#### Reachability in 3-VASS is Elementary

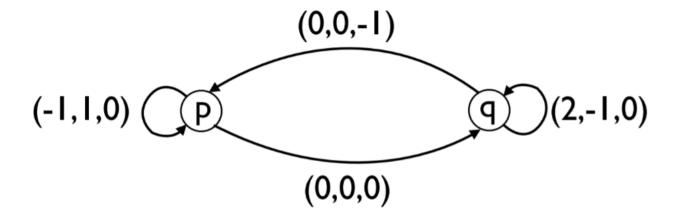
Wojciech Czerwiński<sup>1</sup> Ismaël Jecker<sup>2</sup> Sławomir Lasota<sup>1</sup> Łukasz Orlikowski<sup>1</sup>

> <sup>1</sup>University of Warsaw <sup>2</sup>FEMTO-ST, CNRS, Univ. Franche-Comté

> > 03.09.2025

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## Vector Additions Systems with States



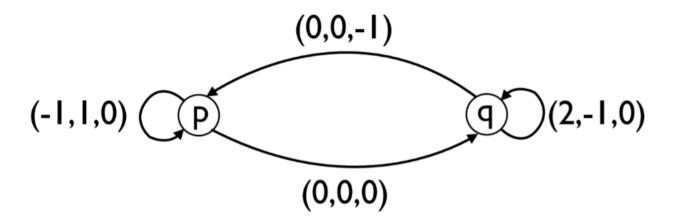
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**Question:** Is there a run from s to t in V?



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- Big gaps between lower bounds and upper bounds
- Hope for elementary complexity bounds

#### Main result

#### Theorem

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The reachability problem in 3-VASS is in 2-ExpSpace, under binary encoding.

- We prove this by induction on the number of strongly connected components
- We show, that length of the shortest run is  $n^{2^{2^{\mathcal{O}(k)}}}$  where n is size of VASS and k is the number of strongly connected components





#### 15h09-16h12 Session 9: Vector Addition Systems

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The Tractability Border of Reachability in Simple Vector Addition Systems with States

15h18-15h27 Shrisha Rao

Continuous Pushdown VASS in One Dimension are Easy

15h27-15h36 A. R. Balasubramanian

Decidability and Complexity of Decision Problems for Affine Continuous VASS

15h36-15h45 Łukasz Orlikowski

Languages of unambiguous vector addition systems with states

15h45-15h54 Roland Guttenberg

Reachability in Priority Vector Addition Systems

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Reachability in Two-Dimensional Branching VASS is decidable

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# Thank You!