Umeå University

Department of Computing Science

Parallel Programming 7.5 p 5DV152

Exercises, Chapter/Topic 3

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1 Introduction

This report is part of the mandatory coursework. It describes the solutions for several chosen exercises from the course book [3].

2 3.2 - Generalization of algorithm for trapezoidal rule

Two functions to adapt the *trapezoidal rule* for calc_local_a and calc_local_b were written and tested with the source code from the book (*mpi_trap.c*).

```
double calc_local_a(int my_rank, double a, double b, int n, int comm_sz) {
 double local_a = 0;
 double h = 0;
 int local_n = 0;
 int rest_n = 0;
 h = (b-a)/n;
 local_n = n/comm_sz;
 rest_n = n%comm_sz;
 if (my_rank < rest_n) {</pre>
    local_a = a + my_rank*local_n*h + my_rank*h;
  } else {
    local_a = a + my_rank*local_n*h + rest_n*h;
    local_a += (my_rank-rest_n) * h;
  }
 return local_a;
}
double calc_local_b(int my_rank, double a, double b, int n, int comm_sz){
 double h;
 int local_n;
 h = (b-a)/n;
 local_n = n/comm_sz;
 if (my_rank == (comm_sz-1)){
   return a + my_rank+1*local_n*h;
  } else {
    return calc_local_a(my_rank+1, a, b, n, comm_sz);
}
```

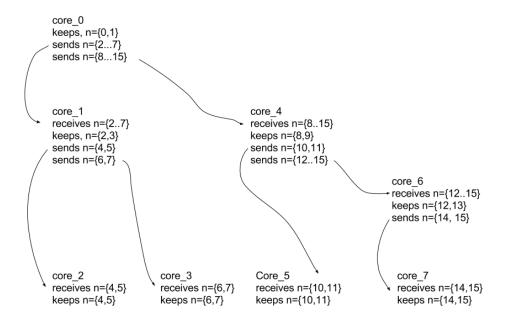


Figure 1: This graph shows a tree based implementation of scatter for comm $_sz = 8$ and n = 16

3 3.6 - Array distributions

Block distribution

Block distribution can be obtained by $b = \lfloor i \div p \rfloor$ where b is the block number, i the index of n and p is the number of processes. This solution is however not fair. An improved, fair expression can be devised using a ternary operator:

$$i < n \mod p \times \lceil n \div p \rceil ? | i \div \lceil n \div p \rceil | : n \mod p + | (i - n \mod p \times \lceil n \div p \rceil) \div | n \mod p | |$$

Cyclic distribution

Cyclic distribution is described by $b = i \mod p$ with b as block number i as index of n and p as number of processes.

Block cyclic distribution

Block cyclic distribution can be expressed as $b = \lfloor i \div l \rfloor mod p$ where b is block index, i index of n, l block length and p number of processes.

4 3.8 - Tree-structured algorithms for scatter and gather

5 3.9 - Vector scaling and dot product

#include <stdio.h>

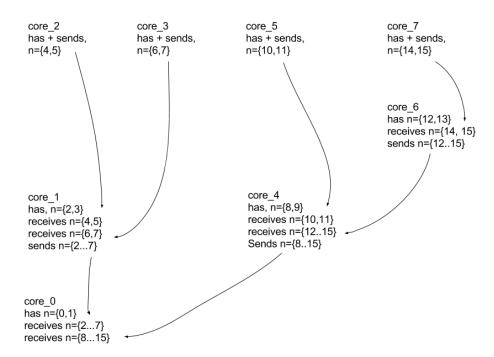


Figure 2: This graph shows a tree based implementation of gather for comm $_sz = 8$ and n = 16.

```
#include <mpi.h>
#include <string.h>
#include <stdlib.h>
#include <time.h>
int main(int argc, char *argv[]) {
 int my_rank, comm_sz;
 int n, local_n, local_dotp_sum = 0, scalar, result_dot;
 int* local_vec1;
 int* local_vec2;
 int* vector1;
 int* vector2;
 /* Initializing */
 MPI_Init(NULL, NULL);
 MPI_Comm_rank(MPI_COMM_WORLD, &my_rank);
 MPI_Comm_size(MPI_COMM_WORLD, &comm_sz);
 srand(time(NULL));
  /* Obtaining Data */
 if (my_rank==0 && argc > 1) {
    if(strcmp(argv[1], "r") == 0){
      printf("using random data, vector length = %d\n", 100*comm_sz);
     n = 100 * comm_sz;
      vector1 = (int *) malloc(100*comm_sz * sizeof(int));
      vector2 = (int *) malloc(100*comm_sz * sizeof(int));
      for (int i = 0; i < n; i++) {
vector1[i] = rand() % 1000;
vector2[i] = rand() % 1000;
     }
     scalar = rand() % 1000;
    }
  } else if (my_rank==0) {
    printf("enter vector length\n");
    scanf("%d", &n);
    printf("enter integer vector 1\n");
    vector1 = (int *) malloc(n * sizeof(int));
    vector2 = (int *) malloc(n * sizeof(int));
    for (int i = 0; i < n; i++) {
     scanf("%d", &vector1[i]);
    }
```

```
printf("enter integer vector 2\n");
  for(int i = 0; i < n; i++){
    scanf("%d", &vector2[i]);
  printf("enter integer scalar\n");
  scanf("%d", &scalar);
/* Distribute Data */
MPI_Bcast(&n, 1, MPI_INT, 0, MPI_COMM_WORLD);
local_n = n/comm_sz;
local_vec1 = (int*) malloc(local_n * sizeof(int));
local_vec2 = (int*) malloc(local_n * sizeof(int));
MPI_Bcast(&scalar, 1, MPI_INT, 0, MPI_COMM_WORLD);
MPI_Scatter(vector1, local_n, MPI_INT, local_vec1, local_n, MPI_INT, 0, MPI_COMM_WORLD)
MPI_Scatter(vector2, local_n, MPI_INT, local_vec2, local_n, MPI_INT, 0, MPI_COMM_WORLD)
/* Calculations */
/* Calculate Dot Product */
for(int i = 0; i < local_n; i++) {</pre>
 local_vec2[i] *=local_vec1[i];
/* Calculate vector-scalar product */
for (int i = 0; i < local n; i++) {
 local_vec1[i] *=scalar;
/* Summing for dot product */
for(int i = 0; i < local_n; i++){
 local_dotp_sum += local_vec2[i];
/* Collect Data */
MPI_Gather(local_vec1, local_n, MPI_INT, vector1, local_n, MPI_INT, 0, MPI_COMM_WORLD);
MPI_Reduce(&local_dotp_sum, &result_dot, 1, MPI_INT, MPI_SUM, 0, MPI_COMM_WORLD);
/* Results */
if(my_rank == 0){
  printf("dot product = %d\n", result_dot);
  printf("vector-scalar product = ");
  for (int i = 0; i < n; i++) {
```

```
printf("%d ", vector1[i]);
}
printf("\n");
}

/* Clean up */
if(my_rank==0) {
  free(vector1);
  free(vector2);
}

free(local_vec1);
free(local_vec2);

MPI_Finalize();

return 0;
}
```

6 3.11 - Prefix sums

Serial Algorithm

Assume vector x of length n and prefix sums vector y to be calculated.

```
y_0 = x_0
for (i = 1, i < n , i++)
y_i = y_{i-1} + n_i
```

Parallel Algorithm

After Blelloch [1], parallel prefix-sum, or 'Sum Scan Algorithm' can be calculated in two steps, the 'Up-sweep' (Reduce) and 'Down-sweep'. Pseudo code from M. Valdez [4] where x is the input data, n the size of the input and d the number of processeors. If $n \neq 2^k$, then n has to be extended with zero's. The 'Up-Sweep':

```
for d = 0 to log2(n) - 1 do
    for all k = 0 to n - 1 by 2^{(d+1)} in parallel do
        x[k + 2^{(d+1)} - 1] = x[k + 2^{d} - 1] + x[k + 2^{(d+1)} - 1]

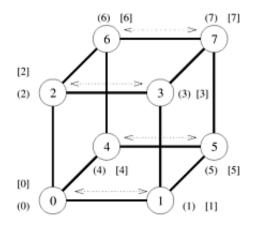
The 'Down-Sweep':

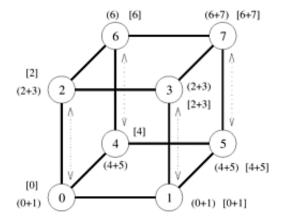
x[n - 1] = 0

for d = log2(n) - 1 down to 0 do
    for all k = 0 to n - 1 by 2^{(d+1)} in parallel do
        t = x[k + 2^{d} - 1]

x[k + 2^{d} - 1] = x[k + 2^{(d+1)} - 1]

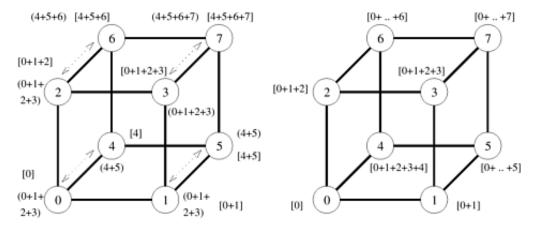
x[k + 2^{(d+1)} - 1] = t + x[k + 2^{(d+1)} - 1]
```





(a) Initial distribution of values

(b) Distribution of sums before second step



- (c) Distribution of sums before third step
- (d) Final distribution of prefix sums

Figure 3: This figure shows the calculation of prefix sums with minimal communication phases. The figure is take from [2, chap 4.3]

Minimum Communication

An algorithm that uses for the prefix-sum of length $n = 2^k$ only k communication steps can be derived from the communication model in a hypercube as shown in figure 3.

Below follows pseudo code for parallel hypercube prefix sum (from [2]):

```
procedure PREFIX_SUMS_HCUBE(my_id, my number, d, result)
begin
    result := my_number;
    msg := result;
    for i := 0 to d - 1 do
        partner := my_id XOR 2i;
        send msg to partner;
        receive number from partner;
        msg := msg + number;
        if (partner < my_id) then result := result + number;
    endfor;</pre>
```

MPI Program using MPI_Scan

A minimal working example using MPI_Scan.

```
#include <stdio.h>
#include <stdlib.h>
#include <mpi.h>
#include <time.h>
int main(void){
 int my_rank, comm_sz;
 int* initial_vector;
 int* prefix_sums;
 MPI_Init(NULL, NULL);
 MPI_Comm_rank(MPI_COMM_WORLD, &my_rank);
 MPI_Comm_size(MPI_COMM_WORLD, &comm_sz);
 srand(time(NULL));
 initial_vector = (int *) malloc(10*sizeof(int));
 prefix_sums = (int *) calloc(10, sizeof(int));
 for (int i = 0; i < 10; i++) {
   initial_vector[i] = rand() % 1000;
 MPI_Scan(initial_vector, prefix_sums, 10,MPI_INT, MPI_SUM, MPI_COMM_WORLD);
 for(int i = 0; i < comm_sz;i++) {
   if (i == my_rank) {
      for (int j = 0; j < 10; j++) {
        printf("%d ", prefix_sums[j]);
      }
    }
  }
 printf("\n");
 MPI_Finalize();
 return 0;
}
```

- 7 3.13 Generalization of vector scaling and dot product
- 8 3.16 Diagram for a butterfly implementation of allgather
- 9 3.18 Derived data types

takes a while to solve requires programming

10 3.20 - Pack and **unpack**

requires programming

11 3.21 - Matrix-vector multiplication

takes a while to solve requires programming requires testing

12 3.22 - Timing the trapezoidal rule

takes a while to solve Requires programming requires testing

13 3.27 - Speedup and efficiencily of odd-even sort

References

- [1] G.E. Blelloch. *Vector Models for Data-Parallel Computing*. MIT Press, Cambridge MA, USA, 1990.
- [2] A Grama, A Gupta, G Karypis, and V Kumar. *Introduction to Parallel Computing*. Pearson, Essex, England, 2003.
- [3] P.S. Pacheco. An Introduction to Parallel Programming. Morgan Kaufman, 2011.
- [4] Stackoverflow, parallel prefix sum fastest implementation. http://stackoverflow.com/questions/10053629/parallel-prefix-sum-fastest-implementation, 2012. accessed: 2017-02-04.