

Tony

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1 Grammar

typ, A, B, C, Char, Bool, ?	::= 	X Int Top $A o B$ $\mu X.A$	types type variable int top type function type recursive type
exp, e, v, u, v _f	::= 	x \uparrow i $\lambda x : A.e$ fix $x : A.e$ $e_1 e_2$ $cast [c]e$	expressions variable top lit abstraction with argument annotation fixpoint applications
mode, m	::= 	\bigoplus_{Θ}	modes positive negative
castop, c	::= 	$\begin{array}{l}id\\ \downarrow_A\\ \uparrow_A\\ c_1 \to c_2 \end{array}$	cast operators id operator unfold operator castdn arrow operator
ctx, Γ	::= 	Γ, $x : A$	term context
$tctx, \ \Delta$::= 	Δ, Χ	type context

2 Subtyping

 $\vdash \Delta$

(Well Formed Type Environment)

$$\begin{array}{ll} \text{WFTyE-cons} \\ \underline{\vdash} \\ \end{array} \qquad \qquad \begin{array}{ll} \text{WFTyE-cons} \\ \underline{\vdash} \\ \Delta \\ \end{array} \qquad \qquad \begin{array}{ll} X \not \in \Delta \\ \\ \underline{\vdash} \\ \Delta \\ \end{array}$$

3 Typing

 $\Delta \vdash \Gamma$

(Well Formed Term Environment)

$$\frac{\text{WFTME-empty}}{\Delta \vdash \cdot} \qquad \qquad \frac{\frac{\text{WFTME-cons}}{\Delta \vdash \Gamma \quad \Delta \vdash A \quad x \notin \Gamma}}{\Delta \vdash \Gamma, x : A}$$

 $\Delta \vdash A$

(Well Formed Type)

 $\Delta \vdash A \hookrightarrow B : c$

(Typing Reduction Rules)

Typing-abs

 Δ ; Γ , x : $A_1 \vdash e$: A_2

$$\begin{array}{c} \text{TCast-arrow} \\ \Delta \vdash A \\ \Delta \vdash A \\ \hline \Delta \vdash A \hookrightarrow A : \text{id} \end{array} \qquad \begin{array}{c} \Delta \vdash A_1 \hookrightarrow A_2 : c_1 \\ \Delta \vdash B_1 \hookrightarrow B_2 : c_2 \\ \hline \Delta \vdash A_1 \rightarrow B_1 \hookrightarrow A_2 \rightarrow B_2 : c_1 \rightarrow c_2 \end{array} \qquad \begin{array}{c} \text{TCast-unfold} \\ \Delta \vdash \mu X.A \\ \hline \Delta \vdash \mu X.A \hookrightarrow A[X \mapsto \mu X.A] : \downarrow_{\mu X.A} \\ \hline \Delta \vdash A[X \mapsto \mu X.A] \hookrightarrow \mu X.A : \uparrow_{\mu X.A} \\ \hline \end{array}$$

 Δ ; $\Gamma \vdash e : A$

Typing-int

 $\vdash \Delta \qquad \Delta \vdash \Gamma$

(Typing rules)

 $\begin{array}{ll} \text{Typing-var} \\ \vdash \Delta & \Delta \vdash \Gamma \end{array} \quad x : \mathbf{A} \subseteq \Gamma$

4 Semantics

value e (Values)

V-LIT	V-abs	V-fold	V-ARROW
V-L11		value e	value e
$\overline{value\ i}$	value $(\lambda x : A.e)$	value $(\mathbf{cast} [\uparrow_A] e)$	value ($\mathbf{cast}\left[c_1 ightarrow c_2] e ight)$

 $c_1 \sim c_2$ (DualCast)

$$\begin{array}{ccc} \mathrm{DCast\text{-}arrow} & & \mathrm{DCast\text{-}arrow} \\ \frac{c_1 \sim c_1' & c_2 \sim c_2'}{c_1 \rightarrow c_2 \sim c_1' \rightarrow c_2'} & & \frac{\mathrm{DCast\text{-}rec}}{\downarrow_A \sim \uparrow_A} \end{array}$$

$$e \hookrightarrow e'$$

(Reduction rules)

$$\overline{(\lambda x : A. e) e' \hookrightarrow e[x \mapsto e']}$$

$$\frac{e_1 \hookrightarrow e_1'}{e_1 e_2 \hookrightarrow e_1' e_2}$$

Red-cast-arr

$$\frac{\text{RED-APPR}}{\text{value } \nu_1} \quad e_2 \hookrightarrow e_2' \\ \hline \nu_1 \, e_2 \hookrightarrow \nu_1 \, e_2' \\ \hline$$

Red-fix

$$\overline{\mathsf{fix}\ x : \mathsf{A.\,e} \hookrightarrow \mathsf{e}[x \mapsto \mathsf{fix}\ x : \mathsf{A.\,e}]}$$

$$\overline{\left(\operatorname{\mathbf{cast}}\left[c_{1} \to c_{2}\right]e_{1}\right)e_{2} \hookrightarrow \operatorname{\mathbf{cast}}\left[c_{2}\right]\left(e_{1}\left(\operatorname{\mathbf{cast}}\left[\neg c_{1}\right]e_{2}\right)\right)}$$

$$\frac{e \hookrightarrow e'}{\text{cast } [c]e \hookrightarrow \text{cast } [c]e'}$$

$$\frac{c_1 \sim c_2 \quad \text{value } \nu}{\mathbf{cast} \, [c_1] (\mathbf{cast} \, [c_2] \nu) \hookrightarrow \nu}$$

$$\frac{\underset{\text{value } v}{\text{rest [id]}} v \hookrightarrow v}{\text{cast [id]} v \hookrightarrow v}$$