
1. Log Analysis System

Scenario: A system that processes log files and finds the most frequent IP addresses.

```
from collections import Counter
```

```
def most_frequent_ips(logs):  
    ip_counts = Counter(logs) # Count occurrences of each IP  
    return ip_counts.most_common(10) # Get top 10 most common IPs
```

Time Complexity Analysis

- Counter(logs) iterates through the logs list ($O(n)$) and constructs a hash map.
- most_common(10) sorts the frequencies using heapq.nlargest(), which is $O(k \log n)$, where $k=10$.
- Total complexity: $O(n + k \log n) \approx O(n)$ (since k is a small constant).

Commercial Use Case: This approach is widely used in log aggregation services like Splunk, ELK Stack, and Datadog.

2. E-Commerce Price Matching System

Scenario: A price comparison tool that matches product prices from different vendors.

```
def find_matching_prices(store_prices, competitor_prices):
```

```
    store_prices.sort() #  $O(n \log n)$ 
```

```
    competitor_prices.sort() #  $O(m \log m)$ 
```

```
    i, j = 0, 0
```

```
    matches = []
```

```
    while i < len(store_prices) and j < len(competitor_prices):
```

```
        if store_prices[i] == competitor_prices[j]:
```

```
            matches.append(store_prices[i])
```

```
            i += 1
```

```
            j += 1
```

```

elif store_prices[i] < competitor_prices[j]:
    i += 1
else:
    j += 1

```

```

return matches

```

Time Complexity Analysis

- Sorting both lists: $O(n \log n + m \log m)$
- Two-pointer search through sorted lists: $O(n + m)$
- Total complexity: $O(n \log n + m \log m)$

Commercial Use Case: Used in price comparison tools like Google Shopping, Amazon, and Walmart's competitive pricing engines.

3. Fraud Detection in Financial Transactions

Scenario: Detect if a transaction pattern is suspicious (e.g., repeated payments to different accounts in a short time).

```

def detect_fraud(transactions, time_window):
    suspicious = set()

    transactions.sort(key=lambda x: x[1]) # Sort by timestamp (O(n log n))

    for i in range(len(transactions)):
        user, timestamp = transactions[i]
        j = i + 1
        while j < len(transactions) and transactions[j][1] - timestamp <= time_window:
            if transactions[j][0] == user:
                suspicious.add(user)
            j += 1

    return list(suspicious)

```

Time Complexity Analysis

- Sorting transactions: $O(n \log n)$
- Nested while loop: $O(n \log n)$ (on average, assuming a balanced time distribution)
- Total complexity: $O(n \log n)$

Commercial Use Case: This technique is used in fraud detection algorithms for payment services like PayPal, Stripe, and banks.

4. Social Media Feed Ranking

Scenario: A recommendation system that merges user-post interactions based on relevance.

```
import heapq
```

```
def merge_feeds(*feeds):
```

```
    merged_feed = []
```

```
    min_heap = []
```

```
    for feed in feeds:
```

```
        if feed:
```

```
            heapq.heappush(min_heap, (feed[0], 0, feed)) # Store (post, index, feed)
```

```
    while min_heap:
```

```
        post, idx, feed = heapq.heappop(min_heap)
```

```
        merged_feed.append(post)
```

```
        if idx + 1 < len(feed):
```

```
            heapq.heappush(min_heap, (feed[idx + 1], idx + 1, feed))
```

```
    return merged_feed
```

Time Complexity Analysis

- Using a min-heap to merge k sorted lists of n elements:
 - **Heap insertion/removal:** $O(\log k)$

- **Total merges:** $O(nk \log k)$, where n is the average feed length and k is the number of feeds.
- Total complexity: $O(nk \log k)$

Commercial Use Case: Used in personalized social media feeds (e.g., Facebook, Instagram, TikTok).

5. Dynamic Pricing in Online Retail

Scenario: A system that dynamically updates product prices based on demand.

```
import bisect
```

```
class DynamicPricing:
```

```
    def __init__(self):
```

```
        self.prices = []
```

```
    def add_price(self, price):
```

```
        bisect.insort(self.prices, price) #  $O(\log n)$  insertion
```

```
    def get_median_price(self):
```

```
        n = len(self.prices)
```

```
        if n % 2 == 0:
```

```
            return (self.prices[n // 2 - 1] + self.prices[n // 2]) / 2
```

```
        else:
```

```
            return self.prices[n // 2]
```

```
# Usage example
```

```
pricing = DynamicPricing()
```

```
pricing.add_price(100)
```

```
pricing.add_price(200)
```

```
pricing.add_price(150)
```

```
print(pricing.get_median_price()) # 150
```

Time Complexity Analysis

- `bisect.insort()`: $O(\log n)$ per insertion.
- Median retrieval: $O(1)$.
- For n price updates, total complexity: $O(n \log n)$.

Commercial Use Case: Used in airline ticket pricing, Uber's surge pricing, and e-commerce flash sales.

Summary of Complexity Results

Use Case	Complexity
Log analysis (top IPs)	$O(n)$
Price matching	$O(n \log n + m \log m)$
Fraud detection	$O(n \log n)$
Social media feed merge	$O(nk \log k)$
Dynamic pricing	$O(n \log n)$