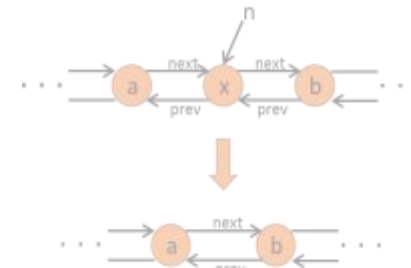
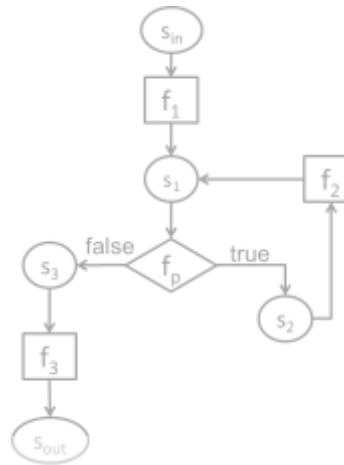


$$\exists c \forall in \ Q(c, in)$$

```

/* Average of x and y without using x+y (avoid overflow)*/
int avg(int x, int y){
  int t = expr({x/2, y/2, x%2, y%2, 2 }, {PLUS, DIV});
  assert t == (x+y)/2;
  return t;
}

```

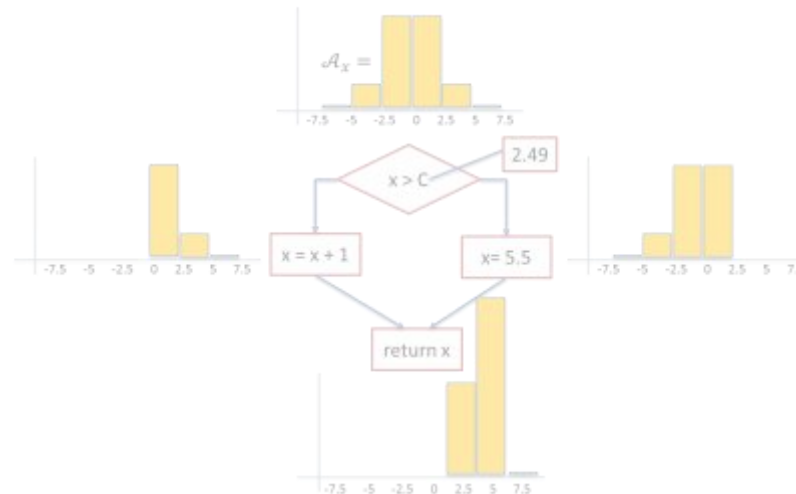
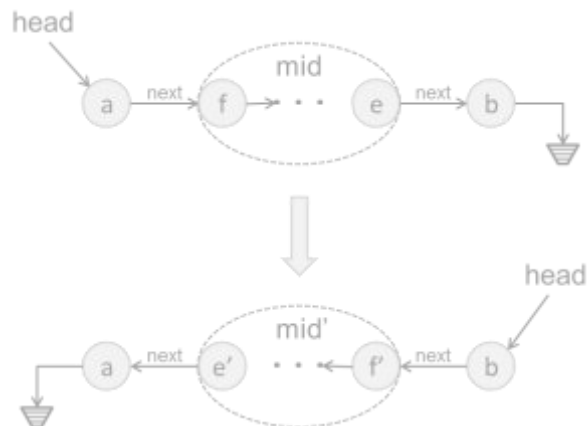


```

{
  s = n.succ;
  p = n.pred;
  p.succ = s;
  s.pred = p;
}

```

# Module II: Synthesizing Complex Programs



$$\varphi(p)$$

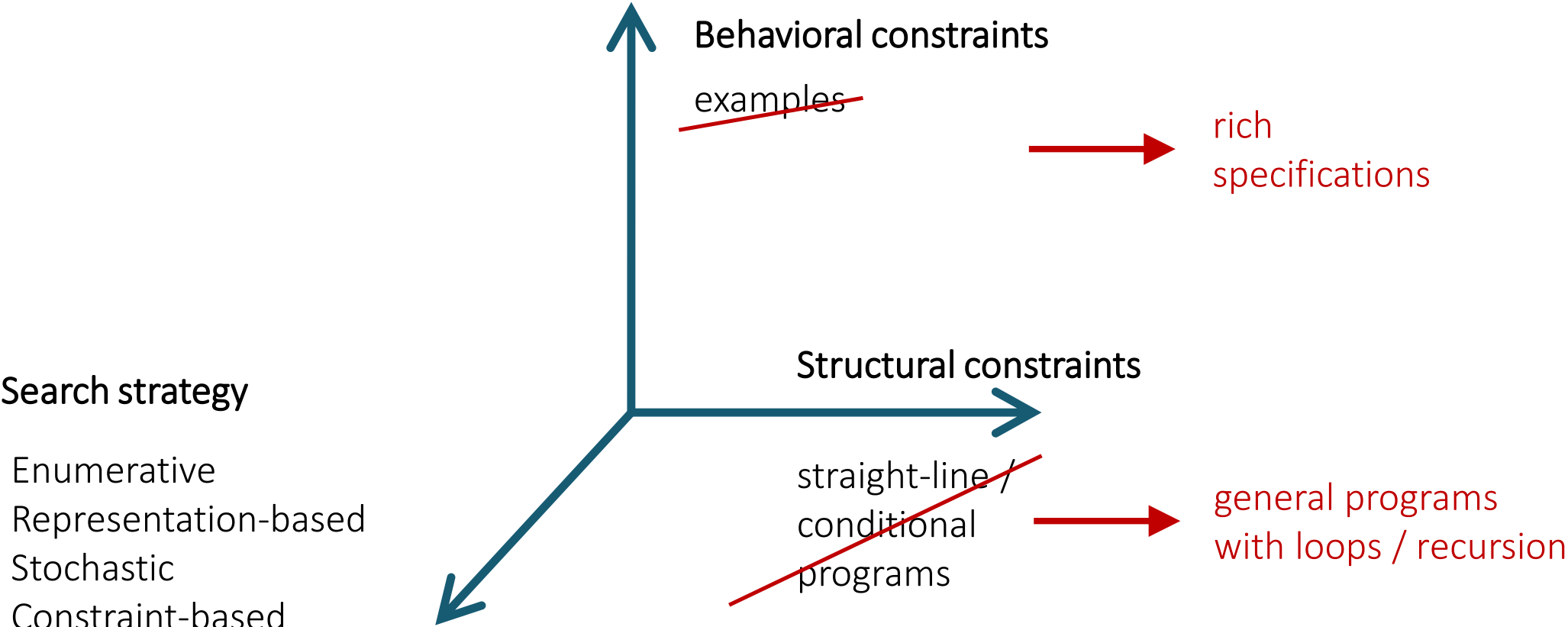
$$Sk[c](in)$$

# Lecture 9

## Specifications

# Module I vs Module II

---



# Examples of rich specifications

---

Reference implementation

Assertions

Pre- and post-condition

Refinement type

# Reference Implementation

---

Easy to compute the result, but hard to compute it efficiently or under structural constraints

```
bit[W] AES_round (bit[W] in, bit[W] rkey)
{
    ... // Transcribe NIST standard
}

bit[W] AES_round _sk (bit[W] in, bit[W] rkey) implements AES_round
{
    ... // Sketch for table lookup
}
```

# Assertions

---

Hard to compute the result, but easy to check its desired properties

```
split_seconds (int totsec) {  
    int h := ??;  
    int m := ??;  
    int s := ??;  
    assert totsec == h*3600 + m*60 + s;  
    assert 0 <= h && 0 <= m < 60 && 0 <= s < 60;  
}
```

# Pre-/post-conditions

---

Hard to compute the result but easy to express its properties in logic

```
sort (int[] in, int n) returns (int[] out)
    requires   $n \geq 0$ 
    ensures   $\forall i\ j. 0 \leq i < j < n \Rightarrow out[i] \leq out[j]$ 
               $\forall i. 0 \leq i < n \Rightarrow \exists j. 0 \leq j < n \wedge in[i] = out[j]$ 
{
    ??
}
```

# Refinement types

Same as pre-/post-conditions but logic goes inside the types

```
data RBT a where
  Empty :: RBT a
  Node  :: x: a ->
    black: Bool ->
    left:  { RBT {a | _v < x} | !black ==> isBlack _v} ->
    right: { RBT {a | x < _v} | (!black ==> isBlack _v) &&
(blackHeight _v == blackHeight left)} ->
  RBT a
```

binary search tree

red nodes have black children

same number of black nodes on every path to leaves

```
insert :: x: a -> t: RBT a -> {RBT a | elems _v == elems t + [x]}
insert = ??
```



# Why go beyond examples?

---

Might need too many

- **Example:** Myth needs 12 for `insert_sorted`, 24 for `list_nth`
- Examples contain *too little* information
- Successful tools use domain-specific ranking

Output difficult to construct

- **Example:** AES cypher, RBT
- Examples also contain *too much* information (concrete outputs)

Need strong guarantees

- **Example:** AES cypher

Reasoning about non-functional properties

- **Example:** security protocols

# Why is this hard?

---

gcd (**int** a, **int** b) **returns** (**int** c)

**requires**  $a > 0 \wedge b > 0$

**ensures**  $a \% c = 0 \wedge b \% c = 0$

$\forall d . c < d \Rightarrow a \% d \neq 0 \vee b \% d \neq 0$

{

**int** x , y := a, b;

**while** (x != y) {

**if** (x > y) x := ??;

**else** y := ??;

}}

infinitely many inputs

cannot validate by testing

infinitely many paths!

hard to generate constraints

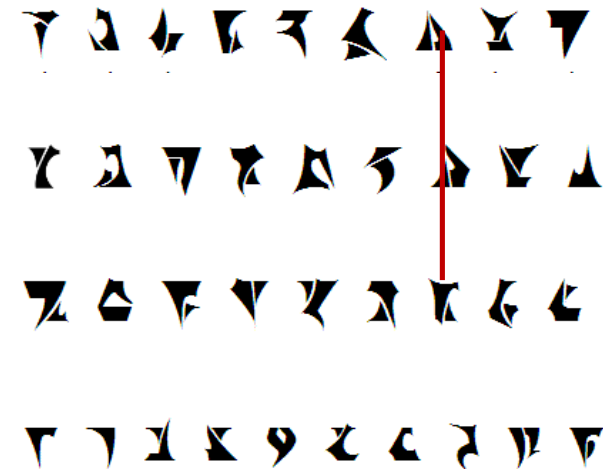
# Why is this hard?

Synthesis from examples



validation was easy!

Synthesis from specifications



SEE IF YOU CAN FIND ANY KLINGON FRUIT!

validation is hard!  
(and search is still hard)

# Module II

---

