## 复旦大学计算机科学技术学院 2021~2022 学年第二学期期末考试试卷 □A 卷 □B 卷

课	<b>怪名</b> 和	陈:_	算法	设计与分析	<u>r̃_</u> 课	程代码:_	_COMP1300	11. 03	<u></u>		
开课院系: 计算机科学技术学院 试形式: 开卷/闭卷/课程论文/											
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	Design and Analysis of Algorithms										
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June 15, 2022											
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<b>Notice:</b> This exam is closed book, no books, papers or recording devices permitted. You may use theorems from class, or the book provided you can recall them correctly. Add some annotation to your algorithm and											
pseudo code when essential.											
Pro				nentals (3%	,		(m)				
1.	Give the tight $\Theta$ -notation bound of recurrences $T(n) = aT\left(\frac{n}{8}\right) + \sqrt[3]{n}$ :						<u>·</u>				
2.	. Consider a modification to QUICKSORT, such that each time PARTITION is called, the median of the										
	partitioned array is found (using the SELECT algorithm) and used as a pivot. The worst-case running time						ng time				
	of this algorithm is										
3.	The	The tight upper bound and lower bound of log(n!) are and									
4.	Any comparison-based algorithm for constructing a binary search tree from an arbitrary list of $n$ elements										
					the worst case						
5.	How	man	y com	parisons oc	curred in Me	ergeSort? _		How	many comp	arisons in	

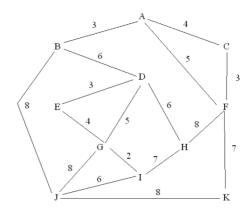
QuickSort? \_\_\_\_\_. (Assume the size of problem is n, express the answer in order of growth)

6.	For a binary search tree, if both <i>pre</i>	<i>order</i> tree v	walk and	post-order (	tree walk a	are known,	then the	binary
	search tree is determined.	True or f	false?					

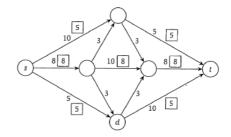
7.	What is the <b>black height</b> of the RB tree that result after successively inserting the keys 39; 32; 27, 19; 8
	into an initially empty red-black tree.

## **Problem 2**. Computation and Justification (44%)

- 1. Let G=(V, E) be a connected graph with *n* vertices, *m* edges, and positive edge costs that you may assume are all distinct. Let T=(V, E') be a spanning tree of G; we define the bottleneck edge of T to be the edge of T with the *greatest* cost. A spanning tree T of G is a *minimum bottleneck spanning tree* if there is no spanning tree T' of G with a cheaper bottleneck edge. Prove or give a counterexample for the following questions.
  - a) Is every minimum-bottleneck spanning tree of G a minimum spanning tree of G? (4%)
  - b) Is every minimum spanning tree of G a minimum bottleneck spanning tree of G? (4%)
- 2. A directed graph is **strongly connected** if and only if a DFS started from *any* vertex will visit every vertex in the graph without needing to be restarted. True or False? Justify your answer (6%)
- 3. Consider the following graph.



- 1) Give the minimum spanning tree of this graph. (4%)
- 2) Using Dijkstra's algorithm, determine the shortest path from node A to I. (4%)
- 4. (8%) The figure shows a flow network on which a flow is defined. The capacity of each edge appears as a label next to the edge and the numbers in boxes give the flow.
  - a) What is the value of this flow? Is this a maximum flow? If not, find the maximum flow. (4%)
  - b) Find a minimum cut in the flow network and say what its capacity is. (4%)



- 5. We have a connected graph G=(V, E), and a specific vertex u in V. Suppose we compute a *depth-first* search tree rooted at u, and obtain a tree T that includes all nodes of G. Suppose we then compute a *breadth-first* search tree rooted at u and obtain the same tree T. Prove that G=T. (8%)
- 6. A sequence of n operations is performed on a data structure. The ith operation costs  $\sqrt{i}$  if i is a perfect square, and 1 otherwise. Using amortized analysis (e.g., aggregate analysis, accounting method, or potential method) to determine the amortized cost per operation (6%)

**Problem 3:** Highways between cities are usually modeled as a directed graph G=(V, E) in which vertices represent cities and edges represent roads between cities. A truck company is planning new routes for shipments from city A (vertex s) to city B (vertex t). It is very costly when a shipment is delayed. For a given road  $e \in E$ , probability  $p(e) \in [0,1]$  is calculated to show that e will close without warning. Give an efficient algorithm for finding a route with the minimum probability of encountering a closed road. You should assume that all road closings are independent.

- 1) Present your algorithm in pseudo-code. (6%)
- 2) Prove and analyze your algorithm. (6%)

**Problem 4.** (12%) Your friends have studied stock patterns and defined something called **a rising trend**. A rising trend in a sequence of stock prices is a subsequence of prices  $p[i_1], p[i_2], ..., p[i_k]$  for days  $i_1 < i_2 < \cdots < i_k$ , so that  $i_1 = 1$ , and  $p[i_j] < p[i_{j+1}]$ , for j = 1, 2, ..., k-1. The problem is to find the longest rising trend in a given sequence of prices. E.g. n=7 and the sequence of prices is 10, 1, 2, 11, 3, 4, 12, then the longest rising trend is prices given on days 1, 4, 7. Give an efficient algorithm that takes a sequence of prices and returns the longest rising trend and its length. Your algorithm should be presented in pseudo-code.

**Problem 5.** (11%) Suppose in activity-selection problem, instead of always selecting the first activity to finish, we select the last activity to start that is compatible with all previously selected activities.

- 1) Show this strategy a greedy property and prove that it yields an optimal solution. (3%)
- 2) Present your algorithm with the selection strategy in pseudo-code. (5%)
- **3**) Analyze the efficiency of your algorithm. (3%)