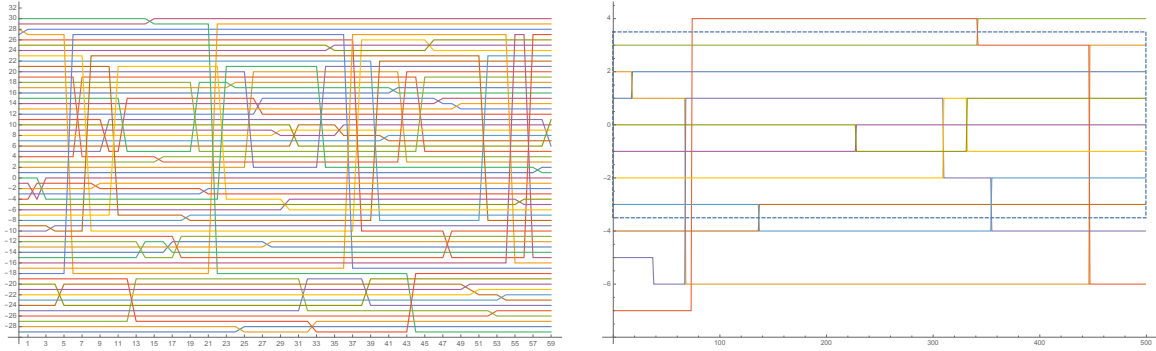


# TRAJECTORIES IN RANDOM MINIMAL TRANSPOSITION FACTORIZATIONS

by

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**Abstract.** — We study random typical minimal factorizations of the  $n$ -cycle, which are factorizations of  $(1, \dots, n)$  as a product of  $n - 1$  transpositions, chosen uniformly at random. Our main result is, roughly speaking, a local convergence theorem for the trajectories of finitely many points in the factorization. The main tool is an encoding of the factorization by an edge and vertex-labelled tree, which is shown to converge to Kesten's infinite Bienaymé-Galton-Watson tree with Poisson offspring distribution, uniform i.i.d. edge labels and vertex labels obtained by a local exploration algorithm.



**Figure 1.** Left: a simulation of the 60 trajectories  $(X_i^{(60)})_{-29 \leq i \leq 30}$  of a uniform minimal factorization of the 60-cycle. Right: a simulation of all the trajectories  $(X_i^{(500)})_{i \in \mathcal{I}_3^{(500)}}$  which cross the dashed rectangle  $[0, 500] \times [-3, 3]$ ; here  $\mathcal{I}_3^{(500)} = \{-7, -5, -4, -3, -2, -1, 0, 1, 2, 3\}$ .

## 1. Introduction

**1.1. Background and informal description of the results.** — We are interested in the combinatorial structure of typical minimal factorizations of the  $n$ -cycle as  $n \rightarrow \infty$ . Specifically, for an integer  $n \geq 1$ , we let  $\mathfrak{S}_n$  be the symmetric group acting on  $[n] := \{1, 2, \dots, n\}$  and we denote by  $\mathfrak{T}_n$  be the set of all transpositions of  $\mathfrak{S}_n$ . Let  $(1, 2, \dots, n)$  be the  $n$ -cycle which maps  $i$  to  $i + 1$  for  $1 \leq i \leq n - 1$ . The elements of the set

$$\mathfrak{M}_n := \{(\tau_1, \dots, \tau_{n-1}) \in \mathfrak{T}_n^{n-1} : \tau_1 \tau_2 \cdots \tau_{n-1} = (1, 2, \dots, n)\}$$

are called *minimal factorizations of  $(1, 2, \dots, n)$  into transpositions* (it is indeed easy to see that at least  $n - 1$  transpositions are required to factorize a  $n$ -cycle). In the sequel, the elements of  $\mathfrak{M}_n$  will be simply called minimal factorizations of size  $n$ . It is known since Dénes [Dén59] that  $|\mathfrak{M}_n| = n^{n-2}$  and bijective proofs were later given by Moszkowski [Mos89], Goulden & Pepper [GP93], Goulden & Yong [GY02].

This work is a companion paper of [FK17] from the same authors: both papers investigate the asymptotic properties of a *uniform random minimal factorization*  $\mathcal{F}^{(n)}$  of  $n$ . One important motivation is the work of Angel, Holroyd, Romik & Virág [AHRV07], who studied uniform random factorization of the reverse permutation by using only nearest-neighbor transpositions (where the reverse permutation  $\rho$  is defined by  $\rho(i) = n + i - 1$  for  $1 \leq i \leq n$ , and such factorizations are usually referred to as *sorting networks*). Further motivation for studying minimal factorizations is given in [FK17, Section 1.1].

Throughout the paper, we write  $\mathcal{F}^{(n)} = (t_1^{(n)}, \dots, t_{n-1}^{(n)})$  for a uniform random element in  $\mathfrak{M}_n$ . We may view  $\mathcal{F}^{(n)}$  as a random permutation-valued process starting at  $\text{id}_n$  and ending at  $(1, 2, \dots, n)$  by considering the partial products

$$(t_1^{(n)} t_2^{(n)} \cdots t_k^{(n)})_{1 \leq k \leq n-1}.$$

Since we view the partial products as a process indexed by  $k$ , the argument  $k$  will be referred to as the time. Each partial product acts on the *space*  $\{1, \dots, n\}$ . The adjectives “global” and “local” below refer to the space.

In [FK17], we considered the global geometry of these partial products and described the one dimensional marginals of this process, i.e., the various possible limiting behaviour for the partial product taken at time  $K_n$  (with  $K_n$  tending to  $+\infty$ ), depending on the behavior of  $K_n/\sqrt{n}$  as  $n \rightarrow \infty$ . In this paper we are interested in some *local* properties, namely in the trajectories of a given element  $i$  (and a fixed number of neighbouring elements) when applying successively the transpositions  $t_k^{(n)}$  ( $1 \leq k \leq n - 1$ ). We prove in Theorem 1 below that these trajectories converge to some random integer-valued step function, after some renormalization in time, but *without* any renormalization in space. Some combinatorial consequences of this result are also discussed.

This *discrete* limit behaviour of the trajectories is in contrast with the model of random sorting networks [AHRV07]: in the latter case, the limiting trajectories are random sine curves, as was conjectured by Angel, Holroyd, Romik and Virág [AHRV07] and recently proved by Dauvergne

[Dau18]; see also [ADHV17, GR17] for some local limit results with no space renormalization and a different time-renormalization.

*Convention:* we shall always multiply permutations from left to right, i.e.  $\tau_1 \dots \tau_{n-1}$  is the permutation obtained by first applying  $\tau_1$  then  $\tau_2$ , and so on (we warn the reader that this is opposite to the standard convention for compositions of functions).

**1.2. Main result: local convergence of the trajectories.** — In order to “zoom-in” around a neighborhood of 1, it is convenient to work with the  $n$ -cycle

$$(-\lfloor (n-1)/2 \rfloor, \dots, 0, 1, \dots, \lfloor n/2 \rfloor).$$

To this end, for every integer  $1 \leq a \leq n-1$ , we set  $\tilde{a} = a$  if  $a \leq n/2$  and  $\tilde{a} = i - n$  otherwise. If  $\tau = (a, b) \in \mathfrak{S}_n$ , we set  $\tilde{\tau} = (\tilde{a}, \tilde{b})$ . Finally, we set  $\tilde{\mathcal{F}}^{(n)} = (\tilde{t}_1^{(n)}, \dots, \tilde{t}_{n-1}^{(n)})$ . Now, for  $|i| \leq n/2$ , we define the *trajectory*  $X_i^{(n)}$  of  $i$  in  $\tilde{\mathcal{F}}^{(n)}$  by  $X_i^{(n)}(0) = i$  and

$$(1) \quad X_i^{(n)}(k) = \tilde{t}_1^{(n)} \tilde{t}_2^{(n)} \dots \tilde{t}_k^{(n)}(i), \quad 1 \leq k \leq n-1,$$

We also set  $X_i^{(n)}(n) = X_i^{(n)}(n-1)$ . See the left part of Fig. 1 for an illustration.

For every  $A \geq 1$ , denote by  $\mathcal{I}_A^{(n)}$  the indices of all the trajectories of  $\tilde{\mathcal{F}}^{(n)}$  that enter in the rectangle  $[0, n] \times [-A, A]$ ; formally,

$$\mathcal{I}_A^{(n)} = \{|i| \leq n/2 : \text{there exists } 0 \leq k \leq n-1, |X_i^{(n)}(k)| \leq A\}.$$

We will consider the vector of trajectories of all  $i$  in  $\mathcal{I}_A^{(n)}$ . Since this is a random set, let us clarify the underlying topology first.

We denote by  $\mathbb{D}([0, 1])$  the set of all real-valued left-continuous functions with right-hand limits (càdlàg functions for short) on  $[0, 1]$ , equipped with Skorokhod  $J_1$  topology (see [JS03, Chap VI] for background, but topologies should not be an issue since we will only work with step functions). For fixed  $I \subset \mathbb{Z}$ ,  $|I| < \infty$ , we equip  $\mathbb{D}([0, 1])^I$  with the product topology. We shall work in the space

$$\bigsqcup_{I \subset \mathbb{Z}, |I| < \infty} \mathbb{D}([0, 1])^I,$$

which is the disjoint union of these topologies, that is  $(f_i^{(n)})_{i \in I^{(n)}}$  tends to  $(f_i)_{i \in I}$  if and only if  $I^{(n)} = I$  for  $n$  large enough and  $f_i^{(n)}$  tends to  $f_i$  in the Skorokhod topology for every  $i$  in  $I$ .

We show that the trajectories of  $\tilde{\mathcal{F}}^{(n)}$  converge locally in distribution as  $n \rightarrow \infty$  in the following sense.

**THEOREM 1.** — *There is a family  $(X_i)_{i \in \mathbb{Z}}$  of integer-valued step functions on  $[0, 1]$  such that the following holds for every  $A \geq 1$ . Let  $\mathcal{I}_A$  denote the indices of all the trajectories  $(X_i)_{i \in \mathbb{Z}}$  that enter the rectangle  $[0, n] \times [-A, A]$ . Then  $\#\mathcal{I}_A < \infty$  almost surely and the convergence*

$$\left( X_i^{(n)}(\lfloor nt \rfloor) : 0 \leq t \leq 1 \right)_{i \in \mathcal{I}_A^{(n)}} \xrightarrow[n \rightarrow \infty]{(d)} (X_i)_{i \in \mathcal{I}_A}$$

*holds in distribution.*



**Figure 2.** The tree  $\text{Tree}(\widetilde{F}_n)$  associated with the minimal factorization  $\widetilde{F}_n = ((-1, -2)(5, -4)(1, 5)(2, 3)(1, -2)(2, 5)(-3, -2)(4, 5)(1, 0))$  of  $(-4, \dots, 0, 1, \dots, 5)$ .

This theorem is illustrated by the right image in Fig. 1, where we see the local nature of the trajectories for a large value of  $n$ .

Let us briefly explain the strategy to establish Theorem 1. The first step is to code  $\widetilde{\mathcal{F}}^{(n)}$  by an edge and vertex labelled tree  $\text{Tree}(\widetilde{\mathcal{F}}^{(n)})$  with vertex set  $\{-\lfloor (n-1)/2 \rfloor, \dots, 0, 1, \dots, \lfloor n/2 \rfloor\}$  and edge set  $\{\tilde{t}_1^{(n)}, \dots, \tilde{t}_{n-1}^{(n)}\}$  (each transposition is seen as a 2-element set). Furthermore, the edge  $\tilde{t}_i^{(n)}$  gets label  $i$ , and the tree is pointed at vertex with label 1 (see Fig. 2 for an example). Then, roughly speaking, the tree obtained by forgetting the vertex labels is simply a Bienaymé–Galton–Watson (BGW) tree with  $\text{Poisson}(1)$  offspring distribution with a uniform order on the edges (Section 2.2). Its local limit is therefore simply described in terms of Kesten’s infinite random BGW tree (Proposition 7). We then prove that the vertex labels can be reconstructed by using a local labelling algorithm (as explained in Section 2.3). Therefore, the edge and vertex labelled tree  $\text{Tree}(\widetilde{\mathcal{F}}^{(n)})$  converges in the local sense to a tree obtained by applying this local labelling algorithm to Kesten’s tree. The convergence of trajectories follows as a consequence.

**1.3. Combinatorial consequences.** — Our approach, based on an explicit relabelling algorithm, also allows us to obtain limit theorems for various “local” statistics of  $\widetilde{\mathcal{F}}^{(n)}$ . In this direction, let  $\widetilde{\mathbb{T}}_i^{(n)} = \{1 \leq j \leq n-1 : i \in \widetilde{t}_j^{(n)}\}$  be the set of indices of all transpositions moving  $i \in \mathbb{Z}$  (transpositions are here again identified with two-element sets), and let  $\widetilde{\mathbb{M}}_i^{(n)}$  be the set  $\{1 \leq j \leq n-1 : \widetilde{t}_1^{(n)} \cdots \widetilde{t}_{j-1}^{(n)}(i) \neq \widetilde{t}_1^{(n)} \cdots \widetilde{t}_j^{(n)}(i)\}$  of indices of all transpositions that affect the trajectory of  $i \in \mathbb{Z}$ . The following results will be deduced from Theorem 1 and from the construction of the limiting trajectories.

COROLLARY 2. — *With the above notation, the following assertions hold.*

- $$\begin{aligned}
(i) \quad & \text{For every } k \geq 1, (\#\widetilde{\mathbb{T}}_i^{(n)}, \#\widetilde{\mathbb{M}}_i^{(n)})_{|i| \leq k} \text{ converges in distribution to a random vector whose} \\
& \text{one dimensional marginal distributions are } 1 + \text{Poisson}(1) \text{ random variables;} \\
(ii) \quad & \text{As } n \rightarrow \infty, \mathbb{P}(X_1^{(n)}(k) \geq 1 \text{ for every } 0 \leq k \leq n-1) \rightarrow 1 - 1/e; \\
(iii) \quad & \text{For every } i, j \geq 1, \mathbb{P}\left(\#\widetilde{\mathbb{T}}_1^{(n)} = i, \#\widetilde{\mathbb{M}}_1^{(n)} = j\right) \xrightarrow{n \rightarrow \infty} \frac{e^{-2}}{(i-1)! \cdot (j-1)!};
\end{aligned}$$

(iv) For every  $i, j \geq 1$ :

$$\mathbb{P}\left(\#\tilde{\mathbb{T}}_1^{(n)} = i, \#\tilde{\mathbb{T}}_2^{(n)} = j\right) \xrightarrow{n \rightarrow \infty} e^{-2} \left( \frac{i+j-2}{(i+j-1)!} + \frac{i+j-1}{i!j!} - \frac{i+j-1}{(i+j)!} \right).$$

The event/statistics considered in items (ii), (iii) and (iv) above are somewhat arbitrary, the purpose is to show on specific examples how our construction allows the explicit computation of some limiting probabilities. More generally, the proof of Corollary 2 gives a means to determine the law of the limiting random vector in (i). The computation becomes however quickly cumbersome.

As for the results, it is interesting to note that  $\tilde{\mathbb{T}}_1^{(n)}$  and  $\tilde{\mathbb{M}}_1^{(n)}$  are asymptotically independent, even though they are not independent for fixed  $n$ . A conjecture on the joint distribution of  $(\tilde{\mathbb{T}}_1^{(n)}, \tilde{\mathbb{M}}_1^{(n)})$  at fixed  $n$  is given at the end of the introduction (Conjecture 1). Finally, we have not recognized any standard bivariate distribution for the limiting probability distribution in (iv).

**1.4. Symmetries.** — Finally, we are interested in symmetry properties that are satisfied by  $\mathcal{F}^{(n)}$ . Indeed, some of these symmetries are not visible on the limiting object, and may therefore help to compute limiting distributions. As a concrete example, the limiting Poisson distribution for  $\#\tilde{\mathbb{M}}_i^{(n)}$  given in Corollary 2(i) is proved as a consequence of an equidistribution result for  $\#\tilde{\mathbb{M}}_i^{(n)}$  and  $\#\tilde{\mathbb{T}}_i^{(n)}$ .

Our result is the following distributional identity (for convenience, we state it with the objects  $\mathbb{T}^{(n)}, \mathbb{M}^{(n)}$  defined exactly as  $\tilde{\mathbb{T}}^{(n)}, \tilde{\mathbb{M}}^{(n)}$  when replacing  $\tilde{\mathcal{F}}^{(n)}$  with  $\mathcal{F}^{(n)}$ ).

**THEOREM 3.** — *For every  $k \geq 1$ , it holds that*

$$\begin{aligned} & (\#\mathbb{T}_1^{(n)}, \#\mathbb{M}_1^{(n)}, \#\mathbb{T}_2^{(n)}, \#\mathbb{M}_2^{(n)}, \dots, \#\mathbb{T}_k^{(n)}, \#\mathbb{M}_k^{(n)}) \\ & \stackrel{(d)}{=} (\#\mathbb{M}_n^{(n)}, \#\mathbb{T}_1^{(n)}, \#\mathbb{M}_1^{(n)}, \#\mathbb{T}_2^{(n)}, \dots, \#\mathbb{M}_{k-1}^{(n)}, \#\mathbb{T}_k^{(n)}) \\ & \stackrel{(d)}{=} (\#\mathbb{M}_k^{(n)}, \#\mathbb{T}_k^{(n)}, \dots, \#\mathbb{M}_2^{(n)}, \#\mathbb{T}_2^{(n)}, \#\mathbb{M}_1^{(n)}, \#\mathbb{T}_1^{(n)}). \end{aligned}$$

**1.5. A conjecture.** — We conclude this Introduction with a conjectural formula for the bivariate probability generating polynomial of  $(\#\mathbb{T}_1^{(n)}, \#\mathbb{M}_1^{(n)})$  for a *fixed value* of  $n$ . Proving this conjecture would give an alternate proof of Corollary 2 (iii), that is of the asymptotic distribution of this pair of statistics.

**CONJECTURE 1.** — *Fix  $n \geq 2$ . Then*

$$\mathbb{E}[x^{\#\mathbb{T}_1^F} y^{\#\mathbb{M}_1^F}] = \left( \frac{xy(n-2+x+y)}{n} \right)^{n-2}.$$

This conjecture has been numerically checked until  $n = 8$ . Note that the right-hand side is known to be the bivariate probability generating polynomial of the degrees of the vertices labelled 1 and 2 in a uniform random Cayley tree (observe that this is different however from that of  $(\#\mathbb{T}_1^{(n)}, \#\mathbb{T}_2^{(n)})$ , i.e. the degrees of the vertices labelled 1 and 2 in  $\text{Tree}(\mathcal{F}^{(n)})$ ). None of the bijections between minimal factorizations and trees we are aware of explain this fact.

## 2. Local convergence of the factorization tree

**Table 1.** Table of the main notation and symbols appearing in Section 2.

$\mathfrak{M}_n$	The set of all minimal factorizations of $(1, 2, \dots, n)$ .
$\text{Tree}(F)$	The edge-labelled vertex-labelled non-plane tree associated with a minimal factorization $F$ .
$\widetilde{F}$	The minimal factorization obtained from $F$ by subtracting $n$ to all edge-labels larger than $n/2$ when $F \in \mathfrak{M}_n$ .
$\mathcal{T}_n$	A BGW tree with Poisson(1) offspring distribution conditioned on having $n$ vertices.
$\mathcal{T}_\infty$	The BGW tree with Poisson(1) offspring distribution conditioned to survive.
$\mathcal{E}(F)$	The pointed edge-labelled non-plane tree associated with a minimal factorization $F$ .
$\tau'$	When $\tau$ is an edge and vertex-labelled tree, $\tau'$ is the vertex-labelled tree obtained from $\tau$ by forgetting the edge labels.

The goal of this section is to establish Theorem 1, which is a local limit theorem for the trajectories of a random uniform minimal factorization  $\mathcal{F}_n$ . The main tool is the encoding of a minimal factorization  $F_n$  as a vertex and edge-labelled tree  $\text{Tree}(\widetilde{F}_n)$ , which was presented in the Introduction (see Fig. 2).

After providing background on trees and local convergence (Section 2.1), we will see that the tree  $\text{Tree}(\mathcal{F}_n)$  without vertex labels is a Bienaymé-Galton-Watson (BGW) tree with a uniform edge-ordering. The local limit of such a tree follows from standard results in the random tree literature (Section 2.2). We then explain (Section 2.3) how to reconstruct the *vertex labels* by a local labelling algorithm, which can also be run on infinite trees (Section 2.4). Continuity and locality properties of the labelling algorithm yield a local limit result for the vertex and edge-labelled tree  $\text{Tree}(\widetilde{F}_n)$  (Theorem 13) Theorem 1 follows easily (Section 2.5).

### 2.1. Preliminaries on trees. —

**2.1.1. Plane trees.** — We use Neveu’s formalism [Nev86] to define (rooted) plane trees: let  $\mathbb{N} = \{1, 2, \dots\}$  be the set of all positive integers, and consider the set of labels  $\mathbb{U} = \bigcup_{n \geq 0} \mathbb{N}^n$  with the convention  $\mathbb{N}^0 = \{\emptyset\}$ . For  $u = (u_1, \dots, u_n) \in \mathbb{U}$ , we denote by  $|u| = n$  the length of  $u$ ; if  $n \geq 1$ , we define  $pr(u) = (u_1, \dots, u_{n-1})$  and for  $i \geq 1$ , we let  $ui = (u_1, \dots, u_n, i)$ ; more generally, for  $v = (v_1, \dots, v_m) \in \mathbb{U}$ , we let  $uv = (u_1, \dots, u_n, v_1, \dots, v_m) \in \mathbb{U}$  be the concatenation of  $u$  and  $v$ . A plane tree is a nonempty subset  $\tau \subset \mathbb{U}$  such that (i)  $\emptyset \in \tau$ ; (ii) if  $u \in \tau$  with  $|u| \geq 1$ , then  $pr(u) \in \tau$ ; (iii) if  $u \in \tau$ , then there exists an integer  $k_u(\tau) \geq 0$  such that  $ui \in \tau$  if and only if  $1 \leq i \leq k_u(\tau)$ . Observe that plane trees are rooted at  $\emptyset$  by definition.

We will view each vertex  $u$  of a tree  $\tau$  as an individual of a population for which  $\tau$  is the genealogical tree. The vertex  $\emptyset$  is called the root of the tree and for every  $u \in \tau$ ,  $k_u(\tau)$

is the number of children of  $u$  (if  $k_u(\tau) = 0$ , then  $u$  is called a leaf, otherwise,  $u$  is called an internal vertex),  $|u|$  is its generation,  $pr(u)$  is its parent and more generally, the vertices  $u, pr(u), pr \circ pr(u), \dots, pr^{|u|}(u) = \emptyset$  are its ancestors. To simplify, we will sometimes write  $k_u$  instead of  $k_u(\tau)$ . A plane tree is said to be locally finite if all its vertices have a finite number of children. Finally, if  $\tau$  is a tree and  $h$  is a nonnegative integer, we let  $[[\tau]]_h = \{u \in \tau : |u| \leq h\}$  denote the tree obtained from  $\tau$  by keeping the vertices in the first  $h$  generations.

**2.1.2. Non-plane trees.** — In parallel to plane trees, since  $\text{Tree}(F_n)$  is by definition a non-plane tree, we will need to consider non-plane trees. By definition, a non-plane tree is a connected graph without cycles. A non-plane tree with a distinguished vertex is called pointed. Equivalently, a pointed non-plane tree is an equivalence class of plane trees under permutation of the order of the children of its vertices, the root of the plane tree being the distinguish vertex. If  $\tau$  is a plane tree, we denote by  $\text{Shape}(\tau)$  the corresponding non-plane tree (informally, we forget the planar structure of  $\tau$ , and point it at the root vertex).

**2.1.3. Labelled trees.** — In this work, we consider trees which are:

- *edge labelled*, in the sense that all edges carry real-valued labels,
- *vertex labelled*, in the sense that some vertices (potentially none, potentially all) carry integer values.

Edge labels will always be assumed to be distinct, and vertex labels as well. We do not impose any relation between vertex and edge labellings. In the figures, we shall use framed labels for vertex labels to avoid confusion with edge labels.

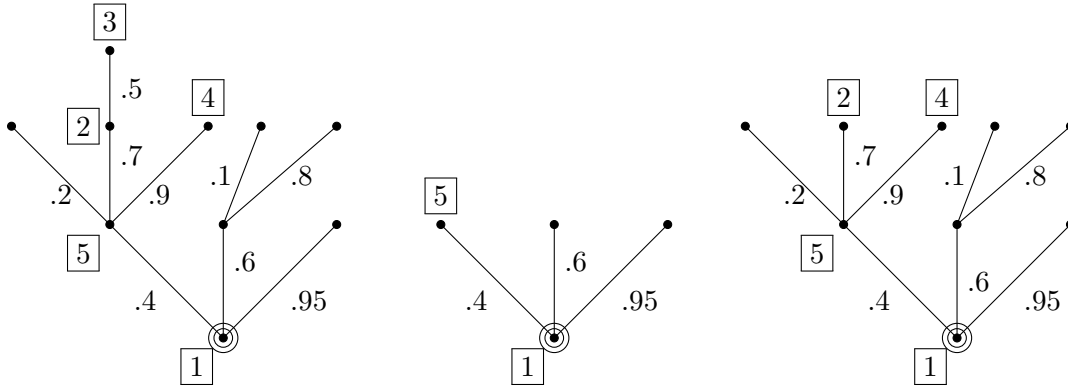
To simplify notation, we say that a tree is **E**-labelled, **V**-labelled or **EV**-labelled if it is respectively edge labelled, vertex labelled, or edge and vertex-labelled.

**2.1.4. Local convergence for labelled trees.** — Let  $\tau$  be a locally finite (potentially infinite) non-plane tree, pointed, **EV**-labelled in the sense defined above. As for plane trees, for every integer  $h \geq 1$ , we denote by  $[[\tau]]_h$  the finite non-plane, pointed, **EV**-labelled tree obtained from  $\tau$  by keeping only the vertices at distance at most  $h$  from the pointed vertex (together with the edges between them and their labels). We say that  $[[\tau]]_h$  is a *labelled ball*, see Figure 3 for an example.

We consider a topology on the set of all locally finite, non-plane, pointed, **EV**-labelled trees such that  $\tau_n \rightarrow \tau$  if and only if for each  $h \geq 1$ , the **V**-labelled ball  $[[\tau_n]]_h$  and  $[[\tau]]_h$  coincide for  $n$  large enough and the edge labels of  $[[\tau_n]]_h$  tends to that of  $[[\tau]]_h$ . It is easy to construct a metric  $d_{\text{loc}}$  for such a topology and to see that the resulting metric space is Polish. Similar metrics can be defined for other families of trees (plane, unlabelled or only **E** or **V**-labelled, etc.) and will be also denoted by  $d_{\text{loc}}$ .

Finally, if  $\tau$  is an **EV**-labelled tree and  $\alpha > 0$ , we denote by  $\alpha \cdot \tau$  the **EV**-labelled tree obtained by multiplying all the *edge labels* by  $\alpha$  (vertex labels are not modified!), and we denote by  $\tau'$  the **V**-labelled tree obtained from  $\tau$  by forgetting the edge labels.

**2.1.5. Random BGW trees.** — Let  $\mu$  be a probability measure on  $\mathbb{Z}_+$  (called the offspring distribution) such that  $\mu(0) > 0$ ,  $\mu(0) + \mu(1) < 1$  (to avoid trivial cases). When  $\sum_{i \geq 0} i\mu(i) \leq 1$ , the BGW measure with offspring distribution  $\mu$  is a probability measure  $\text{BGW}^\mu$  on the set of all finite trees such that  $\text{BGW}^\mu(\tau) = \prod_{u \in \tau} \mu(k_u)$  for every finite tree  $\tau$ . For every integer  $n \geq 1$ ,



**Figure 3.** From left to right: a pointed EV-labelled tree  $\tau$  (considered as non-plane), and its labelled balls  $[[\tau]]_1$  and  $[[\tau]]_2$ .

we denote by  $\text{BGW}_n^\mu$  the conditional probability measure  $\text{BGW}^\mu$  given the set  $\mathbb{A}_n$  of all trees with  $n$  vertices, so that

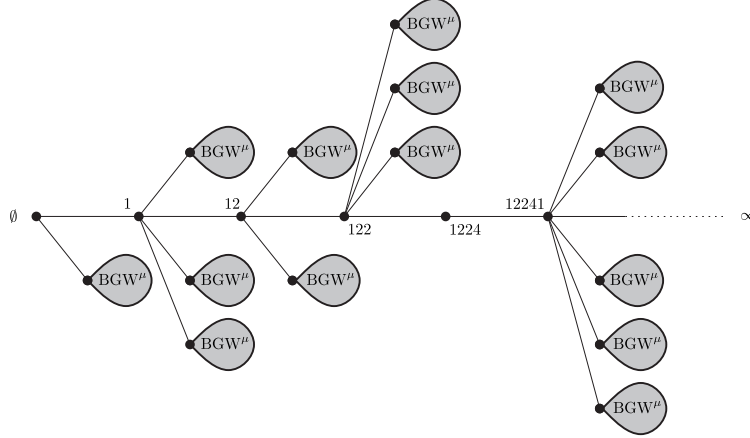
$$\text{BGW}_n^\mu(\tau) = \frac{\text{BGW}^\mu(\tau)}{\text{BGW}^\mu(\mathbb{A}_n)}, \quad \tau \in \mathbb{A}_n$$

(we always implicitly restrict ourselves to the values of  $n$  such that  $\text{BGW}^\mu(\mathbb{A}_n) > 0$ ). This is the distribution of a BGWrandom tree with offspring distribution  $\mu$ , conditioned on having  $n$  vertices.

**2.1.6. The infinite BGW tree.** — Let  $\mu = (\mu_i)_{i \geq 0}$  be a critical offspring distribution (meaning that  $\sum i\mu_i = 1$ ) with  $\mu_1 \neq 1$ . The conditioned random tree  $\text{BGW}_n^\mu$  is known to have a local limit  $\mathcal{T}_\infty$ , which we now present (see [Jan12, Section 5] and [LP16] for a formal definition of  $\mathcal{T}_\infty$ ). Let  $\bar{\mu}$  be the size-biased distribution of  $\mu$  defined by  $\bar{\mu}_k = k\mu_k$  for  $k \geq 0$ . In  $\mathcal{T}_\infty$ , there are two types of nodes: normal nodes and special, with the root being special. Normal nodes have offspring according to independent copies of  $\mu$ , while special nodes have offspring according to independent copies of  $\bar{\mu}$ . Moreover, all children of a normal node are normal; among all the children of a special node one is selected uniformly at random and is special, while all other children are normal. The tree  $\mathcal{T}_\infty$  has a unique infinite path (called the spine) formed by all the special vertices (see Fig. 4 for an illustration of this construction).

In this paper, we shall use the particular case where  $\mu$  is a  $\text{Poisson}(1)$  offspring distribution. Moreover, we consider trees with edge labels. Specifically, let  $\mathcal{T}_n$  a BGW tree with  $\text{Poisson}(1)$  offspring distribution conditioned on having  $n$  vertices, let  $\mathcal{T}_n^e$  be the tree obtained from  $\mathcal{T}_n$  by labelling its edges in a uniform way, using once each integer from 1 to  $n-1$ . Similarly, let  $\mathcal{T}_\infty^e$  be the above constructed random tree  $\mathcal{T}_\infty$ , with edges labelled by i.i.d random variables following the uniform distribution on  $[0, 1]$ . Finally, we recall that  $\frac{1}{n}\text{Shape}(\mathcal{T}_n^e)$  is the tree obtained by multiplying the edge labels of  $\mathcal{T}_n^e$  by  $\frac{1}{n}$  (there are no vertex labels for the moment). The following result is an adaptation to our need of standard local convergence result for BGW trees.





**Figure 4.** An illustration of  $\mathcal{T}_\infty$ .

PROPOSITION 4. — *The convergence*

$$\frac{1}{n} \text{Shape}(\mathcal{T}_n^e) \xrightarrow[n \rightarrow \infty]{(d)} \text{Shape}(\mathcal{T}_\infty^u)$$

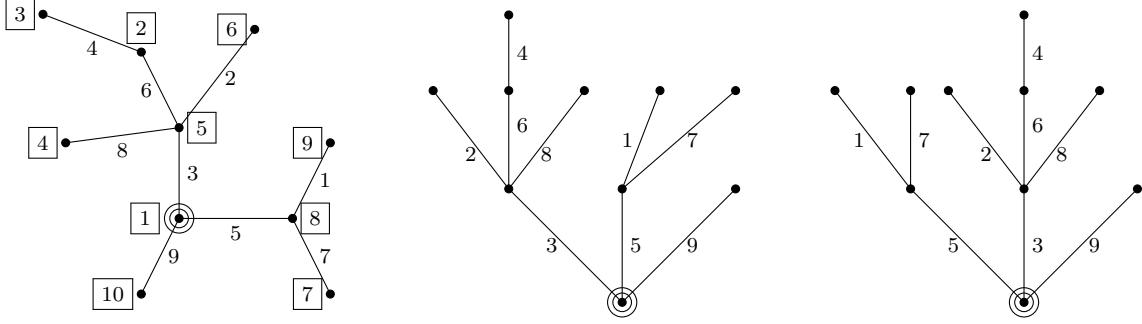
holds in distribution in the set of all locally finite, non-plane pointed,  $E$ -labelled trees equipped with  $d_{\text{loc}}$ .

*Proof.* — The convergence  $\mathcal{T}_n \rightarrow \mathcal{T}_\infty$  of plane non-labelled trees in distribution in the local topology is a classical result in random tree theory, see, e.g., [Jan12, AD14]. The map  $\text{Shape}$  is clearly continuous, implying that  $\text{Shape}(\mathcal{T}_n) \rightarrow \text{Shape}(\mathcal{T}_\infty)$  locally in distribution. We therefore only need to justify that, for each  $h \geq 1$ , the edge labels at height at most  $h$  on the left-hand side jointly converge to those on the right-hand side. This is however obvious, since, on the left-hand side, we have a uniform labeling with the numbers  $1/n, 2/n, \dots, (n-1)/n$ , while on the right we have independent uniform label in  $[0, 1]$ .  $\square$

**2.2. Forgetting the vertex labels yields a random BGW tree.** — In this section, as an intermediate step to study  $\text{Tree}(\mathcal{T}_n)$  and  $\text{Tree}(\widetilde{\mathcal{T}}_n)$ , we consider a variant without vertex labels. Namely, for a minimal factorization  $F_n$ , we introduce  $\mathcal{E}(F_n)$ , the pointed  $E$ -labelled non-plane tree obtained from  $\text{Tree}(F_n)$  by pointing the vertex with label 1 and forgetting other vertex labels. See Fig. 5 for an example. As shown by Moszkowski [Mos89], the map  $\mathcal{E}$  is a bijection. Since the proof is short and elegant, we include it here.

PROPOSITION 5 (MOSZKOWSKI). — *The map  $\mathcal{E}$  is a bijection from  $\mathfrak{M}_n$  to the set of all pointed  $E$ -labelled non-plane trees with  $n$  vertices (where edges are labelled from 1 to  $n-1$ ).*

*Proof.* — Let  $T$  be a non-plane  $EV$ -labelled tree with  $n$  vertices (with vertices labelled from 1 to  $n$  and edges from 1 to  $n-1$ ). We associate with the tree  $T$  the sequence  $\mathcal{S}(T)$  of the  $n-1$  transpositions  $(a_1, b_1) \dots (a_{n-1}, b_{n-1})$ , where  $a_i$  and  $b_i$  are the vertex-labels of the extremities of the edge labelled  $i$ . It is known since the work of Dénes [Dén59] that the function  $\mathcal{S}$  is a bijection between  $EV$ -labelled trees with  $n$  vertices and minimal factorizations of *all* cyclic permutations of length  $n$ .



**Figure 5.** Construction of the trees associated with the factorization  $F_n = ((8, 9) (5, 6) (1, 5) (2, 3) (1, 8) (2, 5) (7, 8) (4, 5) (1, 10))$  of  $(1, 2, \dots, 10)$ . On the left, we have the pointed non-plane EV-labelled tree  $\text{Tree}(F_n)$ . On the middle and on the right, we have represented  $\mathcal{E}(F_n)$ , which is obtained from  $\text{Tree}(F_n)$  by forgetting the vertex labels. To emphasize that  $\mathcal{E}(F_n)$  is non plane, and we have represented the *same* tree  $\mathcal{E}(F_n)$  with two different embeddings.

Observe that if we forget the vertex-labels of  $T$  except 1, then there is a unique way to relabel the  $n - 1$  other vertices to get a factorization of the cycle  $(1, 2, \dots, n)$ . Indeed, relabelling these labels amounts to conjugating the associated cyclic permutation by a permutation fixing 1 and there is always a unique way to conjugate a cyclic permutation by a permutation fixing 1 to get the cycle  $(1, 2, \dots, n)$ .  $\square$

Recall that we denote by  $\mathcal{F}_n$  is a uniform random minimal factorization in  $\mathfrak{M}_n$ . It turns out that the law of  $\mathcal{E}(\mathcal{F}_n)$  can be related to a BGW tree as follows. As in the previous section, let  $\mathcal{T}_n^e$  be the plane E-labelled tree obtained from a  $\text{Poisson}(1)$  BGW tree  $\mathcal{T}_n$  by labelling its edges from 1 to  $n - 1$  in a uniform way. Building on Proposition 5, we can now prove the following result.

**PROPOSITION 6.** — *The random trees  $\mathcal{E}(\mathcal{F}_n)$  and  $\text{Shape}(\mathcal{T}_n^e)$  both follow the uniform distribution on the set of all pointed non-plane E-labelled trees with  $n$  vertices (with edge label set  $\llbracket n - 1 \rrbracket$ ).*

*Proof.* — Since  $\mathcal{E}$  is a bijection from the set  $\mathfrak{M}_n$  to the set of all pointed non-plane E-labelled trees with  $n$  vertices, it is immediate that  $\mathcal{E}(\mathcal{F}_n)$  follows the uniform distribution of the latter set.

Then, note that there is a natural bijection  $\phi$  between pointed non-plane E-labelled trees (with edge label set  $\llbracket n - 1 \rrbracket$ ) and non-plane V-labelled trees (with vertex set  $\llbracket n \rrbracket$ , also known as Cayley trees): label the pointed vertex by 1, and then label every other vertex  $v$  with  $\ell(v^{\leftarrow}) + 1$ , where  $\ell(v^{\leftarrow})$  is the label of the edge adjacent to  $v$  closest to the pointed vertex.

By construction,  $\phi(\text{Shape}(\mathcal{T}_n^e))$  is obtained from  $\mathcal{T}_n$  by labelling the root with 1 and other vertices uniformly with numbers from 1 to  $n$  (and forgetting the root and the planar structure). It is well-known (see the second proof of Theorem 3.17 in [Hof16]) that this is a uniform random

non-plane V-labelled tree with  $n$  vertices. Applying  $\phi^{-1}$ , it follows that  $\text{Shape}(\mathcal{T}_n^e)$  is a uniform random pointed non-plane E-labelled trees with  $n$  vertices. This completes the proof.  $\square$

By combining the local convergence of  $\text{Shape}(\mathcal{T}_n^e)$  (Proposition 4) with Proposition 6, we get the following:

PROPOSITION 7. — *The convergence*

$$(2) \quad \frac{1}{n} \cdot \mathcal{E}(\mathcal{T}_n) \xrightarrow[n \rightarrow \infty]{(d)} \text{Shape}(\mathcal{T}_\infty^u)$$

holds in distribution in the set of all locally finite, non-plane, pointed, E-labelled trees equipped with  $d_{\text{loc}}$ .

We insist on the fact that the above proposition is a convergence of *edge-labelled* trees: indeed, the objects in (2) do not carry vertex labels. The difficulty is now to insert the vertex labels in the above proposition. This is not as easy as for the edge labels, since in  $\text{Tree}(\mathcal{T}_n)$ , the vertex labels are determined by the edge labels (otherwise  $\mathcal{E}$  would not be a bijection). We will see in the next sections that the vertex labels can actually be recovered by a local exploration algorithm. This enables to define a labelling procedure on infinite tree  $\mathcal{T}_\infty^u$ , giving a limit for the EV-labelled tree  $\text{Tree}(\widetilde{\mathcal{T}}_n)$ .

**2.3. The labelling algorithm.** — Consider a locally finite E-labelled tree  $\tau$ , either plane, or non-plane pointed.

If  $k \geq 1$  and  $\tau$  has at least  $k$  vertices, we define a procedure  $\text{Find}_k(\tau, \ell)$  which successively assigns values  $1, 2, \dots, k$  to  $k$  vertices of  $\tau$  as follows. The value 1 is given to the pointed vertex. Assume that  $i$  has been assigned and that we want to assign  $i + 1$ . Among all edges adjacent to  $i$  in  $\tau$ , denote by  $e_1$  the one with the smallest label, and set  $\ell_1 = \ell_{e_1}$ . Call  $j_1$  the other extremity of  $e_1$  and consider all edges  $e$  adjacent to  $j_1$  with  $\ell_e > \ell_1$ .

- If there are none, then assign value  $i + 1$  to  $j_1$ .
- Otherwise, among all such edges, denote by  $e_2$  the one with smallest label, and set  $\ell_2 = \ell_{e_2}$  (in particular  $\ell_2 > \ell_1$ ). Call  $j_2$  the other extremity of  $e_2$ . We then iterate the process: if there is no edge adjacent to  $j_2$  with label greater than  $\ell_2$ , then we assign value  $i + 1$  to the vertex  $j_2$ . Otherwise, among all edges adjacent to  $j_2$  with labels greater than  $\ell_2$ , we consider the one with smallest label, and so one.

We stop the procedure when  $k$  has been assigned. Formally, this is described in Algorithm 1.

EXAMPLE. — We explain in detail how the algorithm runs on the pointed E-labelled tree on the left part of Fig. 6. The result is the middle picture in Fig. 6.

First the root is assigned value  $\boxed{1}$ . The edge adjacent to the root with smallest label is the left-most one, so that in this case  $\ell_{e_1} = .4$  and  $j_1$  is the left-most vertex at height 1. Continuing the process,  $j_1$  has two adjacent edges of labels bigger than  $\ell_{e_1} = .4$ . We pick the one with smallest label  $e_2$ , here  $\ell_{e_2} = .7$ , and call  $j_2$  its other extremity (that is the second left-most vertex at height 2). Now,  $j_2$  has no adjacent edge with label bigger than  $\ell_{e_2} = .7$ , so we assign value  $\boxed{2}$  to  $j_2$ . Starting now from  $\boxed{2}$ , its adjacent edge with smallest label is the edge  $e_3$  with label .5 going to a leaf. This leaf has no other adjacent edges, so in particular none with label

**Algorithm 1** Assign values  $1, 2, \dots, k$  to  $k$  vertices of  $\tau$ 


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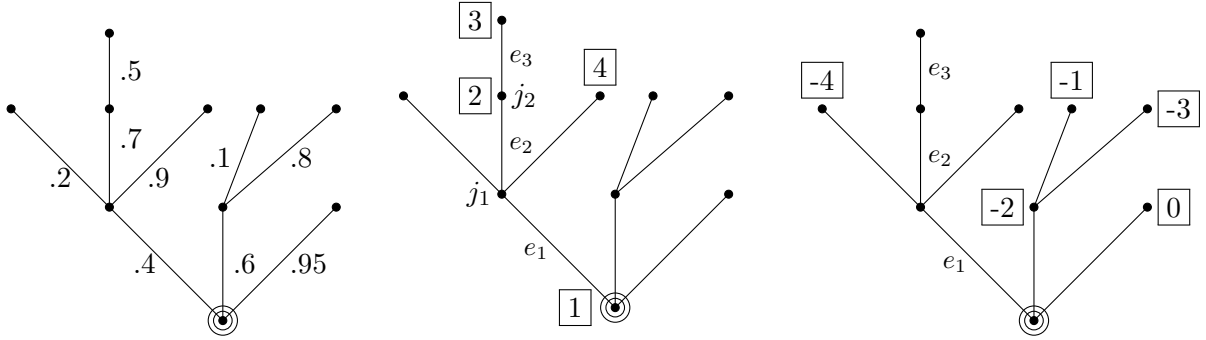
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value(root)  $\leftarrow$  1
 $i \leftarrow 1$ 
while  $i \leq k$  do
  CurrentVertex  $\leftarrow$  vertex with value  $i$ 
  HigherVertices  $\leftarrow$  all neighbours of CurrentVertex
  while HigherVertices  $\neq \emptyset$  do
     $x \leftarrow$  minimal value of  $\ell_{\{\text{CurrentVertex}, v\}}$  for  $v$  in HigherVertices
    CurrentVertex  $\leftarrow$  vertex  $v$  in HigherVertices such that  $\ell_{\{\text{CurrentVertex}, v\}} = x$ 
    HigherVertices  $\leftarrow$  all vertices  $v$  such that  $\ell_{\{\text{CurrentVertex}, v\}} > x$ 
  end while
  value(CurrentVertex)  $\leftarrow i + 1$ 
   $i \leftarrow i + 1$ 
end while

```

---

bigger than .5. We therefore assign value  $\boxed{3}$  to this leaf. Finally, in order to assign the last value, starting from  $\boxed{3}$ , we go through the edge  $e_3$ , go through  $e_2$ , then go through the edge with label .9 and arrive to a leaf which is given value  $\boxed{4}$ . The procedure  $\text{Find}_4$  is over.



**Figure 6.** On the left: a plane tree  $\tau$  with edge-labels  $\ell$ . In the middle:  $\text{Find}_4(\tau, \ell)$ . On the right:  $\overline{\text{Find}_4}(\tau, \ell)$ .

Note that, by construction, this algorithm does not use the planar structure if  $\tau$  is a plane tree. More precisely, if  $(\tau, \ell)$  is a plane E-labelled tree, then we have the commutation relation  $\text{Find}_k(\text{Shape}(\tau), \ell) = \text{Shape}(\text{Find}_k(\tau, \ell))$ .

Recall from Section 2.2 that with a minimal factorization  $F_n$  we have associated two different trees  $\text{Tree}(F_n)$  (which is a non-plane EV-labelled tree) and  $\mathcal{E}(F_n)$  (which is the pointed E-labelled non-plane tree, obtained from  $\text{Tree}(F_n)$  by forgetting vertex labels). The following lemma explains how to go from  $\mathcal{E}(F_n)$  to  $\text{Tree}(F_n)$ , using the above-defined algorithm.

LEMMA 8. — *Let  $F_n$  be a minimal factorization of  $(1, 2, \dots, n)$ . We have  $\text{Find}_n(\mathcal{E}(F_n)) = \text{Tree}(F_n)$ .*

*Proof.* — By construction  $\text{Find}_n(\mathcal{E}(F_n))$  and  $\text{Tree}(F_n)$  might only differ by their vertex labels. We prove by induction that for all  $i \leq n$  the same vertex carries label  $i$  in both trees.

By definition,  $\text{Find}_n$  gives the value 1 to the root, and the root of  $\mathcal{E}(F_n)$  is the vertex that used to have value 1 in  $\text{Tree}(F_n)$ . This proves the base case ( $i = 1$ ) of the induction.

Fix  $i \geq 1$  and assume that  $\text{Find}_n(\mathcal{E}(F_n))$  gives value  $i$  to the vertex  $v_i$  of  $\mathcal{E}(F_n)$  that used to have value  $i$  in  $\text{Tree}(F_n)$ . By construction, to assign value  $i + 1$ ,  $\text{Find}_n(\mathcal{E}(F_n))$  first considers the edge adjacent to  $v_i$  with minimum edge-label. Call  $a$  this label. This means that the transposition  $\tau_a$  in  $F_n$  is of the form  $(i j_1)$  for some  $j_1$ . (This notation is consistent with the construction of  $\text{Find}_n$ .) By minimality of  $a$ , the transpositions  $\tau_1, \dots, \tau_{a-1}$  fix  $i$ . Thus the partial product  $\tau_1 \dots \tau_a$  maps  $i$  on  $j_1$ . We then want to see where  $j_1$  is mapped when we apply the next transpositions  $\tau_{a+1}, \tau_{a+2}, \dots$ . For this, we need to look for an edge adjacent to  $j_1$  with a value bigger than  $a$ , which is exactly what  $\text{Find}_n$  does. If  $b$  is the smallest value of such an edge and  $j_2$  the other extremity of this edge (again the notation is consistent with the one of the construction of  $\text{Find}_n$ ), then  $\tau_1 \dots \tau_b$  maps  $i$  to  $j_2$ . The construction stops at  $j_r$  when there is no edge adjacent to  $j_r$  with a bigger value than the previously considered edge, and then we know that  $\tau_1 \dots \tau_{r-1}$  maps  $i$  to  $j_r$ . But, since  $\tau_1 \dots \tau_{r-1} = (1, 2, \dots, n)$ , we necessarily have  $j_r = i + 1$ . Thus  $\text{Find}_n$  precisely assigns  $i + 1$  to  $j_r$  and this completes our induction step.  $\square$

REMARK 1. — Note that if one replaces the labels  $(\ell_e)$  of  $\mathcal{E}(F_n)$  with labels  $(\ell'_e)$  such that  $\ell_e < \ell_f$  if and only if  $\ell'_e < \ell'_f$  for every edges  $e, f$  (we say that the edge-labellings are *compatible*), one obtains the same vertex labels when running  $\text{Find}_n$  (since  $\text{Find}_n$  only uses the relative order of the labels).

Figs. 5 and 6 illustrate the previous lemma. Indeed the trees in the middle of Fig. 5 and in the left part of Fig. 6 are the same, with compatible edge-labellings. The procedure  $\text{Find}_4$  indeed reassigns the labels 1, 2, 3 and 4 to some vertices of  $\mathcal{E}(F_n)$ , as they are in  $\text{Tree}(F_n)$  (compare the left part of Fig. 5 and the middle picture in Fig. 6).

Lemma 8 explains how to reconstruct  $\text{Tree}(F_n)$  from  $\mathcal{E}(F_n)$ . We are however interested in  $\text{Tree}(\widetilde{F_n})$ , rather than  $\text{Tree}(F_n)$ . We therefore need to introduce a dual labelling procedure  $\overline{\text{Find}}_k$ , which assigns non positive labels. This procedure runs exactly as  $\text{Find}_k$  except that the order of label edges is taken as reversed. Namely we first look at the edge with *largest* label  $\ell_1$  incident to 1, call  $j_1$  its extremity and then look for an edge of *largest* label  $\ell_2 < \ell_1$ , etc. Another difference is that  $\overline{\text{Find}}_k$  now assigns labels  $0, -1, \dots, -k$  to successively found vertices. An example of the outcome of this procedure is shown on Fig. 6.

The following lemma motivates the definition of this dual procedure.

LEMMA 9. — *Let  $F_n$  be a minimal factorization of  $(1, 2, \dots, n)$ . Then  $\overline{\text{Find}}_{n-1}(\mathcal{E}(F_n))$  is obtained from  $\text{Tree}(F_n)$  by subtracting  $n$  to every vertex label.*

*Proof.* — With a minimal factorization  $F_n = (\tau_1, \dots, \tau_{n-1})$  of the full cycle  $(1, 2, \dots, n)$ , we can associate its reversed sequence  $r(F_n) = (\tau_{n-1}, \dots, \tau_1)$ , which is a factorization of the full cycle  $(1, n, n-1, \dots, 2)$ . Reversing the order of the edge labels in  $\text{Tree}(F_n)$  yields  $\text{Tree}(r(F_n))$ . By Lemma 8 (which holds more generally for minimal factorizations of any cycle), the procedure

$\overline{\text{Find}}$  thus assigns label to  $\mathcal{E}(F_n)$  in the order of the cycle  $(1, n, n-1, \dots, 2)$ , i.e.  $n$  gets the first label 0 (after the pointed vertex),  $n-1$  gets the second label  $-1$ , and so on. This proves the lemma.  $\square$

Combining Lemmas 8 and 9, we see that the tree  $\text{Tree}(\widetilde{F}_n)$ , in which we are interested, is obtained by composing both procedures.

COROLLARY 10. — *Let  $F_n$  be a minimal factorization of  $(1, 2, \dots, n)$ . Then*

$$\text{Tree}(\widetilde{F}_n) = \text{Find}_{\lfloor \frac{n}{2} \rfloor} \circ \overline{\text{Find}}_{\lfloor \frac{n-1}{2} \rfloor} (\mathcal{E}(F_n)).$$

**2.4. Relabelling infinite trees.** — We now want to run the procedures  $\text{Find}$  and  $\overline{\text{Find}}$  on infinite, yet locally finite, trees. Note that in general,  $\text{Find}_k$  (and  $\overline{\text{Find}}_k$ ) may be ill-defined (since the inner “while” loop in Algorithm 1 may be infinite).

Let  $T$  be a locally finite tree (either plane, or non-plane pointed) and let  $\ell = (\ell_e)$  be a family of distinct real numbers indexed by the edges of  $T$ . We say that  $\mathbf{T} = (T, \ell)$  satisfies the property  $(\mathcal{I})$  (resp.  $(\mathcal{D})$ ) if there is no infinite increasing (resp. decreasing) path in  $T$ . If  $(T, \ell)$  satisfies  $(\mathcal{I})$  (resp.  $(\mathcal{D})$ ), then it is clear that  $\text{Find}_k(T, \ell)$  (resp.  $\overline{\text{Find}}_k(T, \ell)$ ) is well defined for every  $k \geq 1$  by construction.

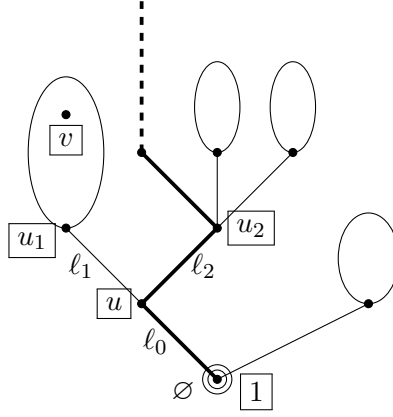
In this case, we can also define a procedure  $\text{Find}_\infty(T, \ell)$  (resp.  $\overline{\text{Find}}_\infty(T, \ell)$ ) that assigns all labels in  $\mathbb{Z}_{>0}$  (resp.  $\mathbb{Z}_{<0}$ ) to the vertices of  $T$ . Under a simple assumption, combining both procedures labels *all* vertices of an infinite tree, as explained in the following lemma, where we say that  $T$  has *one end* if for every  $r \geq 0$ ,  $T \setminus \llbracket T \rrbracket_r$  has a unique infinite connected component.

LEMMA 11. — *Let  $T$  be an infinite locally finite  $E$ -labelled tree with one end (either plane, or non-plane pointed), satisfying both  $(\mathcal{I})$  and  $(\mathcal{D})$ . Then every vertex of the tree is either assigned a label by  $\text{Find}_\infty$  or by  $\overline{\text{Find}}_\infty$ , but not by both.*

*Proof.* — (The reader may want to look at Fig. 7 to visualize the notation in this proof.) Let  $v$  be a vertex of  $T$ . Since  $T$  has one end, there exists a unique infinite injective path starting from the pointed vertex  $\emptyset$ . Denote by  $u$  the vertex of this path which is the closest to  $v$  ( $u$  could be the root vertex, or  $v$  itself). We first assume that  $v \neq u$ , i.e.  $v$  is not on the path from the root to infinity. Then  $u$  has at least two children, one of them, say  $u_1$ , being an ancestor of  $v$  (possibly  $v$  itself) and one other, say  $u_2$ , lying on the infinite path. We set  $e_1 = \{u, u_1\}$  and  $e_2 = \{u, u_2\}$ , both being edges of  $T$ . To simplify the discussion, we also assume that  $u$  is not the root of the tree, and call  $e_0$  the edge joining  $u$  to its parent. The labels of the edges  $e_0$ ,  $e_1$  and  $e_2$  are denoted by  $\ell_0$ ,  $\ell_1$  and  $\ell_2$ , respectively. Whether  $v$  is assigned a label by  $\text{Find}_\infty$  or by  $\overline{\text{Find}}_\infty$  depends on the relative order of  $\ell_0$ ,  $\ell_1$  and  $\ell_2$ , as will be explained below.

Before going into a case distinction, let us make some remarks on the procedures  $\text{Find}_\infty$  and  $\overline{\text{Find}}_\infty$ , using the notion a *fringe subtrees* (a fringe subtree  $S$  of  $T$  is a subtree of  $T$  formed by one of its vertex and *all* its descendants). The following claims are easily checked:

- when the algorithm  $\text{Find}_\infty$  (or  $\overline{\text{Find}}_\infty$ ) enters a finite fringe subtree  $S$  (i.e. `CurrentVertex` is in  $S$  at some stage of Algorithm 1), it does not leave it before having assigned a label to every vertex in  $S$ .



**Figure 7.** Notation of the proof of Lemma 11 (in bold, the infinite path).

- when the algorithm  $\text{Find}_\infty$  (or  $\overline{\text{Find}}_\infty$ ) enters an infinite fringe subtree  $S$ , it never leaves  $S$ .

Denote by respectively  $S_{u_1}$  and  $S_{u_2}$  the fringe subtrees rooted in  $u_1$  and  $u_2$ . By construction,  $S_{u_1}$  is finite and contains  $v$ , while  $S_{u_2}$  is infinite. Determining whether  $v$  is assigned a label by  $\text{Find}_\infty$  (or  $\overline{\text{Find}}_\infty$ ) therefore boils down to determining whether  $\text{Find}_\infty$  (or  $\overline{\text{Find}}_\infty$ ) enters  $S_{u_1}$  or  $S_{u_2}$  first.

From this reformulation it is now easy to see that:

- if  $\ell_0 < \ell_1 < \ell_2$  or  $\ell_1 < \ell_2 < \ell_0$  or  $\ell_2 < \ell_0 < \ell_1$ , then the vertex  $v$  is assigned a label by  $\text{Find}_\infty$  but not by  $\overline{\text{Find}}_\infty$ ;
- if  $\ell_1 < \ell_0 < \ell_2$  or  $\ell_0 < \ell_2 < \ell_1$  or  $\ell_2 < \ell_1 < \ell_0$ , then the vertex  $v$  is assigned a label by  $\overline{\text{Find}}_\infty$  but not by  $\text{Find}_\infty$ .

This proves the lemma in the case  $v \neq u$  and  $u \neq \emptyset$ . If  $v \neq u = \emptyset$ , the above conclusion holds with the convention that  $\ell_0 = \infty$ . If  $v = u \neq \emptyset$ , the same holds with the convention that  $\ell_1 = \infty$ . The only remaining case is that of  $v = \emptyset$ , but it is clear that the root is assigned a label (namely the label 1) by  $\text{Find}_\infty$  and none by  $\overline{\text{Find}}_\infty$ .  $\square$

We now prove the following continuity lemma for the Find procedure, which is crucial to obtain our limit theorem for  $\text{Tree}(\widetilde{\mathcal{F}}_n)$ .

**LEMMA 12.** — *Consider a locally finite  $E$ -labelled tree  $\mathbf{T}$  with one end (either plane, or non-plane pointed), such that both  $(\mathcal{I})$  and  $(\mathcal{D})$  are satisfied. Let  $(\mathbf{T}_n)_{n \geq 1}$  be a sequence of  $E$ -labelled trees, each with  $n$  vertices, such that  $\mathbf{T}_n$  converges to  $\mathbf{T}$  for the local topology on  $E$ -labelled trees. Then the convergence*

$$\text{Find}_{\lfloor \frac{n}{2} \rfloor} \circ \overline{\text{Find}}_{\lfloor \frac{n-1}{2} \rfloor}(\mathbf{T}_n) \xrightarrow{n \rightarrow \infty} \text{Find}_\infty \circ \overline{\text{Find}}_\infty(\mathbf{T})$$

*holds for the local topology on  $EV$ -labelled trees.*

*Proof.* — For an EV-labelled tree  $\tau$ , recall that we denote by  $[\![\tau]\!]'_h$  the V-labelled tree obtained from  $[\![\tau]\!]_h$  by forgetting the edge labels. Fix  $k \geq 1$ . We first prove that, for all  $h \geq 1$ ,

$$(3) \quad [\![\text{Find}_k(\mathbf{T}_n)]\!]_h \xrightarrow{n \rightarrow \infty} [\![\text{Find}_k(\mathbf{T})]\!]_h.$$

It is enough to establish the result for every  $h$  sufficiently large. Since  $\mathbf{T}$  satisfies  $(\mathcal{I})$ , by Lemma 11 and its proof, we may choose  $h \geq 1$  such that  $\text{Find}_k$  does not visit a vertex with height greater than  $h$  in  $\mathbf{T}$ . Therefore  $\text{Find}_k$  only depends on  $[\![\mathbf{T}]\!]_{h+1}$ . By assumption, we can take  $n$  sufficiently large so that  $[\![\mathbf{T}]\!]_{h+1}' = [\![\mathbf{T}_n]\!]_{h+1}'$  and such that the edge labels of  $[\![\mathbf{T}]\!]_{h+1}$  and of  $[\![\mathbf{T}_n]\!]_{h+1}$  are compatible (in the sense of Remark 1). For such  $n$ , the execution of  $\text{Find}_k(\mathbf{T}_n)$  is identical to that of  $\text{Find}_k(\mathbf{T})$ . As a consequence,  $[\![\text{Find}_k(\mathbf{T}_n)]\!]_h' = [\![\text{Find}_k(\mathbf{T})]\!]_h'$  for  $n$  sufficiently large, and (3) follows since the edge-labels converge by assumption. The same holds replacing  $\text{Find}_k$  by  $\overline{\text{Find}}_k$ , and thus also by the composition  $\text{Find}_k \circ \overline{\text{Find}}_k$  by using successively both statements.

We now use the fact that  $\mathbf{T}$  has one end. By Lemma 11, every vertex of  $\mathbf{T}$  is assigned a label by  $\text{Find}_\infty \circ \overline{\text{Find}}_\infty$ . In particular, for every fixed  $h \geq 1$ , there is an integer  $K_h$  such that all vertices at height at most  $h$  are assigned a label with absolute value smaller than  $K_h$ . Then it is clear that

$$[\![\text{Find}_\infty \circ \overline{\text{Find}}_\infty(\mathbf{T})]\!]_h = [\![\text{Find}_{K_h} \circ \overline{\text{Find}}_{K_h}(\mathbf{T})]\!]_h.$$

From the first part of the proof, there exists an integer  $n_h$  such that for every  $n \geq n_h$  we have

$$(4) \quad [\![\text{Find}_{K_h} \circ \overline{\text{Find}}_{K_h}(\mathbf{T}_n)]\!]_h' = [\![\text{Find}_\infty \circ \overline{\text{Find}}_\infty(\mathbf{T})]\!]_h'.$$

This implies that every vertex of height at most  $h$  in  $\mathbf{T}_n$  is assigned a label by  $\text{Find}_k \circ \overline{\text{Find}}_k$  for every  $k \geq K_h$  and  $n \geq n_h$ . Therefore, for  $n > \max(2K_h, n_h)$ , we have

$$[\![\text{Find}_{\lfloor \frac{n}{2} \rfloor} \circ \overline{\text{Find}}_{\lfloor \frac{n-1}{2} \rfloor}(\mathbf{T}_n)]\!]_h' = [\![\text{Find}_\infty \circ \overline{\text{Find}}_\infty(\mathbf{T})]\!]_h'.$$

This completes the proof since this holds for any  $h \geq 1$  and since the edge-labels converge by assumption.  $\square$

**2.5. Local convergence of the random minimal factorization tree.** — We now have all the tools to prove the local convergence as  $n \rightarrow \infty$  of the EV-labelled tree  $\text{Tree}(\widetilde{\mathcal{F}}_n)$ , which will in turn allow us to establish Theorem 1.

**THEOREM 13.** — *The convergence*

$$\frac{1}{n} \cdot \text{Tree}(\widetilde{\mathcal{F}}_n) \xrightarrow[n \rightarrow \infty]{(d)} \text{Shape}(\text{Find}_\infty \circ \overline{\text{Find}}_\infty(\mathcal{T}_\infty^u))$$

*holds in distribution in the set of all locally finite, non-plane, pointed, EV-labelled trees equipped with  $d_{\text{loc}}$ .*

*Proof of Theorem 13.* — In virtue of Skorokhod's representation theorem (see e.g. [Bil99, Theorem 6.7]), we may assume that the convergence of Proposition 7 holds almost surely, so that almost surely, for every  $h \geq 1$ ,

$$\left[ \left[ \frac{1}{n} \cdot \mathcal{E}(\mathcal{F}_n) \right] \right]_h \xrightarrow[n \rightarrow \infty]{} [\![\text{Shape}(\mathcal{T}_\infty^u)]\!]_h.$$



Since  $\mathcal{T}_\infty^u$  has a.s. one end and satisfies a.s. conditions  $(\mathcal{I})$  and  $(\mathcal{D})$ , we can apply Lemma 12. We get that, almost surely, for every  $h \geq 1$ ,

$$\left[ \left[ \text{Find}_{\lfloor \frac{n}{2} \rfloor} \circ \overline{\text{Find}}_{\lfloor \frac{n-1}{2} \rfloor} \left( \frac{1}{n} \cdot \mathcal{E}(\mathcal{F}_n) \right) \right] \right]_h \xrightarrow{n \rightarrow \infty} \left[ \left[ \text{Find}_\infty \circ \overline{\text{Find}}_\infty \left( \text{Shape}(\mathcal{T}_\infty^u) \right) \right] \right]_h.$$

From Corollary 10, the left-hand side has the same distribution as  $\left[ \left[ \frac{1}{n} \text{Tree}(\widetilde{\mathcal{F}}_n) \right] \right]_h$ . Using the commutation between  $\text{Find}$  and  $\text{Shape}$  mentioned in Section 2.3, this completes the proof of the theorem.  $\square$

We are finally in position to establish Theorem 1. Let us first define the limiting trajectories  $(X_i)_{i \in \mathbb{Z}}$ . We consider the limiting tree  $\text{Shape}(\text{Find}_\infty \circ \overline{\text{Find}}_\infty(\mathcal{T}_\infty^u))$  in Theorem 13. For a fixed  $i \in \mathbb{Z}_{>0}$ , we denote by  $V_0^{(i)} = i, \dots, V_{K_i}^{(i)} = i+1$  the labels of the successive vertices visited by  $\text{Find}_\infty$  when assigning the label  $i+1$ ; the number of such vertices  $K_i + 1$  is random, note also that some of these labels might be bigger than  $i$  or negative, so that they are not assigned when we run  $\text{Find}_{i+1}$  on  $\mathcal{T}_\infty^u$ , but are assigned later in the procedure  $\text{Find}_\infty \circ \overline{\text{Find}}_\infty$ . Finally, for  $k \in \{1, \dots, K_i\}$ , we denote  $\ell_k^{(i)}$  the label of the edge between  $V_{k-1}^{(i)}$  and  $V_k^{(i)}$ .

Then, setting  $\ell_0^{(i)} = 0$  and  $\ell_{K_i+1}^{(i)} = 1$ , we define, for  $t \in [0, 1]$ ,

$$(5) \quad X_i(t) = \sum_{k=0}^{K_i} V_k^{(i)} \mathbb{1}_{\ell_k^{(i)} \leq t < \ell_{k+1}^{(i)}}.$$

The construction is similar for  $i \leq 0$ , except that we consider the step where  $\overline{\text{Find}}_\infty$  starts from the vertex labelled  $i+1$  and assigns label  $i$ , and we denote by  $V_{K_i}^{(i)} = i+1, \dots, V_0^{(i)} = i$  the successive visited vertices (note that the order of indices is reversed), with the definition of  $\ell_k^{(i)}$  being unchanged.

*Proof of Theorem 1.* — We first introduce some notation. Fix  $i > 0$ . For  $n \geq |i|$ , run the procedure  $\text{Find}$  on  $\frac{1}{n} \cdot \text{Tree}(\widetilde{\mathcal{F}}_n)$  starting from  $i$  until it assigns label  $i+1$ . Denote by  $V_0^{(n,i)} = i, \dots, V_{K_{n,i}}^{(n,i)} = i+1$  the labels of the successively visited vertices (the number  $K_{n,i} + 1$  of such vertices is a random variable depending on  $n$ ) and by  $\ell_k^{(n,i)}$  the label of the edge between  $V_{k-1}^{(n,i)}$  and  $V_k^{(n,i)}$  (for  $k \in \{1, \dots, K_{n,i}\}$ ; since we consider  $\frac{1}{n} \cdot \text{Tree}(\widetilde{\mathcal{F}}_n)$ , these labels are in  $[0, 1]$ ). Finally set  $\ell_0^{(n,i)} = 0$  and  $\ell_{K_{n,i}+1}^{(n,i)} = 1$ . Recall from (1) the definition of trajectories  $X_i^{(n)}$  of  $i$  in  $\widetilde{\mathcal{F}}^{(n)}$ . From the proof of Lemma 8, we have

$$\forall 0 \leq t \leq 1, \quad X_i^{(n)}(\lfloor nt \rfloor) = \sum_{k=0}^{K_{n,i}} V_k^{(n,i)} \mathbb{1}_{\ell_k^{(n,i)} \leq t < \ell_{k+1}^{(n,i)}}.$$

As above, we use a similar construction for  $i < 0$  and the above relation holds as well in this case.

Fix  $A \geq 1$  and observe that the set of all indices  $\mathcal{I}_A$  of all the trajectories  $(X_i)_{i \in \mathbb{Z}}$  that enter the rectangle  $[0, 1] \times [-A, A]$  satisfies the identity

$$\mathcal{I}_A = \{i \in \mathbb{Z} : \exists 0 \leq t \leq 1 : |X_i(t)| \leq A\}.$$

We note that an element  $i$  can be in  $\mathcal{I}_A$  if and only if there exists an increasing path in  $\mathcal{T}_\infty^u$  from the vertex labelled  $i$  to some vertex label  $i'$  with  $|i'| \leq A$ . Since  $\mathcal{T}_\infty^u$  is locally finite and contains

a single path from the root to infinity, which a.s. contains infinitely many ascents and descents, for a given  $i'$ , the set of such  $i$  is a.s. finite. We conclude that  $\#\mathcal{I}_A < \infty$  almost surely.

Now, by Skorokhod's representation theorem we may assume that the convergence of Theorem 13 holds almost surely. Since  $\#\mathcal{I}_A < \infty$  almost surely, we may fix  $H > 1$  such that

- (i) for every  $i \in \mathcal{I}_A$ , all the vertices visited by the algorithm  $\text{Find}_{|i|}$  and  $\overline{\text{Find}}_{|i|}$  have height at most  $H - 1$ ;
- (ii) for every  $i'$  with  $|i'| \leq A$ , there is no decreasing path from  $i'$  going at height  $H$  or more.

We can find  $N > 1$  such that we have  $\llbracket \frac{1}{n} \cdot \text{Tree}(\widetilde{\mathcal{F}}_n) \rrbracket'_H = \llbracket \text{Shape}(\text{Find}_\infty \circ \overline{\text{Find}}_\infty(\mathcal{T}_\infty^u)) \rrbracket'_H$  for every  $n \geq N$ ; moreover, by possibly increasing  $N$ , we may assume that these  $V$ -labelled balls have compatible edge labellings. Condition (i) above implies that, for  $n \geq N$ , the procedures  $\text{Find}_{|i|}$  and  $\overline{\text{Find}}_{|i|}$  behave similarly on  $\frac{1}{n} \cdot \text{Tree}(\widetilde{\mathcal{F}}_n)$  and  $\mathcal{T}_\infty^u$ . Condition (ii) forces  $\mathcal{I}_A^{(n)}$  to be constituted of labels  $i$  of vertices such that the trajectories of  $i$  “stay” at height at most  $H$  in the tree, so that  $\mathcal{I}_A^{(n)} = \mathcal{I}_A$  for every  $n \geq N$ . As a consequence, for every  $n \geq N$  and  $i \in \mathcal{I}_A$ ,  $K_{n,i} = K_i$  and  $V_k^{(n,i)} = V_k^{(i)}$  for every  $0 \leq k \leq K_n$ . Also, for every  $i \in \mathcal{I}_A$  and  $0 \leq k \leq K_{n,i}$ ,  $\ell_k^{(n,i)} \rightarrow \ell_k^{(i)}$  as  $n \rightarrow \infty$ . The desired result follows.  $\square$

### 3. Combinatorial consequences

The goal of this section is to prove Corollary 2, using the local convergence of  $\text{Tree}(\widetilde{\mathcal{F}}_n)$ . We start in Section 3.1 by item (i), *i.e.* some results on the existence of limiting distributions for “local” statistics. In Section 3.2, we prove items (ii), (iii) and (iv) to illustrate how the explicit construction of the limit of  $\text{Tree}(\widetilde{\mathcal{F}}_n)$ , allows to compute limiting laws of such statistics. As we shall see, explicit computations quickly become quite cumbersome.

**3.1. Existence of distribution limits for local statistics.** — We first need to introduce some notation. As in the Introduction, for a factorization  $F = (\tau_1, \dots, \tau_{n-1}) \in \mathfrak{M}_n$  and an integer  $1 \leq i \leq n$ , let  $\widetilde{T}_i^F = \{1 \leq k \leq n-1 : i \in \widetilde{\tau}_k\}$  be the set of indices of all transpositions moving  $i \in \mathbb{Z}$  (transpositions are as before identified with two-element sets), and let  $\widetilde{M}_i^F = \{1 \leq k \leq n-1 : \widetilde{\tau}_1 \cdots \widetilde{\tau}_{k-1}(i) \neq \widetilde{\tau}_1 \cdots \widetilde{\tau}_k(i)\}$  be the set of all indices of transpositions that affect the trajectory of  $i \in \mathbb{Z}$ .

These sets are easily read on the associated tree  $\text{Tree}(F)$ : in particular,

- the number  $\#\widetilde{T}_i^{(n)}$  of transpositions moving  $i$  in  $F$ , is the degree of the node with label  $i$  in  $\text{Tree}(F)$ ;
- the number  $\#\widetilde{M}_i^{(n)}$  of transpositions that affect the trajectory of  $i$  in  $F$  is the distance between the vertices with labels  $i$  and  $i+1$  in  $\text{Tree}(\widetilde{\mathcal{F}}_n)$ .

As before, taking a factorization  $\mathcal{F}_n$  uniformly at random among all minimal factorizations of size  $n$ , we use the following notation for the corresponding random sets:

$$\widetilde{\mathbb{T}}_i^{(n)} := \widetilde{T}_i^{\mathcal{F}_n}, \quad \widetilde{\mathbb{M}}_i^{(n)} := \widetilde{M}_i^{\mathcal{F}_n}.$$

The local convergence of  $\text{EV}$ -labelled trees implies the (joint) convergence of the degree of the vertex  $i$  and of the distance between the vertices  $i$  and  $i+1$  (for every fixed  $i$  in  $\mathbb{Z}$ ). Therefore the convergence in distribution in Corollary 2 (i) is an immediate consequence of Theorem 13.

The statement of the marginals of the limiting distribution is proved below: in Corollary 15 for  $\#\tilde{\mathbb{T}}_i^{(n)}$  and as a consequence of symmetry considerations in Section 4 for  $\#\tilde{\mathbb{M}}_i^{(n)}$ .

More generally, many other statistics converge jointly in distribution; here is another example.

**COROLLARY 14.** — *Let  $\tilde{I}_1^{(n)}, \dots, \tilde{I}_{M_n}^{(n)}$  be integers such that the transpositions of  $\tilde{\mathcal{F}}_n$  moving 1 are, in this order  $(1, \tilde{I}_1^{(n)}), \dots, (1, \tilde{I}_{M_n}^{(n)})$ . Then  $(\tilde{I}_1^{(n)}, \tilde{I}_2^{(n)}, \dots, \tilde{I}_{M_n}^{(n)})$  converges in distribution.*

**3.2. Some explicit computations.** — In this Section, we compute explicitly some limiting distribution related to the above convergence results. This is based on the explicit construction of the limiting tree in Theorem 13. We start by proving that  $\#\tilde{\mathbb{T}}_i^{(n)}$  converges in distribution to a Poisson size-biased distribution (which is part of Corollary 2 (i)), for which only Proposition 7 is needed (that is the convergence of trees without vertex labels).

**COROLLARY 15.** — *Fix  $i \in \mathbb{Z}$ . Then, for every  $j \geq 1$ ,*

$$\mathbb{P}\left(\#\tilde{\mathbb{T}}_i^{(n)} = j\right) \xrightarrow{n \rightarrow \infty} \frac{e^{-1}}{(j-1)!}.$$

*Proof.* — By using the action by conjugation of  $(1, \dots, n)$  on minimal factorizations, we see that the distribution of the number of transpositions that act on  $i$  is independent from  $i$ . It is therefore enough to establish the result for  $i = 1$ .

By construction,  $\#\tilde{\mathbb{T}}_1^{(n)}$  is the degree of the pointed vertex in  $\mathcal{E}(\mathcal{F}_n)$ . Local convergence of unlabelled or E-labelled tree implies the convergence of the root degree, so that, from Proposition 7,  $\#\tilde{\mathbb{T}}_1^{(n)}$  converges to the root degree in  $[\text{Shape}(\mathcal{T}_\infty^u)]_1$ . By construction of  $\mathcal{T}_\infty^u$ , the degree of its root vertex is a Poisson(1) size-biased distribution, thus giving the desired result.  $\square$

The other parts of Corollary 2 need the full statement of Theorem 13 (that is with vertex labels). We first establish Corollary 2 (ii).

*Proof of Corollary 2 (ii).* — By Theorem 1, as  $n \rightarrow \infty$ , we have

$$\mathbb{P}(X_1^{(n)}(k) \geq 1 \text{ for every } 0 \leq k \leq n) \longrightarrow \mathbb{P}(X_1(t) \geq 1 \text{ for every } t \in [0, 1]).$$

By definition of the limiting trajectories  $X_i$  (Eq. (5)),  $X_1$  takes only positive values if and only if there are only positive labels between the path between 1 and 2 in the limiting tree  $\text{Shape}(\text{Find}_\infty \circ \overline{\text{Find}}_\infty(\mathcal{T}_\infty^u))$ .

We will determine when this happens by distinguishing two cases:

- *Case 1: the edge  $e$  with smallest label adjacent to the root does not belong to the spine.* We call  $v$  its extremity which is not the root. Then the algorithm  $\text{Find}_\infty$  enters first the fringe subtree rooted at  $v$ . The vertex getting label 2 will therefore be in that subtree. Moreover,  $\text{Find}_\infty$  assigns a (positive) label to every vertex in that fringe subtree. We conclude that, in this case, the path between 1 and 2 indeed contains only positive labels.
- *Case 2: the edge  $e$  with smallest label adjacent to the root belongs to the spine.* Then we claim that the path between 1 and 2 contains only positive labels if and only if the second edge-label  $\ell_2$  of the spine is smaller than the first one (call it  $\ell_1$ ). Indeed, if  $\ell_2 > \ell_1$ , then the first nonroot vertex  $v$  on the spine gets a negative label (see the proof of Lemma 11) and lies on the path between 1 and 2. Conversely, if  $\ell_2 < \ell_1$ , this vertex  $v$  gets a positive label.

Moreover, either this label is 2, or  $\text{Find}_\infty$  enters a finite fringe subtree, and will assign only positive labels, including 2, in this fringe subtree. In both cases, the path between 1 and 2 only contain positive labels.

By conditioning on the number of children of the root, we find that the probability of the first event is  $\sum_{k=2}^{\infty} \frac{e^{-1}}{(k-1)!} \frac{k-1}{k} = 1/e$ .

Let us now compute the probability that we are in the second case *and* that  $\ell_2 < \ell_1$ . Let us work conditionally given the number  $k$  of children of the root. The conditional probability that the edge with smallest label adjacent to the root is that on the spine is  $1/k$ . This smallest label has the distribution of the minimum of  $k$  independent uniform random variable in  $[0, 1]$ , that is density  $k(1-x)^{k-1}$ . Since  $\ell_2$  is uniform in  $[0, 1]$ , independently of the number of children of the root and the labels of the corresponding edges, conditionally on  $\ell_1 = x$ , the probability that  $\ell_2 < \ell_1$  is simply  $x$ . Summing up, the probability that we are in the second case *and*  $\ell_2 < \ell_1$  is

$$\sum_{k=1}^{\infty} \frac{e^{-1}}{(k-1)!} \frac{1}{k} \int_0^1 x k (1-x)^{k-1} dx = \int_0^1 x e^{-x} dx = 1 - 2/e,$$

where the first equality follows by exchanging sum and integral. The sum of the two probabilities is  $1 - 1/e$ , and this completes the proof.  $\square$

Finally, we establish Corollary 2 (iii) and (iv), whose proofs are more involved.

*Proof of Corollary 2 (iii) and (iv).* — Consider the limiting tree  $\text{Shape}(\text{Find}_\infty \circ \overline{\text{Find}}_\infty(\mathcal{T}_\infty^u))$  in Theorem 13. We denote  $u_1^\infty$ ,  $u_2^\infty$  and  $d_{1,2}^\infty$  the vertices with labels 1 and 2 in this tree, and their relative distance, respectively.

By Theorem 13 and the discussion in the beginning of Section 3.1 concerning the relation between  $\#\widetilde{\mathbb{T}}_i^{(n)}$ ,  $\#\widetilde{\mathbb{M}}_i^{(n)}$  and  $\text{Tree}(\widetilde{\mathcal{F}}_n)$ , we have, for every  $i, j \geq 1$ ,

$$\mathbb{P}\left(\#\widetilde{\mathbb{T}}_1^{(n)} = i, \#\widetilde{\mathbb{M}}_1^{(n)} = j\right) \xrightarrow{n \rightarrow \infty} \mathbb{P}\left(\deg(u_1^\infty) = i, d_{1,2}^\infty = j\right).$$

and

$$\mathbb{P}\left(\#\widetilde{\mathbb{T}}_1^{(n)} = i, \#\widetilde{\mathbb{T}}_2^{(n)} = j\right) \xrightarrow{n \rightarrow \infty} \mathbb{P}\left(\deg(u_1^\infty) = i, \deg(u_2^\infty) = j\right).$$

We also note that  $u_1^\infty$ ,  $u_2^\infty$  and  $d_{1,2}^\infty$  can be equivalently read on  $\text{Find}_2(\mathcal{T}_\infty^u)$  instead of  $\text{Shape}(\text{Find}_\infty \circ \overline{\text{Find}}_\infty(\mathcal{T}_\infty^u))$ . Therefore, in order to compute the limiting probabilities, we only need to run  $\text{Find}_2$  on  $\mathcal{T}_\infty^u$ .

For integers  $h \geq 1$  and  $0 \leq s \leq h$ , we introduce the probability  $P_h^s(k_0, \dots, k_{h-1}; k_h - 1)$  (resp.  $P_{\geq h}^s(k_0, \dots, k_{h-1}; L_h)$ ) of the following conjunction of events, when running the algorithm  $\text{Find}_2$  on  $\mathcal{T}_\infty^u$ :

- the algorithm goes through exactly  $h$  edges (resp. at least  $h$  edges), i.e.  $u_2^\infty$  is at distance exactly  $h$  (resp. at least  $h$ ) from the root;
- If, as in the description of the algorithm, we call  $j_1, j_2, \dots, j_h$  the vertices successively visited, then  $j_i$  has degree  $k_i$  for each  $i < h$  (by convention,  $j_0$  is the root of the tree);
- $j_0, j_1, \dots, j_s$  are special vertices, while  $j_{s+1}, \dots, j_h$  are not;
- for  $P_h^s(k_0, \dots, k_{h-1}; k_h - 1)$ , we additionally require that  $j_h$  has  $k_h - 1$  children (this shift makes formulas nicer).

- for  $P_{\geq h}^s(k_0, \dots, k_{h-1}; L_h)$ , we also require that the label  $\ell_h$  of the edge from between  $j_{h-1}$  and  $j_h$  lies in  $L_h$ .

Recall that the degree of the root of  $\mathcal{T}_\infty^u$  follows a size-biased Poisson distribution and that edges adjacent to the root are labeled by independent uniform variables in  $[0, 1]$ . If the root degree is  $k_0$ , the minimum among labels of edges adjacent to the root has therefore density  $k_0(1 - \ell_1)^{k_0-1}$ . Moreover,  $j_1$  is uniformly distributed among the children of the roots, and so is the special vertex of height 1, so  $j_1$  has a probability  $1/k_0$  to be a special vertex. Therefore, the probability  $P_{\geq 1}^0(k_0; L_1)$  and  $P_{\geq 1}^1(k_0, L_1)$  are respectively given by

$$P_{\geq 1}^0(k_0; L_1) = \frac{e^{-1}}{(k_0 - 1)!} (k_0 - 1) \int_{\ell_1 \in L_1} (1 - \ell_1)^{k_0-1} d\ell_1;$$

$$P_{\geq 1}^1(k_0; L_1) = \frac{e^{-1}}{(k_0 - 1)!} \int_{\ell_1 \in L_1} (1 - \ell_1)^{k_0-1} d\ell_1.$$

Let us focus, *e.g.*, on the case where  $j_1$  is a special vertex. Then its offspring distribution is again a size-biased Poisson distribution. Conditionally on  $\ell_1$  and on the fact that  $j_1$  has  $k_1$  children, the label  $\ell_2$  has density  $\mathbb{1}_{\ell_2 > \ell_1} k_1(\ell_1 + 1 - \ell_2)^{k_1-1}$ . Again, the probability that  $j_2$  is special is  $1/k_1$ . We therefore have

$$P_{\geq 2}^1(k_0, k_1; L_2) = \frac{e^{-1}}{(k_1 - 1)!} (k_1 - 1) \int_{\ell_2 \in L_2} \int_{\ell_1 < \ell_2} (\ell_1 + 1 - \ell_2)^{k_1-1} P_{\geq 1}^1(k_0; d\ell_1) d\ell_2;$$

$$P_{\geq 2}^2(k_0, k_1; L_2) = \frac{e^{-1}}{(k_1 - 1)!} \int_{\ell_2 \in L_2} \int_{\ell_1 < \ell_2} (\ell_1 + 1 - \ell_2)^{k_1-1} P_{\geq 1}^1(k_0; d\ell_1) d\ell_2.$$

Similarly, if  $j_1$  is not a special vertex, we have

$$P_{\geq 2}^0(k_0, k_1; L_2) = \frac{e^{-1}}{k_1!} k_1 \int_{\ell_2 \in L_2} \int_{\ell_1 < \ell_2} (\ell_1 + 1 - \ell_2)^{k_1-1} P_{\geq 1}^0(k_0; d\ell_1) d\ell_2.$$

Continuing the reasoning, an easy induction proves that we have

$$P_{\geq h}^s(k_0, k_1, \dots, k_{h-1}; L_h) = \left( \prod_{i=0}^{h-1} \frac{e^{-1}}{(k_i - 1)!} \right) (k_s - 1)^*$$

$$\cdot \int_{\ell_h \in L_h} \int_{\ell_1 < \dots < \ell_h} (1 - \ell_1)^{k_0-1} (\ell_1 + 1 - \ell_2)^{k_1-1} \dots (\ell_{h-1} + 1 - \ell_h)^{k_{h-1}-1} d\ell_1 \dots d\ell_h,$$

where  $(k_s - 1)^* = 1$  if  $s = h$  and  $(k_s - 1)^* = k_s - 1$  otherwise.

Conditionally on the fact that the algorithm Find<sub>2</sub> goes through at least  $h$  edges, and conditionally on the label  $\ell_h$  of the last visited edge, Find<sub>2</sub> will stop at height  $h$  if all labels of edges adjacent to  $j_h$  are smaller than  $\ell_h$ , which happens with probability  $\ell_h^{k_h-1}$  (where  $k_h - 1$  is the number of children of  $j_h$ ). Again, the distribution of  $k_h$  depends on whether  $j_h$  is special or not, so that we should consider two cases separately: for  $s < h$

$$(6) \quad P_h^s(k_0, \dots, k_{h-1}; k_h - 1) = \frac{e^{-1}}{(k_h - 1)!} \int_{[0,1]} \ell_h^{k_h-1} P_{\geq h}^s(k_0, k_1, \dots, k_{h-1}; d\ell_h)$$

$$= \left( \prod_{i=0}^h \frac{e^{-1}}{(k_i - 1)!} \right) (k_s - 1) \int_{\ell_1 < \dots < \ell_h} \prod_{i=0}^h (\ell_i + 1 - \ell_{i+1})^{k_i-1} d\ell_1 \dots d\ell_h,$$

with the convention  $\ell_0 = 0$  and  $\ell_{h+1} = 1$ . Similarly, for  $s = h$ , we have

$$\begin{aligned} P_h^h(k_0, \dots, k_{h-1}; k_h - 1) &= \frac{e^{-1}}{(k_h - 2)!} \int_{[0,1]} \ell_h^{k_h-1} P_{\geq h}^s(k_0, k_1, \dots, k_{h-1}; d\ell_h) \\ &= \left( \prod_{i=0}^{h-1} \frac{e^{-1}}{(k_i - 1)!} \right) \frac{e^{-1}}{(k_h - 2)!} \int_{\ell_1 < \dots < \ell_h} \prod_{i=0}^h (\ell_i + 1 - \ell_{i+1})^{k_i-1} d\ell_1 \dots d\ell_h. \end{aligned}$$

Note that this coincides with (6) for  $s = h$ , so that (6) is actually valid for every  $s$  in  $\{0, 1, \dots, h\}$ .

We now come back to the specific probabilities we want to evaluate. For item (i), we fix  $h = j$  and  $k_0 = i$  and sum over  $s$  and over  $k_1, \dots, k_h$ :

$$\mathbb{P}(\deg(u_1^\infty) = k_0, d_{1,2}^\infty = h) = \sum_{s=0}^h \sum_{k_1, \dots, k_h \geq 1} P_h^s(k_0, k_1, \dots, k_{h-1}; k_h - 1)$$

For  $s = 0$ , noting the sum over each  $k_i$  ( $i \geq 1$ ) is the series expansion of an exponential, we have

$$\sum_{k_1, \dots, k_h \geq 1} P_h^0(k_0, k_1, \dots, k_{h-1}; k_h - 1) = \frac{e^{-2}(k_0 - 1)}{(k_0 - 1)!} \int_{\ell_1 < \dots < \ell_h} (1 - \ell_1)^{k_0-1} e^{\ell_1} d\ell_1 \dots d\ell_h.$$

A similar computation for  $s \geq 1$  gives

$$\sum_{k_1, \dots, k_h \geq 1} P_h^s(k_0, k_1, \dots, k_{h-1}; k_h - 1) = \frac{e^{-2}}{(k_0 - 1)!} \int_{\ell_1 < \dots < \ell_h} (1 - \ell_1)^{k_0-1} e^{\ell_1} (\ell_s + 1 - \ell_{s+1}) d\ell_1 \dots d\ell_h.$$

Summing over  $s$  in  $\{0, 1, \dots, h\}$ , we find that

$$\mathbb{P}(\deg(u_1^\infty) = k_0, d_{1,2}^\infty = h) = \frac{e^{-2}}{(k_0 - 1)!} \int_{\ell_1 < \dots < \ell_h} (1 - \ell_1)^{k_0-1} e^{\ell_1} (\ell_1 + h + k_0 - 2) d\ell_1 \dots d\ell_h.$$

The integrand only depends on  $\ell_1$ . Besides,  $\int_{\ell_1 < \dots < \ell_h} d\ell_2 \dots d\ell_h = \frac{(1-\ell_1)^{h-1}}{(h-1)!}$  for any fixed  $\ell_1$ . Thus, we can rewrite the above integral as

$$\begin{aligned} \mathbb{P}(\deg(u_1^\infty) = k_0, d_{1,2}^\infty = h) &= \frac{e^{-2}}{(k_0 - 1)!(h - 1)!} \int_0^1 (1 - \ell_1)^{k_0+h-2} e^{\ell_1} (\ell_1 + h + k_0 - 2) d\ell_1 \\ &= \frac{e^{-2}}{(k_0 - 1)!(h - 1)!}, \end{aligned}$$

where the computation of the last integral is an easy calculus exercise. This shows Corollary 2 (iii).

To establish (iv), we fix  $k_0 = i$  and  $k_h = j$  (for the non-root vertex  $j_h = u_2^\infty$ , having  $k_h - 1$  children means having degree  $k_h$ ) and sum over  $h, s$  and  $k_1, \dots, k_{h-1}$ . Namely, we have

$$\mathbb{P}(\deg(u_1^\infty) = i, \deg(u_2^\infty) = j) = \sum_{h \geq 1} S_h,$$

where

$$S_h = \sum_{s=0}^h \sum_{k_1, \dots, k_{h-1} \geq 1} P_h^s(i, k_1, \dots, k_{h-1}; j - 1).$$

The case  $h = 1$  is somewhat special since the last sum has only one summand corresponding to the empty list. In this case, we may have  $s = 0$  and  $s = 1$  giving

$$(7) \quad S_1 = P_1^0(i; j - 1) + P_1^1(i; j - 1) = \frac{e^{-2}(i + j - 2)}{(i - 1)!(j - 1)!} \int_0^1 (1 - \ell_1)^{i-1} \ell_1^{j-1} d\ell_1 = \frac{e^{-2}(i + j - 2)}{(i + j - 1)!}$$

Consider now the summands corresponding to  $h \geq 2$ . A similar computation as above, starting from (6) and separating the cases  $s = 0$  and  $s = h$ , yields:

$$S_h = \frac{e^{-2}}{(i-1)!(j-1)!} \cdot \int_{\ell_1 < \dots < \ell_h} (1-\ell_1)^{i-1} \ell_h^{j-1} e^{\ell_1 - \ell_h} (\ell_1 - \ell_h + i + j + h - 3) d\ell_1 \cdots d\ell_h.$$

The integrand only involves  $\ell_1$  and  $\ell_h$ . Using the equality  $\int_{\ell_1 < \dots < \ell_h} d\ell_2 \cdots d\ell_{h-1} = \frac{(\ell_h - \ell_1)^{h-2}}{(h-2)!}$  and using variables  $x := \ell_1$  and  $y := \ell_h$ , we get

$$S_h = \frac{e^{-2}}{(i-1)!(j-1)!} \cdot \int_{x < y} \frac{(y-x)^{h-2}}{(h-2)!} (1-x)^{i-1} y^{j-1} e^{x-y} (x-y+i+j+h-3) dx dy.$$

Summing this over  $h \geq 2$  and exchanging sum and integral (the terms are nonnegative), we obtain

$$\sum_{h \geq 2} S_h = \frac{e^{-2}(i+j-1)}{(i-1)!(j-1)!} \cdot \int_{x < y} (1-x)^{i-1} y^{j-1} dx dy = \frac{e^{-2}(i+j-1)}{(i-1)!(j-1)!} \left( \frac{1}{ij} - \frac{(i-1)!(j-1)!}{(i+j)!} \right).$$

Adding  $S_1$ , which was computed in (7), we get

$$\mathbb{P}(\deg(u_1^\infty) = i, \deg(u_2^\infty) = j) = \sum_{h \geq 1} S_h = e^{-2} \left( \frac{i+j-2}{(i+j-1)!} + \frac{i+j-1}{i!j!} - \frac{i+j-1}{(i+j)!} \right).$$

This completes the proof of Corollary 2 (iv).  $\square$

#### 4. A new bijection and duality

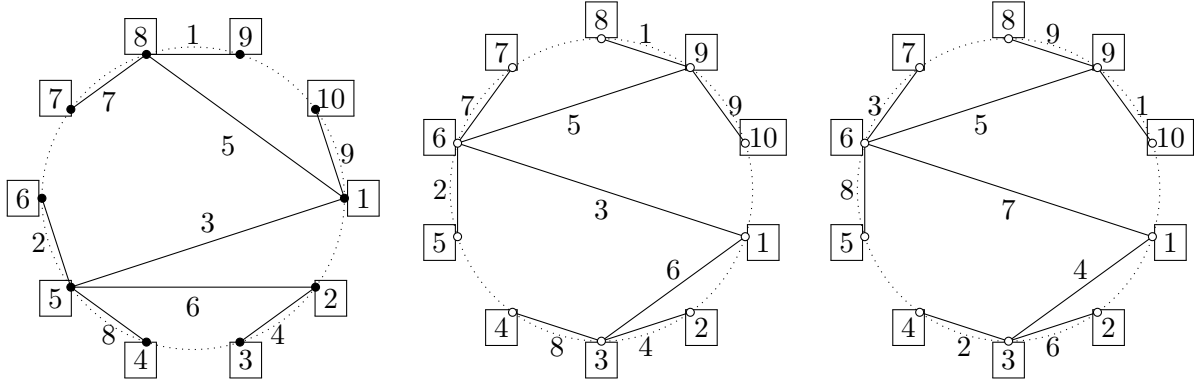
In this section, we construct a new bijection  $\mathcal{B}$  for minimal factorizations with the following property: if  $F' = \mathcal{B}(F)$  with  $F \in \mathfrak{M}_n$ , then, for every  $1 \leq i \leq n$ , the number of transpositions that affect the trajectory of  $i$  in  $F$  is equal to the number of transpositions in  $F'$  containing  $i$ . Combinatorial consequences of this bijection are then discussed.

Our bijection is based on Goulden–Yong’s duality bijection [GY02], which we now present (see Fig. 8 for an example; note that here we multiply from left to right, while in [GY02] the multiplication is done from right to left). Let  $F \in \mathfrak{M}_n$  be a minimal factorization. Recall from Section 2.2 its associated pointed non-plane EV-labelled tree  $\mathcal{E}(F)$ .

First draw  $\mathcal{E}(F)$  inside the complex unit disk  $\mathbb{D}$  by identifying vertex  $j$  (for  $1 \leq j \leq n$ ) with the complex number  $e^{-\frac{2i\pi(j-1)}{n}}$ . A *face* is a connected component of  $\mathbb{D} \setminus \mathcal{E}(F)$ . Then, by [GY02], edges do not cross, and every face contains exactly one arc of  $\mathbb{S}$  of the form  $\widehat{j, j+1}$  for a certain  $1 \leq j \leq n$  (with the convention  $n+1 = 1$  and by identifying  $e^{-\frac{2i\pi(j-1)}{n}}$  with  $j$ , see Fig. 8). Conversely, every arc  $\widehat{j, j+1}$  with  $1 \leq j \leq n$  is contained in a face.

Some combinatorial information is easily read on  $\mathcal{E}(F)$ . Indeed, the number  $\#T_i^F$  of transpositions in  $F$  containing  $i$  is simply the degree of the vertex labelled  $i$  in  $\mathcal{E}(F)$ . Similarly, the number  $\#M_j^F$  of transpositions that affect the trajectory of  $j$  in  $F$  is the number of edges lying around the face of  $\mathcal{E}(F)$  containing the arc  $\widehat{j, j+1}$ . In particular, by applying a horizontal symmetry to  $\mathcal{E}(F)$  we obtain the following identity on generating functions:

$$(8) \quad \sum_{F \in \mathfrak{M}_n} x^{\#T_1^F} y^{\#M_j^F} = \sum_{F \in \mathfrak{M}_n} x^{\#T_1^F} y^{\#M_{n+1-j}^F}.$$



**Figure 8.** Several objects associated with the same minimal factorisation  $F = ((8, 9) (5, 6) (1, 5) (2, 3) (1, 8) (2, 5) (7, 8) (4, 5) (1, 10))$  of  $(1, 2, \dots, 10)$ . From left to right: the EV-labelled tree  $\mathcal{E}(F)$ , its dual tree  $\mathcal{E}^\dagger(F)$  and its symmetrized version  $\bar{\mathcal{E}}^\dagger(F)$  which codes a minimal factorization  $\mathcal{B}(F)$ .

Then, still following Goulden & Yong [GY02], define the “dual” EV-labelled tree  $\mathcal{E}^\dagger(F)$  as follows: the vertices are  $e^{-\frac{2i\pi(j-1)}{n} - \frac{i\pi}{2n}}$  (which is given label  $j$ ) for  $1 \leq j \leq n$ , and two vertices  $e^{-\frac{2i\pi(j-1)}{n} - \frac{i\pi}{2n}}$  and  $e^{-\frac{2i\pi(k-1)}{n} - \frac{i\pi}{2n}}$  are connected if the two faces containing  $\widehat{j, j+1}$  and  $\widehat{k, k+1}$  are adjacent in  $\mathcal{E}(F)$ . Moreover, the corresponding edge gets the label of the edge of  $\mathcal{E}(F)$  separating these two faces.

Finally, define  $\bar{\mathcal{E}}^\dagger(F)$  by changing the edge labels of  $\mathcal{E}^\dagger(F)$  by “symmetrization”: exchange labels  $i$  and  $n - i$  for every  $1 \leq i \leq n/2$ . It turns out that  $\bar{\mathcal{E}}^\dagger(F)$  codes a minimal factorization, which allows to define  $\mathcal{B}(F)$ :

LEMMA 16. — *For every minimal factorization  $F$ , there exists a unique minimal factorisation  $\mathcal{B}(F)$  such that  $\bar{\mathcal{E}}^\dagger(F) = \mathcal{E}(\mathcal{B}(F))$ .*

*Proof.* — We recall [GY02, Theorem 2.2] (adapted to the fact that here we multiply from left to right, while in [GY02] the multiplication is done from right to left): in  $\mathcal{E}(F)$ , when turning along a face in clockwise order starting from the arc of  $\mathbb{S}$  in its boundary, the edge labels are decreasing. Therefore, in  $\bar{\mathcal{E}}^\dagger(F)$ , the edge labels are also decreasing in clockwise order in every face (starting each time from the circle arc contained in the face boundary). By [GY02, Lemma 2.5], the desired result follows.  $\square$

The fact that  $\mathcal{B}$  is a bijection follows by definition of  $\mathcal{B}$ , since  $\mathcal{E}$  is a bijection.

THEOREM 17. — *Let  $F$  be a minimal factorization. For every  $1 \leq i \leq n$ :*

- (i) *the number of transpositions in  $F$  that affect the trajectory of  $i$  is equal to the number of transpositions in  $\mathcal{B}(F)$  containing  $i$ , i.e.  $\#M_i^F = \#T_i^{\mathcal{B}(F)}$ ;*
- (ii) *the number of transpositions in  $F$  containing  $i$  is equal to the number of transpositions in  $\mathcal{B}(F)$  that affect the trajectory of  $i - 1$ , i.e.  $\#T_i^F = \#M_{i-1}^{\mathcal{B}(F)}$ . (We use the convention  $i - 1 = n$  for  $i = 1$ .)*



When  $\mathcal{F}_n$  denotes a minimal factorizations of size  $n$  chosen uniformly at random, as in the Introduction we use the following notation for  $1 \leq i \leq n$  :  $\mathbb{T}_i^{(n)} := T_i^{\mathcal{F}_n}$  and  $\mathbb{M}_i^{(n)} := M_i^{\mathcal{F}_n}$ . Theorem 17 implies that  $\#\mathbb{T}_i^{(n)}$  and  $\#\mathbb{M}_i^{(n)}$  have the same law. Combining with Corollary 15, we conclude that the limiting distribution of  $\#\mathbb{M}_i^{(n)}$  is a size biased Poisson law, as claimed in Corollary 2.

*Proof.* — Fix  $1 \leq i \leq n$ . For (i), observe that, by construction, the following numbers are all equal:

- the number of transpositions that affects the trajectory of  $i$  in  $F$ ;
- the number of edges adjacent to the face of  $\mathcal{E}(F)$  containing the arc  $\widehat{i, i+1}$ ;
- the degree of  $i$  in  $\overline{\mathcal{E}}^\dagger(F) = \mathcal{E}(\mathcal{B}(F))$ ;
- the number of transpositions in  $\mathcal{B}(F)$  containing  $i$ .

The proof of the second assertion is similar and is left to the reader. □

We can now prove the distributional identity stated in Theorem 3.

*Proof of Theorem 3.* — The first equality in distribution is a probabilistic translation of the properties of the bijection  $\mathcal{B}$  (Theorem 17). The second is the result of applying to  $\mathcal{E}(F)$  an axial symmetry around the diameter containing  $e^{\frac{-i\pi(k-1)}{n}}$ . □

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