CyberSecurity: Principle and Practice

BSc Degree in Computer Science 2024-2025

Lesson 12: Intro to Reverse Engineering

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Disclaimer



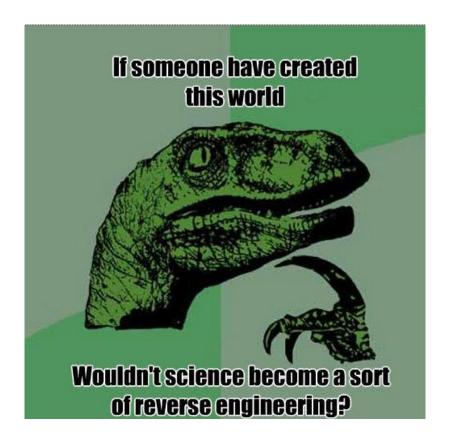
All information presented here has the only purpose of teaching how reverse engineering works.

Use your mad skillz only in CTFs or other situations in which you are legally allowed to do so.

Do not hack the new Playstation. Or maybe do, but be prepared to get legal troubles $\stackrel{\smile}{\smile}$



Reverse Engineering?

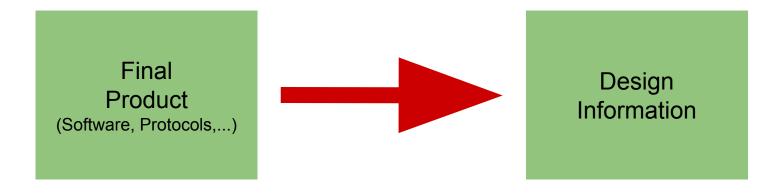


What's Reversing?



"[...] the process of analyzing a subject system to create representations of the system at a higher level of abstraction."

Chikofsky, Cross (1990)



Why?

- Missing or poor documentation
- Opening up proprietary platforms
- Security auditing
- Curiosity

Reversing in CTFs



In reversing challenges you have to understand how a program works, but you don't have its source code.

You typically have to reverse an algorithm (encryption?) to get the flag.

Most of the time, solving a challenge is a bit time consuming but straightforward.

...Unless obfuscation is involved ••

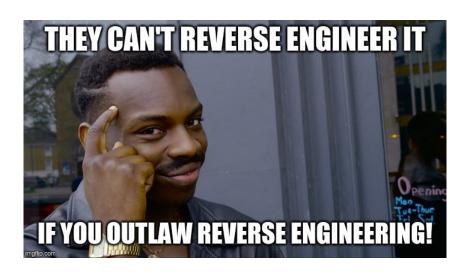
Reversing IRL



Analyze Malwares, remove Ransomwares...

Free Licenses of proprietary software...

A lot of cool stuff, but legally it's a gray area.



(Binary) Software Reverse Engineering



Compiling Software

```
Int main() {
  puts("ILoveCPP");
  return 0;
}
```

COMPILER

00010010010011001

Source Code Binary

(Binary) Software Reverse Engineering



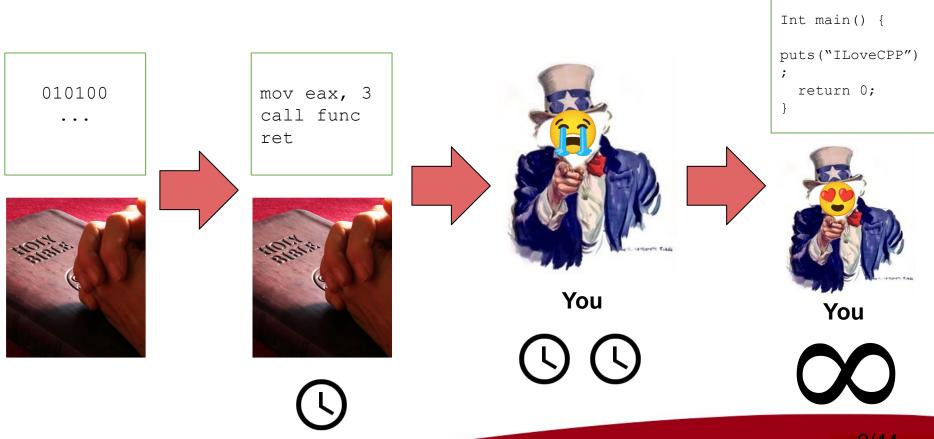
Reversing Software



(Binary) Software Reverse Engineering



Reversing Software (The Truth)



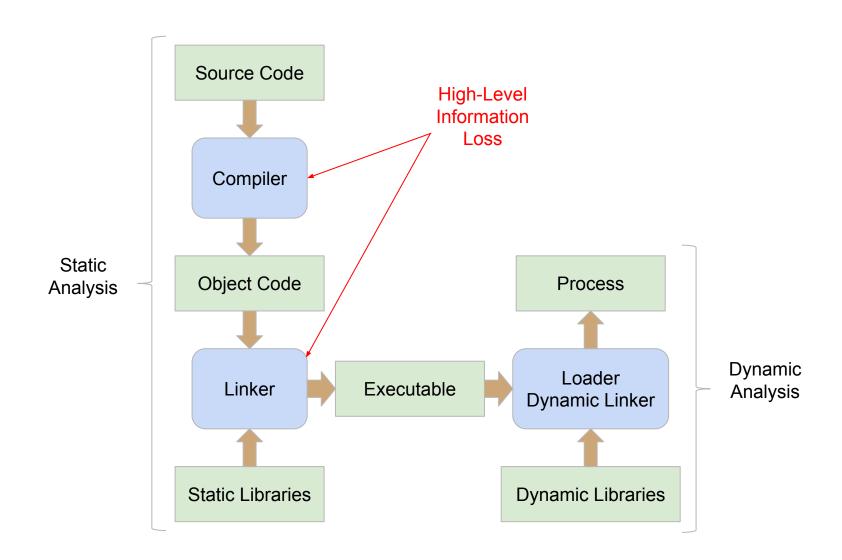
Why is it relevant?



- You don't always have access to source code
- Vulnerability assessment
- Malware analysis
- Pwning
- Algorithm reversing
- Hacking embedded devices
- ...

A program's lifecycle





Executables



- OS-specific format
 - e.g. ELF (*nix), PE (Windows), Mach-O (MacOS, iOS)
- Generally, same format used for programs and libraries
- Made of sections that will be memory-mapped
 - e.g. .text, .(ro)data, .bss
- Specifies imports from dynamic libraries
 - o e.g. GOT/PLT (ELF), IAT (PE)
- Loading methods:
 - Fixed address
 - Relocation
 - Position-independent

Executable and Linkable Format (ELF)



- Introduced in System V Release 4, used by most Unix-like OSes
 - Executables, object code, shared libraries, core dumps
- Designed to be flexible, extensible and cross-platform
- Program headers describe segments (i.e. virtual mappings)
- Section headers describe sections and how to load them into segments
- Supports relocation

Techniques & Tools



- Static analysis doesn't run the executable
 - Disassembly, decompilation
 - Abstract interpretation
 - Symbolic execution
- Dynamic analysis runs the executable
 - Debugging
 - Dynamic binary instrumentation

Disassembler



00010010
Disassembler call func ret

Binary ASM

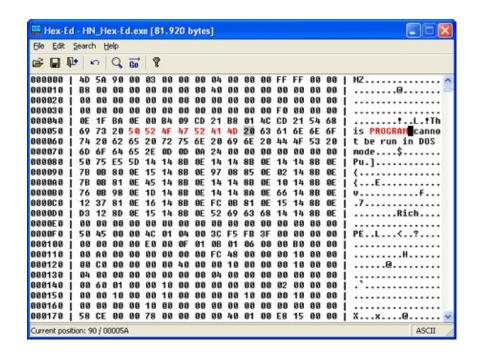
- <u>IDA</u>
 - o GUI, Industry standard
 - Pro Version: \$\$\$\$\$. We will use the <u>free version</u>
- Radare2
 - CLI (experimental GUI <u>here</u>)
 - Opensource
- Ghidra
 - NSA reversing tool (open source!)
- Binary Ninja, Objdump
 - o If you wish...

Hex Editor



- Patch programs
- Inspect file formats
- Change content of files

Many different options here (bless, hexedit, biew, ...)



Other Tools



- Executable information
 - o file, readelf, PEview, ...
- Useful commands
 - o strings, ptrace, ltrace...
- Debuggers
 - o gdb, WinDbg, OllyDbg, Immunity Debugger, qira, ...
- Decompilers

Decompilers



Can't I just use a decompiler?

Can speed up the reversing, but...

Decompiling is (generally) undecidable

Fails in many cases

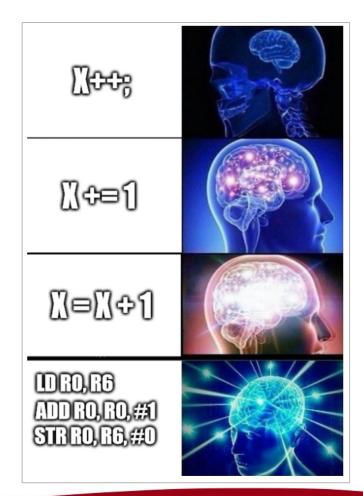
Sometimes you want to work at the ASM level (pwning)

Assembly



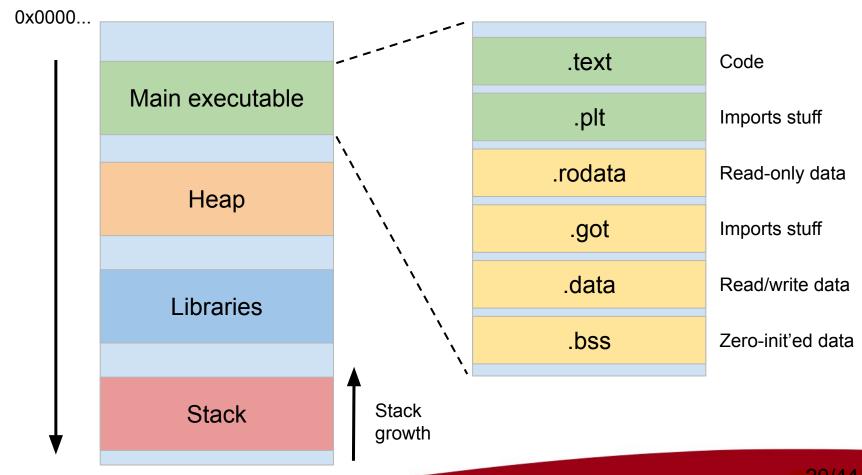
Introduction to x86(_64) ASM

- Your computer probably runs on x86 64
 - x86 still supported
 - 32 bit vs 64 bit
- This is NOT supposed to be a complete ASM lesson





a (Linux) process' memory





x86_64 Registers

General Purpose

Stack Pointer

Base Pointer

Instruction Ptr

64 bit	32 bit 16 bit			
DAY	\$2 E 3 \$2	AX		
RAX	EAX	AH	AL	
DDV	EDV	BX		
RBX EBX	EBA	BH	BL	
DCV	ECY	CX		
RCX	ECX	СН	CL	
RDX EDX	EDV	DX		
	EDA	DH	DL	
RSI	ESI			
RSP	ESP			
RBP	EBP			
RIP	EIP			



Instructions - MOV <dst>, <src>

- Copy <src> into <dst>
- MOV EAX, EBX
 - \circ EAX = EBX
- MOV EAX, 16
 - \circ EAX = 16
- MOV EAX, [ESP+4] [X] = "value at address X"

O EAX = *(ESP+4)

- MOV AL, 'a'
 - \circ AL = 0x61



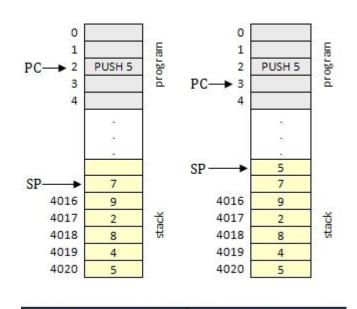
Instructions - LEA <dst>, <src>

- Load Effective Address of <src> into <dst>
- Used to access elements from a buffer/array
- Used to perform simple math operations
- LEA ECX, [EAX+3]
 - ECX = EAX + 3
- LEA EAX, [EBX+2*ESI]
 - EAX = EBX+2*ESI



Instructions - PUSH <src>

- Decrement RSP and put <src> onto the stack (push)
- PUSH EAX
 - ESP -= 4
 - *ESP = (dword) EAX
- PUSH CX
 - ESP -= 2
 - *ESP = (word) CX
- PUSH RAX
 - RSP -= 8
 - *RPS = (qword) RAX



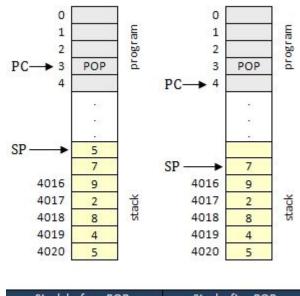
Stack before PUSH 5 Stack after PUSH 5



Instructions - POP <dst>

- <dst> takes the value on top of the stack, RSP gets incremented
- POP EAX

- o ESP += 4
- POP CX
 - CX = *ESP
 - o ESP += 2



Stack before POP

Stack after POP



PUSH/POP example

PUSH EAX
POP EBX
=
MOV EBX, EAX



Instructions - ADD <dst>, <src>

- <dst> += <src>
- ADD EAX, 16
 - o EAX += 16
- ADD AH, AL
 - AH += AL
- ADD ESP, 0x10
 - Remove 16 bytes from the stack



Instructions - SUB <dst>, <src>

- <dst> -= <src>
- SUB EAX, 16
 - EAX -= 16
- SUB AH, AL
 - AH -= AL
- SUB ESP, 0x10
 - Allocate 16 bytes of space on the stack



Flags

- x86 instructions can modify a special register called FLAGS
- FLAGS contains 1-bit flags:
 - Ex: OF, SF, ZF, AF, PF, and CF
- ZF = Zero Flag
 - (set if the result of last operation was zero)
- SF = Sign Flag
 - (set if the result of last operation was negative (dst src < 0)



Flags - Examples

MOV RAX, 555

MOV RAX, 123

SUB RAX, 555

SUB RAX, 555

=>

=>

$$ZF = 1$$

$$ZF = 0$$

$$SF = 0$$



Instructions - CMP <dst>, <src>

- CoMPare
- Perform a SUB but throw away the result
- Used to set flags
- CMP EAX, 13
 - EAX value doesn't change
 - TMP = EAX 13
 - Update the FLAGS according to TMP



Instructions - JMP <dst>

- JuMP to <dst>
- JMP RAX
 - Jump to the address saved in RAX
- JMP 0x1234
 - Jump to address 0x1234



Instructions - Jxx <dst>

- Conditional jump
- Used to control the flow of a program (ex.: IF expressions)
- JZ/JE => jump if ZF = 1
- JNZ/JNE => jump if ZF = 0
- JB, JA => Jump if <dst> Below/Above <src> (unsigned)
- JL, JG => Jump if <dst> Less/Greater than <src> (signed)
- Many others
- See http://unixwiz.net/techtips/x86-jumps.html



Jxx - Example: Password length == 16?

```
MOV RAX, password_length
CMP RAX, 0x10

JZ ok

JMP exit
ok:
...print 'length is correct'...
```



Jxx - Example: Given number >= 11?

```
MOV RAX, integer user input
CMP RAX, 11
JB fail
JMP ok
fail: ...print 'too short'...
ok: ...print 'OK'...
```



Instructions - XOR <dst>, <src>

- Perform a bitwise XOR between <dst> and <src>
- XOR EAX, EBX
 - EAX ^= EBX
- Truth table:

	0	1
0	0	1
1	1	0



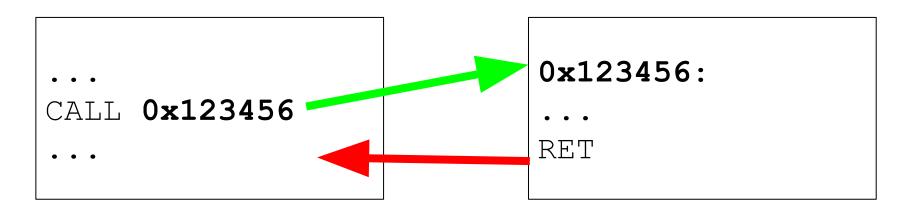
Instructions - CALL <dst>

- CALL a subroutine
- CALL 0x123456
 - Push return address on the stack
 - \circ RIP = 0x123456
- Function parameters passed in many different ways



Instructions - RET

- RETurn from a subroutine
- RET
 - Pop return address from stack
 - Jump to it





How are function parameters passed around?

- On x86, there are many calling conventions
- Sometimes parameters are passed in registers
- Sometimes on the stack
- Return value usually in RAX/EAX
- You should take some time to look at them

https://en.wikipedia.org/wiki/X86 calling conventions



Calling Convention - SystemV AMD64

- Arguments in registers: rdi, rsi, rdx, rcx, r8, r9
- Further args on stack
- Red-zoning: leaf function with frames <= 128 bytes do not need to reserve stack space

```
int callee(int, int, int);
int caller(void)
{
   int ret;

   ret = callee(1, 2, 3);
   ret += 5;
   return ret;
}
```

```
caller:
    ; set up stack frame
    push rbp
    mov rbp, rsp
    ; set up arguments
    mov edi, 1
    mov esi, 2
         edx, 3
    mov
    ; call subroutine 'callee'
    call callee
    ; use subroutine result
    add eax, 5
    : restore old stack frame
    pop rbp
    : return
    ret
```



A very useful instruction

NOP - Single-byte instruction that does nothing

Very useful in patching (to remove CALL, CMP,...)



Now it's Demo Time!

Exercises



- Ex 1 : Can you guess the pin?
- Ex 2: "One of our employees has locked himself out of his account. can you help 'john galt' recover his password? And no snooping around his emails you hear."
- Ex 3: A bomb is going to explode! Defuse the first 4 levels, or go further if you can!

Questions? Feedback? Suggestions?







