Time Complexity

- ·We use a multitape Turing machine
- We count the number of steps until a string is accepted
- •We use the O(k) notation

Example: $L = \{a^n b^n : n \ge 0\}$

Algorithm to accept a string w:

- ·Use a two-tape Turing machine
- \cdot Copy the a on the second tape
- •Compare the a and b

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$$L = \{a^n b^n : n \ge 0\}$$

Time needed:

- •Copy the a on the second tape O(|w|)
- •Compare the a and b O(|w|)

Total time: O(|w|)

 $L = \{a^n b^n : n \ge 0\}$

For string of length n

time needed for acceptance: O(n)

Language class: DTIME(n)

 $\begin{array}{ccc}
DTIME(n) \\
L_1 & L_2 & L_3
\end{array}$

A Deterministic Turing Machine accepts each string of length n in time O(n)

DTIME(n) $\{a^nb^n: n \ge 0\}$ $\{ww\}$

In a similar way we define the class

DTIME(T(n))

for any time function: T(n)

Examples: $DTIME(n^2), DTIME(n^3),...$

Example: The membership problem for context free languages

 $L = \{w : w \text{ is generated by grammar } G\}$

 $L \in DTIME(n^3)$ (CYK - algorithm)

Polynomial time

Theorem: $DTIME(n^{k+1}) \subset DTIME(n^k)$

 $DTIME(n^{k+1})$ $DTIME(n^k)$

Polynomial time algorithms: $DTIME(n^k)$

Represent tractable algorithms: For small $\,k\,$ we can compute the result fast

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The class P

 $P = \bigcup DTIME(n^k)$ for all k

·Polynomial time

·All tractable problems

P CYK-algorithm $\{a^nb^n\}$ \cdots $\{ww\}$

Exponential time algorithms: $DTIME(2^n)$

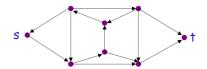
Represent intractable algorithms:

Some problem instances

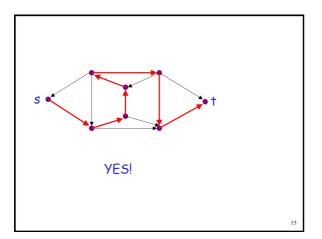
may take centuries to solve

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Example: the Hamiltonian Problem



Question: is there a Hamiltonian path from s to t?



A solution: search exhaustively all paths

L = $\{\langle G, s, t \rangle\}$: there is a Hamiltonian path in G from s to t

 $L \in DTIME(n!) \approx DTIME(2^n)$

Exponential time

Intractable problem

Example: The Satisfiability Problem

Boolean expressions in Conjunctive Normal Form:

$$t_1 \wedge t_2 \wedge t_3 \wedge \cdots \wedge t_k$$

$$t_i = x_1 \lor \overline{x}_2 \lor x_3 \lor \cdots \lor \overline{x}_p$$
Variables

Question: is expression satisfiable?

Example:
$$(\bar{x}_1 \lor x_2) \land (x_1 \lor x_3)$$

Satisfiable:
$$x_1 = 0, x_2 = 1, x_3 = 1$$

$$(\bar{x}_1 \lor x_2) \land (x_1 \lor x_3) = 1$$

Example: $(x_1 \lor x_2) \land \overline{x}_1 \land \overline{x}_2$

Not satisfiable

 $L = \{w : \text{expression } w \text{ is satisfiable}\}$

For n variables: $L \in DTIME(2^n)$

exponential

Algorithm:

search exhaustively all the possible binary values of the variables

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Non-Determinism

Language class: NTIME(n)

 $\begin{array}{c|c}
NTIME(n) \\
L_1 & L_2 \\
\end{array}$

A Non-Deterministic Turing Machine accepts each string of length n in time O(n)

Example: $L = \{ww\}$

Non-Deterministic Algorithm to accept a string ww:

- ·Use a two-tape Turing machine
- •Guess the middle of the string and copy w on the second tape
- •Compare the two tapes

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$L = \{ww\}$

Time needed:

- ·Use a two-tape Turing machine
- •Guess the middle of the string O(|w|) and copy w on the second tape
- •Compare the two tapes O(|w|)

Total time: O(|w|)

NTIME(n) $L = \{ww\}$

In a similar way we define the class

NTIME(T(n))

for any time function: T(n)

Examples: $NTIME(n^2), NTIME(n^3),...$

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Non-Deterministic Polynomial time algorithms:

 $L \in NTIME(n^k)$

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The class NP

 $NP = \bigcup NTIME(n^k)$ for all k

Non-Deterministic Polynomial time

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Example: The satisfiability problem

 $L = \{w : expression \ w \text{ is satisfiable}\}\$

Non-Deterministic algorithm:

•Guess an assignment of the variables

·Check if this is a satisfying assignment

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 $L = \{w : \text{expression } w \text{ is satisfiable}\}$

Time for n variables:

•Guess an assignment of the variables O(n)

•Check if this is a satisfying assignment O(n)

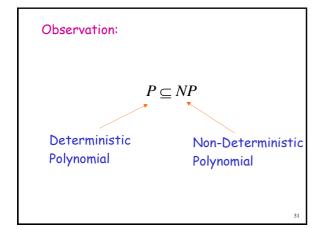
Total time: O(n)

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 $L = \{w : expression \ w \text{ is satisfiable}\}\$

 $L \in NP$

The satisfiability problem is an NP-Problem



Open Problem: P=NP ?

Open Problem: P = NP?

Example: Does the Satisfiability problem have a polynomial time deterministic algorithm?

WE DO NOT KNOW THE ANSWER