

# Exercise 1:

1: First combining Y and I and then adding frequencies and combine it with T and again ~~combining~~ adding and combine with another lowest frequency and so on.

~~E: 4, S: 3, W: 2, N: 2, Y: 1, I: 1, T: 1~~

E: 4, S: 3, W: 2, N: 2, ~~Y: 1, I: 1~~ IT: 2

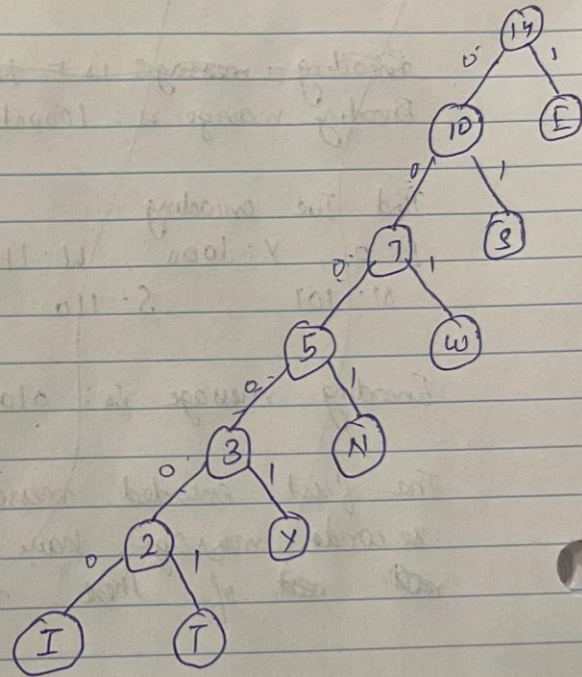
E: 4, S: 3, W: 2, N: 2, ~~IT: 2~~ IY: 3

E: 4, S: 3, W: 2, NITY: 5

E: 4, S: 3, WNITY: 7

E: 4, SWNITY: 10

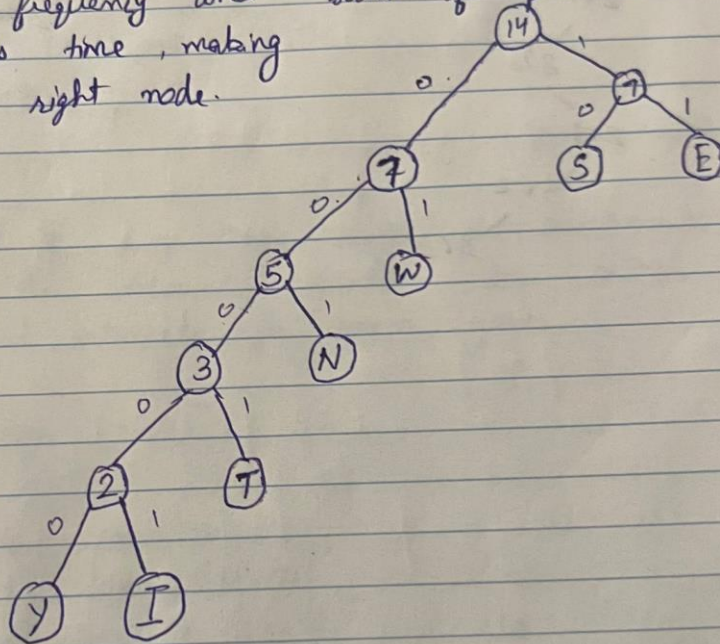
ESWNITY: 14



Frequencies are

E: 4, Y: 1, W: 2, I: 1, T: 1, N: 2, S: 3

Again, combining two nodes with lowest frequencies, Y and I, and sum that frequencies, and again combine one with lowest frequency and sum of previous two trees and this time, making N and S as right node.



Ist.

E: 11, Y: 00000, W: 01, I: 00001, T: 0001, N: 001, S: 10

Message: 11000001101000010001001111010001110110

2nd.

E: 1, Y: 00001, W: 001, I: 000000, T: 000001, N: 0001, S: 01

Message: 1000011001000000000000001000110101000100101

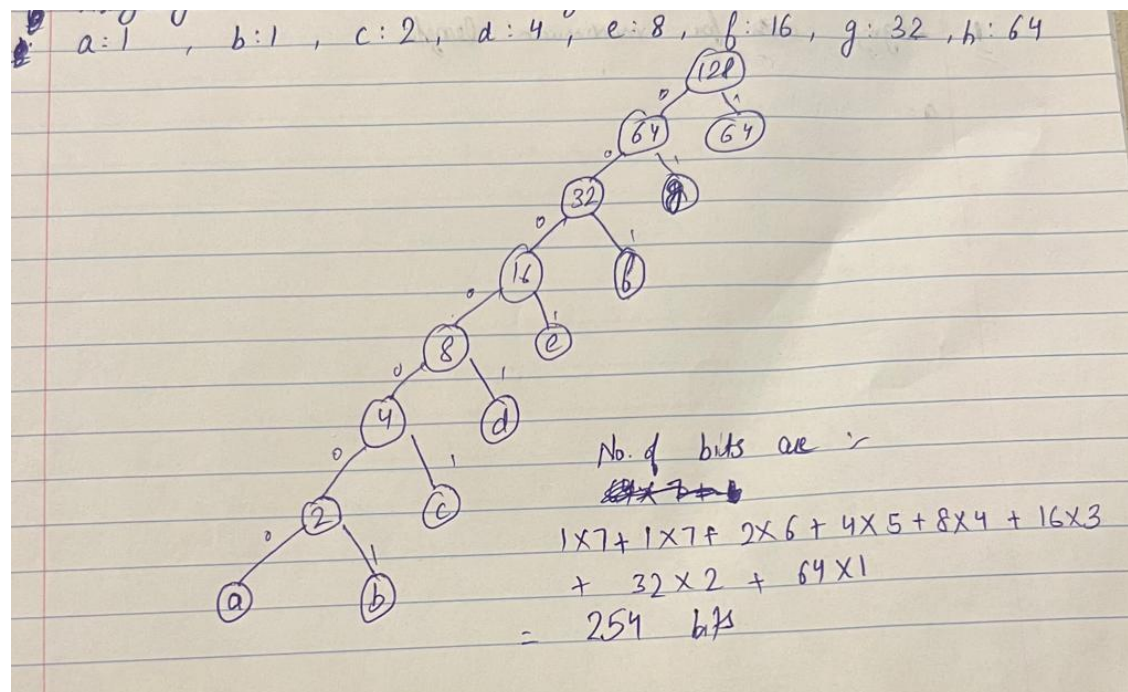
Ist message have 38 bits and 2nd have 41 bits. They are of different size.



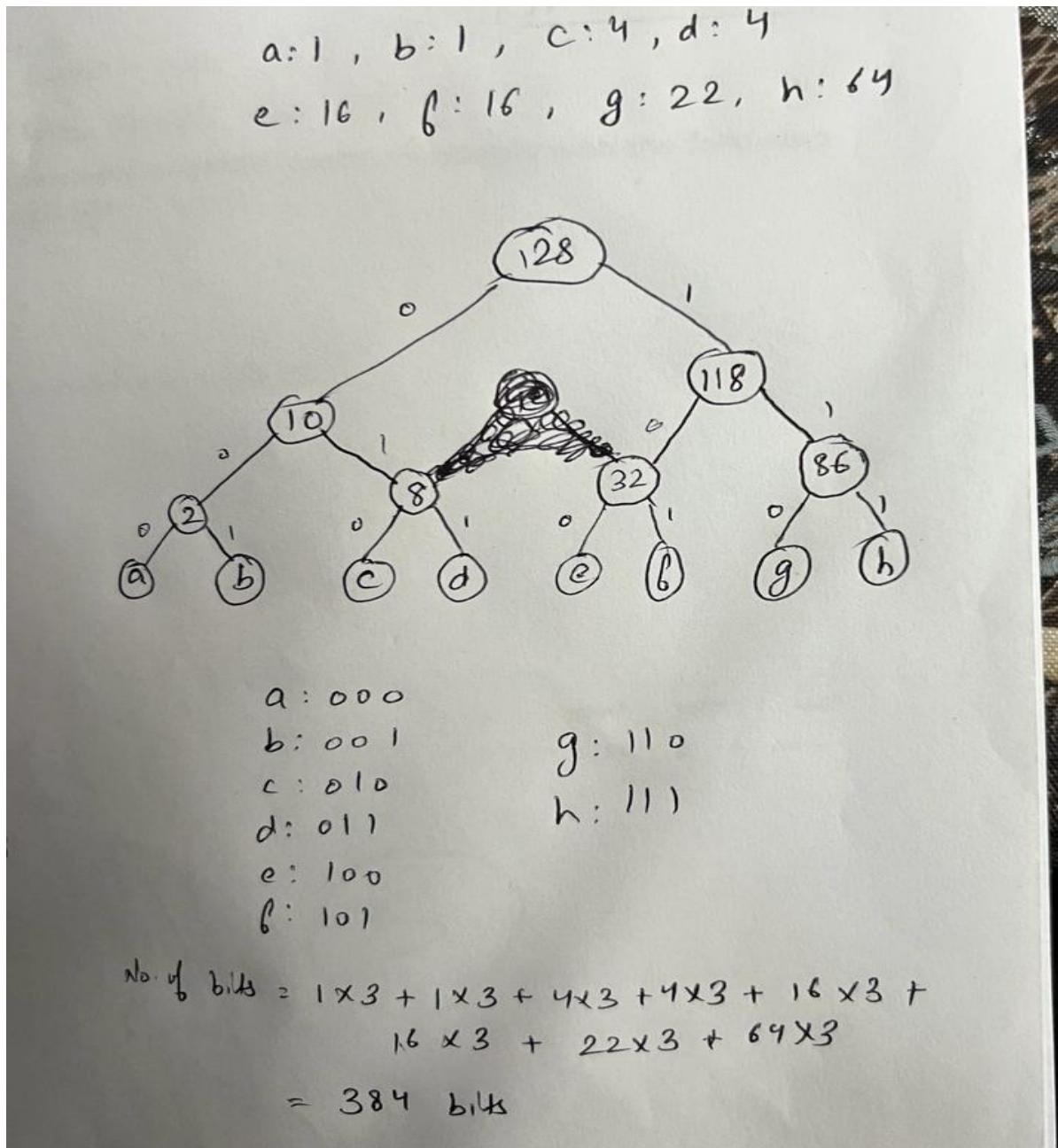
3. ~~The~~ Yes, the length of depends on our choice in building the tree. The best encoding tree depends on the frequencies and the compression ratio of the message. ~~It also depends on~~ If ~~the~~ ~~frequency~~ the frequency is high and make ~~it~~ ~~that~~ in deepest node. ~~and low frequencies~~ ~~with~~ It use lower bits as compared to others.

## Exercise 2:

### 1. For max height



2. For min height



Exercise 4:

Complexity for makeTree()

Here First code runs  $n$  times to get frequencies and another for loop again runs  $n$  times to pushing the nodes.

While loop runs until queue is not 1. Each time it pop 2 elements and push one into the node, it make complexity  $\log n$ .

Overall complexity is  $n \log n$ .

Complexity for print()

Time complexity is  $n$ ; because it goes through each node in a tree and print that node if its left and right node is nullptr.