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Listing 1: Contest

Listing 2: Hash codes

```
#!/usr/bin/env bash
#Hashes a file, ignoring all:
  - whitespace
   - comments
   - asserts
   - includes
   - pragmas
#Use to verify that code was correctly typed.
#usage:
# chmod +x hash.sh
# cat <file> / ./hash.sh
#or just copy this command:
# cat <file> | sed -r '/(assert/include/pragma)/d' | cpp -fpreprocessed -P | tr -d
    \hookrightarrow '[:space:]' | md5sum | cut -c-6
sed -r '/(assert|include|pragma)/d' | cpp -fpreprocessed -P | tr -d '[:space:]' | md5sum
    \hookrightarrow | cut -c-6
```

Listing 3: Test on random inputs

```
#!/usr/bin/env bash
#runs 2 programs against each other on random inputs until they output different results
#usage:
# chmod +x test.sh
# ./test.sh
for((i = 1; ; ++i)); do
    echo $i
    ./test.out > in
    diff --ignore-all-space <(./a.out < in) <(./brute.out < in) || break
done</pre>
```

Listing 4: GRAPHS

Listing 5: Bridges and Cuts

```
//cat bridges_and_cuts.h | ./hash.sh
//34dc49
#pragma once
//with asserts checking correctness of is_bridge and is_cut
//O(n+m) time & space
//2 edge cc and bcc stuff doesn't depend on each other, so delete whatever is not needed
//handles multiple edges
//example initialization of 'adj':
//for (int i = 0; i < m; i++) {
// int u. v:
// cin >> u >> v;
// u--, v--;
// adj[u].emplace_back(v, i);
// adj[v].emplace_back(u, i);
//}
struct info {
```

```
//2 edge connected component stuff (e.g. components split by bridge edges)
         \hookrightarrow https://cp-algorithms.com/qraph/bridge-searching.html
    int num_2_edge_ccs;
    vector<bool> is_bridge;//edge id -> true iff bridge edge
    vector<int> two_edge_ccid; //node -> id of 2 edge component (which are labeled 0, 1,
         \hookrightarrow ..., 'num_2_edge_ccs'-1)
    //bi connected component stuff (e.g. components split by cut/articulation nodes)
         \hookrightarrow https://cp-algorithms.com/graph/cutpoints.html
    int num_bccs;
    vector<bool> is_cut;//node -> true iff cut node
    vector<int> bcc_id; //edge id -> id of bcc (which are labeled 0, 1, ..., 'num_bccs'-1)
info bridge_and_cut(const vector<vector<pair<int/*neighbor*/, int/*edge id*/>>>&

    → adj/*undirected graph*/, int m/*number of edges*/) {
    //stuff for both (always keep)
    int n = adj.size(), timer = 1;
    vector<int> tin(n, 0);
    //2 edge cc stuff (delete if not needed)
    int num_2_edge_ccs = 0;
    vector<bool> is_bridge(m, false);
    vector<int> two_edge_ccid(n), node_stack;
    //bcc stuff (delete if not needed)
    int num_bccs = 0;
    vector<bool> is_cut(n, false);
    vector<int> bcc_id(m), edge_stack;
    auto dfs = [&](auto self, int v, int p_id) -> int {
        int low = tin[v] = timer++, deg = 0;
        node_stack.push_back(v);
        for (auto [to, e_id] : adj[v]) {
            if (e_id == p_id) continue;
            if (!tin[to]) {
                edge_stack.push_back(e_id);
                int low_ch = self(self, to, e_id);
                if (low_ch >= tin[v]) {
                    is cut[v] = true:
                    while (true) {
                         int edge = edge_stack.back();
                         edge_stack.pop_back();
                        bcc_id[edge] = num_bccs;
                        if (edge == e_id) break;
                    num_bccs++;
                }
                low = min(low, low_ch);
                deg++;
            } else if (tin[to] < tin[v]) {</pre>
                edge_stack.push_back(e_id);
                low = min(low, tin[to]);
            }
        if (p_id == -1) is_cut[v] = (deg > 1);
        if (tin[v] == low) {
            if (p_id != -1) is_bridge[p_id] = true;
            while (true) {
                int node = node_stack.back();
                node_stack.pop_back();
                two_edge_ccid[node] = num_2_edge_ccs;
                if (node == v) break;
            num_2_edge_ccs++;
        }
```

//cat block_vertex_tree.h | ./hash.sh

```
return low:
}:
for (int i = 0; i < n; i++)</pre>
    if (!tin[i])
        dfs(dfs, i, -1);
return {num_2_edge_ccs, is_bridge, two_edge_ccid, num_bccs, is_cut, bcc_id};
```

Listing 6: Block Vertex Tree

```
//4052a5
#pragma once
#include "bridges_and_cuts.h"
//(asserts checking correctness of commented-example-usage-loops)
//returns adjacency list of block vertex tree
//usage:
// info cc = bridge_and_cut(adj, m);
// vector<vector<int>> but = block_vertex_tree(adj, cc);
//to loop over each *unique* bcc containing a node v:
// for(int bccid : bvt[v]) {
       bccid -= n:
//
// }
//to loop over each *unique* node inside a bcc:
// for(int v : bvt[bccid + n]) {
// }
vector<vector<int>> block_vertex_tree(const vector<vector<pair<int, int>>>& adj, const
    \hookrightarrow info% cc) {
    int n = adj.size();
    vector<vector<int>> bvt(n + cc.num_bccs);
    vector<bool> vis(cc.num_bccs, false);
    for (int v = 0; v < n; v++) {
        for (auto [_, e_id] : adj[v]) {
            int bccid = cc.bcc_id[e_id];
            if (!vis[bccid]) {
                vis[bccid] = true:
                bvt[v].push_back(bccid + n);//add edge between original node, and bcc
                bvt[bccid + n].push_back(v);
        }
        for (int bccid : bvt[v])
            vis[bccid - n] = false;
    }
    return bvt;
```

Listing 7: Bridge Tree

```
//cat bridge_tree.h | ./hash.sh
//85f56b
#pragma once
#include "bridges and cuts.h"
//never adds multiple edges as bridges_and_cuts.h correctly marks them as non-bridges
//usage:
// info cc = bridge_and_cut(adj, m);
// vector<vector<int>> bt = bridge_tree(adj, cc);
```

```
vector<vector<int>>> bridge_tree(const vector<vector<pair<int, int>>>& adj, const info&
    vector<vector<int>> tree(cc.num_2_edge_ccs);
    for (int i = 0; i < (int)adj.size(); i++)</pre>
        for (auto [to, e_id] : adj[i])
            if (cc.is_bridge[e_id])
                tree[cc.two_edge_ccid[i]].push_back(cc.two_edge_ccid[to]);
    return tree;
```

Listing 8: Centroid

```
//cat centroid.h / ./hash.sh
//4ba5e4
#pragma once
//with asserts checking depth of tree <= log2(n)
//returns array 'par' where 'par[i]' = parent of node 'i' in centroid tree
//'par[root]' is -1
//0-based nodes
//0(n \log n)
//example usage:
// vector<int> parent = get_centroid_tree(adj);
// vector<vector<int>> childs(n);
// int root;
// for (int i = 0; i < n; i++) {
        if (parent[i] == -1)
//
            root = i;
//
        else
//
            childs[parent[i]].push_back(i);
// }
vector<int> get_centroid_tree(const vector<vector<int>>& adj/*unrooted tree*/) {
    int n = adj.size();
    vector<int> sizes(n);
    vector<bool> vis(n, false);
    auto dfs_sz = [&](auto self, int node, int par) -> void {
        sizes[node] = 1;
        for (int to : adj[node]) {
            if (to != par && !vis[to]) {
                self(self, to, node);
                sizes[node] += sizes[to];
        }
    };
    auto find_centroid = [&](int node) -> int {
        dfs_sz(dfs_sz, node, node);
        int size_cap = sizes[node] / 2, par = -1;
        while (true) {
            bool found = false;
            for (int to : adj[node]) {
                if (to != par && !vis[to] && sizes[to] > size_cap) {
                    found = true;
                    par = node;
                    node = to;
                    break;
            if (!found) return node;
    };
    vector<int> parent(n);
```

```
auto dfs = [&](auto self, int node, int par) -> void {
   node = find_centroid(node);
   parent[node] = par;
   vis[node] = true;
   for (int to : adj[node]) {
        if (!vis[to])
            self(self, to, node);
     }
};
dfs(dfs, 0, -1);
return parent;
}
```

Listing 9: Dijkstra

```
//cat dijkstra.h / ./hash.sh
//56a477
#pragma once
//returns array 'len' where 'len[i]' = shortest path from node v to node i
//For example len[v] will always = 0
const long long INF = 1e18;
vector<long long> dijkstra(const vector<vector<pair<int, long long>>>& adj /*directed or
    \hookrightarrow undirected, weighted graph*/, int v) {
    vector<long long> len(adj.size(), INF);
    len[v] = 0:
    set<pair<long long/*weight*/, int/*node*/>> q;
    q.insert({OLL, v});
    while (!q.empty()) {
        auto it = q.begin();
        int node = it->second;
        q.erase(it);
        for (auto [to, weight] : adj[node])
            if (len[to] > weight + len[node]) {
                q.erase({len[to], to});
                len[to] = weight + len[node];
                q.insert({len[to], to});
    }
    return len;
```

Listing 10: Floyd Warshall

```
//cat floyd_warshall.h | ./hash.sh
//84799a
#pragma once
//status: not tested
//**for directed graphs only** if you initialize len[i][i] to infinity, then
//afterward floyds, len[i][i] = length of shortest cycle including node 'i'
//
//another trick: change 'len' to 2d array of *bools* where len[i][j] = true if
//there exists an edge from i -> j in initial graph. Also do:
//'len[i][j] = len[i][j] | (len[i][k] & len[k][j])'
//Then after floyds, len[i][j] = true iff there's exists some path from node
//'i' to node 'j'
//
//Changing the order of for-loops to i-j-k (instead of the current k-i-j)
//results in min-plus matrix multiplication. If adjacency matrix is 'mat', then
//after computing mat^k (with binary exponentiation), mat[i][j] = min length path
```

```
//from i to j with at most k edges.

for (int k = 0; k < n; k++)
    for (int i = 0; i < n; i++)
        for (int j = 0; j < n; j++)
            len[i][j] = min(len[i][j], len[i][k] + len[k][j]);
```

Listing 11: HLD

```
//cat hld.h / ./hash.sh
//8a1639
#pragma once
//source: https://codeforces.com/blog/entry/53170
//assumes a single tree, 1-based nodes is possible by passing in 'root' in range [1, n]
//mnemonic: Heavy Light Decomposition
//NOLINTNEXTLINE(readability-identifier-naming)
struct HLD {
    struct node {
        int sub_sz, par, time_in, next;
   }:
    vector<node> tree:
   HLD(vector<vector<int>>& adj /*single unrooted tree*/, int root) : tree(adj.size(), {
        1, root, (int)adj.size(), root
        dfs1(root, adj);
        int timer = 0:
        dfs2(root, adj, timer);
    void dfs1(int v, vector<vector<int>>& adj) {
       for (int& to : adj[v]) {
            if (to == tree[v].par) continue;
            tree[to].par = v;
            dfs1(to, adj);
            tree[v].sub_sz += tree[to].sub_sz;
            if (tree[to].sub_sz > tree[adj[v][0]].sub_sz || adj[v][0] == tree[v].par)
                swap(to, adj[v][0]);
       }
    void dfs2(int v, const vector<vector<int>>& adj, int& timer) {
        tree[v].time_in = timer++;
        for (int to : adj[v]) {
            if (to == tree[v].par) continue;
            tree[to].next = (timer == tree[v].time_in + 1 ? tree[v].next : to);
            dfs2(to, adj, timer);
       }
   }
    // Returns inclusive-exclusive intervals (of time_in's) corresponding to the path
        \hookrightarrow between u and v, not necessarily in order
    // This can answer queries for "is some node 'x' on some path" by checking if the
        \hookrightarrow tree[x].time_in is in any of these intervals
    vector<pair<int, int>> path(int u, int v) const {
        vector<pair<int, int>> res;
        for (;; v = tree[tree[v].next].par) {
            if (tree[v].time_in < tree[u].time_in) swap(u, v);</pre>
            if (tree[tree[v].next].time_in <= tree[u].time_in) {</pre>
                res.emplace_back(tree[u].time_in, tree[v].time_in + 1);
                return res:
            res.emplace_back(tree[tree[v].next].time_in, tree[v].time_in + 1);
   }
```

//with asserts checking correctness of min vertex cover

//cat hopcroft_karp.h / ./hash.sh

//7d04ba

#pragma once

Listing 12: Hopcroft Karp

```
//source:
    \hookrightarrow https://qithub.com/foreverbell/acm-icpc-cheat-sheet/blob/master/src/graph-algorithm/hopcroft-karp.cpp
//Worst case O(E*sqrt(V)) but faster in practice
struct match {
   //# of edges in matching (which = size of min vertex cover by öKnig's theorem)
    int size_of_matching;
    //an arbitrary max matching is found. For this matching:
    //if l_to_r[node_left] == -1:
    // node_left is not in matching
    // the edge 'node_left' <=> l_to_r[node_left] is in the matching
    //similarly for r_to_l with edge r_to_l[node_right] <=> node_right in matching if
         \hookrightarrow r_to_l[node_right] != -1
    //matchings stored in l_to_r and r_to_l are the same matching
    //provides way to check if any node/edge is in matching
    vector<int> l_to_r, r_to_l;
    //an arbitrary min vertex cover is found. For this mvc: mvc_l[node_left] is true iff
         \hookrightarrow node_left is in the min vertex cover (same for mvc_r)
    //if mvc_l[node_left] is false, then node_left is in the corresponding maximal
         \hookrightarrow independent set
    vector<bool> mvc_1, mvc_r;
//Think of the bipartite graph as having a left side (with size lsz) and a right side
    \hookrightarrow (with size rsz).
//Nodes on left side are indexed 0,1,...,lsz-1
//Nodes on right side are indexed 0,1,...,rsz-1
//'adj' is like a directed adjacency list containing edges from left side -> right side:
//To initialize 'adj': For every edge node_left <=> node_right, do:
    \hookrightarrow adj[node_left].push_back(node_right)
match hopcroft_karp(const vector<vector<int>>& adj/*bipartite graph*/, int rsz/*number
    \hookrightarrow of nodes on right side*/) {
    int size_of_matching = 0, lsz = adj.size();
    vector<int> l_to_r(lsz, -1), r_to_l(rsz, -1);
    while (true) {
        queue<int> q;
        vector<int> level(lsz, -1);
        for (int i = 0; i < lsz; i++) {</pre>
```

```
if (l_to_r[i] == -1) level[i] = 0, q.push(i);
    bool found = false;
    vector<bool> mvc_l(lsz, true), mvc_r(rsz, false);
    while (!q.empty()) {
        int u = q.front();
        q.pop();
        mvc_l[u] = false;
        for (int x : adj[u]) {
            mvc r[x] = true:
            int v = r_{to_1[x]};
            if (v == -1) found = true;
            else if (level[v] == -1) {
                level[v] = level[u] + 1;
                q.push(v);
    if (!found) return {size_of_matching, l_to_r, r_to_l, mvc_l, mvc_r};
    auto dfs = [&](auto self, int u) -> bool {
        for (int x : adj[u]) {
            int v = r_to_1[x];
            if (v == -1 \mid | (level[u] + 1 == level[v] && self(self, v))) {
               l_{to_r[u]} = x;
                r_{to_1[x]} = u;
                return true;
            }
        level[u] = 1e9; //acts as visited array
        return false:
    };
    for (int i = 0; i < lsz; i++)</pre>
        size_of_matching += (l_to_r[i] == -1 && dfs(dfs, i));
}
```

Listing 13: LCA

```
//cat lca.h | ./hash.sh
//22246e
#pragma once
//https://codeforces.com/blog/entry/74847
//assumes a single tree, 1-based nodes is possible by passing in 'root' in range [1, n]
//mnemonic: Least/Lowest Common Ancestor
//NOLINTNEXTLINE(readability-identifier-naming)
struct LCA {
    struct node {
        int jmp, jmp_edges, par, depth;
        long long dist;
   };
   LCA(const vector<vector<pair<int, long long>>>& adj, int root) : tree(adj.size(), {
        root, 1, root, 0, OLL
   }) {
        dfs(root, adj);
    void dfs(int v, const vector<vector<pair<int, long long>>>& adj) {
        int jmp, jmp_edges;
        if (tree[v].depth > 0 && tree[v].jmp_edges == tree[tree[v].jmp].jmp_edges)
            jmp = tree[tree[v].jmp].jmp, jmp_edges = 2 * tree[v].jmp_edges + 1;
```

```
else
            jmp = v, jmp_edges = 1;
        for (auto [ch, w] : adj[v]) {
            if (ch == tree[v].par) continue;
            tree[ch] = {
                jmp,
                jmp_edges,
                1 + tree[v].depth,
                w + tree[v].dist
            dfs(ch, adj);
        }
    //traverse up k edges in O(log(k)). So with k=1 this returns 'v''s parent
    int kth_par(int v, int k) const {
        k = min(k, tree[v].depth);
        while (k > 0) {
            if (tree[v].jmp_edges <= k) {</pre>
                k -= tree[v].jmp_edges;
                v = tree[v].jmp;
            } else {
                k--:
                v = tree[v].par;
        return v;
    }
    int get_lca(int x, int y) const {
        if (tree[x].depth < tree[y].depth) swap(x, y);</pre>
        x = kth_par(x, tree[x].depth - tree[y].depth);
        while (x != y) {
            if (tree[x].jmp != tree[y].jmp)
                x = tree[x].jmp, y = tree[y].jmp;
                x = tree[x].par, y = tree[y].par;
        }
        return x:
    }
    int dist_edges(int x, int y) const {
        return tree[x].depth + tree[y].depth - 2 * tree[get_lca(x, y)].depth;
    }
    long long dist_weight(int x, int y) const {
        return tree[x].dist + tree[y].dist - 2 * tree[get_lca(x, y)].dist;
    }
};
```

Listing 14: SCC

```
vector<int> scc id:
};
//NOLINTNEXTLINE(readability-identifier-naming)
scc_info SCC(const vector<vector<int>>& adj /*directed, unweighted graph*/) {
    int n = adj.size(), timer = 1, num_sccs = 0;
    vector<int> tin(n, 0), scc_id(n, -1), node_stack;
    auto dfs = [&](auto self, int v) -> int {
        int low = tin[v] = timer++;
        node_stack.push_back(v);
        for (int to : adi[v]) {
            if (scc_id[to] < 0)</pre>
                low = min(low, tin[to] ? tin[to] : self(self, to));
        if (tin[v] == low) {
            while (true) {
                int node = node_stack.back();
                node_stack.pop_back();
                scc_id[node] = num_sccs;
                if (node == v) break;
            num_sccs++;
        }
        return low;
    for (int i = 0; i < n; i++) {
        if (!tin[i])
            dfs(dfs, i);
    return {num_sccs, scc_id};
```

Listing 15: RANGE DATA STRUCTURES

Listing 16: Lazy Segment Tree

```
//cat seg_tree.h | ./hash.sh
//0ee39c
#pragma once
//source: https://codeforces.com/blog/entry/18051,
     \hookrightarrow https://github.com/ecnerwala/cp-book/blob/master/src/seg\_tree.hpp
const long long INF = 1e18;
struct seg_tree {
    using dt /*data type*/ = long long;//min
    using ch = long long;
    static dt combine(const dt& 1, const dt& r) {
        return min(1, r);
    struct node {
        dt val:
        ch lazy;
        int 1, r;//[l, r)
    const int N;
    vector<node> tree:
    seg_tree(const vector<long long>& arr) : N(arr.size()), tree(2 * N) {
        for (int i = 0, j = 1 << __lg(2 * N - 1); i < N; i++, <math>j = (j + 1) \% N + N)
            tree[j] = {arr[i], 0, i, i + 1};
        for (int i = N - 1; i >= 1; i--) {
```

```
tree[i] = {
                combine(tree[2 * i].val, tree[2 * i + 1].val),
                tree[2 * i].1.
                tree[2 * i + 1].r
            };
        }
    }
    //what happens when 'change' is applied to every index in range [tree[v].l,
         \hookrightarrow tree[v1.r)?
    void apply(int v, ch change) {
        tree[v].val += change;
        tree[v].lazy += change;
    }
    void push(int v) {
        if (tree[v].lazy) {
            apply(2 * v, tree[v].lazy);
            apply(2 * v + 1, tree[v].lazy);
            tree[v].lazy = 0;
        }
    }
    void build(int v) {
        tree[v].val = combine(tree[2 * v].val, tree[2 * v + 1].val);
    int range_idx(int i) const {
        i += 1 << __lg(2 * N - 1);
        return i < 2 * N ? i : 2 * (i - N);
    }
    //update range [l, r) with 'change'
    void update(int 1, int r, ch change) {
        l = range_idx(l), r = range_idx(r);
        for (int lg = __lg(l); lg > __builtin_ctz(l); lg--) push(l >> lg);
        for (int lg = __lg((l - 1) ^ r); lg > __builtin_ctz(r); lg--) push(r >> lg);
        for (int x = 1, y = r; x < y; x >>= 1, y >>= 1) {
            if (x \& 1) apply(x++, change);
            if (y & 1) apply(--y, change);
        }
        for (int \lg = \_builtin\_ctz(r) + 1; \lg \le \_\lg((1 - 1) \hat{r}); \lg + +) build(r > 
             \hookrightarrow lg);
        for (int lg = __builtin_ctz(l) + 1; lg <= __lg(l); lg++) build(l >> lg);
    }
    //query range [l, r)
    dt query(int 1, int r) {
        l = range_idx(l), r = range_idx(r);
        for (int lg = __lg(1); lg > __builtin_ctz(1); lg--) push(1 >> lg);
        for (int lg = __lg((l - 1) ^ r); lg > __builtin_ctz(r); lg--) push(r >> lg);
        dt resl = INF, resr = INF;
        for (; 1 < r; 1 >>= 1, r >>= 1) {
            if (1 & 1) resl = combine(resl, tree[l++].val);
            if (r & 1) resr = combine(tree[--r].val, resr);
        return combine(resl, resr);
   }
};
```

Listing 17: BIT

```
//cat bit.h | ./hash.sh
//3f62ed
#pragma once
```

```
//mnemonic: Binary Indexed Tree
 //NOLINTNEXTLINE(readability-identifier-namina)
template<class T> struct BIT {
    const int N:
    vector<T> bit;
    BIT(int a_n) : N(a_n), bit(N + 1, 0) {}
    BIT(const vector<T>& a) : N(a.size()), bit(N + 1, 0) {
         for (int i = 1; i <= N; i++) {</pre>
             bit[i] += a[i - 1];
             int i = i + (i \& -i):
             if (j <= N) bit[j] += bit[i];</pre>
    }
    void update(int i, const T& d) {
         assert(0 <= i && i < N);</pre>
         for (i++; i <= N; i += i & -i) bit[i] += d;
    T sum(int r) const {//sum of range [0, r)
         assert(0 \le r \&\& r \le N);
         T ret = 0:
         for (; r; r -= r & -r) ret += bit[r];
         return ret:
    T sum(int 1, int r) const {//sum of range [l, r)
         assert(0 <= 1 && 1 <= r && r <= N);
         return sum(r) - sum(1);
    }
    //Returns min pos such that sum of [0, pos) >= sum
    //Returns N + 1 if no sum is >= sum, or 0 if empty sum is.
    //Doesn't work with negatives
    int lower_bound(T sum) const {
         if (sum <= 0) return 0;</pre>
         int pos = 0:
         for (int pw = 1 << __lg(N | 1); pw; pw >>= 1) {
             if (pos + pw <= N && bit[pos + pw] < sum)</pre>
                 pos += pw, sum -= bit[pos];
        }
         return pos + 1;
    }
};
```

Listing 18: RMQ

```
//cat rmq.h | ./hash.sh
//a90b91
#pragma once
//usage:
// vector<long long> arr;
// ...
// RMQ<long long> st(arr, [&](auto x, auto y) { return min(x,y); });
//
//to also get index of min element, do:
// RMQ<pair<T, int>> st(arr, [&](auto x, auto y) { return min(x,y); });
//and initialize arr[i].second = i (0<=i<n)
//If there are multiple indexes of min element, it'll return the smallest
//(left-most) one
//mnemonic: Range Nin/Max Query
//NOLINTNEXTLINE(readability-identifier-naming)
template <class T> struct RMQ {
    vector<vector<T>> dp;
```

```
function<T(const T&, const T&)> func:
    RMQ(const vector<T>& arr, const function<T(const T&, const T&)>& a_func) : dp(1,
        for (int pw = 1, k = 1, n = arr.size(); 2 * pw \le n; pw *= 2, k++) {
            dp.emplace_back(n - 2 * pw + 1);
           for (int j = 0; j < n - 2 * pw + 1; j++)
               dp[k][j] = func(dp[k - 1][j], dp[k - 1][j + pw]);
       }
   }
    //inclusive-exclusive range [l, r)
   T query(int 1, int r) const {
       assert(0 <= 1 && 1 < r && r <= (int)dp[0].size());
        int \lg = \_\lg(r - 1);
       return func(dp[lg][1], dp[lg][r - (1 << lg)]);</pre>
   }
};
```

Listing 19: Implicit Lazy Segment Tree

```
//cat implicit_seq_tree.h | ./hash.sh
//5f2458
#pragma once
//example initialization:
// implicit_seq_tree<10',000',000> ist(l, r);
const long long INF = 1e18;
template <int N> struct implicit_seg_tree {
    using dt = array<long long, 3>; //sum, max, min
    struct node {
        dt val:
        long long lazy;
        int lch, rch; // children, indexes into 'tree', -1 for null
        node(const dt& a_val) : val(a_val), lazy(0), lch(-1), rch(-1) {}
    } tree[N];
    int ptr, root_1, root_r; //[root_1, root_r) defines range of root node; handles
        \hookrightarrow negatives
    implicit_seg_tree(int 1, int r) : ptr(0), root_l(1), root_r(r) {
        tree[ptr++] = node(dt\{0, 0, 0\});
    //what happens when 'add' is applied to every index in range [tl, tr)?
    void apply(int v, int tl, int tr, long long add) {
        tree[v].val[0] += (tr - tl) * add;
        tree[v].val[1] += add;
        tree[v].val[2] += add;
        tree[v].lazy += add;
   }
    void push(int v, int tl, int tr) {
        if (tr - tl > 1 && tree[v].lch == -1) {
            assert(ptr + 1 < N);</pre>
            tree[v].lch = ptr;
            tree[ptr++] = node(dt{0, 0, 0});
            tree[v].rch = ptr;
            tree[ptr++] = node(dt\{0, 0, 0\});
        if (tree[v].lazy) {
            int tm = tl + (tr - tl) / 2;
            apply(tree[v].lch, tl, tm, tree[v].lazy);
            apply(tree[v].rch, tm, tr, tree[v].lazy);
            tree[v].lazy = 0;
        }
```

```
static dt pull(const dt& 1, const dt& r) {
        return {
            1[0] + r[0],
            \max(1[1], r[1]),
            min(1[2], r[2])
        };
    }
    //update range [l,r) with 'add'
    void update(int 1, int r, long long add) {
        update(0, root_1, root_r, 1, r, add);
    void update(int v, int tl, int tr, int l, int r, long long add) {
        if (r <= tl || tr <= 1)
            return:
        if (1 <= t1 && tr <= r)
            return apply(v, tl, tr, add);
        push(v, tl, tr);
        int tm = tl + (tr - tl) / 2;
        update(tree[v].lch, tl, tm, l, r, add);
        update(tree[v].rch, tm, tr, 1, r, add);
        tree[v].val = pull(tree[tree[v].lch].val,
                           tree[tree[v].rch].val);
    //query range [l,r)
    dt query(int 1, int r) {
        return query(0, root_1, root_r, 1, r);
    dt query(int v, int tl, int tr, int l, int r) {
        if (r <= tl || tr <= 1)
            return {0, -INF, INF};
        if (1 <= t1 && tr <= r)</pre>
            return tree[v].val;
        push(v, tl, tr);
        int tm = tl + (tr - tl) / 2;
        return pull(query(tree[v].lch, tl, tm, l, r),
                    query(tree[v].rch, tm, tr, 1, r));
    }
};
```

Listing 20: Kth Smallest

```
//cat kth_smallest.h / ./hash.sh
//a9f9ed
#pragma once
    struct kth smallest {
   struct node {
      int sum;
      int lch, rch;//children, indexes into 'tree'
   };
   int mn, mx;
   vector<int> roots:
   deque<node> tree;
   kth_smallest(const vector<int>& arr) : mn(INT_MAX), mx(INT_MIN), roots(arr.size() +
      tree.push_back({0, 0, 0}); //acts as null
      for (int val : arr) mn = min(mn, val), mx = max(mx, val + 1);
      for (int i = 0; i < (int)arr.size(); i++)</pre>
```

```
roots[i + 1] = update(roots[i], mn, mx, arr[i]);
    }
    int update(int v, int tl, int tr, int idx) {
        if (tr - tl == 1) {
            tree.push_back(\{\text{tree}[v].sum + 1, 0, 0\});
            return tree.size() - 1;
        int tm = tl + (tr - tl) / 2;
        int lch = tree[v].lch;
        int rch = tree[v].rch;
        if (idx < tm)</pre>
            lch = update(lch, tl, tm, idx);
        else
            rch = update(rch, tm, tr, idx);
        tree.push_back({tree[lch].sum + tree[rch].sum, lch, rch});
        return tree.size() - 1;
    }
    /* find (k+1)th smallest number in range [l, r)
     * k is 0-based, so query(l,r,0) returns the min
    int query(int 1, int r, int k) const {
        assert(0 \le k \&\& k \le r - 1); //note this condition implies <math>l \le r
        assert(0 <= 1 && r < (int)roots.size());</pre>
        return query(roots[1], roots[r], mn, mx, k);
   }
    int query(int vl, int vr, int tl, int tr, int k) const {
        assert(tree[vr].sum > tree[vl].sum);
        if (tr - tl == 1)
            return tl;
        int tm = tl + (tr - tl) / 2;
        int left_count = tree[tree[vr].lch].sum - tree[tree[vl].lch].sum;
        if (left_count > k) return query(tree[v1].lch, tree[vr].lch, t1, tm, k);
        return query(tree[v1].rch, tree[vr].rch, tm, tr, k - left_count);
   }
};
```

Listing 21: Number Distinct Elements

```
//cat distinct_query.h / ./hash.sh
//6dfaad
#pragma once
//source:

→ https://cp-algorithms.com/data_structures/segment_tree.html#preserving-the-history|df-itsmemgle_exonpte_tree/goodneste_queentory/eint>& arr): N(arr.size()), tree(2 * N) {

//works with negatives
//0(n \log n) time and space
struct distinct_query {
    struct node {
        int sum:
        int lch, rch;//children, indexes into 'tree'
   };
    const int N;
    vector<int> roots;
    deque<node> tree;
    distinct_query(const vector<int>& arr) : N(arr.size()), roots(N + 1, 0) {
        tree.push_back({0, 0, 0}); //acts as null
        map<int, int> last_idx;
        for (int i = 0; i < N; i++) {
            roots[i + 1] = update(roots[i], 0, N, last_idx[arr[i]]);
            last idx[arr[i]] = i + 1:
        }
```

```
int update(int v, int tl, int tr, int idx) {
        if (tr - tl == 1) {
            tree.push_back({tree[v].sum + 1, 0, 0});
            return tree.size() - 1;
        int tm = t1 + (tr - t1) / 2:
        int lch = tree[v].lch;
        int rch = tree[v].rch;
        if (idx < tm)
            lch = update(lch, tl, tm, idx);
            rch = update(rch, tm, tr, idx);
        tree.push_back({tree[lch].sum + tree[rch].sum, lch, rch});
        return tree.size() - 1;
    //returns number of distinct elements in range [l,r)
    int query(int 1, int r) const {
        assert(0 <= 1 && 1 <= r && r <= N);
        return query(roots[1], roots[r], 0, N, 1 + 1);
    int query(int vl, int vr, int tl, int tr, int idx) const {
        if (tree[vr].sum == 0 || idx <= tl)</pre>
            return 0;
        if (tr <= idx)</pre>
            return tree[vr].sum - tree[vl].sum;
        int tm = tl + (tr - tl) / 2;
        return query(tree[v1].lch, tree[vr].lch, tl, tm, idx) +
               query(tree[v1].rch, tree[vr].rch, tm, tr, idx);
    }
};
```

Listing 22: Merge Sort Tree

```
//cat merge_sort_tree.h | ./hash.sh
//ffe2d7
#pragma once
//For point updates: either switch to policy based BST, or use sqrt decomposition
struct merge_sort_tree {
    const int N:
   vector<vector<int>> tree:
    //RTE's when 'arr' is empty
        for (int i = 0, j = 1 << __lg(2 * N - 1); i < N; i++, j = (j + 1) % N + N)
            tree[j] = {arr[i]};
        for (int i = N - 1; i >= 1; i--)
            merge(tree[2 * i].begin(), tree[2 * i].end(), tree[2 * i + 1].begin(),

    tree[2 * i + 1].end(), back_inserter(tree[i]));
    int value(int v, int x) const {
        return lower_bound(tree[v].begin(), tree[v].end(), x) - tree[v].begin();
    int range_idx(int i) const {
        i += 1 << __lg(2 * N - 1);
        return i < 2 * N ? i : 2 * (i - N);
    //How many values in range [l, r) are \langle x \rangle
    //O(\log^2(n))
    int query(int 1, int r, int x) const {
        int res = 0;
```

```
for (l = range_idx(l), r = range_idx(r); l < r; l >>= 1, r >>= 1) {
      if (l & 1) res += value(l++, x);
      if (r & 1) res += value(--r, x);
    }
    return res;
}
```

Listing 23: STRINGS

Listing 24: Suffix Array

```
//cat suffix_array.h / ./hash.sh
//46840a
#pragma once
//source: https://judge.yosupo.jp/submission/37410
//0(n)
//mnemonic: Suffix Array Induced Sorting
template<class T> vector<int> sa_is(const T& s, int upper/*max element of 's'; for
    \hookrightarrow std::string, pass in 255*/) {
    int n = (int)s.size();
    if (n == 0) return {};
    if (n == 1) return {0}:
    if (n == 2) {
        if (s[0] < s[1]) {
            return {0, 1};
        } else {
            return {1, 0};
   }
    vector<int> sa(n);
    vector<bool> ls(n);
    for (int i = n - 2; i >= 0; i--)
        ls[i] = (s[i] == s[i + 1]) ? ls[i + 1] : (s[i] < s[i + 1]);
    vector<int> sum_l(upper + 1), sum_s(upper + 1);
    for (int i = 0; i < n; i++) {</pre>
        if (!ls[i])
            sum_s[s[i]]++;
        else
            sum_l[s[i] + 1]++;
   }
   for (int i = 0; i <= upper; i++) {</pre>
        sum_s[i] += sum_l[i];
        if (i < upper) sum_l[i + 1] += sum_s[i];</pre>
    vector<int> buf(upper + 1);
    auto induce = [&](const vector<int>& lms) {
        fill(sa.begin(), sa.end(), -1);
        fill(buf.begin(), buf.end(), 0);
        copy(sum_s.begin(), sum_s.end(), buf.begin());
        for (auto d : lms) {
            if (d == n) continue:
            sa[buf[s[d]]++] = d;
        copy(sum_l.begin(), sum_l.end(), buf.begin());
        sa[buf[s[n-1]]++] = n-1;
        for (int i = 0; i < n; i++) {
            int v = sa[i];
```

```
if (v >= 1 && !ls[v - 1])
            sa[buf[s[v - 1]] ++] = v - 1:
    }
    copy(sum_l.begin(), sum_l.end(), buf.begin());
    for (int i = n - 1; i \ge 0; i--) {
        int v = sa[i];
        if (v >= 1 && ls[v - 1])
            sa[--buf[s[v-1]+1]] = v-1;
    }
}:
vector<int> lms_map(n + 1, -1);
int m = 0:
for (int i = 1; i < n; i++) {</pre>
    if (!ls[i - 1] && ls[i])
        lms_map[i] = m++;
vector<int> lms;
lms.reserve(m);
for (int i = 1; i < n; i++) {
    if (!ls[i - 1] && ls[i])
        lms.push_back(i);
}
induce(lms);
if (m) {
    vector<int> sorted_lms;
    sorted_lms.reserve(m);
    for (int v : sa) {
        if (lms_map[v] != -1) sorted_lms.push_back(v);
    vector<int> rec_s(m);
    int rec_upper = 0;
    rec_s[lms_map[sorted_lms[0]]] = 0;
    for (int i = 1; i < m; i++) {</pre>
        int l = sorted_lms[i - 1], r = sorted_lms[i];
        int end_1 = (lms_map[1] + 1 < m) ? lms[lms_map[1] + 1] : n;</pre>
        int end_r = (lms_map[r] + 1 < m) ? lms[lms_map[r] + 1] : n;
        bool same = true:
        if (end 1 - 1 != end r - r)
            same = false;
        else {
            while (1 < end 1) {
                if (s[1] != s[r])
                    break;
                1++;
                r++;
            if (1 == n || s[1] != s[r]) same = false;
        if (!same) rec_upper++;
        rec_s[lms_map[sorted_lms[i]]] = rec_upper;
    auto rec sa =
        sa_is(rec_s, rec_upper);
    for (int i = 0; i < m; i++)</pre>
        sorted_lms[i] = lms[rec_sa[i]];
    induce(sorted_lms);
return sa;
```

Listing 25: LCP

```
//cat lcp.h | ./hash.sh
//064842
#pragma once
//source: https://judge.yosupo.jp/submission/37410
//mnemonic: Longest Common Prefix
//NOLINTNEXTLINE(readability-identifier-naming)
template<class T> vector<int> LCP(const T& s, const vector<int>& sa) {
   int n = s.size(), k = 0;
   vector<int> lcp(n, 0);
   vector<int> rank(n, 0);
   for (int i = 0; i < n; i++) rank[sa[i]] = i;
   for (int i = 0; i < n; i++, k ? k-- : 0) {
       if (rank[i] == n - 1) {
           k = 0:
            continue;
       }
       int j = sa[rank[i] + 1];
       while (i + k < n \&\& j + k < n \&\& s[i + k] == s[j + k]) k++;
       lcp[rank[i]] = k;
   }
   return lcp;
```

Listing 26: Prefix Function

```
//cat prefix_function.h / ./hash.sh
//aa0518
#pragma once
//source: https://cp-algorithms.com/string/prefix-function.html#implementation
template <class T> vector<int> prefix_function(const T& s) {
    int n = s.size();
    vector<int> pi(n, 0);
    for (int i = 1; i < n; i++) {
        int j = pi[i - 1];
        while (j > 0 && s[i] != s[j]) j = pi[j - 1];
        pi[i] = j + (s[i] == s[j]);
    }
    return pi;
}
```

Listing 27: KMP

```
//cat kmp.h | ./hash.sh
//9d70ad

#pragma once
//mnemonic: Knuth Morris Pratt
#include "prefix_function.h"
//usage:
// string needle;
// ...
// KMP kmp(needle);
//or
// vector<int> needle;
// ...
// KMP kmp(needle);
// kmp doubling trick: to check if 2 arrays are rotationally equivalent: run kmp
//with one array as the needle and the other array doubled (excluding the first
//8 last characters) as the haystack or just use kactl's min rotation code
```

```
//NOLINTNEXTLINE(readability-identifier-naming)
template <class T> struct KMP {
    KMP(const T& a_needle) : pi(prefix_function(a_needle)), needle(a_needle) {}
    // if haustack = "bananas"
    // needle = "ana"
    // then we find 2 matches:
    // bananas
    // _ana___
    // ana
    // 0123456 (indexes)
    // and KMP::find returns {1,3} - the indexes in haystack where
    // each match starts.
    // You can also pass in false for "all" and KMP::find will only
    // return the first match: {1}. Useful for checking if there exists
    // some match:
    //
    // KMP::find(<haystack>,false).size() > 0
    vector<int> find(const T& haystack, bool all = true) const {
        vector<int> matches;
        for (int i = 0, j = 0; i < (int)haystack.size(); i++) {</pre>
            while (j > 0 && needle[j] != haystack[i]) j = pi[j - 1];
            if (needle[j] == haystack[i]) j++;
            if (j == (int)needle.size()) {
                matches.push_back(i - (int)needle.size() + 1);
                if (!all) return matches;
                j = pi[j - 1];
        }
        return matches;
    vector<int> pi;//prefix function
    T needle:
};
```

Listing 28: Trie

```
//cat trie.h | ./hash.sh
//10777e
#pragma once
//status: not tested
//source: https://cp-algorithms.com/string/aho_corasick.html#construction-of-the-trie
//intended to be a base template and to be modified
const int K = 26;//alphabet size
struct trie {
   const char MIN_CH = 'a';//'A' for uppercase, '0' for digits
   struct node {
        int next[K], id, p = -1;
       char ch;
        bool leaf = 0:
        node(int a_p = -1, char a_ch = '#') : p(a_p), ch(a_ch) {
            fill(next, next + K, -1);
       }
   };
   vector<node> t:
   trie(): t(1) {}
   void add_string(const string& s, int id) {
        int c = 0:
       for (char ch : s) {
```

```
int v = ch - MIN CH:
            if (t[c].next[v] == -1) {
                t[c].next[v] = t.size();
                t.emplace_back(c, ch);
            c = t[c].next[v];
        t[c].leaf = 1;
        t[c].id = id;
   }
    void remove_string(const string& s) {
        int c = 0:
        for (char ch : s) {
            int v = ch - MIN_CH;
            if (t[c].next[v] == -1)
                return;
            c = t[c].next[v];
        }
        t[c].leaf = 0;
   }
    int find_string(const string& s) const {
        int c = 0:
        for (char ch : s) {
            int v = ch - MIN_CH;
            if (t[c].next[v] == -1)
                return -1;
            c = t[c].next[v];
        if (!t[c].leaf) return -1;
        return t[c].id;
   }
};
```

Listing 29: Binary Trie

```
//cat binary_trie.h | ./hash.sh
//33aa3a
#pragma once
struct binary_trie {
    const int MX_BIT = 62;
    struct node {
        long long val = -1;
        int sub_sz = 0;//number of inserted values in subtree
        array<int, 2> next = {-1, -1};
   };
    vector<node> t;
   binary_trie() : t(1) {}
    //delta = 1 to insert val, -1 to remove val, 0 to get the # of val's in this data
        \hookrightarrow structure
    int update(long long val, int delta) {
        int c = 0:
        t[0].sub_sz += delta;
        for (int bit = MX_BIT; bit >= 0; bit--) {
            bool v = (val >> bit) & 1:
            if (t[c].next[v] == -1) {
                t[c].next[v] = t.size();
                t.emplace_back();
            c = t[c].next[v];
            t[c].sub_sz += delta;
```

```
}
        t[c].val = val:
        return t[c].sub_sz;
    int size() const {
        return t[0].sub_sz;
    //returns x such that:
    // x is in this data structure
    // value of (x ^val) is minimum
    long long min_xor(long long val) const {
        assert(size() > 0);
        int c = 0;
        for (int bit = MX_BIT; bit >= 0; bit--) {
            bool v = (val >> bit) & 1;
            int ch = t[c].next[v];
            if (ch != -1 && t[ch].sub_sz > 0)
                c = ch:
            else
                c = t[c].next[!v];
        return t[c].val;
    }
};
```

Listing 30: Longest Common Prefix Query

```
//cat lcp_queries.h / ./hash.sh
//ed1840
#pragma once
#include "suffix_array.h"
#include "lcp.h"
#include "../range_data_structures/rmq.h"
//computes suffix array, lcp array, and then sparse table over lcp array
//0(n \log n)
struct lcp_queries {
    lcp_queries(const string& s) : sa(sa_is(s, 255)), inv_sa(s.size()), lcp(LCP(s, sa)),
         \hookrightarrow st(lcp, [](int x, int y) {
        return min(x, y);
    }) {
        for (int i = 0; i < (int)s.size(); i++)</pre>
            inv_sa[sa[i]] = i;
    //length of longest common prefix of suffixes s[idx1 ... n), s[idx2 ... n), 0-based
    //You can check if two substrings s[l1..r1), s[l2..r2) are equal in O(1) by:
    //r1-l1 == r2-l2 \ \theta \theta \ longest\_common\_prefix(l1, l2) >= r1-l1
    int longest_common_prefix(int idx1, int idx2) const {
        if (idx1 == idx2) return (int)sa.size() - idx1;
        idx1 = inv sa[idx1]:
        idx2 = inv_sa[idx2];
        if (idx1 > idx2) swap(idx1, idx2);
        return st.query(idx1, idx2);
    //returns true if suffix s[idx1 ... n) < s[idx2 ... n)
    //(so\ false\ if\ idx1 == idx2)
    bool less(int idx1, int idx2) const {
        return inv_sa[idx1] < inv_sa[idx2];</pre>
```

```
vector<int> sa, inv_sa, lcp;
RMQ<int> st;
};
```

Listing 31: MATH

Listing 32: BIN EXP MOD

```
//cat exp_mod.h / ./hash.sh
//3be256
#pragma once
//returns (base^pw)/mod in O(log(pw)), but returns 1 for 0^0
//What if base doesn't fit in long long?
//Since (base^pw)%mod == ((base%mod)^pw)%mod we can calculate base under mod of 'mod'
//What if pw doesn't fit in long long?
//case 1: mod is prime
//(base^pw)%mod == (base^(pw%(mod-1)))%mod (from Fermat's little theorem)
//so calculate pw under mod of 'mod-1'
//note 'mod-1' is not prime, so you need to be able to calculate 'pw%(mod-1)' without
    \hookrightarrow division
//case 2: non-prime mod
//let t = totient(mod)
//if pw >= log2(mod) then (base^pw)/mod == (base^(t+(pw/t)))/mod (proof)
    ← https://cp-algorithms.com/algebra/phi-function.html#generalization)
//so calculate pw under mod of 't'
//incidentally, totient(p) = p - 1 for every prime p, making this a more generalized
    \hookrightarrow version of case 1
int pow(long long base, long long pw, int mod) {
    assert(0 <= pw && 0 <= base && 1 <= mod);
    int res = 1;
   base %= mod;
    while (pw > 0) {
        if (pw & 1) res = res * base % mod;
        base = base * base % mod;
        pw >>= 1;
   }
    return res;
```

Listing 33: Fibonacci

Listing 34: Matrix Mult and Pow

```
//cat matrix_expo.h / ./hash.sh
//2edd34
#pragma once
//empty matrix -> RTE
vector<vector<int>> mult(const vector<vector<int>>& a, const vector<vector<int>>& b, int
    \hookrightarrow mod) {
    assert(a[0].size() == b.size());
    int n = a.size(), m = b[0].size(), inner = b.size();
    vector<vector<int>> prod(n, vector<int>(m, 0));
    for (int i = 0; i < n; i++) {</pre>
        for (int k = 0; k < inner; k++) {
            for (int j = 0; j < m; j++)
                prod[i][j] = (prod[i][j] + 1LL * a[i][k] * b[k][j]) % mod;
    return prod;
vector<vector<int>> power(vector<int>> mat/*intentional pass by value*/, long
    \hookrightarrow long pw, int mod) {
    int n = mat.size();
    vector<vector<int>> prod(n, vector<int>(n, 0));
    for (int i = 0; i < n; i++)
        prod[i][i] = 1;
    while (pw > 0) {
        if (pw % 2 == 1) prod = mult(prod, mat, mod);
        mat = mult(mat, mat, mod);
        pw /= 2;
   return prod;
```

Listing 35: N Choose K MOD

```
//cat n_choose_k_mod.h / ./hash.sh
//f3a1a9
#pragma once
//for mod inverse
#include "exp_mod.h"
// usage:
       n\_choose\_k \ nk(n, 1e9+7) to use 'choose', 'inv' with inputs < n
//
// or:
       n_choose_k nk(mod, mod) to use 'choose_with_lucas_theorem' with arbitrarily large
     \hookrightarrow inputs
struct n_choose_k {
    n_choose_k(int n, int a_mod) : mod(a_mod), fact(n, 1), inv_fact(n, 1) {
        //this implementation doesn't work if n > mod because n! % mod = 0 when n > mod

→ mod. So 'inv_fact' array will be all 0's
        assert(max(n, 2) <= mod);</pre>
        //assert mod is prime. mod is intended to fit inside an int so that
        //multiplications fit in a longlong before being modded down. So this
        //will take sqrt(2^31) time
        for (int i = 2; i * i <= mod; i++) assert(mod % i);</pre>
        for (int i = 2; i < n; i++)
            fact[i] = 1LL * fact[i - 1] * i % mod;
        inv_fact.back() = pow(fact.back(), mod - 2, mod);
        for (int i = n - 2; i \ge 2; i - -)
            inv_fact[i] = inv_fact[i + 1] * (i + 1LL) % mod;
    //classic n choose k
```

```
//fails when n \ge mod
int choose(int n. int k) const {
    if (k < 0 \mid | k > n) return 0;
    //now we know 0 <= k <= n so 0 <= n
    return 1LL * fact[n] * inv_fact[k] % mod * inv_fact[n - k] % mod;
}
//lucas theorem to calculate n choose k in O(log(k))
//need to calculate all factorials in range [0,mod), so 0(mod) time&space, so need
     \hookrightarrow smallish prime mod (< 1e6 maybe)
//handles n >= mod correctly
int choose_with_lucas_theorem(long long n, long long k) const {
    if (k < 0 \mid | k > n) return 0;
    if (k == 0 || k == n) return 1;
    return 1LL * choose_with_lucas_theorem(n / mod, k / mod) * choose(n % mod, k %
         \hookrightarrow mod) % mod:
}
//returns inverse of n in O(1)
int inv(int n) const {
    assert(1 <= n); //don't divide by 0 :)</pre>
    return 1LL * fact[n - 1] * inv_fact[n] % mod;
}
int mod:
vector<int> fact, inv_fact;
```

Listing 36: Partitions

```
//cat partitions.h / ./hash.sh
//3356f6
#pragma once
//https://oeis.org/A000041
//0 (n sqrt n) time, but small-ish constant factor (there does exist a 0 (n log n)
     \hookrightarrow solution too)
vector<int> partitions(int n/*size of dp array*/, int mod) {
    vector<int> dp(n, 1);
   for (int i = 1; i < n; i++) {
        long long sum = 0;
        for (int j = 1, pent = 1, sign = 1; pent <= i; j++, pent += 3 * j - 2, sign =
            if (pent + j <= i) sum += dp[i - pent - j] * sign + mod;</pre>
            sum += dp[i - pent] * sign + mod;
        dp[i] = sum % mod;
   }
    return dp;
```

Listing 37: Derangements

```
//cat derangements.h | ./hash.sh  
//c221bb  
#pragma once  
//https://oeis.org/A000166  
//
//for a permutation of size i:  
//there are (i-1) places to move 0 to not be at index 0. Let's say we moved 0 to index j  
\hookrightarrow (j>0).  
//If we move value j to index 0 (forming a cycle of length 2), then there are dp[i-2]  
\hookrightarrow derangements of the remaining i-2 elements
```

```
//else there are dp[i-1] derangements of the remaining i-1 elements (including j)
vector<int> derangements(int n/*size of dp array*/, int mod) {
    vector<int> dp(n, 0);
    dp[0] = 1;
    for (int i = 2; i < n; i++)
        dp[i] = 1LL * (i - 1) * (dp[i - 1] + dp[i - 2]) % mod;
    return dp;
}</pre>
```

Listing 38: Prime Sieve Mobius

```
//cat prime_sieve_mobius.h | ./hash.sh
//42052f
#pragma once
//mobius[i] = 0 iff there exists a prime p s.t. i\%(p^2)=0
//mobius[i] = -1 iff i has an odd number of distinct prime factors
//mobius[i] = 1 iff i has an even number of distinct prime factors
const int N = 2e6 + 10;
int mobius[N]:
void calc_mobius() {
    mobius[1] = 1;
    for (int i = 1; i < N; i++)</pre>
        for (int j = i + i; j < N; j += i)
            mobius[j] -= mobius[i];
//a_prime[val] = some random prime factor of 'val'
//to check if 'val' is prime:
// if (a_prime[val] == val)
//to get all prime factors of a number 'val' in O(log(val)):
// while (val > 1) {
//
       int p = a_prime[val];
//
       //p is some prime factor of val
//
        val /= p:
// }
int a prime[N]:
void calc seive() {
    iota(a_prime, a_prime + N, 0);
    for (int i = 2; i * i < N; i++)</pre>
        if (a_prime[i] == i)
            for (int j = i * i; j < N; j += i)
                a_prime[j] = i;
```

Listing 39: Row Reduce

//cat solve_linear_mod.h | ./hash.sh

```
int n = mat.size(), m = mat[0].size(), rank = 0, det = 1;
assert(cols <= m):</pre>
for (int col = 0; col < cols && rank < n; col++) {</pre>
    //find arbitrary pivot and swap pivot to current row
    for (int i = rank; i < n; i++)</pre>
        if (mat[i][col] != 0) {
            if (rank != i) det = det == 0 ? 0 : mod - det;
            swap(mat[i], mat[rank]);
            break;
    if (mat[rank][col] == 0) {
        det = 0:
        continue;
    det = (1LL * det * mat[rank][col]) % mod;
    //make pivot 1 by dividing row by inverse of pivot
    int a_inv = pow(mat[rank][col], mod - 2, mod);
    for (int j = 0; j < m; j++)
        mat[rank][j] = (1LL * mat[rank][j] * a_inv) % mod;
    //zero-out all numbers above & below pivot
    for (int i = 0; i < n; i++)
        if (i != rank && mat[i][col] != 0) {
            int val = mat[i][col];
            for (int j = 0; j < m; j++) {</pre>
                mat[i][j] -= 1LL * mat[rank][j] * val % mod;
                 if (mat[i][j] < 0) mat[i][j] += mod;</pre>
    rank++;
}
assert(rank <= min(n, cols));</pre>
return {rank, det};
```

Listing 40: Solve Linear Equations MOD

```
//44cc6e
#pragma once
#include "row_reduce.h"
struct matrix info {
    int rank, det;
    vector<int> x;
//Solves\ mat\ *\ x\ =\ b\ under\ prime\ mod.
//mat is a n (rows) by m (cols) matrix, b is a length n column vector, x is a length m
    \hookrightarrow vector.
//assumes n.m >= 1. else RTE
//Returns rank of mat, determinant of mat, and x (solution vector to mat *x = b).
//x is empty if no solution. If rank < m, there are multiple solutions and an arbitrary
    \hookrightarrow one is returned.
//Leaves mat in reduced row echelon form (unlike kactl) with b appended.
//O(n * m * min(n.m))
matrix_info solve_linear_mod(vector<vector<int>>& mat, const vector<int>& b, int mod) {
    assert(mat.size() == b.size());
    int n = mat.size(), m = mat[0].size();
   for (int i = 0; i < n; i++)
        mat[i].push_back(b[i]);
    auto [rank, det] = row_reduce(mat, m, mod); //row reduce not including the last column
    //check if solution exists
```

```
for (int i = rank; i < n; i++) {
      if (mat[i].back() != 0) return {rank, det, {} }; //no solution exists
}
      //initialize solution vector ('x') from row-reduced matrix
vector<int> x(m, 0);
for (int i = 0, j = 0; i < rank; i++) {
      while (mat[i][j] == 0) j++; //find pivot column
      x[j] = mat[i].back();
}
return {rank, det, x};
}</pre>
```

Listing 41: Matrix Inverse

```
//cat matrix_inverse.h | ./hash.sh
//3056ad
#pragma once
#include "row_reduce.h"
//returns inverse of square matrix mat, empty if no inverse
vector<vector<int>> matrix_inverse(vector<int>> mat/*intentional pass by value*/,
    \hookrightarrow int mod) {
    int n = mat.size();
    assert(n == (int)mat[0].size());
    //append identity matrix
    for (int i = 0; i < n; i++) {
        mat[i].resize(2 * n, 0);
        mat[i][i + n] = 1;
    auto [rank, det] = row_reduce(mat, n, mod);//row reduce first n columns, leaving
         \hookrightarrow inverse in last n columns
    if (rank < n) return {}: //no inverse
    for (int i = 0; i < n; i++)
        mat[i].erase(mat[i].begin(), mat[i].begin() + n);
    return mat:
```

Listing 42: Euler's Totient Phi Function

```
//cat totient.h / ./hash.sh
//36bd41
#pragma once
//Euler's totient function counts the positive integers
//up to a given integer n that are relatively prime to n.
//
//To improve, use Pollard-rho to find prime factors
int totient(int n) {
   int res = n;
   for (int i = 2; i * i <= n; i++) {
      if (n % i == 0) {
        while (n % i == 0) n /= i;
        res -= res / i;
      }
   }
   if (n > 1) res -= res / n;
   return res;
}
```

Listing 43: MAX FLOW

Listing 44: Dinic

```
//cat dinic.h / ./hash.sh
//21e5b5
#pragma once
//status: not tested
struct max_flow {
    typedef long long 11;
    11 n, s, t;
    max_flow(int a_n, int a_s, int a_t) : n(a_n), s(a_s), t(a_t), d(n), ptr(n), q(n),
        \hookrightarrow g(n) \{\}
    void add_edge(ll a, ll b, ll cap) {
        edge_map[a * n + b] = e.size();
        edge e1 = \{a, b, cap, 0\};
        edge e2 = \{ b, a, 0, 0 \};
        g[a].push_back((ll) e.size());
        e.push_back(e1);
        g[b].push_back((ll) e.size());
        e.push_back(e2);
   }
   11 get_flow() {
        11 \text{ flow} = 0;
        for (::) {
            if (!bfs()) break;
            ptr.assign(ptr.size(), 0);
            while (ll pushed = dfs(s, INF))
                flow += pushed;
        }
        return flow;
   }
    ll get_flow_for_edge(ll a, ll b) {
        return e[edge_map[a * n + b]].flow;
    const 11 INF = 1e18;
    struct edge {
        ll a, b, cap, flow;
    };
    unordered_map<int, 11> edge_map;
    vector<ll> d, ptr, q;
    vector<edge> e;
    vector<vector<ll>> g;
    bool bfs() {
        11 qh = 0, qt = 0;
        q[qt++] = s;
        d.assign(d.size(), -1);
        d[s] = 0;
        while (qh < qt \&\& d[t] == -1) {
            11 v = q[qh++];
            for (size_t i = 0; i < g[v].size(); i++) {</pre>
                11 id = g[v][i],
                   to = e[id].b;
                if (d[to] == -1 && e[id].flow < e[id].cap) {</pre>
                    q[qt++] = to;
                     d[to] = d[v] + 1;
                }
            }
        return d[t] != -1;
   }
    11 dfs(ll v, ll flow) {
        if (!flow) return 0;
```

```
if (v == t) return flow;
    for (; ptr[v] < (l1) g[v].size(); ptr[v]++) {
        ll id = g[v][ptr[v]];
        ll to = e[id].b;
        if (d[to] != d[v] + 1) continue;
        ll pushed = dfs(to, min(flow, e[id].cap - e[id].flow));
        if (pushed) {
            e[id].flow += pushed;
            e[id ^ 1].flow -= pushed;
            return pushed;
        }
    }
    return 0;
}</pre>
```

Listing 45: Hungarian

```
//cat hungarian.h | ./hash.sh
//85c306
#pragma once
//source: https://e-maxx.ru/algo/assignment_hungary
//input: cost[1...n][1...m] with 1 <= n <= m
//n workers, indexed 1, 2, ..., n
//m jobs, indexed 1, 2, ..., m
//it costs 'cost[i][i]' to assign worker i to job j (1<=i<=n, 1<=j<=m)
//this returns *min* total cost to assign each worker to some distinct job
//0(n^2 * m)
//trick 1: set 'cost[i][j]' to INF to say: "worker 'i' cannot be assigned job 'j'"
//trick 2: 'cost[i][j]' can be negative, so to instead find max total cost over all
     \hookrightarrow matchings: set all 'cost[i][j]' to '-cost[i][j]'.
//Now max total cost = - hungarian(cost).min_cost
const long long INF = 1e18;
struct match {
    long long min_cost;
    vector<int> matching; //worker 'i' (1<=i<=n) is assigned to job 'matching[i]'
         \hookrightarrow (1<=matching[i]<=m)
match hungarian(const vector<vector<long long>>& cost) {
    int n = cost.size() - 1, m = cost[0].size() - 1;
    assert(n <= m);</pre>
    vector<int> p(m + 1), way(m + 1);
    vector<long long> u(n + 1), v(m + 1);
    for (int i = 1; i <= n; i++) {
        p[0] = i;
        int j0 = 0;
         vector<long long> minv(m + 1, INF);
         vector<bool> used(m + 1, false);
         do {
             used[j0] = true;
             int i0 = p[j0], j1 = 0;
             long long delta = INF;
             for (int j = 1; j \le m; j++)
                 if (!used[j]) {
                     long long cur = cost[i0][j] - u[i0] - v[j];
                     if (cur < minv[j])</pre>
                         minv[j] = cur, way[j] = j0;
                     if (minv[j] < delta)</pre>
```

```
delta = minv[j], j1 = j;
        for (int j = 0; j \le m; j++)
            if (used[j])
                u[p[j]] += delta, v[j] -= delta;
            else
                minv[j] -= delta;
        j0 = j1;
    } while (p[j0] != 0);
    do {
        int | 1 = way[| 0];
        p[j0] = p[j1];
        j0 = j1;
    } while (j0);
}
vector<int> ans(n + 1);
for (int j = 1; j \le m; j++)
    ans[p[j]] = j;
return {-v[0], ans};
```

Listing 46: Min Cost Max Flow

```
//cat min_cost_max_flow.h / ./hash.sh
//a60fe5
#pragma once
//status: not tested
const long long INF = 1e18;
struct min_cost_max_flow {
   typedef long long 11;
   struct edge {
       ll a, b, cap, cost, flow;
       size_t back;
   };
   vector<edge> e;
   vector<vector<ll>> g;
   11 n, s, t;
   11 k = INF; // max amount of flow allowed
   min_cost_max_flow(int a_n, int a_s, int a_t) : n(a_n), s(a_s), t(a_t) {
       g.resize(n);
   }
   void add_edge(ll a, ll b, ll cap, ll cost) {
       edge e1 = {a, b, cap, cost, 0, g[b].size() };
       edge e2 = {b, a, 0, -cost, 0, g[a].size() };
       g[a].push_back((ll) e.size());
       e.push_back(e1);
       g[b].push_back((ll) e.size());
       e.push_back(e2);
   // returns {flow, cost}
   pair<11, 11> get_flow() {
       11 flow = 0, cost = 0;
       while (flow < k) {</pre>
           vector<ll> id(n, 0), d(n, INF), q(n), p(n);
            vector<size_t> p_edge(n);
           11 qh = 0, qt = 0;
           q[qt++] = s;
            d[s] = 0;
            while (qh != qt) {
               11 v = q[qh++];
```

```
id[v] = 2:
                if (qh == n) qh = 0;
                for (size_t i = 0; i < g[v].size(); i++) {</pre>
                    edge& r = e[g[v][i]];
                    if (r.flow < r.cap && d[v] + r.cost < d[r.b]) {
                        d[r.b] = d[v] + r.cost;
                        if (id[r.b] == 0) {
                             q[qt++] = r.b;
                             if (qt == n) qt = 0;
                        } else if (id[r.b] == 2) {
                             if (--qh == -1) qh = n - 1;
                             q[qh] = r.b;
                        id[r.b] = 1;
                        p[r.b] = v;
                        p_{edge}[r.b] = i;
                    }
                }
            if (d[t] == INF) break;
            11 addflow = k - flow;
            for (11 v = t; v != s; v = p[v]) {
                11 pv = p[v];
                size_t pr = p_edge[v];
                addflow = min(addflow, e[g[pv][pr]].cap - e[g[pv][pr]].flow);
            for (11 v = t; v != s; v = p[v]) {
                11 pv = p[v];
                size_t pr = p_edge[v], r = e[g[pv][pr]].back;
                e[g[pv][pr]].flow += addflow;
                e[g[v][r]].flow -= addflow;
                cost += e[g[pv][pr]].cost * addflow;
            flow += addflow;
        return {flow, cost};
    }
};
```

Listing 47: MISC

Listing 48: DSU

```
//cat dsu.h / ./hash.sh
//gb3c97
#pragma once
//mnemonic: Disjoint Set Union
//NOLINTNEXTLINE(readability-identifier-naming)
struct DSU {
   int num_sets;
   vector<int> par;
   DSU(int n) : num_sets(n), par(n, -1) {}
   DSU(const DSU& rhs) : num_sets(rhs.num_sets), par(rhs.par) {}
   int find(int x) {
      return par[x] < 0 ? x : par[x] = find(par[x]);
   }
   int size_of_set(int x) {
      return -par[find(x)];
}</pre>
```

```
}
bool join(int x, int y) {
    if ((x = find(x)) == (y = find(y))) return false;
    if (par[y] < par[x]) swap(x, y);
    par[x] += par[y];
    par[y] = x;
    num_sets--;
    return true;
}
</pre>
```

Listing 49: PBDS

```
//cat policy_based_data_structures.h / ./hash.sh
//807de9
#pragma once
//status: not tested
//place these includes *before* the '#define int long long' else compile error
//not using <bits/extc++.h> as it compile errors on codeforces c++20 compiler
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
//BST with extra functions https://codeforces.com/blog/entry/11080
//order_of_key - # of elements *strictly* less than given element
//find_by_order - find kth largest element, k is 0 based so find_by_order(0) returns min
    \hookrightarrow element
template<class T> using indexed_set = tree<T, null_type, less<T>, rb_tree_tag,

    tree_order_statistics_node_update>;

//example initialization:
indexed_set<pair<long long, int>> is;
//hash table (apparently faster than unordered_map):
    \hookrightarrow https://codeforces.com/blog/entry/60737
//example initialization:
gp_hash_table<string, long long> ht;
```

Listing 50: Monotonic Stack

```
//cat monotonic stack.h | ./hash.sh
//90f107
#pragma once
//calculates array 'left' with:
//for every index j with left[i] < j < i: arr[j] > arr[i]
//arr[left[i]] \leftarrow arr[i] if left[i] != -1
//trick: pass in vector<pair<T/*value*/, int/*index*/>> with arr[i].second = i (0<=i<n)
     \hookrightarrow to simulate arr[i] >= arr[i]
//0(n)
template<class T> vector<int> monotonic_stack(const vector<T>& arr) {
    int n = arr.size();
    vector<int> left(n);
    for (int i = 0; i < n; i++) {</pre>
        int& j = left[i] = i - 1;
        while (j >= 0 && arr[j] > arr[i]) j = left[j];
   }
    return left;
```

Listing 51: Count Rectangles

```
//cat count_rectangles.h | ./hash.sh
//9873d2
#pragma once
#include "monotonic_stack.h"
//qiven a 2D boolean matrix, calculate cnt[i][j]
//cnt[i][j] = the number of times an i-by-j rectangle appears in the matrix such that
    \hookrightarrow all i*j cells in the rectangle are true
//Note cnt[0][j] and cnt[i][0] will contain garbage values
vector<vector<int>> count_rectangles(const vector<vector<bool>>& grid) {
    int n = grid.size(), m = grid[0].size();
    vector<vector<int>> cnt(n + 1, vector<int>(m + 1, 0));
    vector<int> arr(m, 0);
    auto rv = [&](int j) -> int {//reverse
        return m - 1 - j;
   };
    for (int i = 0; i < n; i++) {</pre>
        vector<pair<int, int>> arr_rev(m);
        for (int j = 0; j < m; j++) {
            arr[j] = grid[i][j] * (arr[j] + 1);
            arr_rev[rv(j)] = {arr[j], j};
        }
        vector<int> left = monotonic_stack(arr);
        vector<int> right = monotonic_stack(arr_rev);
        for (int j = 0; j < m; j++) {
            int l = j - left[j] - 1, r = rv(right[rv(j)]) - j - 1;
            cnt[arr[j]][l + r + 1]++;
            cnt[arr[i]][1]--:
            cnt[arr[i]][r]--;
    for (int i = 1; i <= n; i++)
        for (int k = 0; k < 2; k++)
            for (int j = m; j > 1; j--)
                cnt[i][j - 1] += cnt[i][j];
    for (int j = 1; j <= m; j++)</pre>
        for (int i = n; i > 1; i--)
            cnt[i - 1][j] += cnt[i][j];
   return cnt;
```

Listing 52: LIS

```
//cat lis.h | ./hash.sh
//a243e1
#pragma once
//returns array of indexes representing the longest *strictly* increasing subsequence
//for non-decreasing: pass in a vector<pair<T, int>> with arr[i].second = i (0<=i<n)
//alternatively, there's this https://codeforces.com/blog/entry/13225
//mnemonic: Longest Increasing Subsequence
//NOLINTNEXTLINE(readability-identifier-naming)
template<class T> vector<int> LIS(const vector<T>& arr) {
    if (arr.empty()) return {};
    vector<int> dp{0}/*array of indexes into 'arr'*/, prev(arr.size(), -1);
    for (int i = 1; i < (int)arr.size(); i++) {</pre>
        auto it = lower_bound(dp.begin(), dp.end(), i, [&](int x, int y) -> bool {
            return arr[x] < arr[y];</pre>
        });
        if (it == dp.end()) {
```

```
prev[i] = dp.back();
    dp.push_back(i);
} else {
    prev[i] = it == dp.begin() ? -1 : *(it - 1);
        *it = i;
}
    //here, dp.size() = length of LIS of prefix of arr ending at index i
}
vector<int> res(dp.size());
for (int i = dp.back(), j = dp.size(); i != -1; i = prev[i])
    res[--j] = i;
return res;
}
```

Listing 53: Safe Hash

```
//cat safe_hash.h | ./hash.sh
//d9ea53
#pragma once
//status: not tested
//source: https://codeforces.com/blog/entry/62393
struct custom_hash {
    static uint64_t splitmix64(uint64_t x) {
        // http://xorshift.di.unimi.it/splitmix64.c
        x += 0x9e3779b97f4a7c15;
        x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
        x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
        return x ^ (x >> 31);
   }
   size_t operator()(uint64_t x) const {
        static const uint64_t FIXED_RANDOM =

    chrono::steady_clock::now().time_since_epoch().count();

        return splitmix64(x + FIXED_RANDOM);
   }
};
unordered_map<long long, int, custom_hash> safe_map;
#include "policy_based_data_structures.h"
gp_hash_table<long long, int, custom_hash> safe_hash_table;
```