Data Integration and Large Scale Analysis

Slides credit: Matthias Boehm - Shafaq Siddiqi

09- Distributed Data-Parallel Computation



Lucas Iacono. PhD. - 2024





Part B

Large-Scale Data
Management & Analysis

• LU3. Cloud Computing

- Cloud Computing Fundamentals[Nov 29]
- Cloud Resource Management and Scheduling [Dec 06]
- Distributed Data Storage[Dec 13]

Part B

Large-Scale Data
Management & Analysis

- LU4. Large-Scale Data Analysis
 - Distributed, Data-ParallelComputation [Dec 20]
 - Distributed Stream Processing[Jan 10]
 - Distributed Machine LearningSystems [Jan 17]

Agenda

- Announcements
- Data-Parallel Collection & Processing
- Data-Parallel DataFrame Operations

Announcements

Announcements

Course Evaluation and Exam

- Course evaluation: 20/02/2025
- Exam date: Feb 07, 3:00pm (90 min written exam)
- Oral Exam for Erasmus Students
 - Schedule available in TeachCenter (23/12/2024)

Recap: Distributed Collections

Logical multi-set (bag) of key-value pairs (unsorted collection)

Different physical representations key-value pairs can be stored in various ways (e.g., database, across files, or in memory).

Easy Distribution via Horizontal Partitioning. Data divided into "chunks" (shards or partitions) based on the keys. Each chunk stored on a different machine (easier to handle large-scale data).

How collections are created: from single file with data or a folder of files (even if they're messy and unsorted).

Key	Value
13:00:01	12.1
14:00:05	16.0
13:00:03	12.5
13:00:05	13.0
14:00:04	15.7
14:00:06	16.3
13:00:00	12.1

Recap: Files and Objects

File: large and continuous block of data saved in a specific format (CSV, Binary, etc.).

Object: like a file, but binary and it comes with metadata (Images on S3)

Recap: Object Storage

- 1. Object Storage (e.g. AWS S3):
 - a. Data stored as objects (data, metadata, and UID).
 - b. Ideal for storing unstructured data like media files, backups, or large datasets.
 - c. Objects of a limited size (e.g., 5TB in AWS S3).

Memory Controller M is C C Core Core Q Core Core Q P I Shared L3 Cache I 1

Nehalem Architecture

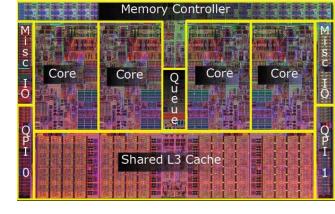
- Integrated Memory Controller: Integrated in chip, -- latency and ++ memory performance.
- Support for DDR3 Memory: Higher memory bandwidth (compared to DDR2).
- QuickPath Interconnect (QPI): High-speed, point-to-point connection (no Front-Side Bus).
- **Enhanced Hyper-Threading:** Each core supports two threads (+++ performance)
- Multi-Core Scalability: 2 to 8 cores per processor (2 threads / core)
- Improved Cache Design: Dedicated L1 and L2 cache p/core shared L3 cache



Michael E.
Thomadakis:
The Architecture of the
Nehalem Processor
and NehalemEP SMP
Platforms, Report,
2010

Nehalem Architecture

- Energy Efficiency: Turbo Boost for dynamic clock speed adjustments.
- Advanced Manufacturing Process: Higher transistor density and better efficiency.
- Integrated Graphics (in later models): Some models included integrated GPUs.
- Foundation for Modern Architectures: Established the groundwork for subsequent Intel architectures like Sandy Bridge and Skylake.

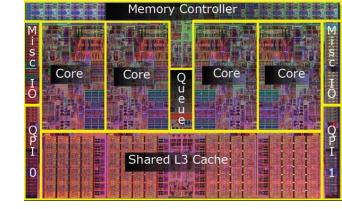




Michael E. Thomadakis: The Architecture of the Nehalem Processor and NehalemEP SMP Platforms, Report, 2010

Nehalem Architecture

- Pipeline
 - o Frontend:
 - Instruction Fetch
 - Pre-Decode
 - Decode CISC 2 uOps (ADD [eax], 5)
 - Load the value from memory.
 - Add 5 to the loaded value.
 - Store the result back to memory.
 - Backend:
 - Rename/Allocate
 - Scheduler
 - Execute

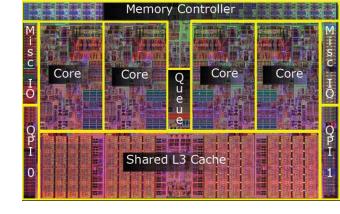




Michael E. Thomadakis: The Architecture of the Nehalem Processor and NehalemEP SMP Platforms, Report, 2010

Nehalem Architecture

- Out-of-Order
 - Instructions are not necessarily executed in the order they appear in the program
- Execution Engine: 4 Inst x Cycle (IPC=4)
- 128-bits Floating-point multiplication
- 128-bits floating-point addition



Goods: Organizing Google's Datasets			
Alon Hallery', File Konv., Natalya F. Nov Sudip Floy , Steen	, Ovisitgher Obton, Necklis Polysolis', in Euljong Whang:		
Youge Resert Ye	out helics of Technology		
alon@recruit.al. (flip, ray obton, repo	lycols, sudpt, swhengi@google.com		
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Michael E.
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The Architecture of the
Nehalem Processor
and NehalemEP SMP
Platforms, Report,
2010

	Single Data	Multiple Data
Single Instruction	SISD (uni-core)	SIMD (vector)
Multiple Instruction	MISD (pipeline)	MIMD (multi-core)

Flynn's Classification

Computer architectures based on how they handle instructions and data.

SISD:

- One task at time one data chunk (e.g. PC running a single program)
- SIMD:
 - One task at time multiple data chunks (e.g. GPUs rendering)
- MISD:
 - Multiple tasks one data chunk (e.g. fault-tolerant computers)
- MIMD:
 - Multiple tasks multiple data chunks (multi-core CPUs 1 Core -> Program)



Michael J. Flynn, Kevin W. Rudd: Parallel Architectures. ACM Comput. Surv. 28(1) 1996

Distributed, Data-Parallel Computation

- Parallel computation of function foo() → single instruction
 - A single function applied to all data items in parallel.
- Collection X of data items (key-value pairs) → multiple data
 - foo() operates on multiple pieces of data (key-value pairs).
- Data parallelism similar to SIMD but more coarse-grained notion of "instruction" and "data" → SPMD (single program, multiple data)

```
Y = X.map(X \rightarrow foo(x)) X = Data Items (e.g. array), .map (operation to each element in X), <math>Y = Output
```



[Frederica Darema: The SPMD Model: Past, Present and Future. PVM/MPI 2001]

SPMD

- Dynamic Work Assignment. Processes can self-schedule, ++
 flexibility & efficiency.
- More General than SIMD. SPMD allows different instruction streams for different data. It can handle more complex tasks.
- Efficient Control. Performed at the application level rather than the OS level (less costly and more efficient than F&J.
- Applications:
 - MPI (Message Passing Interface)
 - PVM (Parallel Virtual Machine)
 - Grid Computing



[Frederica Darema: The SPMD Model : Past, Present and Future. PVM/MPI 2001]

Model	Key Features	Pros	Cons
BSP (Bulk Synchronous)	Global barriers enable synchronization after each phase	+++ Correctness and consistency; simple to implement	Overhead due to waiting at barriers Slow for stragglers
ASP (Asynchronous Parallel)	Processes run independently	Faster execution (no waiting)	Accuracy issues from outdated data
SSP (Stale-Synchronous Parallel)	Controlled staleness allows fastest processes to proceed within a limit	Balances efficiency and consistency; reduces waiting time compared to BSP	Small inaccuracies

Data-Parallel Collection & Processing



Brief Hadoop History

- Google's GFS + MapReduce [ODSI'04] -> Apache Hadoop (2006).
- Apache Hive (SQL), Pig (ETL), Mahout (ML), Giraph (Graph)

Hadoop Architecture / Eco System

- Management (Ambari)
- Coordination / workflows (Zookeeper, Oozie)
- Storage (HDFS)
- Resources (YARN)



Hadoop Ecosytem

- Apache Hive (SQL)
 - 0 What it is:
 - Data warehouse infrastructure built on top of Hadoop.
 - Allows you to query and analyze large datasets stored in Hadoop using a SQL-like language called HiveQL.
 - o Main Purpose:
 - Querying and analysis of big data using familiar SQL syntax.
 - Suitable for batch processing and data summarization.
 - o Use Case:
 - Running SQL queries to analyze log data or generate business reports.







Hadoop Ecosytem

- Apache Pig (ETL)
 - O What it is:
 - **High-level** platform for creating **data processing programs** in Hadoop.
 - Pig language to simplifies the MapReduce jobs writing process.
 - o Main Purpose:
 - ETL operations. Cleaning, transforming, and preparing large datasets for analysis.
 - O Use Case:
 - Processing raw web logs into structured formats for further analysis.



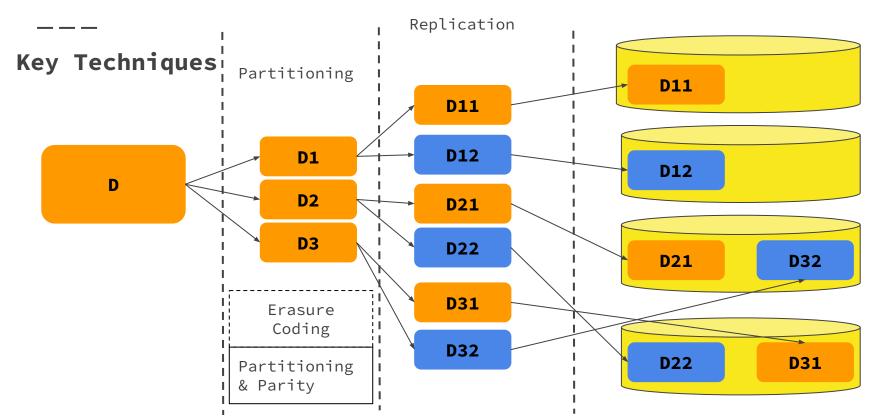




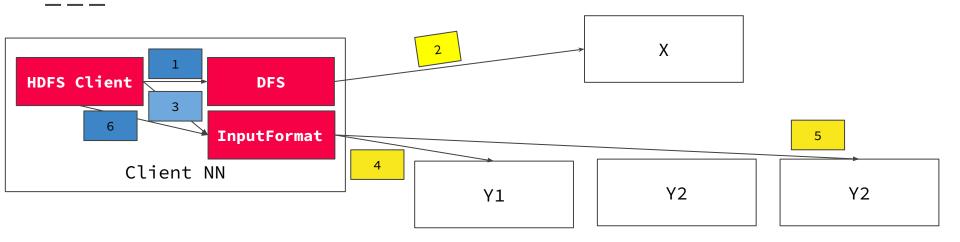
Hadoop Ecosytem

- Apache Mahout (ML)
 - O What it is:
 - A library for building scalable **machine learning** algorithms on top of Hadoop.
 - Focused on distributed or scalable implementations of common ML algorithms.
 - o Main Purpose:
 - Implementing machine learning algorithms like clustering, classification, and recommendation systems on large datasets.
 - o Use Case:
 - Building a recommendation system for an e-commerce platform using collaborative filtering.

HDFS Distribution



Recap: HDFS Read

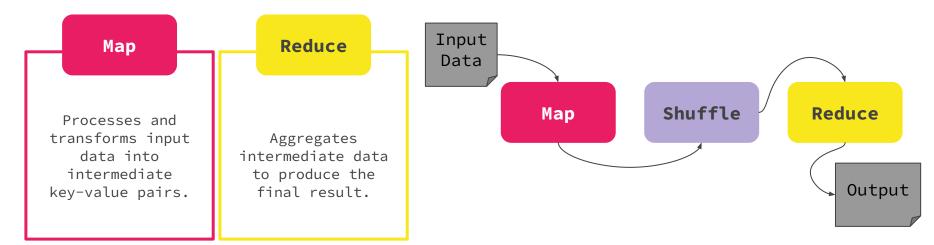


- 1. Open
- 2. Get Block Locations
- 3. Read
- 4. Read
- 5. Read
- 6. Close

MapReduce – Programming Mode

Overview

- MapReduce is a programming model for processing large datasets in parallel, distributed across multiple nodes.
- Developed by Google; popularized by Apache Hadoop.



MapReduce I

Why MapReduce?

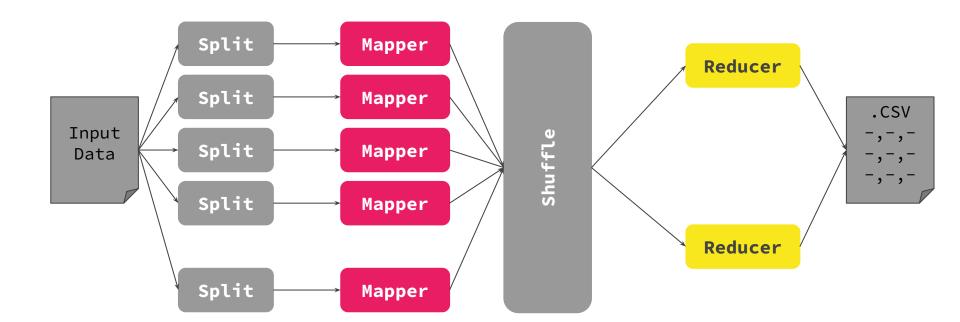
- Handles large-scale data processing efficiently.
- Works on commodity hardware.
- Built-in fault tolerance.
- Suitable for structured, semi-structured, and unstructured data.

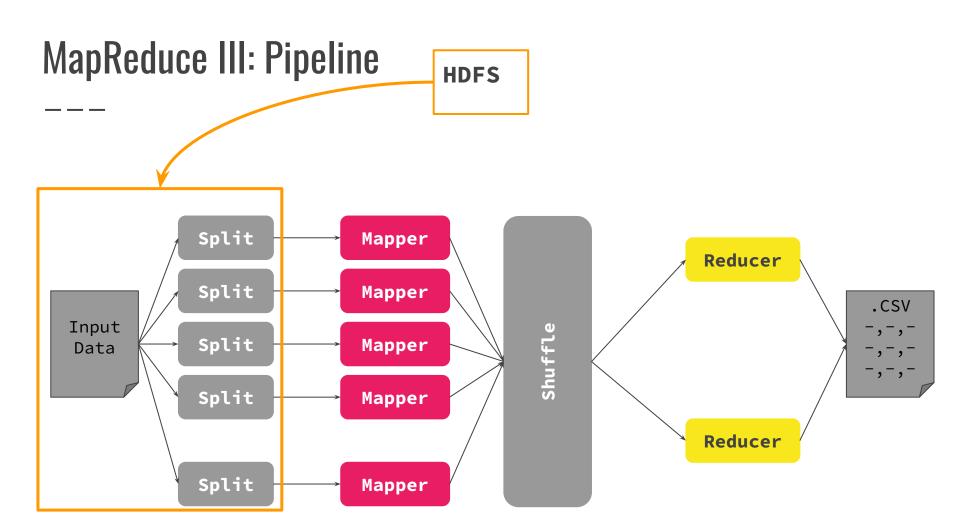
MapReduce II

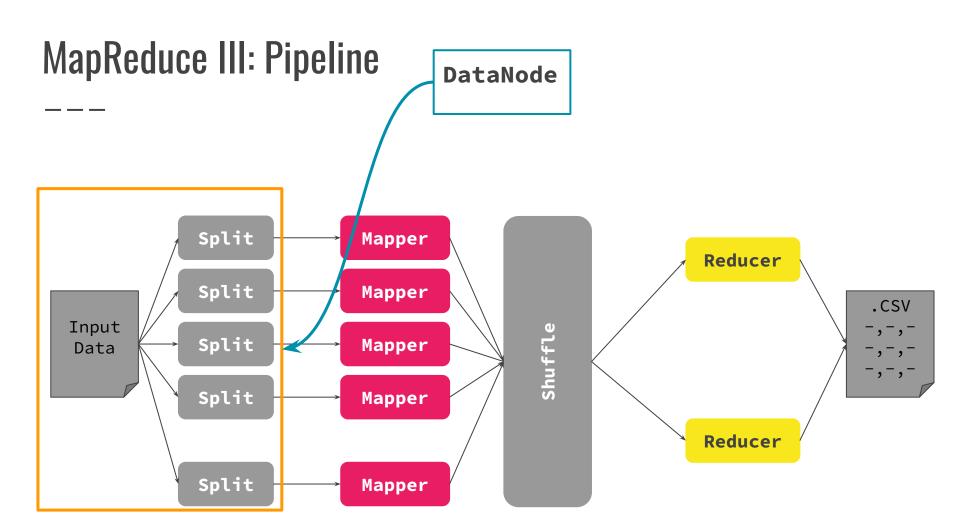
Key Concepts

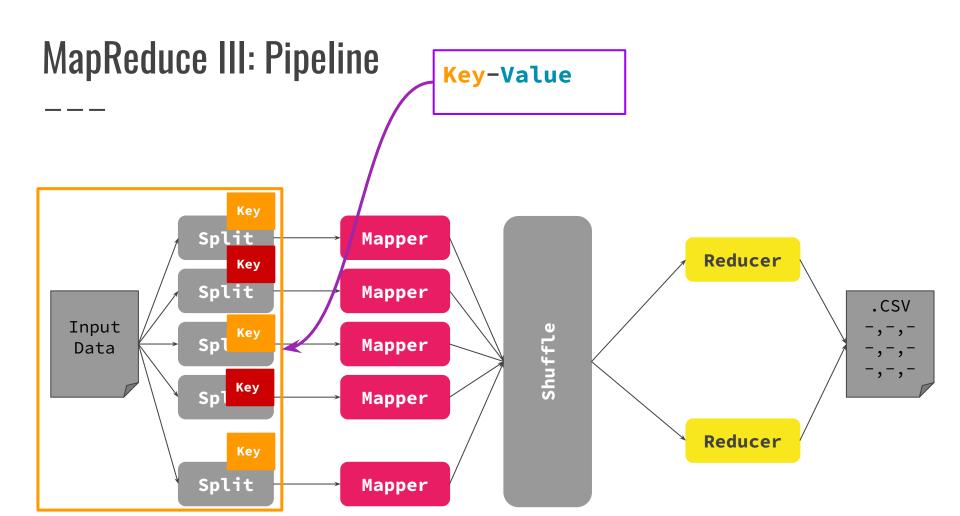
- Distributed Processing: Data is split across multiple nodes for parallel execution.
- Key-Value Pairs: Core data structure in MapReduce.

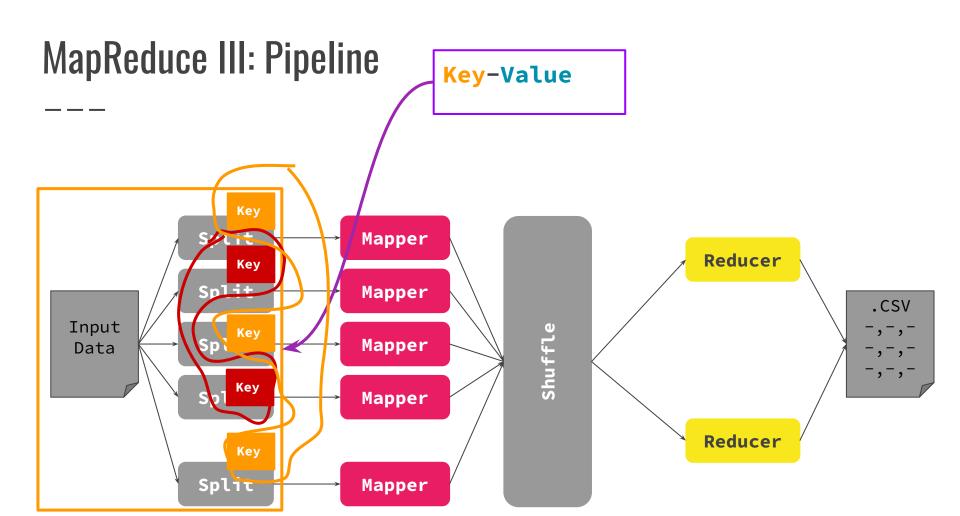
MapReduce III: Pipeline



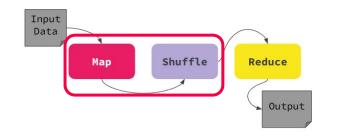


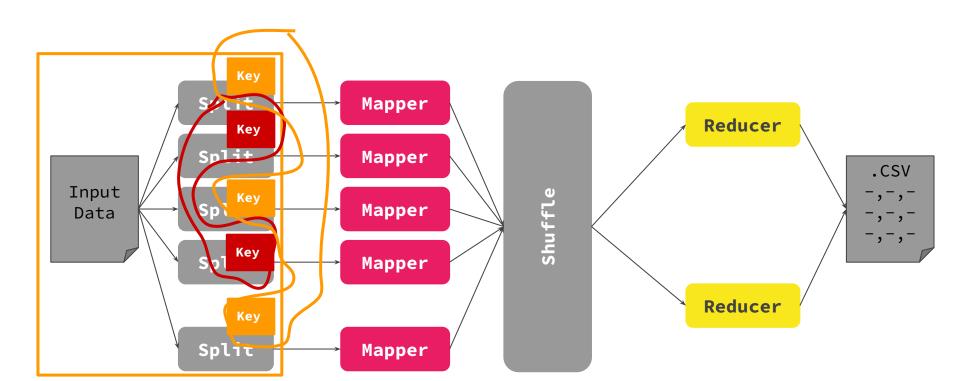






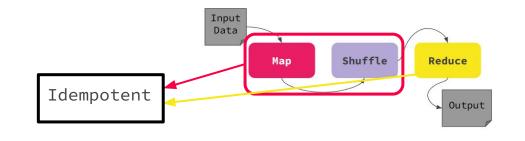
MapReduce III: Pipeline

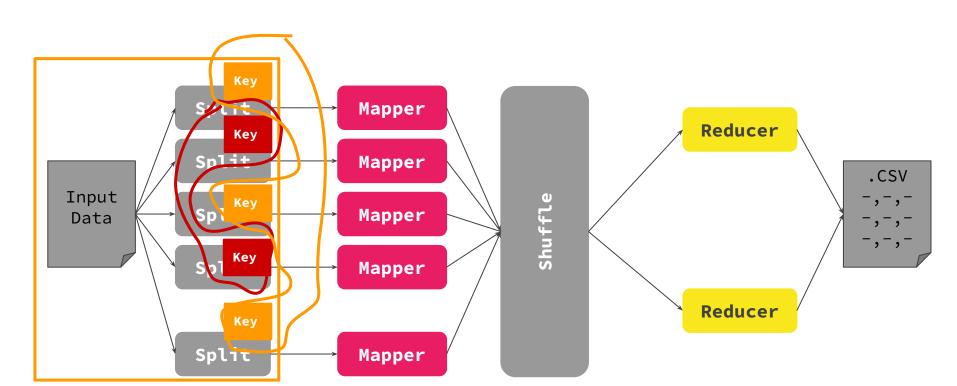




MapReduce III: Pipeline

-





MapReduce IV

```
This - 1
                              is - 1
                              an - 1
               This is
                             apple - 1
This is an
                  an
  apple
                apple
 apple is
  red in
  color
                             apple -
                             is - 1
                apple
                             red - 1
                is red
                             in - 1
                  in
                             color -
                color
```



[Jeffrey Dean, Sanjay Ghemawat: MapReduce: Simplified Data Processing on Large Clusters. OSDI 2004]

MapReduce IV

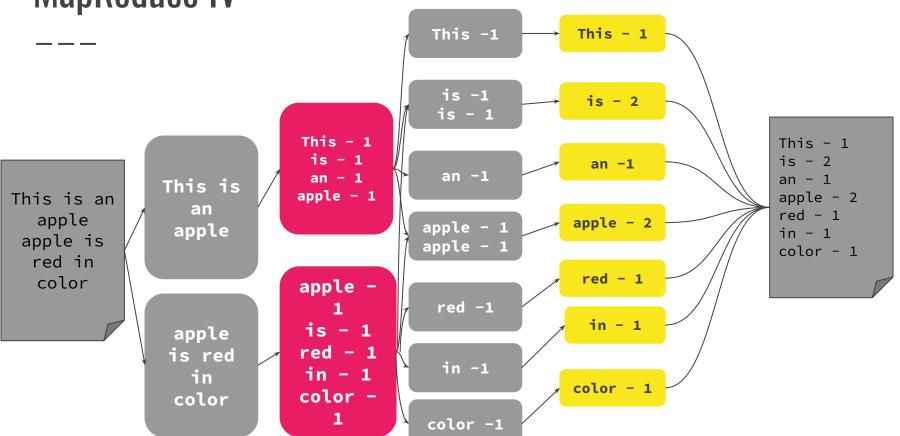
This -1 is -1 is - 1 This - 1 is - 1 an -1 an - 1 This is apple - 1 This is an an apple apple - 1 apple apple is apple - 1 red in color apple red -1 is - <u>1</u> apple red - 1 is red in -1 in - 1 in color color color -1

MapReduce IV

This - 1 This -1 is -1 is - 2 is - 1 This - 1 is - 1 an -1 an -1 an - 1 This is apple - 1 This is an an apple apple - 2 apple - 1 apple apple is apple - 1 red in red - 1apple color red -1 in - 1 is - 1 apple red - 1 is red in -1 in - 1 in color - 1 color color color -1

This - 1

MapReduce IV



MapReduce VI: Hands on Lab

Servers Log



- Use the MapReduce programming model to:
 - Count how many times each page was accessed.
 - Identify the most popular page.
- Calculate the Average Using MapReduce
 - o Given a list [4, 8, 15, 16, 23, 42], compute the average using MapReduce.

MapReduce: Summary (Pros)

- Large-scale processing. Large amounts of data distributed across multiple nodes in a cluster.
- Fault-tolerant. If a node fails, the system can recover and reassign tasks to other nodes.
- User Defined Functions and files. Developers can define their own custom processing logic through UDFs, and the model relies on files to store intermediate and final results.
- Flexibility. Developers can customize processing logic while the system manages distribution and fault recovery automatically.
- Restricted functional APIs. MapReduce relies on a limited set of functional primitives:
 - Map: Transforms input data into key-value pairs.
 - Reduce: Aggregates values associated with the same keys to produce results.
- Implicit parallelism. Developers only need to implement the Map and Reduce functions; the distribution of workload across nodes relies on the system.

MapReduce: Summary (Cons)

- **Performance:** its performance can suffer in complex workloads due to heavy reliance on I/O (writing and reading intermediate data to/from disk).
- Low-level APIs: The API is relatively basic, requiring a lot of manual effort to implement more sophisticated workflows.
- Many different systems: Specialized systems (e.g., Apache Spark, Apache Flink, or distributed database systems) have emerged as alternatives, often being more efficient and user-friendly.





Evolution to Spark (and Flink)

- Spark [HotCloud'10] + Resilent Distributed Data Sets (RDDs) [NSDI'12] → Apache Spark (2014)
- Design 1: Standing executors with in-memory storage:
 - Spark keeps long-running worker processes (executors) active, enabling tasks to run faster by avoiding repeated setup costs.
 - Data is stored in memory whenever possible, minimizing disk I/O for iterative and interactive jobs.
- Design 2: Lazy evaluation:
 - Directed Acyclic Graph of transformations rather than executing them immediately.
 - o Actions (e.g., collect, save, count) trigger DAG's execution, allowing workflow optimization by reordering and combining operations.

Spark History and Architecture



- **Design 3:** Fault tolerance via RDD lineage
 - Data partition lost -> Spark can recompute using **lineage** graph of transformations applied to the data (reliability without heavy replication).

Performance:

- In-memory storage. By keeping intermediate data in memory,
 Spark significantly reduces disk I/O (faster for iterative tasks e.g machine learning).
- Fast job scheduling. Spark's scheduler operates with low overhead, enabling tasks to be scheduled in milliseconds (~100ms), compared to Hadoop's ~10 seconds per job.

Spark History and Architecture



• APIs:

- Richer functional APIs. Wide range of functional operators
 (e.g., map, reduce, filter, groupByKey, flatMap) compared
 to Hadoop -> easier to write complex workflows.
- General computation DAGs. Unlike MapReduce, which forces
 jobs into two rigid phases (map and reduce), Spark supports
 general DAGs for more flexible computation flows.
- **High-level APIs** (**DataFrame/Dataset**). DataFrames and Datasets offer high-level abstractions that simplify working with structured data and enable query optimization.

Spark History and Architecture



- Unified Platform. Multiple workloads into a single platform:
 - Batch processing (similar to MapReduce)
 - Streaming (real-time data)
 - Machine learning (MLlib)
 - Graph processing (GraphX)
 - SQL queries (Spark SQL)

Spark Functionality: Core components



Resilient Distributed Datasets (RDDs):

 Distributed collections of objects (foundation for fault tolerance and parallelism.)

DataFrames and Datasets:

• Higher-level abstractions for structured and semi-structured data (Optimized via Spark's Catalyst engine).

Spark SQL:

Query structured data using SQL.

MLlib:

Machine learning library for scalable algorithms.

GraphX:

Graph processing library.

Spark Functionality: Architecture



Driver Program:

• Defines the application and coordinates tasks.

Cluster Manager:

Allocates resources (YARN, Mesos, Kubernetes).

Executors:

Workers that execute tasks and store data partitions.

DAGs:

• Spark builds a logical execution plan before running tasks.

Spark Functionality: Workflow



- Create RDD/DataFrame: Load data into Spark from HDFS, S3, or other sources.
- Transformations: Apply operations (e.g., map, filter, groupBy).
- Actions: Trigger execution (e.g., collect, save).
- Execution: (a) Splits tasks across nodes, (b) Uses DAG to optimize execution.

Spark: Hands on Lab

Servers Log



- Use the **COLAB** to simulate Spark basic operations
- Let's take a look into Databricks...

Data-Parallel DataFrame Operations

Origins of DataFrames

Recap: Data Preparation Problem

- 80% Argument: 80-90% time for finding, integrating, cleaning data
- Data scientists prefer scripting languages and in-memory libraries

Python DataFrames:

- Python pandas DataFrame for seamless data manipulations (most popular packages/features)
- DataFrame: table with a schema
- Descriptive stats and basic math, reorganization, joins, grouping, windowing
- Limitation: Only in-memory, single-node operations

```
import pandas as pd

df = pd.read_csv('data/tmp1.csv',
index_col=2)

df.head()

df = pd.concat(df, df[['A','C']], axis=0)
```

Spark DataFrames and DataSets

Overview Spark DataFrame

- DataFrame is distributed collection of rows with named/typed columns
- Relational operations (e.g., projection, selection, joins, grouping, aggregation)
- DataSources (e.g., json, jdbc, parquet, hdfs, s3, avro, hbase, csv, cassandra)

DataFrame and Dataset APIs

- DataFrame was introduced as basis for Spark SQL
- DataSets allow more customization and compile-time analysis errors (Spark 2)

DataFrame and Dataset APIs

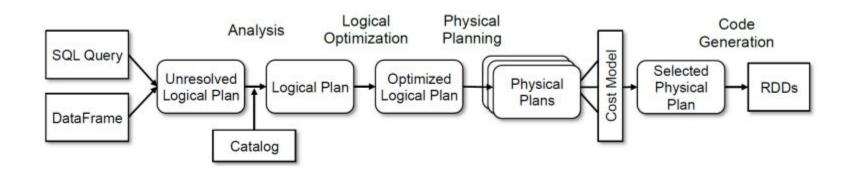
```
logs = spark.read.format("json").open("s3://logs")
logs.groupBy(logs.user_id).agg(sum(logs.time))
.write.format("jdbc").save("jdbc:mysql//...")
```

SparkSQL and DataFrame/Dataset

Overview SparkSQL

- Shark (~2013): academic prototype for SQL on Spark
- SparkSQL (~2015): reimplementation from scratch
- Common IR and compilation of SQL and DataFrame operations

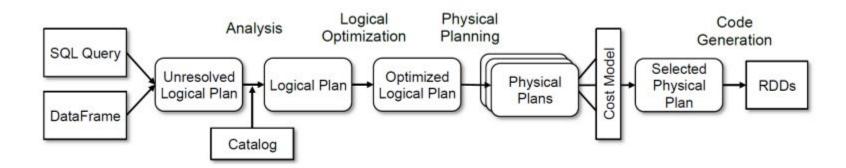
Catalyst: Query Planning



SparkSQL and DataFrame/Dataset

Performance features

- 1. Whole-stage code generation via Janino
- 2. Off-heap memory (sun.misc.Unsafe) for caching and certain operations
- 3. Pushdown of selection, projection, joins into data sources (+ join ordering)



DASK

Overview Dask

- Multi-threaded and distributed operations for arrays, bags, and dataframes
- dask.array: list of numpy n-dim arrays
- dask.dataframe: list of pandas data frames
- dask.bag:unordered list of tuples (second order functions)
- Local and distributed schedulers: threads, processes, YARN, Kubernetes, containers,

HPC, and cloud, GPUs

Execution

- Lazy evaluation
- Limitation: requires static size inference
- Triggered via compute()

```
import dask.array as da
x = da.random.random( (10000,10000),
    chunks=(1000,1000))
y = x + x.T
y.persist() # cache in memory
z = y[::2, 5000:].mean(axis=1) #colMeans ret
= z.compute() # returns NumPy array
```

Summary and Q&A

Summary and Q&A

- Summary and Q&A
 - Motivation and Terminology
 - o Data-Parallel Collection Processing
 - Data-Parallel DataFrame Operations
- Next Lectures
 - Distributed Stream Processing [Jan 10]

Vielen Dank!