Listings

Listing 1: CODE HASHES

Listing 2: hash.sh

Listing 3: GRAPHS

```
% cat graphs/bridges_and_outs.h | ./misc/hash.sh
650062
```

Listing 4: Bridges and Cuts

```
#pragma once
//tested on https://judge.yosupo.jp/problem/two_edge_connected_components and

→ https://judge.yosupo.jp/problem/biconnected_components

//with \ asserts \ checking \ correctness \ of \ is \textit{Bridge} \ and \ is \textit{Cut}
//2 edge cc and bcc stuff doesn't depend on each other, so delete whatever is not needed
//To initialize 'adj':
//eid = 0
//for each edge (a,b):
// adj[a].emplace_back(b, eid);
// adj[b].emplace_back(a, eid++);
struct info {
    /\!/2 \ \textit{edge connected component stuff (e.g. components split by bridge \textit{edges})}
         \hookrightarrow https://cp-algorithms.com/graph/bridge-searching.html
    int num2EdgeCCs;
    vector<bool> isBridge;//edge id -> true iff bridge edge
    vector<int> TwoEdgeCCID; //node -> ID of 2-edge component (which are labeled 0, 1,

    ∴ ∴ `num2EdgeCCs `-1)

    //bi-connected component stuff (e.g. components split by cut/articulation nodes)
         \hookrightarrow https://cp-algorithms.com/graph/cutpoints.html
    int numBCCs;
    vector<bool> isCut;//node -> true iff cut node
    vector<int> bccID;//edge id -> ID of BCC (which are labeled 0, 1, ..., 'numBCCs'-1)
};
info bridge_and_cut(const vector<vector<pair<int/*neiqhbor*/, int/*edqe id*/>>>&

    → adj/*undirected graph*/, int m/*number of edges*/) {

    //stuff for both (always keep)
    int n = adj.size(), timer = 1;
```

```
vector<int> tin(n, 0), low(n, 0);
//2 edge CC stuff (delete if not needed)
int num2EdgeCCs = 0;
vector<bool> isBridge(m, false);
vector<int> TwoEdgeCCID(n), nodeStack;
//BCC stuff (delete if not needed)
int numBCCs = 0:
vector<bool> isCut(n, false);
vector<int> bccID(m), edgeStack;
auto dfs = [&] (auto&& dfsPtr, int v, int pId) -> void {
    tin[v] = low[v] = timer++;
    int deg = 0;
    nodeStack.push_back(v);
    for (auto [to, eId] : adj[v]) {
        if (eId == pId) continue;
        if (!tin[to]) {
            edgeStack.push_back(eId);
            dfsPtr(dfsPtr, to, eId);
            if (low[to] >= tin[v]) {
                isCut[v] = true:
                while (true) {
                    int edge = edgeStack.back();
                    edgeStack.pop_back();
                    bccID[edge] = numBCCs;
                    if (edge == eId) break;
                }
                numBCCs++;
            low[v] = min(low[v], low[to]);
            deg++;
        } else if (tin[to] < tin[v]) {</pre>
            edgeStack.push_back(eId);
            low[v] = min(low[v], tin[to]);
    if (pId == -1) isCut[v] = (deg > 1);
    if (tin[v] == low[v]) {
        if (pId != -1) isBridge[pId] = true;
        while (true) {
            int node = nodeStack.back();
            nodeStack.pop_back();
            TwoEdgeCCID[node] = num2EdgeCCs;
            if (node == v) break;
        num2EdgeCCs++;
};
for (int i = 0; i < n; i++) {</pre>
    if (!tin[i])
        dfs(dfs, i, -1);
}
return {num2EdgeCCs, isBridge, TwoEdgeCCID, numBCCs, isCut, bccID};
```

```
found = true;
                par = node;
                node = to;
                break;
        }
    return node;
void dfs1(int node, int par) {
    removed[node] = true;
    parent[node] = par;
    for (int to : adj[node]) {
        if (!removed[to])
            dfs1(findCentroid(to), node);
//dfs1(findCentroid(1), 0);
% cat graphs/countPathLengths.h | ./misc/hash.sh
470b6f
```

if (to != par && !removed[to] && sizes[to] > sizeCap) {

#pragma once

//status: not tested

vector<int> adj[Max];

bool removed[Max];

const int Max = 2e5 + 2;

sizes[node] = 1;

int sizes[Max], parent[Max];

void dfs2(int node, int par) {

int findCentroid(int node) {

found = false;

dfs2(node, node);

bool found = true;

int par = node; while (found) {

for (int to : adj[node]) {

dfs2(to, node);

int sizeCap = sizes[node] / 2;

for (int to : adj[node]) {

if (to != par && !removed[to]) {

sizes[node] += sizes[to];

% cat graphs/centroid.h | ./misc/hash.sh

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Listing 5: Centroid

Listing 6: Count Path Lengths

```
#pragma once
//status: doesn't compile, but should be correct. Need to import FFT code (like
     \hookrightarrow https://qithub.com/kth-competitive-programming/kactl/blob/main/content/numerical/Fus
```

```
const int Max = 1e6 + 10:
int n. sizes[Max]:
vector<int> adj[Max], cntPathLength[Max];
11 cntTotalPathLengths[Max];
bool removed[Max];
void dfs2(int node, int par, int root, int currDist) {
    while ((int)cntPathLength[root].size() <= currDist)</pre>
        cntPathLength[root].push_back(0);
    cntPathLength[root][currDist]++:
    sizes[node] = 1;
    for (int to : adj[node]) {
        if (to != par && !removed[to]) {
            dfs2(to, node, root, currDist + 1);
            sizes[node] += sizes[to];
   }
int findCentroid(int node) {
    dfs2(node, node, node, 1);
    bool found = true:
    int sizeCap = sizes[node] / 2;
    int par = node;
    while (found) {
        found = false;
        for (int to : adj[node]) {
            if (to != par && !removed[to] && sizes[to] > sizeCap) {
                found = true;
                par = node;
                node = to;
                break;
        }
    }
    return node;
void dfs1(int node, int par) {
    removed[node] = true;
    int maxLength = 1:
    for (int to : adj[node]) {
        if (to != par && !removed[to]) {
            cntPathLength[to].clear();
            cntPathLength[to].push_back(0);
            dfs2(to, to, to, 1);
            maxLength = max(maxLength, (int)cntPathLength[to].size());
    vector<int> temp(maxLength, 0);
    temp[0]++;
    for (int to : adj[node]) {
        if (to != par && !removed[to]) {
            vector<ll> prod = multiply(temp, cntPathLength[to]);
            for (int i = 0; i < (int)prod.size(); ++i)</pre>
                cntTotalPathLengths[i] += prod[i];
            for (int i = 0; i < (int)cntPathLength[to].size(); ++i)</pre>
                temp[i] += cntPathLength[to][i];
    }
    for (int to : adj[node]) {
```

```
% cat graphs/disjointSet.h | ./misc/hash.sh

836916
```

Listing 7: Disjoint Set

```
#pragma once
//status: tested on random inputs, and on https://judge.yosupo.jp/problem/unionfind
struct disjointSet {
    int numSets;
    vector<int> par;
    disjointSet(int n) : numSets(n), par(n, -1) {}
    disjointSet(const disjointSet& rhs) : numSets(rhs.numSets), par(rhs.par) {}
    int find(int x) {
        return par[x] < 0 ? x : par[x] = find(par[x]);</pre>
    int sizeOfSet(int x) {
        return -par[find(x)];
    bool merge(int x, int y) {
        if ((x = find(x)) == (y = find(y))) return false;
        if (par[y] < par[x]) swap(x, y);</pre>
        par[x] += par[y];
        par[y] = x;
        numSets--:
        return true;
};
```

```
% cat graphs/dijkstra.h | ./misc/hash.sh
27560a
```

Listing 8: Dijkstra

```
auto it = q.begin();
int node = it->second;
q.erase(it);
for (auto [to, weight] : adj[node])
    if (len[to] > weight + len[node]) {
        q.erase({len[to], to});
        len[to] = weight + len[node];
        q.insert({len[to], to});
}
return len;
}
```

```
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```

% cat graphs/dsuTree.h | ./misc/hash.sh

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#pragma once

Listing 9: DSU Tree

```
//status: not tested
const int Max = 1e5 + 3:
int color[Max], Time = 1, timeIn[Max], timeOut[Max], ver[Max], Size[Max], cnt[Max],

    heavyChild[Max], Depth[Max] = {0}, answer[Max];
vector<int> adj[Max];
void dfs(int node, int prev) {
   timeIn[node] = Time;
   ver[Time] = node;
   Time++;
   Size[node] = 1;
   int largest = heavyChild[node] = -1;
   Depth[node] = 1 + Depth[prev];
   for (int to : adj[node]) {
       if (to == prev) continue;
       dfs(to, node);
       Size[node] += Size[to];
       if (Size[to] > largest) {
           largest = Size[to];
            heavyChild[node] = to;
       }
   }
   timeOut[node] = Time;
void dfs1(int node, int prev, bool keep = true) {
   for (int to : adj[node]) {
       if (to == prev || to == heavyChild[node]) continue;
       dfs1(to, node, false);
   }
   if (heavyChild[node] != -1)
       dfs1(heavyChild[node], node, true);
   cnt[color[node]]++:
   for (int to : adi[node]) {
       if (to == prev || to == heavyChild[node]) continue;
       for (int i = timeIn[to]; i < timeOut[to]; ++i)</pre>
            cnt[color[ver[i]]]++;
```

```
% cat graphs/floydWarshall.h | ./misc/hash.sh
```

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Listing 10: Floyd Warshall

```
#pragma once
//status: not tested
//**for directed graphs only** if you initialize len[i][i] to infinity, then
//afterward floyds, len[i][i] = length of shortest cycle including node 'i'
//another trick: change 'len' to 2d array of *bools* where len[i][j] = true if
//there exists an edge from i \rightarrow j in initial graph. Also do:
//'len[i][j] /= len[i][k] & len[k][j]'
//Then after floyds, len[i][j] = true iff there's exists some path from node
//'i' to node 'j'
//Changing the order of for-loops to i-j-k (instead of the current k-i-j)
//results in min-plus matrix multiplication. If adjacency matrix is M, then
//after computing M^k (with binary exponentiation), M[i][j] = min length path
//from i to j with at most k edges.
for (int k = 0; k < n; k++)
    for (int i = 0; i < n; i++)
        for (int j = 0; j < n; j++)
            len[i][j] = min(len[i][j], len[i][k] + len[k][j]);
```

% cat graphs/hld.h | ./misc/hash.sh

322dee

Listing 11: HLD

```
vector<int> Size, par, Depth, timeIn, Next, timeInToNode;
hld(vector<vector<int>>& adj /*forest of trees*/, int root = -1/*pass in to specify
    \hookrightarrow root, usually for a single component*/):
    Size(adj.size(), 1), par(adj.size(), -1), Depth(adj.size(), 1),
         int Time = 0;
    auto callDfss = [&](int node) -> void {
        Next[node] = par[node] = node;
        dfs1(node, adj);
        dfs2(node, adj, Time);
    };
    if (root != -1)
        callDfss(root);
    for (int i = 0; i < (int) adj.size(); i++) {</pre>
        if (par[i] == -1) //roots each tree by node with min label
            callDfss(i);
void dfs1(int node, vector<vector<int>>& adj) {
    for (auto& to : adi[node]) {
        if (to == par[node]) continue;
        Depth[to] = 1 + Depth[node];
        par[to] = node;
        dfs1(to, adj);
        Size[node] += Size[to];
        if (Size[to] > Size[adj[node][0]] || adj[node][0] == par[node])
            swap(to, adj[node][0]);
    }
}
void dfs2(int node, const vector<vector<int>>& adj, int& Time) {
    timeIn[node] = Time;
    timeInToNode[Time] = node;
    Time++:
    for (auto to : adj[node]) {
        if (to == par[node]) continue;
        Next[to] = (Time == timeIn[node] + 1 ? Next[node] : to);
        dfs2(to, adj, Time);
    }
}
// Returns intervals (of timeIn's) corresponding to the path between u and v, not
     \hookrightarrow necessarily in order
// This can answer queries for "is some node 'x' on some path" by checking if the
    \hookrightarrow timeIn[x] is in any of these intervals
vector<pair<int, int>> path(int u, int v) const {
    vector<pair<int, int>> res;
    for (;; v = par[Next[v]]) {
        if (timeIn[v] < timeIn[u]) swap(u, v);</pre>
        if (timeIn[Next[v]] <= timeIn[u]) {</pre>
            res.push_back({timeIn[u], timeIn[v]});
            return res;
        res.push_back({timeIn[Next[v]], timeIn[v]});
// Returns interval (of timeIn's) corresponding to the subtree of node i
// This can answer queries for "is some node 'x' in some other node's subtree" by
     \hookrightarrow checking if timeIn[x] is in this interval
pair<int, int> subtree(int i) const {
    return {timeIn[i], timeIn[i] + Size[i] - 1};
// Returns lca of nodes u and v
```

```
int lca(int u, int v) const {
    for (;; v = par[Next[v]]) {
        if (timeIn[v] < timeIn[u]) swap(u, v);
        if (timeIn[Next[v]] <= timeIn[u]) return u;
    }
};</pre>
```

```
% cat graphs/hopcroftKarp.h | ./misc/hash.sh
a3ff4a
```

Listing 12: Hopcroft Karp

#pragma once

```
//Modified from

→ https://qithub.com/foreverbell/acm-icpc-cheat-sheet/blob/master/src/graph-algorithm/
//Worst case O(E*sqrt(V)) but faster in practice
//status: tested on https://judge.yosupo.jp/problem/bipartitematching with asserts
     \hookrightarrow checking correctness of min vertex cover
struct match {
    //# of edges in matching (which = size of min vertex cover by ÖKnig's theorem)
    int sizeOfMatching;
    //an arbitrary max matching is found. For this matching:
    //if \ ml[nodeLeft] == -1:
          'nodeLeft' is not in matching
    //else:
          the edge 'nodeLeft' <=> ml[nodeLeft] is in the matching
    //similarly for mr with edge mr[nodeRight] <=> nodeRight in matching if
         \hookrightarrow mr[nodeRight] != -1
    //matchings stored in ml and mr are the same matching
    //provides way to check if any node is in matching
    vector<int> ml, mr;
    //an arbitrary min vertex cover is found. For this MVC: leftMVC['left node'] is true
         \hookrightarrow iff 'left node' is in the min vertex cover (same for rightMVC)
    //if leftMVC['left node'] is false, then 'left node' is in the corresponding maximal
         \hookrightarrow independent set
    vector<bool> leftMVC, rightMVC;
};
//Think of the bipartite graph as having a left side (with size LSz) and a right side
     \hookrightarrow (with size rSz).
//Nodes on left side are indexed 0,1,\ldots,lSz-1
//Nodes on right side are indexed 0,1,...,rSz-1
//'adj' is like a directed adjacency list containing edges from left side -> right side:
//To initialize 'adj': For every edge nodeLeft <=> nodeRight, do:
     \hookrightarrow adj[nodeLeft].push_back(nodeRight)
match hopcroftKarp(const vector<vector<int>>& adj /*bipartite graph*/, int rSz/*number
     \hookrightarrow of nodes on right side*/) {
    int sizeOfMatching = 0, 1Sz = adj.size();
    vector<int> level(lSz), ml(lSz, -1), mr(rSz, -1);
    vector<bool> visL(lSz, false);
    while (true) {
        queue<int> a:
        for (int i = 0; i < 1Sz; i++) {
```

```
if (ml[i] == -1) level[i] = 0, q.push(i);
        else level[i] = -1:
    }
    while (!q.empty()) {
        int u = q.front();
        q.pop();
        for (int x : adi[u]) {
            int v = mr[x];
            if (v != -1 && level[v] < 0) {</pre>
                level[v] = level[u] + 1:
                q.push(v);
            }
        }
    }
    auto dfs = [&](auto&& dfsPtr, int u) -> bool {
        visL[u] = true;
        for (int x : adj[u]) {
            int v = mr[x]:
            if (v == -1 || (!visL[v] && level[u] < level[v] && dfsPtr(dfsPtr, v))) {</pre>
                mr[x] = u;
                return true:
            }
        return false;
    };
    visL.assign(lSz, false);
    bool found = false;
    for (int i = 0; i < 1Sz; i++)</pre>
        if (ml[i] == -1 && dfs(dfs, i)) {
            found = true;
            sizeOfMatching++;
    if (!found) break;
}
//find min vertex cover
vector<bool> visR(rSz, false);
auto dfs = [&](auto&& dfsPtr. int node) -> void {
    for (int to : adj[node]) {
        if (!visR[to] && mr[to] != -1) {
            visR[to] = true:
            dfsPtr(dfsPtr, mr[to]);
    }
};
for (int i = 0; i < 1Sz; i++) {</pre>
    visL[i] = !visL[i];
    if (ml[i] == -1)
        dfs(dfs, i);
return {sizeOfMatching, ml, mr, visL, visR};
```

```
% cat graphs/lca.h | ./misc/hash.sh
```

eaac25

Listing 13: LCA

```
#pragma once
//https://codeforces.com/blog/entry/74847
//assumes a single tree, 1-based nodes is possible by passing in 'root' in range [1, n]
//status: all functions tested on random trees. 'qetLca' also tested on
    \hookrightarrow https://judge.yosupo.jp/problem/lca
struct lca {
    vector<int> jmp, jmpEdges, par, subSize, depth;
    vector<long long> dist:
    lca(const vector<vector<pair<int, long long>>>& adj, int root) :
        jmp(adj.size(), root),
        jmpEdges(adj.size(), 0),
        par(adj.size(), root),
        subSize(adj.size(), 1),
        depth(adj.size(), 0),
        dist(adj.size(), OLL) {
        dfs(root, -1, adj);
    }
    void dfs(int node, int parent, const vector<vector<pair<int, long long>>>& adj) {
        for (auto [ch, w] : adj[node]) {
            if (ch == parent) continue;
            depth[ch] = 1 + depth[node];
            par[ch] = node;
            dist[ch] = w + dist[node];
            if (parent != -1 && jmpEdges[node] == jmpEdges[jmp[node]])
                jmp[ch] = jmp[jmp[node]], jmpEdges[ch] = 2 * jmpEdges[node] + 1;
                jmp[ch] = node, jmpEdges[ch] = 1;
            dfs(ch, node, adj);
            subSize[node] += subSize[ch]:
    }
    //traverse up k edges in O(\log(k)). So with k=1 this returns 'node''s parent
    int kthPar(int node, int k) const {
        k = min(k, depth[node]);
        while (k > 0) {
            if (impEdges[node] <= k) {</pre>
                k -= jmpEdges[node];
                node = jmp[node];
            } else {
                k--;
                node = par[node];
        return node;
    }
    int getLca(int x, int y) const {
        if (depth[x] < depth[y]) swap(x, y);</pre>
        x = kthPar(x, depth[x] - depth[y]);
        while (x != y) {
            if (jmp[x] == jmp[y])
                x = par[x], y = par[y];
                x = jmp[x], y = jmp[y];
        }
        return x;
```

```
int distEdges(int x, int y) const {
    return depth[x] + depth[y] - 2 * depth[getLca(x, y)];
}

long long distWeight(int x, int y) const {
    return dist[x] + dist[y] - 2 * dist[getLca(x, y)];
}
};
```

```
% cat graphs/scc.h | ./misc/hash.sh
```



Listing 14: SCC

```
#pragma once
//status: tested on https://judge.yosupo.jp/problem/scc
//building of condensation graph tested on https://cses.fi/problemset/task/1686/
struct sccInfo {
    int numSCCs:
    //scc's are labeled 0,1,..., 'numSCCs-1'
    //sccId[i] is the id of the scc containing node 'i'
    vector<int> sccId;
};
sccInfo getSCCs(const vector<vector<int>>& adj /*directed, unweighted graph*/) {
    int n = adj.size(), numSCCs = 0;
    vector<int> sccId(n);
    stack<int> seen;
        vector<bool> vis(n, false);
        auto dfs = [&](auto&& dfsPtr, int curr) -> void {
            vis[curr] = true:
            for (int x : adj[curr]) {
                if (!vis[x])
                    dfsPtr(dfsPtr, x);
            seen.push(curr);
        };
        for (int i = 0; i < n; ++i) {
            if (!vis[i])
                dfs(dfs, i);
        }
    }
    vector<vector<int>> adjInv(n);
    for (int i = 0; i < n; ++i) {
        for (int to : adj[i])
            adjInv[to].push_back(i);
    vector<bool> vis(n, false);
    auto dfs = [&](auto&& dfsPtr. int curr) -> void {
        vis[curr] = true:
        sccId[curr] = numSCCs;
        for (int x : adjInv[curr]) {
            if (!vis[x])
```

```
dfsPtr(dfsPtr, x);
}

};
while (!seen.empty()) {
    int node = seen.top();
    seen.pop();
    if (vis[node])
        continue;
    dfs(dfs, node);
    numSCCs++;
}
return {numSCCs, sccId};
}
```

Listing 15: RANGE DATA STRUCTURES

```
% cat range_data_structures/segTreeBeats.h | ./misc/hash.sh

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```

Listing 16: Segment Tree Beats

```
#pragma once
//status: not tested, used in various problems
struct SegTreeBeats {
    typedef long long 11;
    struct Node {
        11 sum;
        11 mx;
        11 secondMx;
       11 cntMx;
   };
   vector<Node> tree;
   vector<int> lazy;
   int n, size;
   const ll inf = 1e18;
    /*implement these*/
   const Node zero = {0, -inf, -inf, 0};
   Node combine(const Node& L, const Node& R) {
        Node par:
        par.sum = L.sum + R.sum;
        if (L.mx == R.mx)
            par.cntMx = L.cntMx + R.cntMx;
        else if (L.mx > R.mx)
            par.cntMx = L.cntMx;
            par.cntMx = R.cntMx;
        par.mx = max(L.mx, R.mx);
        par.secondMx = -inf;
        for (11 val : {
                    L.mx, R.mx, L.secondMx, R.secondMx
           if (par.mx != val) {
                assert(par.mx > val);
                par.secondMx = max(par.secondMx, val);
```

```
}
    }
    return par;
}
void push(int node, int start, int end) {
    if (start == end) return;
    assert(start < end):</pre>
    for (int child : {
                2 * node, 2 * node + 1
            }) {
        if (tree[child].mx <= tree[node].mx) continue;</pre>
        tree[child].sum -= (tree[child].mx - tree[node].mx) * tree[child].cntMx;
         tree[child].mx = tree[node].mx;
    }
}
SegTreeBeats(const vector<int>& arr) : n((int) arr.size()) {
    size = 1:
    while (size < n) size <<= 1;</pre>
    size <<= 1:
    tree.resize(size);
    lazy.resize(size, 0);
    build(arr, 1, 0, n - 1);
}
void build(const vector<int>& arr, int node, int start, int end) {
    if (start == end) {
        tree[node].sum = arr[start];
        tree[node].mx = arr[start];
        tree[node].secondMx = -inf;
        tree[node].cntMx = 1;
    } else {
        const int mid = (start + end) / 2;
        build(arr, 2 * node, start, mid);
        build(arr, 2 * node + 1, mid + 1, end);
        tree[node] = combine(tree[2 * node], tree[2 * node + 1]);
}
//set \ a[i] = min(a[i], newMn), for i in range: [l,r]
void update(int 1, int r, int newMn) {
    update(1, 0, n - 1, 1, r, newMn);
void update(int node, int start, int end, int 1, int r, int newMn) {
    assert(start <= end);</pre>
    push(node, start, end);
    if (start > r || end < l || tree[node].mx <= newMn) return;</pre>
    if (start >= 1 && end <= r && tree[node].secondMx < newMn) {</pre>
        tree[node].sum -= (tree[node].mx - newMn) * tree[node].cntMx;
        tree[node].mx = newMn;
        return;
    assert(start < end):</pre>
    const int mid = (start + end) / 2:
    update(2 * node, start, mid, 1, r, newMn);
    update(2 * node + 1, mid + 1, end, 1, r, newMn);
    tree[node] = combine(tree[2 * node], tree[2 * node + 1]);
}
//query for sum/max in range [l,r]
Node query(int 1, int r) {
    return query(1, 0, n - 1, 1, r);
Node query(int node, int start, int end, int 1, int r) {
```

```
% cat range_data_structures/implicitSegTree.h | ./misc/hash.sh

920519
```

Listing 17: Implicit Segment Tree

```
#pragma once
//status: tested on https://cses.fi/problemset/task/1144/ and
     \hookrightarrow https://judge.yosupo.jp/problem/point_add_range_sum
struct implicitSegTree {
    struct Node {
        int 1Ch, rCh;//children ptrs, indexes into 'tree'; 0 for null
        long long sum;
    };
    int sz;
    deque<Node> tree;
    implicitSegTree(int _sz): sz(_sz) {
        tree.push_back({0, 0, OLL});//acts as null
        tree.push_back({0, 0, OLL});//root node
    void update(int idx, long long diff) {
        update(1, 0, sz - 1, idx, diff);
    int update(int v, int tl, int tr, int idx, long long diff) {
        if (t1 == tr) {
            if (v == 0) {
                tree.push_back(tree[0]);
                v = tree.size() - 1;
            tree[v].sum += diff;
            return v;
        int tm = (t1 + tr) / 2:
        int 1Ch = tree[v].1Ch:
        int rCh = tree[v].rCh;
        if (idx <= tm)</pre>
            1Ch = update(1Ch, t1, tm, idx, diff);
            rCh = update(rCh, tm + 1, tr, idx, diff);
        if (v == 0) {
            tree.push_back(tree[0]);
            v = tree.size() - 1:
        tree[v] = {1Ch, rCh, tree[1Ch].sum + tree[rCh].sum};
        return v;
```

#pragma once

```
% cat range_data_structures/kth_smallest.h | ./misc/hash.sh
8f3b27
```

Listing 18: Kth Smallest

```
//modified from

→ https://cp-algorithms.com/data_structures/segment_tree.html#preserving-the-history
//tested on https://judge.yosupo.jp/problem/range_kth_smallest
//works for -1e9 <= arr[i] <= 1e9
struct kth_smallest {
    const int mx = 1e9;
    struct Node {
       int 1Ch, rCh;//children, indexes into 'tree'
       int sum;
   };
   deque<Node> tree;
    vector<int> roots;
   kth_smallest(const vector<int>& arr) {
        tree.push_back({0, 0, 0}); //acts as null
       roots.push_back(0);
       for (int i = 0; i < (int)arr.size(); i++) {</pre>
            assert(-mx <= arr[i] && arr[i] <= mx);</pre>
            roots.push_back(update(roots.back(), -mx, mx, arr[i], 1));
       }
   }
    int update(int v, int tl, int tr, int idx, int diff) {
            tree.push_back({0, 0, tree[v].sum + diff});
            return tree.size() - 1;
       int tm = tl + (tr - tl) / 2;
       int 1Ch = tree[v].1Ch;
       int rCh = tree[v].rCh;
        if (idx <= tm)</pre>
            1Ch = update(1Ch, tl, tm, idx, diff);
        else
```

```
rCh = update(rCh, tm + 1, tr, idx, diff);
        tree.push_back({1Ch, rCh, tree[1Ch].sum + tree[rCh].sum});
        return tree.size() - 1;
    /* find kth smallest number among arr[l], arr[l+1], ..., arr[r]
     * k is 1-based, so find_kth(l,r,1) returns the min
    int query(int 1, int r, int k) const {
        assert(1 \le k \&\& k \le r - 1 + 1); //note this condition implies L \le R
        assert(0 <= 1 && r + 1 < (int)roots.size());</pre>
        return query(roots[1], roots[r + 1], -mx, mx, k);
    int query(int vl, int vr, int tl, int tr, int k) const {
        if (tl == tr)
            return tl:
        int tm = tl + (tr - tl) / 2;
        int left_count = tree[tree[vr].lCh].sum - tree[tree[v1].lCh].sum;
        if (left_count >= k) return query(tree[v1].lCh, tree[vr].lCh, t1, tm, k);
        return query(tree[v1].rCh, tree[vr].rCh, tm + 1, tr, k - left_count);
   }
};
```

```
% cat range_data_structures/sparseTable.h | ./misc/hash.sh
f-its-values-persistent-segment-tree
fcbc6d
```

Listing 19: Sparse Table

```
#pragma once
//usage:
// vector<long long> arr;
// ...
// sparseTable < long long > st(arr, [](auto x, auto y) { return min(x,y); });
//to also get index of min element, do:
// sparseTable < pair < long, int >> st(arr, [](auto x, auto y) { return min(x,y); });
//and initialize second to index. If there are multiple indexes of min element,
//it'll return the smallest (left-most) one
//status: tested on random inputs, also on https://judge.yosupo.jp/problem/staticrmq
template <class T>
struct sparseTable {
    vector<vector<T>> dp;
    function<T(const T&, const T&)> func;
    sparseTable(const vector<T>& arr, const function<T(const T&, const T&)>& _func) :
         \hookrightarrow dp(1, arr), func(_func) {
        int n = arr.size();
        for (int pw = 1, k = 1; pw * 2 <= n; pw *= 2, k++) {
            dp.emplace_back(n - pw * 2 + 1);
            for (int j = 0; j < (int)dp[k].size(); j++)</pre>
                dp[k][j] = func(dp[k - 1][j], dp[k - 1][j + pw]);
    }
    //inclusive range [l, r]
    T query(int 1, int r) const {
        assert(0 <= 1 && 1 <= r && r < (int)dp[0].size());
```

```
int lg = 31 - \_builtin\_clz(r - 1 + 1);
        return func(dp[lg][l], dp[lg][r - (1 << lg) + 1]);
   }
};
```

```
% cat range_data_structures/distinct_query.h | ./misc/hash.sh
```

116f83

Listing 20: Number Distinct Elements

```
#pragma once
//modified from
    \hookrightarrow https://cp-algorithms.com/data\_structures/segment\_tree.html*preserving-the-history-of-its-values-persistent-segment-tree
//tested on https://www.spoj.com/problems/DQUERY/ and stress tested
//works with negatives
struct persistentSegTree {
    struct Node {
       int 1Ch, rCh; //children, indexes into 'tree'
        int sum;
   };
   int sz;
    deque<Node> tree;
    vector<int> roots;
   persistentSegTree(const vector<int>& arr) : sz(arr.size() + 1) {
        tree.push_back({0, 0, 0}); //acts as null
       roots.push_back(0);
       map<int, int> lastIdx;
       for (int i = 0; i < (int)arr.size(); i++) {</pre>
            roots.push_back(update(roots.back(), 0, sz - 1, lastIdx[arr[i]], 1));
            lastIdx[arr[i]] = i + 1;
       }
   }
   int update(int v, int tl, int tr, int idx, int diff) {
       if (tl == tr) {
            tree.push_back({0, 0, tree[v].sum + diff});
            return tree.size() - 1;
        int tm = (tl + tr) / 2;
        int 1Ch = tree[v].1Ch;
       int rCh = tree[v].rCh;
       if (idx <= tm)
            1Ch = update(1Ch, tl, tm, idx, diff);
            rCh = update(rCh, tm + 1, tr, idx, diff);
        tree.push_back({lCh, rCh, tree[lCh].sum + tree[rCh].sum});
        return tree.size() - 1:
   }
   //returns number of distinct elements in range [l,r]
    int query(int 1, int r) const {
       return query(roots[1], roots[r + 1], 0, sz - 1, 1 + 1);
   }
```

```
int query(int vl, int vr, int tl, int tr, int idx) const {
        if (tree[vr].sum == 0 || idx <= tl)</pre>
            return 0;
        if (tr < idx)
            return tree[vr].sum - tree[vl].sum;
        int tm = (tl + tr) / 2;
        return query(tree[v1].1Ch, tree[vr].1Ch, tl, tm, idx) +
               query(tree[v1].rCh, tree[vr].rCh, tm + 1, tr, idx);
};
```

```
% cat range_data_structures/buckets.h | ./misc/hash.sh
db53a3
```

Listing 21: Buckets

```
#pragma once
//this code isn't the best. It's meant as a rough start for sqrt-decomposition, and to
     \hookrightarrow be (heavily) modified
//doesn't handle overflow
//status: tested on random inputs, also used in various problems
struct buckets {
    const int BUCKET_SIZE = 50;//TODO: change - small value for testing
    struct bucket {
        int sumLazy = 0;
        int sumBucket = 0;
        int 1, r; //inclusive range of bucket
        int len() const {
            return r - 1 + 1;
        }
    };
    vector<int> values:
    vector<bucket> _buckets;
    buckets(const vector<int>& initial) : values(initial) {
        int numBuckets = ((int) values.size() + BUCKET_SIZE - 1) / BUCKET_SIZE;
        _buckets.resize(numBuckets);
        for (int i = 0; i < numBuckets; i++) {</pre>
            _buckets[i].sumLazy = 0;
            _buckets[i].sumBucket = 0;
            _buckets[i].1 = i * BUCKET_SIZE;
            _buckets[i].r = min((i + 1) * BUCKET_SIZE, (int) values.size()) - 1;
            for (int j = _buckets[i].1; j <= _buckets[i].r; j++)</pre>
                _buckets[i].sumBucket += values[j];
    }
    void pushLazy(int bIdx) {
        bucket& b = _buckets[bIdx];
        if (!b.sumLazy) return;
        for (int i = b.l: i <= b.r: i++)
            values[i] += b.sumLazy;
        b.sumLazy = 0;
    }
```

```
//update range [L,R]
    void update(int L, int R, int diff) {
        int startBucket = L / BUCKET_SIZE;
        int endBucket = R / BUCKET_SIZE;
        if (startBucket == endBucket) { //range contained in same bucket case
            for (int i = L; i <= R; i++) {</pre>
                values[i] += diff;
                 _buckets[startBucket].sumBucket += diff;
            return;
        }
        for (int bIdx : {
                    startBucket, endBucket
                }) { //handle "endpoint" buckets
            bucket& b = _buckets[bIdx];
            for (int i = \max(b.1, L); i \le \min(b.r, R); i++) {
                values[i] += diff;
                b.sumBucket += diff;
            }
        }
        for (int i = startBucket + 1; i < endBucket; i++) {    //handle all n/B buckets</pre>
             \hookrightarrow in middle
            bucket& b = _buckets[i];
            b.sumLazy += diff;
            b.sumBucket += b.len() * diff;
    }
    //sum of range [L,R]
    int query(int L, int R) {
        int startBucket = L / BUCKET_SIZE;
        int endBucket = R / BUCKET_SIZE;
        if (startBucket == endBucket) { //range contained in same bucket case
            pushLazy(startBucket);
            int sum = 0:
            for (int i = L; i <= R; i++)</pre>
                sum += values[i]:
            return sum;
        }
        int sum = 0:
        for (int bIdx : {
                     startBucket, endBucket
                }) { //handle "endpoint" buckets
            bucket& b = _buckets[bIdx];
            pushLazy(bIdx);
            for (int i = max(b.1, L); i <= min(b.r, R); i++)</pre>
                sum += values[i];
        for (int i = startBucket + 1; i < endBucket; i++) //handle all n/B buckets in</pre>
            sum += _buckets[i].sumBucket;
        return sum;
   }
};
```

```
% cat range_data_structures/implicitLazySegTree.h | ./misc/hash.sh
a13f2e
```

Listing 22: Implicit Lazy Segment Tree

```
#pragma once
//status: stress tested && AC's on https://cses.fi/problemset/task/1144
//see TODO for lines of code which usually need to change (not a complete list)
const int N = 1.5e7; //T0D0
struct Node {
    long long val; //could represent max, sum, etc
    long long lazy;
    int 1Ch, rCh; // children, indexes into 'tree', -1 for null
} tree[N];
struct implicitLazySegTree {
    int NEW_NODE, rootL, rootR; //[rootL, rootR] defines range of root node; handles
         \hookrightarrow negatives
    implicitLazySegTree(int 1, int r) : NEW_NODE(0), rootL(1), rootR(r) {
        tree[NEW_NODE++] = {0, 0, -1, -1}; //TODO
    static long long combine(long long val_1, long long val_r) {
        return val_l + val_r; //TODO
    void apply(int v, int tl, int tr, long long add) {
        tree[v].val += (tr - tl + 1) * add; //TODO
        if (tl != tr) {
            tree[tree[v].1Ch].lazy += add; //TODO
            tree[tree[v].rCh].lazy += add;
    }
    void push(int v, int tl, int tr) {
        if (tl != tr && tree[v].lCh == -1) {
            assert(NEW_NODE + 1 < N);</pre>
            tree[v].1Ch = NEW_NODE;
            tree[NEW_NODE++] = \{0, 0, -1, -1\}; //TODO
            tree[v].rCh = NEW_NODE;
            tree [NEW_NODE++] = \{0, 0, -1, -1\};
        if (tree[v].lazy) {
            apply(v, tl, tr, tree[v].lazy);
            tree[v].lazy = 0;
    }
    //update range [l,r] with 'add'
    void update(int 1, int r, long long add) {
        update(0, rootL, rootR, 1, r, add);
    void update(int v, int tl, int tr, int l, int r, long long add) {
        push(v, tl, tr);
        if (tr < 1 || r < t1)
            return;
        if (1 <= t1 && tr <= r)</pre>
            return apply(v, tl, tr, add);
```

```
int tm = tl + (tr - tl) / 2;
        update(tree[v].1Ch, tl, tm, l, r, add);
        update(tree[v].rCh, tm + 1, tr, l, r, add);
        tree[v].val = combine(tree[tree[v].1Ch].val, tree[tree[v].rCh].val);
   }
    //query range [l,r]
    long long query(int 1, int r) {
        return query(0, rootL, rootR, 1, r);
    long long query(int v, int tl, int tr, int l, int r) {
        if (tr < 1 || r < t1)
            return 0; //TODO
        push(v, tl, tr);
        if (1 <= t1 && tr <= r)
            return tree[v].val;
        int tm = tl + (tr - tl) / 2;
        return combine(query(tree[v].lCh, tl, tm, l, r),
                       query(tree[v].rCh, tm + 1, tr, 1, r));
   }
};
```

```
% cat range_data_structures/persistentLazySegTree.h | ./misc/hash.sh
```

87eace

Listing 23: Persistent Lazy Segment Tree

```
#pragma once
//tested on https://codeforces.com/contest/707/problem/D
struct persistentLazySegTree {
   struct Node {
       int 1Ch, rCh;//children, indexes into 'tree'
       int sum;
       bool lazyTog;
   };
   int sz:
   deque<Node> tree;
   vector<int> roots;
   //implicit
   persistentLazySegTree(int _sz) : sz(_sz) {
       tree.push_back({0, 0, 0, 0}); //acts as null
       roots.push_back(0);
   }
   void push(int v, int tl, int tr) {
       if (tl != tr) {
            tree.push_back(tree[tree[v].1Ch]);
            tree[v].lCh = tree.size() - 1:
            tree.push_back(tree[tree[v].rCh]);
            tree[v].rCh = tree.size() - 1;
       if (tree[v].lazyTog) {
            tree[v].sum = (tr - tl + 1) - tree[v].sum;
            tree[v].lazyTog = false;
```

```
if (tl != tr) {
            tree[tree[v].lCh].lazyTog ^= 1;
            tree[tree[v].rCh].lazyTog ^= 1;
    }
}
void set(int idx, int new_val) {
    tree.push_back(tree[roots.back()]);//allocate top down
    roots.push_back(tree.size() - 1);
    set(roots.back(), 0, sz - 1, idx, new_val);
void set(int v, int tl, int tr, int idx, int new_val) {
    push(v, tl, tr);
    if (tr < idx || idx < tl)</pre>
        return;
    if (idx <= tl && tr <= idx) {</pre>
        tree[v].sum = new_val;
        return;
    }
    int tm = (tl + tr) / 2;
    int 1Ch = tree[v].1Ch:
    int rCh = tree[v].rCh;
    set(lCh, tl, tm, idx, new_val);
    set(rCh, tm + 1, tr, idx, new_val);
    tree[v].sum = tree[lCh].sum + tree[rCh].sum;
}
void toggleRange(int 1, int r) {
    tree.push_back(tree[roots.back()]);//allocate top down
    roots.push_back(tree.size() - 1);
    toggleRange(roots.back(), 0, sz - 1, 1, r);
void toggleRange(int v, int tl, int tr, int l, int r) {
    push(v, tl, tr);
    if (tr < 1 || r < t1)
        return:
    int 1Ch = tree[v].1Ch:
    int rCh = tree[v].rCh;
    if (1 <= t1 && tr <= r) {</pre>
        tree[v].sum = (tr - tl + 1) - tree[v].sum:
        if (tl != tr) {
            tree[lCh].lazyTog ^= 1;
            tree[rCh].lazyTog ^= 1;
        return;
    int tm = (tl + tr) / 2;
    toggleRange(1Ch, tl, tm, l, r);
    toggleRange(rCh, tm + 1, tr, l, r);
    tree[v].sum = tree[lCh].sum + tree[rCh].sum;
//let's use implementation trick described here

→ https://codeforces.com/blog/entry/72626
//so that we don't have to propogate lazy vals and thus we don't have to allocate
     \hookrightarrow new nodes
int query(int 1, int r) const {
    int version = roots.size() - 1;
    int root = roots[version];
    return query(root, 0, sz - 1, 1, r, tree[root].lazyTog);
```

```
}
    int query(int v, int tl, int tr, int l, int r, bool tog) const {
        if (v == 0 || tr < 1 || r < t1)</pre>
            return 0:
        if (1 <= t1 && tr <= r) {
            int sum = tree[v].sum;
            if (tree[v].lazyTog) sum = (tr - tl + 1) - sum;
        }
        int tm = (tl + tr) / 2;
        tog ^= tree[v].lazyTog;
        return query(tree[v].1Ch, tl, tm, l, r, tog) +
               query(tree[v].rCh, tm + 1, tr, 1, r, tog);
   }
};
```

```
% cat range_data_structures/MosAlg.h | ./misc/hash.sh
```

//status: not tested, but used in various problems

e7b0af

#include <bits/stdc++.h>

Listing 24: Mos Algorithm

```
using namespace std;
const int Max = 1e6 + 2;
int block, answer[Max], answerToQuery;
struct query {
    int 1, r, index;
};
bool cmp(query x, query y) {
    if (x.1 / block == y.1 / block) return x.r < y.r;</pre>
    return x.1 < y.1;</pre>
void add(int pos) {
void remove(int pos) {
int main() {
    int q;
    cin >> q;
    vector<query> queries(q);
    for (int i = 0; i < q; ++i) {
        cin >> queries[i].l >> queries[i].r;
        queries[i].index = i;
        answer[i] = 0;
    }
    sort(queries.begin(), queries.end(), cmp);
    int left = 0, right = 0;//store inclusive ranges, start at [0,0]
    add(0);
    answerToQuery = 0;
    for (auto& q : queries) {
        while (left > q.1) {
            left--:
```

```
add(left):
    }
    while (right < q.r) {</pre>
        right++;
        add(right);
    while (left < q.1) {</pre>
        remove(left);
        left++;
    while (right > q.r) {
        remove(right);
        right--;
    answer[q.index] = answerToQuery;
for (int i = 0; i < q; ++i) cout << answer[i] << '\n';</pre>
return 0;
```

% cat range_data_structures/persistentSegTree.h | ./misc/hash.sh 47efec

Listing 25: Persistent Segment Tree

```
#pragma once
//modified from
    \hookrightarrow https://cp-algorithms.com/data_structures/segment_tree.html#preserving-the-history
//tested on https://www.spoj.com/problems/PSEGTREE/ and
     \hookrightarrow https://cses.fi/problemset/task/1737/
struct persistentSegTree {
    struct Node {
        int 1Ch, rCh;//children, indexes into 'tree'
        long long sum;
   };
    deque<Node> tree;
    vector<int> roots;
    //implicit
    persistentSegTree(int _sz) : sz(_sz) {
        tree.push_back({0, 0, OLL}); //acts as null
        roots.push_back(0);
    persistentSegTree(const vector<long long>& arr) : sz(arr.size()) {
        tree.push_back({0, 0, OLL}); //acts as null
        roots.push_back(build(arr, 0, sz - 1));
    int build(const vector<long long>& arr, int tl, int tr) {
        if (t1 == tr) {
            tree.push_back({0, 0, arr[t1]});
            return tree.size() - 1;
        int tm = (tl + tr) / 2;
```

```
int 1Ch = build(arr, tl, tm);
        int rCh = build(arr, tm + 1, tr);
       tree.push_back({1Ch, rCh, tree[1Ch].sum + tree[rCh].sum});
       return tree.size() - 1:
   }
    void update(int version, int idx, long long diff) {
        roots.push_back(update(roots[version], 0, sz - 1, idx, diff));
    int update(int v, int tl, int tr, int idx, long long diff) {
       if (tl == tr) {
            tree.push_back({0, 0, tree[v].sum + diff});
            return tree.size() - 1;
       int tm = (tl + tr) / 2;
        int 1Ch = tree[v].1Ch;
       int rCh = tree[v].rCh;
       if (idx \le tm)
            1Ch = update(1Ch, tl, tm, idx, diff);
       else
            rCh = update(rCh, tm + 1, tr, idx, diff);
        tree.push_back({1Ch, rCh, tree[1Ch].sum + tree[rCh].sum});
       return tree.size() - 1:
   }
   long long query(int version, int 1, int r) const {
       return query(roots[version], 0, sz - 1, 1, r);
   long long query(int v, int tl, int tr, int l, int r) const {
        if (tree[v].sum == OLL || tr < 1 || r < t1)</pre>
            return OLL:
       if (1 <= t1 && tr <= r)
            return tree[v].sum:
       int tm = (tl + tr) / 2;
       return query(tree[v].1Ch, t1, tm, 1, r) +
               query(tree[v].rCh, tm + 1, tr, l, r);
   }
};
```

```
\% cat range_data_structures/mergeSortTree.h | ./misc/hash.sh
```

*9cea9*e

Listing 26: Merge Sort Tree

```
#pragma once

//status: stress-tested against persistent seg tree; used in various problems

struct MergeSortTree {
    struct Node {
        vector<int> vals;
        int 1, r;
    };

    vector<Node> tree;

    Node combineChildren(const Node& L, const Node& R) {
```

```
vector<int> par(L.vals.size() + R.vals.size());
        merge(L.vals.begin(), L.vals.end(), R.vals.begin(), R.vals.end(), par.begin());
        return Node{par, L.1, R.r};
    //There's no constructor 'SegmentTree(int size)' because how to initialize l,r in
         \hookrightarrow nodes without calling build?
    //the whole point of this constructor was to be simpler by not calling build
    MergeSortTree(const vector<int>& arr) {
        int n = arr.size(), size = 1:
        while (size < n) size <<= 1;</pre>
        size <<= 1:
        tree.resize(size);
        build(arr, 1, 0, n - 1);
    void build(const vector<int>& arr, int node, int start, int end) {
        if (start == end) {
            tree[node] = Node {
                vector<int>{arr[start]},
                start,
                end
            };
        } else {
            int mid = (start + end) / 2;
            build(arr, 2 * node, start, mid);
            build(arr, 2 * node + 1, mid + 1, end);
            tree[node] = combineChildren(tree[2 * node], tree[2 * node + 1]);
    //inclusive range: [l,r]
    int query(int 1, int r, int x) {
        return query(1, 1, r, x);
    int query(int node, int 1, int r, int x) {
        int start = tree[node].1, end = tree[node].r;
        if (r < start || end < 1) return 0;</pre>
        if (1 <= start && end <= r) {
            vector<int>& v = tree[node].vals;
            return lower_bound(v.begin(), v.end(), x) - v.begin();
        return query(2 * node, 1, r, x) + query(2 * node + 1, 1, r, x);
    }
};
```

```
% cat range_data_structures/segTree.h | ./misc/hash.sh
```

Listing 27: Segment Tree

```
#pragma once

//status: tested on random inputs

const long long inf = 1e18;

struct segTree {
    struct Node {
```

```
long long sum, mx, mn;
    long long lazy;
    int 1, r;
    int len() const {
        return r - 1 + 1;
    //returns 1 + (# of nodes in left child's subtree)
    //https://cp-algorithms.com/data_structures/segment_tree.html#memory-efficient-implementat/com/date range [l,r] with 'add'
    int rCh() const {
        return ((r - 1) & ^1) + 2;
    }
};
vector<Node> tree;
//There's no constructor 'seqTree(int size)' because how to initialize l,r in nodes
     \hookrightarrow without calling build?
//the whole point of 'segTree(int size)' was to be simpler by not calling build
segTree(const vector<long long>& arr) : tree(2 * (int) arr.size() - 1) {
    build(arr, 0, 0, (int) arr.size() - 1);
void build(const vector<long long>& arr, int v, int tl, int tr) {
    if (tl == tr) {
        tree[v] = {
            arr[t1],
            arr[t1],
            arr[t1],
            tl,
            tr
        };
    } else {
        int tm = tl + (tr - tl) / 2;
        build(arr, v + 1, tl, tm):
        build(arr, v + 2 * (tm - tl + 1), tm + 1, tr);
        tree[v] = combine(tree[v + 1], tree[v + 2 * (tm - tl + 1)]);
}
Node combine(const Node& L. const Node& R) {
    return {
        L.sum + R.sum,
        max(L.mx, R.mx),
        min(L.mn, R.mn),
        L.1,
        R.r
    };
}
//what happens when 'add' is applied to every index in range [tree[v].l, tree[v].r]?
void apply(int v, long long add) {
    tree[v].sum += tree[v].len() * add;
    tree[v].mx += add;
    tree[v].mn += add;
    if (tree[v].len() > 1) {
        tree[v + 1].lazy += add;
        tree[v + tree[v].rCh()].lazy += add;
}
```

```
void push(int v) {
        if (tree[v].lazy) {
            apply(v, tree[v].lazy);
            tree[v].lazy = 0;
        }
    }
    void update(int 1, int r, long long add) {
        update(0, 1, r, add);
    void update(int v, int 1, int r, long long add) {
        if (tree[v].r < 1 || r < tree[v].1)</pre>
        if (1 <= tree[v].1 && tree[v].r <= r)</pre>
            return apply(v, add);
        update(v + 1, 1, r, add);
        update(v + tree[v].rCh(), 1, r, add);
        tree[v] = combine(tree[v + 1], tree[v + tree[v].rCh()]);
    //range [l,r]
    Node query(int 1, int r) {
        return query(0, 1, r);
    Node query(int v, int 1, int r) {
        if (tree[v].r < 1 || r < tree[v].1)</pre>
            return {0, -inf, inf, 0, 0, 0};
        push(v);
        if (1 <= tree[v].1 && tree[v].r <= r)</pre>
            return tree[v];
        return combine(query(v + 1, 1, r),
                        query(v + tree[v].rCh(), 1, r));
    }
};
```

```
% cat range_data_structures/fenwickTree.h | ./misc/hash.sh
836i35
```

Listing 28: Fenwick Tree

```
#pragma once

//status: tested on random inputs; also tested on
//https://judge.yosupo.jp/problem/point_add_range_sum, lower_bound tested on
//https://judge.yosupo.jp/problem/predecessor_problem

template<class T>
struct fenwickTree {
    vector<T> bit;
    fenwickTree(int n) : bit(n, 0) {}
    fenwickTree(const vector<T>& a) : bit(a.size()) {
        if (a.empty()) return;
        bit[0] = a[0];
        for (int i = 1; i < (int) a.size(); i++)
            bit[i] = bit[i - 1] + a[i];</pre>
```

```
for (int i = (int) a.size() - 1; i > 0; i--) {
            int lower_i = (i & (i + 1)) - 1;
            if (lower_i >= 0)
                bit[i] -= bit[lower i]:
    }
    void update(int idx, const T& d) {
        for (; idx < (int) bit.size(); idx = idx | (idx + 1))
            bit[idx] += d;
    T sum(int r) const {
        T ret = 0:
        for (; r \ge 0; r = (r \& (r + 1)) - 1)
            ret += bit[r]:
        return ret;
    }
    T sum(int 1, int r) const {
        return sum(r) - sum(l - 1);
    //Returns min pos such that sum of [0, pos] >= sum
    //Returns bit.size() if no sum is >= sum, or -1 if empty sum is.
    //Doesn't work with negatives (since it's greedy), counterexample: array: {1, -1},
         \hookrightarrow sum: 1. this returns 2. but should return 0
    int lower_bound(T sum) const {
        if (sum <= 0) return -1;</pre>
        int pos = 0;
        for (int pw = 1 << (31 - _builtin_clz(bit.size() | 1)); pw; pw >>= 1) {
            if (pos + pw <= (int)bit.size() && bit[pos + pw - 1] < sum)</pre>
                pos += pw, sum -= bit[pos - 1];
        }
        return pos;
    }
};
//status: tested on random inputs
template<class T>
struct rangeUpdatesAndPointQueries {
    fenwickTree<T> ft;
    rangeUpdatesAndPointQueries(int n) : ft(n) {}
    rangeUpdatesAndPointQueries(const vector<T>& arr) : ft(init(arr)) {}
    fenwickTree<T> init(vector<T> arr/*intentional pass by value*/) {
        for (int i = (int) arr.size() - 1; i >= 1; i--)
            arr[i] -= arr[i - 1];
        return fenwickTree<T> (arr);
    //add 'add' to inclusive range [l, r]
    void updateRange(int 1, int r, const T& add) {
        ft.update(1, add);
        if (r + 1 < (int) ft.bit.size())</pre>
            ft.update(r + 1, -add);
    //get value at index 'idx'
    T queryIdx(int idx) const {
        return ft.sum(idx);
    }
};
```

Listing 29: STRINGS

```
% cat strings/suffix_array.h | ./misc/hash.sh
46840a
```

Listing 30: Suffix Array

```
#pragma once
//modified from here: https://judge.yosupo.jp/submission/37410
//status: tested on https://judge.yosupo.jp/problem/suffixarray
// SA-IS, linear-time suffix array construction
// Reference:
// G. Nong, S. Zhang, and W. H. Chan,
// Two Efficient Algorithms for Linear Time Suffix Array Construction
vector<int> sa_is(const T& s, int upper/*max element of 's'; for std::string, pass in
     \hookrightarrow 255*/) {
    int n = (int) s.size();
    if (n == 0) return {};
    if (n == 1) return {0};
    if (n == 2) {
        if (s[0] < s[1]) {
            return {0, 1};
        } else {
            return {1, 0};
    }
    vector<int> sa(n);
    vector<bool> ls(n);
    for (int i = n - 2; i >= 0; i--)
        ls[i] = (s[i] == s[i + 1]) ? ls[i + 1] : (s[i] < s[i + 1]);
    vector<int> sum_l(upper + 1), sum_s(upper + 1);
    for (int i = 0; i < n; i++) {</pre>
        if (!ls[i])
            sum s[s[i]]++:
        else
            sum_l[s[i] + 1]++;
    for (int i = 0; i <= upper; i++) {</pre>
        sum_s[i] += sum_l[i];
        if (i < upper) sum_l[i + 1] += sum_s[i];</pre>
    vector<int> buf(upper + 1);
    auto induce = [&](const vector<int>& lms) {
        fill(sa.begin(), sa.end(), -1);
        fill(buf.begin(), buf.end(), 0);
        copy(sum_s.begin(), sum_s.end(), buf.begin());
        for (auto d : lms) {
            if (d == n) continue;
            sa[buf[s[d]]++] = d;
        copy(sum_l.begin(), sum_l.end(), buf.begin());
        sa[buf[s[n-1]]++] = n-1;
        for (int i = 0: i < n: i++) {
            int v = sa[i];
            if (v >= 1 && !ls[v - 1])
                sa[buf[s[v - 1]] ++] = v - 1;
```

```
}
    copy(sum_l.begin(), sum_l.end(), buf.begin());
    for (int i = n - 1; i >= 0; i--) {
        int v = sa[i]:
        if (v >= 1 \&\& ls[v - 1])
            sa[--buf[s[v - 1] + 1]] = v - 1;
    }
};
vector < int > lms_map(n + 1, -1);
int m = 0;
for (int i = 1; i < n; i++) {
    if (!ls[i - 1] && ls[i])
        lms_map[i] = m++;
}
vector<int> lms:
lms.reserve(m);
for (int i = 1; i < n; i++) {</pre>
    if (!ls[i - 1] && ls[i])
        lms.push_back(i);
}
induce(lms);
if (m) {
    vector<int> sorted_lms;
    sorted_lms.reserve(m);
    for (int v : sa) {
        if (lms_map[v] != -1) sorted_lms.push_back(v);
    vector<int> rec_s(m);
    int rec_upper = 0;
    rec_s[lms_map[sorted_lms[0]]] = 0;
    for (int i = 1; i < m; i++) {</pre>
        int l = sorted_lms[i - 1], r = sorted_lms[i];
        int end_l = (lms_map[l] + 1 < m) ? lms[lms_map[l] + 1] : n;</pre>
        int end_r = (lms_map[r] + 1 < m) ? lms[lms_map[r] + 1] : n;
        bool same = true:
        if (end 1 - 1 != end r - r)
            same = false:
        else {
            while (1 < end_1) {
                if (s[1] != s[r])
                    break:
                1++;
                r++;
            }
            if (1 == n || s[1] != s[r]) same = false;
        if (!same) rec_upper++;
        rec_s[lms_map[sorted_lms[i]]] = rec_upper;
    }
    auto rec_sa =
        sa_is(rec_s, rec_upper);
    for (int i = 0; i < m; i++)
        sorted_lms[i] = lms[rec_sa[i]];
    induce(sorted lms):
}
return sa;
```

396173

Listing 31: Longest Common Prefix Array

```
#pragma once
//modified from here: https://judge.yosupo.jp/submission/37410
//status: tested on https://judge.yosupo.jp/problem/number_of_substrings (answer = (n *
     \hookrightarrow (n+1) / 2) - (sum \ of \ LCP \ array))
// Reference:
// T. Kasai, G. Lee, H. Arimura, S. Arikawa, and K. Park,
// Linear-Time Longest-Common-Prefix Computation in Suffix Arrays and Its
// Applications
template<class T>
vector<int> lcp_array(const T& s, const vector<int>& sa) {
    int n = s.size(), k = 0;
    vector<int> lcp(n, 0);
    vector<int> rank(n, 0);
    for (int i = 0; i < n; i++) rank[sa[i]] = i;</pre>
    for (int i = 0; i < n; i++, k ? k-- : 0) {
        if (rank[i] == n - 1) {
            k = 0;
            continue;
        int j = sa[rank[i] + 1];
        while (i + k < n \&\& j + k < n \&\& s[i + k] == s[j + k]) k++;
        lcp[rank[i]] = k;
    }
    return lcp;
```

% cat strings/rotationally_equivalent.h | ./misc/hash.sh

7a608b

Listing 32: Rotational Equivalence

```
#pragma once

// Checks if two arrays are rotationally equivalent
// uses KMP with doubling trick
// usage:
// string s1, s2;
// ...
// rot_eq(s1, s2)
// or
// vector<int> arr1, arr2;
// ...
// rot_eq(arr1, arr2)
//
//status: tested on random inputs, also on https://open.kattis.com/problems/maze

template <class T>
bool rot_eq(const T& a, const T& b) {
    if (a.size() != b.size()) return false;
    if (a.empty()) return true;
```

```
int n = a.size();
vector<int> fail(n + 1, 0);
auto update = [&](int val, int& p) -> void {
    while (p && val != a[p]) p = fail[p];
    if (val == a[p]) p++;
};
for (int i = 1, p = 0; i < n; i++) {
    update(a[i], p);
    fail[i + 1] = p;
}
for (int i = 0, p = 0; i < 2 * n; i++) {
    update(b[i % n], p);
    if (p == n) return true;
}
return false;</pre>
```

```
% cat strings/rollingHash.h | ./misc/hash.sh
```

// works for -2e9 - 1000 <= arr[i] <= 2e9 + 1000

8e807d

//usage:
// string s;

// Hash h(s);

// vector<int> arr;
// Hash h(arr);

Listing 33: Rolling Hash

```
//status: tested on random inputs, on https://judge.yosupo.jp/problem/zalgorithm, and on

→ https://judge.yosupo.jp/problem/enumerate_palindromes

#pragma once
#include "../misc/random.h"
const unsigned mod = 4294967087; // largest prime p < UINT_MAX such that (p-1)/2 is also
    \hookrightarrow prime
const int mx = 2e9 + 1000;
const vector<unsigned> bases {
   getRand<unsigned>(2u * mx + 2, mod - 1),
    getRand<unsigned>(2u * mx + 2, mod - 1),
    getRand<unsigned>(2u * mx + 2, mod - 1)
template <class T>
    vector<vector<unsigned>> prefix, powB;
   Hash(const T& s):
        prefix(bases.size(), vector<unsigned> (s.size() + 1, 0)),
        powB(bases.size(), vector<unsigned> (s.size() + 1, 1)) {
        for (auto val : s) assert(-mx <= val && val <= mx);</pre>
        for (int i = 0; i < (int) bases.size(); i++) {</pre>
            for (int j = 0; j < (int) s.size(); j++) {</pre>
                powB[i][j + 1] = 1ULL * powB[i][j] * bases[i] % mod;
```

```
prefix[i][j + 1] = (1ULL * prefix[i][j] * bases[i] + s[j] + mx + 1) %
            }
        }
    }
    void debugCollisionProbability() const {
        // (1 - (str_len / mod) ^ #bases) ^ #comparisons
        auto getProb = [&](long long num_comparisons) -> double {
            return pow(1 - pow(powB[0].size() / double(mod), bases.size()),

    num_comparisons);
        };
        long long num_comparisons = 1e5;
        cerr << fixed << setprecision(10) << "Probability of **no** collisions when</pre>
              << num_comparisons << " comparisons is ~ " << getProb(num_comparisons) <<
        long long k = 1e5;
        cerr << fixed << setprecision(10) << "Probability that " << k << " unique</pre>
             \hookrightarrow strings have "
              << k << " unique hashes (if storing hashes in a std::set) is "
              << getProb(k * (k - 1) / 2) << endl;
    }
    //returns hashes of substring/subarray [L,R] inclusive, one hash per base
    vector<unsigned> operator()(int L, int R) const {
        assert(0 <= L && L <= R && R + 1 < (int) prefix[0].size());
        vector<unsigned> res(bases.size());
        for (int i = 0; i < (int) bases.size(); i++) {</pre>
            long long x = 1LL * prefix[i][R + 1] + mod - 1ULL * prefix[i][L] * powB[i][R
                 \hookrightarrow - L + 1] % mod:
            res[i] = x > = mod ? x - mod : x;
        return res;
};
```

```
% cat strings/kmp.h | ./misc/hash.sh
```

509da3

Listing 34: KMP

```
for (int i = 1, p = 0; i < (int) needle.size(); i++) {</pre>
            update(needle[i], p);
            prefixFunction[i + 1] = p;
   };
    // if haystack = "bananas"
    // needle = "ana"
    // then we find 2 matches:
    // bananas
    // _ana___
    // ___ana_
    // 0123456 (indexes)
    // and KMP_Match::find returns {1,3} - the indexes in kaystack where
    // each match starts.
    // You can also pass in false for "all" and KMP_Match::find will only
    // return the first match: {1}. Useful for checking if there exists
    // some match:
    // KMP_Match::find(<haystack>, false).size() > 0
    vector<int> find(const T& haystack, bool all = true) const {
        vector<int> matches;
        for (int i = 0, p = 0; i < (int) haystack.size(); i++) {</pre>
            update(haystack[i], p);
            if (p == (int) needle.size()) {
                matches.push_back(i - (int) needle.size() + 1);
                if (!all) return matches;
                p = prefixFunction[p];
        return matches;
   }
private:
    void update(char val, int& p) const {
        while (p && val != needle[p]) p = prefixFunction[p];
        if (val == needle[p]) p++;
   }
    vector<int> prefixFunction;
    T needle:
};
```

```
% cat strings/trie.h | ./misc/hash.sh
```

45e7a8

Listing 35: Trie

```
#pragma once

//status: not tested, but used on various problems
//intended to be a base template and to be modified

const int K = 26;//character size

struct trie {

   struct node {
```

```
bool leaf = 0:
        int next[K], id, p = -1;
        char pch;
        node(int _p = -1, char ch = '#') : p(_p), pch(ch) {
            fill(next, next + K, -1);
       }
   };
    vector<node> t;
    trie() : t(1) {}
    void add_string(const string& s, int id) {
        int c = 0;
       for (char ch : s) {
            int v = ch - 'a';
            if (t[c].next[v] == -1) {
                t[c].next[v] = t.size();
                t.emplace_back(c, ch);
            c = t[c].next[v];
       }
        t[c].leaf = 1;
        t[c].id = id;
   }
    void remove_string(const string& s) {
        int c = 0;
        for (char ch : s) {
            int v = ch - 'a';
            if (t[c].next[v] == -1)
                return;
            c = t[c].next[v];
        t[c].leaf = 0;
   }
    int find_string(const string& s) {
        int c = 0;
        for (char ch : s) {
            int v = ch - a;
            if (t[c].next[v] == -1)
                return -1;
            c = t[c].next[v];
        if (!t[c].leaf) return -1;
        return t[c].id;
};
```

% cat strings/string_queries.h | ./misc/hash.sh

85ec3a

Listing 36: Longest Common Prefix Query

```
#pragma once
#include "suffix_array.h"
```

```
#include "longest_common_prefix.h"
#include "../range_data_structures/sparseTable.h"
//status: tested on random inputs, and on

→ https://open.kattis.com/problems/automatictrading
//computes suffix array, lcp array, and then sparse table over lcp array
//0(n log n)
struct str_queries {
    str_queries(const string& s) : sa(sa_is(s, 255)), inv_sa(s.size()), lcp(lcp_array(s,
         \hookrightarrow sa)), st(lcp, [](int x, int y) {
        return min(x, y);
   }) {
        for (int i = 0; i < (int) s.size(); i++)</pre>
            inv sa[sa[i]] = i:
    }
    //length of longest common prefix of suffixes s[idx1 ... n-1], s[idx2 ... n-1],
         \hookrightarrow 0-based indexing
    //You can check if two substrings s[L1..R1], s[L2..R2] are equal in O(1) by:
    //R2-L2 == R1-L1 \& U longest_common_prefix(L1, L2) >= R2-L2+1
    int longest_common_prefix(int idx1, int idx2) const {
        if (idx1 == idx2) return (int) inv_sa.size() - idx1;
        idx1 = inv_sa[idx1];
        idx2 = inv_sa[idx2];
        if (idx1 > idx2) swap(idx1, idx2);
        return st.query(idx1, idx2 - 1);
    //returns true if suffix s[idx1 \dots n-1] < s[idx2 \dots n-1]
    //(so\ false\ if\ idx1 == idx2)
    bool less(int idx1, int idx2) const {
        return inv_sa[idx1] < inv_sa[idx2];</pre>
    vector<int> sa, inv_sa, lcp;
    sparseTable<int> st;
```

Listing 37: MATH

```
% cat math/exp_mod.h | ./misc/hash.sh
```

530333

Listing 38: BIN EXP MOD

```
#pragma once

//status: tested on random inputs, and used in misc. problems

//returns a^pw % mod in O(log(pw))
long long fastPow(long long a, long long pw, int mod) {
    long long res = 1;
    a %= mod;
    while (pw > 0) {
        if (pw & 1) res = (res * a) % mod;
    }
}
```

```
a = (a * a) % mod;
    pw >>= 1;
}
return res;
}
```

```
% cat math/fib.h | ./misc/hash.sh

1c28e8
```

Listing 39: Fibonacci

```
% cat math/matrixMultPow.h | ./misc/hash.sh
a37276
```

Listing 40: Matrix Mult and Pow

```
#pragma once
//status: not tested, but used on misc. problems
const int mod = 1e9 + 7;
vector<vector<int>> mult(const vector<vector<int>>& a, const vector<vector<int>>& b) {
    if (a.size() == 0) return {};
    if (a[0].size() == 0) return {};
    if (b.size() == 0) return {};
    if (b[0].size() == 0) return {};
    if (a[0].size() != b.size()) return {};
    int resultRow = a.size(), resultCol = b[0].size(), n = a[0].size();
    vector<vector<int>> product(resultRow, vector<int> (resultCol, 0));
    for (int i = 0; i < resultRow; ++i) {</pre>
        for (int k = 0; k < n; ++k) {
            for (int j = 0; j < resultCol; ++j)</pre>
                product[i][j] = (product[i][j] + 1LL * a[i][k] * b[k][j]) % mod;
    return product;
vector<vector<int>> power(vector<vector<int>> matrix, int b) {
```

vector<vector<int>> res(matrix.size(), vector<int> (matrix.size(), 0));

```
for (int i = 0; i < (int) matrix.size(); i++)
    res[i][i] = 1;
while (b > 0) {
    if (b % 2 == 1)
        res = mult(res, matrix);
    matrix = mult(matrix, matrix);
    b /= 2;
}
return res;
```

```
% cat math/n_choose_k_mod.h | ./misc/hash.sh
fale4f
```

Listing 41: N Choose K MOD

```
#pragma once
//status: tested on random inputs
//for mod inverse
#include "exp_mod.h"
// usage:
       NchooseK nk(n+1, 1e9+7) to use 'choose', 'inv' with inputs <= n
// or:
       NchooseK nk(mod, mod) to use 'chooseWithLucasTheorem'
struct NchooseK {
    // 'factSz' is the size of the factorial array, so only call 'choose', 'inv' with n
   NchooseK(int factSz, int currMod) : mod(currMod), fact(factSz, 1), invFact(factSz) {
        //this implimentation of doesn't work if factSz > mod because n! % mod = 0 when
             \hookrightarrow n >= mod. So 'invFact' array will be all 0's
        assert(factSz <= mod):</pre>
        //assert mod is prime. mod is intended to fit inside an int so that
        //multiplications fit in a longlong before being modded down. So this
        //will take sqrt(2~31) time
        assert(mod >= 2):
        for (int i = 2: i * i \le mod: i++)
            assert(mod % i);
        for (int i = 1; i < factSz; i++)</pre>
            fact[i] = 1LL * fact[i - 1] * i % mod;
        invFact.back() = fastPow(fact.back(), mod - 2, mod);
        for (int i = factSz - 2; i \ge 0; i--)
            invFact[i] = 1LL * invFact[i + 1] * (i + 1) % mod;
   }
    //classic n choose k
    //fails when n \ge mod
    int choose(int n, int k) const {
        if (k < 0 \mid | k > n) return 0;
        //now we know 0 <= k <= n so 0 <= n
        return 1LL * fact[n] * invFact[k] % mod * invFact[n - k] % mod;
   }
    //lucas theorem to calculate n choose k in O(\log(k))
    //need to calculate all factorials in range [0,mod), so O(mod) time&space, so need
        \hookrightarrow smallish prime mod (< 1e6 maybe)
```

```
% cat math/partition.h | ./misc/hash.sh
456850
```

Listing 42: Partition

```
#pragma once
//status: not tested
struct partitionFunction {
    vector<long long> remember;
    //The number of ways you can add to a number
    long long getPartitionsModM(int n, int m) {
        if (n < 0) return 0;
        if (n == 0) return 1;
        if ((int) remember.size() <= n) remember.resize(n + 1, -1);</pre>
        if (remember[n] != -1) return remember[n];
        long long sum = 0;
        long long val = 1;
        for (int i = 1; val <= n; i++) {
            long long multiply = 1;
            if (i % 2 == 0) multiply = -1;
            val = ((3LL * i * i) + i) / 2;
            sum += getPartitionsModM(n - val, m) * multiply % m;
            val = ((3LL * i * i) - i) / 2;
            sum += getPartitionsModM(n - val, m) * multiply % m;
            sum %= m;
            if (sum < 0) sum += m;</pre>
        return remember[n] = sum % m;
};
```

```
% cat math/primeSieveMobius.h | ./misc/hash.sh
e657cb
```

Listing 43: Prime Sieve Mobius

```
#pragma once
//status: not tested, but used on various problems
//mobius[i] = 0 iff there exists a prime p s.t. i\%(p^2)=0
//mobius[i] = -1 iff i has an odd number of distinct prime factors
//mobius[i] = 1 iff i has an even number of distinct prime factors
const int N = 2e6 + 10;
int mobius[N];
void calcMobius() {
   mobius[1] = 1;
    for (int i = 1; i < N; ++i) {
        for (int j = i + i; j < N; j += i)
            mobius[j] -= mobius[i];
   }
}
int minPrime[N];
void calcSeive() {
   fill(minPrime, minPrime + N, N);
   for (int i = N - 1; i \ge 2; --i) {
        for (int j = i; j < N; j += i)
            minPrime[i] = i:
   }
}
```

% cat math/solve linear mod.h | ./misc/hash.sh

971790

#pragma once

Listing 44: Solve Linear Equations MOD

```
//for mod inverse
#include "exp_mod.h"
struct matrixInfo {
    int rank, det;
    vector<int> x;
};
//Solves\ A * x = b\ under\ prime\ mod.
//A is a n (rows) by m (cols) matrix, b is a length n column vector, x is a length m
    \hookrightarrow column vector.
//assumes n,m >= 1, else RTE
//Returns rank of A, determinant of A, and x (solution vector to A * x = b). x is empty
     \hookrightarrow if no solution. If multiple solutions, an arbitrary one is returned.
//Leaves A in reduced row echelon form (unlike kactl).
//O(n * m * min(n.m))
//status: tested on https://judge.yosupo.jp/problem/system_of_linear_equations and

→ https://judge.yosupo.jp/problem/matrix_det

matrixInfo solve_linear_mod(vector<vector<int>>& A, vector<int>& b, const int mod) {
    assert(A.size() == b.size());
    int n = A.size(), m = A[0].size(), rank = 0, det = 1;
    //start of row reduce
    for (int col = 0: col < m && rank < n: ++col) {
        //find arbitrary pivot and swap pivot to current row
```

```
for (int i = rank; i < n; ++i)
        if (A[i][col] != 0) {
            if (rank != i) det = det == 0 ? det : mod - det;
            swap(A[i], A[rank]);
            swap(b[i], b[rank]);
            break;
    if (A[rank][col] == 0) {
        det = 0:
        continue:
    det = (1LL * det * A[rank][col]) % mod;
    //make pivot 1 by dividing row by inverse of pivot
    const int aInv = fastPow(A[rank][col], mod - 2, mod);
    for (int j = 0; j < m; ++j)
        A[rank][j] = (1LL * A[rank][j] * aInv) % mod;
    b[rank] = (1LL * b[rank] * aInv) % mod;
    //zero-out all numbers above & below pivot
    for (int i = 0; i < n; ++i)
        if (i != rank && A[i][col] != 0) {
            const int val = A[i][col];
            for (int j = 0; j < m; ++j) {
                A[i][j] -= 1LL * A[rank][j] * val % mod;
                if (A[i][j] < 0) A[i][j] += mod;</pre>
            b[i] -= 1LL * b[rank] * val % mod;
            if (b[i] < 0) b[i] += mod;</pre>
        }
    ++rank;
//end of row reduce, start of extracting answer ('x') from 'A' and 'b'
assert(rank <= min(n, m));</pre>
//check if solution exists
for (int i = rank; i < n; i++) {</pre>
    if (b[i] != 0) return {rank, det, {} }; //no solution exists
//initialize solution vector ('x')
vector<int> x(m, 0);
for (int i = 0, j = 0; i < rank; i++) {
    while (A[i][j] == 0) j++; //find pivot column
    assert(A[i][j] == 1);
    x[j] = b[i];
return {rank, det, x};
```

```
% cat math/sumFloorArithSeries.h | ./misc/hash.sh
b12ad8
```

Listing 45: Sum Floors of Arithmetic Series

```
#pragma once

//status: used on https://open.kattis.com/problems/itsamodmodmodworld

//computes:
//[p/q] + [2p/q] + [3p/q] + ... + [np/q]
//(p, q, n are natural numbers)
```

```
//[x] = floor(x)
long long cnt(long long p, long long q, long long n) {
    long long t = \_gcd(p, q);
    p = p / t;
    q = q / t;
    long long s = 0;
    long long z = 1;
    while ((q > 0) \&\& (n > 0)) {
        //(point A)
        t = p / q;
        s += z * t * n * (n + 1) / 2;
        p -= q * t;
        //(point B)
        t = n / q;
        s += z * p * t * (n + 1) - z * t * (p * q * t + p + q - 1) / 2;
        n -= q * t;
        //(point C)
        t = n * p / q;
        s += z * t * n;
        n = t;
        swap(p, q);
        z = -z:
    return s;
```

```
% cat math/sumOfKthPowers.h | ./misc/hash.sh
```

5cdedb

#pragma once

Listing 46: Sum of Kth Powers

```
//status: not tested, but used on misc. problems
#define MAX 1000010
#define MOD 100000007
//Faulhaber'the sum of the k-th powers of the first n positive integers
//1^k + 2^k + 3^k + 4^k + \dots + n^k
//0(k*log(k))
//Usage: lgr::lagrange(n, k)
namespace lgr {
short factor[MAX];
int P[MAX], S[MAX], ar[MAX], inv[MAX];
inline int expo(int a, int b) {
   int res = 1;
    while (b) {
        if (b & 1) res = (long long) res * a % MOD;
        a = (long long) a * a % MOD;
        b >>= 1:
   }
   return res;
```

```
int lagrange(long long n, int k) {
    if (!k) return (n % MOD);
    int i, j, x, res = 0;
    if (!inv[0]) {
        for (i = 2, x = 1; i < MAX; i++) x = (long long) x * i % MOD;
        inv[MAX - 1] = expo(x, MOD - 2);
        for (i = MAX - 2; i \ge 0; i--) inv[i] = ((long long) inv[i + 1] * (i + 1)) % MOD;
    }
   k++;
    for (i = 0; i <= k; i++) factor[i] = 0;
    for (i = 4; i <= k; i += 2) factor[i] = 2;
    for (i = 3; (i * i) \le k; i += 2) {
        if (!factor[i]) {
            for (j = (i * i), x = i << 1; j <= k; j += x)
                factor[j] = i;
    for (ar[1] = 1, ar[0] = 0, i = 2; i \le k; i++) {
        if (!factor[i]) ar[i] = expo(i, k - 1);
        else ar[i] = ((long long) ar[factor[i]] * ar[i / factor[i]]) % MOD;
    for (i = 1; i <= k; i++) {
        ar[i] += ar[i - 1];
        if (ar[i] >= MOD) ar[i] -= MOD;
    if (n <= k) return ar[n];</pre>
    P[0] = 1, S[k] = 1;
    for (i = 1; i <= k; i++) P[i] = ((long long) P[i - 1] * ((n - i + 1) % MOD)) % MOD;
    for (i = k - 1; i \ge 0; i--) S[i] = ((long long) S[i + 1] * ((n - i - 1) % MOD)) %
    for (i = 0; i <= k; i++) {</pre>
        x = (long long) ar[i] * P[i] % MOD * S[i] % MOD * inv[k - i] % MOD * inv[i] %
             \hookrightarrow MOD:
        if ((k - i) & 1) {
            res -= x:
            if (res < 0) res += MOD;</pre>
        } else {
            res += x;
            if (res >= MOD) res -= MOD;
    return (res % MOD);
```

```
% cat math/totient.h | ./misc/hash.sh
```

Listing 47: Euler's Totient Phi Function

```
#pragma once

//status: tested on n in range [1, 800]

// Euler's totient function counts the positive integers
// up to a given integer n that are relatively prime to n.
```

Listing 48: MAX FLOW

```
% cat maxflow/dinic.h | ./misc/hash.sh

d421fb
```

Listing 49: Dinic

```
#pragma once
//status: no tests, but used in various problems
struct maxflow {
public:
    typedef long long 11;
    11 n, s, t;
    maxflow(int _n, int _s, int _t) : n(_n), s(_s), t(_t), d(n), ptr(n), q(n), g(n) {}
    void addedge(ll a, ll b, ll cap) {
        edgeMap[a * n + b] = e.size();
        edge e1 = { a, b, cap, 0 };
        edge e2 = \{ b, a, 0, 0 \};
        g[a].push_back((ll) e.size());
        e.push_back(e1);
        g[b].push_back((ll) e.size());
        e.push_back(e2);
   }
   11 getflow() {
        11 \text{ flow} = 0;
        for (;;) {
            if (!bfs()) break;
            ptr.assign(ptr.size(), 0);
            while (ll pushed = dfs(s, inf))
                flow += pushed;
        }
        return flow;
   }
   11 getFlowForEdge(ll a, ll b) {
        return e[edgeMap[a * n + b]].flow;
    }
private:
    const ll inf = 1e18;
    struct edge {
```

```
ll a, b, cap, flow;
    unordered_map<int, 11> edgeMap;
    vector<ll> d, ptr, q;
    vector<edge> e;
    vector<vector<ll>>> g;
    bool bfs() {
        11 qh = 0, qt = 0;
        q[qt++] = s;
        d.assign(d.size(), -1);
        d[s] = 0;
        while (qh < qt && d[t] == -1) {
            11 v = q[qh++];
            for (size_t i = 0; i < g[v].size(); ++i) {</pre>
                ll id = g[v][i],
                    to = e[id].b;
                if (d[to] == -1 && e[id].flow < e[id].cap) {</pre>
                     q[qt++] = to;
                     d[to] = d[v] + 1;
                }
        }
        return d[t] != -1;
    11 dfs(11 v, 11 flow) {
        if (!flow) return 0;
        if (v == t) return flow;
        for (; ptr[v] < (11) g[v].size(); ++ptr[v]) {</pre>
            11 id = g[v][ptr[v]];
            11 \text{ to } = e[id].b;
            if (d[to] != d[v] + 1) continue;
            ll pushed = dfs(to, min(flow, e[id].cap - e[id].flow));
            if (pushed) {
                 e[id].flow += pushed;
                e[id ^ 1].flow -= pushed;
                return pushed;
            }
        }
        return 0;
    }
};
```

```
% cat maxflow/hungarian.h | ./misc/hash.sh
422f16
```

Listing 50: Hungarian

```
//status: tested on https://judge.yosupo.jp/problem/assignment
struct match {
    long long cost;
    vector<int> matching;
};
match HungarianMatch(const vector<vector<long long>>& cost) {
    long long n = cost.size() - 1;
    long long m = cost[0].size() - 1;
    vector\langle int \rangle p(m + 1), way(m + 1);
    vector<long long> u(n + 1), v(m + 1);
    for (int i = 1; i <= n; ++i) {
        p[0] = i;
        int j0 = 0;
        vector<long long> minv(m + 1, inf);
        vector<char> used(m + 1, false);
        do {
            used[j0] = true;
            int i0 = p[j0], j1 = 0;
            long long delta = inf;
            for (int j = 1; j \le m; ++j)
                if (!used[i]) {
                    long long cur = cost[i0][j] - u[i0] - v[j];
                     if (cur < minv[j])</pre>
                         minv[j] = cur, way[j] = j0;
                     if (minv[j] < delta)</pre>
                         delta = minv[j], j1 = j;
            for (int j = 0; j \le m; ++j)
                if (used[j])
                     u[p[j]] += delta, v[j] -= delta;
                else
                     minv[j] -= delta;
            i0 = i1:
        } while (p[j0] != 0);
        do {
            int j1 = way[j0];
            p[j0] = p[j1];
            j0 = j1;
        } while (j0);
    }
    // For each N, it contains the M it selected
    vector<int> ans(n + 1);
    for (int j = 1; j \le m; ++j)
        ans[p[j]] = j;
    return {-v[0], ans};
```

```
% cat maxflow/minCostMaxFlow.h | ./misc/hash.sh
```

b76528

Listing 51: Min Cost Max Flow

```
#pragma once
//status: not tested, but used in various problems
```

```
const long long inf = 1e18;
struct mincostmaxflow {
    typedef long long 11;
    struct edge {
       ll a, b, cap, cost, flow;
        size_t back;
   };
    vector<edge> e;
   vector<vector<ll>>> g;
   11 n, s, t;
   11 k = inf; // The maximum amount of flow allowed
    mincostmaxflow(int _n, int _s, int _t) : n(_n), s(_s), t(_t) {
        g.resize(n);
   }
    void addedge(ll a, ll b, ll cap, ll cost) {
        edge e1 = {a, b, cap, cost, 0, g[b].size() };
        edge e2 = {b, a, 0, -cost, 0, g[a].size() };
        g[a].push_back((ll) e.size());
        e.push_back(e1);
        g[b].push_back((ll) e.size());
        e.push_back(e2);
   }
    // Returns {flow, cost}
    pair<11, 11> getflow() {
        11 \text{ flow} = 0, \text{ cost} = 0;
        while (flow < k) {
            vector<ll> id(n, 0);
            vector<ll> d(n, inf);
            vector<ll> q(n);
            vector<ll> p(n);
            vector<size_t> p_edge(n);
            11 qh = 0, qt = 0;
            q[qt++] = s;
            d[s] = 0;
            while (qh != qt) {
               11 v = q[qh++];
                id[v] = 2;
                if (qh == n) qh = 0;
                for (size_t i = 0; i < g[v].size(); ++i) {</pre>
                    edge& r = e[g[v][i]];
                    if (r.flow < r.cap && d[v] + r.cost < d[r.b]) {</pre>
                        d[r.b] = d[v] + r.cost;
                        if (id[r,b] == 0) {
                            q[qt++] = r.b;
                            if (qt == n) qt = 0;
                        } else if (id[r.b] == 2) {
                            if (--qh == -1) qh = n - 1;
                            q[qh] = r.b;
                        id[r.b] = 1;
                        p[r.b] = v:
                        p_{edge}[r.b] = i;
                }
            }
```

```
if (d[t] == inf) break;
            11 addflow = k - flow:
            for (11 v = t; v != s; v = p[v]) {
                11 pv = p[v];
                size_t pr = p_edge[v];
                addflow = min(addflow, e[g[pv][pr]].cap - e[g[pv][pr]].flow);
            for (11 v = t; v != s; v = p[v]) {
                11 pv = p[v];
                size_t pr = p_edge[v], r = e[g[pv][pr]].back;
                e[g[pv][pr]].flow += addflow;
                e[g[v][r]].flow -= addflow;
                cost += e[g[pv][pr]].cost * addflow;
            flow += addflow;
        }
        return {flow, cost};
   }
};
```

Listing 52: MISC

```
% cat misc/cntRectangles.h | ./misc/hash.sh
```

ca4f09

Listing 53: Count Rectangles

```
#pragma once
//qiven a 2D boolean matrix, calculate cnt[i][j]
//cnt[i][j] = the number of times an (i * j) rectangle appears in the matrix
//such that all cells in the rectangle are false
//Note cnt[0][j] and cnt[i][0] will contain garbage values
//O(R*C)
//
//status: tested on random inputs
vector<vector<int>> getNumRectangles(const vector<vector<bool>>& grid) {
   vector<vector<int>> cnt:
    const int rows = grid.size(), cols = grid[0].size();
    if (rows == 0 || cols == 0) return cnt;
    cnt.resize(rows + 1, vector<int> (cols + 1, 0));
    vector<vector<int>> arr(rows + 2, vector<int> (cols + 1, 0));
   for (int i = 1; i <= rows; ++i) {</pre>
       for (int j = 1; j <= cols; ++j) {
            arr[i][j] = 1 + arr[i][j - 1];
            if (grid[i - 1][j - 1]) arr[i][j] = 0;
   }
   for (int j = 1; j <= cols; ++j) {</pre>
        arr[rows + 1][j] = 0;
        stack<pair<int, int>> st;
       st.push({0, 0});
       for (int i = 1; i <= rows + 1; ++i) {
            pair<int, int> curr = {i, arr[i][j]};
            while (arr[i][j] < st.top().second) {</pre>
                curr = st.top();
                st.pop();
```

```
% cat misc/longest_increasing_subsequence.h | ./misc/hash.sh
e4162a
```

Listing 54: Longest Increasing Subsequence

```
#pragma once
// status: tested on https://open.kattis.com/problems/longincsubseq
//returns array of indexes representing the longest *strictly* increasing subsequence
//for non-decreasing: pass in a vector<pair<T, int>> where second is 0, 1, ..., n-1
template<class T>
vector<int> lis(const vector<T>& arr) {
   int n = arr.size();
   vector<int> dp/*array of indexes into 'arr'*/, prev(n);
   for (int i = 0; i < n; i++) {
        auto it = lower_bound(dp.begin(), dp.end(), i, [&](int x, int y) -> bool {
            return arr[x] < arr[y];</pre>
        if (it == dp.end()) {
            prev[i] = dp.empty() ? -1 : dp.back();
            dp.push_back(i);
            prev[i] = it == dp.begin() ? -1 : *(it - 1);
            *it = i;
   vector<int> res(dp.size());
   int j = dp.size();
   for (int i = dp.back(); i != -1; i = prev[i])
       res[--i] = i;
   return res;
//returns length of longest *strictly* increasing subsequence
//alternatively, there's this https://codeforces.com/blog/entry/13225
template<class T>
int lisSize(const vector<T>& arr) {
   vector<int> dp;
```

```
for (int val : arr) {
   auto it = lower_bound(dp.begin(), dp.end(), val);
   if (it == dp.end())
        dp.push_back(val);
   else
        *it = val;
        //here, 'dp.size()' = length of LIS of prefix of 'arr' so far
}
return dp.size();
```

```
inline T getRand(T 1, T r) {
    assert(1 <= r);
    return uniform_int_distribution<T>(1, r)(rng);
}
inline double getRandReal(double 1, double r) {
    assert(1 < r);
    return uniform_real_distribution(1, r)(rng);
}</pre>
```

```
% cat misc/policy_based_data_structures.h | ./misc/hash.sh
```

807de9

Listing 55: PBDS

```
//status: not tested
//place this include *before* the '#define int long long' else compile error
#include <bits/extc++.h>
using namespace __gnu_pbds;
//BST with extra functions https://codeforces.com/blog/entry/11080
//order_of_key - # of elements *strictly* less than given element
//find_by_order - find kth largest element, k is 0 based so find_by_order(0) returns min
    \hookrightarrow element
template<class T>
using indexed_set = tree<T, null_type, less<T>, rb_tree_tag,

    tree_order_statistics_node_update>;

//example initialization:
indexed_set<pair<long long, int>> is;
//hash table (apparently faster than unordered_map):
    \hookrightarrow https://codeforces.com/blog/entry/60737
//example initialization:
gp_hash_table<string, long long> ht;
```

```
% cat misc/safehash.h | ./misc/hash.sh

a296c3
```

Listing 57: Safe Hash

```
//example initialization:
gp_hash_table<string, long long> ht;

% cat misc/random.h | ./misc/hash.sh

f2bc66

Ligting 56: Dandom Number Consector
```

Listing 56: Random Number Generator