

# Scientific Programming: Algorithms (part B)

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Programming paradigms

# Problems and solutions

Given a problem:

- There are no "general recipes" to solve it
- Nevertheless, we can identify four phases:
  - Problem classification
  - Solution characterization
  - Selection of the algorithmic technique
  - Selection of the data structure
- These phases are not strictly sequential

# Classification of problems

## Decisional problems

- Does the input satisfy a given property?
- Output: the answer is yes/no
- Example: is the graph connected?

## Search problems

- Research space: a set of possible "solutions"
- Admissible solution: a solution that does satisfy some conditions
- Example: position of a substring in the string

# Classification of problems

## Optimization problems

- Each solution is associated with a cost function
- We want to identify the solution with minimum cost
- Example: the shortest path between nodes in a graph

## Approximation problems

- Sometimes, obtaining the optimal solution is computationally infeasible
- We may be satisfied by an approximate solution: low cost, but we are not sure that the cost is the smallest possible
- Example: the traveling salesman problem

# Mathematical characterization

It is important to mathematically define the relationship between input and output

- Very often the mathematical characterization is trivial...
- ... but it could provide a first idea of the solution
- Example: given a sequence of  $n$  elements, a sorted permutation is given by the minimum followed by a sorted permutation of the remaining  $n - 1$  elements (Selection Sort)

The mathematical characterization can suggest a possible technique

- Optimal substructure → Dynamic programming
- Greedy choice → Greedy technique

# Algorithmic techniques

## Divide-et-impera

- The problem is subdivided in **independent** subproblems, that are solved recursively (in a **top-down** approach)
- Area of application: decision problems, search

## Dynamic programming

- The solution is built in a **bottom-up** way from the solution of smaller problems (**potentially repeated**)
- Area of application: optimization problems

## Memoization

- Top-down version of dynamic programming

# Algorithmic techniques

## Greedy

- Greedy approach: select the choice which appears "locally optimal"
- Area of application: optimization problems

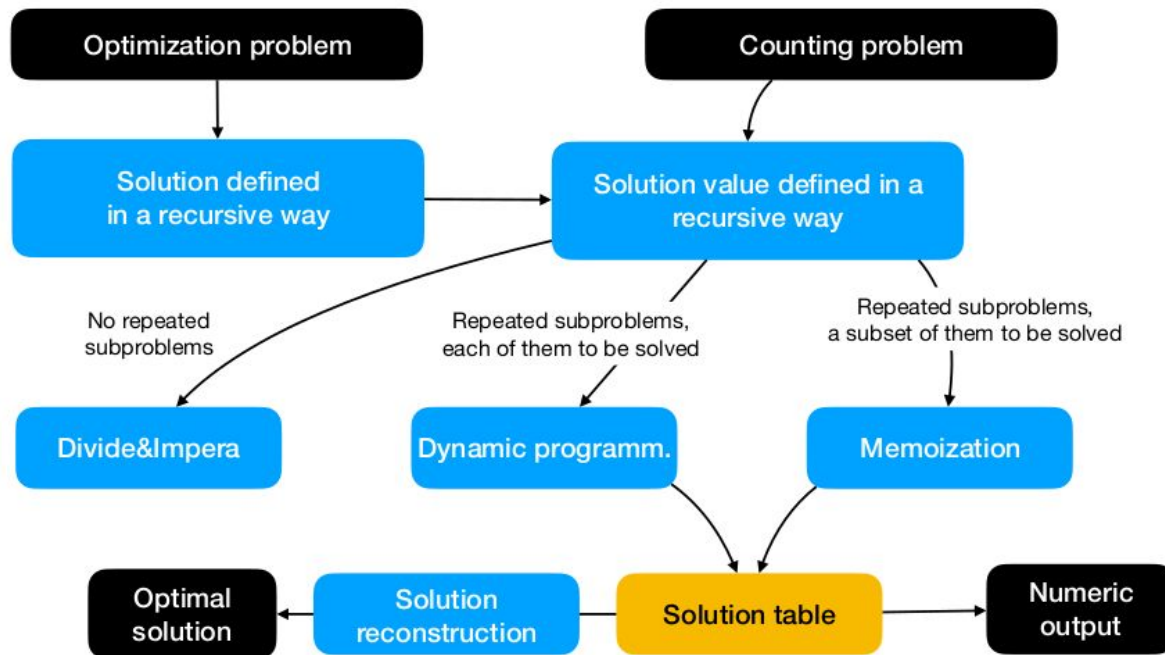
## Backtrack

- Try something, and if does not work, try something else
- Area of application: search problems, optimization problems

## Local search

- The optimal solution can be obtained by continuously improving sub-optimal solutions

# General approach





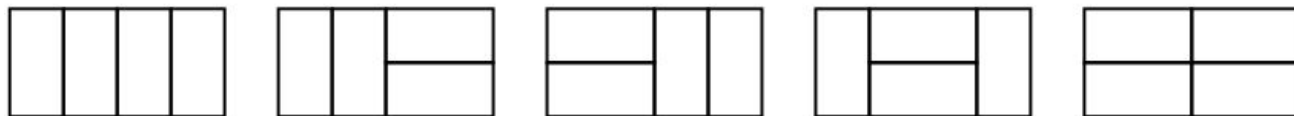
# Dominoes

## Definition

The dominoes game consists of tiles with size  $2 \times 1$ . Let us consider the arrangements of  $n$  tiles inside a rectangle  $2 \times n$ . Write an efficient algorithm that computes the number of possible arrangements and discuss its correctness. Compute an upper bound to its complexity.

## Example

The cases below represent the five possible arrangements in a rectangle  $2 \times 4$ .



Any ideas on how to solve this problem?

# Dominoes

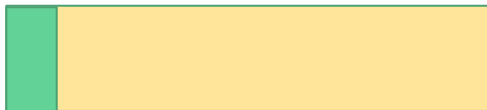
## Recursive definition

Let's define a recursive formula that computes the number of possible arrangements.

- If a vertical tile is placed, the problem of size  $n - 1$  must be solved.
- If an horizontal tile is placed, then another horizontal tile must be placed as well; the problem of size  $n - 2$  must be solved.

$$D(n) = \begin{cases} 1 \\ ? \end{cases}$$

$2 \times n$

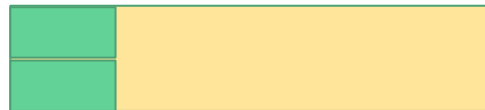


$$n \leq 1$$

$$n > 1$$



$2 \times n$



$n = 0$ , only one possibility: **no tiles**.  
 $n = 1$ , only 1 possibility,  
**vertical tile**

# Dominoes

## Recursive definition

Let's define a recursive formula that computes the number of possible arrangements.

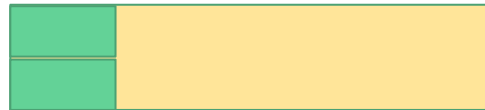
- If a vertical tile is placed, the problem of size  $n - 1$  must be solved.
- If an horizontal tile is placed, then another horizontal tile must be placed as well; the problem of size  $n - 2$  must be solved.

$$D(n) = \begin{cases} 1 & n \leq 1 \\ D(n-2) + D(n-1) & n > 1 \end{cases} \quad \leftarrow$$

2xn



2xn



We sum because the two cases originate different solutions

# Dominoes

$$D(n) = \begin{cases} 1 & n \leq 1 \\ D(n-2) + D(n-1) & n > 1 \end{cases}$$

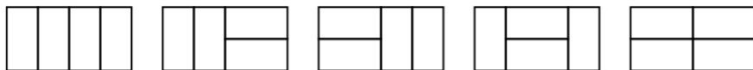
The generated mathematical series is the following:

1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, ...

Does it sound familiar?

Fibonacci's numbers!

$N = 4$  (i.e.  $2 \times 4$ )  $\rightarrow$  5 possible dispositions



# Dominoes: recursive algorithm

$$D(n) = \begin{cases} 1 & n \leq 1 \\ D(n-2) + D(n-1) & n > 1 \end{cases}$$

Write a recursive algorithm that solves the problem

```
def dominoes(n):  
    if n <= 1:  
        return 1  
    else:  
        return dominoes(n-2) + dominoes(n-1)  
  
for i in range(10):  
    print(dominoes(i), end = " ")  
  
1 1 2 3 5 8 13 21 34 55
```

# Complexity

What is the complexity of dominoes?

```
def dominoes(n):  
    if n <= 1:  
        return 1  
    else:  
        return dominoes(n-2) + dominoes(n-1)  
  
for i in range(10):  
    print(dominoes(i), end = " ")  
  
1 1 2 3 5 8 13 21 34 55
```

$$T(n) = \begin{cases} 1 & n \leq 1 \\ T(n-1) + T(n-2) + 1 & n > 1 \end{cases}$$

Theorem not seen:

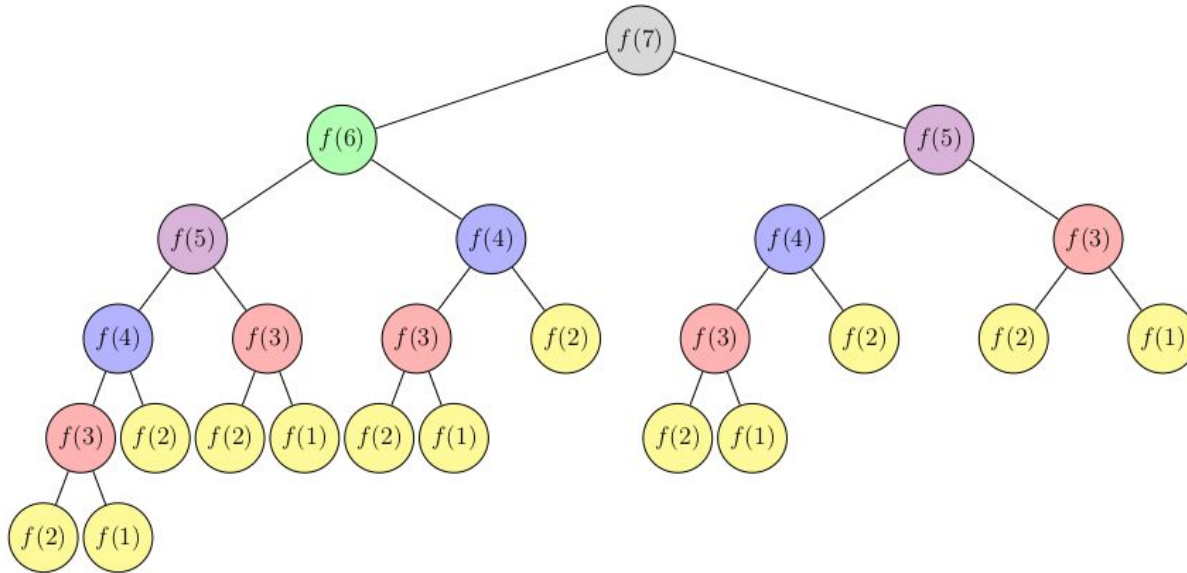
Linear recurrences with constant order:

- $a_1 = 1, a_2 = 1, a = 2, \beta = 0$
- Complexity:  $\Theta(a^n \cdot n^\beta)$

$$T(n) = \Theta(2^n)$$

# Recursive tree

```
def dominoes(n):  
    if n <= 1:  
        return 1  
    else:  
        return dominoes(n-2) + dominoes(n-1)  
  
for i in range(10):  
    print(dominoes(i), end = " ")  
  
1 1 2 3 5 8 13 21 34 55
```



Several sub-problems are repeated!

# How to avoid computing the same thing over and over again

## DP Table

- We use a *DP table* (list, matrix, dictionary, etc.) to store results of sub-problems already solved
- The table contains an entry for each subproblem to be solved
- The table is indexed by a description of the input (e.g., size)
- When the same subproblem has to be solved again, we use the result stored in the table



# How to avoid computing the same thing over and over again

## Base cases

- The base cases do not need to be computed, they can be stored immediately

## Bottom-up iteration

- We start from problems that can be solved using only base cases
- We go up to larger and larger problems...
- ... up to the final goal

$n$	0	1	2	3	4	5	6	7
$DP[]$	1	1	2	3	5	8	13	21

# An iterative solution

$$D(n) = \begin{cases} 1 & n \leq 1 \\ D(n-2) + D(n-1) & n > 1 \end{cases}$$

Write an iterative algorithm that solves the Dominoes problem

```
def dominoes2(n):  
    res = [0]*(n+1)  
    res[0] = 1  
    res[1] = 1  
    for i in range(2,n+1):  
        res[i] = res[i-1] + res[i-2]  
    return res[n]
```

What is the computational complexity of domino2(n)?

$$T(n) = \Theta(n)$$

How about the space complexity?

What is the size of res?

$$S(n) = \Theta(n)$$

**Ideas on how to improve this?**

## Another iterative solution

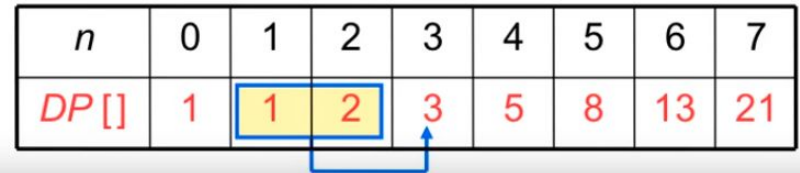
$$D(n) = \begin{cases} 1 & n \leq 1 \\ D(n-2) + D(n-1) & n > 1 \end{cases}$$

```
def dominoes3(n):  
    dp0 = 1  
    dp1 = 1  
    dp2 = 1  
    for i in range(2, n+1):  
        dp0 = dp1  
        dp1 = dp2  
        dp2 = dp0 + dp1  
    return dp2
```

What is the space complexity of domino3(n)?

$$S(n) = \Theta(1)$$

$n$	0	1	2	3	4	5	6	7
$DP[]$	1	1	2	3	5	8	13	21



# Uniform vs Logarithmic cost model

$$D(n) = \begin{cases} 1 & n \leq 1 \\ D(n-2) + D(n-1) & n > 1 \end{cases}$$

Are you sure that our complexity formulas are correct?

## Binet's Formula for Fibonacci's number

$$D(n-1) = F(n) = \frac{\phi^n}{\sqrt{5}} - \frac{(1-\phi)^n}{\sqrt{5}} = \frac{\phi^n - (-\phi)^{-n}}{\sqrt{5}}$$

where

$$\phi = \frac{1 + \sqrt{5}}{2} = 1,6180339887 \dots \quad \text{golden ratio}$$

Watch out: the Fibonacci's number grows exponentially!

How many bits are needed to store  $F(n)$ ?

# Uniform vs Logarithmic cost model

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where

$$\phi = \frac{1 + \sqrt{5}}{2} = 1,6180339887 \dots$$

How many bits are needed to store  $F(n)$ ?

$$\log F(n) = \Theta(n)$$



complexity seen before needs to be multiplied by  $n$

# Uniform vs Logarithmic cost model

$$D(n) = \begin{cases} 1 & n \leq 1 \\ D(n-2) + D(n-1) & n > 1 \end{cases}$$

Under the logarithmic cost model, the three versions have the following complexities:

Function	Time complexity	Space complexity
domino1()	$O(n2^n)$	$O(n^2)$
domino2()	$O(n^2)$	$O(n^2)$
domino3()	$O(n^2)$	$O(n)$

```
s = time.time()
for i in range(1,45):
    print(dominoes(i), end = " ")
e = time.time()
print("Elapsed time: {}s".format(e-s))
s = time.time()
for i in range(1,45):
    print(dominoes2(i), end = " ")
e = time.time()
print("Elapsed time: {}s".format(e-s))

s = time.time()
for i in range(1,45):
    print(dominoes3(i), end = " ")
e = time.time()
print("Elapsed time: {}s".format(e-s))
```

1 2 3 5 8 ... 1134903170  
Elapsed time: 659.3645467758179s

1 2 3 5 8 ... 1134903170  
Elapsed time: 0.0007071495056152344s

1 2 3 5 8 ... 1134903170  
Elapsed time: 0.0011742115020751953s

# Uniform vs Logarithmic cost model

$$D(n) = \begin{cases} 1 & n \leq 1 \\ D(n-2) + D(n-1) & n > 1 \end{cases}$$

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print("Elapsed time: {}s".format(e-s))

s = time.time()
for i in range(1,45):
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1 2 3 5 8 ... 1134903170  
Elapsed time: 0.0011742115020751953s

# Hateville

- Hateville is a strange village, composed of  $n$  houses, numbered  $1$ - $n$  and placed along a single road
- In Hateville, everybody hates his next-door neighbors, on both sides: thus a person living in house  $i$  hates the neighbors living in houses  $i - 1$  and  $i + 1$  (if they exist)
- Hateville wants to organize a festival; your task is to collect money to organize it.
- Each inhabitant  $i$  wants to donate a quantity  $D[i]$  of money, but he will give nothing if any of his neighbors is donating.





# Hateville



Consider the following problems:

- Write an algorithm that returns the largest amount of money that can be collected
- Write an algorithm that returns a subset of indexes  $S \subseteq \{1, \dots, n\}$  such that the total amount  $T = \sum_{i \in S} D[i]$  is maximal.



remember the  
additional constraint  
that indexes must not  
be consecutive

Examples:

- |                                     |                                       |
|-------------------------------------|---------------------------------------|
| • Donation list: $D = [4, 3, 6, 5]$ | • Donation list: $D = [10, 5, 5, 10]$ |
| • Maximum amount: 10                | • Maximum amount: 20                  |
| • Index set: $\{1, 3\}$             | • Index set: $\{1, 4\}$               |

# Hateville

- Donation list:  $D = [4, 3, 6, 5]$
- Maximum amount: 10
- Index set:  $\{1, 3\}$
- Donation list:  $D = [10, 5, 5, 10]$
- Maximum amount: 20
- Index set:  $\{1, 4\}$

How would you solve the problem?

We re-define the problem

- Let  $HV(i)$  be the set of numbers to be selected to obtain the maximum amount of donations from the first  $i$  houses, numbered  $1 \dots n$
- $HV(n)$  is the solution to the original problem

# Hateville

- Donation list:  $D = [4, 3, 6, 5]$
- Maximum amount: 10
- Index set:  $\{1, 3\}$
- Donation list:  $D = [10, 5, 5, 10]$
- Maximum amount: 20
- Index set:  $\{1, 4\}$

Let's compute  $HV(i)$  based on  $HV(0) \dots HV(n-1)$  values

- What happens if I don't accept its donation?

$$HV(i) =$$

# Hateville

- Donation list:  $D = [4, 3, 6, 5]$
- Maximum amount: 10
- Index set:  $\{1, 3\}$
- Donation list:  $D = [10, 5, 5, 10]$
- Maximum amount: 20
- Index set:  $\{1, 4\}$

Let's compute  $HV(i)$  based on  $HV(0) \dots HV(n-1)$  values

- What happens if I don't accept its donation?

$$HV(i) = HV(i-1)$$

# Hateville

- Donation list:  $D = [4, 3, 6, 5]$
- Maximum amount: 10
- Index set:  $\{1, 3\}$
- Donation list:  $D = [10, 5, 5, 10]$
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Let's compute  $HV(i)$  based on  $HV(0) \dots HV(n-1)$  values

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$$HV(i) = HV(i-1)$$

- What happens if I accept its donation?

# Hateville

- Donation list:  $D = [4, 3, 6, 5]$
- Maximum amount: 10
- Index set:  $\{1, 3\}$
- Donation list:  $D = [10, 5, 5, 10]$
- Maximum amount: 20
- Index set:  $\{1, 4\}$

Let's compute  $HV(i)$  based on  $HV(0) \dots HV(n-1)$  values

- What happens if I don't accept its donation?

$$HV(i) = HV(i-1)$$

- What happens if I accept its donation?

$$HV(i) = HV(i-2) + D[i]$$

- How can I choose between the two cases?

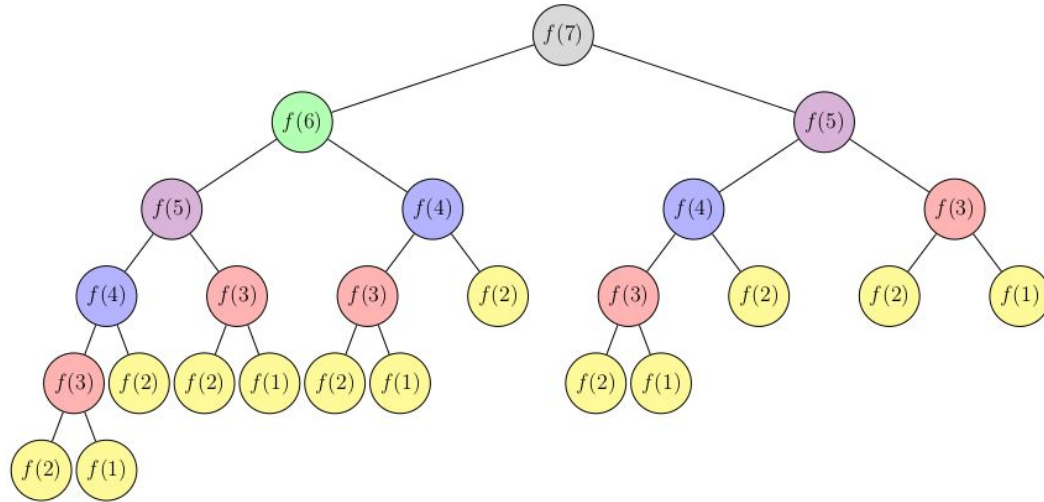
$$\max(HV(i-1), HV(i-2) + D[i])$$

## Hateville: recursive algorithm?

$$\max(HV(i-1), HV(i-2) + D[i])$$

Write a recursive algorithm that solves Hateville?

Would it be a good idea?



# DP Table

$$\max(HV(i-1), HV(i-2) + D[i])$$

## Value of the optimal solution

- Let  $DP(i)$  be the **value** of the maximum amount of donation that we can obtain from the first  $i$  houses of Hateville
- $DP(n)$  is the value of the optimal solution

$$DP(i) = \begin{cases} 0 & \text{if } i = 0 \\ D[1] & \text{if } i = 1 \\ \max(DP(i-1), DP(i-2) + D[i]) & \text{if } n \geq 2 \end{cases}$$



# Iterative solution

$$DP(i) = \begin{cases} 0 & \text{if } i = 0 \\ D[1] & \text{if } i = 1 \\ \max(DP(i-1), DP(i-2) + D[i]) & \text{if } i \geq 2 \end{cases}$$

Write an algorithm that solves the Hateville problem

```
def hateville(D, n):
    dp = [0]*(n+1)
    if n > 0:
        dp[1] = D[0]
    for i in range(2, n+1):
        dp[i] = max(dp[i-1], dp[i-2] + D[i-1])

    return dp[n]
```

```
D = [10,5,5,8,4,7,12]

print("Donations: {}".format(D))
for i in range(len(D)+1):
    print("Solution for {}: {}".format(D[0:i],hateville(D, i)))
```

```
Donations: [10, 5, 5, 8, 4, 7, 12]
Solution for []: 0
Solution for [10]: 10
Solution for [10, 5]: 10
Solution for [10, 5, 5]: 15
Solution for [10, 5, 5, 8]: 18
Solution for [10, 5, 5, 8, 4]: 19
Solution for [10, 5, 5, 8, 4, 7]: 25
Solution for [10, 5, 5, 8, 4, 7, 12]: 31
```

```
D1 = [10,1,1,10,1,1,10]

print("Donations: {}".format(D1))
for i in range(len(D1)+1):
    print("Solution for {}: {}".format(D1[0:i],hateville(D1, i)))
```

```
Donations: [10, 1, 1, 10, 1, 1, 10]
Solution for []: 0
Solution for [10]: 10
Solution for [10, 1]: 10
Solution for [10, 1, 1]: 11
Solution for [10, 1, 1, 10]: 20
Solution for [10, 1, 1, 10, 1]: 20
Solution for [10, 1, 1, 10, 1, 1]: 21
Solution for [10, 1, 1, 10, 1, 1, 10]: 30
```

# Iterative solution

$$DP(i) = \begin{cases} 0 & \text{if } i = 0 \\ D[1] & \text{if } i = 1 \\ \max(DP(i-1), DP(i-2) + D[i]) & \text{if } i \geq 2 \end{cases}$$

<i>i</i>	0	1	2	3	4	5	6	7
<i>D</i>		10	5	5	8	4	7	12
<i>DP</i>	0	10	10	15	18	19	25	31

<i>i</i>	0	1	2	3	4	5	6	7
<i>D</i>		10	1	1	10	1	1	10
<i>DP</i>	0	10	10	11	20	20	21	30

## Problem

- We have the **value** of the optimal solution, but we don't have the solution!
- Look in position  $DP[i]$ . From which cells this value has been computed?
  - If  $DP[i] = DP[i-1]$ , the house  $i$  has not been selected
  - If  $DP[i] = DP[i-2] + D[i-1]$ , house  $i$  has been selected

**Build solution(i) recursively as:**

```
solution(i-2) + index i  
or  
solution(i-1)
```

# Iterative solution

```
def hateville(D, n):
    dp = [0]*(n+1)
    if n > 0:
        dp[1] = D[0]
    for i in range(2, n+1):
        dp[i] = max(dp[i-1], dp[i-2] + D[i-1])

    return build_solution(D, dp, n)

def build_solution(D, dp, i):
    if i == 0:
        return []
    elif i == 1:
        return [0]
    else:
        if dp[i] == dp[i-1]:
            sol = build_solution(D, dp, i-1)
        else:
            sol = build_solution(D, dp, i-2)
            sol.append(i-1)
    return sol
```

```
D = [10,5,5,8,4,7,12]
print("Donations: {}".format(D))
for i in range(len(D)+1):
    HV = hateville(D, i)
    print("Donors for {}: {}. Donations: {}".format(D[0:i], HV, sum([D[x] for x in HV])))
print("\n\n")

D1 = [10,1,1,10,1,1,10]
print("Donations: {}".format(D1))
for i in range(len(D1)+1):
    HV = hateville(D1, i)
    print("Donors for {}: {}. Donations: {}".format(D1[0:i], HV, sum([D1[x] for x in HV])))
```

```
Donations: [10, 5, 5, 8, 4, 7, 12]
Donors for []: []. Donations: 0
Donors for [10]: [0]. Donations: 10
Donors for [10, 5]: [0]. Donations: 10
Donors for [10, 5, 5]: [0, 2]. Donations: 15
Donors for [10, 5, 5, 8]: [0, 3]. Donations: 18
Donors for [10, 5, 5, 8, 4]: [0, 2, 4]. Donations: 19
Donors for [10, 5, 5, 8, 4, 7]: [0, 3, 5]. Donations: 25
Donors for [10, 5, 5, 8, 4, 7, 12]: [0, 2, 4, 6]. Donations: 31
```

```
Donations: [10, 1, 1, 10, 1, 1, 10]
Donors for []: []. Donations: 0
Donors for [10]: [0]. Donations: 10
Donors for [10, 1]: [0]. Donations: 10
Donors for [10, 1, 1]: [0, 2]. Donations: 11
Donors for [10, 1, 1, 10]: [0, 3]. Donations: 20
Donors for [10, 1, 1, 10, 1]: [0, 3, 5]. Donations: 20
Donors for [10, 1, 1, 10, 1, 1]: [0, 3, 5]. Donations: 21
Donors for [10, 1, 1, 10, 1, 1, 10]: [0, 3, 6]. Donations: 30
```

# Complexity

```
def hateville(D, n):
    dp = [0]*(n+1)
    if n > 0:
        dp[1] = D[0]
    for i in range(2, n+1):
        dp[i] = max(dp[i-1], dp[i-2] + D[i-1])

    return build_solution(D, dp, n)

def build_solution(D, dp, i):
    if i == 0:
        return []
    elif i == 1:
        return [0]
    else:
        if dp[i] == dp[i-1]:
            sol = build_solution(D, dp, i-1)
        else:
            sol = build_solution(D, dp, i-2)
            sol.append(i-1)
    return sol
```

What is the complexity of build\_solution?

$$T(n) = O(n)$$

What is the complexity of hate\_ville?

$$T(n) = O(n)$$

**Exercise:**

write hateville with  $S(n) = O(1)$

# Knapsack

## Problem

Given a set of items, each of them characterized by a **weight** and a **value**, determine which items to include in a collection so that the total weight of the collection is less than or equal to a given "knapsack" **capacity** and the total value (or profit) is as large as possible.

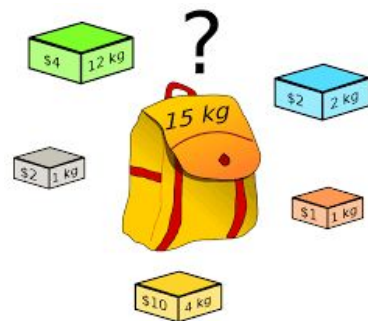
## Input

- List  $w$ , where  $w[i]$  is the weight of the  $i$ -th item
- List  $p$ , where  $p[i]$  is the value (or profit) of the  $i$ -th item
- The capacity  $C$  of the knapsack

## Output

A collection  $S \subseteq \{1, \dots, n\}$  such that:

- Total volume should be smaller or equal than the capacity:  
$$w(S) = \sum_{i \in S} w[i] \leq C$$
- Total profit is maximized:  $p(S) = \sum_{i \in S} p[i]$  is maximal



# Knapsack

## Problem

Given a set of items, each of them characterized by a **weight** and a **value**, determine which items to include in a collection so that the total weight of the collection is less than or equal to a given "knapsack" **capacity** and the total value (or profit) is as large as possible.

Which are the best items for this example?

Item id	1	2	3
Weight	10	4	8
Profit	20	6	12

$$C = 12$$



$$S = \{1\}$$

Item id	1	2	3
Weight	10	4	8
Profit	20	7	15

$$C = 12$$



$$S = \{2,3\}$$

Design an algorithm to solve the Knapsack problem

# Knapsack

## Definition: Sub-problem $DP(i, c)$

Given a knapsack with capacity  $C$  and  $n$  items characterized by weights  $w$  and profits  $p$ , we define  $DP(i, c)$  as the maximal profit we can obtain from the first  $i$  items in a knapsack of capacity  $c$ .

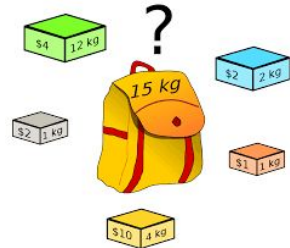
Original problem

The maximal profit of the original problem corresponds to  $DP(n, C)$ .



$$\begin{aligned} i &\leq n \\ c &\leq C \end{aligned}$$

# Knapsack



Let us consider the last item of problem  $DP(i, c)$

- What happens if you don't take it?

$$DP(i, c) = DP(i - 1, c)$$

The capacity and profit do not change

- What happens if you take it?

$$DP(i, c) = DP(i - 1, c - w[i]) + p[i]$$

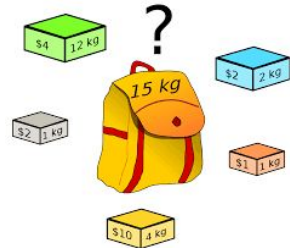
Subtract the weight of the item from the capacity and add its profit

How to select the best solution between the two?

$$DP(i, c) = \max \left\{ DP(i - 1, c), DP(i - 1, c - w[i]) + p[i] \right\}$$



# Knapsack



Let us consider the last item of problem  $DP(i, c)$

- What happens if you don't take it?

$$DP(i, c) = DP(i - 1, c)$$

The capacity and profit do not change

- What happens if you take it?

$$DP(i, c) = DP(i - 1, c - w[i]) + p[i]$$

Subtract the weight of the item from the capacity and add its profit

How to select the best solution between the two?

$$DP(i, c) = \max(DP(i - 1, c - w[i]) + p[i], DP(i - 1, c))$$

# Knapsack



What are the base cases for this recursive definition?

- What happens if you don't have any more items?
- What happens if you don't have any more capacity?
- What happens if your capacity is negative?

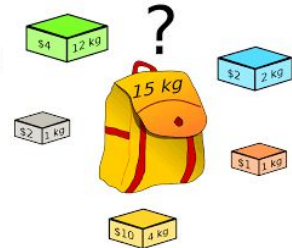
$$DP(i, c) = \begin{cases} 0 & i = 0 \text{ or } c = 0 \\ -\infty & c < 0 \\ \max(DP(i-1, c-w[i]) + p[i], DP(i-1, c)) & \text{otherwise} \end{cases}$$



to enforce NOT  
choosing objects  
that make capacity  
negative

# Knapsack: the code

$$DP(i, c) = \begin{cases} 0 & i = 0 \text{ or } c = 0 \\ -\infty & c < 0 \\ \max(DP(i-1, c-w[i]) + p[i], DP(i-1, c)) & \text{otherwise} \end{cases}$$



```
import numpy as np
import math

def knapsack(w, p, C):
    n = len(w)
    DP = np.zeros((n + 1, C + 1))
    for i in range(1, n+1):
        for c in range(1, C+1):
            not_taken = DP[i-1][c]
            if w[i-1] > c:
                taken = -math.inf
            else:
                taken = DP[i-1][c - w[i-1]] + p[i-1]
            DP[i][c] = max(not_taken, taken)
    #print(DP)
    return DP[n][C]
```

```
w = [4,2,3,4]
p = [10,7,8,6]
C = 9
print(knapsack(w,p,C))
```

25.0

	c									
i	0	1	2	3	4	5	6	7	8	9
0	0	0	0	0	0	0	0	0	0	0
1	0	0	0	0	10	10	10	10	10	10
2	0	0	7	7	10	10	17	17	17	17
3	0	0	7	8	10	15	17	18	18	25
4	0	0	7	8	10	15	17	18	18	25

# Knapsack: the code

What is the computational complexity of function `knapsack()`?

$$T(n) = O(nC)$$

Is it a polynomial algorithm?

No, this is an example of **pseudo-polynomial** algorithm, because  $C$  is not the size of the input, is the input. Thus we need  $k = \log C$  bits to represent it, and thus complexity is equal to:

$$T(n) = O(n2^k)$$

```
import numpy as np
import math

def knapsack(w, p, C):
    n = len(w)
    DP = np.zeros((n + 1, C + 1))
    for i in range(1, n+1):
        for c in range(1, C+1):
            not_taken = DP[i-1][c]
            if w[i-1] > c:
                taken = -math.inf
            else:
                taken = DP[i-1][c - w[i-1]] + p[i-1]
            DP[i][c] = max(not_taken, taken)
    #print(DP)
    return DP[n][C]
```

	c									
i	0	1	2	3	4	5	6	7	8	9
0	0	0	0	0	0	0	0	0	0	0
1	0	0	0	0	10	10	10	10	10	10
2	0	0	7	7	10	10	17	17	17	17
3	0	0	7	8	10	15	17	18	18	25
4	0	0	7	8	10	15	17	18	18	25

# Knapsack: the code

What is the computational complexity of function `knapsack()`?

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$$T(n) = O(n2^k)$$

```
import numpy as np
import math

def knapsack(w, p, C):
    n = len(w)
    DP = np.zeros((n + 1, C + 1))
    for i in range(1, n+1):
        for c in range(1, C+1):
            not_taken = DP[i-1][c]
            if w[i-1] > c:
                taken = -math.inf
            else:
                taken = DP[i-1][c - w[i-1]] + p[i-1]
            DP[i][c] = max(not_taken, taken)
    #print(DP)
    return DP[n][C]
```

	c									
i	0	1	2	3	4	5	6	7	8	9
0	0	0	0	0	0	0	0	0	0	0
1	0	0	0	0	10	10	10	10	10	10
2	0	0	7	7	10	10	17	17	17	17
3	0	0	7	8	10	15	17	18	18	25
4	0	0	7	8	10	15	17	18	18	25

# Memoization

## Note

Not all elements of the table are actually needed to solve our problem.

```
import numpy as np
import math

def knapsack(w, p, C):
    n = len(w)
    DP = np.zeros((n + 1, C + 1))
    for i in range(1, n+1):
        for c in range(1, C+1):
            not_taken = DP[i-1][c]
            if w[i-1] > c:
                taken = -math.inf
            else:
                taken = DP[i-1][c - w[i-1]] + p[i-1]

            DP[i][c] = max(not_taken, taken)
    #print(DP)
    return DP[n][C]
```

```
w = [4,2,3,4]
p = [10,7,8,6]
C = 9
print(knapsack(w,p,C))
```

	c									
i	0	1	2	3	4	5	6	7	8	9
0	0	0	0	0	0	0	0	0	0	0
1	0	0	0	0	10	10	10	10	10	10
2	0	0	7	7	10	10	17	17	17	17
3	0	0	7	8	10	15	17	18	18	25
4	0	0	7	8	10	15	17	18	18	25

$c - w[n-1]$   
 $= 9 - 4$

# Memoization

## Note

Not all elements of the table are actually needed to solve our problem.

```
import numpy as np
import math
```

```
def knapsack(w, p, C):
    n = len(w)
    DP = np.zeros((n + 1, C + 1))
    for i in range(1, n+1):
        for c in range(1, C+1):
            not_taken = DP[i-1][c]
            if w[i-1] > c:
                taken = -math.inf
            else:
                taken = DP[i-1][c - w[i-1]] + p[i-1]

            DP[i][c] = max(not_taken, taken)
    #print(DP)
    return DP[n][C]
```

```
w = [4,2,3,4]
p = [10,7,8,6]
C = 9
print(knapsack(w,p,C))
```

	c									
i	0	1	2	3	4	5	6	7	8	9
0	0	0	0	0	0	0	0	0	0	0
1	0	0	0	0	10	10	10	10	10	10
2	0	0	7	7	10	10	17	17	17	17
3	0	0	7	8	10	15	17	18	18	25
4	0	0	7	8	10	15	17	18	18	25

$5 - w[n-2]$   
 $= 5 - 3$

$9 - w[n-2]$   
 $= 9 - 3$

# Memoization

## Memoization (annotazione)

Programming techniques that merge the tabular aspect of dynamic programming with the top-down approach of divide-et-impera

- Whenever we need to solve a sub-problem, we check in the table first, to see if the problem has already been solved in the past
  - If not, we compute the result and store it in the table
- We use the result stored in the table
- In any case, each subproblem is compute only once as in the bottom-up version



# Memoized-knapsack using a table (np array)

```
import numpy as np
import math

def knapsack_mem(w, p, C):
    n = len(w)
    DP = -np.ones((n+1, C+1))
    return knapsackRec(w, p, DP, n, C)

def knapsackRec(w, p, DP, i, c):
    if c < 0:
        return -math.inf
    if i == 0 or c == 0:
        #DP[i][c] = 0
        return 0
    if DP[i][c] < 0:
        #the solution has not been computed already!
        not_taken = knapsackRec(w,p,DP,i-1,c)
        taken = knapsackRec(w,p,DP,i-1,c-w[i-1]) + p[i-1]
        DP[i][c] = max(not_taken, taken)
    return DP[i][c]
```

```
w = [4,2,3,4]
p = [10,7,8,6]
C = 9
print(knapsack_mem(w,p,C))
```

	c									
i	0	1	2	3	4	5	6	7	8	9
0	-1	-1	-1	-1	-1	-1	-1	-1	-1	-1
1	-1	-1	0	0	10	10	10	10	-1	10
2	-1	-1	7	-1	-1	10	17	-1	-1	17
3	-1	-1	-1	-1	-1	15	-1	-1	-1	25
4	-1	-1	-1	-1	-1	-1	-1	-1	-1	25

**Note:** remember that NOT all elements of the table are actually needed to solve our problem.

# Memoized-knapsack using a table (np array)

	c									
i	0	1	2	3	4	5	6	7	8	9
0	-1	-1	-1	-1	-1	-1	-1	-1	-1	-1
1	-1	-1	0	0	10	10	10	10	-1	10
2	-1	-1	7	-1	-1	10	17	-1	-1	17
3	-1	-1	-1	-1	-1	15	-1	-1	-1	25
4	-1	-1	-1	-1	-1	-1	-1	-1	-1	25

- Table initialization
  - The initialization cost is equal to  $O(nC)$
  - Applied in this way, there is no advantage in using the memoization technique
  - The real advantage is that it makes easy to translate the recursive formula into an algorithm
- Using a dictionary
  - Instead of using a table, we may use a dictionary
  - No pre-initialization is necessary
  - The execution cost is equal to  $O(\min(2^n, nC))$

# Memoized-knapsack using a dictionary

```
import math

def knapsack_mem(w, p, C):
    n = len(w)
    DP = dict()
    return knapsackRec(w, p, DP, n, C)

def knapsackRec(w, p, DP, i, c):
    if c < 0:
        return -math.inf
    if i == 0 or c == 0:
        #DP[(i,c)] = 0
        return 0
    if (i,c) not in DP:
        #the solution has not been computed already!
        not_taken = knapsackRec(w,p,DP,i-1,c)
        taken = knapsackRec(w,p,DP,i-1,c-w[i-1]) + p[i-1]
        DP[(i,c)] = max(not_taken, taken)
    return DP[(i,c)]
```

```
w = [4,2,3,4]
p = [10,7,8,6]
C = 9
print(knapsack_mem(w,p,C))
```

Dictionary:

{(1, 9): 10, (1, 7): 10, (2, 9): 17, (1, 6): 10, (1, 4): 10, (2, 6): 17, (3, 9): 25, (1, 5): 10, (1, 3): 0, (2, 5): 10, (1, 2): 0, (2, 2): 7, (3, 5): 15, (4, 9): 25}

# Longest common **subsequence**

## Definition: subsequence

- A sequence  $P$  is a **subsequence** of  $T$  if  $P$  is obtained from  $T$  by removing one or more of its elements
- Alternatively:  $P$  is defined as the subset of indexes in  $\{0, \dots, n-1\}$  describing the elements of  $T$  that are also in  $P$
- The remaining elements are listed in the same order, although they do not need to be contiguous

## Examples

- $T = \text{"AAAATTGA"}$
- $P = \text{"AAATA"}$

## Note

The empty sequence is a subsequence of every sequence

# Longest common subsequence (LCS)

## Definition: common subsequence

- Given two sequence  $P$  and  $T$ , a sequence  $Z$  is a **common subsequence** of  $P$  and  $T$  if  $Z$  is subsequence of both  $P$  and  $T$
- We write  $Z \in \mathcal{CS}(P, T)$

P: ACAATACT

T: ATCAGTC

Z: ACA

## Definition: longest common subsequence

- Given two sequence  $P$  and  $T$ , a sequence  $Z$  is a **longest common subsequence** of  $P$  and  $T$  if  $Z \in \mathcal{CS}(P, T)$  and there is no other sequence  $W$  such that  $W$  is longer than  $Z$  ( $|W| > |Z|$ ) and  $W$  is common subsequence of  $P$  and  $T$  ( $W \in \mathcal{CS}(P, T)$ ).
- We write  $Z \in \mathcal{LCS}(P, T)$

P: ACAATACT

T: ATCAGTC

Z: ACATC

# Longest common subsequence (LCS)

## Problem: LCS

Given two sequences  $P$  and  $T$  of length  $n$  and  $m$ , respectively, find either the length of the longest common subsequence or one of the longest common subsequences.

Examples:

P: ACAATAT  
T: ATCAGTC  
Out: 4

P: ATATATATAT  
T: ATGATAAT  
Out: 6

P: AAAAA  
T: CTGCTC  
Out: 0

P: ATATATATAT  
T: ATGATAAT  
Out: 6

## Any ideas?

Naive idea (“brute force”): generate all subsequences of  $P$ , all subsequences of  $T$ , compute the common ones and return the longest. **Problem:** all subsequences of a sequence with length  $n$  are  $2^n$  (think about strings of  $n$  0 or 1...)

# Longest common subsequence (LCS)

## Problem: LCS

Given two sequences  $P$  and  $T$  of length  $n$  and  $m$ , respectively, find either the length of the longest common subsequence or one of the longest common subsequences.

Examples:

P: ACAATAT  
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Out: 6

P: AAAAA  
T: CTGCTC  
Out: 0

P: ATATATATAT  
T: ATGATAAT  
Out: 6

## Any ideas?

Naive idea (“brute force”): generate all subsequences of  $P$ , all subsequences of  $T$ , compute the common ones and return the longest. **Problem:** all subsequences of a sequence with length  $n$  are  $2^n$  (think about strings of  $n$  0 or 1...)

# Longest common subsequence (LCS)

## Prefix

Given a sequence  $P$  composed of the characters  $p_1p_2 \dots p_n$ ,  $P(i)$  will denote the **prefix** of  $P$  given by the first  $i$  characters, i.e.:

$$P(i) = p_1p_2 \dots p_i$$

## Examples

- $P = \text{ABDCCAABD}$
- $P(0) = \emptyset$  (empty subsequence)
- $P(3) = \text{ABD}$
- $P(6) = \text{ABDCCA}$



# Longest common subsequence (LCS)

## Goal

Given two sequences  $P$  and  $T$  of length  $n$  and  $m$ , write a recursive formula  $DP(i, j)$  that returns the **length** of the LCS of the prefixes  $P(i)$  and  $T(j)$ .

$$DP(i, j) = \begin{cases} ? & \text{Base case} \\ ? & \text{Recursive case} \end{cases}$$

# Longest common subsequence (LCS)

Goal

Given two sequences  $P$  and  $T$  of length  $n$  and  $m$ , write a recursive formula  $DP(i, j)$  that returns the **length** of the LCS of the prefixes  $P(i)$  and  $T(j)$ .

$$DP(i, j) = \begin{cases} ? & \text{Base case} \\ ? & \text{Recursive case} \end{cases}$$

## Case 1:

Consider the two prefixes  $P(i)$  and  $T(j)$  such that their last characters are the same:  $p_i = t_j$ .

How would you compute  $DF[i, j]$ ?

$$DF[i, j] = DF[i - 1, j - 1] + 1$$

Ex.

P : TACGC**A**

T : ATCG**A**



A is part of the LCS

# Longest common subsequence (LCS)

Goal

Given two sequences  $P$  and  $T$  of length  $n$  and  $m$ , write a recursive formula  $DP(i, j)$  that returns the **length** of the LCS of the prefixes  $P(i)$  and  $T(j)$ .

$$DP(i, j) = \begin{cases} ? & \text{Base case} \\ ? & \text{Recursive case} \end{cases}$$

## Case 2:

Consider the two prefixes  $P(i)$  and  $T(j)$  such that their last characters are different:  $p_i \neq t_j$ .

How would you compute  $DF[i, j]$ ?

Hint: either  $p_i$  or  $t_j$  are useless for the LCS.

$$DF[i, j] = \max(DF[i - 1, j], DF[i, j - 1])$$

Ex.

P : TACGC  
T : ATCG



either C or G is useless (removing C seems the most reasonable choice)

# Longest common subsequence (LCS)

Goal

Given two sequences  $P$  and  $T$  of length  $n$  and  $m$ , write a recursive formula  $DP(i, j)$  that returns the **length** of the LCS of the prefixes  $P(i)$  and  $T(j)$ .

$$DP(i, j) = \begin{cases} ? & \text{Base case} \\ ? & \text{Recursive case} \end{cases}$$

## Base cases:

What if  $i = 0$  or  $j = 0$ ?

$$DP[i, j] = 0$$

Ex.

P : TACGC

T:



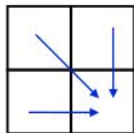
length of LCS is 0

Putting it all together:

$$DP(i, j) = \begin{cases} 0 & i = 0 \text{ or } j = 0 \\ DP(i - 1, j - 1) + 1 & i > 0 \text{ and } j > 0 \text{ and } p_i = t_j \\ \max\{DP(i - 1, j), DP(i, j - 1)\} & i > 0 \text{ and } j > 0 \text{ and } p_i \neq t_j \end{cases}$$

# LCS: example

P: CTCTGT  
T: ACGGCT



arrows specify  
where the values  
come from

$$DP(i, j) = \begin{cases} 0 & i = 0 \text{ or } j = 0 \\ DP(i-1, j-1) + 1 & i > 0 \text{ and } j > 0 \text{ and } p_i = t_j \\ \max\{DP(i-1, j), DP(i, j-1)\} & i > 0 \text{ and } j > 0 \text{ and } p_i \neq t_j \end{cases}$$

i \ j		0	1	2	3	4	5	6
			C	T	C	T	G	T
0		0	0	0	0	0	0	0
1	A	0	↓ 0	↓ 0	↓ 0	↓ 0	↓ 0	↓ 0
2	C	0	↖ 1	→ 1	↖ 1	→ 1	→ 1	→ 1
3	G	0	↓ 1	↓ 1	↓ 1	↓ 1	↖ 2	→ 2
4	G	0	↓ 1	↓ 1	↓ 1	↓ 1	↖ 2	↓ 2
5	C	0	↖ 1	↓ 1	↖ 2	→ 2	↓ 2	↓ 2
6	T	0	↓ 1	↖ 2	↓ 2	↖ 3	→ 3	↖ 3



result

# Memoized LCS

$$DP(i, j) = \begin{cases} 0 & i = 0 \text{ or } j = 0 \\ DP(i-1, j-1) + 1 & i > 0 \text{ and } j > 0 \text{ and } p_i = t_j \\ \max\{DP(i-1, j), DP(i, j-1)\} & i > 0 \text{ and } j > 0 \text{ and } p_i \neq t_j \end{cases}$$

```
def LCSrec(P,T,DP, i,j):
    if i == 0 or j == 0:
        return 0
    if (i,j) not in DP:
        if P[i-1] == T[j-1]:
            DP[(i,j)] = LCSrec(P,T, DP, i-1, j-1) + 1
        else:
            DP[(i,j)] = max(LCSrec(P,T, DP, i-1, j),
                           LCSrec(P,T, DP, i, j-1))
    return DP[(i,j)]

def LCS(P,T):
    n = len(P)
    m = len(T)
    D = dict()
    LCSrec(P,T,D,n,m)
    print(D)
    return D[(n,m)]
```

```
T = "CTCTGT"
P = "ACGGCT"

print(LCS(P,T))
```

		j						
		0	1	2	3	4	5	6
i			C	T	C	T	G	T
	0		0	0	0	0	0	0
1	A	0	↓0	↓0	↓0	↓0	↓0	↓0
2	C	0	↖1	→1	↖1	→1	→1	→1
3	G	0	↓1	↓1	↓1	↓1	↖2	→2
4	G	0	↓1	↓1	↓1	↓1	↖2	↓2
5	C	0	↖1	↓1	↖2	→2	↓2	↓2
6	T	0	↓1	↖2	↓2	↖3	→3	↖3

DP: {(1, 1): 0, (1, 2): 0, (1, 3): 0, (1, 4): 0, (2, 3): 1, (2, 4): 1, (2, 1): 1, (2, 2): 1, (3, 1): 1, (3, 2): 1, (3, 3): 1, (3, 4): 1, (4, 5): 2, (4, 1): 1, (4, 2): 1, (4, 3): 1, (4, 4): 1, (5, 3): 2, (5, 4): 2, (5, 5): 2, (6, 6): 3}

Result:

3

# Memoized LCS: where is my string?

```
def subsequence(DP,P,T,i,j):
    if i == 0 or j == 0:
        return []
    if P[i-1] == T[j-1]:
        S = subsequence(DP,P, T, i-1, j-1)
        S.append(P[i-1]) #or T[j-1]
        return S
    else:
        if DP[(i-1,j)] > DP[(i,j-1)]:
            return subsequence(DP,P,T, i-1, j)
        else:
            return subsequence(DP,P,T,i, j-1)

def LCSrec(P,T,DP, i,j):
    if i == 0 or j == 0:
        return 0
    if (i,j) not in DP:
        if P[i-1] == T[j-1]:
            DP[(i,j)] = LCSrec(P,T, DP, i-1, j-1) + 1
        else:
            DP[(i,j)] = max(LCSrec(P,T, DP, i-1,j),
                           LCSrec(P,T, DP, i,j-1))

    return DP[(i,j)]

def LCS(P,T):
    n = len(P)
    m = len(T)
    D = dict()
    for i in range(n+1):
        D[(0,i)] = 0
    for j in range(m+1):
        D[(j,0)] = 0

    LCSrec(P,T,D,n,m)
    return subsequence(D,P,T,n,m)
```

```
T = "CTCTGT"
P = "ACGGCT"

print("LCS: {}".format("".join(LCS(P,T))))

LCS: CCT
```

i \ j		0	1	2	3	4	5	6
			C	T	C	T	G	T
0		0	0	0	0	0	0	0
1	A	0	↓ 0	↓ 0	↓ 0	↓ 0	↓ 0	↓ 0
2	C	0	↖ 1	→ 1	↖ 1	→ 1	→ 1	→ 1
3	G	0	↓ 1	↓ 1	↓ 1	↓ 1	↖ 2	→ 2
4	G	0	↓ 1	↓ 1	↓ 1	↓ 1	↖ 2	↓ 2
5	C	0	↖ 1	↓ 1	↖ 2	→ 2	↓ 2	↓ 2
6	T	0	↓ 1	↖ 2	↓ 2	↖ 3	→ 3	↖ 3

travel back up to  
build the  
substring...

# Longest common subsequence (LCS)

What is the computational complexity of `subsequence()`?

$$T(n) = O(m + n)$$

we “consume” one element of either of the two sequences at each step

What is the computational complexity of `LCS()`?

$$T(n) = O(mn)$$

that is the size of the matrix



# Automatic memoization in python

```
import time

def fib(n):
    if n<2:
        return 1
    return fib(n-1) + fib(n-2)

s=time.time()
print(fib(45))
e=time.time()
print("elapsed time: {:.3}s".format(e-s))
```

1836311903  
elapsed time: 3.04e+02s

*##Automatic memoization in python*

from functools import wraps

```
def memo(func):
    cache = {} # Stored subproblem solutions
    @wraps(func) # Make wrap look like func
    def wrap(*args): # The memoized wrapper
        if args not in cache: # Not already computed?
            cache[args] = func(*args) # Compute & cache the solution
        return cache[args] # Return the cached solution
    return wrap # Return the wrapper
```

@memo

```
def fib(n):
    if n<2:
        return 1
    return fib(n-1) + fib(n-2)
```

```
s=time.time()
print(fib(45))
e=time.time()
print("elapsed time: {:.3}s".format(e-s))
```

1836311903  
elapsed time: 0.000436s

# Exercise: palindrome

A string is said palindrome if it reads indetically if read from left to right and from right to left.

Write an algorithm that returns the minimum number of characters to be inserted in a string to make it palindrome.

For example, input: "casacca":

- $n = 7$  caratteri: "casaccaACCASAC"
- $n = 6$  caratteri: "casaccaCCASAC"
- $n = 3$  caratteri: "casaccaSAC"
- $n = 2$  caratteri: "ACcasacca"

Please note that characters may be inserted in the middle of the string as well; for example, "anta"  $\rightarrow$  "antNa".

# Exercise: palindrome

- If the string  $s = as'a$  has the same initial and final character “a” then:

$$f(s) = f(s')$$

- If the string  $s = as'b$  has two different initial and final characters:
  - add a  $b$  character at the beginning or a  $a$  character at the end
  - count this added character as an insertion
  - consider the two subproblems given by the first and last (equal) characters removed
  - choose the minim


$$f(s) = \min\{f(as'), f(s'b)\} + 1$$

## Exercise: palindrome

$$DP(i, j) = \begin{cases} 0 & i \geq j \\ DP(i + 1, j - 1) & i < j \wedge s[i] = s[j] \\ \min(DP(i + 1, j), DP(i, j - 1) + 1) & i < j \wedge s[i] \neq s[j] \end{cases}$$

# Shortest common supersequence

- $A$  is a supersequence of  $B$  if  $B$  is a subsequence of  $A$
- The shortest common supersequence (SCS) of  $P, T$  is the shortest sequence that is supersequence of both  $P, T$
- SCS problems appear in the sequencing projects, when the genome of the individual is broken into several pieces. These pieces are sequenced individually, and then the whole genome is constructed as the shortest common supersequence of the sequences
- For the more general problem of finding a string  $S$  which is a supersequence of a set of strings  $S_1, S_2, \dots, S_k$ , the problem is NP-Complete



problems for which there is no polynomial time algorithms known. IF there was, then all NP problems would be solved polynomially