

with 300 MB/s bandwidth we can observe that:

$$t_w = t_{w, \text{parallel}} + t_{w, \text{serial}}$$

time to access disk

stream of new worker all executing this work

Our performance of the stream

$$IS \frac{t_w}{nw} + t_{w, \text{serial}} \cdot nw$$

$$\left(\begin{array}{c} \bigcirc \\ \bigcirc \\ \bigcirc \end{array} \right) \frac{t_w}{nw}$$

questo non può essere diviso me

le devo moltiplicare perché ogni le è non in parallel perché sono risorse condivise

Another version of the algorithm to solve this problem

→ we could use open mp

① pragma section {

pragma section

process image 1

pragma section

process image 2

}

② parallel for is the easiest way

③ pragma omp parallel {

do_img (int i)
for (3 i argc)

to create the thread pool

{ #pragma omp task
do_img(i)

for vs task

I have to do:

pragma omp parallel

pragma omp single

for () {

pragma omp task

do_img

Se Vogliamo usare i pattern?

See.cpp → sequential version

for (i = 3 → argc) {

do_img(i) {

read

process

write result

}

}

Trasformare con pattern usando GRPPI.

Per eseguirlo con GRPPI:

pipeline (

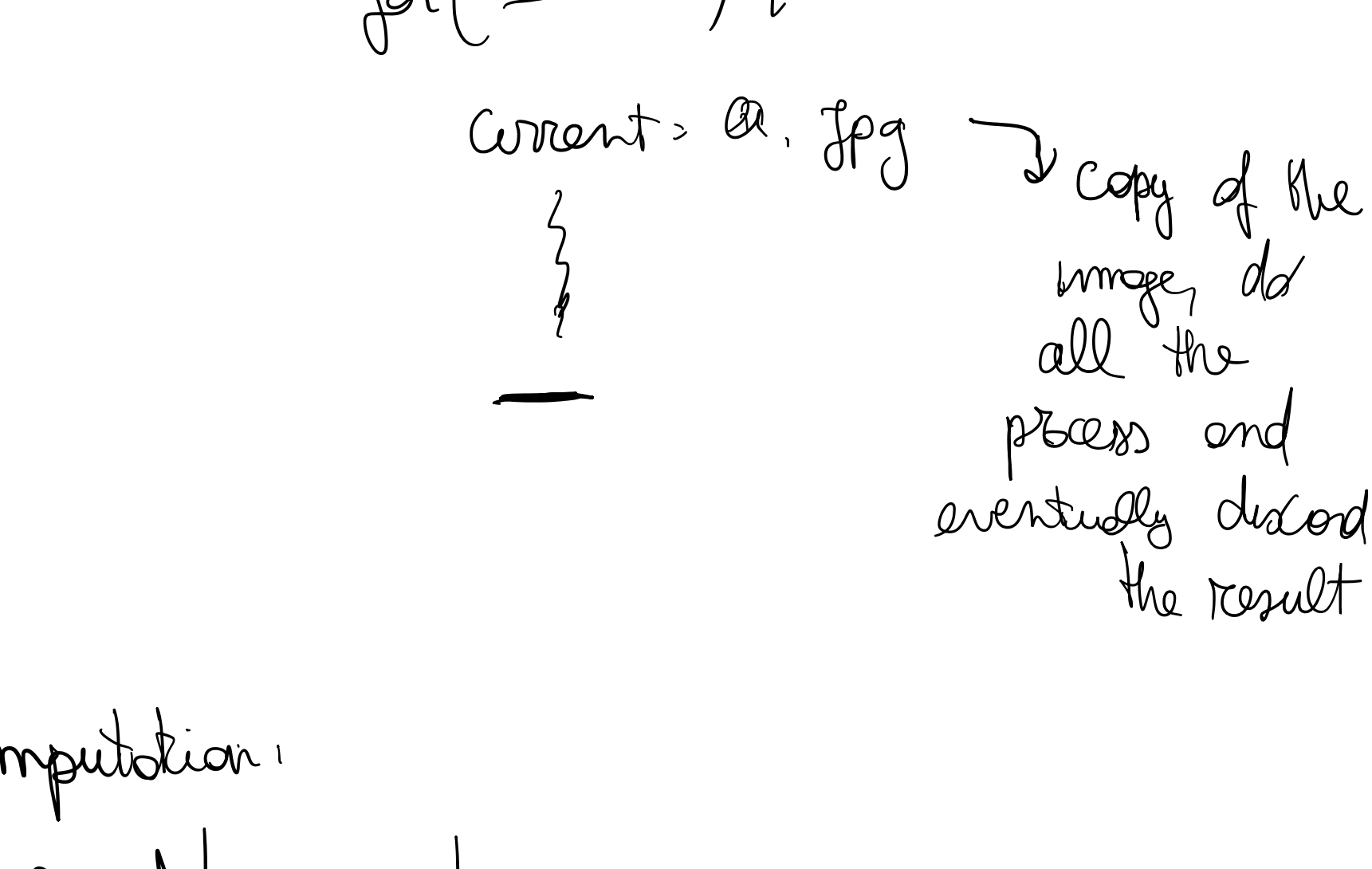
stage 1

stage 2 ← farm

stage 3 ← persona fa in modo che restituisca il numero di elementi già processati.

[GRPPI pipeline and the second stage] is a farm

BYPASS HDD BOTTLENECK



Computation:

load | process | save

farm (pipe (load, process, save))

⇓

pipe (L, P, S)

oppure:

farm (com (L, P, S))

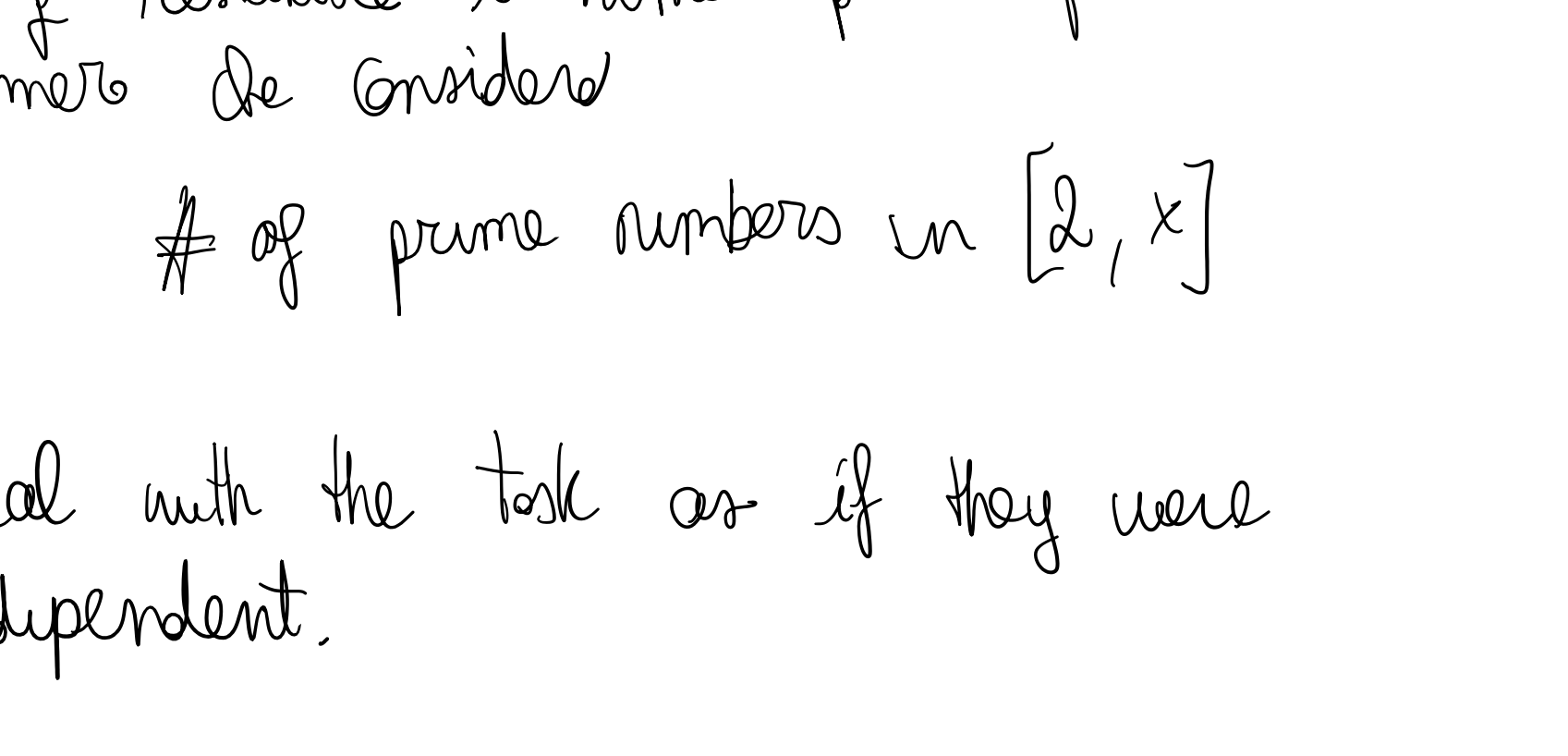
farm (pipe (load, farm (proc), save))

Key point to evaluate this stuffs:

→ You have to evaluate which is the good solution

timer(??)

ASSIGNMENT 2



Le f restituisce i numeri primi prima del numero da considerare

of prime numbers in $[2, x]$

Deal with the task as if they were independent.

parameter:

→ Number of items in shared data structure

Use • Native thread

• Grppi / openmp (done overlo fatto con Native thread).