

# MATH578 - Numerical Analysis 1

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## §1 POLYNOMIAL INTERPOLATION

In general, the goal of interpolation is, given a function  $f(x)$  on  $[a, b]$  and a series of distinct ordered points (often called *nodes* or *collocation points*)  $\{x_j\}_{j=1}^n \subseteq [a, b]$ , to find a polynomial  $P(x)$  such that  $f(x_j) = P(x_j)$  for each  $j$ .

↪ **Theorem 1.1** (Existence and Uniqueness of Lagrange Polynomial): Let  $f \in C[a, b]$  and  $\{x_j\}$  a set of  $n$  distinct points. Then, there exists a unique  $P(x) \in \mathbb{P}_{n-1}$ , the space of  $n - 1$ -degree polynomials, such that  $P(x_j) = f(x_j)$  for each  $j$ .

We call such a  $P$  the *Lagrange polynomial* associated to the points  $\{x_j\}$  for  $f$ .

PROOF. We define the following  $n - 1$  degree “fundamental polynomials” associated to  $\{x_j\}$ ,

$$\ell_j(x) \equiv \prod_{\substack{1 \leq i \leq n \\ i \neq j}} \frac{x - x_i}{x_j - x_i}, \quad j = 1, \dots, n.$$

Then, one readily verifies  $\ell_j(x_i) = \delta_{ij}$ , and that the distinctness of the nodes guarantees the denominator in each term of the product is nonzero. Define

$$P(x) = \sum_{j=1}^n f(x_j) \ell_j(x),$$

which, being a linear combination of  $n - 1$  degree polynomials is also in  $\mathbb{P}_{n-1}$ . Moreover,

$$P(x_i) = \sum f(x_j) \delta_{i,j} = f(x_i),$$

as desired.

For uniqueness, suppose  $\bar{P}$  another  $n - 1$  degree polynomial satisfying the conditions of the theorem. Then,  $q(x) \equiv P(x) - \bar{P}(x)$  is also a degree  $n - 1$  polynomial with  $q(x_i) = 0$  for each  $i = 1, \dots, n$ . Hence,  $q$  a polynomial with more distinct roots than its degree, and thus it must be identically zero, hence  $P = \bar{P}$ , proving uniqueness. ■

↪ **Theorem 1.2** (Interpolation Error): Suppose  $f \in C^n[a, b]$ , and let  $P(x)$  be the Lagrange polynomial for a set of  $n$  points  $\{x_j\}$ , with  $x_1 = a, x_n = b$ . Then, for each  $x \in [a, b]$ , there is a  $\xi \in [a, b]$  such that

$$f(x) - P(x) = \frac{f^{(n)}(\xi)}{n!} (x - x_1) \cdots (x - x_n).$$

Moreover, if we put  $h := \max_i (x_{i+1} - x_i)$ , then

$$\|f - P\|_\infty \leq \frac{h^n}{4n} \|f^{(n)}\|_\infty.$$

PROOF. We prove the first identity, and leave the second “Moreover” as a homework problem. Notice that it holds trivially for  $x = x_j$  for any  $j$ , so assume  $x \neq x_j$  for any  $j$ , and define the function

$$g(t) := f(t) - P(t) - \omega(t) \frac{f(x) - P(x)}{\omega(x)}, \quad \omega(t) := (t - x_1) \cdots (t - x_n) \in \mathbb{P}_n[t].$$

Then, we observe the following:

- $g \in C^n[a, b]$
- $g(x) = 0$
- $g(t = x_j) = 0$  for each  $j$

Recall that by Rolle’s Theorem, if a  $C^1$  function has  $\geq m$  roots, then its derivative has  $\geq m - 1$  roots. Thus, applying this principle inductively to  $g(t)$ , we conclude that  $g^{(n)}(t)$  has at least one root. Take  $\xi$  to be such a root. Then, one readily verifies that  $P^{(n)} \equiv 0$  and  $\omega^{(n)} \equiv n!$  (using polynomial properties), from which we may use the fact that  $g^{(n)}(\xi) = 0$  to simplify to the required identity. ■

**Remark 1.1:** In general, larger  $n$  leads to smaller maximum step size  $h$ . However, it is *not* true that  $n \rightarrow \infty$  implies  $P \rightarrow f$  in  $L^\infty$ . From the previous theorem, one would need to guarantee  $\|f^{(n)}\| \rightarrow 0$  (or at least, doesn’t grow faster than  $\frac{h^n}{4n}$ ), which certainly won’t hold in general; we have no control on the  $n$ th-derivative of an arbitrarily given function. However, we can try to optimize our choice of points  $\{x_j\}$  for a given  $j$ .

We switch notation for convention’s sake to  $n + 1$  points  $x_j$ . Our goal is the optimization problem

$$\min_{x_j} \max_{x \in [a, b]} \left| \prod_j (x - x_j) \right|,$$

the only term in the error bound above that we have control over. Remark that we can expand the product term:

$$\prod_j (x - x_j) = x^n - r(x),$$

where  $r(x) \in \mathbb{P}_n$ . So, really, we equivalently want to solve the problem

$$\min_{r \in \mathbb{P}_n} \|x^{n+1} - r(x)\|_\infty,$$

namely, what  $n$ -degree polynomial minimizes the max difference between  $x^{n+1}$ ?

↪**Theorem 1.3** (De la Vallée-Poussin Oscillation Theorem): Let  $f \in C([a, b])$ , and suppose  $r \in \mathbb{P}_n$  for which there exists  $n + 2$  distinct points  $\{x_j\}$  such that  $a \leq x_0 < \dots < x_{n+1} \leq b$  at which the error  $f(x) - r(x)$  “oscillate” sign, i.e.

$$\text{sign}(f(x_j) - r(x_j)) = -\text{sign}(f(x_{j+1}) - r(x_{j+1})).$$

Then,

$$\min_{P \in \mathbb{P}_n} \|f - P\|_\infty \geq \min_{0 \leq j \leq n+1} |f(x_j) - r(x_j)|.$$