COL726 Homework 4

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1. Since A is diagonal, we have:

$$\lambda_1 = a_{11} \quad \& \quad \lambda_2 = a_{22}$$
 (1)

Since the eigenvectors remain normalized under perturbation i.e. $||q_1 + \delta q_1||_2 = 1$, using this along with the fact that $q_1^T q_1 = 1$, we have

$$||q_1 + \delta q_1||_2^2 = (q_1 + \delta q_1)^T (q_1 + \delta q_1) = 1$$
$$1 + q_1^T \delta q_1 + \delta q_1^T q_1 + \delta q_1^T \delta q_1 = 1$$
$$q_1^T \delta q_1 + \delta q_1^T q_1 \approx 0 \Rightarrow 2q_1^T \delta q_1 \approx 0$$

Or simply,

$$q_1^T \delta q_1 = 0 \tag{2}$$

If we write δq_1 as $\begin{bmatrix} \delta q_{11} \\ \delta q_{21} \end{bmatrix}$, using $q_1 = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$, we get

$$\delta q_{11} = 0 \tag{3}$$

Let the perturbed eigenvalue corresponding to the eigenvector $q_1 + \delta q_1$ be $\lambda_1 + \delta \lambda_1$. We then have

$$(A + \delta A)(q_1 + \delta q_1) = (\lambda_1 + \delta \lambda_1)(q_1 + \delta q_1)$$

Since $Aq_1 = \lambda_1 q_1$, we get

$$A\delta q_1 + \delta Aq_1 + \delta A\delta q_1 = \lambda_1 \delta q_1 + \delta \lambda_1 q_1 + \delta \lambda_1 \delta q_1$$

Dropping the second order terms and replacing the approximation with equality, we get

$$A\delta q_1 + \delta A q_1 = \lambda_1 \delta q_1 + \delta \lambda_1 q_1 \tag{4}$$

Multiply equation 4 with q_1^T to get

$$q_1^T A \delta q_1 + q_1^T \delta A q_1 = \lambda_1 q_1^T \delta q_1 + \delta \lambda_1 q_1^T q_1$$

Using equations 2 and 3, we have

$$\begin{bmatrix} 1 & 0 \end{bmatrix} \begin{bmatrix} a_{11} & 0 \\ 0 & a_{22} \end{bmatrix} \begin{bmatrix} 0 \\ \delta q_{21} \end{bmatrix} + \begin{bmatrix} 1 & 0 \end{bmatrix} \begin{bmatrix} \delta a_{11} & \delta a_{12} \\ \delta a_{21} & \delta a_{22} \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} = 0 + \delta \lambda_1$$

Therefore,

$$\delta \lambda_1 = 0 + \delta a_{11} = \delta a_{11} \tag{5}$$

Put the values obtained in equations 1, 3 and 5 in the equation 4 to get

$$\begin{bmatrix} a_{11} & 0 \\ 0 & a_{22} \end{bmatrix} \begin{bmatrix} 0 \\ \delta q_{21} \end{bmatrix} + \begin{bmatrix} \delta a_{11} & \delta a_{12} \\ \delta a_{21} & \delta a_{22} \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} = a_{11} \begin{bmatrix} 0 \\ \delta q_{21} \end{bmatrix} + \delta a_{11} \begin{bmatrix} 1 \\ 0 \end{bmatrix}$$
$$\begin{bmatrix} 0 \\ \delta q_{21}(a_{22} - a_{11}) \end{bmatrix} = \begin{bmatrix} 0 \\ -\delta a_{21} \end{bmatrix}$$

Therefore using equation 1,

$$\delta q_1 = \begin{bmatrix} \delta q_{11} \\ \delta q_{21} \end{bmatrix} = \begin{bmatrix} 0 \\ \frac{-\delta a_{21}}{a_{22} - a_{11}} \end{bmatrix} = \begin{bmatrix} 0 \\ \frac{\delta a_{21}}{\lambda_1 - \lambda_2} \end{bmatrix}$$

When $\lambda_1 \approx \lambda_2$, we have $||\delta q_1|| = \frac{|\delta a_{21}|}{|\lambda_1 - \lambda_2|} >> 0$. Hence, the eigenvectors are ill-conditioned when the eigenvalues are close to each other.

2. (a) Updating Algorithm 10.1 of Trefethen & Bau for an upper Hessenberg Matrix(H = QR) results in:

$$\begin{split} &\text{i. for } k=1 \text{ to } m-1 \\ &\text{ii. } & x=H_{k:k+1,k} \\ &\text{iii. } & v_k=sign(x_1)||x||_2e_1+x \\ &\text{iv. } & v_k=v_k/||v_k||_2 \\ &\text{v. } & H_{k:k+1,k:m}=H_{k:k+1,k:m}-2v_k(v_k^TH_{k:k+1,k:m}) \\ &\text{vi. } v_m=sign(H_{m,m}) \\ &\text{vii. } H_{m,m}=-H_{m,m} \end{split}$$

H is modified in place and the resultant is R. v_1, \cdots, v_m are the modified reflection vectors. To get the reflection vectors as specified in the book(v_k^b), append zeros at the end of vector v_k , i.e. $v_k^b = [v_k; zeros(min(m-k-1,0))]$. Now, in the for loop, steps ii - iv take constant time - O(1). Step v takes atmost O(m) time per iteration. Therefore total time taken is $O(m*m) = O(m^2)$

(b) Now, suppose we have $(m+1) \times (m+1)$ upper Hessenberg matrix H' as follows:

$$H' = \begin{bmatrix} H & p \\ q^T & r \end{bmatrix}$$

where p and q are m dimensional vectors and ${\bf r}$ is a scalar. Also, q is of the form $\begin{bmatrix} 0 \\ \vdots \\ 0 \\ q_m \end{bmatrix}$.

In the algorithm above, $v_1, v_2, \cdots, v_{m-1}$ are not modified for the QR factorization of H'(H' = Q'R'). Also, if you look at step v, it can be written as $H_{k:k+1,i} = H_{k:k+1,i} - 2(v_k v_k^T) H_{k:k+1,i}$ for each column i. Hence the QR factorization of H' is identical to QR factorization of H for first m-1 columns. So, the update rule i.e. step v is to be performed only for columns m and m+1.

Now, first construct $H' = \begin{bmatrix} R & p \\ q^T & r \end{bmatrix}$ where p,q,r are defined above and R is obtained from the QR factorization of H. Note that $\{m,m\}$ element of R is updated as $R_{m,m} = -R_{m,m}$ before putting in H'. The entire algorithm is stated below:

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\begin{array}{lll} \text{ii.} & for \ k=1 \ \text{to} \ m-1 \\ \text{iii.} & v_k'=v_k \\ \text{iii.} & H_{k:k+1,m+1}'=H_{k:k+1,m+1}'-2v_k'(v_k'^TH_{k:k+1,m+1}') \\ \text{iv.} & x=H_{m:m+1,m}' \\ \text{v.} & v_m'=sign(x_1)||x||_2e_1+x \\ \text{vi.} & v_m'=v_m'/||v_m'||_2 \\ \text{vii.} & H_{m:m+1,m:m+1}'=H_{m:m+1,m:m+1}'-2v_m'(v_m'^TH_{m:m+1,m:m+1}') \\ \text{viii.} & v_{m+1}=sign(H_{m+1,m+1}') \\ \text{ix.} & H_{m+1,m+1}'=-H_{m+1,m+1}' \end{array}
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The resultant H' is R'. Step iii takes O(1) time per iteration. Therefore total time for steps i - iii is O(m). Steps iv - ix take O(1) time. So total time for entire algorithm is O(m).

3. Suppose v be an eigenvector of A with corresponding eigenvalue λ . We have

$$Av = \lambda v \Rightarrow A^{-1}Av = \lambda A^{-1}v \Rightarrow A^{-1}v = \frac{1}{\lambda}v$$

Therefore, A^{-1} has eigenvector v with the corresponding eigenvalue $\frac{1}{\lambda}$.

Also, since A is non-singular and symmetric, A has real eigenvalues and a complete set of orthonormal eigenvectors. Let the eigenvalues of A be $\lambda_1, \dots, \lambda_m$ and the corresponding eigenvectors be u_1, \dots, u_m . Without loss of generality, assume that

$$|\lambda_m| > \cdots > |\lambda_1|$$

.

(a) Now, algorithm 27.1 in Trefethen & Bau(Power Iteration) gives the eigenvector corresponding to the largest eigenvalue of A in absolute terms. Also,

$$v^* = \{v | Av = \lambda_{min}^A v\} = \{v | A^{-1}v = \frac{1}{\lambda_{min}^A} v\} = \{v | A^{-1}v = \lambda_{max}^{A^{-1}} v\}$$

. So, we will apply power iteration on A^{-1} to get the eigenvector corresponding to maximum eigenvalue of A^{-1} . Let the black box be treated as functions f and f^{-1} on some vector x, i.e. f(x) = Ax and $f^{-1}(x) = A^{-1}x$ respectively. Power Iteration algorithm is detailed below using the black box:

- i. $v^{(0)} =$ some vector with $||v_0|| = 1$
- ii. for $k = 1, 2, \cdots$
- iii. $w = f^{-1}(v^{(k-1)})$
- iv. $v^{(k)} = w/||w||$
- $v^{(k)}$ converges to either u_1 or $-u_1$, i.e. $v^{(k)} = \pm u_1$.
- (b) The k^{th} smallest eigenvalue of A in absolute terms is λ_k . Therefore, $\frac{1}{\lambda_k}$ is the k^{th} largest eigenvalue of A^{-1} in absolute terms. Now, we will use Simultaneous Iteration(Algorithm 28.3 Trefethen & Bau) to get the eigenvector corresponding to k^{th} largest eigenvalue of $A^{-1}(k^{th})$ smallest eigenvalue of A).
 - i. Pick $\hat{Q}^{(0)} \in \mathbb{R}^{m \times k}$ with orthonormal columns $q_1^{(0)}, \cdots, q_k^{(0)}$.
 - ii. for $i = 1, 2, \cdots$
 - iii. $Z = [f^{-1}(q_1^{i-1}) \cdots f^{-1}(q_k^{i-1})]$
 - iv. $\hat{Q}^{(i)}\hat{R}^{(i)}=Z$ reduced QR factorization of Z

The vectors $q_1^{(i)}, \dots, q_k^{(i)}$ converge to $\pm u_1, \dots, \pm u_k$ respectively. And u_k is the eigenvector corresponding to k^{th} largest eigenvalue of $A^{-1}(k^{th})$ smallest eigenvalue of A). Step iii takes O(mk) time(copying the vectors to the matrix Z). Step iv is the reduced QR factorization, which can be computed using modified Gram-Schmidt orthogonalization in $O(2mk^2)$ time(Modified Gram-Schmidt gives \hat{Q} directly). Therefore, total time taken is $O(mk^2)$.

4. If two matrices have the same characteristic polynomial, they have same eigenvalues since the roots of the characteristic polynomial are the eigenvalues. Let the characteristic polynomial of a matrix A be denoted as $p_A(\lambda)$. We have

$$p_A(\lambda) = \det(A - \lambda I) = \det((A - \lambda I)^T) = \det(A^T - \lambda I^T) = \det(A^T - \lambda I) = p_{A^T}(\lambda)$$

Therefore, the eigenvalues of A and A^T are same.

For an orthogonal matrix Q, we have $Q^TQ=QQ^T=I\Rightarrow Q^T=Q^{-1}$. Now, let

$$X = Q^T A Q = Q^{-1} A Q = \begin{bmatrix} A_{11} & A_{12} \\ \mathbf{0} & A_{22} \end{bmatrix}$$

Now,

$$p_X(\lambda) = det(Q^{-1}AQ - \lambda I) = det(Q^{-1}(A - \lambda I)Q) = det(Q^{-1})det(A - \lambda I)det(Q) = det(A - \lambda I) = p_A(\lambda)$$

Therefore X and A have the same eigenvalues. Now, suppose v be any n dimensional eigenvector of A_{11} with eigenvalue λ . Define $v' = \begin{bmatrix} v \\ \mathbf{0} \end{bmatrix}$. v' is m dimensional. We have,

$$Xv' = \begin{bmatrix} A_{11} & A_{12} \\ \mathbf{0} & A_{22} \end{bmatrix} \begin{bmatrix} v \\ \mathbf{0} \end{bmatrix} = \begin{bmatrix} A_{11}v + A_{12}\mathbf{0} \\ \mathbf{0}v + A_{22}\mathbf{0} \end{bmatrix} = \begin{bmatrix} A_{11}v \\ \mathbf{0} \end{bmatrix} = \lambda \begin{bmatrix} v \\ \mathbf{0} \end{bmatrix} = \lambda v'$$

Therefore λ is an eigenvalue of X. Since, the choice of v was arbitrary, the eigenvalues of A_{11} are also eigenvalues of X. We proved that X and A have the same eigenvalues. Consequently, the eigenvalues of A_{11} are also eigenvalues of A.

Now, suppose v be any (m-n) dimensional eigenvector of A_{22}^T with eigenvalue λ . Define $v' = \begin{bmatrix} \mathbf{0} \\ v \end{bmatrix}$. v' is m dimensional. We have,

$$X^Tv' = \begin{bmatrix} A_{11}^T & 0 \\ A_{12}^T & A_{22}^T \end{bmatrix} \begin{bmatrix} \mathbf{0} \\ v \end{bmatrix} = \begin{bmatrix} A_{11}^T \mathbf{0} + \mathbf{0}v \\ A_{12}^T \mathbf{0} + A_{22}^T v \end{bmatrix} = \begin{bmatrix} \mathbf{0} \\ A_{12}^T v \end{bmatrix} = \lambda \begin{bmatrix} \mathbf{0} \\ v \end{bmatrix} = \lambda v'$$

Therefore λ is an eigenvalue of X^T . Since, the choice of v was arbitrary, the eigenvalues of A_{22}^T are also eigenvalues of X^T . Since we proved that eigenvalues of any matrix M and M^T are same, we get that the eigenvalues of A_{22} are also eigenvalues of X. We proved that X and X have the same eigenvalues. Consequently, the eigenvalues of X are also eigenvalues of X.

5. We have

$$f(x) = p(x) = x^3 - x^2 - x + 1 = (x+1)(x-1)^2$$

$$f'(x) = p'(x) = 3x^2 - 2x - 1 = (x-1)(3x+1)$$

$$f''(x) = p''(x) = 6x - 2 = 2(3x-1)$$

(a) In Newton's method, we have

$$g(x) = x - \frac{f(x)}{f'(x)} \tag{6}$$

Therefore, the $(k+1)^{th}$ iteration is given by:

$$x_{k+1} = g(x_k) = x_k - \frac{p(x_k)}{p'(x_k)} = x_k - \frac{x_k^3 - x_k^2 - x_k + 1}{3x_k^2 - 2x_k - 1}$$

Now, the 4 iterations performed with $x_0 = -2$ and $x_0 = 2$ give the following iterates:

| $\mathbf{x_0}$ | -2.00000 | 2.00000 |
|------------------|----------|---------|
| $\overline{x_1}$ | -1.40000 | 1.57143 |
| x_2 | -1.10000 | 1.31429 |
| x_3 | -1.00870 | 1.16714 |
| x_4 | -1.00008 | 1.08667 |

(b) Given solution x^* , the error after k^{th} iteration is given by

$$e_{k+1} = x_{k+1} - x^* = g(x_k) - g(x^*)$$

Using Taylor Series, we have

$$e_{k+1} = g(x_k) - g(x^*) = g'(x^*)(x_k - x^*) + g''(x^*)\frac{(x_k - x^*)^2}{2} + \cdots$$

Using $e_k = x_k - x^*$,

$$e_{k+1} = g'(x^*)e_k + g''(x^*)\frac{e_k^2}{2} + \cdots$$
 (7)

Now, using equation 6,

$$g'(x) = \frac{f(x)f''(x)}{(f'(x))^2} = \frac{2(x+1)(3x-1)}{(3x+1)^2}$$

For simple root $x^* = -1$, $g'(x^*) = 0$, therefore the rate of convergence is quadratic (from equation 7):

$$\lim_{k \to \infty} \frac{|e_{k+1}|}{|e_k|^2} = C = \frac{|g''(x^*)|}{2}$$

For double root $x^* = 1$, $g'(x^*) = \frac{1}{2}$, therefore the rate of convergence is linear(from equation 7):

$$\lim_{k \to \infty} \frac{|e_{k+1}|}{|e_k|} = C = |g'(x^*)| = \frac{1}{2} = 0.5$$

(c) Let f(x) be any arbitrary function with double root a. Therefore we can write f(x) as $f(x) = (x-a)^2 h(x)$. We have

$$f'(x) = (x-a)^2 h'(x) + 2(x-a)h(x) = (x-a)((x-a)h'(x) + 2h(x))$$

$$f''(x) = ((x-a)h'(x) + 2h(x)) + (x-a)((x-a)h''(x) + h'(x) + 2h'(x))$$

Now, since $g(x) = x - \frac{f(x)}{f'(x)}$, we get

$$g'(x) = \frac{f(x)f''(x)}{(f'(x))^2} = \frac{(x-a)^2h(x)[(x-a)^2h''(x) + 4(x-a)h'(x) + 2h(x)]}{(x-a)^2((x-a)h'(x) + 2h(x))^2}$$
$$g'(a) = \frac{h(a)[2h(a)]}{(2h(a))^2} = \frac{1}{2} = 0.5$$
 (8)

Now, the rate of convergence starting from a point near a is linear(from equation 7):

$$\lim_{k \to \infty} \frac{|e_{k+1}|}{|e_k|} = C = |g'(a)| = \frac{1}{2} = 0.5$$

In normal newton's method, the update rule is

$$x_{k+1} = g(x_k) = x_k - \frac{f(x_k)}{f'(x_k)} = x_k - (1)\frac{f(x_k)}{f'(x_k)} = x_k - (\alpha)\frac{f(x_k)}{f'(x_k)}$$

Here α can be viewed as the learning rate. If we increase α , we can increase the convergence rate. Ideally, we want the convergence rate to be at least quadratic. Let g_1 be the updated fixed point iteration function($g_1(x) = x - \alpha \frac{f(x_k)}{f'(x_k)}$). For convergence to be at least quadratic near double root a, we have $g'_1(a) = 0$. We have

$$g_1(x) = x - \alpha \frac{f(x)}{f'(x)} = x + \alpha (x - \frac{f(x)}{f'(x)} - x) = x + \alpha (g(x) - x) = \alpha g(x) - (\alpha - 1)x$$

Therefore,

$$g_1'(x) = \alpha g'(x) - (\alpha - 1)$$

Since we want $g'_1(a) = 0$, using equation 8, we get

$$g_1'(a) = \alpha g'(a) - (\alpha - 1) = \frac{\alpha}{2} - (\alpha - 1) = 0 \Rightarrow \alpha = 2$$

Therefore, near a double root use updated fixed point iteration function $g_1(x) = x - 2\frac{f(x)}{f'(x)}$

6. We have the system of equations as $f(x) = \begin{bmatrix} f_1(x) \\ f_2(x) \end{bmatrix} = 0$. The jacobian J(x) is given by $J(x) = \begin{bmatrix} g_1(x) \\ g_2(x) \end{bmatrix}$

The update step in Newton's method is given by

$$x^{(k+1)} = g(x^{(k)}) = x^{(k)} - J_f^{-1}(x^{(k)})f(x^{(k)}) \Rightarrow J_f(x^{(k)})(x^{(k+1)} - x^{(k)}) = -f(x^{(k)})$$

Putting the values of J and f,

$$\begin{bmatrix} g_1(x^{(k)}) \\ g_2(x^{(k)}) \end{bmatrix} (x^{(k+1)} - x^{(k)}) = - \begin{bmatrix} f_1(x^{(k)}) \\ f_2(x^{(k)}) \end{bmatrix}$$

Therefore,

$$f_1(x^{(k)}) + g_1(x^{(k)})(x^{(k+1)} - x^{(k)}) = 0$$

and

$$f_2(x^{(k)}) + g_2(x^{(k)})(x^{(k+1)} - x^{(k)}) = 0$$

which interpreted in terms of $\mathtt{linearize}(\tilde{f})$ is simply $\tilde{f}_1(x^{(k+1)}) = 0$ and $\tilde{f}_2(x^{(k+1)}) = 0$ Hence the next newton iterate(point) is the point which lies at the intersection of the linearized curves of f_1 and f_2 , which in this case is the intersection of lines \tilde{f}_1 and \tilde{f}_2 . After convergence, $x^{(k+1)}$ is the point which lies on both f_1 and f_2 and the linearized curves \tilde{f}_1 and \tilde{f}_2 at this point are tangents to the respective curves.

I used 2 ellipses as conic1 and conic2 with equations $\frac{x_1^2}{16} + \frac{x_2^2}{4} = 1$ and $\frac{x_1^2}{4} + \frac{x_2^2}{16} = 1$ respectively.

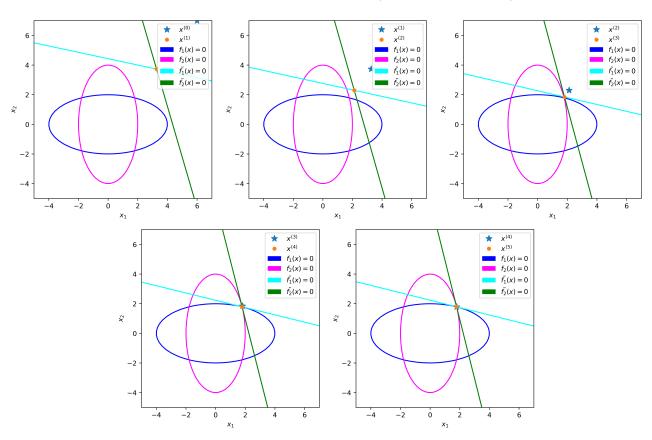


Figure 1: This figure shows the sequence of plots for $x^{(0)}, \dots, x^{(4)}$ with $x^{(k)}$ as the current Newton iterate and $x^{(k+1)}$ as the next Newton iterate along the conics and their respective linearized curves.