PI Measure: A Set-Theoretic View of Partial Information

2020 Spring CS258 Group Project Presentation

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Background

I-Measure

A theory which establishes a one-to-one correspondence between Shannon's information measure and set theory in full generality.

$$H/I \leftrightarrow \mu^* \quad , \quad \leftrightarrow \cup$$

 $; \leftrightarrow \cap \quad | \quad \leftrightarrow -$ (1)

Non-negative Decomposition of Multivariate Information

A measure of *redundancy information* that a set of sources provides about a given variable.

Shannon information can be decomposed into atoms of non-negative partial information.

Our Goal: Using the I-Measure methodology to study the redundancy measure, forming a set-theoretic view of partial information.



Definition

Definition 1 (Field)

The field \mathcal{F}_n generated by sets $\tilde{X}_1, \tilde{X}_2, \cdots, \tilde{X}_n$ is the collection of sets which can be obtained by any sequence of usual set operations (union, intersection, complement, and difference) on $\tilde{X}_1, \tilde{X}_2, \cdots, \tilde{X}_n$

Definition 2 (Atom)

The atoms of \mathcal{F}_n are sets of the form $\bigcap_{i=1}^n Y_i$, where Y_i is either \tilde{X}_i or \tilde{X}_i^c , the complement of \tilde{X}_i

Definition 3 (Signed Measure)

A real function μ defined on \mathcal{F}_n is called a signed measure if it is set-additive, i.e., for disjoint A and B in \mathcal{F}_n

$$\mu(A \cup B) = \mu(A) + \mu(B) \tag{2}$$

Measure Theorem

Use X_G to denote $(X_i, i \in G)$ and \tilde{X}_G to denote $\bigcup_{i \in G} \tilde{X}_i$ for any nonempty subset G of $\mathcal{N}_n = \{1, 2, \cdots, n\}$

Theorem 4 (Completeness)

Let

$$\mathcal{B} = \left\{ \tilde{X}_G : G \text{ is a nonempty subset of } \mathcal{N}_n \right\}$$
 (3)

Then a signed measure μ on \mathcal{F}_n is completely specified by $\{\mu(B), B \in \mathcal{B}\}$ which can be any set of real numbers.

With the terminologies of measure theory above, we can construct a correspondence between set theory and Shannon Information





Correspondence

Build Correspondence

We now construct the I-Measure μ^* on \mathcal{F}_n using Theorem 4 by defining

$$\mu^* \left(\tilde{X}_G \right) = H(X_G) \tag{4}$$

for all nonempty subsets G of \mathcal{N}_n .

The following must hold for all (not necessarily disjoint) subsets G, G', G'' of \mathcal{N}_n :

$$\mu^* \left(\tilde{X}_G \cap \tilde{X}_{G'} - \tilde{X}_{G''} \right) = I(X_G; X_{G'} | X_{G''})$$
 (5)



Applications

Information Diagrams

With the correspondence above, it is reasonable to use an *information diagram*, which is a variation of Venn diagram, to represent the relationship between Shannon's information measures.

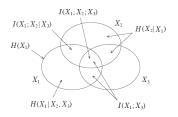


Figure: Generic Information Diagram for X_1, X_2, X_3



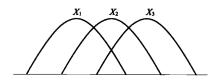
Applications

Information Diagrams for Zero Measure

In Markov Chain $X_1 o X_2 o X_3$, since

$$\mu^* \left(\tilde{X}_1 \cap \tilde{X}_2^c \cap \tilde{X}_3 \right) = I(X_1; X_3 | X_2) = 0 \tag{6}$$

We don't have to plot that region out in the information diagram, making the diagram simpler and more useful when reasoning about Shannon information measures.



In a word, with the set-theoretic view of information diagram, we can discover many useful properties.



Redundancy Measure

Let A_1, A_2, \ldots, A_k be nonempty and potentially overlapping subsets of R, which we call sources. We introduce I_{\min} to measure the information that all sources provide about S.

Definition 5 (Redundancy Measure)

$$I_{\min}\left(S;\left\{\mathbf{A}_{1},\mathbf{A}_{2},\ldots,\mathbf{A}_{k}\right\}\right)=\sum_{s}p(s)\min_{\mathbf{A}_{i}}I\left(S=s;\mathbf{A}_{i}\right)$$
 (7)

where the domain of I_{\min} $\mathcal{A}(\mathbf{R}) = \{\alpha \in \mathcal{P}_1\left(\mathcal{P}_1(\mathbf{R})\right) : \forall \mathbf{A}_i, \mathbf{A}_j \in \alpha, \mathbf{A}_i \notin \mathbf{A}_j\}$

Properties of Redundancy Measure

- I_{min} is non-negative
- I_{min} is less than or equal to $I(S; \mathbf{A}_i)$ for all \mathbf{A}_i 's
- For a given source A, the amount of information redundant with **A** is maximal for $I_{\min}(S; \{\mathbf{A}\}) = I(S; \mathbf{A})$.

Relations between Redundancy Measures

Note that I_{min} can measure redundancy between collections of sources like $\{\{R_1\}, \{R_2, R_3\}\}$ denoted as $\{1\}\{23\}$. Certain relations exists between measures.

Redundancy Order

We can define a partial order over the elements of A(R) such that one is considered to precede another if and only if the latter provides any redundant information that the former provides. Formally,

$$\forall \alpha, \beta \in \mathcal{A}(\mathbf{R}),$$

 $\alpha \preccurlyeq \beta \Leftrightarrow \forall \mathbf{B} \in \beta, \exists \mathbf{A} \in \alpha, \mathbf{A} \subseteq \mathbf{B}$
(8)

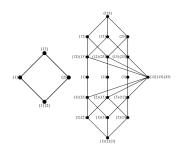


Figure: Redundancy lattice for (A) 3 and (B) 4 variables.

Partial Information Decomposition

We would like to decompose redundancy measure into non-intersecting fragments. For a collection of sources $\alpha \in \mathcal{A}(\mathbf{R})$, the PI-function denoted Π_{R} , is defined implicitly by

$$I_{\min}(S; \alpha) = \sum_{\beta \leq \alpha} \Pi_{\mathcal{R}}(S; \beta)$$
 (9)

Properties of Redundancy Measure

- $\bullet \ II_R$ is non-negative
- ullet IIR can be calculated recursively.
- The relationship between these measures can be shown in a partial information (PI) diagram, which we will discuss later.

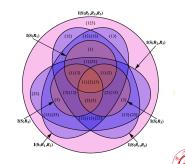


Figure: PI diagram for 4 variables.

Formulation

The sets used to generate fields, and their correspondence to redundancy measures are defined as follows.

Basic Sets Correspondence

Denote the source variables domain as

$$\Gamma = \{ \beta \in \mathcal{P}_1(\mathbf{R}) : \forall \mathbf{A}_i, \mathbf{A}_j \in \beta, \mathbf{A}_i \nsubseteq \mathbf{A}_j \}$$
 (10)

For every random variable collection in the domain Γ , its self-redundancy is defined as $I_{\min}(S; \{A_i, \ldots, A_j\})$. We use $\{\tilde{A}_i, \ldots, \tilde{A}_j\}$, or abbreviation $\tilde{A}_{i\ldots j}$ to represent the basic set its self redundancy corresponds to.



Formulation

To build fields, define the correspondence rules for operations on basic sets.

Field Correspondence

• For every basic sets $\tilde{A}_{i...j}$,

$$\mu^* \left(\tilde{A}_{i...j} \right) = I_{\min} \left(S; \left\{ A_i, \dots, A_j \right\} \right). \tag{11}$$

- For any two sets α in the field, $\mu^*(\alpha^c) = \sum_{\alpha \prec \gamma} \Pi_{\mathbf{R}}(S; \gamma)$
- For any two sets α and β in the field, $\mu^*(\alpha \cap \beta) = \sum_{\gamma \preceq \alpha \text{ and } \gamma \preceq \beta} \Pi_{\mathbf{R}}(S; \gamma)$
- For any two sets α and β , $\mu^*(\alpha \cup \beta) = \sum_{\gamma \preceq \alpha \text{ or } \gamma \preceq \beta} \Pi_{\mathbf{R}}(S; \gamma)$
- For any two sets α and β , $\mu^*(\alpha \beta) = \sum_{\gamma \leq \alpha \text{ and } \beta \prec \gamma} \Pi_{\mathbf{R}}(S; \gamma)$

This definition is based on the partial order given in the redundancy lattice. However, such definition is not very useful in practice.



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We try to eliminate redundant fields in order to avoid the exponential increase of fields.

Example of Two Source Variables

To fix ideas, we take the multivariate information of three variables $I(S; \mathbf{A}_1, \mathbf{A}_2)$ as example, note that $\left\{ \tilde{A}_1 \right\} \subseteq \left\{ \tilde{A}_1, \tilde{A}_2 \right\}$. It follows that

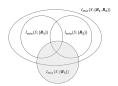
$$\mu^* \left(\left\{ \tilde{A}_1 \right\} \cap \left\{ \tilde{A}_1, \tilde{A}_2 \right\}^C \right) = \mu^* \left(\left\{ \tilde{A}_1 \right\} \right) - \mu^* \left(\left\{ \tilde{A}_1 \right\} \cap \left\{ \tilde{A}_1, \tilde{A}_2 \right\} \right)$$

$$= I_{\min} \left(S; \left\{ \mathbf{A}_1 \right\} \right) - I_{\min} \left(S; \left\{ \mathbf{A}_1 \right\}, \left\{ \mathbf{A}_1, \mathbf{A}_2 \right\} \right) = 0$$
(12)

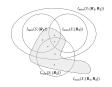
Similarly, with the observation that $\left\{\tilde{A}_2\right\}\subseteq \left\{\tilde{A}_1,\tilde{A}_2\right\}$, we have $\mu^*\left(\left\{\tilde{A}_2\right\}\cap \left\{\tilde{A}_1,\tilde{A}_2\right\}^C\right)=0$. To make our diagram more legible, we can erase these regions out.

Theorem 6 (Field Elimination)

For any two basic sets $\tilde{A}_{i...j}$, $\tilde{A}_{i'...j'} \in \Gamma$, if $\tilde{A}_{i...j} \subseteq \tilde{A}_{i'...j'}$, then the measure of any set intersection with $\tilde{A}_{i...j} \cap \tilde{A}_{i'...i'}^{C}$ is zero.



(a) \tilde{A}_3 produces no extra redundant fields



(b) Adding Basic Set \tilde{A}_{12} produces extra redundant fields as has been starred



(c) After Elimination

Figure: Field Elimination Example for 3 Source Variables



We try to find alternative, useful expressions of fields by means of set theory.

Observation on Atoms

- The measure of atoms in \mathcal{F}_n is a perfect matching to the partial information $\Pi_{\mathbf{R}}$.
- \bullet All redundancy measures can be written as sum of non-negative $\Pi_{R}s.$
- All atoms can be written as sets of $\cap_{i \in \mathcal{N}_n} \tilde{Y}_i(\tilde{Y}_i = \tilde{A}_i \text{ or } \tilde{A}_i^c)$, where the index set $\mathcal{N}_n = \mathcal{P}_1(\{1, 2, \dots, n\})$

In order to completely write measure expressions, we only need to study two operations (complement and intersection) on basic sets $Y_{i \in \mathcal{N}_n}$.



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$$\mu^* \left(\tilde{A}_{N_1} \cap \ldots \cap \tilde{A}_{N_k} \cap \tilde{A}_{M_1}^C \cap \ldots \cap \tilde{A}_{M_j}^C \right) \tag{13}$$

$$= \mu^* \left(\tilde{A}_{N_1} \cap \ldots \cap \tilde{A}_{N_k} \right) - \sum_{i=1}^k \mu^* \left(\tilde{A}_{N_1} \cap \ldots \cap \tilde{A}_{N_k} \cap \tilde{A}_{M_1} \cap \ldots \cap \tilde{A}_{M_i} \right)$$

$$(14)$$

Note that in Equation 14, all the sets being measured are intersections of basic sets. By applying Equation 11, we can write them out in the form of redundancy measure I_{\min} .

Conclusion

- The measure of any atom can be written as a ± 1 combination of some redundancy measures I_{\min} .
- The measure of any field can be written as a linear combination of redundancy measures I_{min} .

Application

An application of our simplification is to calculate partial information in two ways.

Method 1

On the one hand, we can using the definition of $\Pi_{\rm R}$ to calculate, applying Eq.15.

$$\Pi_{\mathbf{R}}(S;\alpha) = I_{\min}(S;\alpha) - \sum_{s} p(s) \max_{\beta \in \alpha^{-}} \min_{\mathbf{B} \in \beta} I(S=s;\mathbf{B})$$
 (15)

Method 2

On the other hand, we can calculate $\Pi_{\rm R}$ by converting it to the linear combination of $I_{\rm min}$.





Example

Consider three equiprobable cases of S as follows:

$$S = 0 : (R_1 = 0, R_2 = 0, R_3 = 0)$$

$$S = 1 : (R_1 = 0, R_2 = 1, R_3 = 1)$$

$$S = 2 : (R_1 = 1, R_2 = 0, R_3 = 0)$$

We want to calculate the value of $\Pi_R(S; \{2\}, \{13\})$.



Method 1

Firstly, we can applying Eq.15 to solve the problem:

$$\Pi_{\mathbf{R}}(S; \{2\}, \{13\}) = I_{\min}(S; \{2\}, \{13\}) - \sum_{s} p(s) \max_{\beta \in \alpha^{-}} \min_{\mathbf{B} \in \beta} I(S = s; \mathbf{B})$$
 (16)

where $\beta \in \{\{\{1\}\{2\}\}, \{\{2\}\{3\}\}\}$. Then we just need to expand the formula to calculate. We have

$$I_{\min}(S; \{2\}, \{13\}) = \log 3 - \frac{2}{3} \log 2$$
 (17)

$$\sum_{s} p(s) \max_{\beta \in \alpha^{-}} \min_{\mathbf{B} \in \beta} I(S = s; \mathbf{B}) = \log 3 - \frac{2}{3} \log 2$$
 (18)

So $\Pi_{\mathbb{R}}(S; \{2\}, \{13\}) = 0$.

Method 2

based on the observations in FIG.3, we can express the field using Eq.19:

$$\mu^* \left(\tilde{A}_{i...j} \cap \ldots \cap \tilde{A}_{m...n} \right) = I_{\min} \left(S; \left\{ \mathbf{A}_i, \ldots, \mathbf{A}_j \right\}, \ldots, \left\{ \mathbf{A}_m, \ldots, \mathbf{A}_n \right\} \right)$$
 (19)

$$\mu^{*}\left(\left\{\tilde{R}_{1}\right\}^{C}\cap\left\{\tilde{R}_{2}\right\}\cap\left\{\tilde{R}_{3}\right\}^{C}\cap\left\{\tilde{R}_{12}\right\}\cap\left\{\tilde{R}_{13}\right\}\cap\left\{\tilde{R}_{23}\right\}\cap\left\{\tilde{R}_{123}\right\}\right)$$

$$=I_{\min}\left(S;\left\{2\right\},\left\{1,2\right\},\left\{1,3\right\},\left\{2,3\right\},\left\{1,2,3\right\}\right)$$

$$-I_{\min}\left(S;\left\{1\right\},\left\{2\right\},\left\{1,2\right\},\left\{1,3\right\},\left\{2,3\right\},\left\{1,2,3\right\}\right)$$

$$-I_{\min}\left(S;\left\{2\right\},\left\{3\right\},\left\{1,2\right\},\left\{1,3\right\},\left\{2,3\right\},\left\{1,2,3\right\}\right)$$

$$+I_{\min}\left(S;\left\{1\right\},\left\{2\right\},\left\{3\right\},\left\{1,2\right\},\left\{1,3\right\},\left\{2,3\right\},\left\{1,2,3\right\}\right)$$

$$=I_{\min}\left(S;\left\{2\right\},\left\{1,3\right\}\right)-I_{\min}\left(S;\left\{1\right\},\left\{2\right\}\right)-I_{\min}\left(S;\left\{2\right\},\left\{3\right\}\right)$$

$$+I_{\min}\left(S;\left\{1\right\},\left\{2\right\},\left\{3\right\}\right)$$

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Method 2

We can also list all I_{\min} of the sets which have the partial order relation with $\{\{2\},\{13\}\}$ to express $\Pi_{\mathbf{R}}$ using the linear combination of I_{\min} .

$$I_{\min}(S; \{2\}, \{13\}) = \Pi_{\mathbb{R}}(S; \{2\}, \{3\}) + \Pi_{\mathbb{R}}(S; \{1\}, \{2\}) + \Pi_{\mathbb{R}}(S; \{2\}, \{13\}) + I_{\min}(S; \{1\}, \{2\}, \{3\})$$
(20)

$$I_{\min}(S; \{2\}, \{3\}) = \Pi_{\mathbb{R}}(S; \{2\}, \{3\}) + I_{\min}(S; \{1\}, \{2\}, \{3\})$$
 (21)

$$I_{\min}(S; \{1\}, \{2\}) = \Pi_{\mathbb{R}}(S; \{1\}, \{2\}) + I_{\min}(S; \{1\}, \{2\}, \{3\})$$
 (22)

Combine Eq.20,21,22, then we have the save result as before.

Calculating by the linear combination of I_{min} , we can also have the same outcome of $\Pi_R(S; \{2\}, \{13\})$.



Summary

Previous Work

- Correspondence between Shannon information and set theory
- Redundancy measure and partial information decomposition methods for multivariate information

Achieved Goals

- Formalize the correspondence between redundancy measure and set theory
- Prove the completeness of the correspondence relation
- Simplify the correspondence expressions
- Propose a top-down strategy to calculate redundancy measure



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