



Lecture IV

Patterns for Parallel Software Development



Outline of This Lecture

The Goals:

- 1) Understand a few basic patterns of sequential algorithms
- 2) Know how to map these onto parallel concepts
- 3) Understand how these scale



What is a Pattern?

Software design pattern

General, reusable solution to a commonly occurring problem in a given context in software design

Parallel pattern

Recurring combination of task distribution and data access that solves a specific problem in parallel algorithm design





Serial Control Flow Patterns

- Before starting with parallelism let's look at what we know about the serial case
- We will have a look at the following ones:
 - Sequence
 - Selection
 - Iteration
- These are all simple concepts, but the vocabulary is important!

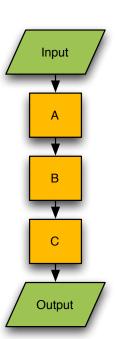


Sequence

- A sequence is an ordered list of tasks/commands to be carried out in a given order
 - The exact dependencies of the commands do not matter
 - Side-effects do not matter
 - There is only one task executed at a time
 - The tasks are executed as defined

Note that

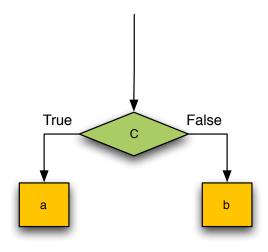
The compiler and the CPU may re-order instructions if they think it optimises runtime





Selection

- In a selection
 - The commands <u>a</u> and <u>b</u> depend on decision of <u>c</u>
 - Always only one of the two sides is being executed



The «if» statement

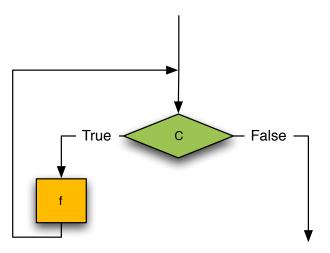
The CPU may apply speculative execution, but it "always" takes care of sanity



Iteration

- In an iteration a certain function <u>f</u> is executed as long as a certain condition c is true.
 - This is the famous while loop

```
while ( c ) {
  f;
}
```





Iteration II

- How do condition and function depend on each other?
 - There must be some dependency, otherwise it is an infinite loop
- Sometimes the dependency is trivial and can be re-formulated as a for loop (a.k.a. counted loop)

```
i = 0;
while ( i < n ) {
   f;
   ++i;
}</pre>
for (i = 0; i < n; ++i) {
   f;
}
```

- In serial code this is mainly just syntactic sugar
 - However, it gives some nice hints to the compiler



Iteration III

- The serial iteration pattern might seem trivially parallelisable but...
 - Beware of dependencies!
- Do multiple iterations depend on each other?
 - Loop-carried dependency
- Different kinds of dependencies translate to different parallelisation possibilities



Iteration IV

```
void doIt( int n, double x[], int a[], int b[], int c[] ) {
  for (int i = 0; i < n; ++i) {
    x[ a[i] ] = x[ a[i] ] * x[ b[i] ] * x[ c[i] ];
  }
}</pre>
```

- Any chance of parallelising this?
- What are the obstacles?
 - i.e. what are the dependencies?



Modern Syntax: An Interlude

- C++ is ever improving with new standards (C++11, C++14, C++17, C++20, C++23)
- Two (not so) recent additions are:
 - auto var = retrieveSomeObject();
 - for (auto & element : myCollection)



- auto: do not specify the type, the compiler finds it out at compile time. Useful to avoid tedious typing also detrimental for readability of the code!
- Range-based loops: build a loop with a concise syntax!

Take advantage of this! ©



Parallel Patterns

- After reminding ourselves about serial control patterns, let's have a look at a few parallel patterns
 - Can help you structure your parallel program
- The serial iteration pattern has many parallel offsprings
 - Map
 - Partition
 - Reduce
 - Scan
- Other useful patterns
 - Pipeline
 - Superscalar Sequences

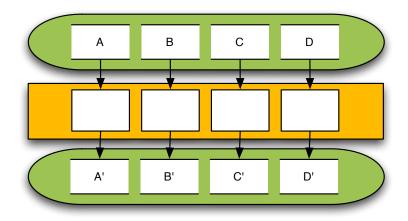
Before implementing yourself, check:

- C++ parallel algorithms library
- Dask
- Spark
- TBB



Map

- The map (or transform / for_each) is the most trivial parallel extension of the serial iteration
 - Apply the same function \underline{f} on all elements of a collection in parallel
 - We hide the loop!
- Requirements:
 - No loop-carried dependency
 - Function \underline{f} is pure, i.e. without side-effects
- Scaling: n (linear w.r.t. the number of elements in the collection)



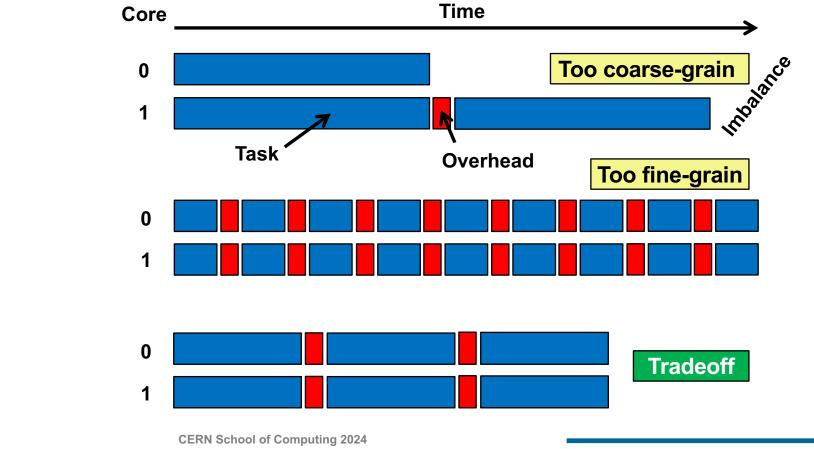


Partition

- The map pattern helps when parallelising on collections
- However, sometimes it is useful to treat multiple items together
 - E.g. for the combination of multithreading and vectorisation
 - Multi-level parallelism!
- Partitioning allows for a custom split of the collection into subcollections or chunks
- A variant of partitioning is called geometric decomposition
 - Update of a partition needs data from other partitions
 - Might require synchronisation



Granularity



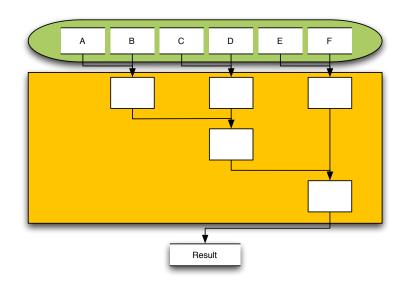


Reduce

- A reduction combines the elements of a collection into a single result using a combiner function
- Requirements:
 - No loop-carried dependency apart from the combined result
 - Combiner function is associative
 - Be careful with floating-point operations!
 - Having a commutative function is beneficial



Reduce II



- Speedup: n / log n
- Counters are a typical example for reduction input
- Before coming to a real example, let's have a look at modern C++ again...



Interlude - Lambdas

- Lambda expressions are anonymous functions and can be assigned to the std::function type
- They can be passed as parameters as if they were regular variables
- When defined, they can capture a specific set of variables (or all)
- Once they have been defined, they can be passed to functions like std::for_each or TBB's parallel_for

```
std::function< double (double, double) >
  f = [ ] (double a, double b) { return a + b; };
std::cout << f ( 23.0, 24.0 );</pre>
```



Interlude - Lambdas II

Using the C++ auto keyword simplifies the syntax, but does not change the behavior auto f = [] (double a, double b) { return a + b; };

Capture the variable globalOffset as a reference and use it in the computation

Capture all variables defined in the current scope by value

```
auto f = [ = ] ( double a, double b )
{ return a + b + globalOffset; };
```

Can you think of the difference in behavior when using capture-by-value instead of capture-by-reference?



Reduce III

 Libraries like Intel's Threading Building Blocks (TBB) or the C++ parallel algorithms library provide already all ingredients for standard patterns like reduce:

```
std::reduce(std::execution::par, // runs parallel reduction
           data.begin(), data.end());
std::reduce(std::execution::par, // runs parallel reduction
           data.begin(), data.end(),
                               // initial element
           1.,
           std::multiplies<>{}
// change reduction operator
           );
```



Map and Reduce Combined

- Usually map and reduce go hand in hand:
 - A function being applied to single elements
 - The results are then passed to a combiner function
- A concrete example:
 - Count the number of times a certain word appears in a text
- Solution:
 - Partition: Split the text in equally-sized chunks
 - Map: Do the word count on each chunk separately
 - Reduce: Add the counts
- Various map/reduce frameworks at your disposal!



The Power of Map-Reduce

- The combination of the Map and Reduce patterns has been extremely successful in massive distributed data processing
- A little bit of history...
 - 2004: Google publishes the MapReduce paper



- 2006: Hadoop is released, inspired by MR
- Nowadays, MR is behind every click on popular web sites or services
 - Facebook, Twitter, Yahoo, ...
 - Analytics to predict user interests, target ads, show recommendations, ... and many more
 - Robust, fault tolerant
 - Scale to crunch large datasets



Map-Reduce and Functional Chains

- Map and reduce were born in functional programming
 - Declare what you want to do, not how
 - No side-effects
- High-level view, based on two main concepts:
 - Data is organised in collections of elements
 - We apply functions to those elements, possibly in a chain

```
histo = events.map(fillHist).reduce(mergeHist)
```



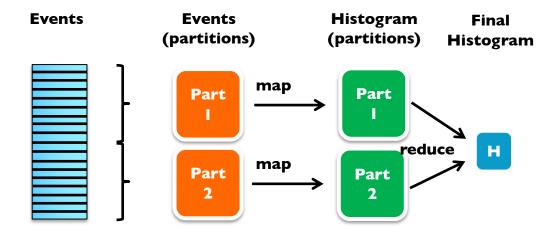


No need to manage parallelisation, just think about opportunities for parallelism!



Map-Reduce and Functional Chains II

- Implementation responsible for producing a parallel execution plan
 - Where are the data?
 - What resources are available?
 - What optimisations can be applied?





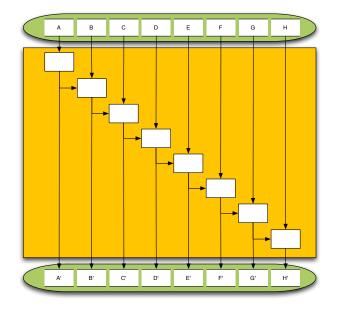
Scan

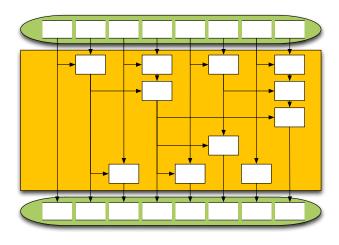
- Scan is another offspring of the iteration pattern with more relaxed boundary conditions
- Requirements:
 - Result of element n depends on n-1
 - Successor function is associative





Scan II





Serial version

Parallel version



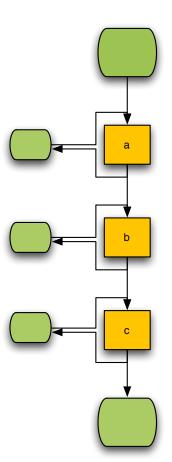
Scan III

- Scan is another offspring of the iteration pattern with more relaxed boundary conditions
- Requirements:
 - Result of element n depends on n-1
 - Successor function is associative
- Already a non-trivial implementation necessary
- Speedup: very limited
 - At most n / log n
 - Number of instructions required is worse (up to x2)



Pipeline

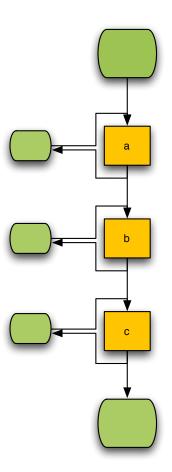
- The pipeline pattern is the good old assembly line
 - Work split into a sequence of operations with a producerconsumer relationship
 - Work items go from one stage to the next
 - The order of steps is important
 - Different operations on different items are independent
 - Stages can be serial or parallel (accept one or more items simultaneously)
- More complex cases can have a directed acyclic graph instead of a purely linear setup
- The speedup of a pipeline is given by Amdahl's Law





Pipeline II

• Intel's TBB offers a feature for implementing a pipeline too:





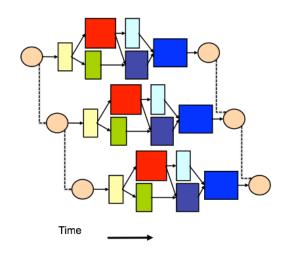
Pipeline III

```
float RootMeanSquare( float* first, float* last, int n ) {
  float sum = 0;
 parallel pipeline (16,
   make filter<void, float*>(
      filter::serial in order,
      [&] (flow control& fc) -> float* {
        if ( first < last ) {</pre>
                                                              Step 1 handles
          return first++;
                                                              the data stream
        } else {
          fc.stop();
          return nullptr;
                                                             Step 2 can run
    make filter<float*, float>(
      filter::parallel,
                                                             in parallel with
      [](float* p) { return (*p)*(*p); }
                                                             itself
   make filter<float, void>(
                                                              Step 3 is not
      filter::serial in order,
      [&] (float x) { sum += x; }
                                                              thread-safe
  return sqrt(sum / n);
```



Superscalar Sequences

- Split work into several tasks and define their data dependencies
- Let a task scheduler do the rest
- Pattern followed by concurrent HEP data processing frameworks



- Assumption of this model is that there are no hidden data dependencies and no side-effects unknown to the scheduler
 - Let's have a look at these assumptions...



Hidden Data Dependencies

```
std::atomic_bool doit(false);

void task1() {
    ...
    if (doit)) {
       eventstore.put(fancystuff);
    }
}
void task2() {
    doit = true;
}
```

- Content of the event store depends on the execution order
- Thread-safe objects don't help at all
- It is a pure logic flaw



Side Effects

- Triggered when a computation modifies some shared state outside of its local environment
 - e.g. a global variable
- They are a major obstacle for parallelism
 - Watch out for them when applying your parallel patterns!
- In general, every non thread-safe resource is an issue
- Remember from previous lectures:
 - Side-effect free resources are the ideal solution
 - If not possible, tell the scheduler about what you need and "reserve" what is unsafe



Take-Away Messages

- There exist design patterns to help you parallelising your programs
 - Check if you can reuse them!
- They all have their origin in serial patterns, but add constraints to the operations allowed
- Map-Reduce is a very successful pattern, used every day for distributed processing of large amounts of data
- High-level features like C++ lambdas, the C++ parallel algorithms library, the TBB library or Spark and Dask make it easier for you to get started with these patterns