

# Writeup CVE-2019-19194

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This is a writeup and *theoretical* Proof-of-Concept of CVE-2019-19194.

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## Summary

This report describes how the *Zero LTK Initialisation* vulnerability (CVE-2019-19194) allows an attacker full communication control over Bluetooth Low Energy (BLE) application by bypassing the *Secure Connections* pairing procedure.

## Vulnerable software and version

This vulnerability affects products using the Telink SMP implementations that support the *Secure Connections* pairing procedure.

## Overview

BLE is a low-consumption, short-range radio communication system. It consists in a set of standardized protocols that provide remote connectivity and security between two devices.

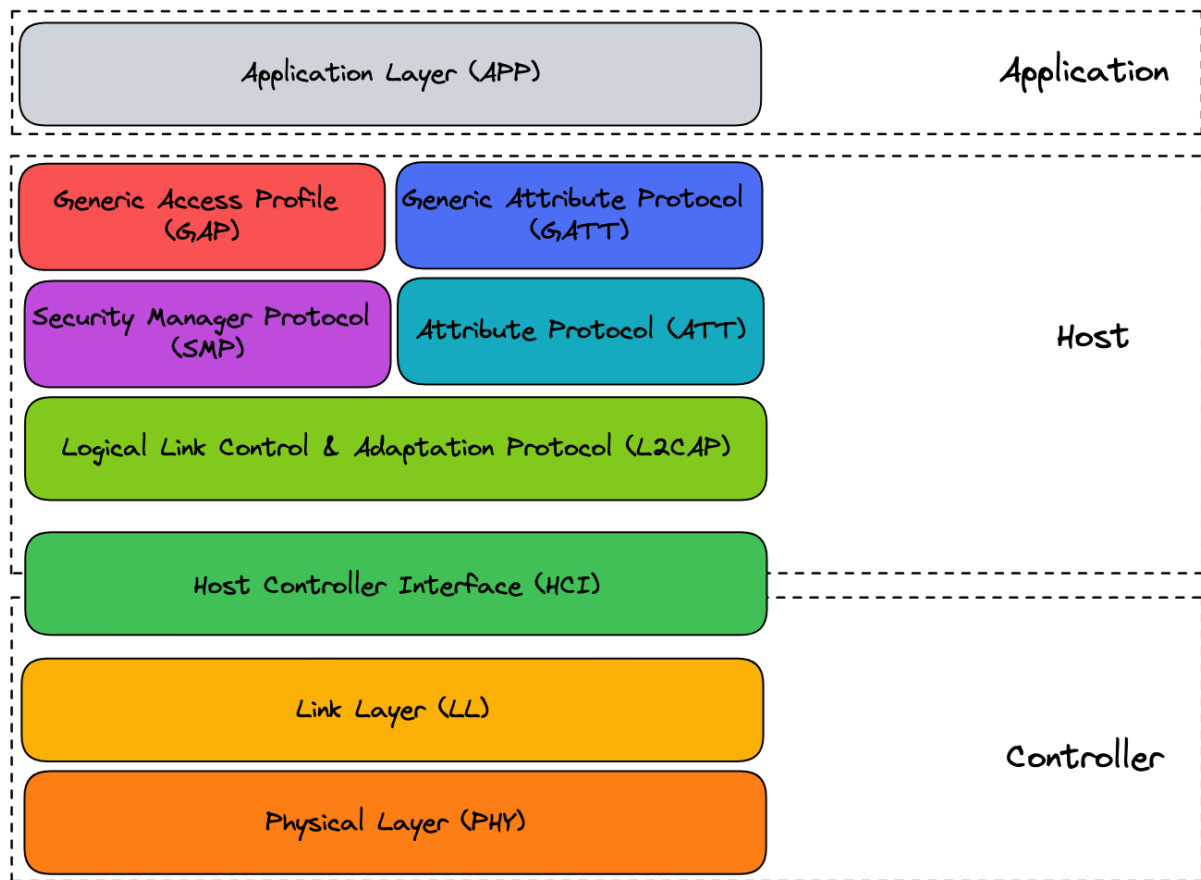
## Protocol Stack and Architecture

The BLE stack is distributed across two architectural blocks: *Host* and *Controller*.

The stack distribution allows to implement each block in physically separate components.

A standard logical interface named the *Host Controller Interface* (HCI) allows communication between the two blocks.

On top of the *Host* block sits the BLE application.



**Figure 1:** Bluetooth Low Energy protocol stack and architecture

### Physical Layer (PHY)

The physical layer operates in the Industrial, Scientific and Medical (ISM) radio band, across the 2.4GHz spectrum. It uses 40 channels: 3 *advertising* channels and 37 *data* channels.

### Link Layer (LL)

The link layer (LL) has many responsibilities which will not be described here. It is governed by a state machine which defines important roles and states:

- An *advertising* device transmits advertisement packets across advertising channels. When advertising, a device informs whether or not it is *connectable*.
- A *scanner* device listens to advertising packets from other devices.
- Once a scanner has received advertising packets from another device, it can initiate a connection procedure if the advertiser is *connectable*.

### **Logical Link Control and Adaptation Protocol (L2CAP)**

The logical link control and adaptation protocol acts as a protocol multiplexing layer. It handles *fragmentation* and *recombination* of packets between the layers below and above.

### **Generic Access Profile (GAP)**

The Generic Access Profile concerns device *discovery* and *connection*. In other words, GAP defines procedures for the transmission of advertising packets and their reception through scanning.

### **Generic Attribute Profile (GATT)**

Once a connection between two BLE devices has been established, GATT uses a client/server model to exchange data between the two devices. And both client and server use the Attribute Protocole (ATT).

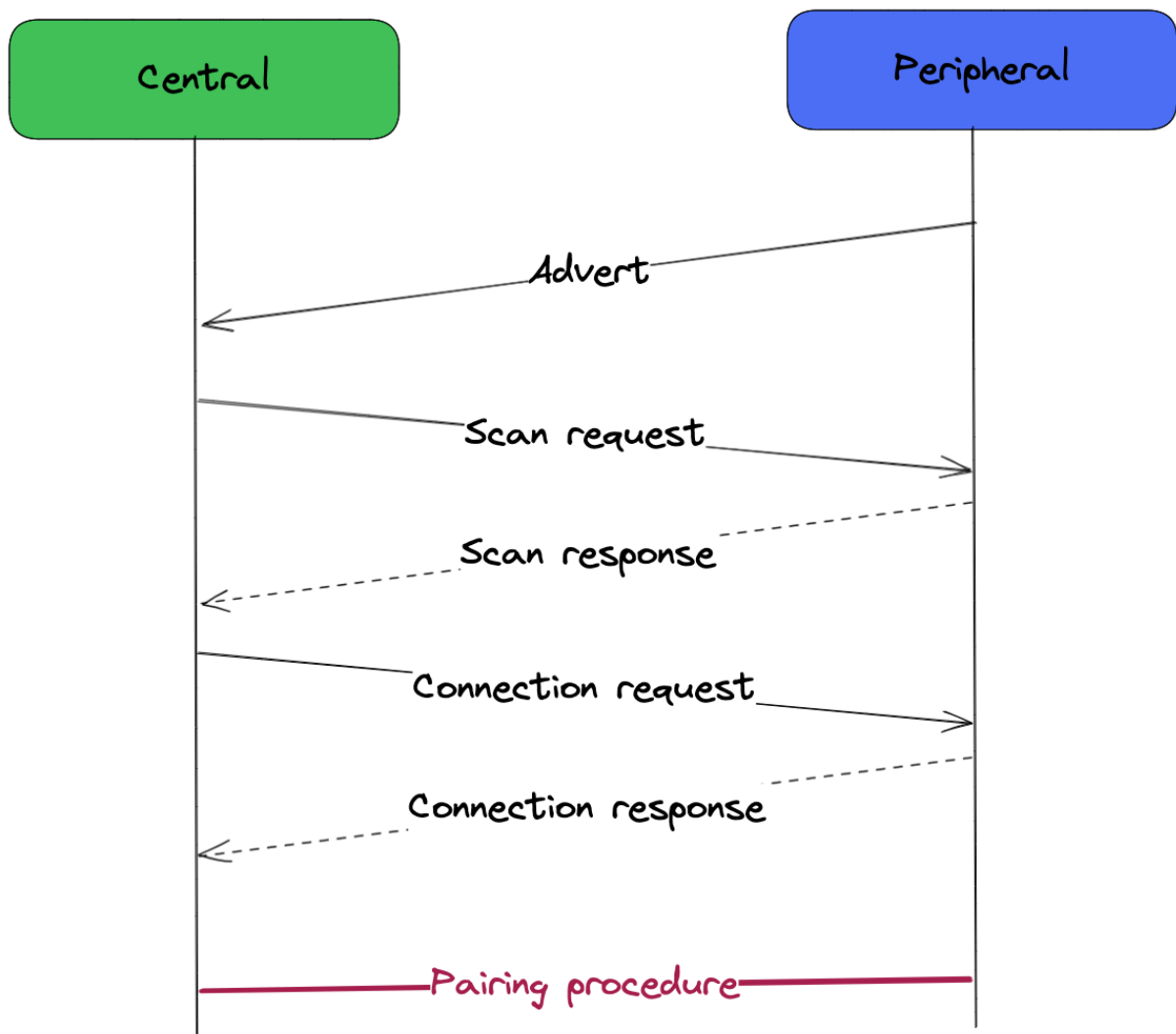
- Server: the device exposing data accepts commands and emits responses, notifications or indications.
- Client: the device requesting the reading of data sends commands, accepts incoming responses, notifications and indications.

### **Security Manager (SM)**

The SM supports security-related procedures such as *pairing*, *bonding* and *key distribution*. Device pairing is considered as the foundation of Bluetooth security: once paired, the two devices can encrypt their communication, authenticate each other, ...

### **Pairing procedure**

Among the different protocols involved in BLE communication, the *zero LTK installation* CVE happens during the pairing procedure, in the *Secure Connections* mode.



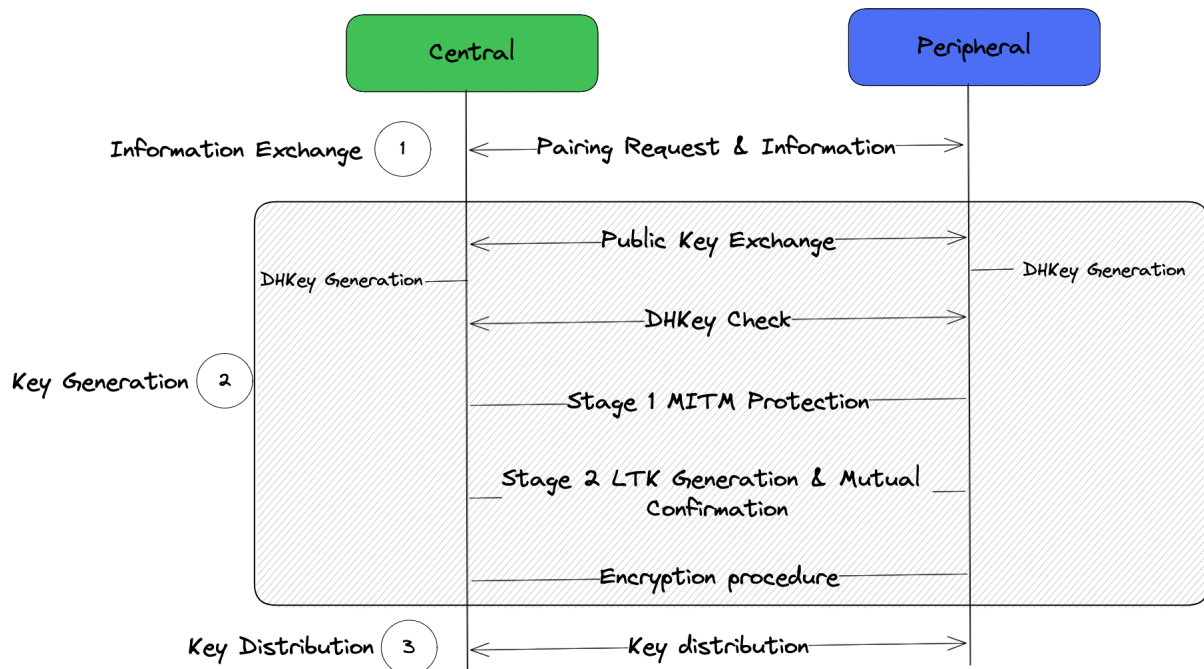
**Figure 2:** Overview pairing steps

### Pairing Modes

- *Legacy* uses a simple process of exchanging secret data to derive a symmetric key with which to encrypt the link during the key distribution phase.
- *Secure Connections* (SC) uses elliptic curve public key cryptography to allow a symmetric key to be derived. That key is used to encrypt the link during the key distribution phase.

The *Secure Connections* pairing mode is the “more secure approach” and was developed to solve the weaknesses of *Legacy* mode. However, the *zero LTK installation* CVE found that bad implementations of SC pairing allow to bypass security.

## Secure Connections pairing process



**Figure 3:** Secure Connections pairing procedure

**Phase 1 - Information Exchange** The central device sends a pairing request and both devices exchange on their security capabilities and requirements. This phase defines the pairing *mode*.

### Phase 2 - Key Generation

- **Public Key exchange:** An exchange of public keys is initiated by the Central device. Both peripheral and central verify that their received key is on the P-256 curve.
- **Calculating the DHKey:** Each device uses its own private key (SK) and the public key of the other device (PK) to calculate their Diffie-Hellman key (DHKey). This way both device possess the same DHKey value.

1	Central: $DHKey = p256(SK_c, PK_p)$
2	Peripheral: $DHKey = p256(SK_p, PK_c)$

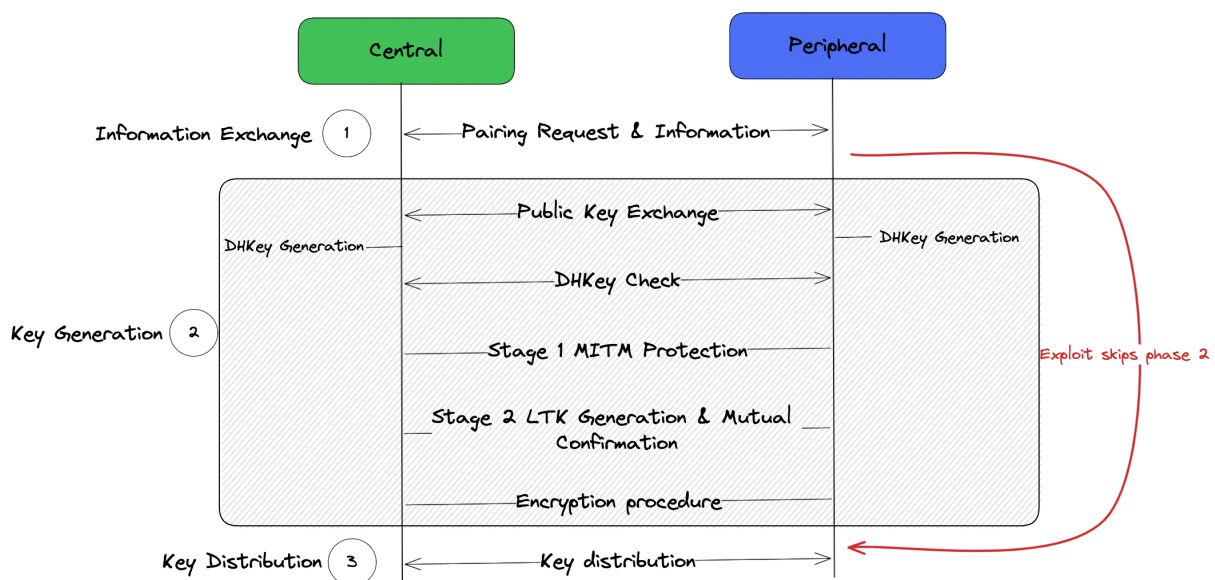
- If MITM protection was requested, an interactive procedure takes place to confirm the authenticity of the pairing devices.
- **Calculating the Long Term Key (LTK) and mutual confirmation:** The devices authenticate each other and calculate a LTK key. A session key is derived from the LTK in order to encrypt the link

before Phase 3.

**Phase 3 - Key Distribution** Through the encrypted link, the devices can distribute keys.

## Proof-of-Concept

The root cause of *zero LTK installation* is that there is no check of the state in which the two devices are in the pairing procedure. As such, a rogue central device can skip the key generation and authentication step. This results in a LTK key set at 0 in the peripheral device, and as such, an easily derivable session key.



**Figure 4:** Exploit skips phase 2 of Secure Connections pairing procedure

This is a fully theoretical Proof-of-Concept as I was not able to get a pcap of BLE discovery, advertisement or pairing procedures, nor did I have a BLE device with Telink SMP implementation.

You can read the code here: <https://github.com/louisabricot/writeup-CVE-2019-19194/blob/main/zero-LTK-installation.py>

## References

- Garbelini, M. E., Wang, C., Chattopadhyay, S., Sun, S., & Kurniawan, E. (2020, July). Sweyntooth: Unleashing mayhem over bluetooth low energy. In Proceedings of the 2020 USENIX Conference on Usenix Annual Technical Conference (pp. 911-925).
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- Bluetooth Core Specification 5.0
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