

Advanced network algorithms, random graph models

You are given four networks in Pajek format (edge list and LNA formats are also available).

- A simple [toy network](#) for testing (tiny)
- The famous [Zachary karate club network](#) (small)
- [iMDB actors collaboration network](#) (medium)
- A part of [Google web graph](#) (large)

I. Number and size of connected components (ii)

1. **(answer)** Study the following algorithm for computing (weakly) connected components $\{C\}$ by any order link traversal. Does the algorithm implement breadth-first or depth-first search? Why? What is the time complexity of the algorithm?

```
input  graph G, nodes N
output network components {C}
1: {C} ← empty list
2: while not N empty do
3:   {C}.add(component(G, N, N.next()))
4: return {C}
```

```
input  graph G, nodes N, node i
output weak component C
1: C ← empty list
2: S ← empty stack
3: N.remove(S.push(i))
4: while not S empty do
5:   C.add(i ← S.pop())
6:   for neighbors j ∈ Γi do
7:     if N.remove(j) then
8:       S.push(j)
9: return C
```

2. **(code)** Implement the algorithm and compute number of (weakly) connected components s and size of the largest (weakly) connected component S of all four networks. Are the results expected or are they surprising?
3. **(answer)** How could you further improve the algorithm to *only* compute s and S ?

II. Average node distance and network diameter

1. **(answer)** Study the following algorithm for computing distances between the nodes $\{d\}$ by level order link traversal. Does the algorithm implement breadth-first or depth-first search? Why? What is the time complexity of the algorithm?

```

input  graph G
output network distances {D}
1: {D} ← empty list
2: for nodes  $i \in N$  do
3:   {D}.add(distances(G, i))
4: return {D}

```

```

input  graph G, node i
output directed distances D
5: ...
6: for successors  $j \in \Gamma_i^{out}$  do
7:   ...

```

```

input  graph G, node i
output undirected distances D
1: D ← empty array
2: Q ← empty queue
3: D[Q.add(i)] ← 0
4: while not Q empty do
5:   i ← Q.remove()
6:   for neighbors  $j \in \Gamma_i$  do
7:     if D[j] undefined then
8:       D[j] ← D[i] + 1
9:       Q.add(j)
10: return D

```

2. **(code)** Implement the algorithm and compute average distance between the nodes $\langle d \rangle$ and maximum distance or diameter d_{\max} of *smaller* networks. Are the results expected or are they surprising?
3. **(answer)** How is the algorithm different from the famous Dijkstra's algorithm? In which case you would necessarily have to use the Dijkstra's algorithm?
4. **(answer)** How could you speed up the algorithm to *only* approximate $\langle d \rangle$ and d_{\max} ?

III. Average node clustering coefficient

1. **(answer)** Study the following algorithm for computing node clustering coefficients $\{C\}$ by link triad counting. Why does the algorithm count triads over the links and not over the nodes? What is the time complexity of the algorithm?

```

input  graph G
output average clustering  $\langle C \rangle$ 
1:  $\langle C \rangle \leftarrow 0$ 
2: for nodes  $i \in N$  do
3:    $\langle C \rangle \leftarrow \text{clustering}(G, i)/n$ 
4: return  $\langle C \rangle$ 

input  graph G, node i
output node clustering C
1: if  $k_i \leq 1$  then
2:   return 0
3: return triads(G, i) · 2 / ( $k_i^2 - k_i$ )

```

```

input  graph G, node i
output node triads t
1: t ← 0
2: for neighbors  $j \in \Gamma_i$  do
3:   if  $|\Gamma_i| \leq |\Gamma_j|$  then
4:     t ← t + triads(G, i, j)/2
5:   else
6:     t ← t + triads(G, j, i)/2
7: return t

```

```

input  graph G, link i, j
output link triads t
1: t ← 0
2: for neighbors  $k \in \Gamma_i$  do
3:   if  $k \in \Gamma_j$  then
4:     t ← t + 1
5: return t

```

2. **(code)** Implement the algorithm and compute average node clustering coefficient $\langle C \rangle$ of all four

networks. Are the results expected or are they surprising?

3. **(answer)** What kind of network representation is required by the algorithm?

IV. Erdős-Rényi random graphs and link indexing

1. **(answer)** Study the following two algorithms for generating Erdős-Rényi random graphs $G(n, m)$ with and without link indexing $\binom{i}{2} + j, i > j$. What is the main difference between the algorithms? What is the time complexity of the algorithms?

```

input  nodes  $n$ , links  $m$ 
output simple random  $G$ 
1:  $H \leftarrow$  empty set
2:  $G \leftarrow n$  isolated nodes
3: while not  $G$  has  $m$  links do
4:    $h \leftarrow \{0, \dots, (n^2 - n)/2 - 1\}.\text{random}()$ 
5:   if  $H.\text{add}(h)$  then
6:      $i \leftarrow 1 + \lfloor -0.5 + \sqrt{0.25 + 2h} \rfloor$ 
7:     add link between  $i$  and  $h - (i^2 - i)/2$ 
8: return  $G$ 

```

```

input  nodes  $n$ , links  $m$ 
output random multi  $G$ 
1:  $G \leftarrow n$  isolated nodes
2: while not  $G$  has  $m$  links do
3:    $i, j \leftarrow \{0, \dots, n - 1\}.\text{random}()$ 
4:   if  $i \neq j$  then
5:     add link between  $i$  and  $j$ 
6: return  $G$ 

```

2. **(code)** Implement one of the algorithms and generate Erdős-Rényi random graphs corresponding to all four networks and compute graphs' S , $\langle d \rangle$ and $\langle C \rangle$. Are the results expected or are they surprising?

V. Configuration model graphs and link rewiring

1. **(answer)** Study the following two algorithms for generating configuration model graphs $G(\{k\})$ with link rewiring and stub matching. What is the main difference between the algorithms? What is the time complexity of the algorithms?

```

input  simple links  $L$ 
output configuration simple  $G$ 
1:  $H \leftarrow$  empty set
2: for links  $\{i, j\} \in L$  do
3:    $H.\text{add}(h_{ij})$ 
4: while not links  $L$  rewired do
5:    $\{i, j\}, \{s, t\} \leftarrow L.\text{random}()$   $\triangleright$  removes links
6:   if  $H.\text{contains}(h_{it} \text{ or } h_{sj})$  or  $i = t$  or  $s = j$  then
7:      $L.\text{add}(\{i, j\})$ 
8:      $L.\text{add}(\{s, t\})$ 
9:   else
10:     $L.\text{add}(\{i, t\})$   $H.\text{add}(h_{it})$   $H.\text{remove}(h_{ij})$ 
11:     $L.\text{add}(\{s, j\})$   $H.\text{add}(h_{sj})$   $H.\text{remove}(h_{st})$ 
12: return  $G$  on links  $L$ 

```

```

input  nodes  $n$ , degrees  $\{k\}$ 
output configuration pseudo  $G$ 
1:  $Q \leftarrow$  empty queue
2:  $G \leftarrow n$  isolated nodes
3: for nodes  $i \in N$  do
4:   for  $k_i$  times do
5:      $Q.\text{add}(i)$ 
6: while not  $Q$  empty do
7:    $i, j \leftarrow Q.\text{random}()$   $\triangleright$  removes nodes
8:   add link between  $i$  and  $j$ 
9: return  $G$ 

```

2. **(code)** Implement one of the algorithms and generate configuration model graphs corresponding to all four networks and compute graphs' S , $\langle d \rangle$ and $\langle C \rangle$. Are the results expected or are they surprising?

