# MATH378 - Nonlinear Optimization

Based on lectures from Fall 2025 by Prof. Tim Hoheisel. Notes by Louis Meunier

### Contents

1 Preliminaries	
1.1 Terminology	
1.2 Convex Sets and Functions	
2 Unconstrained Optimization	4
2.1 Theoretical Foundations	4
2.1.1 Quadratic Approximation	6
2.2 Differentiable Convex Functions	7
2.3 Matrix Norms	9
3 Descent Methods	11
3.1 A General Line-Search Method	11
3.1.1 Global Convergence of Algorithm 3.1	
3.2 The Gradient Method	
3.3 Newton-Type Methods	
3.3.1 Convergence Rates and Landau Notation	
3.3.2 Newton's Method for Nonlinear Equations	
3.3.3 Newton's Method for Optimization Problem	
3.4 Quasi-Newton Methods	21
3.4.1 Direct Methods	22

## §1 Preliminaries

## §1.1 Terminology

We consider problems of the form

minimize 
$$f(x)$$
 subject to  $x \in X$ , (†)

with  $X \subset \mathbb{R}^n$  the feasible region with x a feasible point, and  $f: X \to \mathbb{R}$  the objective (function); more concisely we simply write

$$\min_{x \in X} f(x)$$
.

When  $X = \mathbb{R}^n$ , we say the problem (†) is *unconstrained*, and conversely *constrained* when  $X \subseteq \mathbb{R}^n$ .

**⊗ Example 1.1** (Polynomial Fit): Given  $y_1, ..., y_m \in \mathbb{R}$  measurements taken at m distinct points  $x_1, ..., x_m \in \mathbb{R}$ , the goal is to find a degree  $\leq n$  polynomial  $q : \mathbb{R} \to \mathbb{R}$ , of the form

$$q(x) = \sum_{k=0}^{n} \beta_k x^k,$$

"fitting" the data  $\{(x_i, y_i)\}_i$ , in the sense that  $q(x_i) \approx y_i$  for each i. In the form of (†), we can write this precisely as

$$\min_{\beta \in \mathbb{R}^{n+1}} \frac{1}{2} \sum_{i=0}^{n} \left( \underbrace{\beta_n x_i^n + \dots + \beta_1 x_i + \beta_0}_{q(x_i)} - y_i \right)^2;$$

namely, we seek to minimize the  $\ell^2$ -distance between  $(q(x_i))$  and  $(y_i)$ . If we write

$$X \coloneqq \begin{pmatrix} 1 & x_1 & \dots & x_1^n \\ \vdots & \dots & \dots & \vdots \\ 1 & x_m & \dots & x_m^n \end{pmatrix} \in \mathbb{R}^{m \times (n+1)}, \qquad y \coloneqq \begin{pmatrix} y_1 \\ \vdots \\ y_m \end{pmatrix} \in \mathbb{R}^m,$$

then concisely this problem is equivalent to

$$\min_{\beta \in \mathbb{R}^{n+1}} \frac{1}{2} \|X \cdot \beta - y\|_2^2,$$

a so-called *least-squares* problem.

We have two related tasks:

- 1. Find the optimal value asked for by (†), that is what  $\inf_X f$  is;
- 2. Find a specific point  $\overline{x}$  such that  $f(\overline{x}) = \inf_X f$ , i.e. the value of a point

$$\overline{x} \in \operatorname{argmin}_X f := \left\{ x \in X \mid f(x) = \inf_X f \right\}.$$

(noting that argmin should be viewed as a set-valued function, as there may be multiple admissible minimizers) Notice that if we can accomplish 2., we've accomplished 1. by computing  $f(\overline{x})$ .

1.1 Terminology

Note that  $\overline{x} \in \operatorname{argmin}_X f \Rightarrow f(\overline{x}) = \inf_X f$ , but  $\inf_X f \in \mathbb{R}$  does not necessarily imply  $\operatorname{argmin}_X f \neq \emptyset$ , that is, there needn't be a feasible minimimum; for instance  $\inf_{x \in \mathbb{R}} e^x = 0$ , but  $\operatorname{argmin}_{\mathbb{R}} f = \emptyset$  (there is no x for which  $e^x = 0$ ).

- $\hookrightarrow$  **Definition 1.1** (Minimizers): Let  $X \subset \mathbb{R}^n$  and  $f : \mathbb{R}^n \to \mathbb{R}$ . Then  $\overline{x} \in X$  is called a
- *global minimizer* (of f over X) if  $f(\overline{x}) \le f(x) \forall x \in X$ , or equivalently if  $\overline{x} \in \operatorname{argmin}_X f$ ;
- *local minimizer (of f over X)* if  $f(\overline{x}) \le f(x) \forall x \in X \cap B_{\varepsilon}(\overline{x})$  for some  $\varepsilon > 0$ .

In addition, we have *strict* versions of each by replacing " $\leq$ " with "<".

```
\hookrightarrow Definition 1.2 (Some Geometric Tools): Let f : \mathbb{R}^n \to \mathbb{R}.
```

- gph  $f := \{(x, f(x)) \mid x \in \mathbb{R}^n\} \subseteq \mathbb{R}^n \times \mathbb{R}$
- $f^{-1}(\{c\}) := \{x \mid f(x) = c\} \equiv contour/level \ set \ at \ c$
- $\operatorname{lev}_c f := f^{-1}((-\infty, c]) = \{x \mid f(x) \le c\} \equiv lower \ level/sublevel \ set \ at \ c$

#### Remark 1.1:

- $lev_{inf} f = argmin f$
- assume *f* continuous; then all (sub)level sets are closed (possibly empty)

We recall the following result from calculus/analysis:

```
→Theorem 1.1 (Weierstrass): Let f : \mathbb{R}^n \to \mathbb{R} be continuous and X \subset \mathbb{R}^n compact. Then, \operatorname{argmin}_X f \neq \emptyset.
```

From, we immediately have the following:

**Proposition 1.1**: Let  $f : \mathbb{R}^n \to \mathbb{R}$  continuous. If there exists a  $c \in \mathbb{R}$  such that lev<sub>c</sub>f is nonempty and bounded, then  $\operatorname{argmin}_{\mathbb{R}^n} f \neq \emptyset$ .

PROOF. Since f continuous,  $\operatorname{lev}_c f$  is closed (being the inverse image of a closed set), thus  $\operatorname{lev}_c f$  is compact (and in particular nonempty). By Weierstrass, f takes a minimimum over  $\operatorname{lev}_c f$ , namely there is  $\overline{x} \in \operatorname{lev}_c f$  with  $f(\overline{x}) \leq f(x) \leq c$  for each  $x \in \operatorname{lev}_c f$ . Also, f(x) > c for each  $x \notin \operatorname{lev}_c f$  (by virtue of being a level set), and thus  $f(\overline{x}) \leq f(x)$  for each  $x \in \mathbb{R}^n$ . Thus,  $\overline{x}$  is a global minimizer and so the theorem follows.

#### §1.2 Convex Sets and Functions

**Definition 1.3** (Convex Sets):  $C \subset \mathbb{R}^n$  is *convex* if for any  $x, y \in C$  and  $\lambda \in (0, 1)$ ,  $\lambda x + (1 - \lambda)y \in C$ ; that is, the entire line between x and y remains in C.

1.2 Convex Sets and Functions

 $\hookrightarrow$  **Definition 1.4** (Convex Fucntions): Let  $C \subset \mathbb{R}^n$  be convex. Then,  $f: C \to \mathbb{R}$  is called

1. convex (on C) if

$$f(\lambda x + (1 - \lambda)y) \le \lambda f(x) + (1 - \lambda)f(y),$$

for every  $x, y \in C$  and  $\lambda \in (0, 1)$ ;

- 2. strictly convex (on C) if the inequality  $\leq$  is replaced with  $\leq$ ;
- 3. *strongly convex* (on *C*) if there exists a  $\mu > 0$  such that

$$f(\lambda x + (1 - \lambda)y) + \mu \lambda (1 - \lambda) ||x - y||^2 \le \lambda f(x) + (1 - \lambda)f(y),$$

for every  $x, y \in C$  and  $\lambda \in (0,1)$ ; we call  $\mu$  the modulus of strong convexity.

Remark 1.2:  $3. \Rightarrow 2. \Rightarrow 1.$ 

**Remark 1.3**: A function is convex iff its epigraph is a convex set.

**⊗ Example 1.2**: exp :  $\mathbb{R} \to \mathbb{R}$ , log :  $(0, \infty) \to \mathbb{R}$  are convex. A function  $f : \mathbb{R}^n \to \mathbb{R}^m$  of the form f(x) = Ax - b for  $A \in \mathbb{R}^{m \times n}$ ,  $b \in \mathbb{R}^m$  is called *affine linear*. For m = 1, every affine linear function is convex. All norms on  $\mathbb{R}^n$  are convex.

## $\hookrightarrow$ Proposition 1.2:

- 1. (Positive combinations) Let  $f_i$  be convex on  $\mathbb{R}^n$  and  $\lambda_i > 0$  scalars for i = 1, ..., m, then  $\sum_{i=1}^m \lambda_i f_i$  is convex; as long as one is strictly (resp. strongly) convex, the sum is strictly (resp. strongly) convex as well.
- 2. (Composition with affine mappings) Let  $f : \mathbb{R}^n \to \mathbb{R}$  be convex and  $G : \mathbb{R}^m \to \mathbb{R}^n$  be affine. Then,  $f \circ G$  is convex on  $\mathbb{R}^m$ .

# §2 Unconstrained Optimization

#### §2.1 Theoretical Foundations

We focus on the problem

$$\min_{x\in\mathbb{R}^n} f(x),$$

where  $f: \mathbb{R}^n \to \mathbb{R}$  is continuously differentiable.

**Definition 2.1** (Directional derivative): Let  $D \subset \mathbb{R}^n$  be open and  $f: D \to \mathbb{R}$ . We say f directionally differentiable at  $\overline{x} \in D$  in the direction  $d \in \mathbb{R}^n$  if

$$\lim_{t \to 0^+} \frac{f(\overline{x} + td) - f(\overline{x})}{t}$$

exists, in which case we denote the limit by  $f'(\bar{x}; d)$ .

2.1 Theoretical Foundations

**Lemma 2.1**: Let  $D \subset \mathbb{R}^n$  be open and  $f : D \to \mathbb{R}$  differentiable at  $x \in D$ . Then, f is directionally differentiable at x in every direction d, with

$$f'(x;d) = \nabla f(x)^T d = \langle \nabla f(x), d \rangle.$$

**Example 2.1** (Directional derivatives of the Euclidean norm): Let  $f : \mathbb{R}^n \to \mathbb{R}$  by f(x) = ||x|| the usual Euclidean norm. Then, we claim

$$f'(x;d) = \begin{cases} \frac{x^T d}{\|x\|} & x \neq 0 \\ \|d\| & x = 0 \end{cases}$$

For  $x \neq 0$ , this follows from the previous lemma and the calculation  $\nabla f(x) = \frac{x}{\|x\|}$ . For x = 0, we look at the limit

$$\lim_{t \to 0^+} \frac{f(0+td) - f(0)}{t} = \lim_{t \to 0^+} \frac{t||d|| - 0}{t} = ||d||,$$

using homogeneity of the norm.

**Lemma 2.2** (Basic Optimality Condition): Let *X* ⊂  $\mathbb{R}^n$  be open and  $f: X \to \mathbb{R}$ . If  $\overline{x}$  is a *local minimizer* of f over X and f is directionally differentiable at  $\overline{x}$ , then  $f'(\overline{x};d) \ge 0$  for all  $d \in \mathbb{R}^n$ .

PROOF. Assume otherwise, that there is a direction  $d \in \mathbb{R}^n$  for which the  $f'(\overline{x};d) < 0$ , i.e.

$$\lim_{t \to 0^+} \frac{f(\overline{x} + td) - f(\overline{x})}{t} < 0.$$

Then, for all sufficiently small t > 0, we must have

$$f(\overline{x} + td) < f(\overline{x}).$$

Moreover, since X open, then for t even smaller (if necessary),  $\overline{x} + td$  remains in X, thus  $\overline{x}$  cannot be a local minimizer.

**→Theorem 2.1** (Fermat's Rule): In addition to the assumptions of the previous lemma, assume further that f is differentiable at  $\overline{x}$ . Then,  $\nabla f(\overline{x}) = 0$ .

PROOF. From the previous, we know  $0 \le f'(\overline{x}; d)$  for any d. Take  $d = -\nabla f(\overline{x})$ , then using the representation of a directional derivative for a differentiable function, and the fact that norms are nonnegative,

$$0 \le -\|\nabla f(\overline{x})\|^2 \le 0,$$

which can only hold if  $\|\nabla f(\overline{x})\| = 0$  hence  $\nabla f(\overline{x}) = 0$ 

We recall the following from Calculus:

2.1 Theoretical Foundations

5

**Theorem 2.2** (Taylor's, Second Order): Let  $f : D \subset \mathbb{R}^n \to \mathbb{R}$  be twice continuously differentiable, then for each  $x, y \in D$ , there is an  $\eta$  lying on the line between x and y such that

$$f(y) = f(x) + \nabla f(x)^{T} (y - x) + \frac{1}{2} (y - x)^{T} \nabla^{2} f(\eta) (y - x).$$

**Theorem 2.3** (2nd-order Optimality Conitions): Let  $X \subseteq \mathbb{R}^n$  open and  $f: X \to \mathbb{R}$  twice continuously differentiable. Then, if x a local minimizer of f over X, then the Hessian matrix  $\nabla^2 f(x)$  is positive semi-definite.

PROOF. Suppose not, then there exists a d such that  $d^T \nabla^2 f(x) d < 0$ . By Taylor's, for every t > 0, there is an  $\eta_t$  on the line between x and x + td such that

$$f(x+td) = f(x) + t \underbrace{\nabla f(x)^T}_{=0} d + \frac{1}{2} t^2 d^T \nabla^2 f(\eta_t) d$$
$$= f(x) + \frac{t^2}{d^T} \nabla^2 f(\eta_t) d.$$

As  $t \to 0^+$ ,  $\nabla^2 f(\eta_t) \to \nabla^2 f(x) < 0$ . By continuity, for t sufficiently small,  $\frac{t^2}{2} d^T \nabla^2 f(\eta_t) d < 0$  for t sufficiently small, whence we find

$$f(x+td) < f(x),$$

for sufficiently small t, a contradiction.

**Lemma 2.3**: Let  $X \subset \mathbb{R}^n$  open,  $f: X \to \mathbb{R}$  in  $C^2$ . If  $\overline{x} \in \mathbb{R}^n$  is such that  $\nabla^2 f(\overline{x}) > 0$  (i.e. is positive definite), then there exists  $\varepsilon, \mu > 0$  such that  $B_\varepsilon(\overline{x}) \subset X$  and

$$d^T \nabla^2 f(x) d \geq \mu \|d\|^2, \qquad \forall d \in \mathbb{R}^n, x \in B_\varepsilon(\overline{x}).$$

Combining this and Taylor's Theorem, we can deduce the following (our first "sufficient" result of this section):

**Theorem 2.4** (Sufficient Optimality Condition): Let  $X \subset \mathbb{R}^n$  open and  $f \in C^2(X)$ . Let  $\overline{x}$  be a stationary point of f such that  $\nabla^2 f(\overline{x}) > 0$ . Then,  $\overline{x}$  is a *strict* local minimizer of f.

## 2.1.1 Quadratic Approximation

Let  $f : \mathbb{R}^n \to \mathbb{R}$  be  $C^2$  and  $\overline{x} \in \mathbb{R}^n$ . By Taylor's, we can approximate

$$f(y) \approx g(y) \coloneqq f(\overline{x}) + \nabla f(\overline{x})^T (y - \overline{x}) + \frac{1}{2} (y - \overline{x})^T \nabla^2 f(\overline{x}) (y - \overline{x}).$$

**Example 2.2** (Quadratic Functions): For  $Q \in \mathbb{R}^{n \times n}$  symmetric,  $c \in \mathbb{R}^n$  and  $\gamma \in \mathbb{R}$ , let

$$f: \mathbb{R}^n \to \mathbb{R}, \quad f(x) = \frac{1}{2} x^T Q x + c^T x + \gamma,$$

a typical quadratic function. Then,

$$\nabla f(x) = \frac{1}{2} \big( Q + Q^T \big) x + c = Qx + c, \qquad \nabla^2 f(x) = Q.$$

We find that f has no minimizer if  $c \notin \operatorname{rge}(Q)$  or Q is not positive semi-definite, combining our previous two results. In turn, if Q is positive definite (and thus invertible), there is a unique local minimizer  $\overline{x} = -Q^{-1}c$  (and global minimizer, as we'll see).

## §2.2 Differentiable Convex Functions

 $\hookrightarrow$  Theorem 2.5: Let  $C \subset \mathbb{R}^n$  be open and convex and  $f: C \to \mathbb{R}$  differentiable on C. Then:

1. *f* is convex (on *C*) iff

$$f(x) \ge f(\overline{x}) + \nabla f(\overline{x})^T (x - \overline{x})$$
 \*1

for every  $x, \overline{x} \in C$ ;

- 2. *f* is *strictly* convex iff same inequality as 1. with strict inequality;
- 3. f is *strongly* convex with modulus  $\sigma > 0$  iff

$$f(x) \geq f(\overline{x}) + \nabla f(\overline{x})^T (x - \overline{x}) + \frac{\sigma}{2} \|x - \overline{x}\|^2 \qquad \star_2$$

for every  $x, \overline{x} \in C$ .

PROOF.  $(1., \Rightarrow)$  Let  $x, \overline{x} \in C$  and  $\lambda \in (0, 1)$ . Then,

$$f(\lambda x + (1-\lambda)\overline{x}) - f(\overline{x}) \le \lambda \big(f(x) - f(\overline{x})\big),$$

which implies

$$\frac{f(\overline{x}+\lambda(x-\overline{x}))-f(\overline{x})}{\lambda}\leq f(x)-f(\overline{x}).$$

Letting  $\lambda \to 0^+$ , the LHS  $\to$  the directional derivative of f at  $\overline{x}$  in the direction  $x - \overline{x}$ , which is equal to, by differentiability of f,  $\nabla f(\overline{x})^T(x - \overline{x})$ , thus the result.

$$(1., \Leftarrow)$$
 Let  $x_1, x_2 \in C$  and  $\lambda \in (0, 1)$ . Let  $\overline{x} := \lambda x_1 + (1 - \lambda)x_2$ .  $\star_1$  implies

$$f(x_i) \ge f(\overline{x}) + \nabla f(\overline{x})^T (x_i - \overline{x}),$$

for each of i=1,2. Taking "a convex combination of these inequalities", i.e. multiplying them by  $\lambda$ ,  $1-\lambda$  resp. and adding, we find

$$\lambda f(x_1) + (1-\lambda)f(x_2) \geq f(\overline{x}) + \nabla f(\overline{x})^T \big(\lambda x_1 + (1-\lambda)x_2 - \overline{x}\big) = f\big(\lambda x_1 + (1-\lambda)x_2\big),$$

thus proving convexity.

 $(2., \Rightarrow)$  Let  $x \neq \overline{x} \in C$  and  $\lambda \in (0, 1)$ . Then, by 1., as we've just proven,

$$\lambda \nabla f(\overline{x})^T (x - \overline{x}) \leq f(\overline{x} + \lambda (x - \overline{x})) - f(\overline{x}).$$

But  $f(\overline{x} + \lambda(x - \overline{x})) < \lambda f(x) + (1 - \lambda)f(\overline{x})$  by strict convexity, so we have

$$\lambda \nabla f(\overline{x})^T (x - \overline{x}) < \lambda \big( f(x) - f(\overline{x}) \big),$$

and the result follows by dividing both sides by  $\lambda$ .

- $(2., \Leftarrow)$  Same as  $(1., \Leftarrow)$  replacing " $\leq$ " with "<".
- (3.) Apply 1. to  $f \frac{\sigma}{2} \|\cdot\|^2$ , which is still convex if f  $\sigma$ -strongly convex, as one can check.
- $\hookrightarrow$  Corollary 2.1: Let  $f: \mathbb{R}^n \to \mathbb{R}$  be convex and differentiable. Then,
- a) there exists an *affine function*  $g : \mathbb{R}^n \to \mathbb{R}$  such that  $g(x) \le f(x)$  everywhere;
- b) if f strongly convex, then it is coercive, i.e.  $\lim_{\|x\|\to\infty} f(x) = \infty$ .
- $\hookrightarrow$  Corollary 2.2: Let  $f : \mathbb{R}^n \to \mathbb{R}$  be convex and differentiable, then TFAE:
- 1.  $\bar{x}$  is a global minimizer of f;
- 2.  $\overline{x}$  is a local minimizer of f;
- 3.  $\overline{x}$  is a stationary point of f.

PROOF. 1.  $\Rightarrow$  2. is trivial and 2.  $\Rightarrow$  3. was already proven and 3.  $\Rightarrow$  1. follows from the fact that differentiability gives

$$f(x) \ge f(\overline{x}) + \underline{\nabla(f)(\overline{x})^T(x-\overline{x})}$$

for any  $x \in \mathbb{R}^n$ .

**Corollary 2.3**: (2.2.4)

- **→Theorem 2.6** (Twice Differentiable Convex Functions): Let  $Ω ⊂ \mathbb{R}^n$  open and convex and  $f ∈ C^2(Ω)$ . Then,
- 1. f is convex on  $\Omega$  iff  $\nabla^2 f \ge 0$ ;
- 2. f is strictly convex on  $\Omega \leftarrow \nabla^2 f > 0$ ;
- 2. f is  $\sigma$ -strongly convex on  $\Omega \Leftrightarrow \sigma \leq \lambda_{\min}(\nabla^2 f(x))$  for all  $x \in \Omega$ .
- **Corollary 2.4**: Let  $A \in \mathbb{R}^{n \times n}$  be symmetric,  $b \in \mathbb{R}^n$  and  $f(x) := \frac{1}{2}x^TAx + b^Tx$ . Then,
- 1. f convex  $\Leftrightarrow A \ge 0$ ;
- 2. f strongly convex  $\Leftrightarrow A > 0$ .

**Theorem 2.7** (Convex Optimization): Let  $f : \mathbb{R}^n \to \mathbb{R}$  be convex and continuous,  $X \subset \mathbb{R}^n$  convex (and nonempty), and consider the optimization problem

$$\min f(x)$$
 s.t.  $x \in X$   $(\star)$ .

Then, the following hold:

- 1.  $\overline{x}$  is a global minimizer of  $(\star) \Leftrightarrow \overline{x}$  is a local minimizer of  $(\star)$
- 2.  $\operatorname{argmin}_X f$  is convex (possibly empty)
- 3. f is strictly convex  $\Rightarrow$  argmin<sub>X</sub>f has at *most* one element
- 4. f is strongly convex and differentiable, and X closed,  $\Rightarrow$  argmin<sub>X</sub>f has exactly one element

PROOF.  $(1., \Rightarrow)$  Trivial.  $(1., \Leftarrow)$  Let  $\overline{x}$  be a local minimizer of f over X, and suppose towards a contradiction that there exists some  $\hat{x} \in X$  such that  $f(\hat{x}) < f(\overline{x})$ . By convexity of f, X, we know for  $\lambda \in (0,1)$ ,  $\lambda \overline{x} + (1-\lambda)\hat{x} \in X$  and

$$f(\lambda \overline{x} + (1 - \lambda)\hat{x}) \le \lambda f(\overline{x}) + (1 - \lambda)f(\hat{x}) < f(\overline{x}).$$

Letting  $\lambda \to 1^-$ , we see that  $\lambda \overline{x} + (1 - \lambda)\hat{x} \to \overline{x}$ ; in particular, for any neighborhood of  $\overline{x}$  we can construct a point which strictly lower bounds  $f(\overline{x})$ , which contradicts the assumption that  $\overline{x}$  a local minimizer.

- (2.) and (3.) are left as an exercise.
- (4.) We know that f is strictly convex and level-bounded. By (3.) we know there is at most one minimizer, so we just need to show there exists one. Take  $c \in \mathbb{R}$  such that  $\text{lev}_c(f) \cap X \neq \emptyset$  (which certainly exists by taking, say, f(x) for some  $x \in X$ ). Then, notice that  $(\star)$  and

$$\min_{x \in \text{lev}_c f \cap X} f(x) \qquad (\star \star)$$

have the same solutions i.e. the same set of global minimizers (noting that this remains a convex problem). Since f continuous and  $\text{lev}_c f \cap X$  compact and nonempty, f attains a minimum on  $\text{lev}_c f \cap X$ , as we needed to show.

**Remark 2.1**: Note that level sets of convex functions are convex, this is left as an exercise.

#### §2.3 Matrix Norms

We denote by  $\mathbb{R}^{m \times n}$  the space of real-valued  $m \times n$  matrices (i.e. of linear operators from  $\mathbb{R}^n \to \mathbb{R}^m$ ).

 $\hookrightarrow$  Proposition 2.1 (Operator Norms): Let  $\|\cdot\|_*$  be a norm on  $\mathbb{R}^m$  and  $\mathbb{R}^n$ , resp. Then, the map

$$\mathbb{R}^{m \times n} \ni A \mapsto \|A\|_* \coloneqq \sup_{\substack{x \in \mathbb{R}^n, \\ \|x\|_* \neq 0}} \frac{\|Ax\|_*}{\|x\|_*} \in \mathbb{R}$$

is a norm on  $R^{m \times n}$ . In addition,

$$||A||_* = \sup_{||x||_*=1} ||Ax||_* = \sup_{||x||_* \le 1} ||Ax||_*.$$

2.3 Matrix Norms 9

PROOF. We first note that all of these sup's are truely max's since they are maximizing continuous functions over compact sets.

Let  $A \in \mathbb{R}^{m \times n}$ . The first "In addition" equality follows from positive homogeneity, since  $\frac{x}{\|x\|_*}$  a unit vector. For the second, note that " $\leq$ " is trivial, since we are supping over a larger (super)set. For " $\geq$ ", we have for any x with  $\|x\|_* \leq 1$ ,

$$||Ax||_* = ||x||_* ||A\frac{x}{||x||_*}||_* \le ||A\frac{x}{||x||_*}||.$$

Supping both sides over all such *x* gives the result.

We now check that  $\|\cdot\|_*$  actually a norm on  $\mathbb{R}^{m\times n}$ .

- $1. \ \|A\|_* = 0 \Leftrightarrow \sup_{\|x\|_* = 1} \|Ax\|_* = 0 \Leftrightarrow \|Ax\|_* = 0 \forall \|x\|_* = 1 \Leftrightarrow Ax = 0 \forall \|x\|_* = 1 \Leftrightarrow A = 0$
- 2. For  $\lambda \in \mathbb{R}$ ,  $A \in \mathbb{R}^{m \times n}$ ,  $\|\lambda A\|_* = \sup \|\lambda Ax\|_* = |\lambda| \cdot \sup \|Ax\|_* = |\lambda| \|A\|_*$
- 3. For  $A, B \in \mathbb{R}^{m \times n}$ ,  $||A + B||_* \le ||A||_* + ||B||_*$  using properties of sups of sums

**Proposition 2.2**: Let  $A = (a_{ij})_{i=1,...,m} \in \mathbb{R}^{m \times n}$ , then: j=1,...,n

- 1.  $||A||_1 = \max_{j=1}^n \sum_{i=1}^m |a_{ij}|$
- 2.  $||A||_2 = \sqrt{\lambda_{\max}(A^T A)} = \sigma_{\max}(A)$
- 3.  $||A||_{\infty} = \max_{i=1}^{m} \sum_{i=1}^{n} |a_{ij}|$

 $\hookrightarrow$  Proposition 2.3: Let  $\|\cdot\|_*$  be a norm on  $\mathbb{R}^n$ ,  $\mathbb{R}^m$ , and  $\mathbb{R}^p$ . For  $A \in \mathbb{R}^{m \times n}$  and  $B \in \mathbb{R}^{n \times p}$ ,

- 1.  $||Ax||_* \le ||A||_* \cdot ||x||_*$
- 2.  $||AB||_{*} \leq ||A||_{*} \cdot ||B||_{*}$

**Proposition 2.4** (Banach Lemma): Let  $C \in \mathbb{R}^{n \times n}$  with ||C|| < 1, where  $||\cdot||$  submultiplicative. Then, I + C is invertible, and

$$||(1+C)^{-1}|| \le \frac{1}{1-||C||}.$$

Proof. We have for any m,

$$\left\| \sum_{i=1}^{m} (-C)^{i} \right\| \leq \sum_{i=1}^{m} \|C\|^{i} \underset{m \to \infty}{\longrightarrow} \frac{1}{1 - \|C\|}.$$

Hence,  $A_m := \sum_{i=1}^m (-C)^i$  a sequence of matrices with bounded norm uniformly in m, and thus has a converging subsequence, so wlog  $A_m \to A \in \mathbb{R}^{n \times n}$  (by relabelling). Moreover, observe that

$$A_m \cdot (I+C) = \sum_{i=0}^m (-C)^i (I+C) = \sum_{i=0}^m \left[ (-C)^i - (-C)^{i+1} \right] = (-C)^0 - (-C)^{m+1} = I - (-C)^{m+1}.$$

Now,  $||C^{m+1}|| \le ||C||^{m+1} \to 0$ , since ||C|| < 1, thus  $C \to 0$ . Hence, taking limits in the line above implies

2.3 Matrix Norms 10

$$A(I+C) = \lim_{m \to \infty} A_m(I+C) = I,$$

implying A the inverse of (I + C), proving the proposition.

**Corollary 2.5**: Let  $A, B \in \mathbb{R}^{n \times n}$  with ||I - BA|| < 1 for  $||\cdot||$  submultiplicative. Then, A and B are invertible, and  $||B^{-1}|| \le \frac{||A||}{1 - ||I - BA||}$ .

## §3 DESCENT METHODS

## §3.1 A General Line-Search Method

We deal with the unconstrained problem

$$\min_{x \in \mathbb{R}^n} f(x) \qquad (\star).$$

**Definition 3.1** (Descent Direction): Let  $f : \mathbb{R}^n \to \mathbb{R}$ ,  $x \in \mathbb{R}^n$ .  $d \in \mathbb{R}^n$  is a *descent direction* of f at x if there exists a  $\bar{t} > 0$  such that f(x + td) < f(x) for all  $t \in (0, \bar{t})$ .

**Proposition 3.1**: If  $f : \mathbb{R}^n \to \mathbb{R}$  is directionally differentiable at  $x \in \mathbb{R}^n$  in the direction d with f'(x;d) < 0, then d a descent direction of f at x; in particular if f differentiable at x, then true for d if  $\nabla f(x)^T d < 0$ .

**Corollary 3.1**: Let  $f : \mathbb{R}^n \to \mathbb{R}$  differentiable,  $B \in \mathbb{R}^{n \times n}$  positive definite, and  $x \in \mathbb{R}^n$ . Then  $\nabla f(x) \neq 0 \Rightarrow -B\nabla f(x)$  is a descent direction of f at x.

PROOF. 
$$\nabla f(x)^T (-B\nabla f(x)) = -\nabla f(x)^T B\nabla f(x) < 0.$$

# A generic method/strategy for solving ( $\star$ ):

- S1. (Initialization) Choose  $x^0 \in \mathbb{R}^n$  and set k := 0
- S2. (Termination) If  $x^k$  satisfies a "termination criterion", STOP
- S3. (Search direction) Determine  $d^k$  such that  $\nabla f(x^k)^T d^k < 0$
- S4. (Step-size) Determine  $t_k > 0$  such that  $f(x^k + t_k d^k) < f(x^k)$
- S5. (Update) Set  $x^{k+1} := x^k + t_k d^k$ , iterate k, and go back to step 2.

**Remark 3.1**: a) The generic choice for  $d^k$  in 3. is just  $d^k := -B_k \nabla f(x^k)$  for some  $B_k > 0$ . We focus on:

- $B_k = I$  (gradient-descent)
- $B_k = \nabla^2 f(x^k)^{-1}$  (Newton's method)  $B_k \approx \nabla^2 f(x^k)^{-1}$  (quasi Newton's method)
- b) Step 4. is called *line-search*, since  $t_k > 0$  determined by looking at

$$0 < t \mapsto f(x^k + td^k),$$

i.e. along the (half)line t > 0.

- c) Executing Step 4. is a trade-off between
  - (i) decreasing f along  $x^k + td^k$  as much as possible;
  - (ii) keeping computational efforts low.

For instance, the exact minimization rule  $t_k = \operatorname{argmin}_{t>0} f\left(x_k + td^k\right)$  overemphasizes (i) over (ii).

 $\hookrightarrow$  **Definition 3.2** (Step-size rule): Let  $f \in C^1(\mathbb{R}^n)$  and

$$\mathcal{A}_f \coloneqq \big\{ (x,d) \mid \nabla f(x)^T d < 0 \big\}.$$

A (possible set-valued) map

$$T:(x,d)\in \mathcal{A}_f\mapsto T(x,d)\in \mathbb{R}_+$$

is called a *step-size rule* for *f* .

If T is well-defined for all  $C^1$ -functions, we say T well-defined.

# 3.1.1 Global Convergence of Algorithm 3.1

 $\hookrightarrow$  **Definition 3.3** (Efficient step-size): Let  $f \in C^1(\mathbb{R}^n)$ . The step-size rule T is called *efficient* for *f* if there exists  $\theta > 0$  such that

$$f(x+td) \le f(x) - \theta \left(\frac{\nabla f(x)^T d}{\|d\|}\right)^2, \quad \forall t \in T(x,d), (x,d) \in A_f.$$

**Theorem 3.1**: Let  $f \in C^1(\mathbb{R}^n)$ . Let  $\{x^k\}, \{d^k\}, \{t_k\}$  be generated by Algorithm 3.1. Assume the following:

- 1.  $\exists c > 0$  such that  $-\left(\nabla f(x^k)^T d^k\right) / \left(\|\nabla f(x^k)\| \cdot \|d^k\|\right) \ge c$  for all k (this is called the *angle* condition), and
- 2. there exists  $\theta > 0$  such that  $f(x^k + t_k d^k) \le f(x^k) \theta \cdot (\nabla f(x^k)^T d^k / ||d^k||)^2$  for all k (which is satisfied if  $t_k \in T(x^k, d^k)$  for an efficient T).

Then, every cluster point of  $\{x^k\}$  is a stationary point of f.

Proof. By condition 2., there is  $\theta > 0$  such that

$$f(x^{k+1}) \le f(x^k) - \theta \left(\frac{\nabla f(x^k)^T d^k}{\|d^k\|}\right)^2$$

for all  $k \in \mathbb{N}$ . By 1., we know

$$\left(\frac{\nabla f(x^k)^T d^k}{\|d^k\|}\right)^2 \ge c^2 \|\nabla f(x^k)\|^2.$$

Put  $\kappa := \theta c^2$ , then these two inequalities imply

$$f(x^{k+1}) \le f(x^k) - \kappa \cdot \|\nabla f(x^k)\|^2$$
. (\*)

Let  $\overline{x}$  be a cluster point of  $\{x^k\}$ . As  $\{f(x^k)\}$  is monotonically decreasing (by construction in the algorithm), and has cluster point  $f(\overline{x})$  by continuity, it follows that  $f(x_k) \to f(\overline{x})$  along the whole sequence. In particular,  $f(x^{k+1}) - f(x^k) \to 0$ ; thus, from (\*),

$$0 \le \kappa \left\| \nabla f(x^k) \right\|^2 \le f(x^k) - f(x^{k+1}) \to 0,$$

and thus  $\nabla f(x^k) \to \nabla f(\overline{x}) = 0$ , so indeed  $\overline{x}$  a stationary point of f.

## §3.2 The Gradient Method

We specialize Algorithm 3.1 here. Specifically, we'll take

$$d^k := -\nabla f(x^k);$$

it's know that

$$\frac{-\nabla f(x^k)}{\|\nabla f(x^k)\|} = \operatorname{argmin}_{d:\|d\| \le 1} \nabla f(x^k)^T d,$$

with  $\|\cdot\|$  the 2 norm.

We use a step-size rule called "Armijo rule". Choose parameters  $\beta$ ,  $\sigma \in (0,1)$ . For  $(x,d) \in A_f$ , we define our step-size rule by

$$T_A(x,d) \coloneqq \max_{\ell \in \mathbb{N}_0} \left\{ \beta^\ell \mid \underbrace{f(x+\beta^\ell d) \leq f(x) + \beta^\ell \sigma \nabla f(x)^T d}_{\text{"Armijo condition"}} \right\}.$$

For instance, consider  $f(x) = (x-1)^2 - 1$ . The minimum of this function is  $f^* = -1$ . Choose  $x^k := \frac{1}{k}$ , then

$$f(x^k) = \frac{2k+1}{k^2} \to 0 \neq f^*,$$

even though  $f(x^{k+1}) - f(x^k) < 0$ ; we don't actually reach the right stationary point with our chosen step size.

3.2 The Gradient Method 13

**Example 3.1** (Illustration of Armijo Rule): For (x,d) ∈  $A_f$  and f smooth on  $\mathbb{R}^n$ , defined  $\phi$ :  $\mathbb{R} \to \mathbb{R}$ ,  $\phi(t) := f(x+td)$ . The map  $t \mapsto \sigma \phi'(0)t + \phi(0) = \sigma t \nabla f(x)^T d + \phi(0)$ 

**Proposition 3.2**: Let f :  $\mathbb{R}^n$  →  $\mathbb{R}$  be differentiable with  $\beta$ ,  $\sigma \in (0,1)$ . Then for  $(x,d) \in A_f$ , there exists  $\ell \in \mathbb{N}_0$  such that

$$f(x + \beta^{\ell} d) \le f(x) + \beta^{\ell} \sigma \nabla f(x)^{T} d,$$

i.e.  $T_A(x,d) \neq \emptyset$ .

Proof. Suppose not, i.e.

$$\frac{f(x + \beta^{\ell} d) - f(x)}{\beta^{\ell}} > \sigma \nabla f(x)^{T} d, \forall \ell \in \mathbb{N}_{0}.$$

Letting  $\ell \to \infty$ , the left-hand side converges to  $\nabla f(x)^T d$ , so

$$\nabla f(x)^T d \ge \sigma \nabla f(x)^T d.$$

But  $(x, d) \in A_f$ , so  $\nabla f(x)^T d < 0$  so dividing both sides of this inequality by this quantity, this implies  $\sigma \le 0$ , which is a contradiction.

We now prove convergence of an algorithm based on the Armijo Rule:

# Gradient Descent with Armijo Rule

S0. Choose  $x^0 \in \mathbb{R}^n$ ,  $\sigma$ ,  $\beta \in (0,1)$ ,  $\varepsilon \ge 0$ , and set k := 0

S1. If  $\|\nabla f(x^k)\| \le \varepsilon$ , STOP

S2. Set  $d^k := -\nabla f(x^k)$ 

S3. Determine  $t_k > 0$  by

$$t_k = T_A(x, d)$$

as defined above.

S4. Set  $x^{k+1} = x^k + t_k d^k$ , iterate k and go to S1.

**Lemma 3.1**: Let  $f \in C^1(\mathbb{R}^n)$ ,  $x^k \to x$ ,  $d^k \to d$  and  $t_k \downarrow 0$ . Then

$$\lim_{k \to \infty} \frac{f\left(x^k + t_k d^k\right) - f\left(x^k\right)}{t^k} = \nabla f(x)^T d.$$

Proof. Left as an exercise.

**→Theorem 3.2**: Let  $f \in C^1(\mathbb{R}^n)$ . Then every cluster point of a sequence  $\{x^k\}$  generated by Algorithm 3.2 is a stationary point of f.

PROOF. Let  $\overline{x}$  be a cluster point of  $\left\{x^k\right\}$  and let  $x^k \underset{k \in K}{\to} \overline{x}$ , K an infinite subset of  $\mathbb{N}$ . Assume towards a contradiction  $\nabla f(\overline{x}) \neq 0$ . As  $f\left(x^k\right)$  is monotonically decreasing with cluster point  $f(\overline{x})$ , it must be that  $f\left(x^k\right) \to f(\overline{x})$  along the whole sequence so  $f\left(x^{k+1}\right) - f\left(x^k\right) \to 0$ . Thus,

3.2 The Gradient Method 14

$$0 \le t_k \|\nabla f(x^k)\|^2 \stackrel{\text{S2}}{=} -t_k \nabla f(x^k)^T d^k \stackrel{\text{S3}}{\leq} \frac{f(x^k) - f(x^{k+1})}{\sigma} \to 0.$$

Thus,  $0 = \lim_{k \in K} t_k \|\nabla f(x^k)\| = \|\nabla f(\overline{x})\| \lim_{k \in K} t_k$ . We assumed  $\overline{x}$  not a stationary point, so it follows that  $t_k \underset{k \in K}{\longrightarrow} 0$ . By S3, for  $\beta^{\ell_k} = t_k$ ,

$$\frac{f\left(x^k+\beta^{\ell_k-1}d^k\right)-f\left(x^k\right)}{\beta^{\ell_k-1}}>\sigma\nabla f\left(x^k\right)^Td^k.$$

Letting  $k \to \infty$  along *K*,the LHS converges to, by the previous lemma, to

$$\nabla f(\overline{x})^T d = -\nabla f(\overline{x})^T \nabla f(\overline{x}) = -\|\nabla f(\overline{x})\|^2,$$

and the RHS converges to  $\sigma \|\nabla f(\overline{x})\|^2$ , which implies

$$-\|\nabla f(\overline{x})\|^2 \ge \sigma \|\nabla f(\overline{x})\|^2,$$

which implies  $\sigma$  negative, a contradiction.

**Remark 3.2**: The proof above shows, the following: Let  $\{x^k\}$  such that  $x^{k+1} := x^k + t_k d^k$  for  $d^k \in \mathbb{R}^n$ ,  $t_k > 0$ , and let  $f(x^{k+1}) \le f(x^k)$  and  $x^k \xrightarrow{K} \overline{x}$  such that  $d^k = -\nabla f(x^k)$ ,  $t_k = T_A(x^k, d^k)$  for all  $k \in K$ . Then  $\nabla f(\overline{x}) = 0$ ; i.e., all of the "focus" is on the subsequence along K. The only time we needed the whole sequence was to use the fact that  $f(x^k) \to f(\overline{x})$  along the whole sequence.

## §3.3 Newton-Type Methods

## 3.3.1 Convergence Rates and Landau Notation

 $\hookrightarrow$  **Definition 3.4**: Let  $\{x^k \in \mathbb{R}^n\}$  converge to  $\overline{x}$ . Then,  $\{x^k\}$  converges:

1. *linearly* to  $\overline{x}$  if there exists  $c \in (0,1)$  such that

$$||x^{k+1} - \overline{x}|| \le c||x^k - \overline{x}||, \forall k;$$

2. *superlinearly* to  $\overline{x}$  if

$$\lim_{k \to \infty} \frac{\left\| x^{k+1} - \overline{x} \right\|}{\left\| x^k - \overline{x} \right\|} = 0;$$

3. *quadratically* to  $\overline{x}$  if there exists C > 0 such that

$$||x^{k+1} - \overline{x}|| \le C||x^k - \overline{x}||^2, \forall k.$$

Remark 3.3:  $3. \Rightarrow 2. \Rightarrow 1.$ 

**Remark 3.4**: We needn't assume  $x^k \to \overline{x}$  for the first two definitions; their statements alone imply convergence. However, the last does not; there exists sequences with this property that do not converge.

 $\hookrightarrow$  **Definition 3.5** (Landau Notation): Let {*a<sub>k</sub>*}, {*b<sub>k</sub>*} be positive sequences ↓ 0. Then,

1. 
$$a_k = o(b_k) \Leftrightarrow \lim_{k \to \infty} \frac{a_k}{b_k} = 0$$
;

2.  $a_k = O(b_k) \Leftrightarrow \exists C > 0 : a_k \leq Cb_k$  for all k (sufficiently large).

# **Remark 3.5**: If $x^k \to \overline{x}$ , then

- 1. the convergence is superlinear  $\Leftrightarrow ||x^{k+1} \overline{x}|| = o(||x^k \overline{x}||);$ 2. the convergence is quadratic  $\Leftrightarrow ||x^{k+1} \overline{x}|| = O(||x^k \overline{x}||^2).$

## 3.3.2 Newton's Method for Nonlinear Equations

We consider the nonlinear equation

$$F(x) = 0, \qquad (*)$$

where  $F: \mathbb{R}^n \to \mathbb{R}^n$  is smooth (continuously differentiable). Our goal is to find a numerical scheme that can determine approximate zeros of *F*, i.e. solutions to (\*). The idea of Newton's method for such a problem, is, given  $x^k \in \mathbb{R}^n$ , to consider the (affine) linear approximation of *F* about  $x^k$ ,

$$F_k: x \mapsto F(x^k) + F'(x^k)(x - x^k),$$

where F' the Jacobian of F. Then, we compute  $x^{k+1}$  as a solution of  $F_k(x) = 0$ . Namely, if  $F'(x^k)$ invertible, then solving for  $F_k(x^{k+1}) = 0$ , we find

$$x^{k+1} = x^k - F'(x^k)^{-1}F(x^k).$$

More generally, one solves  $F'(x^k)d = -F(x^k)$  and sets  $x^{k+1} := x^k + d^k$ .

Specifically, we have the following algorithm:

## Newton's Method (Local Version)

S0. Choose  $x^0 \in \mathbb{R}^n$ ,  $\varepsilon > 0$ , and set k := 0.

S1. If  $||F(x^k)|| < \varepsilon$ , STOP.

S2. Compute  $d^k$  as a solution of Newton's equation

$$F'(x^k)d = -F(x^k).$$

S3. Set  $x^{k+1} := x^k + d^k$ , increment k and go to S1.

**Lemma 3.2**: Let  $F : \mathbb{R}^n \to \mathbb{R}^n$  be  $C^1$ , and  $\overline{x} \in \mathbb{R}^n$  such that  $F'(\overline{x})$  is invertible. Then, there exists  $\varepsilon > 0$  such that F'(x) remains invertible for all  $x \in B_{\varepsilon}(\overline{x})$ , and there exists C > 0 such that

$$||F'(x)^{-1}|| \le C, \quad \forall x \in B_{\varepsilon}(\overline{x}).$$

PROOF. Since F' continuous at  $\overline{x}$ , there exists  $\varepsilon > 0$  such that  $\|F'(\overline{x}) - F'(x)\| \le \frac{1}{2\|F'(\overline{x})^{-1}\|}$ for all  $x \in B_{\varepsilon}(\overline{x})$ . Then, for all  $x \in B_{\varepsilon}(\overline{x})$ ,

$$\begin{split} \left\|I-F'(x)F'(\overline{x})^{-1}\right\| &= \left\|\left(F'(\overline{x})-F'(x)\right)F'(\overline{x})^{-1}\right\| \\ &\leq \left\|F'(\overline{x})-F'(x)\right\|\left\|F'(\overline{x})^{-1}\right\| \leq \frac{1}{2} < 1. \end{split}$$

By a corollary of the Banach lemma, F'(x) invertible over  $B_{\varepsilon}(\overline{x})$ , and

$$\left\|F'(x)^{-1}\right\| \leq \frac{\left\|F'(\overline{x})^{-1}\right\|}{1 - \left\|I - F'(x)F'(\overline{x})^{-1}\right\|} \leq 2\left\|F'(\overline{x})^{-1}\right\| =: C.$$

**Remark 3.6**: Observe  $F: \mathbb{R}^n \to \mathbb{R}^m$  is differentiable at  $\overline{x}$  if and only if  $\|F(x^k) - F(\overline{x}) - F'(\overline{x})(x^k - \overline{x})\| = o(\|x^k - \overline{x}\|)$  for every  $x^k \to \overline{x}$ .

This can be sharpened if F' is continuous or even locally Lipschitz.

**Lemma 3.3**: Let  $F: \mathbb{R}^n \to \mathbb{R}$  be continuously differentiable and  $x^k \to \overline{x}$ , then:

1. 
$$\|F(x^k) - F(\overline{x}) - F'(x^k)(x^k - \overline{x})\| = o(\|x^k - \overline{x}\|);$$

2. if 
$$F'$$
 locally Lipschitz at  $\overline{x}$ , then  $\|F(x^k) - F(\overline{x}) - F'(x^k)(x^k - \overline{x})\| = O(\|x^k - \overline{x}\|^2)$ .

Proof.

1. Observe that

$$\begin{split} & \left\| F\left(x^{k}\right) - F(\overline{x}) - F'\left(x^{k}\right)\left(x^{k} - \overline{x}\right) \right\| \\ \leq & \left\| F\left(x^{k}\right) - F(\overline{x}) - F(\overline{x})\left(x^{k} - \overline{x}\right) \right\| + \left\| F'\left(x^{k}\right)\left(x^{k} - \overline{x}\right) - F'(\overline{x})\left(x^{k} - \overline{x}\right) \right\| \\ \leq & \left\| F\left(x^{k}\right) - F(\overline{x}) - F(\overline{x})\left(x^{k} - \overline{x}\right) \right\| + \left\| F'\left(x^{k}\right) - F(\overline{x}) \right\| \left\| x^{k} - \overline{x} \right\|. \end{split}$$

The left-hand term is  $o(\|x^k - \overline{x}\|)$  by our observations previously, and the right-hand term is as well by continuity of F', thus so is the sum.

2. Let L > 0 be a local Lipschitz constant of F' at  $\overline{x}$ . Then,

$$\begin{split} \|F(x^{k}) - F(\overline{x}) - F'(x^{k})(x^{k} - \overline{x})\| &= \left\| \int_{0}^{1} F'(\overline{x} + t(x^{k} - \overline{x})) \, dt(x^{k} - \overline{x}) - F'(x^{k})(x^{k} - \overline{x}) \right\| \\ &\leq \int_{0}^{1} \|F'(\overline{x} + t(x^{k} - \overline{x})) - F'(x^{k})\| \, dt \cdot \|x^{k} - \overline{x}\| \\ &\leq L \int_{0}^{1} |1 - t| \|x^{k} - \overline{x}\| \, dt \cdot \|x^{k} - \overline{x}\| \\ &= L \|x^{k} - \overline{x}\|^{2} \int_{0}^{1} (1 - t) \, dt = \frac{L}{2} \|x^{k} - \overline{x}\|^{2}, \end{split}$$

which implies the result.

3.3.2 Newton's Method for Nonlinear Equations

- **Theorem 3.3**: Let  $F : \mathbb{R}^n \to \mathbb{R}^n$  be continuously differentiable,  $\overline{x} \in \mathbb{R}^n$  such that  $F(\overline{x}) = 0$  and  $F'(\overline{x})$  is invertible. Then, there exists an  $\varepsilon > 0$  such that for every  $x^0 \in B_{\varepsilon}(\overline{x})$ , we have:
- 1. Algorithm 3.3 is well-defined and generates a sequence  $\{x^k\}$  which converges to  $\overline{x}$ ;
- 2. the rate of convergence is (at least) linear;
- 3. if F' is locally Lipschitz at  $\overline{x}$ , then the rate is quadratic.

Proof.

1. By the previous lemma, we know there is  $\varepsilon_1, c > 0$  such that  $\|F'(x)^{-1}\| \le c$  for all  $x \in B_{\varepsilon_1}(x)$ . Further, there exists an  $\varepsilon_2 > 0$  such that  $\|F(x) - F(\overline{x}) - F'(x)(x - \overline{x})\| \le \frac{1}{2c}\|x - \overline{x}\|$  for all  $x \in B_{\varepsilon_2}(\overline{x})$ . Take  $\varepsilon = \min\{\varepsilon_1, \varepsilon_2\}$  and pick  $x^0 \in B_{\varepsilon}(\overline{x})$ . Then,  $x^1$  is well-defined, since  $F'(x^0)$  is invertible, and so

$$||x^{1} - \overline{x}|| = ||x^{0} - F'(x^{0})^{-1}F(x^{0}) - \overline{x}||$$

$$= ||F'(x^{0})^{-1} \left( F(x^{0}) - \underbrace{F(\overline{x})}_{=0} - F'(x^{0})(x^{0} - \overline{x}) \right) ||$$

$$\leq ||F'(x^{0})^{-1}|| ||F(x^{0}) - F(\overline{x}) - F'(x^{0})(x^{0} - \overline{x})||$$

$$\leq c \cdot \frac{1}{2c} ||x^{0} - \overline{x}||$$

$$= \frac{1}{2} ||x^{0} - \overline{x}|| < \frac{\varepsilon}{2},$$

so in particular,  $x^1 \in B_{\varepsilon/2}(\overline{x}) \subset B_{\varepsilon}(\overline{x})$ . Inductively,

$$\left\|x^k - \overline{x}\right\| \le \left(\frac{1}{2}\right)^k \left\|x^0 - \overline{x}\right\|,$$

for every  $k \in \mathbb{N}$ . Thus,  $x^k$  well-defined and converges to  $\overline{x}$ .

2., 3. Analogous to 1.,

$$\begin{aligned} \|x^{k+1} - \overline{x}\| &= \|x^k - d^k - \overline{x}\| \\ &= \|x^k - F'(x^k)^{-1} F(x^k) - \overline{x}\| \\ &\leq \|F'(x^k)^{-1}\| \|F(x^k) - F(\overline{x}) - F'(x^k)(x^k - \overline{x})\| \\ &\leq c \|F(x^k) - F(\overline{x}) - F'(x^k)(x^k - \overline{x})\|. \end{aligned}$$

This final line is little o of  $||x^k - \overline{x}||$  or this quantity squared by the previous lemma, which proves the result depending on the assumptions of 2., 3..

#### 3.3.3 Newton's Method for Optimization Problem

Consider

$$\min_{x \in \mathbb{R}^n} f(x),$$

with  $f: \mathbb{R}^n \to \mathbb{R}$  twice continuously differentiable. Recall that if  $\overline{x}$  a local minimizer of f,  $\nabla f(\overline{x}) = 0$ . We'll now specialize Newton's to  $F := \nabla f$ :

Newton's Method for Optimization (Local Version)

S0. Choose  $x^0 \in \mathbb{R}^n$ ,  $\varepsilon > 0$ , and set k := 0.

S1. If  $\|\nabla f(x^k)\| < \varepsilon$ , STOP. S2. Compute  $d^k$  as a solution of Newton's equation  $\nabla^2 f(x^k) d = -\nabla f(x^k).$ 

$$\nabla^2 f(x^k) d = -\nabla f(x^k).$$

S3. Set  $\underline{x^{k+1}} := \underline{x^k} + d^k$ , increment k and go to S1.

We then have an analogous convergence result to the previous theorem by simply applying F := $\nabla f$ ; in particular, if f thrice continuously differentiable, we have quadratic convergence.

**Example 3.2**: Let  $f(x) := \sqrt{x^2 + 1}$ . Then  $f'(x) = \frac{x}{\sqrt{x^2 + 1}}$ ,  $f''(x) = \frac{1}{(x^2 + 1)^{3/2}}$ . Newton's equation (i.e. Algorithm 3.4, S2) reads in this case:

$$\frac{1}{\left(x_k^2+1\right)^{3/2}}d = -\frac{x_k}{\sqrt{x_k^2+1}}.$$

This gives solution  $d_k = -(x_k^2 + 1)x_k$ , so  $x_{k+1} = -x_k^3$ . Then, notice that if:

$$|x_0| < 1 \Rightarrow x_k \rightarrow 0$$
, quadratically

$$|x_0| > 1 \Rightarrow x_k \text{ diverges}$$

$$|x_0| = 1 \Rightarrow |x_k| = 1 \forall k,$$

so the convergence is truly local; if we start too far from 0, we'll never have convergence.

We can see from this example that this truly a local algorithm. A general globalization strategy is to:

- if Newton's equation has no solution, or doesn't provide sufficient decay, set  $d^k := -\nabla f(x^k)$ ;
- introduce a step-size.

## Newton's Method (Global Version)

S0. Choose  $x^0 \in \mathbb{R}^n$ ,  $\varepsilon > 0$ ,  $\rho > 0$ , p > 2,  $\beta \in (0,1)$ ,  $\sigma \in (0,1/2)$  and set k := 0

S1. If 
$$\|\nabla f(x^k)\| < \varepsilon$$
, STOP

S2. Determine  $d^k$  as a solution of

$$\nabla^2 f(x^k) d = -\nabla f(x^k).$$

If no solution exists, or if  $\nabla f(x^k)^T d^k \le -\rho \|d^k\|^p$ , is violated, set  $d^k := -\nabla f(x^k)$  S3. Determine  $t_k > 0$  by the Armijo back-tracking rule, i.e.

$$t_k \coloneqq \max_{\ell \in \mathbb{N}_0} \Bigl\{ \beta^\ell \, | \, f\bigl(x^k + \beta^\ell d^k\bigr) \leq f\bigl(x^k\bigr) + \beta^\ell \sigma \nabla f\bigl(x^k\bigr)^T d^k \Bigr\}$$

S4. Set  $x^{k+1} := x^k + t_k d^k$ , increment k to k+1, and go back to S1.

**Remark 3.7**: S3. well-defined since in either choice of  $d^k$  in S2., we will have a descent direction so the choice of  $t_k$  in S3. is valid; i.e.  $(x^k, d^k) \in A_f$  for every k.

**Theorem 3.4** (Global convergence of Algorithm 3.5): Let  $f : \mathbb{R}^n \to \mathbb{R}$  be twice continuously differentiable. Then every cluster point of  $\{x^k\}$  generated by Algorithm 3.5 is a stationary point of f.

**Remark 3.8**: Note that we didn't impose any invertibility condition on the Hessian of f; indeed, if say the hessian was nowhere invertible, then Algorithm 3.5 just becomes the gradient method with Armijo back-tracking, for which have already established this result.

PROOF. Let  $\{x^k\}$  be generated by Algorithm 3.5, with  $\{x^k\}_K \to \overline{x}$ . If  $d^k := -\nabla f(x^k)$  for infinitely many  $k \in K$  (i.e. along a subsubsequence of  $\{x^k\}$ ), then we have nothing to prove by the previous remark.

Otherwise, assume wlog (by passing to a subsubsequence again if necessary) that  $d^k$  is determined by the Newton equation for all  $k \in K$ . Suppose towards a contradiction that  $\nabla f(\overline{x}) \neq 0$ . By Newton's equation,

$$\|\nabla f(x^k)\| = \|\nabla^2 f(x^k) d^k\| \le \|\nabla^2 f(x^k)\| \|d^k\|, \quad \forall k \in K$$

By assumption  $\|\nabla^2 f(x^k)\| \neq 0$ ; if it were, then by assumption  $\nabla f(x^k) = 0$ , i.e. we'd have already reached our stationary point, which we assumed doesn't happen. So, we may write  $\frac{\|\nabla f(x^k)\|}{\|\nabla^2 f(x^k)\|} \leq \|d^k\|$  for all  $k \in K$ . We claim that there exists  $c_1, c_2 > 0$  such that

$$0 < c_1 \le \left\| d^k \right\| \le c_2, \qquad \forall k \in K.$$

We have existence of  $c_1$  since, if it didn't, we could find a subsequence of the  $d^k$ 's such that  $d^k \to 0$  along this subsequence; but by our bound above and the fact that  $\|\nabla^2 f(x^k)\|$  uniformly bounded (by continuity), then  $\|\nabla f(x^k)\|$  would converge to zero along the subsequence too, a contradiction.

The existence of  $c_2$  follows from the sufficient decrease condition. Indeed, suppose such a  $c_2$  didn't exist; by the condition

$$\nabla f(x^k)^T \frac{d^k}{\|d^k\|} \le -\rho \|d^k\|^{p-1};$$

the left-hand side is bounded (since  $\nabla f\left(x^k\right) \to \nabla f\left(\overline{x}\right)$  and  $\frac{d^k}{\|d^k\|}$  lives on the unit sphere). Since  $c_2$  assumed not to exist, there is a subsequence  $\|d^k\| \to \infty$ , but then  $-\rho \|d^k\|^{p-1} \to -\infty$ , contradicting the fact that the LHS is bounded. Hence, there also exists such a  $c_2$  as claimed.

As  $\{f(x^k)\}$  is monotonically decreasing (by construction in S3) and converges along a subsequence K to  $f(\overline{x})$ , then  $f(x^k)$  converges along the whole sequence to  $f(\overline{x})$ . In particular,  $f(x^{k+1}) - f(x^k) \to 0$ . Then,

$$\frac{f(x^{k+1}) - f(x^k)}{\sigma} \le t_k \nabla f(x^k)^T d^k \le -\rho t_k \|d^k\|^p \le 0.$$

Taking  $k \to \infty$  along K, we see that  $t_k \|d^k\|^p \to 0$  along K as well. We show now that  $\{t_k\}_K$  actually uniformly bounded away from zero. Suppose not. Then, along a sub(sub)sequence,  $t_k \to 0$ . By the Armijo rule,  $t_k = \beta^{\ell_k}$ , for  $\ell_k \in \mathbb{N}_0$ , uniquely determined. Since  $t_k \to 0$ , then  $\ell_k \to \infty$ . On the other hand, by S3,

$$\frac{f(x^k + \beta^{\ell_k - 1} d^k) - f(x^k)}{\beta^{\ell_k - 1}} > \sigma \nabla f(x^k)^T d^k.$$

Suppose  $d^k \to \overline{d} \neq 0$  (by again passing to a subsequence if necessary), which we may assume by boundedness. Taking  $k \to \infty$ , the LHS converges to  $\nabla f(\overline{x})^T \overline{d}$  and the RHS converges to  $\sigma \nabla f(\overline{x})^T \overline{d}$  so  $\nabla f(\overline{x})^T \overline{d} \geq \sigma \nabla f(\overline{x})^T \overline{d}$ , which implies since  $\sigma \in \left(0, \frac{1}{2}\right)$  that  $\nabla f(\overline{x})^T \overline{d} \geq 0$ . Taking  $k \to \infty$  in the sufficient decrease condition statement shows that this is a contradiction. Hence,  $t_k$  uniformly bounded away from 0. Hence, there exists a  $\overline{t} > 0$  such that  $t_k \geq \overline{t}$  for all  $k \in K$ . But we had that  $t^k \nabla f\left(x^k\right)^T d^k \to 0$ , so by boundedness of  $t_k$  it must be that  $\nabla f\left(x^k\right)^T d^k \to 0$  along the subsequence; by the sufficient decrease condition again, it must be that  $d^k \to 0$ , which it can't, as we showed it was uniformly bounded away, and thus we have a contradiction.

- **Theorem 3.5** (Fast local convergence of Algorithm 3.5): Let  $f: \mathbb{R}^n \to \mathbb{R}$  be twice continuously differentiable,  $\{x^k\}$  generated by Algorithm 3.5. If  $\overline{x}$  is a cluster point of  $\{x^k\}$  with  $\nabla^2 f(\overline{x}) > 0$ . Then:
- 1.  $\{x^k\} \to \overline{x}$  along the *whole* sequence, so  $\overline{x}$  is a strict local minimizer of f;
- 2. for  $k \in \mathbb{N}$  sufficiently large,  $d^k$  wil be determined by the Newton equation in S2;
- 3.  $\{x^k\} \to \overline{x}$  at least superlinearly;
- 4. if  $\nabla^2 f$  locally Lipschitz,  $\{x^k\} \to \overline{x}$  quadratically.

# §3.4 Quasi-Newton Methods

In Newton's, in general we need to find

$$d^k$$
 solving  $\nabla^2 f(x^k)d = -\nabla f(x^k)$ .

3.4 Quasi-Newton Methods

Advantages/disadvantages:

- (+) Global convergence with fast local convergence
- (-) Evaluating  $\nabla^2 f$  can be expensive/impossible.

Dealing with the second, there are two general approaches:

- *Direct Methods:* replace  $\nabla^2 f(x^k)$  with some matrix  $H_k$  approximating it;
- *Indirect Methods:* replace  $\nabla^2 f(x^k)^{-1}$  by  $B_k$  approximating it;

where  $H_k$ ,  $B_k$  reasonably computational, and other convergence results are preserved.

#### 3.4.1 Direct Methods

The typical conditions we put on  $H_{k+1}$  as described above are:

- 1.  $H_{k+1}$  symmetric
- 2.  $H_{k+1}$  satisfies the *Quasi-Newton equation* (QNE)

$$H_{k+1}s^k = y^k$$
,  $s^k \coloneqq x^{k+1} - x^k$ ,  $y^k \coloneqq \nabla f(x^{k+1}) - \nabla f(x^k)$ 

- 3.  $H_{k+1}$  can be achieved from  $H_k$  "efficiently"
- 4. The result method has strong local convergence properties

**Remark 3.9**: Suppose  $x^k$  a current iterate for an algorithm to minimize  $f : \mathbb{R}^n \to \mathbb{R}$  for  $f \in \mathbb{C}^2$ .

- 1.  $\nabla^2 f(x^k)$  does not generally satisfy QNE;
- 2. condition 1 above is motivated by the fact that Hessians are symmetric;
- 3. the QNE is motivated by the mean-value theorem for vector-valued functions,

$$\nabla f(x^{k+1}) - \nabla f(x^k) = \int_0^1 \nabla^2 f(x^k + t(x^{k+1} - x^k)) dt \cdot (x^{k+1} - x^k);$$

we can think of the integrated term as an averaging of the Hessian along the line between  $x^k$ ,  $x^{k+1}$ .

We follow a so-called *symmetric rank-2 approach*; given  $H_k$ , we update

$$H_{k+1} = H_k + \gamma u u^T + \delta v v^T, \qquad \gamma, \delta \in \mathbb{R}; u, v \in \mathbb{R}^n.$$

Note that if we put  $S := uu^T$  for  $u \neq 0$ , rank(S) = 1 and  $S^T = S$ .

3.4.1 Direct Methods