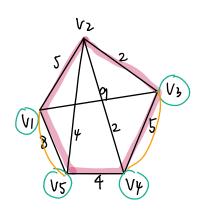
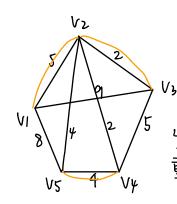
homework b. 16(a) 17(a).

16. 找出图中所有度为奇数的点: V1, V3, V4, V5.



添加边(V//VE), (V3, V4), 使各结点度数均为偶数。

图示圈不满足重边长度之和不超过回路的一半,取该回路共轭重边:

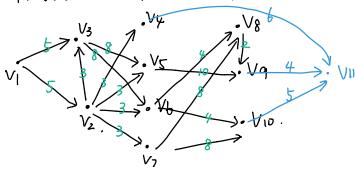


发现该图任一子回路中重边长度之和不超过回路的一半,

∴ 图中-条欧拉回路如(V1, V2, V3, V4, V5, V2, V4, V5, V1, V3, V2, V1)

17.(a) 由表可得, 若将工序2, 工序3交换编号, 则所有工序的前序工序编号均小于该工序编号。

二设工序i对应节点为vi,建立PT图,该图结点排序为拓扑序列。



图中V4,Va, V10均为结束结点,因此增设虚拟超结点 V11并建立W(va, V11)=4  $W(V_{10},V_{11})=5$ ,  $W(V_4,V_{11})=6$ 

转化为求N到VII的关键路径

设π(νi)为νι总开始时间. 初始全π(νi)=0 (ì=1,2... 11)

依次更新π(νj), j=2,3...11

 $\pi(y_j) = \max_{y_i \neq y_j \neq y_i} (\pi(y_i) + w(y_i, y_j)).$ 

 $RIT(v_2)=5$ 

 $\pi (v_3) = \pi(v_2) + 3 = 8$ 

 $\pi(\gamma_4) = \pi(V_2) + 3 = 8$ 

 $\pi (v_5) = \pi (v_3) + 8 = 16$ 

 $\pi (v_b) = \pi (v_3) + \delta = 16$ .

 $\pi(v_2) = \pi(v_2) + 3 = 8$ 

 $\pi (v_8) = \pi (v_b) + 4 = 20$ 

 $\pi(\gamma_q) = \pi(\gamma_c) + 10 = 26.$ 

 $\pi(v_0) = \pi(v_0) + 4 = 20$ 

 $\pi (V_{11}) = \pi (V_{2}) + 4 = 30$ .

.. 关键路径为(Vι, V<sub>2</sub>, V<sub>3</sub>, V5, Vq, V<sub>II</sub>)

下求各工序开始m最晚时间. 初始令 $T(v_i)=\pi(v_n)=30$ , i=1,2... ||

体次更新T(Vj), j=10,9,... 1

 $T(V_j) = \min (T(V_i) - W(V_j, V_i))$ 

Vi为Vjù 直接后继。

 $T(V_{10}) = T(V_{11}) - 5 = 25$ 

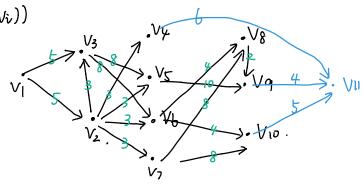
 $T(V_9) = T(V_{11}) - 4 = 2b$ .

 $T(V_8) = T(V_9) - 2 = 24$ 

 $T(V_7) = T(V_{10}) - 8 = 1$ 

 $T(V_6) = T(V_8) - 4 = 20$ .

 $T(V_5) = T(V_9) - 10 = 16.$ 



 $T(V_4) = T(V_{11}) - 6 = 24$   $T(V_3) = T(V_5) - 8 = 8$   $T(V_2) = T(V_3) - 3 = 5$  $T(V_1) = T(V_2) - 5 = 0$ 

.. 工序3(即重编号后加工序2)允许延误时间  $t(V_3')=T(V_2)-\pi(V_3)=0$  工序5 =  $t(V_5')=T(V_5)-\pi(V_5)=0$  工序10 =  $t(V_6')=25-20=5$  .