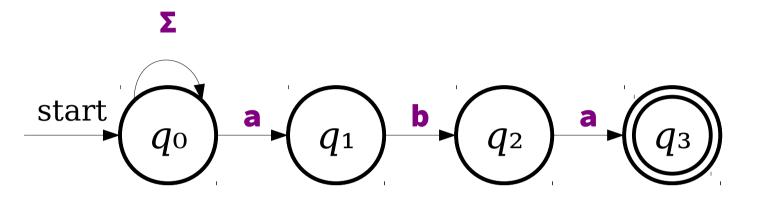
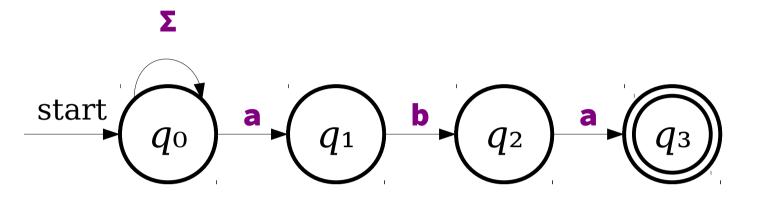
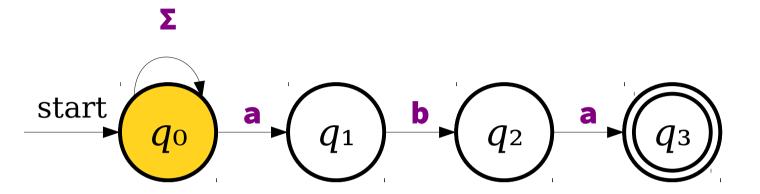
Intuiting Nondeterminism

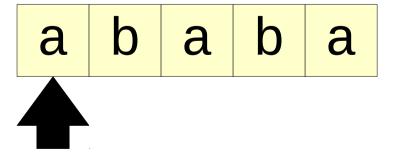
- Nondeterministic machines are a serious departure from physical computers. How can we build up an intuition for them?
- There are two particularly useful frameworks for interpreting nondeterminism:
 - Perfect guessing
 - Massive parallelism

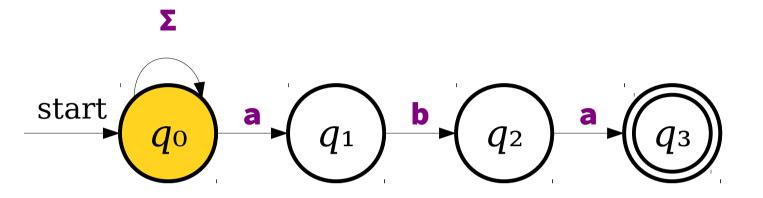


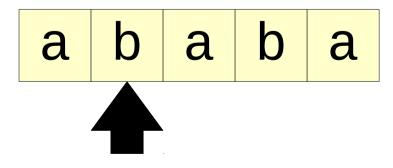


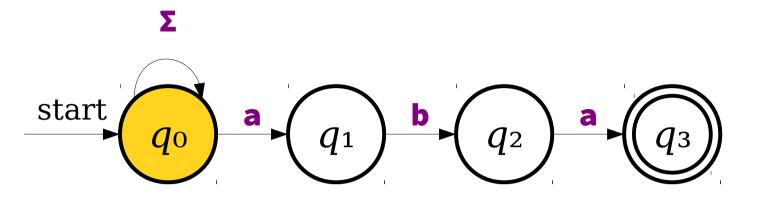
a b a b a

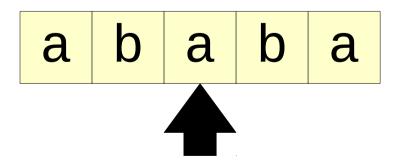


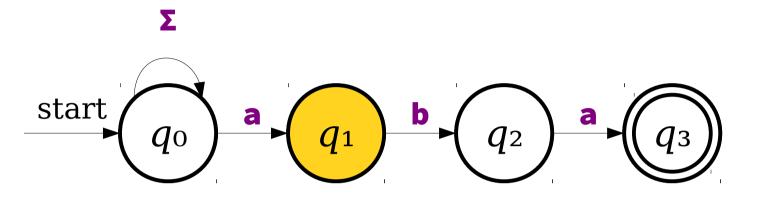


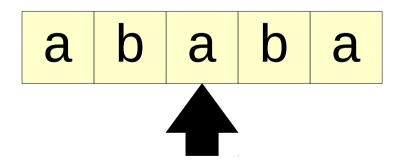


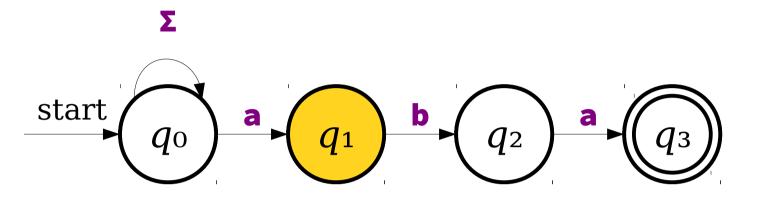


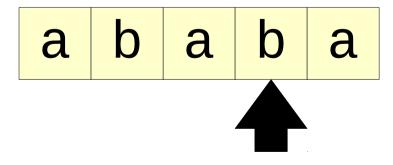


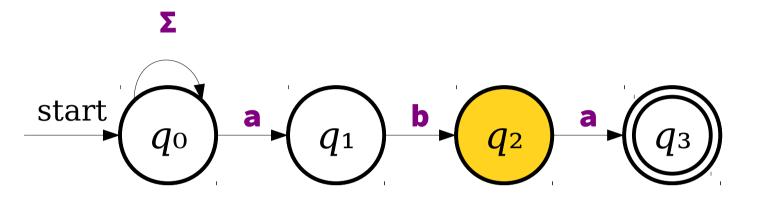


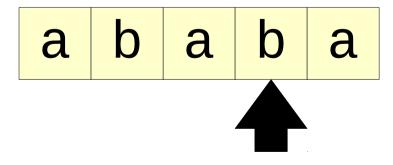


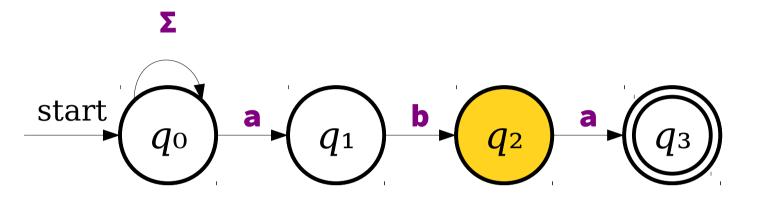


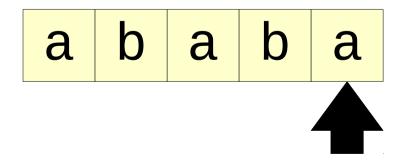


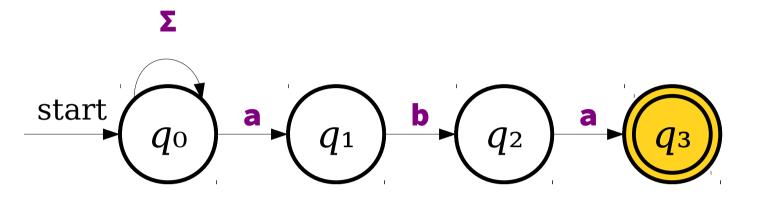


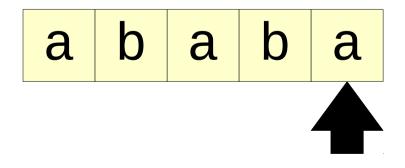


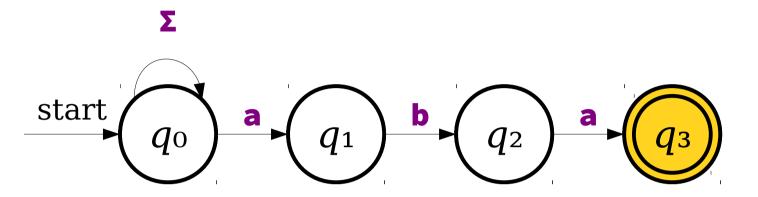






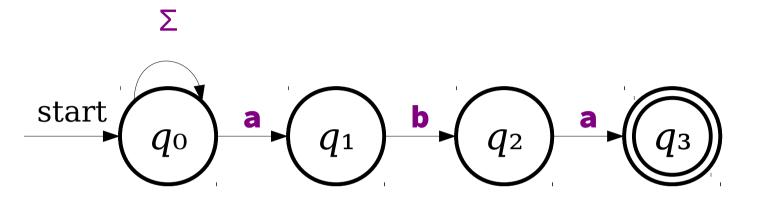


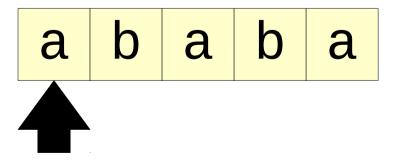


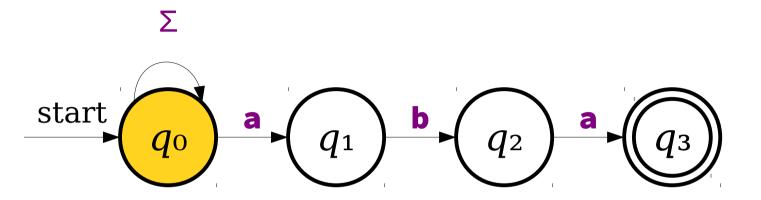


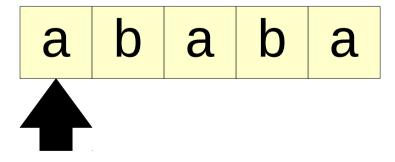
a b a b a

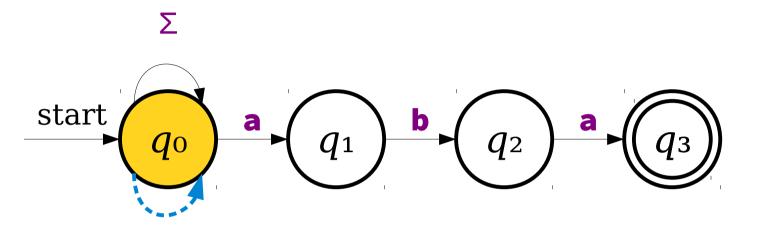
- We can view nondeterministic machines as having *Magic Superpowers* that enable them to guess choices that lead to an accepting state.
 - If there is at least one choice that leads to an accepting state, the machine will guess it.
 - If there are no choices, the machine guesses any one of the wrong guesses.
- No known physical analog for this style of computation – this is totally new!

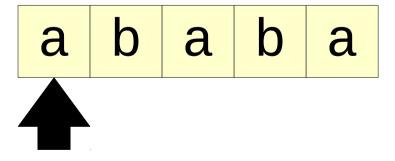


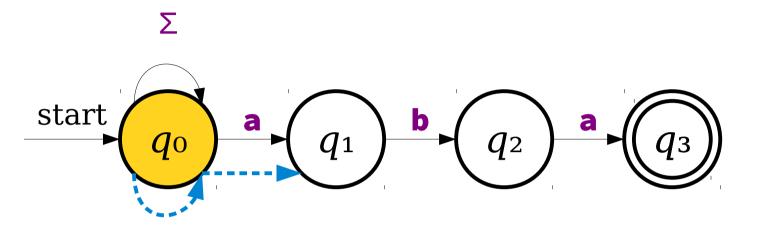


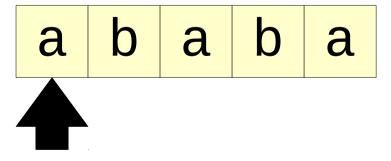


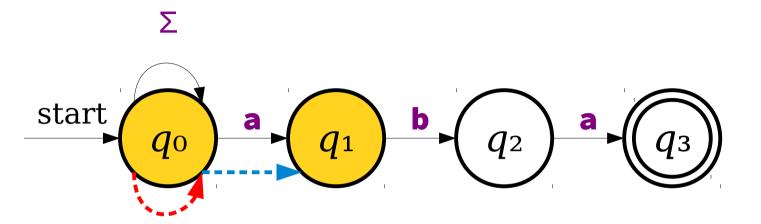


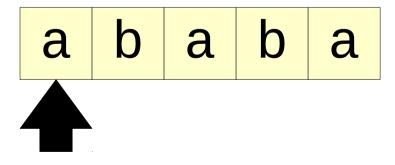


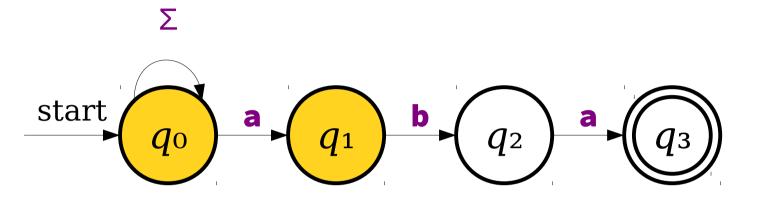


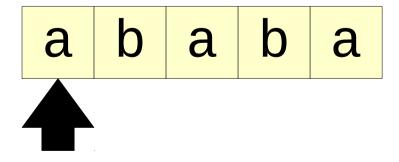


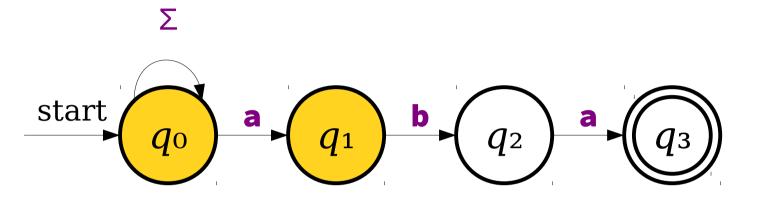


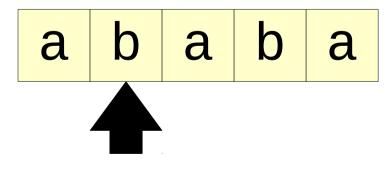


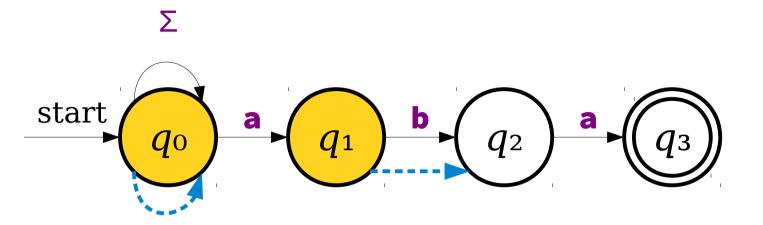


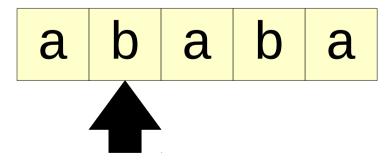


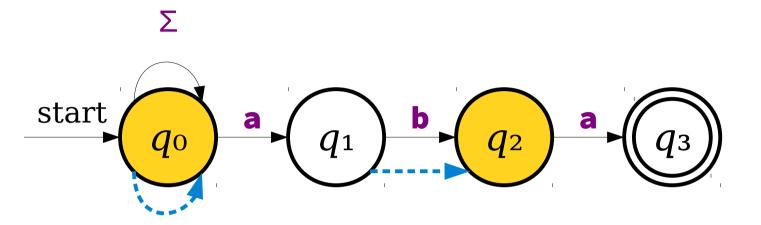


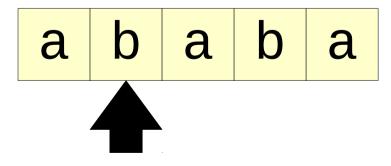


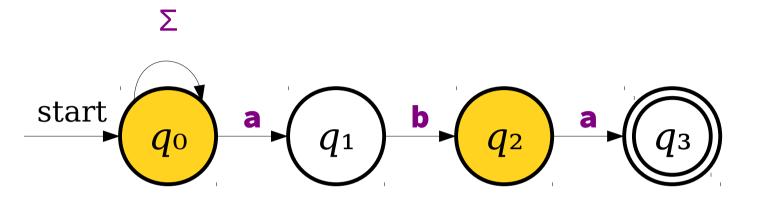


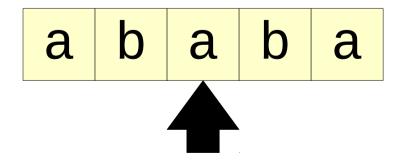


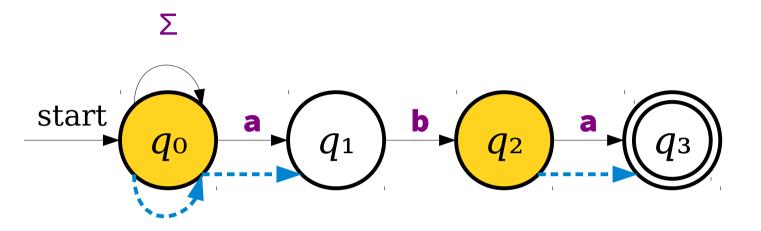


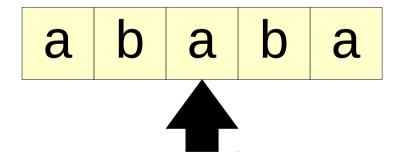


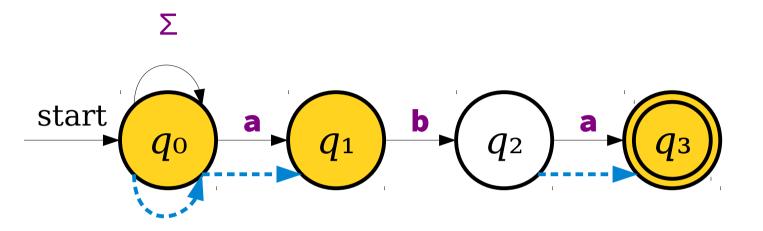


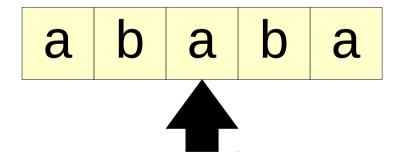


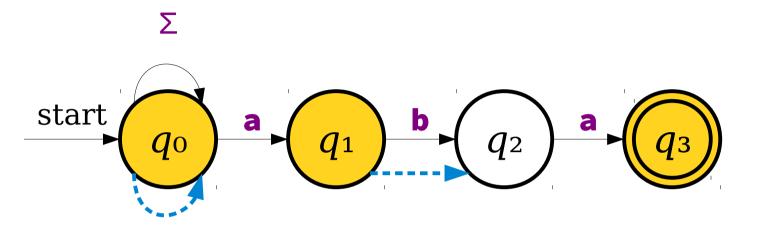


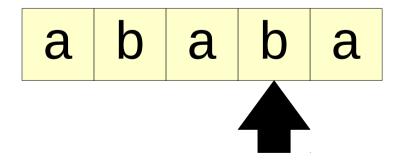


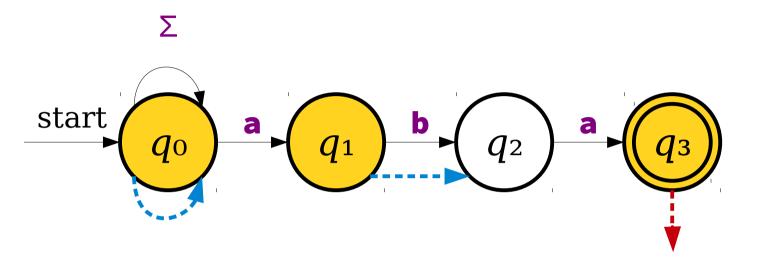


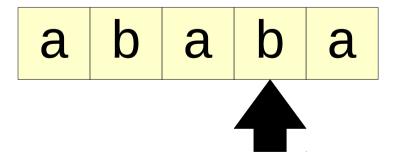


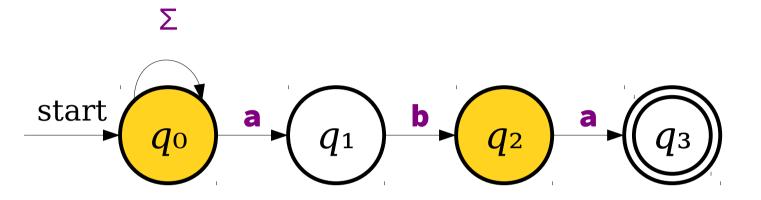


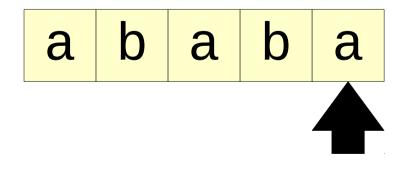


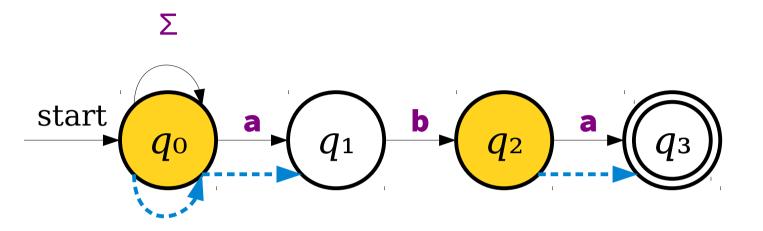


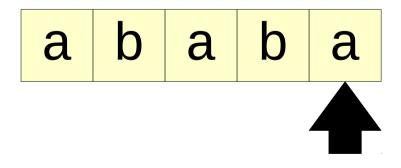


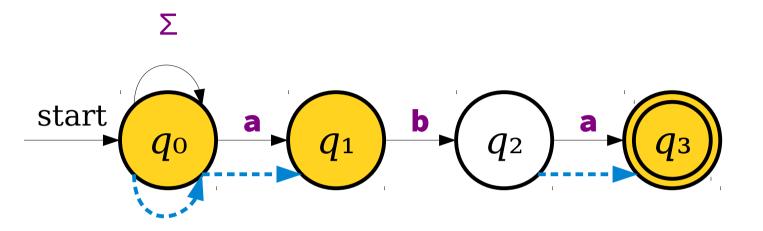


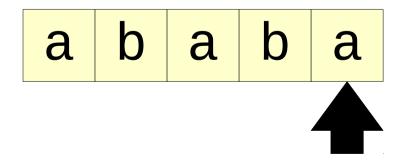


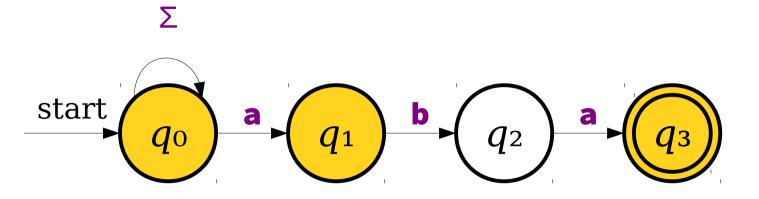




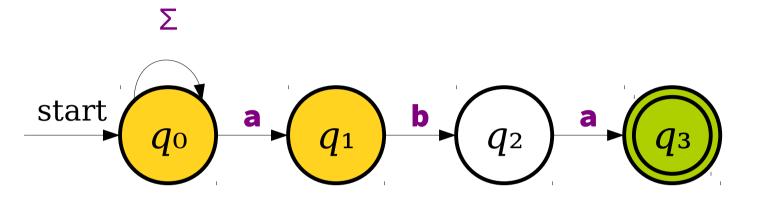








a b a b a



a b a b a

We're in at least one accepting state, so there's some path that gets us to an accepting state.

Therefore, we accept!

- An NFA can be thought of as a DFA that can be in many states at once.
- At each point in time, when the NFA needs to follow a transition, it tries all the options at the same time.
- (Here's a rigorous explanation about how this works; read this on your own time).
 - Start off in the set of all states formed by taking the start state and including each state that can be reached by zero or more ϵ -transitions.
 - When you read a symbol **a** in a set of states *S*:
 - Form the set *S'* of states that can be reached by following a single **a** transition from some state in *S*.
 - Your new set of states is the set of states in S', plus the states reachable from S' by following zero or more ϵ -transitions.

So What?

- Each intuition of nondeterminism is useful in a different setting:
 - Perfect guessing is a great way to think about how to design a machine.
 - Massive parallelism is a great way to test machines and has nice theoretical implications.
- Nondeterministic machines may not be feasible, but they give a great basis for interesting questions:
 - Can any problem that can be solved by a nondeterministic machine be solved by a deterministic machine?
 - Can any problem that can be solved by a nondeterministic machine be solved *efficiently* by a deterministic machine?
- The answers vary from automaton to automaton.

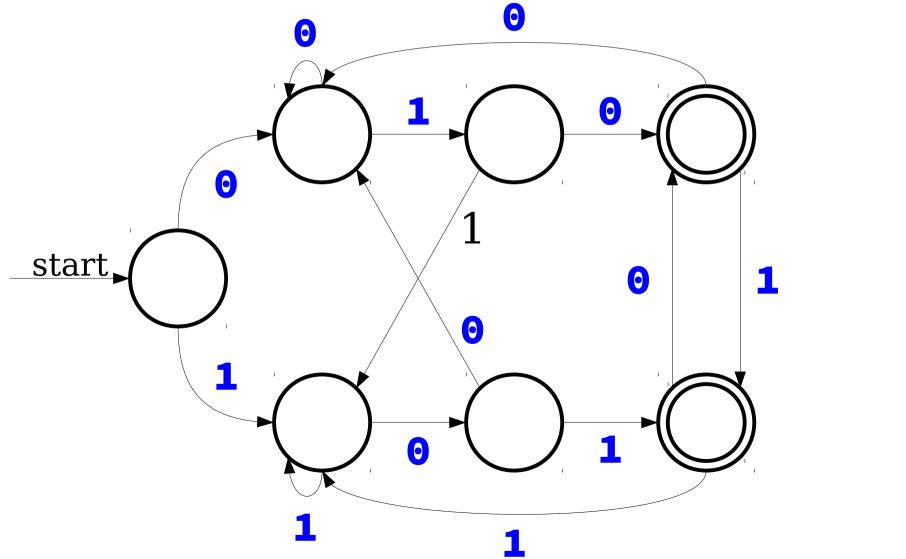
Designing NFAs

Designing NFAs

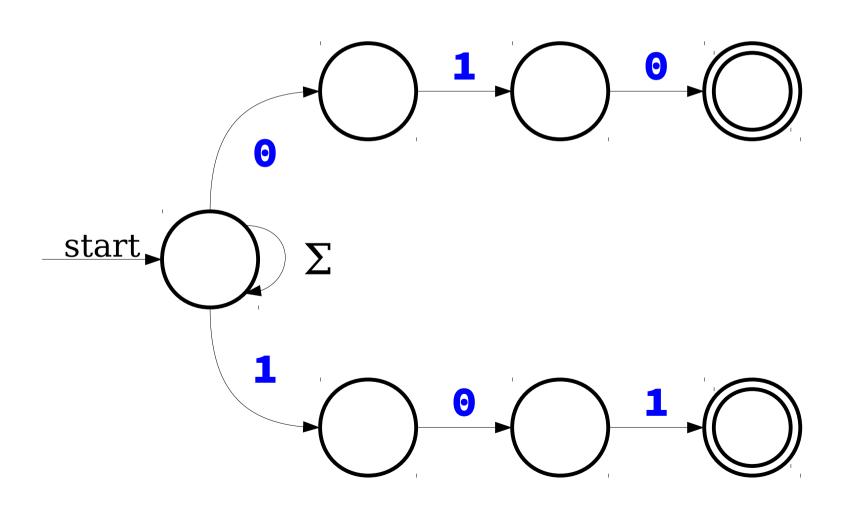
- When designing NFAs, embrace the nondeterminism!
- Good model: *Guess-and-check*:
 - Is there some information that you'd really like to have? Have the machine *nondeterministically guess* that information.
 - Then, have the machine *deterministically check* that the choice was correct.
- The *guess* phase corresponds to trying lots of different options.
- The *check* phase corresponds to filtering out bad guesses or wrong options.

```
L = \{ w \in \{0, 1\}^* \mid w \text{ ends in } 010 \text{ or } 101 \}
```

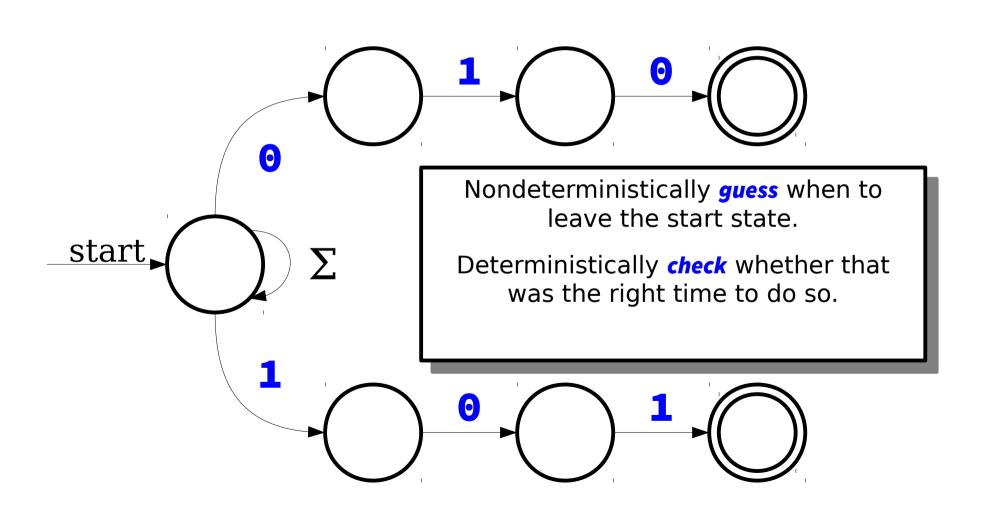
 $L = \{ w \in \{0, 1\}^* \mid w \text{ ends in } 010 \text{ or } 101 \}$



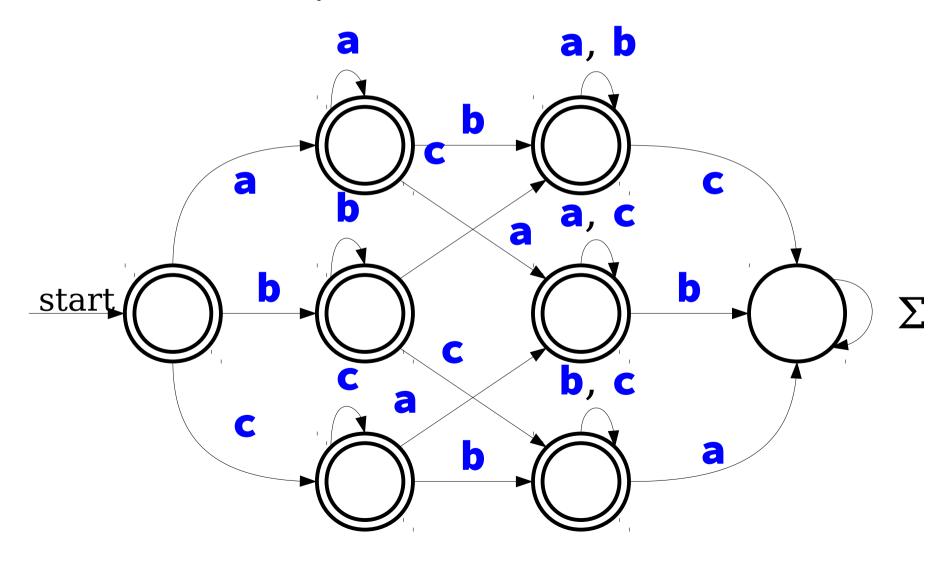
 $L = \{ w \in \{0, 1\}^* \mid w \text{ ends in } 010 \text{ or } 101 \}$



 $L = \{ w \in \{0, 1\}^* \mid w \text{ ends in } 010 \text{ or } 101 \}$

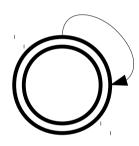


```
L = \{ w \in \{a, b, c\}^* \mid \text{at least one of } a, b, \text{ or } c \text{ is not in } w \}
```

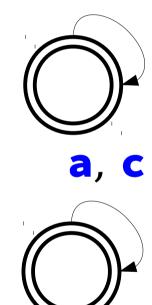


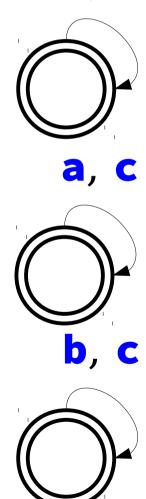
```
L = \{ w \in \{a, b, c\}^* \mid \text{at least one of } a, b, \text{ or } c \text{ is not in } w \}
```

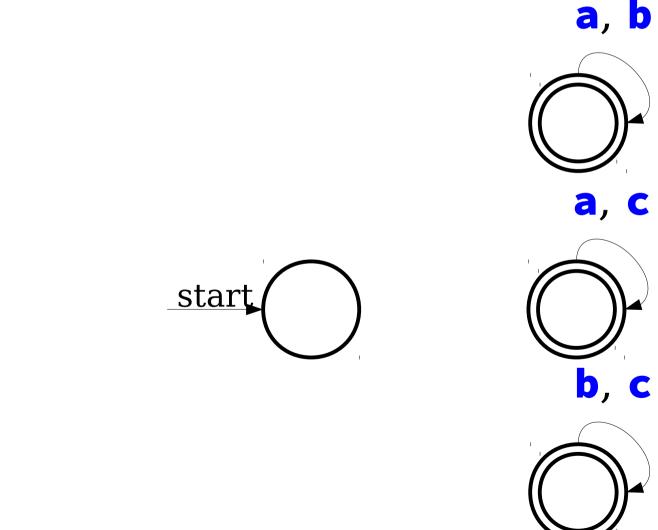
 $L = \{ w \in \{a, b, c\}^* \mid \text{at least one of a, b, or c is not in } w \}$

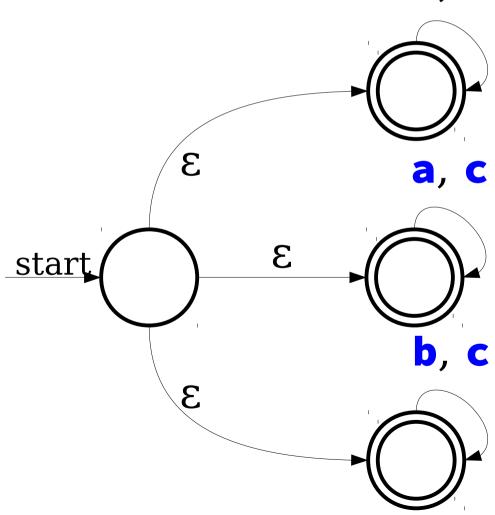


 $L = \{ w \in \{a, b, c\}^* \mid \text{at least one of a, b, or c is not in } w \}$

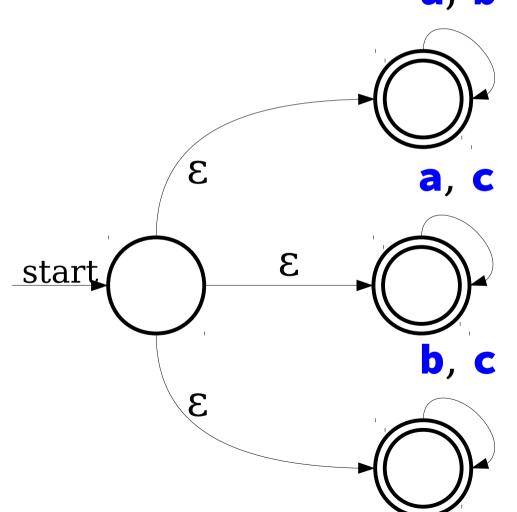








 $L = \{ w \in \{a, b, c\}^* \mid \text{at least one of } a, b, \text{ or } c \text{ is not in } w \}$



Nondeterministically **guess** which character is missing.

Deterministically *check* whether that character is indeed missing.

Just how powerful are NFAs?