Contents

Chapter 1

Geometry of probability distributions

Quantum states, as we will see, can be thought of as a generalization of probability distributions. In this chapter, we will study probability distributions from a geometrical point of view, and in this framework, we shall prove the Cramer-Rao bound. We will only consider probability distributions defined on finite sample spaces since this is all we need for finite-dimensional pure quantum states.

1.1 Manifolds of probability distributions

1.1.1 Space of probability distributions

Consider a *random process* and the set $\mathcal{X}$ of all its possible outcomes. We call this set the *sample space*, and we will only consider random processes for which it is finite. Then, a *probability distribution* on $\mathcal{X}$ is a function $p \in \mathbb{R}^{\mathcal{X}} := \{f \mid f \colon \mathcal{X} \to \mathbb{R}\}$ which satisfies

$$p(x) \geq 0 \quad \forall x \in \mathcal{X} \quad \text{and} \quad \sum_{x \in \mathcal{X}} p(x) = 1 \tag{1.1}$$

where $p(x)$ represents the probability of the outcome x.

Further, every function $A \in \mathbb{R}^{\mathcal{X}}$ represents a real *random variable*, as it maps every outcome of a random process to a number. Then the *expectation value* of A when the underlying probability distribution is p is expressed by

$$\mathrm{E}_p[A] := \sum_{x \in \mathcal{X}} p(x) A(x) \tag{1.2}$$

Also, given two random variables A, B their *covariance* is

$$\text{Cov}_p[A, B] := \text{E}_p[(A - \text{E}_p[A])(B - \text{E}_p[B])] \tag{1.3}$$

and so the *variance* of a random variable A is

$$\text{V}_p[A] := \text{Cov}_p[A, A] = \text{E}_p[(A - \text{E}_p[A])^2] \tag{1.4}$$

Let now N be the cardinality of $\mathcal{X}$. To have a picture of $\mathbb{R}^{\mathcal{X}}$ we can index the outcomes and consider the natural isomorphism between $\mathbb{R}^{\mathcal{X}}$ and $\mathbb{R}^N$

$$f \leftrightarrow (f(x_1), \ldots, f(x_N)) \tag{1.5}$$

then it's easy to recognize that the *space of probability distributions* is a convex subset of the affine subspace $\mathcal{A}_1 := \{f \in \mathbb{R}^{\mathcal{X}} \mid \sum_{x \in \mathcal{X}} f(x) = 1\}$. In particular, it is the set resulting from the convex mixing of the trivial probability distributions $f_k(x_i) = \delta_{ik}$, represented by the unit vectors of $\mathbb{R}^N$. Finally, it's also interesting to consider the inner product induced on $\mathbb{R}^{\mathcal{X}}$ by the Euclidean one of $\mathbb{R}^N$. Let p be a probability distribution and A a random variable, then

$$p \cdot A = \sum_{x \in \mathcal{X}} p(x) A(x) = \text{E}_p[A] \tag{1.6}$$

1.1.2 Statistical models and manifolds

We call an n-dimensional *statistical model* on $\mathcal{X}$ a family of probability distributions that are globally parametrized by n real-valued variables. Formally this is a subset $\mathcal{S}$ of the space of probability distributions with an invertible function $\psi \colon \mathcal{S} \to \Xi \subseteq \mathbb{R}^n$, so that we may write

$$\mathcal{S} = \left\{ p_\xi \mid \exists \xi = (\xi^{(1)}, \ldots, \xi^{(n)}) \in \Xi \colon p_\xi = \psi^{-1}(\xi) \right\} \tag{1.7}$$

where $p_\xi(x)$ may be equivalently written as $p(x; \xi)$ or $p(x; \xi^{(1)}, \ldots, \xi^{(n)})$. This definition of a statistical model reflects the act of hypothesizing an underlying model, that may depend on some parameters, for the generation of the random variable's samples. Then only a subset, here represented by $\mathcal{S}$, of all the possible probability distributions is considered as a candidate of the underlying probability distribution, and every candidate probability distribution is identified uniquely by the corresponding parameters, here represented by ξ.

We now introduce some additional requirements to statistical models so that we may define well-behaved manifolds from them. Firstly we regard S

as a subset of $\mathcal{A}_1$ equipped with the topology induced by the standard one of $\mathbb{R}^N$. Then we assume that

$$\begin{aligned}&\Xi \text{ is an open set}\\&\psi \text{ is a } C^\infty \text{ diffeomorphism from } \mathcal{S} \text{ to } \Xi\end{aligned} \qquad (1.8)$$

This allows us to differentiate the probability distributions with respect to the parameters so that $\partial_i p(x;\xi)$ is well defined, where we wrote $\partial_i := \frac{\partial}{\partial \xi^{(i)}}$. These conditions also imply that the pair $\mathcal{S}$ and ψ form a chart of $\mathcal{S}$. Then for any another statistical model on $\mathcal{S}$ with parametrization $\psi' \colon \mathcal{S} \to \Xi' \subseteq \mathbb{R}^n$ that follows eq. (1.8), the composed function $\psi' \circ \psi^{-1} \colon \Xi \to \Xi'$ will be a C^∞ diffeomorphism. By considering all the possible parametrizations of this kind we may treat $\mathcal{S}$ as a C^∞-differentiable manifold, where statistical models are the charts and the different parametrizations are the coordinate systems; we call manifolds like these *statistical manifolds*.

From our definitions, it is clear that the maximal dimension of a model is $n = N - 1$ and that every statistical manifold is a submanifold of

$$\mathcal{P} := \{p \in \mathbb{R}^{\mathcal{X}} \mid p(x) > 0 \quad \forall x \in \mathcal{X} \quad \text{and} \quad \sum_{x \in \mathcal{X}} p(x) = 1\} \qquad (1.9)$$

that we call the *manifold of probability distributions*. Notice that $\mathcal{P}$ is the interior of the space of probability distributions, this is because from our definitions follows that every $(N-1)$-dimensional statistical manifold must be an open subset of $\mathcal{A}_1$. (Examples?)

1.1.3 The tangent space and its representations

We will now study tangent vectors of statistical manifolds looking for useful statistical interpretations of them. To do this we will use the fact that, as explained in section 1.1.1, $\mathcal{P}$ can be embedded in the space of random variables $\mathbb{R}^{\mathcal{X}}$. Then we can try to also embed the tangent spaces in $\mathbb{R}^{\mathcal{X}}$ in some meaningful ways, thus linking tangent vectors and random variables.

The mixture representation

Since $\mathcal{P}$ is an open subset of the affine space $\mathcal{A}_1$ we can naturally identify the tangent space at every point with the displacement vector space $\mathcal{A}_0 := \{A \in \mathbb{R}^{\mathcal{X}} \mid \sum_{x \in \mathcal{X}} A(x) = 0\}$. This is the natural embedding of $T_p(\mathcal{P})$ that arises from the trivial embedding of $\mathcal{P}$ in $\mathbb{R}^{\mathcal{X}}$, in fact for any $X \in T_p(\mathcal{P})$ we can define

$$X^{(m)}(x) := X(p(x)) \qquad (1.10)$$

then by considering a parametrization $\{\xi^{(i)}\}$ of $\mathcal{P}$ and its relative coordinate basis $\{\partial_i\}$ we have that

$$\partial_i^{(m)}(x) = \partial_i p(x; \xi) \in \mathcal{A}_0 \tag{1.11}$$

since

$$\sum_{x \in \mathcal{X}} \partial_i p(x; \xi) = \partial_i \sum_{x \in \mathcal{X}} p(x; \xi) = 0 \tag{1.12}$$

Finally, from eq. (1.8) follows that $\partial_i p(x; \xi)$ are $N - 1$ linearly independent functions and thus

$$X^{(m)} \leftrightarrow X \tag{1.13}$$

is an isomorphism and

$$T_p(\mathcal{P}) \sim T_p^{(m)}(\mathcal{P}) := \{X^{(m)} \mid X \in T_p(\mathcal{P})\} = \mathcal{A}_0 \qquad \forall p \in \mathcal{P} \tag{1.14}$$

We call $X^{(m)}$ the *mixture representation* or *m-representation* of X.

The exponential representation

Since $T_p(\mathcal{P})$ is an $(N - 1)$-dimensional vector space, for every p we may try to identify it to the subspace of $\mathbb{R}^{\mathcal{X}}$ orthogonal to p with respect to the inner product defined in eq. (1.6). This is interesting given the statistical meaning of the inner product between a generic element of $\mathbb{R}^{\mathcal{X}}$ and a probability distribution. For every $p \in \mathcal{P}$ the orthogonal space is

$$\mathcal{A}_p^{\perp} := \{A \in \mathbb{R}^{\mathcal{X}} \mid p \cdot A = \mathrm{E}_p[A] = 0\} \tag{1.15}$$

that is the space of random variables with null expectation value when the underlying probability distribution is p.

Now we wish to find a natural isomorphism between $T_p(\mathcal{P})$ and $\mathcal{A}_p^{\perp}$. One way to do this that will prove to be useful is to consider the following alternative embedding of $\mathcal{P}$ in $\mathbb{R}^{\mathcal{X}}$

$$p \mapsto \ln p \in \mathbb{R}^{\mathcal{X}} \tag{1.16}$$

then, for any $X \in T_p(\mathcal{P})$ we can define

$$X^{(e)}(x) := X(\ln p(x)) = \frac{X(p(x))}{p(x)} \tag{1.17}$$

and by considering a parametrization $\{\xi^{(i)}\}$ of $\mathcal{P}$ and its relative coordinate basis $\{\partial_i\}$ we have that

$$\partial_i^{(e)}(x) = \partial_i \ln p(x; \xi) \in \mathcal{A}_p^{\perp} \tag{1.18}$$

since

$$\mathrm{E}_p[\partial_i \ln p(x;\xi)] = \sum_{x \in \mathcal{X}} p(x;\xi) \frac{\partial_i p(x;\xi)}{p(x;\xi)} = \sum_{x \in \mathcal{X}} \partial_i p(x;\xi) = 0 \qquad (1.19)$$

It's easy to prove that the linear independence of $\partial_i \ln p(x;\xi)$ follows from the one of $\partial_i p(x;\xi)$ and thus

$$X^{(e)} \leftrightarrow X \qquad (1.20)$$

is an isomorphism and

$$T_p(\mathcal{P}) \sim T_p^{(e)}(\mathcal{P}) := \{X^{(e)} \mid X \in T_p(\mathcal{P})\} = \mathcal{A}_p^\perp \qquad \forall p \in \mathcal{P} \qquad (1.21)$$

We call $X^{(m)}$ the *exponential representation* or *e-representation* of X.

From their definitions, we have that the two representations of a tangent vector $X \in T_p(\mathcal{P})$ are related as follows

$$X^{(m)}(x) = X^{(e)}(x)p(x) \qquad (1.22)$$

and that while $T_p^{(m)}(\mathcal{P})$ is the same for every p, $T_p^{(e)}(\mathcal{P})$ varies since the set of random variables with null expectation value will differ depending on the underlying probability distribution.

1.2   The information metric

1.2.1   Relative entropy

Given a statistical manifold, we may ask ourselves if a certain metric can be naturally defined on it. Such a metric would give rise to a Riemannian connection and consequently to a geodesic distance between elements of the manifold. For this reason, we should first find a statistical meaning to the notion of distance between probability distributions and only then try to find a metric coherent with it.

One natural way to proceed is to consider how hard it is to distinguish a probability distribution from another one by extracting some samples. More precisely let's assume that a random process has an underlying probability distribution q and that N_s samples are generated. Then we can consider the probability that the resulting frequencies f_i of the samples correspond to the probabilities p_i of another probability distribution p.

For simplicity, let's consider the $N = 2$ case, i.e. the case of binomial distributions. Let $\mathbf{q} = (t, 1 - t)$ and $\mathbf{p} = (r, 1 - r)$ be two probability

distributions. Then if N_s samples are drawn with underlying probability q, the probability $P_{N_s}(\mathbf{p})$ that the obtained frequencies correspond to $\mathbf{p}$ is given by

$$P_{N_s}(\mathbf{p}) = \binom{N_s}{N_s \cdot r} t^{N_s \cdot r}(1-t)^{N_s \cdot (1-r)}$$

then by assuming $r \neq 0, 1$ and using Stirling's formula, we obtain the following asymptotic behavior for $N_s \to \infty$

$$P_{N_s}(\mathbf{p}) \sim \exp\left\{-N_s\left[r\ln\left(\frac{r}{t}\right) + (1-r)\ln\left(\frac{1-r}{1-t}\right)\right]\right\} \qquad (1.23)$$

so the probability decreases exponentially with N_s times the factor in the square parenthesis. This factor only depends on the probability distributions p and q, and one may recognize it from statistics as the *relative entropy* of the two distributions. For a generic finite sample space $\mathcal{X}$, the relative entropy is defined as

$$S(p \parallel q) := \sum_{x \in \mathcal{X}} p(x)\ln\left(\frac{p(x)}{q(x)}\right) \qquad (1.24)$$

More generally it can be shown that the following theorem holds.

Theorem 1 (Sanov's Theorem). *Let $\mathcal{E} \subset \mathcal{P}$ be a closed set of probability distributions without isolated points. Then if N_s samples are drawn with underlying probability distribution $q \in \mathcal{P}$, the probability $P_{N_s}(\mathcal{E})$ that the obtained frequencies correspond to an element in $\mathcal{E}$ has the following asymptotic behavior*

$$P_{N_s}(\mathcal{E}) \sim e^{-N_s S(p_* \parallel q)} \qquad \text{for} \quad N_s \to \infty \qquad (1.25)$$

where p_* is the element of $\mathcal{E}$ for which $S(p_* \parallel q)$ is smallest.

Roughly speaking, this shows that the greater the relative entropy $S(p \parallel q)$ the faster the probability of obtaining frequencies in a small neighborhood of p decreases with the number of samples drawn with underlying probability q. In this view, relative entropy can serve as a kind of distance between probability distributions, but with some caveats. From the definition in eq. (1.24) relative entropy has the following properties

$$S(p \parallel q) \geq 0 \quad \forall p, q \in \mathcal{P} \qquad (1.26)$$
$$S(p \parallel q) = 0 \iff p = q \qquad (1.27)$$

but it is not symmetric and it doesn't follow the triangle inequality, so it is not a metric distance. We may ask ourselves if the asymmetry is an accident of our definition of relative entropy or if it is inherent in the distinguishability of probability distributions. The latter turns out to be true, as shown by the following example.

Example. *Consider two coins, one fair and one with heads on both sides. We want to pick one and guess which one it is just by tossing it multiple times. Clearly, the game is not symmetric in the choice of the coin; in fact, if we pick the fair coin the first time we will get a tail we will be sure that we picked the fair one, while if we pick the double-head one we will only get heads but this result will always be also compatible with a fair coin that, by chance, is only giving heads.*

This game is precisely a problem of distinguishability of probability distributions. In fact, we have the $N = 2$ sample space and two probability distributions: $\mathbf{p} = (0.5, 0.5)$ (the fair coin) and $\mathbf{q} = (1, 0)$ (the double-head coin), so the two relative entropies are $S(p \parallel q) \to \infty$ and $S(q \parallel p) = \ln 2$. If we pick the fair coin for large N_s the obtained frequencies will approach $\mathbf{p}$; then we consider the probability of obtaining these frequencies if the underlying distribution was q. This probability is identically 0, and it is coherent with Sanov's theorem since $S(p \parallel q) \to \infty$ and so $P_{N_s}(p) \sim e^{-N_s \cdot \infty} = 0$. Otherwise, if we pick the double-head coin the frequencies will always match exactly $\mathbf{q}$, and the probability of getting this if the underlying distribution was q is 0.5^{N_s}. This is coherent with Sanov's theorem since $S(q \parallel p) = \ln 2$ and so $P_{N_s}(q) \sim e^{-N_s \cdot \ln 2} = 0.5^{N_s}$.

Even though relative entropy doesn't follow triangle inequality, it does follow a generalization of the Pythagorean theorem stated as follows

Theorem 2 (Generalized Pythagorean theorem). *Let $\mathcal{E} \subset \mathcal{P}$ be a convex set and consider $p \in \mathcal{E}$ and $q \in \mathcal{P} \setminus \mathcal{E}$. Then*

$$S(p \parallel q) \geq S(p \parallel p_*) + S(p_* \parallel q) \tag{1.28}$$

where p_* is the element of $\partial \mathcal{E}$ for which $S(p_* \parallel q)$ is smallest.

This is a generalization of the Pythagorean theorem in the sense that if it was stated in terms of the Euclidean distance squared, the angle between $\overline{pp_*}$ and $\overline{p_*q}$ would be obtuse and so eq. (1.28) would be the corollary of the Pythagorean theorem for obtuse triangles. This suggests that the relative entropy may be regarded as an asymmetric distance squared and as we will see this is enough to define a metric on the manifold.

1.2.2 Squared Riemannian distances

Now that we have some notion of distance, we shall explore how to define a coherent metric on the manifold. First, we shall study this for a Riemannian distance.

Let (M, g) be a Riemannian manifold where g is the metric tensor, then consider a point $p \in M$. We define the *exponential map* in p as follows

$$\text{Exp}_p : T_p M \to M, \quad \text{Exp}_p(\mathbf{v}) := \gamma_{\mathbf{v}}(1) \tag{1.29}$$

where $\mathbf{v} \in T_p$ and $\gamma_{\mathbf{v}} : [0, 1] \to M$ is the unique geodesic tangent to v in p, i.e. satisfying $\gamma_{\mathbf{v}}(0) = p$ and $\gamma_{\mathbf{v}}'(0) = \mathbf{v}$. In general, this map will be well-defined only from a neighborhood of the origin of T_p to a neighborhood of p, since only locally the uniqueness of the geodesic curve is guaranteed. By eventually further restricting the neighborhood, this map will also be 1-1 since locally the geodesic curves don't cross.

Then in this neighborhood, we have the inverse of the exponential map that maps $q \mapsto \mathbf{v}_q \in T_p M$ so that $\gamma_{\mathbf{v}_q}(1) = q$. Since $\gamma_{\mathbf{v}_q}$ is the only geodesic connecting p and q, the geodesic distance between them will be

$$L(p, q) = \int_0^1 \sqrt{g_{\gamma_{\mathbf{v}_q}(\lambda)}(\gamma_{\mathbf{v}_q}'(\lambda), \gamma_{\mathbf{v}_q}'(\lambda))} d\lambda \tag{1.30}$$

Because $\gamma_{\mathbf{v}_q}$ is a geodesic, by definition we have that $\gamma_{\mathbf{v}_q}'(\lambda)$ is parallel transported along the curve and so

$$g_{\gamma_{\mathbf{v}_q}(\lambda)}(\gamma_{\mathbf{v}_q}'(\lambda), \gamma_{\mathbf{v}_q}'(\lambda)) = g_p(\mathbf{v}_q, \mathbf{v}_q) \quad \forall \lambda \in [0, 1] \tag{1.31}$$

Finally then, we get

$$L(p, q) = \int_0^1 \sqrt{g_p(\mathbf{v}_q, \mathbf{v}_q)} d\lambda = \sqrt{g_p(\mathbf{v}_q, \mathbf{v}_q)} = \|\mathbf{v}_q\| \tag{1.32}$$

and so $\hat{\mathbf{v}}_q = \mathbf{v}_q \backslash L(p, q)$.

Now chose a vector $\mathbf{dp} \in T_p M$ and let $q = \text{Exp}_p(\mathbf{dp})$. Then, let $\Gamma_{\hat{\mathbf{dp}}}$ be the unique geodesic curve such that $\Gamma_{\hat{\mathbf{dp}}}(0) = p$ and $\Gamma_{\hat{\mathbf{dp}}}'(0) = \hat{\mathbf{dp}}$. Clearly, this is the following reparametrization of $\gamma_{\mathbf{dp}}$

$$\Gamma_{\hat{\mathbf{dp}}}(l) = \gamma_{\mathbf{dp}}\left(\frac{l}{\|\mathbf{dp}\|}\right) = \gamma_{\mathbf{dp}}\left(\frac{l}{L(p, q)}\right) \tag{1.33}$$

and so $q = \Gamma_{\hat{\mathbf{dp}}}(L(p, q))$.

Now let $\{x^{(i)}\}$ be a coordinate system for the neighborhood; from the Taylor expansion of $\Gamma_{\hat{\mathbf{dp}}}^{(i)}(l)$ in $l = 0$ we get

$$\Gamma_{\hat{\mathbf{dp}}}^{(i)}(l) = \Gamma_{\hat{\mathbf{dp}}}^{(i)}(0) + \left.\frac{d\Gamma_{\hat{\mathbf{dp}}}^{(i)}(l)}{dl}\right|_{l=0} \cdot l + O(l^2) \tag{1.34}$$

then, from $\Gamma_{\hat{\mathbf{dp}}}(0) = p$, $\Gamma'_{\hat{\mathbf{dp}}}(0) = \hat{\mathbf{dp}}$ and $q = \Gamma_{\hat{\mathbf{dp}}}(L(p,q))$ we get that

$$q^{(i)} = p^{(i)} + \hat{dp}^i \, L(p,q) + O(L^2(p,q)) = p^{(i)} + dp^i + O(\|\mathbf{dp}\|^2) \qquad (1.35)$$

for every q in the image of $\Gamma_{\hat{\mathbf{dp}}}$. Then, combining eq. (1.32) and eq. (1.35) we get

$$L^2(\{p^{(i)}\}, \{p^{(i)} + dp^i\}) \to L^2(p,q) = g_{ij} dp^i dp^j \quad \text{for} \quad dp^i \to 0 \qquad (1.36)$$

We define $L_p^2 : M \to [0, +\infty)$, $L_p^2(q) \coloneqq L^2(p,q)$ and consider its Taylor expansion in $q = p$

$$L_p^2(\{p^{(i)} + dp^i\}) = \frac{1}{2} \Big[\partial_i \partial_j L_p^2(q) \Big]_{q=p} dp^i dp^j + O(dp^3) \qquad (1.37)$$

where the first derivative vanishes because p is a minimum of L_p^2. Then we get

$$g_{ij}^{(p)} = \frac{1}{2} \Big[\partial_i \partial_j L_p^2(q) \Big]_{q=p} \qquad (1.38)$$

so the metric tensor in p is proportional to the Hessian matrix of L_p^2 in p. This is possible because the Hessian matrix of any C^2 function of a differentiable manifold evaluated in a critical point is a $\binom{0}{2}$ tensor.

We have thus found a way to recover the metric tensor from its squared Riemannian distance.

1.2.3   Divergences

As argued in the last paragraph of section 1.2.1, relative entropy has some properties of squared distances and so we may try to find a metric tensor coherent with it as we did in section 1.2.2.

Let M be a differentiable manifold and $D(\,\cdot\,\|\,\cdot\,) : M \times M \to [0, +\infty)$ a C^2 function (possibly asymmetric) satisfying

$$D(p \,\|\, q) \geq 0 \quad \text{and} \quad D(p \,\|\, q) = 0 \iff p = q \qquad \forall p, q \in M \qquad (1.39)$$

Then, given a coordinate system $\{\xi^{(i)}\}$ on M we have that every pair of points $(q, q') \in M \times M$ has coordinates $(\{\xi^{(i)}\}, \{\xi^{(i)'}\})$ and we use the following notation for partial derivatives in one of the two terms on the diagonal (p, p)

$$D[\,\partial_i \,\|\, \cdot\,] \colon p \mapsto [\partial_i D(q \,\|\, q')]_{(q,q')=(p,p)}$$
$$D[\,\cdot\, \|\, \partial_i\,] \colon p \mapsto [\partial_i' D(q \,\|\, q')]_{(q,q')=(p,p)}$$

From the fact that the diagonal (p, p) is a constant surface of minima of D follows that

$$D[\partial_i \parallel \cdot] = D[\cdot \parallel \partial_i] \equiv 0 \qquad (1.40)$$

so the diagonal is also a constant surface of the derivatives of D. Then by further deriving parallel to the diagonal, we get

$$(\partial_i + \partial_i')D[\cdot \parallel \partial_j] = D[\cdot \parallel \partial_i\partial_j] + D[\partial_i \parallel \partial_j] \equiv 0$$
$$(\partial_j + \partial_j')D[\partial_i \parallel \cdot] = D[\partial_i\partial_j \parallel \cdot] + D[\partial_i \parallel \partial_j] \equiv 0$$

where we used the fact that since D is C^2 we can swap second-order derivatives. Finally, we get

$$D[\partial_i\partial_j \parallel \cdot] = D[\cdot \parallel \partial_i\partial_j] = -D[\partial_i \parallel \partial_j] =: g_{ij}^{(D)} \qquad (1.41)$$

where from the fact that the diagonal is a surface of minima, it follows that the previous expression defines a (symmetric) positive semi-definite tensor. From eq. (1.40) and eq. (1.41) it follows that denoting

$$D_p^{(R)} : q \mapsto D(p \parallel q) \quad \text{and} \quad D_p^{(L)} : q \mapsto D(q \parallel p)$$

to second order, we get

$$D_p^{(R)}(q) = \frac{1}{2}g_{ij}^{(D)}d\xi^i d\xi^j + O(d\xi^2) \quad \text{and} \quad D_p^{(L)}(q) = \frac{1}{2}g_{ij}^{(D)}d\xi^i d\xi^j + O(d\xi^2)$$
$$\text{where } d\xi^i := \xi_p^{(i)} - \xi_q^{(i)}$$

and so to the lowest order, the asymmetry is not present.

Finally, if $g_{ij}^{(D)}$ is positive definite we say that D is a *divergence*, then $\frac{1}{2}g_{ij}^{(D)}$ defines a metric tensor and so a unique Riemannian structure on the manifold. The induced squared Riemannian distance coincides at the lowest order near the diagonal with the divergence.

1.2.4 The Fisher information metric

We now go back to the relative entropy, we report the definition given in section 1.2.1

$$S(p \parallel q) := \sum_{x \in \mathcal{X}} p(x) \ln\left(\frac{p(x)}{q(x)}\right)$$

this is a C^2 function of $\mathcal{P} \times \mathcal{P}$ and it follows eq. (1.39). So for any model $\mathcal{S}$ with parameters $\{\xi^{(i)}\}$ we can define

$$g_{ij}^{(S)} = S[\cdot \parallel \partial_i\partial_j] = \partial_i'\partial_j' \sum_{x \in \mathcal{X}} p(x;\xi) \ln\left(\frac{p(x;\xi)}{q(x;\xi')}\right)\Bigg|_{\xi=\xi'} \qquad (1.42)$$

And so we have

$$g_{ij}^{(S)} = -\sum_{x \in \mathcal{X}} p(x) \partial_i \partial_j \ln p(x; \xi) = -\mathrm{E}_p[\partial_i \partial_j \ln p(x; \xi)] \qquad (1.43)$$

and equivalently

$$g_{ij}^{(S)} = \mathrm{E}_p\left[\frac{1}{p^2(x)} \partial_i p(x; \xi) \partial_j p(x; \xi)\right] = \mathrm{E}_p[\partial_i \ln p(x; \xi) \partial_j \ln p(x; \xi)] \qquad (1.44)$$

$$= \sum_{x \in \mathcal{X}} \frac{\partial_i p(x; \xi) \partial_j p(x; \xi)}{p(x)} \qquad (1.45)$$

where we used the fact that

$$\sum_{x \in \mathcal{X}} \partial_i \partial_j p(x; \xi) = \partial_i \partial_j \sum_{x \in \mathcal{X}} p(x; \xi) = 0 \qquad (1.46)$$

Since we know that any $g_{ij}^{(D)}$ is positive semi-definite, $g_{ij}^{(S)}$ will be positive definite if and only if it is invertible. It can be easily shown that if the functions $\partial_i p(x; \xi)$ are linearly independent, then $g_{ij}^{(S)}$ is invertible and thus positive definite (proof?) . Thus, relative entropy is a divergence and in fact, it is also known as *Kullback-Leibler divergence* or *information divergence*.

Finally, then, $G_F := \{g_{ij}^{(S)}\}$ defines a metric tensor at every point and it is known as the *Fisher information metric*. This is, up to a constant factor, the unique metric induced by the relative entropy, and it plays a focal role in the geometrical modeling and interpretation of statistics.

1.2.5 The geometry of $\mathcal{P}$

The Fisher metric defines the inner product $\langle \cdot , \cdot \rangle_p$ between tangent vectors of a point $p \in \mathcal{P}$. Let us now express this inner product through the representations we defined in section 1.1.3. Given two tangent vectors $X, Y \in T_p\mathcal{P}$ from eq. (1.44) we find that

$$\langle X, Y \rangle_p = \mathrm{E}_p[X^{(e)} Y^{(e)}] \qquad (1.47)$$

while from eq. (1.45) we get

$$\langle X, Y \rangle_p = \sum_{x \in \mathcal{X}} \frac{X^{(m)}(x) Y^{(m)}(x)}{p(x)} \qquad (1.48)$$

$$= \sum_{x \in \mathcal{X}} X^{(m)}(x) Y^{(e)}(x) = \sum_{x \in \mathcal{X}} X^{(e)}(x) Y^{(m)}(x) \qquad (1.49)$$

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These expressions will prove to be very useful in section 1.3.

We notice that in neither representation the inner product is the Euclidean one induced by $\mathbb{R}^{\mathcal{X}}$ on the respective embeddings. For such a representation we would have that

$$\langle X, Y \rangle_p = \sum_{x \in \mathcal{X}} X^{(0)}(x) Y^{(0)}(x) \tag{1.50}$$

then it's easy to guess that the embedding

$$p \mapsto 2\sqrt{p} =: p_{(0)} \tag{1.51}$$

is the one whose representation of the tangent spaces

$$X^{(0)} := X(2\sqrt{p}) = \frac{X(p)}{\sqrt{p}}, \qquad X \in T_p(\mathcal{P}) \tag{1.52}$$

follows eq. (1.50). This means that the information geometry of $\mathcal{P}$, i.e. the geometry induced on it by the Fisher metric, is that of an N-dimensional round sphere (of radius 2) since

$$\sum_{x \in \mathcal{X}} p(x) = 1 \implies \sum_{x \in \mathcal{X}} p_{(0)}^2(x) = 4 \tag{1.53}$$

(Talk about e and m connections?)

1.3 Parameter estimation

1.3.1 Unbiased estimators

Consider a random process and an n-dim statistical model $\mathcal{S} = \{p_\xi \mid \xi \in \Xi\}$ of it, as defined in section 1.1.2; it is often the case that from a measured sample $x \in \mathcal{X}$ we want to estimate the parameters ξ of the underlying probability distribution, that we assume to be in $\mathcal{S}$.

The estimation is represented by a function

$$\hat{\xi} = (\hat{\xi}^{(1)}, \ldots, \hat{\xi}^{(n)}) \colon \mathcal{X} \to \Xi \subseteq \mathbb{R}^n \tag{1.54}$$

that we call *estimator*. Each component $\hat{\xi}^{(i)}$ is a random variable, and we say that $\hat{\xi}$ is an *unbiased estimator* if

$$\mathrm{E}_{p_\xi}\left[\hat{\xi}\right] = \left(\mathrm{E}_{p_\xi}\left[\hat{\xi}^{(1)}\right], \ldots, \mathrm{E}_{p_\xi}\left[\hat{\xi}^{(n)}\right]\right) = \xi \qquad \forall \xi \in \Xi \tag{1.55}$$

i.e. if for each $p_\xi \in \mathcal{S}$ the expectation value of the estimator is the correct parameter ξ.

Then for an unbiased estimator, we may represent the deviation from the true parameters with the variance-covariance matrix of the estimator $V_\xi \left[\hat{\xi} \right] := \left\{ v_\xi^{ij} \right\}$ where

$$v_\xi^{ij} := \operatorname{Cov}_{p_\xi} \left[\hat{\xi}^{(i)}, \hat{\xi}^{(j)} \right] = \operatorname{E}_{p_\xi} \left[\left(\hat{\xi}^{(i)}(x) - \xi^i \right) \left(\hat{\xi}^{(j)}(x) - \xi^j \right) \right] \qquad (1.56)$$

In particular, the elements on the diagonal are the variances of the components of the estimator. (Talk about the covariance ellipse?)

1.3.2 Variance and expectation value

For a generic random variable $A \in \mathbb{R}^{\mathcal{X}}$ we may define a real function $\operatorname{E}[A]$ on $\mathcal{P}$ that maps every probability distribution p to the expectation value of A when the sample is generated with p as underlying probability distribution

$$\operatorname{E}[A] \colon \mathcal{P} \to \mathbb{R} \qquad p \mapsto \operatorname{E}_p[A] \qquad (1.57)$$

Since this is a function of the manifold $\mathcal{P}$, at every point p we may consider its differential $(d\operatorname{E}[A])_p$. This is the element of the cotangent space $T_p^*(\mathcal{P})$ such that for any tangent vector $X \in T_p(\mathcal{P})$ we have

$$(d\operatorname{E}[A])_p(X) = X(\operatorname{E}[A]) \qquad (1.58)$$

Also, because a metric is defined on $\mathcal{P}$ we have a natural isomorphism between tangent and cotangent vectors, thus the gradient of $\operatorname{E}[A]$ is the tangent vector defined by

$$\langle (grad\operatorname{E}[A])_p, X \rangle_p = (d\operatorname{E}[A])_p(X) = X(\operatorname{E}[A]) \qquad \forall X \in T_p(\mathcal{P}) \qquad (1.59)$$

We now state and prove the following theorem that relates the deviation of a random variable to the gradient of its expectation value

Theorem 3. *For any random variable $A \in \mathbb{R}^{\mathcal{X}}$ we have that*

$$(grad\operatorname{E}[A])_p^{(e)} = A - \operatorname{E}_p[A] \qquad \forall p \in \mathcal{P} \qquad (1.60)$$

where the gradient is the dual tangent vector of the differential with respect to the Fisher metric.

Proof. For every $X \in T_p$ we have

$$X(\mathrm{E}[A]) = \sum_{x \in \mathcal{X}} X(p(x))A(x) = \sum_{x \in \mathcal{X}} X^{(m)}(x)A(x)$$
$$= \mathrm{E}_p[X^{(e)}A] = \mathrm{E}_p[X^{(e)}(A - \mathrm{E}_p[A])]$$

where in the last equation we used the fact that $\mathrm{E}_p[X^{(e)}] = 0$. We notice that

$$\mathrm{E}_p[A - \mathrm{E}_p[A]] = 0 \implies (A - \mathrm{E}_p[A]) \in T_p^{(e)}(\mathcal{P})$$

and so there must exist a tangent vector $Y \in T_p(\mathcal{P})$ such that $Y^{(e)} = (A - \mathrm{E}_p[A])$. Then

$$X(\mathrm{E}[A]) = \mathrm{E}_p[X^{(e)}Y^{(e)}] = \langle X, Y \rangle_p$$

and so from eq. (1.59) we have that $Y = (grad\mathrm{E}[A])_p$ and so

$$(grad\mathrm{E}[A])_p^{(e)} = A - \mathrm{E}_p[A]$$

$\square$

We also get the following

Corollary 3.1. *For any random variable A*

$$\mathrm{V}_p[A] = \|(d\mathrm{E}[A])_p\|_p^2 \tag{1.61}$$

Proof. Follows from theorem 3 noticing that

$$\|(d\mathrm{E}[A])_p\|_p^2 = \langle (grad\mathrm{E}[A])_p, (grad\mathrm{E}[A])_p \rangle_p$$
$$= \mathrm{E}_p[(A - \mathrm{E}_p[A])^2]$$

$\square$

Let now $\mathcal{S}$ be an n-dim statistical manifold, since it is a submanifold of $\mathcal{P}$ we have that $T_p(\mathcal{S})$ is a linear subspace of $T_p(\mathcal{P})$ for every $p \in \mathcal{S}$. Then the gradient of the restriction on $\mathcal{S}$ of a function of $\mathcal{P}$ is its orthogonal projection on $T_p(\mathcal{S})$. In particular, we have that

$$T_p(\mathcal{P}) = T_p(\mathcal{S}) \oplus T_p(\mathcal{S})^\perp \tag{1.62}$$

and so we may uniquely decompose

$$(grad\mathrm{E}[A])_p = v_\| + v_\perp \qquad v_\| \in T_p(\mathcal{S}), v_\perp \in T_p(\mathcal{S})^\perp \tag{1.63}$$

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then we find that $\forall X \in T_p(\mathcal{S})$

$$\langle (grad\mathrm{E}[A]|_{\mathcal{S}})_p, X \rangle_p = X(\mathrm{E}[A]) = \langle v_\parallel + v_\perp, X \rangle_p = \langle v_\parallel, X \rangle_p \qquad (1.64)$$

and so $(grad\mathrm{E}[A]|_{\mathcal{S}})_p = v_\parallel$.

Finally, we have the following theorem relating the variance of a random variable and the sensitivity of its expectation value to the changes in the model parameters

Theorem 4. *Given a statistical manifold S, for any random variable A we have that*

$$\mathrm{V}_p[A] \geq \|(d\mathrm{E}[A]|_{\mathcal{S}})_p\|_p^2 \qquad (1.65)$$

where the equality holds if and only if

$$A - \mathrm{E}_p[A] \in T_p^{(e)}(\mathcal{S}) \qquad (1.66)$$

Proof. Follows immediately from corollary 3.1 and eq. (1.64) $\qquad \square$

1.3.3 The Cramér-Rao bound

We are now in the position to state and prove an important result of parameter estimation theory

Theorem 5 (Cramér-Rao bound). *Let $\mathcal{S} = \{p_\xi \mid \xi \in \Xi\}$ be an n-dim statistical model of $\mathcal{P}$. Then, for any unbiased estimator $\hat{\xi}$ the variance-covariance matrix $\mathrm{V}_\xi[\hat{\xi}]$ satisfies*

$$\mathrm{V}_\xi[\hat{\xi}] \geq G_F^{-1}(p_\xi) \qquad (1.67)$$

in the sense that $\mathrm{V}_\xi[\hat{\xi}] - G_F^{-1}(p_\xi)$ is positive semi-definite.

Proof. Let $A = c_i \hat{\xi}^{(i)}$ where c is an arbitrary element of $\mathbb{R}^n$. Then A is a random variable with $\mathrm{E}_{p_\xi}[A] = c_i \xi^i$ and

$$\begin{aligned}
\mathrm{V}_{p_\xi}[A] &= \mathrm{E}_{p_\xi}[(c_i \hat{\xi}^{(i)}(x) - c_i \xi^i)(c_j \hat{\xi}^{(j)}(x) - c_j \xi^j)] \\
&= \mathrm{E}_{p_\xi}[c_i(\hat{\xi}^{(i)}(x) - \xi^i)(\hat{\xi}^{(j)}(x) - \xi^j)c_j] \\
&= c_i v_\xi^{ij} c_j
\end{aligned}$$

Then letting $p = p_\xi$ from theorem 4 we get

$$c_i v_\xi^{ij} c_j \geq \|(d\mathrm{E}[A]|_{\mathcal{S}})_p\|_p^2$$

where, in the coordinate basis of $\{\xi^{(i)}\}$

$$\|(d\mathrm{E}[A]|_{\mathcal{S}})_p\|_p^2 = (\partial_i\mathrm{E}[A])_p g^{ij}(p)(\partial_j\mathrm{E}[A])_p$$
$$= c_i g^{ij}(p) c_j$$

and so, finally

$$c_i(v_\xi^{ij} - g^{ij}(p_\xi))c_j \geq 0$$

$\square$

An unbiased estimator that saturates eq. (1.67) is called an *efficient estimator* and is the best unbiased estimator in the sense that its variance is minimum between all unbiased estimators. It's important to notice that an efficient estimator doesn't always exist. (talk about asymptotically efficient estimators?)

This result shows that the efficiency with which we can infer the underlying probability distribution of a process is deeply linked with the information geometry of the model. (Examples)

Chapter 2

Geometry of quantum states