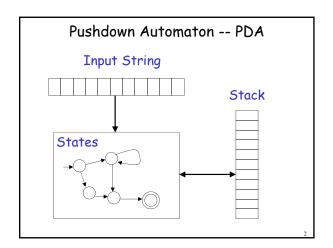
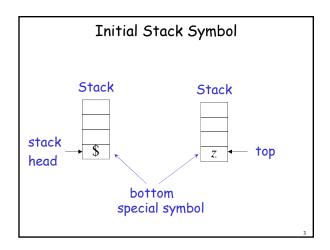
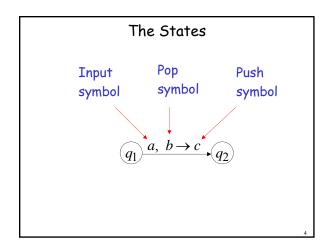
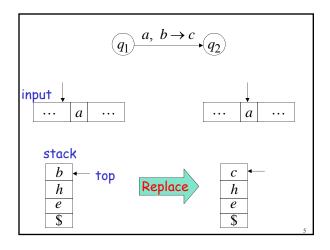
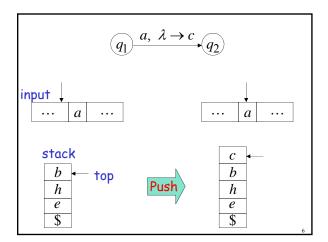
Pushdown Automata PDAs class 10

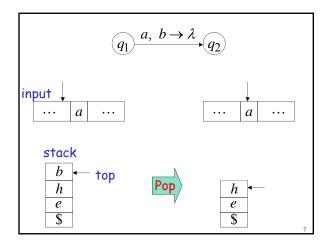


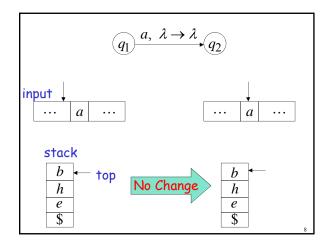


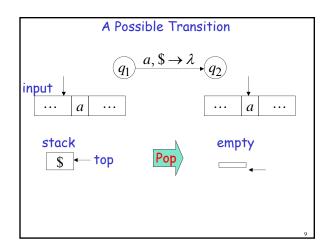


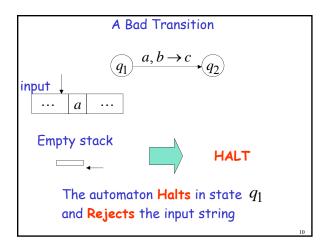


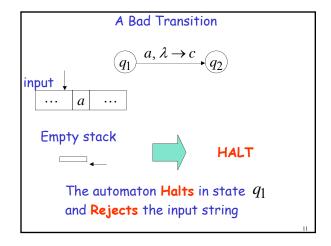


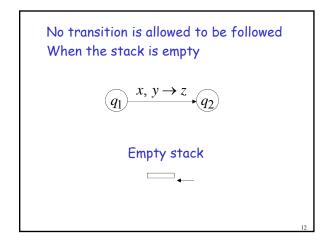


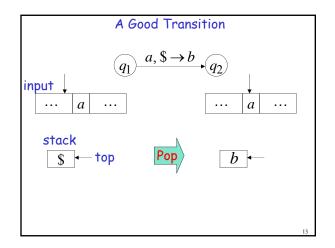


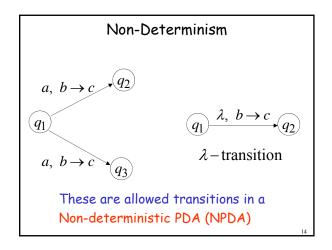


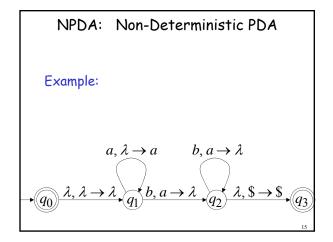


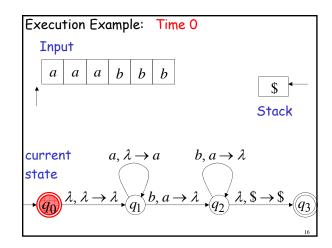


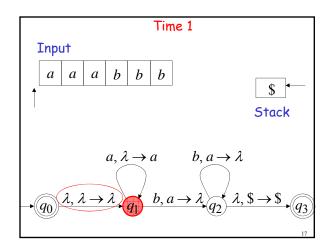


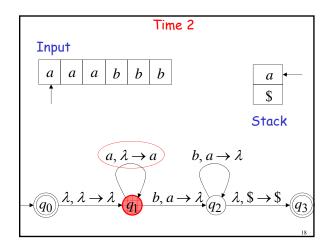


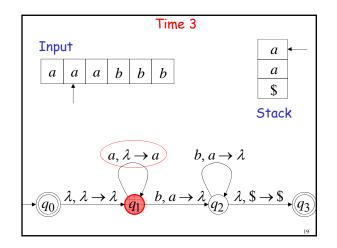


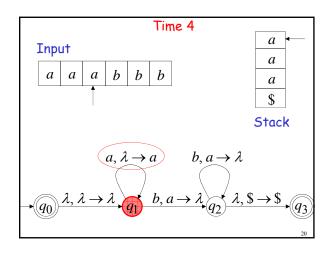


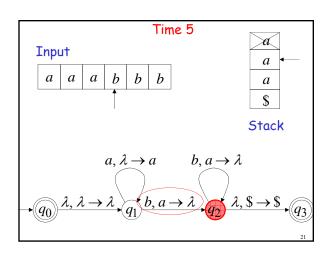


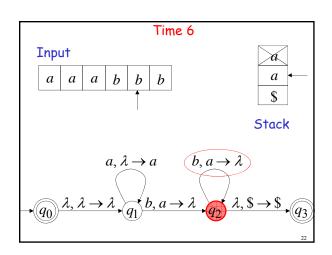


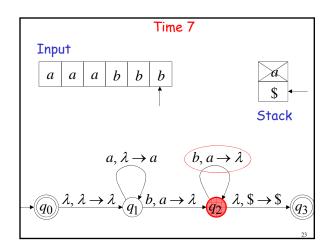


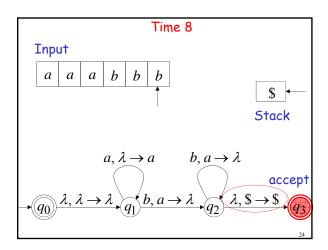












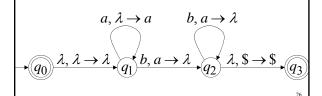
A string is accepted if there is a computation such that:

All the input is consumed AND

The last state is a final state

At the end of the computation, we do not care about the stack contents

The input string aaabbb is accepted by the NPDA:



In general,

$$L = \{a^n b^n : n \ge 0\}$$

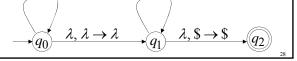
is the language accepted by the NPDA:

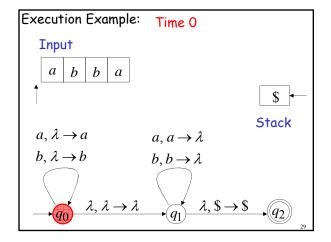
Another NPDA example

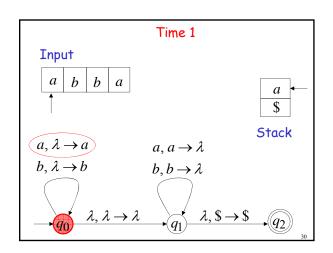
NPDA M

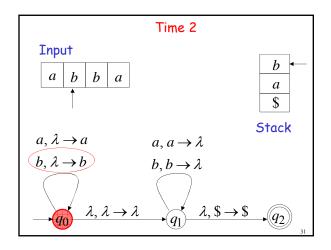
$$L(M) = \{ww^R\}$$

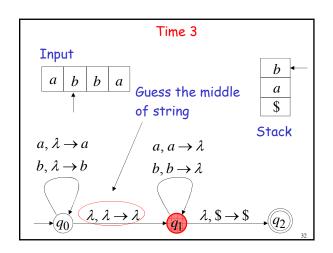
 $a, \lambda \to a$ $a, a \to \lambda$ $b, \lambda \to b$ $b, b \to \lambda$

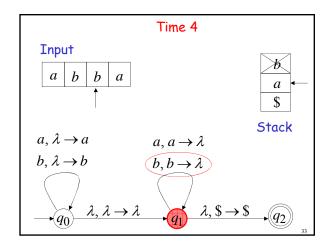


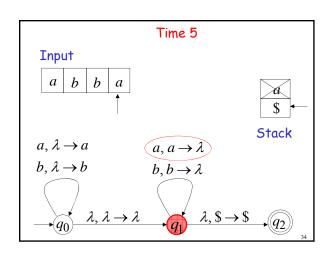


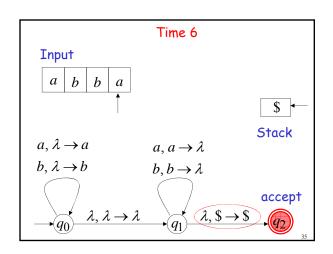


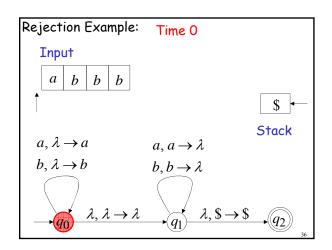


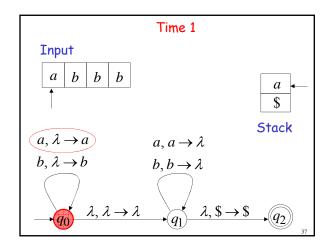


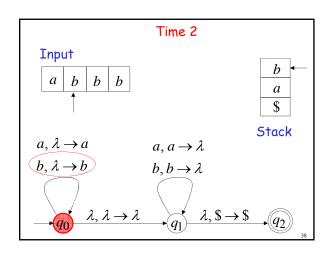


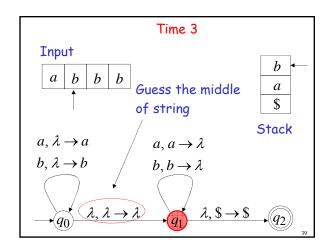


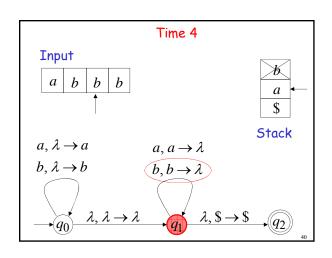


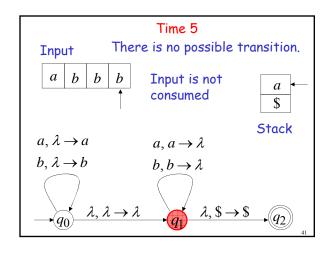


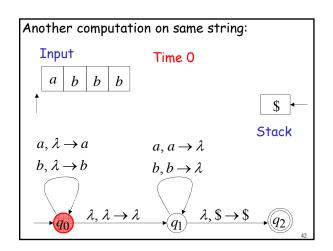


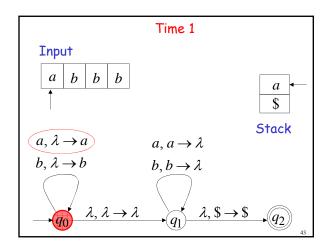


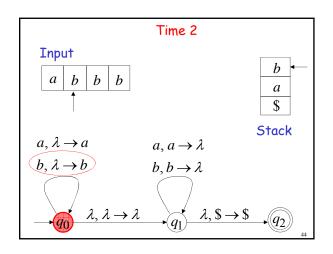


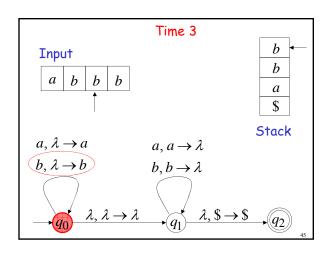


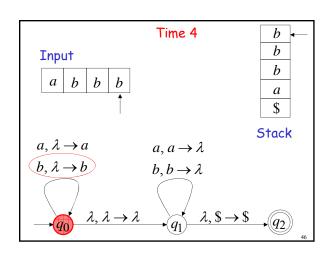


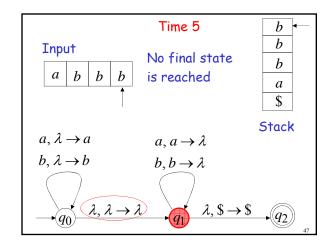


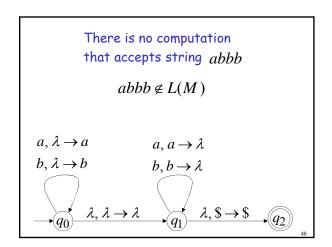












A string is rejected if there is no computation such that:

All the input is consumed AND

The last state is a final state

At the end of the computation, we do not care about the stack contents

In other words, a string is rejected if in every computation with this string:

The input cannot be consumed

OR

The input is consumed and the last state is not a final state

OR

The stack head moves below the bottom of the stack

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Another NPDA example

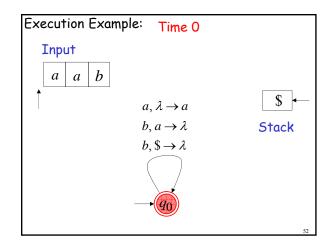
NPDA M

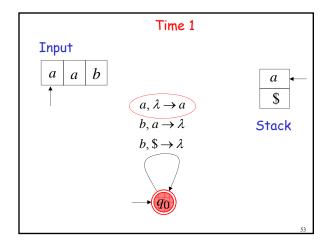
$$L(M) = \{a^n b^m : n \ge m - 1\}$$

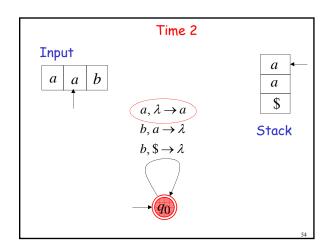
$$a, \lambda \to a$$

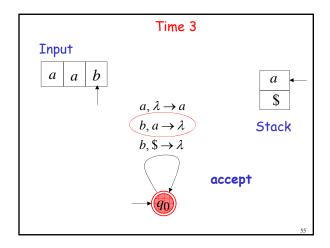
$$b, a \to \lambda$$

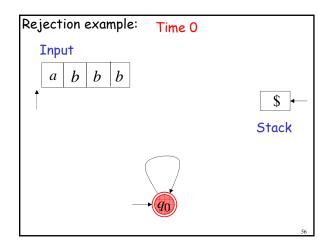
$$q_0$$

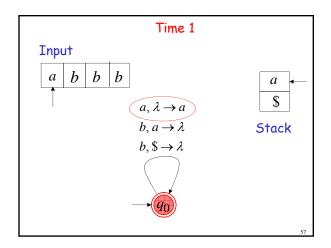


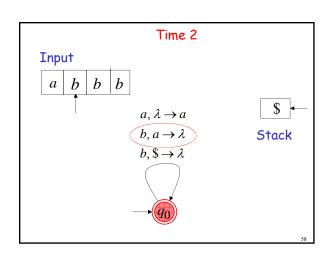


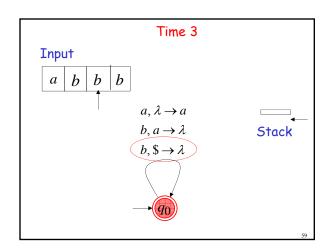


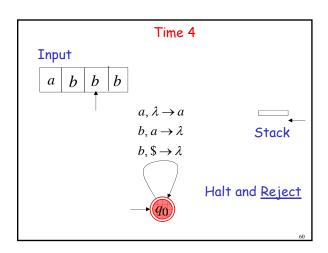


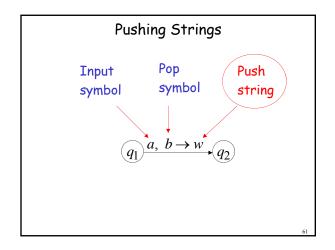


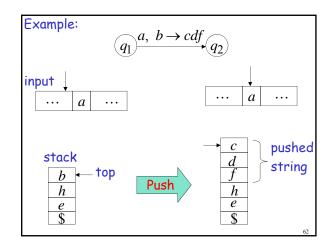


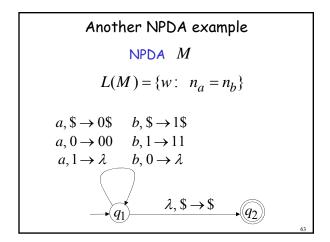


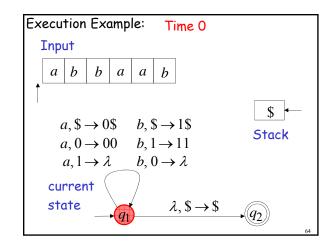


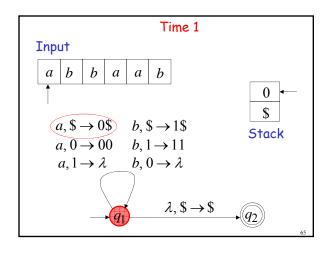


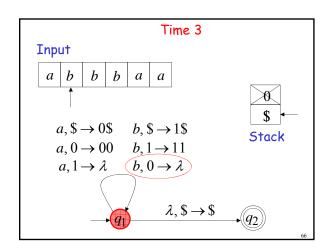


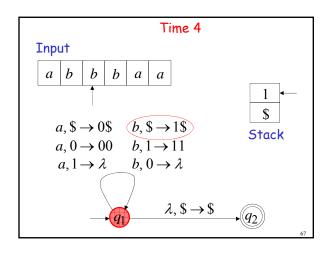


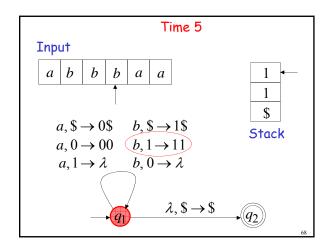


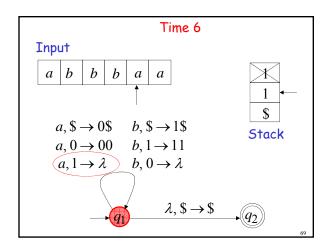


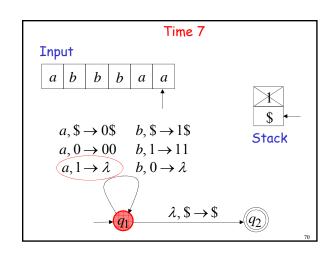


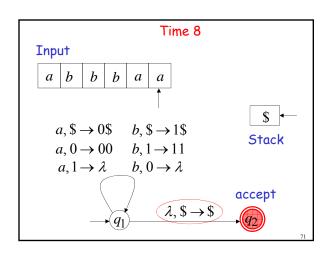


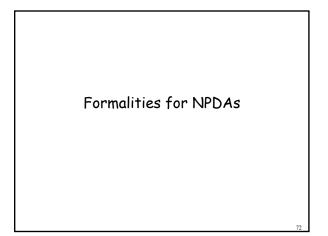


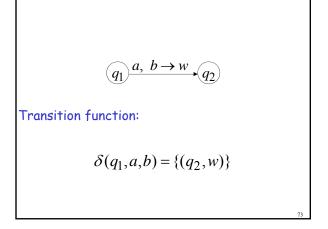


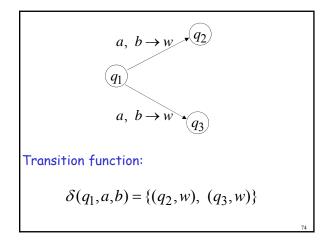


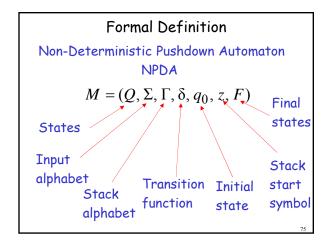


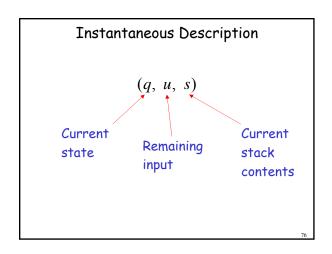


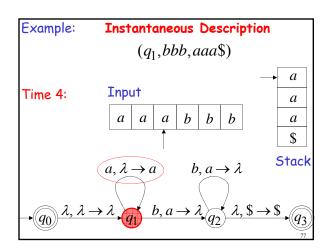


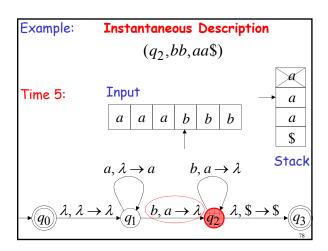












We write:

$$(q_1,bbb,aaa\$) \succ (q_2,bb,aa\$)$$
 Time 4 Time 5

A computation:

 $(q_0, aaabbb,\$) \succ (q_1, aaabbb,\$) \succ$ $(q_1, aabbb, a\$) \succ (q_1, abbb, aa\$) \succ (q_1, bbb, aaa\$) \succ$ $(q_2, bb, aa\$) \succ (q_2, b, a\$) \succ (q_2, \lambda,\$) \succ (q_3, \lambda,\$)$

$$a, \lambda \to a \qquad b, a \to \lambda$$

$$\downarrow q_0 \qquad \lambda, \lambda \to \lambda \qquad \downarrow q_1 \qquad b, a \to \lambda \qquad \downarrow q_2 \qquad \lambda, \$ \to \$ \qquad \downarrow q_3$$

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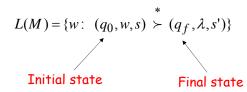
 $(q_0,aaabbb,\$) \succ (q_1,aaabbb,\$) \succ$ $(q_1,aabbb,a\$) \succ (q_1,abbb,aa\$) \succ (q_1,bbb,aaa\$) \succ$ $(q_2,bb,aa\$) \succ (q_2,b,a\$) \succ (q_2,\lambda,\$) \succ (q_3,\lambda,\$)$

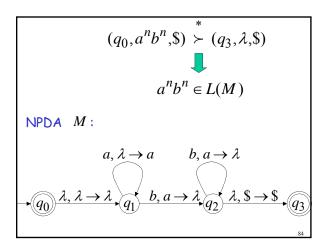
For convenience we write:

$$(q_0,aaabbb,\$) \stackrel{*}{\succ} (q_3,\lambda,\$)$$

Formal Definition

Language L(M) of NPDA M:





Therefore:
$$L(M) = \{a^n b^n : n \ge 0\}$$

NPDA $M:$

$$a, \lambda \to a \qquad b, a \to \lambda$$

$$\downarrow q_0 \qquad \lambda, \lambda \to \lambda \qquad q_1 \qquad b, a \to \lambda \qquad q_2 \qquad \lambda, \$ \to \$ \qquad q_3 \qquad q_3 \qquad q_4 \qquad q_5 \qquad$$