MATH378 - Nonlinear Optimization

Based on lectures from Fall 2025 by Prof. Tim Hoheisel. Notes by Louis Meunier

Contents

1 Preliminaries	
1.1 Terminology	
1.2 Convex Sets and Functions	3
2 Unconstrained Optimization	4
2.1 Theoretical Foundations	
2.1.1 Quadratic Approximation	6
2.2 Differentiable Convex Functions	
2.3 Matrix Norms	
3 Descent Methods	11
3.1 A General Line-Search Method	
3.1.1 Global Convergence of Algorithm 3.1	
3.2 The Gradient Method	
3.3 Newton-Type Methods	
3.3.1 Convergence Rates and Landau Notation	
3.3.2 Newton's Method for Nonlinear Equations	
3.3.3. Newton's Method for Ontimization Problem	18

§1 Preliminaries

§1.1 Terminology

We consider problems of the form

minimize
$$f(x)$$
 subject to $x \in X$, (†)

with $X \subset \mathbb{R}^n$ the feasible region with x a feasible point, and $f: X \to \mathbb{R}$ the objective (function); more concisely we simply write

$$\min_{x \in X} f(x)$$
.

When $X = \mathbb{R}^n$, we say the problem (†) is *unconstrained*, and conversely *constrained* when $X \subseteq \mathbb{R}^n$.

⊗ Example 1.1 (Polynomial Fit): Given $y_1, ..., y_m \in \mathbb{R}$ measurements taken at m distinct points $x_1, ..., x_m \in \mathbb{R}$, the goal is to find a degree $\leq n$ polynomial $q : \mathbb{R} \to \mathbb{R}$, of the form

$$q(x) = \sum_{k=0}^{n} \beta_k x^k,$$

"fitting" the data $\{(x_i, y_i)\}_i$, in the sense that $q(x_i) \approx y_i$ for each i. In the form of (†), we can write this precisely as

$$\min_{\beta \in \mathbb{R}^{n+1}} \frac{1}{2} \sum_{i=0}^{n} \left(\underbrace{\beta_n x_i^n + \dots + \beta_1 x_i + \beta_0}_{q(x_i)} - y_i \right)^2;$$

namely, we seek to minimize the ℓ^2 -distance between $(q(x_i))$ and (y_i) . If we write

$$X \coloneqq \begin{pmatrix} 1 & x_1 & \dots & x_1^n \\ \vdots & \dots & \dots & \vdots \\ 1 & x_m & \dots & x_m^n \end{pmatrix} \in \mathbb{R}^{m \times (n+1)}, \qquad y \coloneqq \begin{pmatrix} y_1 \\ \vdots \\ y_m \end{pmatrix} \in \mathbb{R}^m,$$

then concisely this problem is equivalent to

$$\min_{\beta \in \mathbb{R}^{n+1}} \frac{1}{2} ||X \cdot \beta - y||_2^2,$$

a so-called *least-squares* problem.

We have two related tasks:

- 1. Find the optimal value asked for by (†), that is what $\inf_X f$ is;
- 2. Find a specific point \overline{x} such that $f(\overline{x}) = \inf_X f$, i.e. the value of a point

$$\overline{x} \in \operatorname{argmin}_X f := \left\{ x \in X \mid f(x) = \inf_X f \right\}.$$

(noting that argmin should be viewed as a set-valued function, as there may be multiple admissible minimizers) Notice that if we can accomplish 2., we've accomplished 1. by computing $f(\overline{x})$.

1.1 Terminology

Note that $\overline{x} \in \operatorname{argmin}_X f \Rightarrow f(\overline{x}) = \inf_X f$, but $\inf_X f \in \mathbb{R}$ does not necessarily imply $\operatorname{argmin}_X f \neq \emptyset$, that is, there needn't be a feasible minimimum; for instance $\inf_{x \in \mathbb{R}} e^x = 0$, but $\operatorname{argmin}_{\mathbb{R}} f = \emptyset$ (there is no x for which $e^x = 0$).

- \hookrightarrow **Definition 1.1** (Minimizers): Let $X \subset \mathbb{R}^n$ and $f : \mathbb{R}^n \to \mathbb{R}$. Then $\overline{x} \in X$ is called a
- *global minimizer* (of f over X) if $f(\overline{x}) \le f(x) \forall x \in X$, or equivalently if $\overline{x} \in \operatorname{argmin}_X f$;
- *local minimizer (of f over X)* if $f(\overline{x}) \le f(x) \forall x \in X \cap B_{\varepsilon}(\overline{x})$ for some $\varepsilon > 0$.

In addition, we have *strict* versions of each by replacing " \leq " with "<".

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\hookrightarrow Definition 1.2 (Some Geometric Tools): Let f : \mathbb{R}^n \to \mathbb{R}.
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- gph $f := \{(x, f(x)) \mid x \in \mathbb{R}^n\} \subseteq \mathbb{R}^n \times \mathbb{R}$
- $f^{-1}(\{c\}) := \{x \mid f(x) = c\} \equiv contour/level \ set \ at \ c$
- $\operatorname{lev}_c f := f^{-1}((-\infty, c]) = \{x \mid f(x) \le c\} \equiv lower \ level/sublevel \ set \ at \ c$

Remark 1.1:

- $lev_{inf} f = argmin f$
- assume *f* continuous; then all (sub)level sets are closed (possibly empty)

We recall the following result from calculus/analysis:

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→Theorem 1.1 (Weierstrass): Let f : \mathbb{R}^n \to \mathbb{R} be continuous and X \subset \mathbb{R}^n compact. Then, \operatorname{argmin}_X f \neq \emptyset.
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From, we immediately have the following:

Proposition 1.1: Let $f : \mathbb{R}^n \to \mathbb{R}$ continuous. If there exists a $c \in \mathbb{R}$ such that lev_cf is nonempty and bounded, then $\operatorname{argmin}_{\mathbb{R}^n} f \neq \emptyset$.

PROOF. Since f continuous, $\operatorname{lev}_c f$ is closed (being the inverse image of a closed set), thus $\operatorname{lev}_c f$ is compact (and in particular nonempty). By Weierstrass, f takes a minimimum over $\operatorname{lev}_c f$, namely there is $\overline{x} \in \operatorname{lev}_c f$ with $f(\overline{x}) \leq f(x) \leq c$ for each $x \in \operatorname{lev}_c f$. Also, f(x) > c for each $x \notin \operatorname{lev}_c f$ (by virtue of being a level set), and thus $f(\overline{x}) \leq f(x)$ for each $x \in \mathbb{R}^n$. Thus, \overline{x} is a global minimizer and so the theorem follows.

§1.2 Convex Sets and Functions

Definition 1.3 (Convex Sets): $C \subset \mathbb{R}^n$ is *convex* if for any $x, y \in C$ and $\lambda \in (0, 1)$, $\lambda x + (1 - \lambda)y \in C$; that is, the entire line between x and y remains in C.

1.2 Convex Sets and Functions

 \hookrightarrow **Definition 1.4** (Convex Fucntions): Let $C \subset \mathbb{R}^n$ be convex. Then, $f: C \to \mathbb{R}$ is called

1. convex (on C) if

$$f(\lambda x + (1 - \lambda)y) \le \lambda f(x) + (1 - \lambda)f(y),$$

for every $x, y \in C$ and $\lambda \in (0, 1)$;

- 2. strictly convex (on C) if the inequality \leq is replaced with \leq ;
- 3. *strongly convex* (on *C*) if there exists a $\mu > 0$ such that

$$f(\lambda x + (1 - \lambda)y) + \mu \lambda (1 - \lambda) ||x - y||^2 \le \lambda f(x) + (1 - \lambda)f(y),$$

for every $x, y \in C$ and $\lambda \in (0,1)$; we call μ the modulus of strong convexity.

Remark 1.2: $3. \Rightarrow 2. \Rightarrow 1.$

Remark 1.3: A function is convex iff its epigraph is a convex set.

⊗ Example 1.2: exp : $\mathbb{R} \to \mathbb{R}$, log : $(0, \infty) \to \mathbb{R}$ are convex. A function $f : \mathbb{R}^n \to \mathbb{R}^m$ of the form f(x) = Ax - b for $A \in \mathbb{R}^{m \times n}$, $b \in \mathbb{R}^m$ is called *affine linear*. For m = 1, every affine linear function is convex. All norms on \mathbb{R}^n are convex.

\hookrightarrow Proposition 1.2:

- 1. (Positive combinations) Let f_i be convex on \mathbb{R}^n and $\lambda_i > 0$ scalars for i = 1, ..., m, then $\sum_{i=1}^m \lambda_i f_i$ is convex; as long as one is strictly (resp. strongly) convex, the sum is strictly (resp. strongly) convex as well.
- 2. (Composition with affine mappings) Let $f : \mathbb{R}^n \to \mathbb{R}$ be convex and $G : \mathbb{R}^m \to \mathbb{R}^n$ be affine. Then, $f \circ G$ is convex on \mathbb{R}^m .

§2 Unconstrained Optimization

§2.1 Theoretical Foundations

We focus on the problem

$$\min_{x\in\mathbb{R}^n} f(x),$$

where $f: \mathbb{R}^n \to \mathbb{R}$ is continuously differentiable.

Definition 2.1 (Directional derivative): Let $D \subset \mathbb{R}^n$ be open and $f: D \to \mathbb{R}$. We say f directionally differentiable at $\overline{x} \in D$ in the direction $d \in \mathbb{R}^n$ if

$$\lim_{t \to 0^+} \frac{f(\overline{x} + td) - f(\overline{x})}{t}$$

exists, in which case we denote the limit by $f'(\bar{x}; d)$.

2.1 Theoretical Foundations

Lemma 2.1: Let $D \subset \mathbb{R}^n$ be open and $f : D \to \mathbb{R}$ differentiable at $x \in D$. Then, f is directionally differentiable at x in every direction d, with

$$f'(x;d) = \nabla f(x)^T d = \langle \nabla f(x), d \rangle.$$

Example 2.1 (Directional derivatives of the Euclidean norm): Let $f : \mathbb{R}^n \to \mathbb{R}$ by f(x) = ||x|| the usual Euclidean norm. Then, we claim

$$f'(x;d) = \begin{cases} \frac{x^T d}{\|x\|} & x \neq 0 \\ \|d\| & x = 0 \end{cases}$$

For $x \neq 0$, this follows from the previous lemma and the calculation $\nabla f(x) = \frac{x}{\|x\|}$. For x = 0, we look at the limit

$$\lim_{t \to 0^+} \frac{f(0+td) - f(0)}{t} = \lim_{t \to 0^+} \frac{t||d|| - 0}{t} = ||d||,$$

using homogeneity of the norm.

Lemma 2.2 (Basic Optimality Condition): Let *X* ⊂ \mathbb{R}^n be open and $f: X \to \mathbb{R}$. If \overline{x} is a *local minimizer* of f over X and f is directionally differentiable at \overline{x} , then $f'(\overline{x};d) \ge 0$ for all $d \in \mathbb{R}^n$.

PROOF. Assume otherwise, that there is a direction $d \in \mathbb{R}^n$ for which the $f'(\overline{x};d) < 0$, i.e.

$$\lim_{t \to 0^+} \frac{f(\overline{x} + td) - f(\overline{x})}{t} < 0.$$

Then, for all sufficiently small t > 0, we must have

$$f(\overline{x} + td) < f(\overline{x}).$$

Moreover, since X open, then for t even smaller (if necessary), $\overline{x} + td$ remains in X, thus \overline{x} cannot be a local minimizer.

→Theorem 2.1 (Fermat's Rule): In addition to the assumptions of the previous lemma, assume further that f is differentiable at \overline{x} . Then, $\nabla f(\overline{x}) = 0$.

PROOF. From the previous, we know $0 \le f'(\overline{x}; d)$ for any d. Take $d = -\nabla f(\overline{x})$, then using the representation of a directional derivative for a differentiable function, and the fact that norms are nonnegative,

$$0 \le -\|\nabla f(\overline{x})\|^2 \le 0,$$

which can only hold if $\|\nabla f(\overline{x})\| = 0$ hence $\nabla f(\overline{x}) = 0$

We recall the following from Calculus:

2.1 Theoretical Foundations

5

Theorem 2.2 (Taylor's, Second Order): Let $f : D \subset \mathbb{R}^n \to \mathbb{R}$ be twice continuously differentiable, then for each $x, y \in D$, there is an η lying on the line between x and y such that

$$f(y) = f(x) + \nabla f(x)^{T} (y - x) + \frac{1}{2} (y - x)^{T} \nabla^{2} f(\eta) (y - x).$$

Theorem 2.3 (2nd-order Optimality Conitions): Let $X \subseteq \mathbb{R}^n$ open and $f: X \to \mathbb{R}$ twice continuously differentiable. Then, if x a local minimizer of f over X, then the Hessian matrix $\nabla^2 f(x)$ is positive semi-definite.

PROOF. Suppose not, then there exists a d such that $d^T \nabla^2 f(x) d < 0$. By Taylor's, for every t > 0, there is an η_t on the line between x and x + td such that

$$f(x+td) = f(x) + t \underbrace{\nabla f(x)^T}_{=0} d + \frac{1}{2} t^2 d^T \nabla^2 f(\eta_t) d$$
$$= f(x) + \frac{t^2}{d^T} \nabla^2 f(\eta_t) d.$$

As $t \to 0^+$, $\nabla^2 f(\eta_t) \to \nabla^2 f(x) < 0$. By continuity, for t sufficiently small, $\frac{t^2}{2} d^T \nabla^2 f(\eta_t) d < 0$ for t sufficiently small, whence we find

$$f(x+td) < f(x),$$

for sufficiently small t, a contradiction.

Lemma 2.3: Let $X \subset \mathbb{R}^n$ open, $f: X \to \mathbb{R}$ in C^2 . If $\overline{x} \in \mathbb{R}^n$ is such that $\nabla^2 f(\overline{x}) > 0$ (i.e. is positive definite), then there exists $\varepsilon, \mu > 0$ such that $B_\varepsilon(\overline{x}) \subset X$ and

$$d^T \nabla^2 f(x) d \geq \mu \|d\|^2, \qquad \forall d \in \mathbb{R}^n, x \in B_\varepsilon(\overline{x}).$$

Combining this and Taylor's Theorem, we can deduce the following (our first "sufficient" result of this section):

Theorem 2.4 (Sufficient Optimality Condition): Let $X \subset \mathbb{R}^n$ open and $f \in C^2(X)$. Let \overline{x} be a stationary point of f such that $\nabla^2 f(\overline{x}) > 0$. Then, \overline{x} is a *strict* local minimizer of f.

2.1.1 Quadratic Approximation

Let $f : \mathbb{R}^n \to \mathbb{R}$ be C^2 and $\overline{x} \in \mathbb{R}^n$. By Taylor's, we can approximate

$$f(y) \approx g(y) \coloneqq f(\overline{x}) + \nabla f(\overline{x})^T (y - \overline{x}) + \frac{1}{2} (y - \overline{x})^T \nabla^2 f(\overline{x}) (y - \overline{x}).$$

Example 2.2 (Quadratic Functions): For $Q \in \mathbb{R}^{n \times n}$ symmetric, $c \in \mathbb{R}^n$ and $\gamma \in \mathbb{R}$, let

$$f: \mathbb{R}^n \to \mathbb{R}, \quad f(x) = \frac{1}{2} x^T Q x + c^T x + \gamma,$$

a typical quadratic function. Then,

$$\nabla f(x) = \frac{1}{2} \big(Q + Q^T \big) x + c = Qx + c, \qquad \nabla^2 f(x) = Q.$$

We find that f has no minimizer if $c \notin \operatorname{rge}(Q)$ or Q is not positive semi-definite, combining our previous two results. In turn, if Q is positive definite (and thus invertible), there is a unique local minimizer $\overline{x} = -Q^{-1}c$ (and global minimizer, as we'll see).

§2.2 Differentiable Convex Functions

 \hookrightarrow Theorem 2.5: Let $C \subset \mathbb{R}^n$ be open and convex and $f: C \to \mathbb{R}$ differentiable on C. Then:

1. *f* is convex (on *C*) iff

$$f(x) \ge f(\overline{x}) + \nabla f(\overline{x})^T (x - \overline{x})$$
 *1

for every $x, \overline{x} \in C$;

- 2. *f* is *strictly* convex iff same inequality as 1. with strict inequality;
- 3. f is *strongly* convex with modulus $\sigma > 0$ iff

$$f(x) \geq f(\overline{x}) + \nabla f(\overline{x})^T (x - \overline{x}) + \frac{\sigma}{2} \|x - \overline{x}\|^2 \qquad \star_2$$

for every $x, \overline{x} \in C$.

PROOF. $(1., \Rightarrow)$ Let $x, \overline{x} \in C$ and $\lambda \in (0, 1)$. Then,

$$f(\lambda x + (1-\lambda)\overline{x}) - f(\overline{x}) \le \lambda \big(f(x) - f(\overline{x})\big),$$

which implies

$$\frac{f(\overline{x}+\lambda(x-\overline{x}))-f(\overline{x})}{\lambda}\leq f(x)-f(\overline{x}).$$

Letting $\lambda \to 0^+$, the LHS \to the directional derivative of f at \overline{x} in the direction $x - \overline{x}$, which is equal to, by differentiability of f, $\nabla f(\overline{x})^T(x - \overline{x})$, thus the result.

$$(1., \Leftarrow)$$
 Let $x_1, x_2 \in C$ and $\lambda \in (0, 1)$. Let $\overline{x} := \lambda x_1 + (1 - \lambda)x_2$. \star_1 implies

$$f(x_i) \ge f(\overline{x}) + \nabla f(\overline{x})^T (x_i - \overline{x}),$$

for each of i=1,2. Taking "a convex combination of these inequalities", i.e. multiplying them by λ , $1-\lambda$ resp. and adding, we find

$$\lambda f(x_1) + (1-\lambda)f(x_2) \geq f(\overline{x}) + \nabla f(\overline{x})^T \big(\lambda x_1 + (1-\lambda)x_2 - \overline{x}\big) = f\big(\lambda x_1 + (1-\lambda)x_2\big),$$

thus proving convexity.

 $(2., \Rightarrow)$ Let $x \neq \overline{x} \in C$ and $\lambda \in (0, 1)$. Then, by 1., as we've just proven,

$$\lambda \nabla f(\overline{x})^T (x - \overline{x}) \leq f(\overline{x} + \lambda (x - \overline{x})) - f(\overline{x}).$$

But $f(\overline{x} + \lambda(x - \overline{x})) < \lambda f(x) + (1 - \lambda)f(\overline{x})$ by strict convexity, so we have

$$\lambda \nabla f(\overline{x})^T (x - \overline{x}) < \lambda \big(f(x) - f(\overline{x}) \big),$$

and the result follows by dividing both sides by λ .

- $(2., \Leftarrow)$ Same as $(1., \Leftarrow)$ replacing " \leq " with "<".
- (3.) Apply 1. to $f \frac{\sigma}{2} \|\cdot\|^2$, which is still convex if f σ -strongly convex, as one can check.
- \hookrightarrow Corollary 2.1: Let $f: \mathbb{R}^n \to \mathbb{R}$ be convex and differentiable. Then,
- a) there exists an *affine function* $g : \mathbb{R}^n \to \mathbb{R}$ such that $g(x) \le f(x)$ everywhere;
- b) if f strongly convex, then it is coercive, i.e. $\lim_{\|x\|\to\infty} f(x) = \infty$.
- \hookrightarrow Corollary 2.2: Let $f : \mathbb{R}^n \to \mathbb{R}$ be convex and differentiable, then TFAE:
- 1. \bar{x} is a global minimizer of f;
- 2. \overline{x} is a local minimizer of f;
- 3. \overline{x} is a stationary point of f.

PROOF. 1. \Rightarrow 2. is trivial and 2. \Rightarrow 3. was already proven and 3. \Rightarrow 1. follows from the fact that differentiability gives

$$f(x) \ge f(\overline{x}) + \underline{\nabla(f)(\overline{x})^T(x-\overline{x})}$$

for any $x \in \mathbb{R}^n$.

Corollary 2.3: (2.2.4)

- **→Theorem 2.6** (Twice Differentiable Convex Functions): Let $Ω ⊂ \mathbb{R}^n$ open and convex and $f ∈ C^2(Ω)$. Then,
- 1. f is convex on Ω iff $\nabla^2 f \ge 0$;
- 2. f is strictly convex on $\Omega \leftarrow \nabla^2 f > 0$;
- 2. f is σ -strongly convex on $\Omega \Leftrightarrow \sigma \leq \lambda_{\min}(\nabla^2 f(x))$ for all $x \in \Omega$.
- **Corollary 2.4**: Let $A \in \mathbb{R}^{n \times n}$ be symmetric, $b \in \mathbb{R}^n$ and $f(x) := \frac{1}{2}x^TAx + b^Tx$. Then,
- 1. f convex $\Leftrightarrow A \ge 0$;
- 2. f strongly convex $\Leftrightarrow A > 0$.

Theorem 2.7 (Convex Optimization): Let $f : \mathbb{R}^n \to \mathbb{R}$ be convex and continuous, $X \subset \mathbb{R}^n$ convex (and nonempty), and consider the optimization problem

$$\min f(x)$$
 s.t. $x \in X$ (\star) .

Then, the following hold:

- 1. \overline{x} is a global minimizer of $(\star) \Leftrightarrow \overline{x}$ is a local minimizer of (\star)
- 2. $\operatorname{argmin}_X f$ is convex (possibly empty)
- 3. f is strictly convex \Rightarrow argmin_Xf has at *most* one element
- 4. f is strongly convex and differentiable, and X closed, \Rightarrow argmin_Xf has exactly one element

PROOF. $(1., \Rightarrow)$ Trivial. $(1., \Leftarrow)$ Let \overline{x} be a local minimizer of f over X, and suppose towards a contradiction that there exists some $\hat{x} \in X$ such that $f(\hat{x}) < f(\overline{x})$. By convexity of f, X, we know for $\lambda \in (0,1)$, $\lambda \overline{x} + (1-\lambda)\hat{x} \in X$ and

$$f(\lambda \overline{x} + (1 - \lambda)\hat{x}) \le \lambda f(\overline{x}) + (1 - \lambda)f(\hat{x}) < f(\overline{x}).$$

Letting $\lambda \to 1^-$, we see that $\lambda \overline{x} + (1 - \lambda)\hat{x} \to \overline{x}$; in particular, for any neighborhood of \overline{x} we can construct a point which strictly lower bounds $f(\overline{x})$, which contradicts the assumption that \overline{x} a local minimizer.

- (2.) and (3.) are left as an exercise.
- (4.) We know that f is strictly convex and level-bounded. By (3.) we know there is at most one minimizer, so we just need to show there exists one. Take $c \in \mathbb{R}$ such that $\text{lev}_c(f) \cap X \neq \emptyset$ (which certainly exists by taking, say, f(x) for some $x \in X$). Then, notice that (\star) and

$$\min_{x \in \text{lev}_c f \cap X} f(x) \qquad (\star \star)$$

have the same solutions i.e. the same set of global minimizers (noting that this remains a convex problem). Since f continuous and $\text{lev}_c f \cap X$ compact and nonempty, f attains a minimum on $\text{lev}_c f \cap X$, as we needed to show.

Remark 2.1: Note that level sets of convex functions are convex, this is left as an exercise.

§2.3 Matrix Norms

We denote by $\mathbb{R}^{m \times n}$ the space of real-valued $m \times n$ matrices (i.e. of linear operators from $\mathbb{R}^n \to \mathbb{R}^m$).

 \hookrightarrow Proposition 2.1 (Operator Norms): Let $\|\cdot\|_*$ be a norm on \mathbb{R}^m and \mathbb{R}^n , resp. Then, the map

$$\mathbb{R}^{m \times n} \ni A \mapsto \|A\|_* \coloneqq \sup_{\substack{x \in \mathbb{R}^n, \\ \|x\|_* \neq 0}} \frac{\|Ax\|_*}{\|x\|_*} \in \mathbb{R}$$

is a norm on $R^{m \times n}$. In addition,

$$||A||_* = \sup_{||x||_*=1} ||Ax||_* = \sup_{||x||_* \le 1} ||Ax||_*.$$

2.3 Matrix Norms 9

PROOF. We first note that all of these sup's are truely max's since they are maximizing continuous functions over compact sets.

Let $A \in \mathbb{R}^{m \times n}$. The first "In addition" equality follows from positive homogeneity, since $\frac{x}{\|x\|_*}$ a unit vector. For the second, note that " \leq " is trivial, since we are supping over a larger (super)set. For " \geq ", we have for any x with $\|x\|_* \leq 1$,

$$||Ax||_* = ||x||_* ||A\frac{x}{||x||_*}||_* \le ||A\frac{x}{||x||_*}||.$$

Supping both sides over all such *x* gives the result.

We now check that $\|\cdot\|_*$ actually a norm on $\mathbb{R}^{m\times n}$.

- $1. \ \|A\|_* = 0 \Leftrightarrow \sup_{\|x\|_* = 1} \|Ax\|_* = 0 \Leftrightarrow \|Ax\|_* = 0 \forall \|x\|_* = 1 \Leftrightarrow Ax = 0 \forall \|x\|_* = 1 \Leftrightarrow A = 0$
- 2. For $\lambda \in \mathbb{R}$, $A \in \mathbb{R}^{m \times n}$, $\|\lambda A\|_* = \sup \|\lambda Ax\|_* = |\lambda| \cdot \sup \|Ax\|_* = |\lambda| \|A\|_*$
- 3. For $A, B \in \mathbb{R}^{m \times n}$, $||A + B||_* \le ||A||_* + ||B||_*$ using properties of sups of sums

Proposition 2.2: Let $A = (a_{ij})_{i=1,...,m} \in \mathbb{R}^{m \times n}$, then: j=1,...,n

- 1. $||A||_1 = \max_{j=1}^n \sum_{i=1}^m |a_{ij}|$
- 2. $||A||_2 = \sqrt{\lambda_{\max}(A^T A)} = \sigma_{\max}(A)$
- 3. $||A||_{\infty} = \max_{i=1}^{m} \sum_{i=1}^{n} |a_{ij}|$

 \hookrightarrow Proposition 2.3: Let $\|\cdot\|_*$ be a norm on \mathbb{R}^n , \mathbb{R}^m , and \mathbb{R}^p . For $A \in \mathbb{R}^{m \times n}$ and $B \in \mathbb{R}^{n \times p}$,

- 1. $||Ax||_* \le ||A||_* \cdot ||x||_*$
- 2. $||AB||_{*} \leq ||A||_{*} \cdot ||B||_{*}$

Proposition 2.4 (Banach Lemma): Let $C \in \mathbb{R}^{n \times n}$ with ||C|| < 1, where $||\cdot||$ submultiplicative. Then, I + C is invertible, and

$$||(1+C)^{-1}|| \le \frac{1}{1-||C||}.$$

Proof. We have for any m,

$$\left\| \sum_{i=1}^{m} (-C)^{i} \right\| \leq \sum_{i=1}^{m} \|C\|^{i} \underset{m \to \infty}{\longrightarrow} \frac{1}{1 - \|C\|}.$$

Hence, $A_m := \sum_{i=1}^m (-C)^i$ a sequence of matrices with bounded norm uniformly in m, and thus has a converging subsequence, so wlog $A_m \to A \in \mathbb{R}^{n \times n}$ (by relabelling). Moreover, observe that

$$A_m \cdot (I+C) = \sum_{i=0}^m (-C)^i (I+C) = \sum_{i=0}^m \left[(-C)^i - (-C)^{i+1} \right] = (-C)^0 - (-C)^{m+1} = I - (-C)^{m+1}.$$

Now, $||C^{m+1}|| \le ||C||^{m+1} \to 0$, since ||C|| < 1, thus $C \to 0$. Hence, taking limits in the line above implies

2.3 Matrix Norms 10

$$A(I+C) = \lim_{m \to \infty} A_m(I+C) = I,$$

implying A the inverse of (I + C), proving the proposition.

Corollary 2.5: Let $A, B \in \mathbb{R}^{n \times n}$ with ||I - BA|| < 1 for $||\cdot||$ submultiplicative. Then, A and B are invertible, and $||B^{-1}|| \le \frac{||A||}{1 - ||I - BA||}$.

§3 DESCENT METHODS

§3.1 A General Line-Search Method

We deal with the unconstrained problem

$$\min_{x \in \mathbb{R}^n} f(x) \qquad (\star).$$

Definition 3.1 (Descent Direction): Let $f : \mathbb{R}^n \to \mathbb{R}$, $x \in \mathbb{R}^n$. $d \in \mathbb{R}^n$ is a *descent direction* of f at x if there exists a $\bar{t} > 0$ such that f(x + td) < f(x) for all $t \in (0, \bar{t})$.

Proposition 3.1: If $f : \mathbb{R}^n \to \mathbb{R}$ is directionally differentiable at $x \in \mathbb{R}^n$ in the direction d with f'(x;d) < 0, then d a descent direction of f at x; in particular if f differentiable at x, then true for d if $\nabla f(x)^T d < 0$.

Corollary 3.1: Let $f : \mathbb{R}^n \to \mathbb{R}$ differentiable, $B \in \mathbb{R}^{n \times n}$ positive definite, and $x \in \mathbb{R}^n$. Then $\nabla f(x) \neq 0 \Rightarrow -B\nabla f(x)$ is a descent direction of f at x.

PROOF.
$$\nabla f(x)^T (-B\nabla f(x)) = -\nabla f(x)^T B\nabla f(x) < 0.$$

A generic method/strategy for solving (\star):

- S1. (Initialization) Choose $x^0 \in \mathbb{R}^n$ and set k := 0
- S2. (Termination) If x^k satisfies a "termination criterion", STOP
- S3. (Search direction) Determine d^k such that $\nabla f(x^k)^T d^k < 0$
- S4. (Step-size) Determine $t_k > 0$ such that $f(x^k + t_k d^k) < f(x^k)$
- S5. (Update) Set $x^{k+1} := x^k + t_k d^k$, iterate k, and go back to step 2.

Remark 3.1: a) The generic choice for d^k in 3. is just $d^k := -B_k \nabla f(x^k)$ for some $B_k > 0$. We focus on:

- $B_k = I$ (gradient-descent)
- $B_k = \nabla^2 f(x^k)^{-1}$ (Newton's method) $B_k \approx \nabla^2 f(x^k)^{-1}$ (quasi Newton's method)
- b) Step 4. is called *line-search*, since $t_k > 0$ determined by looking at

$$0 < t \mapsto f(x^k + td^k),$$

i.e. along the (half)line t > 0.

- c) Executing Step 4. is a trade-off between
 - (i) decreasing f along $x^k + td^k$ as much as possible;
 - (ii) keeping computational efforts low.

For instance, the exact minimization rule $t_k = \operatorname{argmin}_{t>0} f\left(x_k + td^k\right)$ overemphasizes (i) over (ii).

 \hookrightarrow **Definition 3.2** (Step-size rule): Let $f \in C^1(\mathbb{R}^n)$ and

$$\mathcal{A}_f \coloneqq \big\{ (x,d) \mid \nabla f(x)^T d < 0 \big\}.$$

A (possible set-valued) map

$$T:(x,d)\in \mathcal{A}_f\mapsto T(x,d)\in \mathbb{R}_+$$

is called a *step-size rule* for *f* .

If T is well-defined for all C^1 -functions, we say T well-defined.

3.1.1 Global Convergence of Algorithm 3.1

 \hookrightarrow **Definition 3.3** (Efficient step-size): Let $f \in C^1(\mathbb{R}^n)$. The step-size rule T is called *efficient* for *f* if there exists $\theta > 0$ such that

$$f(x+td) \le f(x) - \theta \left(\frac{\nabla f(x)^T d}{\|d\|}\right)^2, \quad \forall t \in T(x,d), (x,d) \in A_f.$$

Theorem 3.1: Let $f \in C^1(\mathbb{R}^n)$. Let $\{x^k\}, \{d^k\}, \{t_k\}$ be generated by Algorithm 3.1. Assume the following:

- 1. $\exists c > 0$ such that $-\left(\nabla f(x^k)^T d^k\right) / \left(\|\nabla f(x^k)\| \cdot \|d^k\|\right) \ge c$ for all k (this is called the *angle* condition), and
- 2. there exists $\theta > 0$ such that $f(x^k + t_k d^k) \le f(x^k) \theta \cdot (\nabla f(x^k)^T d^k / ||d^k||)^2$ for all k (which is satisfied if $t_k \in T(x^k, d^k)$ for an efficient T).

Then, every cluster point of $\{x^k\}$ is a stationary point of f.

Proof. By condition 2., there is $\theta > 0$ such that

$$f(x^{k+1}) \le f(x^k) - \theta \left(\frac{\nabla f(x^k)^T d^k}{\|d^k\|}\right)^2$$

for all $k \in \mathbb{N}$. By 1., we know

$$\left(\frac{\nabla f(x^k)^T d^k}{\|d^k\|}\right)^2 \ge c^2 \|\nabla f(x^k)\|^2.$$

Put $\kappa := \theta c^2$, then these two inequalities imply

$$f(x^{k+1}) \le f(x^k) - \kappa \cdot \|\nabla f(x^k)\|^2$$
. (*)

Let \overline{x} be a cluster point of $\{x^k\}$. As $\{f(x^k)\}$ is monotonically decreasing (by construction in the algorithm), and has cluster point $f(\overline{x})$ by continuity, it follows that $f(x_k) \to f(\overline{x})$ along the whole sequence. In particular, $f(x^{k+1}) - f(x^k) \to 0$; thus, from (*),

$$0 \le \kappa \left\| \nabla f(x^k) \right\|^2 \le f(x^k) - f(x^{k+1}) \to 0,$$

and thus $\nabla f(x^k) \to \nabla f(\overline{x}) = 0$, so indeed \overline{x} a stationary point of f.

§3.2 The Gradient Method

We specialize Algorithm 3.1 here. Specifically, we'll take

$$d^k := -\nabla f(x^k);$$

it's know that

$$\frac{-\nabla f(x^k)}{\|\nabla f(x^k)\|} = \operatorname{argmin}_{d:\|d\| \le 1} \nabla f(x^k)^T d,$$

with $\|\cdot\|$ the 2 norm.

We use a step-size rule called "Armijo rule". Choose parameters β , $\sigma \in (0,1)$. For $(x,d) \in A_f$, we define our step-size rule by

$$T_A(x,d) \coloneqq \max_{\ell \in \mathbb{N}_0} \left\{ \beta^\ell \mid \underbrace{f(x+\beta^\ell d) \leq f(x) + \beta^\ell \sigma \nabla f(x)^T d}_{\text{"Armijo condition"}} \right\}.$$

For instance, consider $f(x) = (x-1)^2 - 1$. The minimum of this function is $f^* = -1$. Choose $x^k := \frac{1}{k}$, then

$$f(x^k) = \frac{2k+1}{k^2} \to 0 \neq f^*,$$

even though $f(x^{k+1}) - f(x^k) < 0$; we don't actually reach the right stationary point with our chosen step size.

3.2 The Gradient Method 13

Example 3.1 (Illustration of Armijo Rule): For (x,d) ∈ A_f and f smooth on \mathbb{R}^n , defined ϕ : $\mathbb{R} \to \mathbb{R}$, $\phi(t) := f(x+td)$. The map $t \mapsto \sigma \phi'(0)t + \phi(0) = \sigma t \nabla f(x)^T d + \phi(0)$

Proposition 3.2: Let f : \mathbb{R}^n → \mathbb{R} be differentiable with β , $\sigma \in (0,1)$. Then for $(x,d) \in A_f$, there exists $\ell \in \mathbb{N}_0$ such that

$$f(x + \beta^{\ell} d) \le f(x) + \beta^{\ell} \sigma \nabla f(x)^{T} d,$$

i.e. $T_A(x,d) \neq \emptyset$.

Proof. Suppose not, i.e.

$$\frac{f(x + \beta^{\ell} d) - f(x)}{\beta^{\ell}} > \sigma \nabla f(x)^{T} d, \forall \ell \in \mathbb{N}_{0}.$$

Letting $\ell \to \infty$, the left-hand side converges to $\nabla f(x)^T d$, so

$$\nabla f(x)^T d \ge \sigma \nabla f(x)^T d.$$

But $(x, d) \in A_f$, so $\nabla f(x)^T d < 0$ so dividing both sides of this inequality by this quantity, this implies $\sigma \le 0$, which is a contradiction.

We now prove convergence of an algorithm based on the Armijo Rule:

Gradient Descent with Armijo Rule

S0. Choose $x^0 \in \mathbb{R}^n$, σ , $\beta \in (0,1)$, $\varepsilon \ge 0$, and set k := 0

S1. If $\|\nabla f(x^k)\| \le \varepsilon$, STOP

S2. Set $d^k := -\nabla f(x^k)$

S3. Determine $t_k > 0$ by

$$t_k = T_A(x, d)$$

as defined above.

S4. Set $x^{k+1} = x^k + t_k d^k$, iterate k and go to S1.

Lemma 3.1: Let $f \in C^1(\mathbb{R}^n)$, $x^k \to x$, $d^k \to d$ and $t_k \downarrow 0$. Then

$$\lim_{k \to \infty} \frac{f\left(x^k + t_k d^k\right) - f\left(x^k\right)}{t^k} = \nabla f(x)^T d.$$

Proof. Left as an exercise.

→Theorem 3.2: Let $f \in C^1(\mathbb{R}^n)$. Then every cluster point of a sequence $\{x^k\}$ generated by Algorithm 3.2 is a stationary point of f.

PROOF. Let \overline{x} be a cluster point of $\left\{x^k\right\}$ and let $x^k \underset{k \in K}{\to} \overline{x}$, K an infinite subset of \mathbb{N} . Assume towards a contradiction $\nabla f(\overline{x}) \neq 0$. As $f\left(x^k\right)$ is monotonically decreasing with cluster point $f(\overline{x})$, it must be that $f\left(x^k\right) \to f(\overline{x})$ along the whole sequence so $f\left(x^{k+1}\right) - f\left(x^k\right) \to 0$. Thus,

3.2 The Gradient Method 14

$$0 \le t_k \|\nabla f(x^k)\|^2 \stackrel{\text{S2}}{=} -t_k \nabla f(x^k)^T d^k \stackrel{\text{S3}}{\leq} \frac{f(x^k) - f(x^{k+1})}{\sigma} \to 0.$$

Thus, $0 = \lim_{k \in K} t_k \|\nabla f(x^k)\| = \|\nabla f(\overline{x})\| \lim_{k \in K} t_k$. We assumed \overline{x} not a stationary point, so it follows that $t_k \underset{k \in K}{\longrightarrow} 0$. By S3, for $\beta^{\ell_k} = t_k$,

$$\frac{f\left(x^k+\beta^{\ell_k-1}d^k\right)-f\left(x^k\right)}{\beta^{\ell_k-1}}>\sigma\nabla f\left(x^k\right)^Td^k.$$

Letting $k \to \infty$ along *K*,the LHS converges to, by the previous lemma, to

$$\nabla f(\overline{x})^T d = -\nabla f(\overline{x})^T \nabla f(\overline{x}) = -\|\nabla f(\overline{x})\|^2,$$

and the RHS converges to $\sigma \|\nabla f(\overline{x})\|^2$, which implies

$$-\|\nabla f(\overline{x})\|^2 \ge \sigma \|\nabla f(\overline{x})\|^2,$$

which implies σ negative, a contradiction.

Remark 3.2: The proof above shows, the following: Let $\{x^k\}$ such that $x^{k+1} := x^k + t_k d^k$ for $d^k \in \mathbb{R}^n$, $t_k > 0$, and let $f(x^{k+1}) \le f(x^k)$ and $x^k \xrightarrow{K} \overline{x}$ such that $d^k = -\nabla f(x^k)$, $t_k = T_A(x^k, d^k)$ for all $k \in K$. Then $\nabla f(\overline{x}) = 0$; i.e., all of the "focus" is on the subsequence along K. The only time we needed the whole sequence was to use the fact that $f(x^k) \to f(\overline{x})$ along the whole sequence.

§3.3 Newton-Type Methods

3.3.1 Convergence Rates and Landau Notation

 \hookrightarrow **Definition 3.4**: Let $\{x^k \in \mathbb{R}^n\}$ converge to \overline{x} . Then, $\{x^k\}$ converges:

1. *linearly* to \overline{x} if there exists $c \in (0,1)$ such that

$$||x^{k+1} - \overline{x}|| \le c||x^k - \overline{x}||, \forall k;$$

2. *superlinearly* to \overline{x} if

$$\lim_{k \to \infty} \frac{\left\| x^{k+1} - \overline{x} \right\|}{\left\| x^k - \overline{x} \right\|} = 0;$$

3. *quadratically* to \overline{x} if there exists C > 0 such that

$$||x^{k+1} - \overline{x}|| \le C||x^k - \overline{x}||^2, \forall k.$$

Remark 3.3: $3. \Rightarrow 2. \Rightarrow 1.$

Remark 3.4: We needn't assume $x^k \to \overline{x}$ for the first two definitions; their statements alone imply convergence. However, the last does not; there exists sequences with this property that do not converge.

 \hookrightarrow **Definition 3.5** (Landau Notation): Let {*a_k*}, {*b_k*} be positive sequences ↓ 0. Then,

1.
$$a_k = o(b_k) \Leftrightarrow \lim_{k \to \infty} \frac{a_k}{b_k} = 0$$
;

2. $a_k = O(b_k) \Leftrightarrow \exists C > 0 : a_k \leq Cb_k$ for all k (sufficiently large).

Remark 3.5: If $x^k \to \overline{x}$, then

- 1. the convergence is superlinear $\Leftrightarrow ||x^{k+1} \overline{x}|| = o(||x^k \overline{x}||);$ 2. the convergence is quadratic $\Leftrightarrow ||x^{k+1} \overline{x}|| = O(||x^k \overline{x}||^2).$

3.3.2 Newton's Method for Nonlinear Equations

We consider the nonlinear equation

$$F(x) = 0, \qquad (*)$$

where $F: \mathbb{R}^n \to \mathbb{R}^n$ is smooth (continuously differentiable). Our goal is to find a numerical scheme that can determine approximate zeros of *F*, i.e. solutions to (*). The idea of Newton's method for such a problem, is, given $x^k \in \mathbb{R}^n$, to consider the (affine) linear approximation of *F* about x^k ,

$$F_k: x \mapsto F(x^k) + F'(x^k)(x - x^k),$$

where F' the Jacobian of F. Then, we compute x^{k+1} as a solution of $F_k(x) = 0$. Namely, if $F'(x^k)$ invertible, then solving for $F_k(x^{k+1}) = 0$, we find

$$x^{k+1} = x^k - F'(x^k)^{-1}F(x^k).$$

More generally, one solves $F'(x^k)d = -F(x^k)$ and sets $x^{k+1} := x^k + d^k$.

Specifically, we have the following algorithm:

Newton's Method (Local Version)

S0. Choose $x^0 \in \mathbb{R}^n$, $\varepsilon > 0$, and set k := 0.

S1. If $||F(x^k)|| < \varepsilon$, STOP.

S2. Compute d^k as a solution of Newton's equation

$$F'(x^k)d = -F(x^k).$$

S3. Set $x^{k+1} := x^k + d^k$, increment k and go to S1.

Lemma 3.2: Let $F : \mathbb{R}^n \to \mathbb{R}^n$ be C^1 , and $\overline{x} \in \mathbb{R}^n$ such that $F'(\overline{x})$ is invertible. Then, there exists $\varepsilon > 0$ such that F'(x) remains invertible for all $x \in B_{\varepsilon}(\overline{x})$, and there exists C > 0 such that

$$||F'(x)^{-1}|| \le C, \quad \forall x \in B_{\varepsilon}(\overline{x}).$$

PROOF. Since F' continuous at \overline{x} , there exists $\varepsilon > 0$ such that $\|F'(\overline{x}) - F'(x)\| \le \frac{1}{2\|F'(\overline{x})^{-1}\|}$ for all $x \in B_{\varepsilon}(\overline{x})$. Then, for all $x \in B_{\varepsilon}(\overline{x})$,

$$\begin{split} \left\|I-F'(x)F'(\overline{x})^{-1}\right\| &= \left\|\left(F'(\overline{x})-F'(x)\right)F'(\overline{x})^{-1}\right\| \\ &\leq \left\|F'(\overline{x})-F'(x)\right\|\left\|F'(\overline{x})^{-1}\right\| \leq \frac{1}{2} < 1. \end{split}$$

By a corollary of the Banach lemma, F'(x) invertible over $B_{\varepsilon}(\overline{x})$, and

$$\left\|F'(x)^{-1}\right\| \leq \frac{\left\|F'(\overline{x})^{-1}\right\|}{1 - \left\|I - F'(x)F'(\overline{x})^{-1}\right\|} \leq 2\left\|F'(\overline{x})^{-1}\right\| =: C.$$

Remark 3.6: Observe $F: \mathbb{R}^n \to \mathbb{R}^m$ is differentiable at \overline{x} if and only if $\|F(x^k) - F(\overline{x}) - F'(\overline{x})(x^k - \overline{x})\| = o(\|x^k - \overline{x}\|)$ for every $x^k \to \overline{x}$.

This can be sharpened if F' is continuous or even locally Lipschitz.

Lemma 3.3: Let $F: \mathbb{R}^n \to \mathbb{R}$ be continuously differentiable and $x^k \to \overline{x}$, then:

1.
$$\|F(x^k) - F(\overline{x}) - F'(x^k)(x^k - \overline{x})\| = o(\|x^k - \overline{x}\|);$$

2. if
$$F'$$
 locally Lipschitz at \overline{x} , then $\|F(x^k) - F(\overline{x}) - F'(x^k)(x^k - \overline{x})\| = O(\|x^k - \overline{x}\|^2)$.

Proof.

1. Observe that

$$\begin{split} & \left\| F\left(x^{k}\right) - F(\overline{x}) - F'\left(x^{k}\right)\left(x^{k} - \overline{x}\right) \right\| \\ \leq & \left\| F\left(x^{k}\right) - F(\overline{x}) - F(\overline{x})\left(x^{k} - \overline{x}\right) \right\| + \left\| F'\left(x^{k}\right)\left(x^{k} - \overline{x}\right) - F'(\overline{x})\left(x^{k} - \overline{x}\right) \right\| \\ \leq & \left\| F\left(x^{k}\right) - F(\overline{x}) - F(\overline{x})\left(x^{k} - \overline{x}\right) \right\| + \left\| F'\left(x^{k}\right) - F(\overline{x}) \right\| \left\| x^{k} - \overline{x} \right\|. \end{split}$$

The left-hand term is $o(\|x^k - \overline{x}\|)$ by our observations previously, and the right-hand term is as well by continuity of F', thus so is the sum.

2. Let L > 0 be a local Lipschitz constant of F' at \overline{x} . Then,

$$\begin{split} \|F(x^{k}) - F(\overline{x}) - F'(x^{k})(x^{k} - \overline{x})\| &= \left\| \int_{0}^{1} F'(\overline{x} + t(x^{k} - \overline{x})) \, dt(x^{k} - \overline{x}) - F'(x^{k})(x^{k} - \overline{x}) \right\| \\ &\leq \int_{0}^{1} \|F'(\overline{x} + t(x^{k} - \overline{x})) - F'(x^{k})\| \, dt \cdot \|x^{k} - \overline{x}\| \\ &\leq L \int_{0}^{1} |1 - t| \|x^{k} - \overline{x}\| \, dt \cdot \|x^{k} - \overline{x}\| \\ &= L \|x^{k} - \overline{x}\|^{2} \int_{0}^{1} (1 - t) \, dt = \frac{L}{2} \|x^{k} - \overline{x}\|^{2}, \end{split}$$

which implies the result.

3.3.2 Newton's Method for Nonlinear Equations

- **Theorem 3.3**: Let $F : \mathbb{R}^n \to \mathbb{R}^n$ be continuously differentiable, $\overline{x} \in \mathbb{R}^n$ such that $F(\overline{x}) = 0$ and $F'(\overline{x})$ is invertible. Then, there exists an $\varepsilon > 0$ such that for every $x^0 \in B_{\varepsilon}(\overline{x})$, we have:
- 1. Algorithm 3.3 is well-defined and generates a sequence $\{x^k\}$ which converges to \overline{x} ;
- 2. the rate of convergence is (at least) linear;
- 3. if F' is locally Lipschitz at \overline{x} , then the rate is quadratic.

Proof.

1. By the previous lemma, we know there is $\varepsilon_1, c > 0$ such that $\|F'(x)^{-1}\| \le c$ for all $x \in B_{\varepsilon_1}(x)$. Further, there exists an $\varepsilon_2 > 0$ such that $\|F(x) - F(\overline{x}) - F'(x)(x - \overline{x})\| \le \frac{1}{2c}\|x - \overline{x}\|$ for all $x \in B_{\varepsilon_2}(\overline{x})$. Take $\varepsilon = \min\{\varepsilon_1, \varepsilon_2\}$ and pick $x^0 \in B_{\varepsilon}(\overline{x})$. Then, x^1 is well-defined, since $F'(x^0)$ is invertible, and so

$$||x^{1} - \overline{x}|| = ||x^{0} - F'(x^{0})^{-1}F(x^{0}) - \overline{x}||$$

$$= ||F'(x^{0})^{-1} \left(F(x^{0}) - \underbrace{F(\overline{x})}_{=0} - F'(x^{0})(x^{0} - \overline{x}) \right) ||$$

$$\leq ||F'(x^{0})^{-1}|| ||F(x^{0}) - F(\overline{x}) - F'(x^{0})(x^{0} - \overline{x})||$$

$$\leq c \cdot \frac{1}{2c} ||x^{0} - \overline{x}||$$

$$= \frac{1}{2} ||x^{0} - \overline{x}|| < \frac{\varepsilon}{2},$$

so in particular, $x^1 \in B_{\varepsilon/2}(\overline{x}) \subset B_{\varepsilon}(\overline{x})$. Inductively,

$$\left\|x^k - \overline{x}\right\| \le \left(\frac{1}{2}\right)^k \left\|x^0 - \overline{x}\right\|,$$

for every $k \in \mathbb{N}$. Thus, x^k well-defined and converges to \overline{x} .

2., 3. Analogous to 1.,

$$\begin{aligned} \|x^{k+1} - \overline{x}\| &= \|x^k - d^k - \overline{x}\| \\ &= \|x^k - F'(x^k)^{-1} F(x^k) - \overline{x}\| \\ &\leq \|F'(x^k)^{-1}\| \|F(x^k) - F(\overline{x}) - F'(x^k)(x^k - \overline{x})\| \\ &\leq c \|F(x^k) - F(\overline{x}) - F'(x^k)(x^k - \overline{x})\|. \end{aligned}$$

This final line is little o of $||x^k - \overline{x}||$ or this quantity squared by the previous lemma, which proves the result depending on the assumptions of 2., 3..

3.3.3 Newton's Method for Optimization Problem

Consider

$$\min_{x \in \mathbb{R}^n} f(x),$$

with $f: \mathbb{R}^n \to \mathbb{R}$ twice continuously differentiable. Recall that if \overline{x} a local minimizer of f, $\nabla f(\overline{x}) = 0$. We'll now specialize Newton's to $F := \nabla f$:

Newton's Method for Optimization (Local Version)

S0. Choose $x^0 \in \mathbb{R}^n$, $\varepsilon > 0$, and set k := 0.

S1. If $\|\nabla f(x^k)\| < \varepsilon$, STOP. S2. Compute d^k as a solution of Newton's equation $\nabla^2 f(x^k) d = -\nabla f(x^k).$

$$\nabla^2 f(x^k) d = -\nabla f(x^k).$$

S3. Set $\underline{x^{k+1}} := \underline{x^k} + d^k$, increment k and go to S1.

We then have an analogous convergence result to the previous theorem by simply applying F := ∇f ; in particular, if f thrice continuously differentiable, we have quadratic convergence.

Example 3.2: Let $f(x) := \sqrt{x^2 + 1}$. Then $f'(x) = \frac{x}{\sqrt{x^2 + 1}}$, $f''(x) = \frac{1}{(x^2 + 1)^{3/2}}$. Newton's equation (i.e. Algorithm 3.4, S2) reads in this case:

$$\frac{1}{\left(x_k^2+1\right)^{3/2}}d = -\frac{x_k}{\sqrt{x_k^2+1}}.$$

This gives solution $d_k = -(x_k^2 + 1)x_k$, so $x_{k+1} = -x_k^3$. Then, notice that if:

$$|x_0| < 1 \Rightarrow x_k \rightarrow 0$$
, quadratically

$$|x_0| > 1 \Rightarrow x_k \text{ diverges}$$

$$|x_0| = 1 \Rightarrow |x_k| = 1 \forall k,$$

so the convergence is truly local; if we start too far from 0, we'll never have convergence.

We can see from this example that this truly a local algorithm. A general globalization strategy is to:

- if Newton's equation has no solution, or doesn't provide sufficient decay, set $d^k := -\nabla f(x^k)$;
- introduce a step-size.

Newton's Method (Global Version)

- S0. Choose $x^0 \in \mathbb{R}^n$, $\varepsilon > 0$, $\rho > 0$, p > 2, $\beta \in (0,1)$, $\sigma \in (0,1/2)$ and set $k \coloneqq 0$
- S1. If $\|\nabla f(x^k)\| < \varepsilon$, STOP
- S2. Determine d^k as a solution of

$$\nabla^2 f(x^k) d = -\nabla f(x^k).$$

If no solution exists, or if $\nabla f(x^k)^T d^k \le -\rho \|d^k\|^p$, is violated, set $d^k := -\nabla f(x^k)$

S3. Determine $t_k > 0$ by the Armijo back-tracking rule, i.e.

$$t_k \coloneqq \max_{\ell \in \mathbb{N}_0} \Bigl\{ \beta^\ell \, | \, f\bigl(x^k + \beta^\ell d^k\bigr) \le f\bigl(x^k\bigr) + \beta^\ell \sigma \nabla f\bigl(x^k\bigr)^T d^k \Bigr\}$$

S4. Set $x^{k+1} := x^k + t_k d^k$, increment k to k+1, and go back to S1.

Remark 3.7: S3. well-defined since in either choice of d^k in S2., we will have a descent direction so the choice of t_k in S3. is valid; i.e. $(x^k, d^k) \in A_f$ for every k.

Theorem 3.4 (Global convergence of Algorithm 3.5): Let $f : \mathbb{R}^n \to \mathbb{R}$ be twice continuously differentiable. Then every cluster point of $\{x^k\}$ generated by Algorithm 3.5 is a stationary point of f.

Remark 3.8: Note that we didn't impose any invertibility condition on the Hessian of f; indeed, if say the hessian was nowhere invertible, then Algorithm 3.5 just becomes the gradient method with Armijo back-tracking, for which have already established this result.

- **Theorem 3.5** (Fast local convergence of Algorithm 3.5): Let $f: \mathbb{R}^n \to \mathbb{R}$ be twice continuously differentiable, $\{x^k\}$ generated by Algorithm 3.5. If \overline{x} is a cluster point of $\{x^k\}$ with $\nabla^2 f(\overline{x}) > 0$. Then:
- 1. $\{x^k\} \to \overline{x}$ along the *whole* sequence, so \overline{x} is a strict local minimizer of f;
- 2. for $k \in \mathbb{N}$ sufficiently large, d^k wil be determined by the Newton equation in S2;
- 3. $\{x^k\} \to \overline{x}$ at least superlinearly;
- 4. if $\nabla^2 f$ locally Lipschitz, $\{x^k\} \to \overline{x}$ quadratically.