Lempel-Ziv Compression: Design

Lucais Sanderson

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1 Introduction

The purpose of this program is to implement the LZ Compression algorithm. This algorithm saves space by taking sequences of characters repeated frequently and uses pointers to reference those repeated sequences.

Specifically, we will use a Trie ADT to organize the sequences of characters and a Word Table ADT for the decompression so we know what each code refers to each prefix.

2 Pseudocode

2.1 Compression

See section A.1 in the assignment document.

```
compress(infile, outfile):
   root = TrieCreate()
   curr_node = root
   prev_node = NULL
   curr_sym = 0
   prev_sym = 0
   # START_CODE is 1
   next_code = START_CODE
   while read_sym(infile, &curr_sym) is true:
       next_node = trie_step(curr_node, curr_sym)
       if next_node is not NULL:
          prev_node = curr_node
          curr_node = next_node
          write_pair(outfile, curr_node.code, curr_sym, bit-length(next_code))
           curr_node.children[curr_sym] = trie_node_create(next_node)
           curr_node = root
          next_code = next_code + 1
       if next_code is MAX_CODE:
          trie_reset(root)
          curr_node = root
          next_code = START_CODE
       prev_sym = curr_sym
   if curr_node is not root:
       write_pair(outfile, prev_code.code, prev_sym, bit-length(next_code))
       next_code = (next_code + 1) % MAX_CODE
   write_pair(outfile, STOP_CODE, 0, bit-length(next_code))
```

2.2 Trie ADT

```
# Below is the structure for a Trie Node
TrieNode:
   # ALPHABET is 256, number of ASCII characters
   TrieNode *children[ALPHABET]
   code
TrieNode *trie_node_create(code):
   # might not need to allocate?
   node = allocate(sizeof(TrieNode))
   node.code = code
   for i in ALPHABET:
       node.children[i] = NULL
   return node
trie_node_delete(TrieNode n):
   free(n)
TrieNode *trie_create():
   root = allocate(sizeof(TrieNode))
   root = TrieNode
   root.code = EMPTY_CODE
   return root
trie_reset(TrieNode root):
   for i in ALPHABET:
       trie_delete(root.children[i])
       root.children[i] = NULL
trie_delete(TrieNode n):
   # base case
   if n == NULL:
       terminate
   # recursive case
       trie_delete(next child)
   # delete current node
   trie_node_delete(n)
   n = NULL
TrieNode *trie_step(TrieNode n, sym):
   for n's children:
       if child == sym:
          return n.children[sym]
   # NULL if invalid child
```

2.3 Decompression

See section A.2 in the assignment document.

```
decompression(infile, outfile):
    table = wt_create()
    curr_sym = 0
    curr_node = 0
    next_code = START_CODE
    while read_pair(infile, &curr_code, &curr_sym, bit-length(next_code)) is true:
        table[next_code] = word_append_sym(table[curr_code], curr_sym)
        write_word(outfile, table[next_code])
        next_code = next_code + 1
        if next_code is MAX_CODE
            wt_reset(table)
            next_code = START_CODE
        flush_words(outfile)
```

2.4 Work Table ADT

```
# Word structure
Word:
   # array of "symbols", a word
   # length of the array syms
# WordTable is an array of Words
# (sorry for C syntax, not sure how else to describe it)
typedef Word * WordTable
Word *word_create(*syms, len):
   Word w = allocate(sizeof(Word))
   # might need to copy memory instead of just setting a pointer to same location
   w.syms = syms
   w.len = len
   return w
Word *word_append_sym(Word *w, sym):
   alloc new block for new word, nw
   copy contents of w to nw
   nw.length = nw.length + 1
   add one more byte to nw.syms[]
   add sym to nw.syms[]
   return nw
```

word_delete(Word *w):

```
if w isn't empty word:
       free w->syms
   free(w)
WordTable *wt_create():
   WordTable wt = allocate(sizeof(Word *) * MAX_CODE)
   # empty word. can be null, just accommodate in other functions.
   wt[EMPTY_CODE].syms = NULL
   return wt
wt_reset(WordTables *wt):
   # start at index 2 so as not to delete empty code
   for i in range(2(START_CODE), MAX_CODE):
       word_delete(wt[i])
       wt[i] = NULL
void wt_delete(WordTable *wt):
   word_delete(wt[EMPTY_CODE])
   free(wt[EMPTY_CODE])
2.5 I/O
FileHeader:
   # magic number is OxBAADBAAC ... 32 bit unsigned
   # protection bits for infile to use when writing outfile
   protection
int read_bytes(int infile, uint8_t *buf, int to_read):
   # bytes to keep track of how many bytes are read
   # and which part or the array to continue reading
   bytes = 0
   # temp is for the current bytes read from read()
   while temp != 0 and bytes < to_read:</pre>
       temp = read (to_read - bytes) bytes into buf from infile
       bytes += temp
   return bytes
int write_bytes(int outfile, uint7_t *buf, int to_write):
   counter = 0
   # temp is for the current bytes written from write()
   while temp != 0 and counter < to_write:</pre>
       temp = write (to_write - counter) bytes into buf from infile
       counter += temp
   return counter
```

```
read_header(int infile, FileHeader *header):
   read in sizeof(FileHeader) into header
   swap endianness if byte order not little endian
   verify magic number
write_header(int outfile, FileHeader *header):
   write sizeof(FileHeader) to output file supplied from header
bool read_sym(int infile, uint8_t *sym):
   # terminal, char_buf, and char_index are static, global variables within i/o
   # if at beginning or end of buffer, fill it back up
   if char_index not 0 or BLOCK:
       terminal = read_bytes(infile, char_buf, BLOCK)
       char_index = 0
   if less than BLOCK bytes are read (terminal) and index is terminal:
       # (both could be 0 for example)
       return false
   # set character and increment
   *sym = char_buf(char_index)
   return true
write_pair(int outfile, uint16_t code, uint8_t sym, int bitlen):
   # first write code with bitlen bits
   for i = 0, bitlen:
       # if current bit in code is a 1, set bit_buf as such
       if bitmask code with 1 << i:
          bit_buf[bit_index / 8] |= (1 << (bit_index / 8))
       bit_index = bit_index + 1
       # if buffer full, flush and reset
       if buffer full:
          flush_pairs(outfile)
          bit_index = 0
          reset_buffer(bit_buf)
   # now do the same but for sym, which will always be 8 bits
   for i = 0, 8:
       # if current bit in code is a 1, set bit_buf as such
       if bitmask code with 1 << i:
          bit_buf[bit_index / 8] |= (1 << (bit_index / 8))
       bit_index = bit_index + 1
       # if buffer full, flush and reset
       if buffer full:
          flush_pairs(outfile)
          bit_index = 0
          reset_buffer(bit_buf)
flush_pairs(int outfile):
   write_bytes(outfile, bit_buf, bit_index / 8)
# essentially inverse of write_pair()
bool read_pair(int infile, uint16_t *code, uint8_t *sym, int bitlen):
```

code = 0

```
sym = 0
   for i = 0, bitlen:
       if buffer full or empty (initial state):
           read_bytes(infile, bit_buf, BLOCK)
           bit_index = 0
       if bitmask current bytes in bit_buf with 1 << (bit_index/8):</pre>
           code |= 1 << i
       bit_index++
   for i = 0, 8:
       if buffer full or empty (initial state):
           read_bytes(infile, bit_buf, BLOCK)
           bit_index = 0
       if bitmask current bytes in bit_buf with 1 << (bit_index/8):</pre>
           code |= 1 << i
       bit_index++
   if code is STOP_CODE:
       return false
   return true
write_word(int outfile, Word *w):
   for i = 0, w->len:
       if at end of buffer:
           flush_words(outfile)
           char_index = 0
       char_buf[char_index] = w->syms[i]
       char_index++
flush_words(int outfile):
   write_bytes(outfile, char_buf, char_index)
```

3 Citations

- Darrel Long from the assignment document.
- Thanks to Audrey Ostrom for this template!