# CSE 1729:Principles of Programming

# Lecture 10: Lists in Scheme

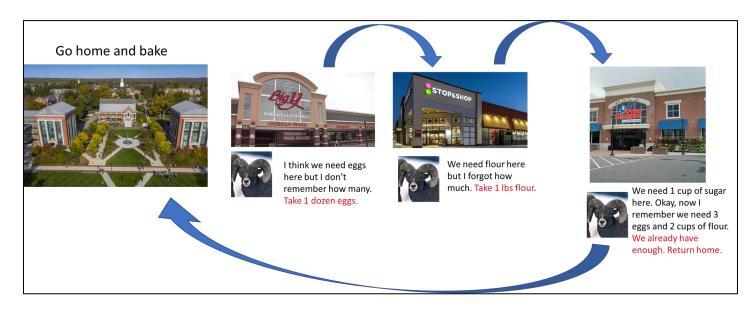
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# Previously in CSE 1729, two concepts...

#### Rams need Cake: Tail Recursion using a ridiculous example







# Previously in CSE 1729, two concepts...

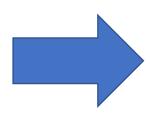
#### Pairs in Scheme:

```
PAIRS

• Construction

(define z (cons x y))
```

```
1 (define z (cons 2 3))
2 3 (car z)
4 5 (cdr z)
```



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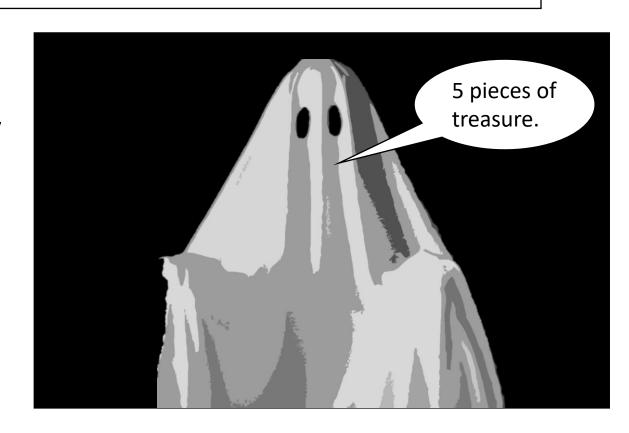
# **Ghost Guarded Treasure!**



- You stumble upon a cavern of riches.
- The treasure is guarded by a ghost who says you can only take a pouch in.
- The pouch can hold two items.
   Once you enter the cavern one time and leave, you can never return again.
- How do you recover all the treasure in one trip?

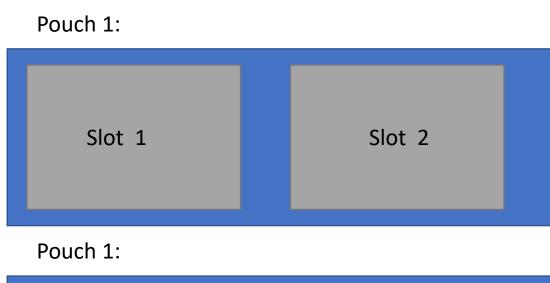
#### Solution to the Ghost Guarded Treasure problem

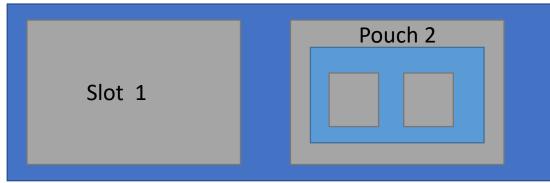
- The ghost says you can only take in a pouch that can hold two items.
- However the ghost did <u>NOT</u> specify the pouch had to be empty.
- Do the following:
- 1. Ask the ghost how many pieces of treasure they are guarding.
- 2. Get your first pouch. Keep the first slot empty. Put a SECOND pouch in the second slot of the first pouch.

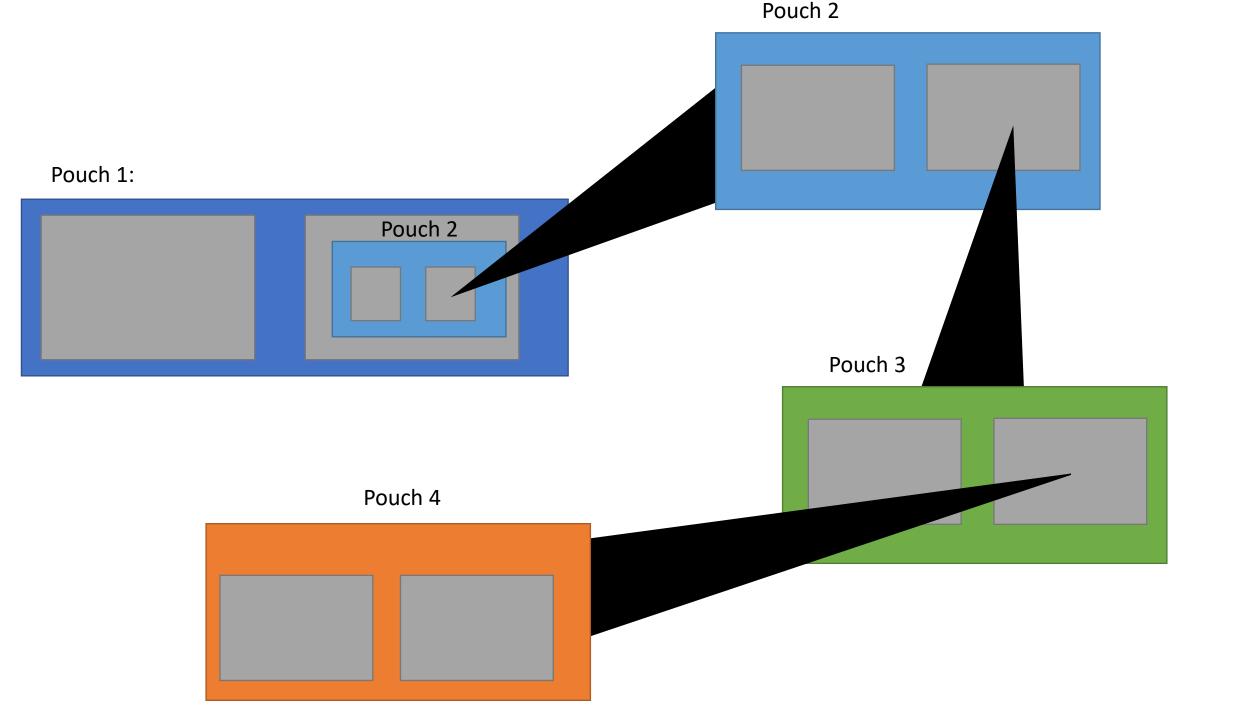


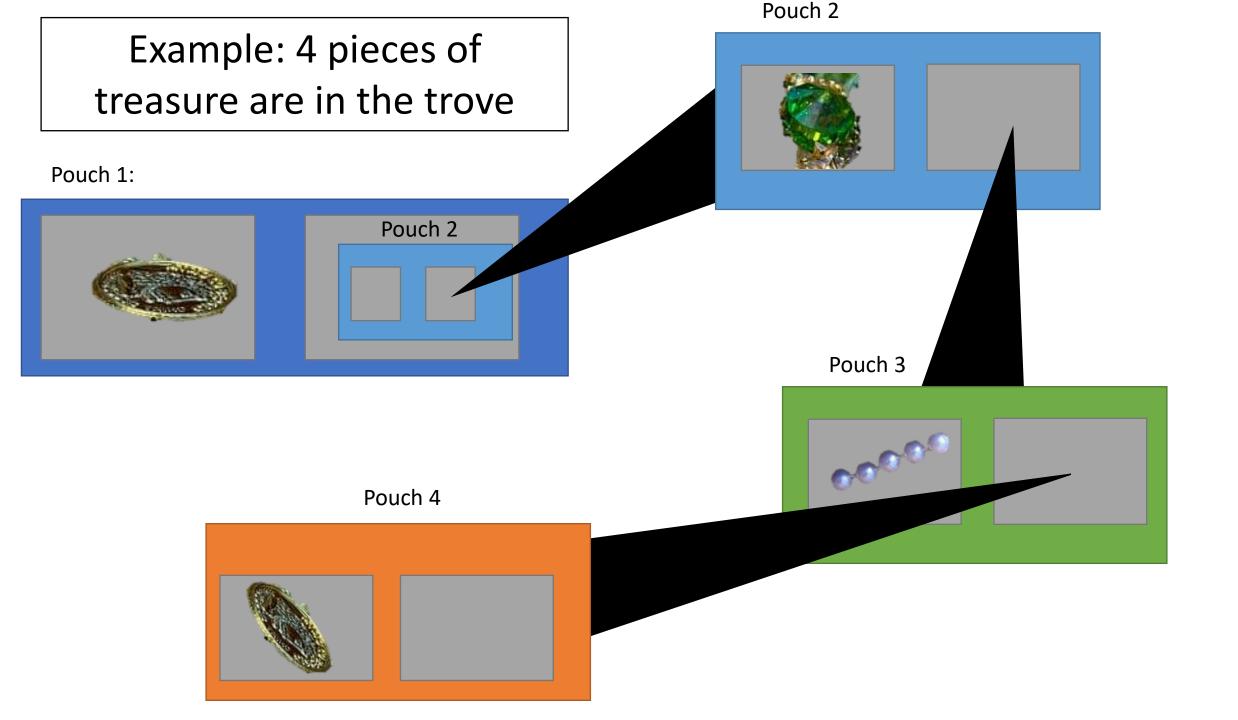
#### Solution to the Ghost Guarded Treasure problem

- Do the following:
- 1. Ask the ghost how many pieces of treasure they are guarding.
- 2. Get your first pouch. Keep the first slot empty. Put a SECOND pouch in the second slot of the first pouch.
- 3. In the second pouch, second slot, put another pouch.
- 4. Repeat this process until you have a number empty slots equal to the amount of treasured items in the cave.
- 5. Profit.











**Question**: In Scheme we learned pairs (an object that can hold two items). Can we build objects that can hold more than two items out of pairs?

# **Lists in Scheme**

- A list is an extremely flexible data structure that maintains an ordered list of objects, for example:
  - Ceres, Pluto, Makemake, Haumea, Eris, a list of 5 extrasolar planets.
- Scheme implements lists in terms of the pair structure you have already met.
  - However, pairs have only 2 slots, so we need a mechanism for using pairs to represent lists of arbitrary length.
- Roughly, Scheme uses the following recursive convention: the list of k objects  $a_1, \ldots, a_k$  is represented as a pair where...
  - The first element of the pair is the first element of the list  $a_1$ .
  - The second element of the list is...a list containing the rest of the elements.

### **BUILDING UP LISTS WITH PAIRS**

- To be more precise: A *list* is either
  - the empty list, or
  - a pair, whose first coordinate is the first element of the list, and whose second coordinate is a list containing the remainder of the elements.
- Note: the second element of the pair must be a list.

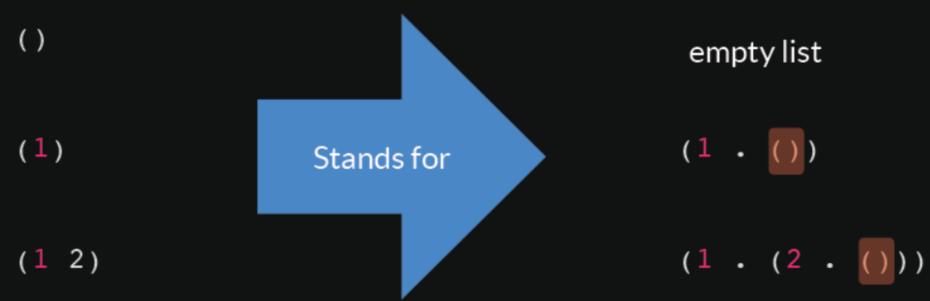
 For example, if • denotes the empty list, then... (1) $(1 \ 2)$ 

# A GENERAL LIST; SCHEME NOTATION

• Thus, a list has the form: Pair

First element List of remaining elements

Since lists are used so frequently, Scheme provides special notation for them:



Note: In Scheme, lists are always terminated with the empty list.

# **QUOTATION; ENTERING LISTS IN THE SCHEME INTERPRETER**

Recall the Scheme evaluation rule for compound (list!) objects.

(1)

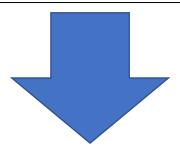
 This means that the natural way to enter a list doesn't work: Scheme wants to apply evaluation:

```
1 > ()
2 . #%app: missing procedure expression; probably originally (), which is an illegal empty application in: (#%app)
3 > (1 2)
4 . . procedure application: expected procedure, given: 1; arguments were: 2
```

Scheme provides the (quote <expr>) (or '<expr>) form, which evaluates to
 <expr> without further evaluation:

## List Syntax

```
1 (define x (list 1 2 3 ))
2 x
```



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# How do we know a list is also technically a pair?

Try using the car and cdr functions on the list.

```
1 (define x (list 1 2 3))
2 (car x)
3 (cdr x)
```

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# Can you put a list in a list?



```
1 (define y (list 1 (list 2 3) 4 5 ))
2 y
```

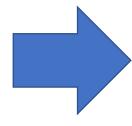


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# How do you check if a list is empty?

Just ask!

```
1  ; Empty List
2  (define x (list ))
3  ; Non-empty List
4  (define y (list 1 2 3 ))
5
6  (null? x)
7  (null? y)
```



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# LIST PROCESSING:

# HANDLE THE FIRST ELEMENTS AND, THEN,...HANDLE THE REST

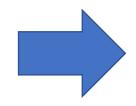
- (null? x) returns #t if x is the empty list.
- list processing:
  - handle the first element (the car) and, then,
  - handle the remaining list (the cdr).
  - Notice that these have different "types."
- Computing the length, for example...

```
(define (nlength xyz)
  (if (null? xyz)
      0
       (+ 1 (nlength (cdr xyz)))))
                 Then ...
> (nlength '(1 2 3))
3
  (nlength '())
0
  (nlength '((1 2) (3 4)))
2
```

# How do we find the end of an arbitrary list?

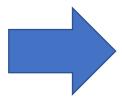
What if we repeatedly use the cdr function to reduce the list...

```
1 ;Non-empty List
2 (define x (list 1 2 3 ))
3 (cdr x)
```



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```
1 ;Non-empty List
2 (define x (list 1 2 3 ))
3 (cdr (cdr x))
```



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# Calling cdr one more time...

```
1 ; Non-empty List
2 (define x (list 1 2 3 ))
3 (cdr (cdr (cdr x)))

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Language: R5RS; memory limit: 128 MB
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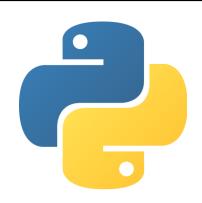
What happened? Scheme uses the empty list to indicate termination (that you final element of a list)

# Getting to the last element in a list...

- If list is a list, it is easy to get to the first element: (car list).
- The last element, however, takes more work to find.
  - This is an inherent feature (and, sometimes, shortcoming) of this "data structure."

# A side note...How did we get the last element of a list in Python?





```
(define (last-element 1)
  (if (null? (cdr 1))
        (car 1)
        (last-element (cdr 1))))
```

## **ANOTHER EXAMPLE: SUMMING THE NUMBERS OF A LIST**

Adding the elements of a list:

• Then...

```
1 > (sum-list '())
2 0
3 > (sum-list '(1 3 5 7))
4 16
```

#### How do we combine lists?

Use the built in append function...

```
(define x (list 1 2 3))
(define y (list 4 5 6))
(define z (append x y))
                                (1 2 3 4 5 6)
```

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# Append: What is going on under the hood?

Basic operation on lists: place one after the other:

# **HOW LONG DOES THIS TAKE?**

- A good measure of the "time taken" by a Scheme function (without looping constructs, which we will discuss later) is simply the number of recursive calls it generates.
- (append list1 list2) involves a total of length (lista) recursive calls.
   (Why? It needs to find the end of the list.)

```
(append list1 list2)

(append (cdr list1) list2)

(append (cdr (cdr list1)) list2)
```

# A LIST USER'S GUIDE...

- Suppose that L is a list in Scheme;
  - then you can tell if it is empty by testing (null? L); if not...
  - its first element is (car L);
  - the "rest" of the elements are (cdr L) (this is a list, and might be empty).

- Suppose that L is a list in Scheme and x is a value;
  - '() or (list) is the empty list.
  - (cons x L) is a new list—its first element is x; the rest of the elements are those of L.
  - The list containing only the value x? Same idea, but use the empty list for L: (cons x '()) or (list x).

# A new foe has appeared

CHALLENGER APPROACHING



# The Ordered Square List Challenge

• Hmm... it lists them backwards!

> (square-list 4)
(16 9 4 1 0)

# **SQUARES IN THE RIGHT ORDER**

It's easy if both ends of a range are given: (why did this make it easy?)

• We can wrap this in a definition that starts at zero:

```
(define (forward-squares k)
  (define (square x) (* x x))
  (define (squares start finish)
     (if (> start finish) '()
            (cons (square start) (squares (+ start 1) finish))))
  (squares 0 k))
```

# Figure Sources

- <a href="https://cdn.vox-cdn.com/thumbor/pr3jD5sfTRKpPinnYym\_4A0gJaQ=/0x27:4415x333">https://cdn.vox-cdn.com/thumbor/pr3jD5sfTRKpPinnYym\_4A0gJaQ=/0x27:4415x333</a> 8/1400x1400/filters:focal(0x27:4415x3338):format(jpeg)/cdn.vox-cdn.com/uploads/chorus image/image/43170476/ghosts.0.0.jpg
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