



Calcul Scientifique

Projet de Calcul Scientifique

Groupe EF06

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1.

1.1.

| Matrix dimension | Matrix type | Exec. time for <code>eig</code> (s) | Exec. time for <code>power_v11</code> , (s) |
|------------------|-------------|-------------------------------------|---|
| 200×200 | Type 1 | 9.000e-02 | 1.510e+00 |
| 400×400 | Type 1 | 4.000e-02 | 1.831e+01 |
| 600×600 | Type 1 | 6.000e-02 | 6.021e+01 |
| 200×200 | Type 2 | 3.000e-02 | 3.000e-02 |
| 400×400 | Type 2 | 4.000e-02 | 4.000e-02 |
| 600×600 | Type 2 | 7.000e-02 | 1.700e-01 |
| 200×200 | Type 3 | 1.000e-02 | 5.000e-02 |
| 400×400 | Type 3 | 3.000e-02 | 5.200e-01 |
| 600×600 | Type 3 | 7.000e-02 | 1.270e+00 |
| 200×200 | Type 4 | 2.000e-02 | 1.670e+00 |
| 400×400 | Type 4 | 3.000e-02 | 2.094e+01 |
| 600×600 | Type 4 | 6.000e-02 | 5.456e+01 |

Table 1: Execution time for different sizes and types of matrices

We can see that the `power_v11` algorithm is generally slower than the `eigen` function especially for the type 2 and 4 matrices.

1.2.

```

1  nb_it = 1;
2  norme = norm(beta*v - z, 2)/norm(beta,2);
3
4  while(norme > eps && nb_it < maxit)
5      beta_old = beta;
6      v = z/norm(z, 2);
7      z = A*v;
8      beta = (v'*z)/(v'*v);
9      norme = abs(beta-beta_old)/abs(beta_old);
10     nb_it = nb_it + 1;
11 end

```

Listing 1: Inner loop of the new algorithm

| Matrix dimension | Matrix type | Exec. time for <code>power_v11</code> , (s) | Exec. time for <code>power_v12</code> , (s) |
|------------------|-------------|---|---|
| 200×200 | Type 1 | 1.960e+00 | 3.200e-01 |
| 400×400 | Type 1 | 1.888e+01 | 2.660e+00 |
| 600×600 | Type 1 | 5.031e+01 | 7.070e+00 |
| 200×200 | Type 2 | 1.000e-02 | 1.000e-02 |
| 400×400 | Type 2 | 7.000e-02 | 1.000e-02 |
| 600×600 | Type 2 | 1.800e-01 | 4.000e-02 |
| 200×200 | Type 3 | 3.000e-02 | 1.000e-02 |
| 400×400 | Type 3 | 6.100e-01 | 1.100e-01 |
| 600×600 | Type 3 | 1.270e+00 | 2.600e-01 |

| | | | |
|------------------|--------|-----------|-----------|
| 200×200 | Type 4 | 1.530e+00 | 2.900e-01 |
| 400×400 | Type 4 | 2.113e+01 | 3.060e+00 |
| 600×600 | Type 4 | 5.914e+01 | 6.480e+00 |

We can see that the `power_v12` algorithm is globally faster than the `power_v11`.

1.3.

The main drawback of the deflated power method is the numerous matrix-vector products required to compute the eigenvectors as well as the fact that each iteration compute only one eigenvalue which can be slow if a lot of eigenvalues are desired.

1.4.

If we apply Algorithm 1 to m vectors, there is no reason for the columns of V to converge to a base. Each vector will converge toward a different projection of the dominant eigenvalue.

1.5.

In Algorithm 2, the matrix H is a smaller matrix, with dimension $n \times m$, therefore, even for larger matrices A , computing the spectral decomposition of H will not be computationally expensive.

1.6.

1.7.

```

1: function SUBSPACE ITER V1 (RALEIGH-RITZ PROJECTION)
2:   Input :  $A \in \mathbb{R}^{n \times n}, \epsilon, \text{MaxIter}, \text{PercentTrace}$ 
3:   Output :  $n_{\text{ev}}$  dominant eigenvectors  $V_{\text{out}}$  and the corresponding eigenvalues  $\Lambda_{\text{out}}$ 
4:   Generate an initial set of  $m$  orthonormal vectors  $V \in \mathbb{R}^{n \times m}; k = 0; \text{PercentReached} = 0$ 
5:   repeat until (line 52)  $\text{PercentReached} > \text{PercentTrace} \vee n_{\text{ev}} = m \vee k > \text{MaxIter}$ 
6:      $k \leftarrow k + 1$ 
7:     Compute  $Y$  such that  $Y = A \cdot V$ 
8:      $V \leftarrow$  orthonormalisation of the columns of  $Y$ 
9:     Rayleigh-Ritz projection (line 60) applied on matrix  $A$  and orthonormal vectors  $V$ 
10:    Convergence analysis step (line 70) : save eigenpairs that converged and update PercentReached
```

1.8.

There is $A \in \mathbb{R}^{n \times n}$, a square matrix, for $A \times A$ there are $\mathcal{O}(n^3)$ Flops. Since repeating p times the operation, there are $\mathcal{O}(p \times n^3)$ Flops to compute A^p . After, there is $V \in \mathbb{R}^{n \times m}$ which leads to $\mathcal{O}(p \times n^3 \times m)$ Flops to compute $A^p \times V$.

1.9.

1.10.

| Matrix dimension | Matrix type | Iterations for subspace_iter0 | Iterations for subspace_iter1 | Iterations for subspace_iter2 | p (A^p) |
|------------------|-------------|-------------------------------|-------------------------------|-------------------------------|-------------|
| 200×200 | Type 1 | 2309 | 263 | 132 | 2 |
| 200×200 | Type 1 | 2309 | 263 | 88 | 3 |
| 200×200 | Type 1 | 2309 | 263 | 53 | 5 |
| 200×200 | Type 1 | 2309 | 263 | 27 | 10 |

Table 2: After computing

When increasing the value of p to compute A^p in `subspace_iter2`, the number of flops to compute the results is : $\text{Iterations}(\text{iter2}) \simeq \frac{\text{Iterations}(\text{iter2})}{p}$. This is due to the eigenvalues of A^p raised to the p power also. Increasing the conditionment of the matrix leading to a faster convergence.

1.11.

The accuracy differs because eigenpairs are computed from the columns of new matrix V. However, the firsts columns are recalculated at each step, that will lead to have different approximate size for the new eigenpairs computed. By recomputing them, the quality reduces and diverge from the approximate size of the first one.

1.12.

By freezing the converged columns, the algorithm will not have to recalculate them everytime. Which means that the accuracy for the eigenpairs will be more equal. The first and last will have the same approximate size.

1.13.

1.14.

1.15.

2.

2.1.

Σ_k is of size (k, k)

U_k is of size (q, k)

V_k is of size (p, k)

2.2.

| | |
|--------------|-----------|
| eps | 10^{-8} |
| maxit | 10000 |
| search_space | 400 |
| percentage | 0.995 |
| puiss | 1 |

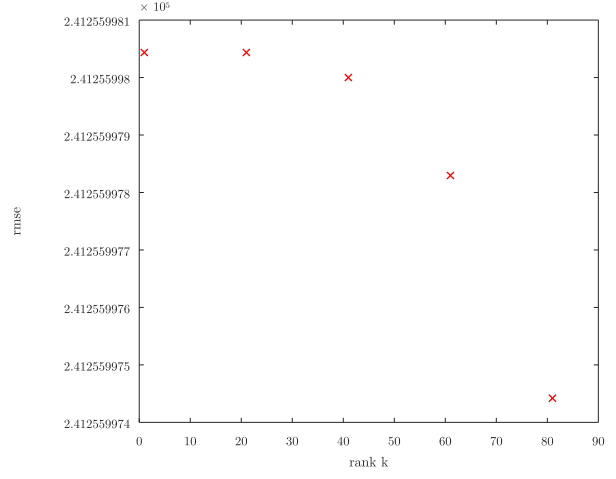


Figure 1: eig method

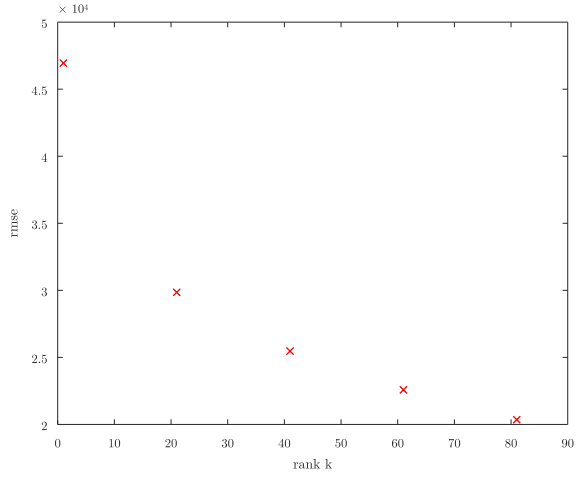


Figure 2: power method

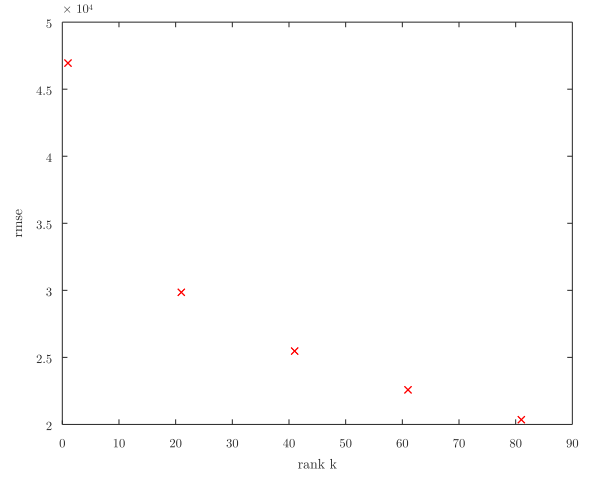


Figure 3: subspace_iter0 method

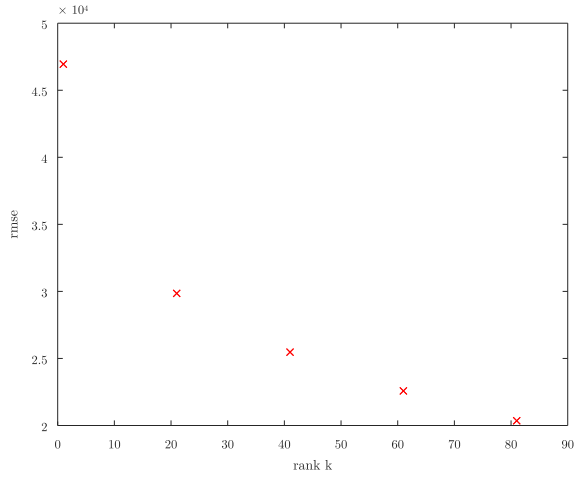


Figure 4: subspace_iter1 method

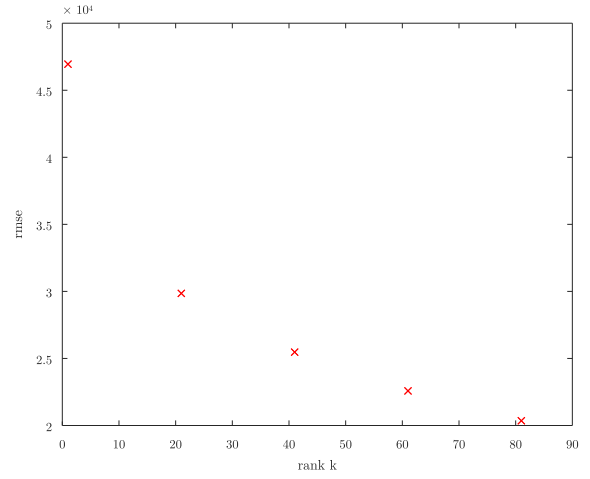


Figure 5: subspace_iter2 method

| | |
|--------------|-----------|
| eps | 10^{-8} |
| maxit | 3000 |
| search_space | 500 |
| percentage | 0.995 |
| puiss | 1 |

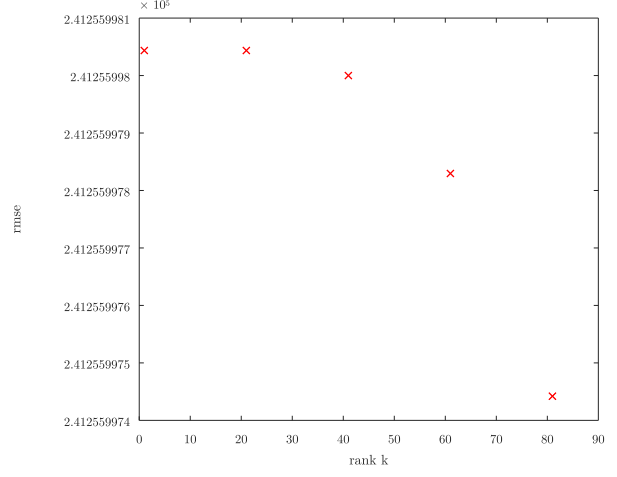


Figure 6: eig method

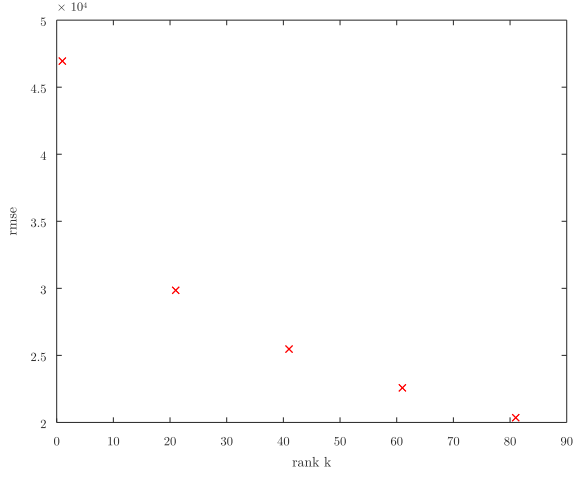


Figure 7: power method

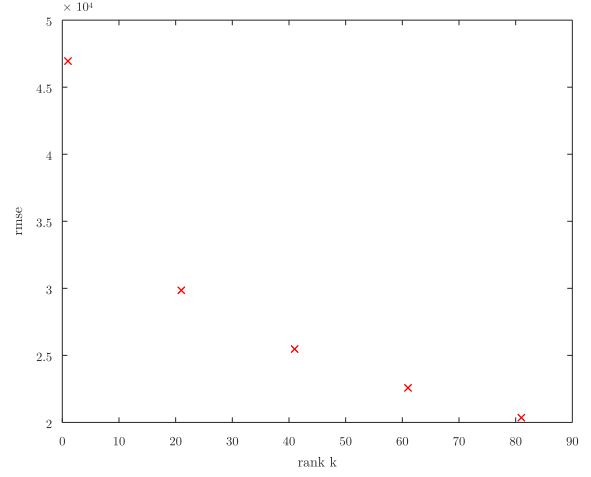


Figure 8: subspace_iter0 method

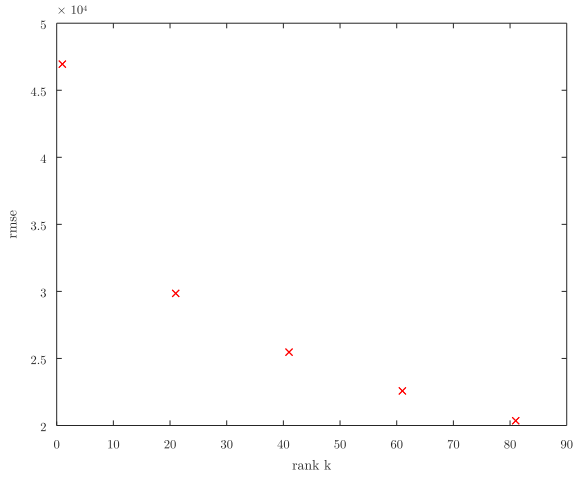


Figure 9: subspace_iter1 method

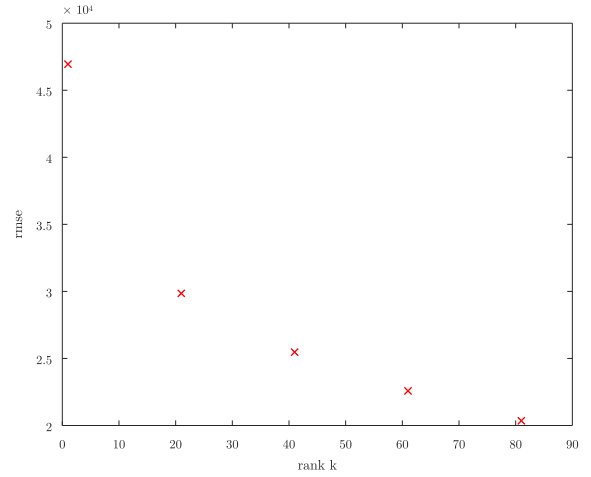


Figure 10: subspace_iter2 method

We might had an issue when calculating the RMSE with the eig function. Despite this error, we can see that using different method or changing the parameter do not affect the style of the curve.

In the first part when reconstructing the comic, we saw that the image definition looks like more and more of the original when reconstructing it for k increasing.

However, at a certain point, the RMSE tend to have little variation for some k values (less than $\frac{1}{2} \times 10^4$ between $k = 60$ and $k = 80$). But what is impacted is the time for computation, taking more time when k improve.

3. Conclusion

The RMSE is a good method to measure the reconstruction of an image after compression. All algorithms compute the same result but can take longer time to render the result. We may take that time in argument to choose which algorithm to pick, alongside with the value for k (or other factors do not see here).