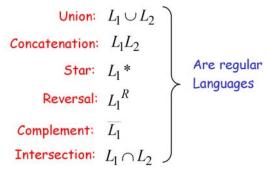
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Note: The purpose of the following questions is:

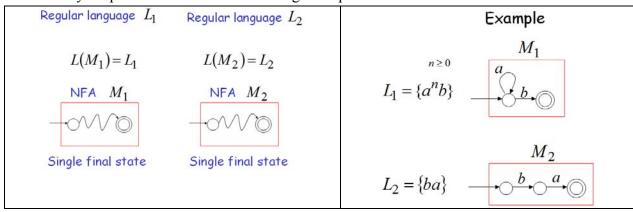
| Enhance learning | Summarized points | Analyze abstract ideas |
|------------------|-------------------|------------------------|
|------------------|-------------------|------------------------|

Class 4: Properties of Regular Languages

- 1. [Slide 2-4] Prove that any NFA can be converted to an equivalent NFA with a single final state.
- 2. [Slide 5] Extreme case: show that an NFA *without* final state be converted to an equivalent NFA with a single final state.
- 3. [Slide 7-22] For regular languages L_1 and L_2 prove that:



Show that your proof works for the following example:



4. [Slide 24-30] One way of describing regular languages is via the notation of regular expressions. This notation involves a combination of strings and symbols from some alphabet ∑, parentheses, and the operators +, ., and *. The simplest case is the language {a}, which will be denoted by a regular expression a. We will use + do denote union. We use . for concatenation and * for a star-closure. Complete the following table

| | Language | Corresponding regular expression |
|----|-----------|----------------------------------|
| 1. | {a} | а |
| 2. | {a, b, c} | |
| 3. | | (a + (b.c))* |

5. A regular language over an alphabet Σ is one that can be obtained from the very simplest languages over Σ , those containing a single string of length 0 or 1, using only the operations of union, concatenation, and Kleene *. A regular language can therefore be described by an explicit formula.

a) Find the corresponding regular expression for each of the following languages:

| | Language | Corresponding regular expression |
|-----|---|----------------------------------|
| 1. | {a} | |
| 2. | {0} | |
| 3. | {001} (i.e., {0}{0}{1}) | |
| 4. | {0, 1} (i.e., {0} ∪ {1}) | |
| 5. | {0, 10} (i.e., {0} ∪ {10}) | |
| 6. | $\{1, \lambda\}\{001\}$ | |
| 7. | {110} [*] {0, 1} | |
| 8. | {1} [*] {10} | |
| 9. | {10, 111, 11010} [*] | |
| 10. | $\{0, 10\}^*(\{11\}^* \bigcup \{001, \lambda \})$ | |

6. [Slide 31-34] Complete the following table

| | Language | Corresponding regular expression |
|----|----------|----------------------------------|
| 1. | : | $r = (a + b)^*(a + bb)$ |
| 2. | | r = (aa)*(bb)*b |
| 3. | | $r = (0+1)^*00(0+1)^*$ |
| 4. | | $r = (1+01)^*(0+\lambda)$ |

7. [Slide 33] For $\Sigma = \{0, 1\}$, give the regular expression r such that

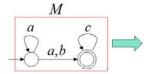
$$L(r) = \{w \in \sum^* : w \text{ has at least one pair of consecutive zeros}\}.$$

8. [Slide 34] Find a regular expression for the language

$$L(r) = \{w \in \{0,1\}^* : w \text{ has no pair of consecutive zeros}\}.$$

9. [Slide 38-50] Prove the following **Theorem:** Let r be a regular expression. Then there exists some nondeterministic finite acceptor that accepts L(r). Consequently, L(r) is a regular language.

10. [Slide 51] From *M* construct the equivalent **Generalized Transition Graph (GTG)** in which transition labels are regular expressions:



11. [Slide 52-56] Find regular expressions for the languages accepted by the following automata.

