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Listing 1: Contest

Listing 2: Template

```
#include <bits/stdc++.h>
using namespace std;
int main() {
    cin.tie(0)->sync_with_stdio(0);
    return 0;
}
```

Listing 3: Hash codes

```
#!/usr/bin/env bash
#Hashes a file, ignoring all:
   - whitespace
   - comments
   - asserts
   - includes
   - pragmas
#Use to verify that code was correctly typed.
#usage:
# chmod +x hash.sh
# cat <file> / ./hash.sh
#or just copy this command:
   cat <file> | sed -r '/(assert|include|pragma)/d' | cpp -fpreprocessed -P | tr -d
    \hookrightarrow '[:space:]' | md5sum | cut -c-6
sed -r '/(assert|include|pragma)/d' | cpp -fpreprocessed -P | tr -d '[:space:]' | md5sum
    \hookrightarrow | cut -c-6
```

Listing 4: Test on random inputs

```
#!/usr/bin/env bash
#runs 2 programs against each other on random inputs until they output different results
#source: https://github.com/Errichto/youtube/blob/master/testing/s.sh
#usage:
# chmod +x test.sh
# ./test.sh
for((i = 1; ; ++i)); do
    echo $i
    ./test.out > in
    diff --ignore-all-space <(./a.out < in) <(./brute.out < in) || break
done</pre>
```

Listing 5: GRAPHS

Listing 6: Bridges and Cuts

```
//cat bridges_and_cuts.h | ./hash.sh

//1310ef

#pragma once

//0(n+m) time & space

//2 edge cc and bcc stuff doesn't depend on each other, so delete whatever is not needed
```

```
//handles multiple edges
//example initialization of 'adj':
//for (int i = 0; i < m; i++) {
// int u, v;
// cin >> u >> v:
// u--. v--:
// adj[u].emplace_back(v, i);
// adj[v].emplace_back(u, i);
//}
struct info {
    //2 edge connected component stuff (e.g. components split by bridge edges)
         \hookrightarrow https://cp-algorithms.com/graph/bridge-searching.html
    int num_2_edge_ccs;
    vector<bool> is_bridge;//edge id -> 1 iff bridge edge
    vector<int> two_edge_ccid;//node -> id of 2 edge component (which are labeled 0, 1,
         \hookrightarrow ..., 'num_2_edge_ccs'-1)
    //bi connected component stuff (e.g. components split by cut/articulation nodes)
         \hookrightarrow https://cp-algorithms.com/graph/cutpoints.html
    int num_bccs;
    vector<bool> is_cut;//node -> 1 iff cut node
    vector<int> bcc_id; //edge id -> id of bcc (which are labeled 0, 1, ..., 'num_bccs'-1)
};
info bridge_and_cut(const vector<vector<pair<int/*neighbor*/, int/*edge id*/>>>&

    → adj/*undirected graph*/, int m/*number of edges*/) {
    //stuff for both (always keep)
    int n = adj.size(), timer = 1;
    vector<int> tin(n, 0);
    //2 edge cc stuff (delete if not needed)
    int num_2_edge_ccs = 0;
    vector<bool> is_bridge(m, 0);
    vector<int> two_edge_ccid(n), node_stack;
    //bcc stuff (delete if not needed)
    int num_bccs = 0;
    vector<bool> is_cut(n, 0);
    vector<int> bcc_id(m), edge_stack;
    auto dfs = [&](auto self, int v, int p_id) -> int {
        int low = tin[v] = timer++, deg = 0;
        node_stack.push_back(v);
        for (auto [to, e_id] : adj[v]) {
            if (e_id == p_id) continue;
            if (!tin[to]) {
                 edge_stack.push_back(e_id);
                int low_ch = self(self, to, e_id);
                if (low_ch >= tin[v]) {
                     is_cut[v] = 1;
                     while (1) {
                         int edge = edge_stack.back();
                         edge_stack.pop_back();
                         bcc_id[edge] = num_bccs;
                         if (edge == e_id) break;
                    }
                     num_bccs++;
                }
                low = min(low, low_ch);
                 deg++;
            } else if (tin[to] < tin[v]) {</pre>
                 edge_stack.push_back(e_id);
                 low = min(low, tin[to]);
        }
```

```
if (p_id == -1) is_cut[v] = (deg > 1);
    if (tin[v] == low) {
        if (p_id != -1) is_bridge[p_id] = 1;
        while (1) {
            int node = node_stack.back();
            node_stack.pop_back();
            two_edge_ccid[node] = num_2_edge_ccs;
            if (node == v) break;
        num_2_edge_ccs++;
    }
    return low;
};
for (int i = 0; i < n; i++)
    if (!tin[i])
        dfs(dfs, i, -1);
return {num_2_edge_ccs, is_bridge, two_edge_ccid, num_bccs, is_cut, bcc_id};
```

Listing 7: Block Vertex Tree

```
//cat block_vertex_tree.h / ./hash.sh
//ea8ef1
#pragma once
#include "bridges_and_cuts.h"
//returns adjacency list of block vertex tree
// info cc = bridge_and_cut(adj, m);
// vector<vector<int>> but = block_vertex_tree(adj, cc);
//to loop over each *unique* bcc containing a node v:
// for (int bccid : bvt[v]) {
    bccid -= n:
//
11 }
//to loop over each *unique* node inside a bcc:
// for (int v : bvt[bccid + n]) {
vector<vector<int>> block_vertex_tree(const vector<vector<pair<int, int>>>& adj, const
    \hookrightarrow info% cc) {
    int n = adj.size();
    vector<vector<int>> bvt(n + cc.num bccs);
    vector<bool> vis(cc.num bccs. 0):
    for (int v = 0; v < n; v++) {
        for (auto [_, e_id] : adj[v]) {
            int bccid = cc.bcc_id[e_id];
            if (!vis[bccid]) {
                vis[bccid] = 1:
                bvt[v].push_back(bccid + n); //add edge between original node, and bcc
                     \hookrightarrow node
                bvt[bccid + n].push_back(v);
           }
        for (int bccid : bvt[v]) vis[bccid - n] = 0;
   }
    return bvt:
```

Listing 8: Bridge Tree

Listing 9: Frequency Table of Tree Distance

```
//cat tree_freq_dist.h | ./hash.sh
//4e0d39
#pragma once
#include "../../kactl/content/numerical/FastFourierTransform.h"
//returns array 'cnt_paths' where 'cnt_paths[i]' = # of paths in tree with 'i' edges
//centroid decomposition + FFT
//0(n \log^2 n)
vector<long long> tree_freq_dist(const vector<vector<int>& adj/*unrooted, connected
    \hookrightarrow tree*/) {
    int n = adj.size();
    vector<int> vis(n, 0), sizes(n);
    auto dfs_sz = [&] (auto self, int node, int par) -> void {
        sizes[node] = 1;
        for (int ch : adj[node]) {
            if (ch == par || vis[ch]) continue;
            self(self, ch, node);
            sizes[node] += sizes[ch];
       }
    auto find_centroid = [&](int node) -> int {
        dfs_sz(dfs_sz, node, -1);
        int size_cap = sizes[node] / 2, par = -1;
        while (1) {
            bool found = 0;
            for (int ch : adj[node]) {
                if (ch != par && !vis[ch] && sizes[ch] > size_cap) {
                    found = 1;
                    par = node;
                    node = ch;
                    break;
                }
            if (!found) return node;
       }
   };
    vector<long long> cnt_paths(n, 0);
    auto dfs = [&](auto self. int node) -> void {
        node = find_centroid(node);
        vis[node] = 1:
        vector<double> total_depth(1, 1.0);
```

```
for (int to : adj[node]) {
        if (vis[to]) continue:
        vector<double> cnt_depth(1, 0.0);
        for (queue<pair<int, int>> q({{to, node}}); !q.empty();) {
            cnt_depth.push_back(q.size());
            queue<pair<int, int>> new_q;
            while (!q.empty()) {
                auto [curr, par] = q.front();
                q.pop();
                for (int ch : adj[curr]) {
                   if (ch == par || vis[ch]) continue;
                   new_q.emplace(ch, curr);
            }
            swap(q, new_q);
            vector<double> prod = conv(total_depth, cnt_depth);
            for (int i = 1; i < (int)prod.size(); i++) cnt_paths[i] +=</pre>
                if (total_depth.size() < cnt_depth.size())</pre>

    total_depth.resize(cnt_depth.size(), 0.0);
        for (int i = 1; i < (int)cnt_depth.size(); i++) total_depth[i] +=</pre>
            self(self, to);
    }
};
dfs(dfs, 0);
return cnt_paths;
```

Listing 10: Dijkstra

```
//cat dijkstra.h | ./hash.sh
//56a477
#pragma once
//returns array 'len' where 'len[i]' = shortest path from node v to node i
//For\ example\ len[v]\ will\ always = 0
const long long INF = 1e18;
vector<long long> dijkstra(const vector<vector<pair<int, long long>>>& adj /*directed or
    vector<long long> len(adj.size(), INF);
   len[v] = 0;
   set<pair<long long/*weiqht*/, int/*node*/>> q;
   q.insert({OLL, v});
   while (!q.empty()) {
       auto it = q.begin();
       int node = it->second;
       q.erase(it);
       for (auto [to, weight] : adi[node])
           if (len[to] > weight + len[node]) {
               q.erase({len[to], to});
               len[to] = weight + len[node];
               q.insert({len[to], to});
   }
   return len:
```

Listing 11: HLD

```
//cat hld.h / ./hash.sh
//499032
#pragma once
//source: https://codeforces.com/blog/entry/53170
//assumes a single tree, 1-based nodes is possible by passing in 'root' in range [1, n]
//mnemonic: Heavy Light Decomposition
//NOLINTNEXTLINE(readability-identifier-naming)
struct HLD {
    struct node {
        int sub_sz, par, time_in, next;
    };
    vector<node> tree;
    HLD(vector<vector<int>>& adj /*single unrooted tree*/, int root) : tree(adj.size(), {
        1, root, (int)adj.size(), root
    }) {
        dfs1(root, adj);
        int timer = 0;
        dfs2(root, adj, timer);
    void dfs1(int v, vector<vector<int>>& adj) {
        auto par = find(adj[v].begin(), adj[v].end(), tree[v].par);
        if (par != adj[v].end()) adj[v].erase(par);
        for (int& to : adj[v]) {
            tree[to].par = v;
            dfs1(to, adj);
            tree[v].sub_sz += tree[to].sub_sz;
            if (tree[to].sub_sz > tree[adj[v][0]].sub_sz)
                swap(to, adj[v][0]);
    void dfs2(int v, const vector<vector<int>>& adj, int& timer) {
        tree[v].time_in = timer++;
        for (int to : adj[v]) {
            tree[to].next = (timer == tree[v].time_in + 1 ? tree[v].next : to);
            dfs2(to, adj, timer);
    }
    // Returns inclusive-exclusive intervals (of time_in's) corresponding to the path
         \hookrightarrow between u and v, not necessarily in order
    // This can answer queries for "is some node 'x' on some path" by checking if the
         \hookrightarrow tree[x].time_in is in any of these intervals
    vector<pair<int, int>> path(int u, int v) const {
        vector<pair<int, int>> res;
        for (;; v = tree[tree[v].next].par) {
            if (tree[v].time_in < tree[u].time_in) swap(u, v);</pre>
            if (tree[tree[v].next].time_in <= tree[u].time_in) {</pre>
                res.emplace_back(tree[u].time_in, tree[v].time_in + 1);
                return res;
            res.emplace_back(tree[tree[v].next].time_in, tree[v].time_in + 1);
    // Returns interval (of time_in's) corresponding to the subtree of node i
    // This can answer queries for "is some node 'x' in some other node's subtree" by
         \hookrightarrow checking if tree[x].time_in is in this interval
    pair<int, int> subtree(int i) const {
        return {tree[i].time_in, tree[i].time_in + tree[i].sub_sz};
    // Returns lca of nodes u and v
```

```
int lca(int u, int v) const {
    for (;; v = tree[tree[v].next].par) {
        if (tree[v].time_in < tree[u].time_in) swap(u, v);
        if (tree[tree[v].next].time_in <= tree[u].time_in) return u;
    }
}</pre>
```

Listing 12: Hopcroft Karp

```
//cat hopcroft_karp.h / ./hash.sh
//de75d7
#pragma once
//source:
     \hookrightarrow https://github.com/foreverbell/acm-icpc-cheat-sheet/blob/master/src/graph-algorithm/hopcroft-karp.cpp
//Worst case O(E*sqrt(V)) but faster in practice
struct match {
    //# of edges in matching (which = size of min vertex cover by öKnig's theorem)
    int size_of_matching;
    //an arbitrary max matching is found. For this matching:
    //if l_to_r[node_left] == -1:
    // node_left is not in matching
    // the edge 'node_left' <=> l_to_r[node_left] is in the matching
    //similarly for r_to_l with edge r_to_l[node_right] <=> node_right in matching if
         \hookrightarrow r_to_l[node_right] != -1
    //matchings stored in l_to_r and r_to_l are the same matching
    //provides way to check if any node/edge is in matching
    vector<int> l_to_r, r_to_l;
    //an arbitrary min vertex cover is found. For this mvc: mvc_l[node_left] is 1 iff
         \hookrightarrow node_left is in the min vertex cover (same for mvc_r)
    //if muc_l[node_left] is 0, then node_left is in the corresponding maximal
         \hookrightarrow independent set
    vector<bool> mvc_l, mvc_r;
//Think of the bipartite graph as having a left side (with size lsz) and a right side
    \hookrightarrow (with size rsz).
//Nodes on left side are indexed 0,1,...,lsz-1
//Nodes on right side are indexed 0,1,...,rsz-1
//'adj' is like a directed adjacency list containing edges from left side -> right side:
//To initialize 'adj': For every edge node_left <=> node_right, do:
    \hookrightarrow adj[node_left].push_back(node_right)
match hopcroft_karp(const vector<vector<int>>& adj/*bipartite graph*/, int rsz/*number
    \hookrightarrow of nodes on right side*/) {
    int size_of_matching = 0, lsz = adj.size();
    vector<int> l_to_r(lsz, -1), r_to_l(rsz, -1);
    while (1) {
        queue<int> q;
        vector<int> level(lsz, -1);
        for (int i = 0; i < lsz; i++)</pre>
            if (l_to_r[i] == -1)
                level[i] = 0, q.push(i);
        bool found = 0;
        vector<bool> mvc_l(lsz, 1), mvc_r(rsz, 0);
        while (!q.empty()) {
            int u = q.front();
            q.pop();
            mvc_1[u] = 0;
```

```
for (int x : adj[u]) {
            mvc r[x] = 1:
            int v = r_{to_1[x]};
            if (v == -1) found = 1:
            else if (level[v] == -1) {
                level[v] = level[u] + 1;
                q.push(v);
            }
    }
    if (!found) return {size_of_matching, l_to_r, r_to_l, mvc_l, mvc_r};
    auto dfs = [&](auto self, int u) -> bool {
        for (int x : adj[u]) {
            int v = r_to_1[x];
            if (v == -1 || (level[u] + 1 == level[v] && self(self, v))) {
               l_{to_r[u]} = x;
                r_{to_1[x]} = u;
                return 1;
        level[u] = 1e9; //acts as visited array
   };
    for (int i = 0; i < lsz; i++)</pre>
        size_of_matching += (l_to_r[i] == -1 && dfs(dfs, i));
}
```

Listing 13: LCA

```
//cat lca.h / ./hash.sh
//22246e
#pragma once
//https://codeforces.com/blog/entry/74847
//assumes a single tree, 1-based nodes is possible by passing in 'root' in range [1, n]
//mnemonic: Least/Lowest Common Ancestor
//NOLINTNEXTLINE(readability-identifier-naming)
struct LCA {
   struct node {
        int jmp, jmp_edges, par, depth;
       long long dist;
   vector<node> tree;
   LCA(const vector<vector<pair<int, long long>>>& adj, int root) : tree(adj.size(), {
        root, 1, root, 0, OLL
   }) {
        dfs(root, adj);
   void dfs(int v, const vector<vector<pair<int, long long>>>& adj) {
        int jmp, jmp_edges;
        if (tree[v].depth > 0 && tree[v].jmp_edges == tree[tree[v].jmp].jmp_edges)
            jmp = tree[tree[v].jmp].jmp, jmp_edges = 2 * tree[v].jmp_edges + 1;
        else
            jmp = v, jmp_edges = 1;
       for (auto [ch, w] : adj[v]) {
           if (ch == tree[v].par) continue;
            tree[ch] = {
                jmp,
                jmp_edges,
               v.
```

```
1 + tree[v].depth,
                w + tree[v].dist
            };
            dfs(ch, adj);
        }
    }
    //traverse up k edges in O(log(k)). So with k=1 this returns 'v''s parent
    int kth_par(int v, int k) const {
        k = min(k, tree[v].depth);
        while (k > 0) {
            if (tree[v].jmp_edges <= k) {</pre>
                k -= tree[v].jmp_edges;
                v = tree[v].jmp;
            } else {
                k--:
                v = tree[v].par;
        }
        return v;
    }
    int get_lca(int x, int y) const {
        if (tree[x].depth < tree[y].depth) swap(x, y);</pre>
        x = kth_par(x, tree[x].depth - tree[y].depth);
        while (x != y) {
            if (tree[x].jmp != tree[y].jmp)
                x = tree[x].jmp, y = tree[y].jmp;
                x = tree[x].par, y = tree[y].par;
        }
        return x;
    }
    int dist_edges(int x, int y) const {
        return tree[x].depth + tree[y].depth - 2 * tree[get_lca(x, y)].depth;
    }
    long long dist_weight(int x, int y) const {
        return tree[x].dist + tree[y].dist - 2 * tree[get_lca(x, y)].dist;
   }
};
```

Listing 14: Kth Node on Path

```
//cat kth_node_on_path.h / ./hash.sh
//7a4c3c
#pragma once
#include "lca.h"
struct kth_node_on_path {
    LCA lca:
    kth_node_on_path(const vector<vector<pair<int, long long>>>& adj, int root) :
         \hookrightarrow lca(adj, root) {}
    //consider path \{u, u's par, \ldots, LCA(u, v), \ldots, v's par, v\}. This returns the node
         \hookrightarrow at index k
    //assumes 0 \le k \le number of edges on path from u to v
    int query(int u, int v, int k) const {
        int lca_uv = lca.get_lca(u, v);
        int u_lca = lca.tree[u].depth - lca.tree[lca_uv].depth;
        int v_lca = lca.tree[v].depth - lca.tree[lca_uv].depth;
        assert(0 <= k && k <= u_lca + v_lca);
        return k <= u_lca ? lca.kth_par(u, k) : lca.kth_par(v, u_lca + v_lca - k);</pre>
    }
};
```

Listing 15: SCC

```
//cat scc.h / ./hash.sh
//ee9331
#pragma once
//source:
     \hookrightarrow https://qithub.com/kth-competitive-programming/kactl/blob/main/content/graph/SCC.h
//mnemonic: Strongly Connected Component
struct scc_info {
    int num_sccs;
    //scc's are labeled 0,1,..., 'num_sccs-1'
    //scc_id[i] is the id of the scc containing node 'i'
    //for each edge i \rightarrow j: scc_id[i] >= scc_id[j] (topo order of scc's)
    vector<int> scc_id;
//NOLINTNEXTLINE(readability-identifier-naming)
scc_info SCC(const vector<vector<int>>& adj /*directed, unweighted graph*/) {
    int n = adj.size(), timer = 1, num_sccs = 0;
    vector<int> tin(n, 0), scc_id(n, -1), node_stack;
    auto dfs = [&](auto self, int v) -> int {
        int low = tin[v] = timer++:
        node_stack.push_back(v);
        for (int to : adj[v]) {
            if (scc_id[to] < 0)</pre>
                low = min(low, tin[to] ? tin[to] : self(self, to));
        if (tin[v] == low) {
            while (1) {
                int node = node_stack.back();
                node_stack.pop_back();
                scc_id[node] = num_sccs;
                if (node == v) break;
            num_sccs++;
        }
        return low:
    for (int i = 0; i < n; i++) {
        if (!tin[i])
            dfs(dfs, i);
    }
    return {num_sccs, scc_id};
```

Listing 16: RANGE DATA STRUCTURES

Listing 17: Lazy Segment Tree

```
}
static const dt INF = 1e18:
struct node {
    dt val:
    ch lazy;
    int 1, r;//[l, r)
};
const int N, S/*smallest power of 2 >= N*/;
vector<node> tree;
//doesn't work with empty array
seg_tree(const \ vector < dt > \ arr) : N(arr.size()), S(1 << __lg(2 * N - 1)), tree(2 * N - 1)), tree(2 * N - 1))
    for (int i = 0; i < N; i++)</pre>
        tree[i + N] = {arr[i], 0, i, i + 1};
    rotate(tree.rbegin(), tree.rbegin() + S - N, tree.rbegin() + N);
    for (int i = N - 1; i >= 1; i --) {
        tree[i] = {
            combine(tree[2 * i].val, tree[2 * i + 1].val),
            tree[2 * i].1,
            tree[2 * i + 1].r
        };
    }
}
void apply(int v, ch change) {
    tree[v].val += change;
    tree[v].lazy += change;
}
void push(int v) {
    if (tree[v].lazy) {
        apply(2 * v, tree[v].lazy);
        apply(2 * v + 1, tree[v].lazy);
        tree[v].lazy = 0;
    }
}
void build(int v) {
    tree[v].val = combine(tree[2 * v].val, tree[2 * v + 1].val);
}
int to_leaf(int i) const {
    return i < 2 * N ? i : 2 * (i - N);
}
//update range [l, r)
void update(int 1, int r, ch change) {
    assert(0 \le 1 \&\& 1 \le r \&\& r \le N);
    1 = to_leaf(1), r = to_leaf(r);
    int lca_l_r = __lg((l - 1) ^ r);
    for (int lg = __lg(l); lg > __builtin_ctz(l); lg--) push(l >> lg);
    for (int lg = lca_l_r; lg > __builtin_ctz(r); lg--) push(r >> lg);
    for (int x = 1, y = r; x < y; x >>= 1, y >>= 1) {
        if (x & 1) apply(x++, change);
        if (y & 1) apply(--y, change);
    }
    for (int lg = __builtin_ctz(r) + 1; lg <= lca_l_r; lg++) build(r >> lg);
    for (int lg = __builtin_ctz(1) + 1; lg <= __lg(1); lg++) build(1 >> lg);
}
//query range [l, r)
dt query(int 1, int r) {
    assert(0 <= 1 && 1 < r && r <= N);
    1 = to_leaf(1), r = to_leaf(r);
    int lca_l_r = __lg((l - 1) ^ r);
```

```
for (int lg = __lg(l); lg > __builtin_ctz(l); lg--) push(l >> lg);
    for (int lg = lca_l_r; lg > __builtin_ctz(r); lg--) push(r >> lg);
    dt resl = INF, resr = INF;
    for (; l < r; l >>= 1, r >>= 1) {
        if (l & 1) resl = combine(resl, tree[l++].val);
        if (r & 1) resr = combine(tree[--r].val, resr);
    }
    return combine(resl, resr);
}
```

Listing 18: BIT

```
//cat bit.h / ./hash.sh
//83059d
#pragma once
//mnemonic: Binary Indexed Tree
//NOLINTNEXTLINE(readability-identifier-naming)
template<class T> struct BIT {
    const int N;
    vector<T> bit;
    BIT(int a_n) : N(a_n), bit(N, 0) {}
    BIT(const vector<T>& a) : BIT(a.size()) {
        for (int i = 0; i < N; i++) {</pre>
            bit[i] += a[i];
            int j = i | (i + 1);
            if (j < N) bit[j] += bit[i];</pre>
    void update(int i, const T& d) {
        assert(0 <= i && i < N);
        for (; i < N; i |= i + 1) bit[i] += d;
    T sum(int r) const {//sum of range [0, r)
        assert(0 <= r && r <= N);
        T ret = 0;
        for (; r > 0; r \&= r - 1) ret += bit[r - 1];
        return ret:
    T sum(int 1, int r) const {//sum of range [l, r)
        assert(0 <= 1 && 1 <= r && r <= N);
        return sum(r) - sum(1);
    //Returns min pos (0 \le pos \le N+1) such that sum of [0, pos) \ge sum
    //Returns N + 1 if no sum is >= sum, or 0 if empty sum is.
    //Doesn't work with negatives
    int lower_bound(T sum) const {
        if (sum <= 0) return 0;</pre>
        int pos = 0;
        for (int pw = 1 << __lg(N | 1); pw; pw >>= 1)
            if (pos + pw <= N && bit[pos + pw - 1] < sum)</pre>
                pos += pw, sum -= bit[pos - 1];
        return pos + 1;
    }
};
```

Listing 19: RMQ

//082180

```
#pragma once
//source:
    // vector<long long> arr;
// RMQ<long long> rmg(arr, [8](auto x, auto y) \{ return min(x,y); \});
//to also get index of min element, do:
// RMQ<pair<T, int>> rmq(arr, [8](auto x, auto y) \{ return min(x,y); \});
//and initialize arr[i].second = i (0 <= i < n)
//If there are multiple indexes of min element, it'll return the smallest
//(left-most) one
//mnemonic: Range Min/Max Query
//NOLINTNEXTLINE(readability-identifier-naming)
template <class T> struct RMQ {
   vector<vector<T>> dp;
   function<T(const T&, const T&)> op;
   RMQ(const vector<T>& arr, const function<T(const T&, const T&)>& a_op) : dp(1, arr),
       for (int pw = 1, k = 1, n = arr.size(); 2 * pw <= n; pw *= 2, k++) {
           dp.emplace_back(n - 2 * pw + 1);
           for (int j = 0; j < n - 2 * pw + 1; j++)
              dp[k][j] = op(dp[k - 1][j], dp[k - 1][j + pw]);
   }
   //inclusive-exclusive range [l, r)
   T query(int 1, int r) const {
       assert(0 <= 1 && 1 < r && r <= (int)dp[0].size());
       int lg = __lg(r - 1);
       return op(dp[lg][l], dp[lg][r - (1 << lg)]);</pre>
   }
};
```

Listing 20: Implicit Lazy Segment Tree

```
//cat implicit_seg_tree.h / ./hash.sh
//d5be85
#pragma once
//example initialization:
// implicit_seg_tree<10'000'000> ist(l, r);
template <int N> struct implicit_seg_tree {
    using dt = array<long long, 2>;//min, number of mins
    using ch = long long;
    static dt combine(const dt& 1, const dt& r) {
        if (1[0] == r[0]) return {1[0], 1[1] + r[1]};
        return min(1, r);
    static constexpr dt UNIT{(long long)1e18, OLL);
    struct node {
        dt val;
        int lch, rch; // children, indexes into 'tree', -1 for null
        node() {}
        node(const dt& a_val) : val(a_val), lazy(0), lch(-1), rch(-1) {}
    int ptr, root_1, root_r;//[root_1, root_r) defines range of root node; handles
        \hookrightarrow negatives
    implicit_seg_tree(int 1, int r) : ptr(0), root_l(1), root_r(r) {
```

```
tree[ptr++] = node(dt\{0, r - 1\});
    void apply(int v, ch add) {
        tree[v].lazy += add;
    void push(int v, int tl, int tr) {
        if (tr - tl > 1 && tree[v].lch == -1) {
            int tm = tl + (tr - tl) / 2;
            assert(ptr + 1 < N):
            tree[v].lch = ptr;
            tree[ptr++] = node(dt{0, tm - tl});
            tree[v].rch = ptr;
            tree[ptr++] = node(dt{0, tr - tm});
        if (tree[v].lazy) {
            apply(tree[v].lch, tree[v].lazy);
            apply(tree[v].rch, tree[v].lazy);
            tree[v].lazy = 0;
        }
    //update range [l,r)
    void update(int 1, int r, ch add) {
        update(0, root_1, root_r, 1, r, add);
    }
    void update(int v, int tl, int tr, int l, int r, ch add) {
        if (r <= tl || tr <= 1)
            return;
        if (1 <= t1 && tr <= r)</pre>
            return apply(v, add);
        push(v, tl, tr);
        int tm = tl + (tr - tl) / 2;
        update(tree[v].lch, tl, tm, l, r, add);
        update(tree[v].rch, tm, tr, 1, r, add);
        tree[v].val = combine(tree[tree[v].lch].val,
                               tree[tree[v].rch].val);
    //query range [l,r)
    dt query(int 1, int r) {
        return query(0, root_1, root_r, 1, r);
    dt query(int v, int tl, int tr, int l, int r) {
        if (r <= tl || tr <= 1)
            return UNIT;
        if (1 <= t1 && tr <= r)</pre>
            return tree[v].val;
        push(v, tl, tr);
        int tm = tl + (tr - tl) / 2;
        return combine(query(tree[v].lch, tl, tm, l, r),
                        query(tree[v].rch, tm, tr, 1, r));
    }
};
```

Listing 21: Kth Smallest

```
//cat kth_smallest.h | ./hash.sh
//f8ce8b
#include <bits/stdc++.h>
using namespace std;
#pragma once
```

```
//source:
     \hookrightarrow https://cp-algorithms.com/data\_structures/segment\_tree.html#preserving-the-history-df-itsstandatesnoders(istent-segment-tree)
struct kth_smallest {
    struct node {
        int sum;
        int lch, rch;//children, indexes into 'tree'
    };
    int mn, mx;
    vector<int> roots;
    deque<node> tree:
    kth_smallest(const vector<int>& arr) : mn(INT_MAX), mx(INT_MIN), roots(arr.size() +
        tree.push_back({0, 0, 0}); //acts as null
        for (int val : arr) mn = min(mn, val), mx = max(mx, val + 1);
        for (int i = 0; i < (int)arr.size(); i++)</pre>
            roots[i + 1] = update(roots[i], mn, mx, arr[i]);
    }
    int update(int v, int tl, int tr, int idx) {
        if (tr - tl == 1) {
            tree.push_back({tree[v].sum + 1, 0, 0});
            return tree.size() - 1;
        }
        int tm = tl + (tr - tl) / 2:
        int lch = tree[v].lch;
        int rch = tree[v].rch;
        if (idx < tm)</pre>
            lch = update(lch, tl, tm, idx);
            rch = update(rch, tm, tr, idx);
        tree.push_back({tree[lch].sum + tree[rch].sum, lch, rch});
        return tree.size() - 1;
    }
    /* find (k+1)th smallest number in range [l, r)
     * k is 0-based, so query(l,r,0) returns the min
    int query(int 1, int r, int k) const {
        assert(0 <= k \&\& k < r - 1); //note this condition implies l < r
        assert(0 <= 1 && r < (int)roots.size());</pre>
        return query(roots[1], roots[r], mn, mx, k);
    }
    int query(int vl, int vr, int tl, int tr, int k) const {
        assert(tree[vr].sum > tree[vl].sum);
        if (tr - tl == 1)
            return tl;
        int tm = tl + (tr - tl) / 2;
        int left_count = tree[tree[vr].lch].sum - tree[tree[v1].lch].sum;
        if (left_count > k) return query(tree[v1].lch, tree[vr].lch, t1, tm, k);
        return query(tree[v1].rch, tree[vr].rch, tm, tr, k - left_count);
   }
};
```

Listing 22: Number Distinct Elements

```
//cat distinct_query.h / ./hash.sh
//6dfaad
#pragma once

→ https://cp-algorithms.com/data_structures/segment_tree.html#preserving-the-history| of-itsneugleusoptertried(groundstegment/bried(groundstegment/bried)), S(1 << __lg(2 * N - 1)),
</p>
//works with negatives
//0(n \log n) time and space
```

```
struct distinct_query {
        int sum;
        int lch, rch;//children, indexes into 'tree'
    };
    const int N;
    vector<int> roots:
    deque<node> tree;
    distinct_query(const vector<int>& arr) : N(arr.size()), roots(N + 1, 0) {
        tree.push_back({0, 0, 0}); //acts as null
        map<int, int> last_idx;
        for (int i = 0; i < N; i++) {</pre>
            roots[i + 1] = update(roots[i], 0, N, last_idx[arr[i]]);
            last_idx[arr[i]] = i + 1;
        }
    }
    int update(int v, int tl, int tr, int idx) {
        if (tr - tl == 1) {
            tree.push_back({tree[v].sum + 1, 0, 0});
            return tree.size() - 1;
        int tm = tl + (tr - tl) / 2;
        int lch = tree[v].lch:
        int rch = tree[v].rch;
        if (idx < tm)</pre>
            lch = update(lch, tl, tm, idx);
            rch = update(rch, tm, tr, idx);
        tree.push_back({tree[lch].sum + tree[rch].sum, lch, rch});
        return tree.size() - 1;
    //returns number of distinct elements in range [l,r)
    int query(int 1, int r) const {
        assert(0 <= 1 && 1 <= r && r <= N);
        return query(roots[1], roots[r], 0, N, 1 + 1);
    int query(int vl, int vr, int tl, int tr, int idx) const {
        if (tree[vr].sum == 0 || idx <= t1)</pre>
            return 0;
        if (tr <= idx)
            return tree[vr].sum - tree[vl].sum:
        int tm = tl + (tr - tl) / 2;
        return query(tree[v1].lch, tree[vr].lch, tl, tm, idx) +
               query(tree[v1].rch, tree[vr].rch, tm, tr, idx);
    }
};
```

Listing 23: Merge Sort Tree

```
//cat merge_sort_tree.h / ./hash.sh
//a84032
#pragma once
//For point updates: either switch to policy based BST, or use sqrt decomposition
struct merge_sort_tree {
    const int N, S/*smallest power of 2 >= N*/;
    vector<vector<int>> tree:
    //doesn't work with empty array
        \hookrightarrow tree(2 * N) {
        for (int i = 0; i < N; i++)
```

```
tree[i + N] = {arr[i]};
        rotate(tree.rbegin(), tree.rbegin() + S - N, tree.rbegin() + N);
        for (int i = N - 1; i >= 1; i--) {
            const auto& 1 = tree[2 * i];
            const auto& r = tree[2 * i + 1];
            tree[i].reserve(l.size() + r.size());
            merge(1.begin(), 1.end(), r.begin(), r.end(), back_inserter(tree[i]));
        }
    }
    int value(int v, int x) const {
        return lower_bound(tree[v].begin(), tree[v].end(), x) - tree[v].begin();
   }
    int to_leaf(int i) const {
        i += S:
        return i < 2 * N ? i : 2 * (i - N);
    //How many values in range [l, r) are < x?
    //0(log^2(n))
    int query(int 1, int r, int x) const {
        int res = 0:
        for (l = to_leaf(l), r = to_leaf(r); l < r; l >>= 1, r >>= 1) {
            if (1 & 1) res += value(1++, x);
            if (r & 1) res += value(--r, x);
        }
        return res;
   }
};
```

Listing 24: STRINGS

Listing 25: Suffix Array and LCP Array

```
//cat string.hpp / ./hash.sh
//67378f
#ifndef ATCODER_STRING_HPP
#define ATCODER_STRING_HPP 1
#include <algorithm>
#include <cassert>
#include <numeric>
#include <string>
#include <vector>
namespace atcoder {
namespace internal {
std::vector<int> sa_naive(const std::vector<int>& s) {
    int n = int(s.size());
    std::vector<int> sa(n);
    std::iota(sa.begin(), sa.end(), 0);
    std::sort(sa.begin(), sa.end(), [&](int 1, int r) {
        if (1 == r) return false;
        while (1 < n \&\& r < n) {
            if (s[1] != s[r]) return s[1] < s[r];</pre>
            1++;
            r++;
        }
```

```
return 1 == n:
   }):
    return sa;
std::vector<int> sa_doubling(const std::vector<int>& s) {
    int n = int(s.size());
    std::vector<int> sa(n), rnk = s, tmp(n);
    std::iota(sa.begin(), sa.end(), 0);
    for (int k = 1; k < n; k *= 2) {
        auto cmp = [\&] (int x, int y) {
            if (rnk[x] != rnk[y]) return rnk[x] < rnk[y];</pre>
            int rx = x + k < n ? rnk[x + k] : -1;
            int ry = y + k < n ? rnk[y + k] : -1;
            return rx < ry;</pre>
        std::sort(sa.begin(), sa.end(), cmp);
        tmp[sa[0]] = 0;
        for (int i = 1; i < n; i++) {
            tmp[sa[i]] = tmp[sa[i - 1]] + (cmp(sa[i - 1], sa[i]) ? 1 : 0);
        std::swap(tmp, rnk);
   }
    return sa;
// SA-IS, linear-time suffix array construction
// Reference:
// G. Nong, S. Zhang, and W. H. Chan,
// Two Efficient Algorithms for Linear Time Suffix Array Construction
template <int THRESHOLD_NAIVE = 10, int THRESHOLD_DOUBLING = 40>
std::vector<int> sa_is(const std::vector<int>& s, int upper) {
    int n = int(s.size());
    if (n == 0) return {};
    if (n == 1) return {0};
    if (n == 2) {
        if (s[0] < s[1]) {</pre>
            return {0, 1};
       } else {
            return {1, 0};
    if (n < THRESHOLD_NAIVE) {</pre>
        return sa_naive(s);
    if (n < THRESHOLD DOUBLING) {</pre>
        return sa_doubling(s);
    std::vector<int> sa(n);
    std::vector<bool> ls(n):
    for (int i = n - 2; i \ge 0; i--) {
        ls[i] = (s[i] == s[i + 1]) ? ls[i + 1] : (s[i] < s[i + 1]);
    std::vector<int> sum_l(upper + 1), sum_s(upper + 1);
    for (int i = 0; i < n; i++) {
        if (!ls[i]) {
            sum_s[s[i]]++;
       } else {
            sum_l[s[i] + 1]++;
       }
```

```
}
for (int i = 0: i <= upper: i++) {</pre>
    sum_s[i] += sum_l[i];
    if (i < upper) sum_l[i + 1] += sum_s[i];</pre>
}
auto induce = [&](const std::vector<int>& lms) {
    std::fill(sa.begin(), sa.end(), -1);
    std::vector<int> buf(upper + 1);
    std::copy(sum_s.begin(), sum_s.end(), buf.begin());
    for (auto d : lms) {
        if (d == n) continue;
        sa[buf[s[d]]++] = d;
    std::copy(sum_l.begin(), sum_l.end(), buf.begin());
    sa[buf[s[n-1]]++] = n-1;
    for (int i = 0; i < n; i++) {</pre>
        int v = sa[i]:
        if (v >= 1 && !ls[v - 1]) {
            sa[buf[s[v - 1]] ++] = v - 1;
    }
    std::copy(sum_l.begin(), sum_l.end(), buf.begin());
    for (int i = n - 1; i \ge 0; i--) {
        int v = sa[i];
        if (v >= 1 && ls[v - 1]) {
            sa[--buf[s[v-1]+1]] = v-1;
};
std::vector<int> lms_map(n + 1, -1);
int m = 0:
for (int i = 1; i < n; i++) {</pre>
    if (!ls[i - 1] && ls[i]) {
        lms_map[i] = m++;
    }
}
std::vector<int> lms;
lms.reserve(m);
for (int i = 1: i < n: i++) {
    if (!ls[i - 1] && ls[i]) {
        lms.push_back(i);
}
induce(lms);
if (m) {
    std::vector<int> sorted_lms;
    sorted lms.reserve(m):
    for (int v : sa) {
        if (lms_map[v] != -1) sorted_lms.push_back(v);
    std::vector<int> rec_s(m);
    int rec_upper = 0;
    rec s[lms map[sorted lms[0]]] = 0:
    for (int i = 1; i < m; i++) {</pre>
        int l = sorted_lms[i - 1], r = sorted_lms[i];
        int end_l = (lms_map[l] + 1 < m) ? lms[lms_map[l] + 1] : n;</pre>
        int end_r = (lms_map[r] + 1 < m) ? lms[lms_map[r] + 1] : n;
```

```
bool same = true:
            if (end 1 - 1 != end r - r) {
                same = false;
            } else {
                while (1 < end_1) {
                     if (s[1] != s[r]) {
                        break:
                    }
                    1++;
                    r++:
                if (1 == n || s[1] != s[r]) same = false;
            if (!same) rec_upper++;
            rec_s[lms_map[sorted_lms[i]]] = rec_upper;
        auto rec_sa =
            sa_is<THRESHOLD_NAIVE, THRESHOLD_DOUBLING>(rec_s, rec_upper);
        for (int i = 0; i < m; i++) {</pre>
            sorted lms[i] = lms[rec sa[i]]:
        induce(sorted_lms);
    }
    return sa;
} // namespace internal
std::vector<int> suffix_array(const std::vector<int>& s, int upper) {
    assert(0 <= upper);</pre>
    for (int d : s) {
        assert(0 <= d && d <= upper);</pre>
    auto sa = internal::sa_is(s, upper);
    return sa;
template <class T> std::vector<int> suffix_array(const std::vector<T>& s) {
    int n = int(s.size());
    std::vector<int> idx(n);
    iota(idx.begin(), idx.end(), 0);
    sort(idx.begin(), idx.end(), [&](int 1, int r) { return s[1] < s[r]; });</pre>
    std::vector<int> s2(n);
    int now = 0:
    for (int i = 0; i < n; i++) {</pre>
        if (i && s[idx[i - 1]] != s[idx[i]]) now++;
        s2[idx[i]] = now;
    return internal::sa_is(s2, now);
std::vector<int> suffix_array(const std::string& s) {
    int n = int(s.size());
    std::vector<int> s2(n);
    for (int i = 0; i < n; i++) {</pre>
        s2[i] = s[i];
    return internal::sa_is(s2, 255);
```

```
// Reference:
// T. Kasai, G. Lee, H. Arimura, S. Arikawa, and K. Park,
// Linear-Time Longest-Common-Prefix Computation in Suffix Arrays and Its
// Applications
template <class T>
std::vector<int> lcp_array(const std::vector<T>& s,
                           const std::vector<int>& sa) {
    int n = int(s.size());
    assert(n >= 1):
    std::vector<int> rnk(n);
    for (int i = 0; i < n; i++) {</pre>
        rnk[sa[i]] = i;
    std::vector<int> lcp(n - 1);
    int h = 0;
   for (int i = 0; i < n; i++) {</pre>
        if (h > 0) h--;
        if (rnk[i] == 0) continue;
        int j = sa[rnk[i] - 1];
        for (; j + h < n && i + h < n; h++) {
            if (s[j + h] != s[i + h]) break;
        lcp[rnk[i] - 1] = h;
   }
    return lcp;
std::vector<int> lcp_array(const std::string& s, const std::vector<int>& sa) {
   int n = int(s.size());
    std::vector<int> s2(n);
    for (int i = 0; i < n; i++) {</pre>
        s2[i] = s[i];
   }
    return lcp_array(s2, sa);
// Reference:
// D. Gusfield,
// Algorithms on Strings, Trees, and Sequences: Computer Science and
// Computational Biology
template <class T> std::vector<int> z_algorithm(const std::vector<T>& s) {
    int n = int(s.size());
    if (n == 0) return {};
    std::vector<int> z(n);
   z[0] = 0:
    for (int i = 1, j = 0; i < n; i++) {
        int & k = z[i];
        k = (j + z[j] \le i) ? 0 : std::min(j + z[j] - i, z[i - j]);
        while (i + k < n \&\& s[k] == s[i + k]) k++;
        if (j + z[j] < i + z[i]) j = i;
   }
    z[0] = n;
    return z:
std::vector<int> z_algorithm(const std::string& s) {
    int n = int(s.size());
    std::vector<int> s2(n);
    for (int i = 0; i < n; i++) {</pre>
        s2[i] = s[i];
```

```
return z_algorithm(s2);
} // namespace atcoder
#endif // ATCODER STRING HPP
```

Listing 26: KMP

```
//cat kmp.h / ./hash.sh
//73f1be
#pragma once
//mnemonic: Knuth Morris Pratt
#include "../../kactl/content/strings/KMP.h"
// string needle;
// ...
// KMP kmp(needle);
1/or
// vector<int> needle;
// ...
// KMP kmp(needle);
//kmp doubling trick: to check if 2 arrays are rotationally equivalent: run kmp
//with one array as the needle and the other array doubled (excluding the first
//8 last characters) as the haystack or just use kactl's min rotation code
//NOLINTNEXTLINE(readability-identifier-naming)
template <class T> struct KMP {
    KMP(const T& a_needle) : needle(a_needle), pf(pi(needle)) {}
    // if haystack = "bananas"
    // needle = "ana"
    // then we find 2 matches:
    // bananas
    // _ana___
    // ___ana_
    // 0123456 (indexes)
    // and KMP::find returns {1,3} - the indexes in haystack where
    // each match starts.
    // You can also pass in 0 for "all" and KMP::find will only
    // return the first match: {1}. Useful for checking if there exists
     // some match:
    // KMP::find(<haystack>,0).size() > 0
    vector<int> find(const T& haystack, bool all = 1) const {
        vector<int> matches;
        for (int i = 0, j = 0; i < (int)haystack.size(); i++) {</pre>
            while (j > 0 && needle[j] != haystack[i]) j = pf[j - 1];
            if (needle[j] == haystack[i]) j++;
            if (j == (int)needle.size()) {
                matches.push_back(i - (int)needle.size() + 1);
                if (!all) return matches;
                j = pf[j - 1];
            }
        return matches;
    }
    T needle:
    vector<int> pf;//prefix function
```

|};

Listing 27: Trie

```
//cat trie.h | ./hash.sh
//fd9c8d
#pragma once
//source: https://cp-algorithms.com/string/aho_corasick.html#construction-of-the-trie
const int K = 26://alphabet size
struct trie {
    const char MIN_CH = 'A';//'a' for lowercase, '0' for digits
    struct node {
        int next[K], cnt_words = 0, par = -1;
        char ch;
        node(int a_par = -1, char a_ch = '#') : par(a_par), ch(a_ch) {
            fill(next, next + K, -1);
        }
    };
    vector<node> t;
    trie() : t(1) {}
    void add_string(const string& s) {
        int v = 0:
        for (char ch : s) {
            int let = ch - MIN_CH;
            if (t[v].next[let] == -1) {
                t[v].next[let] = t.size();
                t.emplace_back(v, ch);
            v = t[v].next[let];
        }
        t[v].cnt_words++;
   }
   bool find_string(const string& s) const {
        int v = 0:
        for (char ch : s) {
            int let = ch - MIN_CH;
            if (t[v].next[let] == -1) return 0:
            v = t[v].next[let];
        return t[v].cnt words:
   }
};
```

Listing 28: Binary Trie

```
int update(long long val, int delta) {
        int c = 0:
        t[0].sub_sz += delta;
        for (int bit = MX_BIT; bit >= 0; bit--) {
            bool v = (val >> bit) & 1;
            if (t[c].next[v] == -1) {
                t[c].next[v] = t.size();
                t.emplace_back();
            c = t[c].next[v]:
            t[c].sub_sz += delta;
        t[c].val = val;
        return t[c].sub_sz;
    int size() const {
        return t[0].sub_sz;
    //returns x such that:
    // x is in this data structure
    // value of (x ^val) is minimum
    long long min_xor(long long val) const {
        assert(size() > 0);
        int c = 0;
        for (int bit = MX_BIT; bit >= 0; bit--) {
            bool v = (val >> bit) & 1;
            int ch = t[c].next[v];
            if (ch != -1 && t[ch].sub_sz > 0)
                c = ch;
            else
                c = t[c].next[!v];
        return t[c].val;
    }
};
```

Listing 29: Longest Common Prefix Query

```
//cat lcp_query.h / ./hash.sh
//951676
#pragma once
#include "../../ac-library/atcoder/string.hpp"
#include "../range_data_structures/rmq.h"
//computes suffix array, lcp array, and then sparse table over lcp array
//0(n \log n)
struct lcp_query {
    const int N;
    vector<int> sa, lcp, inv_sa;
    lcp_query(const string& s) : N(s.size()), sa(atcoder::suffix_array(s)),
         \hookrightarrow lcp(atcoder::lcp_array(s, sa)), inv_sa(N), st(lcp, [](int x, int y) {
        return min(x, y);
    }) {
        for (int i = 0; i < N; i++) inv_sa[sa[i]] = i;</pre>
    //length of longest common prefix of suffixes s[idx1 ... n), s[idx2 ... n), 0-based
    //You can check if two substrings s[l1..r1), s[l2..r2) are equal in O(1) by:
    //r1-l1 == r2-l2 \& longest_common_prefix(l1, l2) >= r1-l1
```

```
int longest_common_prefix(int idx1, int idx2) const {
    if (idx1 == idx2) return N - idx1;
    idx1 = inv_sa[idx1];
    idx2 = inv_sa[idx2];
    if (idx1 > idx2) swap(idx1, idx2);
    return st.query(idx1, idx2);
}

//returns 1 if suffix s[idx1 ... n) < s[idx2 ... n)
//(so 0 if idx1 == idx2)
bool less(int idx1, int idx2) const {
    return inv_sa[idx1] < inv_sa[idx2];
}
};</pre>
```

Listing 30: Palindrome Query

```
//cat palindrome_query.h / ./hash.sh
//7326d0
#pragma once
#include "../../kactl/content/strings/Manacher.h"
struct pal_query {
   const int N;
    array<vi, 2> pal_len;
    pal_query(const string& s) : N(s.size()), pal_len(manacher(s)) {}
    //returns 1 if substring s[l...r) is a palindrome
    //(returns\ 1\ when\ l == r)
   bool is_pal(int 1, int r) const {
        assert(0 <= 1 && 1 <= r && r <= N);
        int len = r - 1;
        return pal_len[len & 1][l + len / 2] >= len / 2;
   }
};
```

Listing 31: MATH

Listing 32: BIN EXP MOD

```
//cat exp_mod.h / ./hash.sh
//3be256
#pragma once
//returns (base^pw)/mod in O(log(pw)), but returns 1 for 0^0
//What if base doesn't fit in long long?
//Since (base^pw) %mod == ((base%mod)^pw) %mod we can calculate base under mod of 'mod'
//What if pw doesn't fit in long long?
//case 1: mod is prime
//(base^pw)%mod == (base^(pw%(mod-1)))%mod (from Fermat's little theorem)
//so calculate pw under mod of 'mod-1'
//note 'mod-1' is not prime, so you need to be able to calculate 'pw%(mod-1)' without
     \hookrightarrow division
//case 2: non-prime mod
//let t = totient(mod)
//if pw >= log2(mod) then (base^pw)%mod == (base^(t+(pw%t)))%mod (proof)

→ https://cp-algorithms.com/algebra/phi-function.html#generalization)

//so calculate pw under mod of 't'
```

Listing 33: Fibonacci

Listing 34: Matrix Mult and Pow

```
//cat matrix_expo.h / ./hash.sh
//2edd34
#pragma once
//empty matrix -> RTE
vector<vector<int>> mult(const vector<vector<int>>& a, const vector<vector<int>>& b, int
    \hookrightarrow mod) {
    assert(a[0].size() == b.size());
    int n = a.size(), m = b[0].size(), inner = b.size();
    vector<vector<int>> prod(n, vector<int>(m, 0));
    for (int i = 0; i < n; i++) {
        for (int k = 0; k < inner; k++) {
            for (int j = 0; j < m; j++)
                prod[i][j] = (prod[i][j] + 1LL * a[i][k] * b[k][j]) % mod;
   }
    return prod;
vector<vector<int>> power(vector<int>> mat/*intentional pass by value*/, long
     \hookrightarrow long pw, int mod) {
    int n = mat.size();
    vector<vector<int>> prod(n, vector<int>(n, 0));
    for (int i = 0; i < n; i++)</pre>
        prod[i][i] = 1;
    while (pw > 0) {
        if (pw % 2 == 1) prod = mult(prod, mat, mod);
        mat = mult(mat, mat, mod);
        pw /= 2;
```

```
return prod;
}
```

Listing 35: N Choose K MOD

```
//cat n choose k mod.h | ./hash.sh
//f3a1a9
#pragma once
//for mod inverse
#include "exp_mod.h"
// usage:
       n_choose_k \ nk(n, 1e9+7) to use 'choose', 'inv' with inputs strictly < n
// or:
       n_choose_k nk(mod, mod) to use 'choose_with_lucas_theorem' with arbitrarily large
struct n_choose_k {
    n_choose_k(int n, int a_mod) : mod(a_mod), fact(n, 1), inv_fact(n, 1) {
        //this implementation doesn't work if n > mod because n! % mod = 0 when n > = 0
             \hookrightarrow mod. So 'inv_fact' array will be all 0's
        assert(max(n, 2) \le mod);
        //assert mod is prime. mod is intended to fit inside an int so that
        //multiplications fit in a longlong before being modded down. So this
        //will take sqrt(2^31) time
        for (int i = 2; i * i <= mod; i++) assert(mod % i);</pre>
        for (int i = 2; i < n; i++)
            fact[i] = 1LL * fact[i - 1] * i % mod;
        inv_fact.back() = pow(fact.back(), mod - 2, mod);
        for (int i = n - 2; i \ge 2; i--)
            inv_fact[i] = inv_fact[i + 1] * (i + 1LL) % mod;
   }
    //classic n choose k
    //fails when n \ge mod
    int choose(int n, int k) const {
        if (k < 0 \mid k > n) return 0:
        //now we know 0 <= k <= n so 0 <= n
        return 1LL * fact[n] * inv_fact[k] % mod * inv_fact[n - k] % mod;
    }
    //lucas theorem to calculate n choose k in O(\log(k))
    //need to calculate all factorials in range [0,mod), so O(mod) time&space, so need
         \hookrightarrow smallish prime mod (< 1e6 maybe)
    //handles n >= mod correctly
    int choose_with_lucas_theorem(long long n, long long k) const {
        if (k < 0 \mid k > n) return 0;
        if (k == 0 || k == n) return 1;
        return 1LL * choose_with_lucas_theorem(n / mod, k / mod) * choose(n % mod, k %
             \hookrightarrow mod) % mod;
   }
    //returns \ x \ such \ that \ x * n % \ mod == 1
    int inv(int n) const {
        assert(1 <= n); //don't divide by 0 :)</pre>
        return 1LL * fact[n - 1] * inv_fact[n] % mod;
   }
    int mod:
    vector<int> fact, inv_fact;
```

Listing 36: Partitions

Listing 37: Derangements

```
//cat derangements.h / ./hash.sh
//c221bb
#pragma once
//https://oeis.org/A000166
//
//for a permutation of size i:
//there are (i-1) places to move 0 to not be at index 0. Let's say we moved 0 to index j

-> (j>0).
//If we move value j to index 0 (forming a cycle of length 2), then there are dp[i-2]
-> derangements of the remaining i-2 elements
//else there are dp[i-1] derangements of the remaining i-1 elements (including j)
vector<int> derangements(int n, int mod) {
    vector<int> derangements(int n, int mod) {
    vector<int> dp(n, 0);
    dp[0] = 1;
    for (int i = 2; i < n; i++)
        dp[i] = 1LL * (i - 1) * (dp[i - 1] + dp[i - 2]) % mod;
    return dp;
}
```

Listing 38: Prime Sieve

```
//cat prime_sieve.h / ./hash.sh
//45fc23
#pragma once
//a_prime[val] = some random prime factor of 'val'
//to check if 'val' is prime:
// if (a_prime[val] == val)
//to get all prime factors of a number 'val' in O(log(val)):
// while (val > 1) {
//
        int p = a_prime[val];
//
        //p is some prime factor of val
        val /= p;
// }
const int N = 1e6 + 10:
int a_prime[N];
```

Listing 39: Mobius Inversion

```
//cat mobius_inversion.h / ./hash.sh
//811515
#pragma once
//mobius[i] = 0 iff there exists a prime p s.t. i%(p^2)=0
//mobius[i] = -1 iff i has an odd number of distinct prime factors
//mobius[i] = 1 iff i has an even number of distinct prime factors
const int N = 1e6 + 10;
int mobius[N];
void calc_mobius() {
   mobius[1] = 1;
   for (int i = 1; i < N; i++)
        for (int j = i + i; j < N; j += i)
        mobius[j] -= mobius[i];
}</pre>
```

Listing 40: Row Reduce

```
//cat row_reduce.h | ./hash.sh
//1d7c3e
#pragma once
//for mod inverse
#include "exp_mod.h"
//First 'cols' columns of mat represents a matrix to be left in reduced row echelon form
//Row operations will be performed to all later columns
//example usage:
// row_reduce(mat, mat[0].size(), mod) //row reduce matrix with no extra columns
pair<int/*rank*/, int/*determinant*/> row_reduce(vector<vector<int>>& mat, int cols, int
    \hookrightarrow mod) {
    int n = mat.size(), m = mat[0].size(), rank = 0, det = 1;
    assert(cols <= m);</pre>
    for (int col = 0; col < cols && rank < n; col++) {</pre>
        //find arbitrary pivot and swap pivot to current row
        for (int i = rank; i < n; i++)</pre>
            if (mat[i][col] != 0) {
                if (rank != i) det = det == 0 ? 0 : mod - det;
                swap(mat[i], mat[rank]);
                break;
        if (mat[rank][col] == 0) {
            det = 0;
            continue;
        det = (1LL * det * mat[rank][col]) % mod;
        //make pivot 1 by dividing row by inverse of pivot
        int a_inv = pow(mat[rank][col], mod - 2, mod);
        for (int j = 0; j < m; j++)
            mat[rank][j] = (1LL * mat[rank][j] * a_inv) % mod;
        //zero-out all numbers above & below pivot
```

```
for (int i = 0; i < n; i++)
    if (i != rank && mat[i][col] != 0) {
        int val = mat[i][col];
        for (int j = 0; j < m; j++) {
            mat[i][j] -= 1LL * mat[rank][j] * val % mod;
            if (mat[i][j] < 0) mat[i][j] += mod;
        }
    }
    rank++;
}
assert(rank <= min(n, cols));
return {rank, det};
}</pre>
```

Listing 41: Solve Linear Equations MOD

```
//cat solve_linear_mod.h / ./hash.sh
//44cc6e
#pragma once
#include "row_reduce.h"
struct matrix_info {
   int rank, det;
    vector<int> x;
};
//Solves\ mat\ *\ x\ =\ b\ under\ prime\ mod.
//mat is a n (rows) by m (cols) matrix, b is a length n column vector, x is a length m
    \hookrightarrow vector.
//assumes n,m >= 1, else RTE
//Returns rank of mat, determinant of mat, and x (solution vector to mat * x = b).
//x is empty if no solution. If rank < m, there are multiple solutions and an arbitrary
    \hookrightarrow one is returned.
//Leaves mat in reduced row echelon form (unlike kactl) with b appended.
//0(n * m * min(n.m))
matrix_info solve_linear_mod(vector<vector<int>>& mat, const vector<int>& b, int mod) {
    assert(mat.size() == b.size());
    int n = mat.size(), m = mat[0].size();
    for (int i = 0; i < n; i++)
        mat[i].push_back(b[i]);
    auto [rank, det] = row_reduce(mat, m, mod);//row reduce not including the last column
    //check if solution exists
    for (int i = rank; i < n; i++) {</pre>
        if (mat[i].back() != 0) return {rank, det, {} }; //no solution exists
    //initialize solution vector ('x') from row-reduced matrix
    vector<int> x(m, 0);
    for (int i = 0, j = 0; i < rank; i++) {</pre>
        while (mat[i][j] == 0) j++; //find pivot column
        x[j] = mat[i].back();
    return {rank, det, x};
```

Listing 42: Matrix Inverse

```
//cat matrix_inverse.h | ./hash.sh

//3056ad

#pragma once

#include "row_reduce.h"

//returns inverse of square matrix mat, empty if no inverse
```

Listing 43: Euler's Totient Phi Function

```
//cat totient.h | ./hash.sh
//36bd41
#pragma once
//Euler's totient function counts the positive integers
//up to a given integer n that are relatively prime to n.
//To improve, use Pollard-rho to find prime factors
int totient(int n) {
   int res = n;
   for (int i = 2; i * i <= n; i++) {
       if (n % i == 0) {
            while (n \% i == 0) n /= i;
            res -= res / i;
       }
   }
   if (n > 1) res -= res / n;
   return res;
```

Listing 44: MAX FLOW

Listing 45: Dinic

```
//cat dinic.h / ./hash.sh
//e6a3d5
#pragma once
struct max_flow {
    using ll = long long;
    const ll INF = le18;
    struct edge {
        int a, b;
        ll cap, flow;
    };
    vector<edge> e;
    vector<vector<int>> g;
    vector<int> q, d, ptr;
    max_flow(int n) : g(n), q(n), d(n), ptr(n) {}
    void add_edge(int a, int b, ll cap) {
```

```
edge e1 = { a, b, cap, 0 };
        edge e2 = \{ b, a, 0, 0 \};
        g[a].push_back(e.size());
        e.push_back(e1);
        g[b].push_back(e.size());
        e.push_back(e2);
    11 get_flow(int s, int t) {
        11 \text{ flow} = 0;
        for (::) {
            if (!bfs(s, t)) break;
            ptr.assign(ptr.size(), 0);
            while (ll pushed = dfs(s, INF, t))
                flow += pushed;
        }
        return flow;
    }
    bool bfs(int s, int t) {
        int qh = 0, qt = 0;
        q[qt++] = s;
        d.assign(d.size(), -1);
        d[s] = 0:
        while (qh < qt && d[t] == -1) {
            int v = q[qh++];
            for (int i = 0; i < (int)g[v].size(); i++) {</pre>
                 int id = g[v][i], to = e[id].b;
                if (d[to] == -1 && e[id].flow < e[id].cap) {</pre>
                     q[qt++] = to;
                     d[to] = d[v] + 1;
                }
            }
        return d[t] != -1;
    11 dfs(int v, ll flow, int t) {
        if (!flow) return 0;
        if (v == t) return flow;
        for (; ptr[v] < (int)g[v].size(); ptr[v]++) {</pre>
            int id = g[v][ptr[v]], to = e[id].b;
            if (d[to] != d[v] + 1) continue;
            ll pushed = dfs(to, min(flow, e[id].cap - e[id].flow), t);
            if (pushed) {
                 e[id].flow += pushed;
                e[id ^ 1].flow -= pushed;
                return pushed;
        }
        return 0;
    }
};
```

Listing 46: Hungarian

```
//cat hungarian.h | ./hash.sh
//625431
#pragma once
//source: https://e-maxx.ru/algo/assignment_hungary
//
//input: cost[1...n][1...m] with 1 <= n <= m
//n workers, indexed 1, 2, ..., n
```

```
//m jobs, indexed 1, 2, ..., m
//it costs 'cost[i][j]' to assign worker i to job j (1<=i<=n, 1<=j<=m)
//this returns *min* total cost to assign each worker to some distinct job
//0(n^2 * m)
//trick 1: set 'cost[i][j]' to INF to say: "worker 'i' cannot be assigned job 'j'"
//trick 2: 'cost[i][i]' can be negative, so to instead find max total cost over all
     \hookrightarrow matchings: set all 'cost[i][j]' to '-cost[i][j]'.
//Now max total cost = - hungarian(cost).min_cost
const long long INF = 1e18:
struct match {
    long long min_cost;
    vector<int> matching://worker 'i' (1<=i<=n) is assigned to job 'matching[i]'
         \hookrightarrow (1<=matching[i]<=m)
match hungarian(const vector<vector<long long>>& cost) {
    int n = cost.size() - 1, m = cost[0].size() - 1;
    assert(n <= m):
    vector\langle int \rangle p(m + 1), way(m + 1);
    vector<long long> u(n + 1), v(m + 1);
    for (int i = 1; i <= n; i++) {</pre>
        p[0] = i;
        int j0 = 0;
        vector<long long> minv(m + 1, INF);
        vector<bool> used(m + 1, 0);
        do {
             used[j0] = 1;
            int i0 = p[j0], j1 = 0;
            long long delta = INF;
            for (int j = 1; j <= m; j++)
                 if (!used[j]) {
                     long long cur = cost[i0][j] - u[i0] - v[j];
                     if (cur < minv[j])</pre>
                         minv[j] = cur, way[j] = j0;
                     if (minv[j] < delta)</pre>
                         delta = minv[j], j1 = j;
                 }
            for (int j = 0; j <= m; j++)
                 if (used[i])
                     u[p[j]] += delta, v[j] -= delta;
                 else
                     minv[j] -= delta;
             j0 = j1;
        } while (p[j0] != 0);
        do {
             int j1 = way[j0];
            p[j0] = p[j1];
             j0 = j1;
        } while (j0);
    vector<int> ans(n + 1);
    for (int j = 1; j <= m; j++)
        ans[p[j]] = j;
    return {-v[0], ans};
```

Listing 47: Min Cost Max Flow

```
//cat min_cost_max_flow.h / ./hash.sh
//6d926c
```

```
#pragma once
const long long INF = 1e18;
struct min_cost_max_flow {
    using 11 = long long;
    struct edge {
        int a, b;
        11 cap, cost, flow;
        int back;
   };
    const int N:
    vector<edge> e;
    vector<vector<int>> g;
    min_cost_max_flow(int a_n) : N(a_n), g(N) {}
    void add_edge(int a, int b, ll cap, ll cost) {
        edge e1 = {a, b, cap, cost, 0, (int)g[b].size() };
        edge e2 = {b, a, 0, -cost, 0, (int)g[a].size() };
        g[a].push_back(e.size());
        e.push_back(e1);
        g[b].push_back(e.size());
        e.push_back(e2);
    //returns minimum cost to send 'total_flow' flow through the graph, or -1 if
        \hookrightarrow impossible
   ll get_flow(int s, int t, ll total_flow) {
        11 flow = 0, cost = 0;
        while (flow < total_flow) {</pre>
            vector<ll> d(N, INF);
            vector<int> p_edge(N), id(N, 0), q(N), p(N);
            int qh = 0, qt = 0;
            q[qt++] = s;
            d[s] = 0;
            while (qh != qt) {
                int v = q[qh++];
                id[v] = 2;
                if (ah == N) ah = 0:
                for (int i = 0; i < (int)g[v].size(); i++) {</pre>
                    const edge& r = e[g[v][i]];
                    if (r.flow < r.cap && d[v] + r.cost < d[r.b]) {
                        d[r.b] = d[v] + r.cost;
                        if (id[r.b] == 0) {
                            a[at++] = r.b:
                            if (qt == N) qt = 0;
                        } else if (id[r.b] == 2) {
                            if (--qh == -1) qh = N - 1;
                            q[qh] = r.b;
                        id[r.b] = 1;
                        p[r.b] = v;
                        p_{edge}[r.b] = i;
                }
            if (d[t] == INF) break;
            11 addflow = total flow - flow:
            for (int v = t; v != s; v = p[v]) {
                int pv = p[v], pr = p_edge[v];
                addflow = min(addflow, e[g[pv][pr]].cap - e[g[pv][pr]].flow);
            for (int v = t; v != s; v = p[v]) {
                int pv = p[v], pr = p_edge[v], r = e[g[pv][pr]].back;
                e[g[pv][pr]].flow += addflow;
```

```
}
    //here, dp.size() = length of LIS of prefix of arr ending at index i
}
vector<int> res(dp.size());
for (int i = dp.back(), j = dp.size(); i != -1; i = prev[i])
    res[--j] = i;
return res;
```

Listing 48: MISC

Listing 49: PBDS

```
//cat policy_based_data_structures.h / ./hash.sh
//807de9
#pragma once
//place these includes *before* the '#define int long long' else compile error
//not using <bits/extc++.h> as it compile errors on codeforces c++20 compiler
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
//BST with extra functions https://codeforces.com/bloq/entry/11080
//order_of_key - # of elements *strictly* less than given element
//find_by_order - find kth largest element, k is 0 based so find_by_order(0) returns min
    \hookrightarrow element
template<class T> using indexed_set = tree<T, null_type, less<T>, rb_tree_tag,

    tree_order_statistics_node_update>;

//example initialization:
indexed_set<pair<long long, int>> is;
//hash table (apparently faster than unordered_map):

→ https://codeforces.com/blog/entry/60737

//example initialization:
gp_hash_table<string, long long> ht;
```

Listing 50: LIS

```
//cat lis.h / ./hash.sh
//a243e1
#pragma once
//returns array of indexes representing the longest *strictly* increasing subsequence
//for non-decreasing: pass in a vectorvector(T, int>> with arr[i].second = i (0<=i<n)</pre>
//alternatively, there's this https://codeforces.com/blog/entry/13225
//mnemonic: Longest Increasing Subsequence
//NOLINTNEXTLINE(readability-identifier-naming)
template<class T> vector<int> LIS(const vector<T>& arr) {
   if (arr.empty()) return {};
   vector<int> dp{0}/*array of indexes into 'arr'*/, prev(arr.size(), -1);
   for (int i = 1; i < (int)arr.size(); i++) {</pre>
        auto it = lower_bound(dp.begin(), dp.end(), i, [&](int x, int y) -> bool {
            return arr[x] < arr[y];</pre>
       }):
       if (it == dp.end()) {
            prev[i] = dp.back();
            dp.push_back(i);
       } else {
            prev[i] = it == dp.begin() ? -1 : *(it - 1);
```

Listing 51: Number of Distinct Subsequences DP

```
//cat num_distinct_subsequences.h | ./hash.sh
//9542f5
#pragma once
//returns number of distinct subsequences
//the empty subsequence is counted
int num_subsequences(const vector<int>& arr, int mod) {
    int n = arr.size():
    vector < int > dp(n + 1, 1);
    map<int, int> last;
    for (int i = 0; i < n; i++) {</pre>
        int& curr = dp[i + 1] = 2 * dp[i];
        if (curr >= mod) curr -= mod;
        auto it = last.find(arr[i]);
        if (it != last.end()) {
            curr -= dp[it->second];
            if (curr < 0) curr += mod;</pre>
            it->second = i;
        } else last[arr[i]] = i;
   }
    return dp[n];
```

Listing 52: Safe Hash

```
//cat safe_hash.h | ./hash.sh
//d9ea53
#pragma once
//source: https://codeforces.com/blog/entry/62393
struct custom_hash {
    static uint64_t splitmix64(uint64_t x) {
       // http://xorshift.di.unimi.it/splitmix64.c
       x += 0x9e3779b97f4a7c15;
       x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
       x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
       return x ^ (x >> 31);
    size_t operator()(uint64_t x) const {
       static const uint64_t FIXED_RANDOM =
            return splitmix64(x + FIXED_RANDOM);
   }
};
//usage:
unordered_map<long long, int, custom_hash> safe_map;
#include "policy_based_data_structures.h"
gp_hash_table<long long, int, custom_hash> safe_hash_table;
```

//cat count_rectangles.h | ./hash.sh

Listing 53: Monotonic Stack

```
//cat monotonic_stack.h | ./hash.sh
//4c7a40
#pragma once
//usages:
// vector<int> left = monotonic_stack<int>(arr, less()); //(or replace 'less' with:
    \hookrightarrow less_equal, greater, greater_equal
// vector < int > left = monotonic_stack < int > (arr, [8] (int x, int y) {return x < y;});
//
//returns array 'left' where 'left[i]' = max index such that:
// 'left[i]' \langle i \& \& !op(arr[left[i]], arr[i])
//or -1 if no index exists
//0(n)
template<class T> vector<int> monotonic_stack(const vector<T>& arr, const
    int n = arr.size():
   vector<int> left(n);
   for (int i = 0; i < n; i++) {
       int \& i = left[i] = i - 1:
       while (j >= 0 && op(arr[j], arr[i])) j = left[j];
   }
   return left;
```

Listing 54: Count Rectangles

```
//b2cced
#pragma once
#include "monotonic_stack.h"
//given a 2D boolean matrix, calculate cnt[i][j]
//cnt[i][j] = the number of times an i-by-j rectangle appears in the matrix such that
    \hookrightarrow all i*j cells in the rectangle are 1
//Note cnt[0][j] and cnt[i][0] will contain garbage values
//0(n*m)
vector<vector<int>> count_rectangles(const vector<vector<bool>>& grid) {
    int n = grid.size(), m = grid[0].size();
    vector<vector<int>> cnt(n + 1, vector<int>(m + 1, 0));
    vector<int> arr(m, 0);
    auto rv /*reverse*/ = [&](int j) -> int {
        return m - 1 - j;
   };
    for (int i = 0; i < n; i++) {
        for (int j = 0; j < m; j++)
            arr[j] = grid[i][j] * (arr[j] + 1);
        vector<int> left = monotonic_stack<int>(arr, greater());
        vector<int> right = monotonic_stack<int>(vector<int>(arr.rbegin(), arr.rend()),
             \hookrightarrow greater_equal());
        for (int j = 0; j < m; j++) {
            int l = j - left[j] - 1, r = rv(right[rv(j)]) - j - 1;
            cnt[arr[j]][l + r + 1]++;
            cnt[arr[j]][1]--;
            cnt[arr[j]][r]--;
   }
    for (int i = 1: i \le n: i++)
        for (int k = 0: k < 2: k++)
            for (int j = m; j > 1; j--)
                cnt[i][j - 1] += cnt[i][j];
    for (int j = 1; j <= m; j++)
```

Listing 55: Cartesian Tree

```
//cat cartesian_tree.h | ./hash.sh
//0b95bc
#pragma once
#include "monotonic stack.h"
//min cartesian tree
vector<int> cartesian tree(const vector<int>& arr) {
    int n = arr.size();
    auto rv /*reverse*/ = [&](int i) -> int {
        return n - 1 - i:
   };
    vector<int> left = monotonic_stack<int>(arr, greater());
    vector<int> right = monotonic_stack<int>(vector<int>(arr.rbegin(), arr.rend()),
         \hookrightarrow greater());
    vector<int> par(n);
    for (int i = 0; i < n; i++) {
        int l = left[i], r = rv(right[rv(i)]);
        if (1 >= 0 \&\& r < n) par[i] = arr[l] > arr[r] ? 1 : r;
        else if (1 >= 0) par[i] = 1;
        else if (r < n) par[i] = r;
        else par[i] = i;
   }
   return par;
```

Listing 56: Max Rectangle in Histogram

```
//cat max_rect_histogram.h | ./hash.sh
//4e6291
#pragma once
#include "monotonic_stack.h"
long long max_rect_histogram(const vector<int>& arr) {
    int n = arr.size();
    auto rv /*reverse*/ = [&](int i) -> int {
        return n - 1 - i;
   };
    vector<int> left = monotonic_stack<int>(arr, greater_equal());
    vector<int> right = monotonic_stack<int>(vector<int>(arr.rbegin(), arr.rend()),
        long long max_area = 0;
    for (int i = 0; i < n; i++) {
        int l = left[i], r = rv(right[rv(i)]); //arr[i] is the max of range (l, r)
        \max_{\text{area}} = \max(\max_{\text{area}}, 1LL * arr[i] * (r - 1 - 1));
   }
   return max_area;
```