



Gradiance Online Accelerated Learning

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Based on Sections 10.2 and 10.3 of HMU.

1. In the following expressions, - represents negation of a variable. For example, -x stands for "NOT x", + represents logical OR, and juxtaposition represents logical AND (e.g., (x+y)(y+z) represents (x OR y) AND (y OR z)).

Identify the expression that is satisfiable, from the list below.

- a) $(x)(-x+y)(-x+-y)(y+z)$
- b) $(-x+y)(-x+-y)(y+z)(x)$
- c) $(x)(-x+y)(-y+z)(-z+-x)$
- d) $(x)(-x+-y)(-y+z)(-z+-x)$

Answer submitted: **d)**

You have answered the question correctly.

Question Explanation:

All choices fall into one of four categories (possibly with clauses reordered):

$(x)(-x+-y)(-y+z)(-z+-x)$ is satisfiable. Let $x=1$ and $y=z=0$.

$(x+y)(x+-y)(-x+y)(-x+-y)$ is not satisfiable. If $x=1$ and $y=1$, the fourth clause is false. If $x=1$ and $y=0$, the third clause is false. If $x=0$ and $y=1$, the second clause is false. If $x=0$ and $y=0$, the first clause is false.

$(x)(-x+y)(-y+z)(-z+-x)$ is not satisfiable. The first clause forces x to be true if the whole expression is to be satisfied. Then, the second clause forces y to be true and the third clause forces z to be true. But then the fourth clause is false.

$(x)(-x+y)(-x+-y)(y+z)$ is not satisfiable. The first clause forces x to be true if the whole expression is to be satisfied. Then, the second clause forces y to be true. But then the third clause is false.

2. The polynomial-time reduction from SAT to CSAT, as described in Section 10.3.3 (p. 452), needs to introduce new variables. The reason is

that the obvious manipulation of a boolean expression into an equivalent CNF expression could exponentiate the size of the expression, and therefore could not be polynomial time.

Suppose we apply this construction to the expression $(u+(vw))+x$, with the parse implied by the parentheses. Suppose also that when we introduce new variables, we use y_1, y_2, \dots

After constructing the corresponding CNF expression, identify one of its clauses from the list below. Note: logical OR is represented by $+$, logical AND by juxtaposition, and logical NOT by $-$.

- a) (y_2+y_1+u)
- b) (x)
- c) $(y_3+-y_2+-y_1+w)$
- d) (y_2+u)

Answer submitted: **c)**

Your answer is incorrect.

Possible error: you may have used the OR construction for AND as well. Hint: remember that when you take the AND of two CNF expressions (i.e., AND's of clauses) you simply take the product of the clauses for each expression. You should consult the reduction from SAT to CSAT in Section 10.3.3 (p. 452).

Question Explanation:

The first subexpression to which we apply the transformation is vw . The AND rule is simple: take the AND of the clauses for each side. That gives us $(v)(w)$ as the CNF expression.

Next, we work on $u+(vw)$. The rule for OR requires us to introduce variable y_1 . It is added positively to all the clauses on the left side and negatively to all clauses on the right side. That gives us $(y_1+u)(-y_1+v)(-y_1+w)$.

Finally, we apply the same transformation to $(u+(vw))+x$, introducing y_2 . The final answer is $(y_2+y_1+u)(y_2+-y_1+v)(y_2+-y_1+w)(-y_2+x)$.

The correct choice is: **a)**

3. The NOT-ALL-EQUAL 3SAT problem is defined as follows: Given a 3-CNF formula F , is there a truth assignment for the variables such that each clause has at least one true literal and at least one false literal? The NOT-ALL-EQUAL 3SAT problem is NP-complete.

This question is about trying to reduce the NOT-ALL-EQUAL 3SAT problem to the MAX-CUT problem defined below to show the latter to be NP-complete.

A cut in an undirected graph $G=(V,E)$ is a partitioning of the set of nodes V into two disjoint subsets V_1 and V_2 . The size of a cut is the number of edges $e=(u,v)$ where u is in V_1 and v is in V_2 . The MAX-CUT problem is defined as follows: Given an undirected graph $G=(V,E)$ and a positive integer k , does G have a cut of size k or more?

Given a 3CNF expression E , we create the graph $G=(V,E)$ using the transformation given by Theorem 10.18 in Section 10.4.2 on p. 460 of the text. Then given an assignment A , create a cut C in G by partitioning the set of nodes V as follows: the nodes corresponding to the

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uncomplemented literals are in set V1 and those corresponding to the complemented variables are in set V2.

For variable a , let a' denote NOT(a). Let

$$E = (a + b + c)(a + b' + c)(a' + b' + d)(c' + d' + e)$$

be an instance of NOT-ALL-EQUAL 3SAT. Suppose a cut separates the true nodes from false nodes according to some truth assignment applied to E . How many edges between nodes corresponding to the literals in the same clause are cut? How many other edges are cut? Find out how the cut-size can be computed for an arbitrary instance of NOT-ALL-EQUAL 3SAT. Then for the instance E , determine in which of the cases below, the cut-size C corresponds to the satisfiable assignment given.

- a) $a = F, b = F, c = T, d = T, e = T, C = 15$
- b) $a = T, b = F, c = F, d = T, e = F, C = 8$
- c) $a = F, b = T, c = T, d = F, e = T, C = 15$
- d) $a = T, b = F, c = T, d = F, e = T, C = 15$

Answer submitted: **d)**

Your answer is incorrect.

Hint: The edges in the cut come from two sets: the edges between complemented and uncomplemented literals for the same variable and from the edges between literals in the same clause. Clause $(a + b' + c)$ has all literals true and hence contributes nothing to the second set.

Question Explanation:

If each clause has at least two literals with different truth values, they end up in different partitions of the cut. The maximum contribution from a single clause to the total cut-size is hence 2. For 4 clauses, the total is 8. Complemented and uncomplemented literals for the same variable will necessarily be in different partitions of the cut and hence each such pair contributes 1 to the cut-size. There are 7 such pairs in E with a total contribution of 7. Hence maximum cut-size = $8 + 7 = 15$.

The correct choice is: **c)**

4. The Boolean expression $wxyz+u+v$ is equivalent to an expression in 3-CNF (a product of clauses, each clause being the sum of exactly three literals). Find the simplest such 3-CNF expression and then identify one of its clauses in the list below. Note: $-e$ denotes the negation of e . Also note: we are looking for an expression that involves only u, v, w, x, y , and z , no other variables. Not all boolean expressions can be converted to 3-CNF without introducing new variables, but this one can.
- a) $(y+z+v)$
 - b) $(z+u+v)$
 - c) $(u+v+-x)$
 - d) $(u+v+-w)$

Answer submitted: **c)**

Your answer is incorrect.

Hint: If C is a clause in a 3-CNF expression equivalent to some boolean expression E , then E must logically imply C (because the 3-CNF

expression can only be true when C is true). However, $wxyz+u+v$ can be true when $(u+v+x)$ is not; let $u=v=0$ and $w=x=y=z=1$.

The definition of 3-CNF is in Section 10.3.1 (p. 448). Note that the reduction of SAT to CSAT in Theorem 10.13 (p. 452) and of CSAT to 3-SAT in Theorem 10.15 (p. 457) involve converting arbitrary boolean expressions to 3-CNF expressions, but usually introduce new variables, as they would in this case.

Question Explanation:

The simplest way to proceed is to use the distributing law of OR over AND, three times, to distribute $u+v$ over $wxyz$. The result is $(w+u+v)(x+u+v)(y+u+v)(z+u+v)$.

The correct choice is: **b)**

5. Use the construction from Theorem 10.15 (p. 457) to convert the following clauses:

1. $(a+b)$
2. $(c+d+e+f)$
3. $(g+h+i+j+k+l+m)$

to products of 3 literals per clause. In each case, the new clauses must be satisfiable if and only if the original clause is satisfiable. For the first clause, introduce variables x_1, x_2, \dots in that order from the left; for the second introduce y_1, y_2, \dots in that order from the left, and for the third introduce z_1, z_2, \dots in that order from the left. Use $\neg w$ as shorthand for NOT w . Then identify, in the list below, the one clause that would appear among the clauses generated by the construction.

- a) $(c+\neg d+y_1)$
- b) $(a+b+\neg x_1)$
- c) $(e+y_1+\neg y_2)$
- d) $(m+z_4+\neg z_5)$

Answer submitted: **b)**

You have answered the question correctly.

Question Explanation:

$(a+b)$ becomes $(a+b+x_1)(a+b+\neg x_1)$.

$(c+d+e+f)$ becomes $(c+d+\neg y_1)(e+f+y_1)$.

$(g+h+i+j+k+l+m)$ becomes $(g+h+\neg z_1)(i+z_1+\neg z_2)(j+z_2+\neg z_3)(k+z_3+\neg z_4)(l+m+z_4)$.