

Graph

$$G = (V, E)$$

vertex set $\rightarrow E \subseteq V \times V$

m edges

n vertices

undirected:

$$(u, v) \in E \iff (v, u) \in E$$

directed:

adjacency matrix

$$m_{ij} = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}$$

adjacency list

1: 3, 7, 10

2: 1, 6, 9

3

	matrix	list
$(u, v) \in E$	$O(1)$	$O(n)$ $O(d)$
$\text{neigh}(v)$	$O(n)$	$O(n)$ $O(d)$
space	$O(n^2)$	$O(m)$

mark vertices

reachable from v .

Connectivity.

Explore (G, v)

depth-first search (DFS)

visited $[v] = \text{true}$.

DFS (G) :

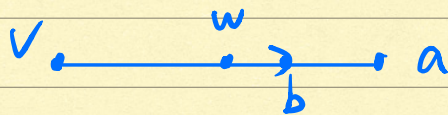
for each $(v, u) \in E$:

- $\forall v \in V, \text{visited}[v] = \text{false}.$
- $\forall v \in V; \text{if not visited}[v]$
 $\text{explore}(G, v)$

if not visited[u];
 $\text{explore}(G, u)$

claim: $\text{explore}(G, v)$ visits every node reachable from v :

Proof: suppose not. let a be a vertex that is not visited and is reachable



Runtime of DFS: $O(n+m)$

Connectivity in undirected graph.

$G \rightarrow$ connected components
 $u, v \in V$ are connected.

$\begin{pmatrix} u \rightarrow v \\ v \rightarrow u \end{pmatrix}$

$CC(G)$:

- $\forall v \in V, \text{visited}[v] = \text{false}.$
 $CCNum[v] = \text{null}.$
- $cc := 0.$

$\text{Explore}(G, v)$

- $\text{visited}[v] = \text{true}.$
- $CCNum[v] = cc.$
- $\forall (v, u) \in E:$

- $V \in V$;
if not visited.

C++

explore(h, v)

if not visited[u]:
 Explore(h, u)

Maintain a global clock for time on stack.

PFS (n):

- $U \cup \{v\}$, $visited[v] = \text{false}$.

- $\boxed{\text{prev}[u] = \text{null}, \text{post}[u] = \text{null},}$

- clock = 0.

- $\forall v \in V$, if not visited $[v]$:
 $\text{explore}(h, v)$

$\text{Explore}(G, u) :$

- visited[u] = true.

- `prev = clock`.
- `clock++`.

- $\forall (v, u) \in E$, if not visited u :
 $\text{explore}(G, u)$

- $post[0] = \text{clock}$
- $\text{clock} + t_i$

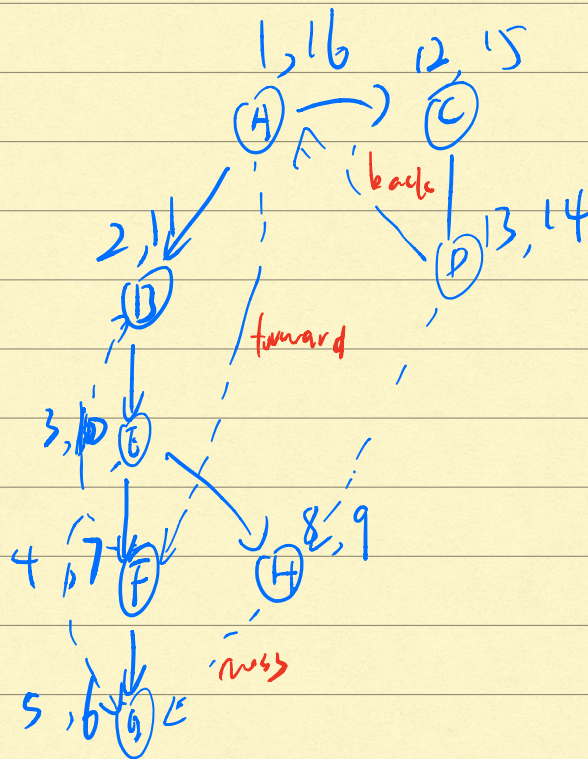
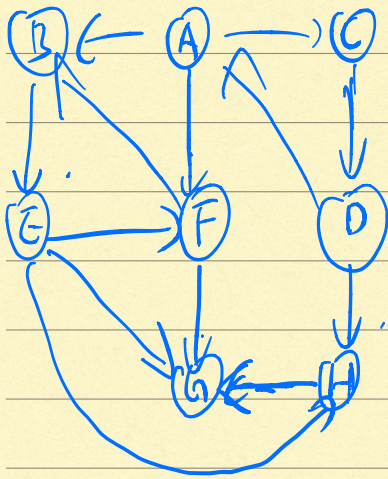
Read off cc :

Type of intervals:

$(u, v) \in E \rightarrow$

[illegible]

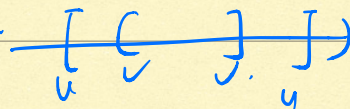
3 connected components.



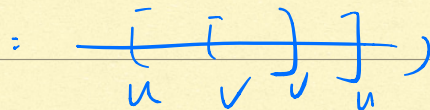
Type of edges:

$$(u, v) \in E$$

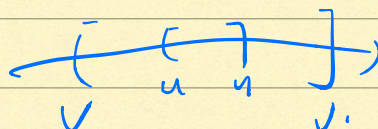
• tree edge (solid).



• forward edge



• back edge



• cross edge.

