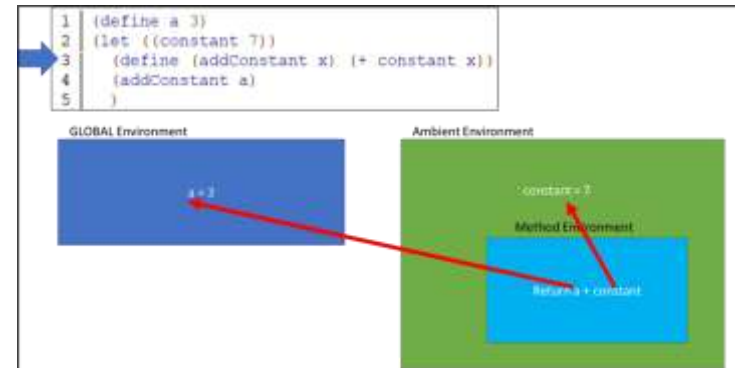


Lecture 6: Let There Be Scheme



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University of Connecticut

Previously in CSE 1729... *So many examples of recursion!!!!*

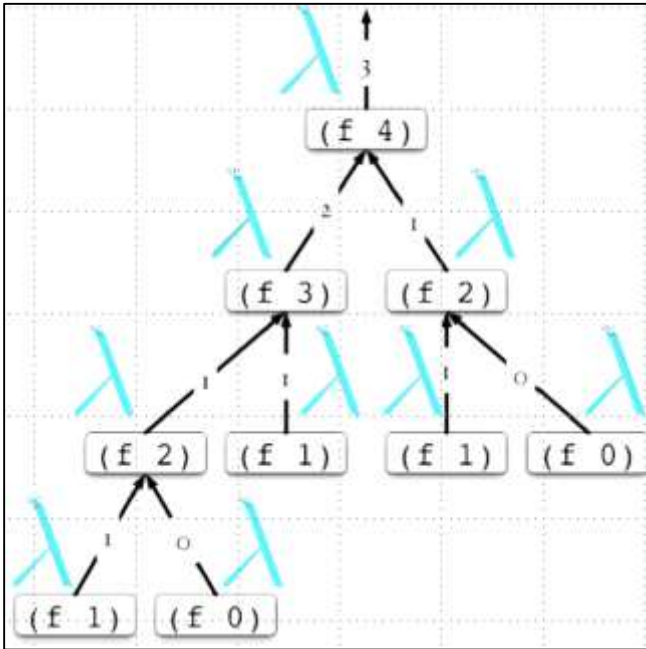
2. Multiplication through addition

EXAMPLE: MULTIPLICATION IN TERMS OF ADDITION

- Consider the definition of multiplication as repeated addition:

$$a \times b = \underbrace{b + b + \dots + b}_{a \text{ times}}$$

1. Fibonacci Function



2. Guessing a number using binary search



COMPUTING SQUARE ROOTS BY AVERAGING

- One simple way to compute an approximation to the square root of a number x is to
 - Start with two guesses, a and b , with the property that

$$a < \sqrt{x} < b$$

- (For example, if $x > 1$, we could start with $a = 1$, $b = x$.) Thus we know that the actual square root is between a and b .
- If $\frac{(a+b)}{2}$ is larger than the square root (which we can check by comparing $[\frac{(a+b)}{2}]^2$ with x) we know the real square root lies between a and $\frac{(a+b)}{2}$.
- Otherwise, the real square root lies between $\frac{(a+b)}{2}$ and b .

IN SCHEME

```
(define (average a b) (/ (+ a b) 2))  
(define (square a) (* a a))  
  
(define (sqrt-converge x a b)  
  (if (< (abs (- a b)) .000001)  
      a  
      (if (> (square (average a b)) x)  
          (sqrt-converge x a (average a b))  
          (sqrt-converge x (average a b) b)))))
```

Now, we might like to define a more attractive square root function that does not require choosing a and b:

```
(define (new-sqrt x) (sqrt-converge x 1 x))
```

LOCAL VARIABLES

- (average a b) is referred to several times in sqrt-converge.

Wouldn't it be nice if we could temporarily bind a "local" variable to this value?

- The let construct does exactly this:
- **Semantics:**
 - Evaluate each $\langle \text{expr}_i \rangle$, yielding a value v_i .
 - Create a new environment by starting with the current one and binding each x_i to v_i .
 - Then return the value of $\langle \text{body-expr} \rangle$ in this environment.

```
(let ((x1 <expr1>)
      (x2 <expr2>)
  ...
      (xk <exprk>))
<body-expr>)
```


Now I know what everyone is thinking.
This new “let” statement is so much cooler than “define”.

UConn CSE Students

New “let”
statement

Old “define”
statement



LOCAL VARIABLES

```
(define (sqrt-converge x a b)
  (let ((avg (/ (+ a b) 2)))
    (if (< (abs (- a b)) .000001)
      a
      (if (> (square avg) x)
          (sqrt-converge x a avg)
          (sqrt-converge x avg b))))))
```

The let statement binds avg to $\frac{(a+b)}{2}$ for the shaded block of code

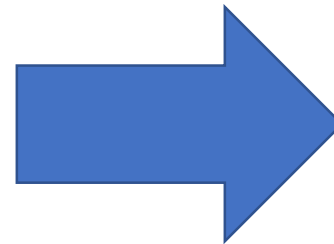


"No one man should have all that power"

-Kanye West,
probably referring to the
let statement in Scheme

How about we try using the let statement outside of a method call?

```
1 (define a 3)
2 (define constant 5)
3
4 (define (addConstant x)
5   (+ constant x)
6 )
7
8 (addConstant a)
```



```
Welcome to DrRacket, version 8.3 [cs].
Language: R5RS; memory limit: 128 MB.
8
> |
```

Can we rewrite this statement using “let”?

Rewriting the “add constant” method with the let statement.

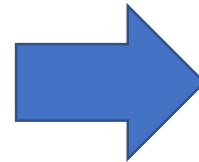
```
1 (define a 3)
2 (define constant 5)
3
4 (define (addConstant x)
5   (+ constant x)
6 )
7
8 (let ((constant 7))
9   (addConstant a))
```

What should the output of this code be?

10 right?

Rewriting the “add constant” method with the let statement.

```
1 (define a 3)
2 (define constant 5)
3
4 (define (addConstant x)
5   (+ constant x)
6 )
7
8 (let ((constant 7))
9   (addConstant a))
```



```
Welcome to DrRacket, version 8.3 [cs].
Language: R5RS; memory limit: 128 MB.
8
>
```

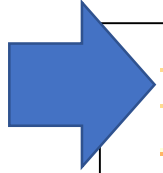
When Scheme gives you 8 but
you wanted 10:



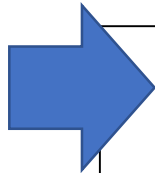
We need to understand the rules of scoping in Scheme



- The “let” statement behaves differently depending on WHERE in the code it is called.
- We will now trace through that example pictorial to understand.



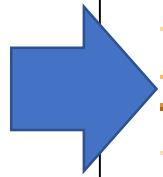
```
1 (define a 3)
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4 (define (addConstant x)
5   (+ constant x)
6 )
7
8 (let ((constant 7))
9   (addConstant a))
```



```
1 (define a 3)
2 (define constant 5)
3
4 (define (addConstant x)
5   (+ constant x)
6 )
7
8 (let ((constant 7))
9   (addConstant a))
```

GLOBAL Environment

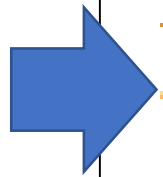
a = 3



```
1 (define a 3)
2 (define constant 5)
3
4 (define (addConstant x)
5   (+ constant x)
6 )
7
8 (let ((constant 7))
9   (addConstant a))
```

GLOBAL Environment

a = 3
constant = 5



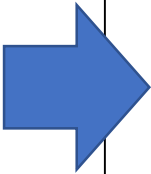
```
1 (define a 3)
2 (define constant 5)
3
4 (define (addConstant x)
5   (+ constant x)
6 )
7
8 (let ((constant 7))
9   (addConstant a))
```

GLOBAL Environment

a = 3
constant = 5

GLOBAL Environment

a = 3
constant = 5



```
1 (define a 3)
2 (define constant 5)
3
4 (define (addConstant x)
5   (+ constant x)
6 )
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9   (addConstant a))
```

```
1 (define a 3)
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3
4 (define (addConstant x)
5   (+ constant x)
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8 (let ((constant 7))
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```

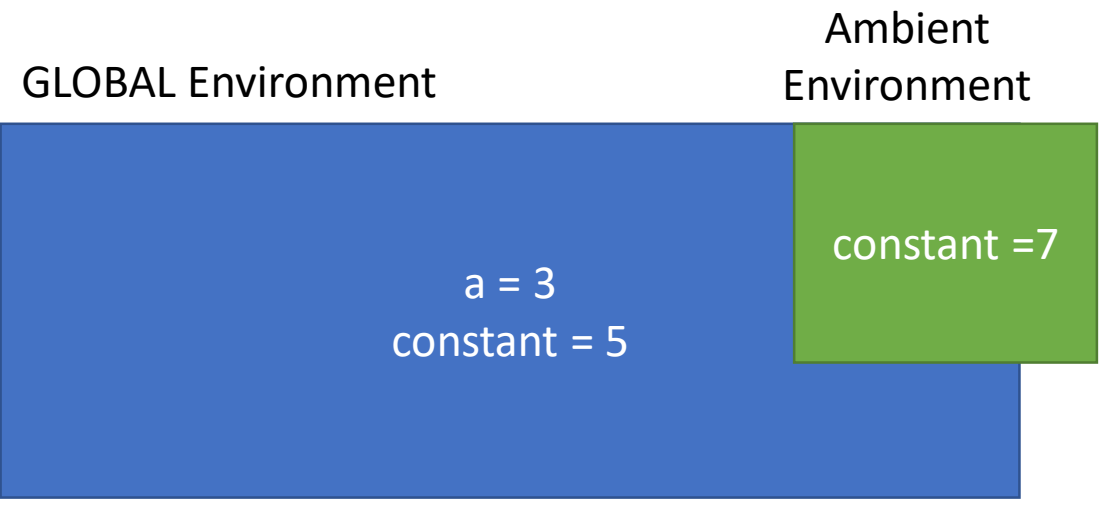
GLOBAL Environment

Ambient
Environment

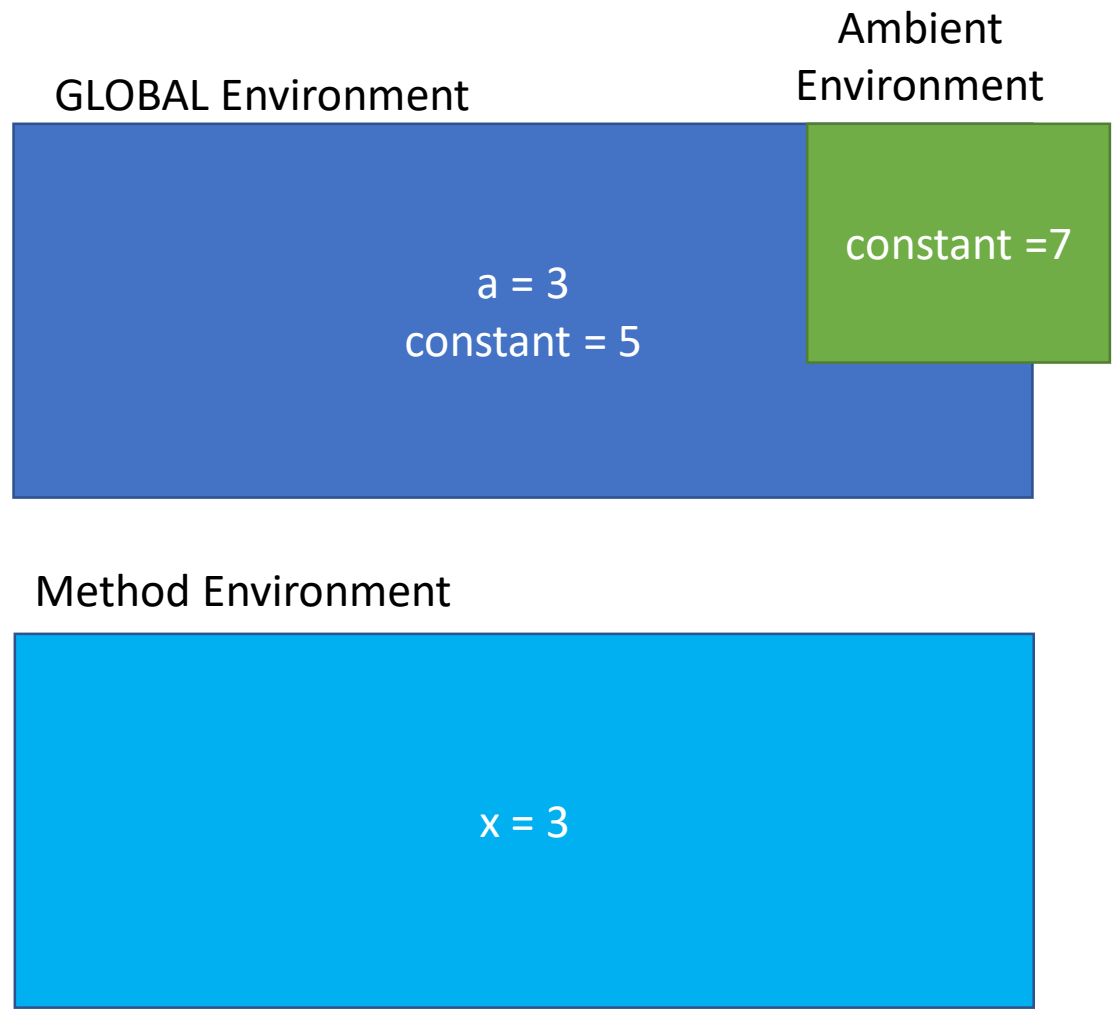
a = 3
constant = 5

constant = 7

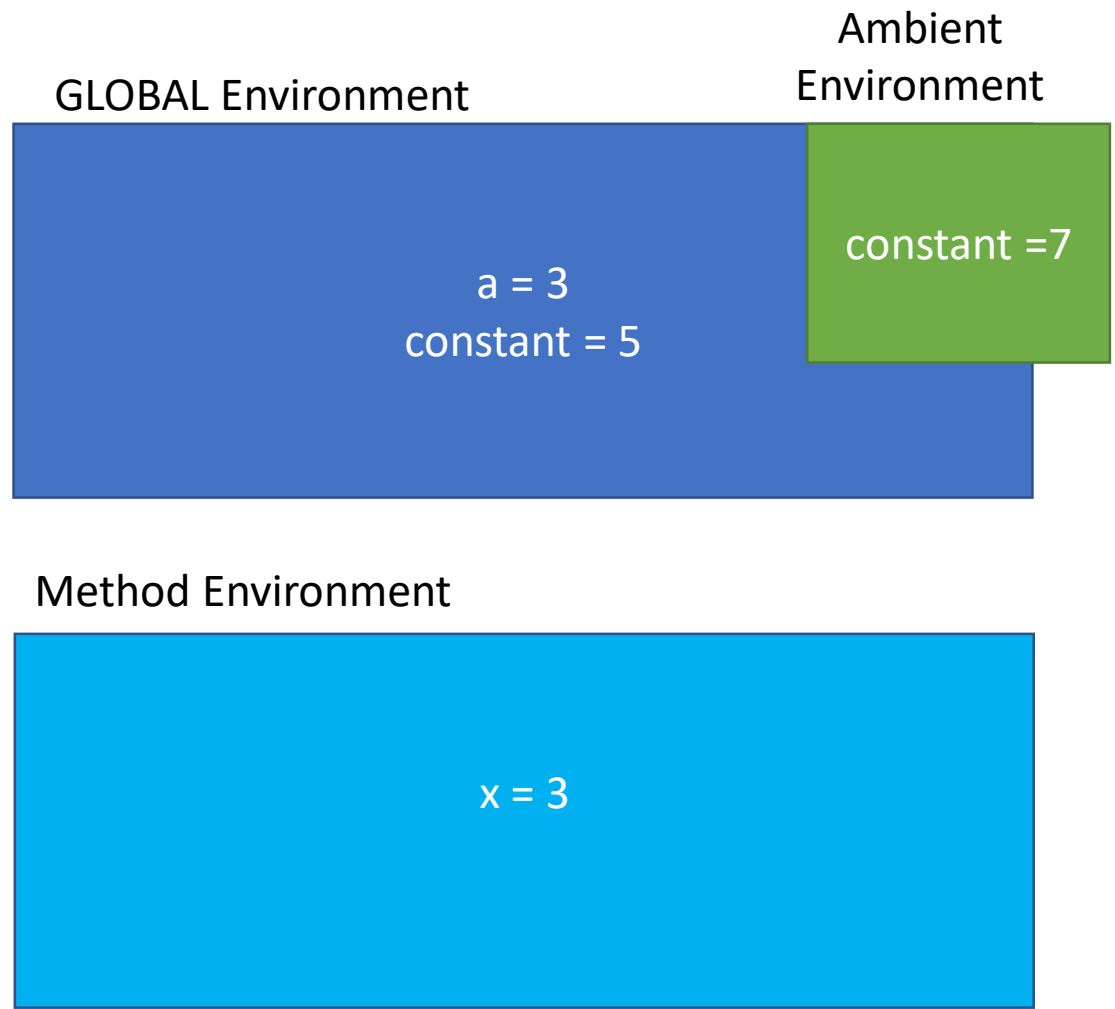
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1 (define a 3)
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5   (+ constant x)
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```




```
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4 (define (addConstant x)
5   (+ constant x)
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```

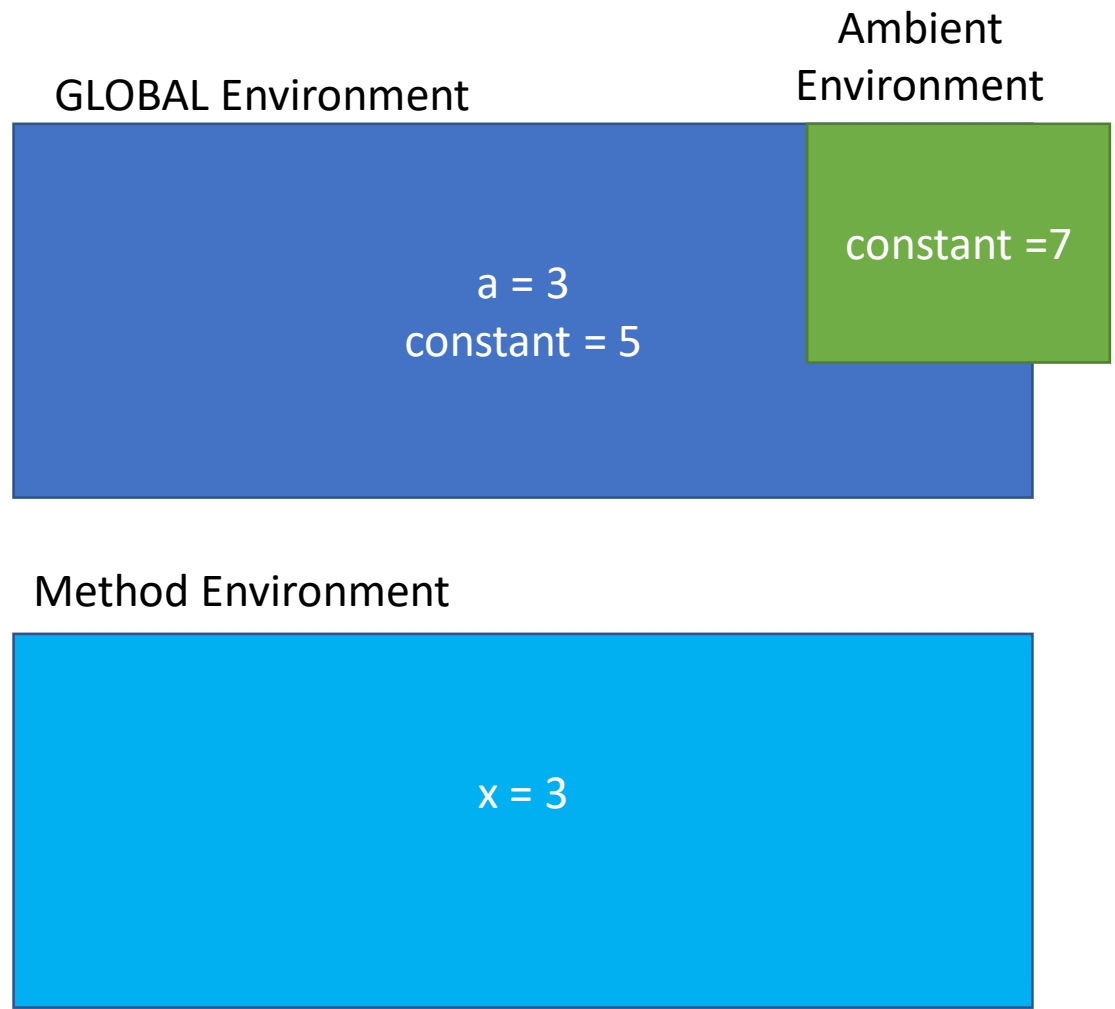
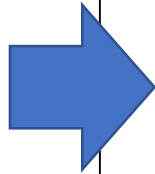


```
1 (define a 3)
2 (define constant 5)
3
4 (define (addConstant x)
5   (+ constant x)
6 )
7
8 (let ((constant 7))
9   (addConstant a))
```



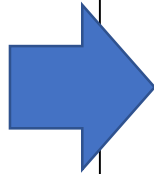
Where is constant?

```
1 (define a 3)
2 (define constant 5)
3
4 (define (addConstant x)
5   (+ constant x)
6 )
7
8 (let ((constant 7))
9   (addConstant a))
```

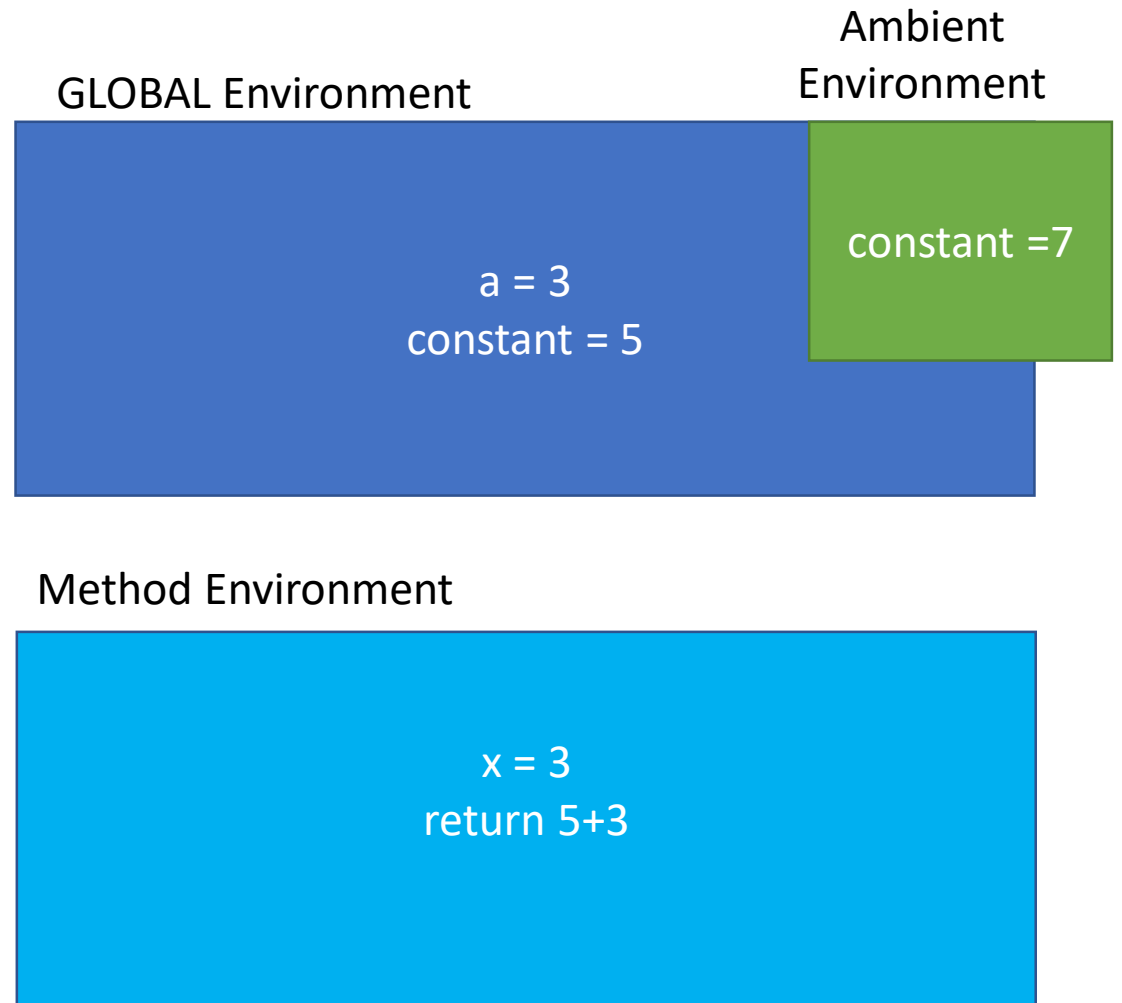


Where is constant?

First look in the method environment (missing)



```
1 (define a 3)
2 (define constant 5)
3
4 (define (addConstant x)
5   (+ constant x)
6 )
7
8 (let ((constant 7))
9   (addConstant a))
```

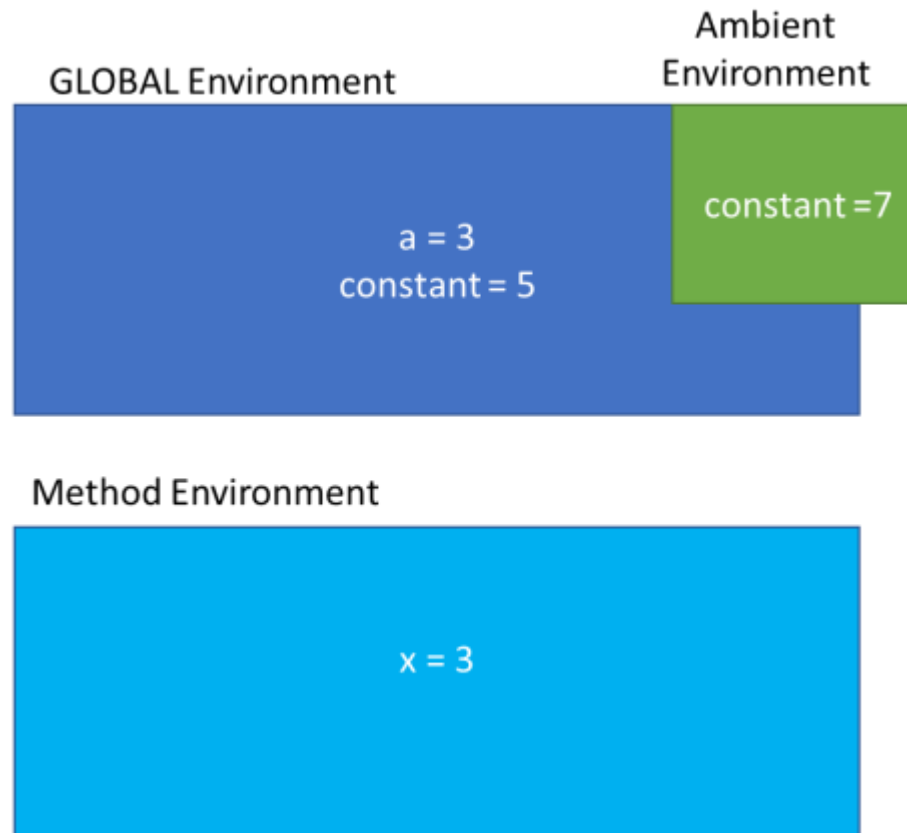


Where is constant?

Next look in the GLOBAL environment

Question: What happens if we don't make constant a global variable?

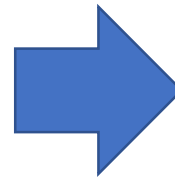
- Does that mean the interpreter will look in the “ambient environment” next?



Question: What happens if we don't make constant a global variable?

- Does that mean the interpreter will look in the “ambient” environment next?

```
1 (define a 3)
2 (define constant 5)
3
4 (define (addConstant x)
5   (+ constant x)
6 )
7
8 (let ((constant 7))
9   (addConstant a))
10
```

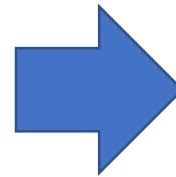


```
1 (define a 3)
2 ;(define constant 5)
3
4 (define (addConstant x)
5   (+ constant x)
6 )
7
8 (let ((constant 7))
9   (addConstant a))
```



Question: What happens if we don't make constant a global variable?

- Does that mean the interpreter will look in the “ambient” environment next?

```
1 (define a 3)
2 ;(define constant 5)
3
4 (define (addConstant x)
5   (+ constant x)
6 )
7
8 (let ((constant 7))
9   (addConstant a))
```



Welcome to [DrRacket](#), version 8.3 [cs].
Language: **R5RS**; memory limit: 128 MB.

  *constant: undefined;
cannot reference an identifier before its definition*

>

Answer: Nope!

How can ambient variables defined OUTSIDE of methods be used INSIDE of methods?

1. By passing them as METHOD PARAMETERS:

```
1 (define a 3)
2 (define constant 5)
3
4 (define (addConstant x constant)
5   (+ constant x)
6 )
7
8 (let ((constant 7))
9   (addConstant a constant))
```



```
Welcome to DrRacket, version 8.3 [cs].
Language: R5RS; memory limit: 128 MB.
10
>
```

OR what if we put the method definition inside the ambient environment?

```
1 (define a 3)
2 ;(define constant 5)
3
4 (define (addConstant x)
5   (+ constant x)
6 )
7
8 (let ((constant 7))
9   (addConstant a))
```



```
1 (define a 3)
2 (let ((constant 7))
3   (define (addConstant x) (+ constant x))
4   (addConstant a)
5 )
```

Result:

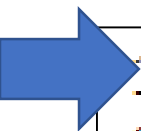
```
Welcome to DrRacket, version 8.3 [cs].
Language: R5RS; memory limit: 128 MB.
10
>
```

Why did that work?



```
1 | (define a 3)
2 | (let ((constant 7))
3 |   (define (addConstant x) (+ constant x))
4 |   (addConstant a)
5 | )
```

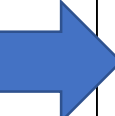
Let's draw a picture to understand...



```
1 | (define a 3)
2 | (let ((constant 7))
3 |     (define (addConstant x) (+ constant x))
4 |     (addConstant a)
5 | )
```

GLOBAL Environment

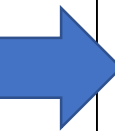
a = 3



```
1 | (define a 3)
2 | (let ((constant 7))
3 |     (define (addConstant x) (+ constant x))
4 |     (addConstant a)
5 | )
```

GLOBAL Environment

a = 3



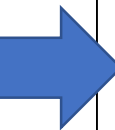
```
1 | (define a 3)
2 | (let ((constant 7))
3 |   (define (addConstant x) (+ constant x))
4 |   (addConstant a)
5 | )
```

GLOBAL Environment

a = 3

Ambient Environment

constant = 7



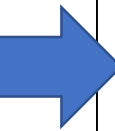
```
1 | (define a 3)
2 | (let ((constant 7))
3 |     (define (addConstant x) (+ constant x))
4 |     (addConstant a)
5 | )
```

GLOBAL Environment

a = 3

Ambient Environment

constant = 7



```
1 | (define a 3)
2 | (let ((constant 7))
3 |   (define (addConstant x) (+ constant x))
4 |   (addConstant a)
5 | )
```

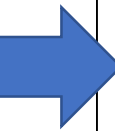
GLOBAL Environment

a = 3

Ambient Environment

constant = 7

Method Environment



```
1 | (define a 3)
2 | (let ((constant 7))
3 |   (define (addConstant x) (+ constant x))
4 |   (addConstant a)
5 | )
```

GLOBAL Environment

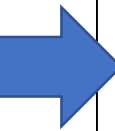
a = 3

Ambient Environment

constant = 7

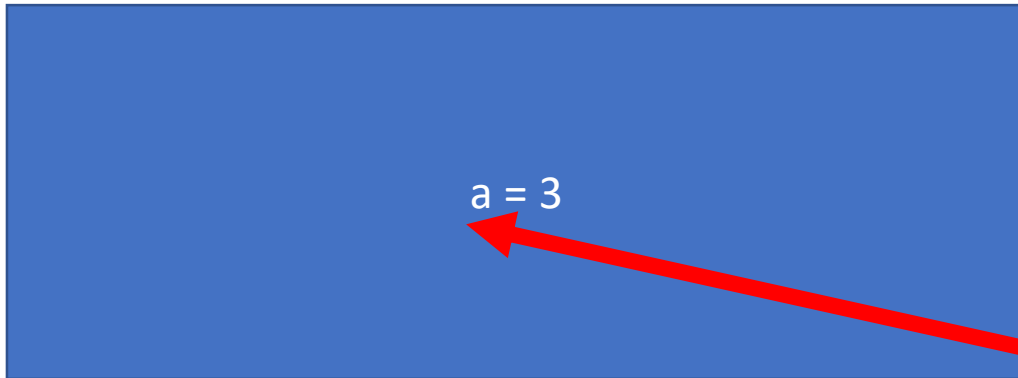
Method Environment

Return a + constant

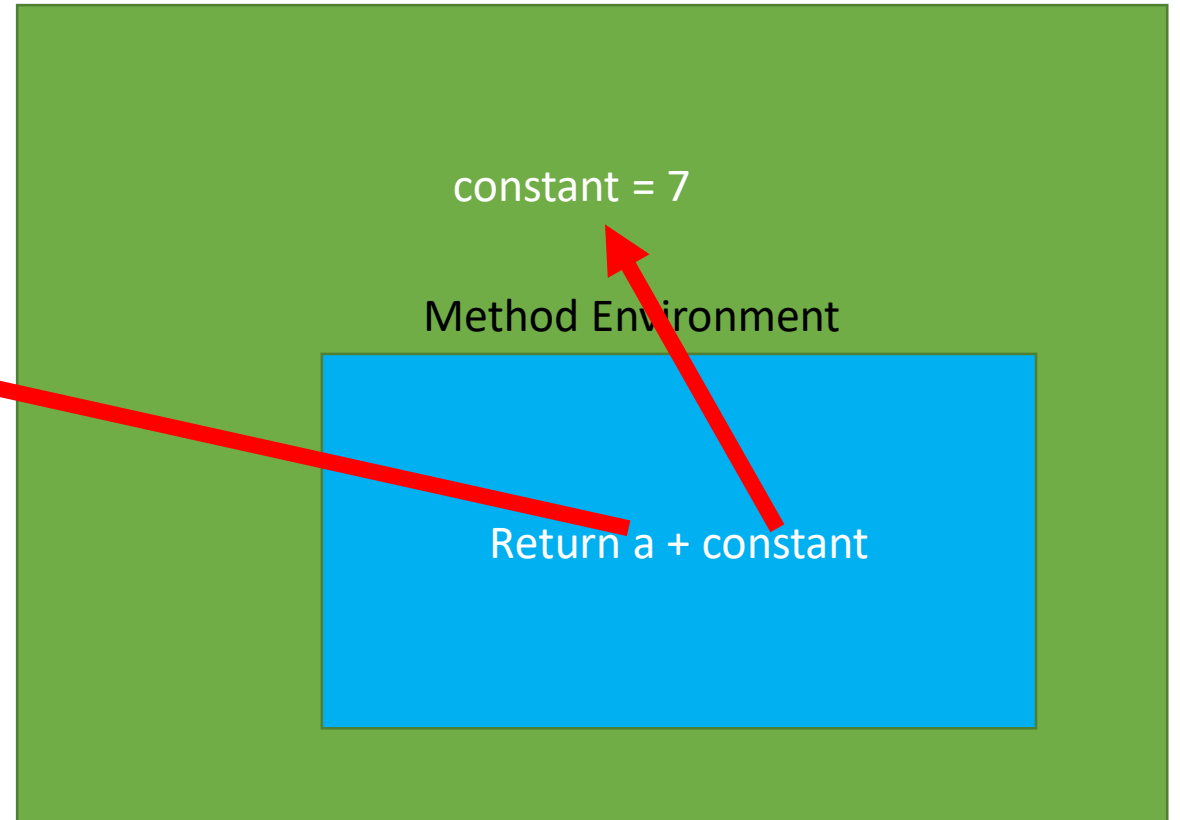


```
1 | (define a 3)
2 | (let ((constant 7))
3 |   (define (addConstant x) (+ constant x))
4 |   (addConstant a)
5 | )
```

GLOBAL Environment



Ambient Environment



Conclusion 1: How can methods access ambient variables?

1. Pass the ambient variables as method input parameters.
2. Put the method declaration INSIDE of the ambient environment.

Other important points:

- The “let” statement creates an ambient environment.
- Methods can always access variables from the global environment.

One last twist...

```
1 (define a 0)
2 (define constant 0)
3 (let ((constant 7)
4      (a 3))
5     (define (addConstant x) (+ constant x))
6     (addConstant a)
7 )
```

*"Ha ha you'll never
defeat me and my
memory questions!"*

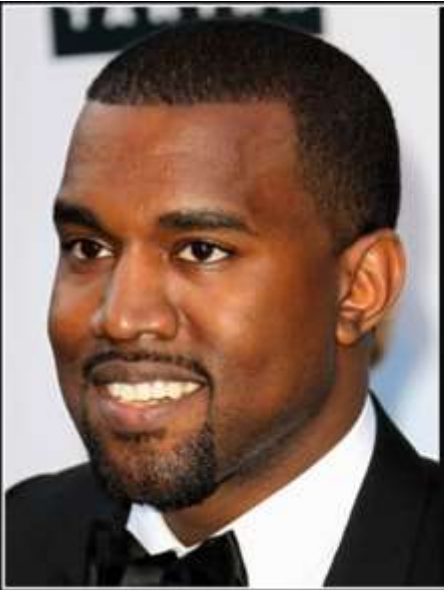


The Scheme Schemer

What should the output of this code be?

```
Welcome to DrRacket, version 8.3 [cs].
Language: R5RS; memory limit: 128 MB.
10
>
```

Conclusion 2: Why did Kanye West say this quote?

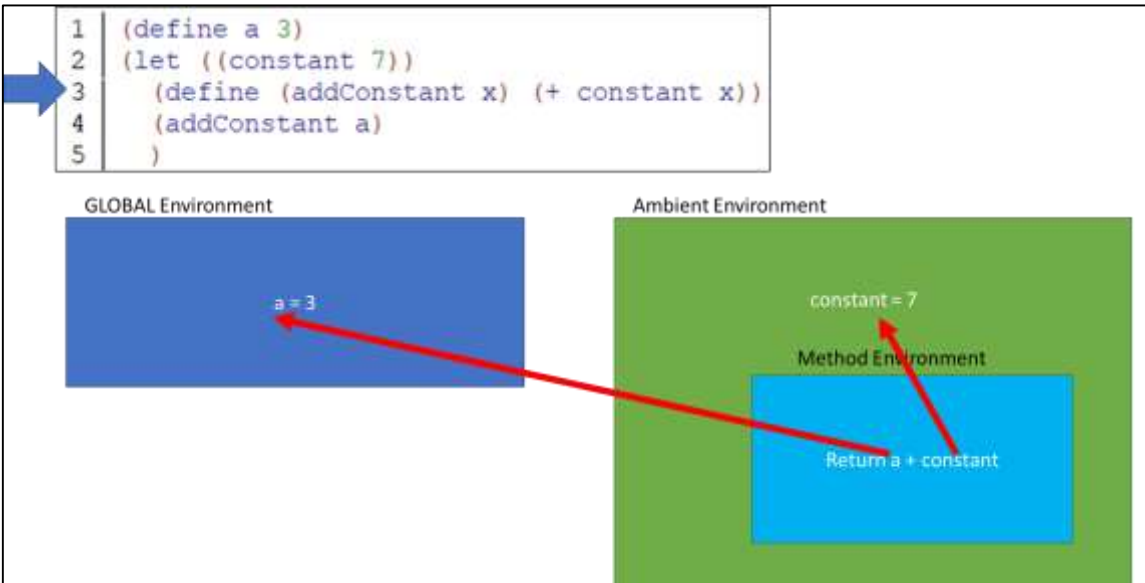


*"No one man should have all
that power"*

-Kanye West,

- Two possible answers:
 1. He was rapping.
 2. He was warning you to be very careful when using the `let` statement and variable scoping in Scheme.

Conclusion 3: WHERE Methods are defined MATTERS



- If the method is defined (notice here we say defined, not called) inside a let statement, it will use the ambient environment variables.
- If a method is called in the global environment it cannot use ambient variables unless they are passed as parameters to the method.

Figure Sources

- <https://external-preview.redd.it/fudpTkeL0ltbeklZ6lRujiTelPlye0823LKwAMkuyYY.jpg?auto=webp&s=4ac03fb1cc98a547d9b578938442686c117f0ca1>
- <https://uploads.dailydot.com/04c/00/a72e177e45dc0788.jpg?auto=compress%2Cformat&fit=scale&h=300&ixlib=php-3.3.0&w=600&wpsize=medium>
- <https://www.azquotes.com/picture-quotes/quote-no-one-man-should-have-all-that-power-kanye-west-85-46-66.jpg>
- https://i.kym-cdn.com/entries/icons/facebook/000/038/208/Anime_Girl_Punching_Wall_banner.jpg
- <https://datinginmiddleearth.files.wordpress.com/2015/01/reading.jpg>
- <https://external-preview.redd.it/fudpTkeL0ltbeklZ6lRujiTelPlye0823LKwAMkuyYY.jpg?auto=webp&s=4ac03fb1cc98a547d9b578938442686c117f0ca1>
- <https://freesvg.org/img/1317604469.png>
- https://static.wikia.nocookie.net/characterprofile/images/f/fe/Lex_Luthor_Gods_Among_Us_001.png/revision/latest?cb=20160104030401