

Next generation of Tencent OLAP Engine

Tencent TEG LongYue





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01 Background

O2 Storage: Various Columns and Indexes

O3 Computation: Integrated with Presto

O4 Benchmark and Applications



Background

Data :

- 1. Thousand Columns, 10 Billions Rows
- 2. Index on arbitrary column(s)
- 3. Real-Time Write
- 4. Both Row-oriented and Column-oriented
- 5. Different Indexes

MercsDB

Performance :

- 1. Second response on 10 Billions rows query
- 2. MPP: both ad-hoc query and real-time query
- 3. Real-time Write: 100 Billion rows / day

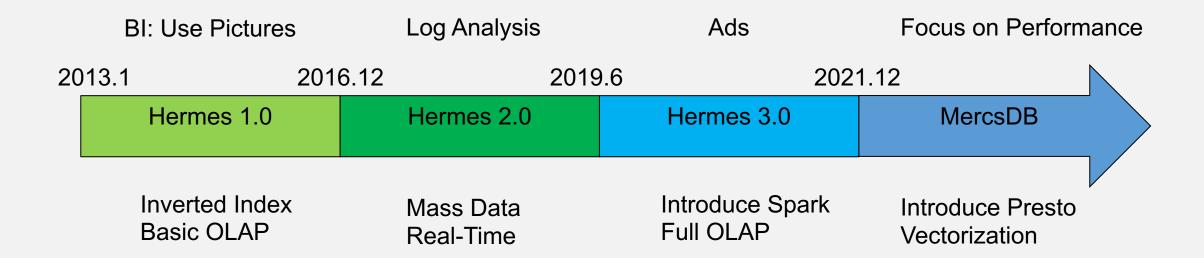
ElasticSearch

= Presto/Impala

ClickHouse



History





Current Status

Clusters

Query

Storage

Peak IO

5k+ Nodes

10M / Day

Total: 100 PB

Daily: 1PB

100M rows / s















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01 Background

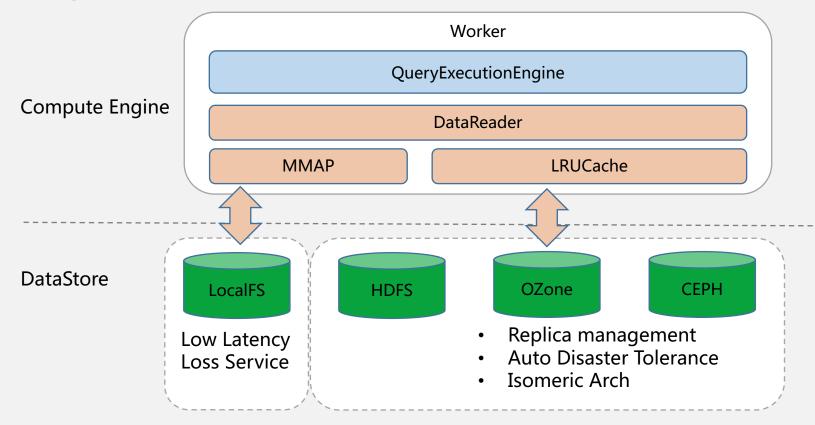
O2 Storage: Various Columns and Indexes

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O4 Benchmark and Applications



Basic Arch



- Districted vs Local
- HA vs Low Latency
- Hot Data vs Cold Data



Column-Oriented & Index

Q1: High QPS?

Q2: Second response for 10Bilion rows?

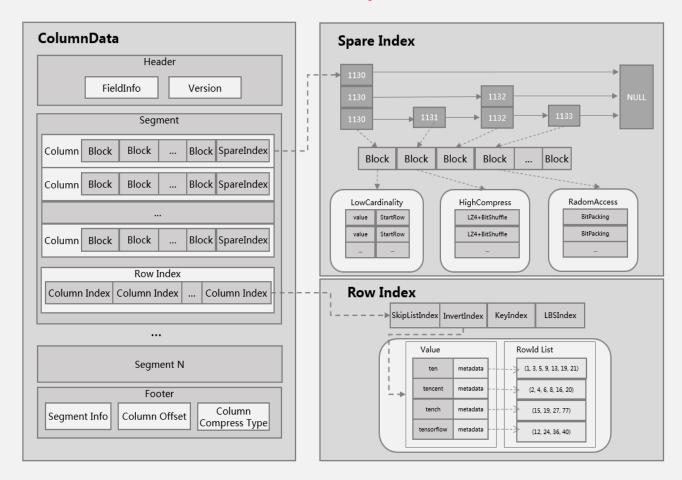
Q3: Cost-friendly for mass data?

Columns:

- Retrieval —— Low latency
- Sorted Mass Data
- 3. Compressed —— Cost-Friendly
- 4. Nested —— Support parquet

Indexes:

- 1. SpareIndex
- 2. SkipListIndex
- 3. InvertedIndex
- 4. KeyIndex
- 5. LBSIndex

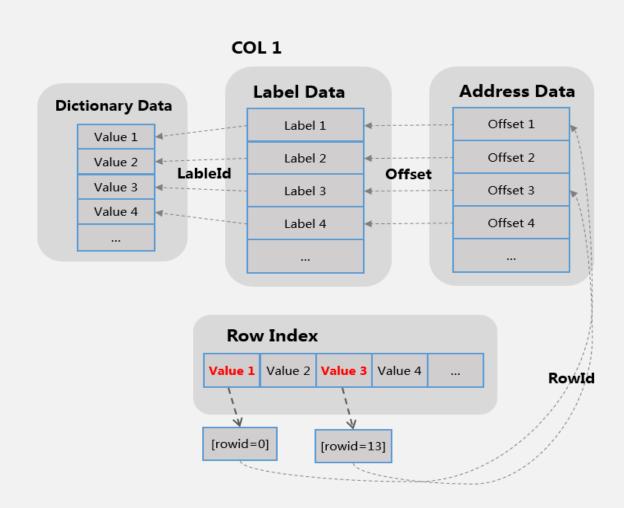




Retrieval Column

- Application
- 1. Low latency
- 2. Medium data
- 3. Simple Query
- Implementation
- 1. Storage -> Time
- 2. Dictionary Index

Index Size/Origin Data = 40%

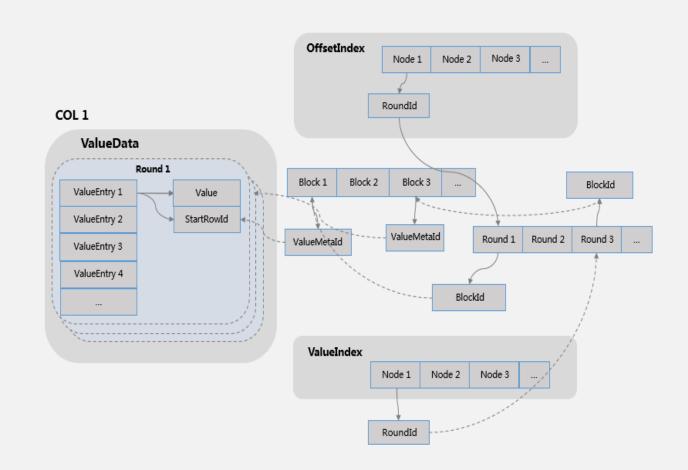




Sorted Column

- Application
- 1. Mass Data (much bigger than memory)
- 2. No other accelerations
- Implementation
- Sorted
- 2. Index both on data and offset

Improvement: 10x speed up

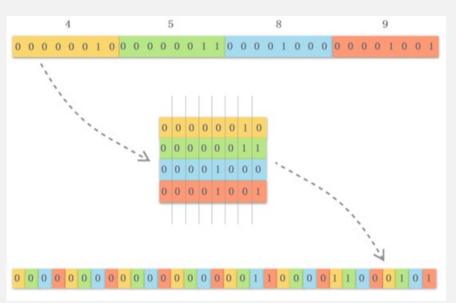


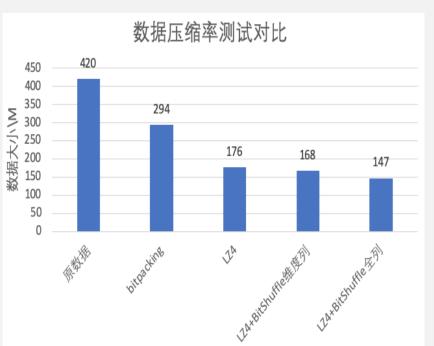


Compressed Column

- Application
- 1. High Cardinal
- 2. Direct compression: low performance
- Implementation
- 1. BitShuffle + LZ4

Compression rate: 50% up

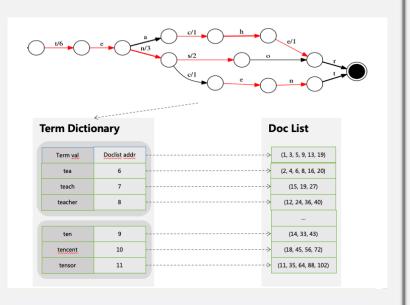




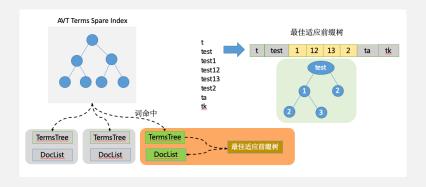


Indexs

Inverted Index

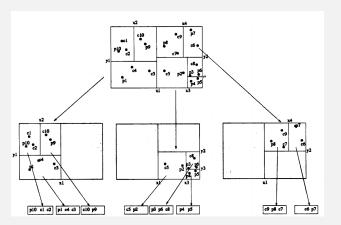


FST KeyIndex



	完全基于 FST 的索引的 key 值索引	基于平衡二叉树的 key 值索引	Lucene 倒排索引
十万次数据点查耗时	490 ms	243 ms	736 ms
相对于 Lucene 倒排索引的存储效果	1.3 倍	1.013 倍	基准 1
千次随机范围查询耗时	142706 ms	13737 ms	24163 ms

LBS(KDB) Index





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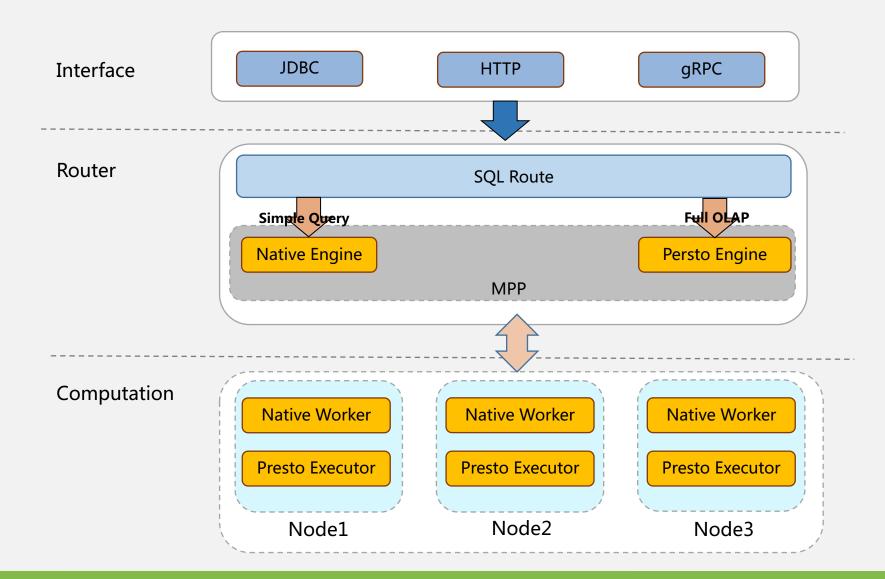
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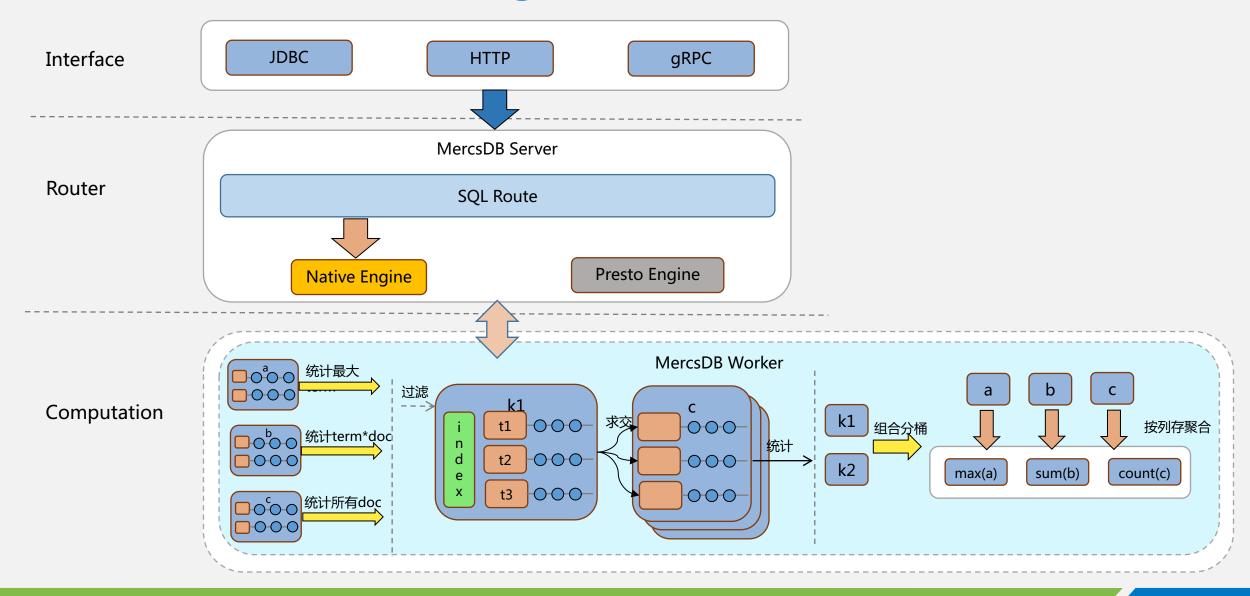


MercsDB — Dynamic Router

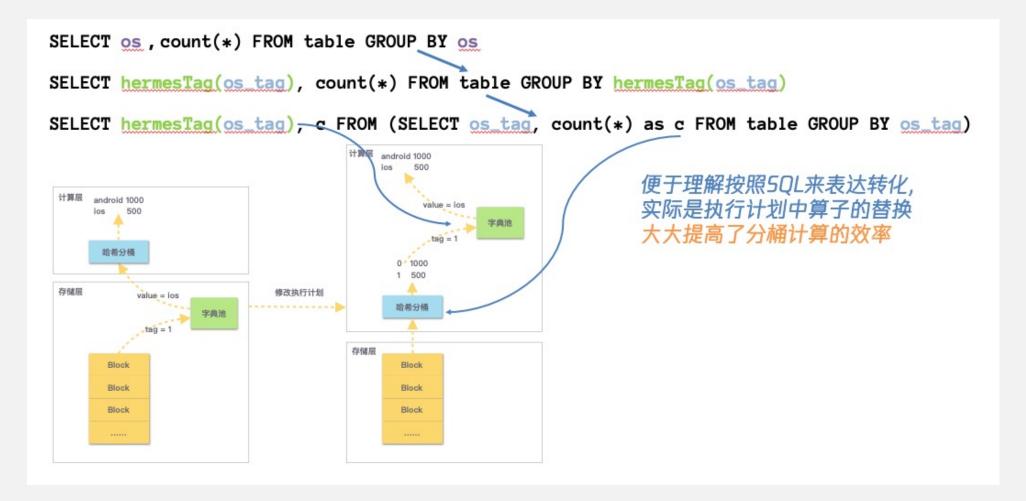




MercsDB Native Engine

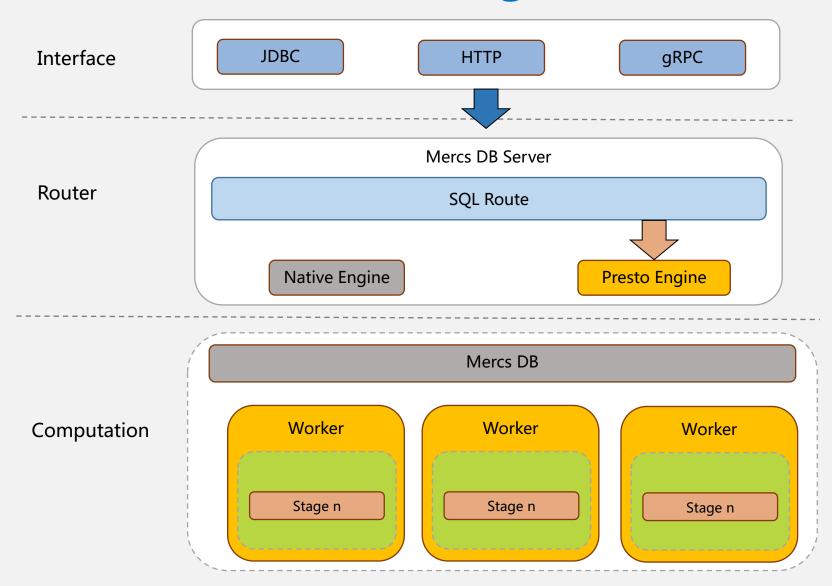


Late Materialization





MercsDB Presto Engine





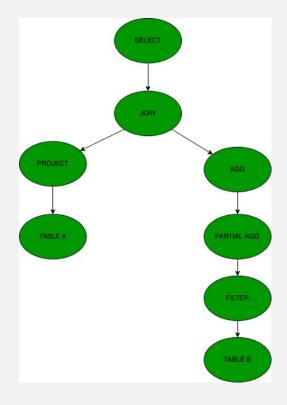
Push down into MercsDB

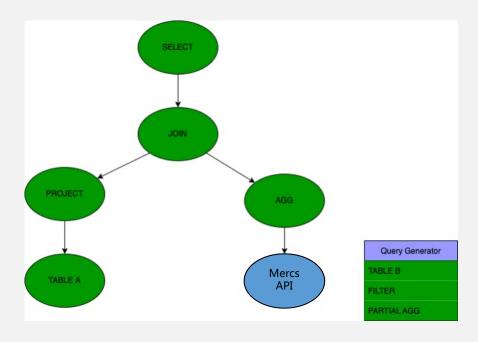
```
SELECT * FROM (SELECT FROM A)

JOIN

(SELECT agg FROM B WHERE X)

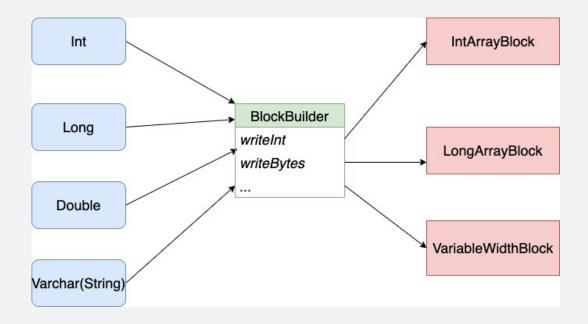
ON C = D
```



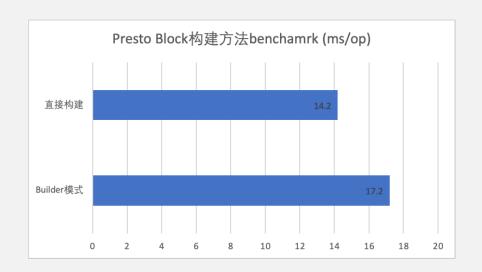




Data transformation



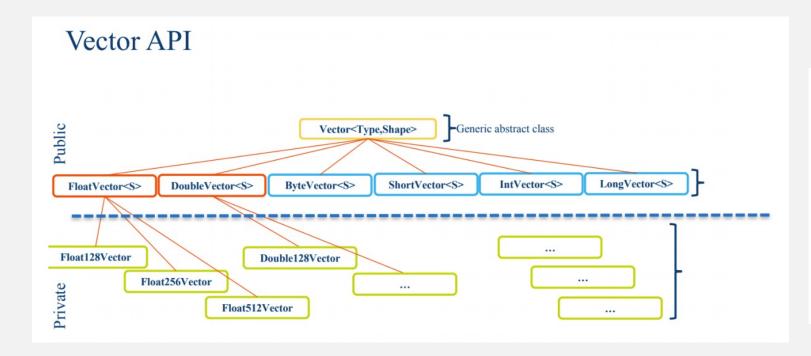
Presto Block	MercsDB Data
long[] Values	Result of Mercs API
bool[] isNulls	Vectorized





Java Vector API

- Incubator since JDK 16
- Tencent <u>KONA JDK 17</u>

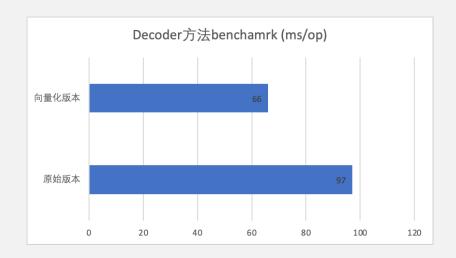


```
void scalarComputation(float[] a, float[] b, float[] c) {
  for (int i = 0; i < a.length; i++) {
      c[i] = (a[i] * a[i] + b[i] * b[i]) * -1.0f;
  }
}</pre>
```



Vectorization with Vector API

- Opt1: Unroll Loop
- Opt2: Unify Vector Species
- Opt3:
 - No Boxing & Unboxing
 - No Object creation
 - No function call



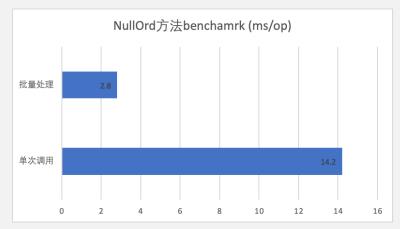
```
void decode(long[] input, long[] output) {
  for (int i = 0; i < input.length; ++ i) {
    long index = input[i];
    long offset = (index * 20) >>> 3;
    int v = readInt(offset) >>> 8;
    int shift = (int) ( ((index + 1) & 1) << 2);
    output[i] = (v >>> shift) & 0xFFFFF;
}
```

```
c void decodeVector(long[] input, long[] output) {
int len = input.length - input.length % (4*LONG_SPECIES.length());
for (i = 0; i < len; i += 4 * LONG_SPECIES.length()) {</pre>
   LongVector v1 = LongVector.fromArray(LONG_SPECIES, input, i);
   v1.lanewise(VectorOperators.LSHL, e: 4).add(v1.lanewise(VectorOperators.LSHL, e: 2))
           .lanewise(VectorOperators.LSHR, e: 3).add(currentOffset).intoArray(values, offset: 0);
   LongVector v2 = LongVector.fromArray(LONG_SPECIES, input, offset: i + LONG_SPECIES.length());
   v2.lanewise(VectorOperators.LSHL, e: 4).add(v2.lanewise(VectorOperators.LSHL, e: 2))
            .lanewise(VectorOperators.LSHR, e: 3).add(currentOffset).intoArray(values, LONG SPECIES.length());
   LongVector v3 = LongVector.fromArray(LONG_SPECIES, input, offset: i + 2 * LONG_SPECIES.length());
   v3.lanewise(VectorOperators.LSHL, e: 4).add(v3.lanewise(VectorOperators.LSHL, e: 2))
   LongVector v4 = LongVector.fromArray(LONG_SPECIES, input, offset: i + 3 * LONG_SPECIES.length());
   v4.lanewise(VectorOperators.LSHL, e: 4).add(v4.lanewise(VectorOperators.LSHL, e: 2))
            .lanewise(VectorOperators.LSHR, e: 3).add(currentOffset).intoArray(values, offset: 3 * LONG_SPECIES.length());
   LongVector.fromArray(LONG_SPECIES, values, offset: 0)
           .lanewise(VectorOperators.LSHR, e: 8)
           .lanewise(VectorOperators.LSHR, v1.add(1).and(1).lanewise(VectorOperators.LSHL, e: 2))
           .lanewise(VectorOperators.LSHR, @: 8)
            .lanewise(VectorOperators.LSHR, v2.add(1).and(1).lanewise(VectorOperators.LSHL, e: 2))
            .and(@xFFFFF).intoArray(output, offset: i + LONG_SPECIES.length())
   LongVector.fromArray(LONG_SPECIES, values, offset: 2 * LONG_SPECIES.length())
           .lanewise(VectorOperators.LSHR, @: 8)
           .lanewise(VectorOperators.LSHR, v3.add(1).and(1).lanewise(VectorOperators.LSHL, e: 2))
           .and(0xFFFFF).intoArray(output, offset: i + 2 * LONG_SPECIES.length());
   LongVector.fromArray(LONG_SPECIES, values, offset: 3 * LONG_SPECIES.length())
           .lanewise(VectorOperators.LSHR, e: 8)
            .lanewise(VectorOperators.LSHR, v4.add(1).and(1).lanewise(VectorOperators.LSHL, e: 2))
            .and(0xFFFFF).intoArray(output, offset: i + 3 * LONG_SPECIES.length());
```

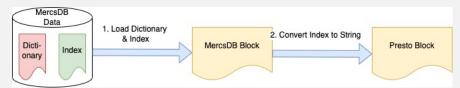


Others

Batch Vectorization



• Sequential memory access





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SSB Table: flat_lineorder

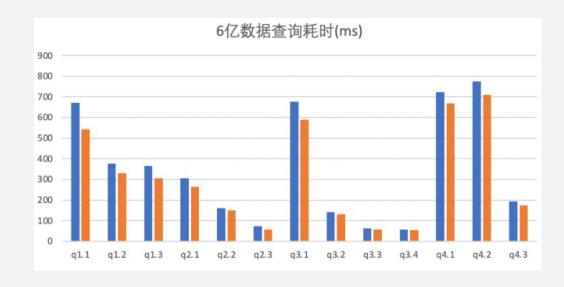


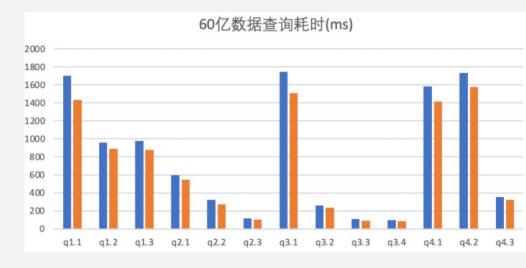
SSB Benchmark 1

•	Rows	Origin Size	MercsDB (with index)	ClickHouse (with spare and primary-key index)
	600M	200GB	89GB	60GB
	6B	2TB	882GB	593GB

60B rows







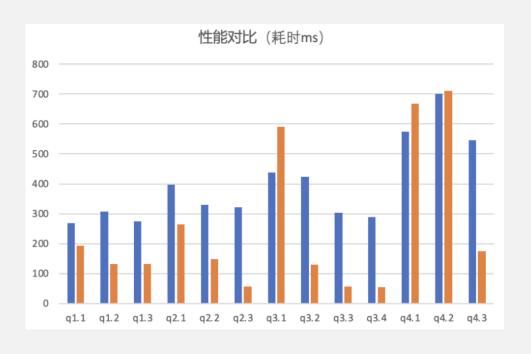
MercsDB

Optimized MercsDB



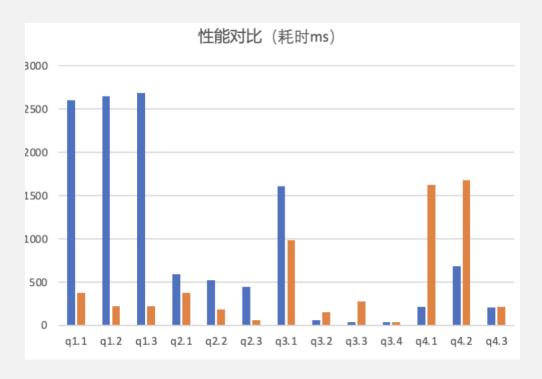
SSB Benchmark 2: QPS 1

600M rows



CK

60B rows

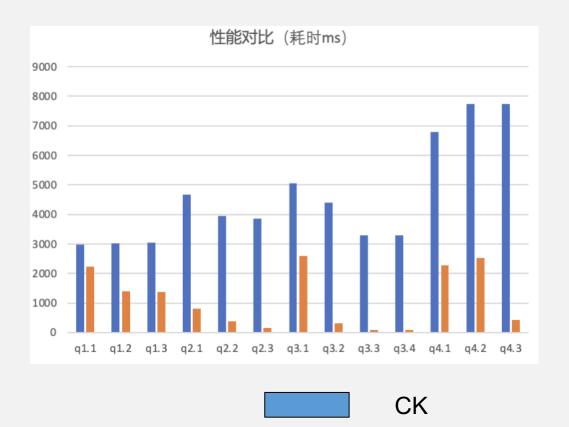


MercsDB

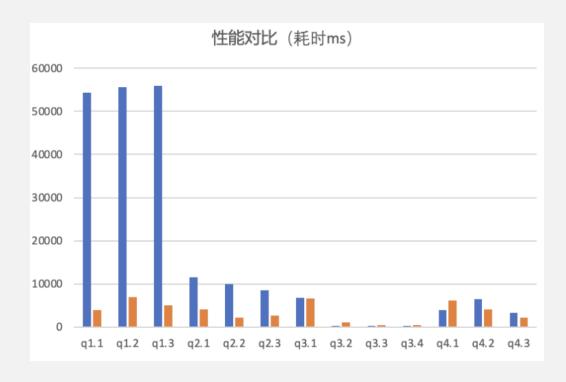


SSB Benchmark 3: QPS 20

600M rows



60B rows



MercsDB



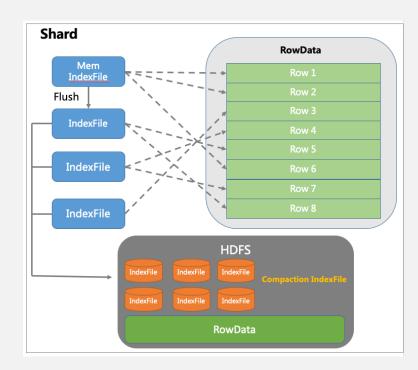
Log Retrieval in WeChat Pay

Background:

- 1. Real Time Write: 100 Billion Rows / day
- 2. Retrieval both Global and Specific
- 3. Mass Data

Solutions:

- 1. Write with TubeMQ
- 2. Participle and Index on Data
- 3. Separation of storage:
 - I. IndexFile: Local Disk
 - II. RowData: HDFS





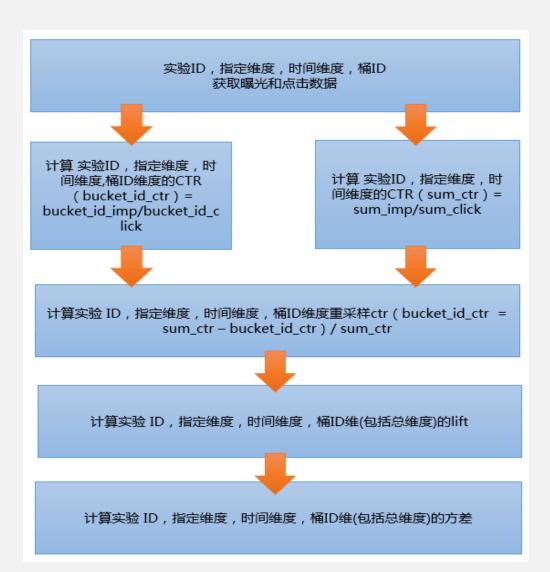
AB Test for Ads

Background:

- 1. Second response for Mass Data
- 2. Thousand columns in single query
- 3. Arbitrary Join

Solutions:

- Sorted by Primary Key (ad_id)
- 2. Compressions on metric column, 40% storage of origin data
- 3. Presto/Spark supported with Cache





Future Plan

- Cloud Service & Open-Source
- Vectorization
- Fault-Tolerance
- Memory Management
- ...



腾讯大数据



Thanks

