South Dakota Mines				Pa	
Listings		38	Max Rectangle in Histogram	13	
			39	Monotonic Stack	. 13
1	CONTEST	1	40	GCD Convolution	. 13
2	Tips and Tricks	1	41	Iterate Chooses	. 13
3	Hash codes	1	42	Iterate Submasks	. 14
4	Test on random inputs	2	43	Iterate Supermasks	. 14
5	MAX FLOW	2	44	Number of Distinct Subsequences DP	. 14
6	Hungarian	2	45	PBDS	. 14
7	Min Cost Max Flow	2	46	Random	. 14
8	GRAPHS	3	47	Safe Hash	. 15
9	Block Vertex Tree	3	48	RANGE DATA STRUCTURES	. 15
10	Bridge Tree	3	49	Number Distinct Elements	. 15
11	Bridges and Cuts	4	50	Implicit Lazy Segment Tree	. 15
12	Strongly Connected Components	4	51	Kth Smallest	. 16
13	Centroid Decomposition	5	52	Merge Sort Tree	. 17
14	Frequency Table of Tree Distance	5	53	BIT	. 17
15	Count Paths Per Node	6	54	RMQ	. 17
16	Dijkstra	6	55	Disjoint RMQ	. 18
17	HLD	6	56	Lazy Segment Tree	. 18
18	Hopcroft Karp	7	57	STRINGS	. 19
19	Kth Node on Path	8	58	Binary Trie	. 19
20	LCA	8	59	Prefix Function	
21	Rooted Tree Isomorphism	9	60	KMP String Matching	. 20
22	MATH	9	61	Suffix and LCP Arrays	. 20
23	Derangements	9	62	Suffix Array Related Queries	. 21
24	Binary Exponentiation MOD	9	63	Palindrome Query	. 22
25	Fibonacci	9	64	Trie	. 22
26	Matrix Multiplication	9			
27	Mobius Inversion	10			
28	N Choose K MOD	10			
29	Partitions	10			
30	Prime Sieve	10			
31	Row Reduce	11			
32	Solve Linear Equations MOD	11			
33	Euler's Totient Phi Function				
34	Tetration MOD	12			
35	MISC				
36	Cartesian Tree	12			

### CONTEST

# Tips and Tricks

```
## Tips and Tricks
- [C++ tips and tricks](https://codeforces.com/blog/entry/74684)
- invokes RTE (Run Time Error) upon integer overflow
#pragma GCC optimize "trapv"
- invoke RTE for input error (e.g. reading a long long into an int)
cin.exceptions(cin.failbit);
- use pramgas for C++ speed boost
#pragma GCC optimize("03,unroll-loops")
#pragma GCC target("avx2,bmi,bmi2,lzcnt,popcnt")
### Troubleshooting
/* stuff you should look for
   * int overflow, array bounds
   * special cases (n=1?)
   * do smth instead of nothing and stay organized
   * WRITE STUFF DOWN
   * DON'T GET STUCK ON ONE APPROACH
- - -
Author: Beng
refer to [KACTL
    → Troubleshoot](https://github.com/kth-competitive-programming/kactl/blob/main/content/contest/toubleshoot.txt)
## Sources
- [[Tutorial] GCC Optimization Pragmas](https://codeforces.com/blog/entry/96344)
- [Don't use rand(): a guide to random number generators in
```

### Hash codes

```
#! /bin/bash
#Hashes a file, ignoring all:
# - whitespace
# - comments
# - asserts
# - includes
# - pragmas
#Use to verify that code was correctly typed.

#usage:
# chmod +x hash.sh
# cat a.cpp | ./hash.sh
#or just copy this command:
```

### Test on random inputs

```
#! /bin/bash
#runs 2 programs against each other on random inputs until they output different results
#source: https://github.com/Errichto/youtube/blob/master/testing/s.sh
#usage:
# chmod +x test.sh
# ./test.sh
for((i = 1; ; ++i)); do
    echo $i
    ./test.out > in
    diff --ignore-all-space <(./a.out < in) <(./brute.out < in) || break
done</pre>
```

#### MAX FLOW

### Hungarian

```
//cat hungarian.hpp | ./hash.sh
//935a16
#pragma once
const long long INF = 1e18;
struct weighted_match {
    long long min_cost;
    vector<int> matching;//worker `i` (1<=i<=n) is assigned to job `matching[i]`</pre>
         \hookrightarrow (1<=matching[i]<=m)
 * Given cost[1...n][1...m] with 1 <= n <= m.
* n workers, indexed 1, 2, ..., n
* m jobs, indexed 1, 2, ..., m
 * It costs `cost[i][j]` to assign worker i to job j (1<=i<=n, 1<=j<=m). This
* returns *min* sum of costs to assign each worker to some distinct job.
* Set `cost[i][j]` to INF to say: "worker i cannot be assigned job j".
* This works for negatives, so negating cost array gives max matching.
* @see https://e-maxx.ru/algo/assignment_hungary
 * @time O(n^2 * m)
* @memory O(n * m)
*/
weighted_match hungarian(const vector<vector<long long>>& cost) {
    int n = ssize(cost) - 1, m = ssize(cost[0]) - 1;
    assert(n <= m);</pre>
    vector<int> p(m + 1), way(m + 1);
    vector<long long> u(n + 1), v(m + 1);
    for (int i = 1; i \le n; i++) {
        p[0] = i;
        int i0 = 0;
        vector<long long> minv(m + 1, INF);
        vector<bool> used(m + 1, 0);
```

```
do {
        used[j0] = 1;
        int i0 = p[j0], j1 = 0;
        long long delta = INF;
        for (int j = 1; j \le m; j++)
            if (!used[j]) {
                long long cur = cost[i0][j] - u[i0] - v[j];
                if (cur < minv[j])</pre>
                     minv[j] = cur, way[j] = j0;
                if (minv[j] < delta)</pre>
                     delta = minv[j], j1 = j;
        for (int j = 0; j \le m; j++)
            if (used[j])
                u[p[j]] += delta, v[j] -= delta;
            else
                minv[j] -= delta;
        j0 = j1;
    } while (p[j0] != 0);
    do {
        int j1 = way[j0];
        p[j0] = p[j1];
        j0 = j1;
    } while (j0);
vector<int> ans(n + 1);
for (int j = 1; j \le m; j++)
    ans[p[j]] = j;
return {-v[0], ans};
```

### Min Cost Max Flow

```
//cat min_cost_max_flow.hpp | ./hash.sh
//9dd6b6
#pragma once
const long long INF = 1e18;
* @see https://e-maxx.ru/algo/min_cost_flow
struct mcmf {
   using 11 = long long;
   struct edge {
       int a, b;
       11 cap, cost, flow;
       int back;
   };
   const int N;
   vector<edge> e;
   vector<vector<int>> g;
   mcmf(int a_n) : N(a_n), g(N) {}
   void add_edge(int a, int b, ll cap, ll cost) {
       edge e1 = \{a, b, cap, cost, 0, ssize(g[b])\};
       edge e2 = {b, a, 0, -cost, 0, ssize(g[a])};
       g[a].push_back(ssize(e));
       e.push_back(e1);
       g[b].push_back(ssize(e));
        e.push_back(e2);
```

```
pair<ll, ll> get_flow(int s, int t, ll total_flow) {
        11 flow = 0, cost = 0;
        while (flow < total_flow) {</pre>
            vector<ll> d(N, INF);
            vector\langle int \rangle p_edge(N), id(N, 0), q(N), p(N);
            int qh = 0, qt = 0;
            q[qt++] = s;
            d[s] = 0;
            while (qh != qt) {
                int v = q[qh++];
                id[v] = 2;
                if (qh == N) qh = 0;
                for (int i = 0; i < ssize(g[v]); i++) {
                    const edge& r = e[g[v][i]];
                    if (r.flow < r.cap && d[v] + r.cost < d[r.b]) {
                        d[r.b] = d[v] + r.cost;
                        if (id[r.b] == 0) {
                             q[qt++] = r.b;
                             if (qt == N) qt = 0;
                        } else if (id[r.b] == 2) {
                             if (--qh == -1) qh = N - 1;
                             q[qh] = r.b;
                        }
                        id[r.b] = 1;
                        p[r.b] = v;
                        p_edge[r.b] = i;
                }
            if (d[t] == INF) break;
            11 addflow = total_flow - flow;
            for (int v = t; v != s; v = p[v]) {
                int pv = p[v], pr = p_edge[v];
                addflow = min(addflow, e[g[pv][pr]].cap - e[g[pv][pr]].flow);
            for (int v = t; v != s; v = p[v]) {
                int pv = p[v], pr = p_edge[v], r = e[g[pv][pr]].back;
                e[g[pv][pr]].flow += addflow;
                e[g[v][r]].flow -= addflow;
                cost += e[g[pv][pr]].cost * addflow;
            flow += addflow;
        return {flow, cost};
    }
};
```

### **GRAPHS**

#### Block Vertex Tree

```
//cat block_vertex_tree.hpp | ./hash.sh
//a5c2b9
#pragma once
#include "bridges_and_cuts.hpp"
/**
```

```
* Returns adjacency list of block vertex tree.
* @code{.cpp}
      graph_info cc = bridge_and_cut(adj, m);
      vector<vector<int>>> bvt = block_vertex_tree(adj, cc);
      //to loop over each *unique* bcc containing a node v:
      for (int bccid : bvt[v]) {
          bccid -= n:
      //to loop over each *unique* node inside a bcc:
      for (int v : bvt[bccid + n]) {}
* @endcode
* @time O(n + m)
* @memory O(n + m)
*/
vector<vector<int>> block_vertex_tree(const vector<vector<pair<int, int>>>& adj, const
    int n = ssize(adj);
   vector<vector<int>> bvt(n + cc.num_bccs);
   vector<bool> vis(cc.num_bccs, 0);
   for (int v = 0; v < n; v++) {
       for (auto [_, e_id] : adj[v]) {
           int bccid = cc.bcc_id[e_id];
           if (!vis[bccid]) {
               vis[bccid] = 1;
               bvt[v].push_back(bccid + n);//add edge between original node, and bcc node
               bvt[bccid + n].push_back(v);
           }
       for (int bccid : bvt[v]) vis[bccid - n] = 0;
   }
   return bvt;
```

# Bridge Tree

```
//cat bridge_tree.hpp | ./hash.sh
//8eb014
#pragma once
#include "bridges_and_cuts.hpp"
* Never adds multiple edges as bridges_and_cuts.hpp correctly marks them as
* non-bridges.
* @code{.cpp}
       graph_info cc = bridge_and_cut(adj, m);
       vector<vector<int>>> bt = bridge_tree(adj, cc);
* @endcode
* @time O(n + m)
* @memory O(n + m)
vector<vector<int>>> bridge_tree(const vector<vector<pair<int, int>>>& adj, const graph_info&
    \hookrightarrow cc) {
   vector<vector<int>> tree(cc.num_2_edge_ccs);
    for (int i = 0; i < ssize(adj); i++)
        for (auto [to, e_id] : adj[i])
            if (cc.is_bridge[e_id])
                tree[cc.two_edge_ccid[i]].push_back(cc.two_edge_ccid[to]);
   return tree;
```

### Bridges and Cuts

```
//cat bridges_and_cuts.hpp | ./hash.sh
//3f21b9
#pragma once
struct graph_info {
    //2 edge connected component stuff (e.g. components split by bridge edges)
    int num_2_edge_ccs;
    vector<bool> is_bridge;//edge id -> 1 iff bridge edge
    vector<int> two_edge_ccid;//node -> id of 2 edge component (which are labeled 0, 1, ...,
         \hookrightarrow `num_2_edge_ccs `-1)
    //bi connected component stuff (e.g. components split by cut/articulation nodes)
    int num_bccs;
    vector<bool> is_cut;//node -> 1 iff cut node
    vector<int> bcc_id;//edge id -> id of bcc (which are labeled 0, 1, ..., `num_bccs`-1)
};
/**
* @code{.cpp}
       //example initialization of `adj`:
       for (int i = 0; i < m; i++) {
           int u, v;
           cin >> u >> v;
           u--, v--;
           adj[u].emplace_back(v, i);
           adj[v].emplace_back(u, i);
 * @endcode
 * @see https://cp-algorithms.com/graph/bridge-searching.html
      https://cp-algorithms.com/graph/cutpoints.html
 * @time O(n + m)
* @memory O(n + m)
graph_info bridge_and_cut(const vector<vector<pair<int/*neighbor*/, int/*edge id*/>>>&
     → adj/*undirected graph*/, int m/*number of edges*/) {
    //stuff for both (always keep)
    int n = ssize(adj), timer = 1;
    vector<int> tin(n, 0);
    //2 edge cc stuff (delete if not needed)
    int num_2_edge_ccs = 0;
    vector<bool> is_bridge(m, 0);
    vector<int> two_edge_ccid(n), node_stack;
    node_stack.reserve(n);
    //bcc stuff (delete if not needed)
    int num_bccs = 0;
    vector<bool> is_cut(n, 0);
    vector<int> bcc_id(m), edge_stack;
    edge_stack.reserve(m);
    auto dfs = [&](auto self, int v, int p_id) -> int {
        int low = tin[v] = timer++, deg = 0;
        node_stack.push_back(v);
        for (auto [to, e_id] : adj[v]) {
            if (e_id == p_id) continue;
            if (!tin[to]) {
                edge_stack.push_back(e_id);
                int low_ch = self(self, to, e_id);
                if (low_ch >= tin[v]) {
                    is_cut[v] = 1;
                    while (1) {
                        int edge = edge_stack.back();
```

```
edge_stack.pop_back();
                    bcc_id[edge] = num_bccs;
                    if (edge == e_id) break;
                }
                num_bccs++;
            low = min(low, low_ch);
            deg++;
        } else if (tin[to] < tin[v]) {</pre>
            edge_stack.push_back(e_id);
            low = min(low, tin[to]);
        }
    if (p_id == -1) is_cut[v] = (deg > 1);
    if (tin[v] == low) {
        if (p_id != -1) is_bridge[p_id] = 1;
        while (1) {
            int node = node_stack.back();
            node_stack.pop_back();
            two_edge_ccid[node] = num_2_edge_ccs;
            if (node == v) break;
        num_2_edge_ccs++;
    }
    return low;
};
for (int i = 0; i < n; i++)
    if (!tin[i])
        dfs(dfs, i, -1);
return {num_2_edge_ccs, is_bridge, two_edge_ccid, num_bccs, is_cut, bcc_id};
```

# Strongly Connected Components

```
//cat strongly_connected_components.hpp | ./hash.sh
//f6f849
#pragma once
//source: https://github.com/kth-competitive-programming/
// kactl/blob/main/content/graph/SCC.h
struct scc_info {
   int num_sccs;
   //scc's are labeled 0,1,..., `num_sccs-1`
   //scc_id[i] is the id of the scc containing node `i`
   //for each edge i -> j: scc_id[i] >= scc_id[j] (reverse topo order of scc's)
   vector<int> scc_id;
//NOLINTNEXTLINE(readability-identifier-naming)
scc_info SCC(const vector<vector<int>>& adj/*directed, unweighted graph*/) {
    int n = ssize(adj), timer = 1, num_sccs = 0;
   vector<int> tin(n, 0), scc_id(n, -1), node_stack;
   node_stack.reserve(n);
   auto dfs = [&](auto self, int v) -> int {
       int low = tin[v] = timer++;
       node_stack.push_back(v);
        for (int to : adj[v])
            if (scc_id[to] < 0)</pre>
               low = min(low, tin[to] ? tin[to] : self(self, to));
       if (tin[v] == low) {
```

# Centroid Decomposition

```
//cat centroid_decomposition.hpp | ./hash.sh
//71e2e5
#pragma once
* Given an unweighted, undirected forest and a function, centroid_decomp runs
 * the function on the centroid of every decomposition.
 * @code{.cpp}
      //example usage
       centroid_decomp decomp(adj, [&](const vector<vector<int>>& adj_removed_edges, int cent)
      \hookrightarrow -> void {
      });
 * @endcode
 * @time O(n log n)
 * @memory O(n)
template <typename F> struct centroid_decomp {
    vector<vector<int>> adj;
    F func:
    vector<int> sub_sz;
    centroid_decomp(const vector<vector<int>>& a_adj/*undirected forest*/, const F& a_func)
        : adj(a_adj), func(a_func), sub_sz(ssize(adj), -1) {
        for (int i = 0; i < ssize(adj); i++)</pre>
            if (sub_sz[i] == -1)
                dfs(i);
    }
    void calc_subtree_sizes(int u, int p = -1) {
        sub_sz[u] = 1;
        for (int v : adj[u]) {
            if (v == p) continue;
            calc_subtree_sizes(v, u);
            sub_sz[u] += sub_sz[v];
        }
    void dfs(int u) {
        calc_subtree_sizes(u);
        for (int p = -1, sz_root = sub_sz[u];;) {
            auto big_ch = find_if(adj[u].begin(), adj[u].end(), [&](int v) -> bool {
                return v != p && 2 * sub_sz[v] > sz_root;
            if (big_ch == adj[u].end()) break;
```

# Frequency Table of Tree Distance

```
//cat count_paths_per_length.hpp | ./hash.sh
//182b70
#pragma once
#include "../../kactl/content/numerical/FastFourierTransform.h"
#include "centroid_decomposition.hpp"
* Returns array `num_paths` where `num_paths[i]` = # of paths in tree with `i`
* edges.
* @time O(n log^2 n)
vector<long long> count_paths_per_length(const vector<vector<int>>& adj/*unrooted, connected
    \hookrightarrow tree*/) {
   vector<long long> num_paths(ssize(adj), 0);
   centroid_decomp decomp(adj, [&](const vector<vector<int>>& adj_removed_edges, int cent) ->
        vector<vector<double>> child_depths;
        for (int to : adj_removed_edges[cent]) {
            child_depths.emplace_back(1, 0.0);
            for (queue<pair<int, int>> q({{to, cent}}); !q.empty();) {
                child_depths.back().push_back(ssize(q));
                queue<pair<int, int>> new_q;
                while (!q.empty()) {
                    auto [curr, par] = q.front();
                    q.pop();
                    for (int ch : adj_removed_edges[curr]) {
                        if (ch == par) continue;
                        new_q.emplace(ch, curr);
                    }
                swap(q, new_q);
           }
        sort(child_depths.begin(), child_depths.end(), [&](const auto & x, const auto & y) {
           return x.size() < y.size();</pre>
       });
        vector<double> total_depth(1, 1.0);
        for (const auto& cnt_depth : child_depths) {
            auto prod = conv(total_depth, cnt_depth);
            for (int i = 1; i < ssize(prod); i++)
                num_paths[i] += llround(prod[i]);
            total_depth.resize(ssize(cnt_depth), 0.0);
            for (int i = 1; i < ssize(cnt_depth); i++)</pre>
                total_depth[i] += cnt_depth[i];
   });
```

```
return num_paths;
}
```

#### Count Paths Per Node

```
//cat count_paths_per_node.hpp | ./hash.sh
//2077d2
#pragma once
#include "centroid_decomposition.hpp"
* Returns array `num_paths` where `num_paths[i]` = number of paths with k
 * edges where node `i` is on the path. 0-based nodes.
 * @time O(n log n)
vector<long long> count_paths_per_node(const vector<vector<int>>& adj/*unrooted tree*/, int k)
    vector<long long> num_paths(ssize(adj));
    centroid_decomp decomp(adj, [&](const vector<vector<int>>& adj_removed_edges, int cent) ->
         \hookrightarrow void {
        vector<int> pre_d(1, 1), cur_d(1);
        auto dfs = [&](auto self, int u, int p, int d) -> long long {
            if (d > k) return 0;
            if (ssize(cur_d) <= d) cur_d.push_back(0);</pre>
            cur_d[d]++;
            long long cnt = 0;
            if (k - d < ssize(pre_d)) cnt += pre_d[k - d];</pre>
            for (int v : adj_removed_edges[u])
                if (v != p)
                     cnt += self(self, v, u, d + 1);
            num_paths[u] += cnt;
            return cnt;
        auto dfs_child = [&](int child) -> long long {
            long long cnt = dfs(dfs, child, cent, 1);
            pre_d.resize(ssize(cur_d));
            for (int i = 1; i < ssize(cur_d) && cur_d[i]; i++)</pre>
                pre_d[i] += cur_d[i], cur_d[i] = 0;
            return cnt;
        for (int child : adj_removed_edges[cent])
            num_paths[cent] += dfs_child(child);
        pre_d = vector<int>(1);
        cur_d = vector<int>(1);
        for (auto it = adj_removed_edges[cent].rbegin(); it != adj_removed_edges[cent].rend();
             \hookrightarrow it++)
            dfs_child(*it);
    });
    return num_paths;
```

# Dijkstra

```
//cat dijkstra.hpp | ./hash.sh
//aa6eda
#pragma once
const long long INF = 1e18;
/**
```

```
* Returns array 'len' where 'len[i]' = shortest path from node 'start' to node
* 'i'. For example 'len[start]' = 0.
* @time O((n + m) \log n) Note \log(m) < \log(n^2) = 2*\log(n), so O(\log n) = 2*\log(n)
* @memory O(n + m)
*/
vector<long long> dijkstra(const vector<vector<pair<int, long long>>>& adj /*directed or

    undirected, weighted graph*/, int start) {

   using node = pair<long long, int>;
   vector<long long> len(ssize(adj), INF);
   len[start] = 0;
   priority_queue<node, vector<node>, greater<node>> q;
   q.emplace(0, start);
   while (!q.empty()) {
       auto [curr_len, v] = q.top();
       q.pop();
       if (len[v] < curr_len) continue;//important check: O(n*m) without it
       for (auto [to, weight] : adj[v])
           if (len[to] > weight + len[v]) {
               len[to] = weight + len[v];
               q.emplace(len[to], to);
           }
   }
   return len;
```

### HLD

```
//cat hld.hpp | ./hash.sh
//d30c4a
#pragma once
* @see https://codeforces.com/blog/entry/53170
//NOLINTNEXTLINE(readability-identifier-naming)
struct HLD {
   struct node {
       int sub_sz = 1, par = -1, time_in = -1, next = -1;
   };
   vector<node> tree;
   /**
    * @time O(n)
    * @memory O(n)
   HLD(vector<vector<int>>& adj/*forest of unrooted trees*/) : tree(ssize(adj)) {
       for (int i = 0, timer = 0; i < ssize(adj); i++) {
           if (tree[i].next == -1) {//lowest indexed node in each tree becomes root
               tree[i].next = i;
               dfs1(i, adj);
               dfs2(i, adj, timer);
           }
       }
   }
   void dfs1(int v, vector<vector<int>>& adj) {
       auto par = find(adj[v].begin(), adj[v].end(), tree[v].par);
       if (par != adj[v].end()) adj[v].erase(par);
       for (int& to : adj[v]) {
```

```
tree[to].par = v;
            dfs1(to, adj);
            tree[v].sub_sz += tree[to].sub_sz;
            if (tree[to].sub_sz > tree[adj[v][0]].sub_sz)
                swap(to, adj[v][0]);
        }
    }
    void dfs2(int v, const vector<vector<int>>& adj, int& timer) {
        tree[v].time_in = timer++;
        for (int to : adi[v]) {
            tree[to].next = (timer == tree[v].time_in + 1 ? tree[v].next : to);
            dfs2(to, adj, timer);
    }
    /**
     * Returns inclusive-exclusive intervals (of time_in's) corresponding to
     * the path between u and v, not necessarily in order. u, v must be in the
     * same component.
     * @time O(log n)
     * @memory O(log n)
    vector<pair<int, int>> path(int u, int v) const {
        vector<pair<int, int>> res;
        for (;; v = tree[tree[v].next].par) {
            if (tree[v].time_in < tree[u].time_in) swap(u, v);</pre>
            if (tree[tree[v].next].time_in <= tree[u].time_in) {</pre>
                res.emplace_back(tree[u].time_in, tree[v].time_in + 1);
                return res;
            }
            res.emplace_back(tree[tree[v].next].time_in, tree[v].time_in + 1);
    }
    pair<int, int> subtree(int i) const {
        return {tree[i].time_in, tree[i].time_in + tree[i].sub_sz};
    /**
     * u, v must be in the same component.
     * @time O(log n)
     * @memory 0(1)
    int lca(int u, int v) const {
        for (;; v = tree[tree[v].next].par) {
            if (tree[v].time_in < tree[u].time_in) swap(u, v);</pre>
            if (tree[tree[v].next].time_in <= tree[u].time_in) return u;</pre>
    }
};
```

# Hopcroft Karp

```
//cat hopcroft_karp.hpp | ./hash.sh
//5d1682
#pragma once
struct match {
    //# of edges in matching (which = size of min vertex cover by öKnig's theorem)
    int size_of_matching;
    //an arbitrary max matching is found. For this matching:
    //if l_to_r[node_left] == -1:
```

```
// node_left is not in matching
   // the edge `node_left` <=> l_to_r[node_left] is in the matching
   //
   //similarly for r_to_l with edge r_to_l[node_right] <=> node_right in matching if
        \hookrightarrow r_to_l[node_right] != -1
   //matchings stored in l_to_r and r_to_l are the same matching
   //provides way to check if any node/edge is in matching
   vector<int> l_to_r, r_to_l;
   //an arbitrary min vertex cover is found. For this mvc: mvc_l[node_left] is 1 iff

→ node_left is in the min vertex cover (same for mvc_r)
   //if mvc_l[node_left] is 0, then node_left is in the corresponding maximal independent set
   vector<bool> mvc_1, mvc_r;
* Think of the bipartite graph as having a left side (with size lsz) and a
* right side (with size rsz).
* Nodes on left side are indexed 0,1,...,lsz-1.
* Nodes on right side are indexed 0,1,...,rsz-1.
* To initialize `adj`: For every edge node_left <=> node_right, do:
* adj[node_left].push_back(node_right)
* @see https://github.com/foreverbell/acm-icpc-cheat-sheet/
      blob/master/src/graph-algorithm/hopcroft-karp.cpp
* @time O(m * sqrt(n))
* @memory O(n ^ (3/2)) Some note about the complexity.
match hopcroft_karp(const vector<vector<int>>& adj/*bipartite graph*/, int rsz/*number of
    \hookrightarrow nodes on right side*/) {
   int size_of_matching = 0, lsz = ssize(adj);
   vector<int> l_to_r(lsz, -1), r_to_l(rsz, -1);
   while (1) {
       queue<int> q;
        vector<int> level(lsz, -1);
        for (int i = 0; i < lsz; i++)
           if (l_to_r[i] == -1)
               level[i] = 0, q.push(i);
       bool found = 0:
       vector<bool> mvc_l(lsz, 1), mvc_r(rsz, 0);
       while (!q.empty()) {
           int u = q.front();
           q.pop();
           mvc_1[u] = 0;
           for (int x : adj[u]) {
               mvc_r[x] = 1;
               int v = r_to_1[x];
               if (v == -1) found = 1;
               else if (level[v] == -1) {
                   level[v] = level[u] + 1;
                    q.push(v);
               }
           }
       if (!found) return {size_of_matching, l_to_r, r_to_l, mvc_l, mvc_r};
       auto dfs = [&](auto self, int u) -> bool {
            for (int x : adj[u]) {
               int v = r to l[x]:
               if (v == -1 || (level[u] + 1 == level[v] && self(self, v))) {
                   l_{to_r[u]} = x;
```

#### Kth Node on Path

```
//cat kth_node_on_path.hpp | ./hash.sh
//c59307
#pragma once
#include "lca.hpp"
struct kth_node_on_path {
    LCA lca;
    kth_node_on_path(const vector<vector<pair<int, long long>>>& adj/*forest of weighted
         \hookrightarrow trees*/) : lca(adj) {}
     * Consider path {u, u's par, ..., LCA(u,v), ..., v's par, v}. This returns
     * the node at index k. So k=0 returns u, k=#path_edges returns v.
     * u, v must be in the same component.
     * @time O(log n)
     */
    int query(int u, int v, int k) const {
        int lca_uv = lca.get_lca(u, v);
        int u_lca = lca.tree[u].depth - lca.tree[lca_uv].depth;
        int v_lca = lca.tree[v].depth - lca.tree[lca_uv].depth;
        assert(0 <= k && k <= u_lca + v_lca);
        return k <= u_lca ? lca.kth_par(u, k) : lca.kth_par(v, u_lca + v_lca - k);</pre>
    }
};
```

#### LCA

```
//cat lca.hpp | ./hash.sh
//b28532
#pragma once
/**
* @see https://codeforces.com/blog/entry/74847
//NOLINTNEXTLINE(readability-identifier-naming)
struct LCA {
    struct node {
        int jmp = -1, jmp\_edges = 0, par = -1, depth = 0;
        long long dist = 0LL;
    };
    vector<node> tree:
    /**
     * @time O(n)
     * @memory O(n)
    LCA(const vector<vector<pair<int, long long>>>& adj/*forest of weighted trees*/) :
         \hookrightarrow tree(ssize(adj)) {
```

```
for (int i = 0; i < ssize(adj); i++) {</pre>
        if (tree[i].jmp == -1) {//lowest indexed node in each tree becomes root
            tree[i].jmp = i;
            dfs(i, adj);
        }
    }
}
void dfs(int v, const vector<vector<pair<int, long long>>>& adj) {
    int jmp, jmp_edges;
    if (tree[v].jmp != v && tree[v].jmp_edges == tree[tree[v].jmp].jmp_edges)
        jmp = tree[tree[v].jmp].jmp, jmp_edges = 2 * tree[v].jmp_edges + 1;
    else
        jmp = v, jmp\_edges = 1;
    for (auto [ch, w] : adj[v]) {
        if (ch == tree[v].par) continue;
        tree[ch] = {
            jmp,
            jmp_edges,
            1 + tree[v].depth,
            w + tree[v].dist
        };
        dfs(ch, adj);
}
 * Traverse up k edges. So with k=1 this returns v's parent.
 * @time O(log k)
int kth_par(int v, int k) const {
    k = min(k, tree[v].depth);
    while (k > 0) {
        if (tree[v].jmp_edges <= k) {</pre>
            k -= tree[v].jmp_edges;
            v = tree[v].jmp;
        } else {
            k--;
            v = tree[v].par;
        }
    }
    return v;
}
/**
 * x, y must be in the same component.
 * @time O(log n)
int get_lca(int x, int y) const {
    if (tree[x].depth < tree[y].depth) swap(x, y);</pre>
    x = kth_par(x, tree[x].depth - tree[y].depth);
    while (x != y) {
        if (tree[x].jmp != tree[y].jmp)
            x = tree[x].jmp, y = tree[y].jmp;
        else
            x = tree[x].par, y = tree[y].par;
    }
    return x;
}
int dist_edges(int x, int y) const {
    return tree[x].depth + tree[y].depth - 2 * tree[get_lca(x, y)].depth;
```

```
}
long long dist_weight(int x, int y) const {
    return tree[x].dist + tree[y].dist - 2 * tree[get_lca(x, y)].dist;
}
};
```

### Rooted Tree Isomorphism

```
//cat subtree_isomorphism.hpp | ./hash.sh
//455aef
#pragma once
struct iso_info {
    int num_distinct_subtrees; //0 <= id[i] < num_distinct_subtrees</pre>
    vector<int> id; //id[u] == id[v] iff rooted subtree u is isomorphic to rooted subtree v
};
/**
* Classifies each rooted subtree by isomorphism.
* @time O(n log n)
* @memory O(n)
iso_info subtree_iso(const vector<vector<int>>& adj) {
    vector<int> id(ssize(adj), -1);
    map<vector<int>, int> hashes;
    auto dfs = [&](auto self, int u, int p) -> int {
        vector<int> ch_ids;
        ch_ids.reserve(ssize(adj[u]));
        for (int v : adj[u]) {
            if (v != p)
                ch_ids.push_back(self(self, v, u));
        sort(ch_ids.begin(), ch_ids.end());
        auto it = hashes.find(ch_ids);
        if (it == hashes.end())
            return id[u] = hashes[ch_ids] = ssize(hashes);
        return id[u] = it->second;
    for (int i = 0; i < ssize(adj); i++)</pre>
        if (id[i] == -1)
            dfs(dfs, i, i);
    return {ssize(hashes), id};
```

### MATH

# Derangements

```
//cat derangements.hpp | ./hash.sh
//64d325
#pragma once
/**
    * @see https://oeis.org/A000166
    * @time O(n)
    */
vector<long long> derangements(int n, long long mod) {
    vector<long long> dp(n, 0);
    dp[0] = 1;
```

### Binary Exponentiation MOD

```
//cat binary_exponentiation_mod.hpp | ./hash.sh
//92a3ef
#pragma once
* Returns (base^pw)%mod; returns 1 for 0^0.
* @time O(log pw)
*/
long long bin_exp(long long base, long long pw, long long mod) {
   assert(0 <= pw && 0 <= base && 1 <= mod);</pre>
   long long res = 1;
   base %= mod;
   while (pw > 0) {
       if (pw & 1) res = res * base % mod;
       base = base * base % mod;
       pw >>= 1;
   }
   return res;
```

### Fibonacci

# Matrix Multiplication

#### Mobius Inversion

#### N Choose K MOD

```
//cat n_choose_k_mod.hpp | ./hash.sh
//daaaa8
#pragma once
#include "binary_exponentiation_mod.hpp"
* @code{.cpp}
      n_choose_k nk(n, 1e9+7); // to use `choose`, `inv` with inputs strictly < n</pre>
      n_choose_k nk(mod, mod); // to use `choose_lucas` with arbitrarily large inputs
* @endcode
 */
struct n_choose_k {
    long long mod;
    vector<long long> fact, inv_fact;
     * Only works for `n <= mod` and prime mod.
     * @time O(n + sqrt(mod)) The sqrt is only to assert mod is prime.
     * @memory O(n)
    n_choose_k(int n, long long a_mod) : mod(a_mod), fact(n, 1), inv_fact(n, 1) {
        assert(max(n, 2) <= mod);</pre>
        for (int i = 2; i * i <= mod; i++) assert(mod % i);</pre>
        for (int i = 2; i < n; i++)
            fact[i] = fact[i - 1] * i % mod;
        inv_fact.back() = bin_exp(fact.back(), mod - 2, mod);
        for (int i = n - 2; i >= 2; i--)
```

```
inv_fact[i] = inv_fact[i + 1] * (i + 1) % mod;
   }
    /**
    * n choose k for n.k < ssize(fact). Fails when n >= mod.
    * @time O(1)
     * @memory O(n) precomp
    long long choose(int n, int k) const {
        if (k < 0 \mid | k > n) return 0;
        //now we know 0 <= k <= n so 0 <= n
        return fact[n] * inv_fact[k] % mod * inv_fact[n - k] % mod;
   }
    /**
    * Lucas theorem - n choose k for n, k up to LLONG_MAX. Handles n>=mod
    * correctly.
     * @time O(log(k))
     * @memory O(mod) precomp, so can't use 1e9 + 7.
     */
    long long choose_lucas(long long n, long long k) const {
        if (k < 0 \mid | k > n) return 0;
        if (k == 0 | | k == n) return 1;
        return choose_lucas(n / mod, k / mod) * choose(int(n % mod), int(k % mod)) % mod;
   }
   /**
    * Returns x such that x * n \% \mod == 1.
    * @time O(1)
    */
   long long inv(int n) const {
        assert(1 <= n); //don't divide by 0 :)</pre>
        return fact[n - 1] * inv_fact[n] % mod;
   }
};
```

# \_\_\_\_\_ Row Redu

```
Partitions
//cat partitions.hpp | ./hash.sh
//e7ae42
#pragma once
//https://oeis.org/A000041
//O(n sqrt n) time, but small-ish constant factor (there does exist a O(n log n) solution as
vector<long long> partitions(int n, long long mod) {
   vector<long long> dp(n, 1);
   for (int i = 1; i < n; i++) {
       long long sum = 0;
       for (int j = 1, pent = 1, sign = 1; pent <= i; j++, pent += 3 * j - 2, sign = -sign) {
           if (pent + j \le i) sum += dp[i - pent - j] * sign + mod;
           sum += dp[i - pent] * sign + mod;
       dp[i] = sum \% mod;
   }
   return dp;
```

### Prime Sieve

```
//cat prime_sieve.hpp | ./hash.sh
```

```
//25a877
#pragma once
bool is_prime(int val, const vector<int>& sieve) {
    assert(val < ssize(sieve));</pre>
    return val >= 2 && sieve[val] == val;
vector<int> get_prime_factors(int val, const vector<int>& sieve) {
    assert(val < ssize(sieve));</pre>
    vector<int> factors;
    while (val > 1) {
        int p = sieve[val];
        factors.push_back(p);
        val /= p;
    }
    return factors;
/**
* Returns array `sieve` where `sieve[i]` = some prime factor of `i`.
* @time O(n * log(logn))
*/
vector<int> get_sieve(int n) {
    vector<int> sieve(n);
    iota(sieve.begin(), sieve.end(), 0);
    for (int i = 2; i * i < n; i++)
        if (sieve[i] == i)
            for (int j = i * i; j < n; j += i)
                sieve[i] = i;
    return sieve;
```

### Row Reduce

```
//cat row reduce.hpp | ./hash.sh
//c812f1
#pragma once
#include "binary_exponentiation_mod.hpp"
* First `cols` columns of mat represents a matrix to be left in reduced row
 * echelon form. Row operations will be performed to all later columns.
* @code{.cpp}
       auto [rank, det] = row_reduce(mat, ssize(mat[0]), mod);
* @endcode
* @time O(n * m * min(cols, n))
* @memorv O(n * m)
pair<int/*rank*/, long long/*determinant*/> row_reduce(vector<vector<long long>>& mat, int
     \hookrightarrow cols, long long mod) {
    int n = ssize(mat), m = ssize(mat[0]), rank = 0;
    long long det = 1;
    assert(cols <= m);</pre>
    for (int col = 0; col < cols && rank < n; col++) {</pre>
        //find arbitrary pivot and swap pivot to current row
        for (int i = rank; i < n; i++)
            if (mat[i][col] != 0) {
                if (rank != i) det = det == 0 ? 0 : mod - det;
                swap(mat[i], mat[rank]);
                break;
```

```
if (mat[rank][col] == 0) {
        det = 0:
        continue:
    det = det * mat[rank][col] % mod;
    //make pivot 1 by dividing row by inverse of pivot
    long long a_inv = bin_exp(mat[rank][col], mod - 2, mod);
    for (int j = 0; j < m; j++)
        mat[rank][j] = mat[rank][j] * a_inv % mod;
    //zero-out all numbers above & below pivot
    for (int i = 0; i < n; i++)
        if (i != rank && mat[i][col] != 0) {
            long long val = mat[i][col];
            for (int j = 0; j < m; j++) {
                mat[i][j] -= mat[rank][j] * val % mod;
                if (mat[i][j] < 0) mat[i][j] += mod;</pre>
        }
    rank++;
}
assert(rank <= min(n, cols));</pre>
return {rank, det};
```

# Solve Linear Equations MOD

```
//cat solve_linear_mod.hpp | ./hash.sh
//eba3e8
#pragma once
#include "row_reduce.hpp"
struct matrix_info {
   int rank;
   long long det:
   //solution vector
   //empty if no solution exists
   //if there are multiple solutions, this is an arbitrary one
   vector<long long> x;
* Solves mat * x = b under prime mod.
* `mat` is a n (rows) by m (cols) matrix, b is a length n column vector, x is
* Assumes n,m >= 1, else RTE. Leaves mat in reduced row echelon form with b
* appended.
* Number of unique solutions = (size of domain) ^ (# of free variables).
* (# of free variables) is generally equivalent to n - rank.
* @time O(n * m * min(n, m))
* @memorv O(n * m)
*/
matrix_info solve_linear_mod(vector<vector<long long>>& mat, const vector<long long>& b, long
    \hookrightarrow long mod) {
   assert(ssize(mat) == ssize(b));
   int n = ssize(mat), m = ssize(mat[0]);
   for (int i = 0; i < n; i++)
       mat[i].push_back(b[i]);
   auto [rank, det] = row_reduce(mat, m, mod);//row reduce not including the last column
   //check if solution exists
```

```
for (int i = rank; i < n; i++)
    if (mat[i].back() != 0)
        return {rank, det, {} }; //no solution exists
    //initialize solution vector (`x`) from row-reduced matrix
    vector<long long> x(m, 0);
    for (int i = 0, j = 0; i < rank; i++) {
        while (mat[i][j] == 0) j++; //find pivot column
        x[j] = mat[i].back();
    }
    return {rank, det, x};
}</pre>
```

### Euler's Totient Phi Function

```
//cat totient.hpp | ./hash.sh
//36bd41
#pragma once
* Euler's totient function counts the positive integers up to a given integer
* n that are relatively prime to n. To improve, pre-calc prime factors or use
* Pollard-rho to find prime factors.
* @time O(sqrt n)
int totient(int n) {
    int res = n;
    for (int i = 2; i * i <= n; i++) {
        if (n % i == 0) {
            while (n \% i == 0) n /= i;
            res -= res / i:
    }
    if (n > 1) res -= res / n;
    return res;
```

### Tetration MOD

```
//cat tetration_mod.hpp | ./hash.sh
//e2153e
#pragma once
#include "binary_exponentiation_mod.hpp"
#include "totient.hpp"
* Returns base ^ (base ^ (base ^ ... )) % mod, where the height of the tower
* is pw.
* Let t = totient(mod).
* If log2(mod) \le pw then (base^pw)^mod == (base^(t+(pw^t)))^mod
* @see https://cp-algorithms.com/algebra/phi-function.html#generalization
* @time O(sqrt(mod) * log(mod))
long long tetration(long long base, long long pw, long long mod) {
   if (mod == 1)
       return 0;
    if (base == 0)
        return (pw + 1) % 2 % mod;
    if (base == 1 || pw == 0)
        return 1;
```

```
if (pw == 1)
    return base % mod;
if (base == 2 && pw == 2)
    return 4 % mod;
if (base == 2 && pw == 3)
    return 16 % mod;
if (base == 3 && pw == 2)
    return 27 % mod;
//need enough base cases such that the following is true
//log2(mod) <= tetration(base, pw - 1) (before modding)
int t = totient(int(mod));
long long exp = tetration(base, pw - 1, t);
return bin_exp(base, exp + t, mod);</pre>
```

#### MISC

### Cartesian Tree

```
//cat cartesian_tree.hpp | ./hash.sh
//204c45
#pragma once
#include "monotonic_stack.hpp"
* Min cartesian tree - root stores min.
* @time O(n)
* @memory O(n)
*/
vector<int> cartesian_tree(const vector<int>& arr) {
   int n = ssize(arr);
   auto rv /*reverse*/ = [&](int i) -> int {
       return n - 1 - i;
   };
   vector<int> left = monotonic_stack<int>(arr, greater());
   vector<int> right = monotonic_stack<int>(vector<int>(arr.rbegin(), arr.rend()), greater());
   vector<int> par(n);
   for (int i = 0; i < n; i++) {
       int le = left[i], ri = rv(right[rv(i)]);
       if (le >= 0 && ri < n) par[i] = arr[le] > arr[ri] ? le : ri;
       else if (le >= 0) par[i] = le;
       else if (ri < n) par[i] = ri;</pre>
       else par[i] = i; //true only for root
   }
   return par;
```

# Count Rectangles

```
//cat count_rectangles.hpp | ./hash.sh
//8990f7
#pragma once
#include "monotonic_stack.hpp"
/**
 * Given a n-by-m boolean matrix, calculate cnt[i][j]. cnt[i][j] = the number
 * of times an i-by-j sub rectangle appears in the matrix such that all i*j
 * cells in the sub rectangle are 1.
```

```
* cnt[0][j] and cnt[i][0] will contain garbage values.
* @time O(n * m)
* @memory O(n * m)
vector<vector<int>> count_rectangles(const vector<vector<bool>>& grid) {
    int n = ssize(grid), m = ssize(grid[0]);
    vector<vector<int>> cnt(n + 1, vector<math><int>(m + 1, 0));
    vector<int> arr(m, 0);
    auto rv /*reverse*/ = [&](int j) -> int {
        return m - 1 - j;
   };
    for (int i = 0; i < n; i++) {
        for (int j = 0; j < m; j++)
            arr[j] = grid[i][j] * (arr[j] + 1);
        vector<int> left = monotonic_stack<int>(arr, greater());
        vector<int> right = monotonic_stack<int>(vector<int>(arr.rbegin(), arr.rend()),

    greater_equal());
        for (int j = 0; j < m; j++) {
            int le = j - left[j] - 1, ri = rv(right[rv(j)]) - j - 1;
            cnt[arr[j]][le + ri + 1]++;
            cnt[arr[j]][le]--;
            cnt[arr[j]][ri]--;
       }
    for (int i = 1; i \le n; i++)
        for (int k = 0; k < 2; k++)
            for (int j = m - 1; j >= 1; j--)
               cnt[i][j] += cnt[i][j + 1];
    for (int j = 1; j \le m; j++)
        for (int i = n - 1; i >= 1; i--)
            cnt[i][j] += cnt[i + 1][j];
    return cnt;
```

## Max Rectangle in Histogram

```
//cat max_rect_histogram.hpp | ./hash.sh
//95288f
#pragma once
#include "monotonic_stack.hpp"
/**
* @time O(n)
* @memory O(n)
long long max_rect_histogram(const vector<int>& arr) {
    auto rv /*reverse*/ = [&](int i) -> int {
        return ssize(arr) - 1 - i;
    vector<int> left = monotonic_stack<int>(arr, greater_equal());
    vector<int> right = monotonic_stack<int>(vector<int>(arr.rbegin(), arr.rend()),
         \hookrightarrow greater_equal());
    long long max_area = 0;
    for (int i = 0; i < ssize(arr); i++) {</pre>
        int le = left[i], ri = rv(right[rv(i)]);//arr[i] is the max of range (le, ri)
        max_area = max(max_area, 1LL * arr[i] * (ri - le - 1));
    return max_area;
```

```
}
```

#### Monotonic Stack

```
//cat monotonic_stack.hpp | ./hash.sh
//35c95c
#pragma once
* Returns array 'le' where 'le[i]' = max number such that: 'le[i]' < i and
* !op(arr[le[i]], arr[i]). Returns -1 if no number exists.
* @code{.cpp}
      vector<int> le = monotonic_stack<int>(arr, less()); //(or replace `less` with:
     vector<int> le = monotonic_stack<int>(arr, [&](int x, int y) {return x < y;});</pre>
* @endcode
* @time O(n)
* @memorv O(n)
template <typename T> vector<int> monotonic_stack(const vector<T>& arr, const

    function<bool(const T&, const T&)>& op) {

   vector<int> le(ssize(arr));
   for (int i = 0; i < ssize(arr); i++) {</pre>
      le[i] = i - 1;
       while (le[i] >= 0 && op(arr[le[i]], arr[i])) le[i] = le[le[i]];
   }
   return le;
```

## GCD Convolution

```
//cat gcd_convolution.hpp | ./hash.sh
//d92c44
#pragma once
* Returns array 'c' where 'c[k]' = the sum for all pairs where gcd(i,j) == k
* of a[i] * b[j].
* @time O(n log n)
* @memory O(n)
template<int MOD> vector<int> gcd_convolution(const vector<int>& a, const vector<int>& b) {
   assert(ssize(a) == ssize(b));
   int n = ssize(a);
   vector<int> c(n);
   for (int gcd = n - 1; gcd >= 1; gcd--) {
       int sum_a = 0, sum_b = 0;
       for (int i = gcd; i < n; i += gcd) {
           sum_a = (sum_a + a[i]) \% MOD, sum_b = (sum_b + b[i]) \% MOD;
           c[gcd] = (c[gcd] - c[i] + MOD) % MOD;
       c[gcd] = int((c[gcd] + 1LL * sum_a * sum_b) % MOD);
   }
   return c;
```

### Iterate Chooses

```
//cat iterate_chooses.hpp | ./hash.sh
//c79083
#pragma once
int next_subset(int mask) {
    int c = mask & -mask, r = mask + c;
    return r | (((r ^ mask) >> 2) / c);
}
/**

* Iterates over all bitmasks of size n with k bits set.

* @see https://github.com/kth-competitive-programming/

* kactl/blob/main/content/various/chapter.tex

* @time O(n choose k)

* @memory O(1)

*/
void iterate_chooses(int n, int k, const function<void(int)>& func) {
    for (int mask = (1 << k) - 1; mask < (1 << n); mask = next_subset(mask))
        func(mask);
}</pre>
```

#### Iterate Submasks

```
//cat iterate_submasks.hpp | ./hash.sh
//084c05
#pragma once
/**
 * Iterates over all submasks of mask.
 * @time O(3^n) to iterate every submask of every mask of size n
 * @memory O(1)
 */
void iterate_submasks(int mask, const function<void(int)>& func) {
   for (int submask = mask; submask; submask = (submask - 1) & mask)
     func(submask);
}
```

# Iterate Supermasks

```
//cat iterate_supermasks.hpp | ./hash.sh
//76b38f
#pragma once
/**

* Iterates over all supermasks of mask.

* @time O(3^n) to iterate every supermask of every mask of size n

* @memory O(1)

*/

void iterate_supermasks(int mask, int n, const function<void(int)>& func) {
   for (int supermask = mask; supermask < (1 << n); supermask = (supermask + 1) | mask)
      func(supermask);
}</pre>
```

# Number of Distinct Subsequences DP

```
//cat num_distinct_subsequences.hpp | ./hash.sh
//d94bdc
#pragma once
/**
```

```
* Returns the number of distinct subsequences of `arr`. The empty subsequence
* @time O(n log n)
* @memory O(n)
int num_subsequences(const vector<int>& arr, int mod) {
   vector<int> dp(ssize(arr) + 1, 1);
   map<int, int> last;
   for (int i = 0; i < ssize(arr); i++) {</pre>
       int& curr = dp[i + 1] = 2 * dp[i];
       if (curr >= mod) curr -= mod;
       auto it = last.find(arr[i]);
       if (it != last.end()) {
           curr -= dp[it->second];
           if (curr < 0) curr += mod;</pre>
            it->second = i;
       } else last[arr[i]] = i;
   }
   return dp.back();
```

#### **PBDS**

```
//cat policy_based_data_structures.hpp | ./hash.sh
//a777d7
#pragma once
//place these includes *before* the `#define int long long` else compile error
//not using <bits/extc++.h> as it compile errors on codeforces c++20 compiler
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
//BST with extra functions https://codeforces.com/blog/entry/11080
//order_of_key - # of elements *strictly* less than given element
//find_by_order - find kth largest element, k is 0 based so find_by_order(0) returns min
template <typename T> using indexed_set = tree<T, null_type, less<T>, rb_tree_tag,

    tree_order_statistics_node_update>;
//example initialization:
indexed_set<pair<long long, int>> is;
//hash table (apparently faster than unordered_map): https://codeforces.com/blog/entry/60737
//example initialization:
gp_hash_table<string, long long> ht;
```

#### Random

```
//cat random.hpp | ./hash.sh
//ab9111
#pragma once

//source: https://codeforces.com/blog/entry/61675

mt19937 rng(chrono::steady_clock::now().time_since_epoch().count());

//intended types: int, unsigned, long long
//returns a random number in range [le, ri)
template <typename T> inline T get_rand(T le, T ri) {
    assert(le < ri);</pre>
```

```
return uniform_int_distribution<T>(le, ri - 1)(rng);
}
//vector<int> a;
//shuffle(a.begin(), a.end(), rng);
```

#### Safe Hash

```
//cat safe_hash.hpp | ./hash.sh
//d9ea53
#pragma once
//source: https://codeforces.com/blog/entry/62393
struct custom_hash {
    static uint64_t splitmix64(uint64_t x) {
        // http://xorshift.di.unimi.it/splitmix64.c
        x += 0x9e3779b97f4a7c15;
        x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
        x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
        return x ^ (x >> 31);
    size_t operator()(uint64_t x) const {
        static const uint64_t FIXED_RANDOM =

    chrono::steady_clock::now().time_since_epoch().count();

        return splitmix64(x + FIXED_RANDOM);
};
//usage:
unordered_map<long long, int, custom_hash> safe_map;
#include "policy_based_data_structures.hpp"
gp_hash_table<long long, int, custom_hash> safe_hash_table;
```

### RANGE DATA STRUCTURES

#### Number Distinct Elements

```
//cat distinct_query.hpp | ./hash.sh
//79534f
#pragma once
/**
* Can't handle updates; works with negatives.
 * @see https://cp-algorithms.com/data_structures/segment_tree.html
       #preserving-the-history-of-its-values-persistent-segment-tree
*/
struct distinct_query {
    struct node {
        int sum;
        int lch, rch;//children, indexes into `tree`
        node(int a_sum, int a_lch, int a_rch) : sum(a_sum), lch(a_lch), rch(a_rch) {}
    };
    const int N:
    vector<int> roots;
    deque<node> tree;
    * @time O(n log n)
     * @memory O(n log n)
```

```
distinct_query(const vector<int>& arr) : N(ssize(arr)), roots(N + 1, 0) {
        tree.emplace_back(0, 0, 0); //acts as null
        map<int, int> last_idx;
        for (int i = 0; i < N; i++) {
            roots[i + 1] = update(roots[i], 0, N, last_idx[arr[i]]);
            last_idx[arr[i]] = i + 1;
   }
    int update(int v, int tl, int tr, int idx) {
        if (tr - tl == 1) {
            tree.emplace_back(tree[v].sum + 1, 0, 0);
            return ssize(tree) - 1;
        int tm = tl + (tr - tl) / 2;
        int lch = tree[v].lch;
        int rch = tree[v].rch;
        if (idx < tm) lch = update(lch, tl, tm, idx);</pre>
        else rch = update(rch, tm, tr, idx);
        tree.emplace_back(tree[lch].sum + tree[rch].sum, lch, rch);
        return ssize(tree) - 1;
   }
    /**
    * Returns number of distinct elements in range [le, ri).
    * @time O(log n)
     */
    int query(int le, int ri) const {
        assert(0 <= le && le <= ri && ri <= N);
        return query(roots[le], roots[ri], 0, N, le + 1);
   }
   int query(int vl, int vr, int tl, int tr, int idx) const {
        if (tree[vr].sum == 0 || idx <= tl) return 0;</pre>
        if (tr <= idx) return tree[vr].sum - tree[vl].sum;</pre>
        int tm = tl + (tr - tl) / 2;
        return query(tree[v1].lch, tree[vr].lch, tl, tm, idx) +
               query(tree[v1].rch, tree[vr].rch, tm, tr, idx);
   }
};
```

# Implicit Lazy Segment Tree

```
//cat implicit_seg_tree.hpp | ./hash.sh
//ab5eb9
#pragma once
//example initialization:
// implicit_seg_tree<10'000'000> ist(le, ri);
template <int N> struct implicit_seg_tree {
   using dt = array<long long, 2>;//min, number of mins
   using ch = long long;
   static dt combine(const dt& le, const dt& ri) {
       if (le[0] == ri[0]) return {le[0], le[1] + ri[1]};
        return min(le, ri);
   }
   static constexpr dt UNIT{(long long)1e18, 0LL);
   struct node {
       dt val;
       int lch = -1, rch = -1; // children, indexes into 'tree', -1 for null
   } tree[N];
```

```
int ptr = 0, root_l, root_r;//[root_l, root_r) defines range of root node; handles
    implicit_seg_tree(int le, int ri) : root_l(le), root_r(ri) {
        tree[ptr++].val = {0, ri - le};
    void apply(int v, ch add) {
        tree[v].val[0] += add;
        tree[v].lazy += add;
    void push(int v, int tl, int tr) {
        if (tr - tl > 1 && tree[v].lch == -1) {
            int tm = tl + (tr - tl) / 2;
            assert(ptr + 1 < N);
            tree[v].lch = ptr;
            tree[ptr++].val = {0, tm - tl};
            tree[v].rch = ptr;
            tree[ptr++].val = {0, tr - tm};
        if (tree[v].lazy) {
            apply(tree[v].lch, tree[v].lazy);
            apply(tree[v].rch, tree[v].lazy);
            tree[v].lazy = 0;
        }
    }
    //update range [le,ri)
    void update(int le, int ri, ch add) {
        update(0, root_l, root_r, le, ri, add);
    void update(int v, int tl, int tr, int le, int ri, ch add) {
        if (ri <= tl || tr <= le)
            return;
        if (le <= tl && tr <= ri)
            return apply(v, add);
        push(v, tl, tr);
        int tm = tl + (tr - tl) / 2;
        update(tree[v].lch, tl, tm, le, ri, add);
        update(tree[v].rch, tm, tr, le, ri, add);
        tree[v].val = combine(tree[tree[v].lch].val,
                              tree[tree[v].rch].val);
    //query range [le,ri)
    dt query(int le, int ri) {
        return query(0, root_l, root_r, le, ri);
    dt query(int v, int tl, int tr, int le, int ri) {
        if (ri <= tl || tr <= le)
            return UNIT;
        if (le <= tl && tr <= ri)
            return tree[v].val;
        push(v, tl, tr);
        int tm = tl + (tr - tl) / 2;
        return combine(query(tree[v].lch, tl, tm, le, ri),
                       query(tree[v].rch, tm, tr, le, ri));
};
```

```
//cat kth_smallest.hpp | ./hash.sh
//b86dd0
#pragma once
* Can't handle updates.
 * @see https://cp-algorithms.com/data_structures/segment_tree.html
       preserving-the-history-of-its-values-persistent-segment-tree
*/
struct kth_smallest {
   struct node {
        int sum;
        int lch, rch;//children, indexes into `tree`
        node(int a_sum, int a_lch, int a_rch) : sum(a_sum), lch(a_lch), rch(a_rch) {}
   };
   int mn = INT_MAX, mx = INT_MIN;
    vector<int> roots;
    deque<node> tree;
    * @time O(n log max)
     * @memory O(n log max)
    kth_smallest(const vector<int>& arr) : roots(ssize(arr) + 1, 0) {
        tree.emplace_back(0, 0, 0); //acts as null
        for (int val : arr) mn = min(mn, val), mx = max(mx, val + 1);
        for (int i = 0; i < ssize(arr); i++)
            roots[i + 1] = update(roots[i], mn, mx, arr[i]);
   }
    int update(int v, int tl, int tr, int idx) {
        if (tr - tl == 1) {
            tree.emplace_back(tree[v].sum + 1, 0, 0);
            return ssize(tree) - 1;
        int tm = tl + (tr - tl) / 2;
        int lch = tree[v].lch;
        int rch = tree[v].rch;
        if (idx < tm) lch = update(lch, tl, tm, idx);</pre>
        else rch = update(rch, tm, tr, idx);
        tree.emplace_back(tree[lch].sum + tree[rch].sum, lch, rch);
        return ssize(tree) - 1;
   }
    /**
    * Returns (k+1)th smallest number in range [le, ri). k is 0-based, so
     * query(le,ri,0) returns the min.
     * @time O(log max)
    int query(int le, int ri, int k) const {
        assert(0 <= k && k < ri - le); //note this condition implies le < ri</pre>
        assert(0 <= le && ri < ssize(roots));</pre>
        return query(roots[le], roots[ri], mn, mx, k);
   int query(int vl, int vr, int tl, int tr, int k) const {
        assert(tree[vr].sum > tree[vl].sum);
        if (tr - tl == 1)
            return tl;
        int tm = tl + (tr - tl) / 2;
        int left_count = tree[tree[vr].lch].sum - tree[tree[vl].lch].sum;
        if (left_count > k) return query(tree[v1].lch, tree[vr].lch, tl, tm, k);
        return query(tree[v1].rch, tree[vr].rch, tm, tr, k - left_count);
```

```
]};
```

### Merge Sort Tree

```
//cat merge_sort_tree.hpp | ./hash.sh
//e110ed
#pragma once
//For point updates: either switch to policy based BST, or use sgrt decomposition
struct merge_sort_tree {
    const int N, S/*smallest power of 2 >= N*/;
    vector<vector<int>> tree;
    merge_sort_tree(const vector<int>& arr) : N(ssize(arr)), S(N ? 1 << __lg(2 * N - 1) : 0),
         \hookrightarrow tree(2 * N) {
        for (int i = 0; i < N; i++)
            tree[i + N] = {arr[i]};
        rotate(tree.rbegin(), tree.rbegin() + S - N, tree.rbegin() + N);
        for (int i = N - 1; i >= 1; i--) {
            const auto& le = tree[2 * i];
            const auto& ri = tree[2 * i + 1];
            tree[i].reserve(ssize(le) + ssize(ri));
            merge(le.begin(), le.end(), ri.begin(), ri.end(), back_inserter(tree[i]));
        }
    }
    int value(int v, int x) const {
        return int(lower_bound(tree[v].begin(), tree[v].end(), x) - tree[v].begin());
    int to_leaf(int i) const {
        i += S;
        return i < 2 * N ? i : 2 * (i - N);
     * Returns the number of values in range [le, ri) which are < x.
     * @time O(log^2(n))
    int query(int le, int ri, int x) const {
        int res = 0;
        for (le = to_leaf(le), ri = to_leaf(ri); le < ri; le >>= 1, ri >>= 1) {
            if (le & 1) res += value(le++, x);
            if (ri & 1) res += value(--ri, x);
        return res;
    }
};
```

### BIT

```
//cat bit.hpp | ./hash.sh
//ab7995
#pragma once
//NOLINTNEXTLINE(readability-identifier-naming)
template <typename T> struct BIT {
    vector<T> bit;
    BIT(int n) : bit(n, 0) {}
    BIT(const vector<T>& a) : bit(a) {
        for (int i = 0; i < ssize(a); i++) {
            int j = i | (i + 1);
            if (j < ssize(a)) bit[j] += bit[i];</pre>
```

```
}
    }
    void update(int i, const T& d) {
        assert(0 <= i && i < ssize(bit));</pre>
        for (; i < ssize(bit); i |= i + 1) bit[i] += d;
    }
    /**
     * sum of range [0, ri)
     * @time O(log n)
    T sum(int ri) const {
        assert(0 <= ri && ri <= ssize(bit));</pre>
        T ret = 0;
        for (; ri > 0; ri &= ri - 1) ret += bit[ri - 1];
        return ret;
    }
    /**
     * sum of range [le, ri)
     * @time O(log n)
     */
    T sum(int le, int ri) const {
        assert(0 <= le && le <= ri && ri <= ssize(bit));
        return sum(ri) - sum(le);
    }
    /**
     * Returns min pos such that sum of range [0, pos) >= sum. Returns
     * ssize(bit) + 1 if no sum is >= sum.
     * Doesn't work if BIT::sum(i, i + 1) < 0
     * @time O(log n)
    int lower_bound(T sum) const {
        if (sum <= 0) return 0;</pre>
        int pos = 0;
        for (int pw = 1 << __lg(ssize(bit) | 1); pw; pw >>= 1)
            if (pos + pw <= ssize(bit) && bit[pos + pw - 1] < sum)</pre>
                pos += pw, sum -= bit[pos - 1];
        return pos + 1;
    }
};
```

# RMQ

```
//cat rmq.hpp | ./hash.sh
//59ae14
#pragma once
/**
    * @code{.cpp}
    * vector<long long> arr;
    * RMQ<long long> rmq(arr, [&](auto x, auto y) { return min(x, y); });
    *
    * //To get index of min element:
    * vector<pair<long long, int>> arr; //initialize arr[i].second = i
    * RMQ<pair<long long, int>> rmq(arr, [&](auto x, auto y) { return min(x, y); });
    * @endcode
    * @see https://github.com/kth-competitive-programming/
    * kactl/blob/main/content/data-structures/RMQ.h
    */
//NOLINTNEXTLINE(readability-identifier-naming)
```

```
template <typename T> struct RMQ {
    vector<vector<T>> dp;
    function<T(const T&, const T&)> op;
     * @time O(n log n)
     * @memory O(n log n)
    RMQ(const vector<T>& arr, const function<T(const T&, const T&)>& a_op) : dp(1, arr),
         \hookrightarrow op(a_op) {
        for (int i = 0; (2 << i) <= ssize(arr); i++) {
            dp.emplace_back(ssize(arr) - (2 << i) + 1);
            for (int j = 0; j < ssize(dp.back()); j++)
                dp[i + 1][j] = op(dp[i][j], dp[i][j + (1 << i)]);
        }
    }
    * range [le, ri)
     * @time O(1)
    T query(int le, int ri) const {
        assert(0 <= le && le < ri && ri <= ssize(dp[0]));
        int \lg = _{-}\lg(ri - \lg);
        return op(dp[lg][le], dp[lg][ri - (1 << lg)]);</pre>
    }
};
```

### Disjoint RMQ

```
//cat disjoint_rmq.hpp | ./hash.sh
//848e2b
#pragma once
 * Disjoint RMQ is like normal RMQ except these ranges never overlap. It is
 * useful for:
 * - min and # of mins.
 * - product under composite mod
 * - 2-by-2 matrix multiply
 * @code{.cpp}
       //usage for min and # of mins:
       vector<pair<long long, int>> arr; //initialize arr[i].second = 1
       disjoint_rmq<pair<long long, int>> rmq(arr, {LLONG_MAX, 0}, [&](auto x, auto y) {
           if (x.first == y.first) return make_pair(x.first, x.second + y.second);
           return min(x, y);
       });
 * @endcode
 * @see https://codeforces.com/blog/entry/87940,
       https://github.com/sgtlaugh/algovault/blob/
       master/code_library/disjoint_sparse_table.cpp
template <typename T> struct disjoint_rmg {
    const int N:
    vector<vector<T>> dp:
    function<T(const T&, const T&)> op; // any associative operation
     * @time O(n log n)
     * @memory O(n log n)
    disjoint_rmq(const vector<T>& arr, const T& identity, const function<T(const T&, const
```

```
\hookrightarrow T&)>& a_op) : N(ssize(arr)), op(a_op) {
        for (int i = 0, len = 1; len <= N; i++, len *= 2) {
            dp.emplace_back(N + 1, identity);
            for (int center = len; center < N + len; center += 2 * len) {</pre>
                for (int j = center + 1; j <= min(N, center + len); j++)
                     dp[i][j] = op(dp[i][j - 1], arr[j - 1]);
                for (int j = min(N, center) - 1; j >= center - len; j--)
                     dp[i][j] = op(arr[j], dp[i][j + 1]);
            }
        }
    }
    /**
     * range [le, ri)
     * @time O(1)
     */
    T query(int le, int ri) const {
        assert(0 <= le && le < ri && ri <= N);
        int lg = __lg(le ^ ri);
        return op(dp[lg][le], dp[lg][ri]);
   }
};
```

## Lazy Segment Tree

```
//cat lazy_segment_tree.hpp | ./hash.sh
//96535f
#pragma once
* Internal nodes are [1, n), leaf nodes are [n, 2 * n).
* Rotating leaves makes it a single complete binary tree (instead of a set of
* perfect binary trees). So now, even for non-power of 2 size:
* - recursive seg tree works
* - recursive tree walks AKA binary search works
* - root is at tree[1]
* @see https://codeforces.com/blog/entry/18051
      https://github.com/ecnerwala/cp-book/blob/master/src/seg_tree.hpp
      https://github.com/yosupo06/Algorithm/blob/master/src/datastructure/segtree.hpp
*/
struct seg_tree {
   using dt = long long;
   using ch = long long;
   static dt combine(const dt& le, const dt& ri) {
        return min(le, ri);
   }
   static const dt UNIT = 1e18;
   struct node {
       dt val;
       ch lazy;
       int le, ri;//[le, ri)
   };
   const int N, S/*smallest power of 2 >= N*/;
   vector<node> tree:
   seg\_tree(const\ vector< dt>\&\ arr) : N(ssize(arr)), S(N ? 1 << __lg(2 * N - 1) : 0), tree(2 * N - 1)
        \hookrightarrow N) {
       for (int i = 0; i < N; i++)
            tree[i + N] = {arr[i], 0, i, i + 1};
       rotate(tree.rbegin(), tree.rbegin() + S - N, tree.rbegin() + N);
        for (int i = N - 1; i >= 1; i--) {
```

```
tree[i] = {
            combine(tree[2 * i].val, tree[2 * i + 1].val),
            tree[2 * i].le,
            tree[2 * i + 1].ri
        };
    }
void apply(int v, ch change) {
    tree[v].val += change;
    tree[v].lazy += change;
}
void push(int v) {
    if (tree[v].lazy) {
        apply(2 * v, tree[v].lazy);
        apply(2 * v + 1, tree[v].lazy);
        tree[v].lazy = 0;
    }
void build(int v) {
    tree[v].val = combine(tree[2 * v].val, tree[2 * v + 1].val);
int to_leaf(int i) const {
    return i < 2 * N ? i : 2 * (i - N);
}
//update range [le, ri)
void update(int le, int ri, ch change) {
    assert(0 <= le && le <= ri && ri <= N);
    le = to_leaf(le), ri = to_leaf(ri);
    int lca_l_r = __lg((le - 1) ^ ri);
    for (int lg = __lg(le); lg > __builtin_ctz(le); lg--) push(le >> lg);
    for (int lg = lca_l_r; lg > __builtin_ctz(ri); lg--) push(ri >> lg);
    for (int x = le, y = ri; x < y; x >>= 1, y >>= 1) {
        if (x & 1) apply(x++, change);
        if (y & 1) apply(--y, change);
    for (int lg = __builtin_ctz(ri) + 1; lg <= lca_l_r; lg++) build(ri >> lg);
    for (int lg = \_builtin\_ctz(le) + 1; lg <= \__lg(le); lg++) build(le >> lg);
void update(int v/* = 1*/, int le, int ri, ch change) {
    if (ri <= tree[v].le || tree[v].ri <= le) return;</pre>
    if (le <= tree[v].le && tree[v].ri <= ri) return apply(v, change);</pre>
    push(v);
    update(2 * v, le, ri, change);
    update(2 * v + 1, le, ri, change);
    build(v);
}
//query range [le, ri)
dt query(int le, int ri) {
    assert(0 <= le && le <= ri && ri <= N);
    le = to_leaf(le), ri = to_leaf(ri);
    int lca_l_r = __lg((le - 1) ^ ri);
    for (int lg = __lg(le); lg > __builtin_ctz(le); lg--) push(le >> lg);
    for (int lg = lca_l_r; lg > __builtin_ctz(ri); lg--) push(ri >> lg);
    dt resl = UNIT, resr = UNIT;
    for (; le < ri; le >>= 1, ri >>= 1) {
        if (le & 1) resl = combine(resl, tree[le++].val);
        if (ri & 1) resr = combine(tree[--ri].val, resr);
```

```
    return combine(resl, resr);
}
dt query(int v/* = 1*/, int le, int ri) {
    if (ri <= tree[v].le || tree[v].ri <= le) return UNIT;
    if (le <= tree[v].le && tree[v].ri <= ri) return tree[v].val;
    push(v);
    return combine(query(2 * v, le, ri), query(2 * v + 1, le, ri));
}
};
</pre>
```

### **STRINGS**

## Binary Trie

```
//cat binary_trie.hpp | ./hash.sh
//88fa9c
#pragma once
struct binary_trie {
   const int MX_BIT = 62;
   struct node {
       long long val = -1;
       int sub_sz = 0;//number of inserted values in subtree
       array < int, 2 > next = \{-1, -1\};
   };
   vector<node> t;
   binary_trie() : t(1) {}
   /**
    * Pass delta = 1 to insert val, -1 to remove val, 0 to get the # of val's
    * in this data structure.
    * @time O(MX_BIT)
   int update(long long val, int delta) {
       int c = 0;
        t[0].sub_sz += delta;
        for (int bit = MX_BIT; bit >= 0; bit--) {
           bool v = (val >> bit) & 1;
           if (t[c].next[v] == -1) {
                t[c].next[v] = ssize(t);
                t.emplace_back();
           }
           c = t[c].next[v];
            t[c].sub_sz += delta;
       t[c].val = val;
       return t[c].sub_sz;
   int size() const {
       return t[0].sub_sz;
   }
    * Returns x such that x is in this data structure, and the value of
    * (x^val) is minimum.
    * @time O(MX_BIT)
   long long min_xor(long long val) const {
       assert(size() > 0);
```

```
int c = 0;
for (int bit = MX_BIT; bit >= 0; bit--) {
    bool v = (val >> bit) & 1;
    int ch = t[c].next[v];
    if (ch != -1 && t[ch].sub_sz > 0)
        c = ch;
    else
        c = t[c].next[!v];
    }
    return t[c].val;
}
```

#### **Prefix Function**

```
//cat prefix_function.hpp | ./hash.sh
//65fea7
#pragma once
/**
    * @see https://cp-algorithms.com/string/prefix-function.html#implementation
    * @time O(n)
    */

template <typename T> vector<int> prefix_function(const T& s) {
    vector<int> pi(ssize(s), 0);
    for (int i = 1; i < ssize(s); i++) {
        int j = pi[i - 1];
        while (j > 0 && s[i] != s[j]) j = pi[j - 1];
        pi[i] = j + (s[i] == s[j]);
    }
    return pi;
}
```

### KMP String Matching

```
//cat kmp.hpp | ./hash.sh
//1edfb2
#pragma once
#include "prefix_function.hpp"
* @code{.cpp}
      string s;
      KMP kmp(s);
       vector<int> a;
       KMP kmp(a);
 * @endcode
 * KMP doubling trick: to check if 2 arrays are rotationally equivalent: run
 * kmp with one array as the needle and the other array doubled (excluding the
* first & last characters) as the haystack or just use kactl's min rotation
//NOLINTNEXTLINE(readability-identifier-naming)
template <typename T> struct KMP {
    T needle;
    vector<int> pi;
     * @time O(|needle|)
     * @memory O(|needle|)
```

```
KMP(const T& a_needle) : needle(a_needle), pi(prefix_function(needle)) {}
     * Returns array `matches` where:
     * haystack.substr(matches[i], ssize(needle)) == needle
     * @time O(|needle| + |haystack|)
    vector<int> find(const T& haystack) const {
        vector<int> matches;
        for (int i = 0, j = 0; i < ssize(haystack); i++) {
            while (j > 0 \& needle[j] != haystack[i]) j = pi[j - 1];
            if (needle[j] == haystack[i]) j++;
            if (j == ssize(needle)) {
                matches.push_back(i - ssize(needle) + 1);
                j = pi[j - 1];
            }
        }
        return matches;
   }
};
```

# Suffix and LCP Arrays

```
//cat suffix_array.hpp | ./hash.sh
//0272f8
#pragma once
* suffixes of "banana":
* 0 banana
* 1 anana
* 2 nana
* 3 ana
* 4 na
* 5 a
* sorted, lcp
* 5 a
*
* 3 ana
* |||
           3
* 1 anana
           0
* 0 banana
* 4 na
* ||
           2
* 2 nana
* suffix array = [5, 3, 1, 0, 4, 2]
* rank array satisfies sa[rank[i]] == i and rank[sa[i]] == i
* lcp array = [1, 3, 0, 0, 2, 0] - last 0 is a dummy value
* @code{.cpp}
      string s;
      suffix_array info(s, 128);
      vector<int> arr;
      suffix_array info(arr, 1e5 + 1);
* @endcode
```

```
* @see https://github.com/kth-competitive-programming/kactl
 * /blob/main/content/strings/SuffixArray.h
*/
template <typename T> struct suffix_array {
    const int N:
    vector<int> sa, rank, lcp;
     * Assumes 0 <= s[i] < max_val.
     * @time O((nlogn) + max_val)
     * @memory O(n + max_val)
    suffix_array(const T& s, int max_val) : N(ssize(s)), sa(N), rank(s.begin(), s.end()),
         \hookrightarrow lcp(N)  {
        iota(sa.begin(), sa.end(), 0);
        vector<int> tmp(N + 1, -1);
        for (int len = 0; len < N; len = max(1, 2 * len)) {//suffix array
            iota(tmp.begin(), tmp.begin() + len, N - len);
            for (int i = 0, j = len; i < N; i++)
                if (sa[i] >= len)
                    tmp[j++] = sa[i] - len;
            vector<int> freq(max_val, 0);
            for (int val : rank) freq[val]++;
            partial_sum(freq.begin(), freq.end(), freq.begin());
            for (int i = N - 1; i \ge 0; i--)
                sa[--freq[rank[tmp[i]]] = tmp[i];
            copy(rank.begin(), rank.end(), tmp.begin());
            max_val = 1, rank[sa[0]] = 0;
            for (int i = 1; i < N; i++) {
                int a = sa[i - 1], b = sa[i];
                rank[b] = (tmp[a] == tmp[b] && tmp[a + len] == tmp[b + len]) ? max_val - 1 :
                     \hookrightarrow max_val++;
            }
            if (max_val == N) break;
        for (int i = 0, k = 0; i < N; i++) {//lcp array
            if (k > 0) k--;
            if (rank[i] == 0) continue;
            int j = sa[rank[i] - 1];
            while (\max(i, j) + k < N \&\& s[j + k] == s[i + k]) k++;
            lcp[rank[i] - 1] = k;
    }
};
```

# Suffix Array Related Queries

```
//cat suffix_array_query.hpp | ./hash.sh
//ae4931
#pragma once
#include "suffix_array.hpp"
#include "../range_data_structures/rmq.hpp"
/**
    * Various queries you can do based on suffix array.
    * @see https://github.com/yosupo06/Algorithm/blob
    * /master/src/string/suffixarray.hpp
    */
template <typename T> struct sa_query {
        T s;
```

```
suffix_array<T> info;
RMQ<int> rmq_lcp, rmq_sa;
/**
 * Assumes 0 \le s[i] \le \max_{val}.
 * @time O((nlogn) + max_val)
 * @memory O((nlogn) + max_val) An O(max_val) size freq array is used
 * temporarily during suffix array construction. Only O(n log n) memory is
 * permanently stored by this struct.
sa_query(const T& a_s, int max_val) :
    info(suffix_array(s, max_val)),
    rmq_lcp(info.lcp, [](int i, int j) -> int {return min(i, j);}),
    rmq_sa(info.sa, [](int i, int j) -> int {return min(i, j);}) {}
 * Returns length of longest common prefix of suffixes s[idx1...N),
 * s[idx2...N), 0-based indexing.
 * To check if two substrings s[l1..r1), s[l2..r2) are equal:
 * r1-l1 == r2-l2 \&\& get_lcp(l1, l2) >= r1-l1
 * @time O(1)
int get_lcp(int idx1, int idx2) const {
    if (idx1 == idx2) return ssize(s) - idx1;
    auto [le, ri] = minmax(info.rank[idx1], info.rank[idx2]);
    return rmq_lcp.query(le, ri);
}
/**
 * Returns 1 if suffix s[idx1 ... N) < s[idx2 ... N).
 * Returns 0 if idx1 == idx2.
 * @time O(1)
*/
bool less(int idx1, int idx2) const {
    return info.rank[idx1] < info.rank[idx2];</pre>
}
/**
 * Returns range [le, ri) such that:
 * - for all i in [le, ri): t == s.substr(info.sa[i], ssize(t))
 * - `ri - le` is the # of matches of t in s.
 * @time O(|t| * log(|s|))
pair<int, int> find(const T& t) const {
    auto cmp = [&](int i, int cmp_val) -> bool {
        return s.compare(i, ssize(t), t) < cmp_val;</pre>
    };
    auto le = lower_bound(info.sa.begin(), info.sa.end(), 0, cmp);
    auto ri = lower_bound(le, info.sa.end(), 1, cmp);
    return {le - info.sa.begin(), ri - info.sa.begin()};
}
/**
 * Returns min i such that t == s.substr(i, ssize(t)) or -1. For example,
 * replace RMQ with kth-smallest PST/Wavelet to solve
 * https://open.kattis.com/problems/anothersubstringqueryproblem
 * 0(|t| * \log(|s|))
int find_first(const T& t) const {
    auto [le, ri] = find(t);
    if (le == ri) return -1;
    return rmq_sa.query(le, ri);
```

```
}
};
```

### Palindrome Query

```
//cat palindrome_query.hpp | ./hash.sh
//68c8e1
#pragma once
#include "../../kactl/content/strings/Manacher.h"
struct pal_query {
    const int N;
    array<vi, 2> pal_len;
     * @time O(n)
     * @memory O(n)
    pal_query(const string& s) : N(ssize(s)), pal_len(manacher(s)) {}
     * Returns 1 if substring s[le...ri) is a palindrome.
     * Returns 1 when le == ri.
     * @time O(1)
    bool is_pal(int le, int ri) const {
        assert(0 <= le && le <= ri && ri <= N);
        int len = ri - le;
        return pal_len[len & 1][le + len / 2] >= len / 2;
    }
};
```

#### Trie

```
//cat trie.hpp | ./hash.sh
//2aa8c6
#pragma once
//source: https://cp-algorithms.com/string/aho_corasick.html#construction-of-the-trie
const int K = 26;//alphabet size
struct trie {
    const char MIN_CH = 'A';//'a' for lowercase, '0' for digits
    struct node {
        int next[K], cnt_words = 0, par = -1;
        char ch;
        node(int a_par = -1, char a_ch = '#') : par(a_par), ch(a_ch) {
            fill(next, next + K, -1);
    };
    vector<node> t;
    trie() : t(1) {}
    void insert(const string& s) {
        int v = 0;
        for (char ch : s) {
            int let = ch - MIN_CH;
            if (t[v].next[let] == -1) {
                t[v].next[let] = ssize(t);
                t.emplace_back(v, ch);
            v = t[v].next[let];
```

```
t[v].cnt_words++;
}
int find(const string& s) const {
   int v = 0;
   for (char ch : s) {
      int let = ch - MIN_CH;
      if (t[v].next[let] == -1) return 0;
      v = t[v].next[let];
   }
   return t[v].cnt_words;
}
};
```