

Umeå University
Department of Computing Science

Parallel Programming 7.5 p
5DV152

Exercises, Chapter/Topic 1

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1 Introduction

This report is part of the mandatory coursework. It describes the solutions for several chosen exercises from the course book [1].

2 1.1 - Formulas for block partitioning

The overwhelming idea is to load balance p number of cores with n tasks. Here, we use two functions to obtain block partitioning using a `for` loop:

```
for (my_i = my_first_i; my_i < my_last_i; my_i++)
```

The functions `my_first_i` and `my_last_i` are used to set the limits in the loop. Besides n , i and p we also need an index for the actual core: k . It is understood that indices i and k start at 0. The book text hints to start with the case when n is evenly divisible by p :

```
my_first_i = k * n / p
my_last_i = (k + 1) * n / p
```

Testing this expression for $n = 10$, $p = 5$, $k = \{0, 1, \dots, 4\}$ seems to be correct. Now when n is not even divisible by p , one has to distribute the $n \bmod p$ tasks for example on the first $n \bmod p$ cores:

```
my_first_i = k * n / p + ( k < n mod p ? k : n mod p )
my_last_i = (k + 1) * n / p + (k + 1 < n mod p ? k + 1 : n mod p)
```

Testing this expression for $n = 9$, $p = 5$, $k = \{0, 1, \dots, 4\}$ gives the correct results.

3 1.2 - Modify 1.1 with non-uniform costs

4 1.3 - Tree-structured global sum

5 1.4 - Alternative algorithm for 1.3

6 1.5 - Generalization of 1.3 and 1.4

7 1.6 - Cost analysis of global sum algorithms

References

[1] P.S. Pacheco. *An Introduction to Parallel Programming*. Morgan Kaufman, 2011.