x86

bananaapple

Before we start

Add architecture

- dpkg --add-architecture i386

Update repository

- apt-get update

Install library

- apt-get install ia32-libs
- apt-get install gcc-multilib

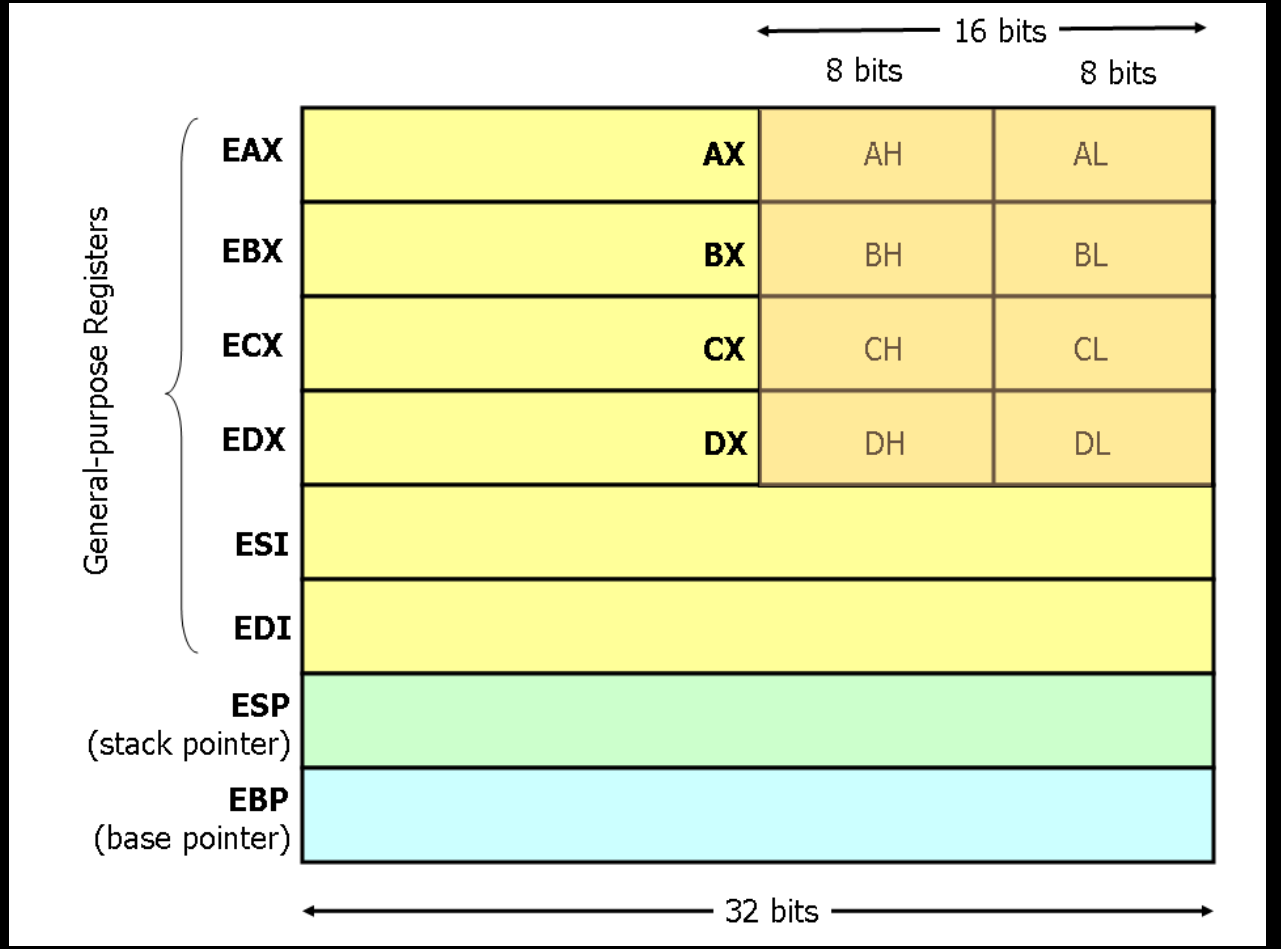
Who am I?

- ID : bananaapple
- 學校科系 ： 交通大學資工系
- 年級 ： 大三升大四
- 目前為 Bamboofox 中的一員

Outline

- Registors
- Flags
- Modes
- Common Instructions
- Intel and AT&T Syntax
- System Call
- Practice
- Example

Registors


16 bits
8 bits
8 bits
General-purpose Registers
EAX
AX
AH
AL
EBX
BX
BH
BL
ECX
CX
CH
CL
EDX
DX
DH
DL
ESI
EDI
ESP
(stack pointer)
EBP
(base pointer)
32 bits

Registors

- eax : accumulator
- ebx : base registor
- ecx : loop counter
- edx : data registor
- esi, edi : index registor
- esp : stack pointer
- ebp : stack base pointer

- eip : instruction pointer

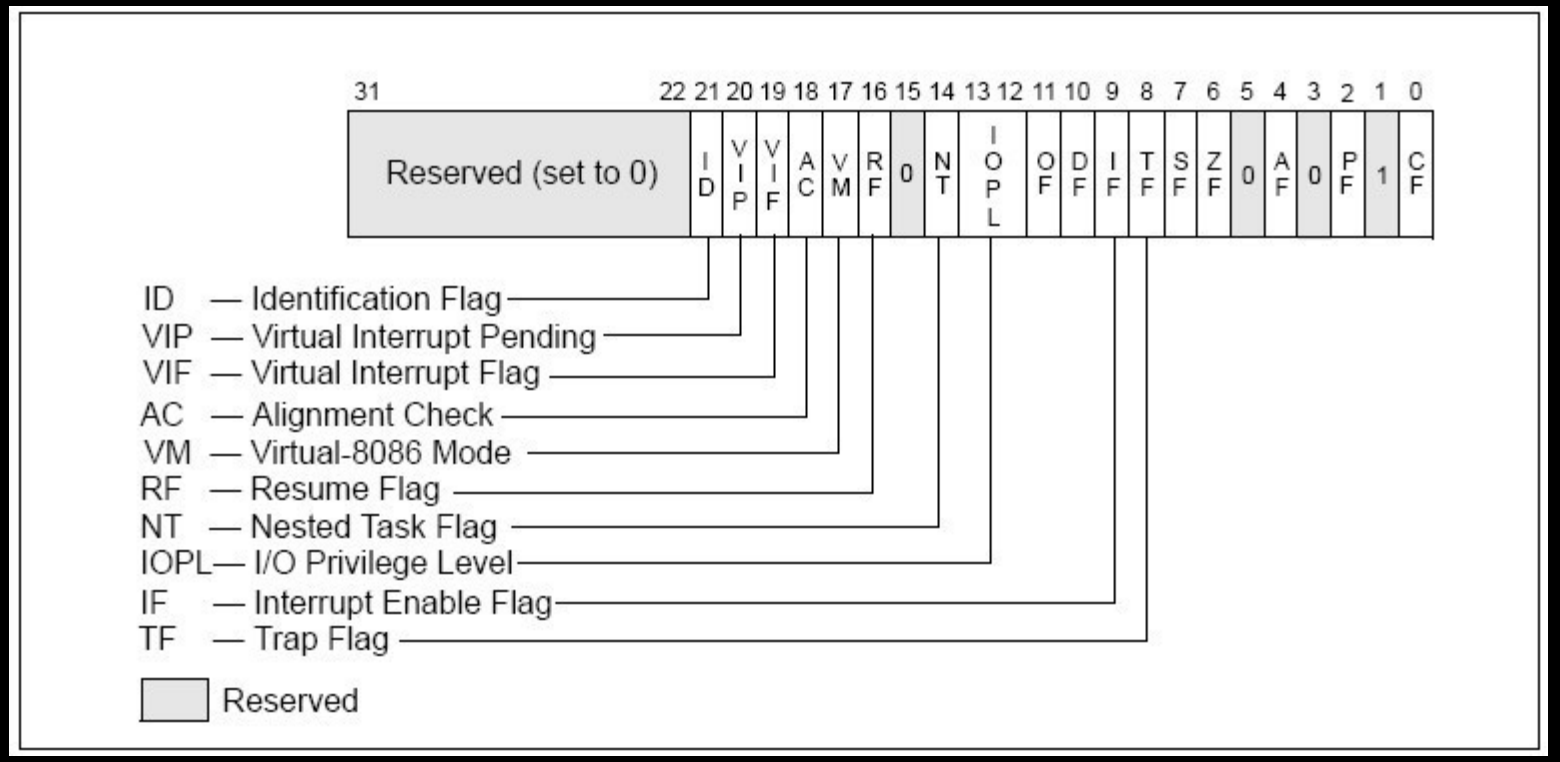
Segment Registers

- cs : code segment
- ds : data segment
- ss : stack segment
- es, fs, gs : additional segment

flags

- Status flag
- Each flag is one bit

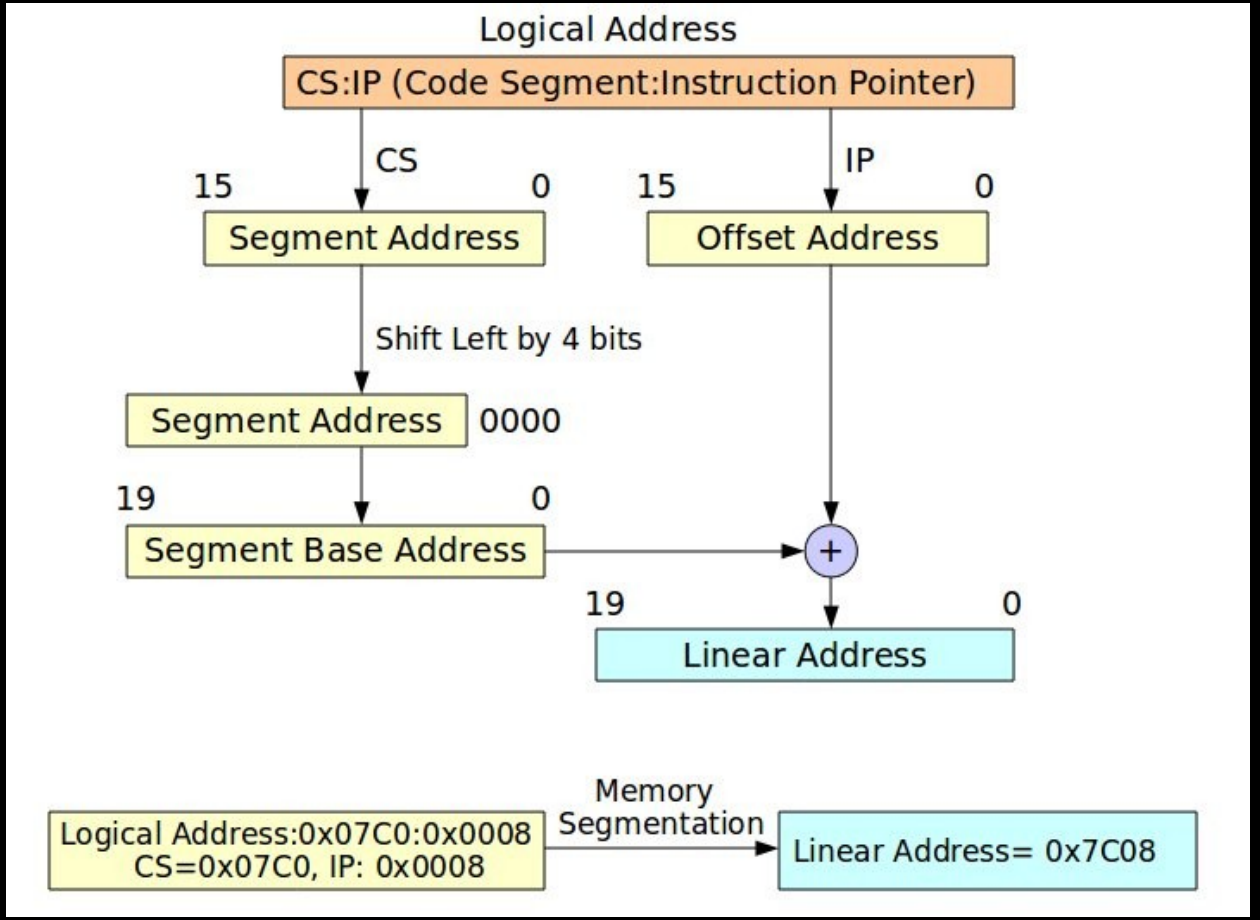
Flags


31 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0
Reserved (set to 0)
ID VIP VIF AC VM RF 0 NT IOPL OF DF IF TF SF ZF 0 AF 0 PF 1 CF
ID — Identification Flag
VIP — Virtual Interrupt Pending
VIF — Virtual Interrupt Flag
AC — Alignment Check
VM — Virtual-8086 Mode
RF — Resume Flag
NT — Nested Task Flag
IOPL — I/O Privilege Level
IF — Interrupt Enable Flag
TF — Trap Flag
Reserved

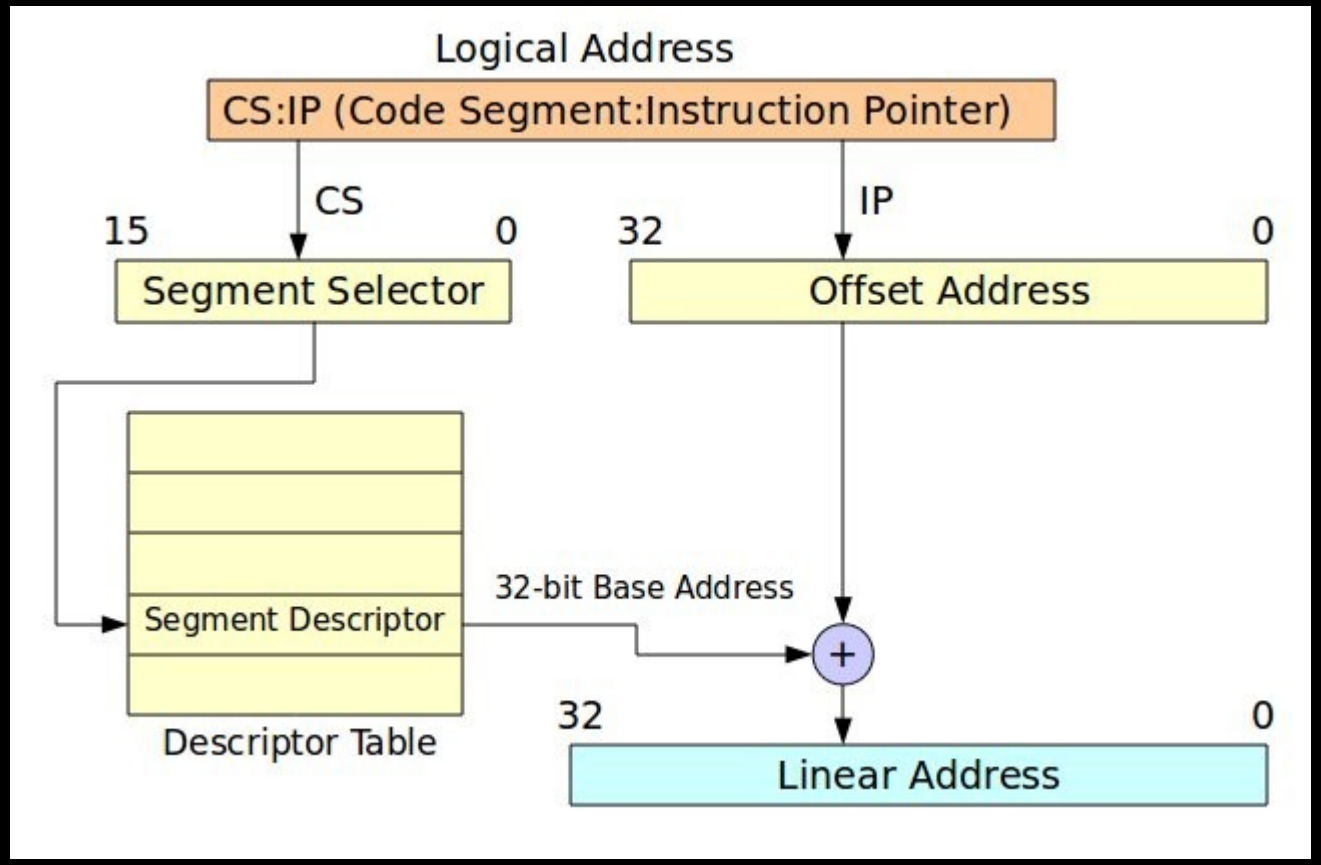
Modes

- Two Modes, Real Mode and Protect Mode
- Real Mode use two 16 bit register to represent 20 bit address space
  - segment:offset => segment << 4 + offset
- segment:offset => segment << 4 + offset
- Can use up 1MB memory (1MB = 2^{20})
- Protect Mode
  - segment:offset => Segment Descriptor + offset
- segment:offset => Segment Descriptor + offset

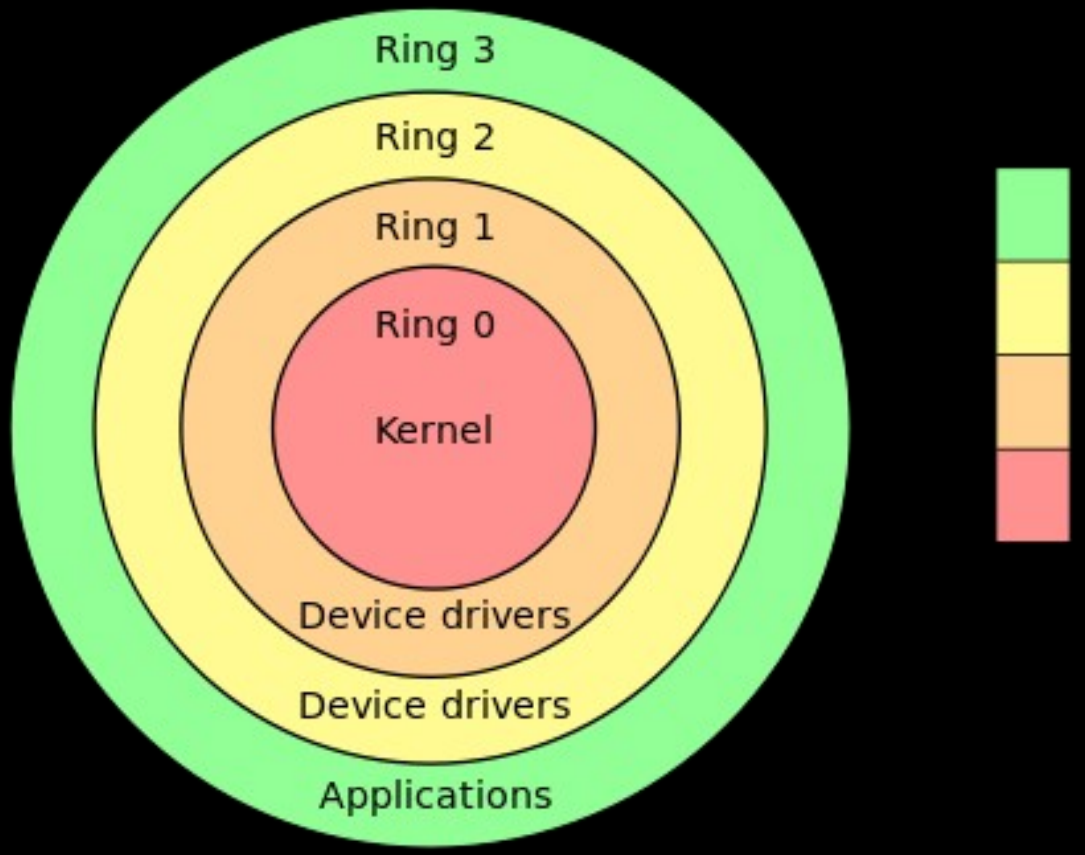
Real Mode


Logical Address
CS:IP (Code Segment:Instruction Pointer)
CS
IP
15
0
15
0
Segment Address
Offset Address
Shift Left by 4 bits
Segment Address
0000
19
0
Segment Base Address
+
19
0
Linear Address
Logical Address:0x07C0:0x0008
CS=0x07C0, IP: 0x0008
Memory
Segmentation
Linear Address= 0x7C08

Protect Mode



Kernel Mode User Mode



Common Instructions

mov - Move

Syntax

- mov dest, source

Example

- mov eax, [ebx]
- mov eax, [ebp - 4]
- mov [var], ebx

Common Instructions

push - Push stack

pop - Pop stack

Example

- push eax
- push 0
- pop eax
- pop [ebx]

Common Instructions

lea - Load effective address

Syntax

- lea <reg32>, <mem>

Example

- lea ebx, [ebx+eax*8]
- lea eax, [ebp-0x44]

Common Instructions

add, sub, mul, div - Arithmetic

inc ,dec - Increment, Decrement

Syntax

- add dest, source
- inc <reg> or <mem>

Example

- add eax, 10
- inc eax

Common Instructions

jmp – Jump

- je <label> (jump when equal)
- jne <label> (jump when not equal)
- jz <label> (jump when last result was zero)
- jg <label> (jump when greater than)
- jge <label> (jump when greater than or equal to)
- jl <label> (jump when less than)
- jle <label> (jump when less than or equal to)

Common Instructions

cmp – Compare

Example

- cmp DWORD PTR [eax], 10
- je loop
- cmp eax, ebx
- jle done
- jmp DWORD PTR [eax]

Intel and AT&T Syntax

- Prefixes
- Direction of Operands
- Memory Operands
- Suffixes

Prefixes

Intex Syntax

- mov eax,1
- mov ebx,0ffh
- int 80h

AT&T Syntax

- movl $1,%eax
- movl $0xff,%ebx
- int $0x80

Direction of Operands

Intex Syntax

- instr dest,source
- mov eax,[ecx]

AT&T Syntax

- instr source,dest
- movl (%ecx),%eax

Memory Operands

Intex Syntax

- mov eax,[ebx]
- mov eax,[ebx+3]

AT&T Syntax

- movl (%ebx),%eax
- movl 3(%ebx),%eax

Suffixes

Intel Syntax

- Instr foo,segreg:[base+index* scale+disp]
- mov eax,[ebx+20h]
- add eax,[ebx+ecx*2h]
- lea eax,[ebx+ecx]
- sub eax,[ebx+ecx*4h-20h]

AT&T Syntax

- Instr %segreg:disp(base,index,scale),foo
- movl 0x20(%ebx),%eax
- addl (%ebx,%ecx,0x2),%eax
- leal (%ebx,%ecx),%eax
- subl -0x20(%ebx,%ecx,0x4),%eax

System Call

- Syscalls are the interface between user programs and the Linux kernel
- Put value on registers eax, ebx
- eax represent system call number
- ebx, ecx ..... represent arguments
- Finally, execute **int 0x80** instruction
- Return value will put on eax register
- If you want to know more about system call, type man 2 system_call (ex:open)
- http://docs.cs.up.ac.za/programming/asm/derick_tut/syscalls.html

Practice

```
section     .text
global      _start

_start:
    ;You are going to practice system call
    ;What you should do?
    ;put system call number in %eax
    ;put fd number in %ebx
    ;put string address in %ecx
    ;put string length in %edx
    ;interrupt

section     .data

msg     db  'Hello, world!',0xa
len     equ $ - msg
```

wget http://people.cs.nctu.edu.tw/~wpchen/x86/practice.asm

nasm -f elf practice.asm

ld -m elf_i386 -s -o practice practice.o

./practice

//Hello, world!

Answer

```
section     .text
global      _start

_start:

    mov     edx,len
    mov     ecx,msg
    mov     ebx,1
    mov     eax,4
    int     0x80

    mov     eax,1
    int     0x80

section     .data

msg     db  'Hello, world!',0xa
len     equ $ - msg
```

wget http://people.cs.nctu.edu.tw/~wpchen/x86/hello.asm

nasm -f elf hello.asm

ld -m elf_i386 -s -o hello hello.o

./hello

//Hello, world!

Not enough?

Try this one:

http://secprog.cs.nctu.edu.tw/problems/3

Open your terminal and type:

nc secprog.cs.nctu.edu.tw 10003

Hint : open /home/rop/flag -> read from fd -> write to stdout

Have fun!!!

Example

```c
#include<stdio.h>
int sum(int i,int j)
{
    int sum;
    sum=i+j;
    return sum;
}
int main(void)
{
    int i;
    int j;
    int k;
    scanf("%d%d",&i,&j);
    k=sum(i,j);
    printf("Sum:%d\n",k);
    return 0;
}
```

wget http://people.cs.nctu.edu.tw/~wpchen/x86/sum.c

gcc -m32 -o sum sum.c

//or just download it

wget http://people.cs.nctu.edu.tw/~wpchen/sum

objdump -d sum | less

Example

```
08048482 <main>:
 8048482:       55                      push   %ebp
 8048483:       89 e5                   mov    %esp,%ebp
 8048485:       83 e4 f0                and    $0xfffffff0,%esp
 8048488:       83 ec 20                sub    $0x20,%esp
 804848b:       8d 44 24 14             lea    0x14(%esp),%eax
 804848f:       89 44 24 08             mov    %eax,0x8(%esp)
 8048493:       8d 44 24 18             lea    0x18(%esp),%eax
 8048497:       89 44 24 04             mov    %eax,0x4(%esp)
 804849b:       c7 04 24 70 85 04 08    movl   $0x8048570,(%esp)
 80484a2:       e8 c9 fe ff ff          call   8048370 <__isoc99_scanf@plt>
 80484a7:       8b 54 24 14             mov    0x14(%esp),%edx
 80484ab:       8b 44 24 18             mov    0x18(%esp),%eax
 80484af:       89 54 24 04             mov    %edx,0x4(%esp)
 80484b3:       89 04 24                mov    %eax,(%esp)
 80484b6:       e8 b1 ff ff ff          call   804846c <sum>
```

Why?

Answer

This code makes sure that the stack is aligned to 16 bytes. After t
his operation esp will be less than or equal to what it was before
this operation, so the stack may grow, which protects anything t
hat might already be on the stack. This is sometimes done in ma
in just in case the function is called with an unaligned stack, whi
ch can cause things to be really slow (16 byte is a cache line widt
h on x86, I think, though 4 byte alignment is what is really impor
tant here). If main has a unaligned stack the rest of the program
will too.

http://stackoverflow.com/questions/4228261/understanding-th
e-purpose-of-some-assembly-statements

Example

```
80484a2:    e8 c9 fe ff ff          call   8048370 <__isoc99_scanf@plt>
80484a7:    8b 54 24 14             mov    0x14(%esp),%edx
80484ab:    8b 44 24 18             mov    0x18(%esp),%eax
80484af:    89 54 24 04             mov    %edx,0x4(%esp)
80484b3:    89 04 24                mov    %eax,(%esp)
80484b6:    e8 b1 ff ff ff          call   804846c <sum>
80484bb:    89 44 24 1c             mov    %eax,0x1c(%esp)
80484bf:    8b 44 24 1c             mov    0x1c(%esp),%eax
80484c3:    89 44 24 04             mov    %eax,0x4(%esp)
80484c7:    c7 04 24 75 85 04 08    movl   $0x8048575,(%esp)
80484ce:    e8 6d fe ff ff          call   8048340 <printf@plt>
```

Function call

Example

```
0804846c <sum>:
 804846c:       55                      push   %ebp
 804846d:       89 e5                   mov    %esp,%ebp
 804846f:       83 ec 10                sub    $0x10,%esp
 8048472:       8b 45 0c                mov    0xc(%ebp),%eax
 8048475:       8b 55 08                mov    0x8(%ebp),%edx
 8048478:       01 d0                   add    %edx,%eax
 804847a:       89 45 fc                mov    %eax,-0x4(%ebp)
 804847d:       8b 45 fc                mov    -0x4(%ebp),%eax
 8048480:       c9                      leave
 8048481:       c3                      ret
```

0804846c is address of function

each line represent one command

each command has different length

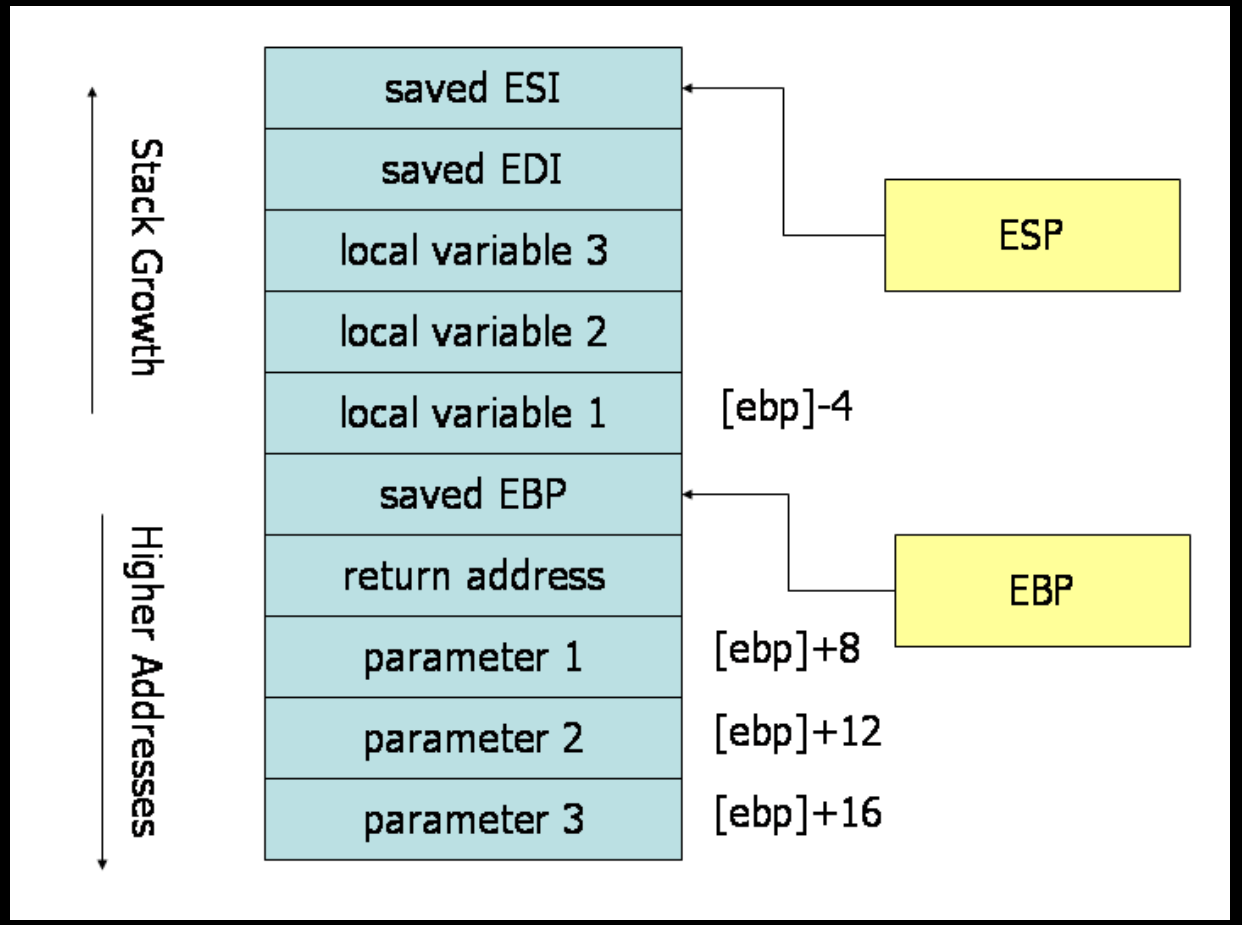
Example

```
0804846c <sum>:
 804846c:       55                      push   %ebp
 804846d:       89 e5                   mov    %esp,%ebp
 804846f:       83 ec 10                sub    $0x10,%esp
 8048472:       8b 45 0c                mov    0xc(%ebp),%eax
 8048475:       8b 55 08                mov    0x8(%ebp),%edx
 8048478:       01 d0                   add    %edx,%eax
 804847a:       89 45 fc                mov    %eax,-0x4(%ebp)
 804847d:       8b 45 fc                mov    -0x4(%ebp),%eax
 8048480:       c9                      leave
 8048481:       c3                      ret
```

Function prologue

Function epilogue

Example


Stack Growth
Higher Addresses
saved ESI
saved EDI
local variable 3
local variable 2
local variable 1
saved EBP
return address
parameter 1
parameter 2
parameter 3
ESP
EBP
[ebp]-4
[ebp]+8
[ebp]+12
[ebp]+16

Example

```
0804846c <sum>:
 804846c:       55                      push   %ebp
 804846d:       89 e5                   mov    %esp,%ebp
 804846f:       83 ec 10                sub    $0x10,%esp
 8048472:       8b 45 0c                mov    0xc(%ebp),%eax
 8048475:       8b 55 08                mov    0x8(%ebp),%edx
 8048478:       01 d0                   add    %edx,%eax
 804847a:       89 45 fc                mov    %eax,-0x4(%ebp)
 804847d:       8b 45 fc                mov    -0x4(%ebp),%eax
 8048480:       c9                      leave
 8048481:       c3                      ret
```

Why?

I only use 4bytes

Answer

Sometimes , compiler will optimize the code by adding some pa dding to make it align to word boundary

You have to inspect the assembly code to know the exactly stack position
There are special instructions called SSE2 on x86 CPUs *do* requir e the data to be 128-bit (16-byte) aligned

Most of the SSE2 instructions implement the integer vector oper ations also found in MMX
https://en.wikipedia.org/wiki/Data_structure_alignment

Example

```
0804846c <sum>:
 804846c:       55                      push   %ebp
 804846d:       89 e5                   mov    %esp,%ebp
 804846f:       83 ec 10                sub    $0x10,%esp
 8048472:       8b 45 0c                mov    0xc(%ebp),%eax
 8048475:       8b 55 08                mov    0x8(%ebp),%edx
 8048478:       01 d0                   add    %edx,%eax
 804847a:       89 45 fc                mov    %eax,-0x4(%ebp)
 804847d:       8b 45 fc                mov    -0x4(%ebp),%eax
 8048480:       c9                      leave
 8048481:       c3                      ret
```

second argument

first argument

return value on eax

Example

- Intel and AT&T Syntax

http://asm.sourceforge.net/articles/linasm.html

- hello.asm

http://asm.sourceforge.net/intro/hello.html

- Stack overflow

http://stackoverflow.com/questions/4228261/understanding-the-purpose-of-some-assembly-statements

Reference

- x86 Assembly Guide (recommended)

http://www.cs.virginia.edu/~evans/cs216/guides/x86.html

- Linux System Call Table

http://
docs.cs.up.ac.za/programming/asm/derick_tut/syscalls.html

- Wiki

https://en.wikipedia.org/wiki/X86_assembly_language

https://
en.wikibooks.org/wiki/X86_Assembly/Interfacing_with_Linux