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Listing 1: CODE HASHES

Listing 2: hash.sh

Listing 3: **GRAPHS**

Listing 4: Bridges and Cuts

```
//cat bridges_and_cuts.h / ./hash.sh
//3c13d9
#pragma once
//library checker tests: https://judge.yosupo.jp/problem/biconnected_components,

→ https://judge.yosupo.jp/problem/two_edge_connected_components

//with asserts checking correctness of is_bridge and is_cut
//O(n+m) time & space
//2 edge cc and bcc stuff doesn't depend on each other, so delete whatever is not needed
//handles multiple edges
//example initialization of 'adj':
//for (int i = 0; i < m; i++) {
// int u, v:
// cin >> u >> v:
// u--, v--;
// adj[u].emplace_back(v, i);
// adj[v].emplace_back(u, i);
struct info {
    //2 edge connected component stuff (e.g. components split by bridge edges)
         \hookrightarrow https://cp-algorithms.com/qraph/bridge-searching.html
    int num_2_edge_ccs;
    vector<bool> is_bridge;//edge id -> true iff bridge edge
    vector<int> two_edge_ccid; //node -> ID of 2-edge component (which are labeled 0, 1,
        \hookrightarrow ..., 'num_2_edge_ccs'-1)
    //bi-connected component stuff (e.g. components split by cut/articulation nodes)
         \hookrightarrow https://cp-algorithms.com/qraph/cutpoints.html
    int num_bccs;
    vector<bool> is_cut;//node -> true iff cut node
    vector<int> bcc_id;//edge id -> ID of BCC (which are labeled 0, 1, ..., 'num_bccs'-1)
info bridge_and_cut(const vector<vector<pair<int/*neighbor*/, int/*edge id*/>>>&

    → adj/*undirected graph*/, int m/*number of edges*/) {
    //stuff for both (always keep)
    int n = adj.size(), timer = 1;
    vector<int> tin(n, 0);
    //2 edge CC stuff (delete if not needed)
    int num_2_edge_ccs = 0;
```

```
vector<bool> is_bridge(m, false);
vector<int> two_edge_ccid(n), node_stack;
//BCC stuff (delete if not needed)
int num_bccs = 0;
vector<bool> is_cut(n, false);
vector<int> bcc_id(m), edge_stack;
auto dfs = [&](auto self, int v, int p_id) -> int {
    int low = tin[v] = timer++;
    int deg = 0;
    node_stack.push_back(v);
    for (auto [to, e_id] : adj[v]) {
        if (e_id == p_id) continue;
        if (!tin[to]) {
            edge_stack.push_back(e_id);
            int low_ch = self(self, to, e_id);
            if (low_ch >= tin[v]) {
                is_cut[v] = true;
                while (true) {
                    int edge = edge_stack.back();
                    edge_stack.pop_back();
                    bcc_id[edge] = num_bccs;
                    if (edge == e_id) break;
                }
                num_bccs++;
            }
            low = min(low, low_ch);
            deg++;
        } else if (tin[to] < tin[v]) {</pre>
            edge_stack.push_back(e_id);
            low = min(low, tin[to]);
    if (p_id == -1) is_cut[v] = (deg > 1);
    if (tin[v] == low) {
        if (p_id != -1) is_bridge[p_id] = true;
        while (true) {
            int node = node_stack.back();
            node_stack.pop_back();
            two_edge_ccid[node] = num_2_edge_ccs;
            if (node == v) break;
        num_2_edge_ccs++;
    }
    return low;
};
for (int i = 0; i < n; i++) {
    if (!tin[i])
        dfs(dfs, i, -1);
return {num_2_edge_ccs, is_bridge, two_edge_ccid, num_bccs, is_cut, bcc_id};
```

Listing 5: Block Vertex Tree

}

```
//cat block_vertex_tree.h / ./hash.sh
//567ff1
#pragma once
#include "bridges_and_cuts.h"
//library checker tests: https://judge.yosupo.jp/problem/biconnected_components
//(asserts checking correctness of commented-example-usage-loops)
```

```
//returns adjacency list of block vertex tree
//usage:
// info cc = bridge_and_cut(adj, m);
// vector<vector<int>> but = block_vertex_tree(adj, cc);
//to loop over each *unique* BCC containing a node i:
// for(int bccid : bvt[i]) {
        bccid -= n:
//
// }
//to loop over each *unique* node inside a BCC:
// for(int i : bvt[bccid + n]) {
//
// }
vector<vector<int>> block_vertex_tree(const vector<vector<pair<int, int>>>& adj, const
     \hookrightarrow info% cc) {
    int n = adj.size();
    vector<vector<int>> tree(n + cc.num_bccs);
    vector<int> cnt(cc.num_bccs, 0);
    for (int i = 0; i < n; i++) {</pre>
        for (auto [_, e_id] : adj[i]) {
            int bcc = cc.bcc_id[e_id];
            if (cnt[bcc]++ == 0) {
                tree[i].push_back(bcc + n); // Add edge between original node, and BCC
                tree[bcc + n].push_back(i);
        }
        for (auto [_, e_id] : adj[i])
            cnt[cc.bcc_id[e_id]]--;
    }
    return tree;
```

Listing 6: Centroid

```
//cat centroid.h / ./hash.sh
//4ba5e4
#pragma once
//library checker tests: https://judge.yosupo.jp/problem/frequency_table_of_tree_distance
//with asserts checking depth of tree <= log2(n)
//returns array 'par' where 'par[i]' = parent of node 'i' in centroid tree
//'par[root]' is -1
//0-based nodes
//0(n \log n)
//example usage:
// vector<int> parent = get_centroid_tree(adj);
// vector<vector<int>> childs(n);
// int root:
// for (int i = 0; i < n; i++) {
//
       if (parent[i] == -1)
           root = i:
//
       else
//
            childs[parent[i]].push_back(i);
11 }
vector<int> get_centroid_tree(const vector<vector<int>>& adj/*unrooted tree*/) {
   int n = adj.size();
   vector<int> sizes(n):
   vector<bool> vis(n, false);
   auto dfs_sz = [&] (auto self, int node, int par) -> void {
        sizes[node] = 1;
```

```
for (int to : adj[node]) {
        if (to != par && !vis[to]) {
            self(self, to, node);
            sizes[node] += sizes[to];
   }
}:
auto find_centroid = [&](int node) -> int {
    dfs_sz(dfs_sz, node, node);
    int size_cap = sizes[node] / 2, par = -1;
    while (true) {
        bool found = false;
        for (int to : adj[node]) {
            if (to != par && !vis[to] && sizes[to] > size_cap) {
                found = true;
                par = node;
                node = to;
                break;
        if (!found) return node;
   }
};
vector<int> parent(n);
auto dfs = [&](auto self, int node, int par) -> void {
    node = find_centroid(node);
    parent[node] = par;
    vis[node] = true;
    for (int to : adj[node]) {
        if (!vis[to])
            self(self, to, node);
   }
};
dfs(dfs, 0, -1);
return parent;
```

Listing 7: Dijkstra

```
//cat dijkstra.h | ./hash.sh
//6b6195
#pragma once
//library checker tests: https://judqe.yosupo.jp/problem/shortest_path
//returns\ array\ 'len'\ where\ 'len[i]' = shortest path from node 'startNode' to node i
//For example len[startNode] will always = 0
const long long inf = 1e18;
vector<long long> dijkstra(const vector<vector<pair<int, long long>>>& adj /*directed or

    undirected, weighted graph*/, int v) {
    vector<long long> len(adj.size(), inf);
    len[v] = 0;
    set<pair<long long/*weight*/, int/*node*/>> q;
    q.insert({OLL, v});
    while (!q.empty()) {
        auto it = q.begin();
        int node = it->second;
        q.erase(it);
        for (auto [to, weight] : adj[node])
            if (len[to] > weight + len[node]) {
                q.erase({len[to], to});
                len[to] = weight + len[node];
```

```
q.insert({len[to], to});
}
return len;
}
```

Listing 8: Floyd Warshall

```
//cat floyd_warshall.h / ./hash.sh
//84799a
#pragma once
//status: not tested
//**for directed graphs only** if you initialize len[i][i] to infinity, then
//afterward floyds, len[i][i] = length of shortest cycle including node 'i'
//another trick: change 'len' to 2d array of *bools* where len[i][j] = true if
//there exists an edge from i -> j in initial graph. Also do:
//'len[i][j] |= len[i][k] & len[k][j]'
//Then after floyds, len[i][j] = true iff there's exists some path from node
//'i' to node 'j'
//Changing the order of for-loops to i-j-k (instead of the current k-i-j)
//results in min-plus matrix multiplication. If adjacency matrix is M, then
//after computing M^k (with binary exponentiation), M[i][j] = min length path
//from i to j with at most k edges.
for (int k = 0; k < n; k++)
   for (int i = 0; i < n; i++)
       for (int j = 0; j < n; j++)
           len[i][j] = min(len[i][j], len[i][k] + len[k][j]);
```

Listing 9: HLD

```
//cat hld.h / ./hash.sh
//103dab
#pragma once
//library checker tests: https://judge.yosupo.jp/problem/lca,

→ https://judge.yosupo.jp/problem/vertex_add_path_sum,

    \hookrightarrow https://judge.yosupo.jp/problem/vertex_add_subtree_sum
//source: https://codeforces.com/blog/entry/53170
//assumes a single tree, 1-based nodes is possible by passing in 'root' in range [1, n]
struct HLD { //NOLINT(readability-identifier-naming)
   int n;
   vector<int> sub_sz, par, time_in, next;
   HLD(vector<vector<int>>& adj /*single unrooted tree*/, int root) :
       n(adj.size()), sub_sz(n, 1), par(n, root), time_in(n), next(n, root) {
       dfs1(root, adj);
       int timer = 0;
       dfs2(root, adj, timer);
   }
   void dfs1(int node, vector<vector<int>>& adj) {
       for (int& to : adj[node]) {
            if (to == par[node]) continue;
            par[to] = node;
            dfs1(to, adj);
            sub sz[node] += sub sz[to]:
            if (sub_sz[to] > sub_sz[adj[node][0]] || adj[node][0] == par[node])
                swap(to, adj[node][0]);
   }
```

```
void dfs2(int node, const vector<vector<int>>& adj, int& timer) {
        time in[node] = timer++:
        for (int to : adj[node]) {
            if (to == par[node]) continue;
            next[to] = (timer == time_in[node] + 1 ? next[node] : to);
            dfs2(to, adj, timer);
    }
    // Returns intervals (of time_in's) corresponding to the path between u and v, not
         \hookrightarrow necessarily in order
    // This can answer queries for "is some node 'x' on some path" by checking if the
         \hookrightarrow time_in[x] is in any of these intervals
    vector<pair<int, int>> path(int u, int v) const {
        vector<pair<int, int>> res;
        for (;; v = par[next[v]]) {
            if (time_in[v] < time_in[u]) swap(u, v);</pre>
            if (time_in[next[v]] <= time_in[u]) {</pre>
                 res.emplace_back(time_in[u], time_in[v]);
                 return res;
            res.emplace_back(time_in[next[v]], time_in[v]);
    }
    // Returns interval (of time_in's) corresponding to the subtree of node i
    // This can answer queries for "is some node 'x' in some other node's subtree" by
         \hookrightarrow checking if time_in[x] is in this interval
    pair<int, int> subtree(int i) const {
        return {time_in[i], time_in[i] + sub_sz[i] - 1};
    // Returns lca of nodes u and v
    int lca(int u, int v) const {
        for (;; v = par[next[v]]) {
            if (time_in[v] < time_in[u]) swap(u, v);</pre>
            if (time_in[next[v]] <= time_in[u]) return u;</pre>
    }
};
```

Listing 10: Hopcroft Karp

```
//cat hopcroft_karp.h / ./hash.sh
//e42443
#pragma once
//library checker tests: https://judge.yosupo.jp/problem/bipartitematching
//with asserts checking correctness of min vertex cover
//Modified from
     \hookrightarrow https://github.com/foreverbell/acm-icpc-cheat-sheet/blob/master/src/graph-algorithm/
//Worst case O(E*sqrt(V)) but faster in practice
struct match {
    //# of edges in matching (which = size of min vertex cover by öKnig's theorem)
    int size of matching:
    //an arbitrary max matching is found. For this matching:
    //if match_l[nodeLeft] == -1:
          'nodeLeft' is not in matching
    //else:
          the edge 'nodeLeft' <=> match l[nodeLeft] is in the matching
    //similarly for match_r with edge match_r[nodeRight] <=> nodeRight in matching if
         \hookrightarrow match_r[nodeRight] != -1
    //matchings stored in match_l and match_r are the same matching
```

```
//provides way to check if any node is in matching
    vector<int> match 1. match r:
    //an arbitrary min vertex cover is found. For this MVC: mvc_l['left node'] is true
        \hookrightarrow iff 'left node' is in the min vertex cover (same for muc r)
    //if mvc_l['left node'] is false, then 'left node' is in the corresponding maximal
        \hookrightarrow independent set
   vector<bool> mvc 1. mvc r:
//Think of the bipartite graph as having a left side (with size l_sz) and a right side
    \hookrightarrow (with size r sz).
//Nodes on left side are indexed 0,1,\ldots,l_sz-1
//Nodes on right side are indexed 0,1,\ldots,r_sz-1
//'adj' is like a directed adjacency list containing edges from left side -> right side:
//To initialize 'adj': For every edge nodeLeft <=> nodeRight, do:

    adj[nodeLeft].push_back(nodeRight)

match hopcroft_karp(const vector<vector<int>>& adj/*bipartite qraph*/, int r_sz/*number
    \hookrightarrow of nodes on right side*/) {
    int size_of_matching = 0, l_sz = adj.size();
   vector<int> match_l(l_sz, -1), match_r(r_sz, -1);
    while (true) {
       queue<int> q:
       vector<int> level(l_sz, -1);
       for (int i = 0; i < 1_sz; i++) {</pre>
            if (match_1[i] == -1) level[i] = 0, q.push(i);
        bool found = false;
       vector<bool> mvc_l(l_sz, true), mvc_r(r_sz, false);
        while (!q.empty()) {
           int u = q.front();
           q.pop();
            mvc_l[u] = false;
            for (int x : adj[u]) {
                mvc_r[x] = true;
                int v = match r[x]:
                found |= v == -1:
                if (v != -1 && level[v] < 0) {</pre>
                    level[v] = level[u] + 1:
                    q.push(v);
                }
       if (!found) return {size_of_matching, match_l, match_r, mvc_l, mvc_r};
        auto dfs = [&](auto self, int u) -> bool {
           for (int x : adj[u]) {
                int v = match r[x]:
                if (v == -1 || (level[u] + 1 == level[v] && self(self, v))) {
                    match_1[u] = x;
                    match_r[x] = u;
                    return true;
                }
            level[u] = 1e9; //acts as visited array
            return false:
       for (int i = 0; i < l_sz; i++)</pre>
            }
```

Listing 11: LCA

```
//cat lca.h / ./hash.sh
//90ab04
#pragma once
//library checker tests: https://judge.yosupo.jp/problem/lca
//https://codeforces.com/blog/entru/74847
//assumes a single tree, 1-based nodes is possible by passing in 'root' in range [1, n]
struct LCA { //NOLINT(readability-identifier-naming)
    int n:
    vector<int> jmp, jmp_edges, par, depth;
    vector<long long> dist;
    LCA(const vector<vector<pair<int, long long>>>& adj, int root) :
        n(adj.size()), jmp(n, root), jmp_edges(n, 1), par(n, root), depth(n, 0), dist(n,
        dfs(root, adj);
    void dfs(int node, const vector<vector<pair<int, long long>>>& adj) {
        for (auto [ch, w] : adj[node]) {
            if (ch == par[node]) continue;
            par[ch] = node:
            depth[ch] = 1 + depth[node];
            dist[ch] = w + dist[node];
            if (depth[node] > 0 && jmp_edges[node] == jmp_edges[jmp[node]])
                jmp[ch] = jmp[jmp[node]], jmp_edges[ch] = 2 * jmp_edges[node] + 1;
                imp[ch] = node;
            dfs(ch, adj);
        }
    //traverse up k edges in O(\log(k)). So with k=1 this returns 'node''s parent
    int kth_par(int node, int k) const {
        k = min(k, depth[node]);
        while (k > 0) {
            if (jmp_edges[node] <= k) {</pre>
                k -= jmp_edges[node];
                node = jmp[node];
            } else {
                k--;
                node = par[node];
        }
        return node;
    int get_lca(int x, int y) const {
        if (depth[x] < depth[y]) swap(x, y);</pre>
        x = kth_par(x, depth[x] - depth[y]);
        while (x != y) {
            if (jmp[x] != jmp[y])
                x = jmp[x], y = jmp[y];
                x = par[x], y = par[y];
        }
        return x;
    int dist_edges(int x, int y) const {
        return depth[x] + depth[y] - 2 * depth[get_lca(x, y)];
    long long dist_weight(int x, int y) const {
        return dist[x] + dist[y] - 2 * dist[get_lca(x, y)];
```

```
Listing 12: SCC
```

|};

```
//cat scc.h / ./hash.sh
//8fa337
#pragma once
//library checker tests: https://judge.yosupo.jp/problem/scc
struct scc info {
   int num_sccs;
    //scc's are labeled 0,1,..., 'numSCCs-1'
    //scc_id[i] is the id of the scc containing node 'i'
    //for each edge i \rightarrow j: scc_id[i] >= scc_id[j]
    vector<int> scc_id;
scc_info SCC(const vector<vector<int>>& adj /*directed, unweighted graph*/) {

→ //NOLINT(readability-identifier-naming)
    int n = adj.size(), timer = 1, num_sccs = 0;
    vector<int> tin(n, 0), scc_id(n, -1), node_stack;
    auto dfs = [&](auto self, int v) -> int {
        int low = tin[v] = timer++;
        node_stack.push_back(v);
        for (int to : adj[v]) {
            if (scc_id[to] < 0)</pre>
                low = min(low, tin[to] ? tin[to] : self(self, to));
        if (tin[v] == low) {
            while (true) {
                int node = node_stack.back();
                node_stack.pop_back();
                scc_id[node] = num_sccs;
                if (node == v) break;
           }
            num_sccs++;
        }
        return low;
   for (int i = 0; i < n; i++) {
        if (!tin[i])
            dfs(dfs, i):
   }
    return {num_sccs, scc_id};
```

Listing 13: RANGE DATA STRUCTURES

Listing 14: Segment Tree

```
//cat seg_tree.h | ./hash.sh
//28f96e
#pragma once
//stress tests: tests/stress_tests/range_data_structures/seg_tree.cpp
const long long inf = 1e18;
struct seg tree {
    struct node {
       long long sum, mx, mn;
       long long lazy;
```

```
int 1. r:
    int len() const {
        return r - 1 + 1;
    //returns 1 + (# of nodes in left child's subtree)
    //https://cp-algorithms.com/data_structures/segment_tree.html#memory-efficient-imple
    int rch() const { //right child
        return ((r - 1) & ~1) + 2;
   }
};
vector<node> tree;
//There's no constructor 'seqTree(int size)' because how to initialize l,r in nodes
     \hookrightarrow without calling build?
//the whole point of 'seqTree(int size)' was to be simpler by not calling build
seg_tree(const vector<long long>& arr) : tree(2 * (int)arr.size() - 1) {
    build(arr, 0, 0, (int)arr.size() - 1);
void build(const vector<long long>& arr, int v, int tl, int tr) {
    if (t1 == tr) {
        tree[v] = {
            arr[t1],
            arr[tl].
            arr[tl].
            tl,
        };
   } else {
        int tm = tl + (tr - tl) / 2;
        build(arr, v + 1, tl, tm);
        build(arr, v + 2 * (tm - tl + 1), tm + 1, tr);
        tree[v] = combine(tree[v + 1], tree[v + 2 * (tm - tl + 1)]);
    }
}
static node combine(const node& 1, const node& r) {
    return {
        1.sum + r.sum,
        max(1.mx. r.mx).
        min(1.mn, r.mn),
        1.1
        r.r
   };
//what happens when 'add' is applied to every index in range [tree[v].l, tree[v].r]?
void apply(int v, long long add) {
    tree[v].sum += tree[v].len() * add;
    tree[v].mx += add;
    tree[v].mn += add;
    if (tree[v].len() > 1) {
        tree[v + 1].lazy += add;
        tree[v + tree[v].rch()].lazy += add;
void push(int v) {
    if (tree[v].lazy) {
        apply(v, tree[v].lazy);
        tree[v].lazy = 0;
//update range [l,r] with 'add'
```

```
void update(int 1, int r, long long add) {
        update(0, 1, r, add);
    }
    void update(int v, int 1, int r, long long add) {
        if (tree[v].r < 1 || r < tree[v].1)</pre>
            return:
        if (1 <= tree[v].1 && tree[v].r <= r)</pre>
            return apply(v, add);
        update(v + 1, 1, r, add):
        update(v + tree[v].rch(), 1, r, add);
        tree[v] = combine(tree[v + 1], tree[v + tree[v].rch()]);
    }
    //range [l,r]
    node query(int 1, int r) {
        return query(0, 1, r);
   }
    node query(int v, int l, int r) {
        if (tree[v].r < 1 || r < tree[v].1)</pre>
            return {0, -inf, inf, 0, 0, 0};
        if (1 <= tree[v].1 && tree[v].r <= r)</pre>
            return tree[v];
        return combine(query(v + 1, 1, r),
                       query(v + tree[v].rch(), 1, r));
   }
};
```

Listing 15: Fenwick Tree

```
//cat bit.h | ./hash.sh
//516197
#pragma once
//library checker tests: https://judge.yosupo.jp/problem/point_add_range_sum,

→ https://judge.yosupo.jp/problem/vertex_add_path_sum,

→ https://judge.yosupo.jp/problem/vertex_add_subtree_sum,

    \hookrightarrow https://judge.yosupo.jp/problem/predecessor_problem
template<class T>
struct BIT { //NOLINT(readability-identifier-naming)
    vector<T> bit:
    BIT(int n) : bit(n, 0) {}
    BIT(const vector<T>& a) : bit(a.size()) {
        if (a.empty()) return;
        bit[0] = a[0];
        for (int i = 1; i < (int)a.size(); i++)</pre>
            bit[i] = bit[i - 1] + a[i];
        for (int i = (int)a.size() - 1; i > 0; i--) {
            int lower_i = (i & (i + 1)) - 1;
            if (lower_i >= 0)
                bit[i] -= bit[lower_i];
        }
   }
    void update(int idx, const T& d) {
        for (; idx < (int)bit.size(); idx = idx | (idx + 1))</pre>
            bit[idx] += d;
   }
   T sum(int r) const {
        T ret = 0;
        for (; r \ge 0; r = (r \& (r + 1)) - 1)
            ret += bit[r];
```

```
return ret:
    T sum(int 1, int r) const {
        return sum(r) - sum(l - 1);
    //Returns min pos such that sum of [0, pos] >= sum
    //Returns bit.size() if no sum is >= sum, or -1 if empty sum is.
    //Doesn't work with negatives (since it's greedy), counterexample: array: {1, -1},
         \hookrightarrow sum: 1, this returns 2, but should return 0
    int lower bound(T sum) const {
        if (sum <= 0) return -1;</pre>
        int pos = 0;
        for (int pw = 1 << (31 - __builtin_clz(bit.size() | 1)); pw; pw >>= 1) {
            if (pos + pw <= (int)bit.size() && bit[pos + pw - 1] < sum)</pre>
                pos += pw, sum -= bit[pos - 1];
        return pos;
    }
};
```

Listing 16: Sparse Table

```
//cat rmq.h | ./hash.sh
//dfa815
#pragma once
//library checker tests: https://judge.yosupo.jp/problem/staticrmg,

→ https://judge.yosupo.jp/problem/zalgorithm,

→ https://judge.yosupo.jp/problem/enumerate_palindromes,

     \hookrightarrow https://judge.yosupo.jp/problem/cartesian_tree
//usage:
// vector<long long> arr;
// sparseTable < long long > st(arr, [](auto x, auto y) { return min(x,y); });
//
//to also get index of min element, do:
// sparseTable < pair < long, int >> st(arr, [](auto x, auto y) { return min(x,y); });
//and initialize '.second''s to index. If there are multiple indexes of min element,
//it'll return the smallest (left-most) one
template <class T>
struct RMQ { //NOLINT(readability-identifier-naming)
    vector<vector<T>> dp;
    function<T(const T&, const T&)> func:
    RMQ(const vector<T>& arr, const function<T(const T&, const T&)>& a_func) : dp(1,
         \hookrightarrow arr), func(a_func) {
        int n = arr.size();
        for (int pw = 1, k = 1; pw * 2 <= n; pw *= 2, k++) {
            dp.emplace_back(n - pw * 2 + 1);
            for (int j = 0; j < (int)dp[k].size(); j++)</pre>
                dp[k][j] = func(dp[k - 1][j], dp[k - 1][j + pw]);
    }
    //inclusive range [l, r]
    T query(int 1, int r) const {
        int lg = 31 - \_builtin\_clz(r - 1 + 1);
        return func(dp[lg][l], dp[lg][r - (1 << lg) + 1]);</pre>
};
```

Listing 17: Implicit Lazy Segment Tree

//cat implicit_seg_tree.h | ./hash.sh

//f20253

#pragma once

```
//stress tests: tests/stress_tests/range_data_structures/implicit_seq_tree.cpp
//see TODO for lines of code which usually need to change (not a complete list)
const int sz = 1.5e7; //T0D0
struct node {
   long long val; //could represent max, sum, etc
    long long lazy;
    int lch. rch: // children. indexes into 'tree'. -1 for null
} tree[sz];
struct implicit_seg_tree {
    int ptr, root_1, root_r;//[root_1, root_r] defines range of root node; handles
         \hookrightarrow negatives
    implicit_seg_tree(int 1, int r) : ptr(0), root_1(1), root_r(r) {
        tree[ptr++] = {0, 0, -1, -1}; //TODO
   }
    static long long combine(long long val_1, long long val_r) {
        return val_l + val_r; //TODO
   }
    void apply(int v, int tl, int tr, long long add) {
        tree[v].val += (tr - tl + 1) * add; //TODO
        if (tl != tr) {
            tree[tree[v].lch].lazy += add; //TODO
            tree[tree[v].rch].lazy += add;
   }
    void push(int v, int tl, int tr) {
        if (tl != tr && tree[v].lch == -1) {
            assert(ptr + 1 < sz);</pre>
            tree[v].lch = ptr;
            tree[ptr++] = {0, 0, -1, -1}; //TODO
            tree[v].rch = ptr;
            tree[ptr++] = \{0, 0, -1, -1\};
        if (tree[v].lazy) {
            apply(v, tl, tr, tree[v].lazy);
            tree[v].lazy = 0;
    //update range [l,r] with 'add'
    void update(int 1, int r, long long add) {
        update(0, root_1, root_r, 1, r, add);
   }
    void update(int v, int tl, int tr, int l, int r, long long add) {
        push(v, tl, tr);
        if (tr < 1 || r < t1)
            return;
        if (1 <= t1 && tr <= r)
            return apply(v, tl, tr, add);
        int tm = tl + (tr - tl) / 2;
        update(tree[v].lch, tl, tm, l, r, add);
        update(tree[v].rch, tm + 1, tr, l, r, add);
        tree[v].val = combine(tree[tree[v].lch].val, tree[tree[v].rch].val);
   }
    //query range [l,r]
    //for more complicated query which doesn't allocate new nodes, see:
    //https://qithub.com/lrvideckis/Programming-Team-Code/blob/dc659297850440b65af2550a834kd42af8dctCffc&indirargo/rsnge_data_structures/implicitLazySegTree.h
    long long query(int 1, int r) {
        return query(0, root_1, root_r, 1, r);
   }
```

```
long long query(int v, int tl, int tr, int l, int r) {
        if (tr < 1 || r < t1)
            return 0; //TODO
        push(v, tl, tr);
        if (1 <= t1 && tr <= r)</pre>
            return tree[v].val;
        int tm = tl + (tr - tl) / 2:
        return combine(query(tree[v].lch, tl, tm, l, r),
                        query(tree[v].rch, tm + 1, tr, 1, r));
    }
};
```

Listing 18: Range Updates, Point Queries

```
//cat fenwick_inv.h | ./hash.sh
//e1114e
//library checker tests: https://judge.yosupo.jp/problem/vertex_add_subtree_sum,

→ https://judge.yosupo.jp/problem/point_add_range_sum
#include "../bit.h"
template<class T>
struct fenwick_inv {
    BIT<T> ft:
    fenwick inv(int n) : ft(n) {}
    fenwick_inv(const vector<T>& arr) : ft(init(arr)) {}
    BIT<T> init(vector<T> arr/*intentional pass by value*/) const {
        for (int i = (int)arr.size() - 1; i >= 1; i--)
            arr[i] -= arr[i - 1];
        return BIT<T>(arr);
    //add 'add' to inclusive range [l, r]
    void update(int 1, int r, const T& add) {
        ft.update(1, add);
        if (r + 1 < (int)ft.bit.size())</pre>
            ft.update(r + 1, -add);
    //get value at index 'idx'
    T query(int idx) const {
        return ft.sum(idx);
};
```

Listing 19: Kth Smallest

```
//cat kth_smallest.h | ./hash.sh
//8783cf
#pragma once
//library checker tests: https://judge.yosupo.jp/problem/range_kth_smallest
     \hookrightarrow https://cp-algorithms.com/data_structures/segment_tree.html#preserving-the-history-
struct kth smallest {
    struct node {
        int sum;
        int lch. rch://children. indexes into 'tree'
    };
    int mn, mx;
    deque<node> tree;
    kth_smallest(const vector<int>& arr) : mn(INT_MAX), mx(INT_MIN), roots(arr.size() +
         \hookrightarrow 1, 0) {
```

```
tree.push_back({0, 0, 0}); //acts as null
        for (int val : arr) mn = min(mn, val), mx = max(mx, val);
        for (int i = 0; i < (int)arr.size(); i++)</pre>
            roots[i + 1] = update(roots[i], mn, mx, arr[i]);
   }
    int update(int v, int tl, int tr, int idx) {
        if (tl == tr) {
            tree.push_back({tree[v].sum + 1, 0, 0});
            return tree.size() - 1;
        int tm = tl + (tr - tl) / 2;
        int lch = tree[v].lch;
        int rch = tree[v].rch;
        if (idx <= tm)</pre>
            lch = update(lch, tl, tm, idx);
            rch = update(rch, tm + 1, tr, idx);
        tree.push_back({tree[lch].sum + tree[rch].sum, lch, rch});
        return tree.size() - 1;
   }
    /* find (k+1)th smallest number among arr[l], arr[l+1], ..., arr[r]
     * k is 0-based, so query(l,r,0) returns the min
    int query(int 1, int r, int k) const {
        assert(0 \le k \&\& k \le r - 1 + 1); //note this condition implies L <= R
        assert(0 \le 1 \&\& r + 1 \le (int)roots.size());
        return query(roots[1], roots[r + 1], mn, mx, k);
   }
    int query(int vl, int vr, int tl, int tr, int k) const {
        if (t1 == tr)
            return tl;
        int tm = tl + (tr - tl) / 2;
        int left_count = tree[tree[vr].lch].sum - tree[tree[vl].lch].sum;
        if (left_count > k) return query(tree[v1].lch, tree[vr].lch, t1, tm, k);
        return query(tree[v1].rch, tree[vr].rch, tm + 1, tr, k - left_count);
   }
};
```

Listing 20: Number Distinct Elements

```
//cat distinct_query.h / ./hash.sh
//6bdf2f
#pragma once
//stress tests: tests/stress_tests/range_data_structures/distinct_query.cpp
//modified from

→ https://cp-algorithms.com/data_structures/segment_tree.html#preserving-the-history|-df-its-valuets-sources/segment-tree

https://cp-algorithms.com/data_structures/segment_tree.html#preserving-the-history|-df-its-valuets-sources/segment-tree

https://cp-algorithms.com/data_structures/segment_tree.html#preserving-the-history|-df-its-valuets-sources/segment_tree

https://cp-algorithms.com/data_structures/segment_tree

https://cp-algorithms.com/data_structures/segment_tr
 //works with negatives
//0(n \log n) time and space
struct distinct_query {
              struct node {
                              int lch, rch; //children, indexes into 'tree'
             };
              vector<int> roots:
              deque<node> tree;
              distinct_query(const vector<int>& arr) : roots(arr.size() + 1, 0) {
                              tree.push_back({0, 0, 0}); //acts as null
                             map<int, int> last_idx;
                             for (int i = 0; i < (int)arr.size(); i++) {</pre>
                                              roots[i + 1] = update(roots[i], 0, arr.size(), last_idx[arr[i]]);
```

```
last_idx[arr[i]] = i + 1;
        }
    }
    int update(int v, int tl, int tr, int idx) {
        if (t1 == tr) {
            tree.push_back({tree[v].sum + 1, 0, 0});
            return tree.size() - 1;
        int tm = (tl + tr) / 2;
        int lch = tree[v].lch:
        int rch = tree[v].rch;
        if (idx <= tm)
            lch = update(lch, tl, tm, idx);
            rch = update(rch, tm + 1, tr, idx);
        tree.push_back({tree[lch].sum + tree[rch].sum, lch, rch});
        return tree.size() - 1;
    //returns number of distinct elements in range [l,r]
    int query(int 1, int r) const {
        return query(roots[1], roots[r + 1], 0, (int)roots.size() - 1, 1 + 1);
    int query(int vl, int vr, int tl, int tr, int idx) const {
        if (tree[vr].sum == 0 || idx <= t1)</pre>
            return 0;
        if (tr < idx)
            return tree[vr].sum - tree[vl].sum;
        int tm = (tl + tr) / 2;
        return query(tree[v1].lch, tree[vr].lch, tl, tm, idx) +
               query(tree[v1].rch, tree[vr].rch, tm + 1, tr, idx);
    }
};
```

Listing 21: Buckets

```
//cat buckets.h | ./hash.sh
//6d8842
#pragma once
//stress tests: tests/stress_tests/range_data_structures/buckets.cpp
//this code isn't the best. It's meant as a rough start for sqrt-decomposition, and to
    \hookrightarrow be modified
//doesn't handle overflow
struct buckets {
    const int bucket_size = 50;//TODO: change - small value for testing
    struct node {
        int sum_bucket = 0;
        int 1, r; //inclusive range of bucket
        int len() const {
            return r - 1 + 1;
   };
    vector<int> values;
    vector<node> bucket:
    buckets(const vector<int>& initial) : values(initial) {
        int numbucket = ((int)values.size() + bucket_size - 1) / bucket_size;
        bucket.resize(numbucket):
        for (int i = 0; i < numbucket; i++) {</pre>
            bucket[i].sum lazv = 0:
            bucket[i].sum_bucket = 0;
```

```
bucket[i].l = i * bucket size:
        bucket[i].r = min((i + 1) * bucket_size, (int)values.size()) - 1;
        for (int j = bucket[i].1; j <= bucket[i].r; j++)</pre>
            bucket[i].sum_bucket += values[j];
}
void push(int b_idx) {
    node& b = bucket[b_idx];
    if (!b.sum_lazy) return;
    for (int i = b.1: i \le b.r: i++)
        values[i] += b.sum_lazy;
    b.sum_lazy = 0;
}
//update range [l,r]
void update(int 1, int r, int diff) {
    int start_bucket = 1 / bucket_size;
    int end_bucket = r / bucket_size;
    if (start_bucket == end_bucket) { //range contained in same bucket case
        for (int i = 1; i <= r; i++) {</pre>
            values[i] += diff;
            bucket[start_bucket].sum_bucket += diff;
        }
        return;
    for (int b_idx : {
                start_bucket, end_bucket
            }) { //handle "endpoint" buckets
        node& b = bucket[b_idx];
        for (int i = \max(b.1, 1); i \le \min(b.r, r); i ++) {
            values[i] += diff;
            b.sum_bucket += diff;
    }
    for (int i = start_bucket + 1; i < end_bucket; i++) { //handle all n/B buckets</pre>
         \hookrightarrow in middle
        node& b = bucket[i]:
        b.sum_lazy += diff;
        b.sum bucket += b.len() * diff:
    }
//sum of range [l.r]
int query(int 1, int r) {
    int start_bucket = 1 / bucket_size;
    int end_bucket = r / bucket_size;
    if (start_bucket == end_bucket) { //range contained in same bucket case
        push(start_bucket);
        int sum = 0;
        for (int i = 1; i <= r; i++)
            sum += values[i]:
        return sum;
    int sum = 0:
    for (int b_idx : {
                start bucket, end bucket
            }) { //handle "endpoint" buckets
        node& b = bucket[b_idx];
        push(b idx):
        for (int i = \max(b.1, 1); i \le \min(b.r, r); i++)
            sum += values[i];
    for (int i = start_bucket + 1; i < end_bucket; i++) //handle all n/B buckets in
```

```
Listing 22: Persistent Lazy Segment Tree
//cat persistent lazu seg tree.h | ./hash.sh
//5f187b
#pragma once
//status: not tested
struct persistent_lazy_seg_tree {
   struct node {
        int lch, rch;//children, indexes into 'tree'
        bool lazy_tog;
   };
   int sz;
   deque<node> tree;
   vector<int> roots;
   //implicit
   persistent_lazy_seg_tree(int a_sz) : sz(a_sz) {
        tree.push_back({0, 0, 0, 0}); //acts as null
        roots.push_back(0);
   void push(int v, int tl, int tr) {
        if (t1 != tr) {
            tree.push_back(tree[tree[v].lch]);
            tree[v].lch = tree.size() - 1;
            tree.push_back(tree[tree[v].rch]);
            tree[v].rch = tree.size() - 1;
        if (tree[v].lazy_tog) {
            tree[v].sum = (tr - tl + 1) - tree[v].sum;
           tree[v].lazy_tog = false;
            if (tl != tr) {
                tree[tree[v].lch].lazy_tog ^= 1;
                tree[tree[v].rch].lazy_tog ^= 1;
       }
   void set(int idx, int new_val) {
        tree.push_back(tree[roots.back()]);//allocate top down
        roots.push_back(tree.size() - 1);
        set(roots.back(), 0, sz - 1, idx, new_val);
   void set(int v, int tl, int tr, int idx, int new_val) {
        push(v, tl, tr);
        if (tr < idx || idx < tl)</pre>
            return:
        if (idx <= tl && tr <= idx) {</pre>
           tree[v].sum = new_val;
            return:
        int tm = (t1 + tr) / 2:
        int lch = tree[v].lch:
        int rch = tree[v].rch;
        set(lch, tl, tm, idx, new_val);
        set(rch, tm + 1, tr, idx, new_val);
```

```
tree[v].sum = tree[lch].sum + tree[rch].sum;
    }
    void toggle_range(int 1, int r) {
        tree.push_back(tree[roots.back()]);//allocate top down
        roots.push_back(tree.size() - 1);
        toggle_range(roots.back(), 0, sz - 1, 1, r);
    void toggle_range(int v, int tl, int tr, int l, int r) {
        push(v, tl, tr);
        if (tr < 1 || r < t1)
            return;
        int lch = tree[v].lch;
        int rch = tree[v].rch;
        if (1 <= t1 && tr <= r) {
            tree[v].sum = (tr - tl + 1) - tree[v].sum;
            if (tl != tr) {
                tree[lch].lazy_tog ^= 1;
                tree[rch].lazy_tog ^= 1;
            }
            return;
        int tm = (t1 + tr) / 2:
        toggle_range(lch, tl, tm, l, r);
        toggle_range(rch, tm + 1, tr, 1, r);
        tree[v].sum = tree[lch].sum + tree[rch].sum;
    //let's use implementation trick described here
         \hookrightarrow https://codeforces.com/blog/entry/72626
    //so that we don't have to propagate lazy vals and thus we don't have to allocate
        \hookrightarrow new nodes
    int query(int 1, int r) const {
        int version = roots.size() - 1;
        int root = roots[version];
        return query(root, 0, sz - 1, 1, r, tree[root].lazy_tog);
    }
    int query(int v, int tl, int tr, int l, int r, bool tog) const {
        if (v == 0 || tr < 1 || r < t1)
            return 0:
        if (1 <= t1 && tr <= r) {</pre>
            int sum = tree[v].sum;
            if (tree[v].lazv tog) sum = (tr - tl + 1) - sum;
            return sum;
        int tm = (tl + tr) / 2;
        tog ^= tree[v].lazy_tog;
        return query(tree[v].lch, tl, tm, l, r, tog) +
               query(tree[v].rch, tm + 1, tr, 1, r, tog);
   }
};
```

Listing 23: Merge Sort Tree

```
//cat merge_sort_tree.h | ./hash.sh
//55d2b5
#pragma once
//library checker tests: https://judge.yosupo.jp/problem/static_range_frequency,

→ https://judge.yosupo.jp/problem/range_kth_smallest
//For point updates: either switch to policy based BST, or use sqrt decomposition
struct merge_sort_tree {
   struct node {
```

```
vector<int> vals:
        //returns 1 + (# of nodes in left child's subtree)
        //https://cp-algorithms.com/data_structures/segment_tree.html#memory-efficient-impl
        int rch() const {
            return ((r - 1) & ~1) + 2;
    };
    vector<node> tree;
    //RTE's when 'arr' is emptu
    merge_sort_tree(const vector<int>& arr) : tree(2 * (int)arr.size() - 1) {
        build(arr, 0, 0, (int)arr.size() - 1);
    void build(const vector<int>& arr, int v, int tl, int tr) {
        if (t1 == tr) {
            tree[v] = {
                {arr[t1]},
                tl.
            };
        } else {
            int tm = tl + (tr - tl) / 2:
            build(arr. v + 1. tl. tm):
            build(arr, v + 2 * (tm - tl + 1), tm + 1, tr);
            tree[v] = combine(tree[v + 1], tree[v + 2 * (tm - tl + 1)]);
    }
    node combine(const node& 1, const node& r) const {
        vector<int> par(l.vals.size() + r.vals.size());
        merge(1.vals.begin(), 1.vals.end(), r.vals.begin(), r.vals.end(), par.begin());
        return {par, 1.1, r.r};
    //How many of arr[l], arr[l+1], ..., arr[r] are < x?
    //0(log^2(n))
    int query(int 1, int r, int x) const {
        return query(0, 1, r, x);
    int query(int v, int 1, int r, int x) const {
        if (tree[v].r < 1 || r < tree[v].1)</pre>
            return 0;
        if (1 <= tree[v].1 && tree[v].r <= r) {</pre>
            const vector<int>& vals = tree[v].vals;
            return lower_bound(vals.begin(), vals.end(), x) - vals.begin();
        return query(v + 1, 1, r, x) +
               query(v + tree[v].rch(), 1, r, x);
    }
};
```

Listing 24: STRINGS

Listing 25: Suffix Array

```
//cat suffix_array.h / ./hash.sh
//46840a
//library checker tests: https://judge.yosupo.jp/problem/suffixarray,

→ https://judge.yosupo.jp/problem/zalgorithm,
```

```
→ https://judge.yosupo.jp/problem/number_of_substrings,

→ https://judge.yosupo.jp/problem/enumerate_palindromes

//modified from here: https://judge.yosupo.jp/submission/37410
// SA-IS, linear-time suffix array construction
// Reference:
// G. Nong, S. Zhang, and W. H. Chan,
// Two Efficient Algorithms for Linear Time Suffix Array Construction
template<class T>
vector<int> sa is(const T& s. int upper/*max element of 's': for std::string. pass in
    \hookrightarrow 255*/) {
    int n = (int)s.size();
    if (n == 0) return {};
    if (n == 1) return {0};
   if (n == 2) {
        if (s[0] < s[1]) {
            return {0, 1};
        } else {
            return {1, 0};
    vector<int> sa(n);
    vector<bool> ls(n):
    for (int i = n - 2; i >= 0; i--)
        ls[i] = (s[i] == s[i + 1]) ? ls[i + 1] : (s[i] < s[i + 1]);
    vector<int> sum_l(upper + 1), sum_s(upper + 1);
    for (int i = 0; i < n; i++) {
        if (!ls[i])
            sum_s[s[i]]++;
        else
            sum_l[s[i] + 1]++;
   }
   for (int i = 0; i <= upper; i++) {</pre>
        sum_s[i] += sum_l[i];
        if (i < upper) sum_l[i + 1] += sum_s[i];</pre>
    vector<int> buf(upper + 1);
    auto induce = [&](const vector<int>& lms) {
        fill(sa.begin(), sa.end(), -1);
        fill(buf.begin(), buf.end(), 0);
        copy(sum_s.begin(), sum_s.end(), buf.begin());
        for (auto d : lms) {
            if (d == n) continue;
            sa[buf[s[d]]++] = d;
        copy(sum_l.begin(), sum_l.end(), buf.begin());
        sa[buf[s[n-1]]++] = n-1;
        for (int i = 0; i < n; i++) {</pre>
            int v = sa[i]:
            if (v >= 1 && !ls[v - 1])
                sa[buf[s[v - 1]] ++] = v - 1;
        copy(sum_l.begin(), sum_l.end(), buf.begin());
        for (int i = n - 1; i \ge 0; i--) {
            int v = sa[i];
            if (v >= 1 \&\& ls[v - 1])
                sa[--buf[s[v-1] + 1]] = v - 1:
        }
    vector < int > lms_map(n + 1, -1);
    int m = 0;
```

```
for (int i = 1: i < n: i++) {
    if (!ls[i - 1] && ls[i])
        lms_map[i] = m++;
vector<int> lms;
lms.reserve(m);
for (int i = 1: i < n: i++) {
    if (!ls[i - 1] && ls[i])
        lms.push_back(i);
induce(lms);
if (m) {
    vector<int> sorted_lms;
    sorted_lms.reserve(m);
    for (int v : sa) {
        if (lms_map[v] != -1) sorted_lms.push_back(v);
    vector<int> rec_s(m);
    int rec_upper = 0;
    rec_s[lms_map[sorted_lms[0]]] = 0;
    for (int i = 1; i < m; i++) {</pre>
        int l = sorted_lms[i - 1], r = sorted_lms[i];
        int end_1 = (lms_map[1] + 1 < m) ? lms[lms_map[1] + 1] : n;</pre>
        int end_r = (lms_map[r] + 1 < m) ? lms[lms_map[r] + 1] : n;
        bool same = true;
        if (end_l - l != end_r - r)
            same = false;
        else {
            while (1 < end_1) {</pre>
                if (s[l] != s[r])
                    break;
                1++;
            if (1 == n \mid | s[1] != s[r]) same = false:
        if (!same) rec_upper++;
        rec_s[lms_map[sorted_lms[i]]] = rec_upper;
    }
    auto rec_sa =
        sa_is(rec_s, rec_upper);
    for (int i = 0; i < m; i++)
        sorted_lms[i] = lms[rec_sa[i]];
    induce(sorted_lms);
}
return sa;
```

Listing 26: Longest Common Prefix Array

```
// Linear-Time Longest-Common-Prefix Computation in Suffix Arrays and Its
// Applications
template<class T>
vector<int> LCP(const T& s, const vector<int>& sa) {
    \hookrightarrow //NOLINT(readability-identifier-naming)
    int n = s.size(), k = 0;
   vector<int> lcp(n, 0);
   vector<int> rank(n, 0);
   for (int i = 0; i < n; i++) rank[sa[i]] = i;
   for (int i = 0: i < n: i++, k ? k-- : 0) {
       if (rank[i] == n - 1) {
            k = 0:
            continue;
       int j = sa[rank[i] + 1];
        while (i + k < n \&\& j + k < n \&\& s[i + k] == s[j + k]) k++;
       lcp[rank[i]] = k;
   }
   return lcp;
```

Listing 27: Prefix Function

```
//cat prefix_function.h | ./hash.sh
//aa0518
#pragma once
//library checker tests: https://judge.yosupo.jp/problem/zalgorithm
//stress tests: tests/stress_tests/strings/kmp.cpp
//source: https://cp-algorithms.com/string/prefix-function.html#implementation
template <class T>
vector<int> prefix_function(const T& s) {
   int n = s.size();
   vector<int> pi(n, 0);
   for (int i = 1; i < n; i++) {
       int j = pi[i - 1];
       while (j > 0 \&\& s[i] != s[j]) j = pi[j - 1];
       pi[i] = j + (s[i] == s[j]);
   }
   return pi;
```

Listing 28: KMP

```
//cat kmp.h | ./hash.sh
//9d70ad
#pragma once
//stress tests: tests/stress_tests/strings/kmp.cpp
#include "prefix_function.h"
//usage:
// string needle;
// ...
// KMP_Match kmp(needle);
//or
// vector<int> needle;
// ...
// KMP_Match kmp(needle);
// wethor<int> needle;
// ...
// KMP_outh kmp(needle);
// wethor<int> needle;
// ...
// KMP_outh kmp(needle);
// kmp-doubling-trick: to check if 2 arrays are rotationally equivalent: run kmp
//with one array as the needle and the other array doubled as the haystack
//or just use kactl's min rotation code
```

```
template <class T>
struct KMP { //NOLINT(readability-identifier-naming)
    KMP(const T& a_needle) : pi(prefix_function(a_needle)), needle(a_needle) {}
    // if haustack = "bananas"
    // needle = "ana"
    // then we find 2 matches:
    // bananas
    // _ana___
    // ana
    // 0123456 (indexes)
    // and KMP_Match::find returns {1,3} - the indexes in haystack where
    // each match starts.
    // You can also pass in false for "all" and KMP_Match::find will only
    // return the first match: {1}. Useful for checking if there exists
    // some match:
    //
    // KMP_Match::find(<haystack>,false).size() > 0
    vector<int> find(const T& haystack, bool all = true) const {
        vector<int> matches;
        for (int i = 0, j = 0; i < (int)haystack.size(); i++) {</pre>
            while (j > 0 && needle[j] != haystack[i]) j = pi[j - 1];
            if (needle[j] == haystack[i]) j++;
            if (j == (int)needle.size()) {
                matches.push_back(i - (int)needle.size() + 1);
                if (!all) return matches;
                j = pi[j - 1];
        }
        return matches;
    vector<int> pi;//prefix function
    T needle:
};
```

Listing 29: Trie

```
//cat trie.h | ./hash.sh
//57ce2c
#pragma once
//status: not tested
//source: https://cp-algorithms.com/string/aho_corasick.html#construction-of-the-trie
//intended to be a base template and to be modified
const int k = 26;//character size
struct trie {
   const char min_ch = 'a';//'A' for uppercase, '0' for digits
   struct node {
        int next[k], id, p = -1;
       char ch;
        bool leaf = 0:
        node(int a_p = -1, char a_ch = '#') : p(a_p), ch(a_ch) {
            fill(next, next + k, -1);
       }
   };
   vector<node> t:
   trie(): t(1) {}
   void add_string(const string& s, int id) {
        int c = 0:
       for (char ch : s) {
```

```
int v = ch - min ch:
            if (t[c].next[v] == -1) {
                t[c].next[v] = t.size();
                t.emplace_back(c, ch);
            c = t[c].next[v];
        t[c].leaf = 1;
        t[c].id = id;
   }
    void remove_string(const string& s) {
        int c = 0:
        for (char ch : s) {
            int v = ch - min_ch;
            if (t[c].next[v] == -1)
                return;
            c = t[c].next[v];
        }
        t[c].leaf = 0;
   }
    int find_string(const string& s) {
        int c = 0:
        for (char ch : s) {
            int v = ch - min_ch;
            if (t[c].next[v] == -1)
                return -1;
            c = t[c].next[v];
        if (!t[c].leaf) return -1;
        return t[c].id;
   }
};
```

Listing 30: Binary Trie

```
//cat binary_trie.h | ./hash.sh
//d13f33
#pragma once
//library checker tests: https://judge.yosupo.jp/problem/set_xor_min
struct binary_trie {
   const int mx_bit = 62;
   struct node {
       long long val = -1;
       int sub_sz = 0;//number of inserted values in subtree
       int next[2] = \{-1, -1\};
   };
   vector<node> t;
   binary_trie() : t(1) {}
    //delta = 1 to insert val, -1 to remove val
   void update(long long val, int delta) {
       int c = 0:
       t[0].sub_sz += delta;
       for (int bit = mx_bit; bit >= 0; bit--) {
            bool v = (val >> bit) & 1:
            if (t[c].next[v] == -1) {
                t[c].next[v] = t.size();
                t.emplace_back();
            c = t[c].next[v];
            t[c].sub_sz += delta;
```

```
t[c].val = val:
   }
    //returns # of val's in this data structure
    int cnt(long long val) const {
        int c = 0;
        for (int bit = mx bit; bit >= 0; bit--) {
            bool v = (val >> bit) & 1;
            if (t[c].next[v] == -1)
                return 0:
            c = t[c].next[v];
        return t[c].sub_sz;
    int size() const {
        return t[0].sub_sz;
    //returns x such that:
    // x is in this data structure
    // value of (x ^val) is minimum
    long long min_xor(long long val) const {
        assert(size() > 0);
        int c = 0:
        for (int bit = mx_bit; bit >= 0; bit--) {
            bool v = (val >> bit) & 1;
            int ch = t[c].next[v];
            if (ch != -1 && t[ch].sub_sz > 0)
                c = ch;
            else
                c = t[c].next[!v];
        return t[c].val;
};
```

Listing 31: Longest Common Prefix Query

```
//cat lcp_queries.h / ./hash.sh
//a4013c
#pragma once
//library checker tests: https://judge.yosupo.jp/problem/zalqorithm,
     \hookrightarrow https://judge.yosupo.jp/problem/enumerate_palindromes
#include "suffix_array.h"
#include "lcp.h"
#include "../range_data_structures/rmq.h"
//computes suffix array, lcp array, and then sparse table over lcp array
//O(n \log n)
struct lcp_queries {
    lcp_queries(const string& s) : sa(sa_is(s, 255)), inv_sa(s.size()), lcp(LCP(s, sa)),
         \hookrightarrow st(lcp, [](int x, int y) {
        return min(x, y);
    }) {
        for (int i = 0; i < (int)s.size(); i++)</pre>
            inv_sa[sa[i]] = i;
    //length of longest common prefix of suffixes s[idx1 \dots n-1], s[idx2 \dots n-1],
         \hookrightarrow 0-based indexing
    //You can check if two substrings s[L1..R1], s[L2..R2] are equal in O(1) by:
    //R2-L2 == R1-L1 \&\& longest_common_prefix(L1, L2) >= R2-L2+1
```

//cat exp_mod.h / ./hash.sh

```
int longest_common_prefix(int idx1, int idx2) const {
    if (idx1 == idx2) return (int)sa.size() - idx1;
    idx1 = inv_sa[idx1];
    idx2 = inv_sa[idx2];
    if (idx1 > idx2) swap(idx1, idx2);
    return st.query(idx1, idx2 - 1);
}

//returns true if suffix s[idx1 ... n-1] < s[idx2 ... n-1]

//(so false if idx1 == idx2)

bool less(int idx1, int idx2) const {
    return inv_sa[idx1] < inv_sa[idx2];
}

vector<int> sa, inv_sa, lcp;
RMQ<int> st;
};
```

Listing 32: MATH

Listing 33: BIN EXP MOD

```
//f6e76e
#pragma once
//library checker tests: https://judge.yosupo.jp/problem/system_of_linear_equations,

→ https://judge.yosupo.jp/problem/binomial_coefficient,

→ https://judge.wosupo.jp/problem/matrix det.

    \hookrightarrow https://judge.yosupo.jp/problem/inverse_matrix
//stress tests: tests/stress_tests/math/exp_mod.cpp
//returns (base^pw)/mod in O(log(pw)), but returns 1 for 0^0
//What if base doesn't fit in long long?
//Since (base^pw) // mod == ((base/mod)^pw) // mod we can calculate base under mod of 'mod'
//What if pw doesn't fit in long long?
//case 1: mod is prime
//(base^pw)%mod == (base^(pw%(mod-1)))%mod (from Fermat's little theorem)
//so calculate pw under mod of 'mod-1'
//case 2: non-prime mod
//let t = totient(mod)
//if pw >= log2(mod) then (base^pw)%mod == (base^(t+(pw%t)))%mod (proof)
    \hookrightarrow https://cp-algorithms.com/algebra/phi-function.html#generalization)
//so calculate pw under mod of 't'
//incidentally, totient(p) = p - 1 for every prime p, making this a more generalized
    \hookrightarrow version of case 1
int pow(long long base, long long pw, int mod) {
    assert(0 <= pw && 0 <= base && 1 <= mod);</pre>
    int res = 1;
   base %= mod;
    while (pw > 0) {
        if (pw & 1) res = res * base % mod;
        base = base * base % mod;
        pw >>= 1;
   }
    return res;
```

Listing 34: Fibonacci

Listing 35: Matrix Mult and Pow

```
//cat matrix_expo.h | ./hash.sh
//424a12
#pragma once
//library checker tests: https://judge.yosupo.jp/problem/matrix_product
//stress tests: tests/stress_tests/math/fib_matrix_expo.cpp
//empty matrix -> RTE
vector<vector<int>> mult(const vector<vector<int>>& a, const vector<vector<int>>& b, int
    \hookrightarrow mod) {
    assert(a[0].size() == b.size());
    int n = a.size(), m = b[0].size(), inner = b.size();
    vector<vector<int>> prod(n, vector<int>(m, 0));
    for (int i = 0; i < n; i++) {
        for (int k = 0; k < inner; k++) {
            for (int j = 0; j < m; j++)
                prod[i][j] = (prod[i][j] + 1LL * a[i][k] * b[k][j]) % mod;
    return prod;
vector<vector<int>> power(vector<int>> mat/*intentional pass by value*/, long
    \hookrightarrow long pw, int mod) {
    int n = mat.size();
    vector<vector<int>> prod(n, vector<int>(n, 0));
    for (int i = 0; i < n; i++)</pre>
        prod[i][i] = 1;
    while (pw > 0) {
        if (pw % 2 == 1) prod = mult(prod, mat, mod);
        mat = mult(mat, mat, mod);
        pw /= 2:
   }
   return prod;
```

Listing 36: N Choose K MOD

```
//cat n_choose_k_mod.h | ./hash.sh
//2c3f33
#pragma once
//library checker tests: https://judge.yosupo.jp/problem/binomial_coefficient
//only the tests with prime mod
//for mod inverse
#include "exp_mod.h"
```

```
// usage:
       n_{choose} = k \ nk(n, 1e9+7) to use 'choose', 'inv' with inputs < n
// or:
       n choose k nk(mod. mod) to use 'chooseWithLucasTheorem'
struct n_choose_k {
    n_choose_k(int n, int a_mod) : mod(a_mod), fact(n, 1), inv_fact(n, 1) {
        //this implementation doesn't work if n > mod because n! % mod = 0 when n > =
             \hookrightarrow mod. So 'inv_fact' array will be all 0's
        assert(max(n, 2) <= mod);</pre>
        //assert mod is prime. mod is intended to fit inside an int so that
        //multiplications fit in a longlong before being modded down. So this
        //will take sqrt(2~31) time
        for (int i = 2; i * i <= mod; i++) assert(mod % i);</pre>
        for (int i = 2; i < n; i++)</pre>
            fact[i] = 1LL * fact[i - 1] * i % mod;
        inv_fact.back() = pow(fact.back(), mod - 2, mod);
        for (int i = n - 2; i \ge 2; i--)
             inv_fact[i] = 1LL * inv_fact[i + 1] * (i + 1) % mod;
    //classic n choose k
    //fails when n \ge mod
    int choose(int n, int k) const {
        if (k < 0 \mid k > n) return 0:
        //now we know 0 <= k <= n so 0 <= n
        return 1LL * fact[n] * inv_fact[k] % mod * inv_fact[n - k] % mod;
    //lucas theorem to calculate n choose k in O(\log(k))
    //need to calculate all factorials in range [0,mod), so O(mod) time&space, so need
         \hookrightarrow smallish prime mod (< 1e6 maybe)
    //handles n >= mod correctly
    int choose_with_lucas_theorem(long long n, long long k) const {
        if (k < 0 \mid \mid k > n) return 0;
        if (k == 0 || k == n) return 1;
        return 1LL * choose_with_lucas_theorem(n / mod, k / mod) * choose(n % mod, k %
             \hookrightarrow mod) % mod:
    }
    //returns inverse of n in O(1)
    int inv(int n) const {
        assert(1 <= n); //don't divide by 0 :)</pre>
        return 1LL * fact[n - 1] * inv_fact[n] % mod;
    }
    int mod;
    vector<int> fact, inv_fact;
};
```

Listing 37: Partitions

Listing 38: Derangements

```
//cat derangements.h | ./hash.sh
//c221bb
#pragma once
//library checker tests: https://judge.yosupo.jp/problem/montmort_number_mod
//https://oeis.org/A000166
//for a permutation of size i:
//there are (i-1) places to move 0 to not be at index 0. Let's say we moved 0 to index j
//If we move value j to index 0 (forming a cycle of length 2), then there are dp[i-2]
     \hookrightarrow derangements of the remaining i-2 elements
//else there are dp[i-1] derangements of the remaining i-1 elements (including j)
vector<int> derangements(int n/*size of dp array*/, int mod) {
    vector<int> dp(n, 0);
    dp[0] = 1:
    for (int i = 2; i < n; i++)
        dp[i] = 1LL * (i - 1) * (dp[i - 1] + dp[i - 2]) % mod;
    return dp;
}
```

Listing 39: Prime Sieve Mobius

```
//cat prime_sieve_mobius.h | ./hash.sh
//4986da
#pragma once
//stress tests: tests/stress_tests/math/prime_sieve_mobius.cpp
//mobius[i] = 0 iff there exists a prime p s.t. i\%(p^2) = 0
//mobius[i] = -1 iff i has an odd number of distinct prime factors
//mobius[i] = 1 iff i has an even number of distinct prime factors
const int sz = 2e6 + 10:
int mobius[sz];
void calc_mobius() {
    mobius[1] = 1;
    for (int i = 1; i < sz; i++)
        for (int j = i + i; j < sz; j += i)
            mobius[j] -= mobius[i];
//a_prime[val] = some random prime factor of 'val'
//to check if 'val' is prime:
// if (a_prime[val] == val)
//
//to get all prime factors of a number 'val':
// while(val > 1) {
//
        int p = a_prime[val];
//
        //p is some prime factor of val
//
        val /= p;
// }
int a_prime[sz];
```

Listing 40: Row Reduce

```
//cat row_reduce.h / ./hash.sh
//a5406d
#pragma once
//library checker tests: https://judge.yosupo.jp/problem/system_of_linear_equations,

→ https://judge.yosupo.jp/problem/matrix_det,
     \hookrightarrow https://judge.yosupo.jp/problem/inverse_matrix
//for mod inverse
#include "exp_mod.h"
//First 'cols' columns of A represents a matrix to be left in reduced row echelon form
//Row operations will be performed to all later columns
//example usage:
// row_reduce(A, A[0].size(), mod) //row reduce matrix with no extra columns
pair<int/*rank*/, int/*determinant*/> row_reduce(vector<vector<int>>& mat, int cols, int
    int n = \text{mat.size}(), m = \text{mat}[0].\text{size}(), rank = 0, det = 1;
    assert(cols <= m);</pre>
    for (int col = 0; col < cols && rank < n; col++) {</pre>
        //find arbitrary pivot and swap pivot to current row
        for (int i = rank; i < n; i++)</pre>
            if (mat[i][col] != 0) {
                if (rank != i) det = det == 0 ? 0 : mod - det;
                 swap(mat[i], mat[rank]);
                 break;
        if (mat[rank][col] == 0) {
            det = 0:
             continue:
        }
        det = (1LL * det * mat[rank][col]) % mod;
        //make pivot 1 by dividing row by inverse of pivot
        const int a_inv = pow(mat[rank][col], mod - 2, mod);
        for (int j = 0; j < m; j++)</pre>
            mat[rank][j] = (1LL * mat[rank][j] * a_inv) % mod;
        //zero-out all numbers above & below pivot
        for (int i = 0; i < n; i++)
            if (i != rank && mat[i][col] != 0) {
                 const int val = mat[i][col];
                 for (int j = 0; j < m; j++) {
                     mat[i][j] -= 1LL * mat[rank][j] * val % mod;
                     if (mat[i][j] < 0) mat[i][j] += mod;</pre>
        rank++:
    assert(rank <= min(n, cols));</pre>
    return {rank, det};
```

Listing 41: Solve Linear Equations MOD

```
//cat solve_linear_mod.h | ./hash.sh
//867590
#pragma once
//library checker tests: https://judge.yosupo.jp/problem/system_of_linear_equations
#include "row_reduce.h"
struct matrix_info {
    int rank, det;
    vector<int> x;
};
//Solves\ A * x = b\ under\ prime\ mod.
//A is a n (rows) by m (cols) matrix, b is a length n column vector, x is a length m
    \hookrightarrow column vector.
//assumes n.m >= 1. else RTE
//Returns rank of A, determinant of A, and x (solution vector to A * x = b). x is empty
     \hookrightarrow if no solution. If multiple solutions, an arbitrary one is returned.
//Leaves A in reduced row echelon form (unlike kactl) with b appended.
//0(n * m * min(n.m))
matrix_info solve_linear_mod(vector<vector<int>>& mat, const vector<int>& b, int mod) {
    assert(mat.size() == b.size());
    int n = mat.size(), m = mat[0].size();
    for (int i = 0; i < n; i++)</pre>
        mat[i].push_back(b[i]);
    auto [rank, det] = row_reduce(mat, m, mod); //row reduce not including the last column
    //check if solution exists
    for (int i = rank; i < n; i++) {</pre>
        if (mat[i].back() != 0) return {rank, det, {} }; //no solution exists
    //initialize solution vector ('x') from row-reduced matrix
    vector<int> x(m, 0);
    for (int i = 0, j = 0; i < rank; i++) {
        while (mat[i][j] == 0) j++; //find pivot column
        x[j] = mat[i].back();
    return {rank, det, x};
```

Listing 42: Matrix Inverse

```
//cat matrix_inverse.h | ./hash.sh
//c32af9
#pragma once
//library checker tests: https://judqe.yosupo.jp/problem/inverse_matrix
#include "row_reduce.h"
//returns inverse of square matrix mat, empty if no inverse
vector<vector<int>> matrix_inverse(vector<vector<int>> mat/*intentional pass by value*/,
     \hookrightarrow int mod) {
    int n = mat.size();
    assert(n == (int)mat[0].size());
    //append identity matrix
    for (int i = 0; i < n; i++) {
        mat[i].resize(2 * n, 0);
        mat[i][i + n] = 1;
    auto [rank, det] = row_reduce(mat, n, mod);//row reduce first n columns, leaving
         \hookrightarrow inverse in last n columns
    if (rank < n) return {}: //no inverse
    for (int i = 0; i < n; i++)
        mat[i].erase(mat[i].begin(), mat[i].begin() + n);
    return mat;
```

```
Listing 43: Euler's Totient Phi Function
```

```
//cat totient.h | ./hash.sh
//36bd41
#pragma once
//stress tests: tests/stress_tests/math/totient.cpp
//Euler's totient function counts the positive integers
//up to a given integer n that are relatively prime to n.
//To improve, use Pollard-rho to find prime factors
int totient(int n) {
   int res = n:
   for (int i = 2; i * i <= n; i++) {
       if (n % i == 0) {
            while (n \% i == 0) n /= i;
            res -= res / i;
       }
   }
    if (n > 1) res -= res / n;
   return res:
```

Listing 44: MAX FLOW

Listing 45: Dinic

```
//cat dinic.h | ./hash.sh
//23e871
#pragma once
//status: not tested
struct max_flow {
    typedef long long 11;
    11 n, s, t;
    max_flow(int a_n, int a_s, int a_t) : n(a_n), s(a_s), t(a_t), d(n), ptr(n), q(n),
    void add_edge(ll a, ll b, ll cap) {
        edge_map[a * n + b] = e.size();
        edge e1 = \{a, b, cap, 0\};
        edge e2 = \{ b, a, 0, 0 \};
        g[a].push_back((ll) e.size());
        e.push_back(e1);
        g[b].push_back((11) e.size());
        e.push_back(e2);
   }
   11 get_flow() {
        11 \text{ flow} = 0:
        for (;;) {
            if (!bfs()) break;
            ptr.assign(ptr.size(), 0);
            while (ll pushed = dfs(s, inf))
                flow += pushed;
        }
        return flow;
    }
    ll get_flow_for_edge(ll a, ll b) {
```

```
return e[edge_map[a * n + b]].flow;
    const ll inf = 1e18;
    struct edge {
        11 a, b, cap, flow;
    };
    unordered_map<int, 11> edge_map;
    vector<1l> d, ptr, q;
    vector<edge> e;
    vector<vector<ll>>> g;
    bool bfs() {
        11 qh = 0, qt = 0;
        q[qt++] = s;
        d.assign(d.size(), -1);
        d[s] = 0;
        while (qh < qt && d[t] == -1) {
            11 v = q[qh++];
            for (size_t i = 0; i < g[v].size(); i++) {</pre>
                11 id = g[v][i],
                    to = e[id].b;
                if (d[to] == -1 && e[id].flow < e[id].cap) {</pre>
                     q[qt++] = to;
                     d[to] = d[v] + 1;
                }
            }
        return d[t] != -1;
    }
    11 dfs(ll v, ll flow) {
        if (!flow) return 0;
        if (v == t) return flow;
        for (; ptr[v] < (11) g[v].size(); ptr[v]++) {</pre>
            11 id = g[v][ptr[v]];
            11 to = e[id].b;
            if (d[to] != d[v] + 1) continue;
            11 pushed = dfs(to, min(flow, e[id].cap - e[id].flow));
            if (pushed) {
                e[id].flow += pushed;
                e[id ^ 1].flow -= pushed;
                return pushed;
        }
        return 0;
    }
};
```

Listing 46: Hungarian

```
//cat hungarian.h | ./hash.sh
//8ea32c
#pragma once
//library checker tests: https://judge.yosupo.jp/problem/assignment
//source: https://e-maxx.ru/algo/assignment_hungary
//
//input: cost[1...n][1...m] with n <= m
//n workers, indexed 1, 2, ..., n
//m jobs, indexed 1, 2, ..., m
//it costs 'cost[i][j]' to assign worker i to job j (1<=i<=n, 1<=j<=m)
//this returns *min* total cost to assign each worker to some distinct job
//0(n^2 * m)
```

```
//trick 1: set 'cost[i][j]' to INF to say: "worker 'i' cannot be assigned job 'j'"
//trick 2: 'cost[i][j]' can be negative, so to instead find max total cost over all
     \hookrightarrow matchings: set all 'cost[i][j]' to '-cost[i][j]'.
//Now max total cost = - HungarianMatch(cost).minCost
const long long inf = 1e18;
struct match {
    long long min_cost;
    vector<int> matching; //worker 'i' (1<=i<=n) is assigned to job 'matching[i]'
         \hookrightarrow (1<=matchina[i]<=m)
match hungarian(const vector<vector<long long>>& cost) {
    int n = cost.size() - 1, m = cost[0].size() - 1;
    assert(n <= m);</pre>
    vector<int> p(m + 1), way(m + 1);
    vector<long long> u(n + 1), v(m + 1);
    for (int i = 1; i <= n; i++) {
        p[0] = i;
        int j0 = 0;
        vector<long long> minv(m + 1, inf);
        vector<bool> used(m + 1, false);
        do {
            used[j0] = true;
             int i0 = p[i0], i1 = 0;
            long long delta = inf;
            for (int j = 1; j \le m; j++)
                 if (!used[i]) {
                     long long cur = cost[i0][j] - u[i0] - v[j];
                     if (cur < minv[j])</pre>
                         minv[j] = cur, way[j] = j0;
                     if (minv[j] < delta)</pre>
                         delta = minv[j], j1 = j;
            for (int j = 0; j \le m; j++)
                 if (used[j])
                     u[p[j]] += delta, v[j] -= delta;
                 else
                     minv[i] -= delta:
            j0 = j1;
        } while (p[j0] != 0);
            int j1 = wav[j0];
            p[j0] = p[j1];
             j0 = j1;
        } while (j0);
    vector<int> ans(n + 1);
    for (int j = 1; j \le m; j++)
        ans[p[j]] = j;
    return {-v[0], ans};
```

Listing 47: Min Cost Max Flow

```
//cat min_cost_max_flow.h | ./hash.sh
//805596
#pragma once
//status: not tested
const long long inf = 1e18;
struct min_cost_max_flow {
```

```
typedef long long 11;
struct edge {
    ll a, b, cap, cost, flow;
    size t back:
};
vector<edge> e;
vector<vector<1l>>> g;
11 n. s. t:
11 k = inf; // The maximum amount of flow allowed
min_cost_max_flow(int a_n, int a_s, int a_t) : n(a_n), s(a_s), t(a_t) {
    g.resize(n);
void add_edge(ll a, ll b, ll cap, ll cost) {
    edge e1 = {a, b, cap, cost, 0, g[b].size() };
    edge e2 = {b, a, 0, -cost, 0, g[a].size() };
    g[a].push_back((ll) e.size());
    e.push_back(e1);
    g[b].push_back((ll) e.size());
    e.push_back(e2);
}
// Returns {flow, cost}
pair<11, 11> get_flow() {
    11 \text{ flow} = 0, \text{ cost} = 0;
    while (flow < k) {
        vector<ll> id(n, 0), d(n, inf), q(n), p(n);
        vector<size_t> p_edge(n);
        11 qh = 0, qt = 0;
        q[qt++] = s;
        d[s] = 0;
        while (qh != qt) {
            11 v = q[qh++];
            id[v] = 2;
            if (qh == n) qh = 0;
            for (size_t i = 0; i < g[v].size(); i++) {</pre>
                edge& r = e[g[v][i]];
                if (r.flow < r.cap && d[v] + r.cost < d[r.b]) {
                    d[r.b] = d[v] + r.cost;
                    if (id[r.b] == 0) {
                        q[qt++] = r.b;
                        if (qt == n) qt = 0;
                    } else if (id[r.b] == 2) {
                        if (--qh == -1) qh = n - 1;
                        q[qh] = r.b;
                    id[r.b] = 1;
                    p[r.b] = v;
                    p_{edge}[r.b] = i;
            }
        if (d[t] == inf) break;
        11 addflow = k - flow:
        for (11 v = t; v != s; v = p[v]) {
            11 pv = p[v];
            size_t pr = p_edge[v];
            addflow = min(addflow, e[g[pv][pr]].cap - e[g[pv][pr]].flow);
        for (11 v = t; v != s; v = p[v]) {
            11 pv = p[v];
            size_t pr = p_edge[v], r = e[g[pv][pr]].back;
            e[g[pv][pr]].flow += addflow;
```

Listing 48: MISC

Listing 49: Disjoint Set

```
//cat dsu.h | ./hash.sh
//9b3c97
#pragma once
//library checker tests: https://judge.yosupo.jp/problem/unionfind
//stress tests: tests/stress_tests/misc/disjoint_set.cpp
struct DSU { //NOLINT(readability-identifier-naming)
    int num_sets;
    vector<int> par;
    DSU(int n) : num_sets(n), par(n, -1) {}
    DSU(const DSU& rhs) : num_sets(rhs.num_sets), par(rhs.par) {}
    int find(int x) {
        return par[x] < 0 ? x : par[x] = find(par[x]);</pre>
    }
    int size_of_set(int x) {
        return -par[find(x)];
   }
   bool join(int x, int y) {
        if ((x = find(x)) == (y = find(y))) return false;
        if (par[y] < par[x]) swap(x, y);</pre>
        par[x] += par[y];
        par[y] = x;
        num_sets--:
        return true:
   }
};
```

```
Listing 50: PBDS
//cat policy_based_data_structures.h / ./hash.sh
//807de9
#pragma once
//status: not tested
//place these includes *before* the '#define int long long' else compile error
//not using <bits/extc++.h> as it compile errors on codeforces c++20 compiler
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
//BST with extra functions https://codeforces.com/bloq/entry/11080
//order_of_key - # of elements *strictly* less than given element
//find_by_order - find kth largest element, k is 0 based so find_by_order(0) returns min
    \hookrightarrow element
template<class T>
using indexed_set = tree<T, null_type, less<T>, rb_tree_tag,

    tree_order_statistics_node_update>;
```

Listing 51: Monotonic Stack

```
//cat monotonic_stack.h | ./hash.sh
//90f107
#pragma once
//library checker tests: https://judge.yosupo.jp/problem/cartesian_tree
//stress tests: tests/stress_tests/misc/count_rectangles.cpp
//calculates array 'left' with:
//for\ every\ index\ j\ with\ left[i] < j < i:\ arr[j] > arr[i]
//and
//arr[left[i]] <= arr[i] if left[i] != -1
//trick: pass in vector<pair<T/*value*/, int/*index*/>> with arr[i].second = i (0<=i<n)
     \hookrightarrow to simulate arr[i] >= arr[i]
//0(n)
template<class T>
vector<int> monotonic stack(const vector<T>& arr) {
    int n = arr.size();
    vector<int> left(n);
    for (int i = 0; i < n; i++) {
        int& j = left[i] = i - 1;
        while (j >= 0 && arr[j] > arr[i]) j = left[j];
    return left;
```

Listing 52: Count Rectangles

```
//cat count_rectangles.h | ./hash.sh
//69a1b0
#pragma once
#include "monotonic stack.h"
//stress tests: tests/stress_tests/misc/count_rectangles.cpp
//given a 2D boolean matrix, calculate cnt[i][j]
//cnt[i][j] = the number of times an (i * j) rectangle appears in the matrix such that
    \hookrightarrow all i*j cells in the rectangle are true
//Note cnt[0][j] and cnt[i][0] will contain garbage values
vector<vector<int>> count_rectangles(const vector<vector<bool>>& grid) {
    const int n = grid.size(), m = grid[0].size();
    vector<vector<int>> cnt(n + 1, vector<int>(m + 1, 0));
    vector<int> arr(m, 0);
    auto rv = [&](int j) -> int {//reverse
        return m - 1 - j;
   for (int i = 0; i < n; i++) {</pre>
        vector<pair<int, int>> arr_rev(m);
        for (int j = 0; j < m; j++) {
            arr[j] = grid[i][j] * (arr[j] + 1);
            arr_rev[rv(j)] = {arr[j], j};
       }
```

```
vector<int> left = monotonic stack(arr):
    vector<int> right = monotonic_stack(arr_rev);
    for (int j = 0; j < m; j++) {
        int l = j - left[j] - 1, r = rv(right[rv(j)]) - j - 1;
        cnt[arr[j]][1 + r + 1]++;
        cnt[arr[j]][1]--;
        cnt[arr[j]][r]--;
    }
}
for (int i = 1; i <= n; i++)
    for (int k = 0; k < 2; k++)
        for (int j = m; j > 1; j--)
            cnt[i][j - 1] += cnt[i][j];
for (int j = 1; j <= m; j++)</pre>
    for (int i = n; i > 1; i--)
        cnt[i - 1][j] += cnt[i][j];
return cnt;
```

```
static uint64_t splitmix64(uint64_t x) {
       // http://xorshift.di.unimi.it/splitmix64.c
       x += 0x9e3779b97f4a7c15;
       x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
       x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
       return x ^ (x >> 31);
   size_t operator()(uint64_t x) const {
       static const uint64_t fixed_random =
           return splitmix64(x + fixed_random);
   }
};
//usage:
unordered_map<long long, int, custom_hash> safe_map;
#include "policy_based_data_structures.h"
gp_hash_table<long long, int, custom_hash> safe_hash_table;
```

Listing 53: Longest Increasing Subsequence

```
//cat lis.h | ./hash.sh
//a243e1
#pragma once
//library checker tests: https://judqe.yosupo.jp/problem/static_range_lis_query
//returns array of indexes representing the longest *strictly* increasing subsequence
//for non-decreasing: pass in a vector\langle pair \langle T, int \rangle \rangle where second is 0, 1, ..., n-1
//alternatively, there's this https://codeforces.com/blog/entry/13225
template<class T>
vector<int> LIS(const vector<T>& arr) { //NOLINT(readability-identifier-naming)
    if (arr.empty()) return {};
    vector<int> dp{0}/*array of indexes into 'arr'*/, prev(arr.size(), -1);
    for (int i = 1; i < (int)arr.size(); i++) {</pre>
        auto it = lower_bound(dp.begin(), dp.end(), i, [&](int x, int y) -> bool {
            return arr[x] < arr[y];</pre>
        });
        if (it == dp.end()) {
            prev[i] = dp.back();
            dp.push_back(i);
        } else {
            prev[i] = it == dp.begin() ? -1 : *(it - 1);
        }
        //here, dp.size() = length of LIS of prefix of arr ending at index i
   }
    vector<int> res(dp.size());
    for (int i = dp.back(), j = dp.size(); i != -1; i = prev[i])
        res[--j] = i;
    return res;
```

Listing 54: Safe Hash

```
//cat safe_hash.h | ./hash.sh
//e837ee
#pragma once
//status: not tested
//source: https://codeforces.com/blog/entry/62393
struct custom_hash {
```