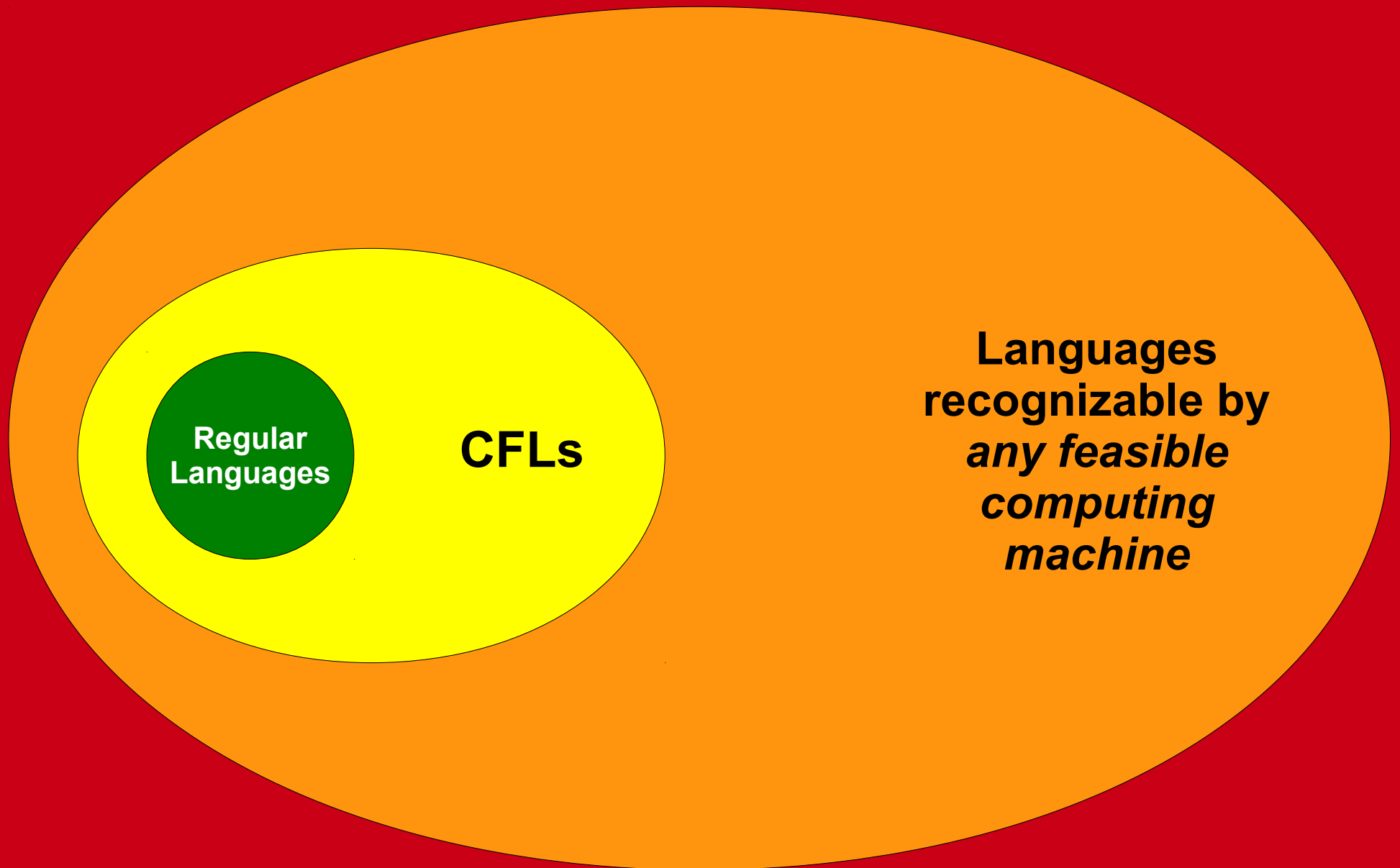


Turing Machines

Part One

What problems can we solve with a computer?



All Languages

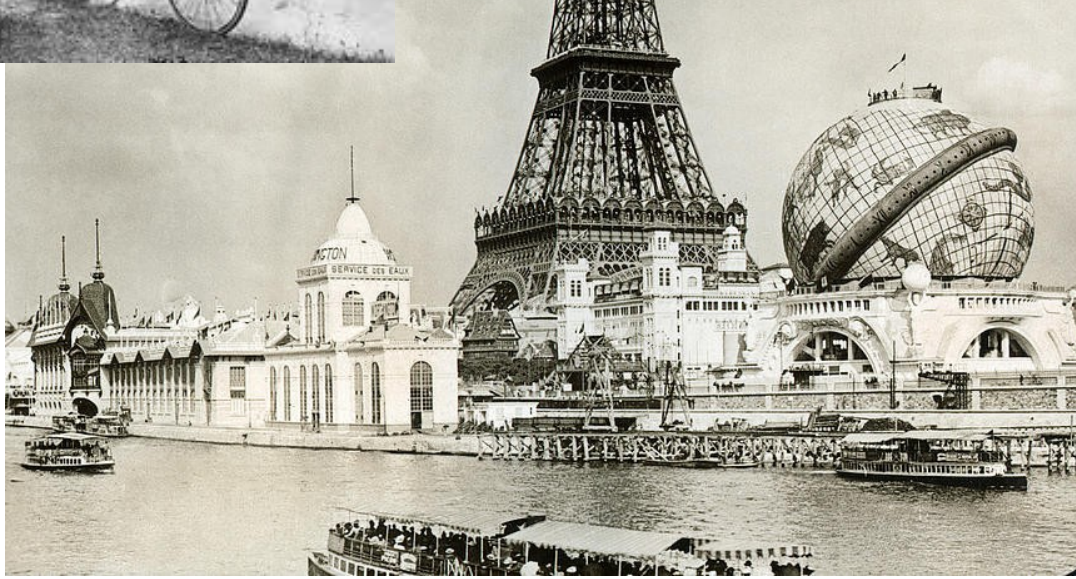
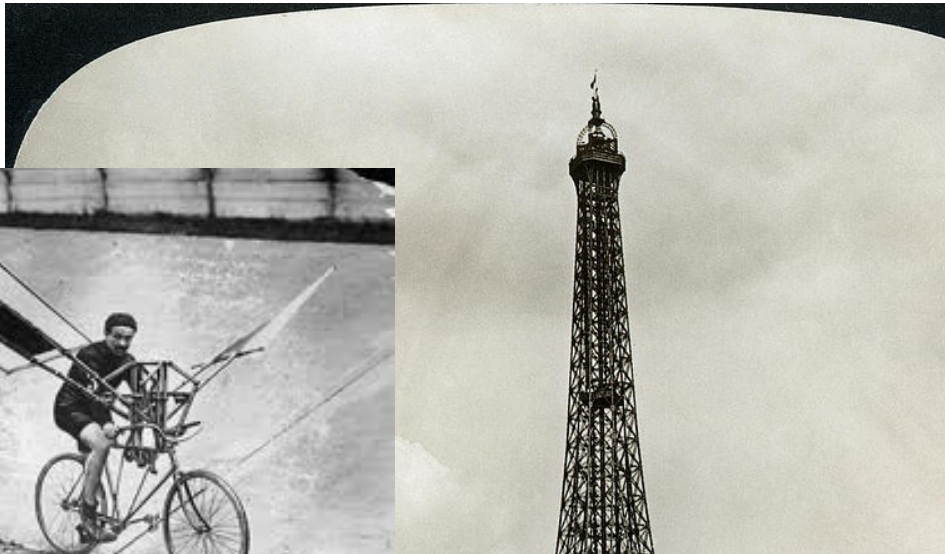
That same drawing, to scale.

The Problem

- Finite automata accept precisely the regular languages.
- We may need unbounded memory to recognize context-free languages.
 - e.g. $\{ \mathbf{a}^n \mathbf{b}^n \mid n \in \mathbb{N} \}$ requires unbounded counting.
- How do we build an automaton with finitely many states but unbounded memory?

A Brief History Lesson

Technology has solved all of mankind's ^{<ahem>} problems! No more wars or sad ever!



Hilbert's Vision

- 1900: International Congress of Mathematicians meeting in Paris
- Proposes 23 unsolved problems as the agenda for the coming years
- An important theme is not simply proving more theorems, but achieving *automation* of theorem-proving, even theorem generation.
- Humanity lives in leisure while all Truth flows effortlessly into our hands on a ticker tape!

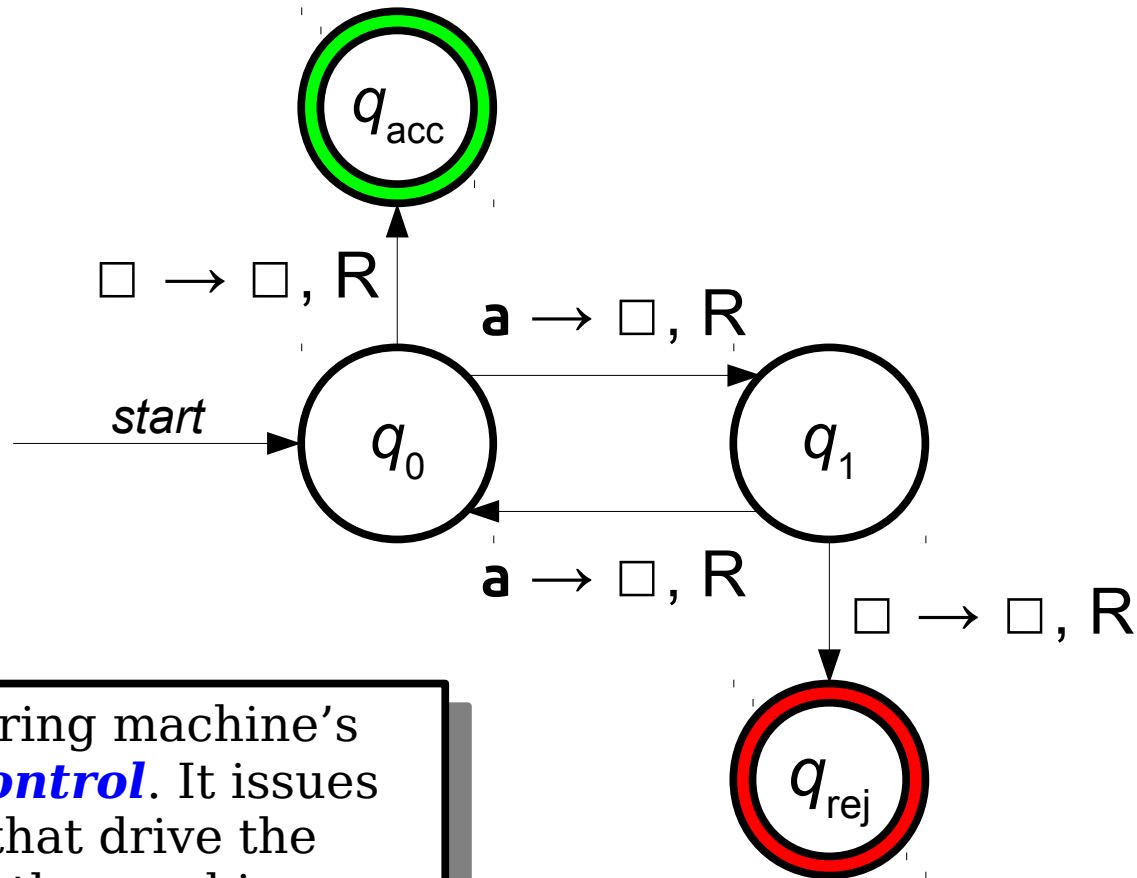
“No one shall expel us from the Paradise that Cantor has created!” -David Hilbert

Hi there, Reality!

- Hilbert's agenda is both a spectacular success and a spectacular failure
- Inspires some of the most impactful theoretical work in mathematical history, human history
- Brings us heroes like **Alan Turing** and **Kurt Gödel!**
- These incredible results consist of utterly demolishing all the pillars of Hilbert's vision of automated knowledge creation, within just a few years

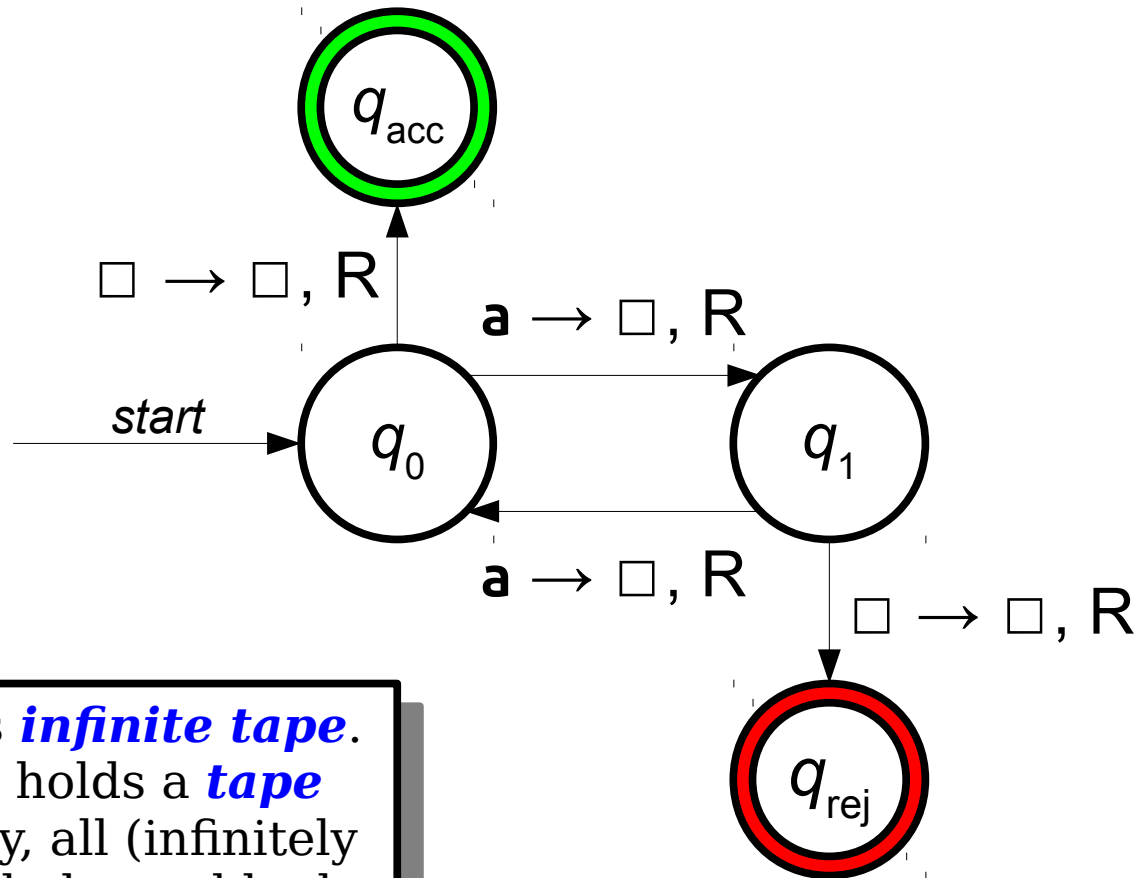
“No one shall expel us from the Paradise that Cantor has created!” -David Hilbert

A Simple Turing Machine

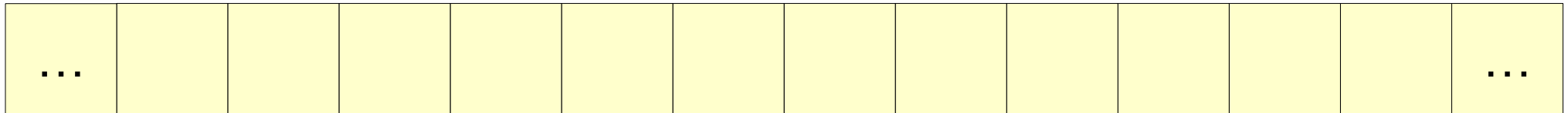


This is the Turing machine's ***finite state control***. It issues commands that drive the operation of the machine.

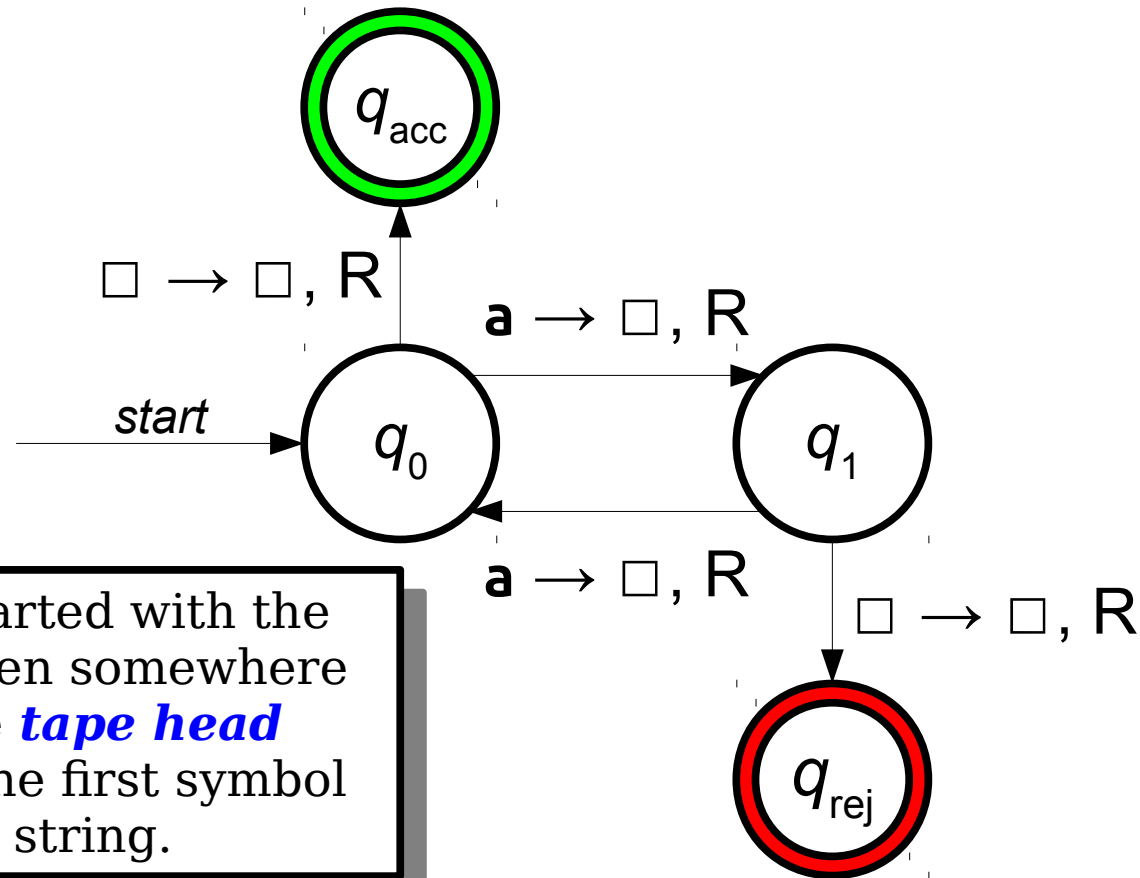
A Simple Turing Machine



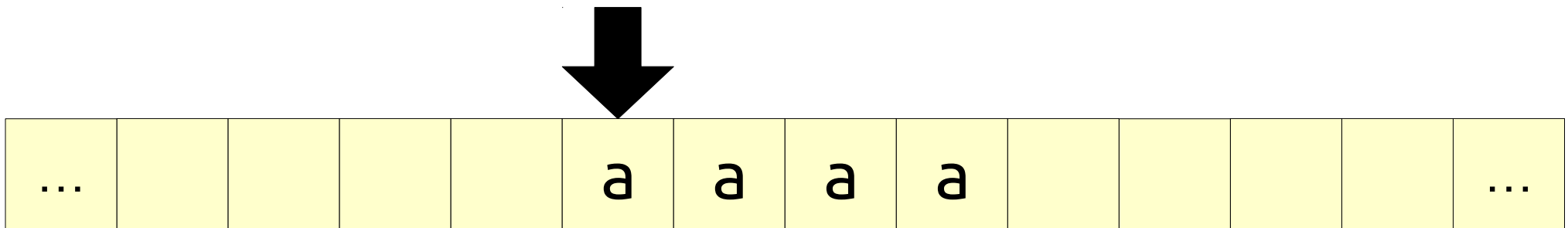
This is the TM's **infinite tape**.
Each tape cell holds a **tape symbol**. Initially, all (infinitely many) tape symbols are blank.



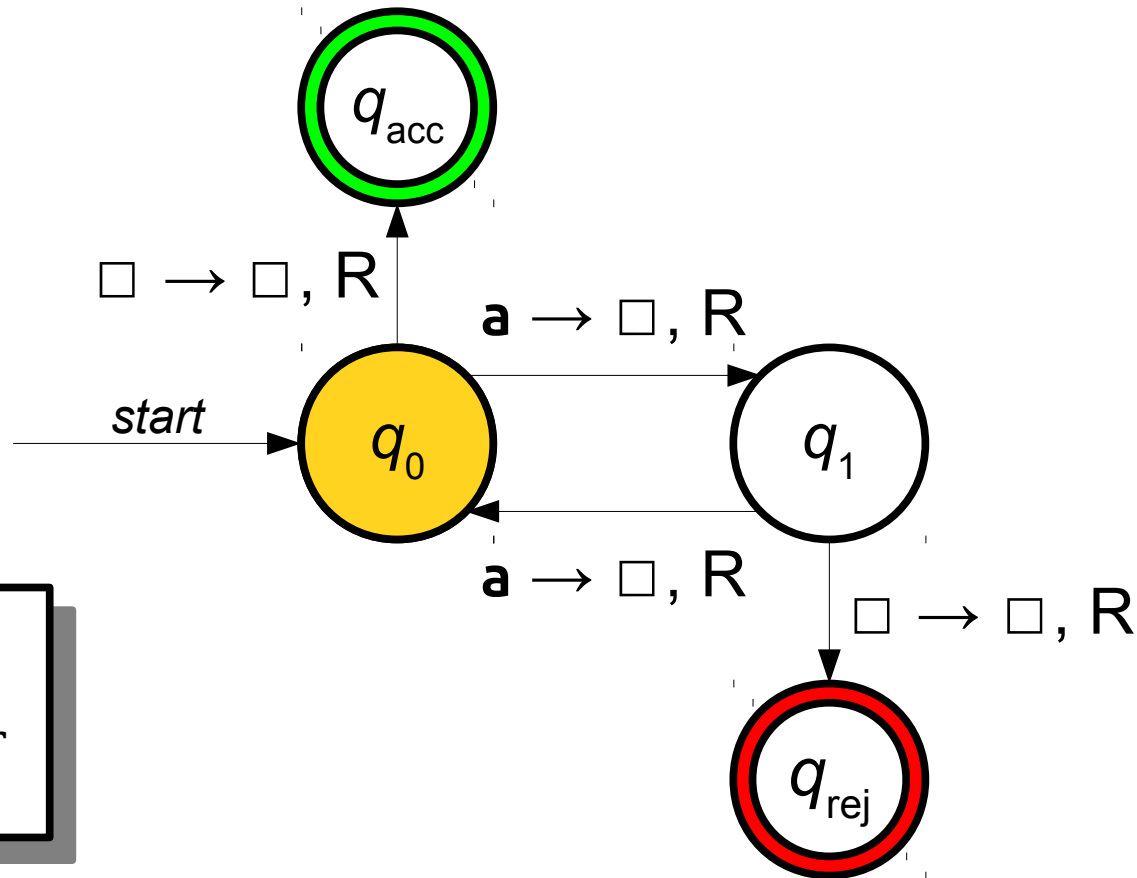
A Simple Turing Machine



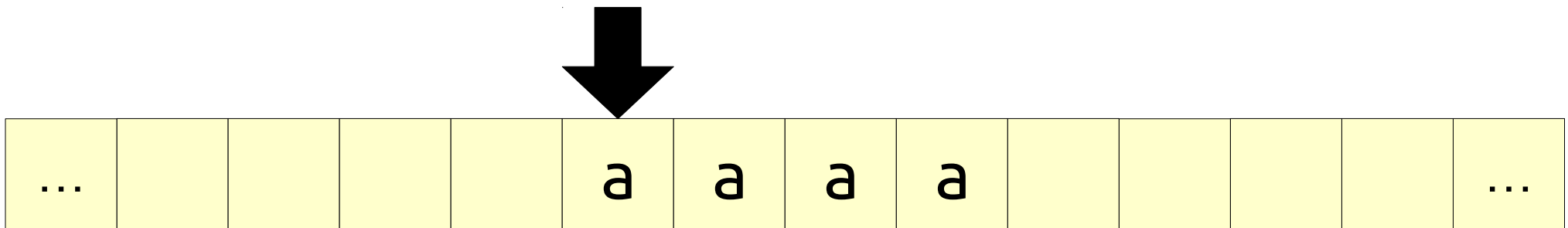
The machine is started with the **input string** written somewhere on the tape. The **tape head** initially points to the first symbol of the input string.



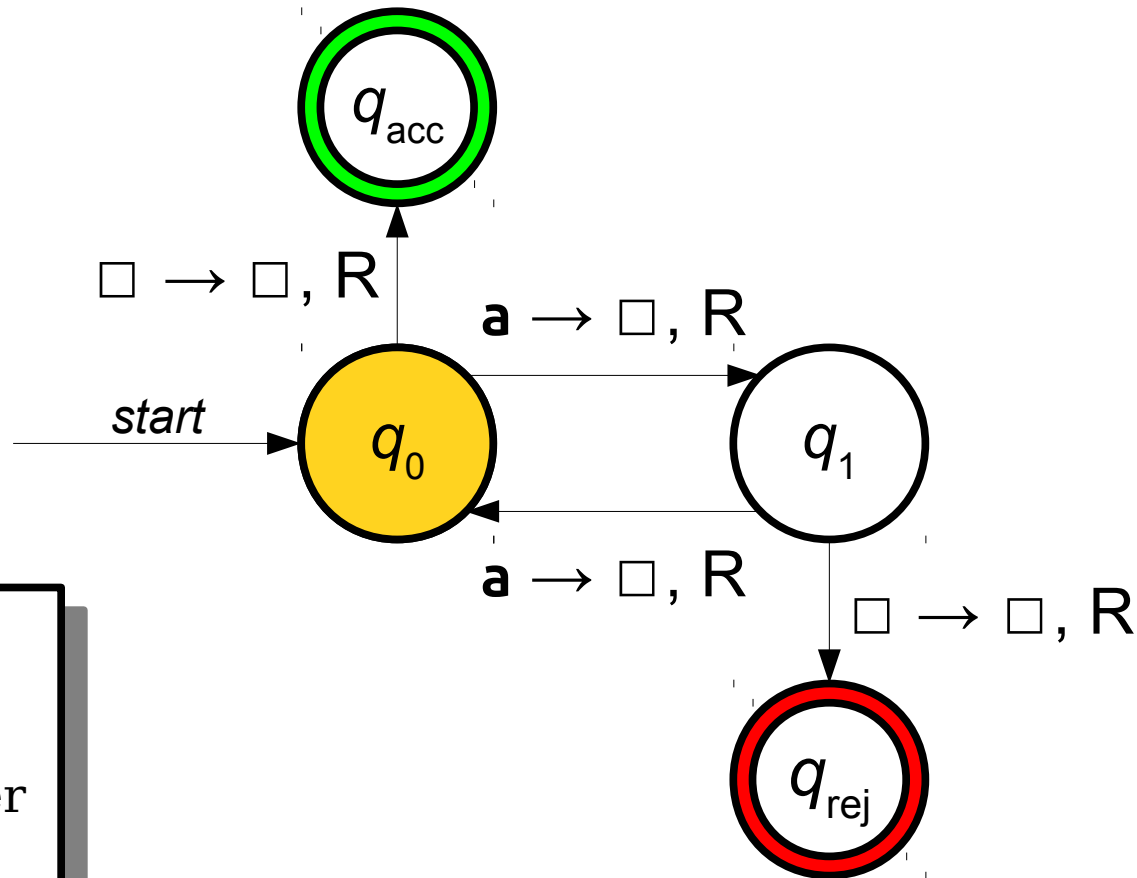
A Simple Turing Machine



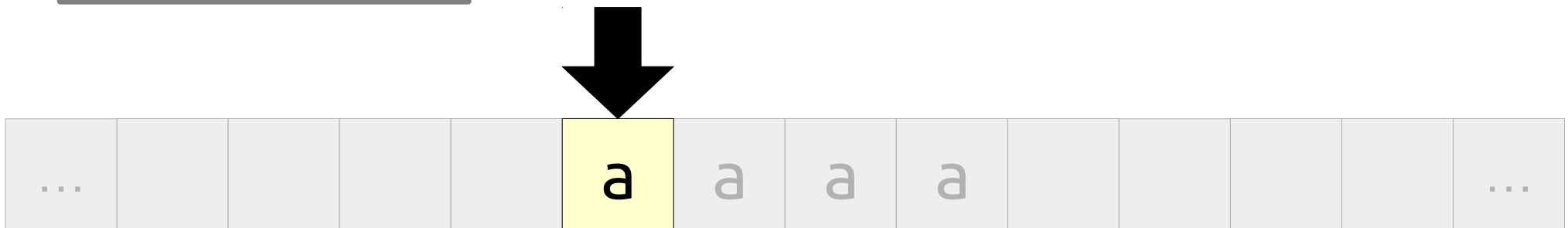
Like DFAs and NFAs, TMs begin execution in their ***start state***.



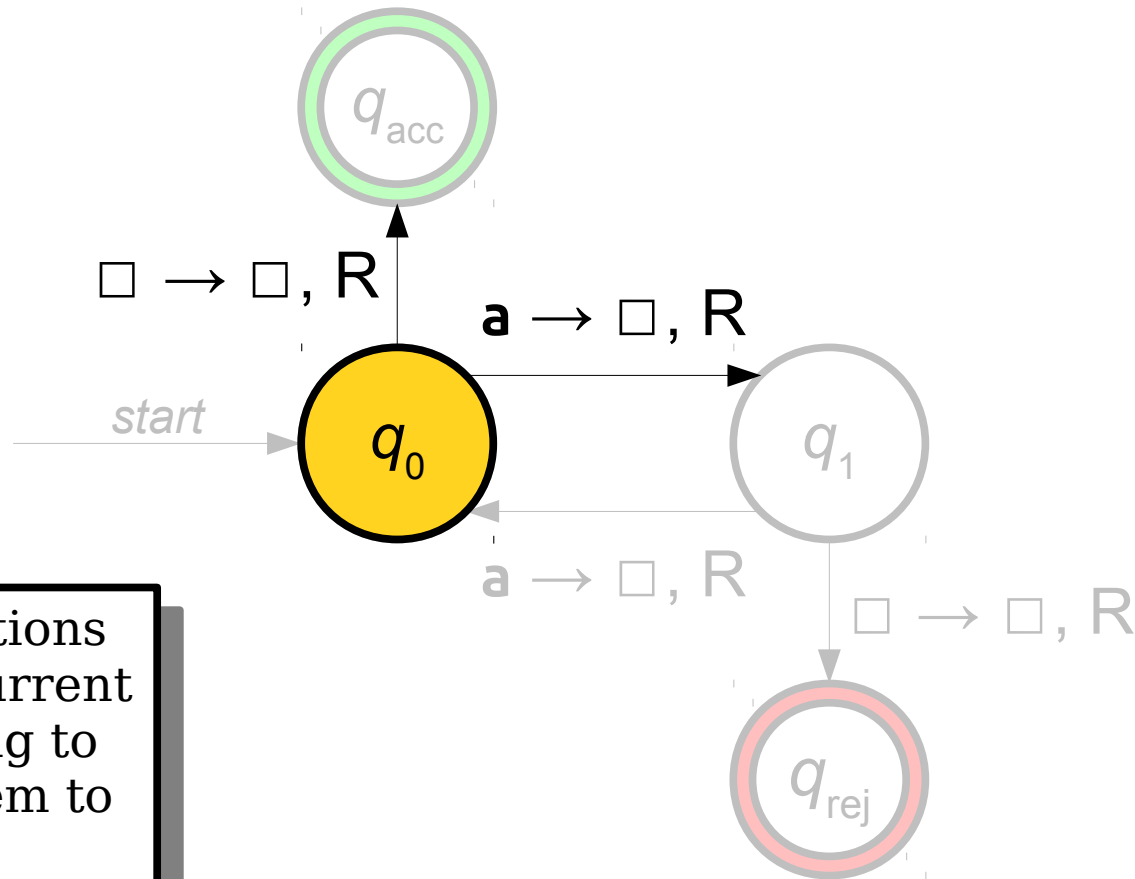
A Simple Turing Machine



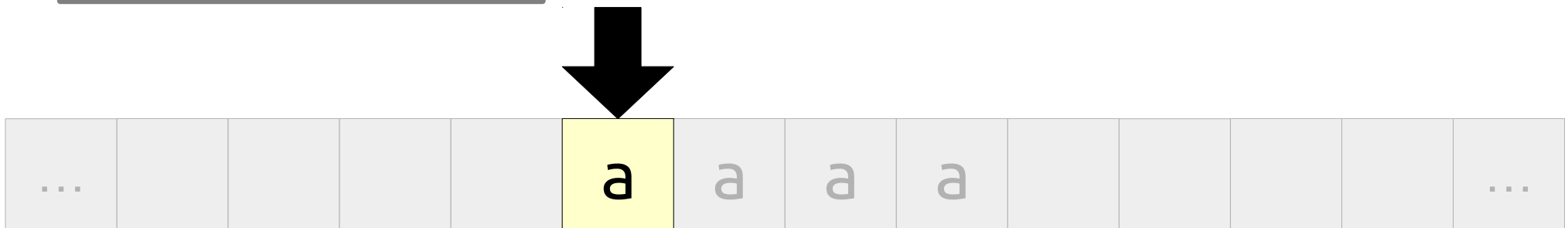
At each step, the TM only looks at the symbol immediately under the **tape head**.



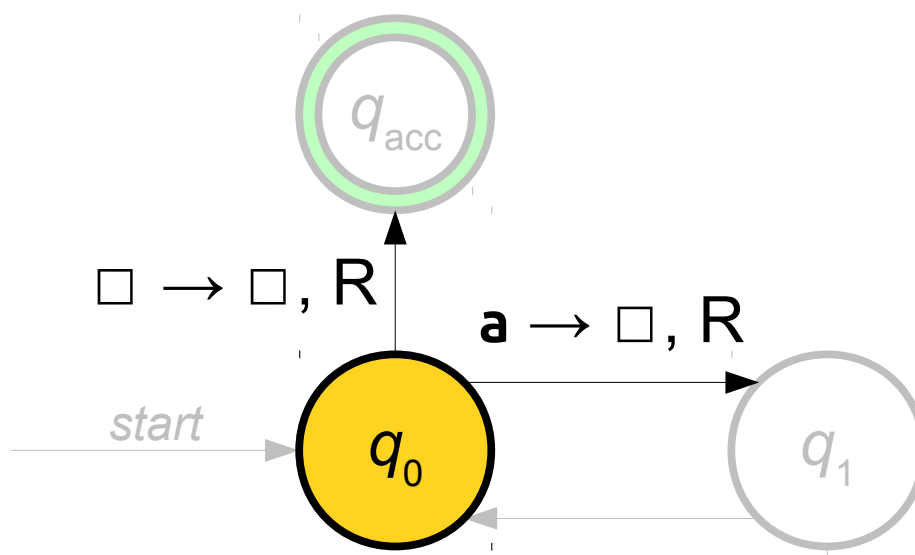
A Simple Turing Machine



These two transitions originate at the current state. We're going to choose one of them to follow.



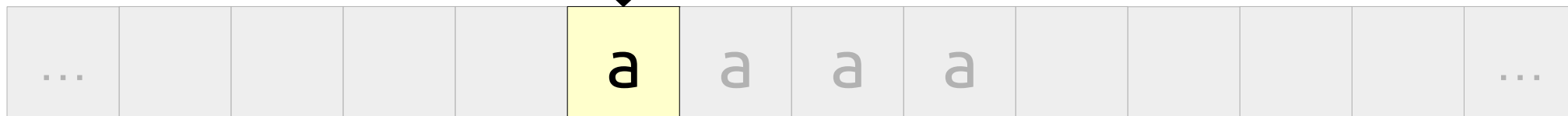
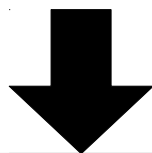
A Simple Turing Machine



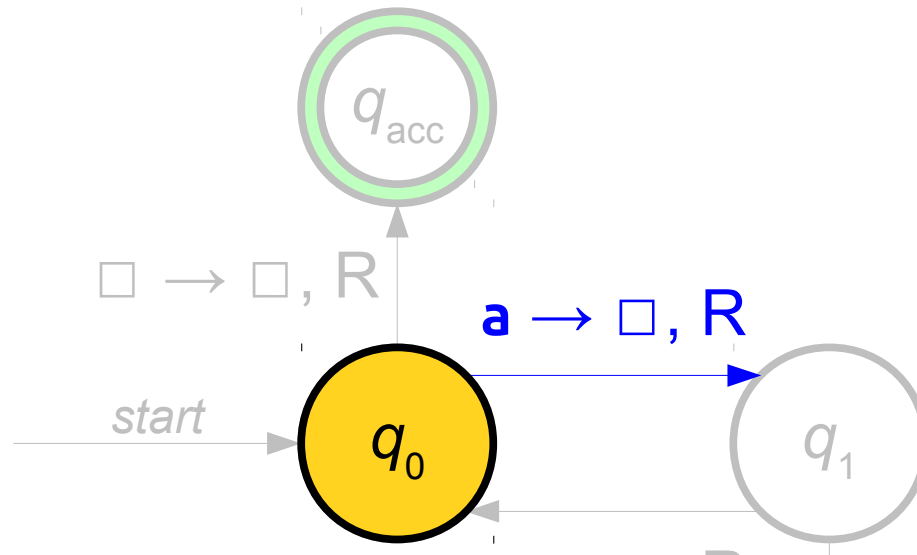
Each transition has the form

read \rightarrow ***write***, ***dir***

and means “if symbol ***read*** is under the tape head, replace it with ***write*** and move the tape head in direction ***dir*** (L or R). The \square symbol denotes a blank cell.



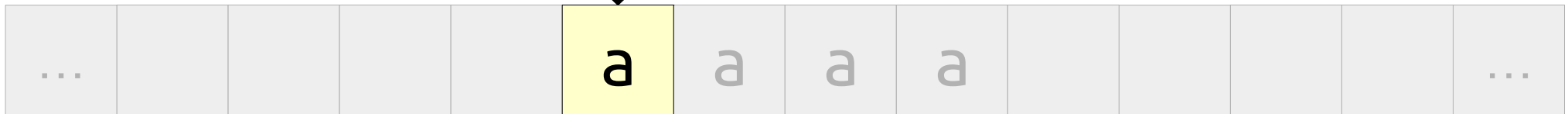
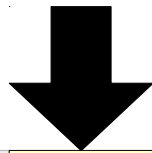
A Simple Turing Machine



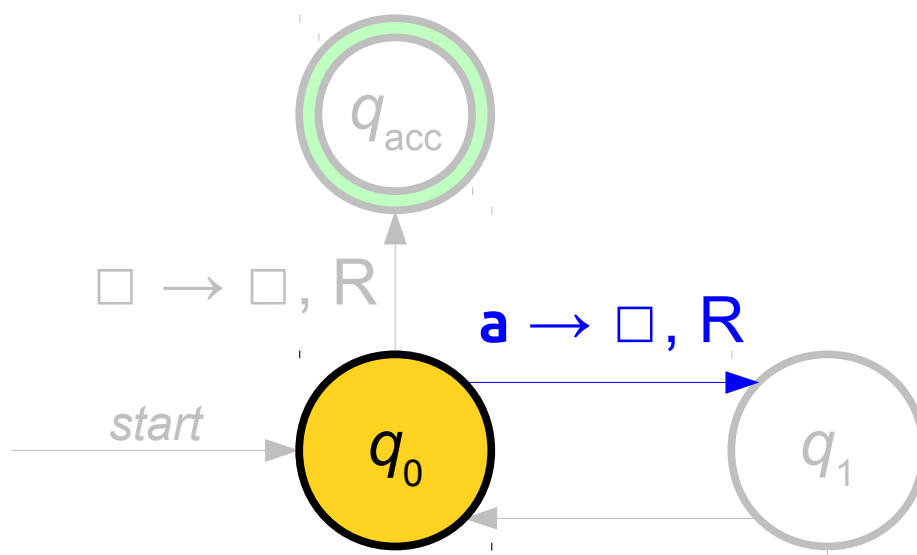
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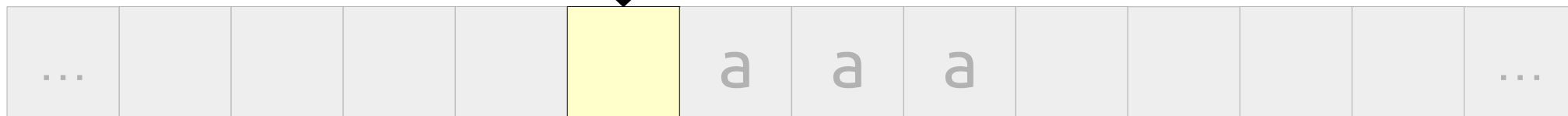
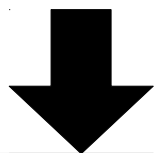
A Simple Turing Machine



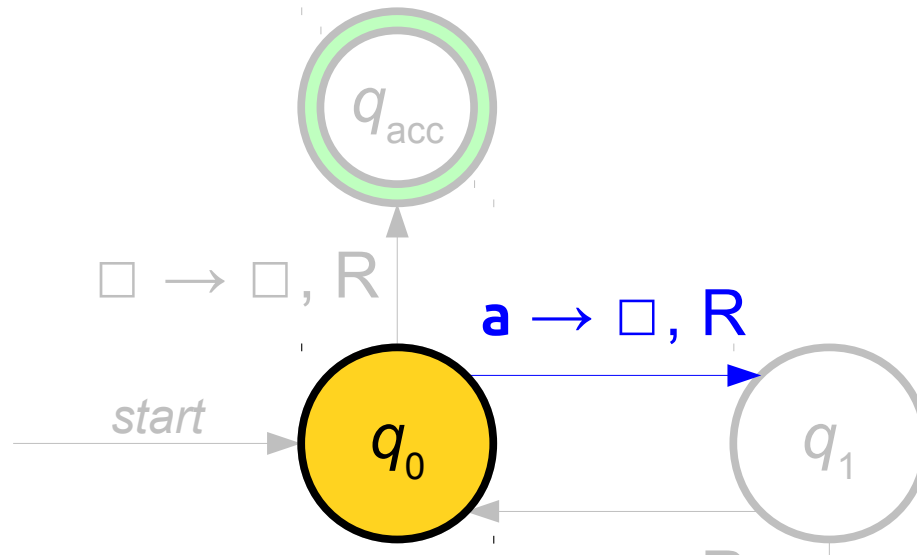
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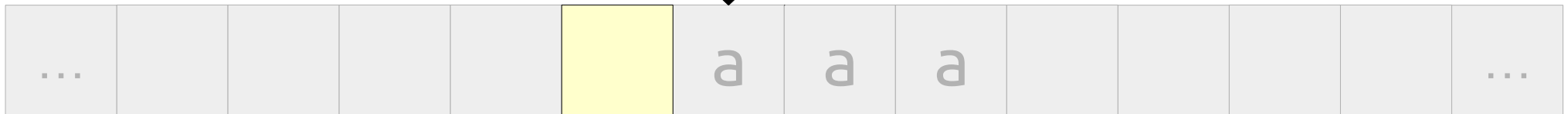
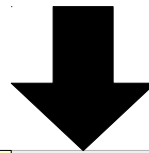
A Simple Turing Machine



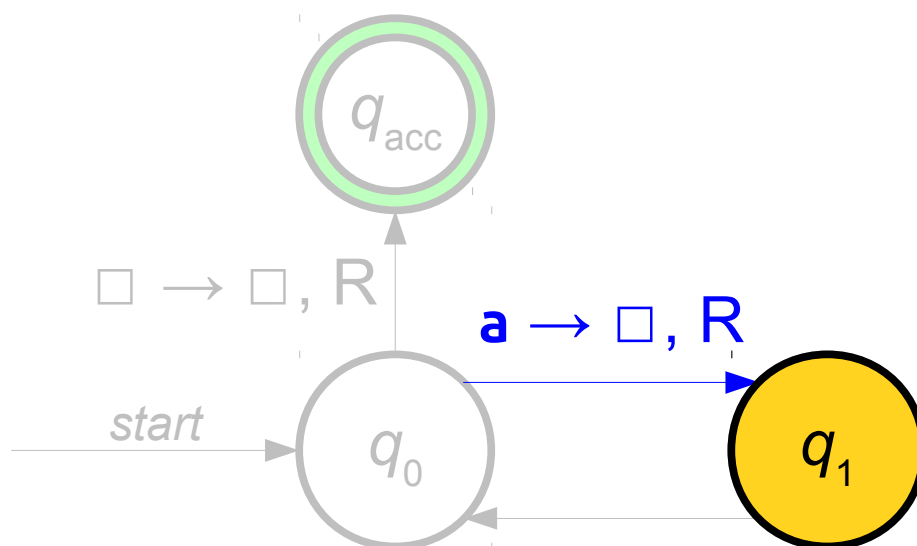
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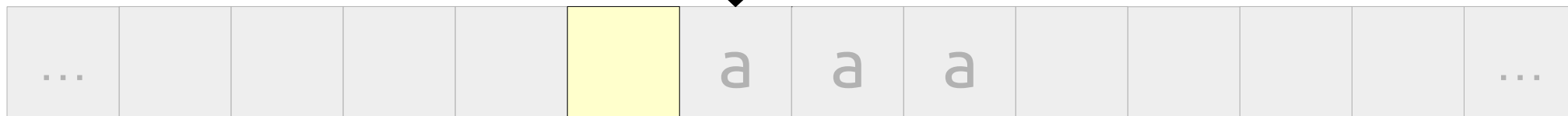
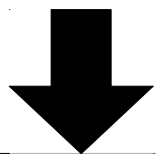
A Simple Turing Machine



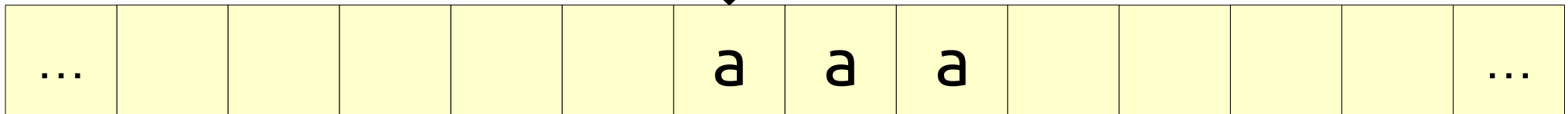
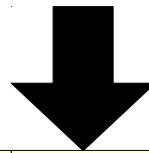
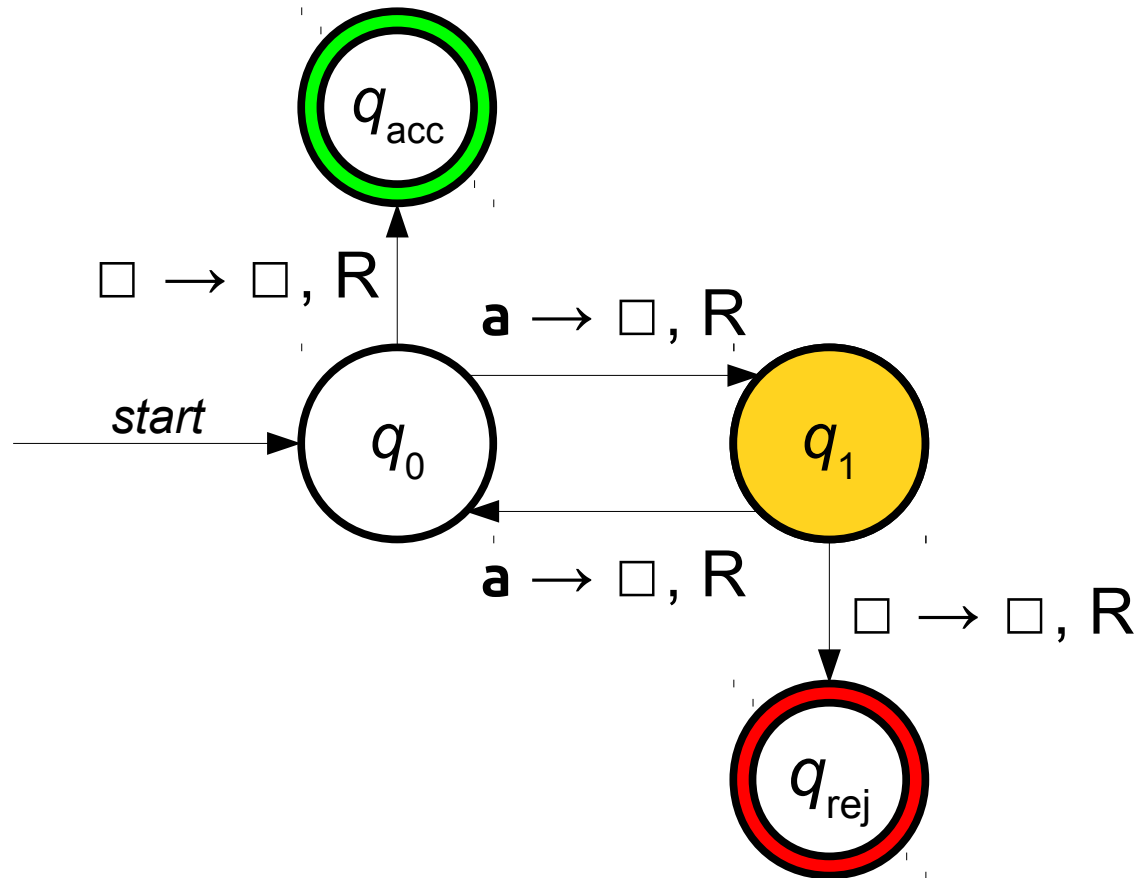
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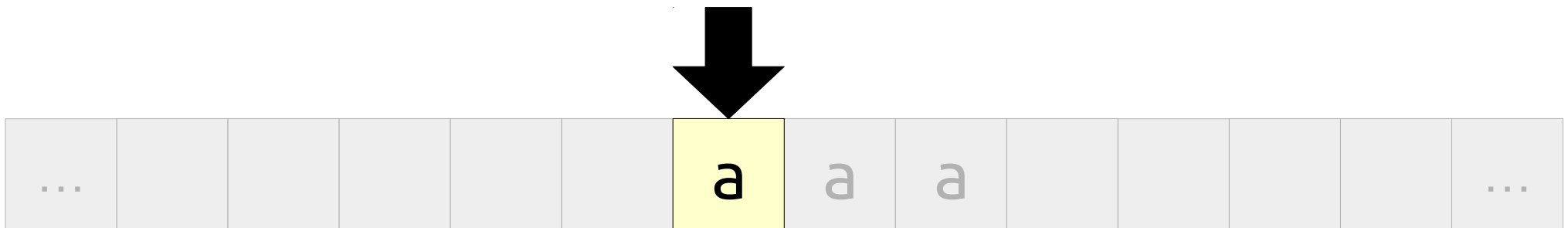
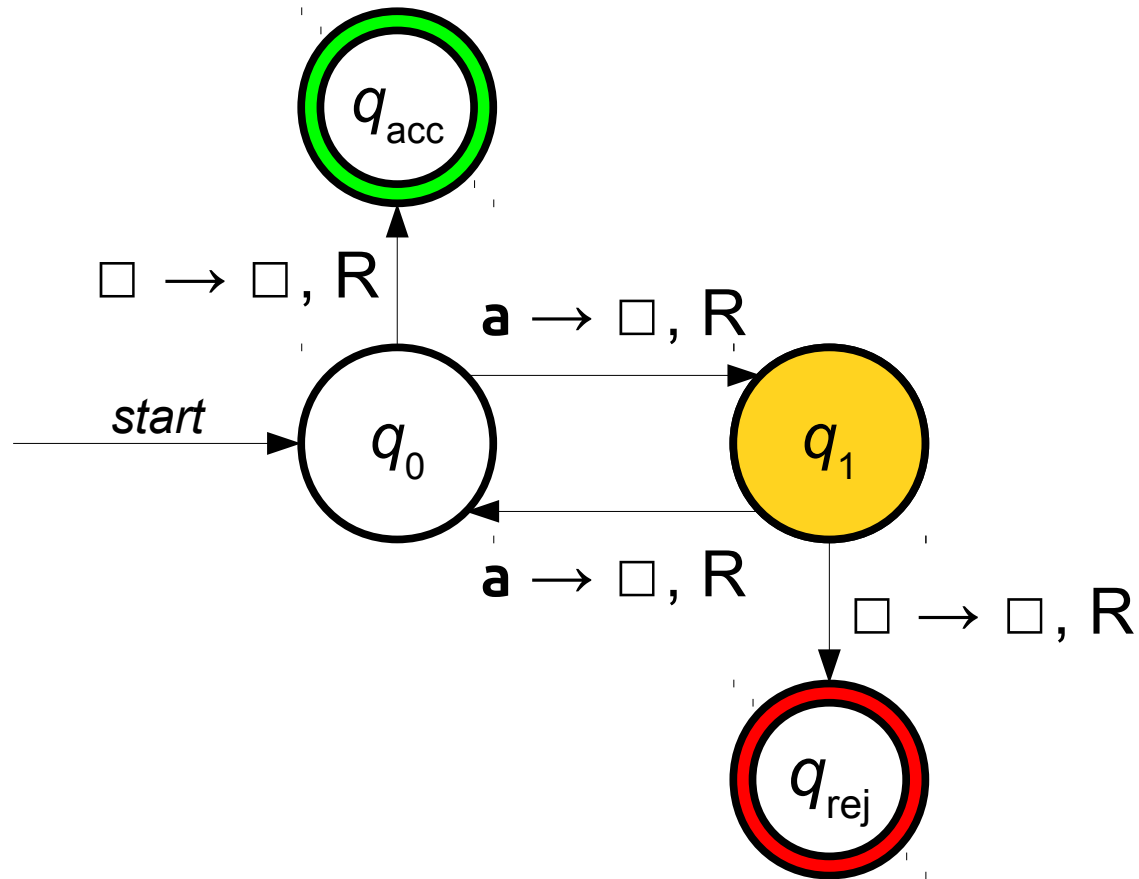
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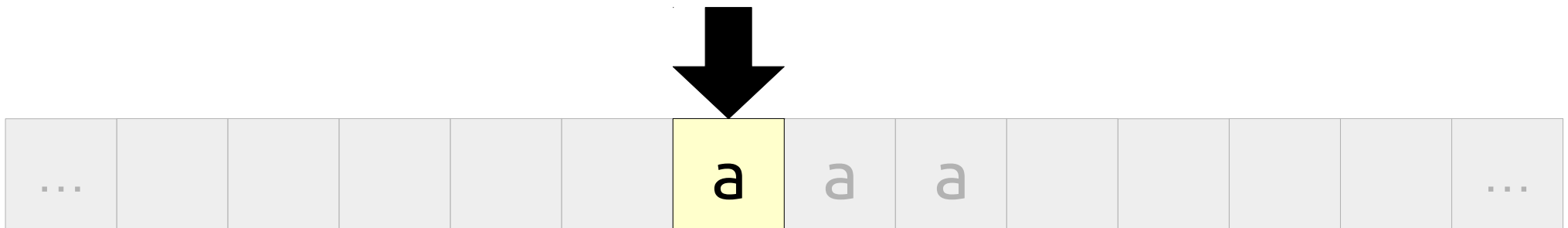
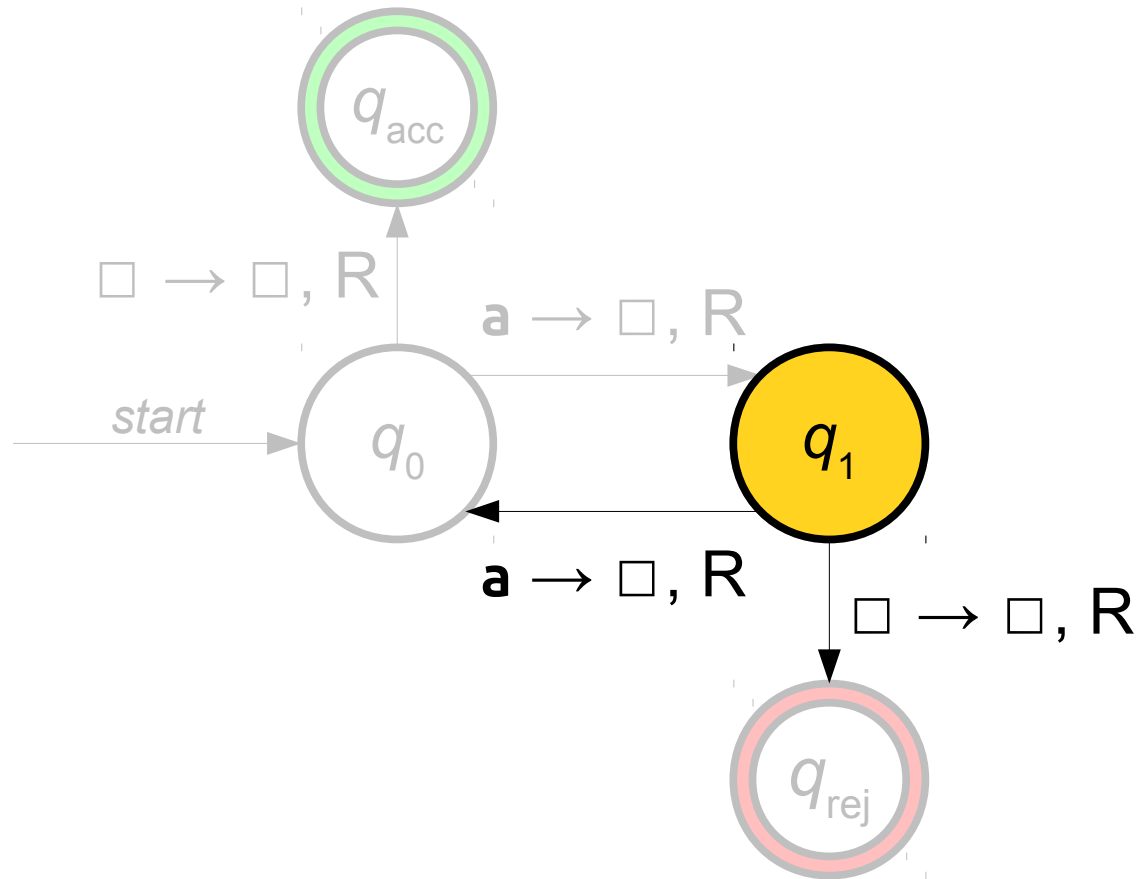
A Simple Turing Machine



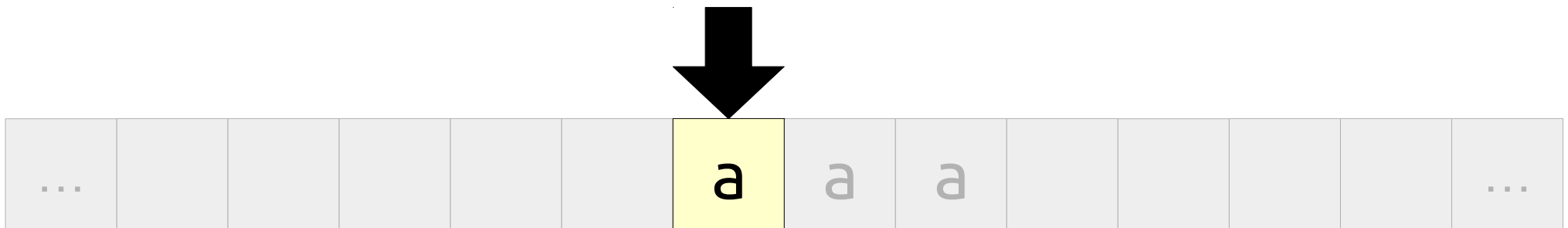
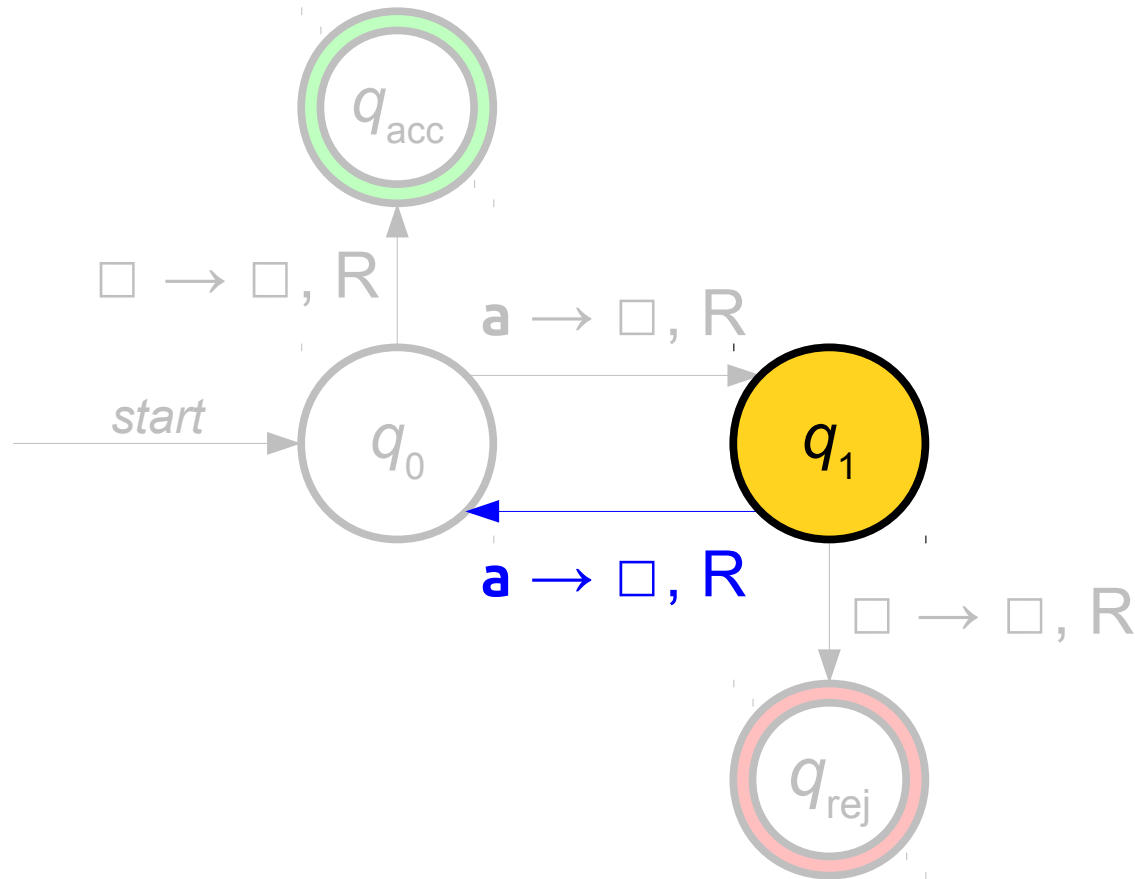
A Simple Turing Machine



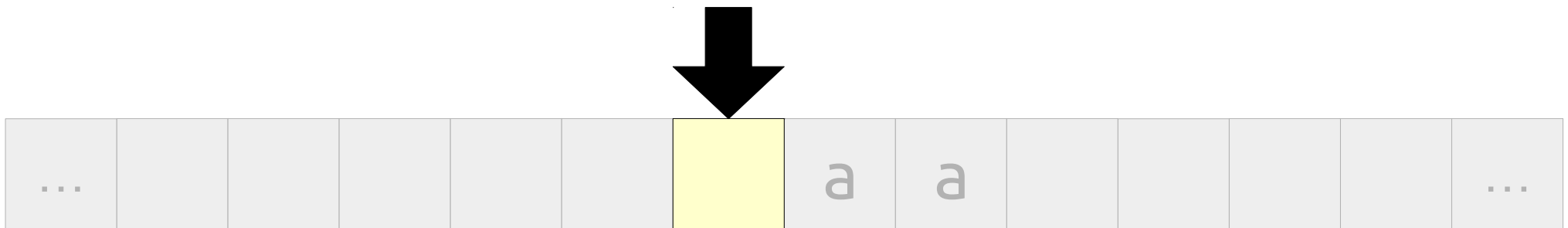
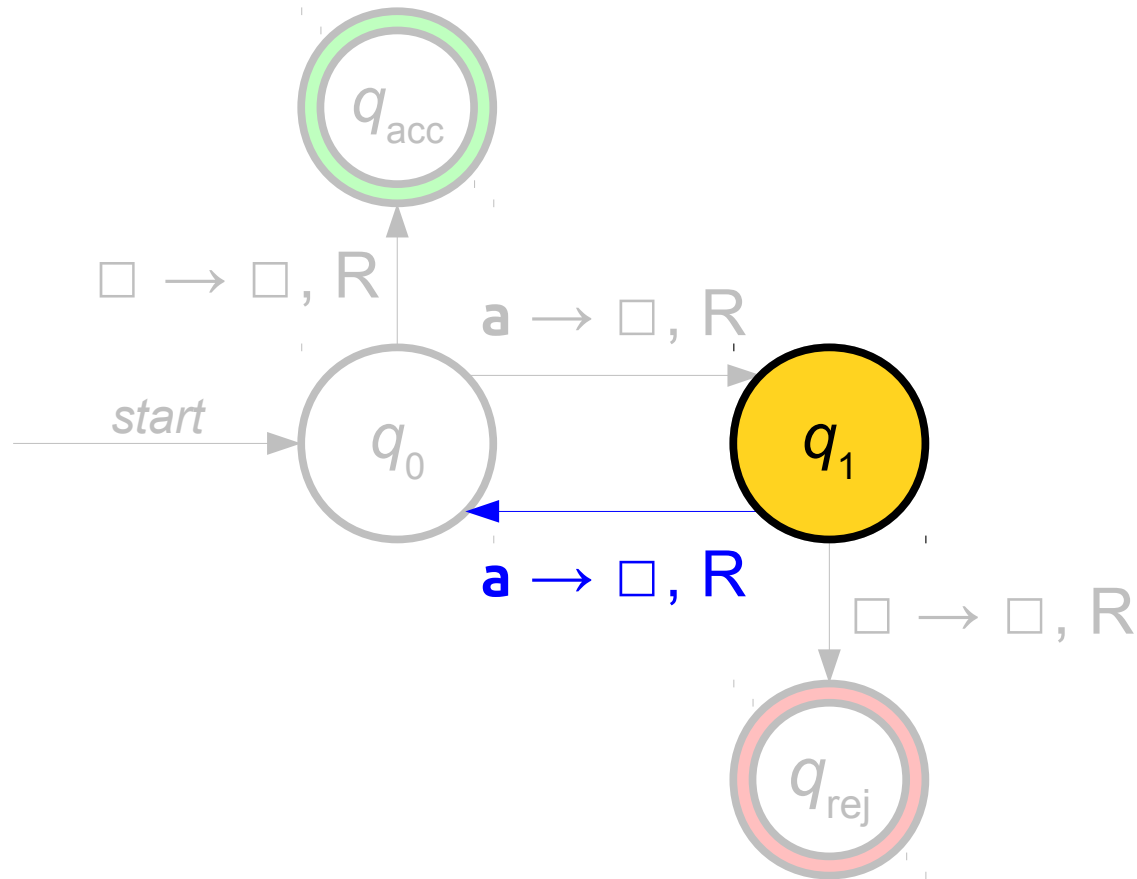
A Simple Turing Machine



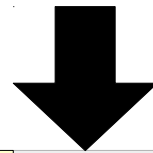
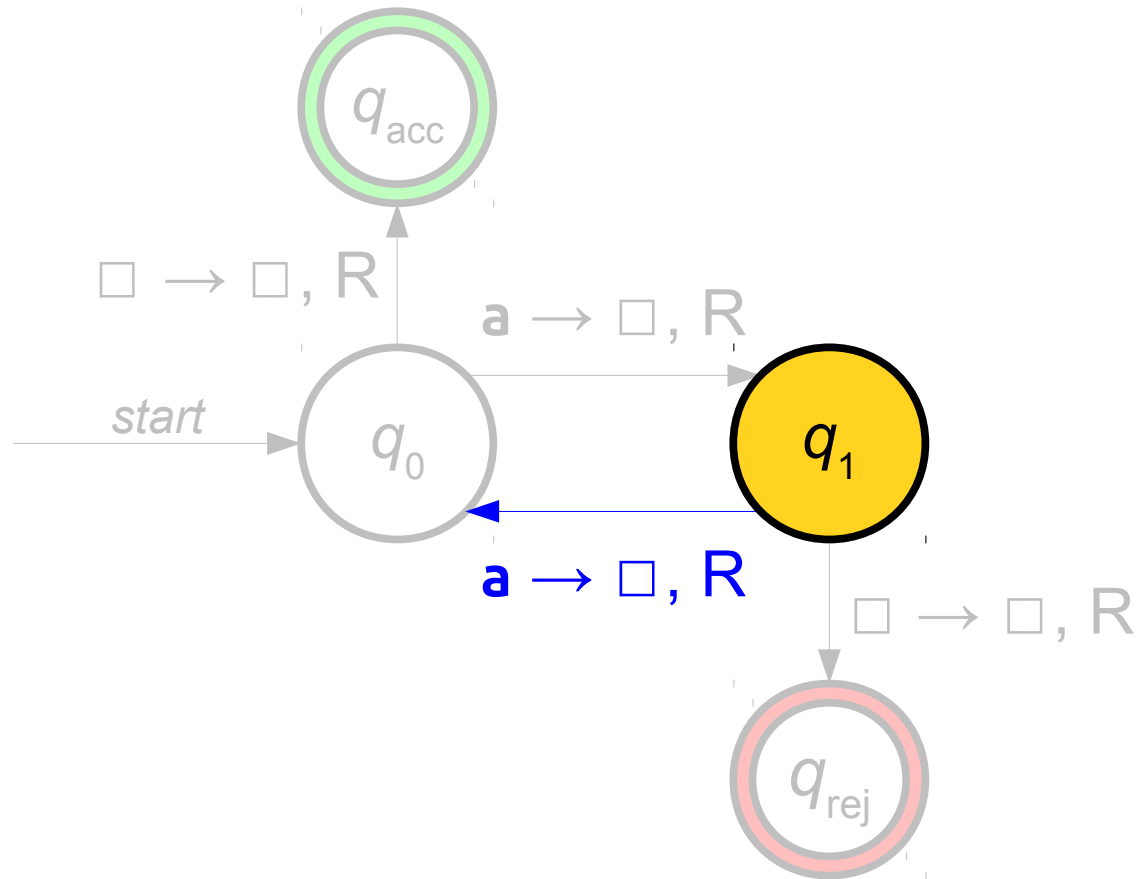
A Simple Turing Machine



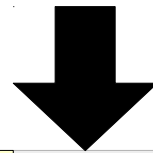
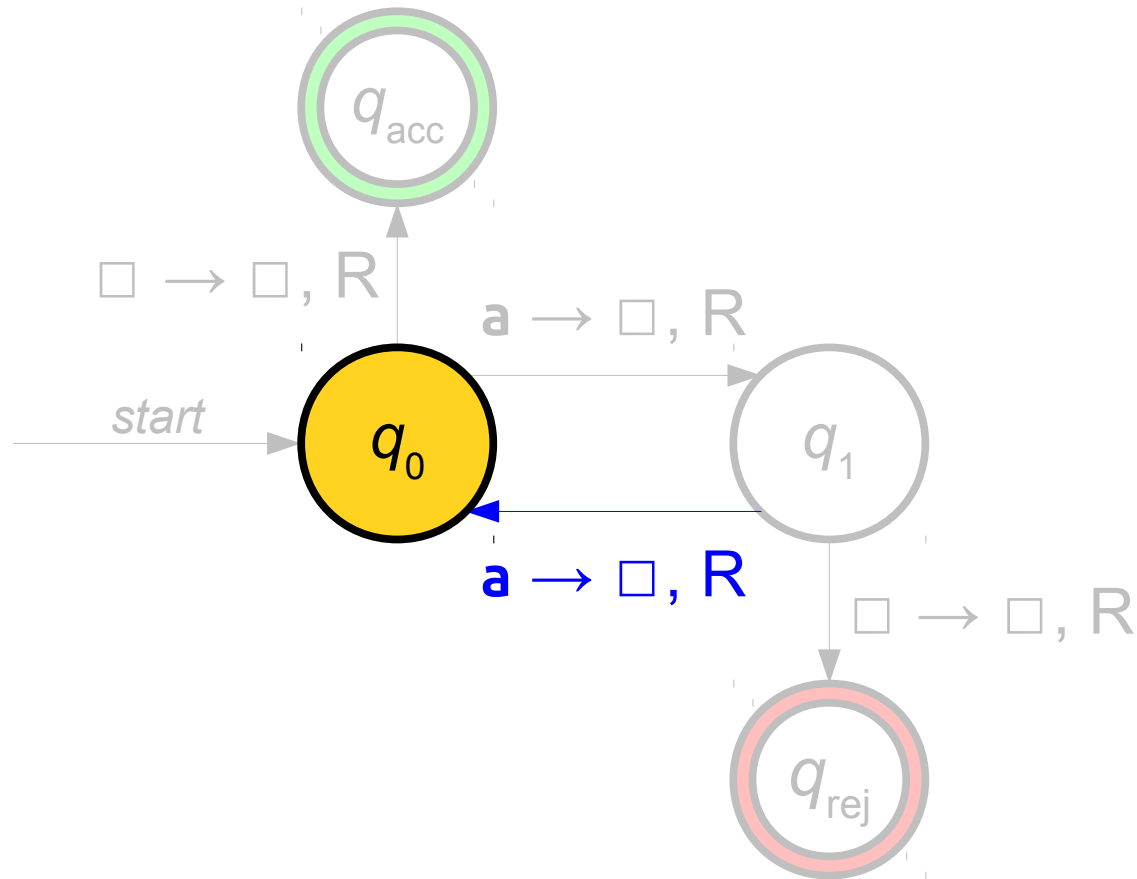
A Simple Turing Machine



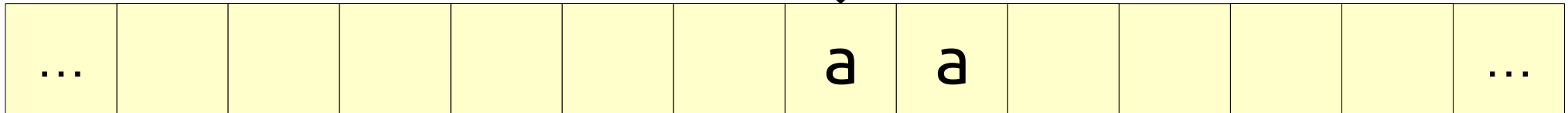
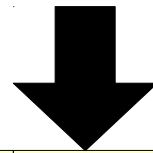
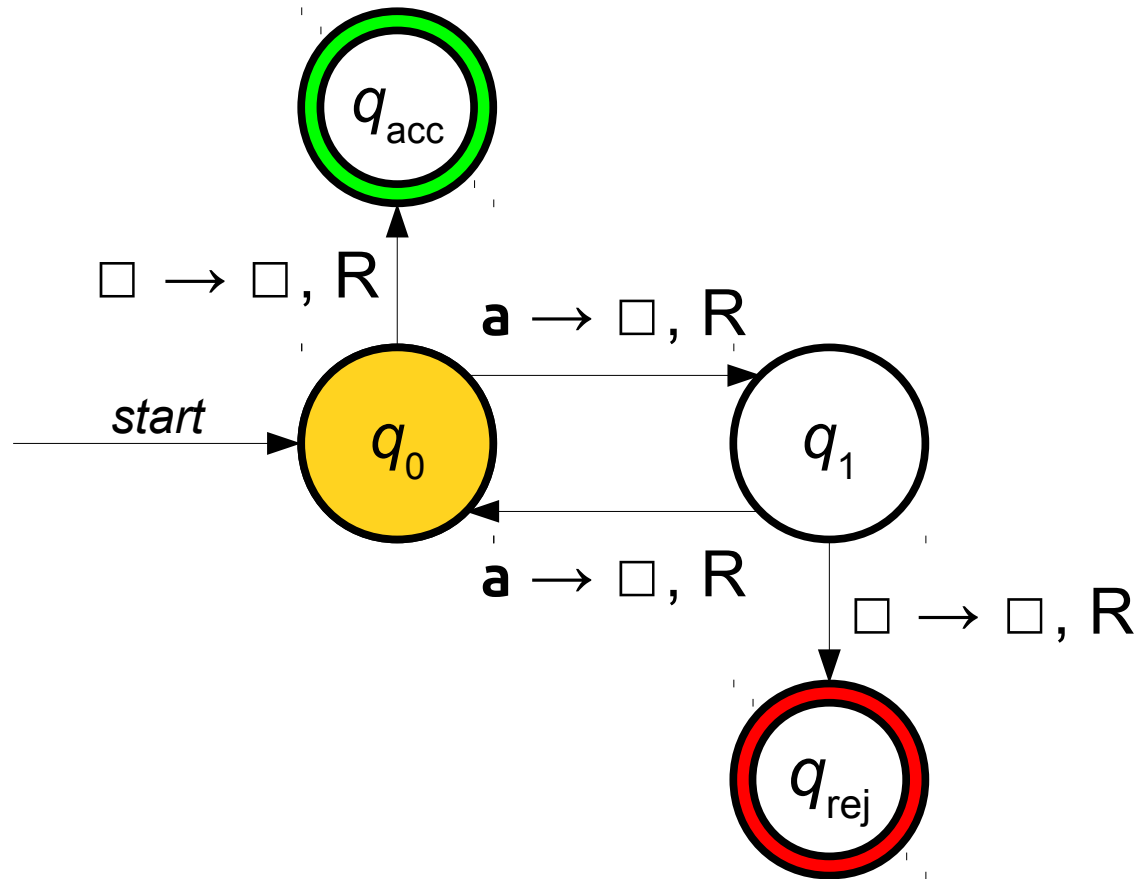
A Simple Turing Machine



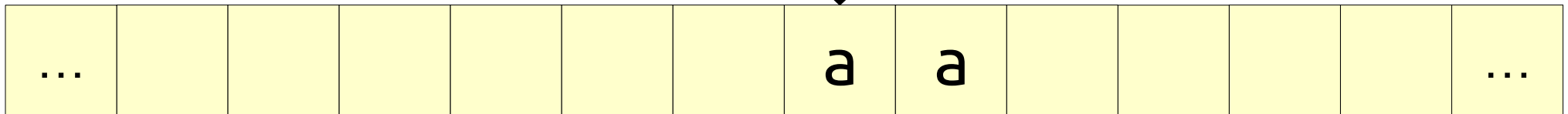
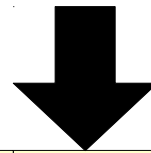
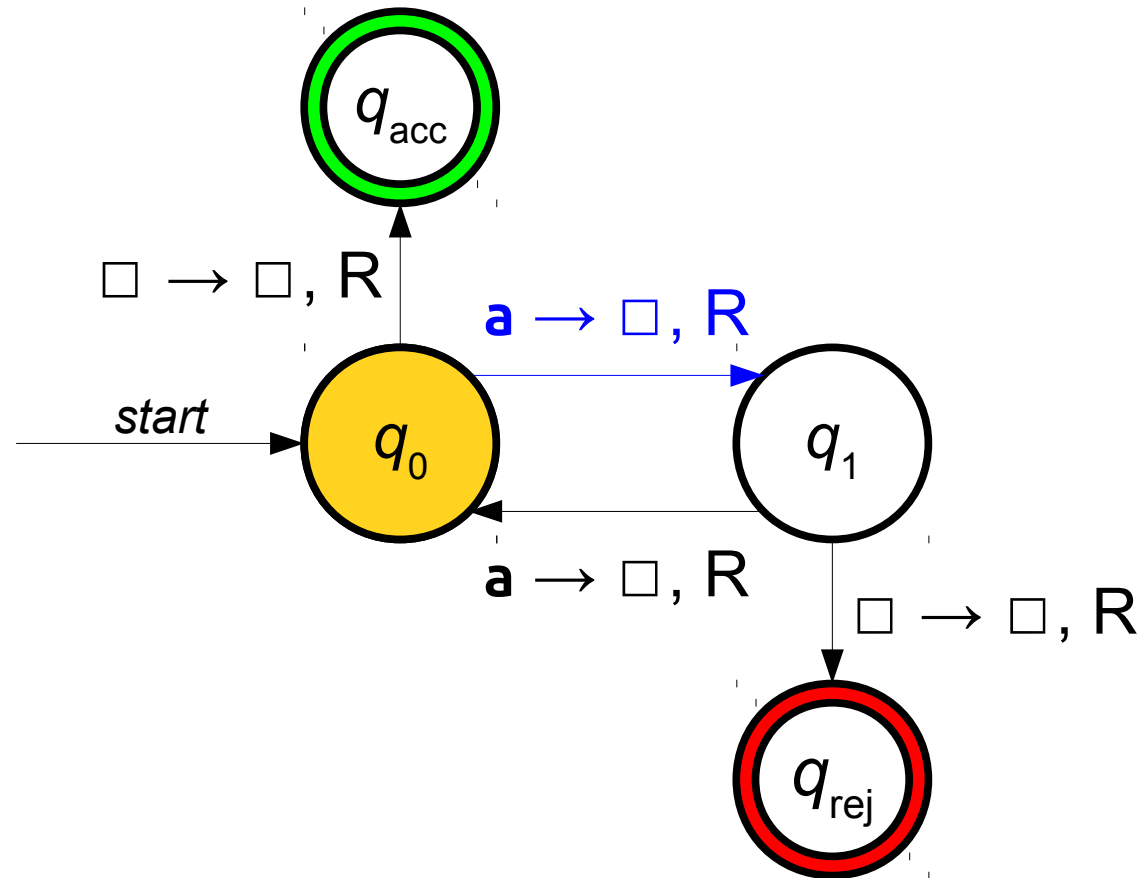
A Simple Turing Machine



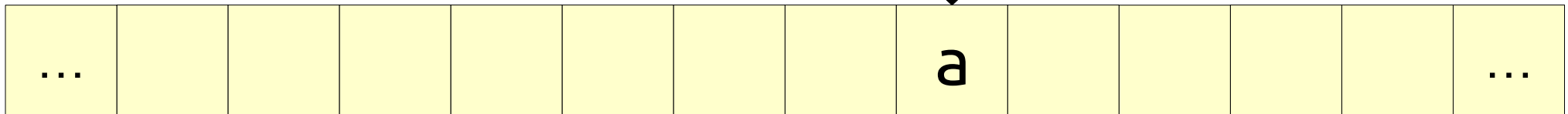
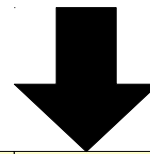
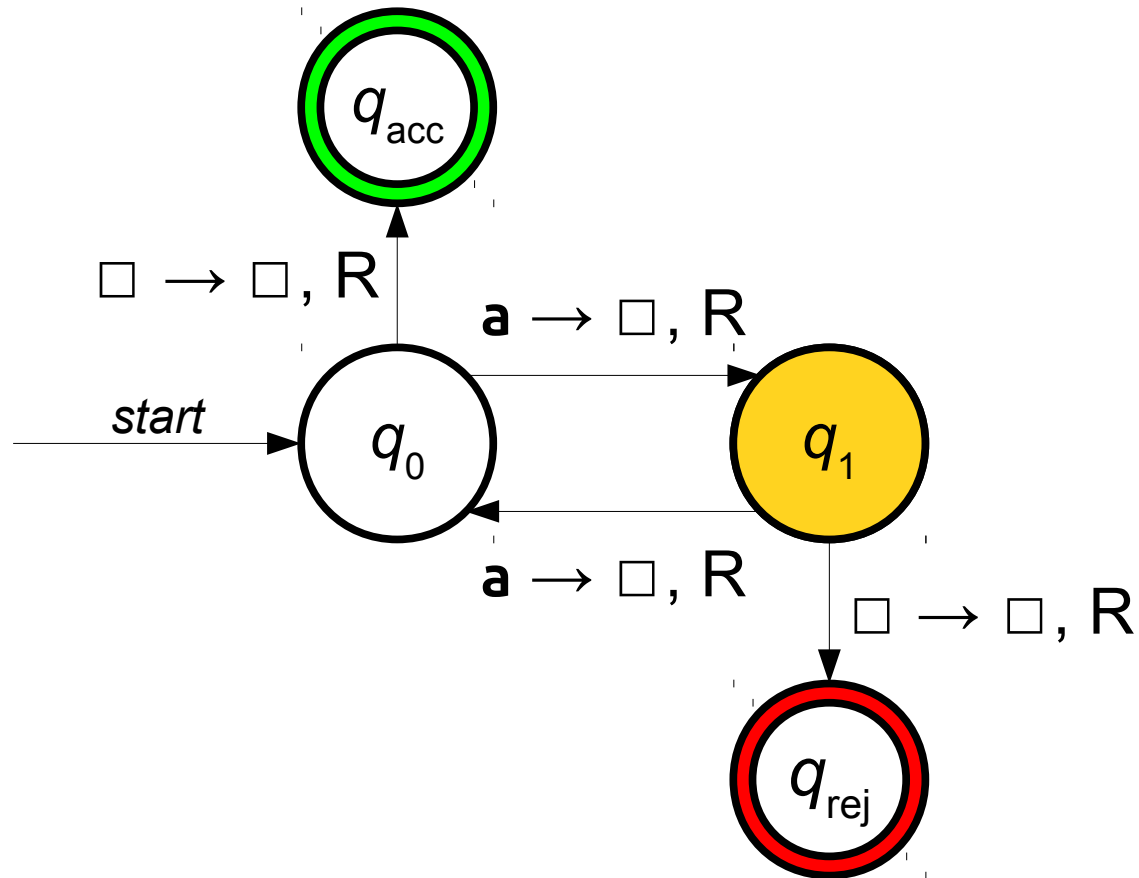
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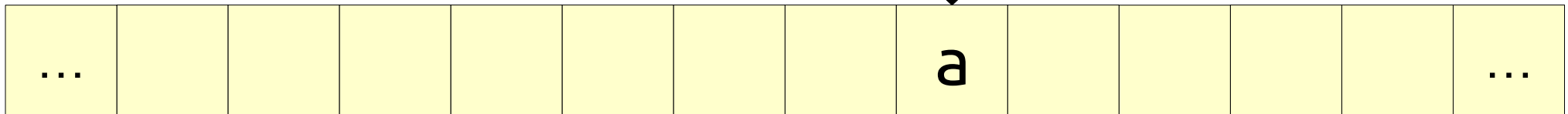
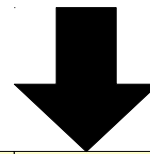
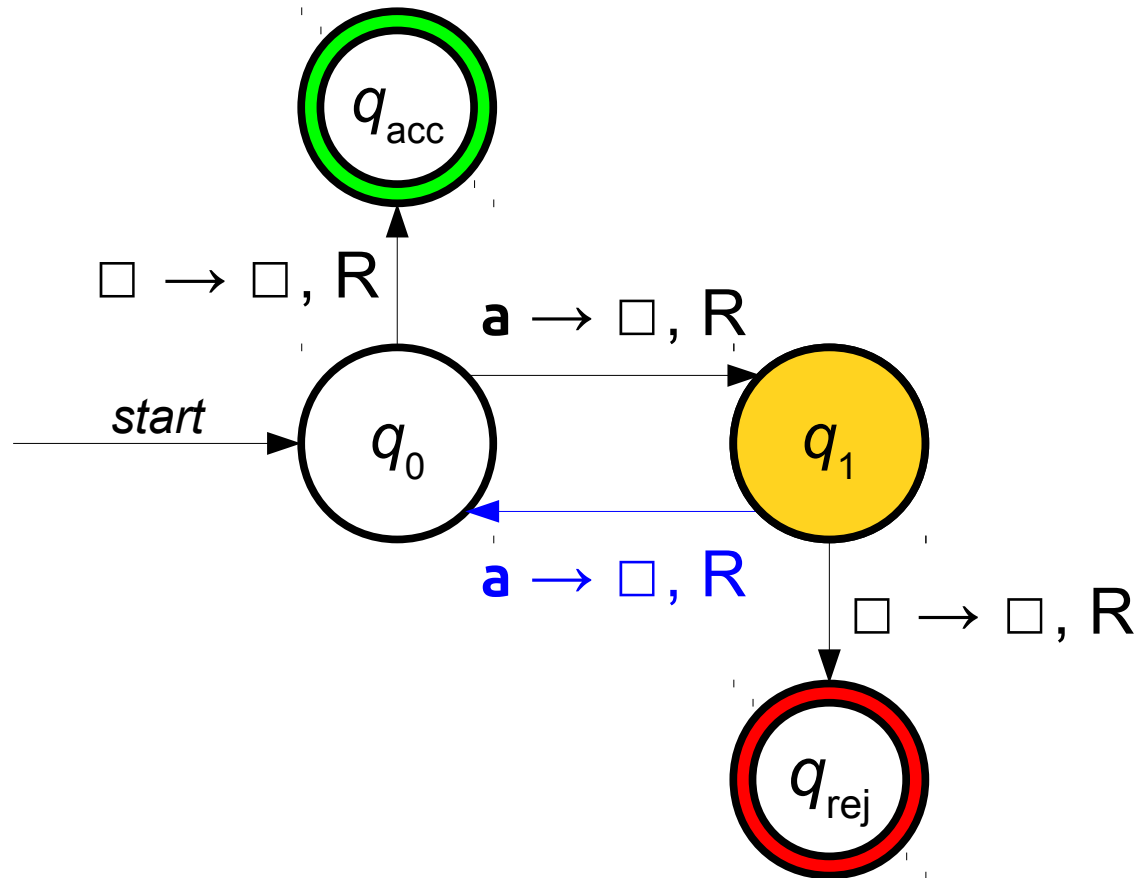
A Simple Turing Machine



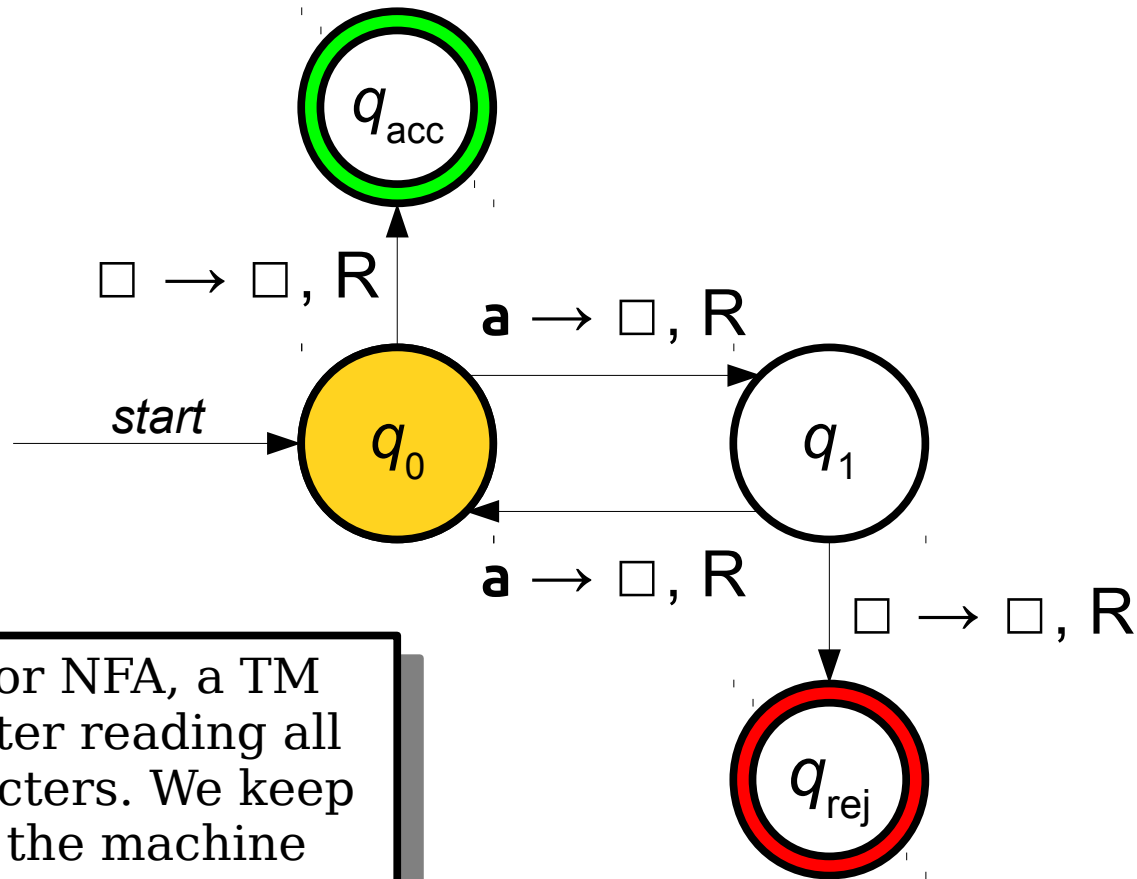
A Simple Turing Machine



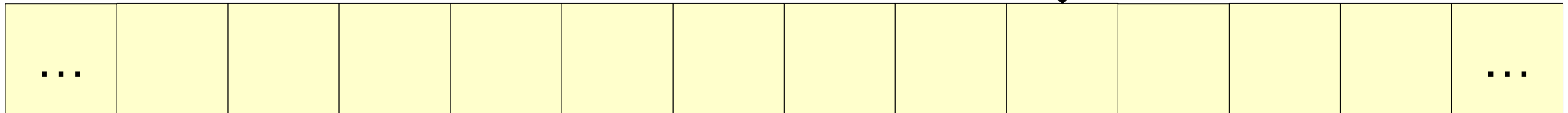
A Simple Turing Machine



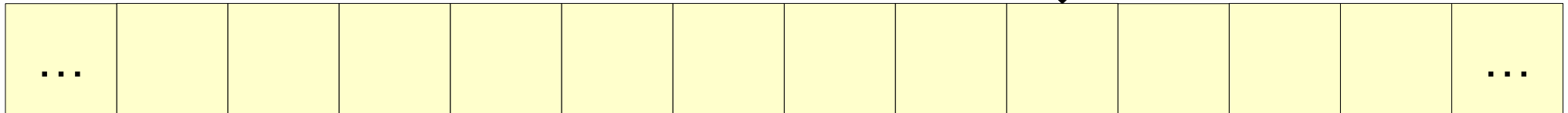
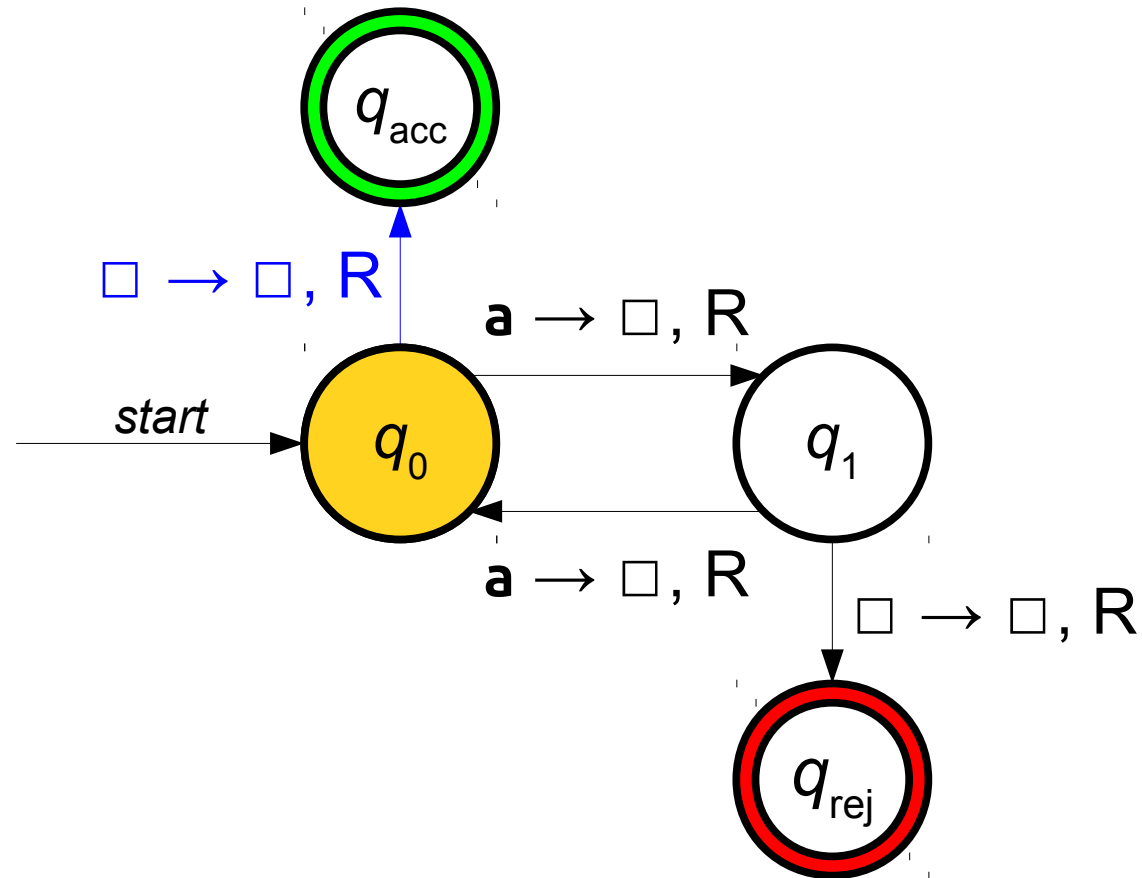
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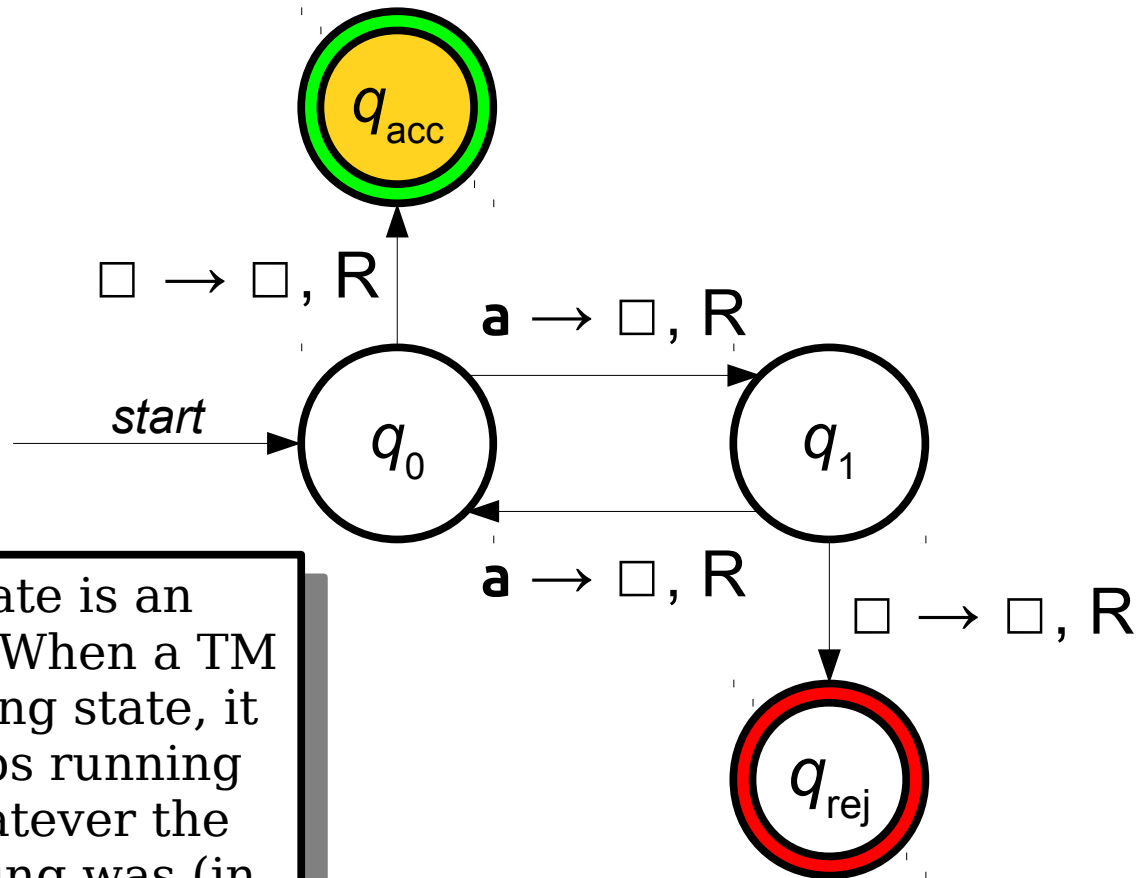
Unlike a DFA or NFA, a TM doesn't stop after reading all the input characters. We keep running until the machine explicitly says to stop.



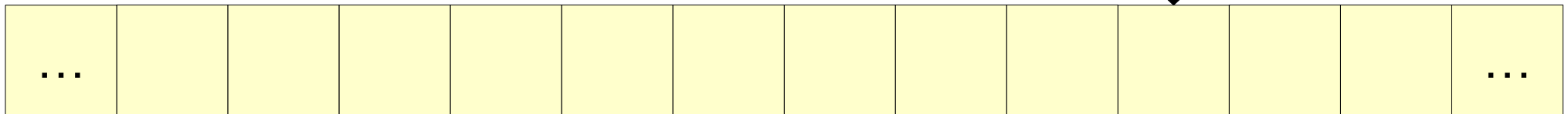
A Simple Turing Machine



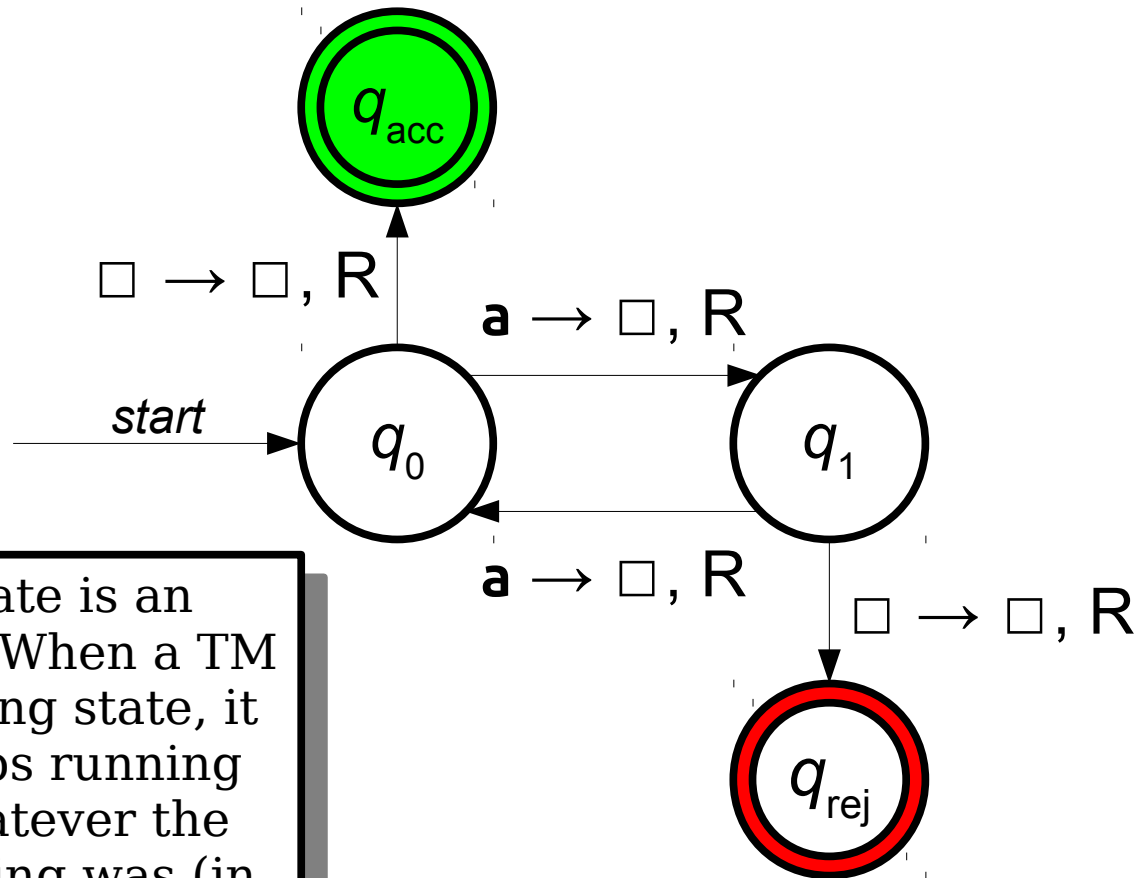
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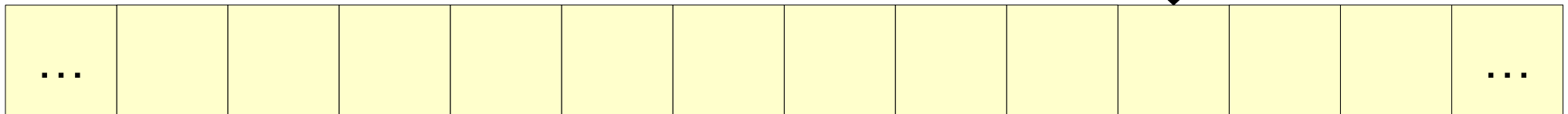
This special state is an **accepting state**. When a TM enters an accepting state, it *immediately* stops running and accepts whatever the original input string was (in this case, **aaaa**).



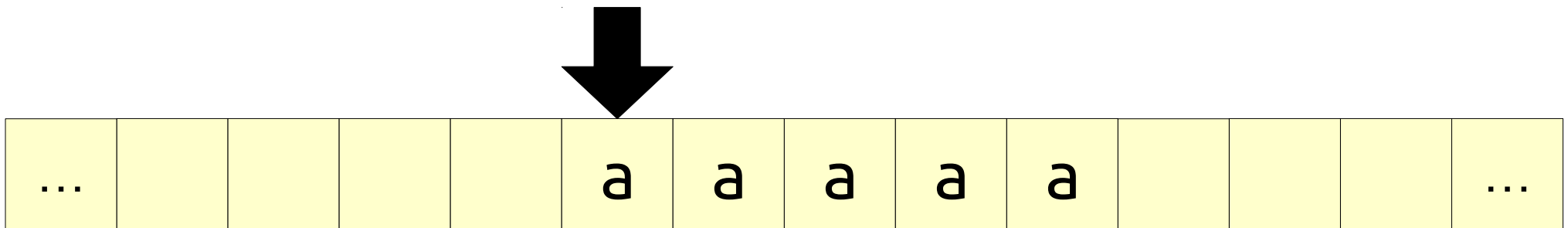
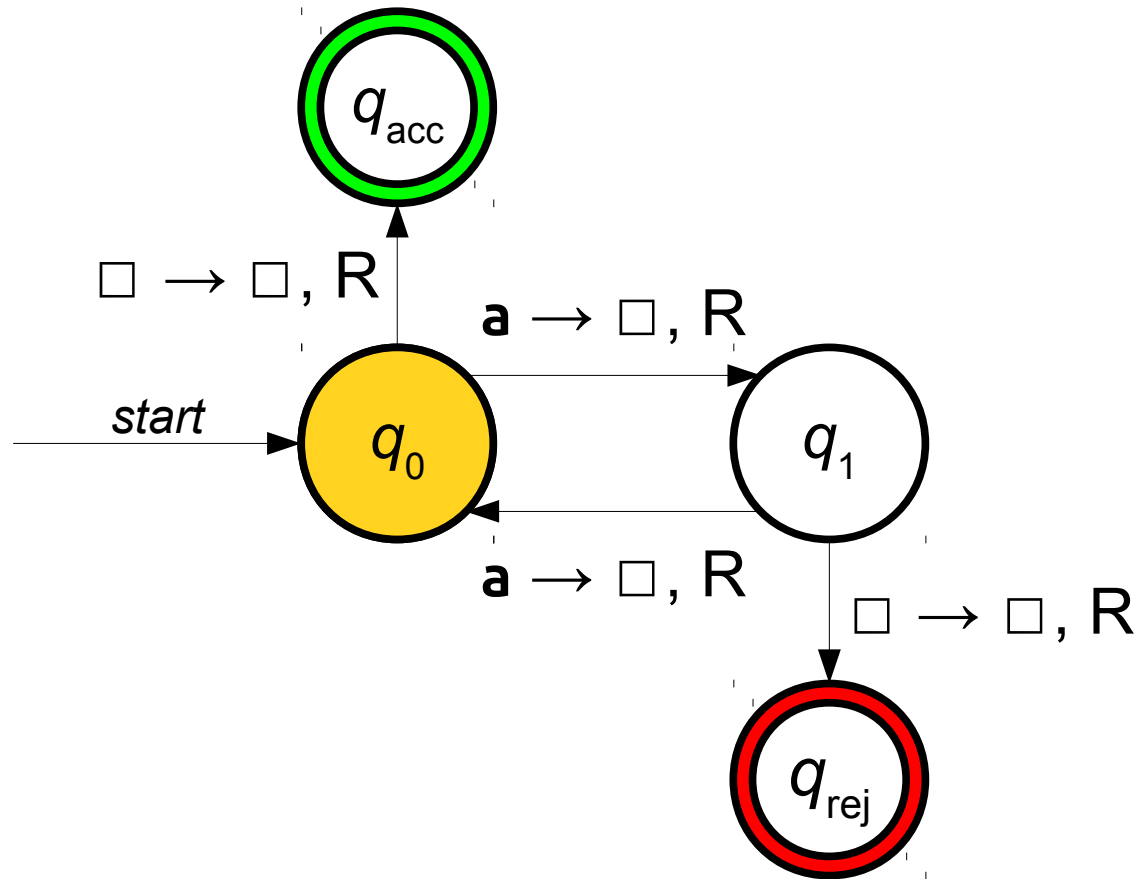
A Simple Turing Machine



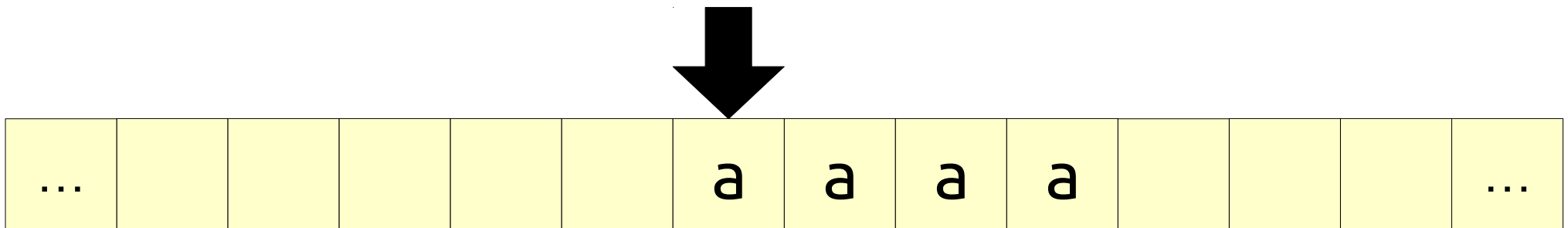
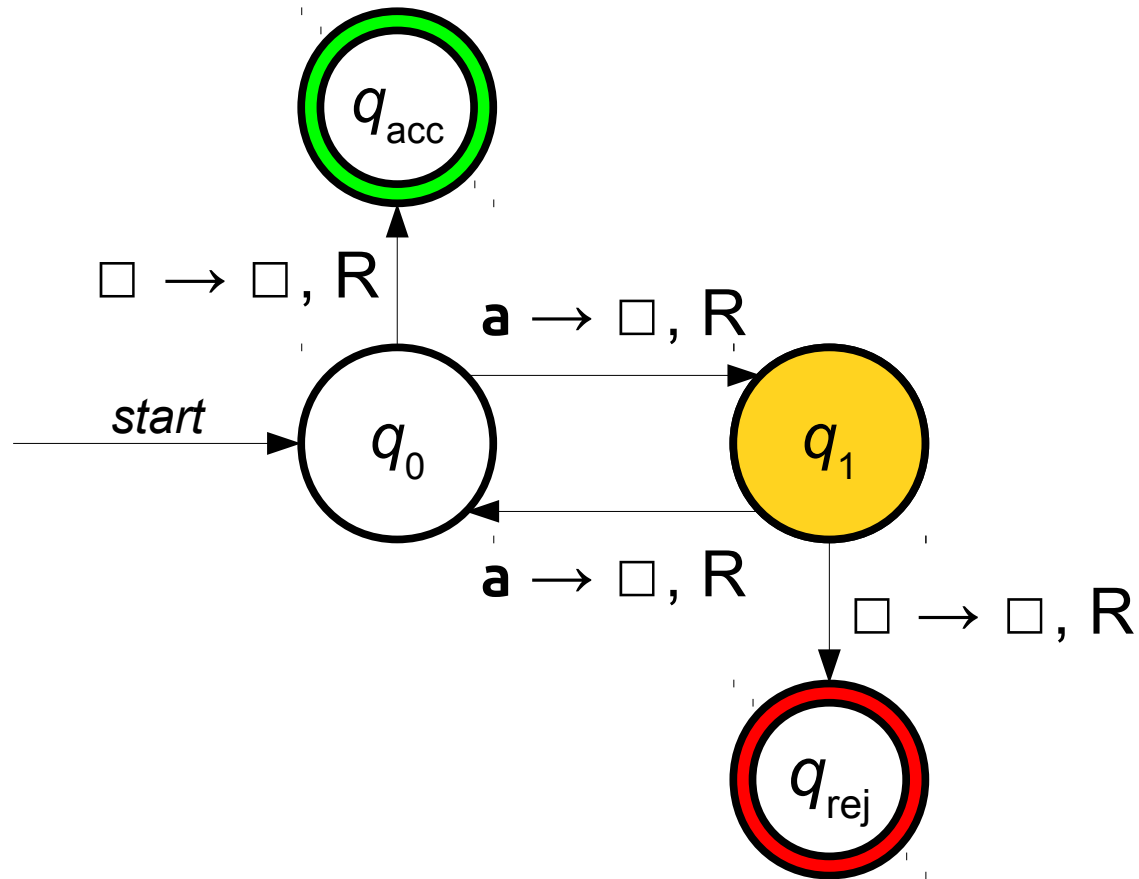
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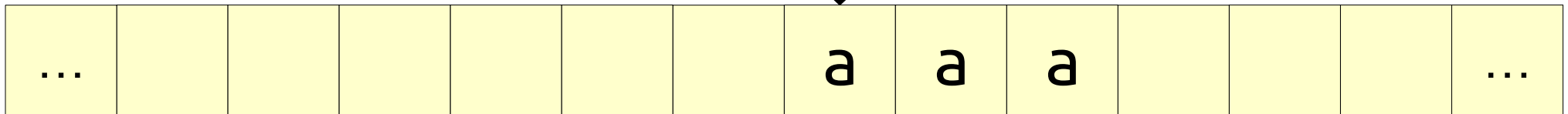
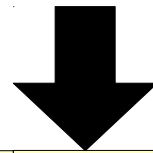
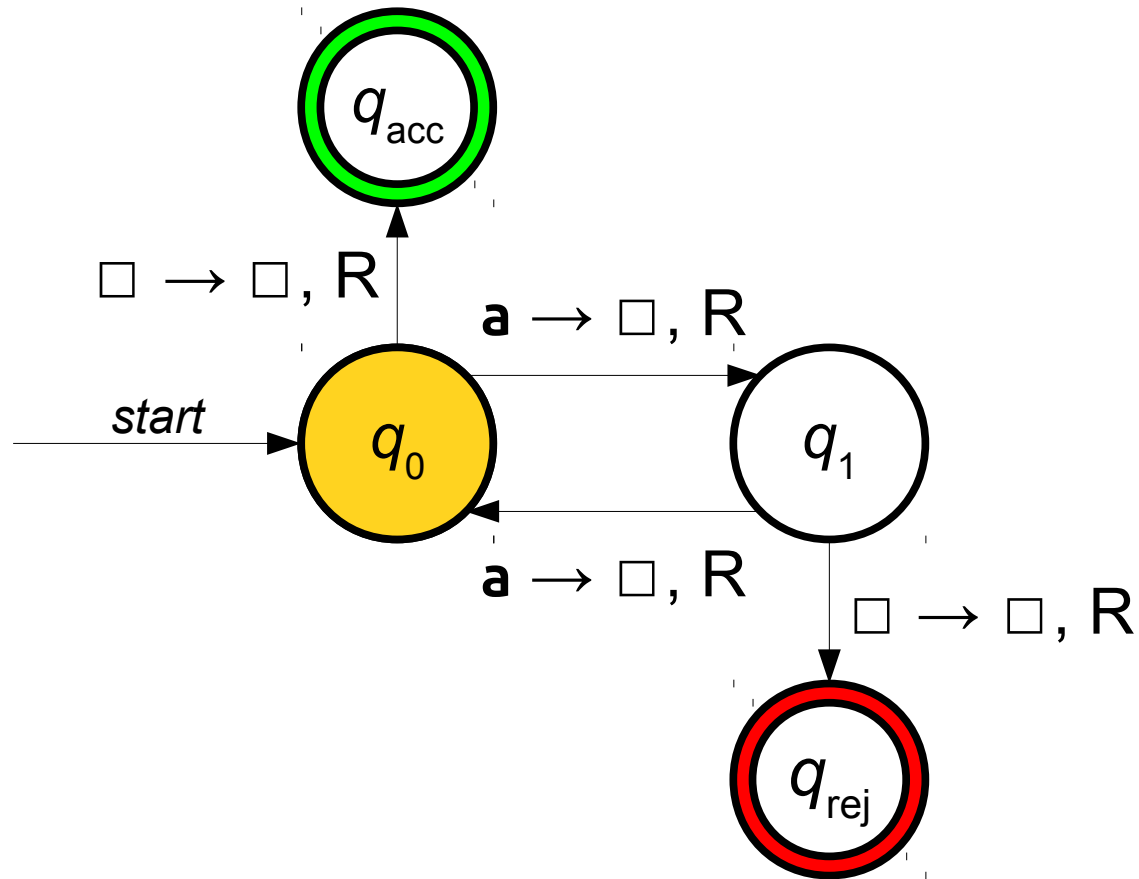
A Simple Turing Machine



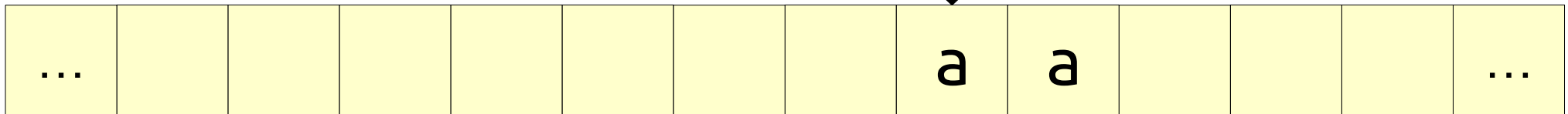
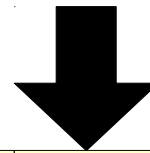
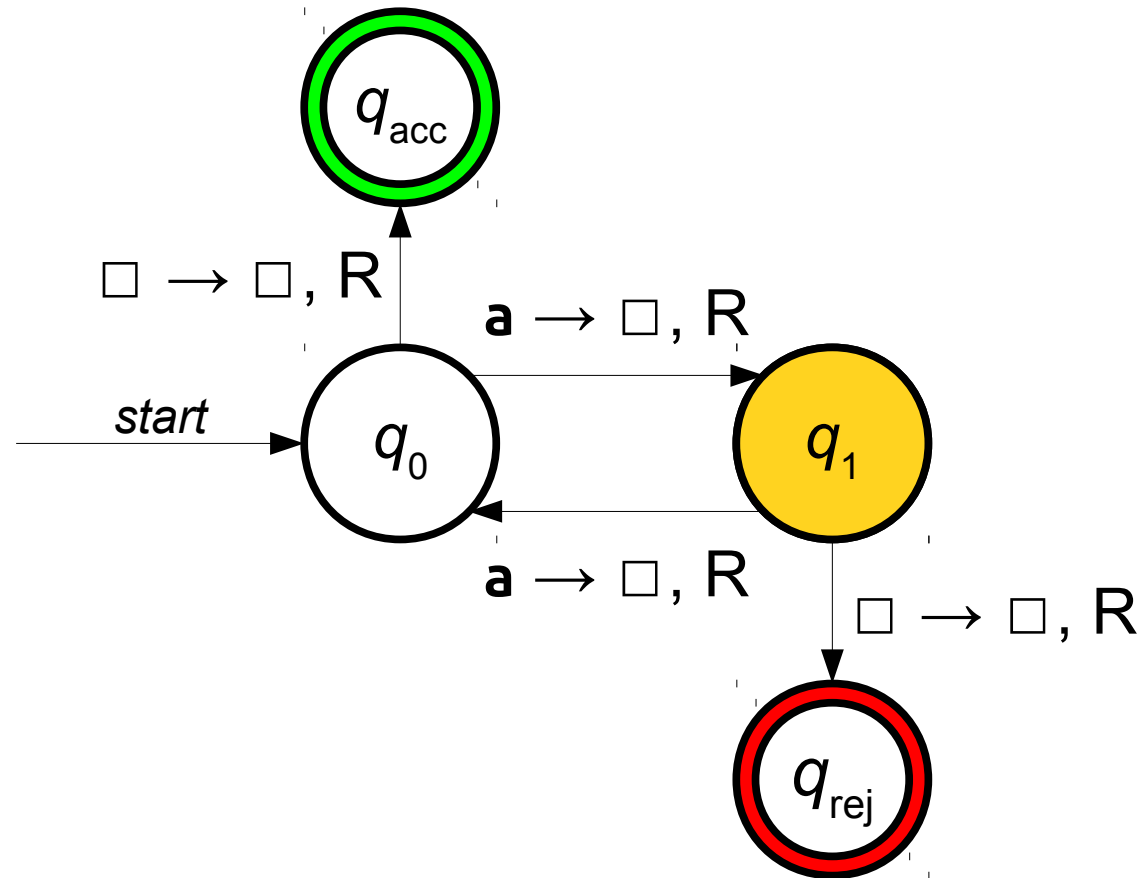
A Simple Turing Machine



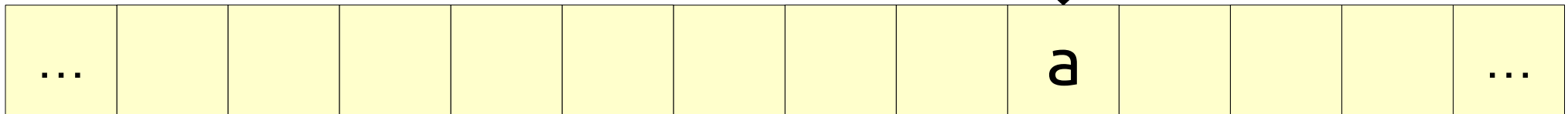
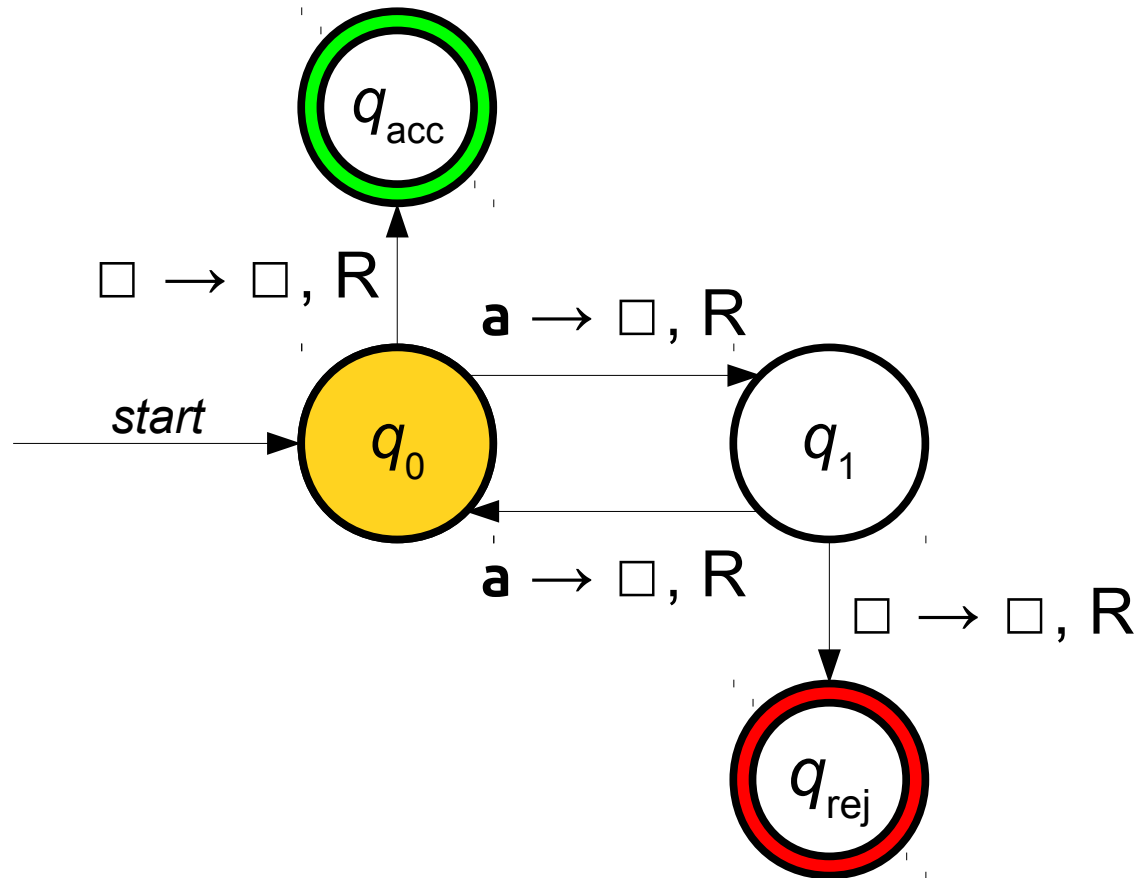
A Simple Turing Machine



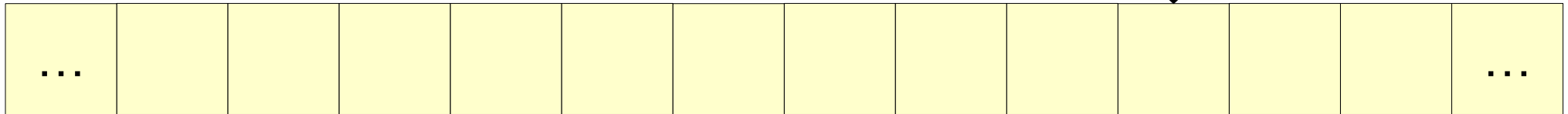
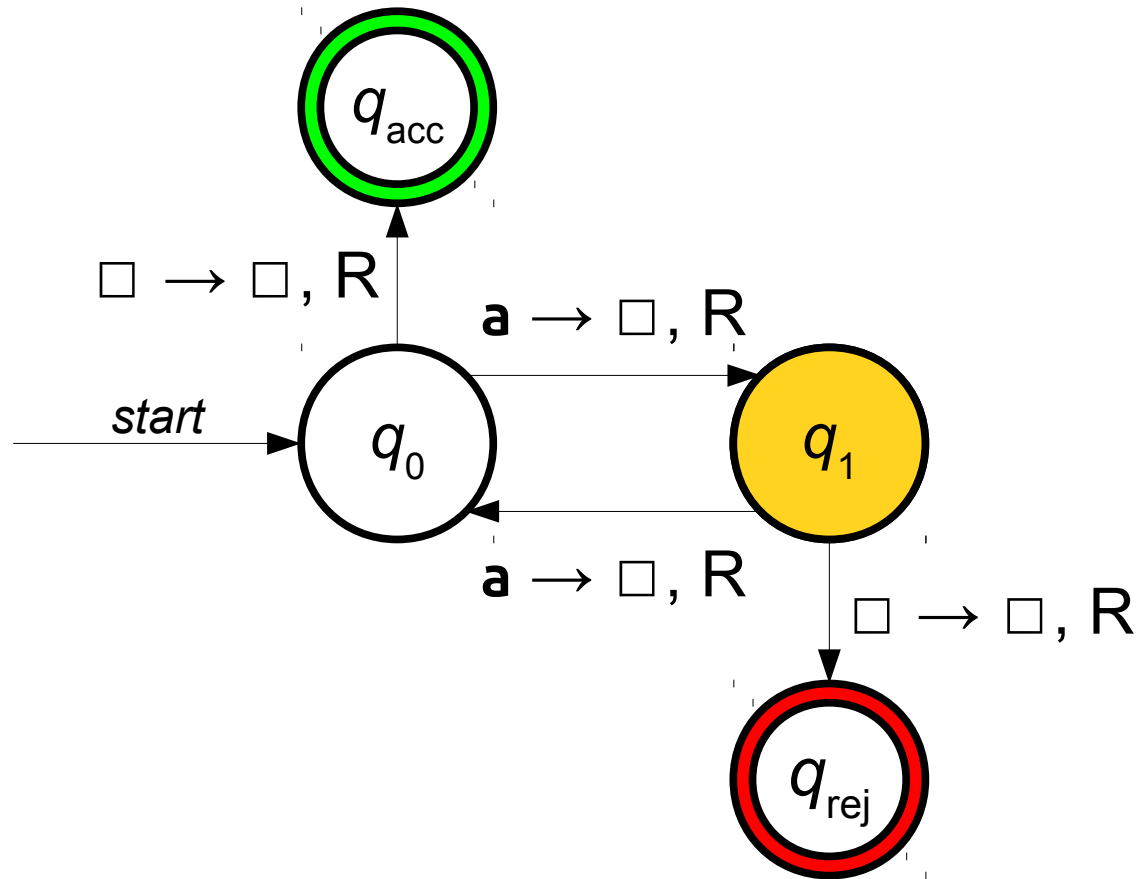
A Simple Turing Machine



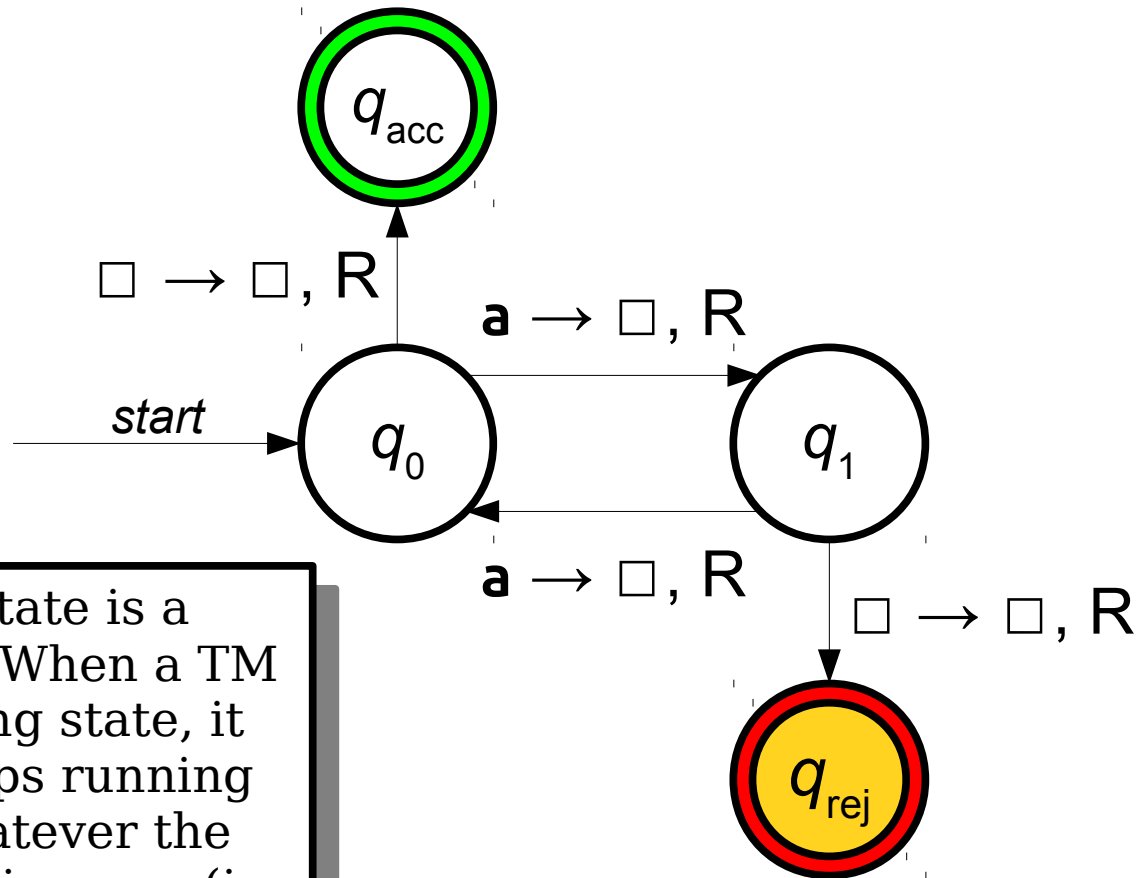
A Simple Turing Machine



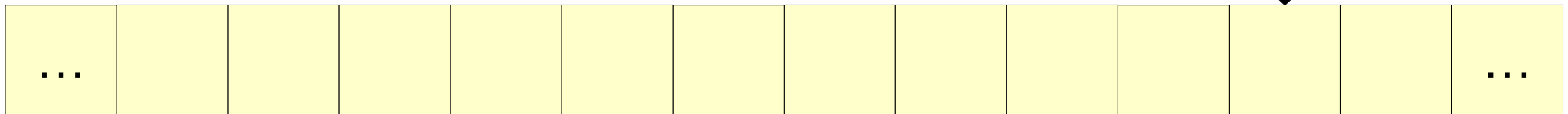
A Simple Turing Machine



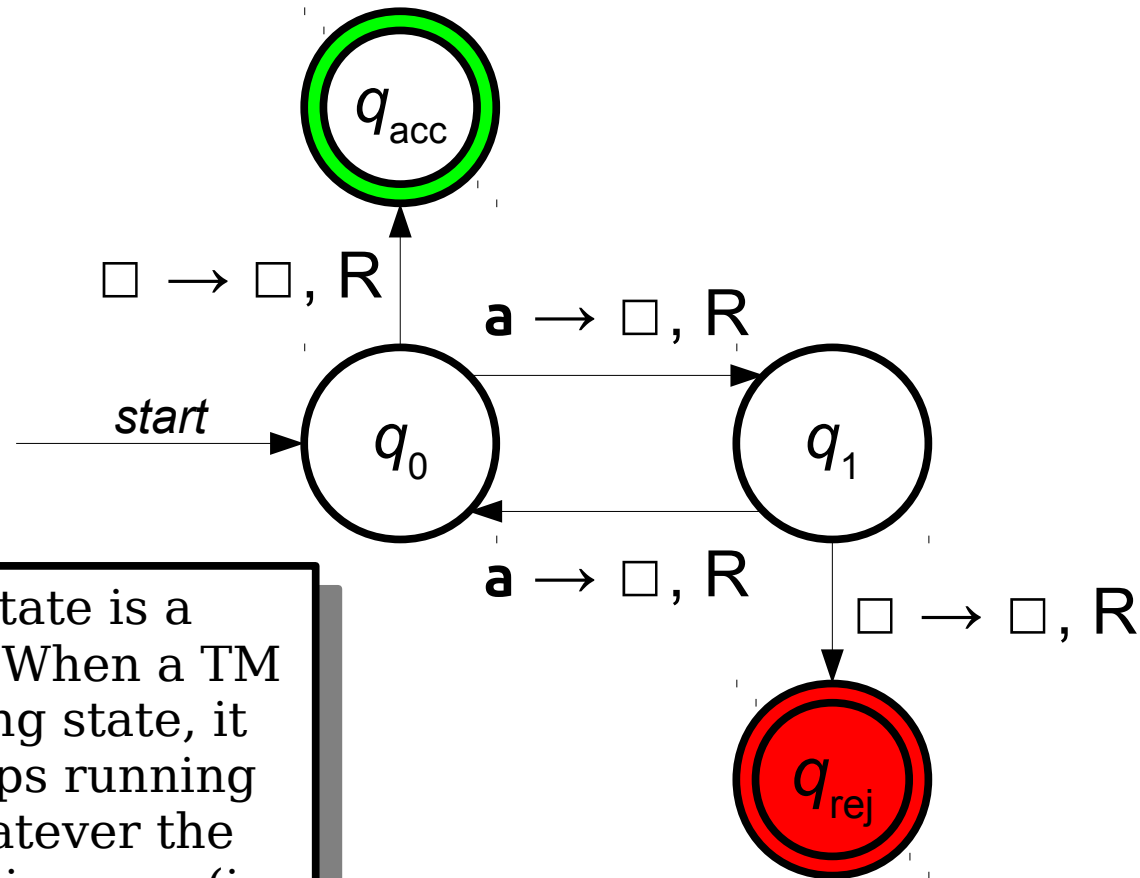
A Simple Turing Machine



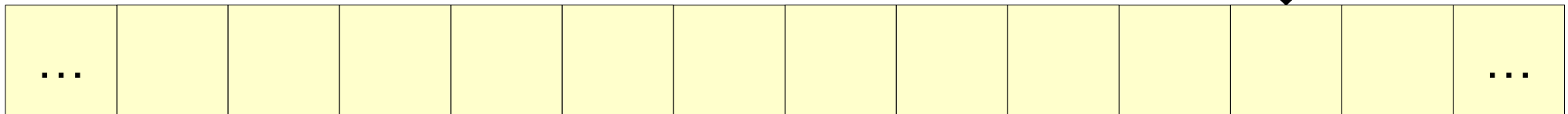
This special state is a **rejecting state**. When a TM enters a rejecting state, it *immediately* stops running and rejects whatever the original input string was (in this case, **aaaaa**).



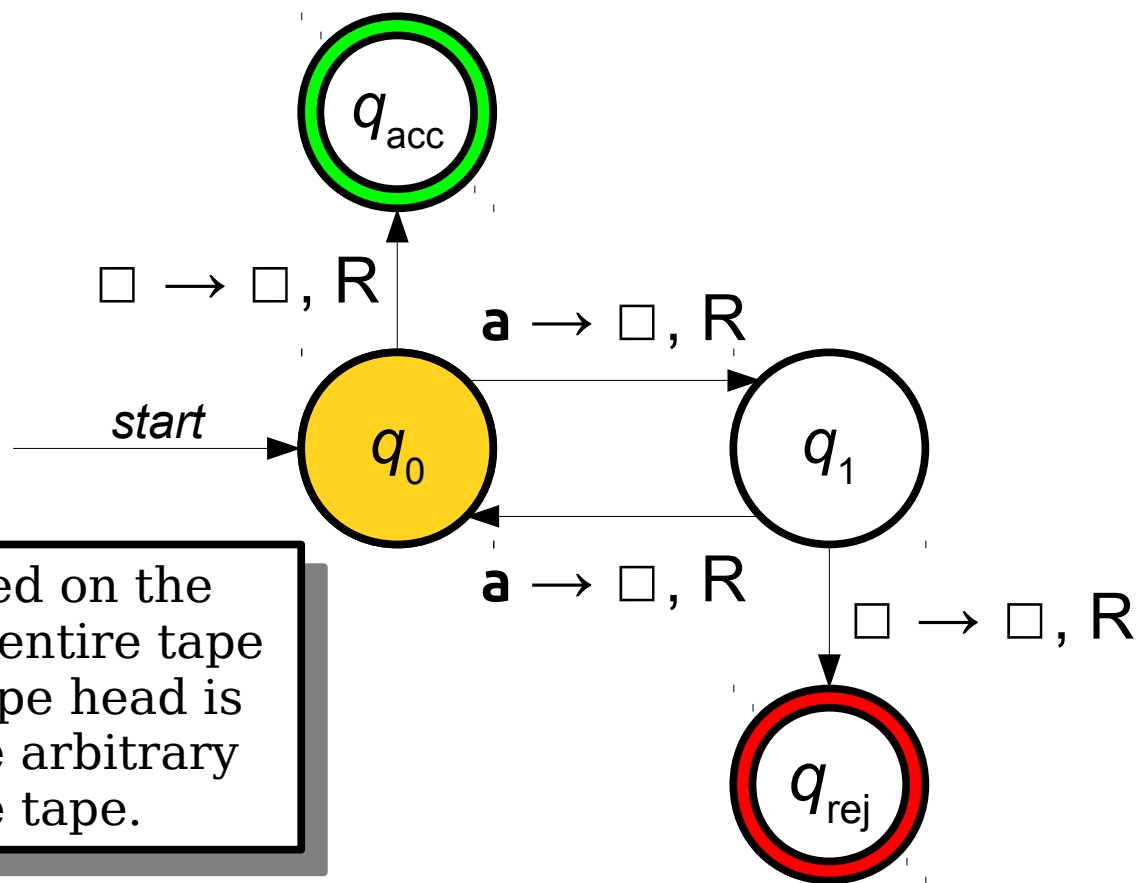
A Simple Turing Machine



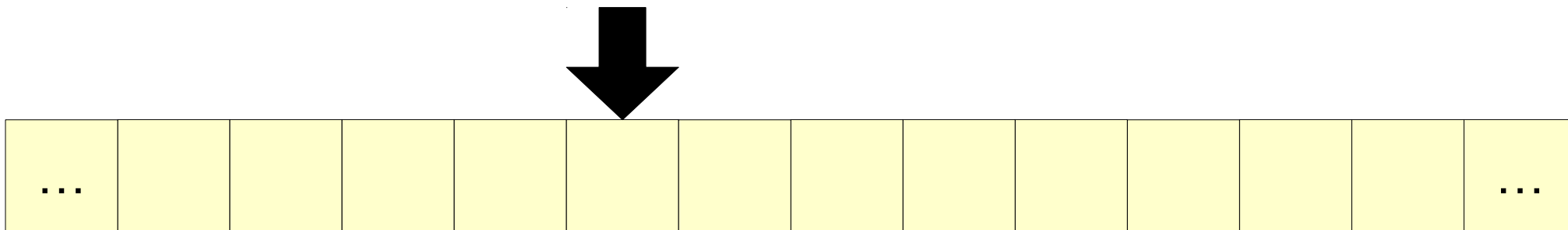
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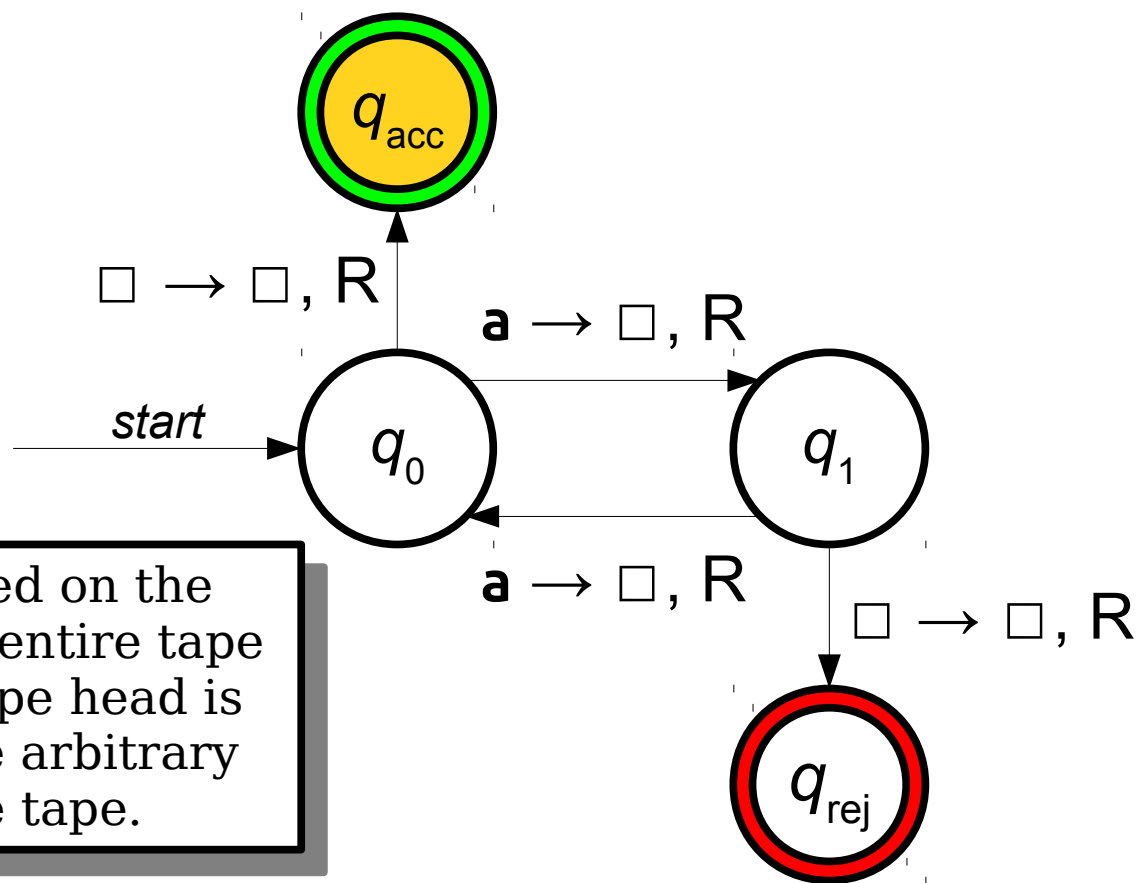
A Simple Turing Machine



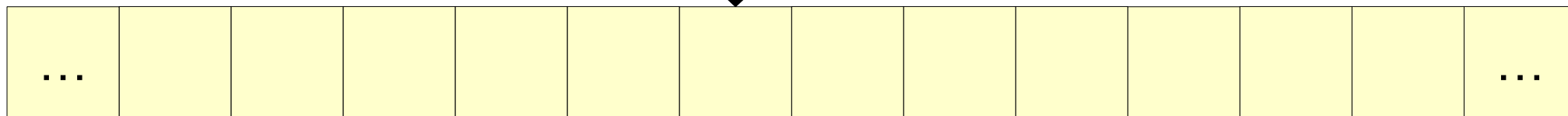
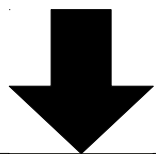
If the TM is started on the empty string ε , the entire tape is blank and the tape head is positioned at some arbitrary location on the tape.



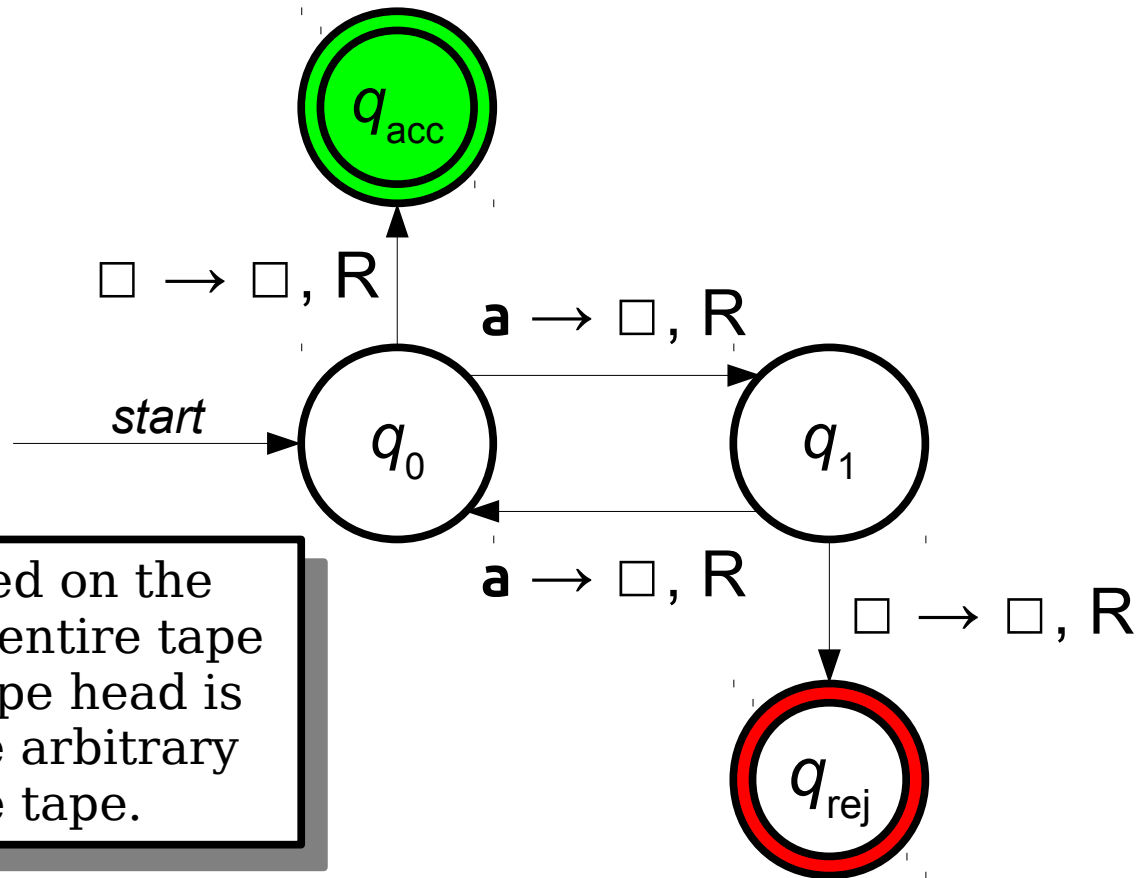
A Simple Turing Machine



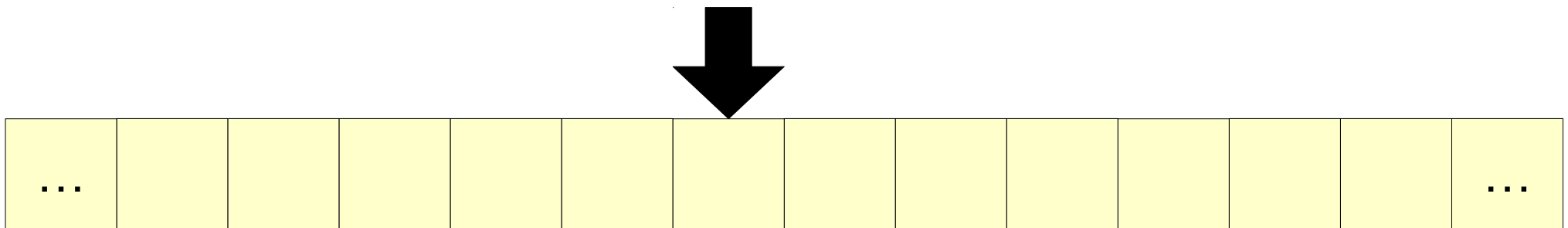
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A Simple Turing Machine



If the TM is started on the empty string ε , the entire tape is blank and the tape head is positioned at some arbitrary location on the tape.



The Turing Machine

- A Turing machine consists of three parts:
 - A *finite-state control* that issues commands,
 - an *infinite tape* for input and scratch space, and
 - a *tape head* that can read and write a single tape cell.
- At each step, the Turing machine
 - writes a symbol to the tape cell under the tape head,
 - changes state, and
 - moves the tape head to the left or to the right.

Input and Tape Alphabets

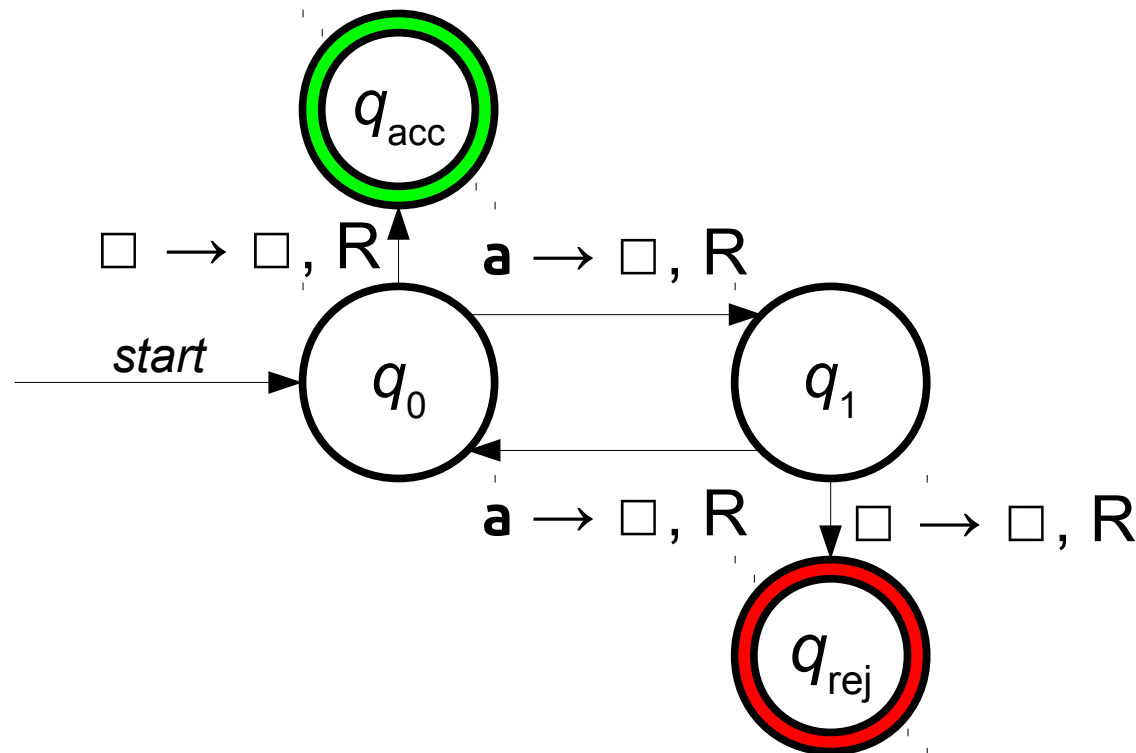
- A Turing machine has two alphabets:
 - An **input alphabet** Σ . All input strings are written in the input alphabet.
 - A **tape alphabet** Γ , where $\Sigma \subseteq \Gamma$. The tape alphabet contains all symbols that can be written onto the tape.
- The tape alphabet Γ can contain any number of symbols, but always contains at least one **blank symbol**, denoted \square . You are guaranteed $\square \notin \Sigma$.
- At startup, the Turing machine begins with an infinite tape of \square symbols with the input written at some location. The tape head is positioned at the start of the input.

Accepting and Rejecting States

- Unlike DFAs, Turing machines do not stop processing the input when they finish reading it.
- Turing machines decide when (and if!) they will accept or reject their input.
- Turing machines can enter infinite loops and never accept or reject; more on that later...

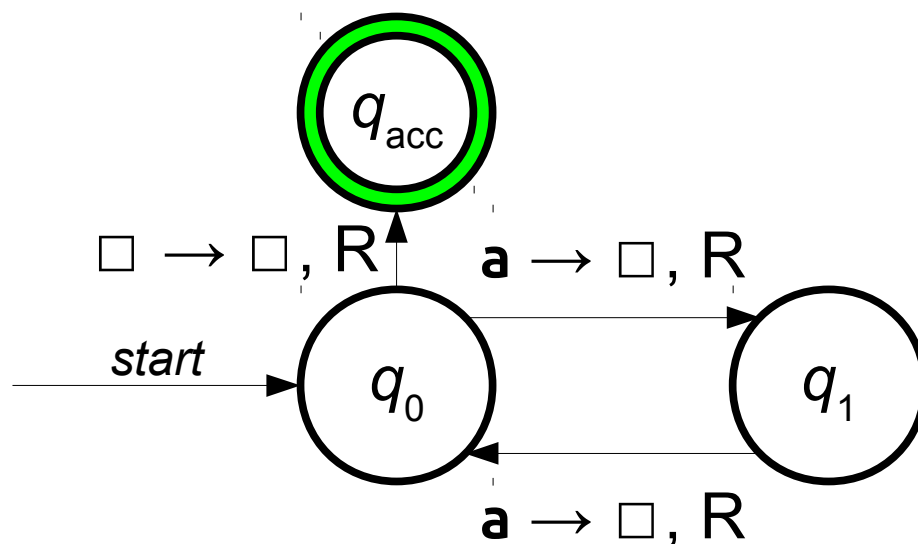
Determinism

- Turing machines are **deterministic**: for every combination of a (non-accepting, non-rejecting) state q and a tape symbol $a \in \Gamma$, there must be exactly one transition defined for that combination of q and a .
- Any transitions that are missing implicitly go straight to a rejecting state. We'll use this later to simplify our designs.

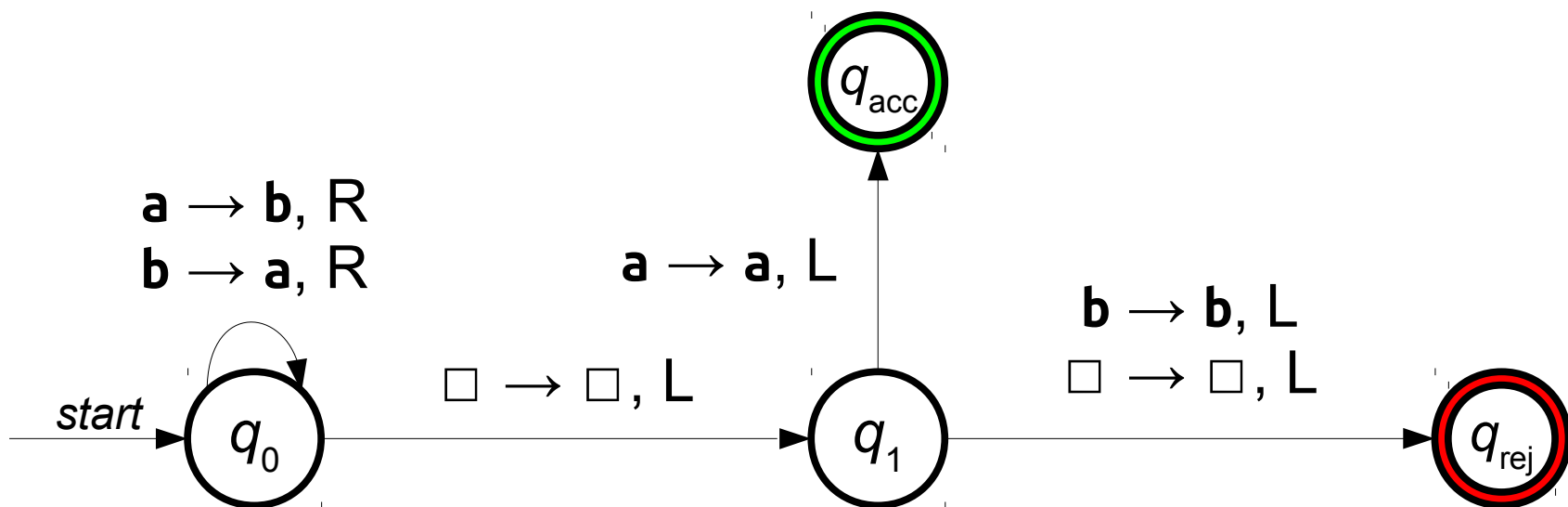


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This machine is exactly the same as the previous one.



Run the TM shown above on the input string **bba**.
What will the tape look like when the TM finishes running?

A.

...		b	b	a		...
-----	--	---	---	---	--	-----

B.

...		a	a	b		...
-----	--	---	---	---	--	-----

C.

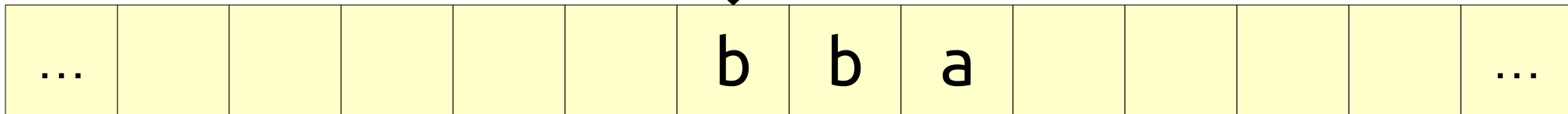
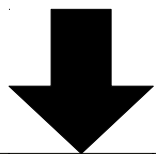
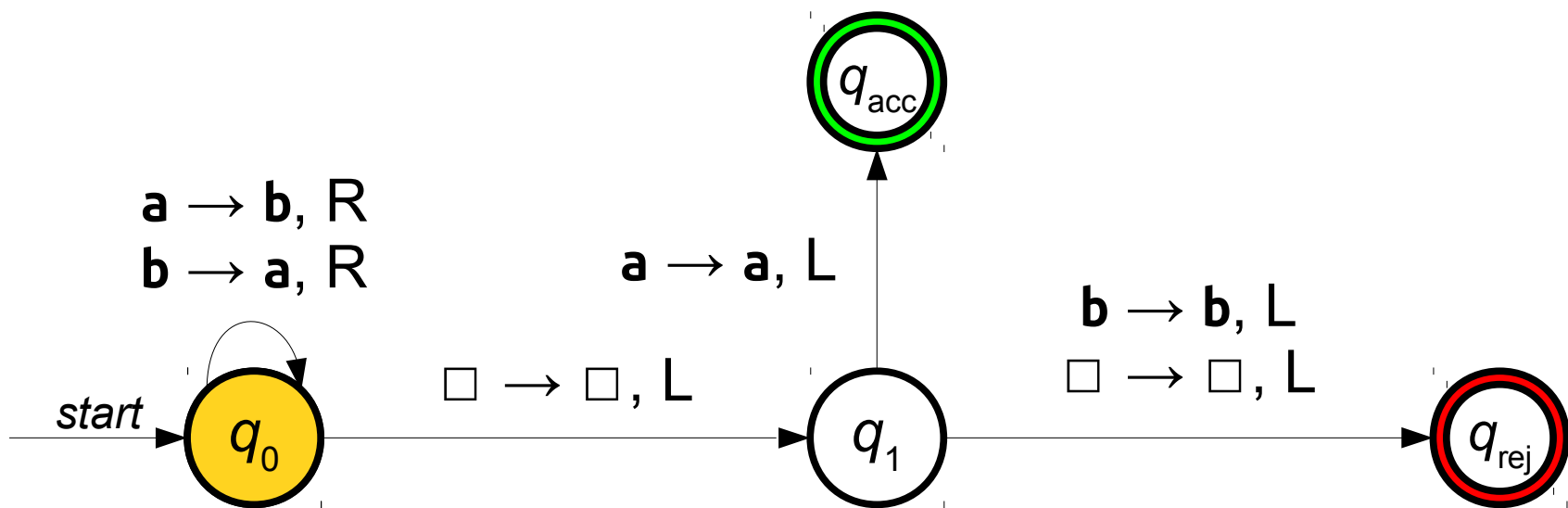
...		b	b	a		...
-----	--	---	---	---	--	-----

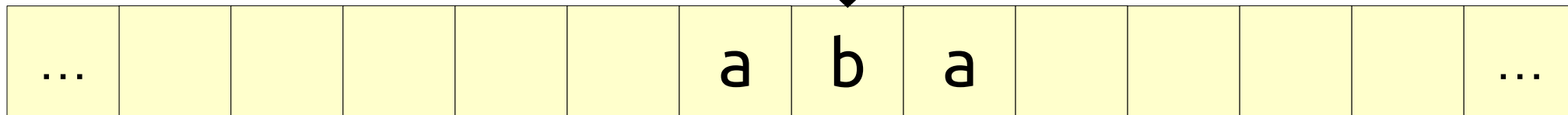
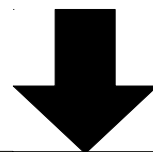
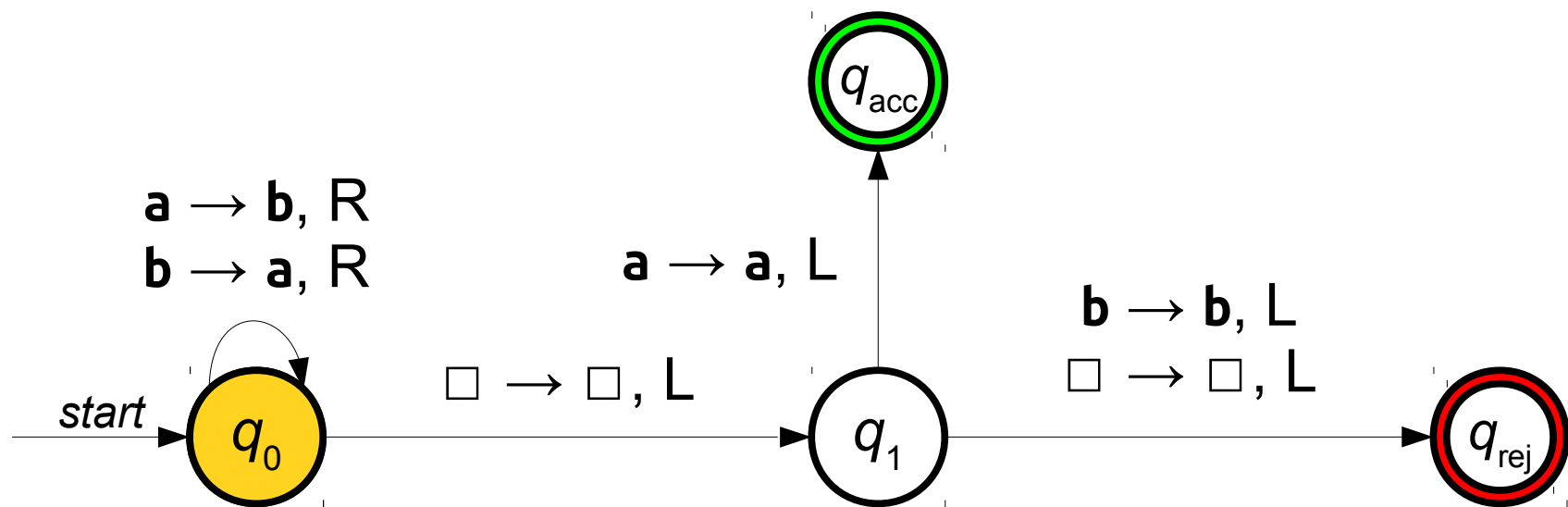
D.

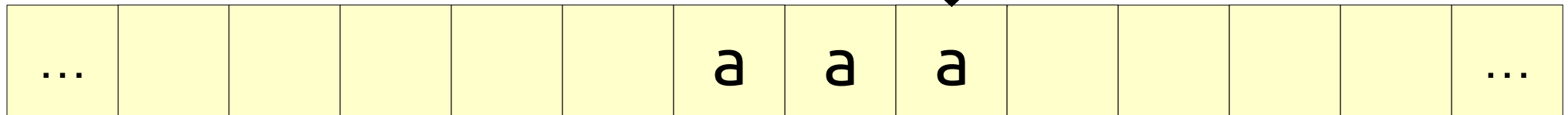
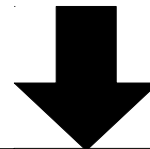
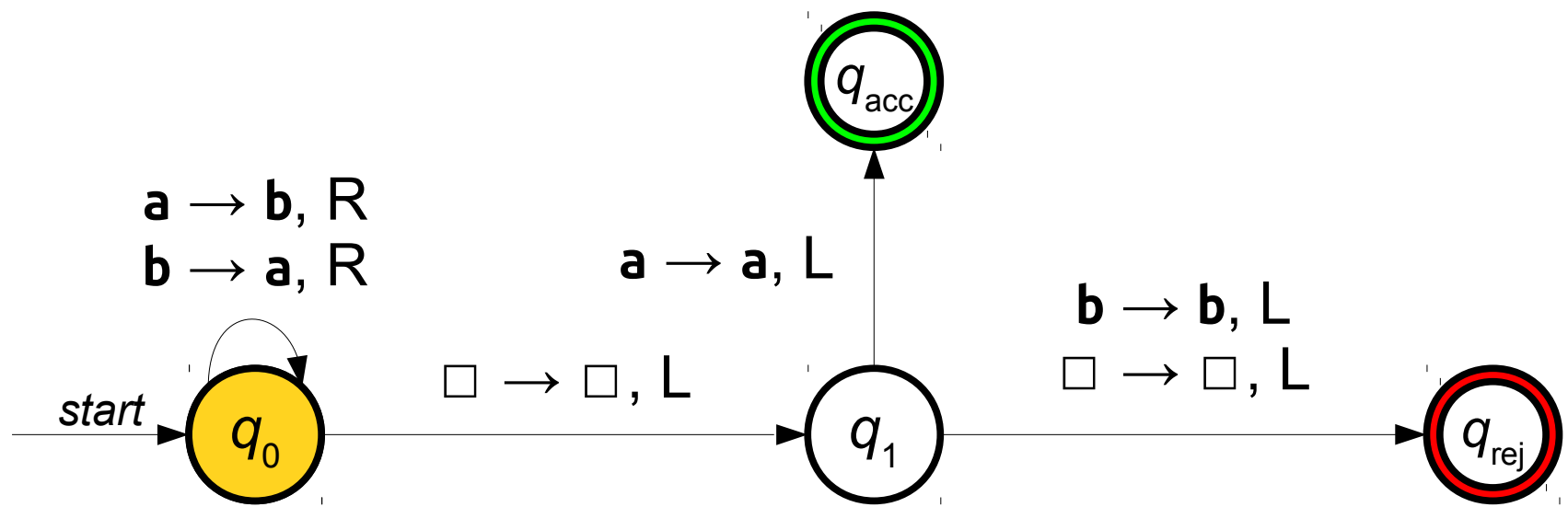
...		a	a	b		...
-----	--	---	---	---	--	-----

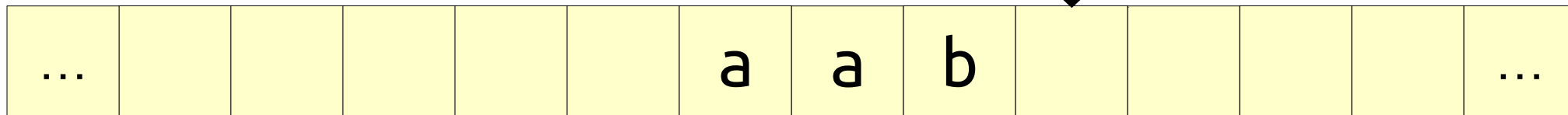
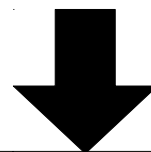
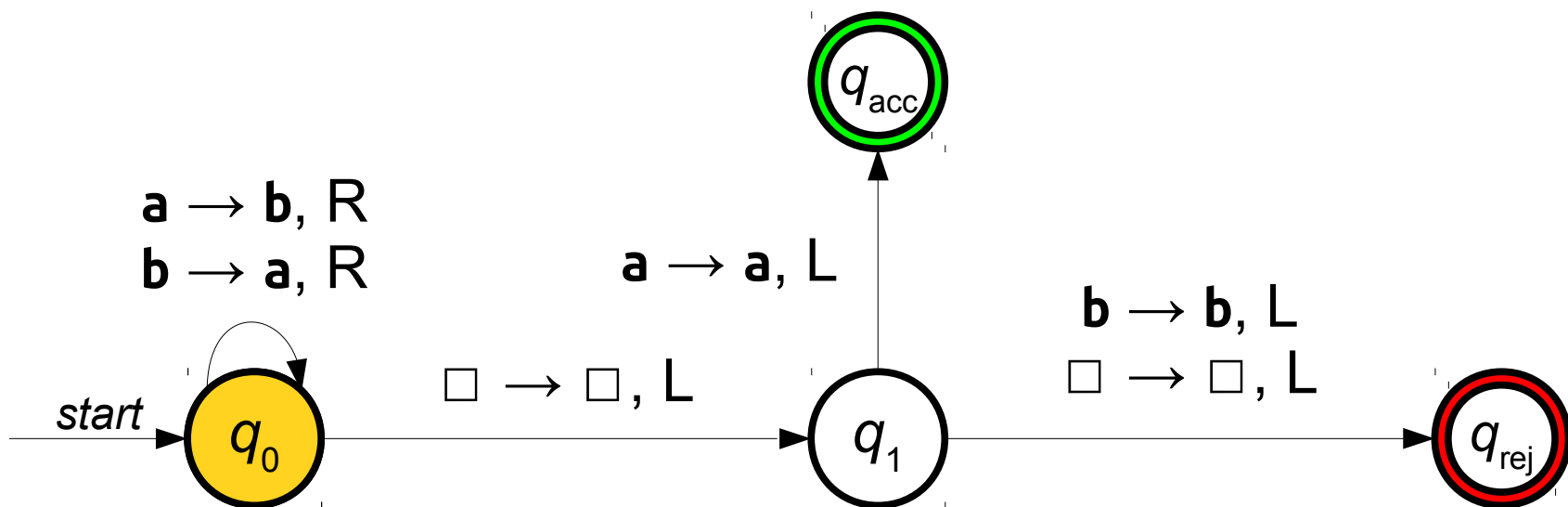
E. None of these, or two or more of these.

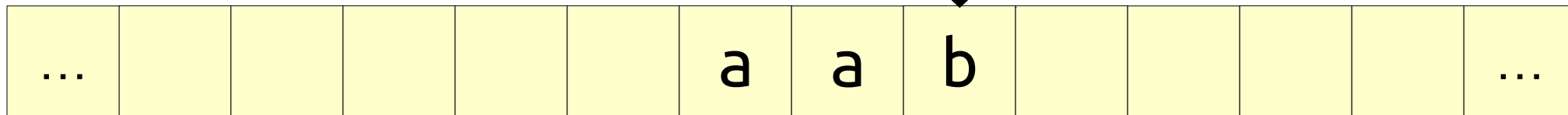
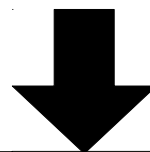
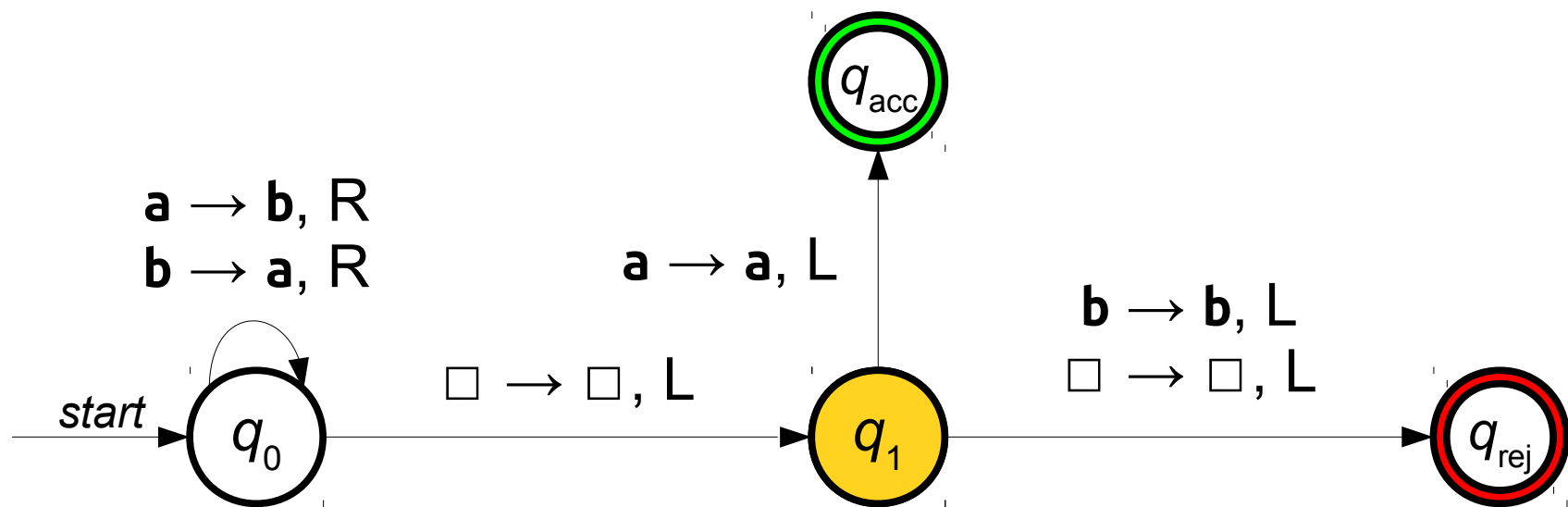
Answer at [PollEv.com/cs103](https://pollev.com/cs103) or
text **CS103** to **22333** once to join, then **A**, **B**, **C**, **D**, or **E**.

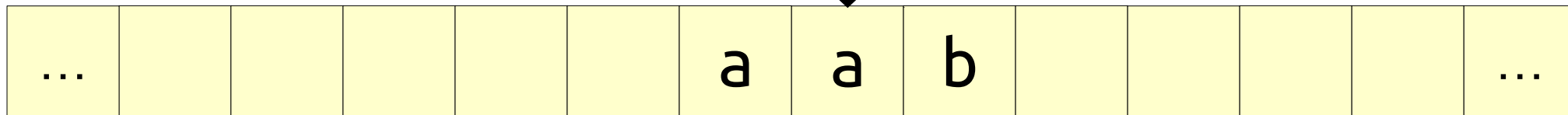
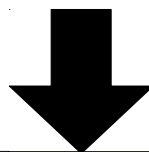
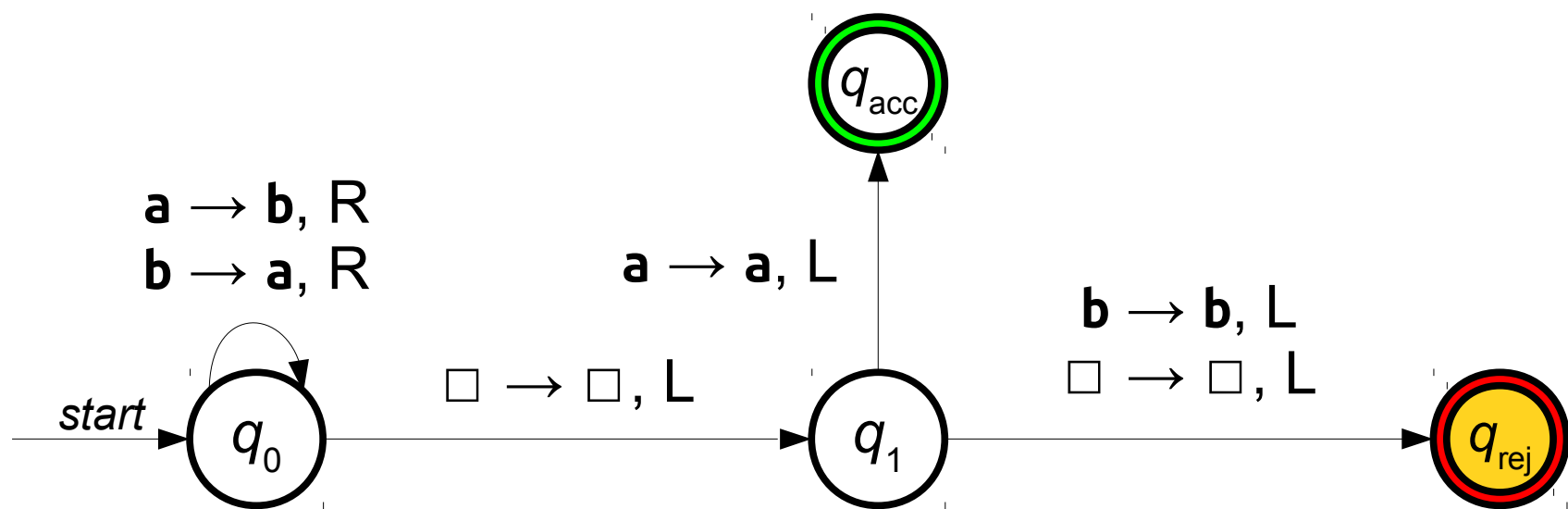


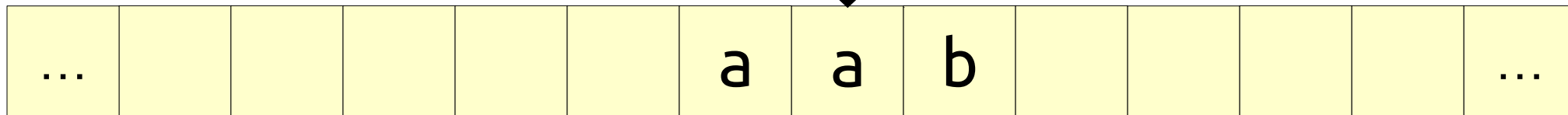
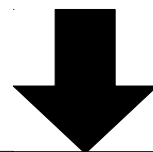
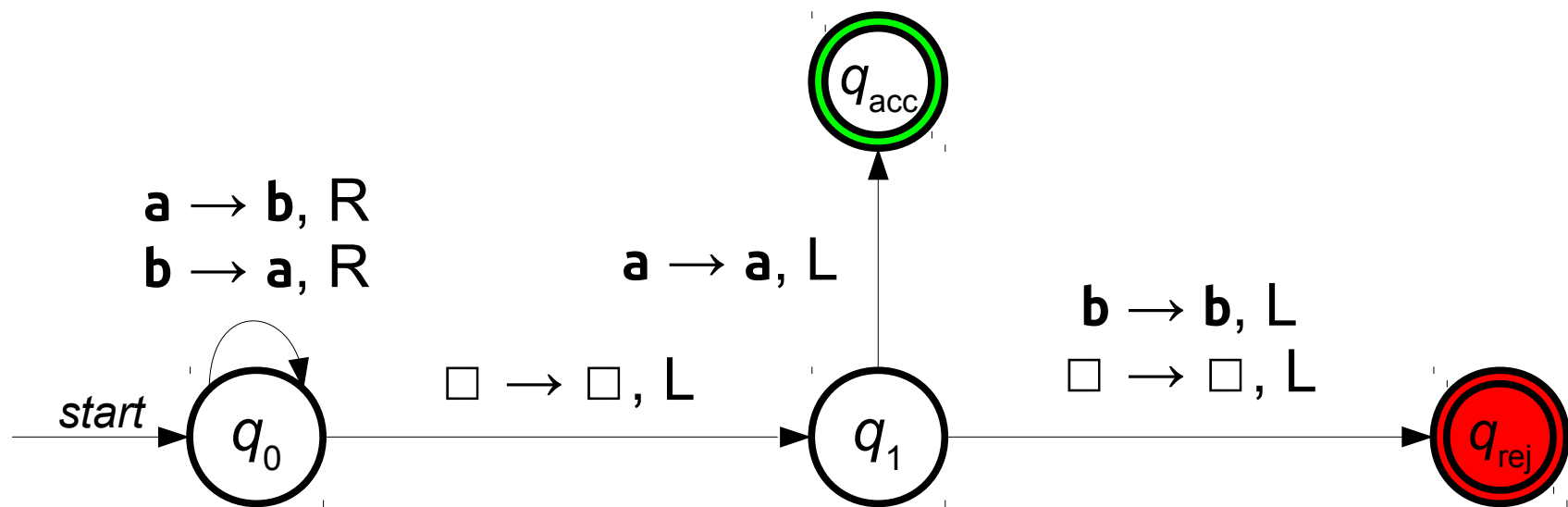


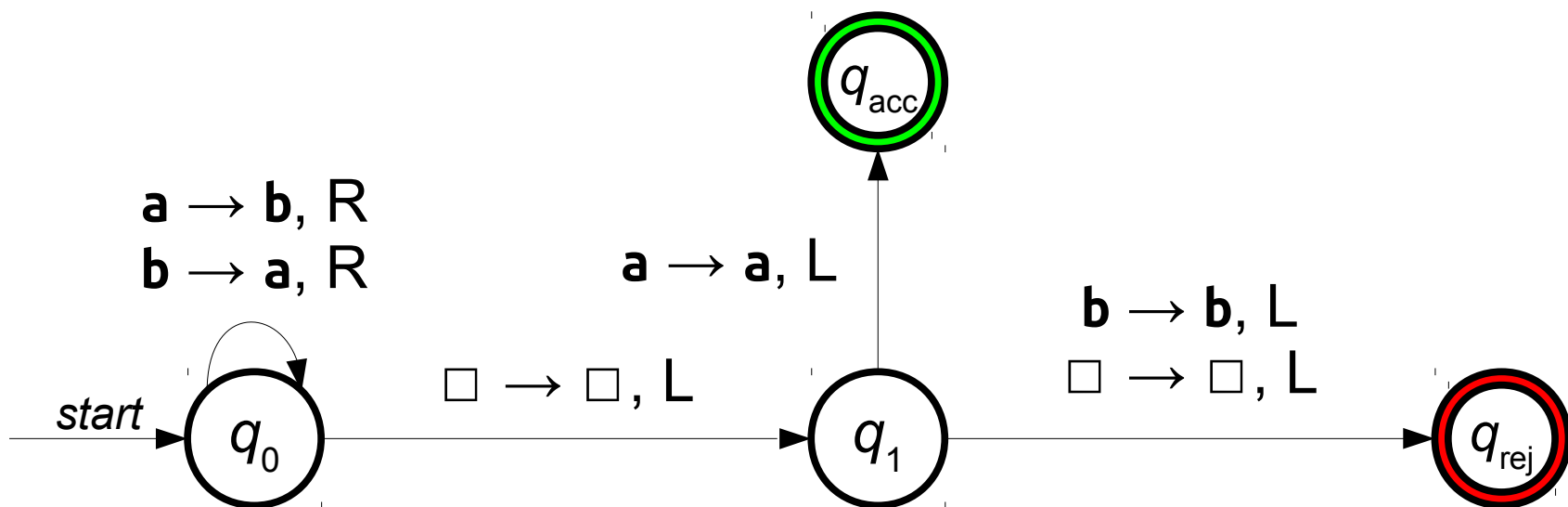






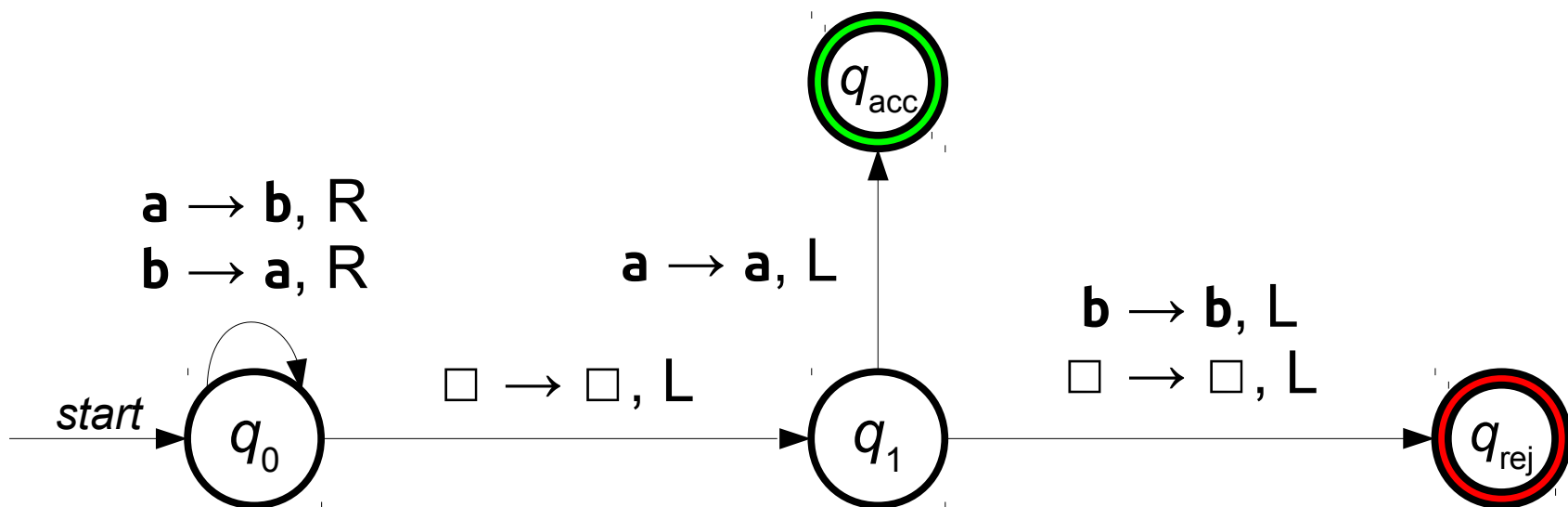






If M is a Turing machine with input alphabet Σ , then the **language of M** , denoted $\mathcal{L}(M)$, is the set

$$\mathcal{L}(M) = \{ w \in \Sigma^* \mid M \text{ accepts } w \}$$



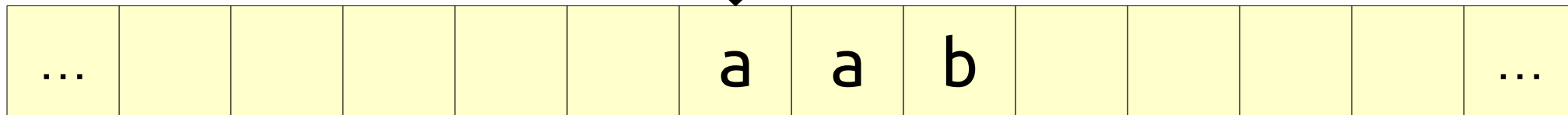
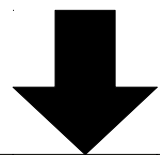
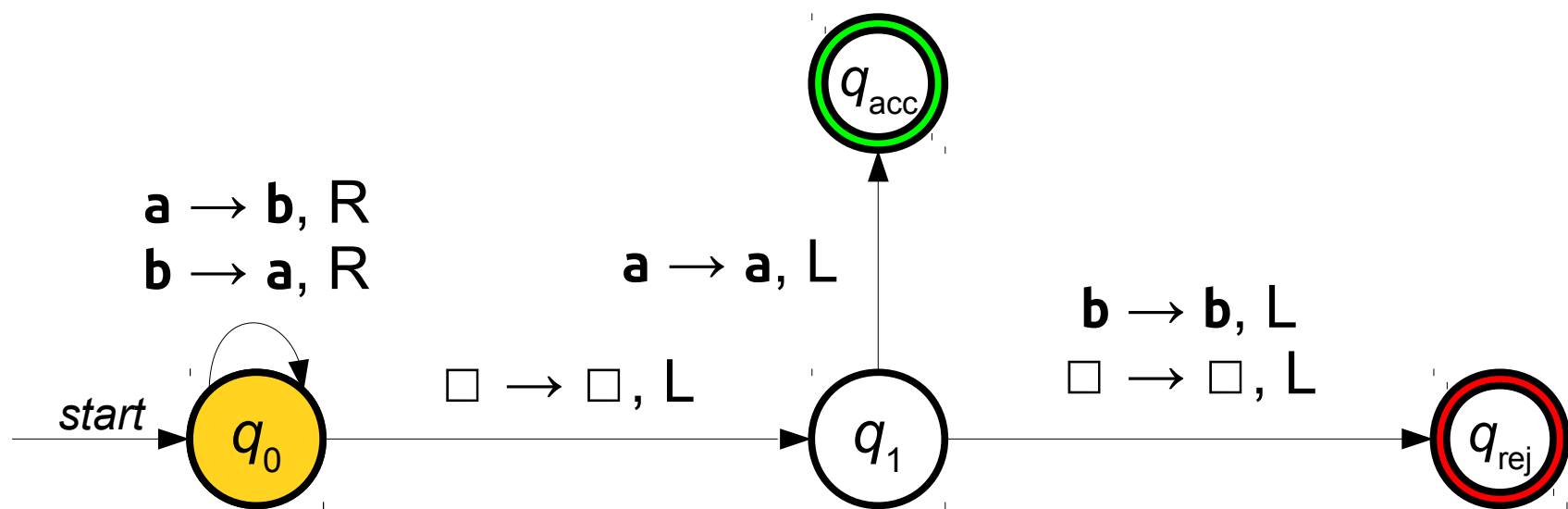
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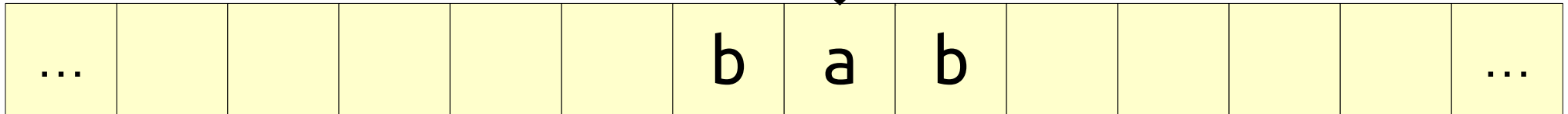
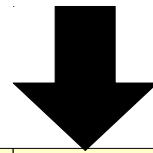
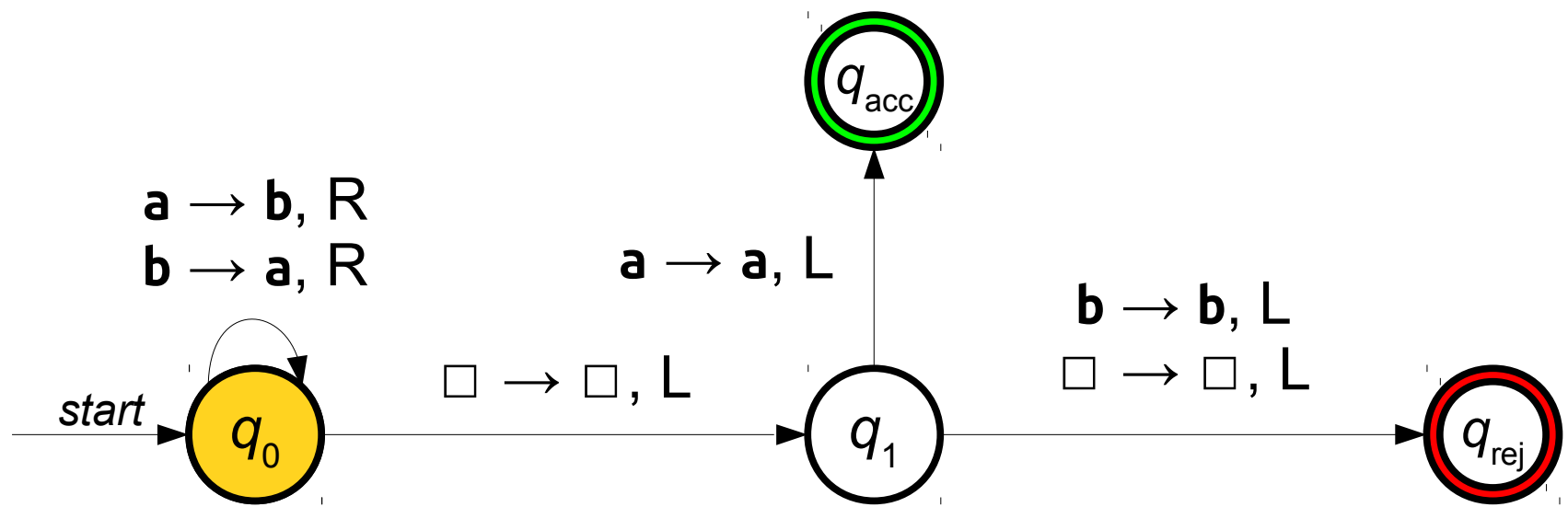
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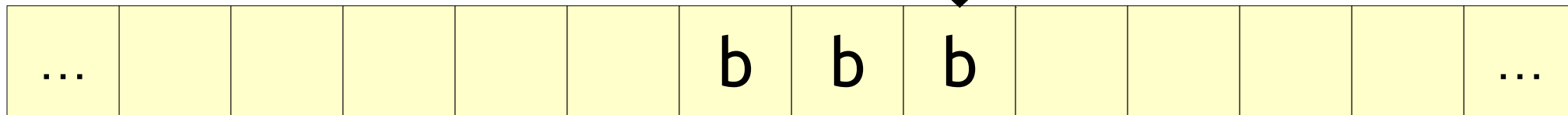
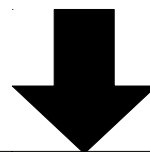
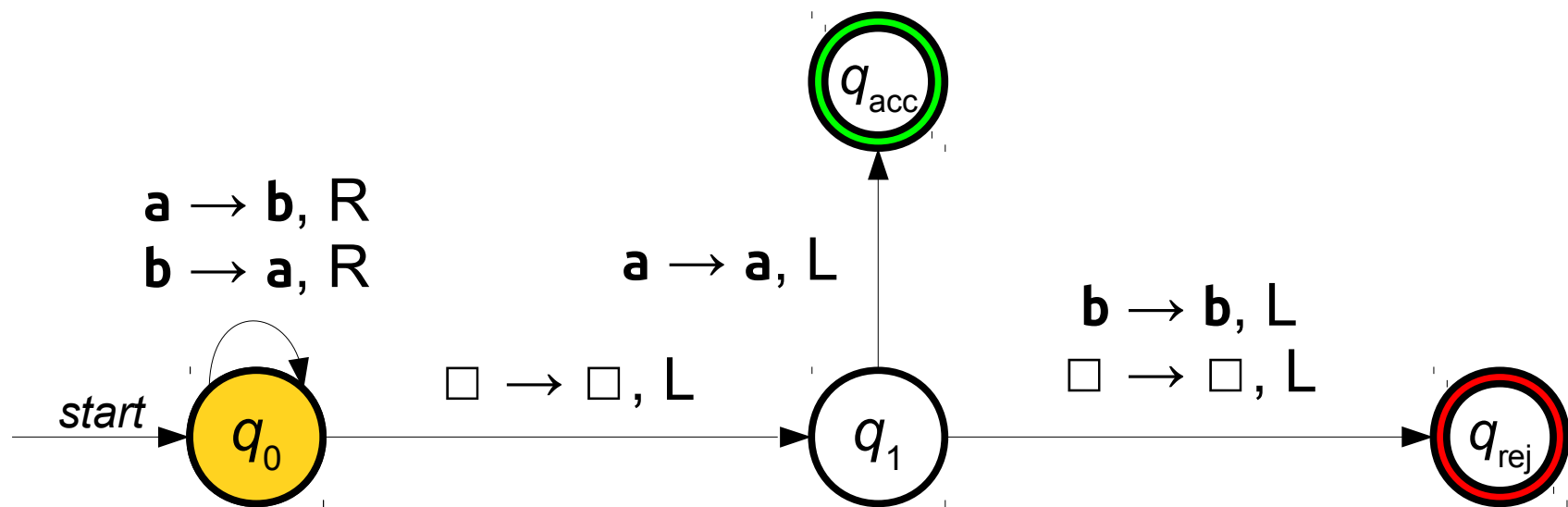
Let M be the above TM, and assume its input alphabet is $\{a, b\}$. What is $\mathcal{L}(M)$?

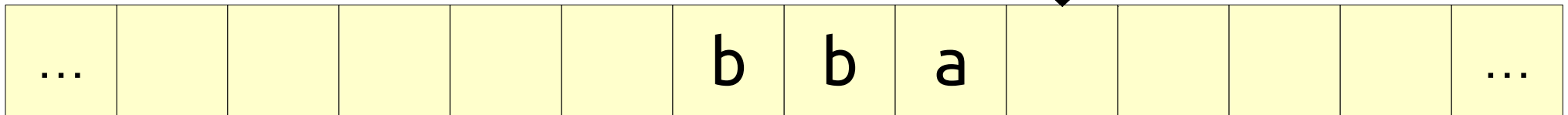
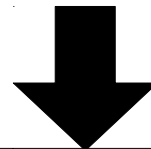
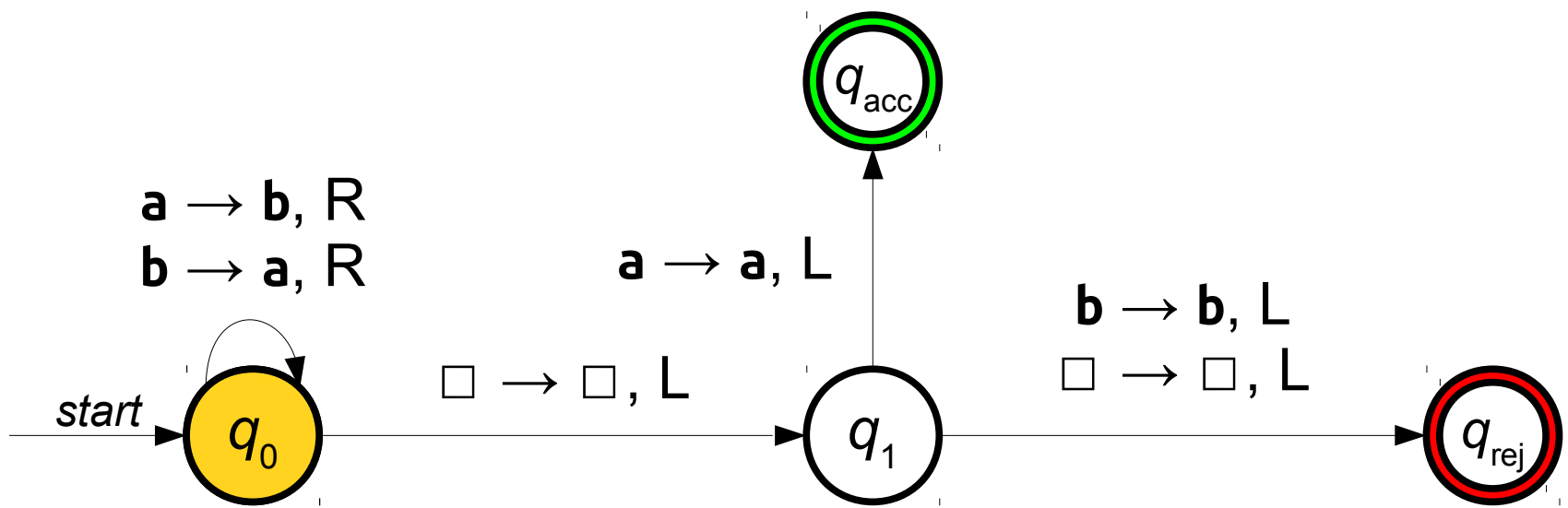
- A. $\{ w \in \{a, b\}^* \mid w \text{ ends in } a \}$
- B. $\{ w \in \{a, b\}^* \mid w \text{ ends in } b \}$
- C. \emptyset
- D. None of these, or two or more of these.

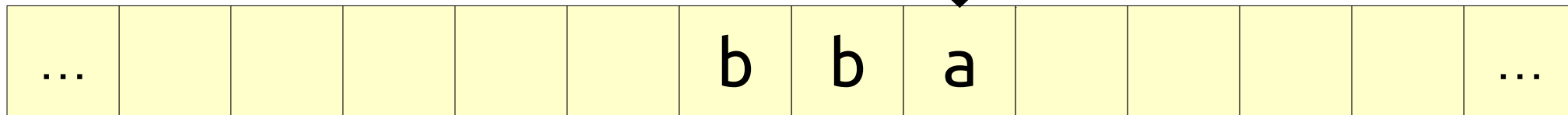
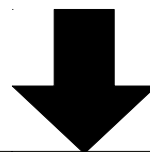
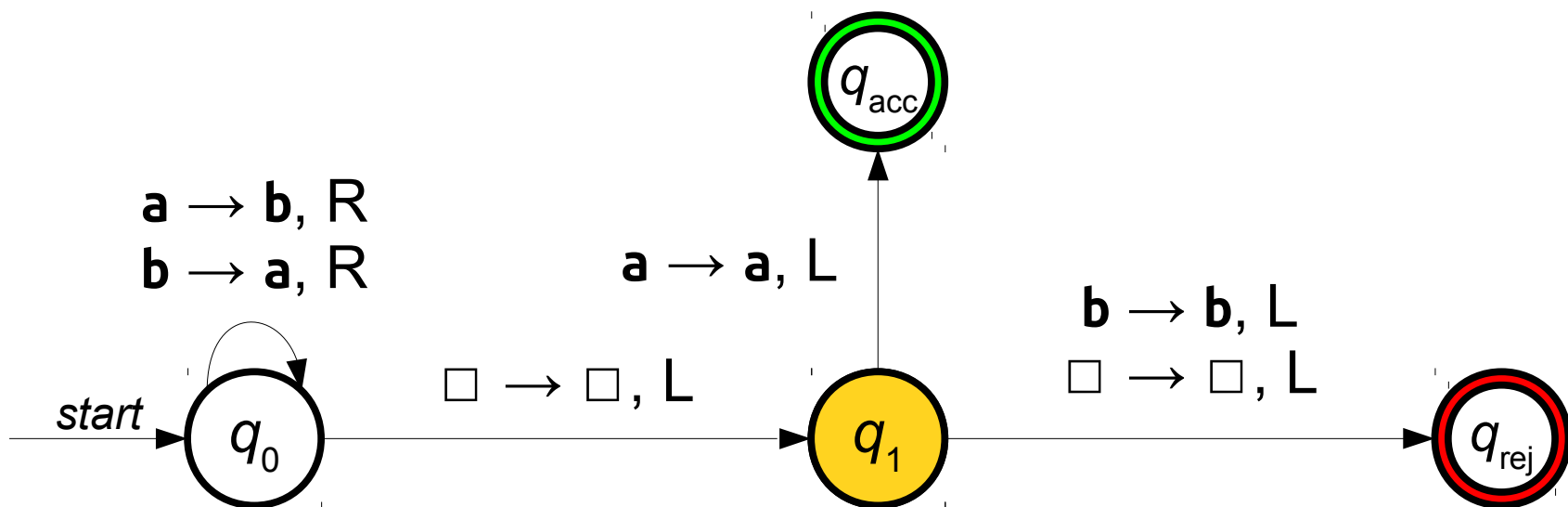
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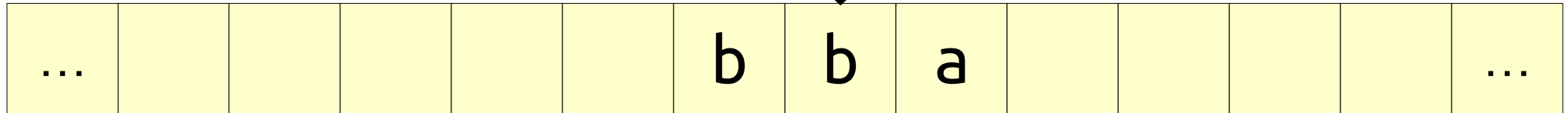
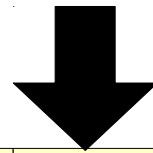
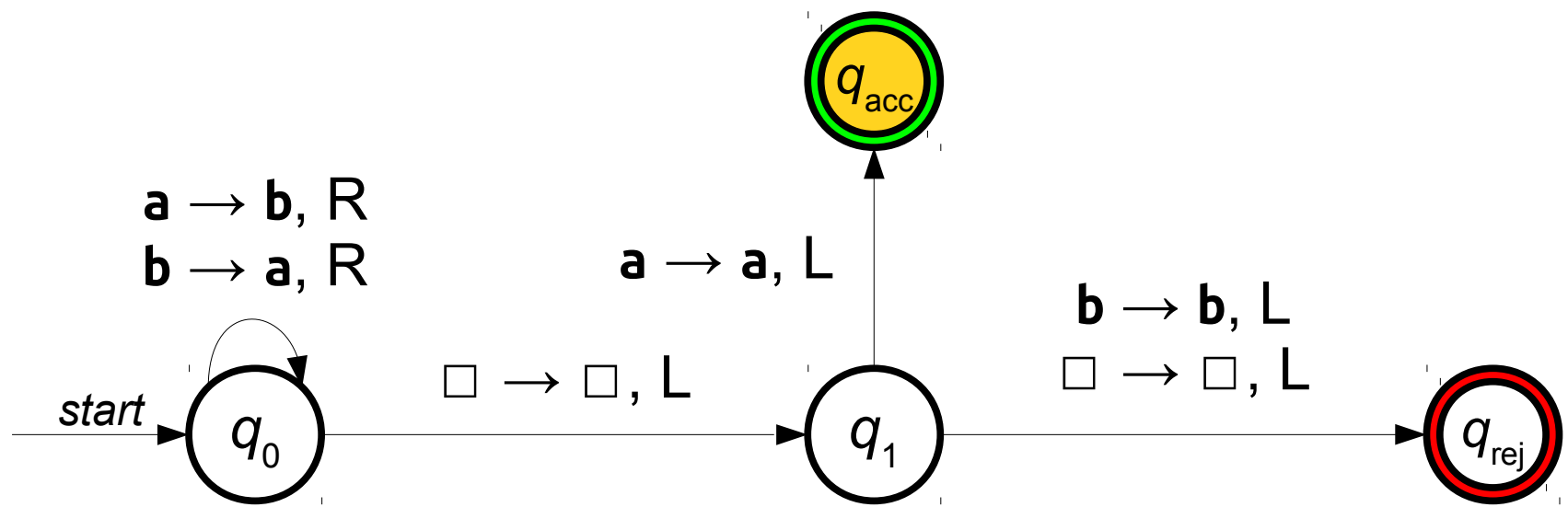


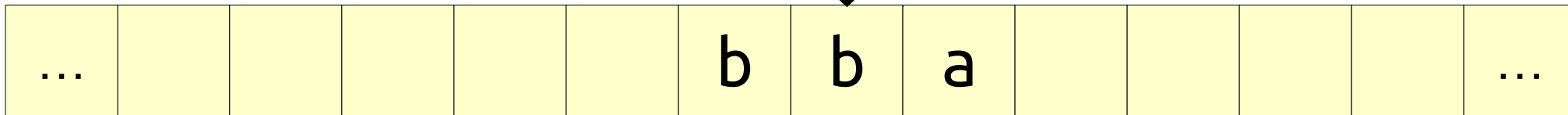
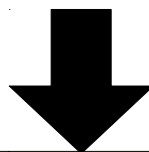
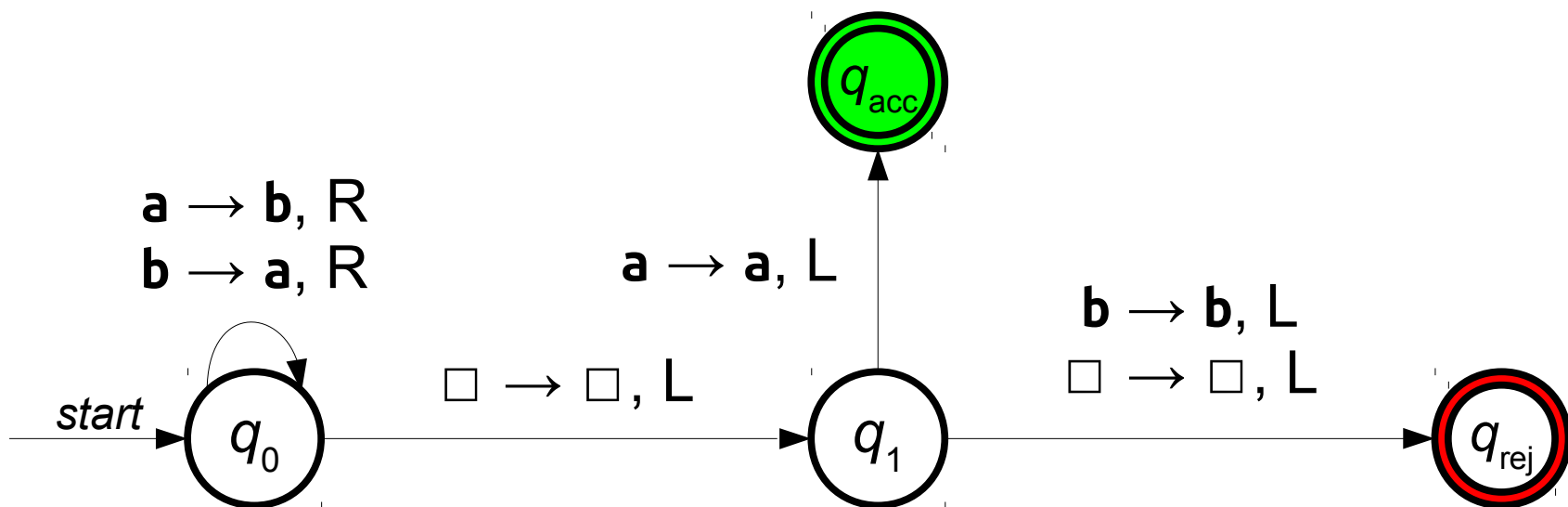


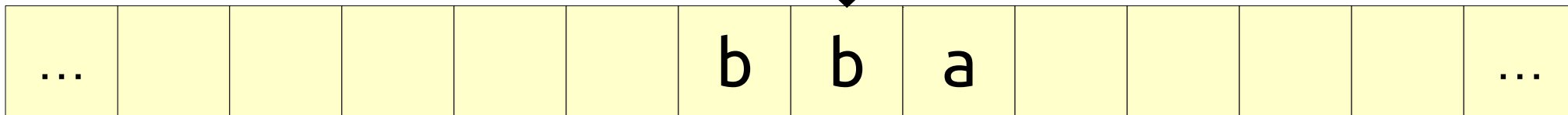
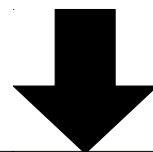
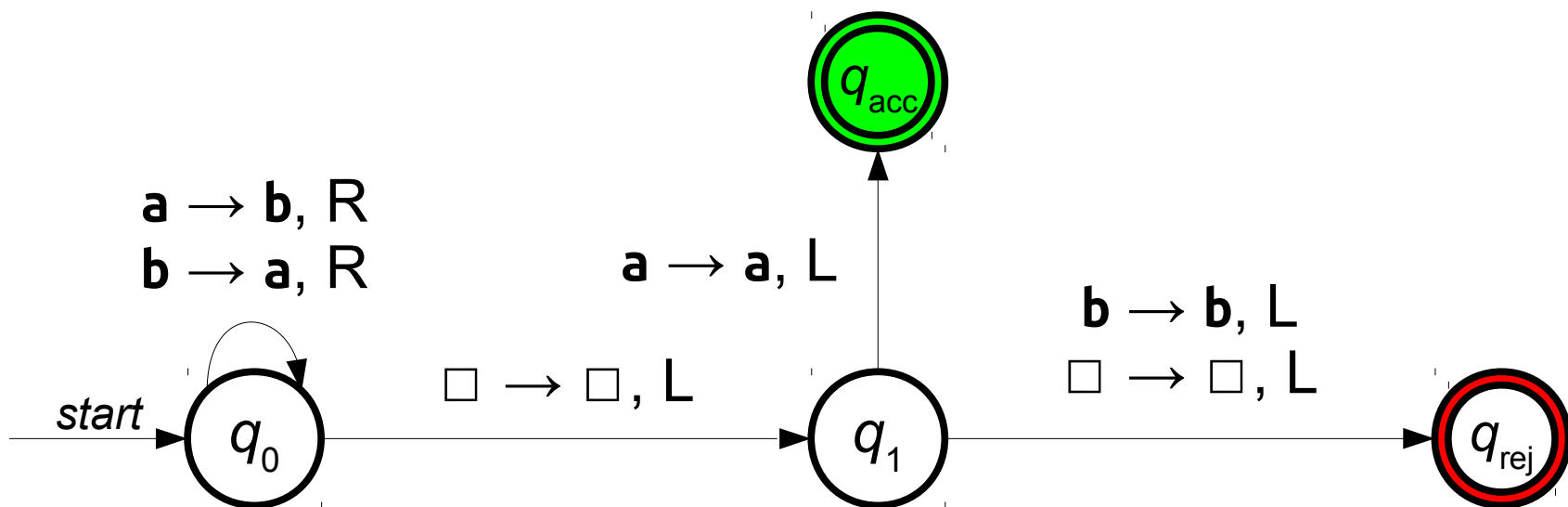












Although the tape ends with **bba** written on it, the original input string was **aab**. This shows that the TM accepts **aab**, not **bba**.

$$\text{So } \mathcal{L}(M) = \{ w \in \{ \mathbf{a}, \mathbf{b} \}^* \mid w \text{ ends in } \mathbf{b} \}$$

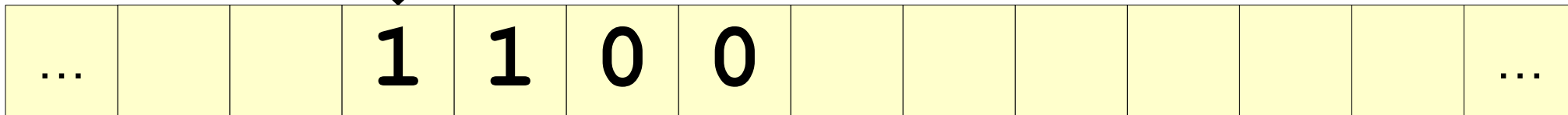
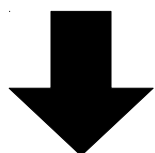
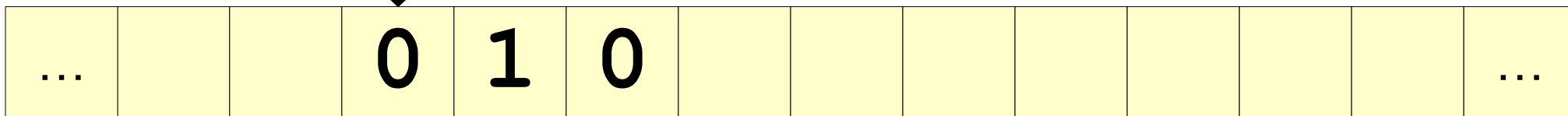
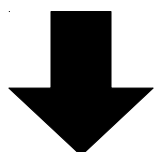
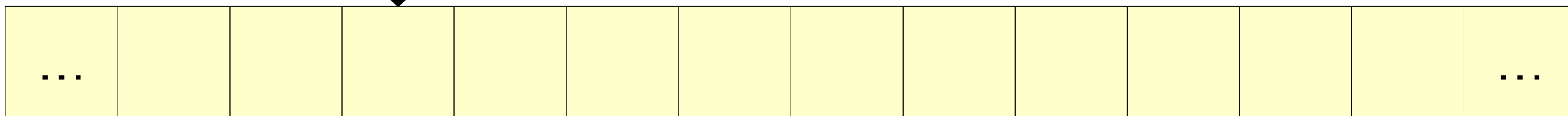
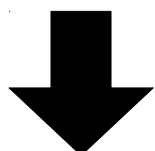
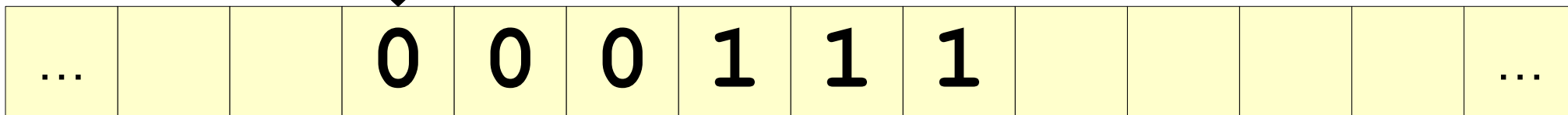
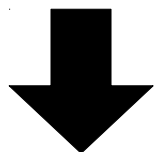
Designing Turing Machines

- Despite their simplicity, Turing machines are very powerful computing devices.
- Today's lecture explores how to design Turing machines for various languages.

Designing Turing Machines

- Let $\Sigma = \{0, 1\}$ and consider the language $L = \{0^n 1^n \mid n \in \mathbb{N}\}$.
- We know that L is context-free.
- How might we build a Turing machine for it?

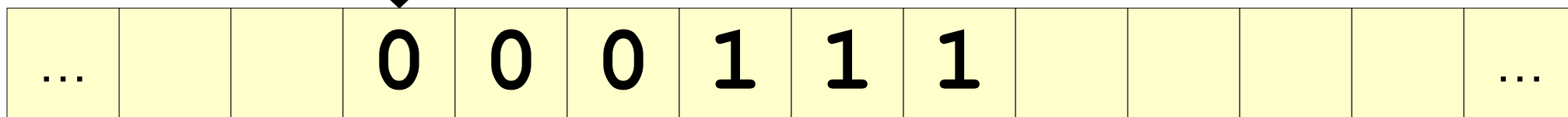
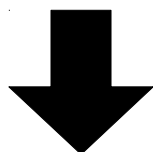
$$L = \{ \mathbf{0}^n \mathbf{1}^n \mid n \in \mathbb{N} \}$$



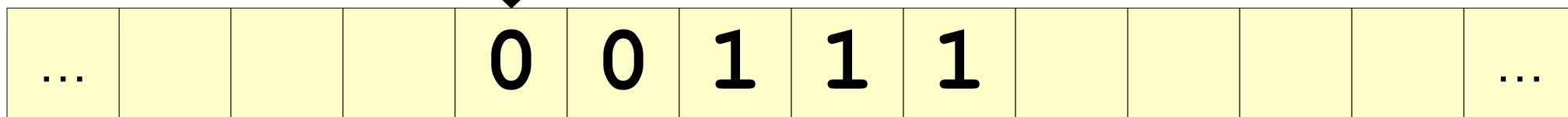
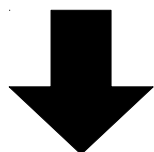
A Recursive Approach

- The string ε is in L .
- The string $0w1$ is in L iff w is in L .
- Any string starting with 1 is not in L .
- Any string ending with 0 is not in L .

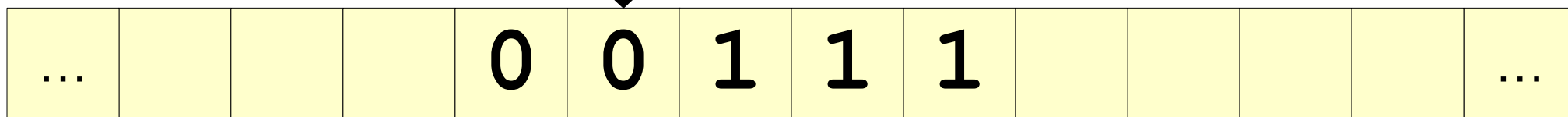
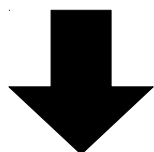
A Sketch of the TM



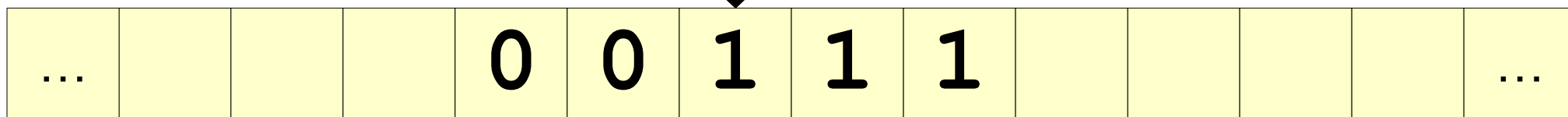
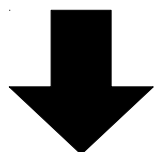
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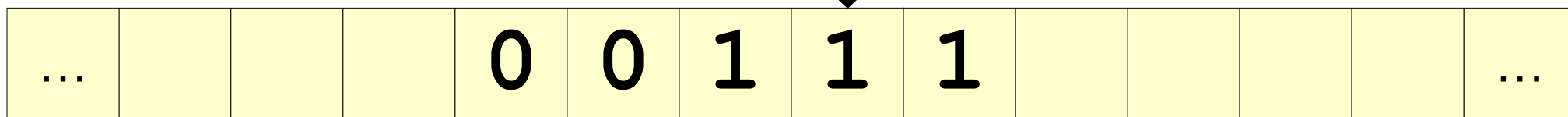
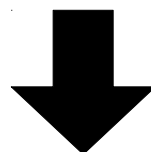
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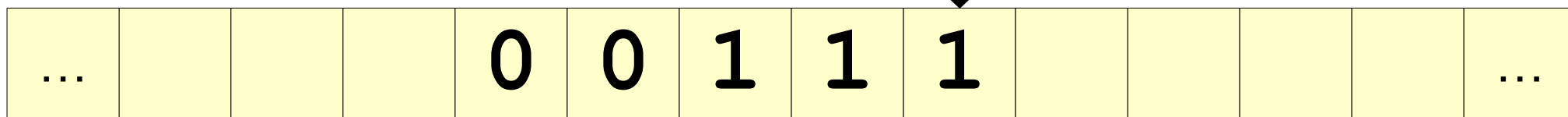
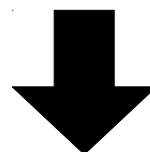
A Sketch of the TM



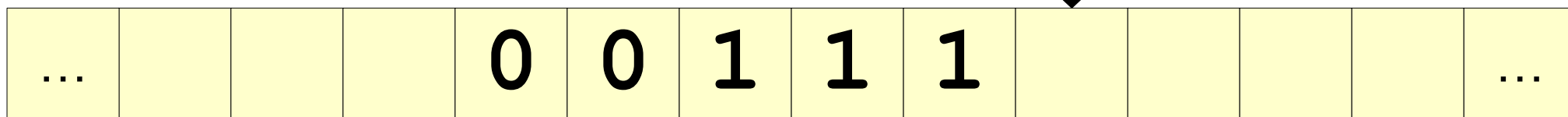
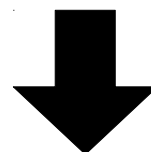
A Sketch of the TM



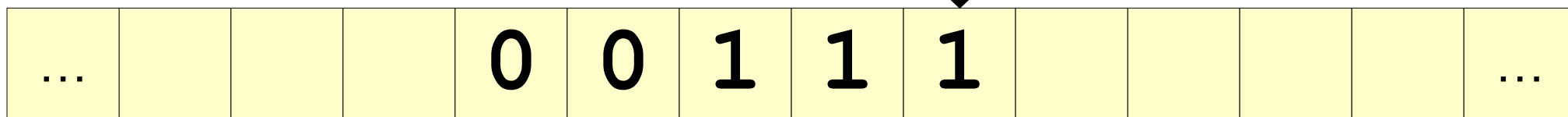
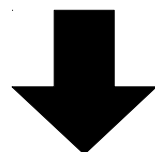
A Sketch of the TM



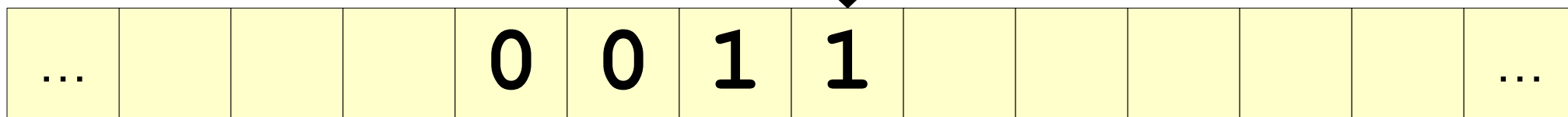
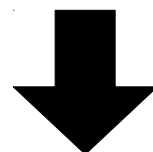
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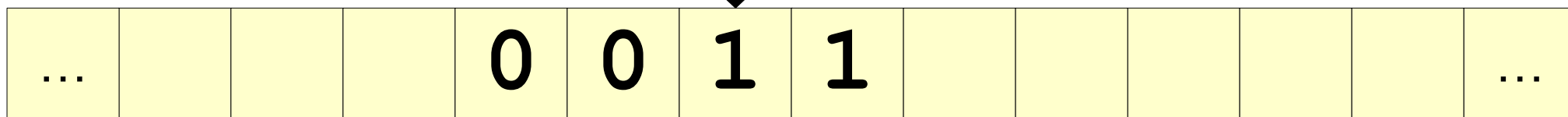
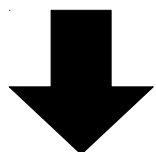
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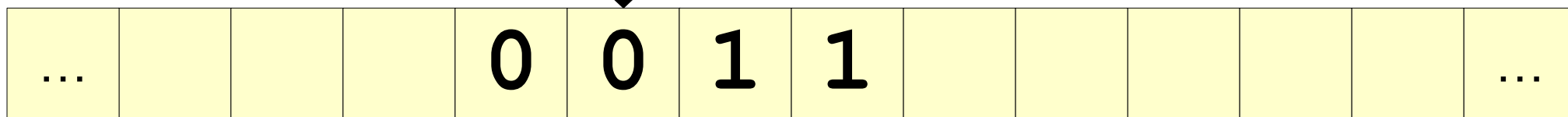
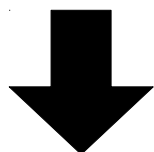
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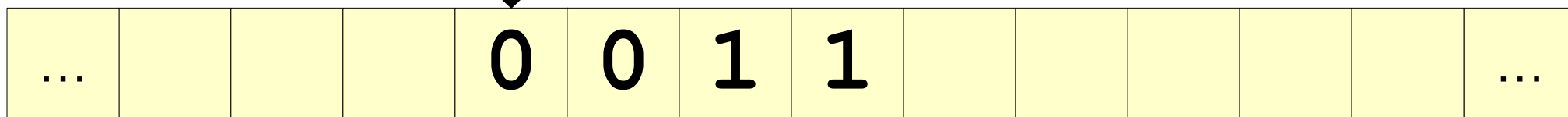
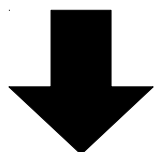
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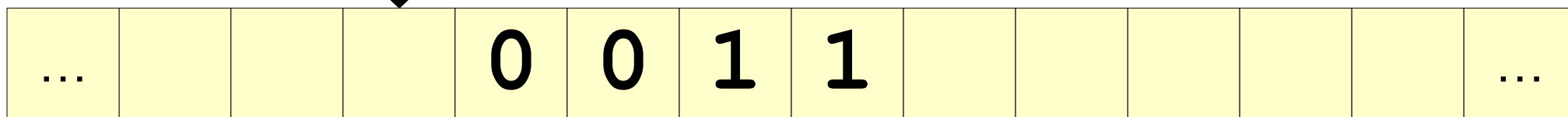
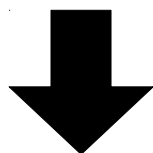
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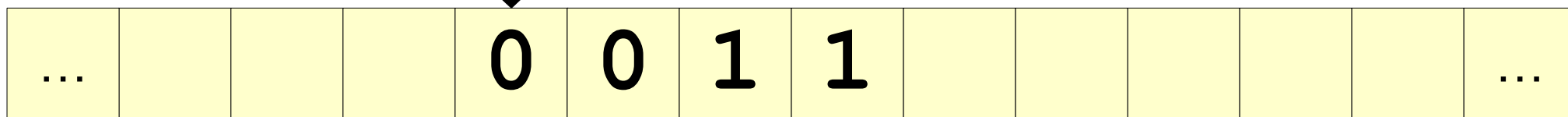
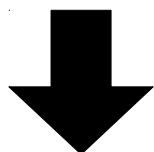
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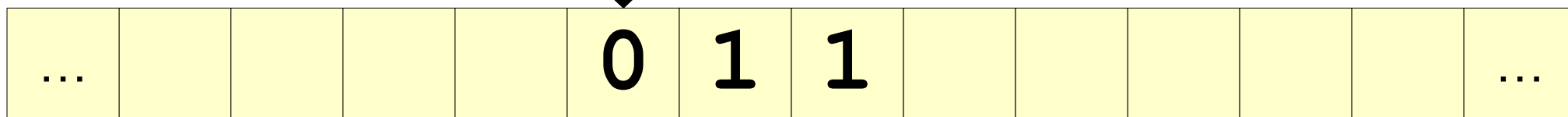
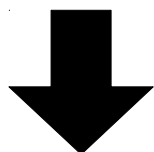
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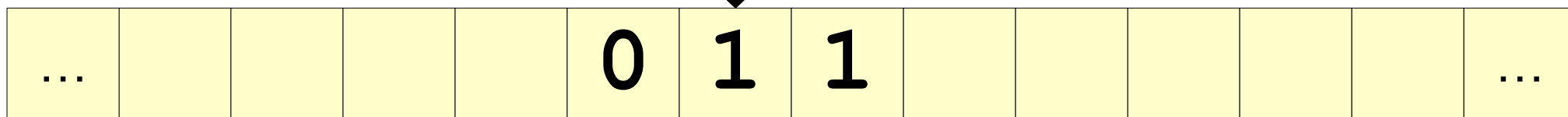
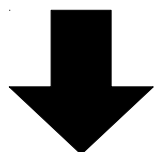
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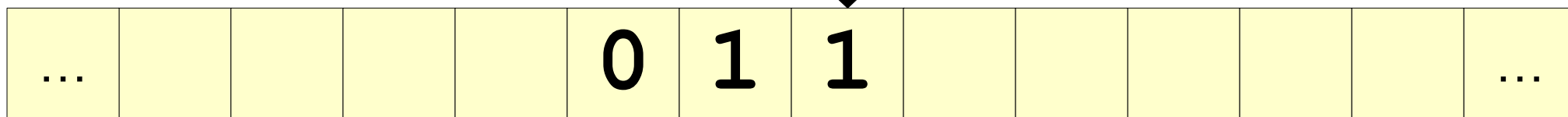
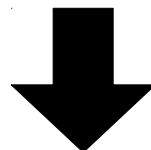
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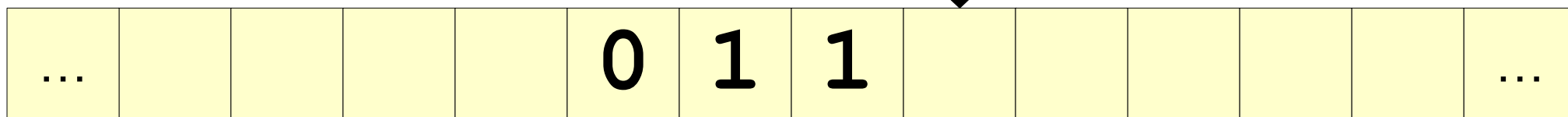
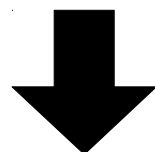
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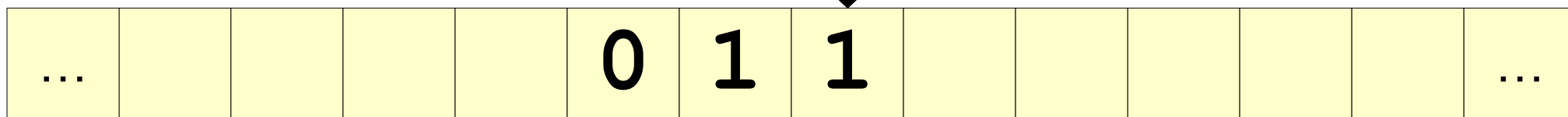
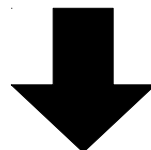
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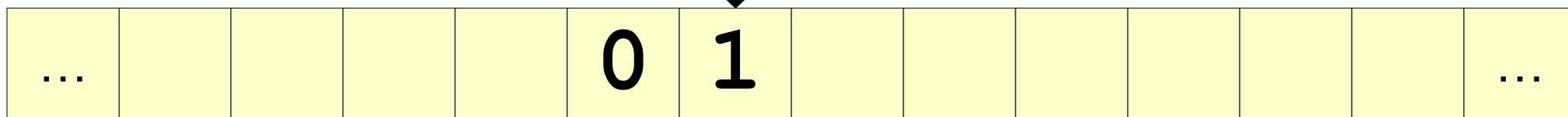
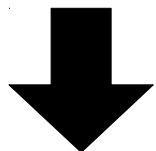
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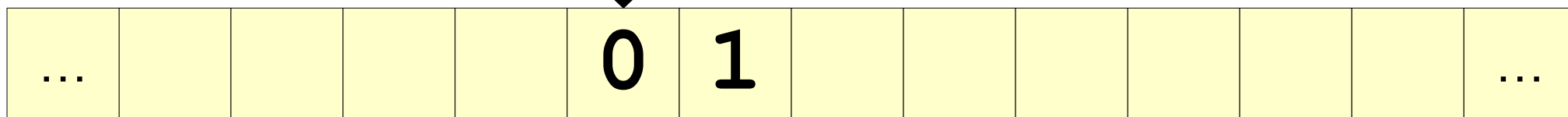
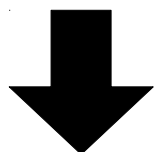
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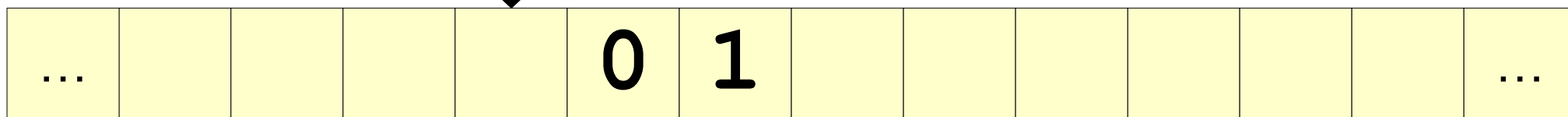
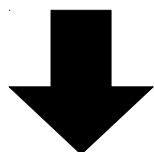
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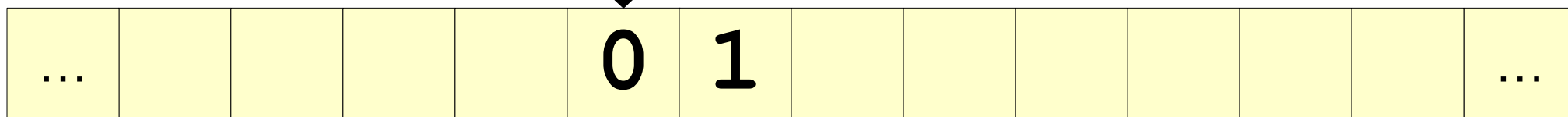
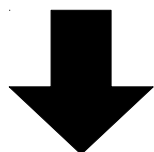
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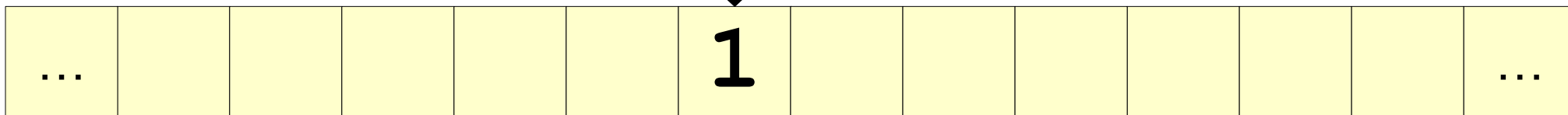
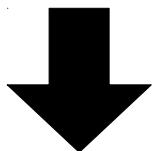
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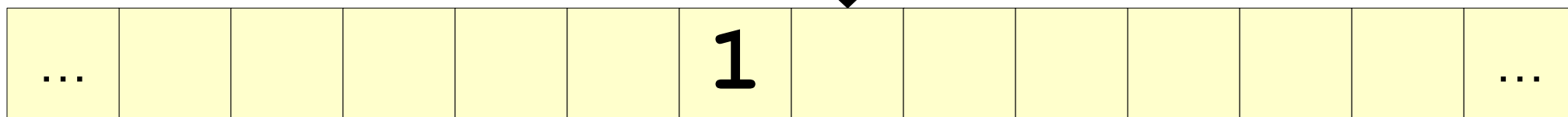
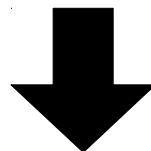
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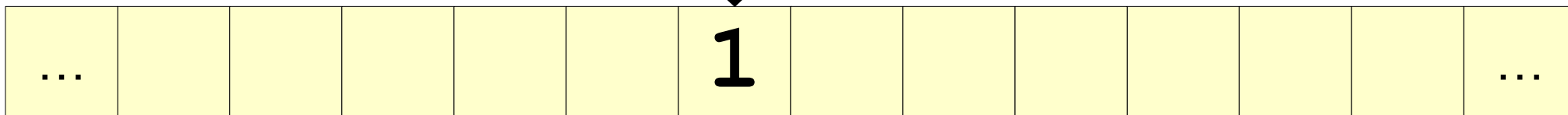
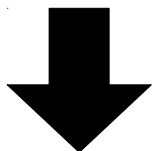
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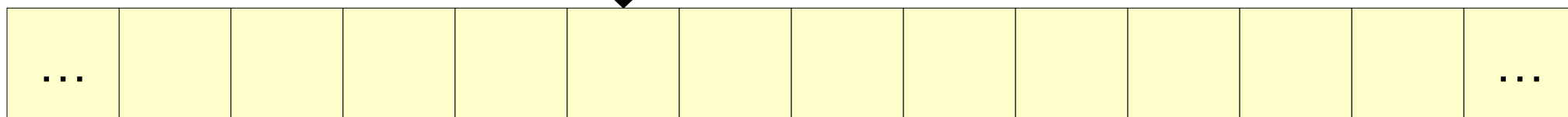
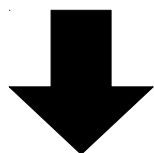
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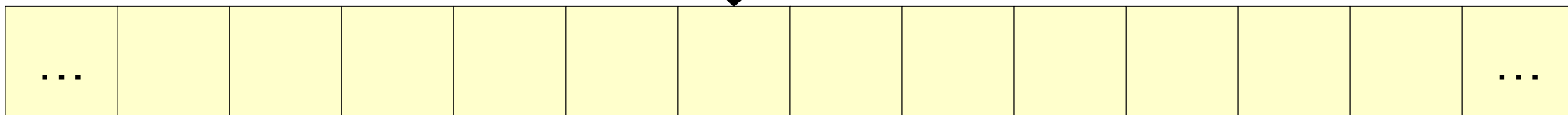
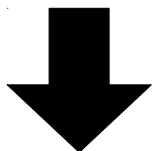
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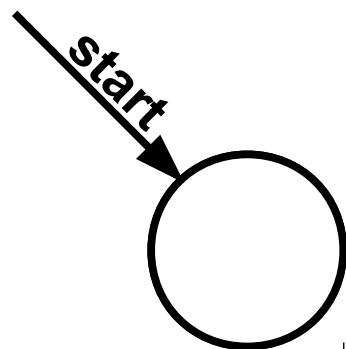


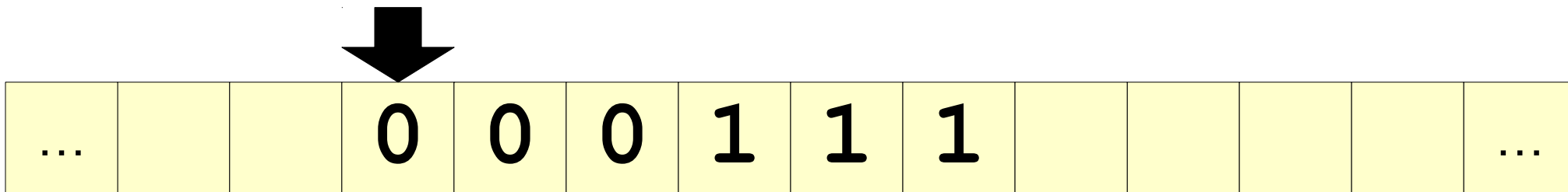
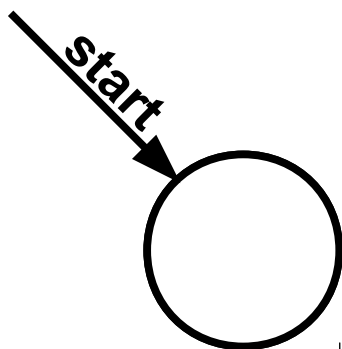
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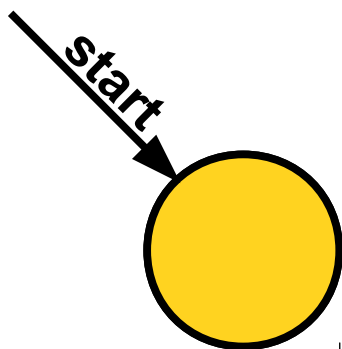


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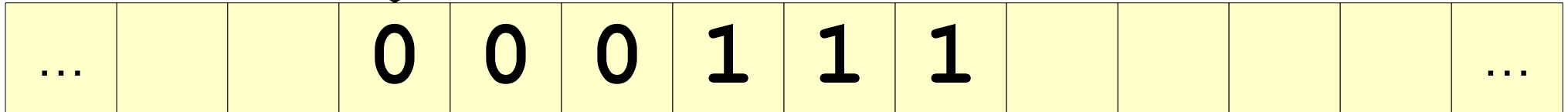
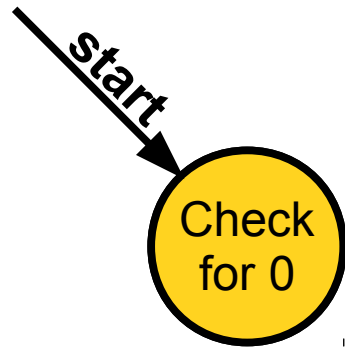


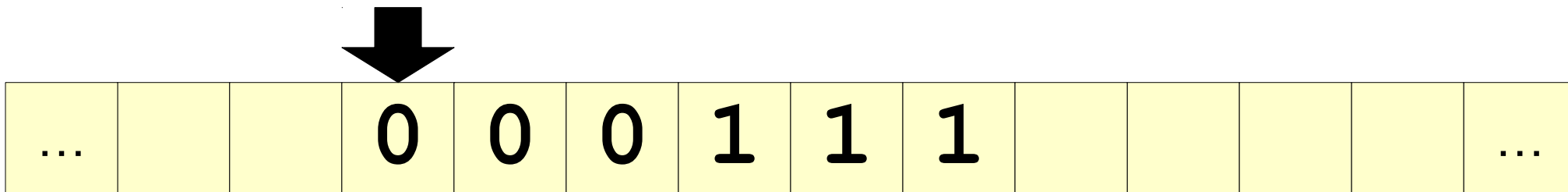
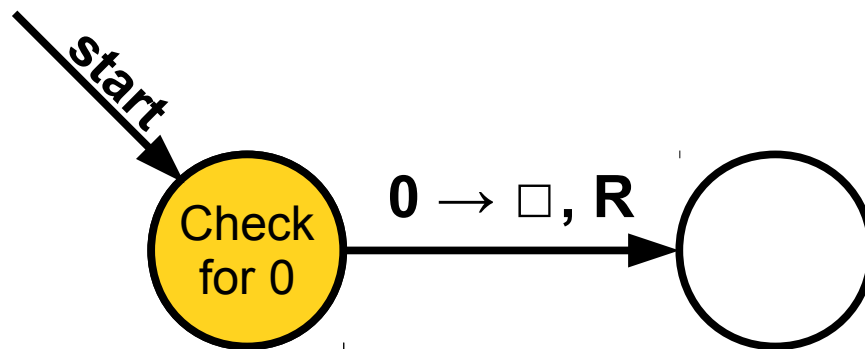


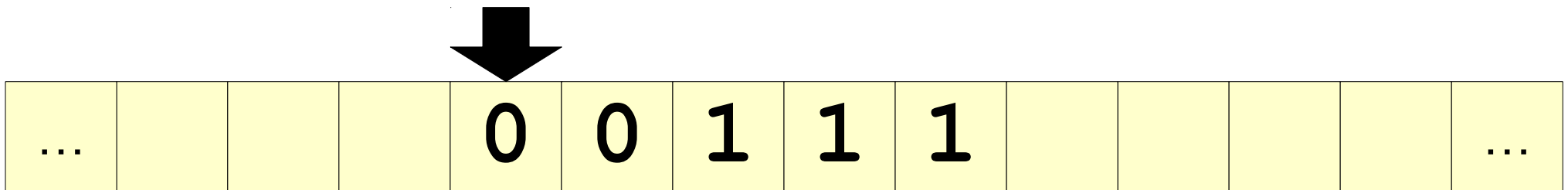
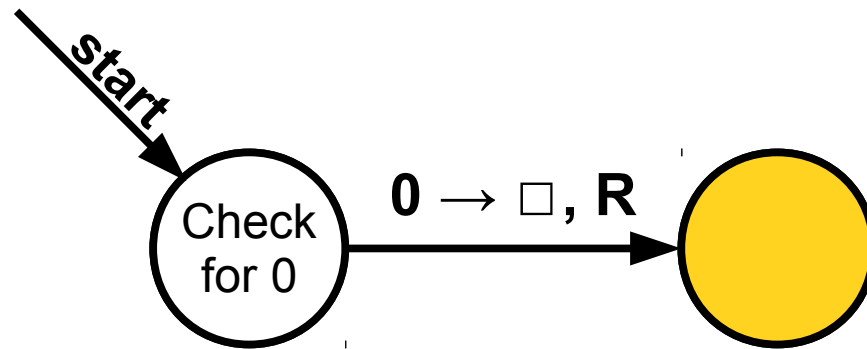


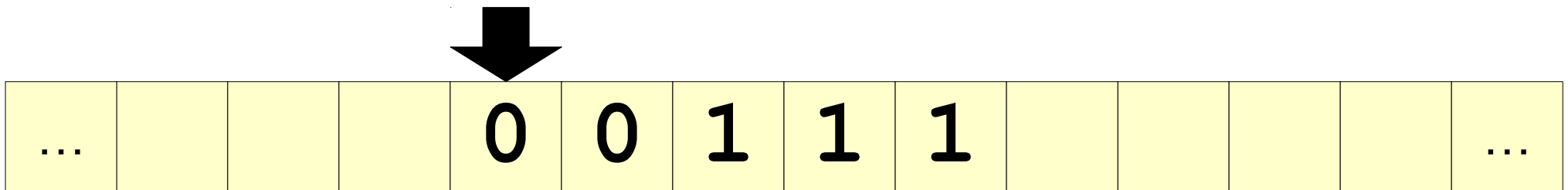
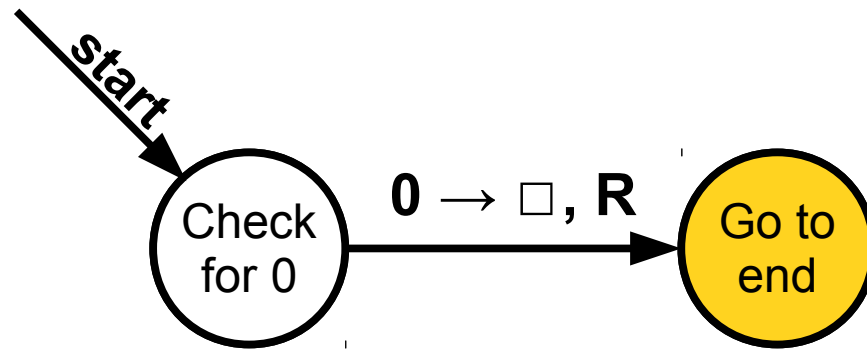


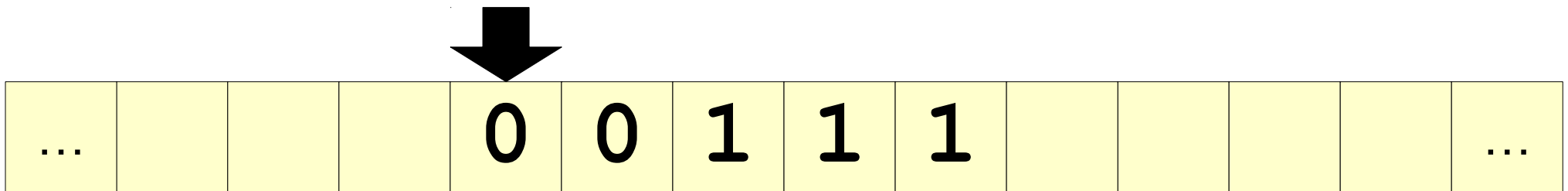
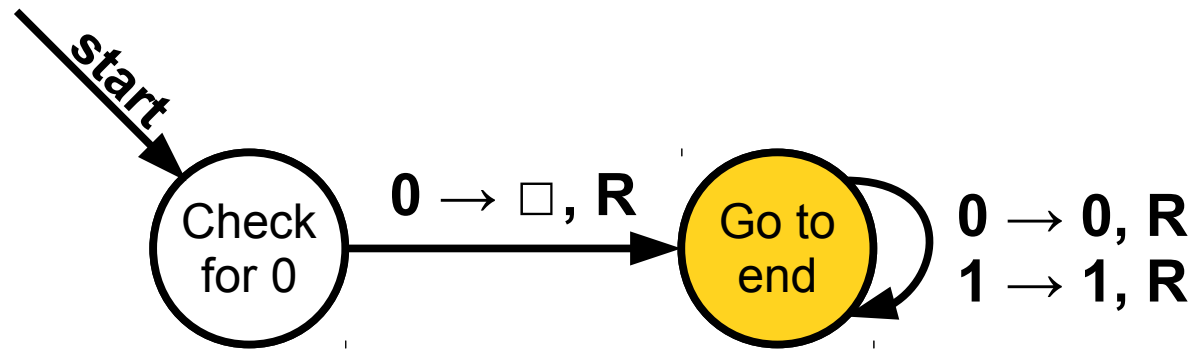
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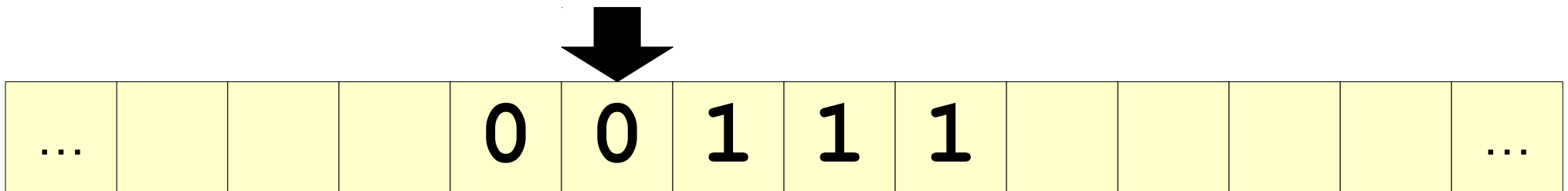
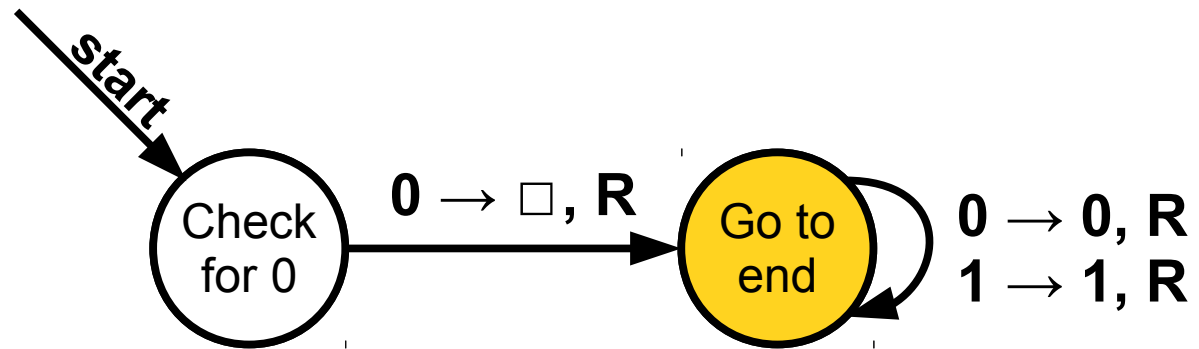


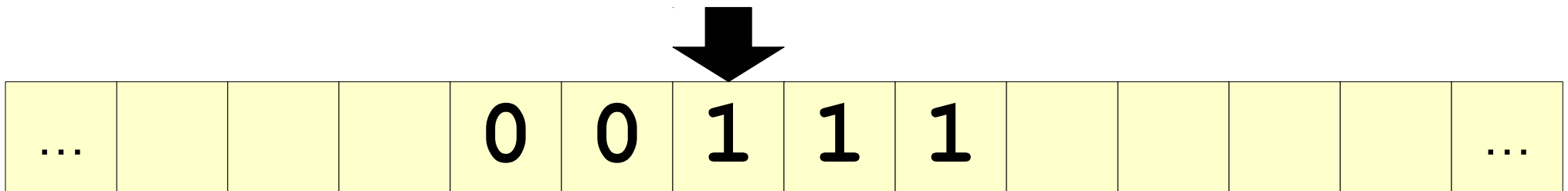
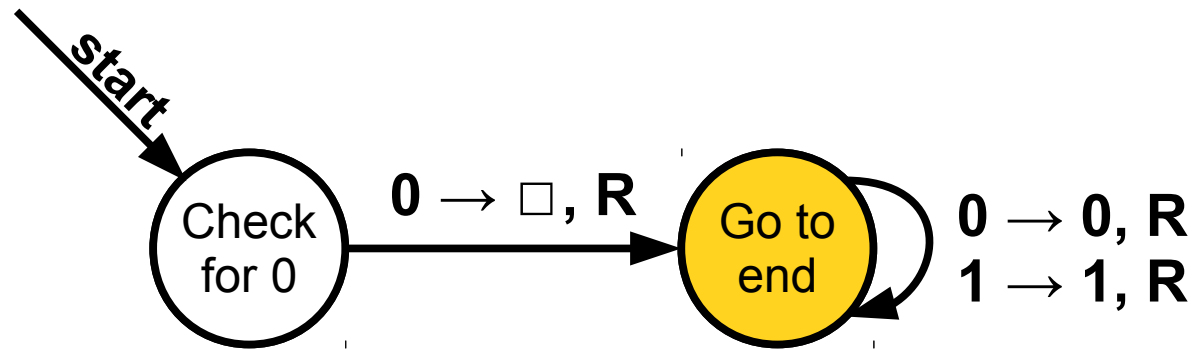


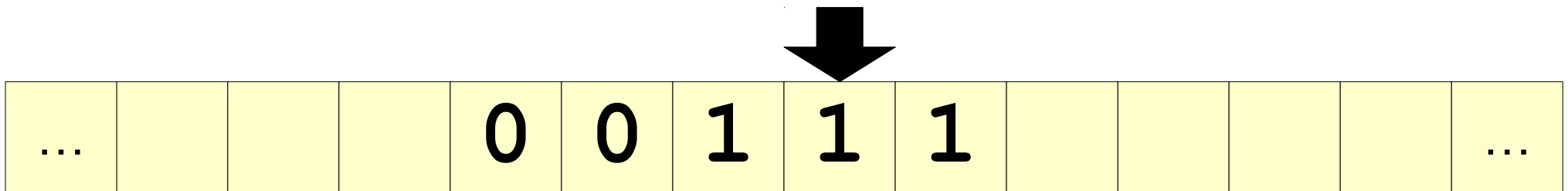
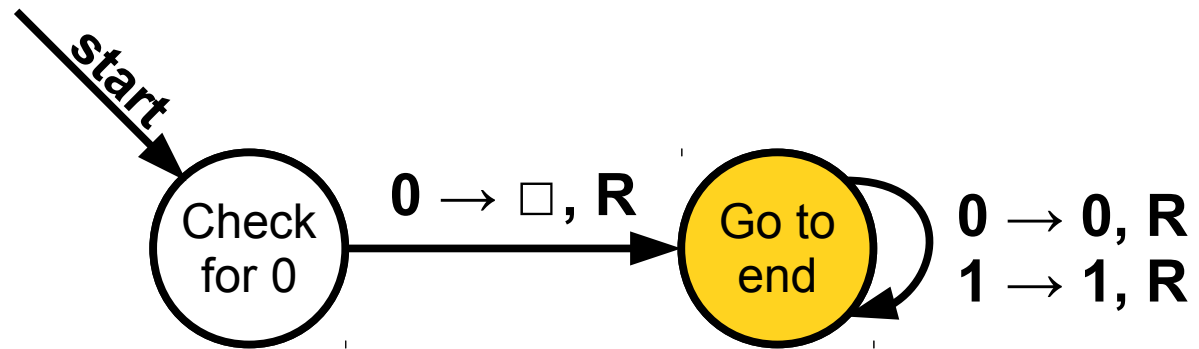


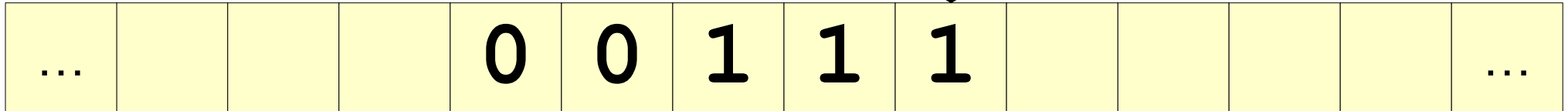
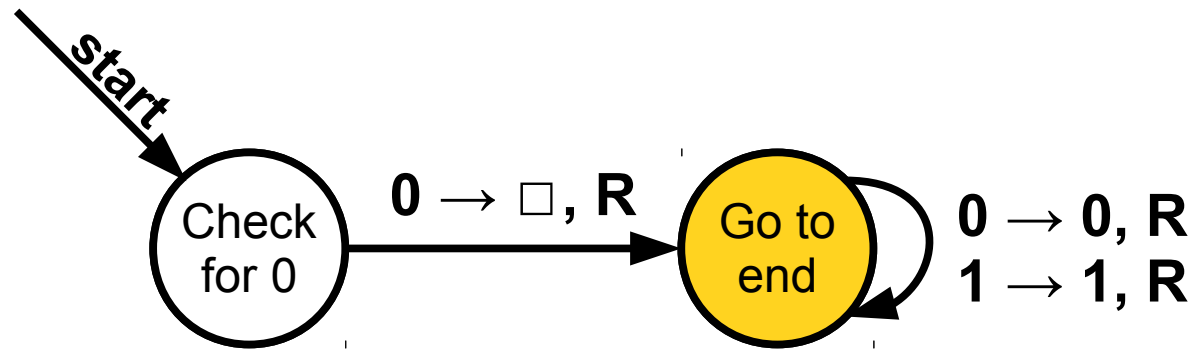


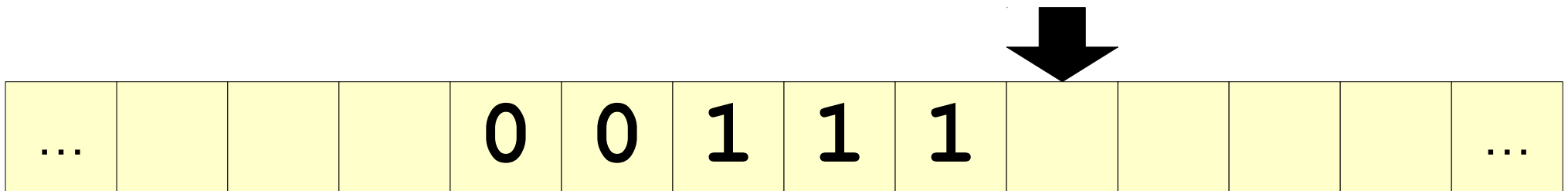
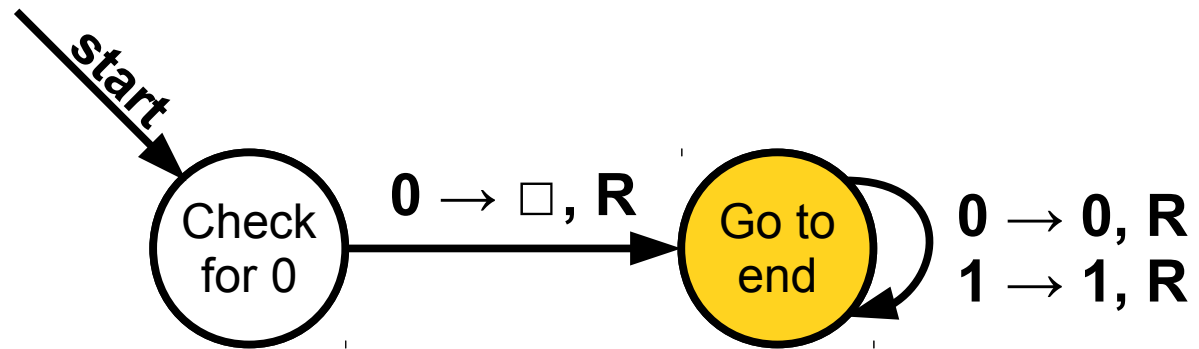


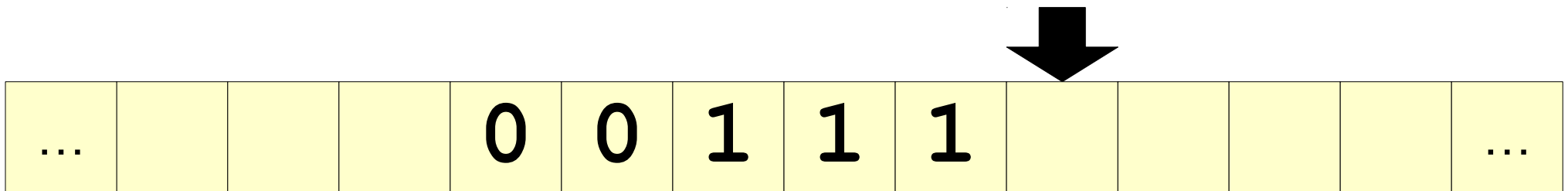
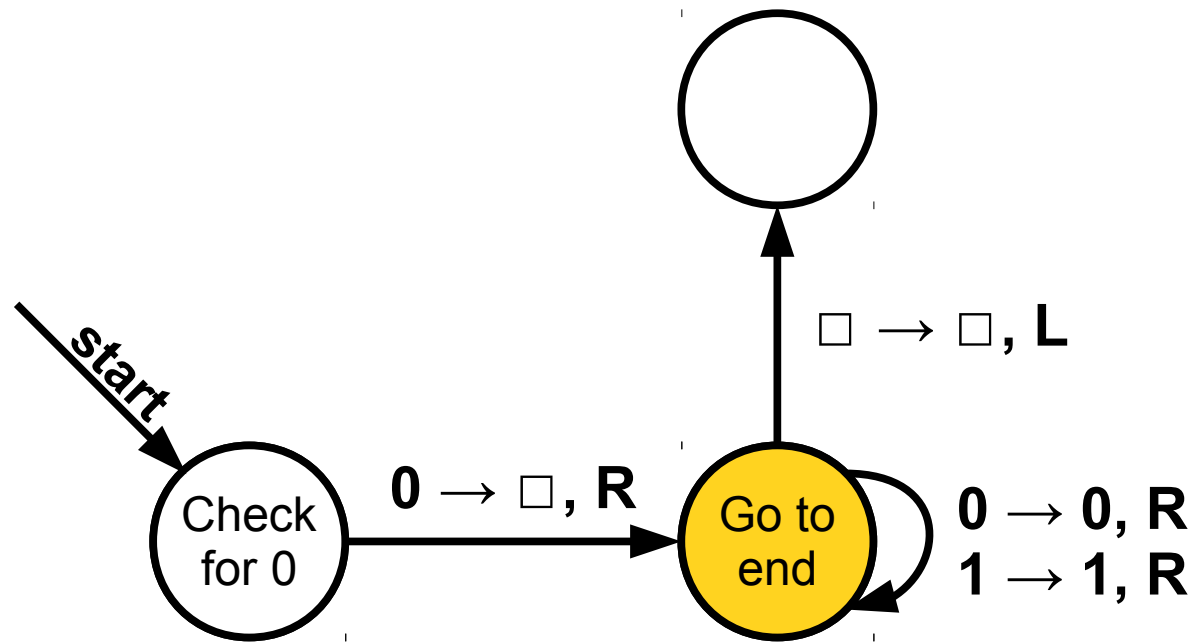


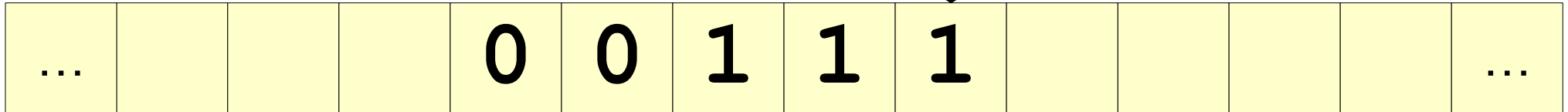
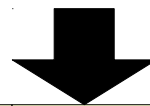
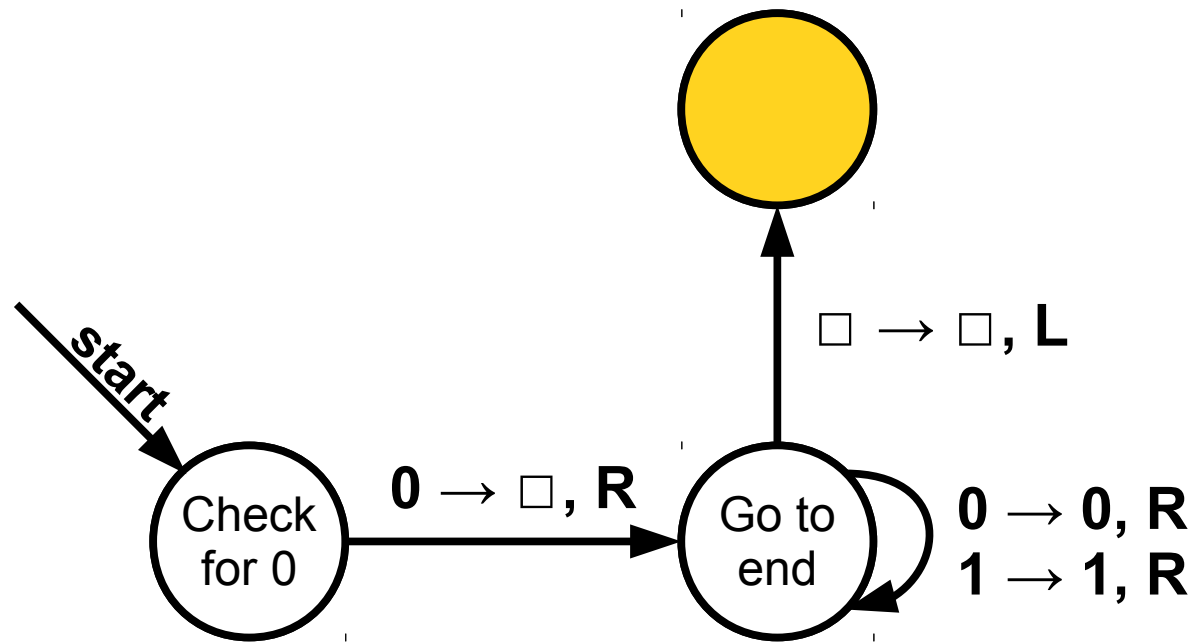


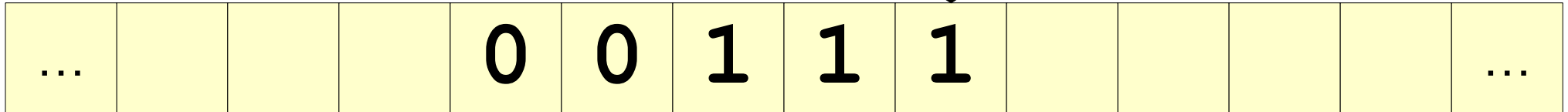
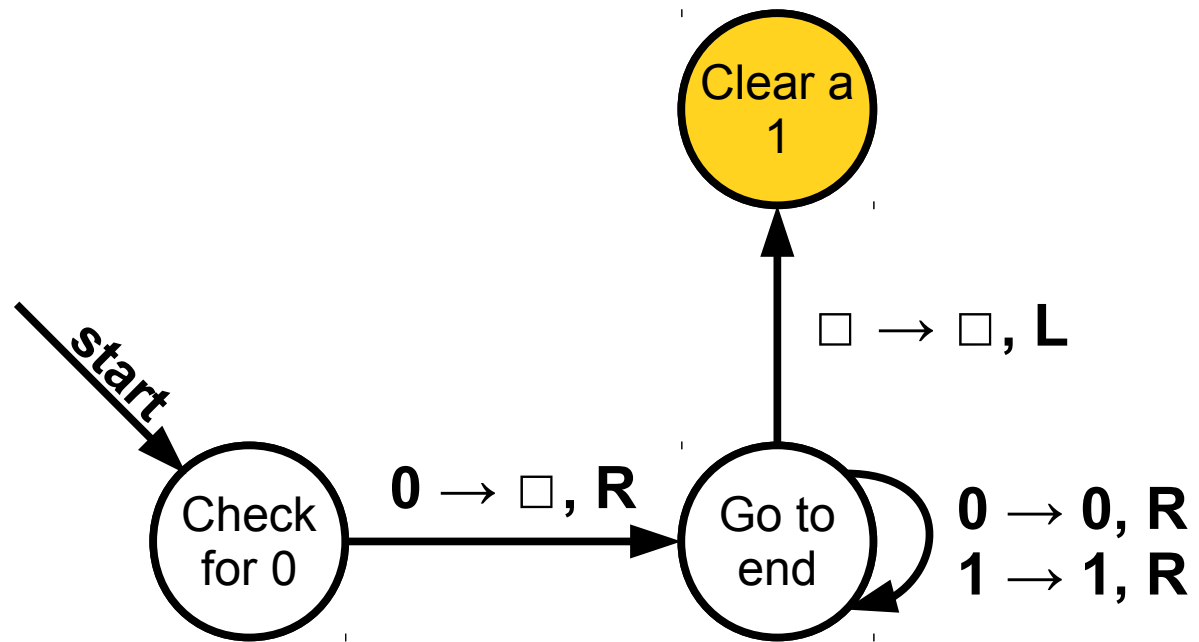


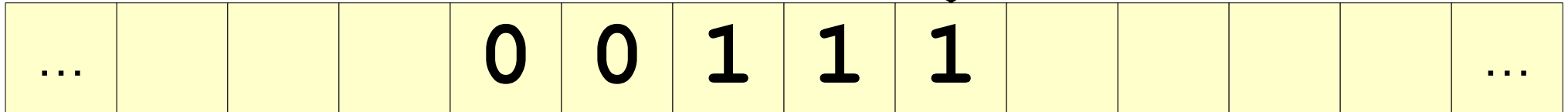
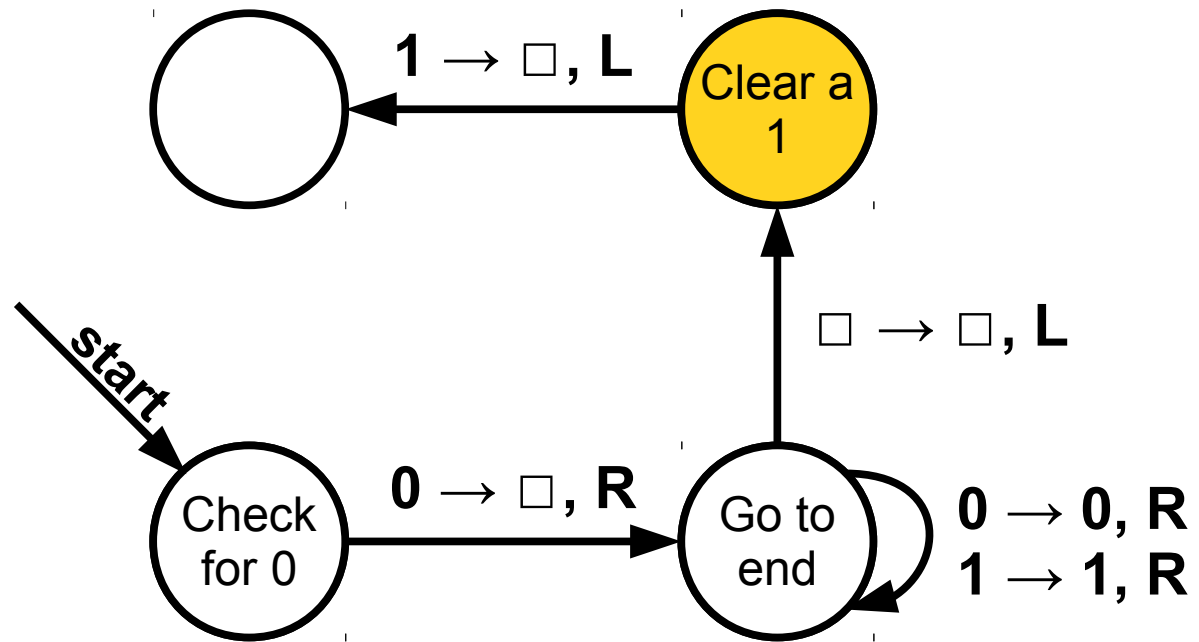


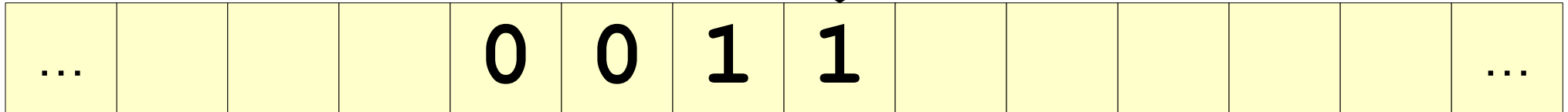
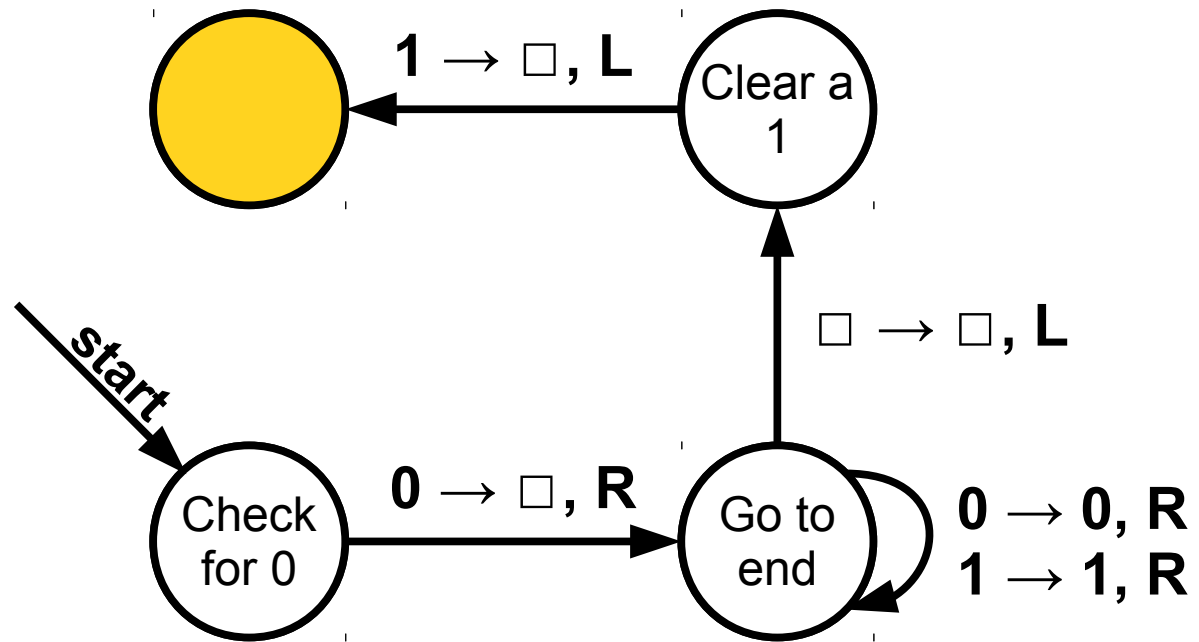


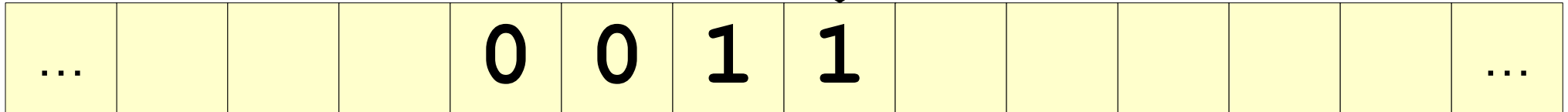
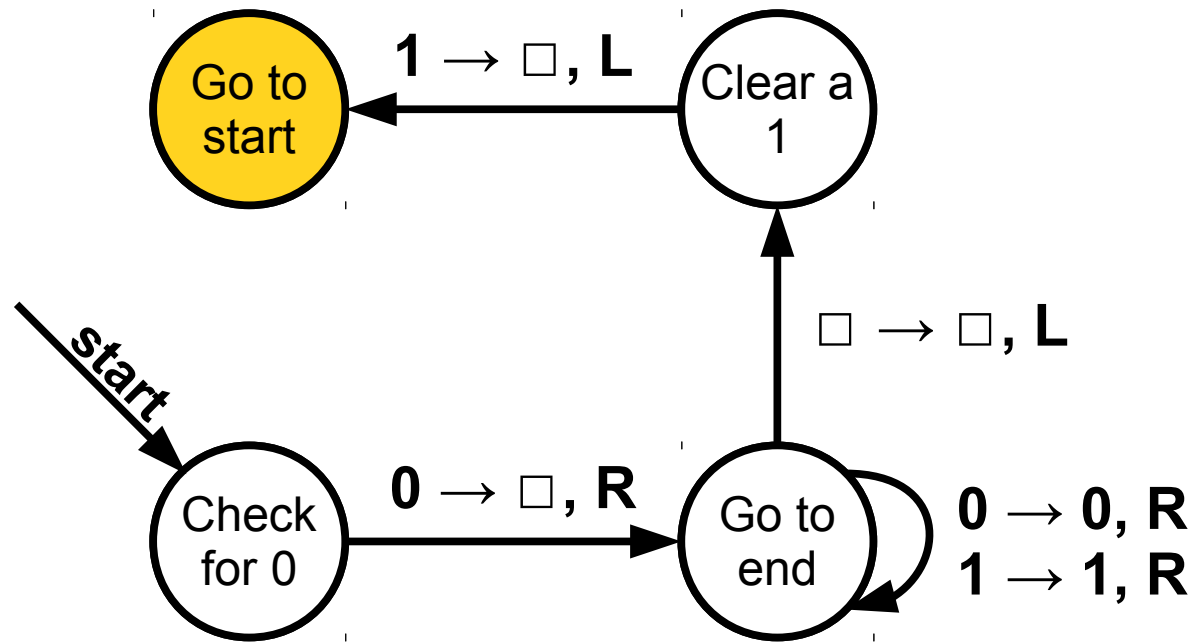


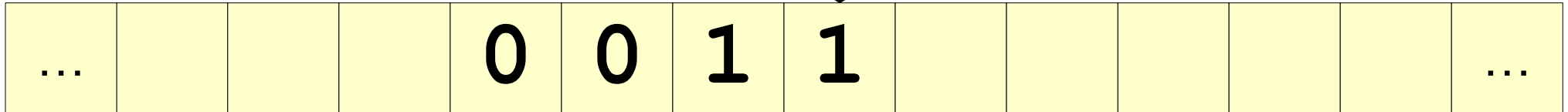
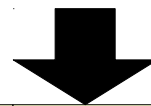
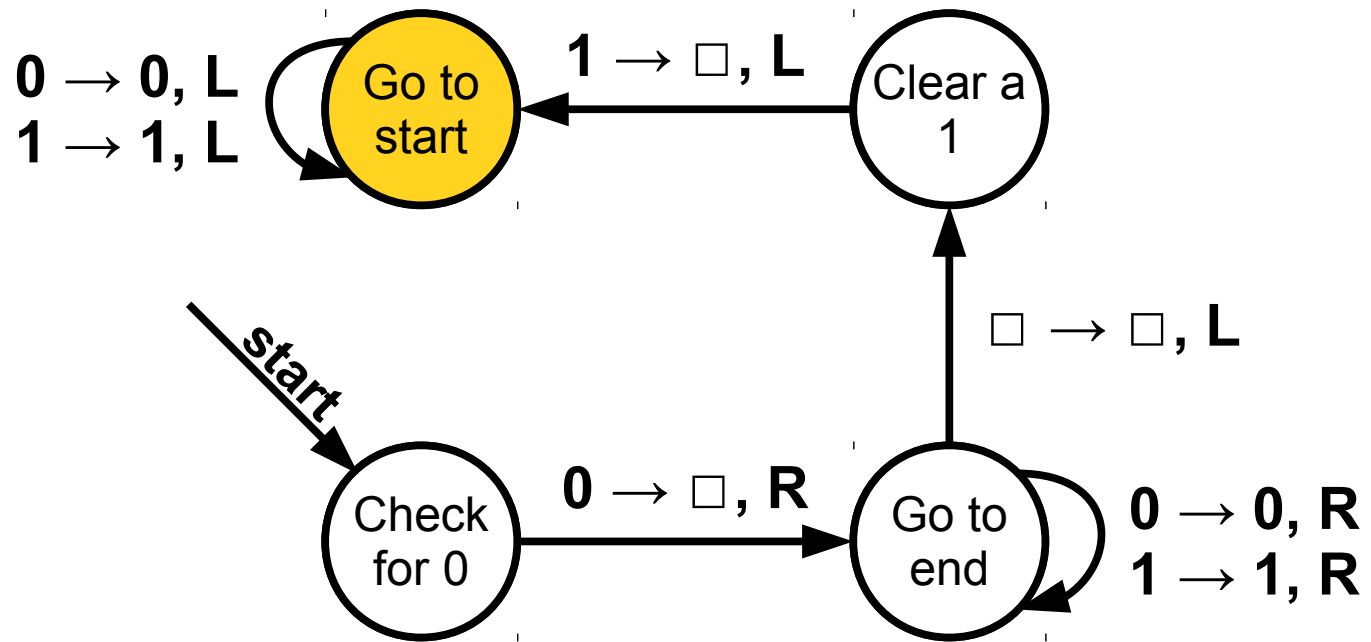


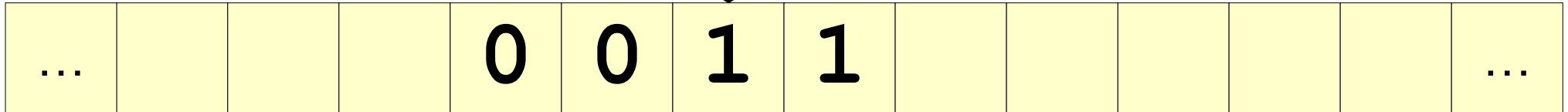
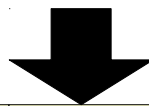
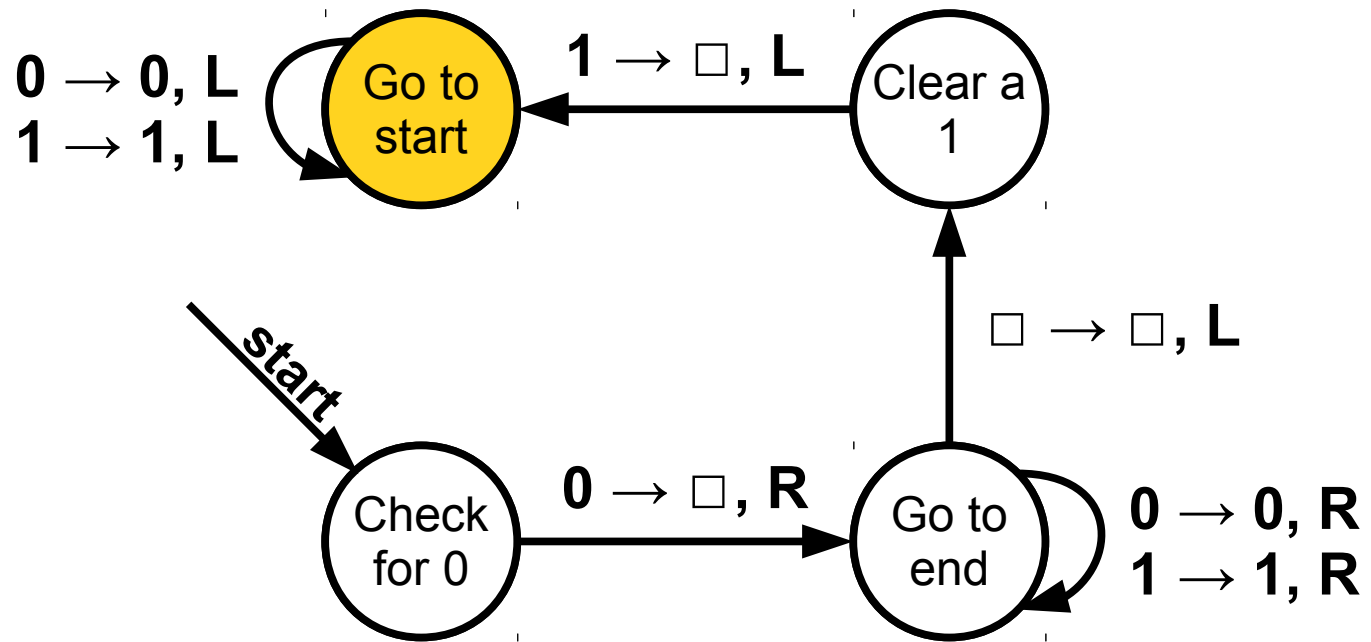


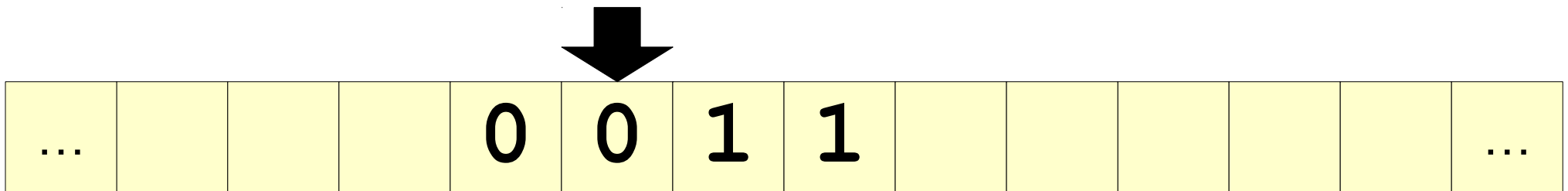
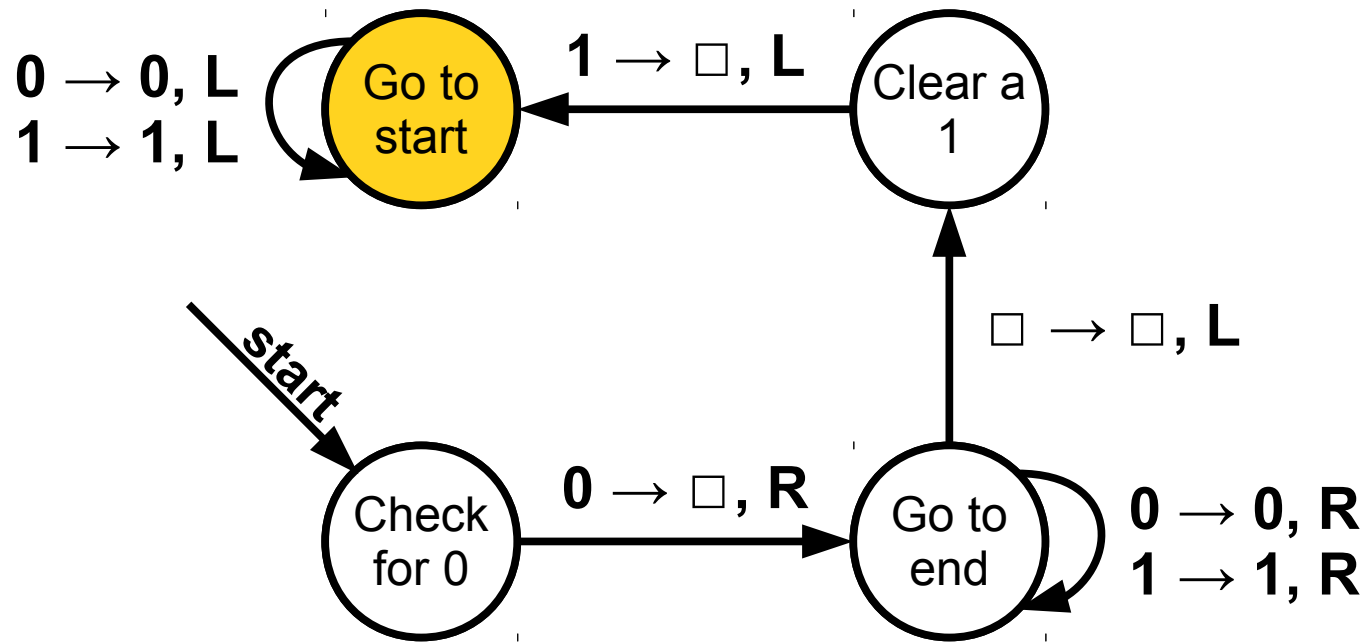


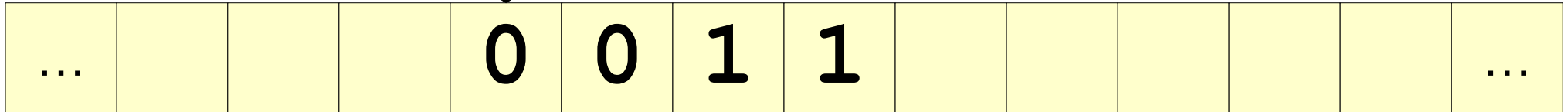
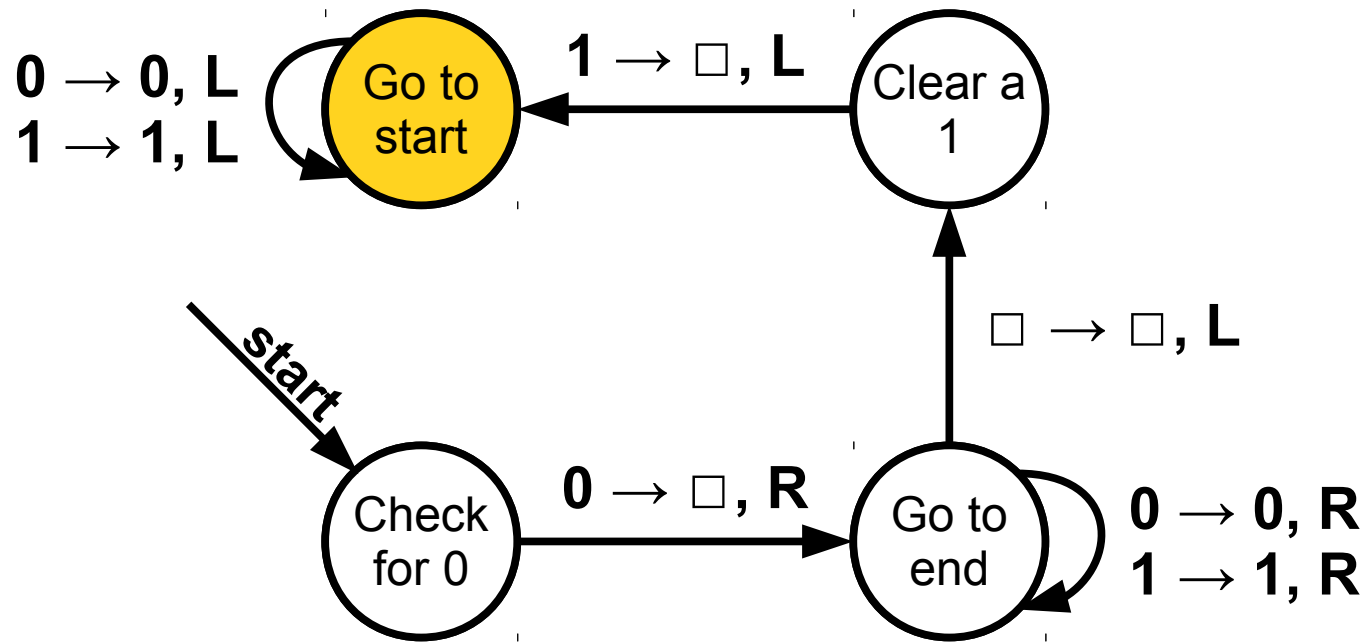


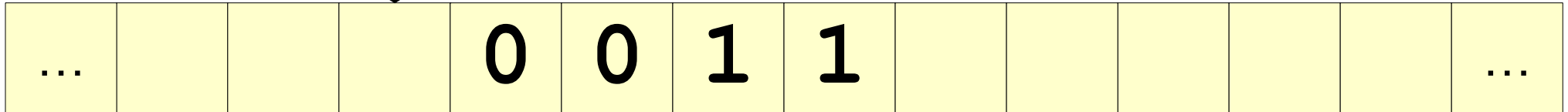
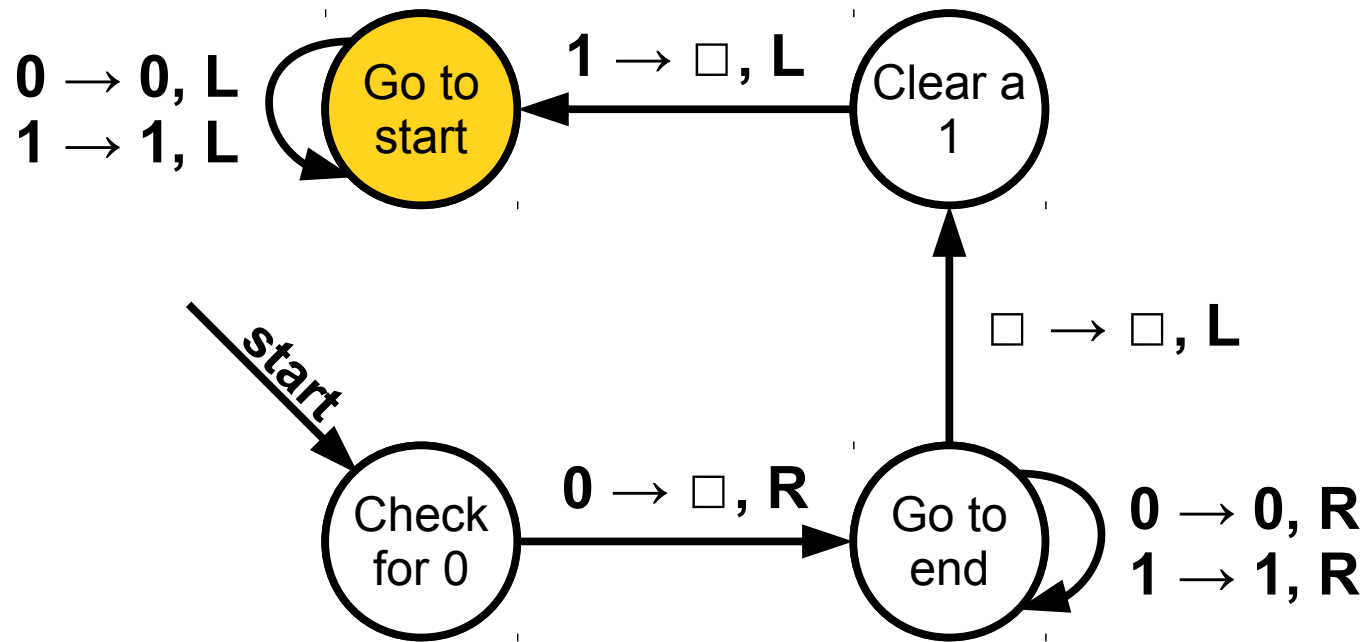


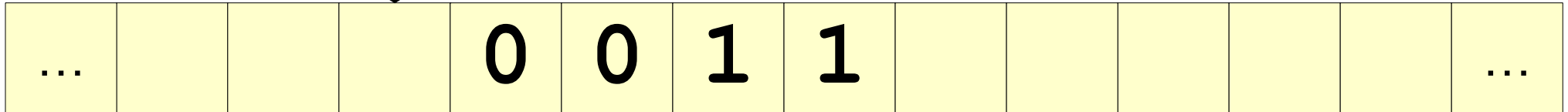
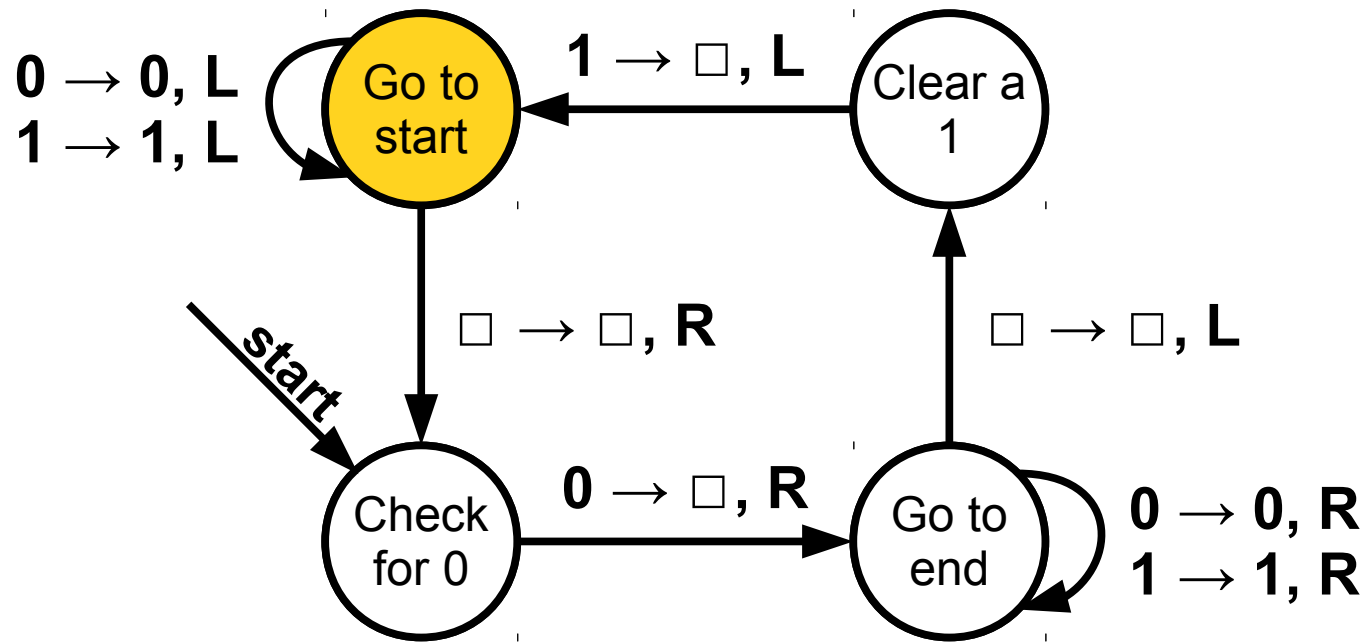


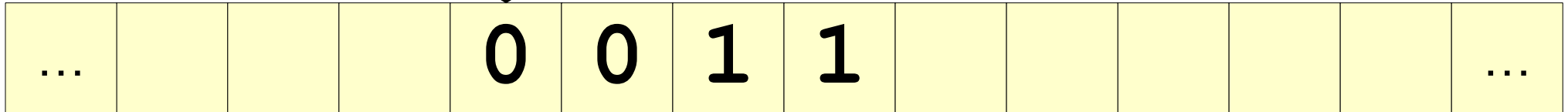
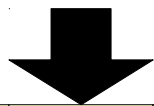
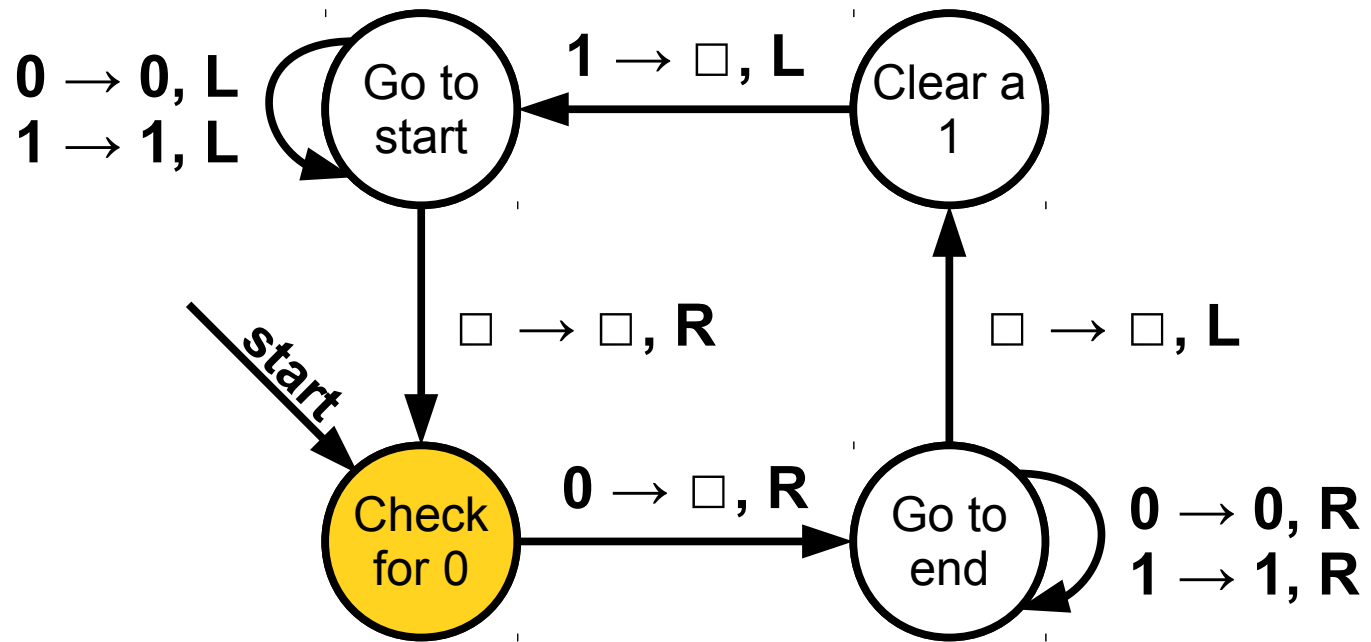


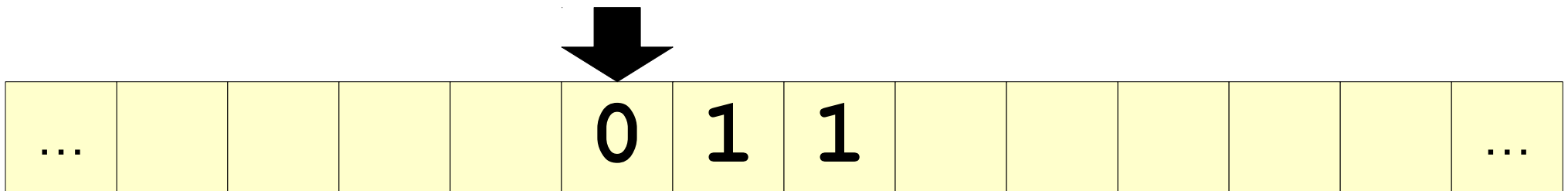
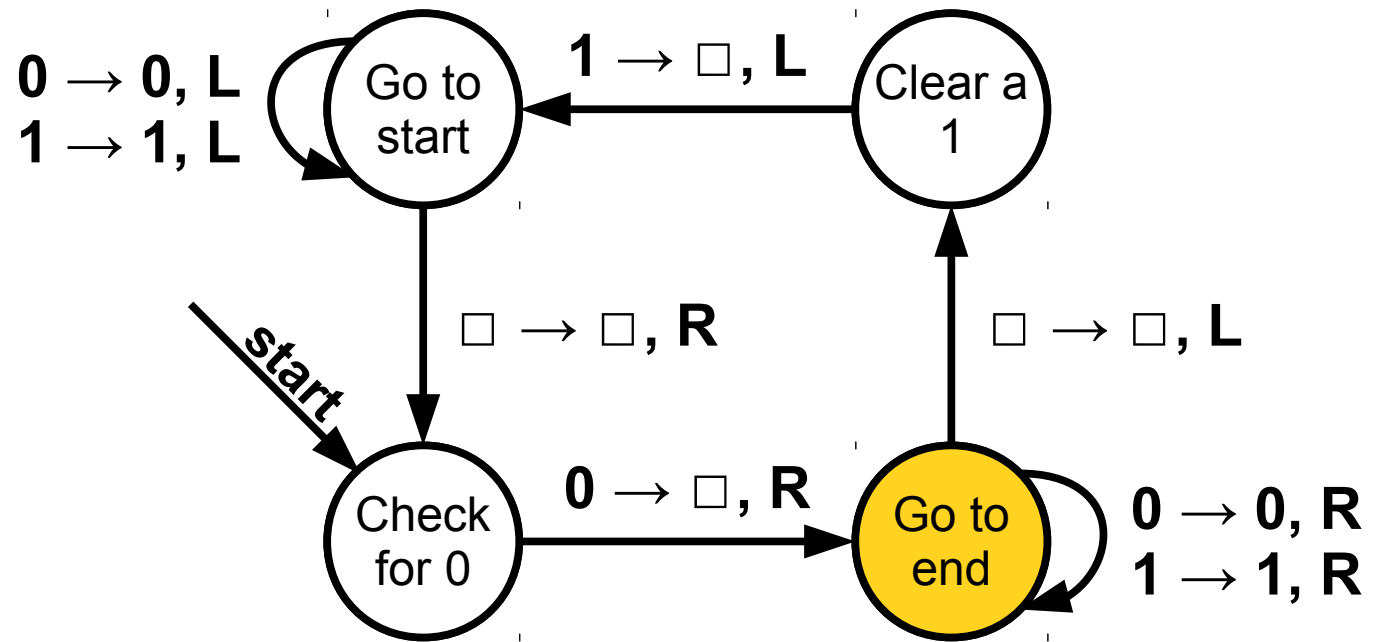


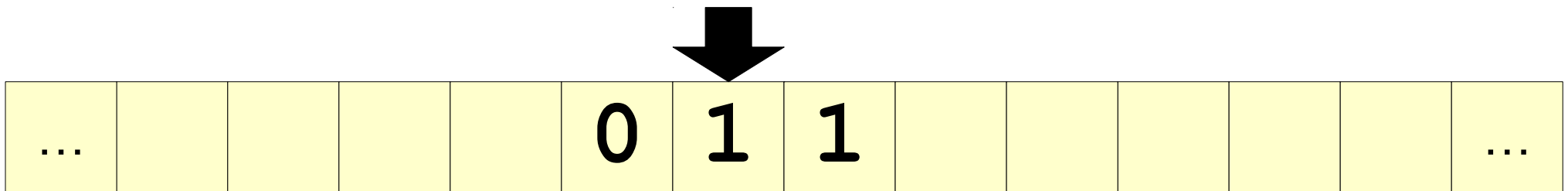
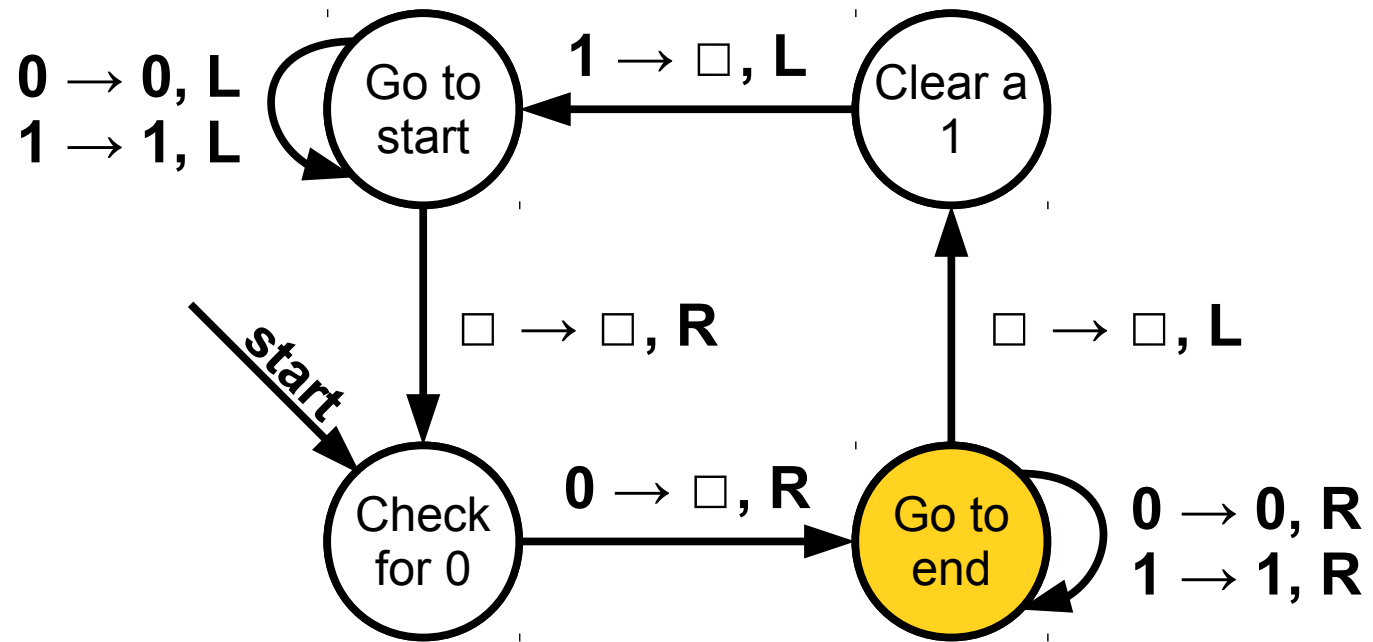


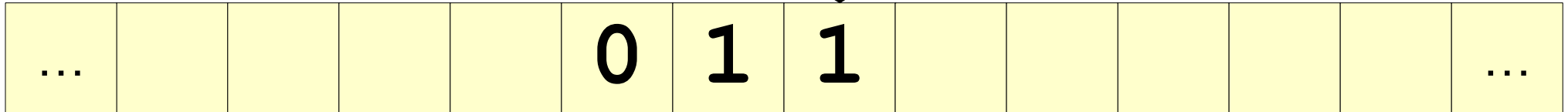
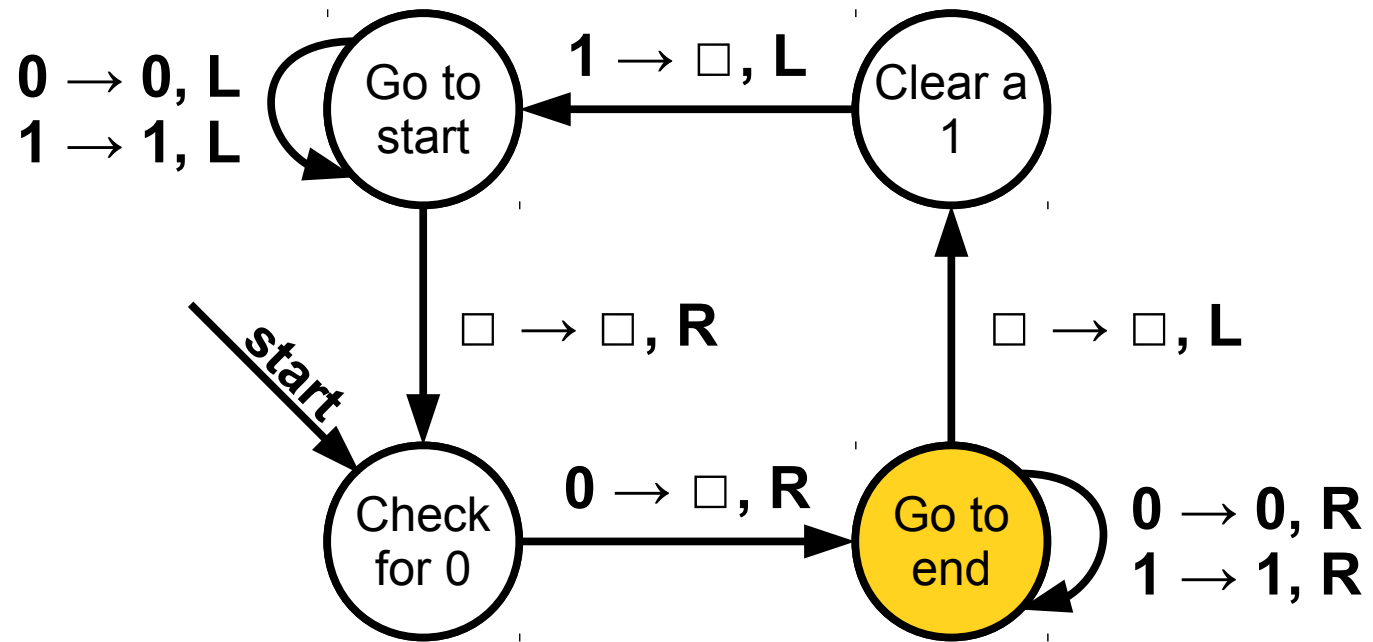


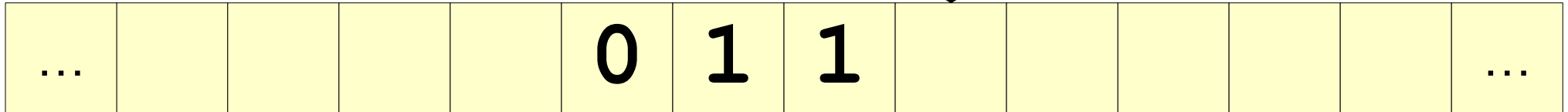
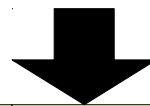
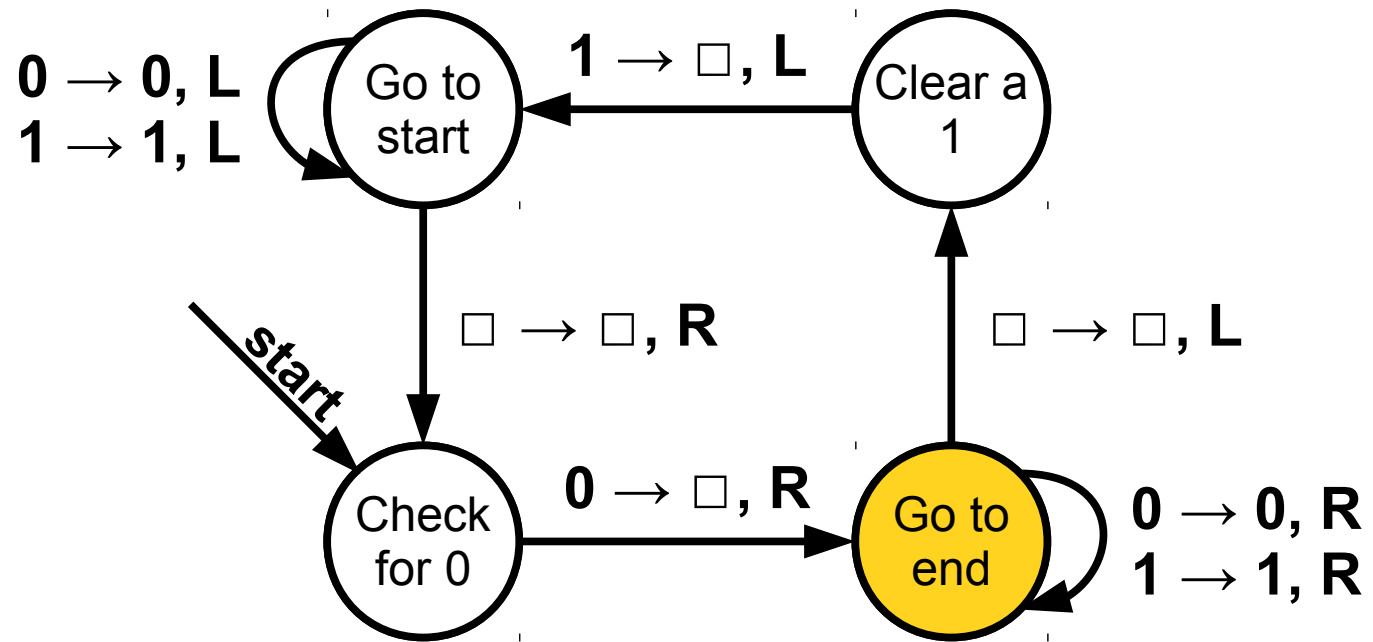


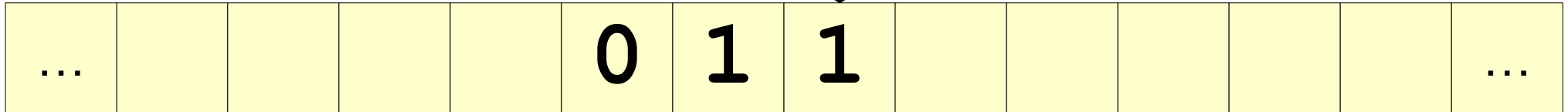
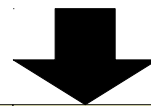
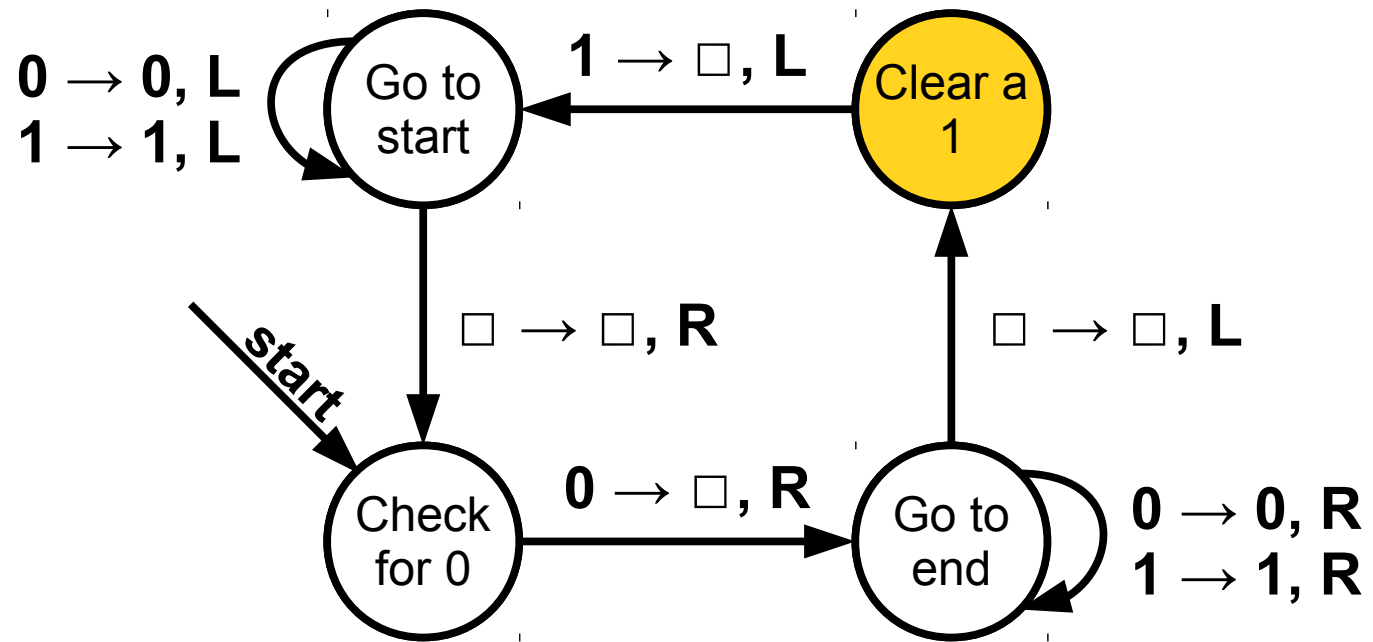


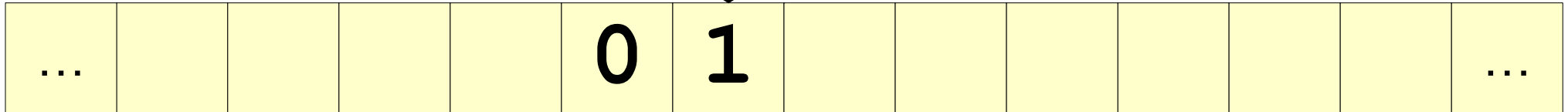
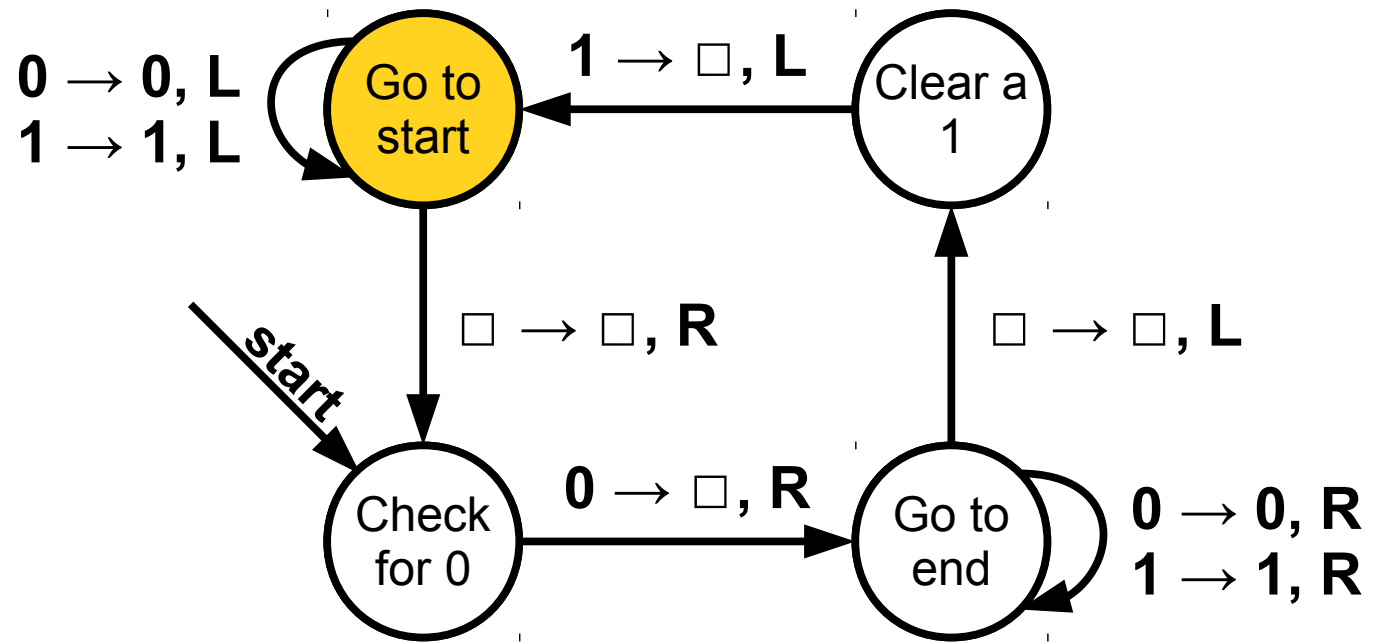


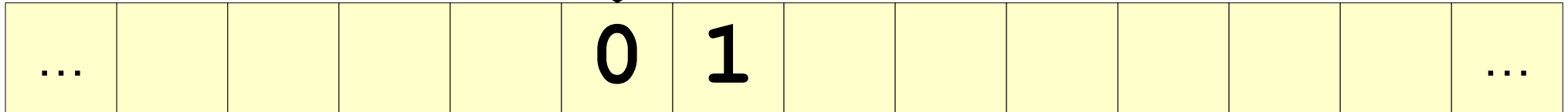
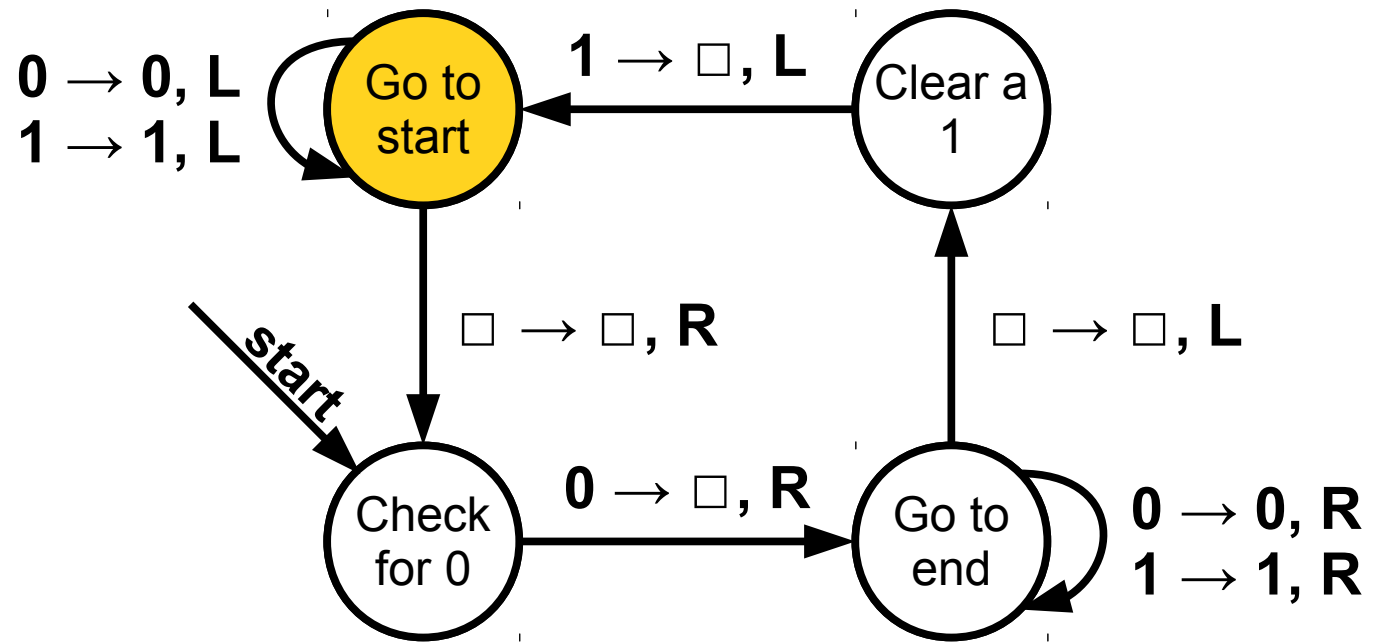


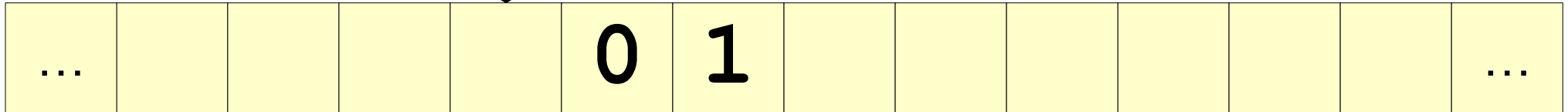
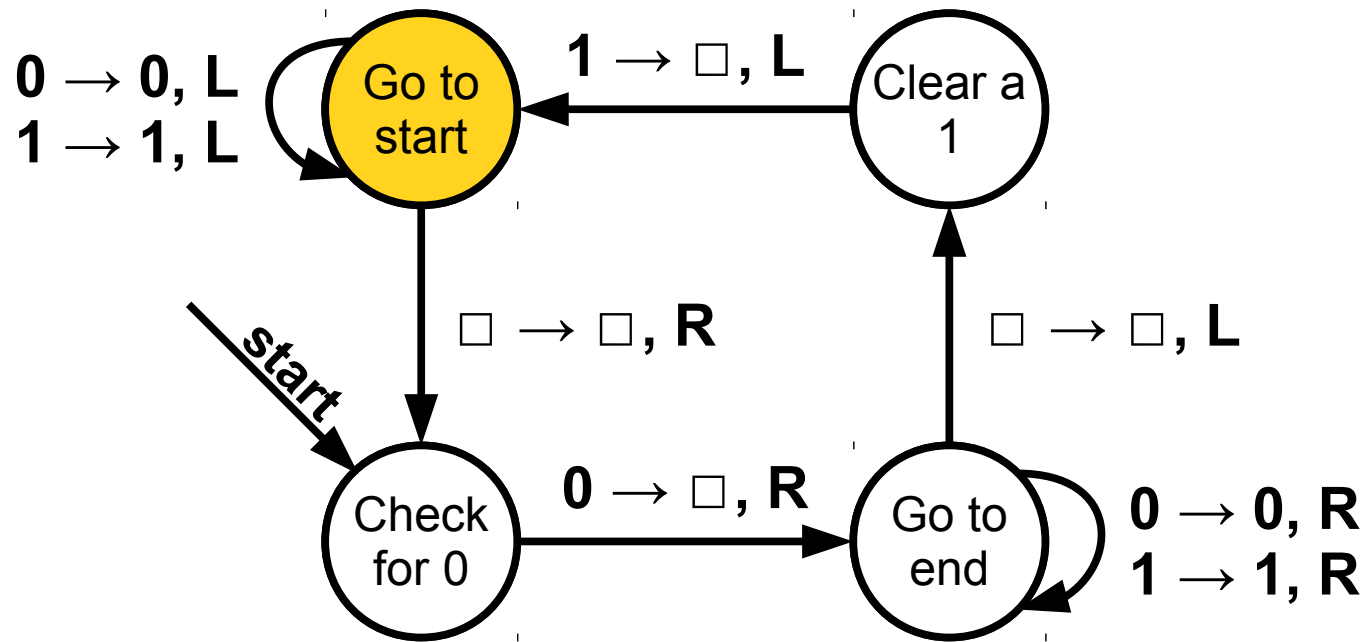


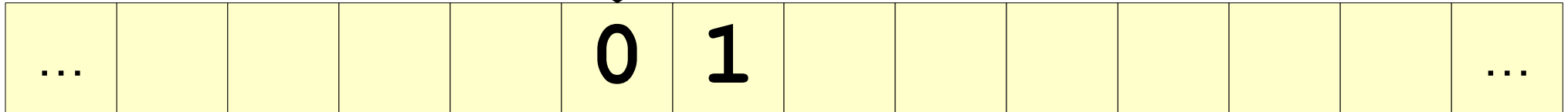
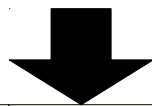
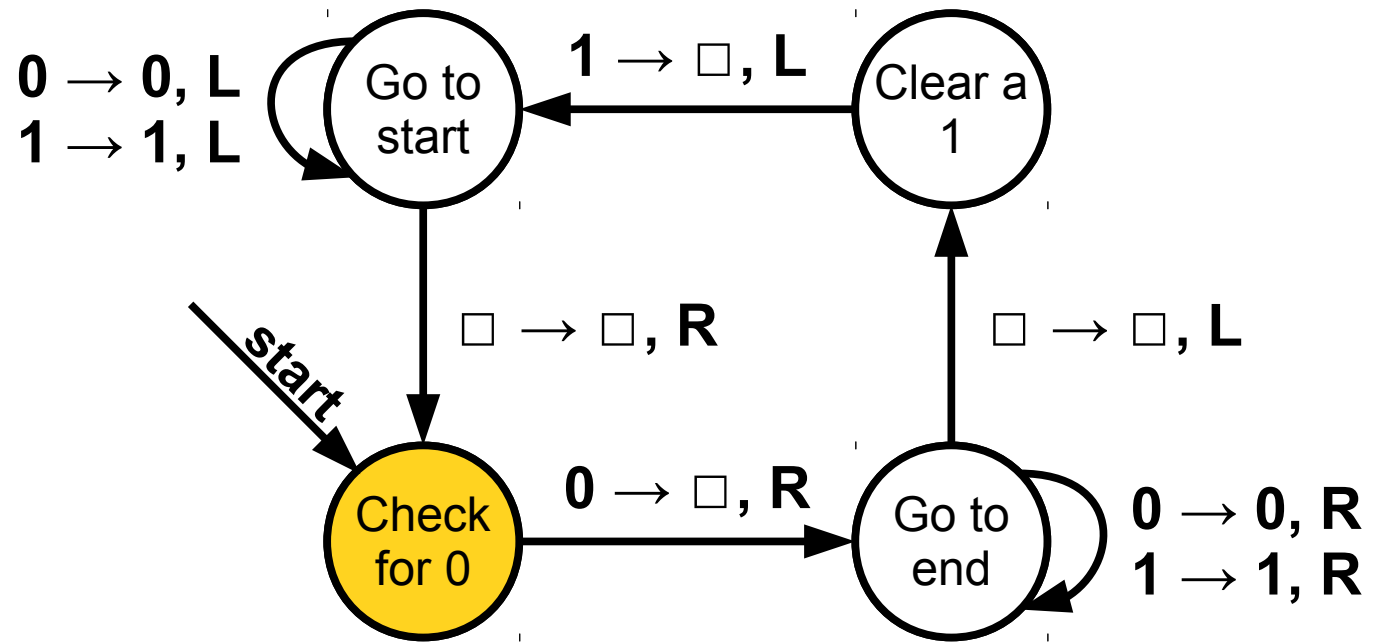


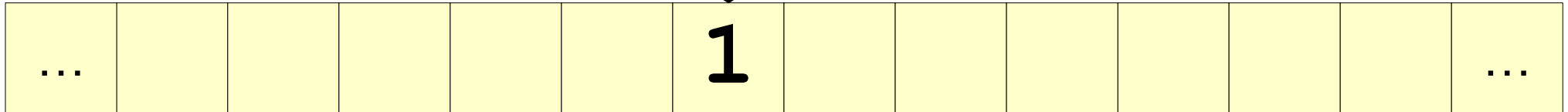
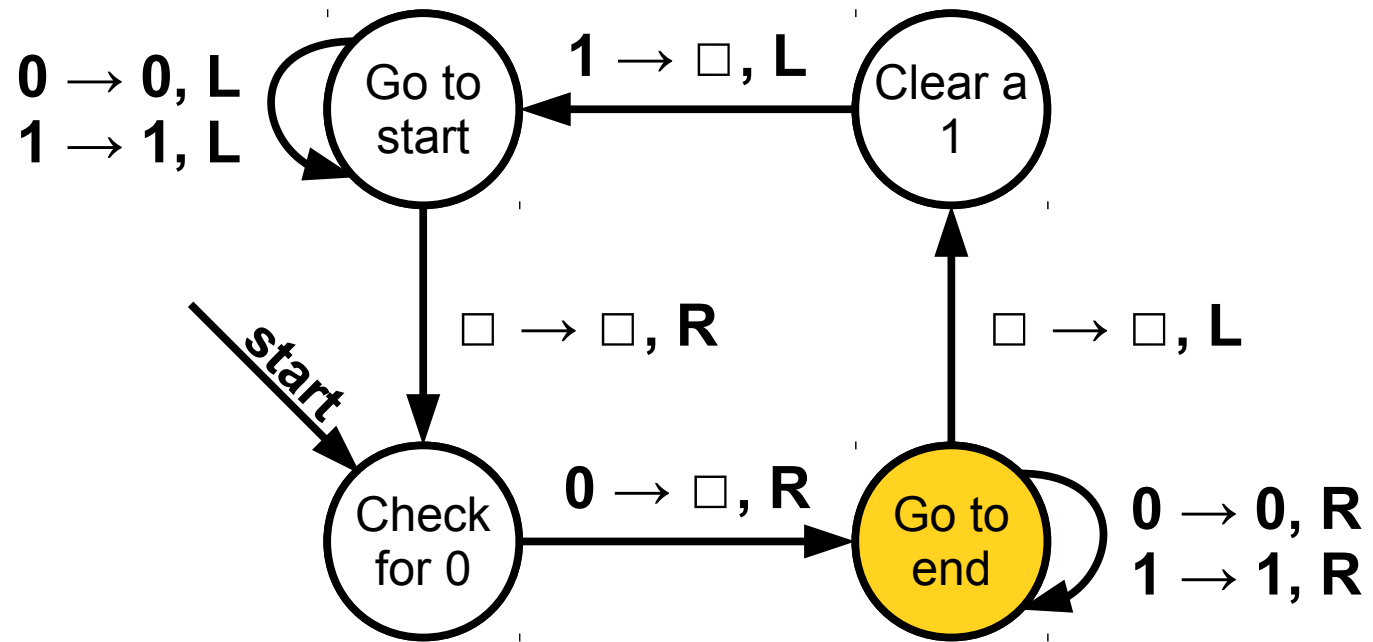


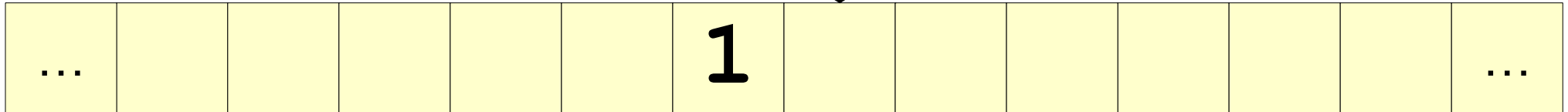
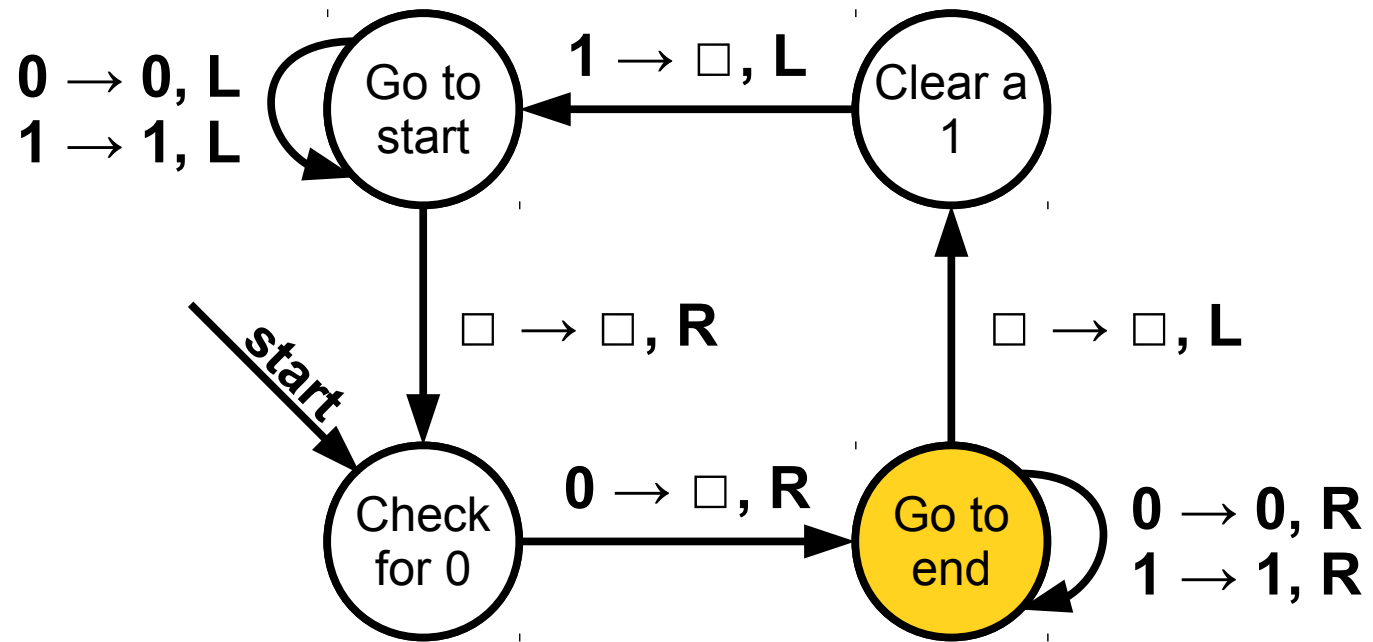


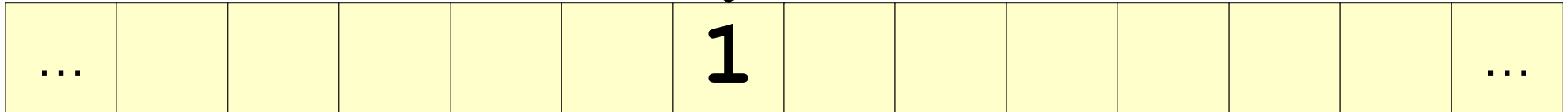
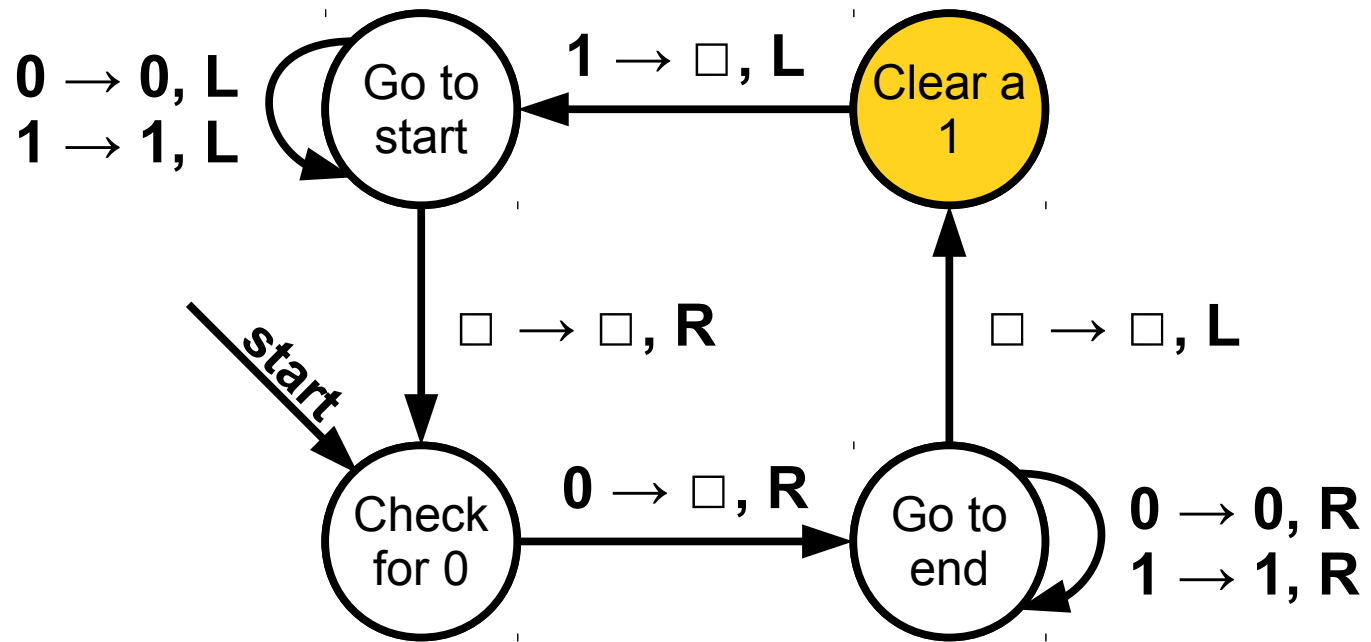


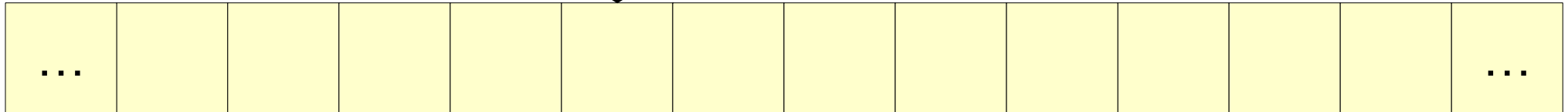






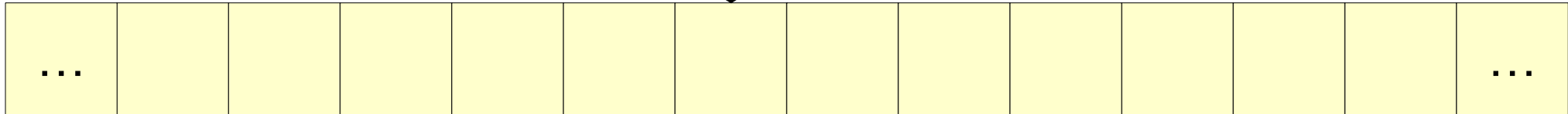


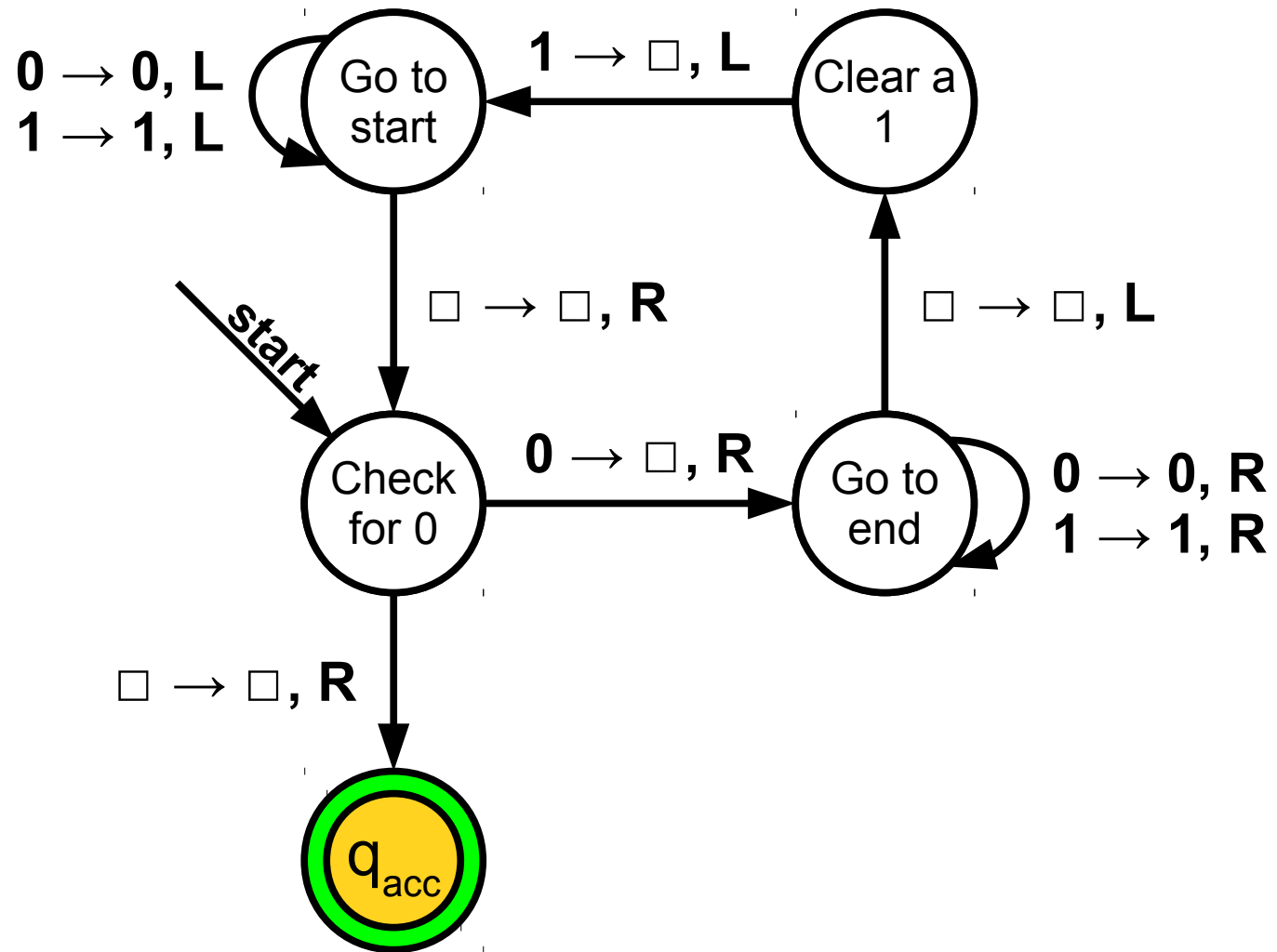


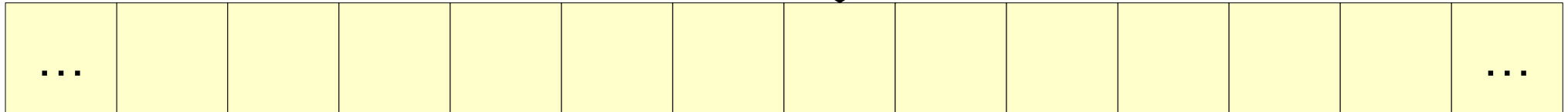
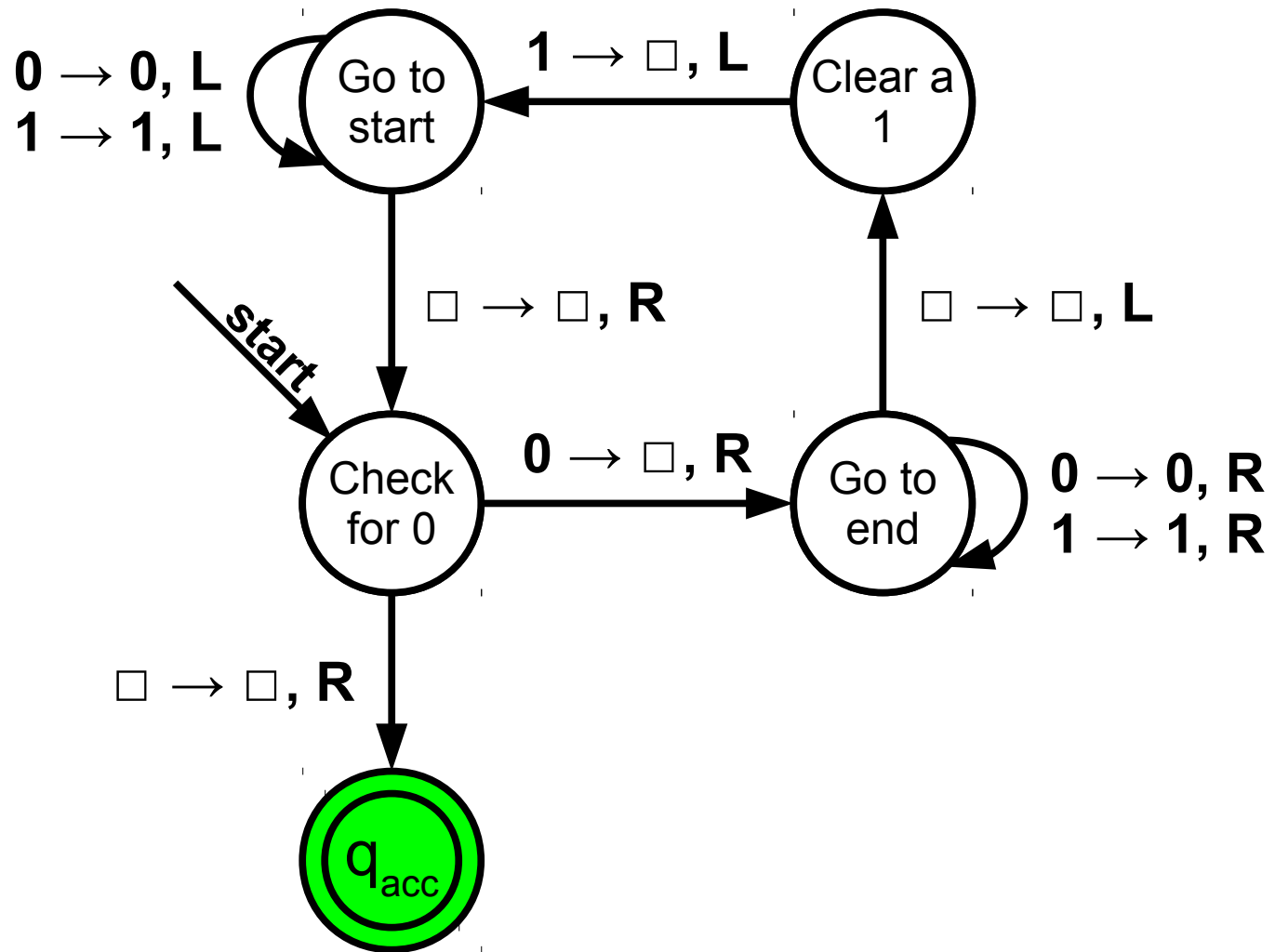


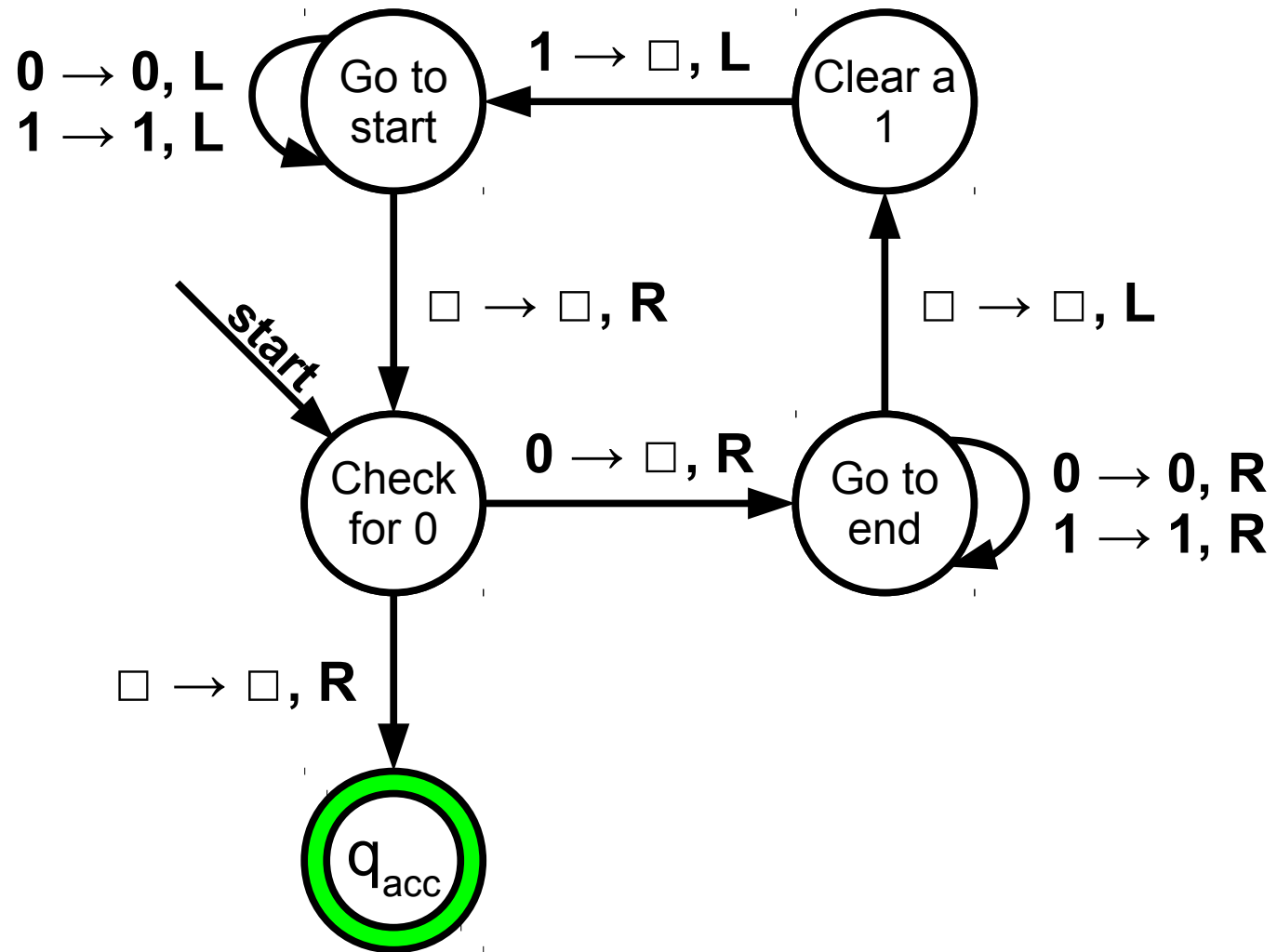


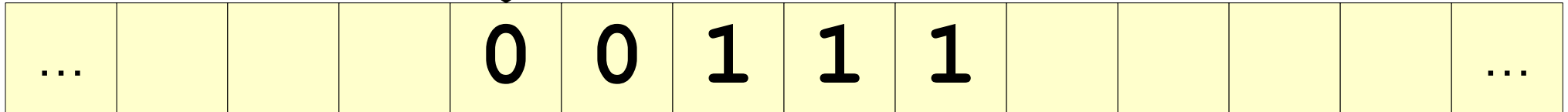
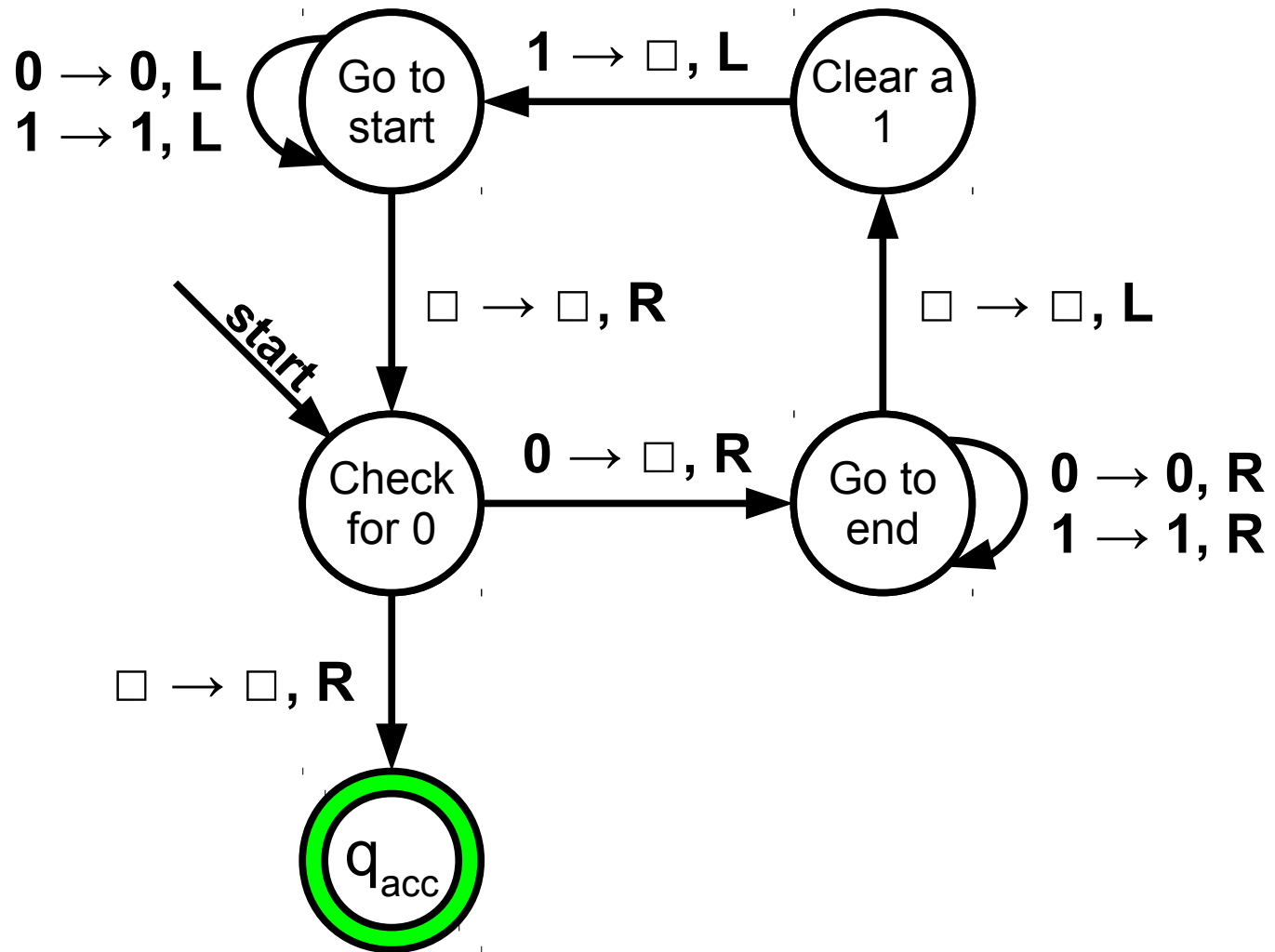
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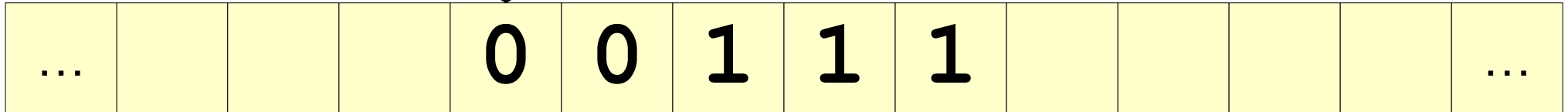
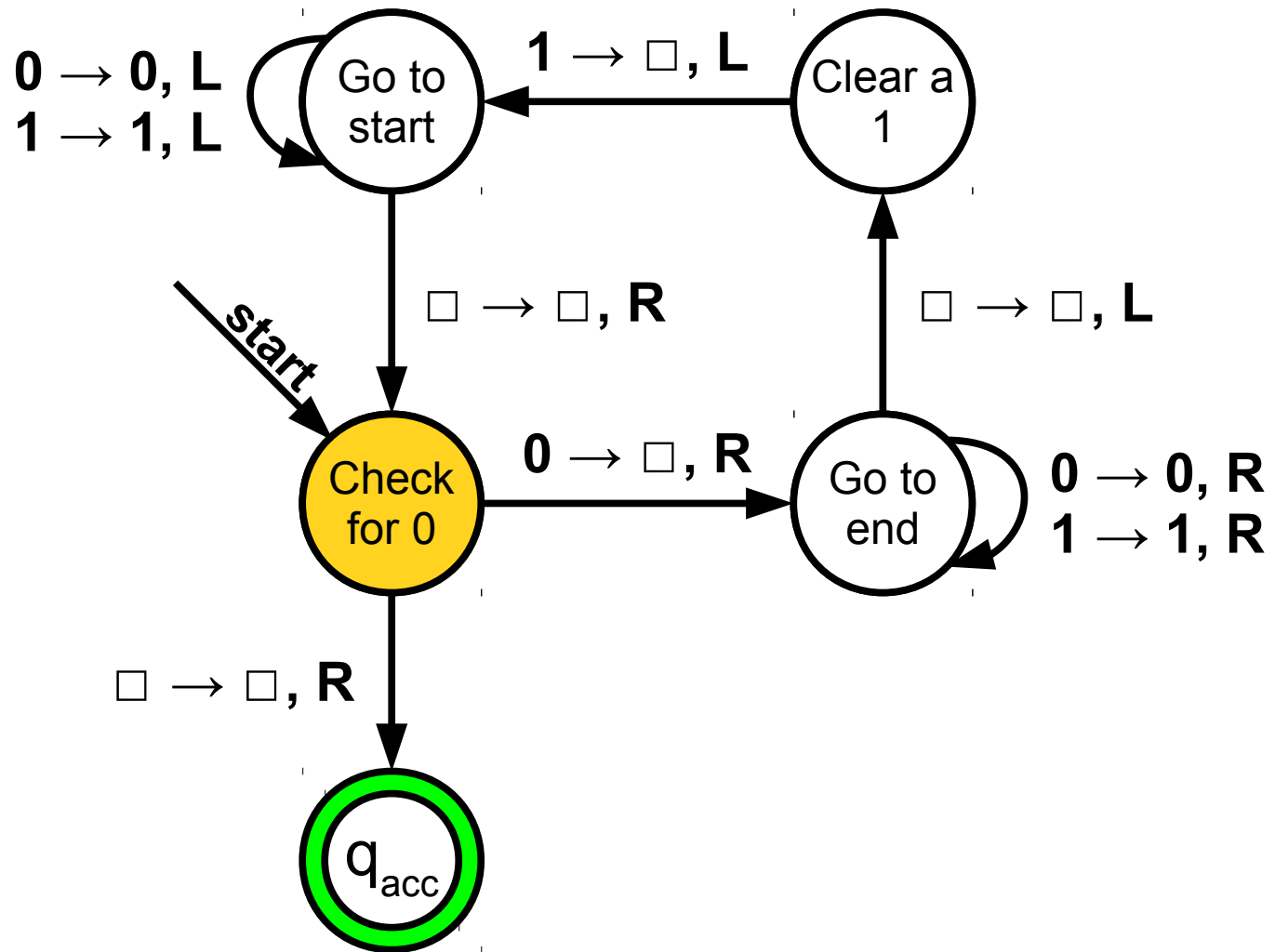


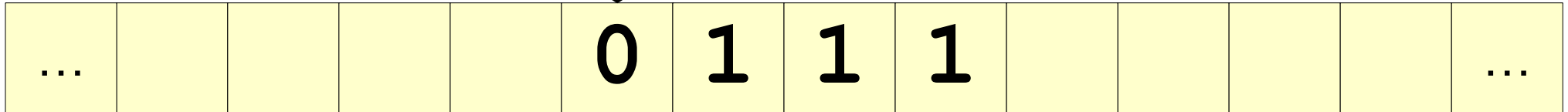
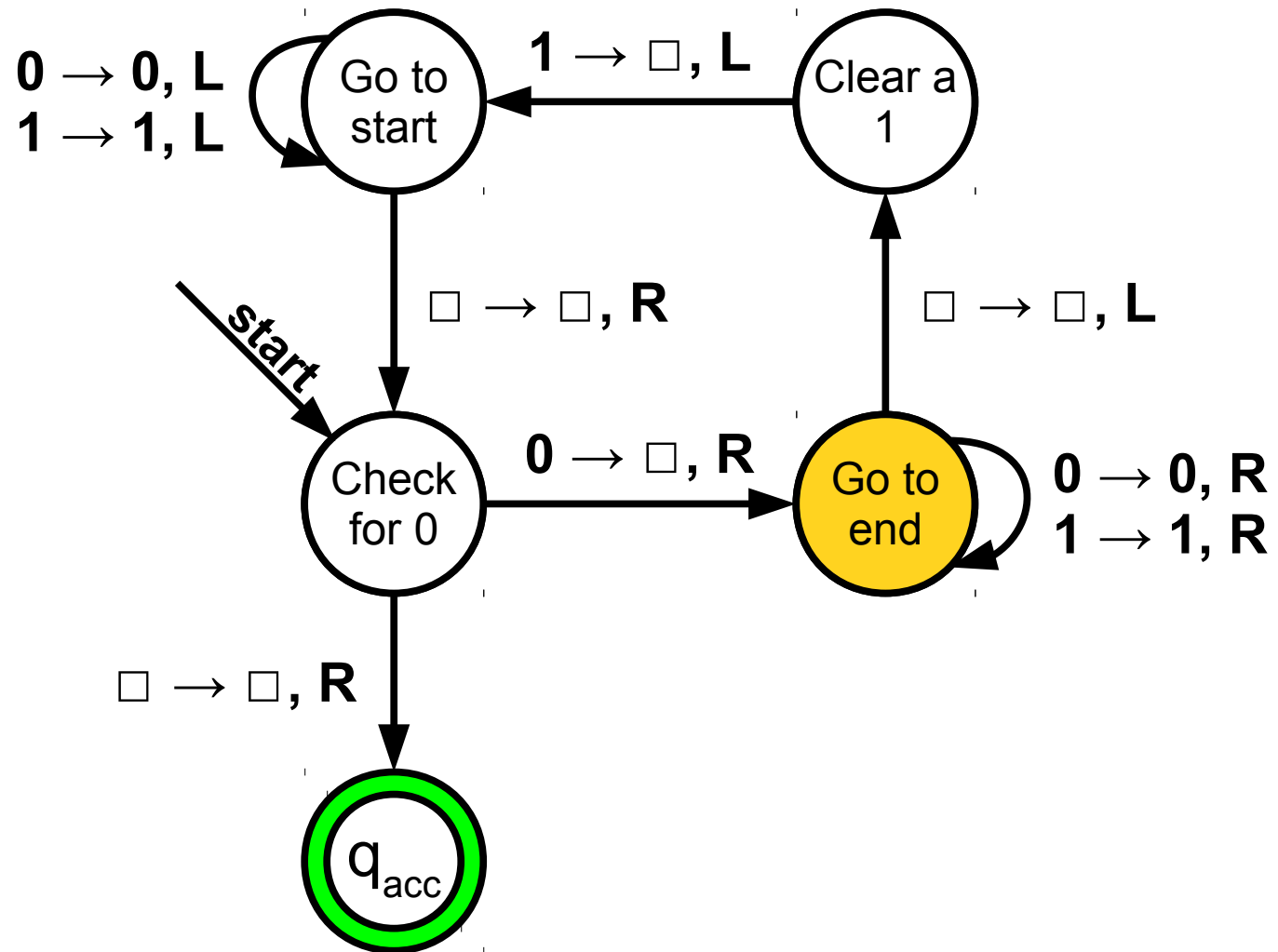


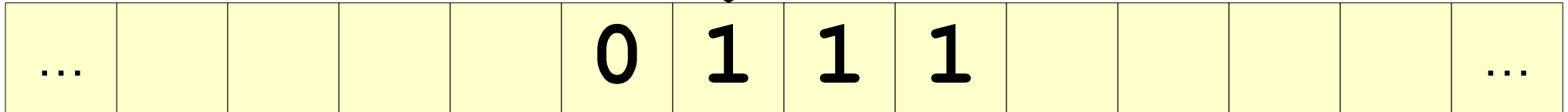
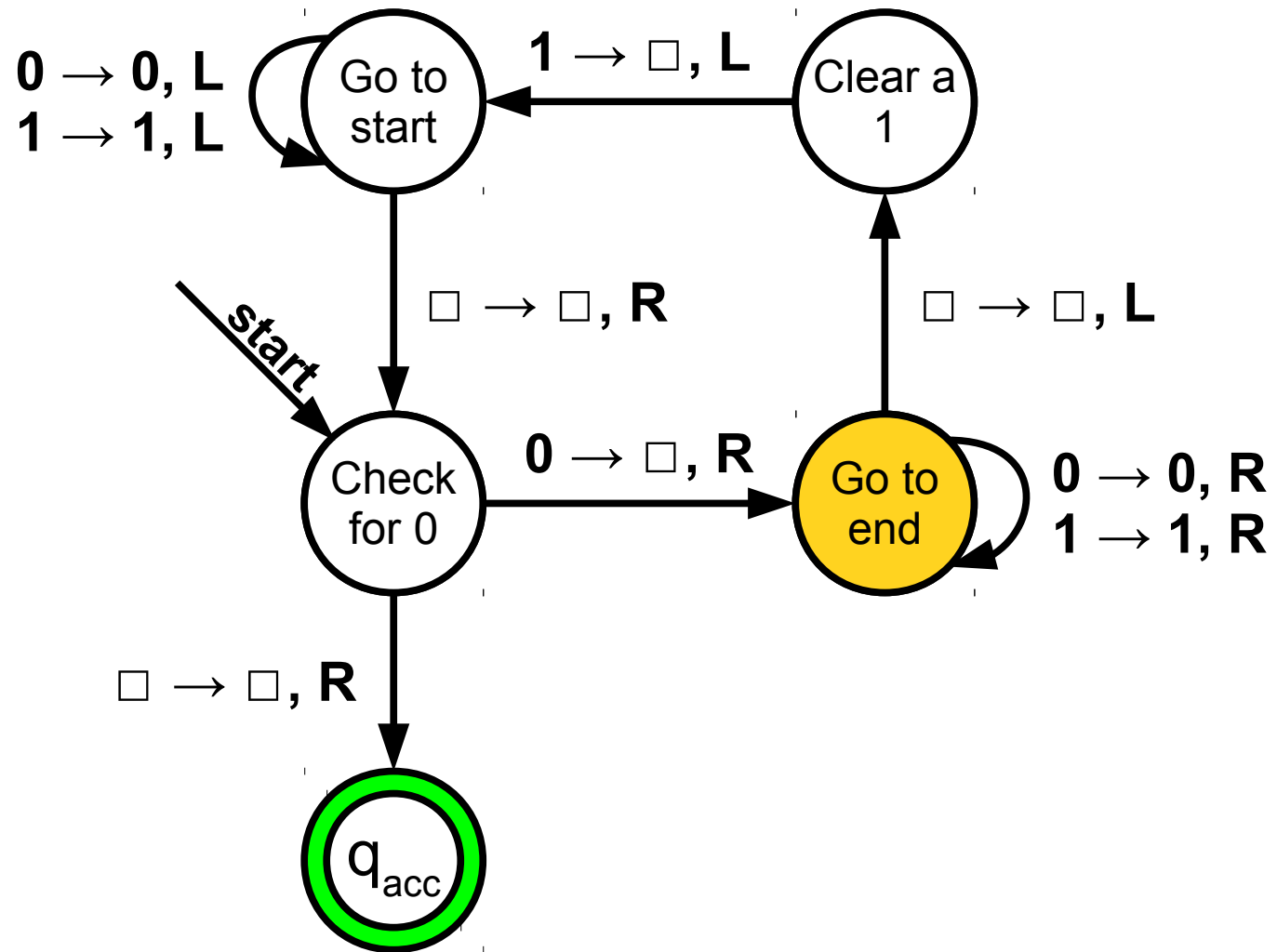


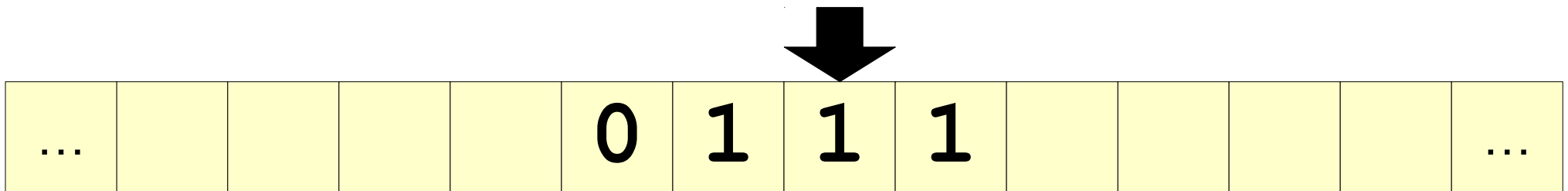
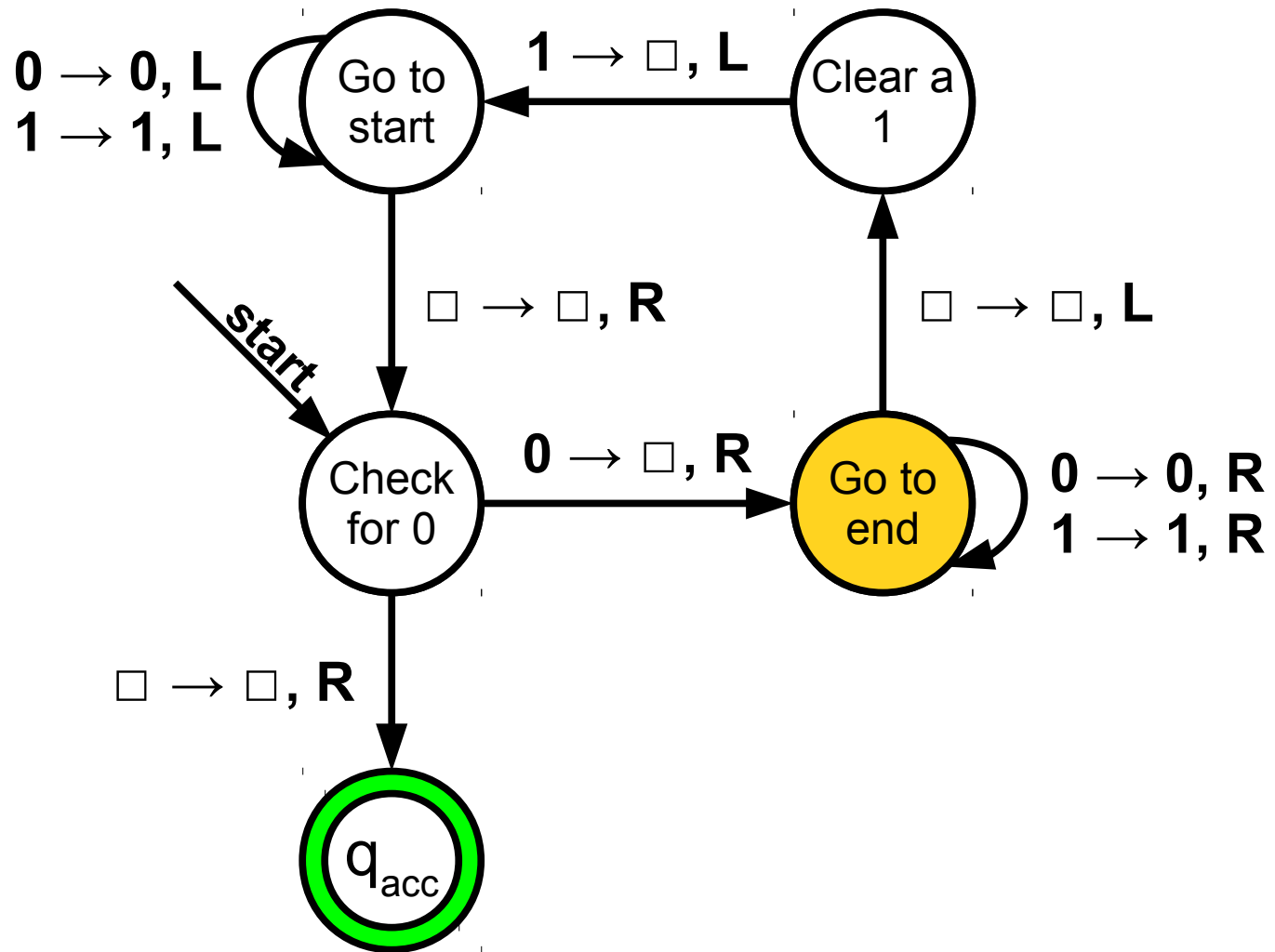


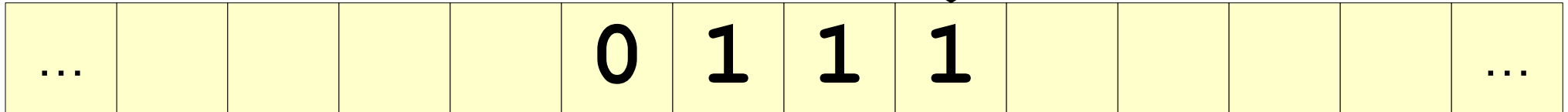
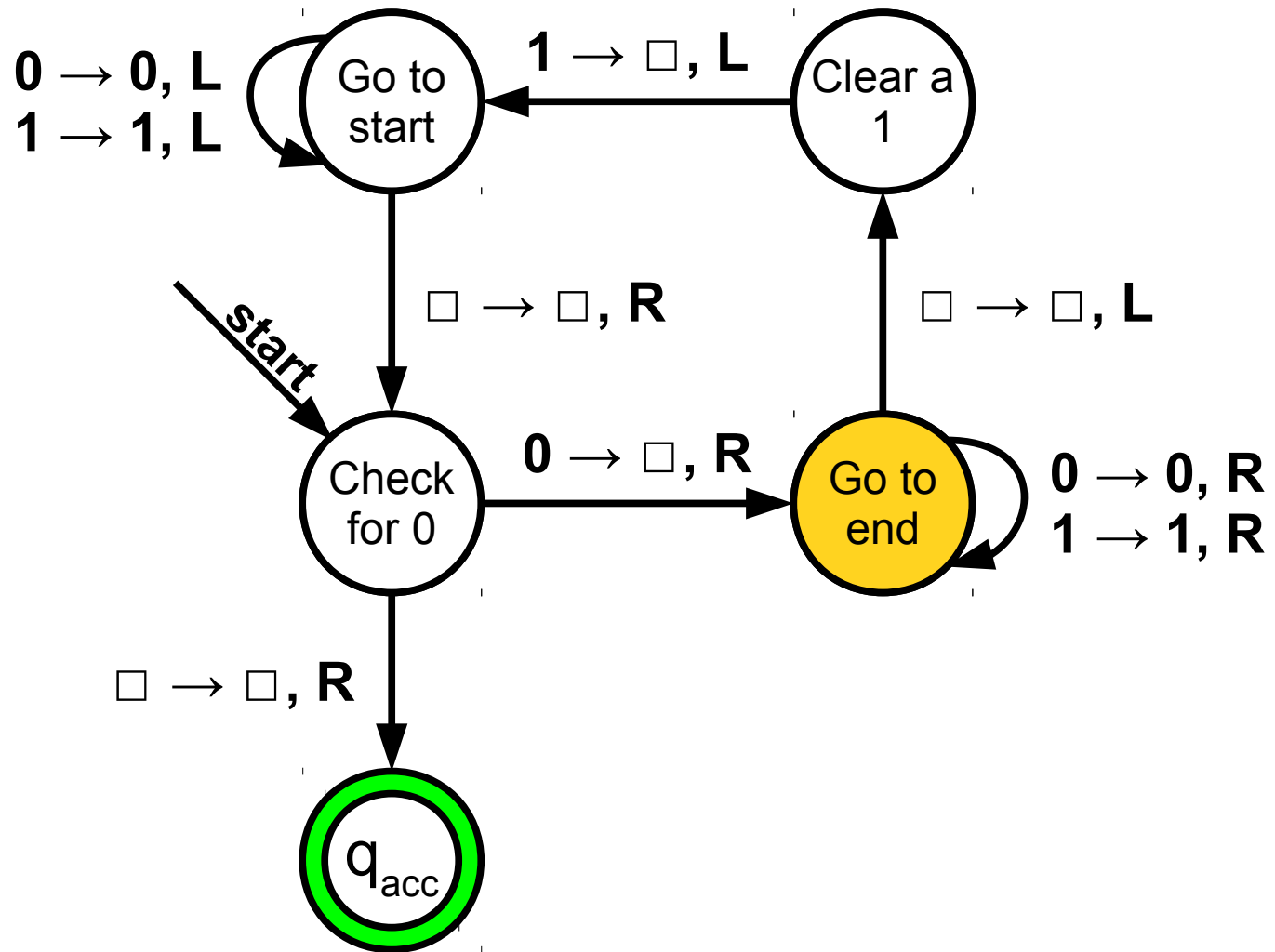


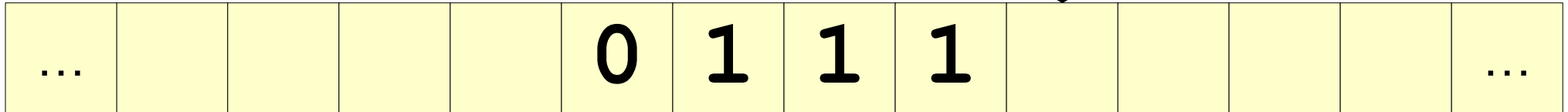
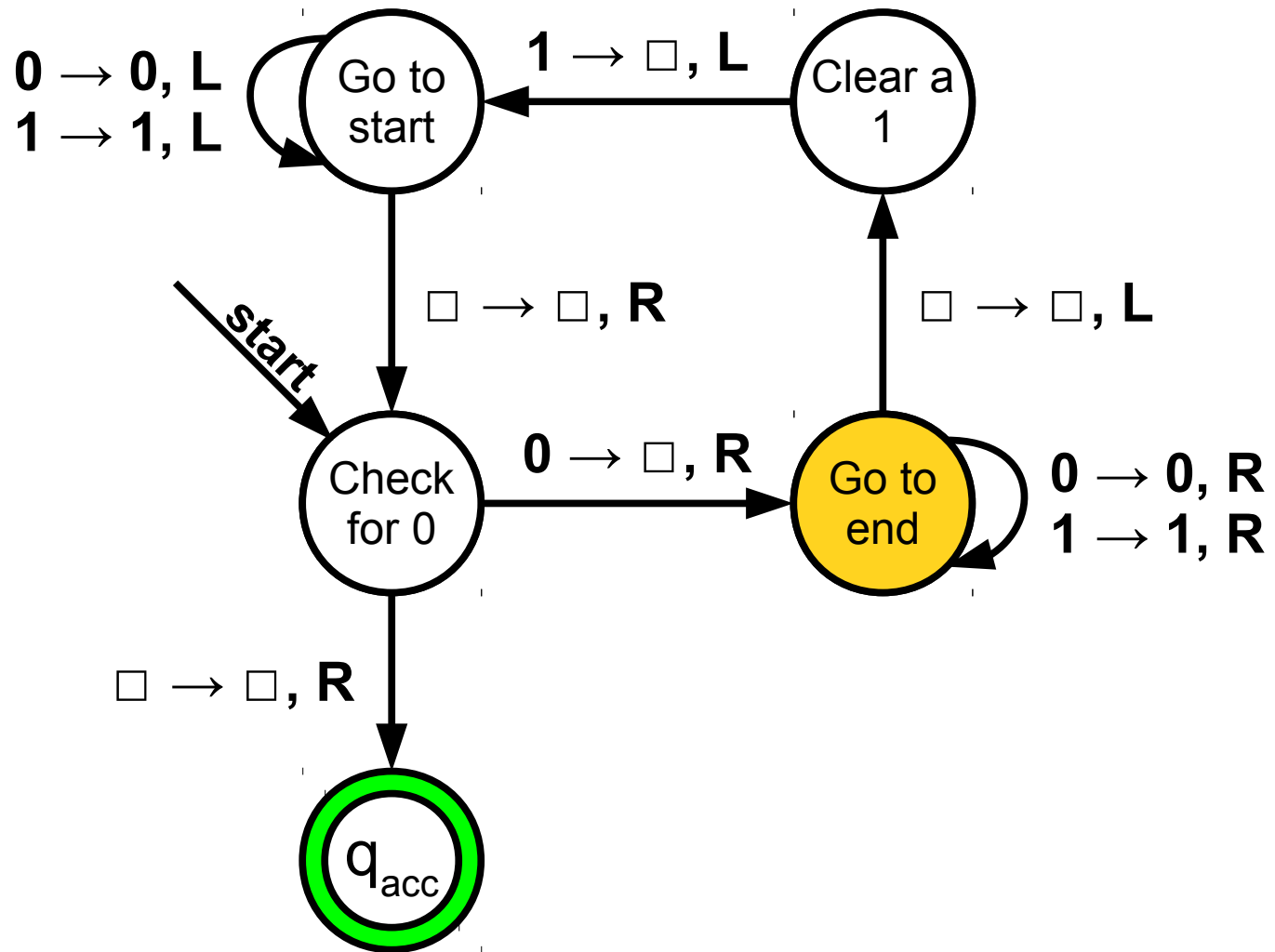


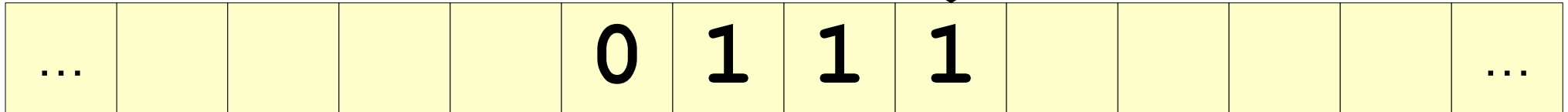
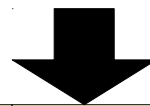
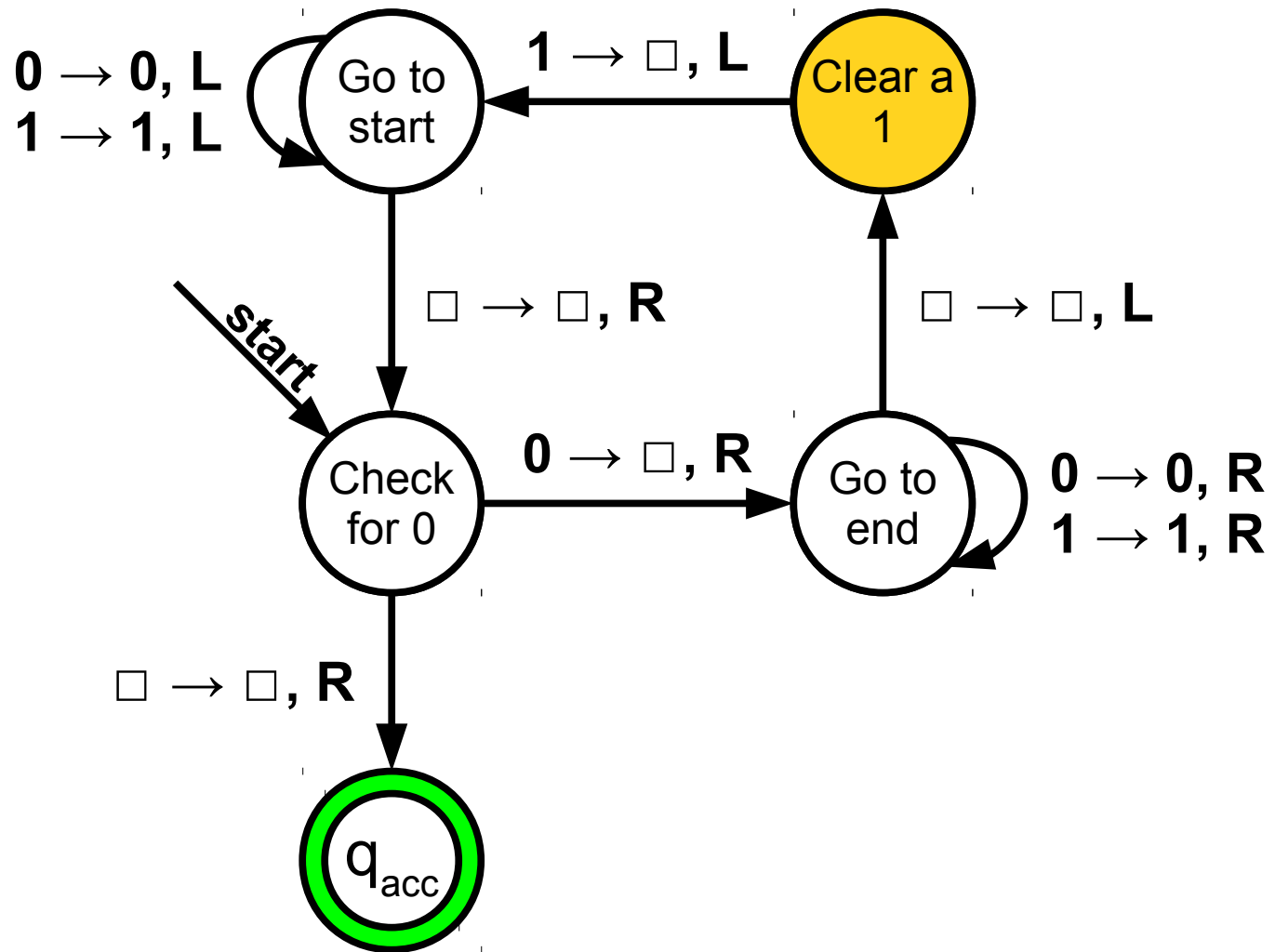


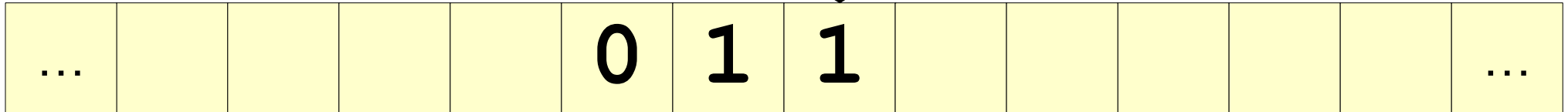
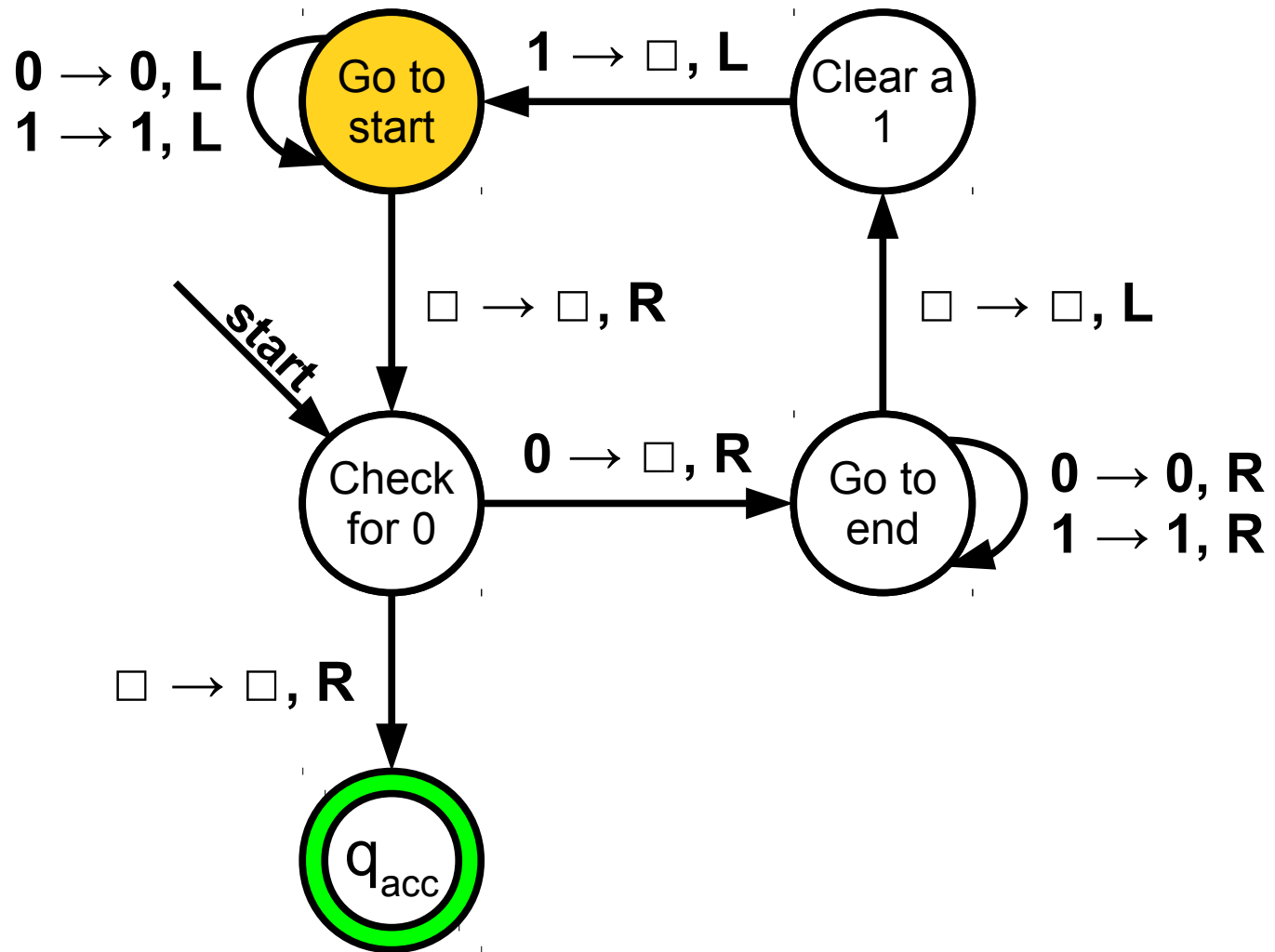


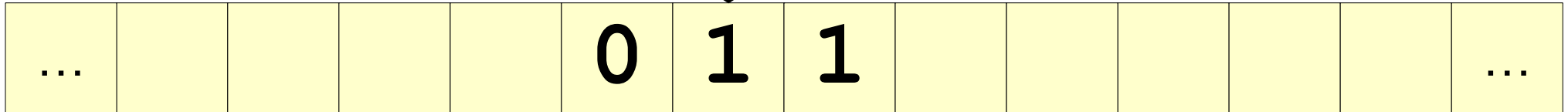
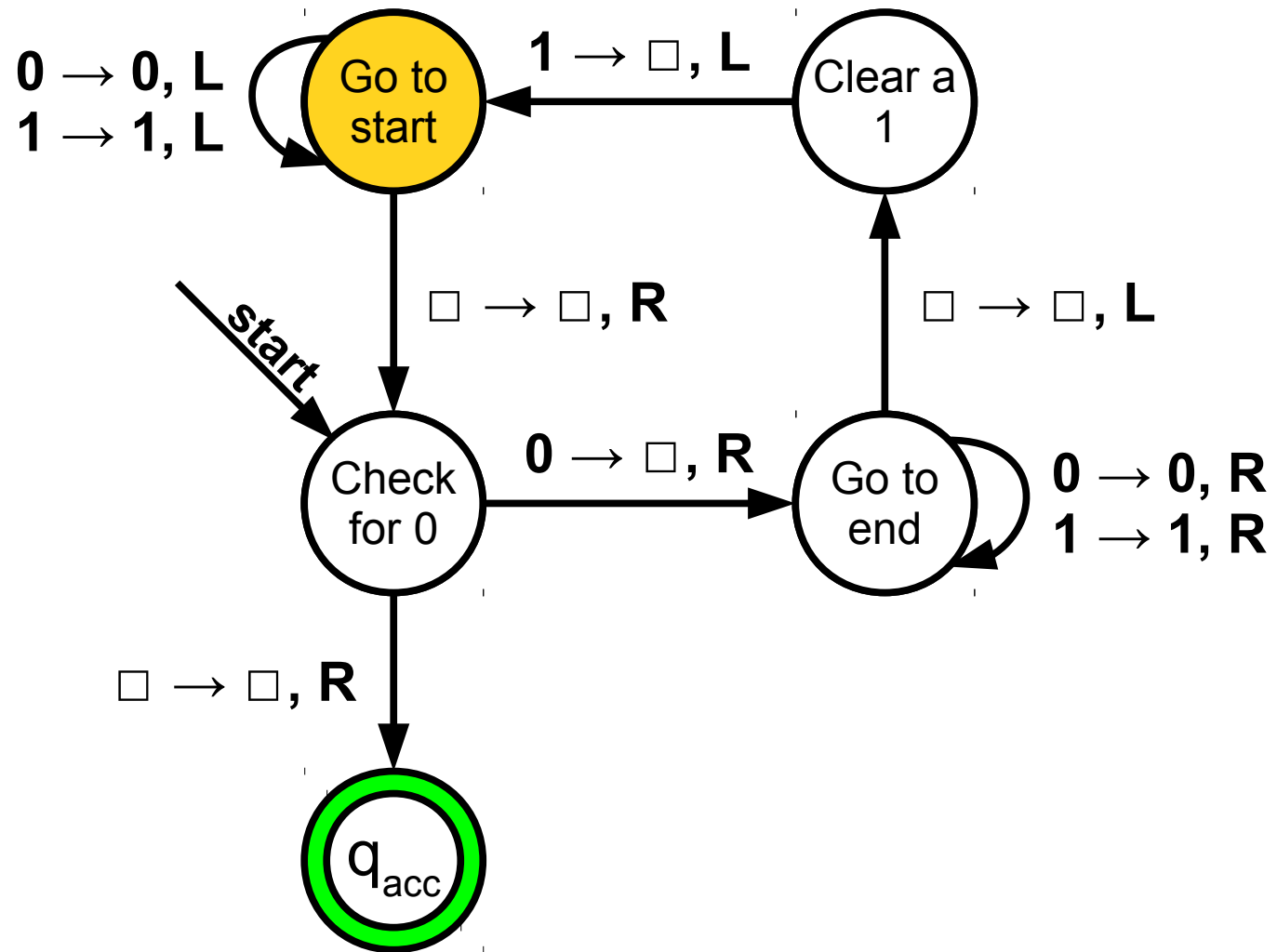


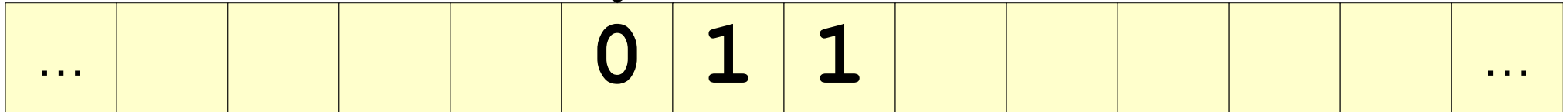
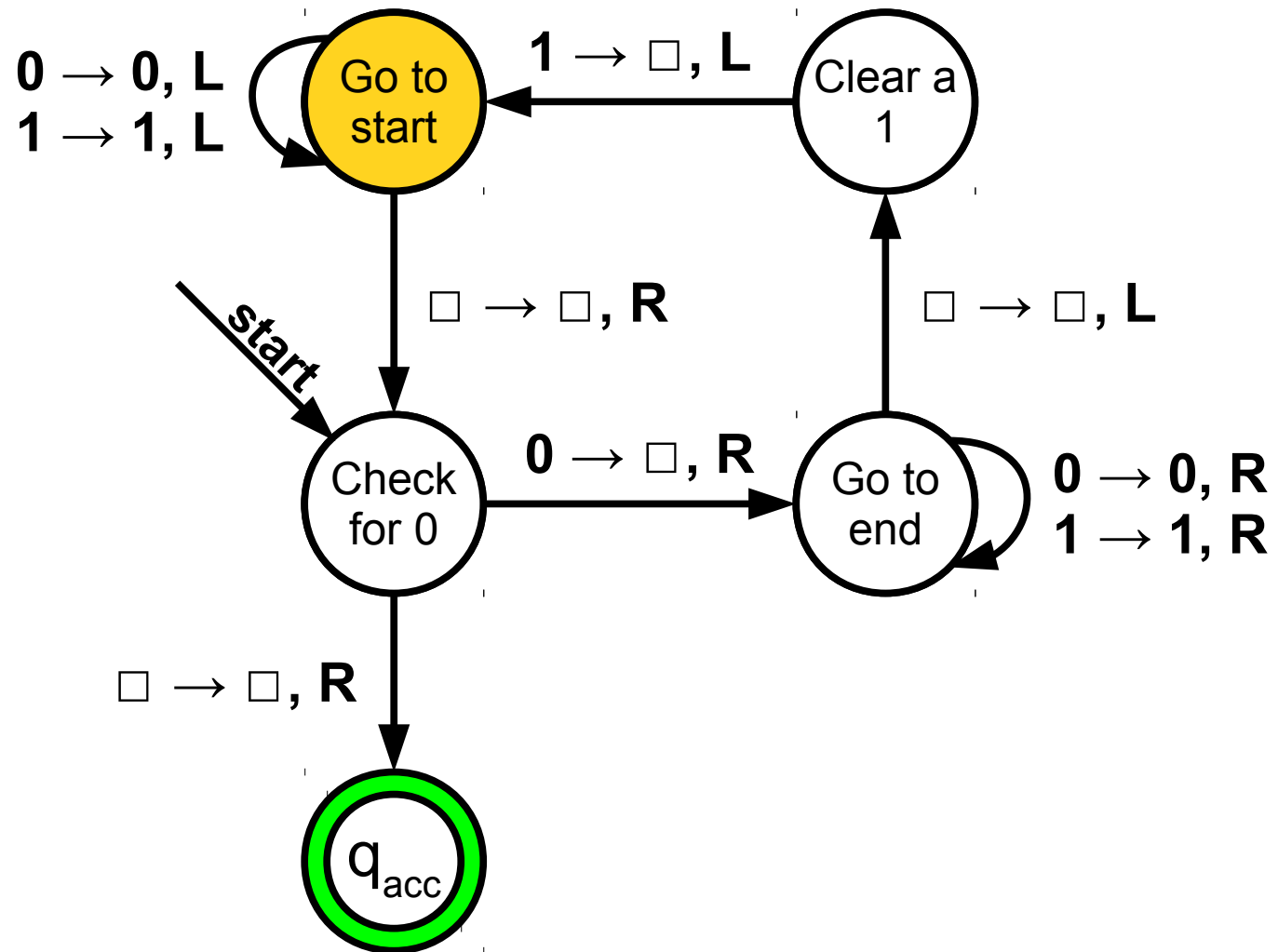


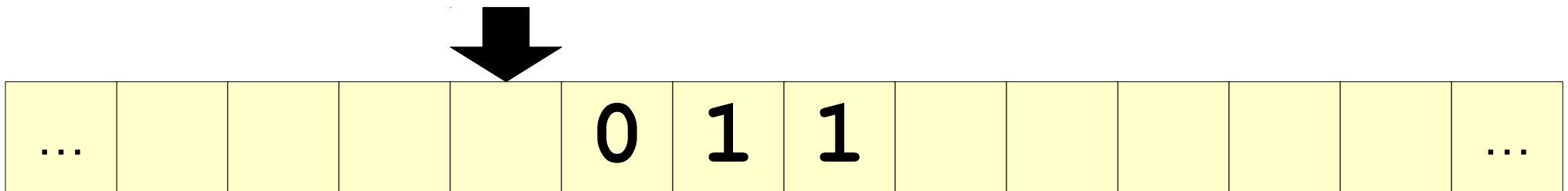
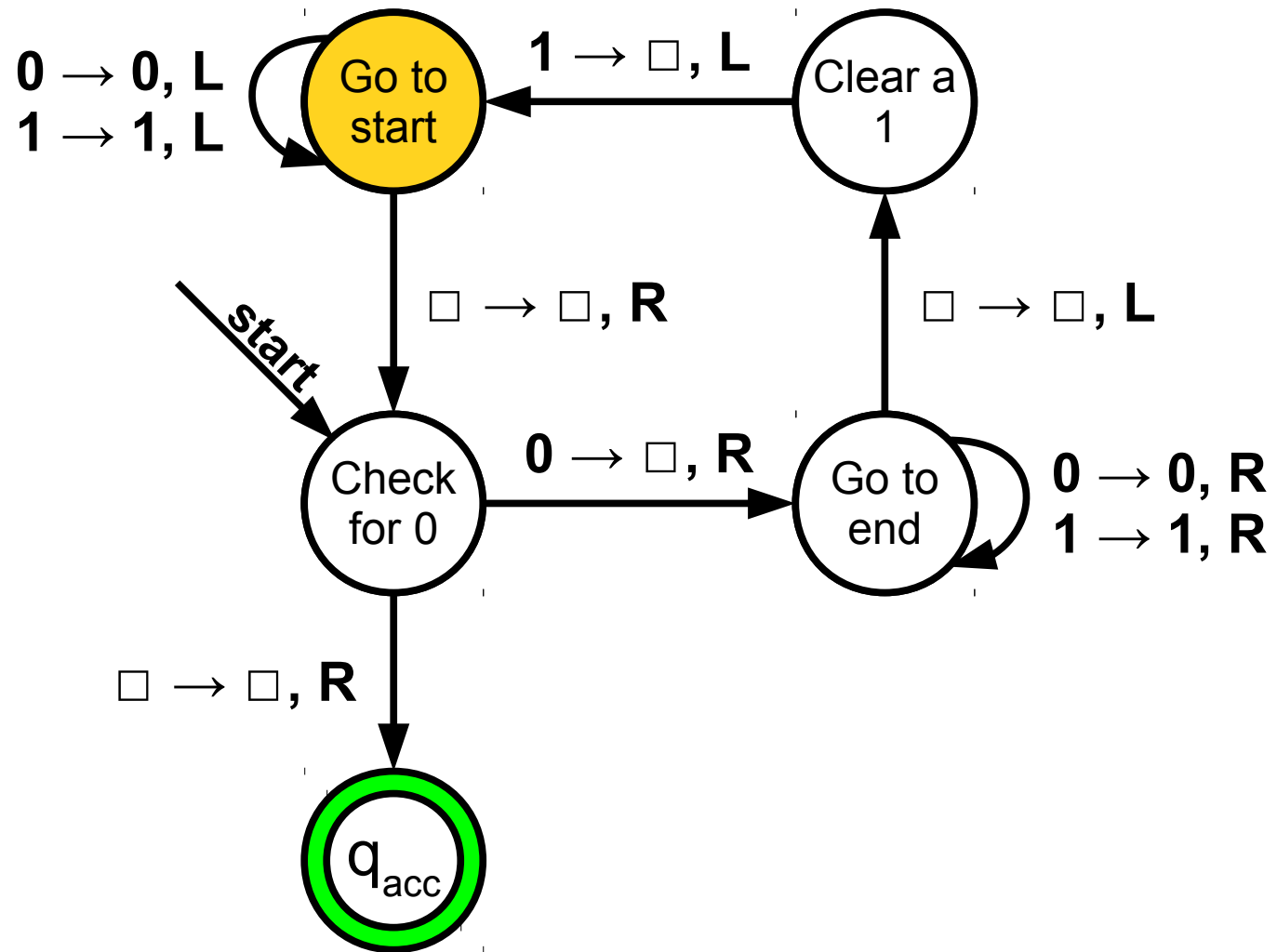


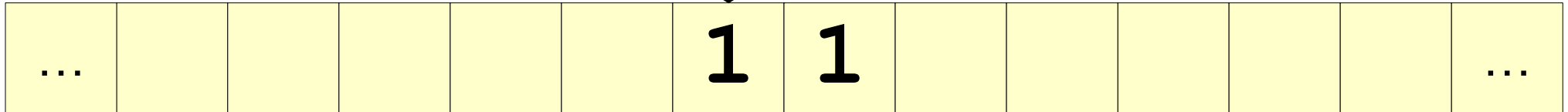
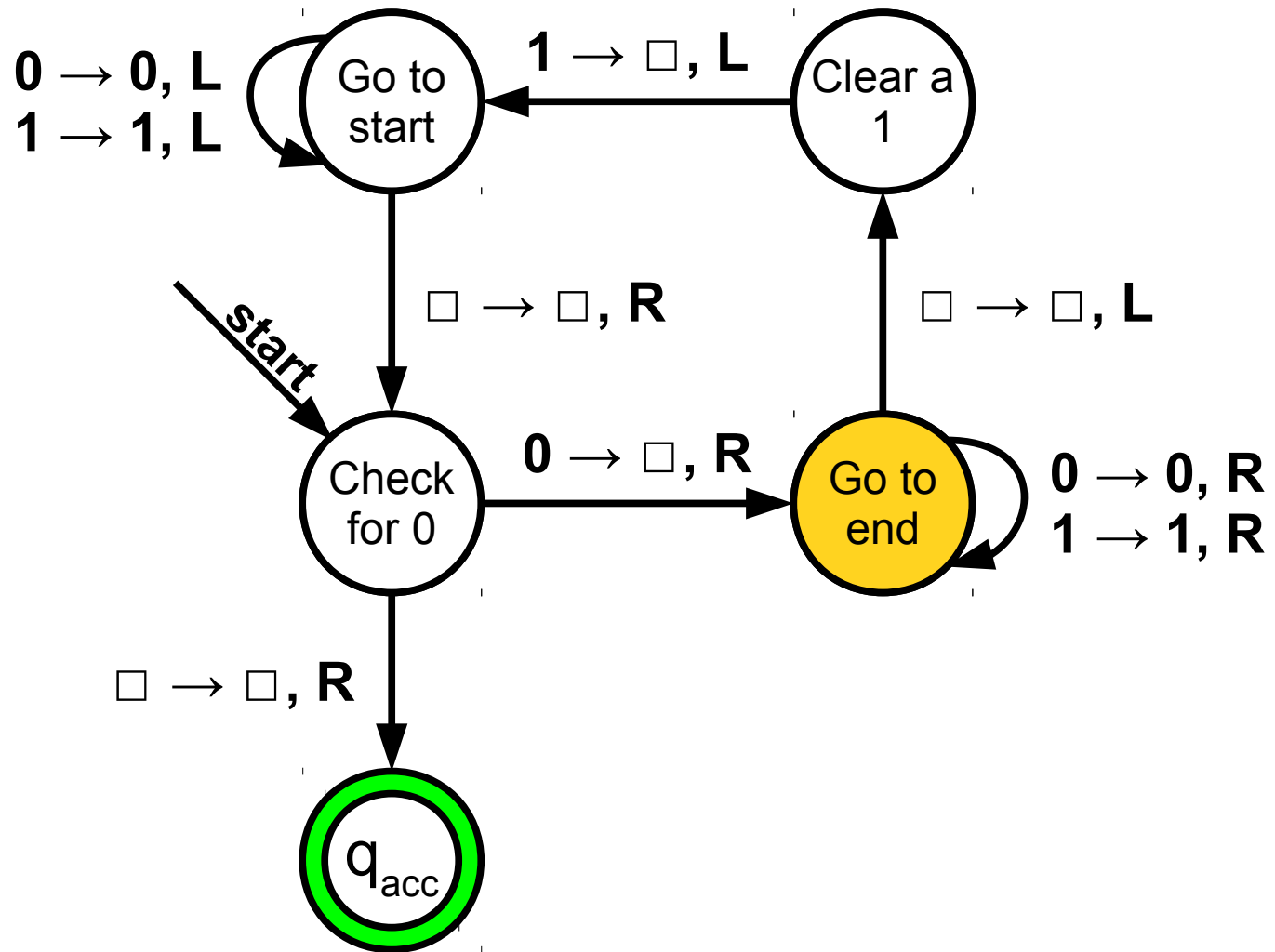


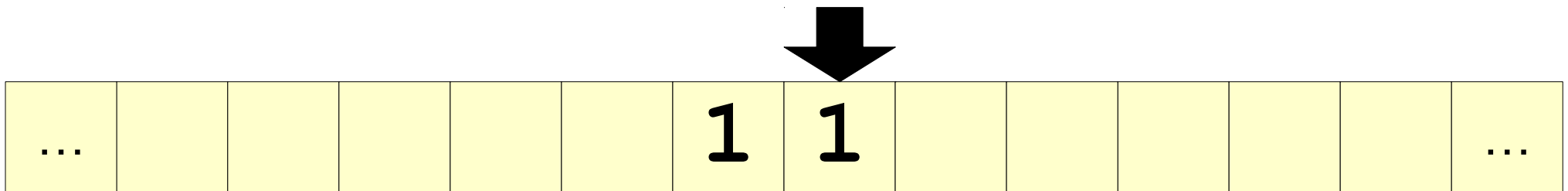
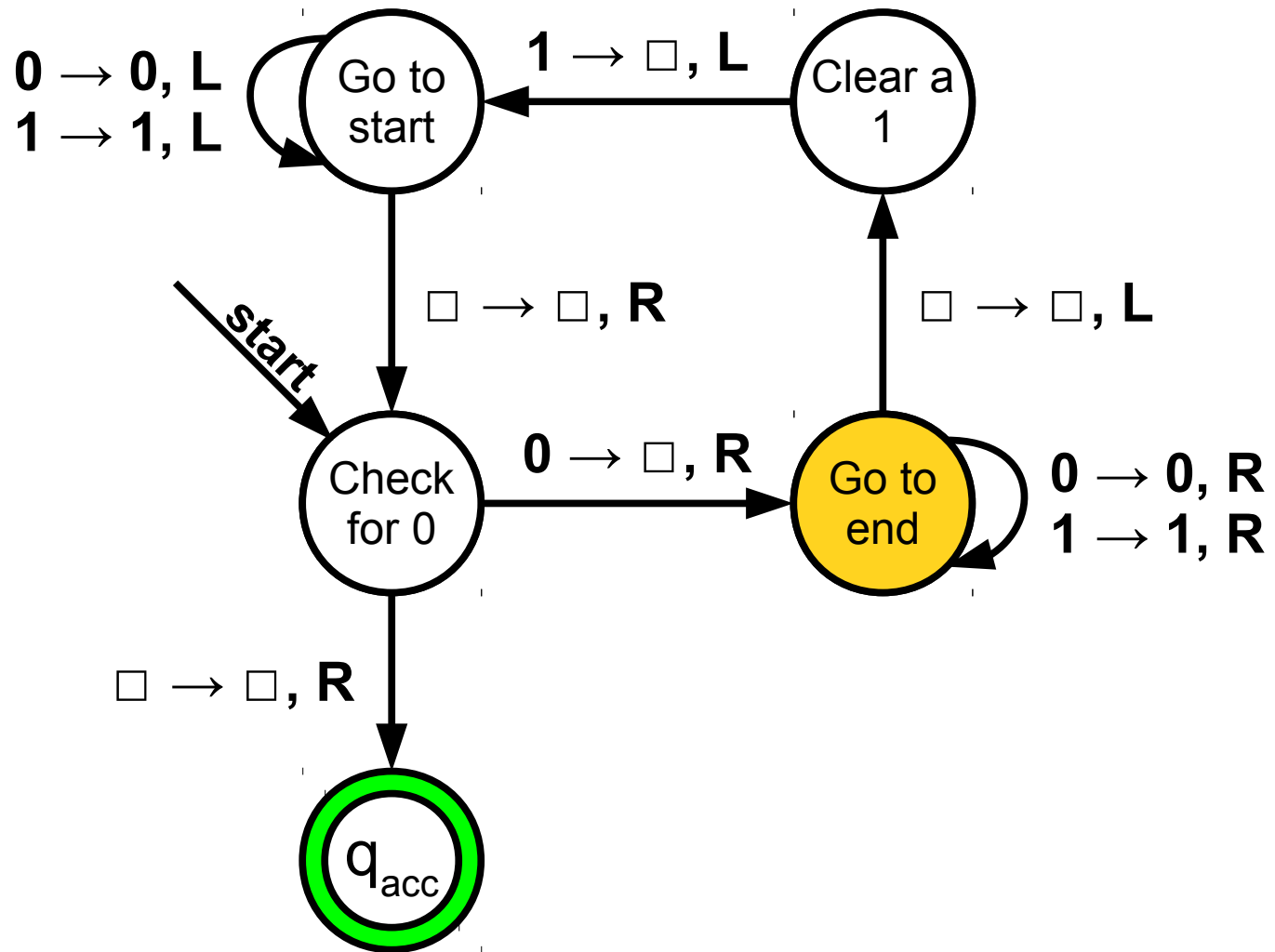


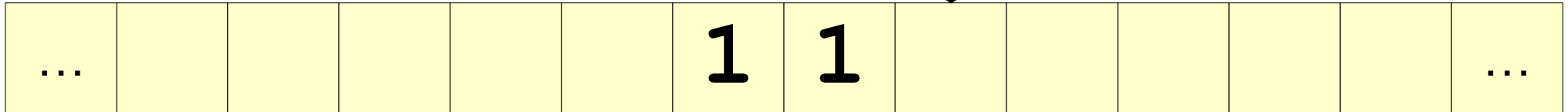
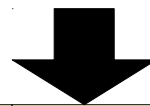
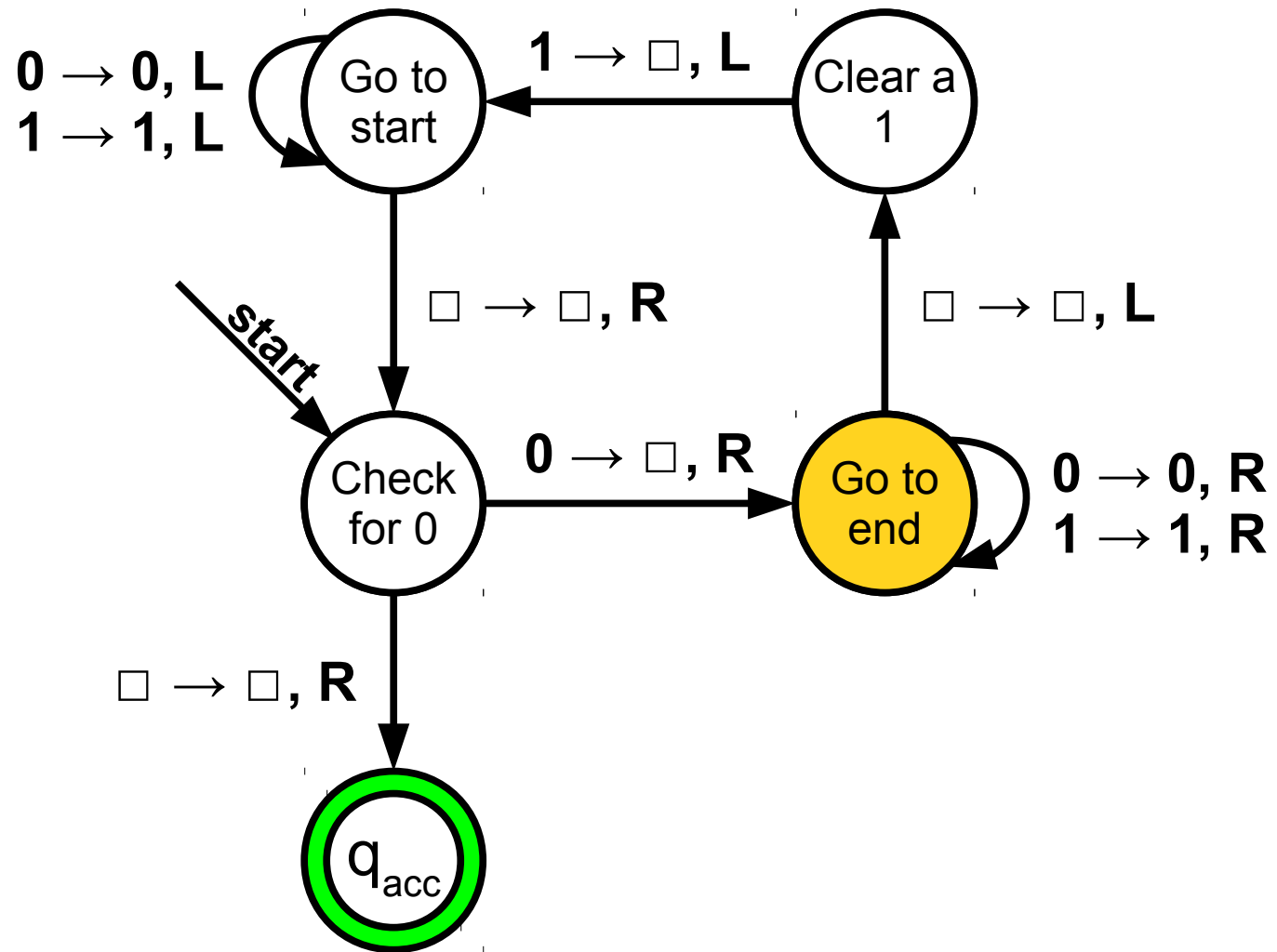


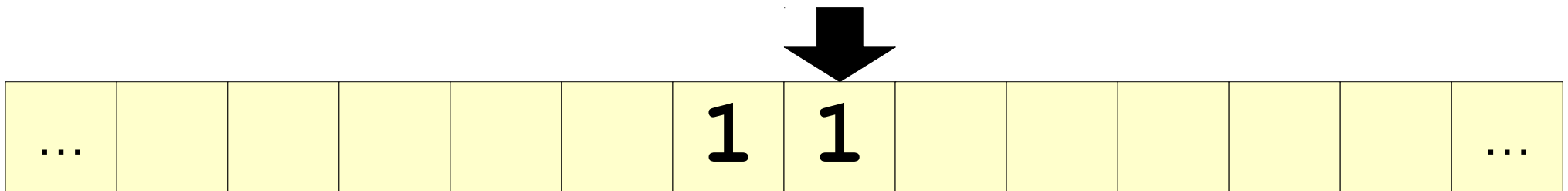
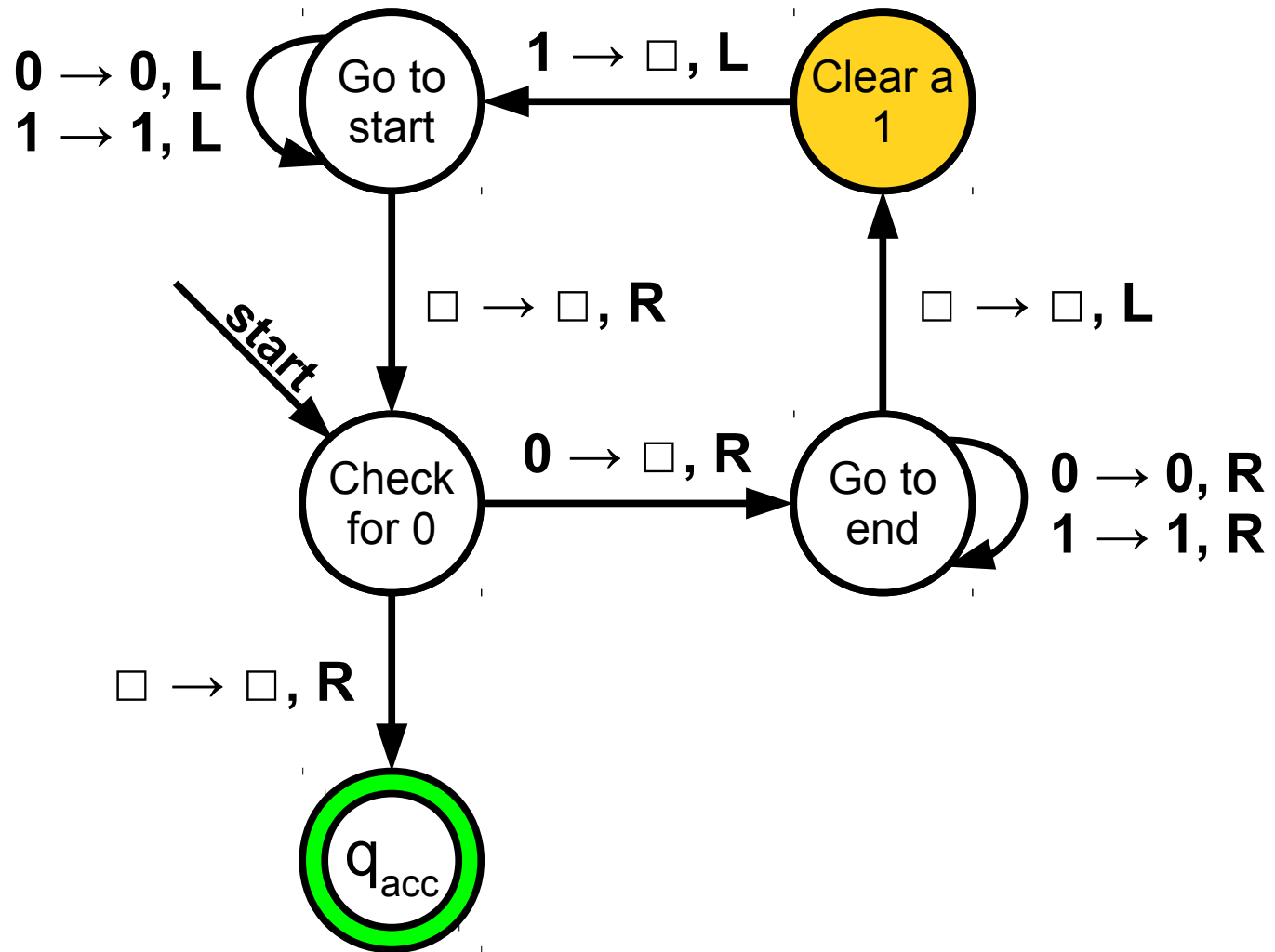


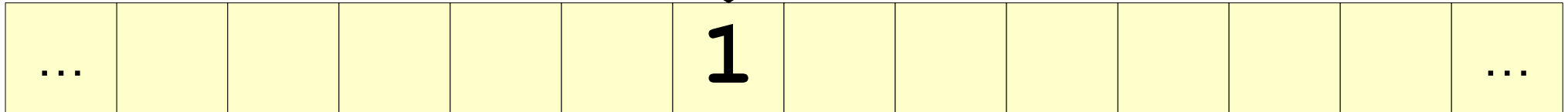
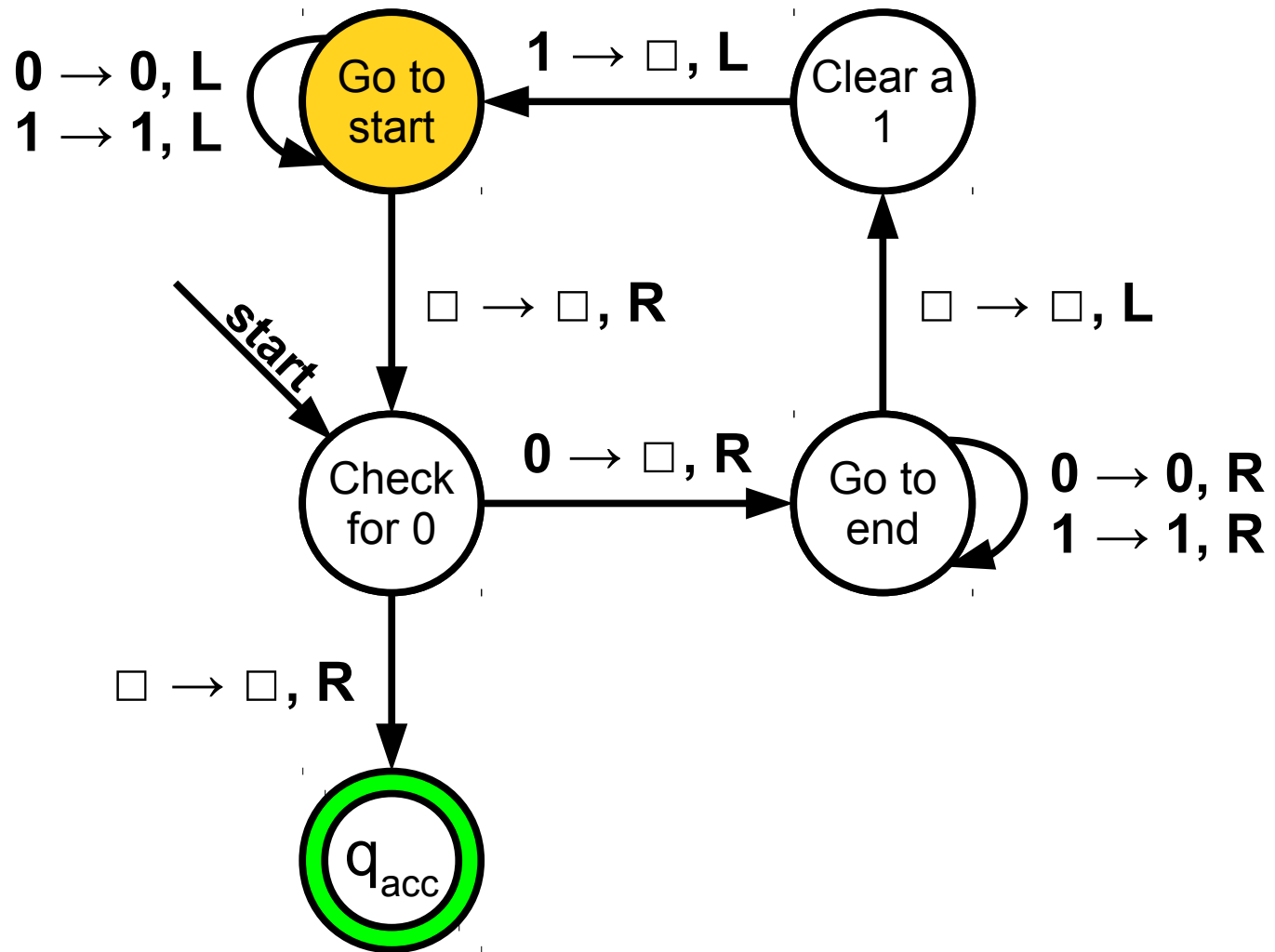


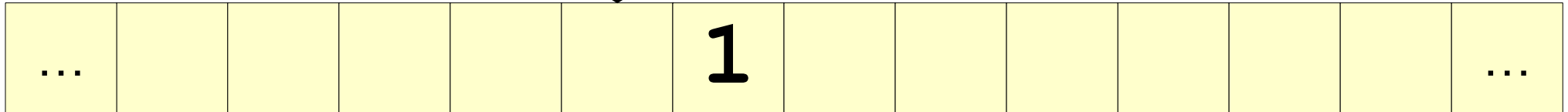
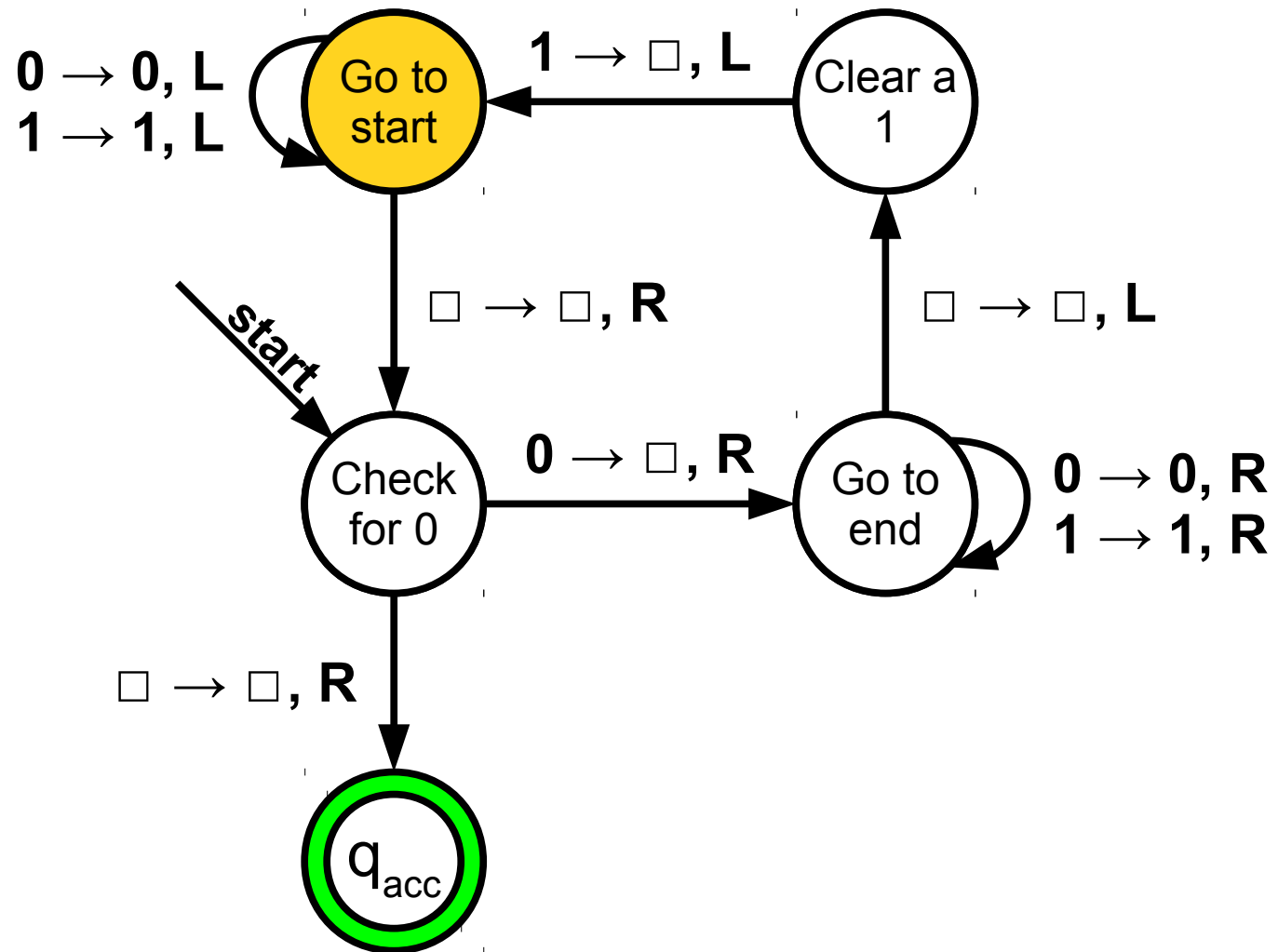


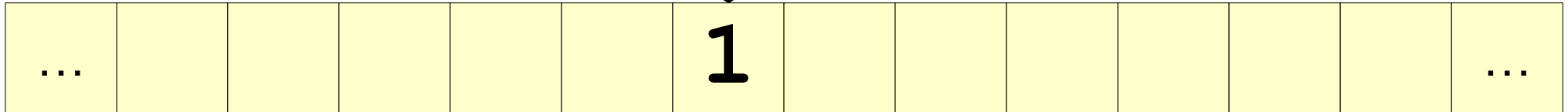
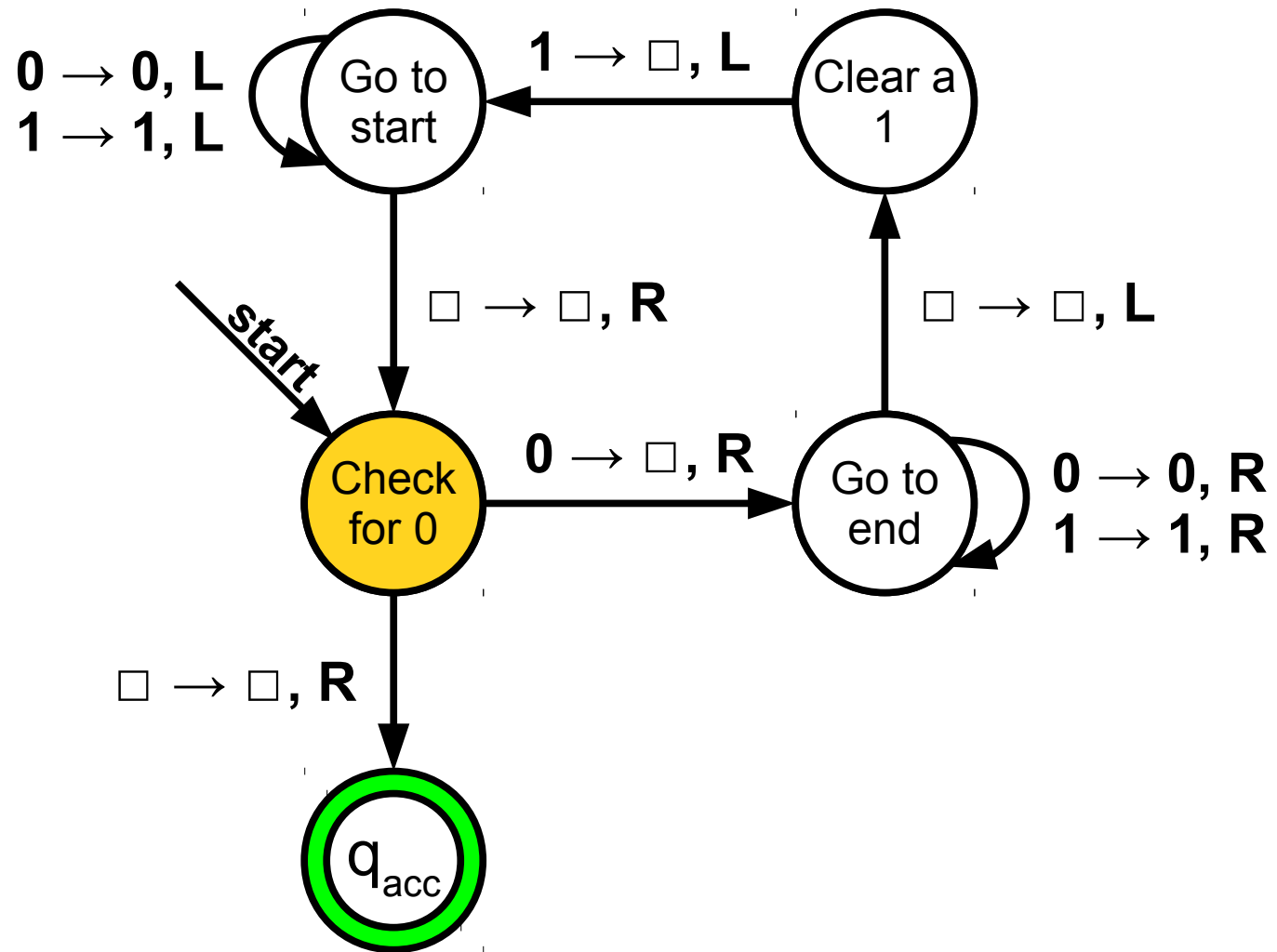


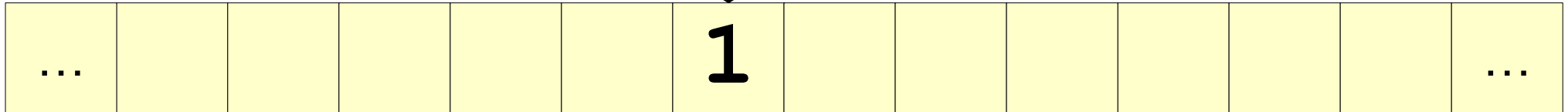
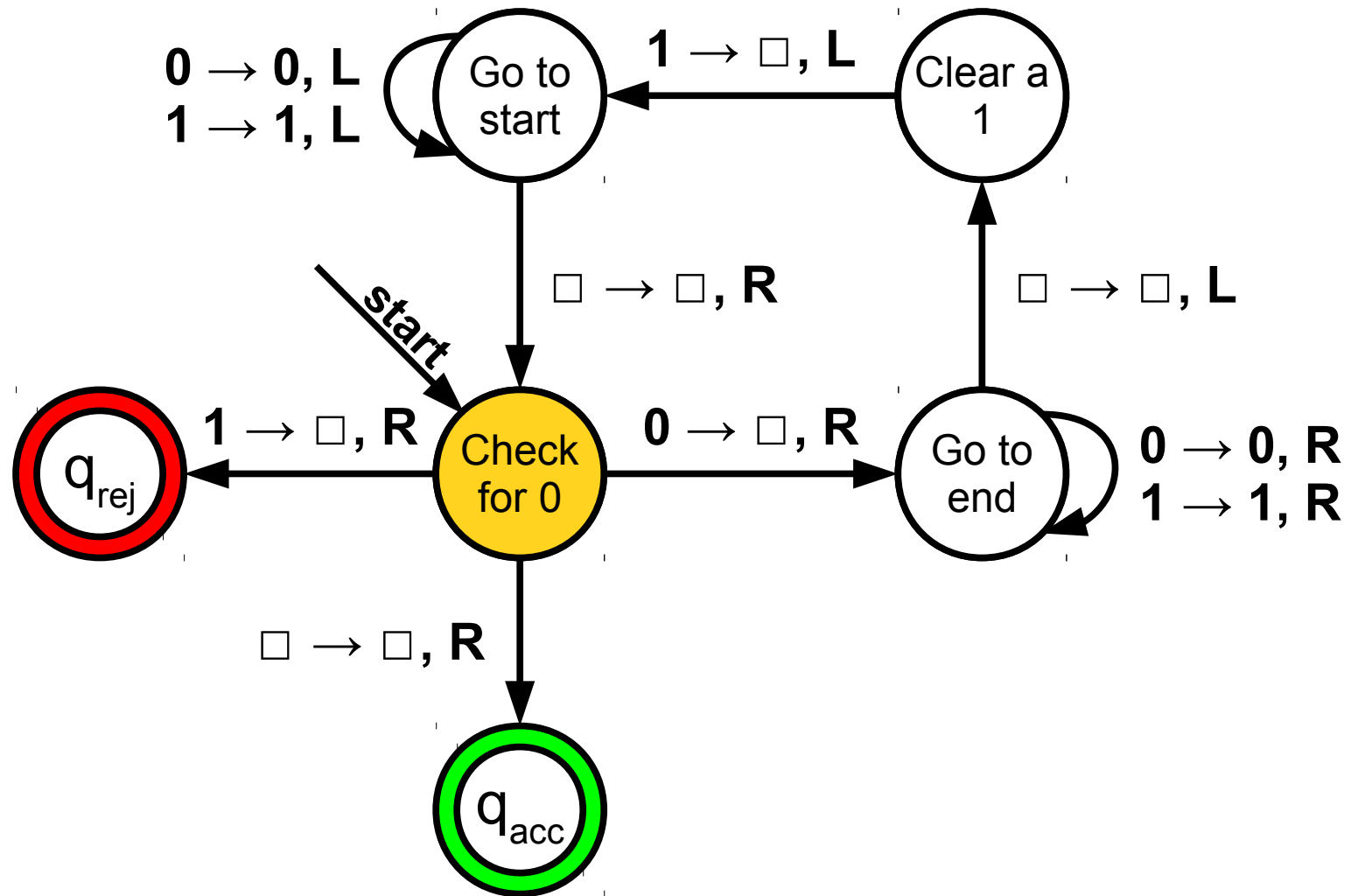


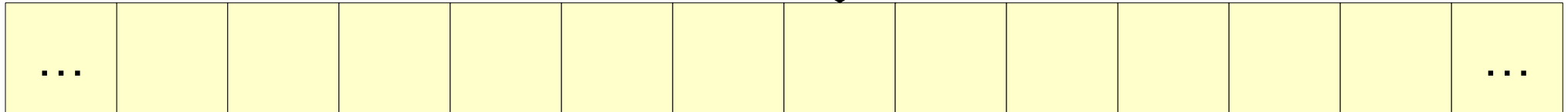
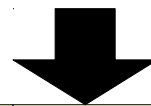
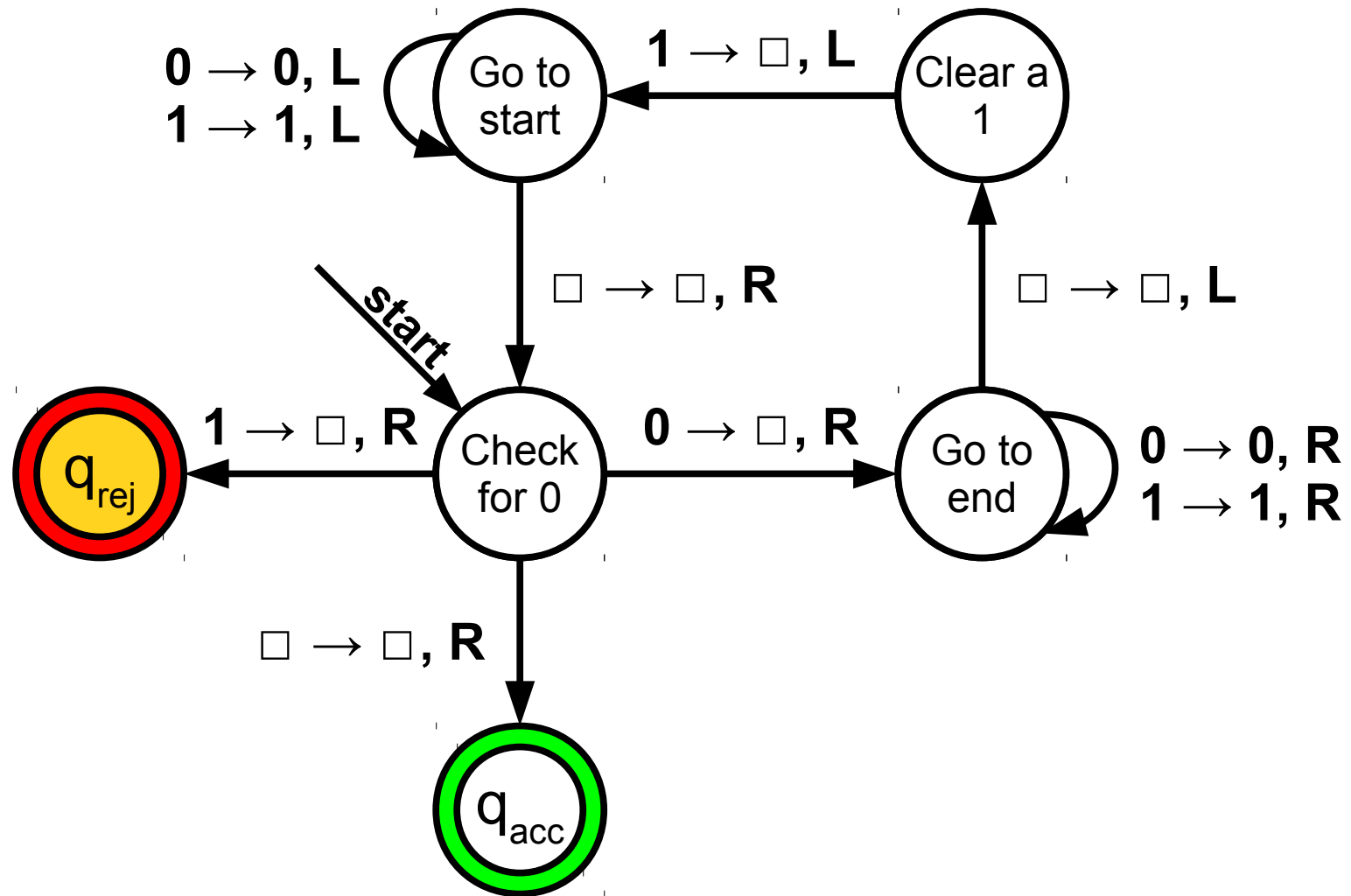


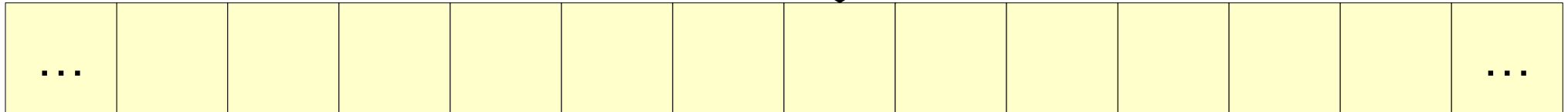
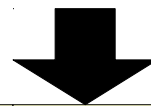
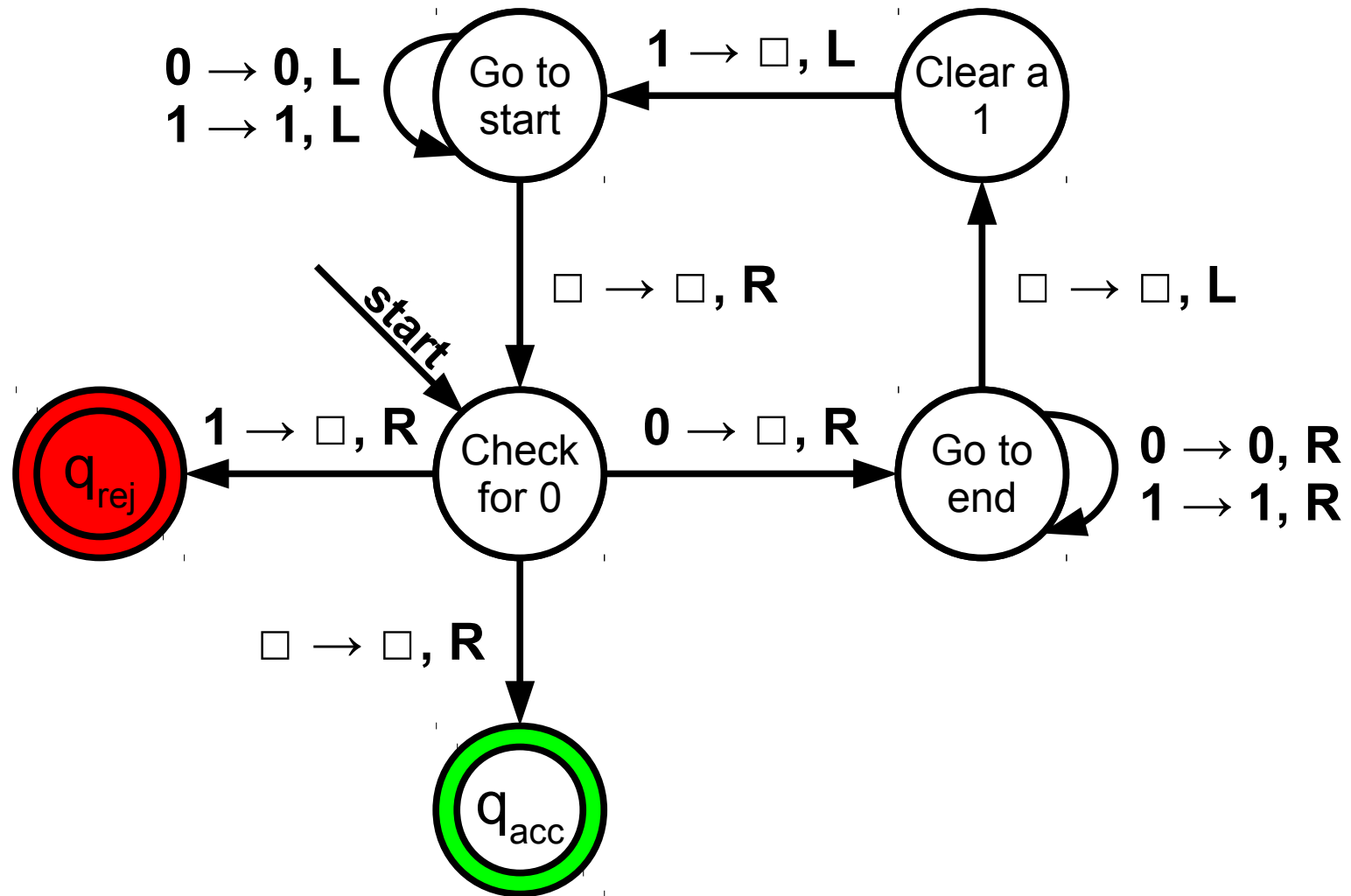


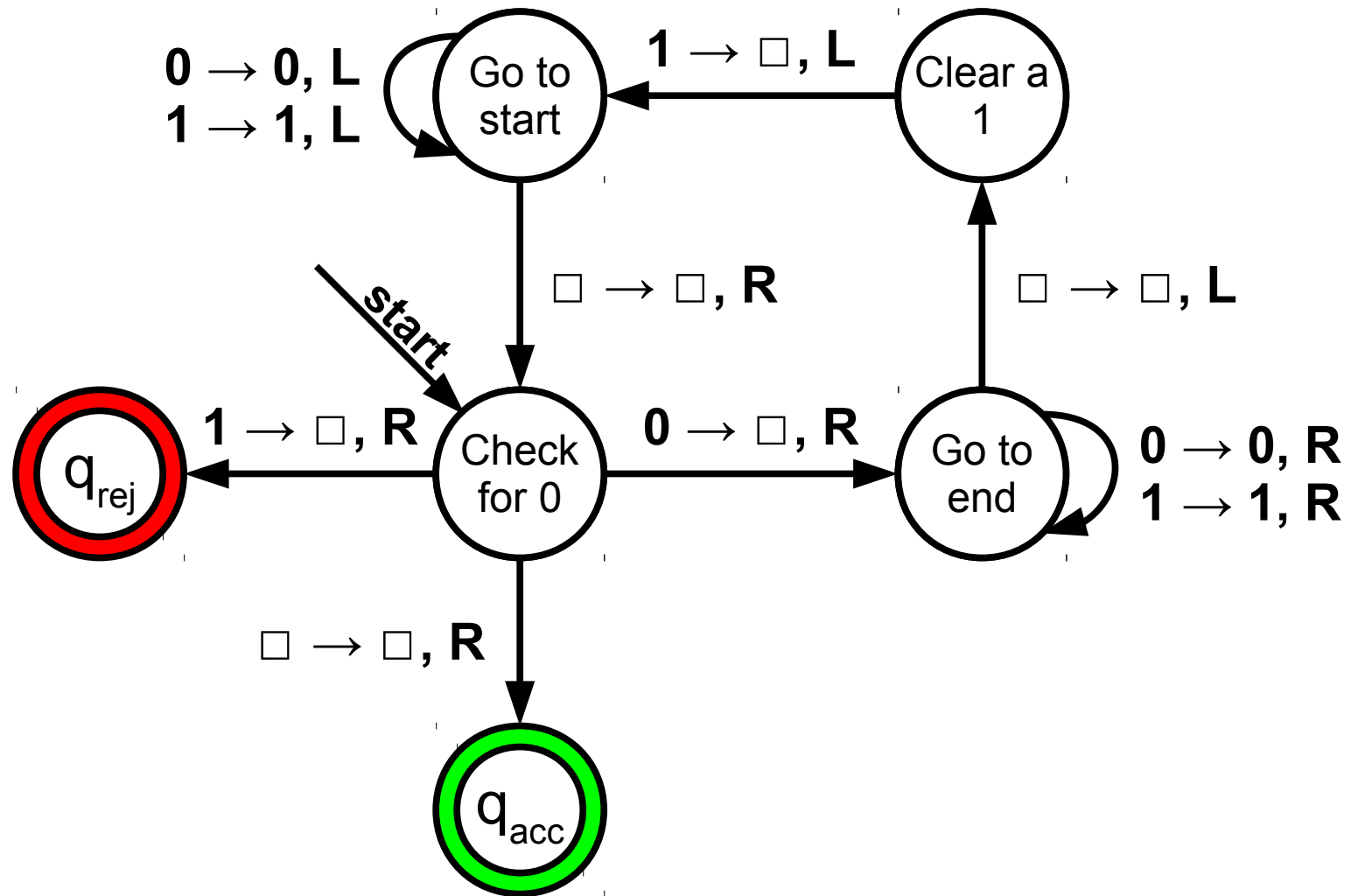


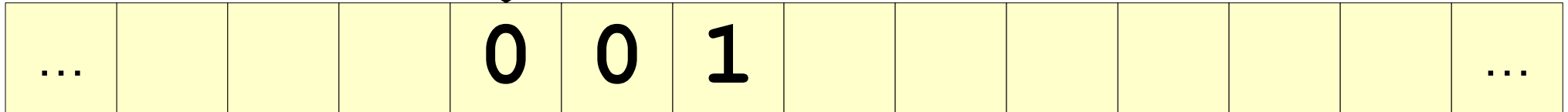
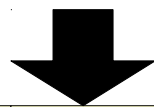
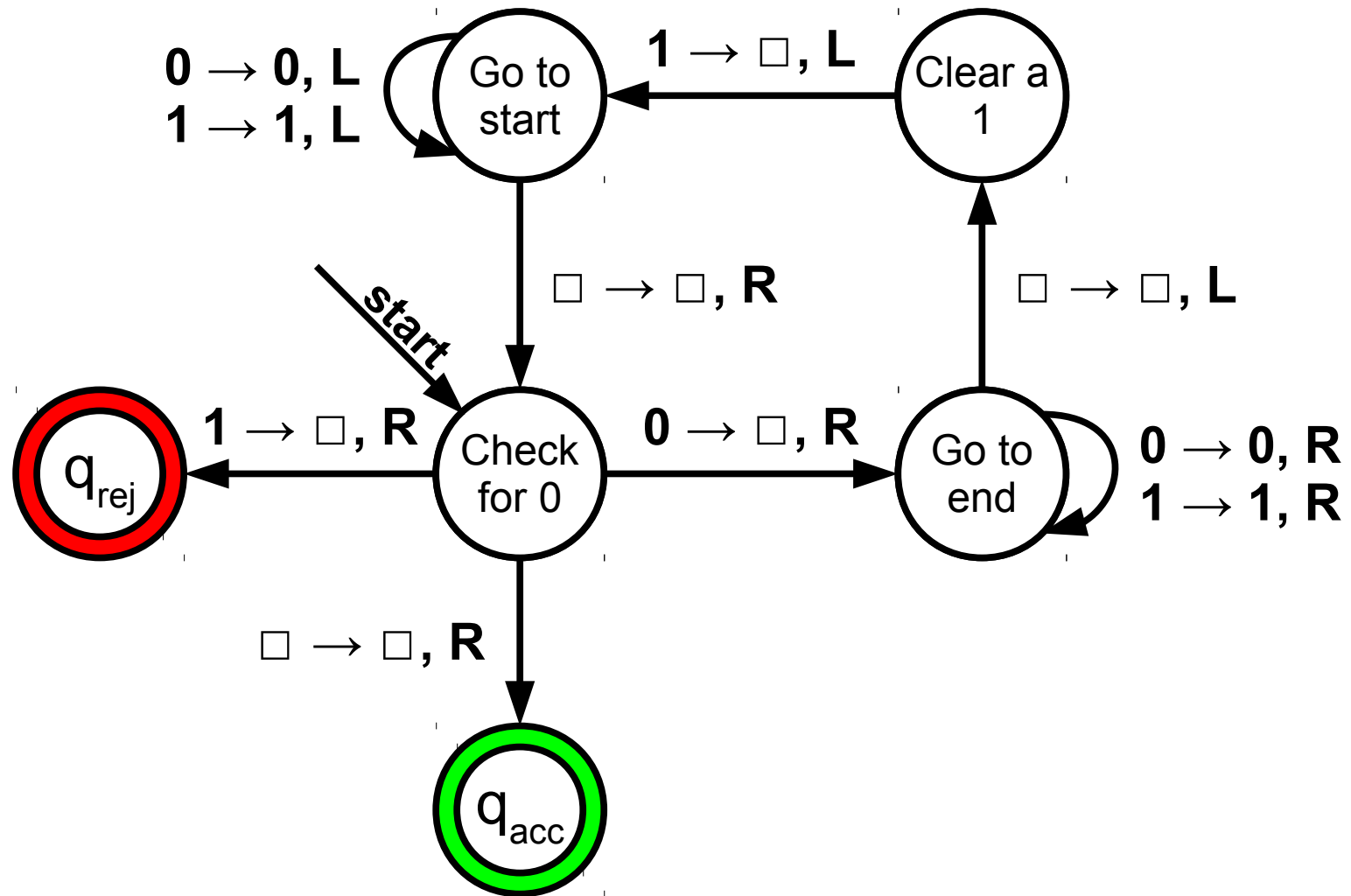


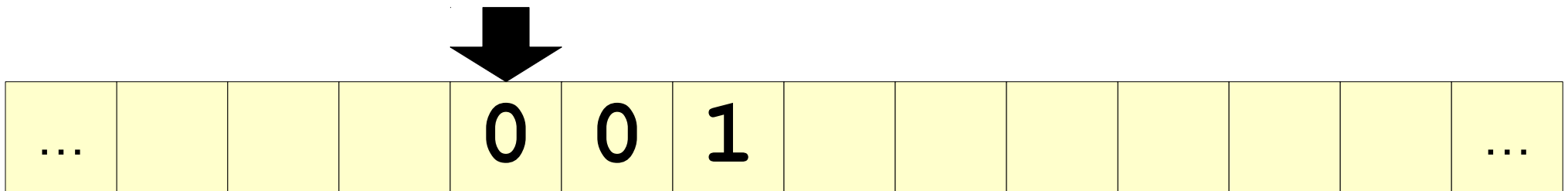
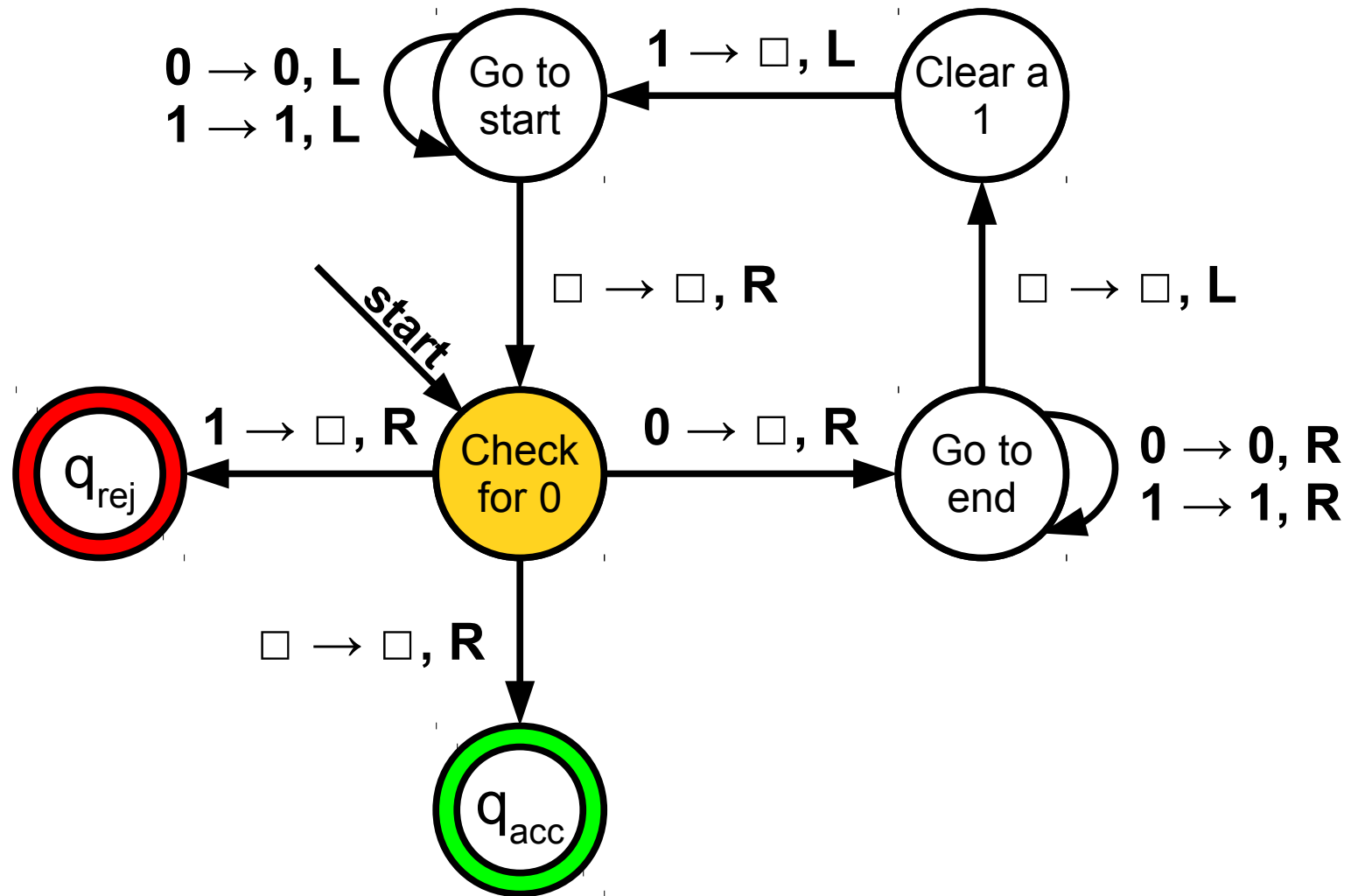


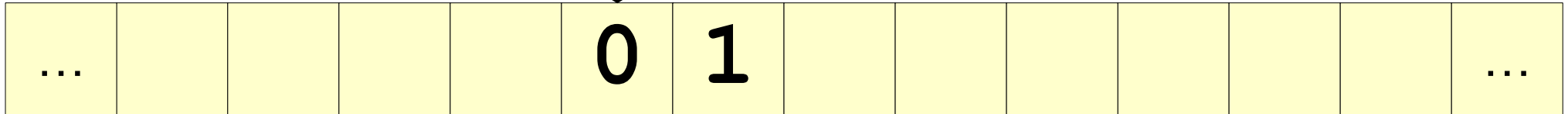
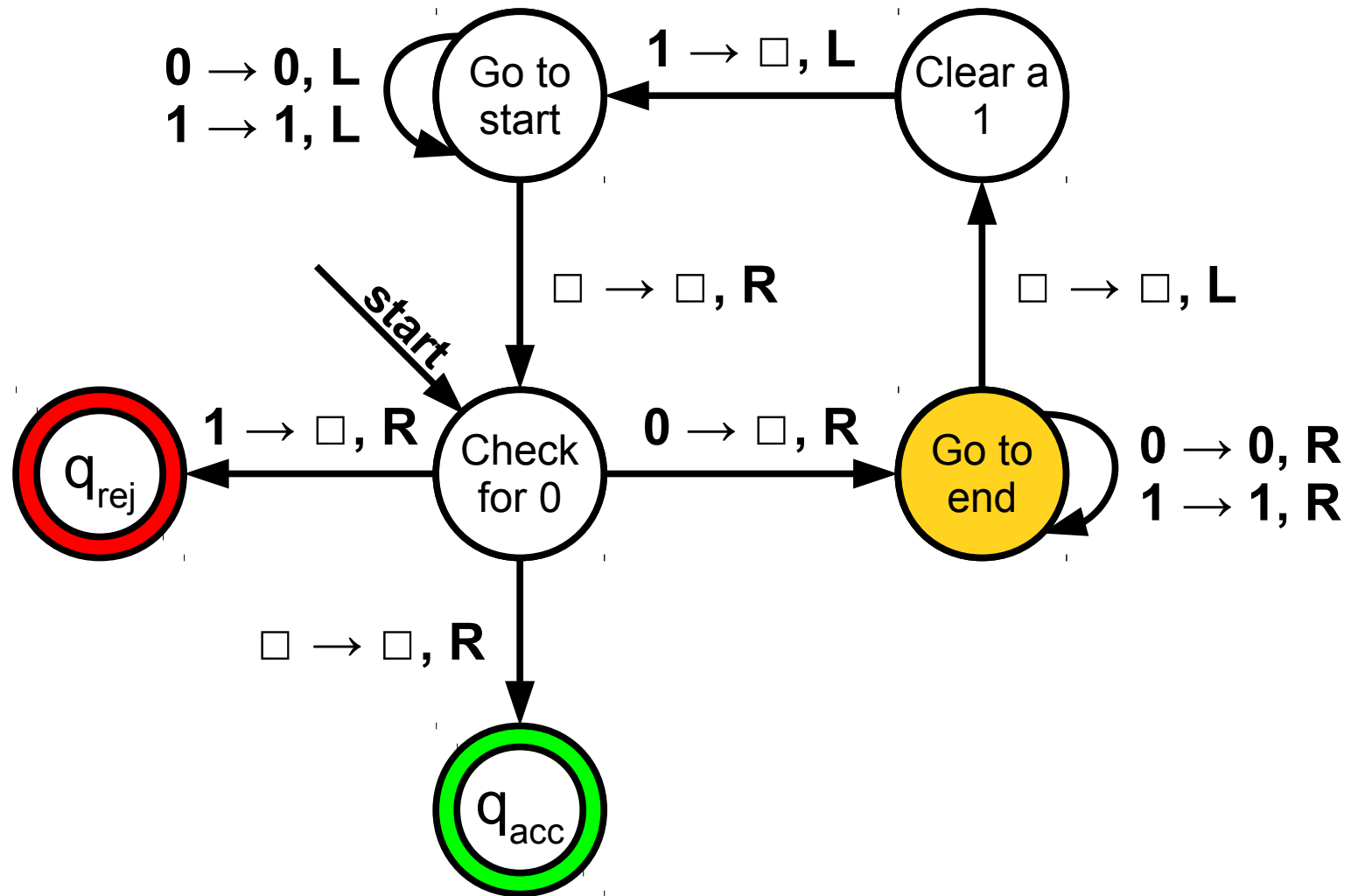


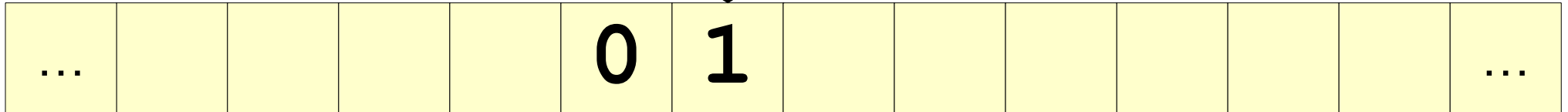
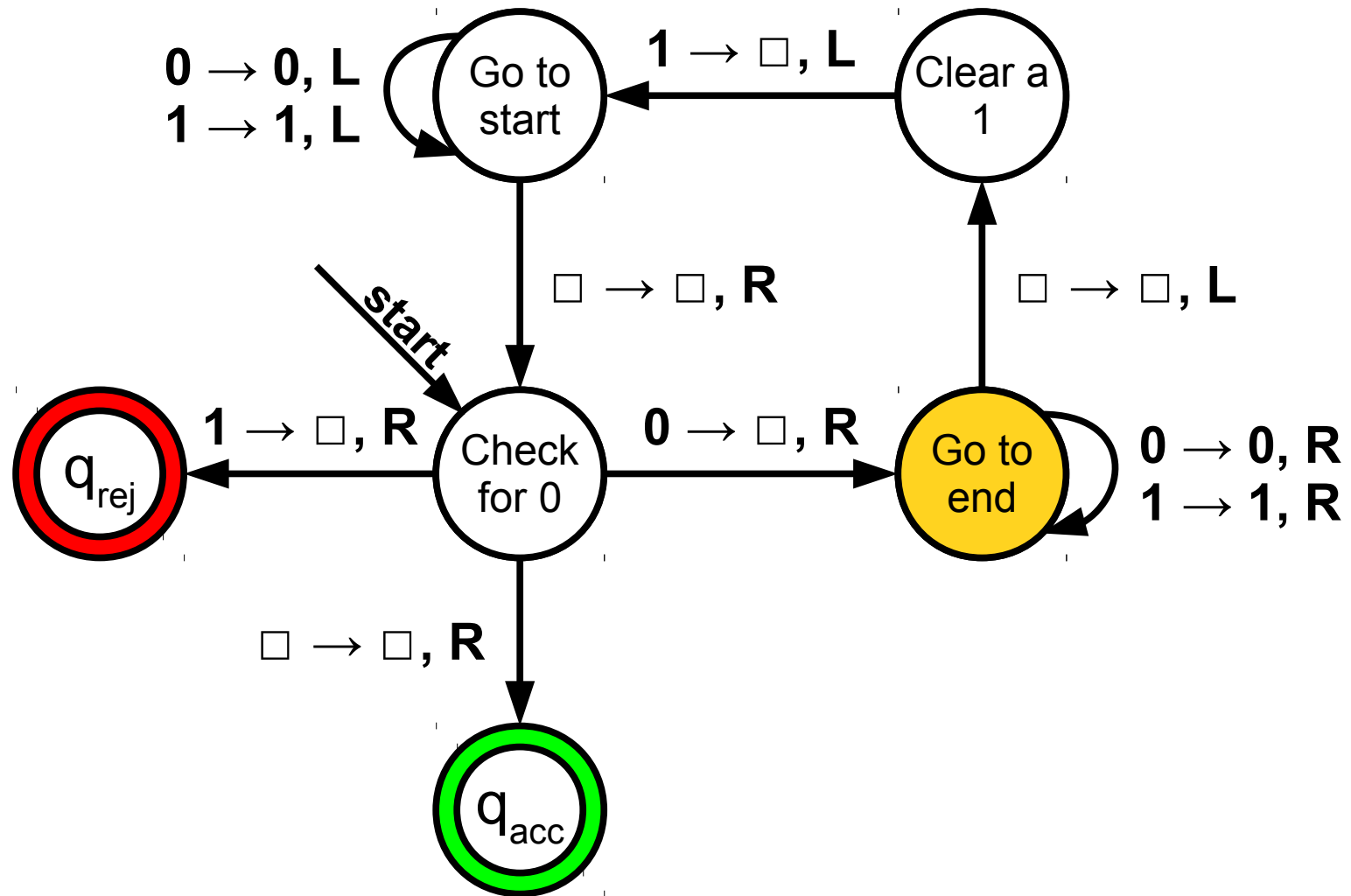


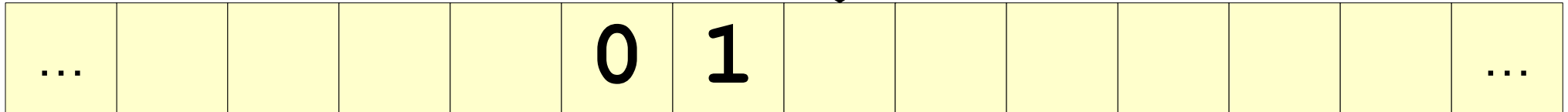
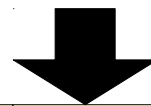
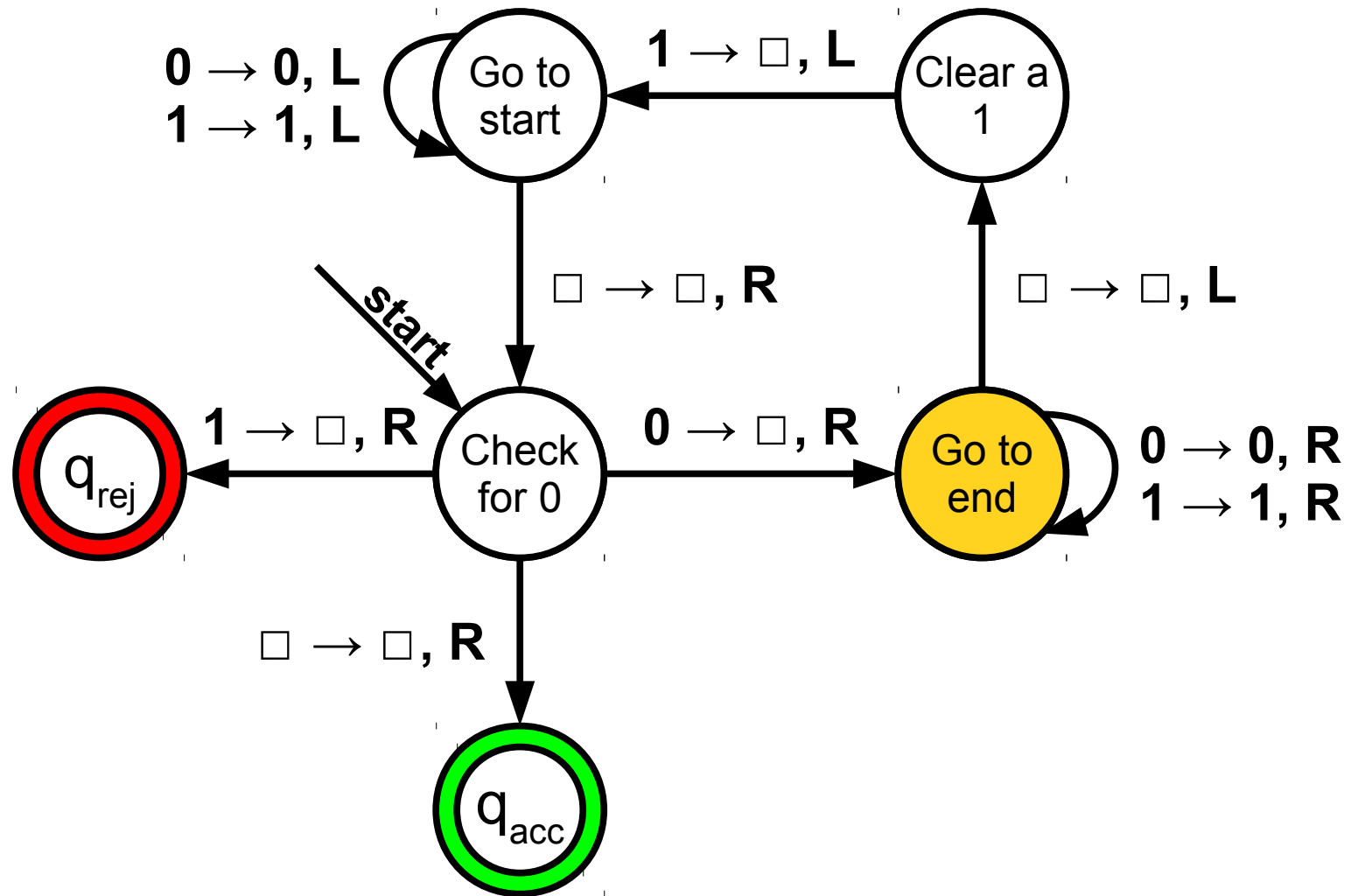


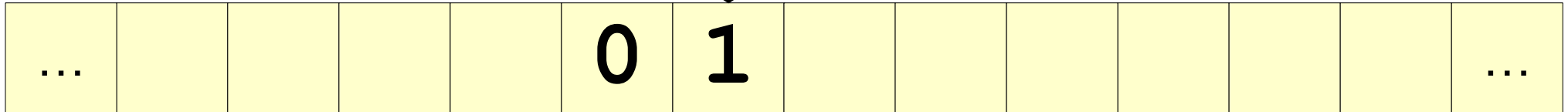
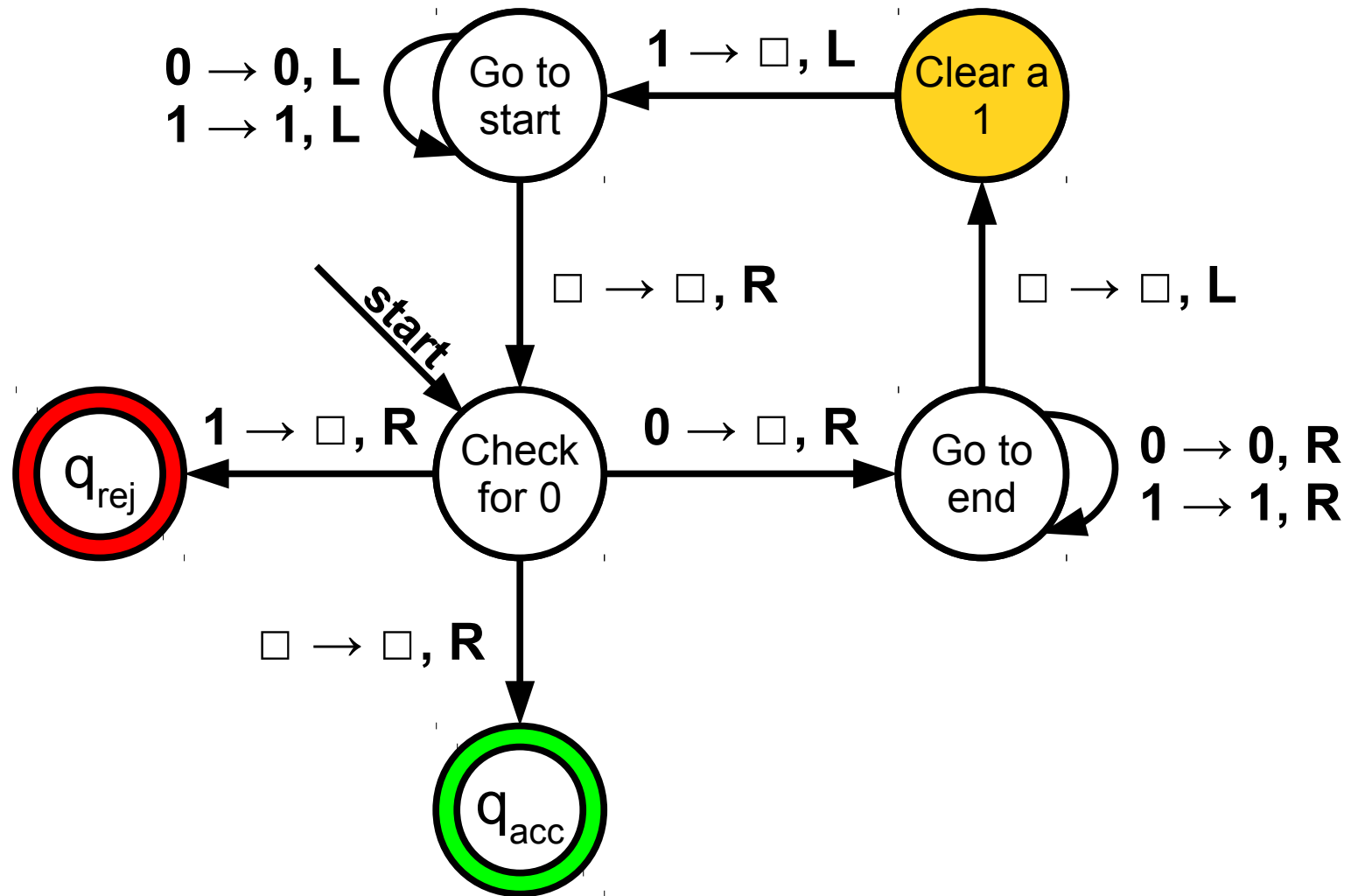


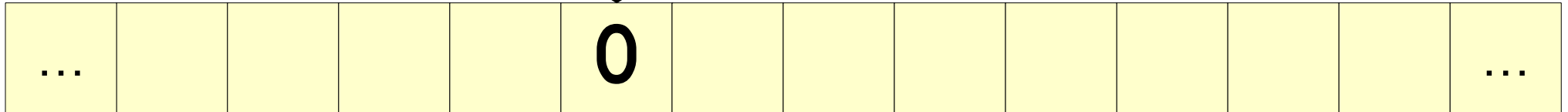
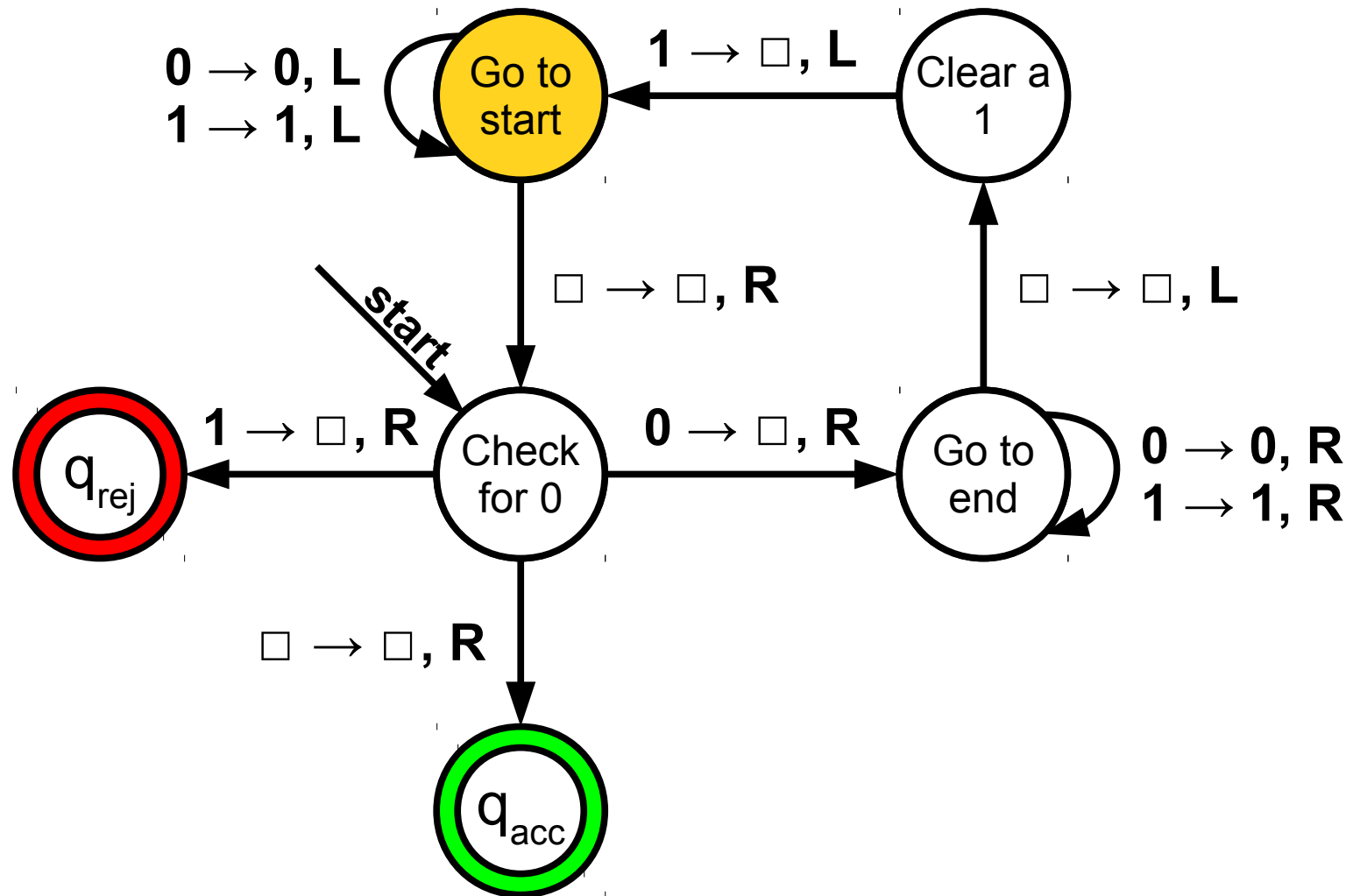


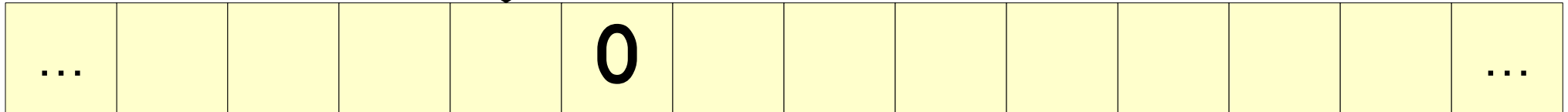
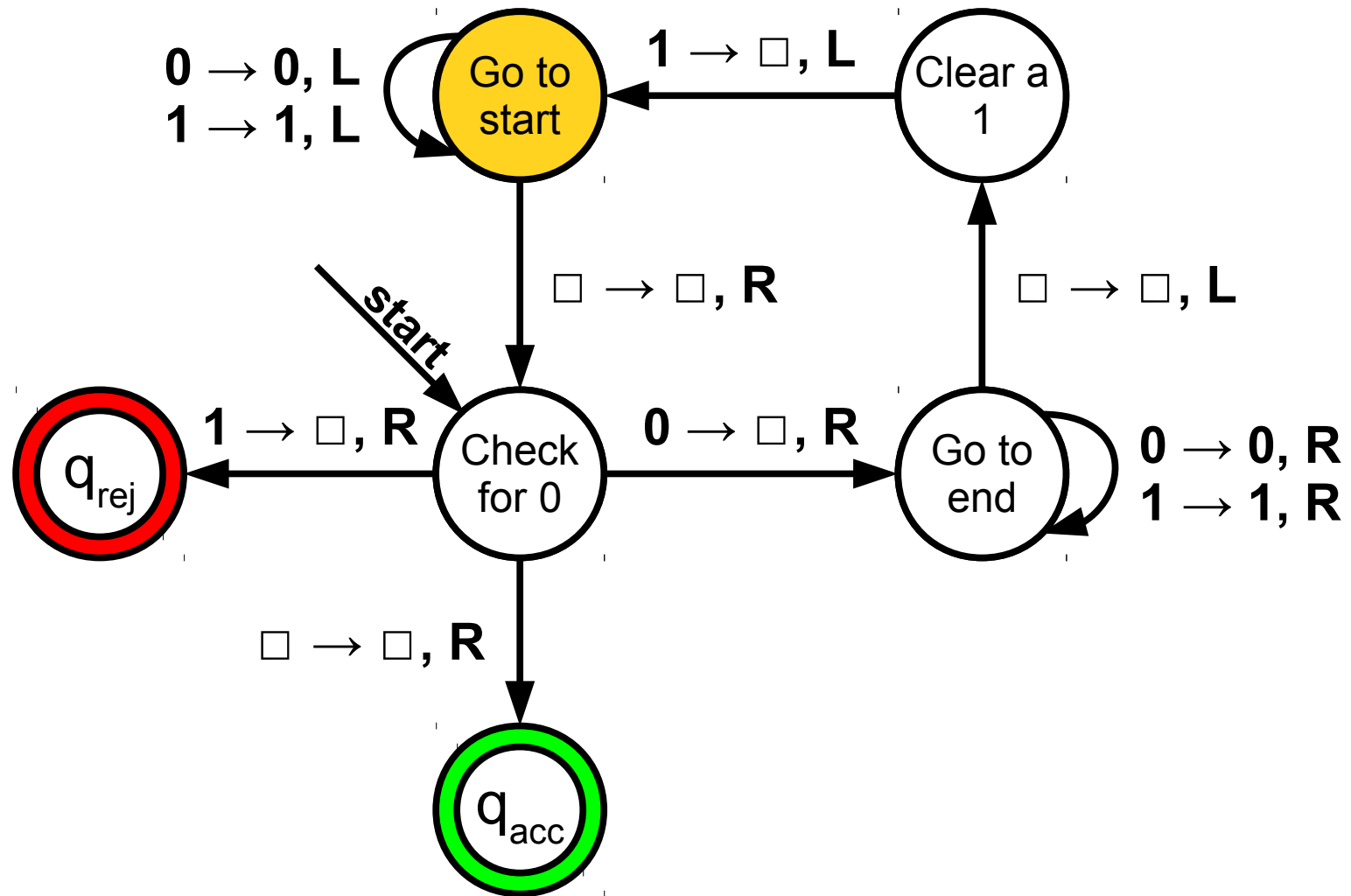


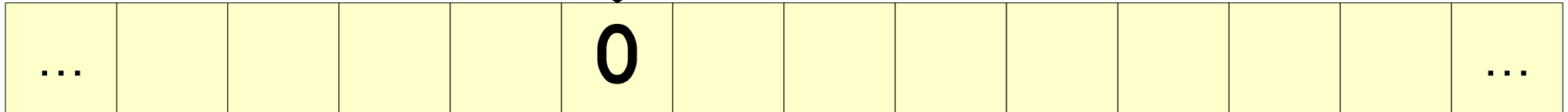
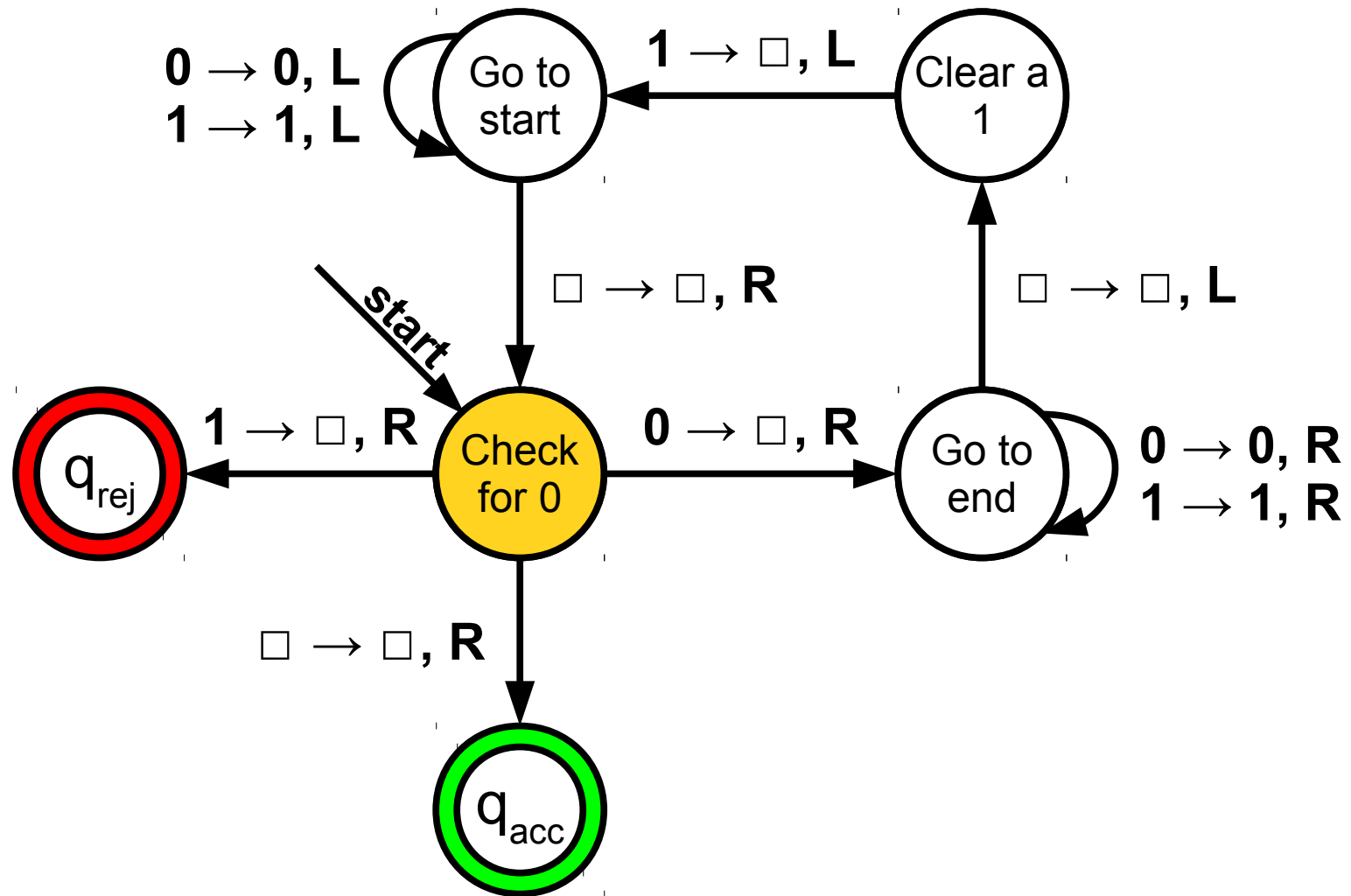


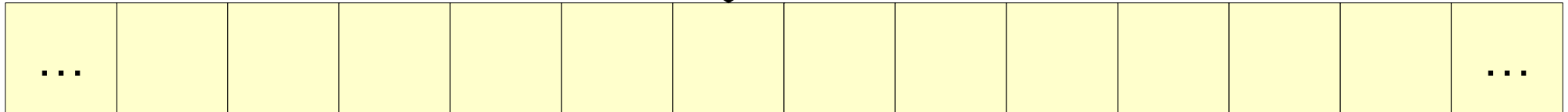
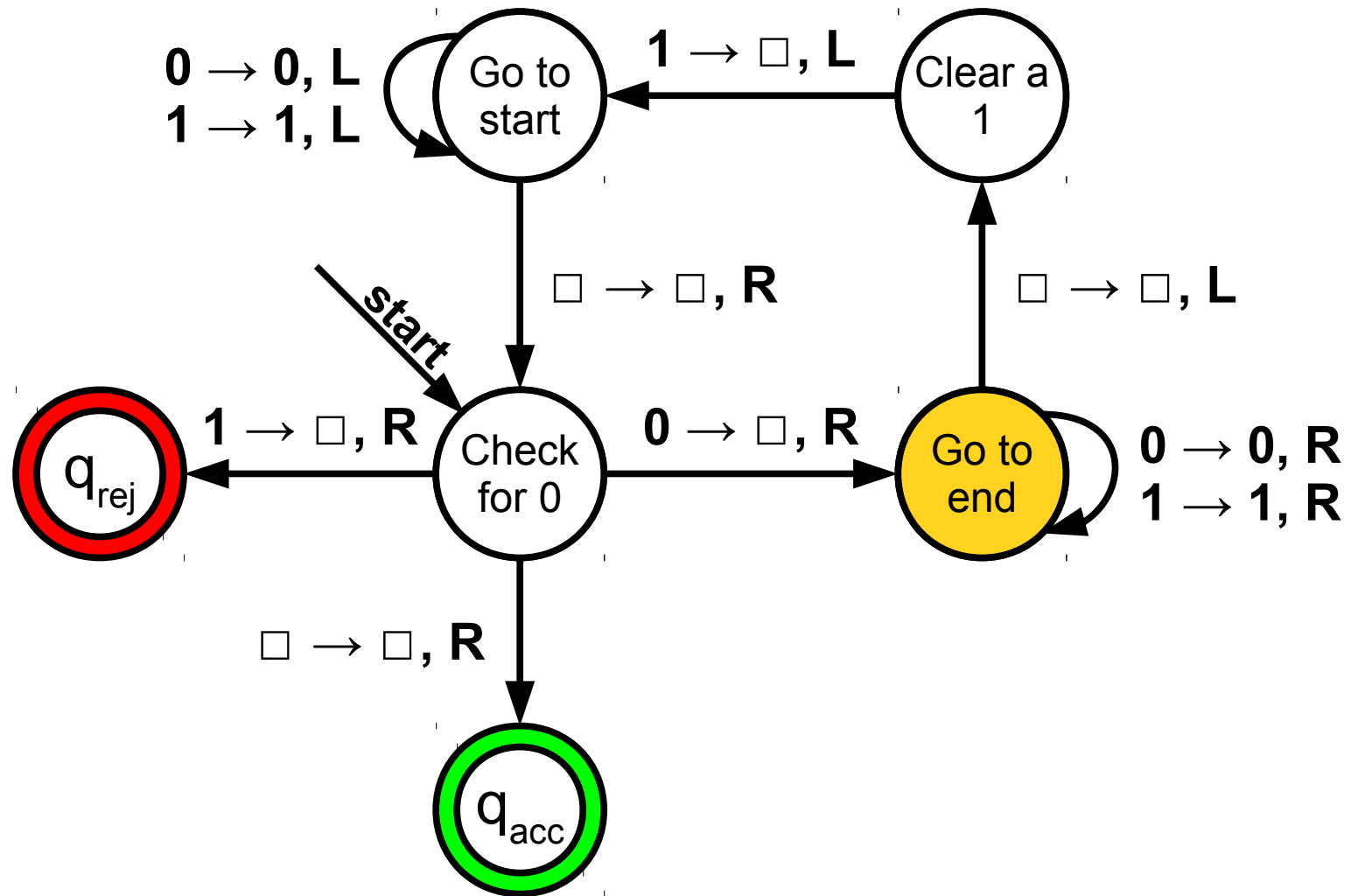


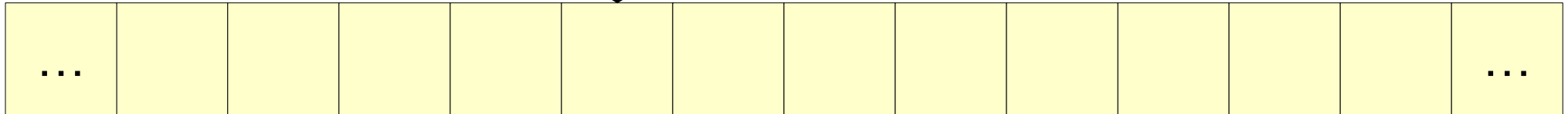
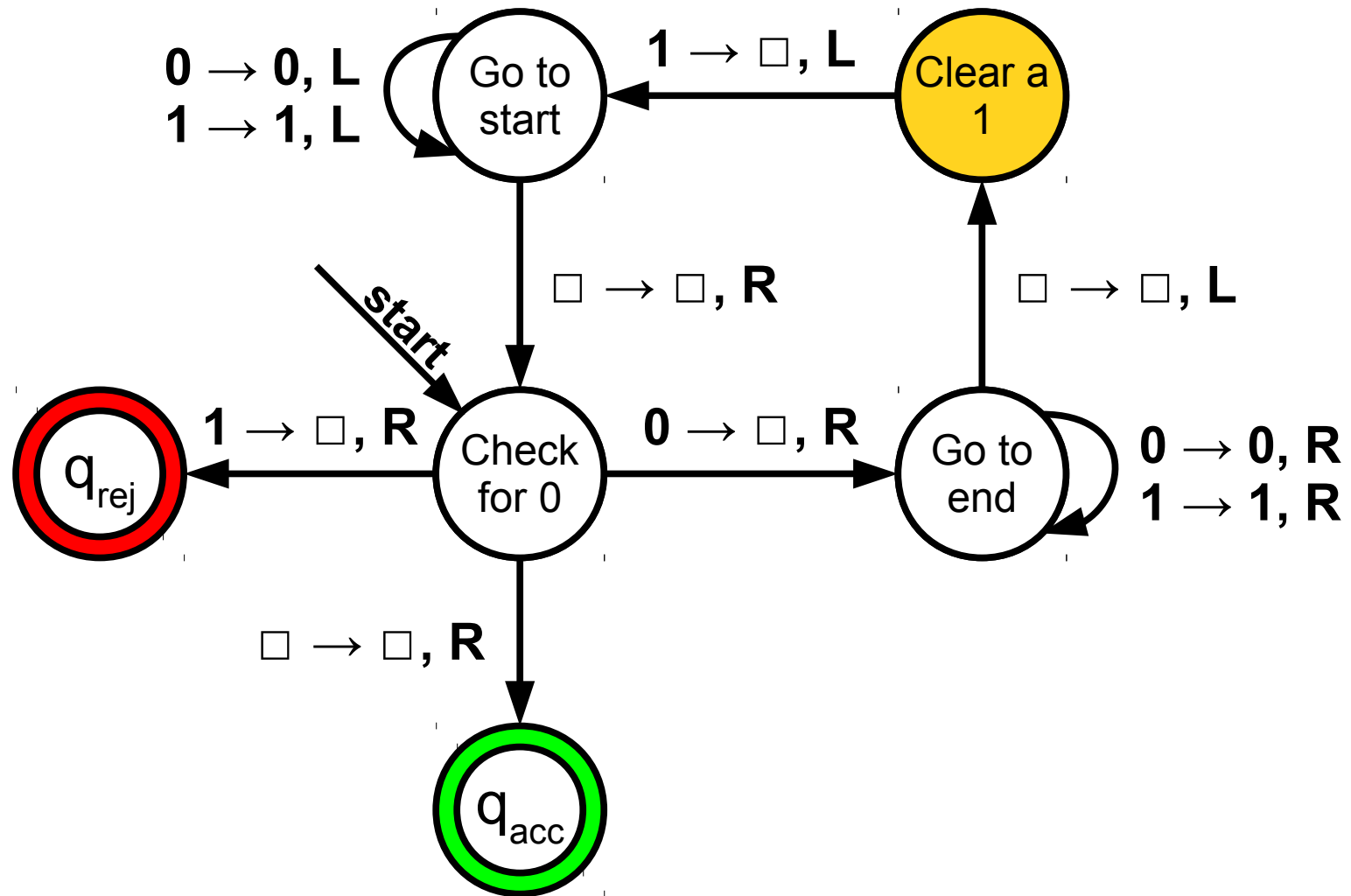


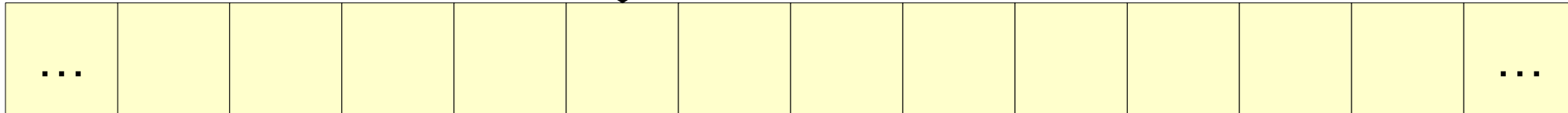
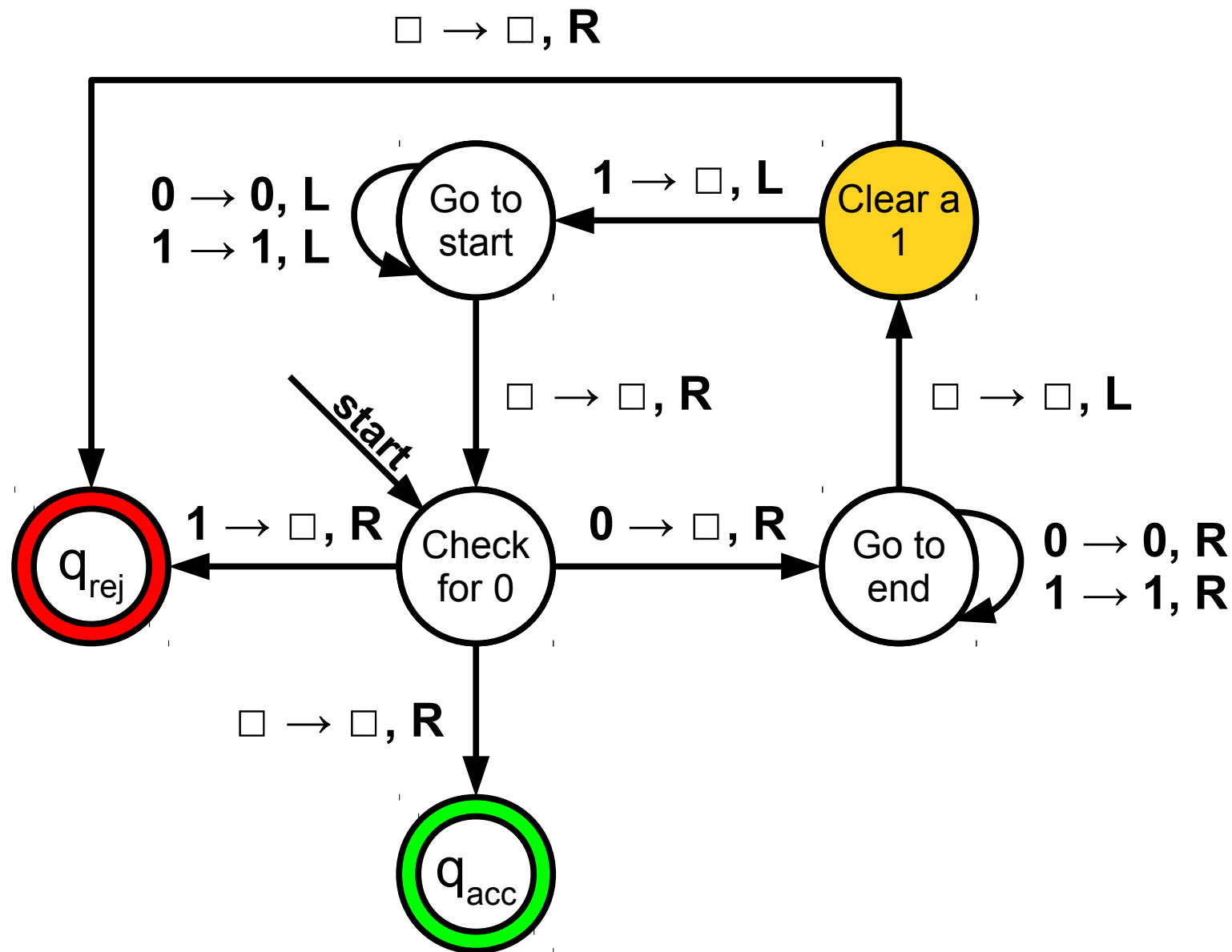


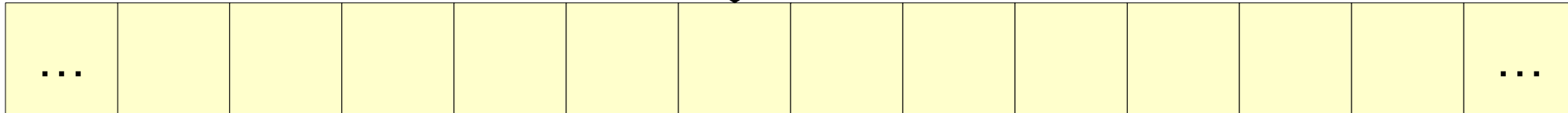
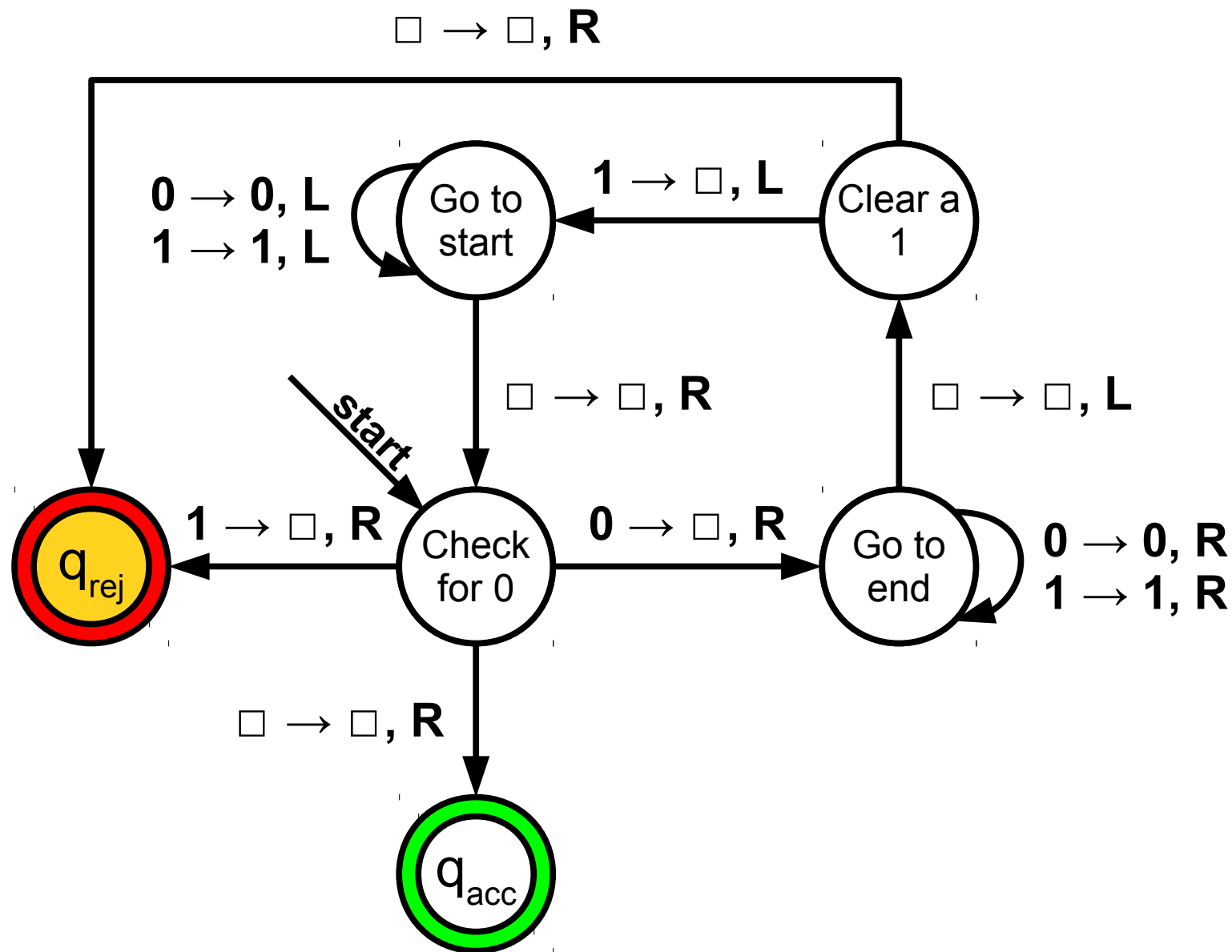


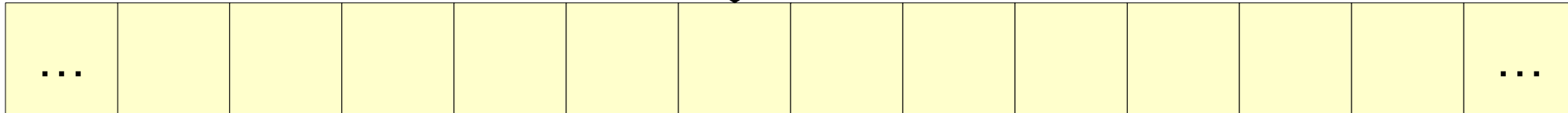
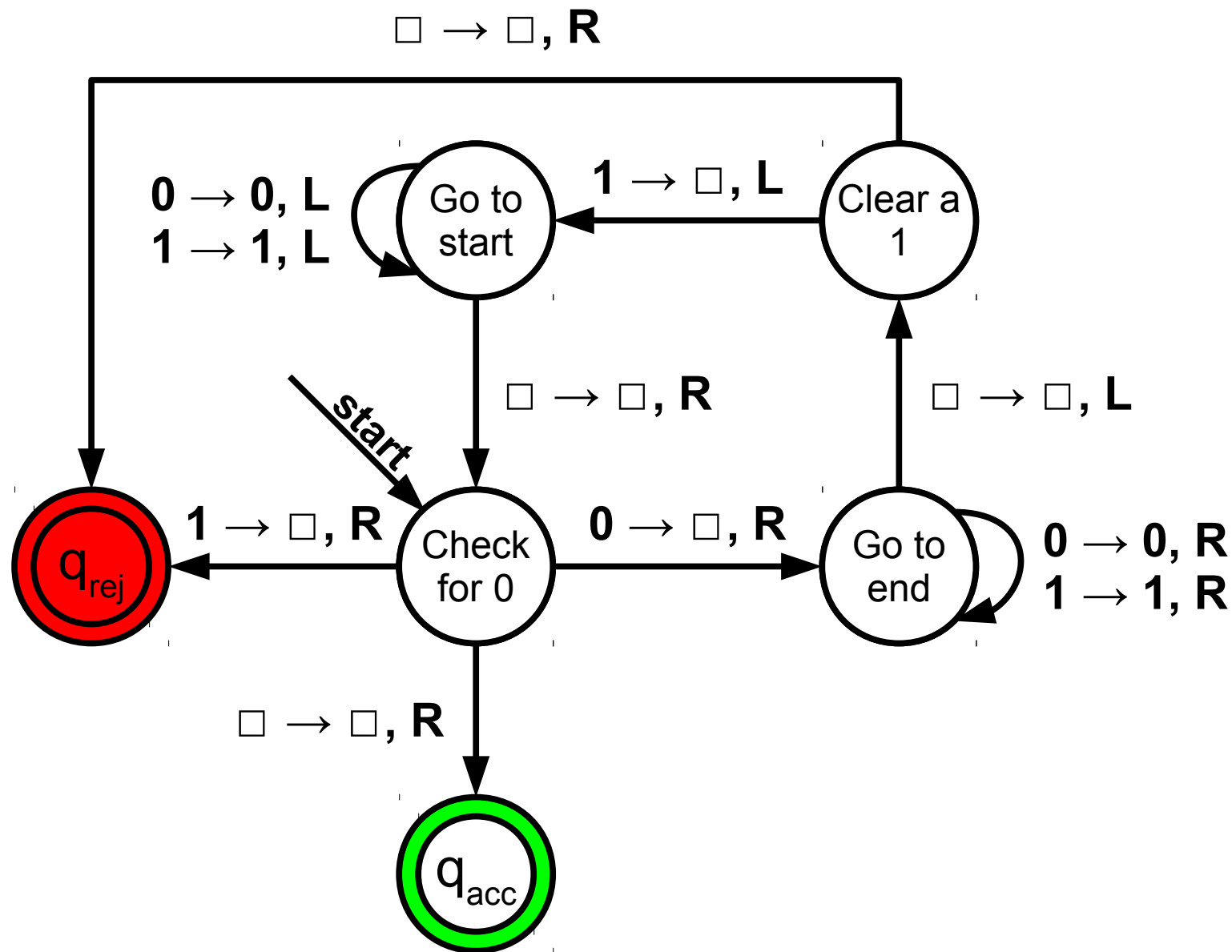


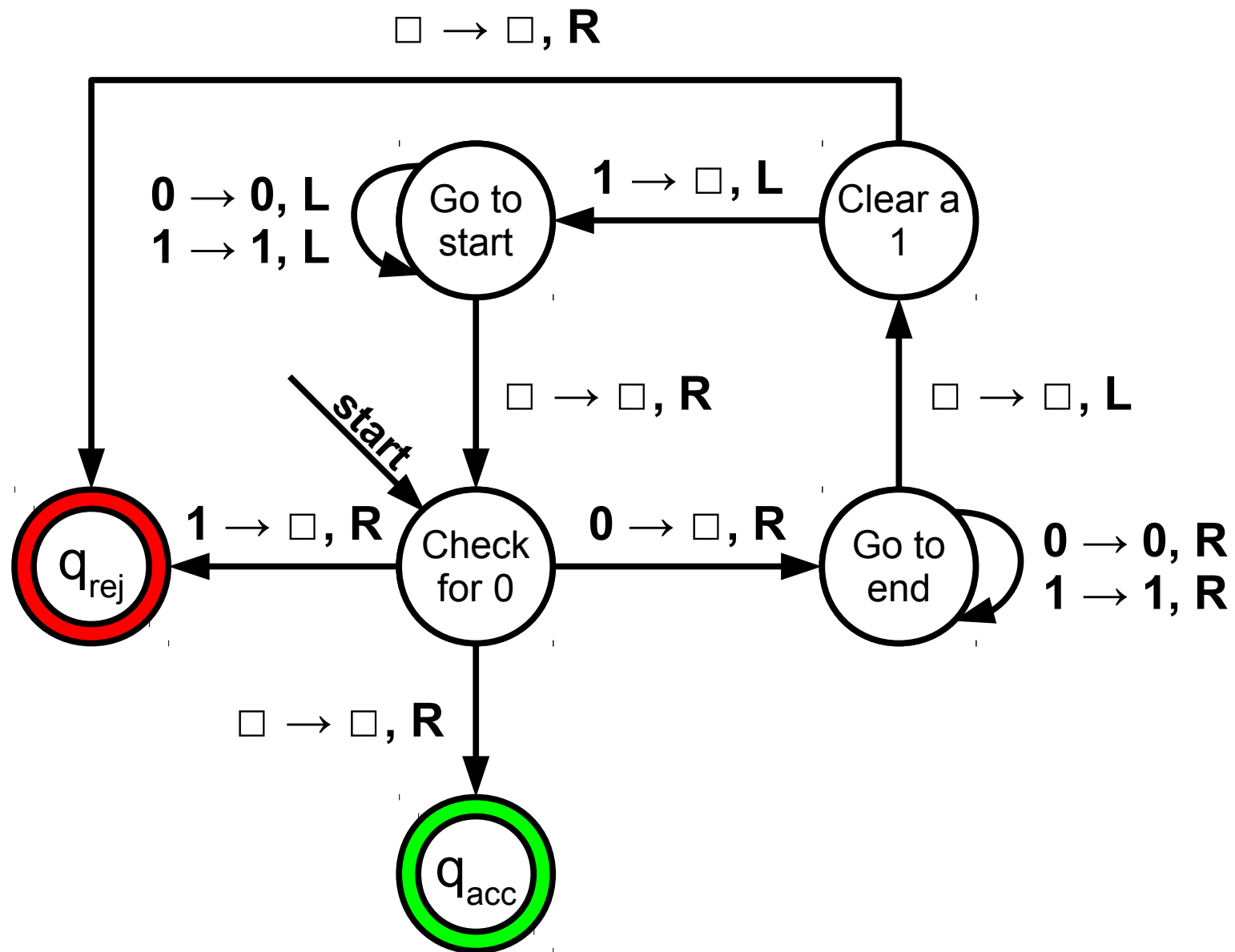


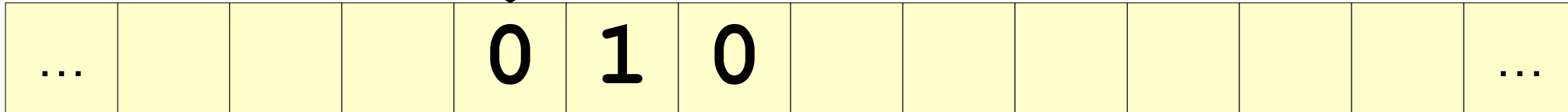
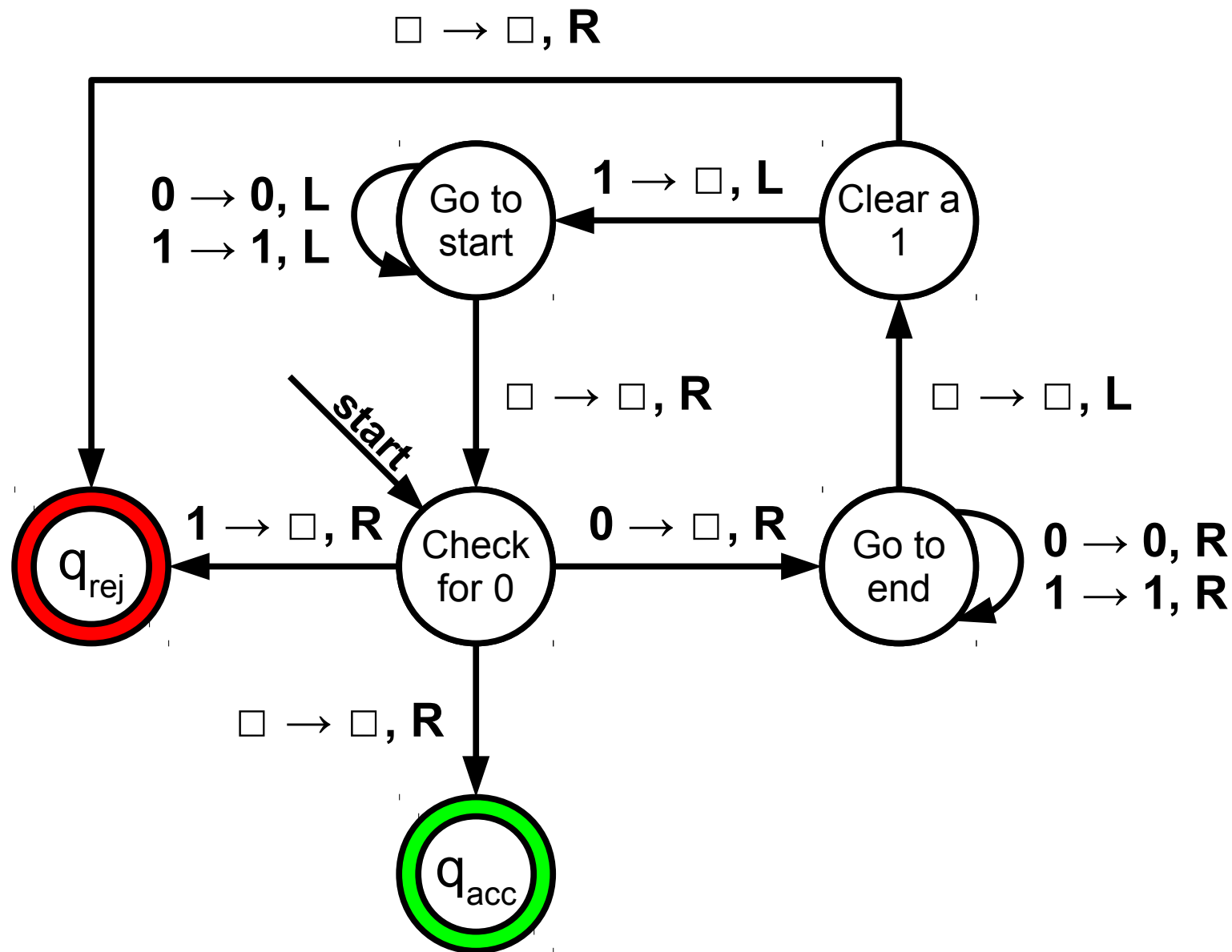


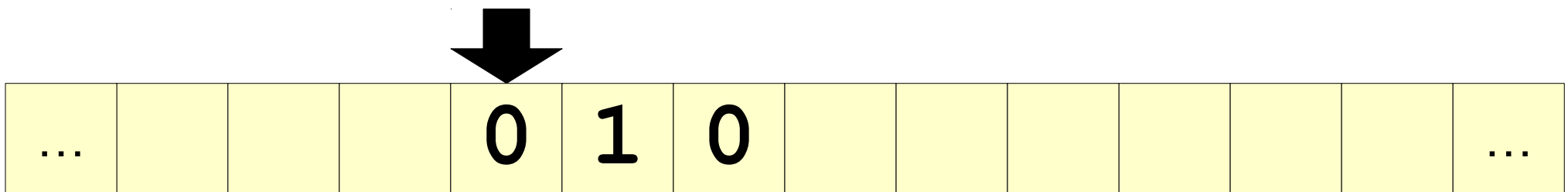
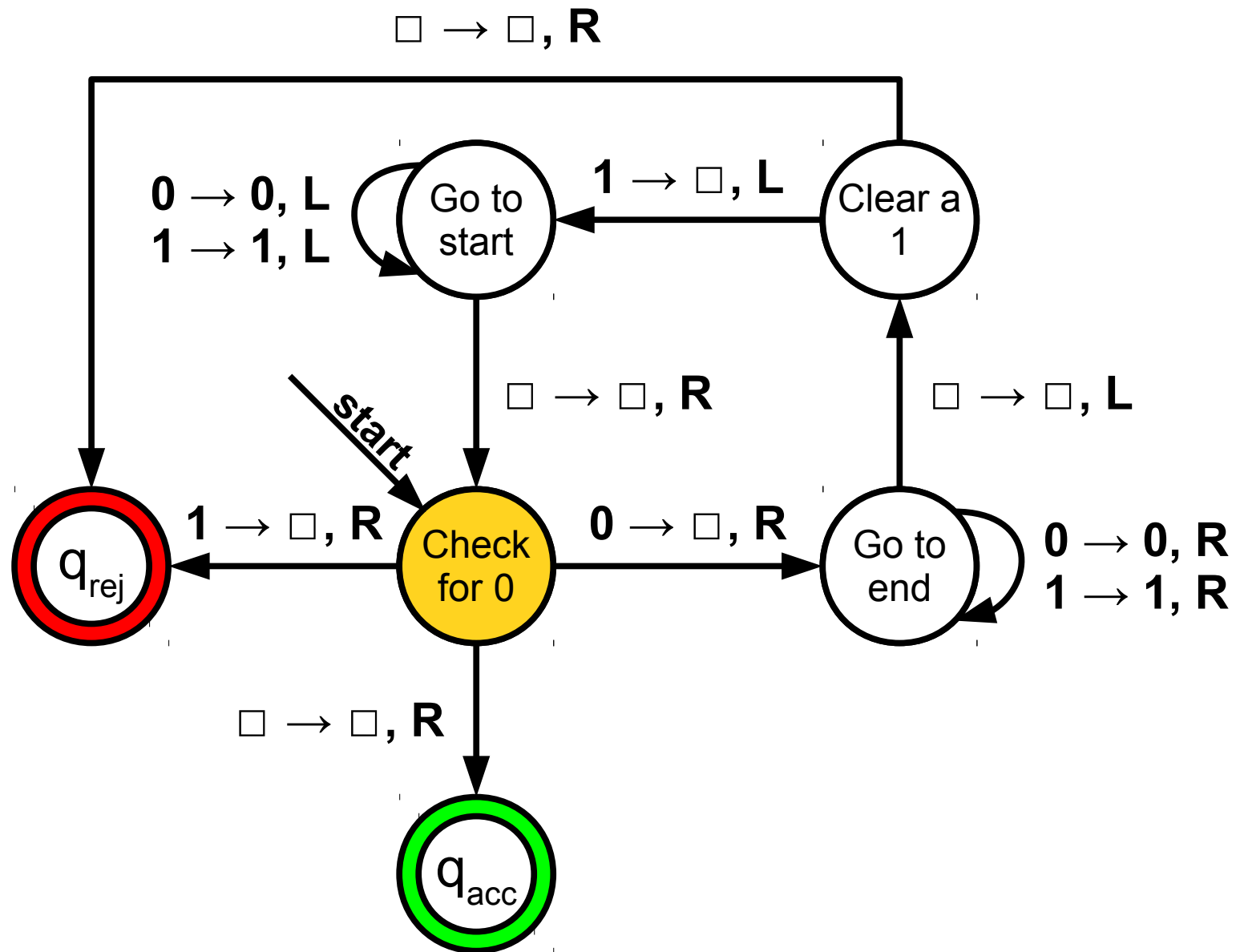


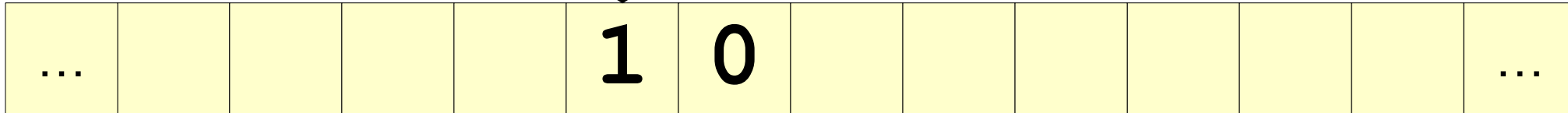
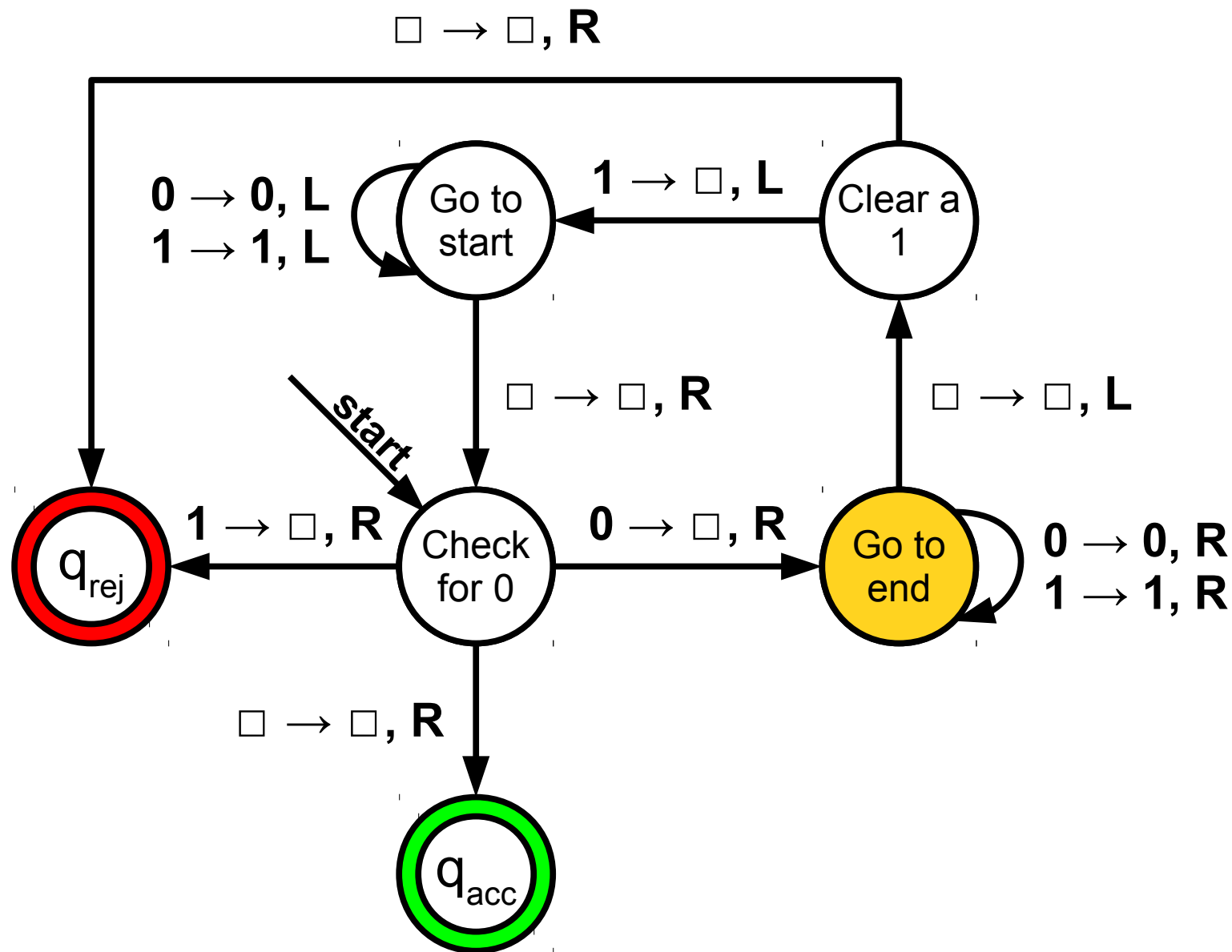


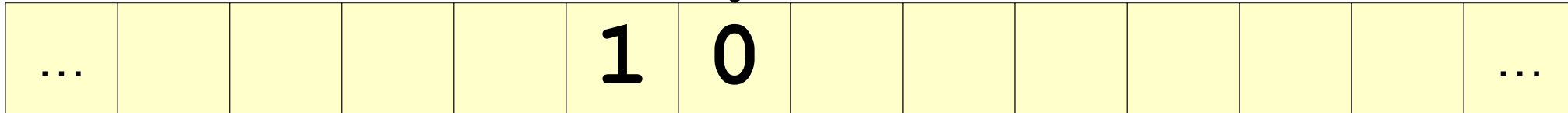
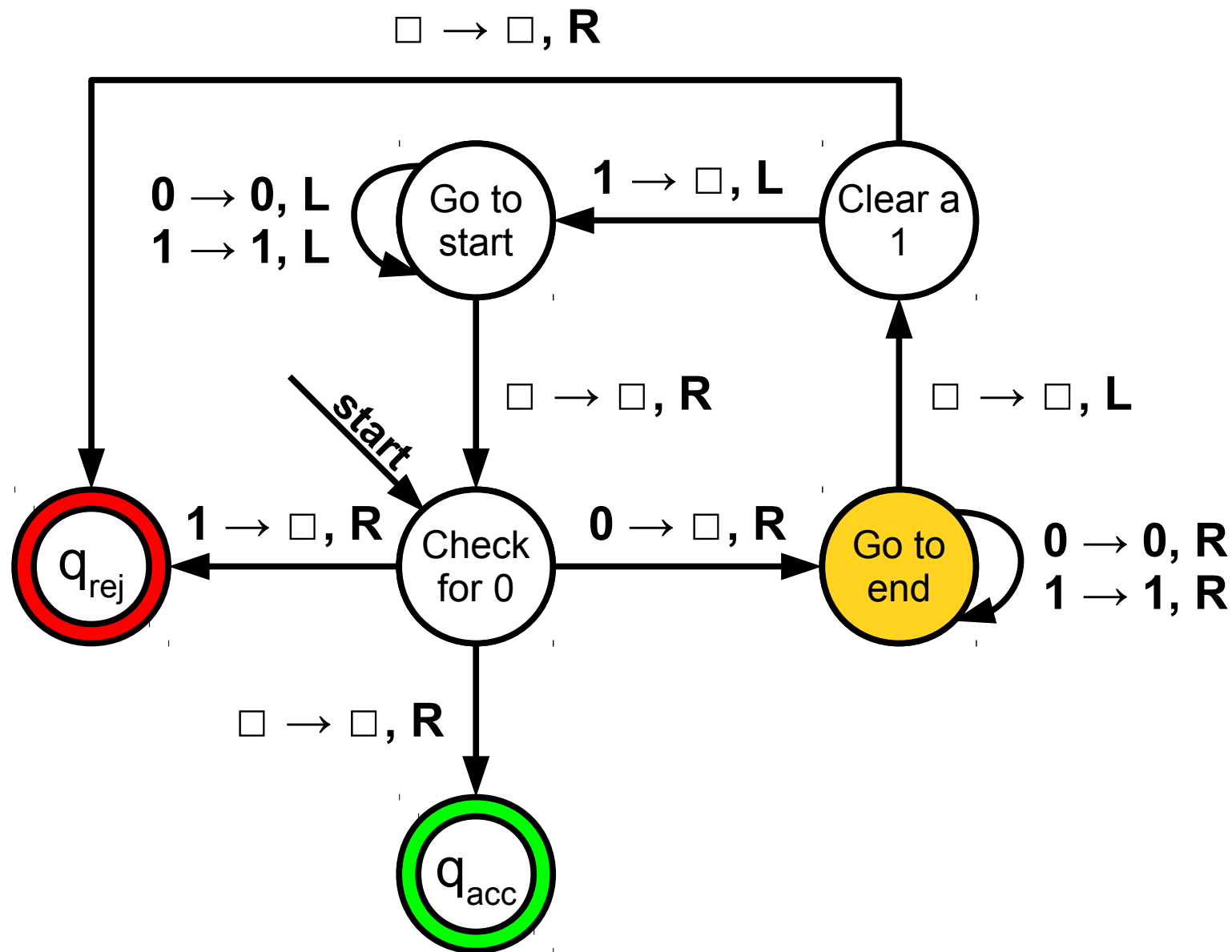


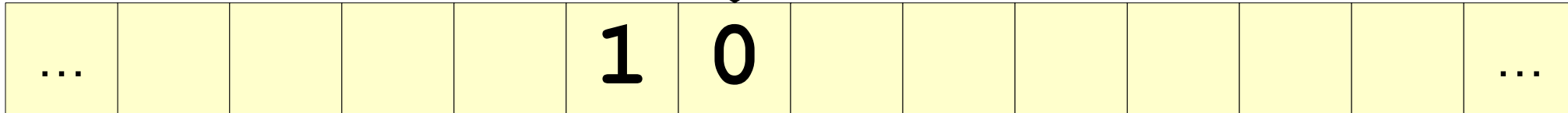
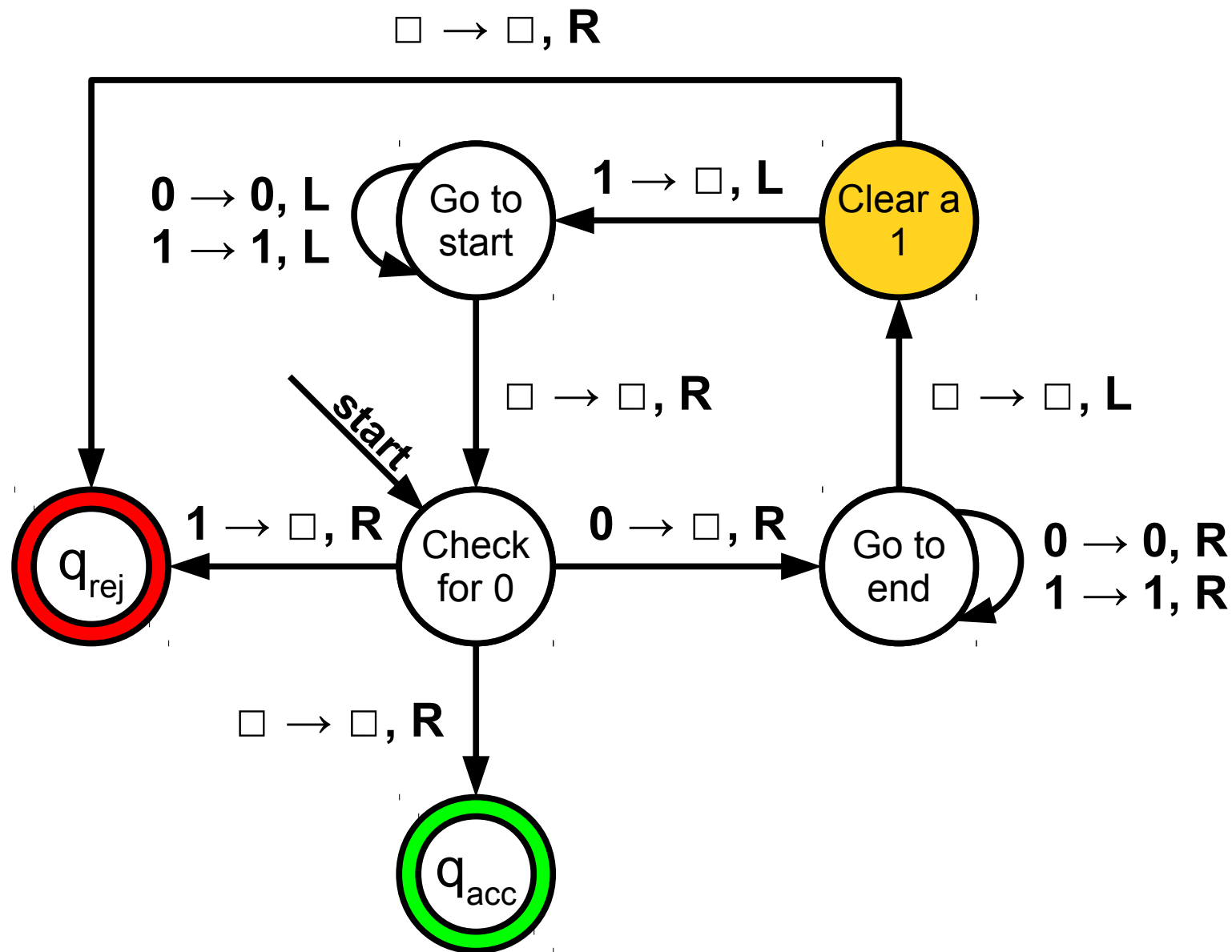


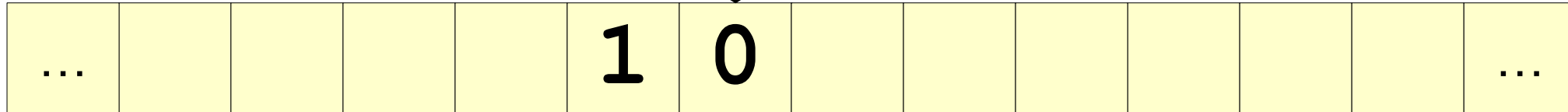
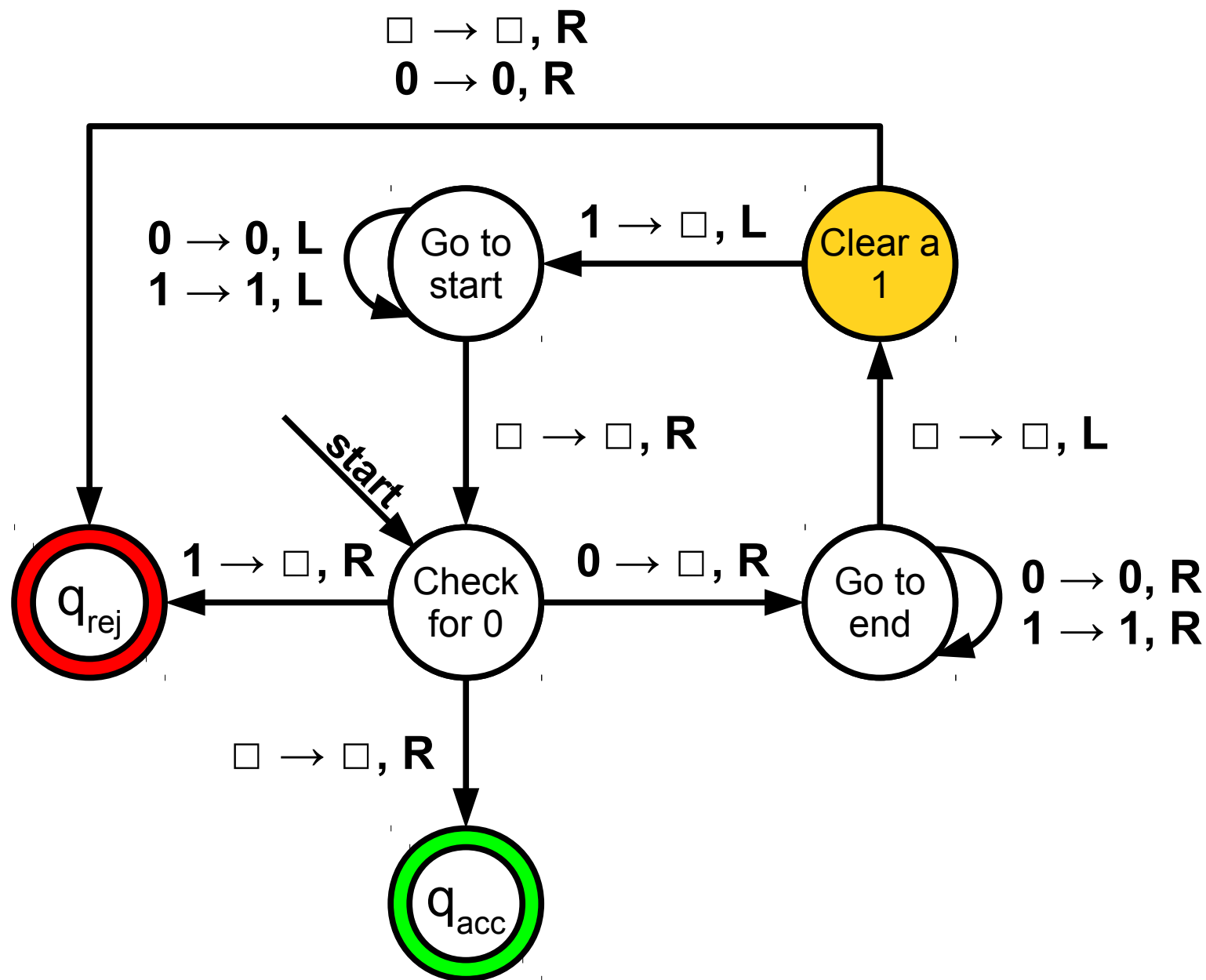


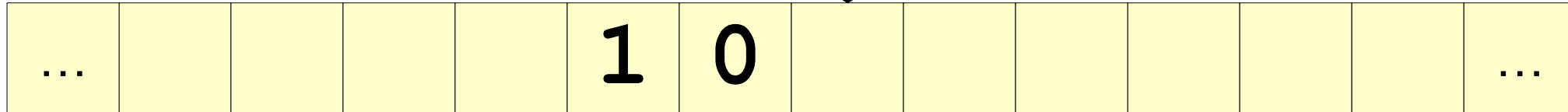
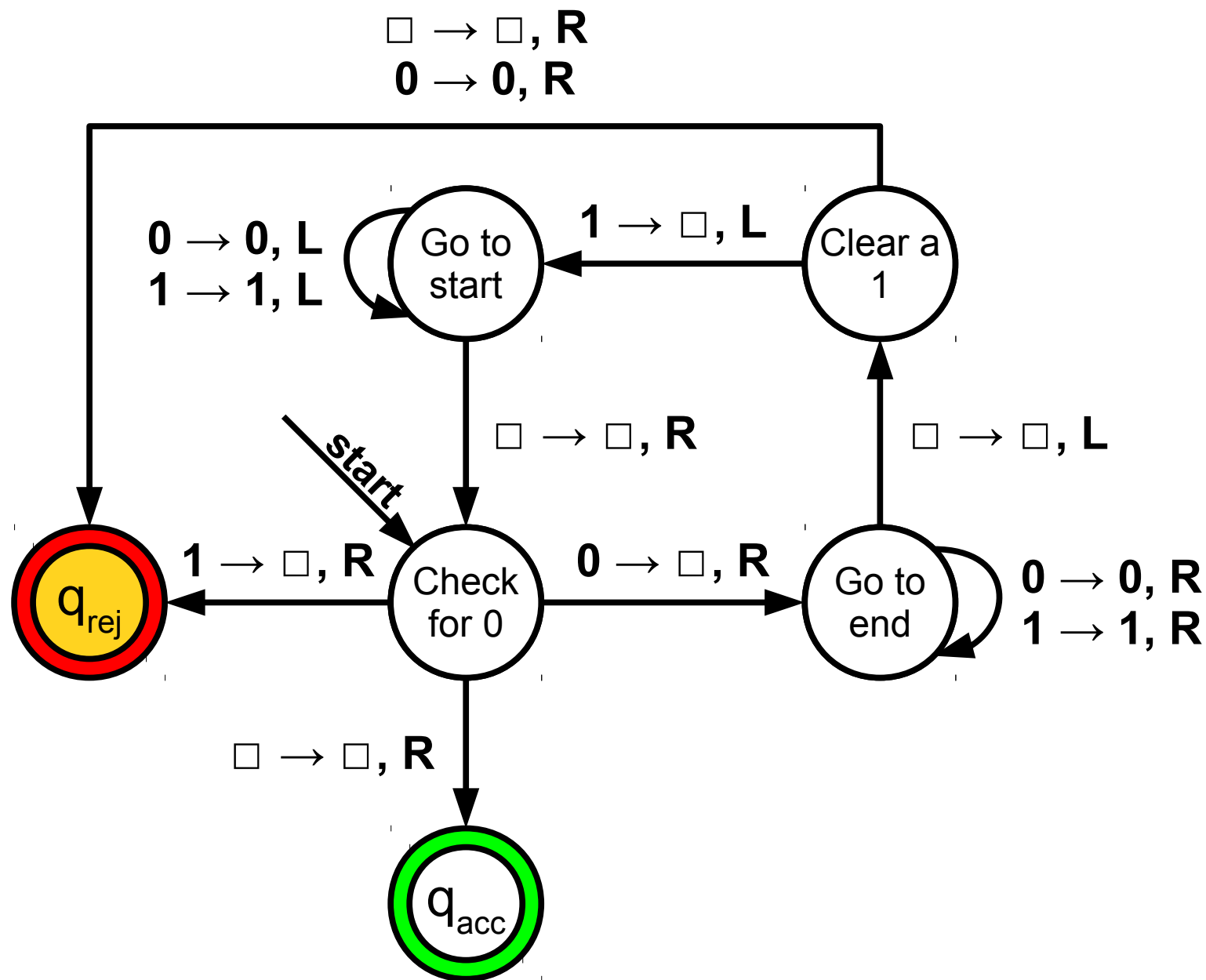


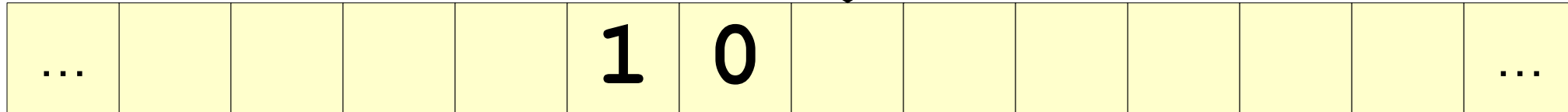
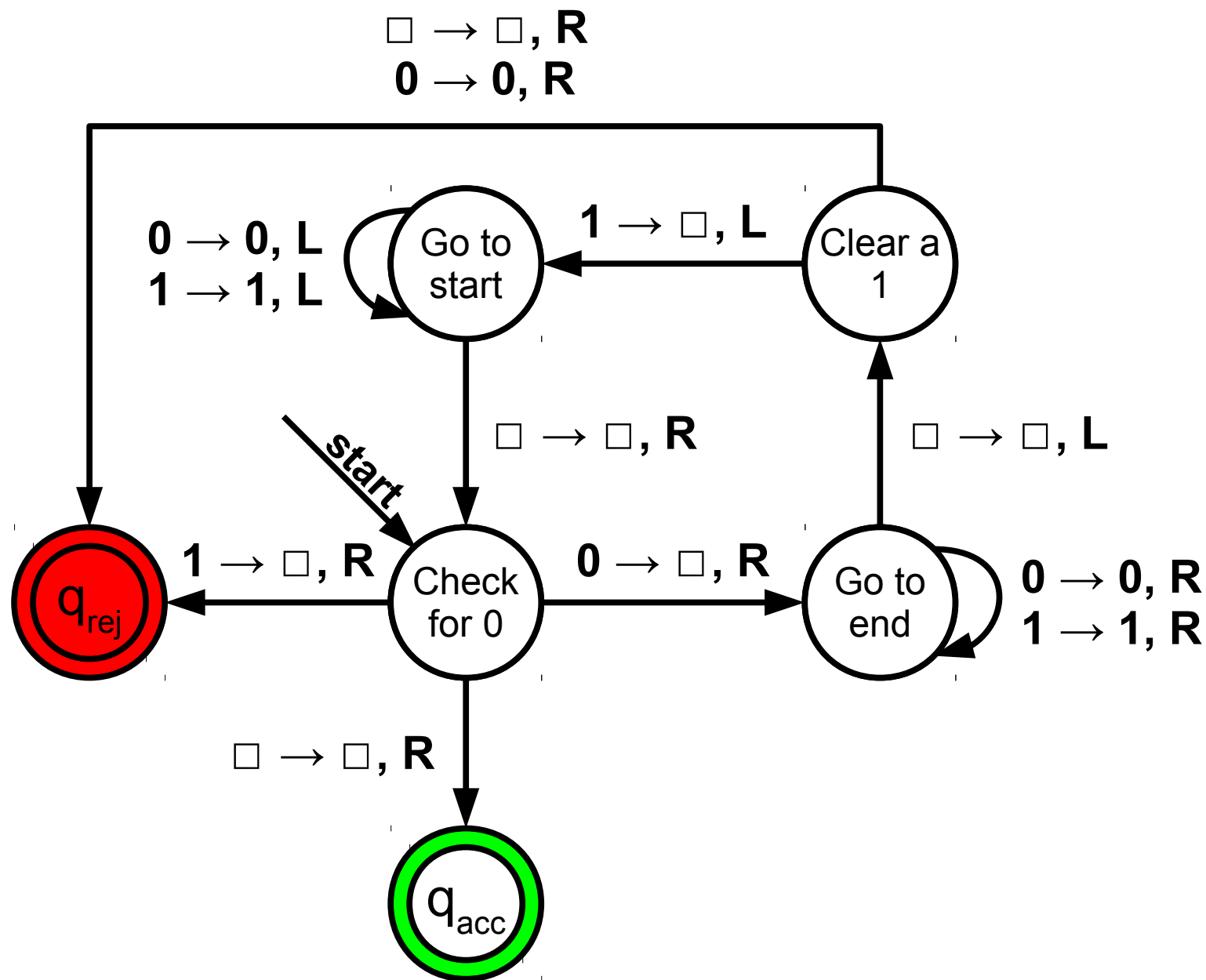


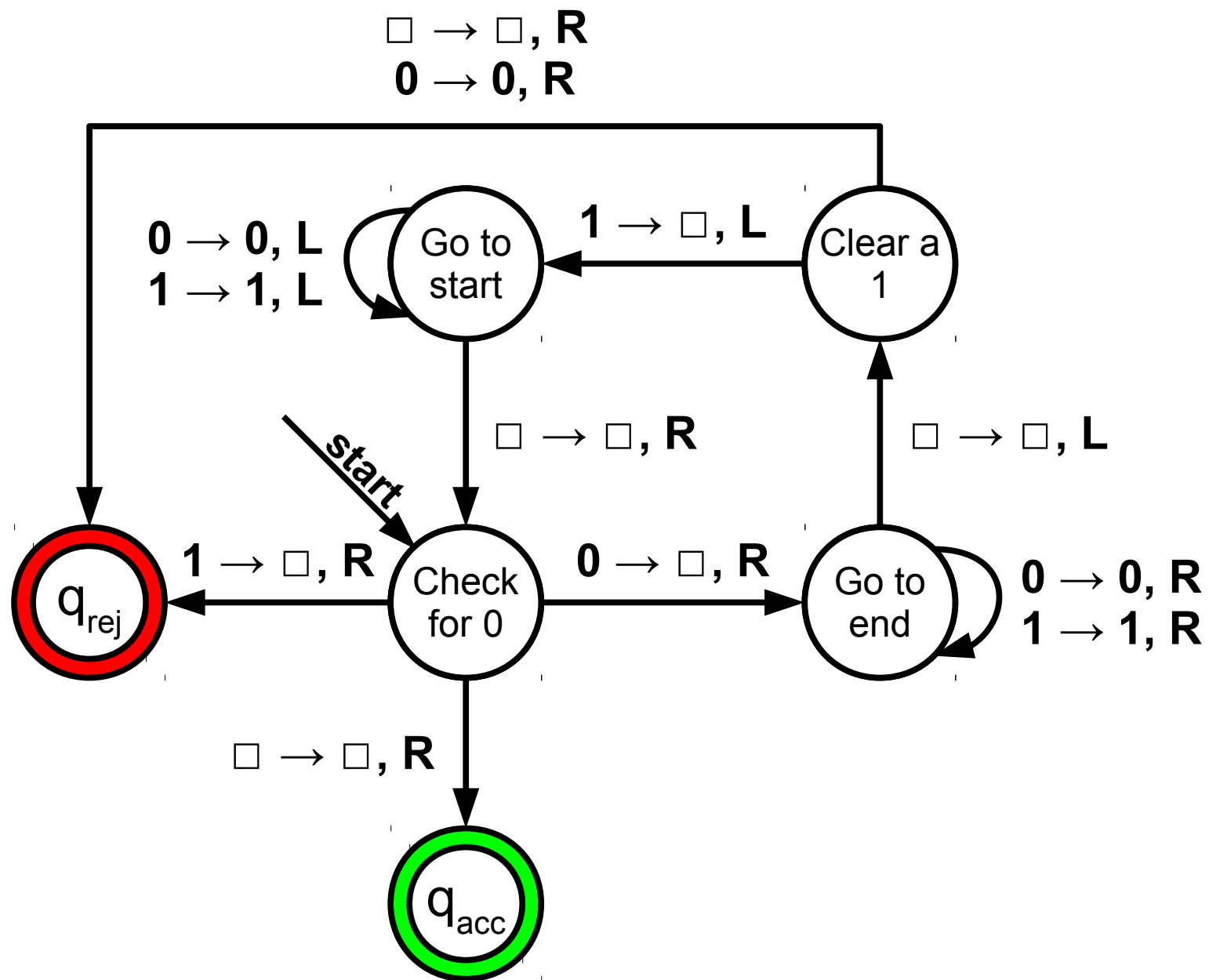


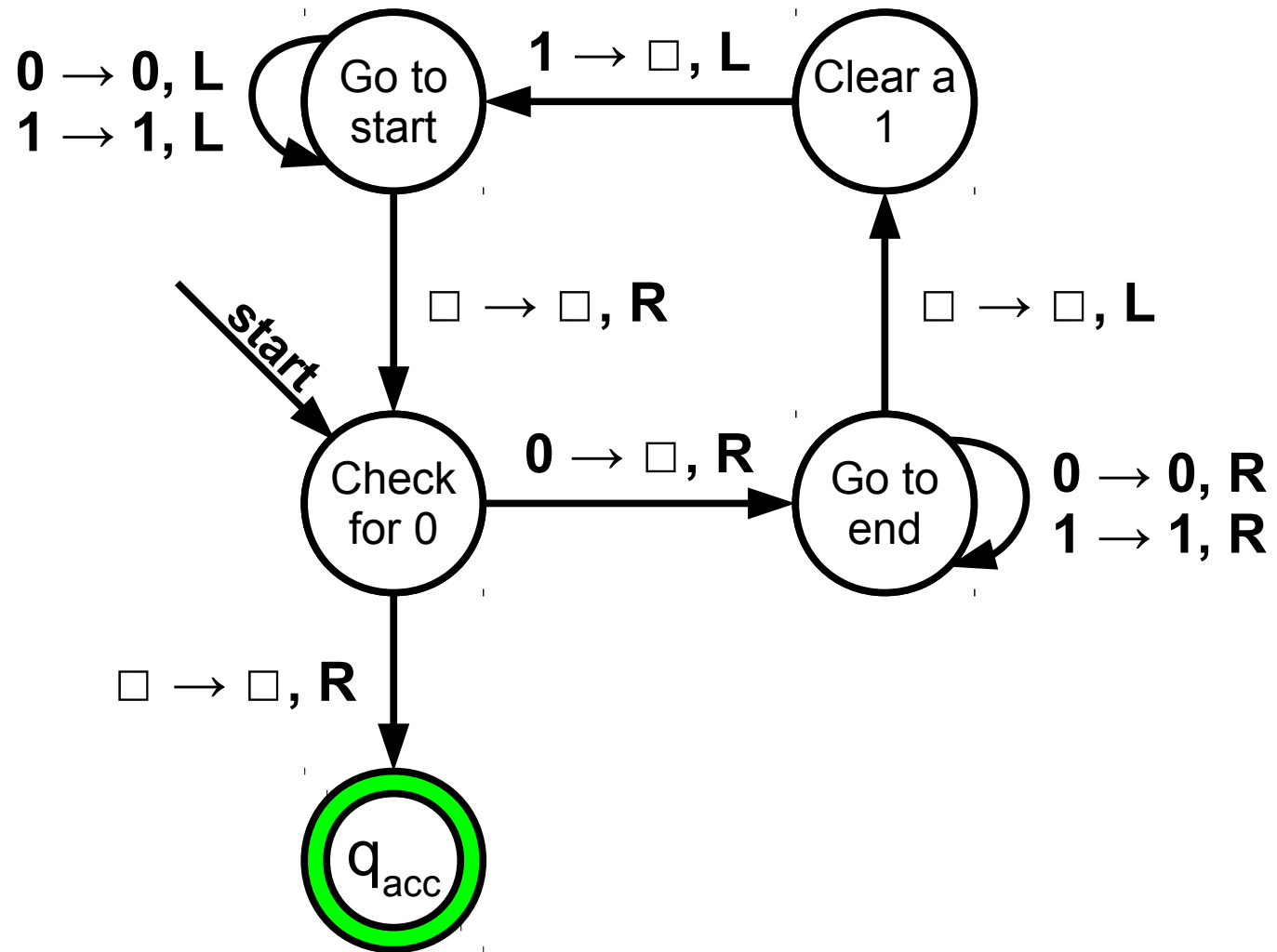












Another TM Design

- We've designed a TM for $\{0^n1^n \mid n \in \mathbb{N}\}$.
- Consider this language over $\Sigma = \{0, 1\}$:
$$L = \{ w \in \Sigma^* \mid w \text{ has the same number of 0s and 1s} \}$$
- This language is also not regular, but it is context-free.
- How might we design a TM for it?

A Caveat



...			0	0	0	1	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

A Caveat



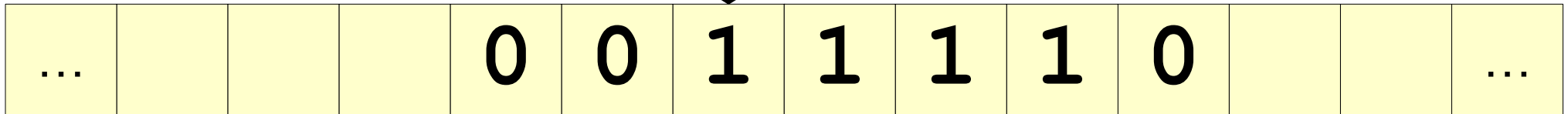
...				0	0	1	1	1	1	0			...
-----	--	--	--	---	---	---	---	---	---	---	--	--	-----

A Caveat



...				0	0	1	1	1	1	0			...
-----	--	--	--	---	---	---	---	---	---	---	--	--	-----

A Caveat



A Caveat



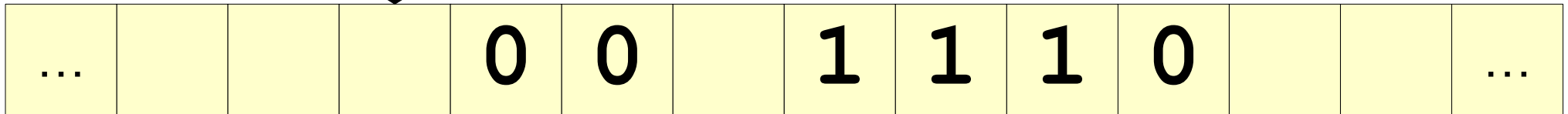
...				0	0		1	1	1	0			...
-----	--	--	--	---	---	--	---	---	---	---	--	--	-----

A Caveat



...				0	0		1	1	1	0			...
-----	--	--	--	---	---	--	---	---	---	---	--	--	-----

A Caveat



A Caveat



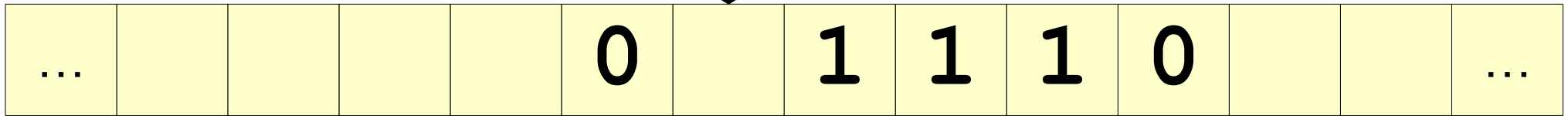
...				0	0		1	1	1	0			...
-----	--	--	--	---	---	--	---	---	---	---	--	--	-----

A Caveat

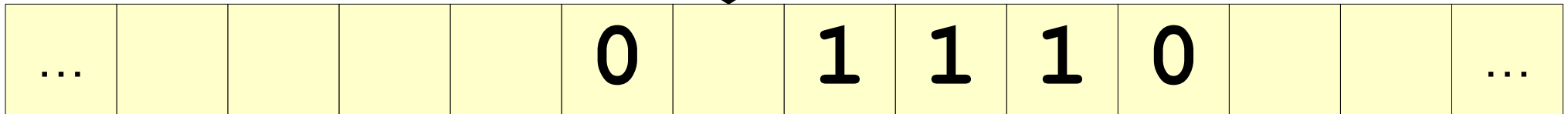


...					0		1	1	1	0			...
-----	--	--	--	--	---	--	---	---	---	---	--	--	-----

A Caveat

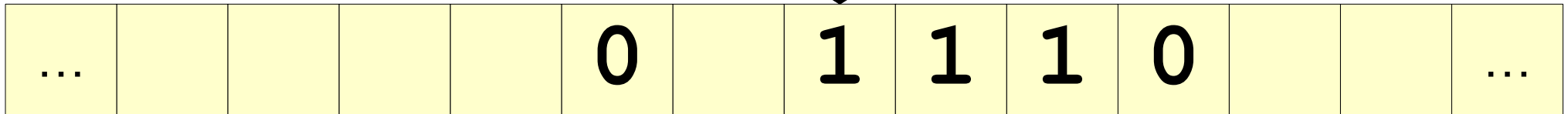
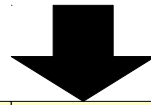


A Caveat

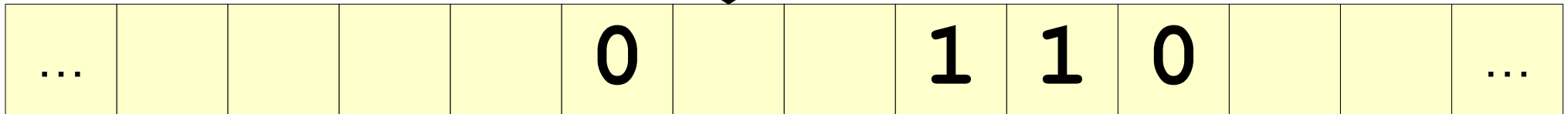


How do we know that
this blank isn't one of
the infinitely many
blanks after our input
string?

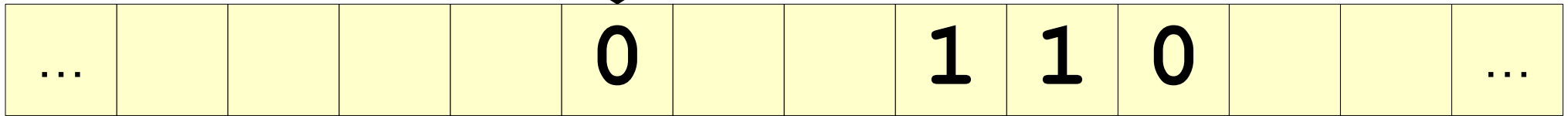
A Caveat



A Caveat



A Caveat



A Caveat



...					0			1	1	0			...
-----	--	--	--	--	---	--	--	---	---	---	--	--	-----

A Caveat



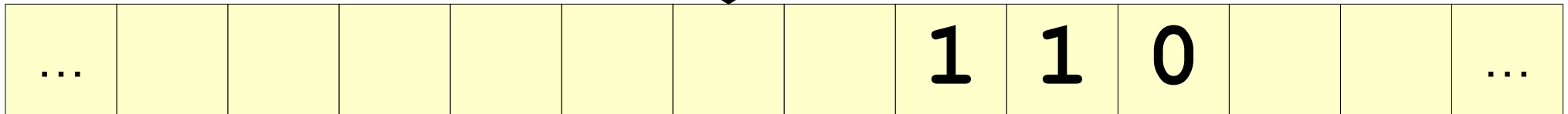
...					0			1	1	0			...
-----	--	--	--	--	---	--	--	---	---	---	--	--	-----

A Caveat



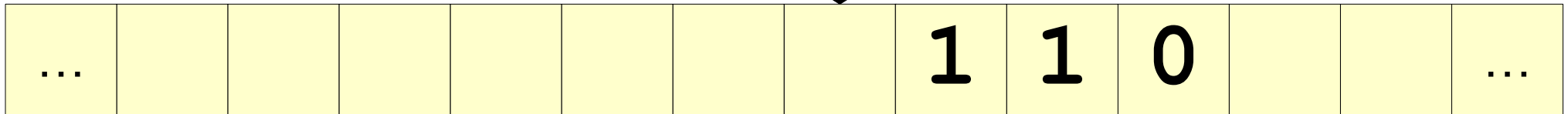
...								1	1	0			...
-----	--	--	--	--	--	--	--	---	---	---	--	--	-----

A Caveat

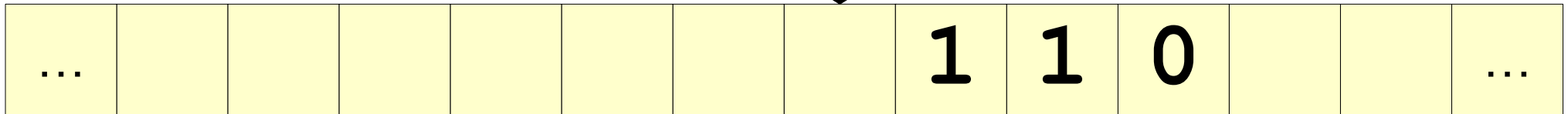


How do we know that
this blank isn't one of
the infinitely many
blanks after our input
string?

A Caveat



A Caveat



How do we know that
this blank isn't one of
the infinitely many
blanks after our input
string?

One Solution



...			0	0	0	1	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

One Solution



...			×	0	0	1	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

One Solution



...			×	0	0	1	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

One Solution



...			×	0	0	1	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

One Solution



...			x	0	0	x	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

One Solution



...			x	0	0	x	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

One Solution



...			x	0	0	x	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

One Solution



...			x	0	0	x	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

One Solution



...			x	0	0	x	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

One Solution



...			x	0	0	x	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

One Solution



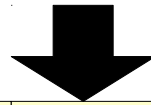
...			x	x	0	x	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

One Solution



...			x	x	0	x	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

One Solution



...			x	x	0	x	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

One Solution



...			x	x	0	x	x	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

One Solution



...			x	x	0	x	x	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

One Solution



...			x	x	0	x	x	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

One Solution



...			x	x	0	x	x	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

One Solution



...			x	x	0	x	x	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

One Solution



...			x	x	0	x	x	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

One Solution



...			x	x	0	x	x	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

One Solution



...			x	x	0	x	x	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

One Solution



...			x	x	x	x	x	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

One Solution



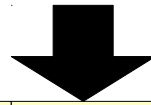
...			x	x	x	x	x	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

One Solution



...			x	x	x	x	x	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

One Solution

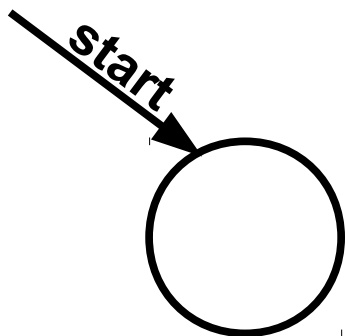


...			x	x	x	x	x	x	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

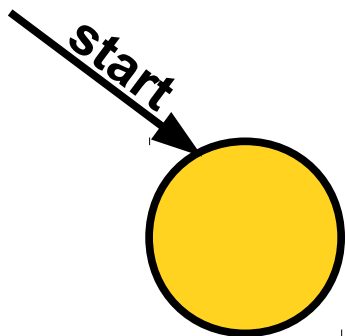
One Solution



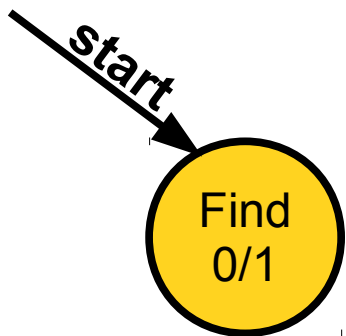
...			x	x	x	x	x	x	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----



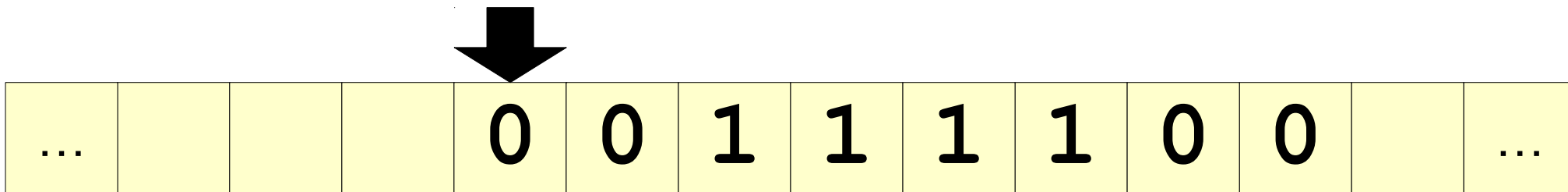
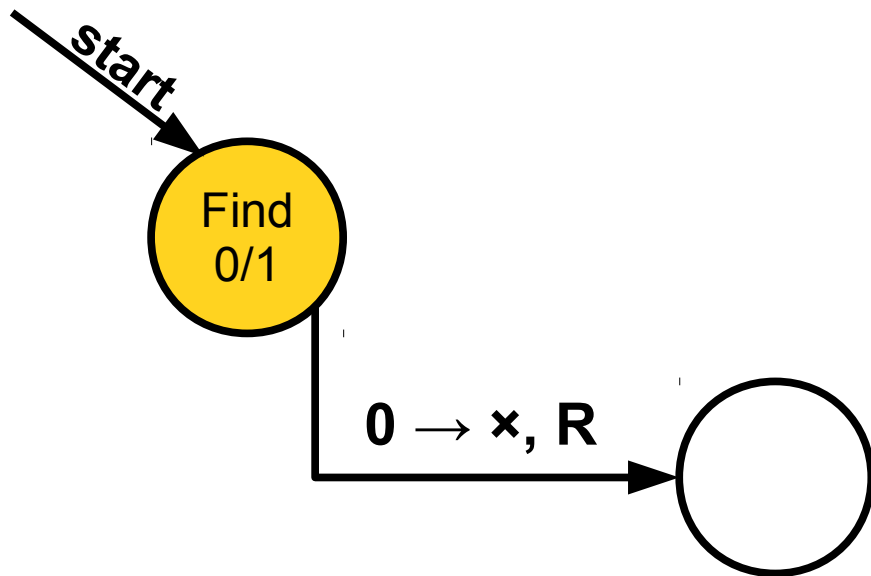
...				0	0	1	1	1	1	0	0		...
-----	--	--	--	---	---	---	---	---	---	---	---	--	-----

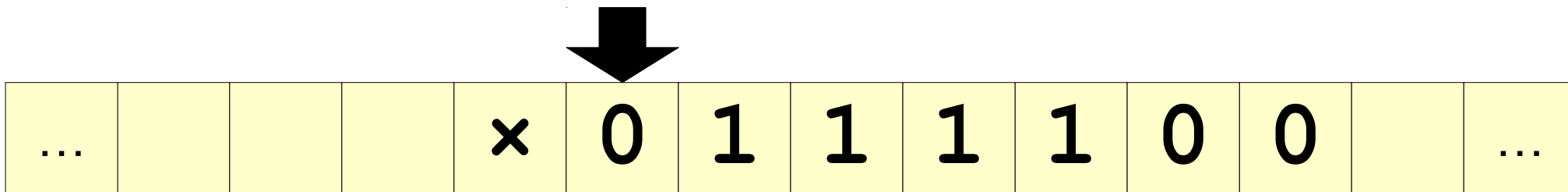
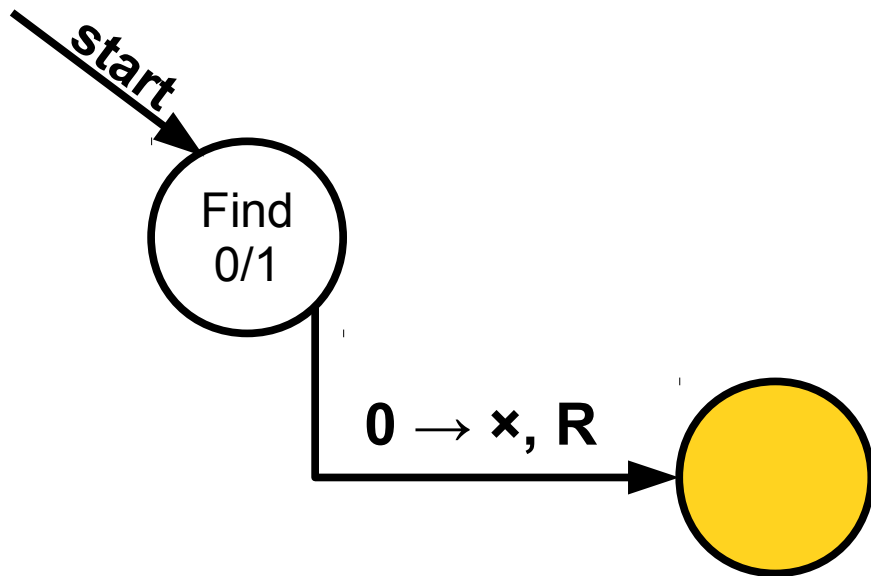


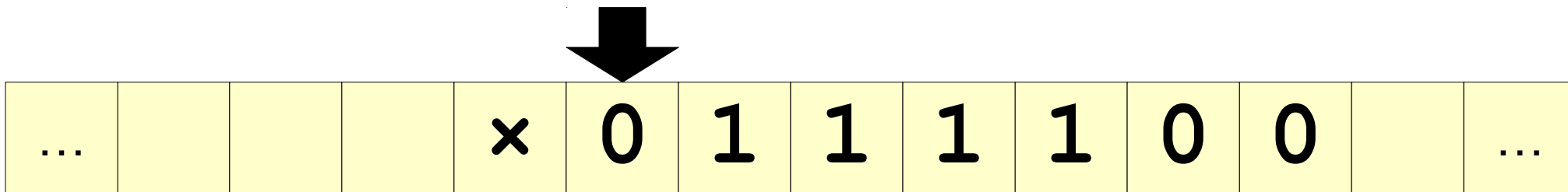
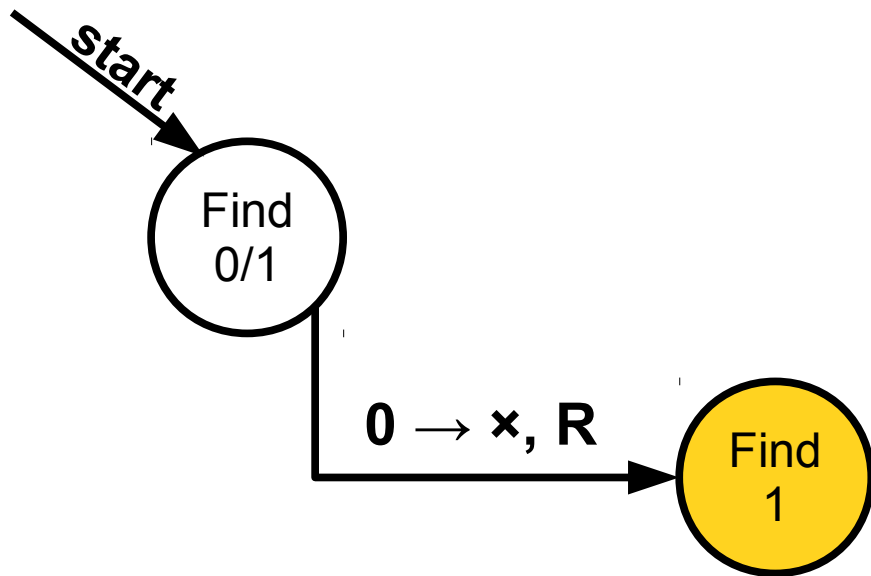
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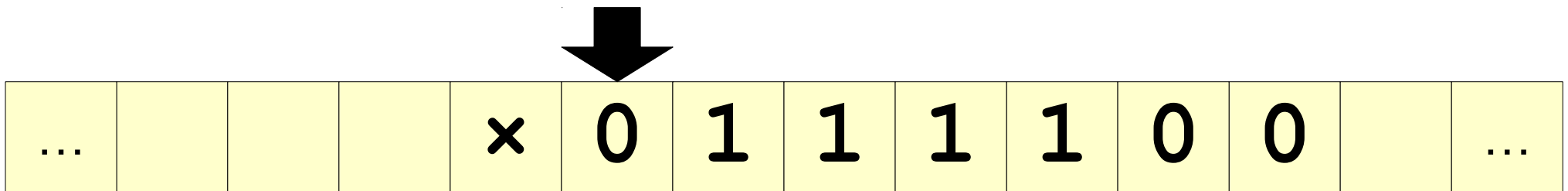
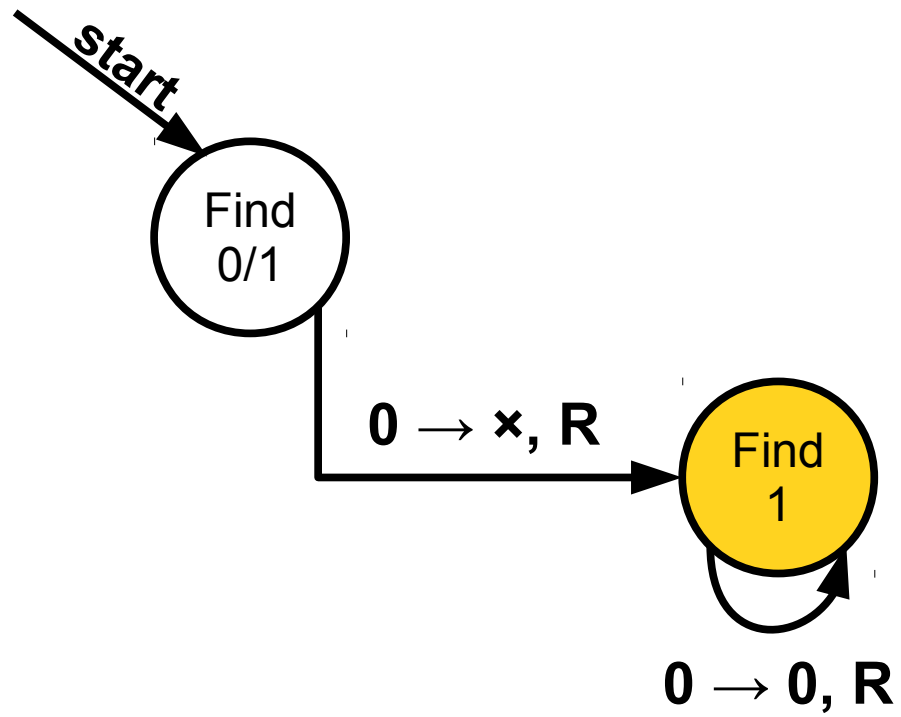


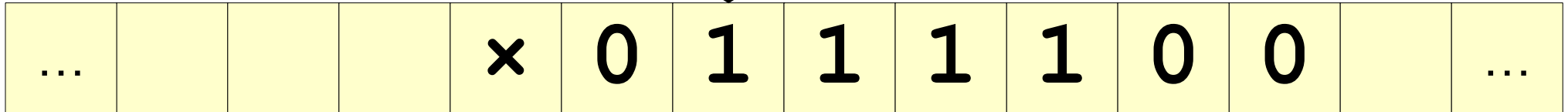
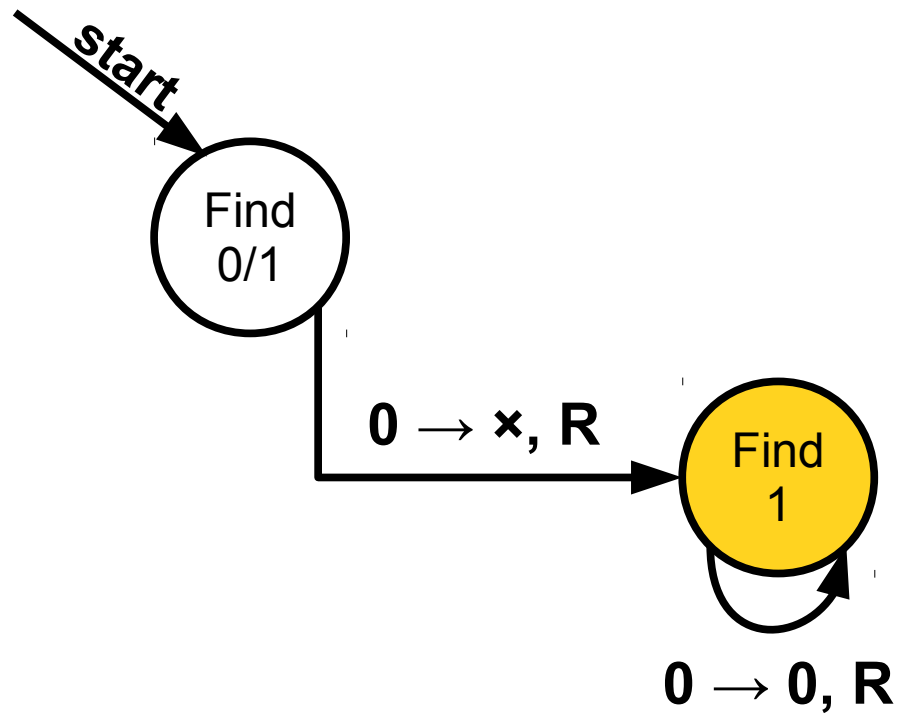
...				0	0	1	1	1	1	0	0		...
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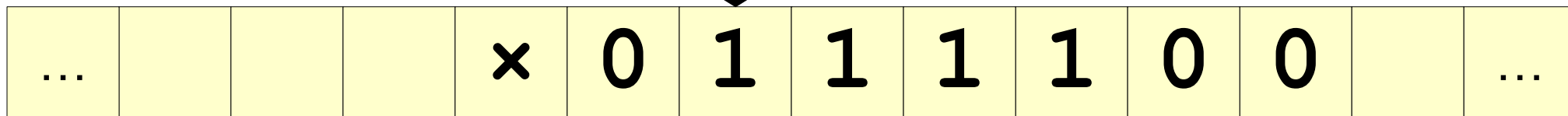
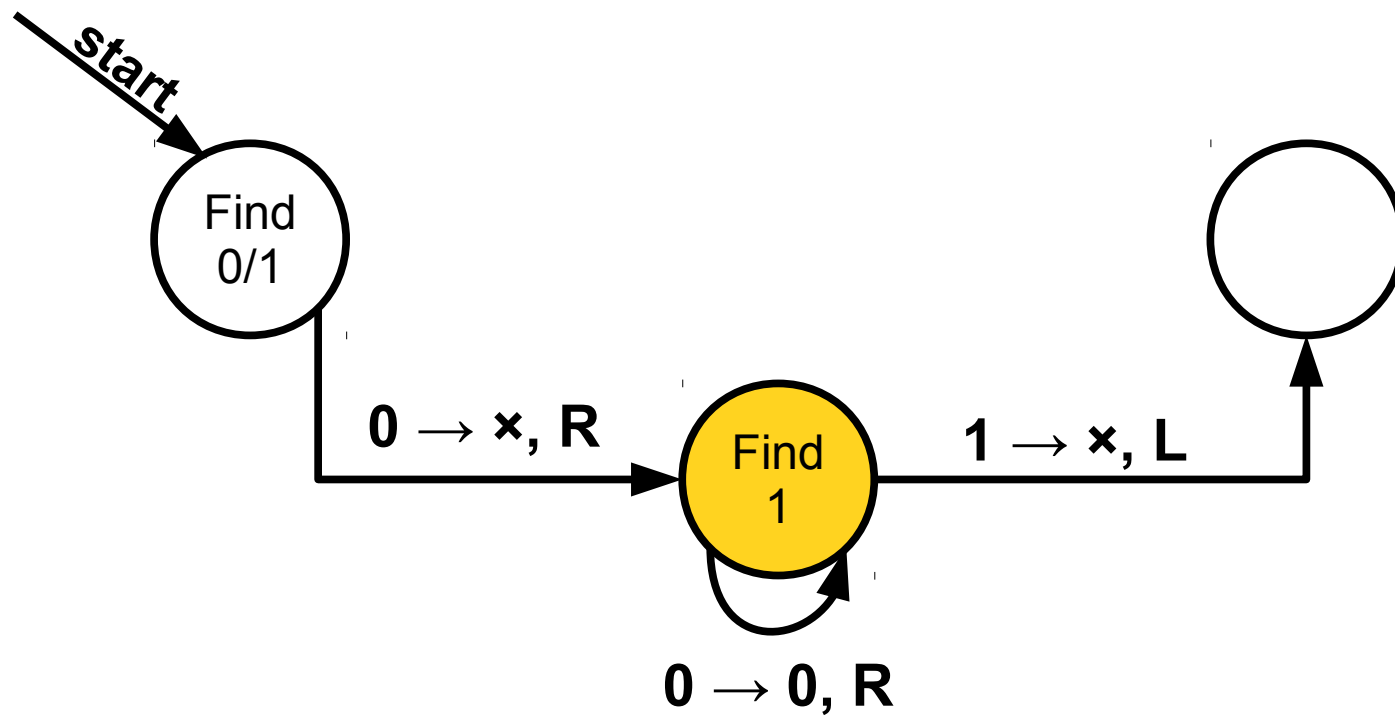


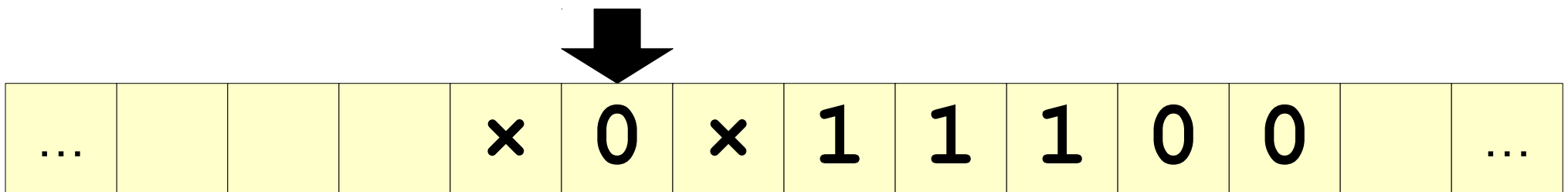
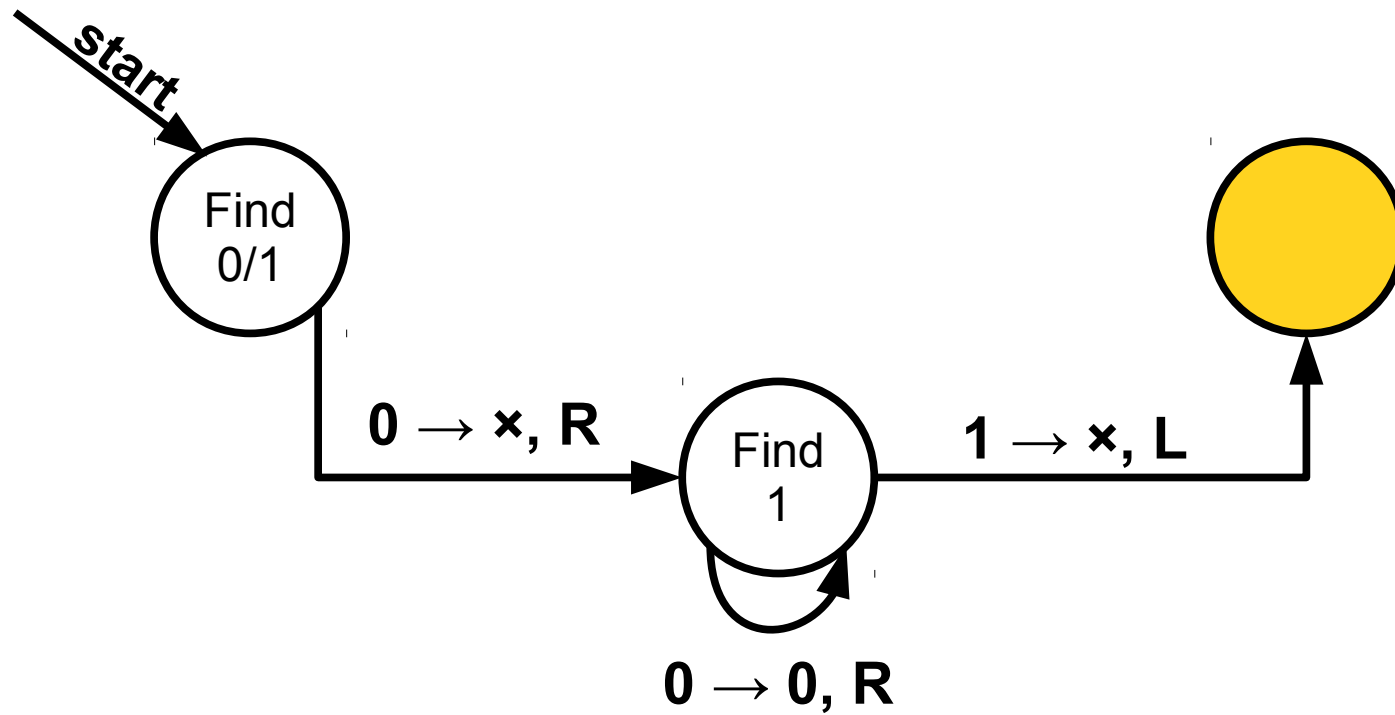


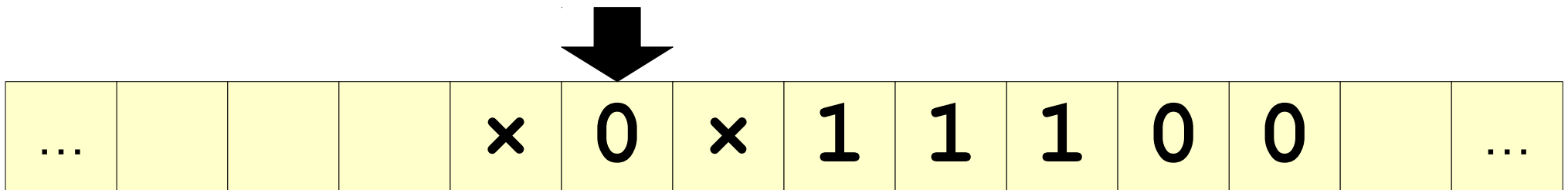
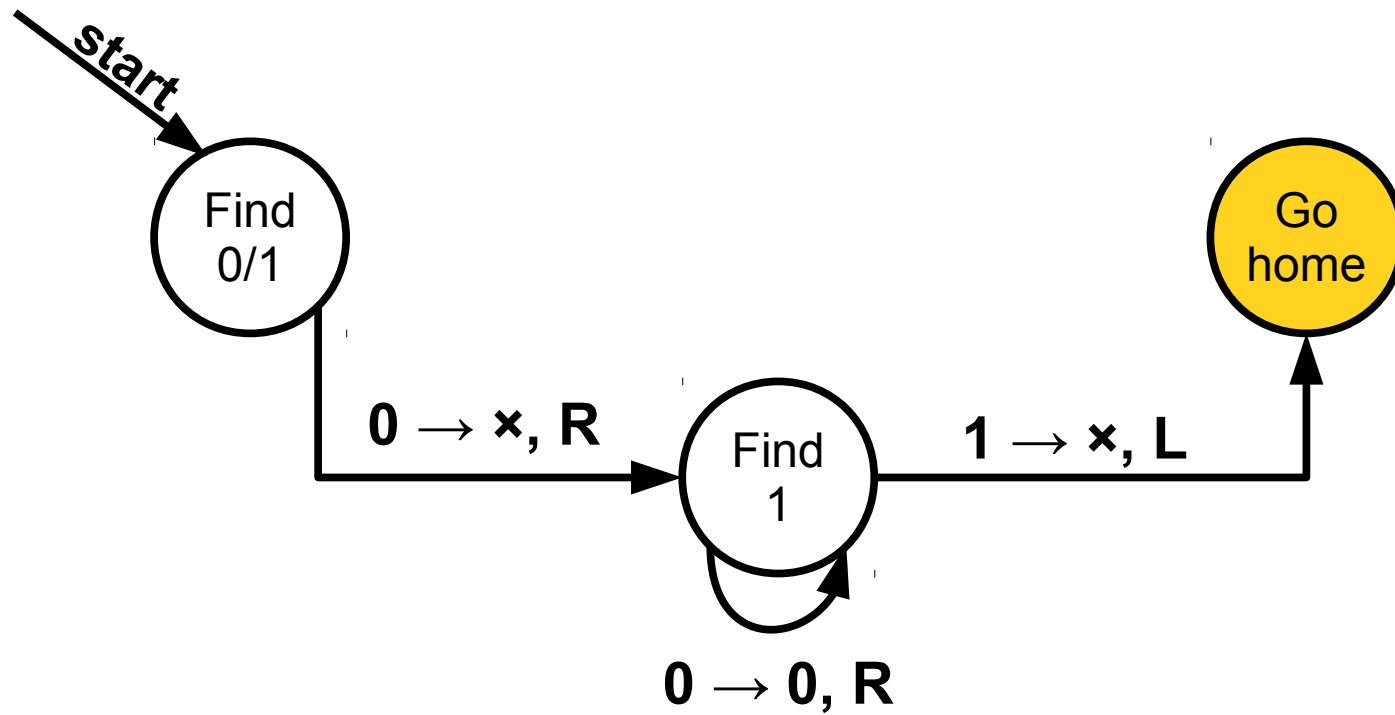


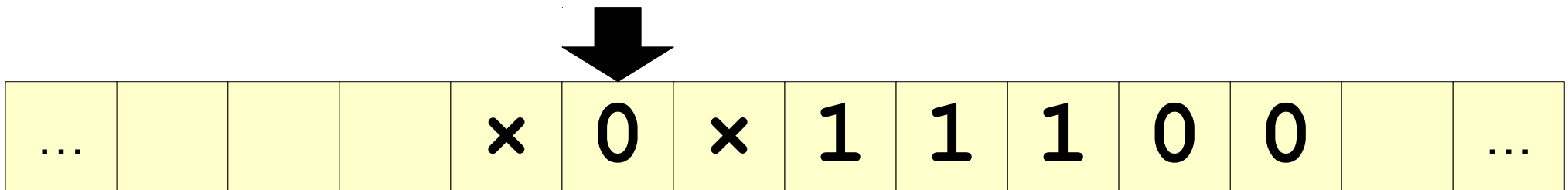
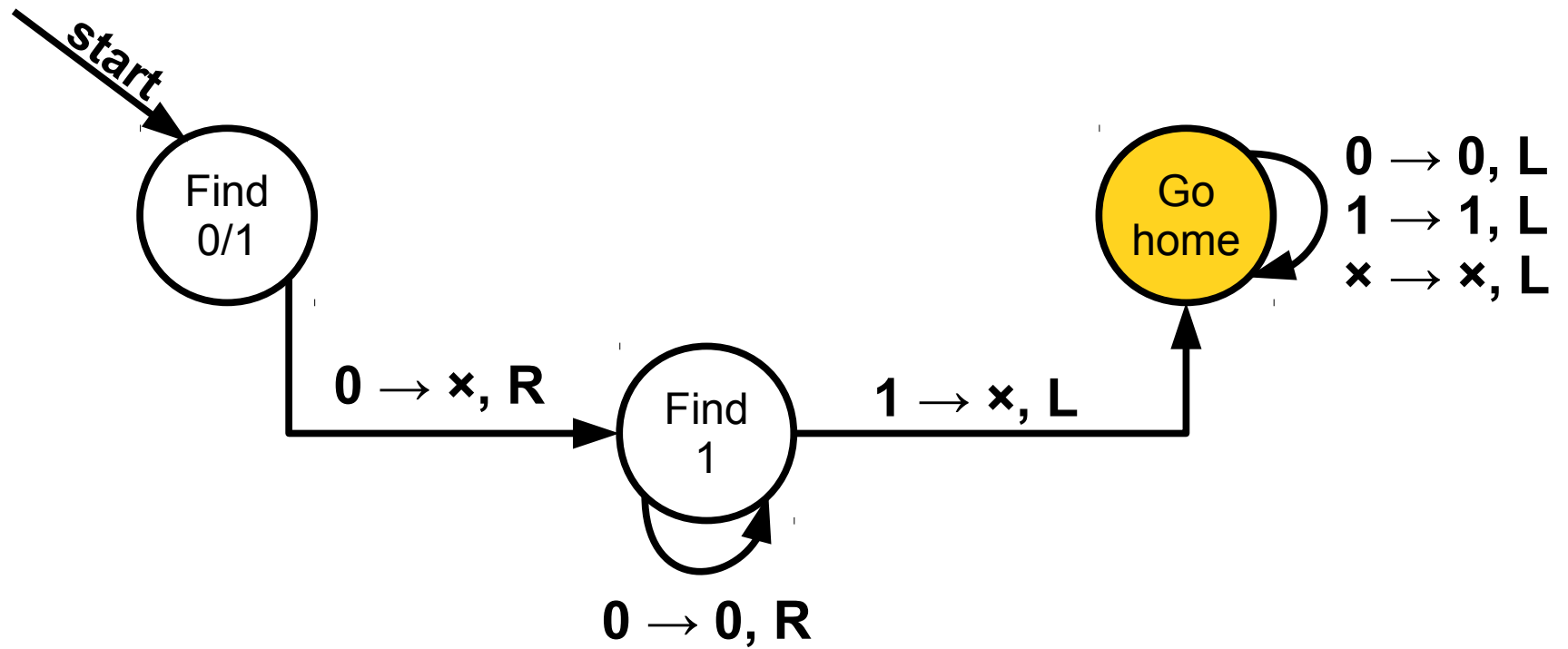


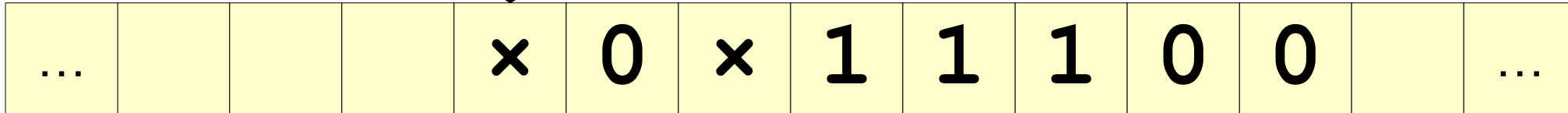
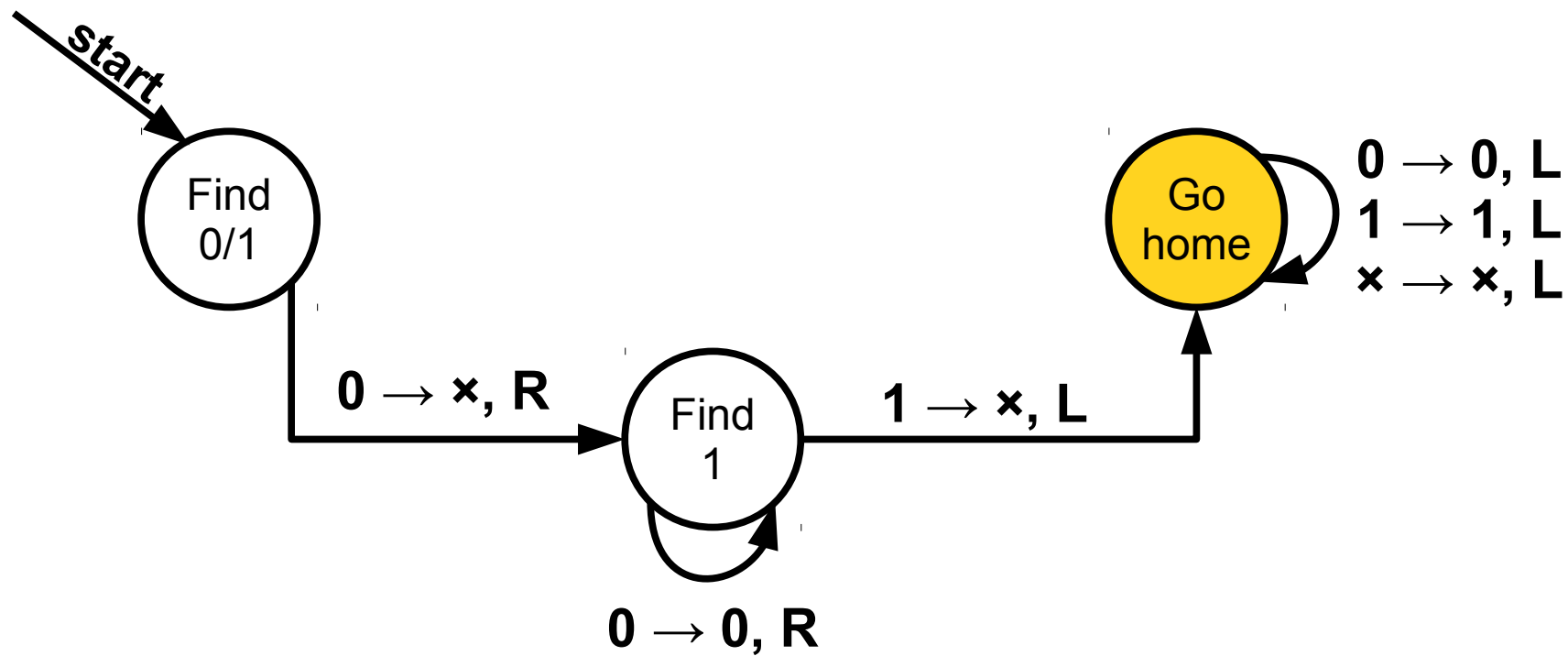


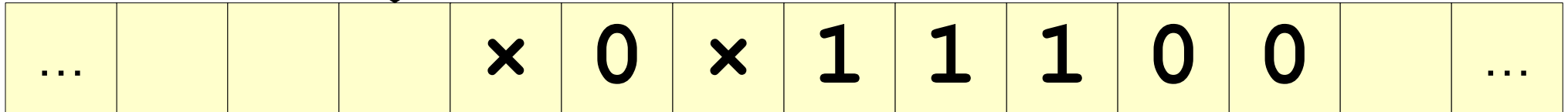
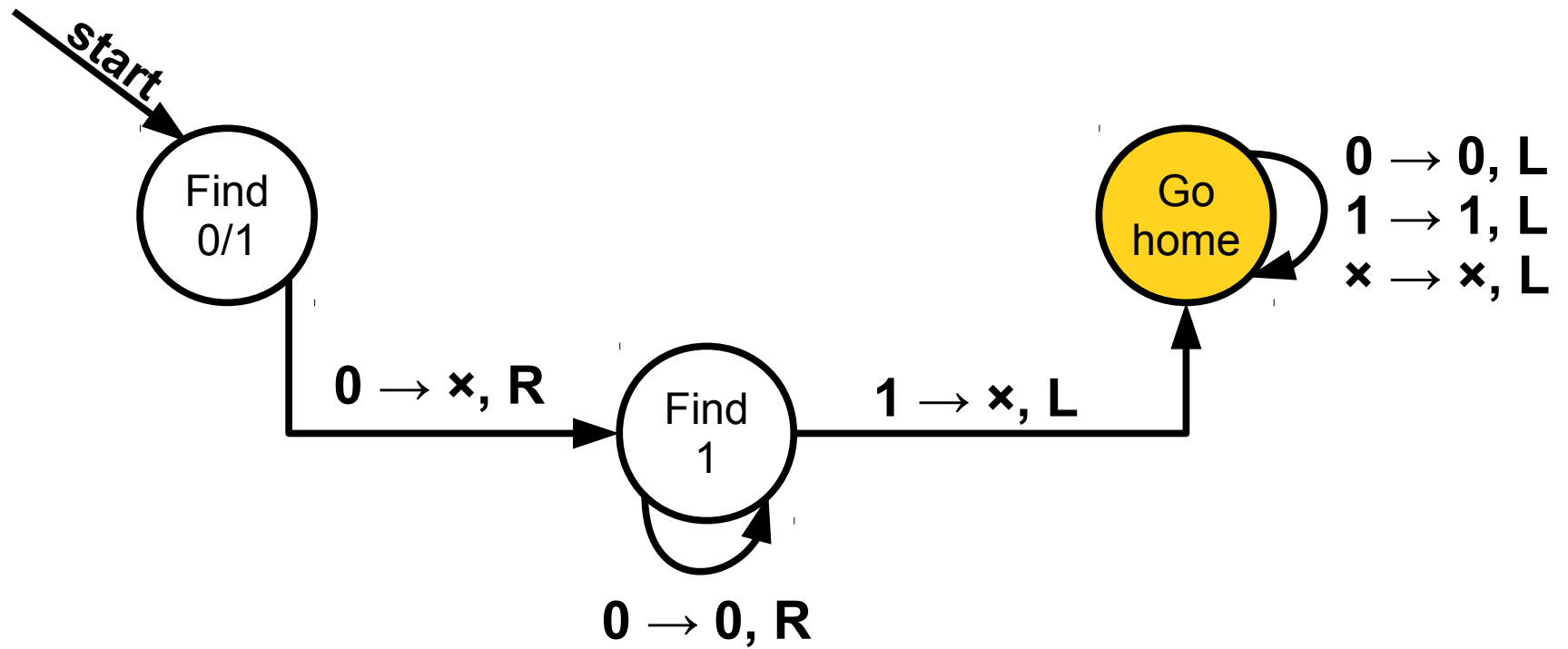


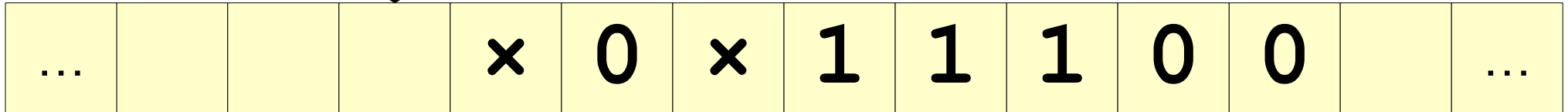
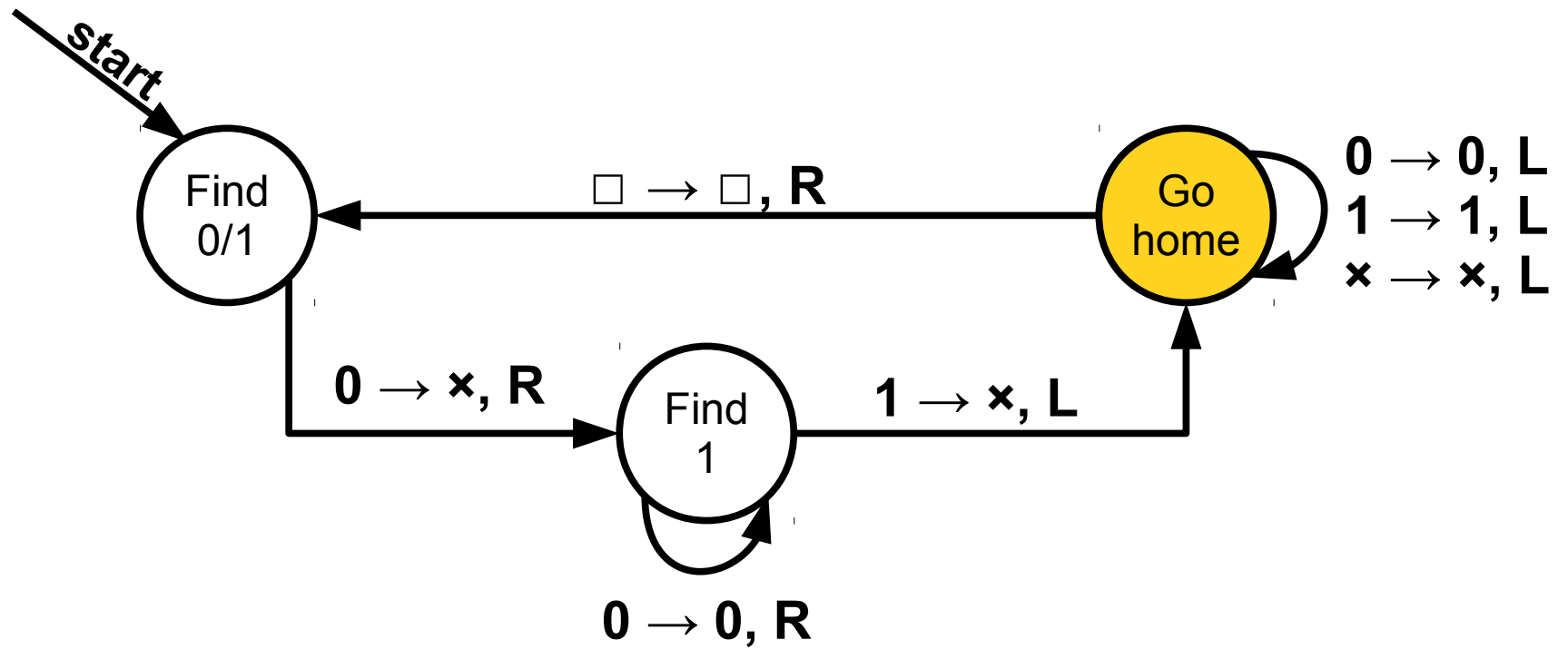


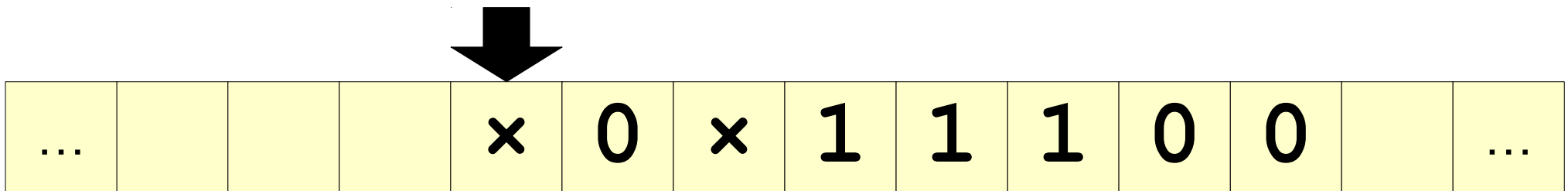
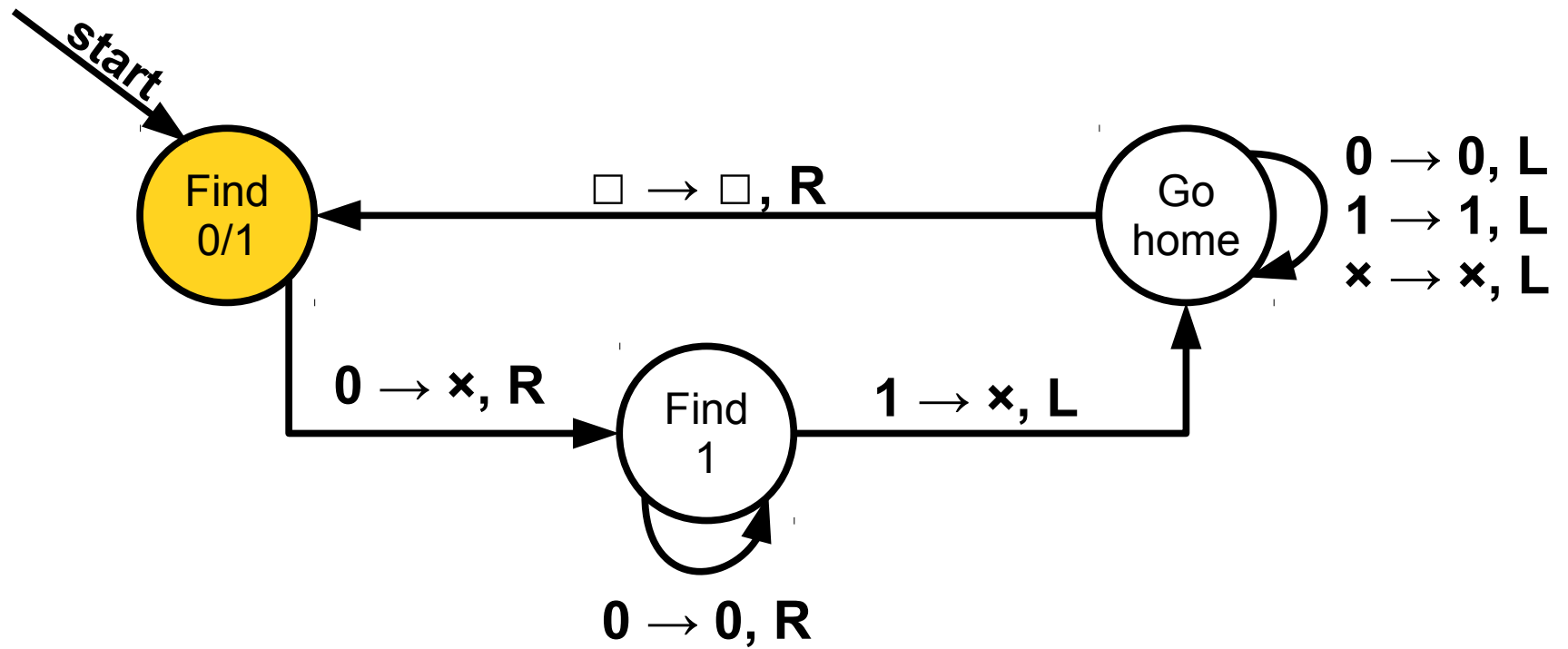


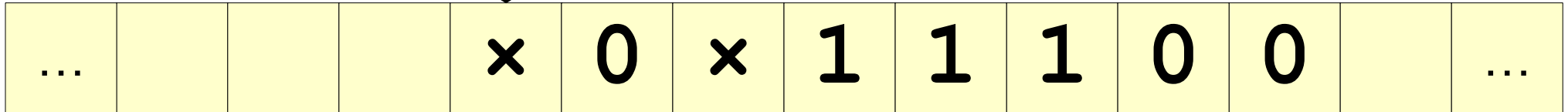
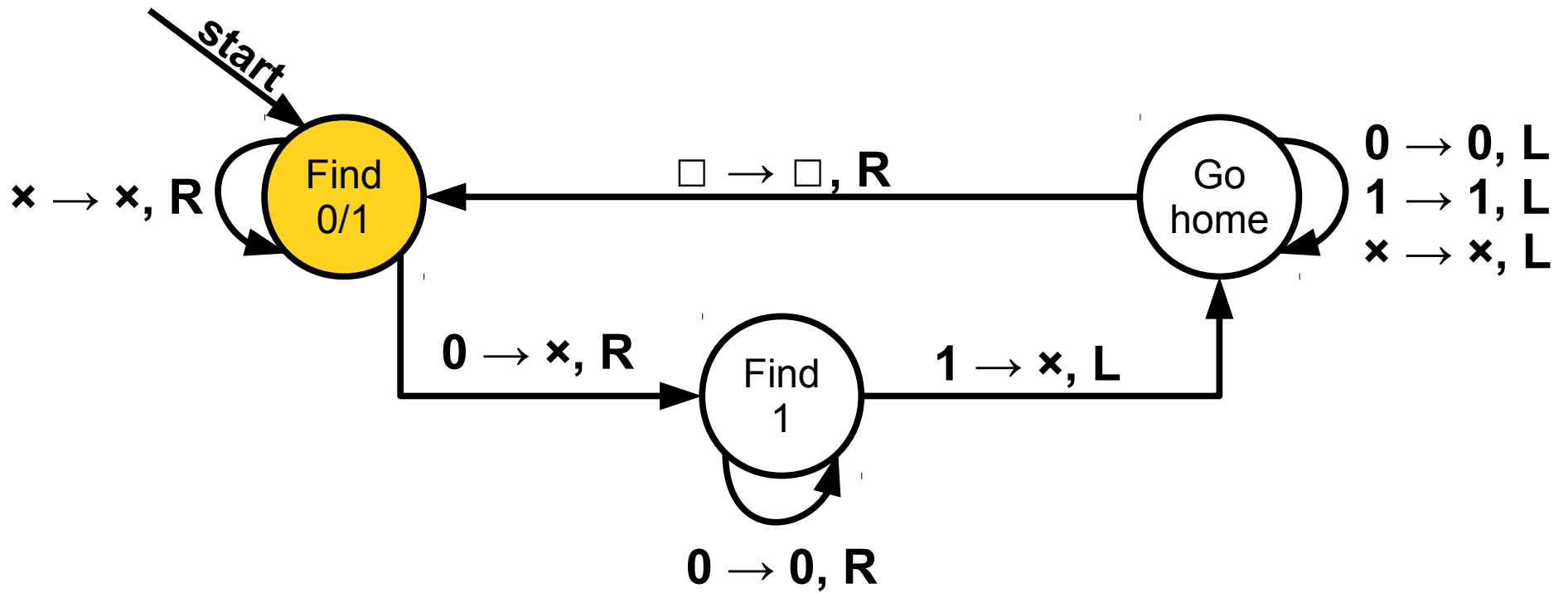


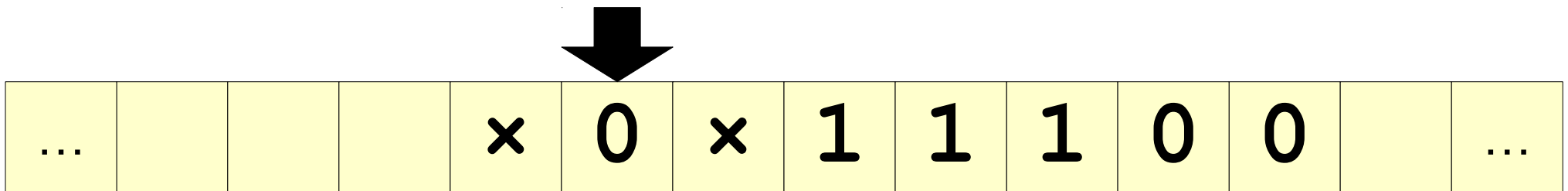
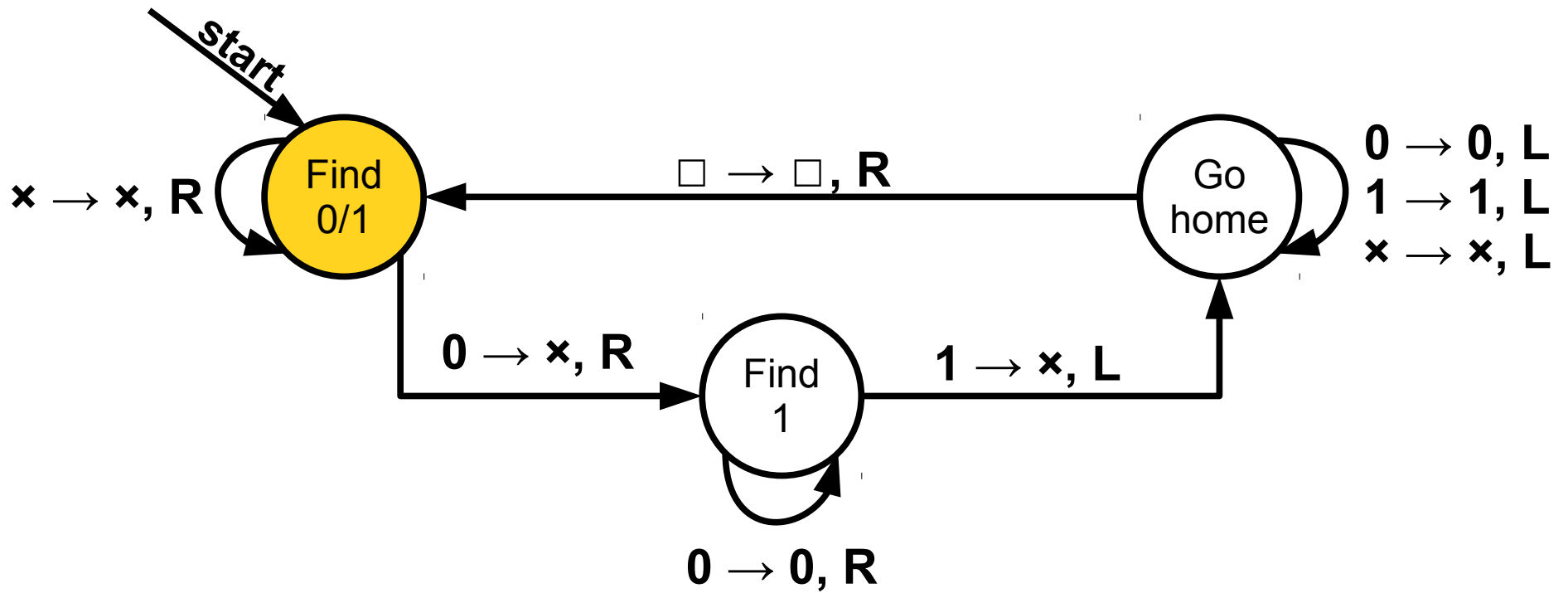


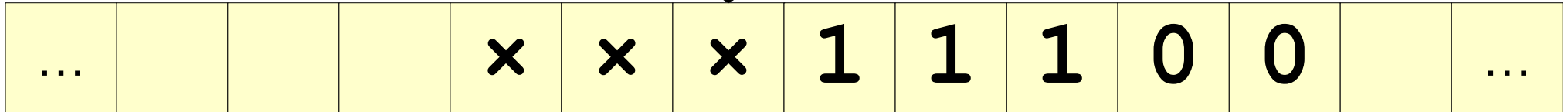
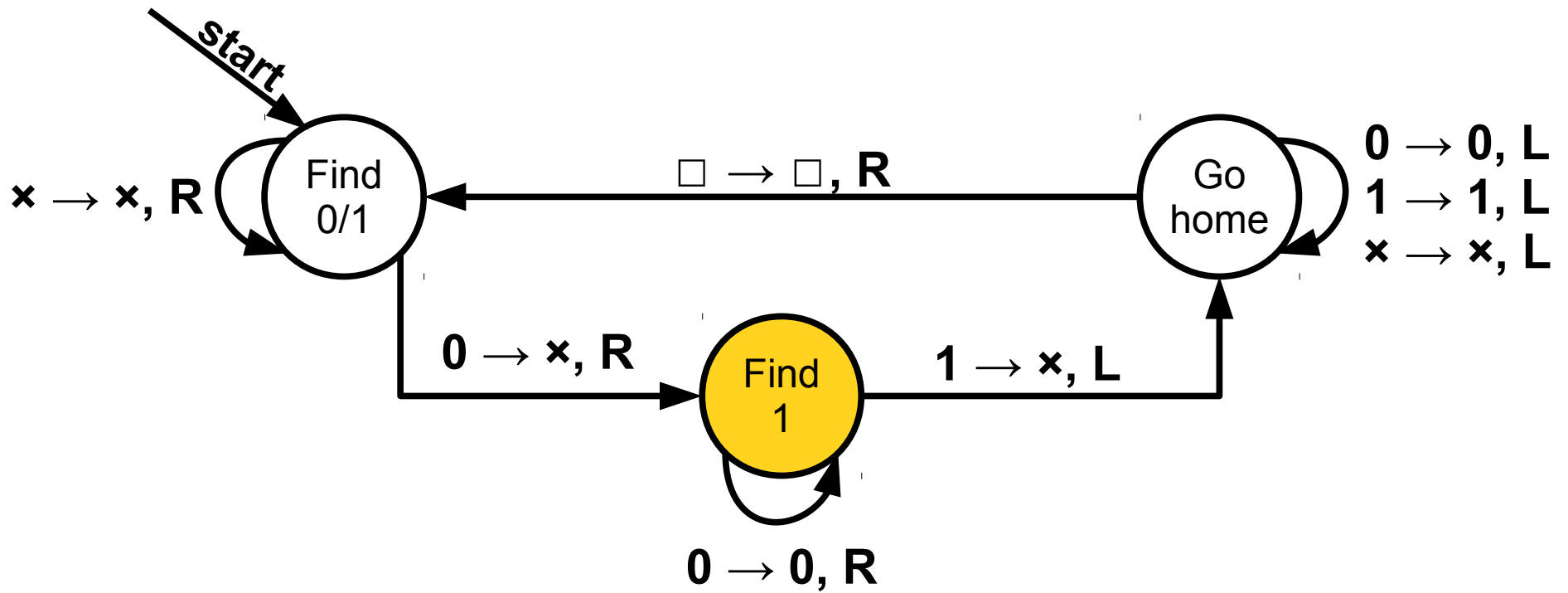


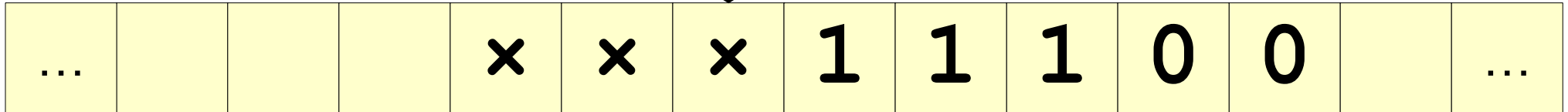
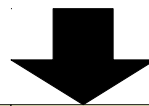
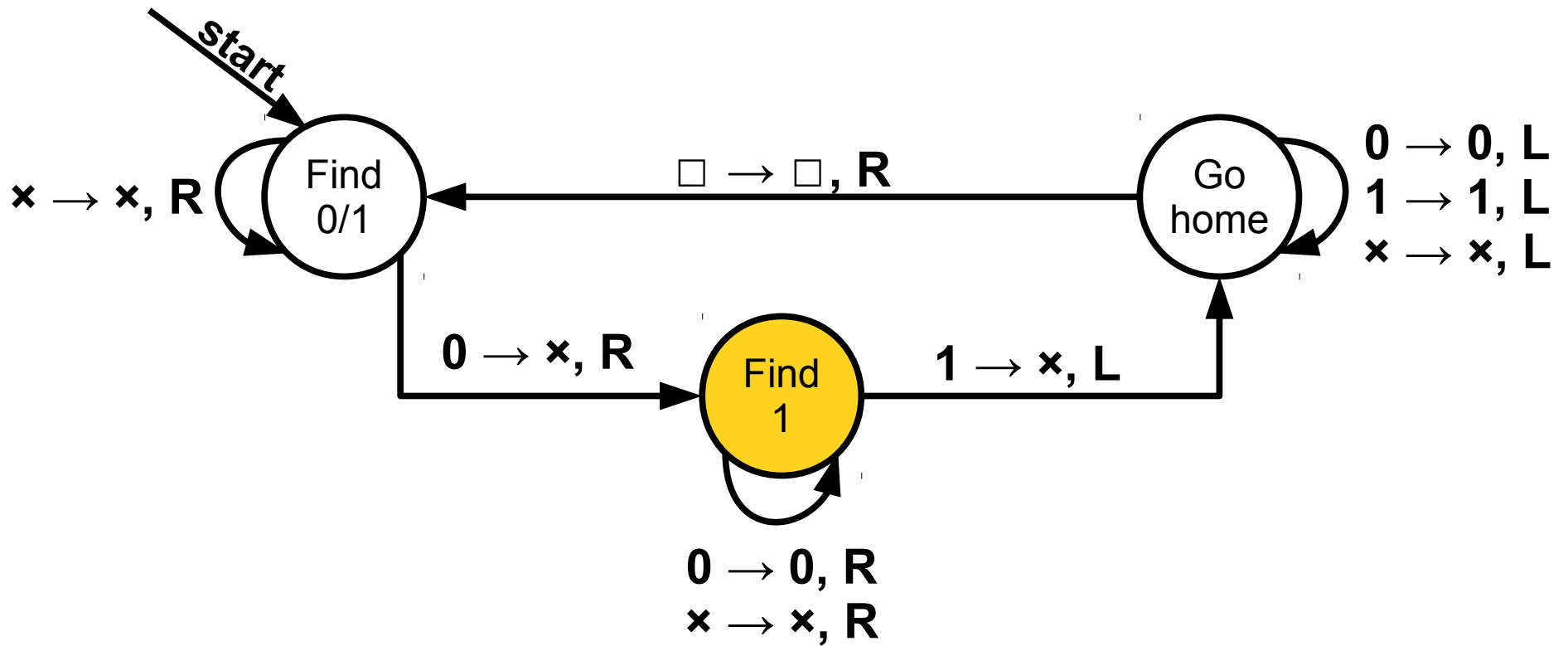


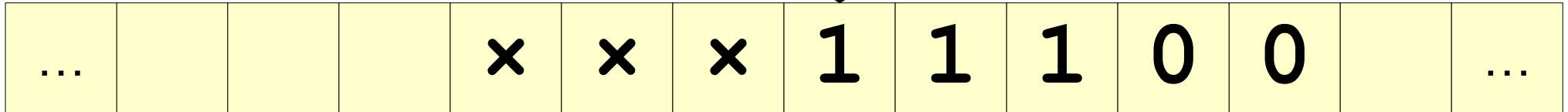
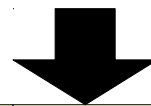
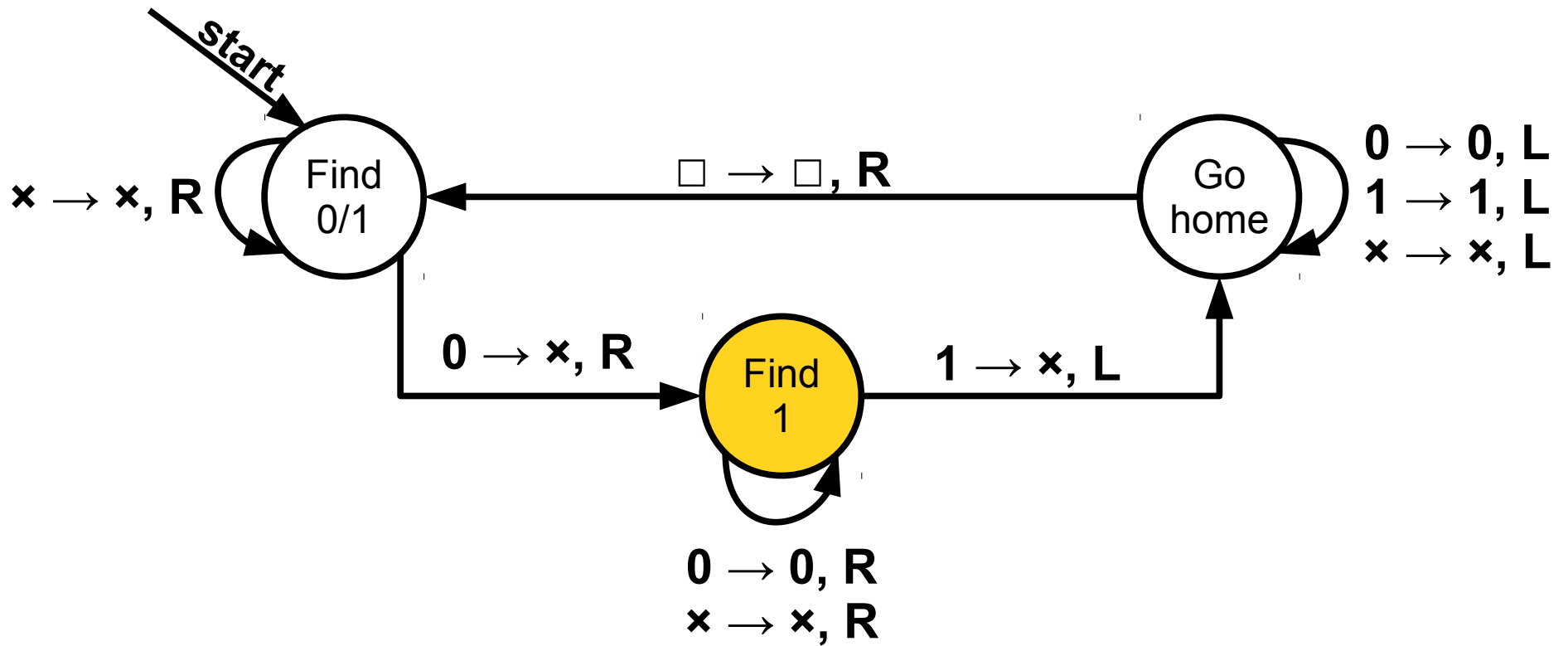


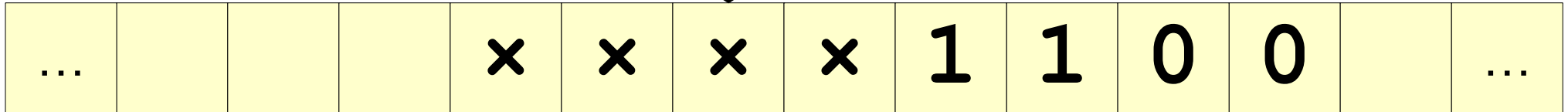
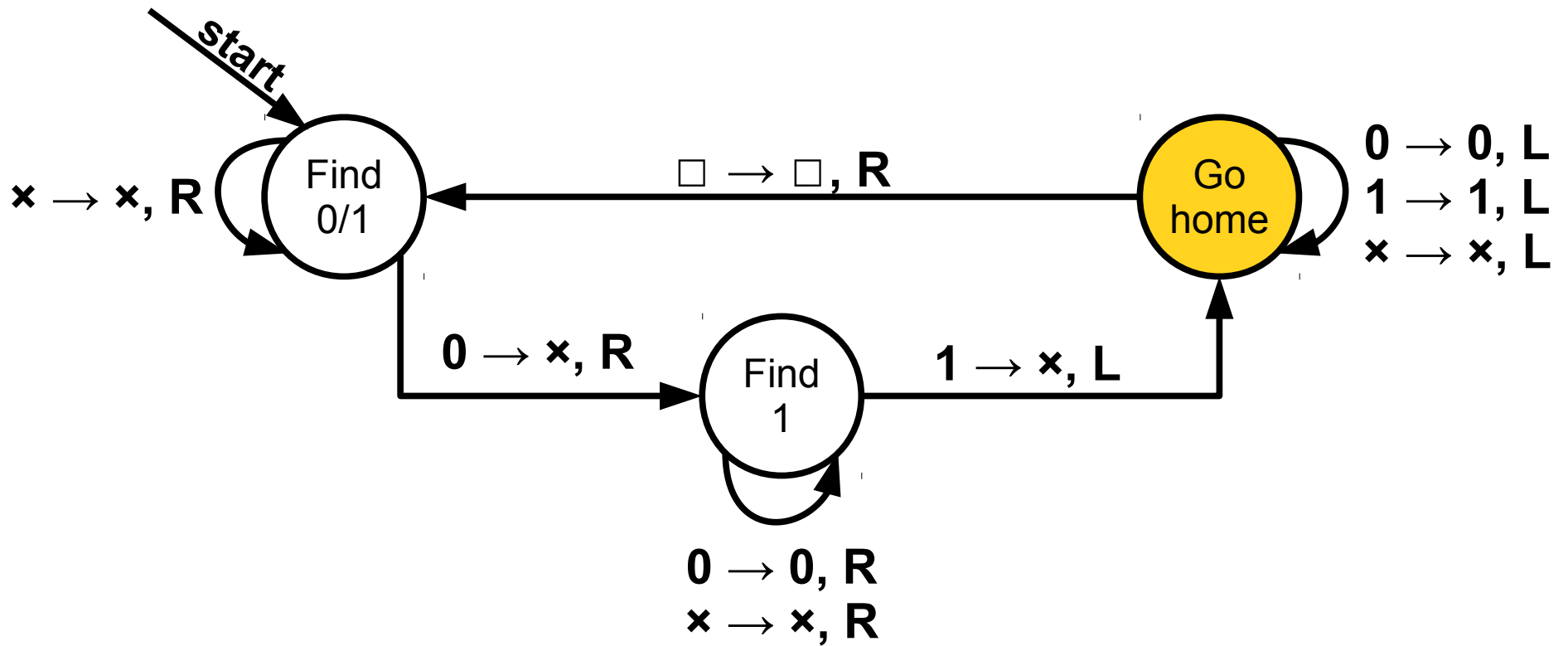


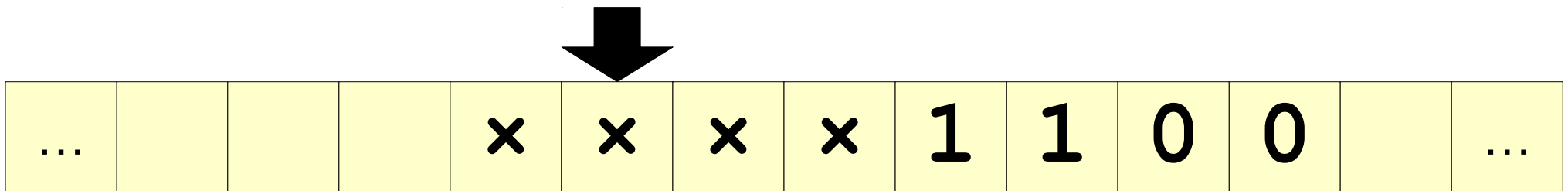
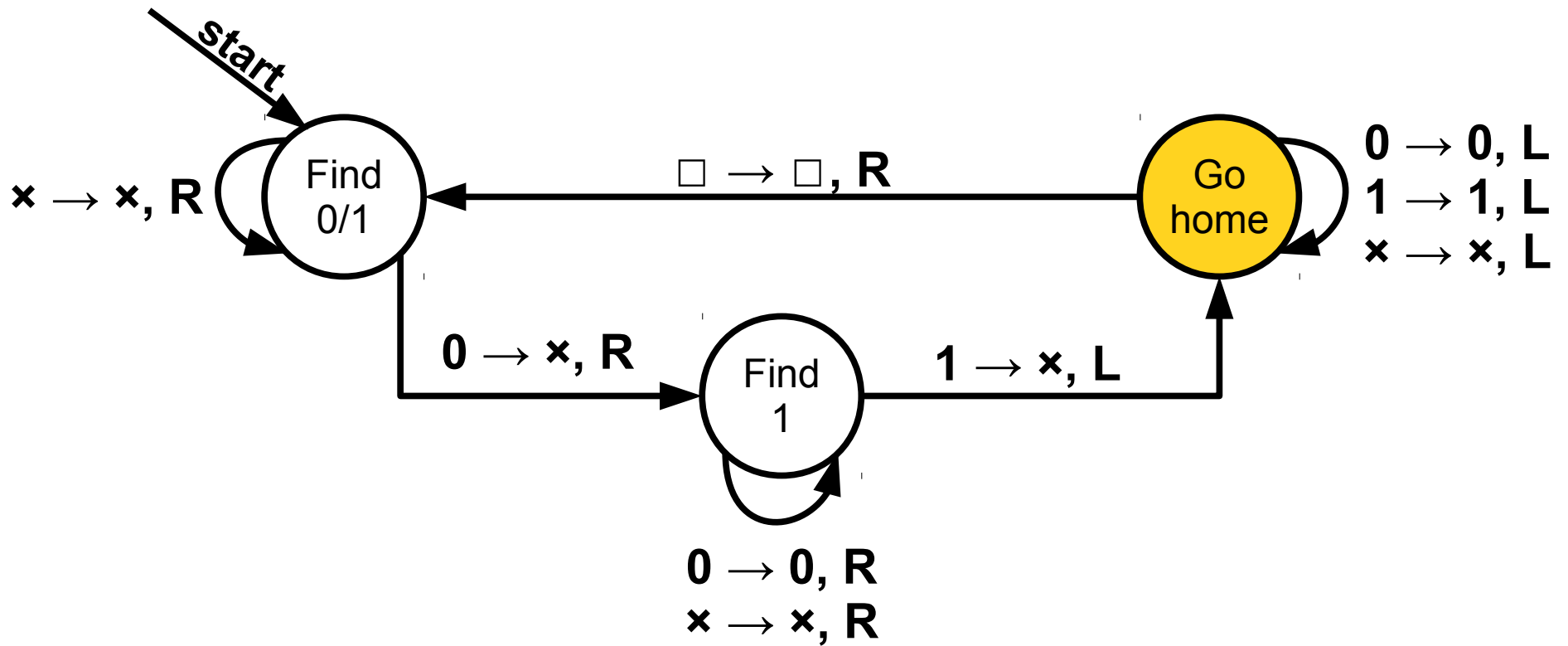


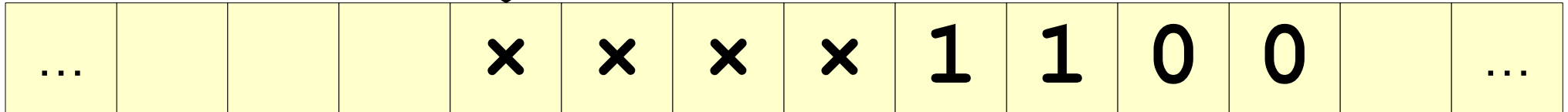
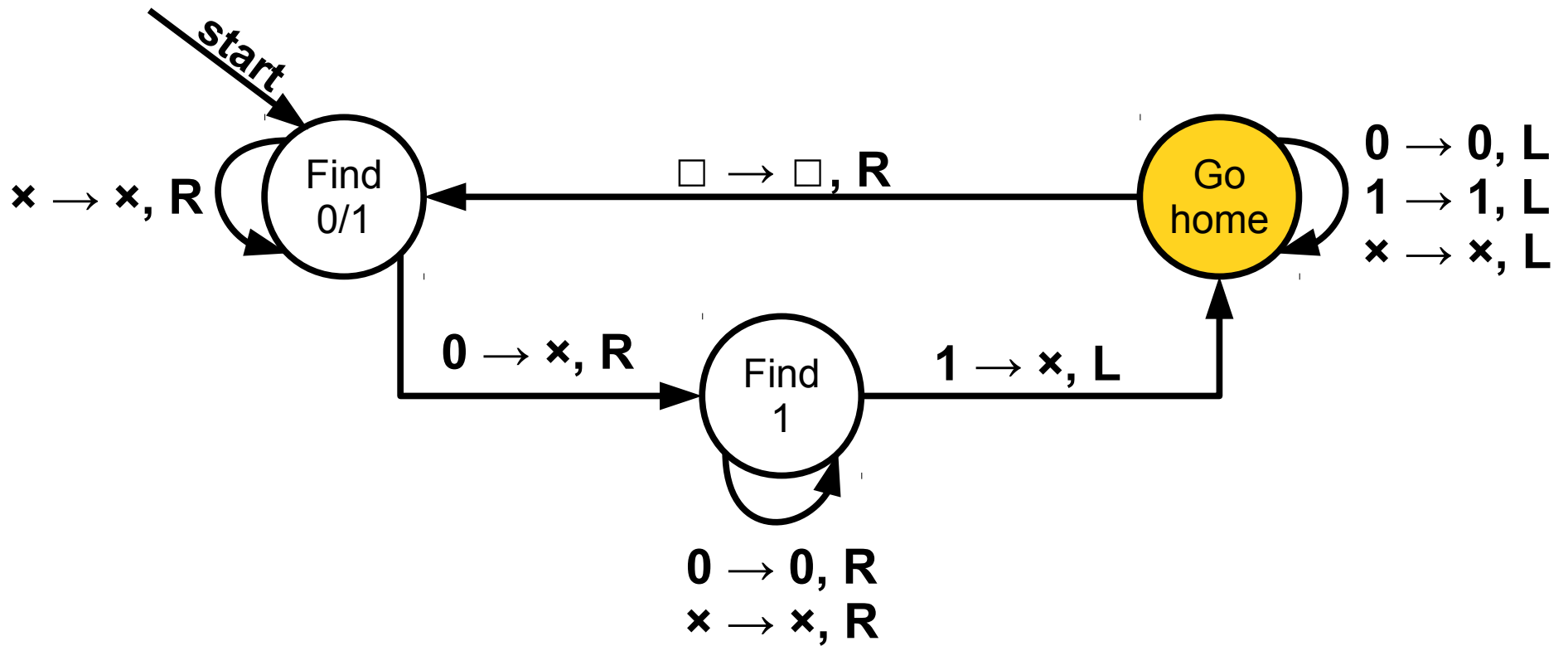


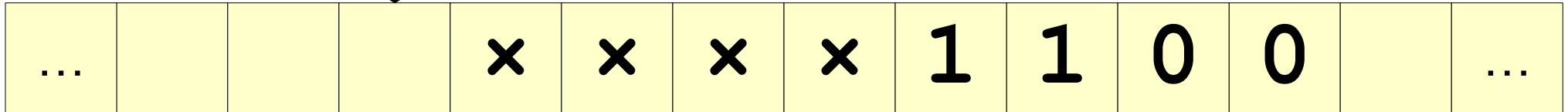
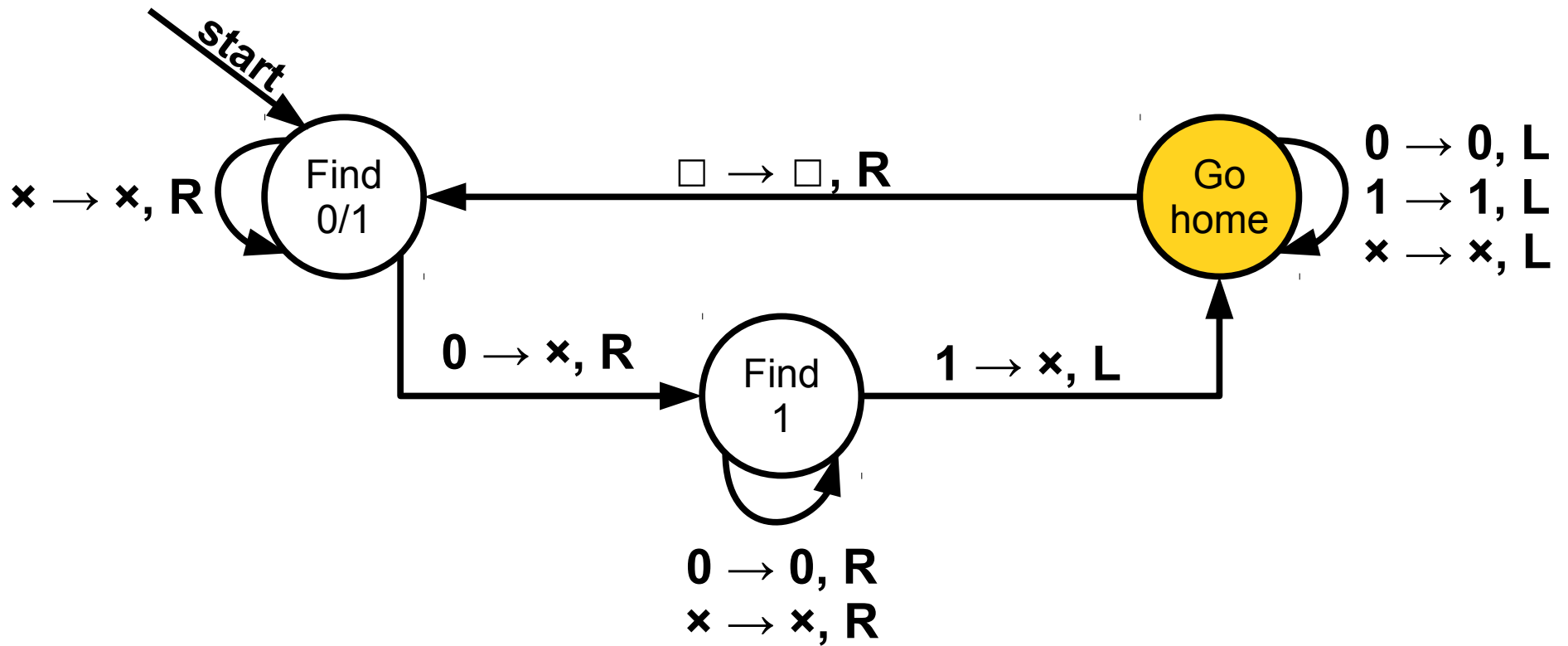


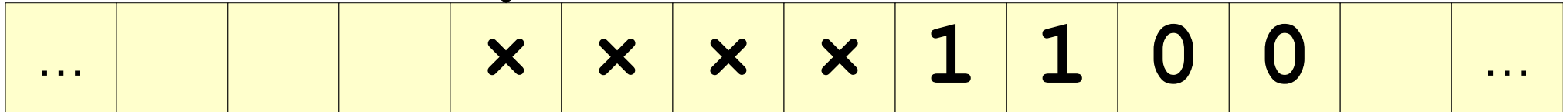
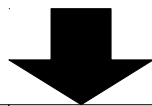
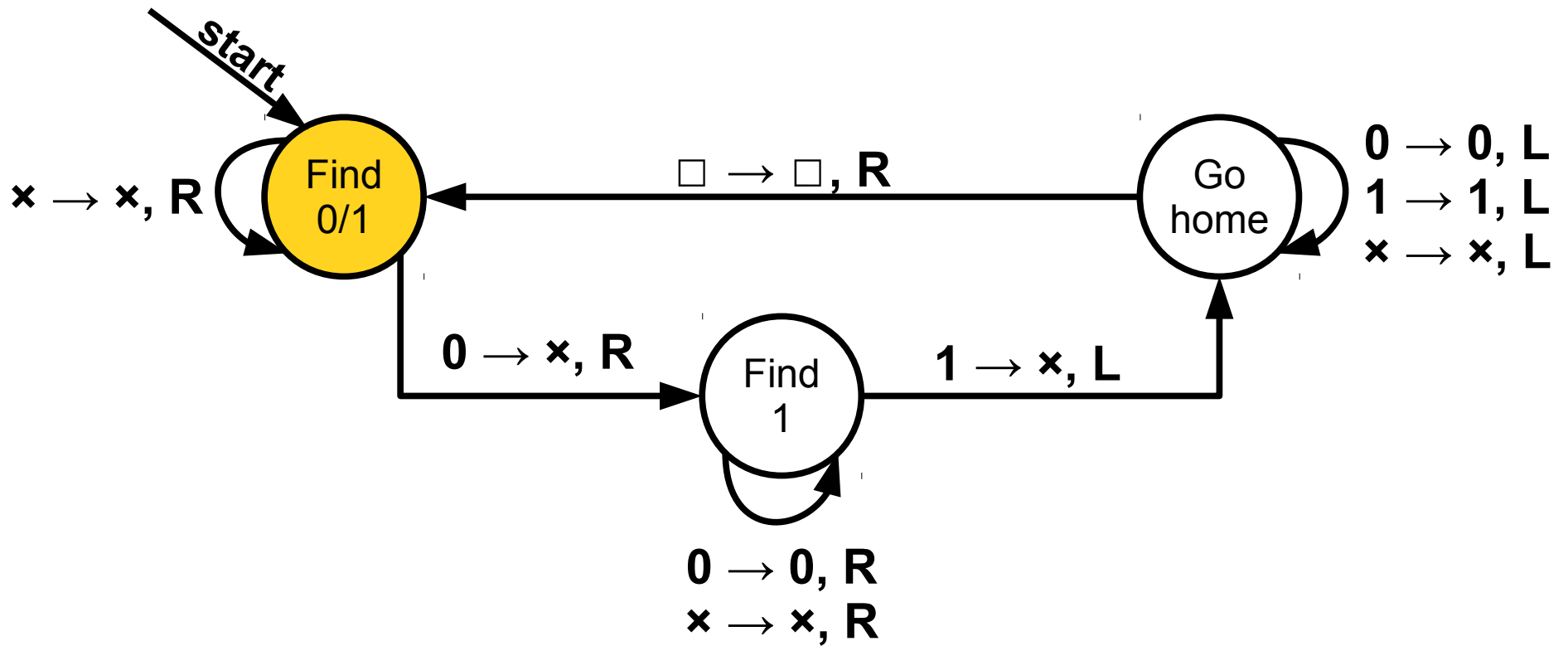


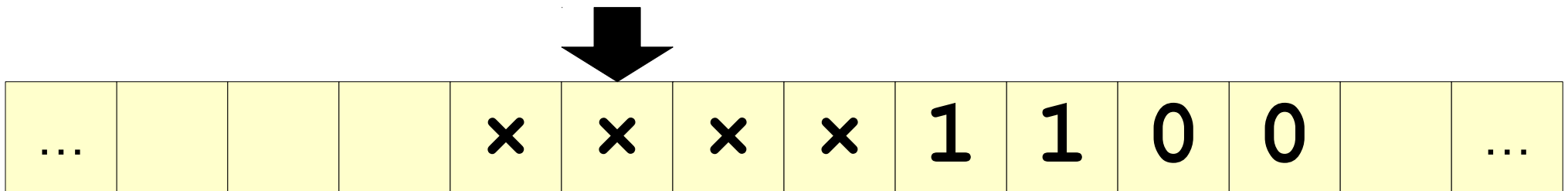
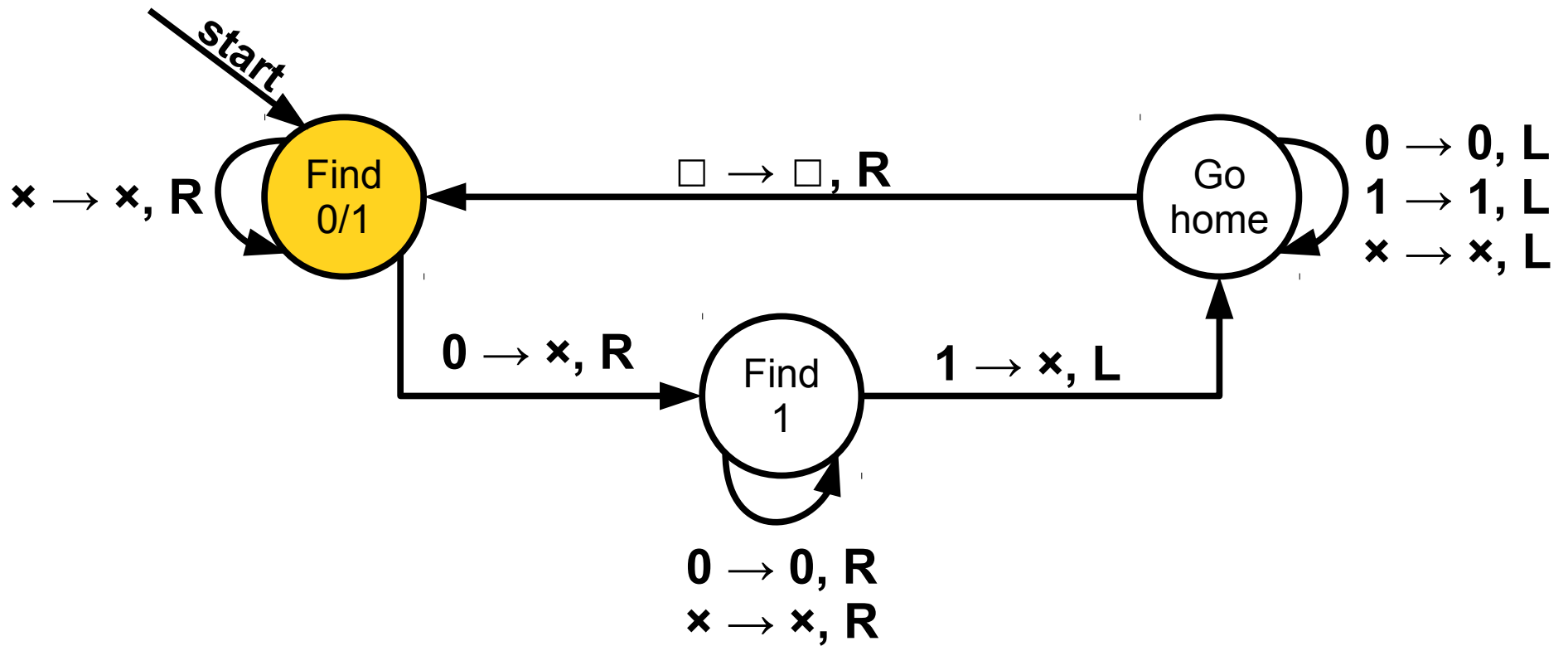


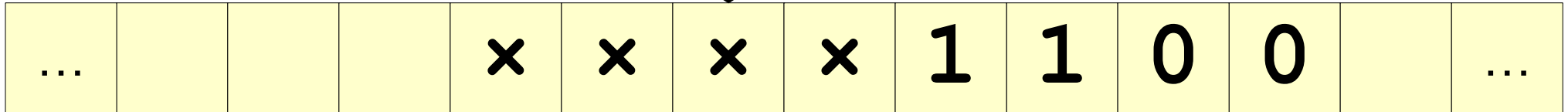
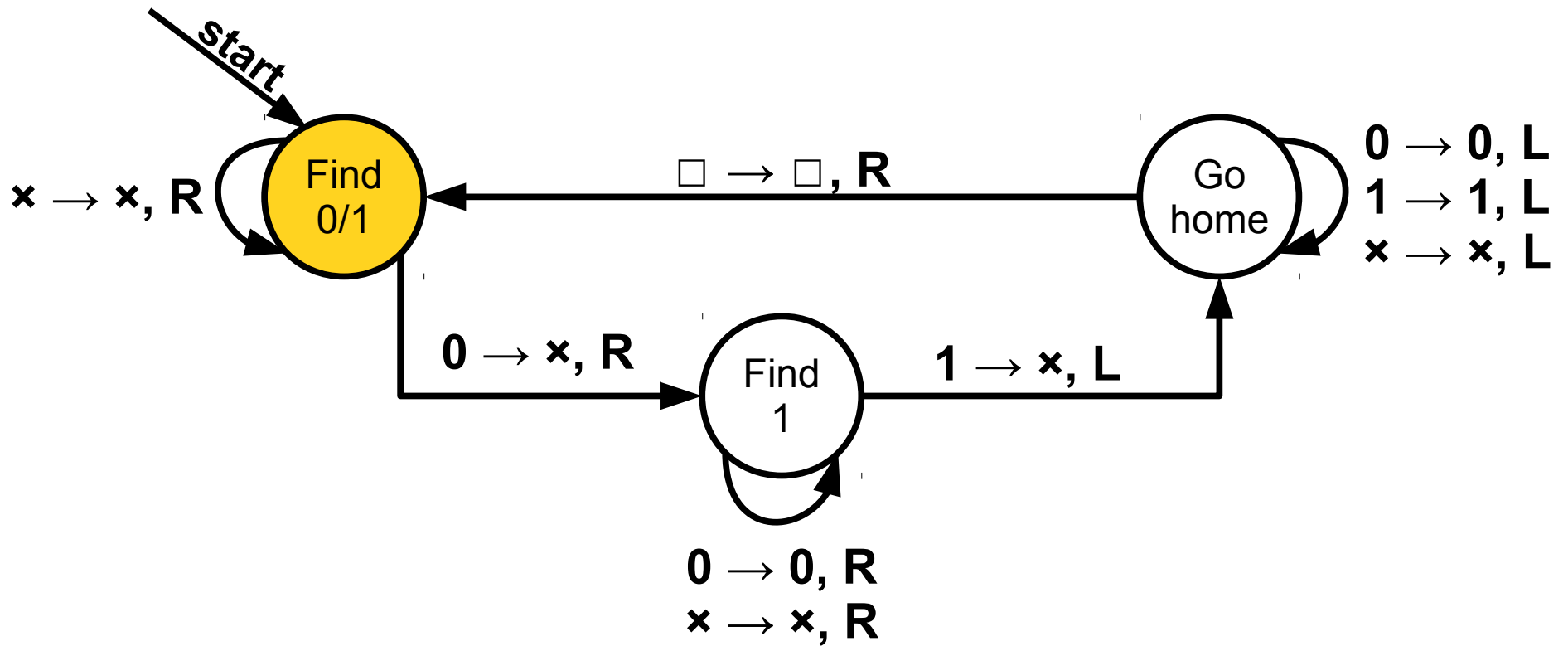


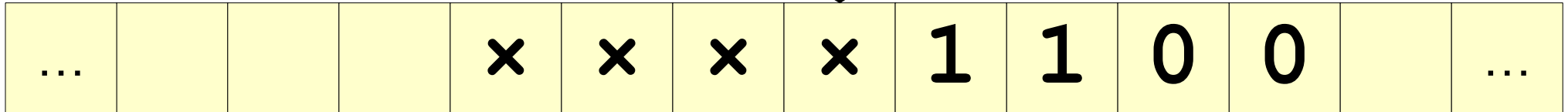
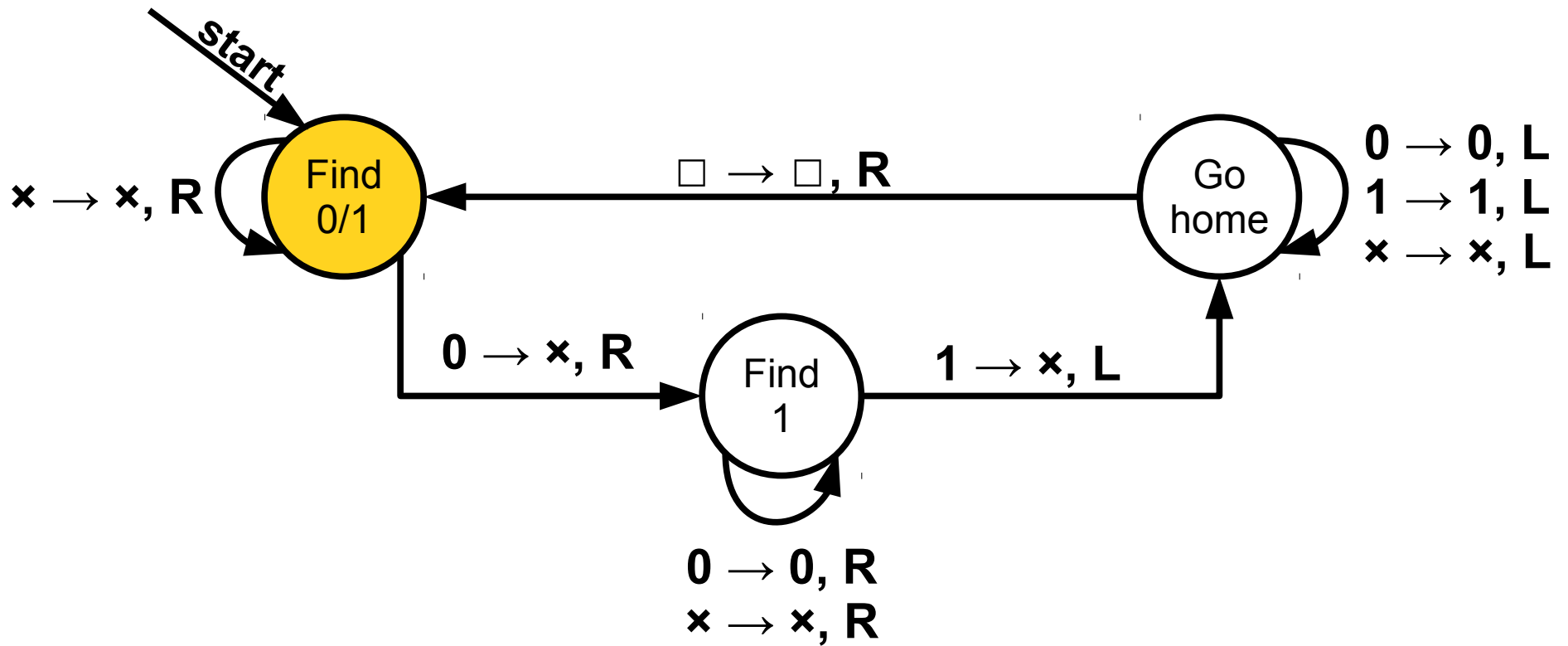


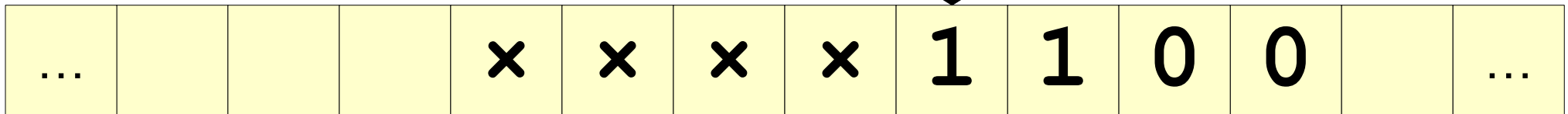
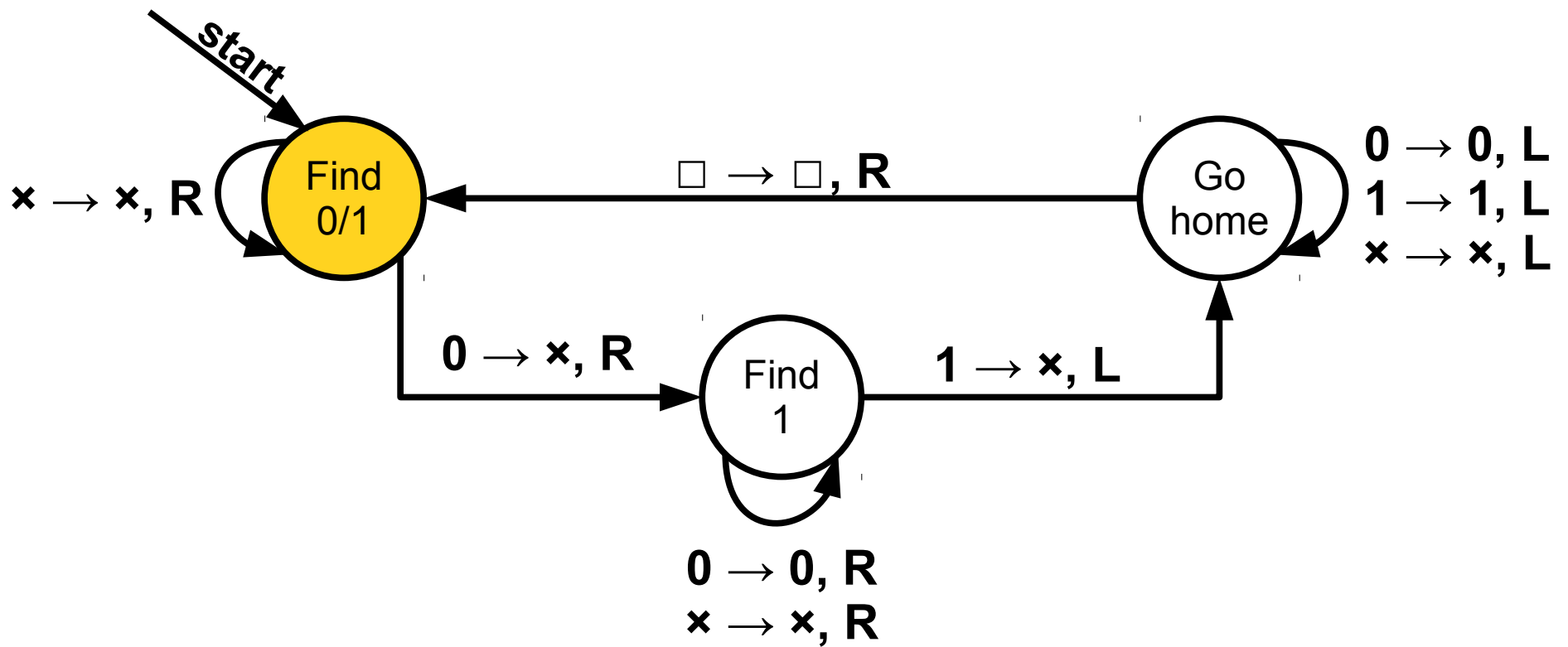


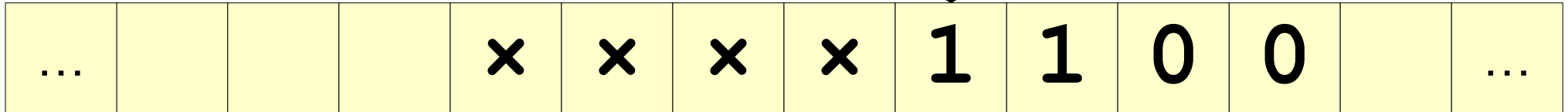
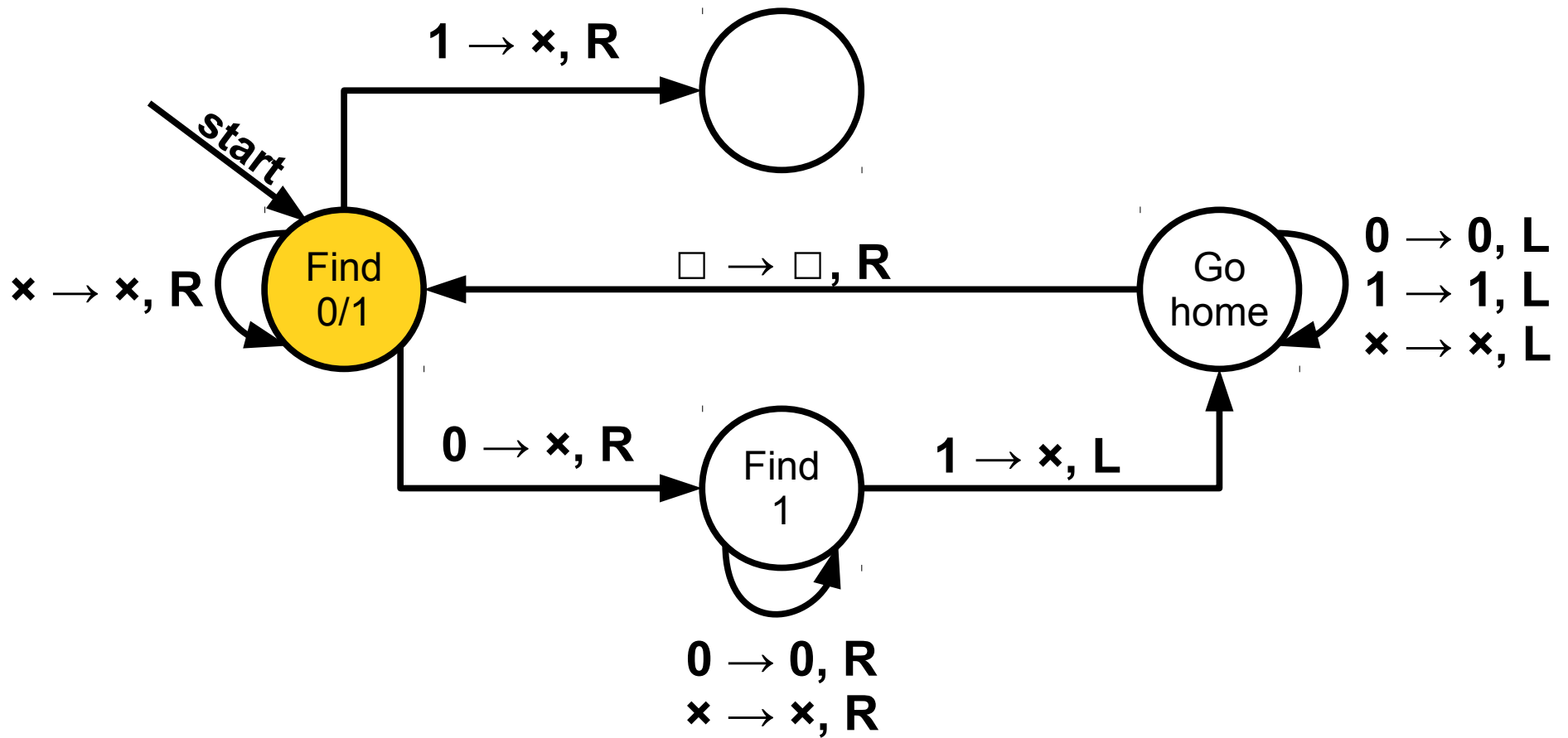


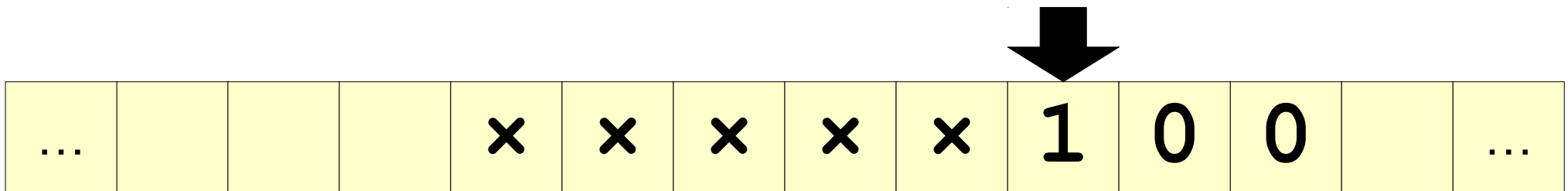
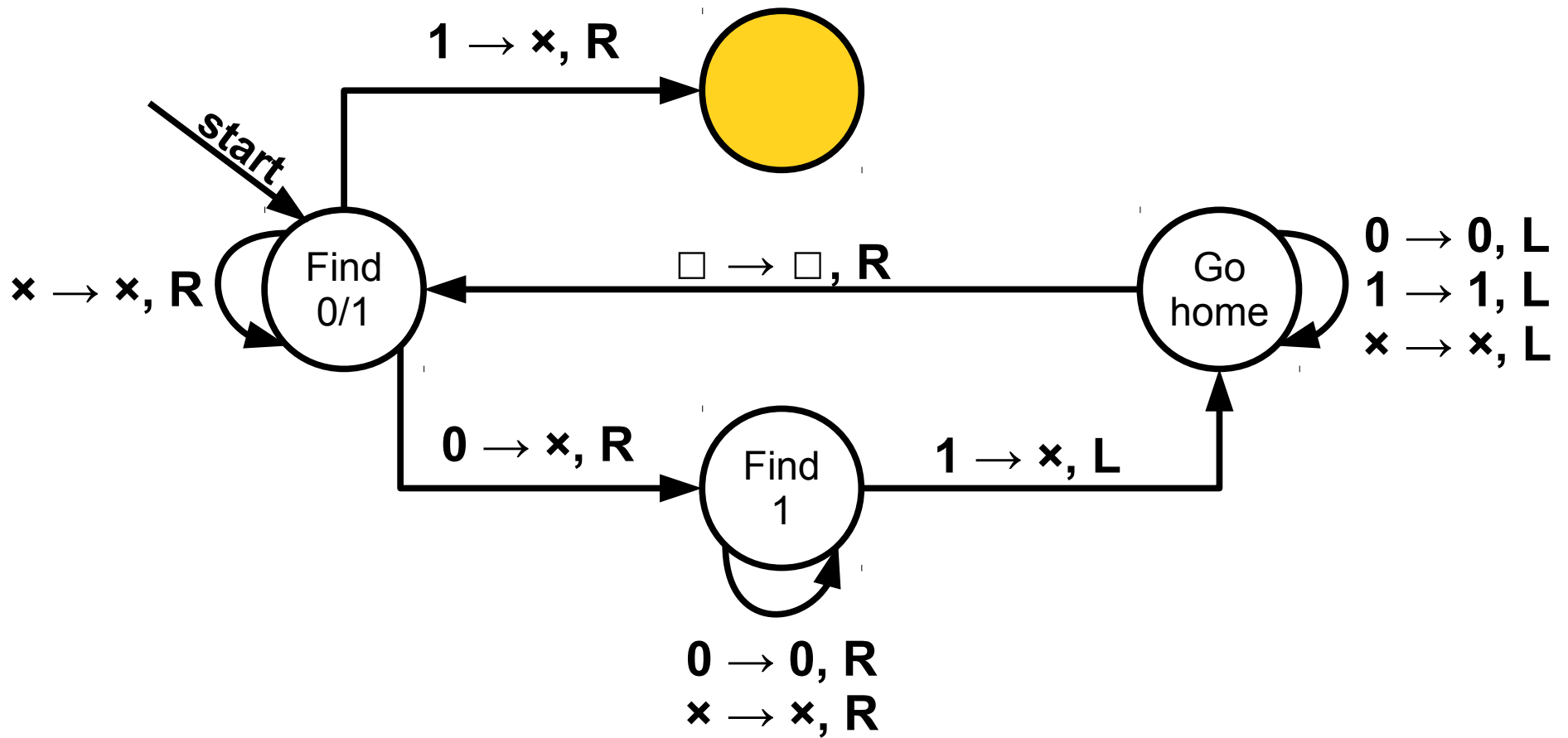


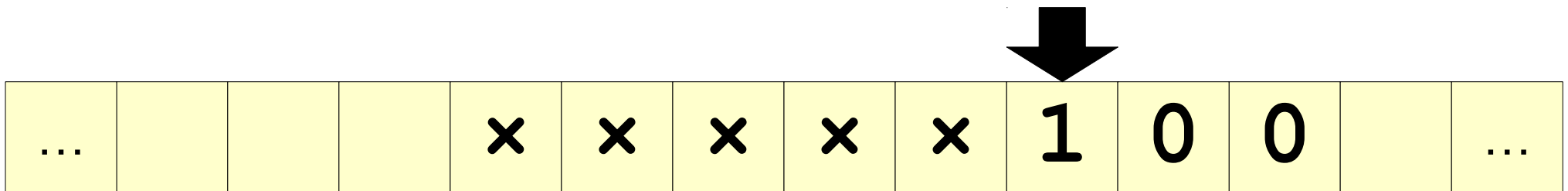
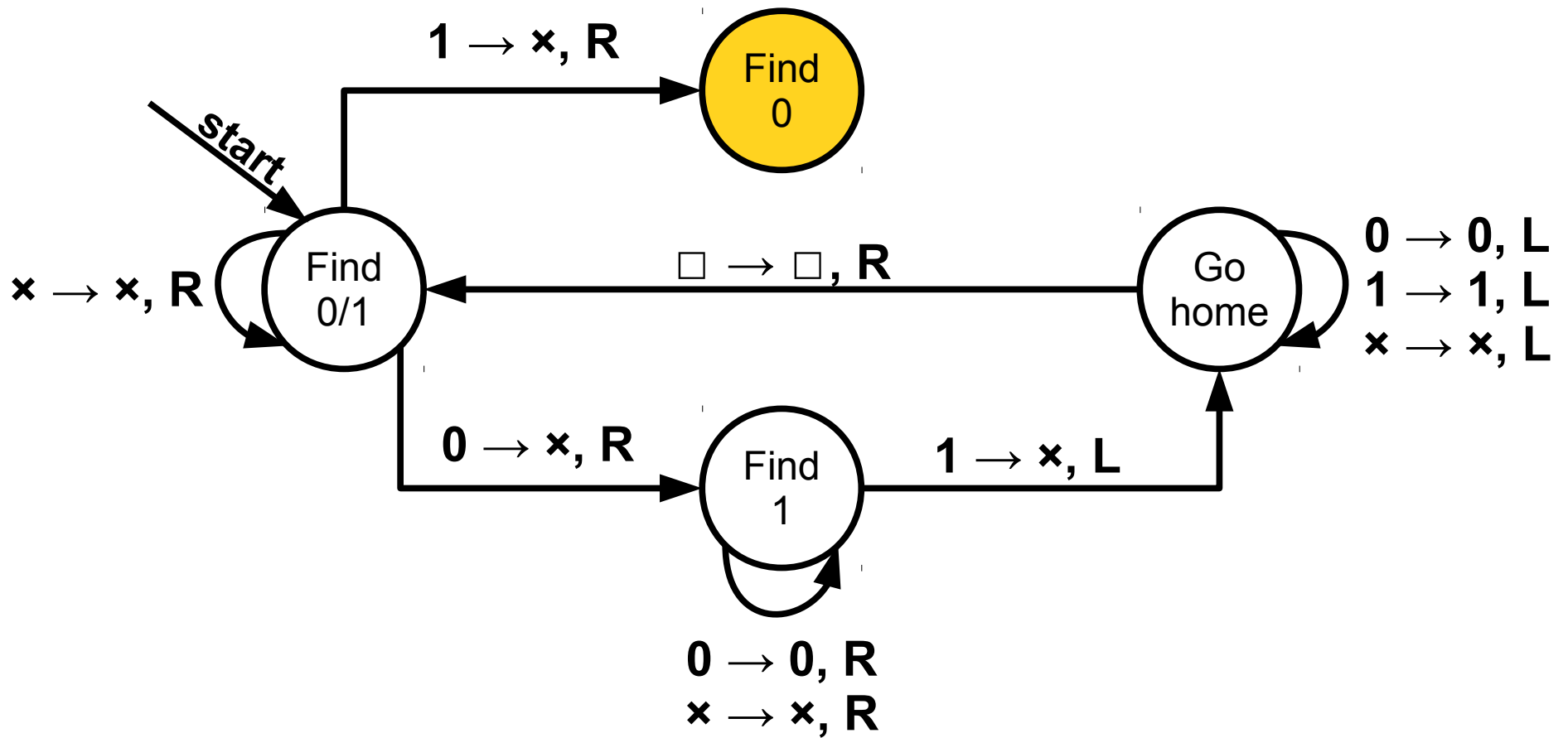


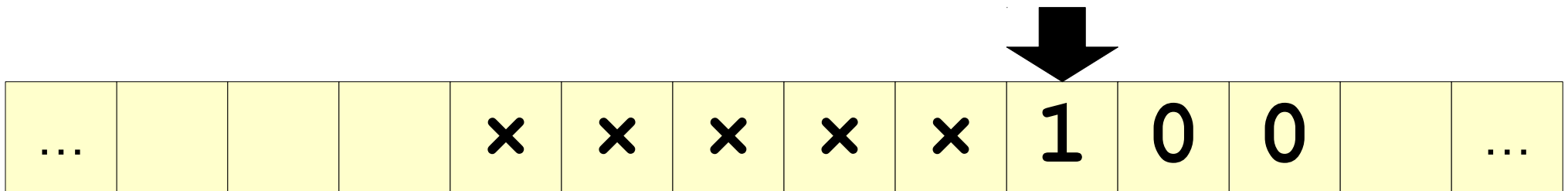
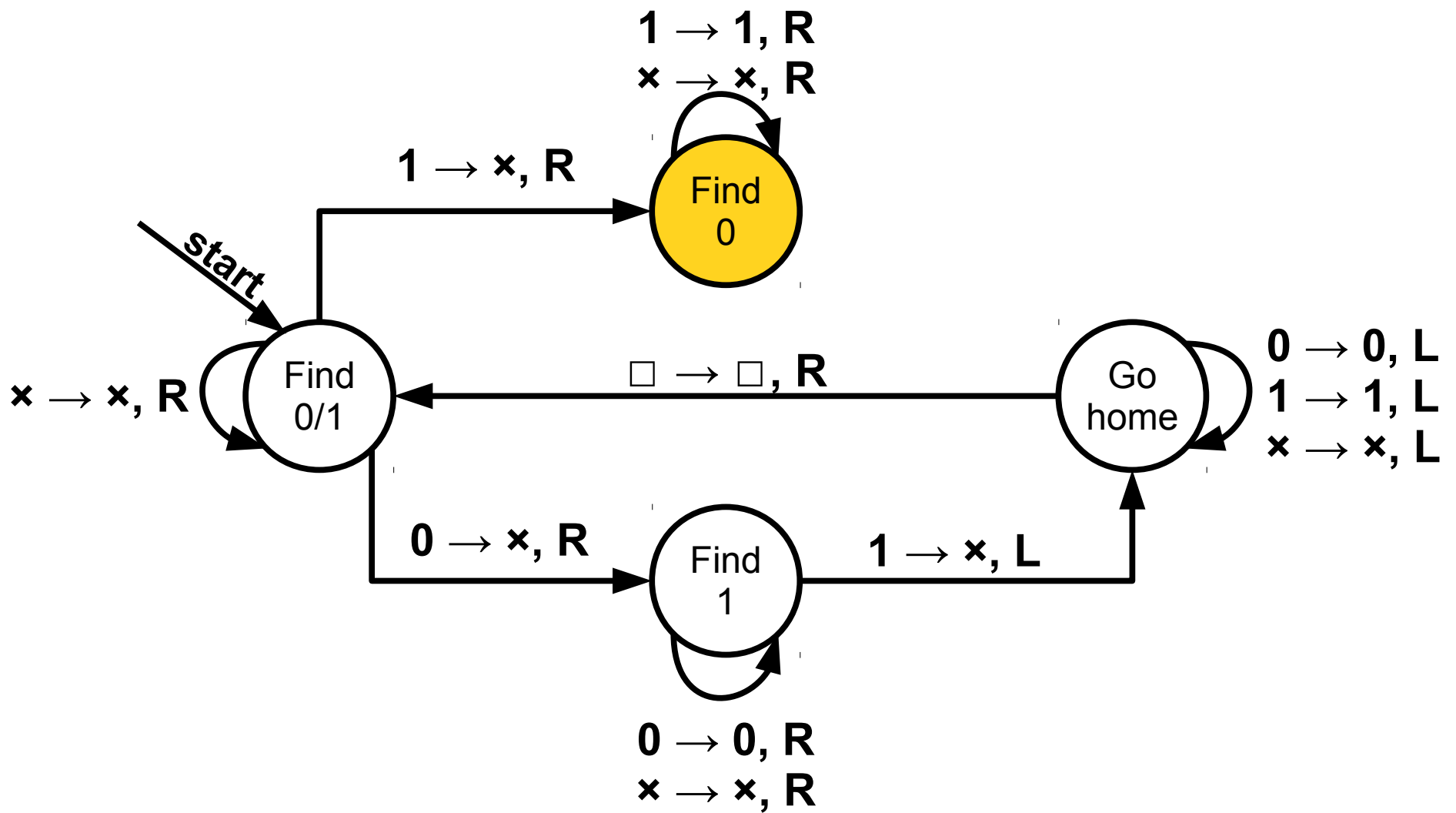


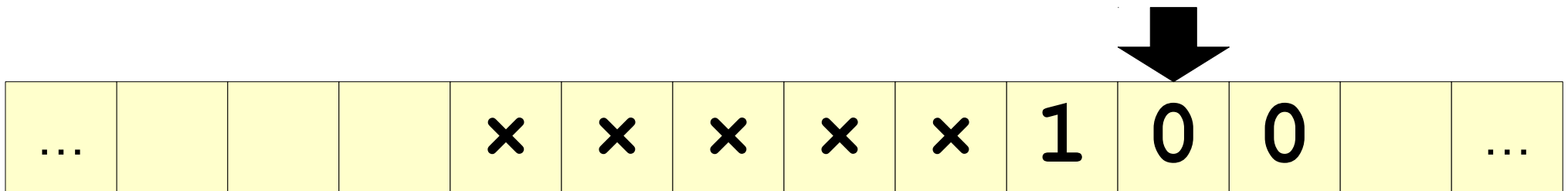
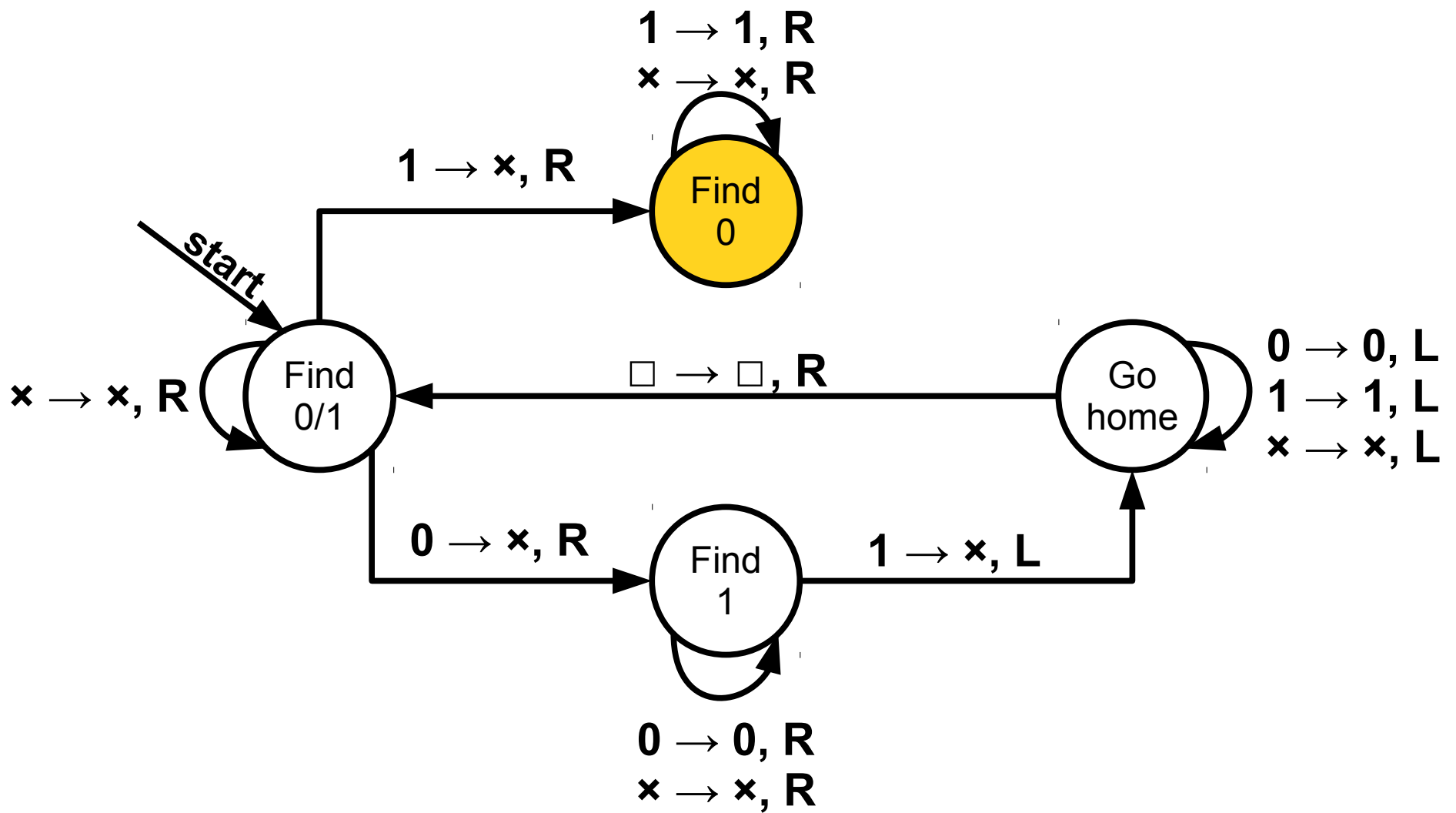


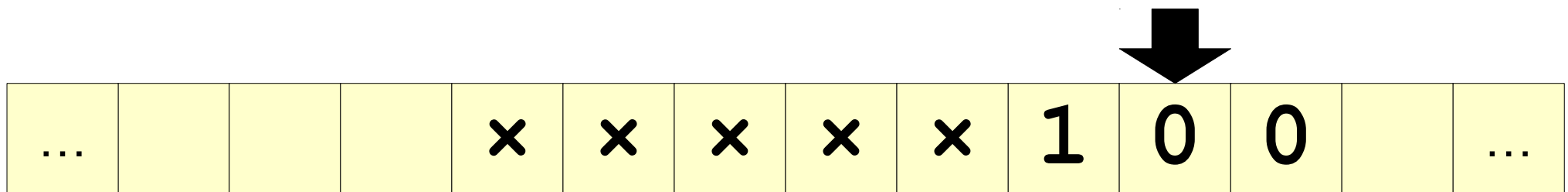
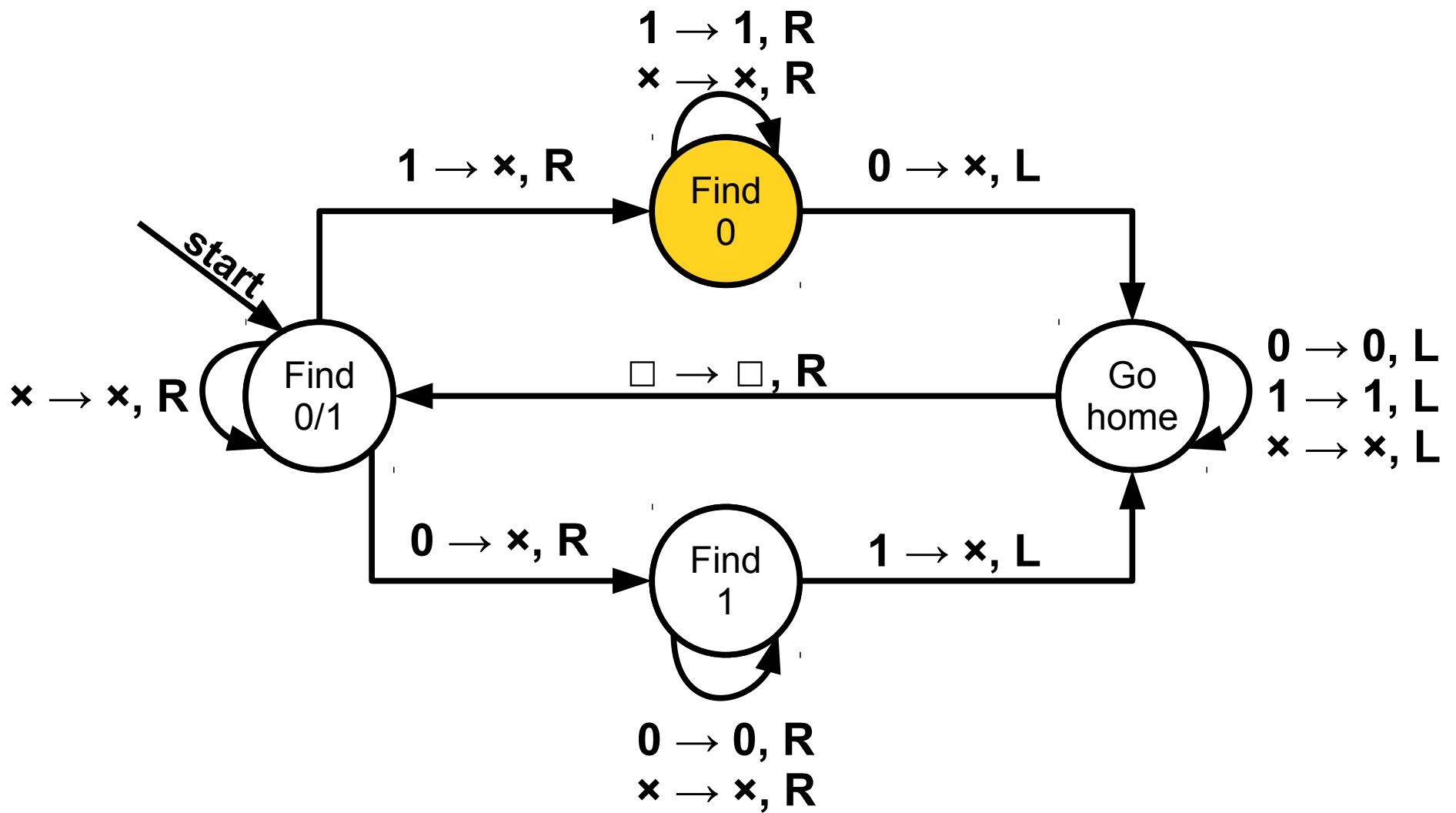


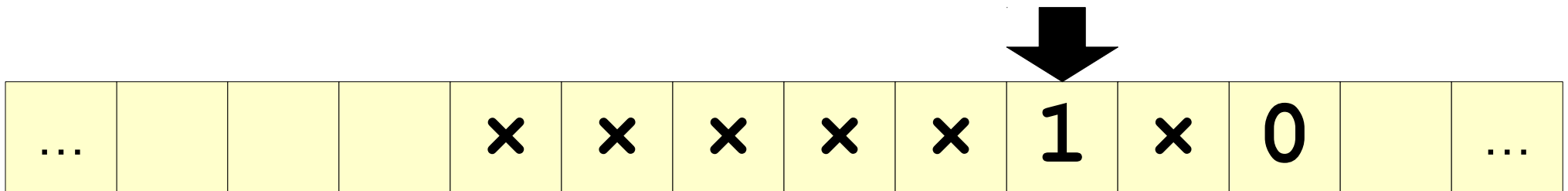
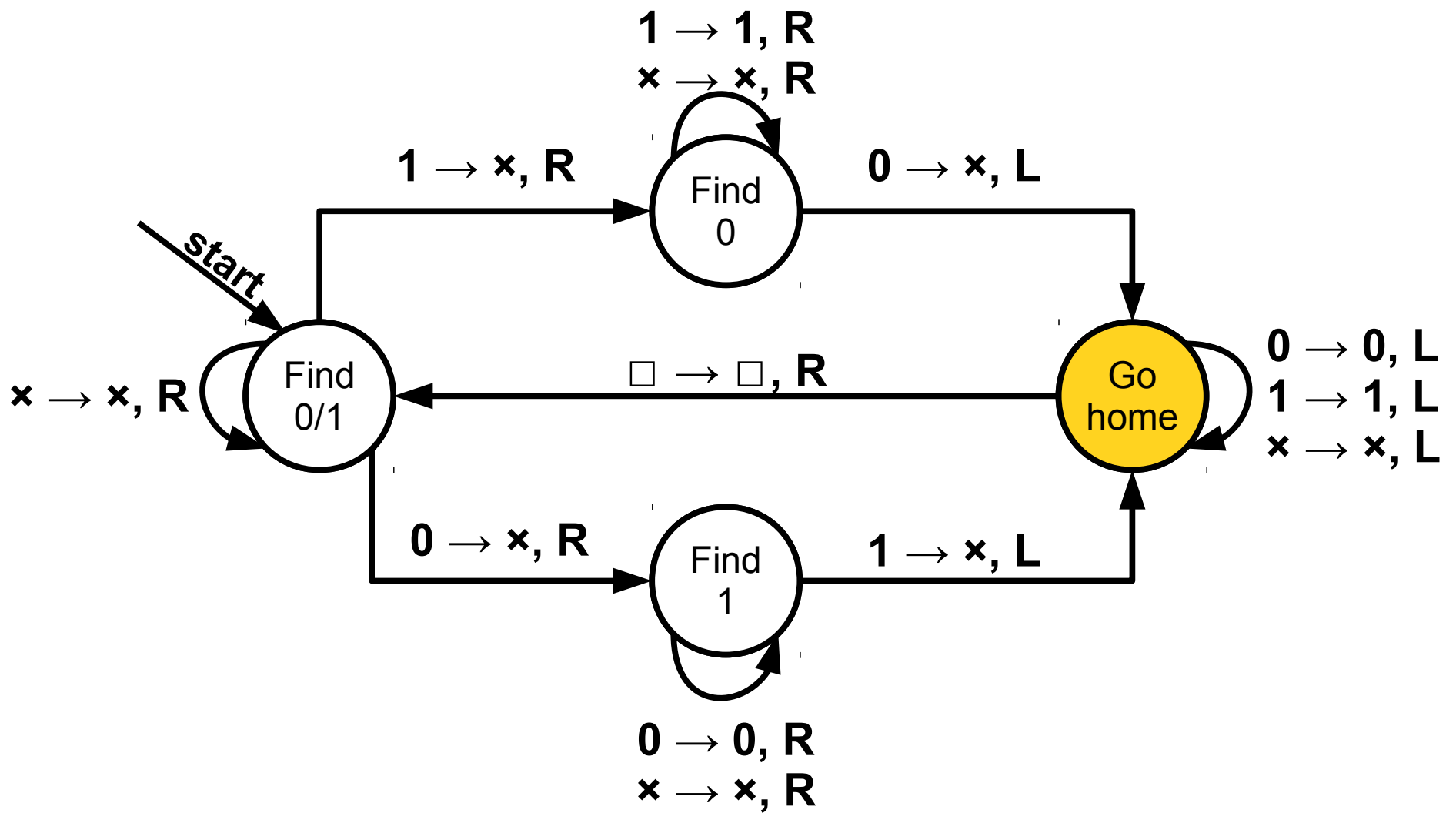


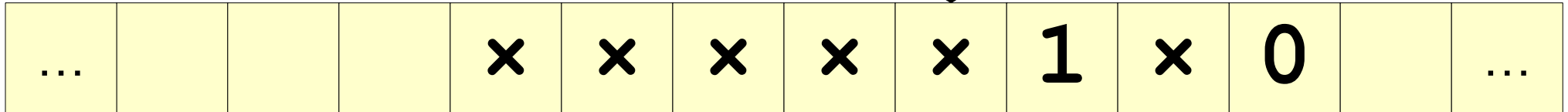
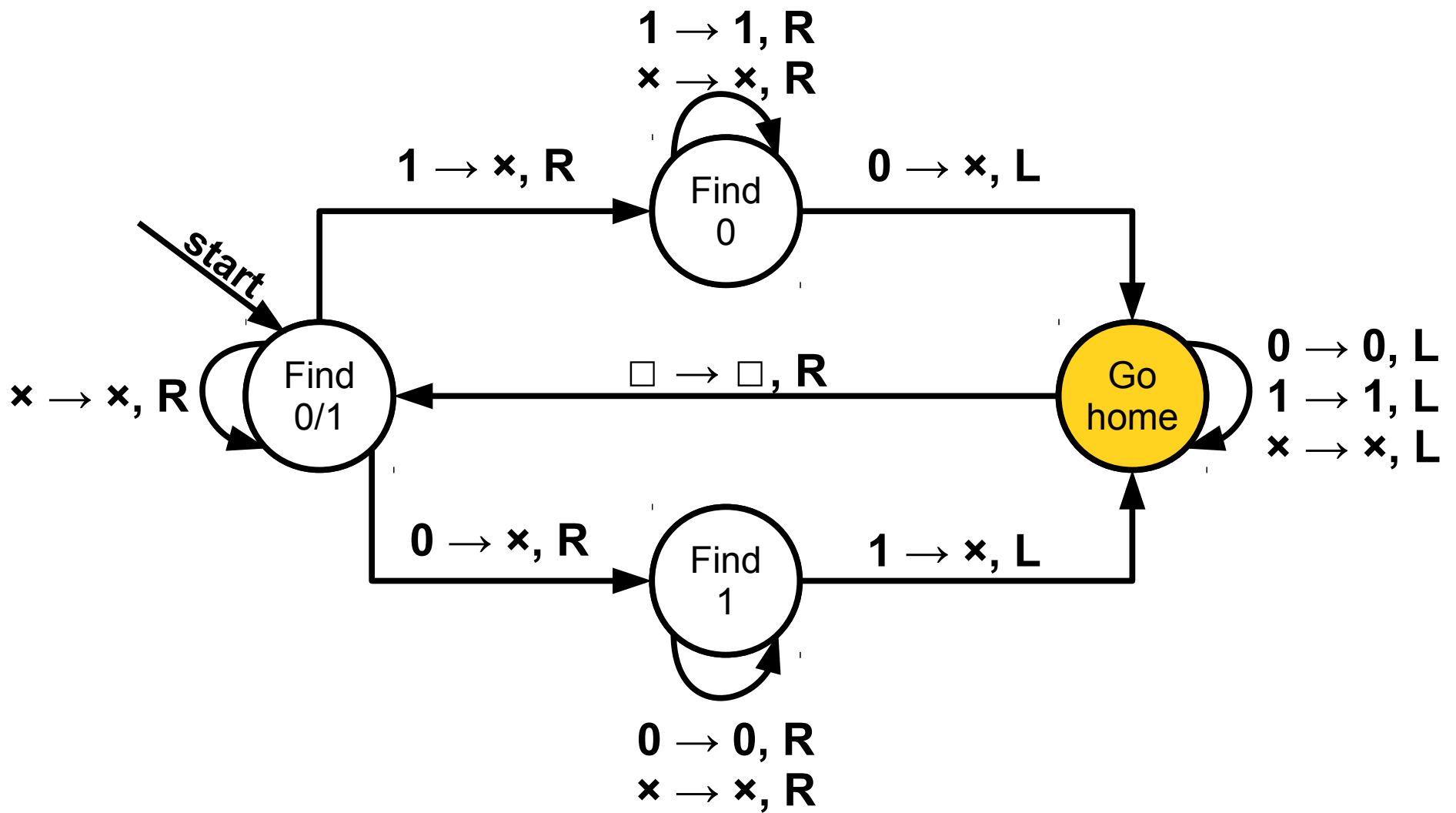


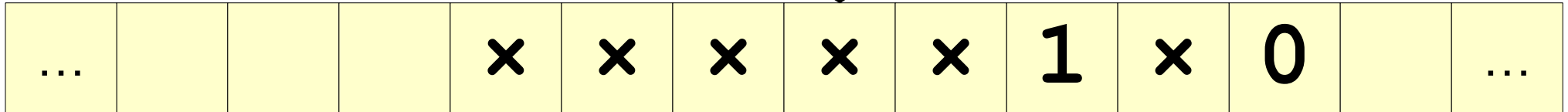
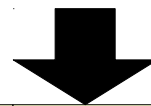
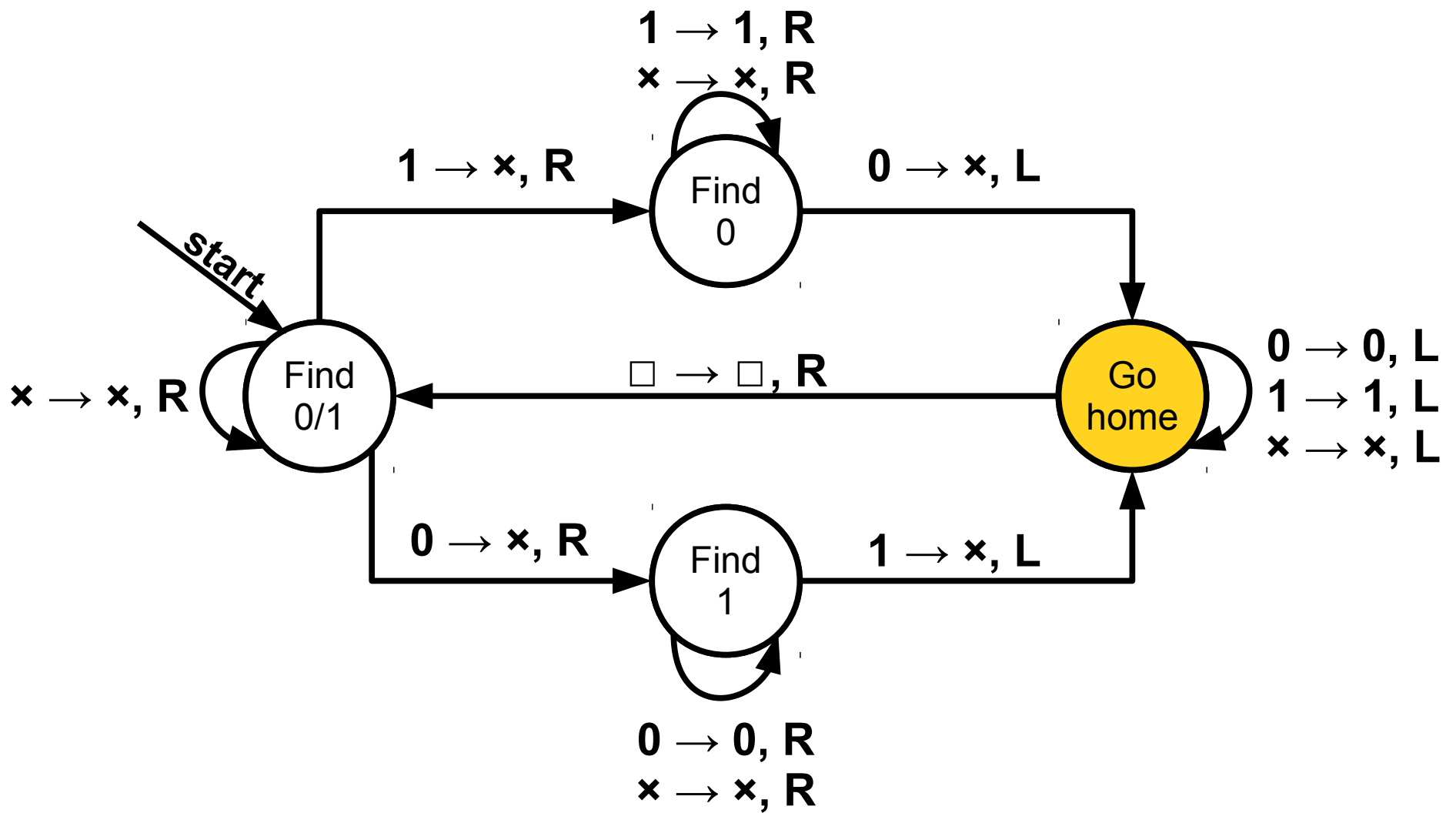


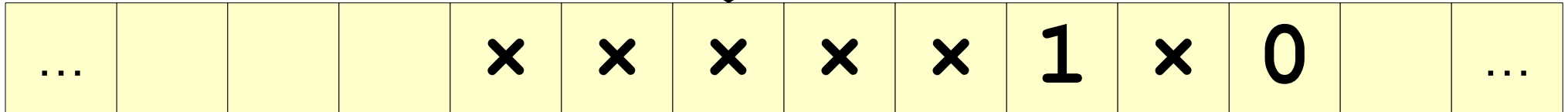
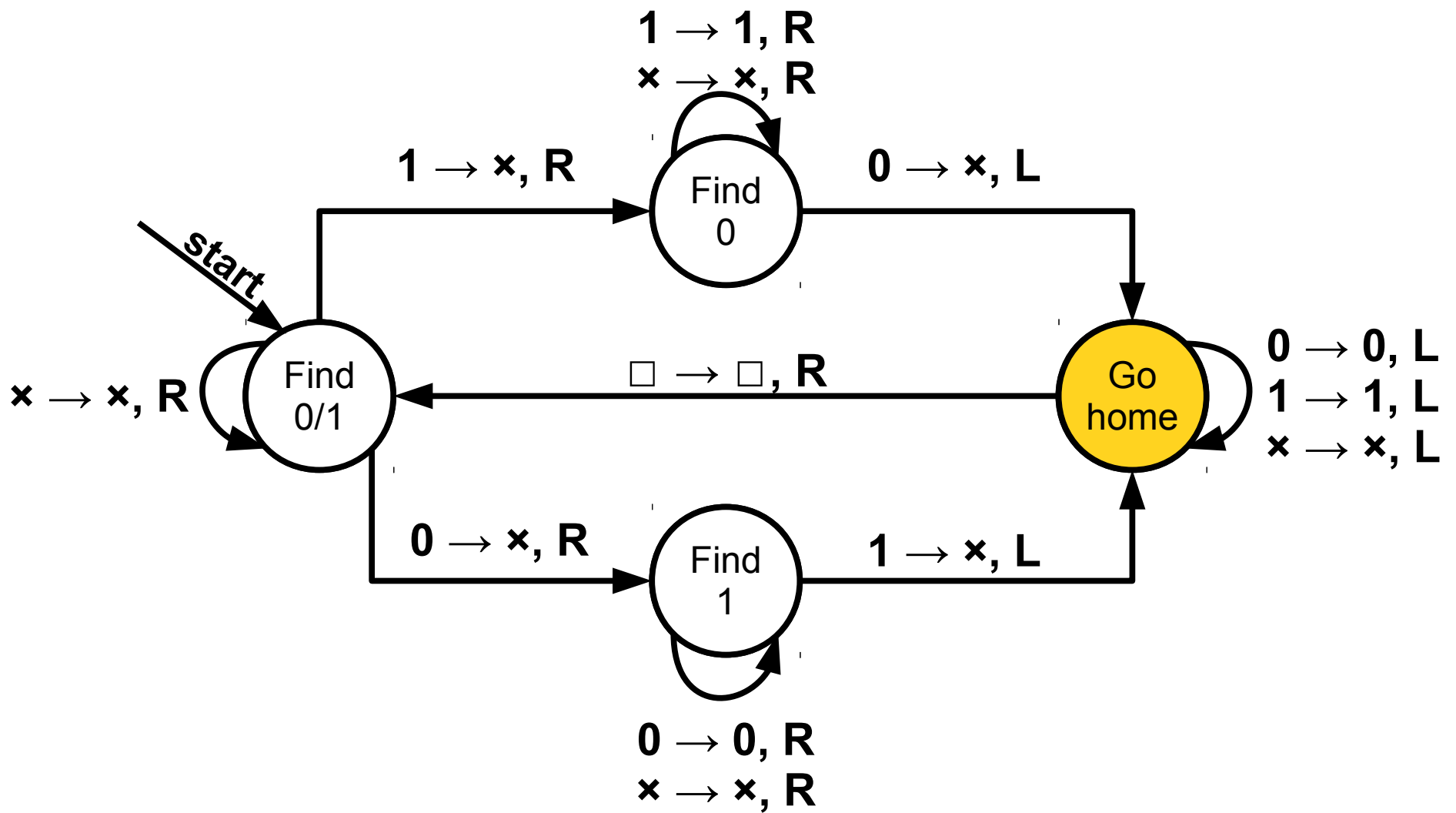


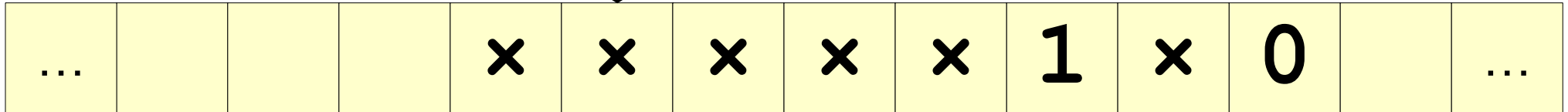
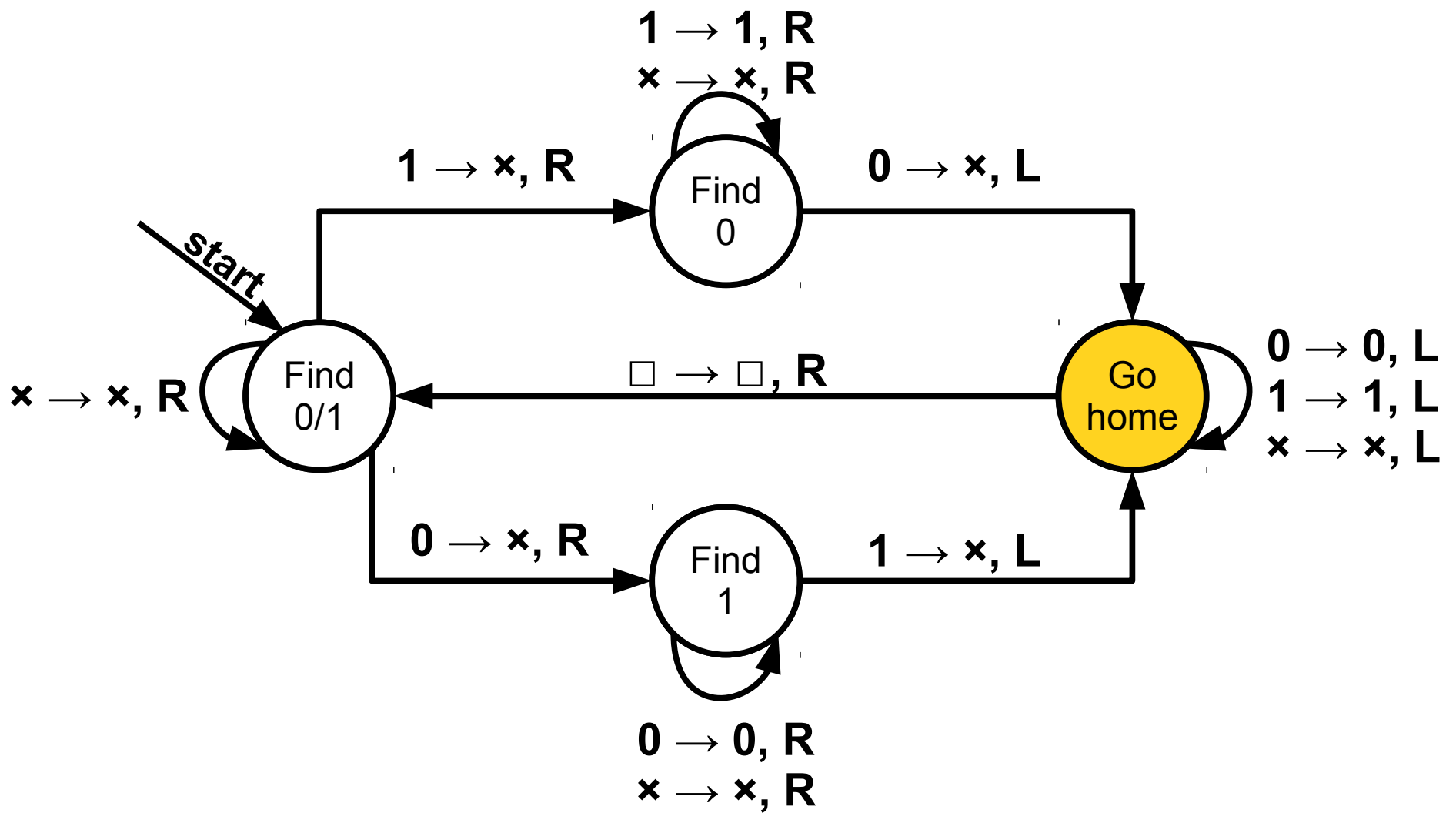


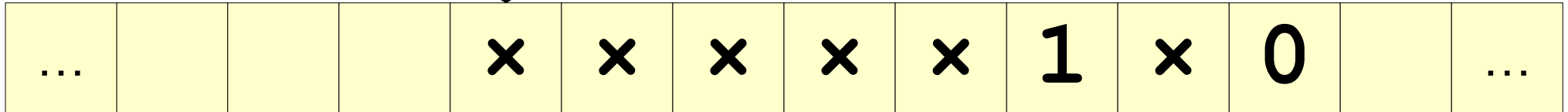
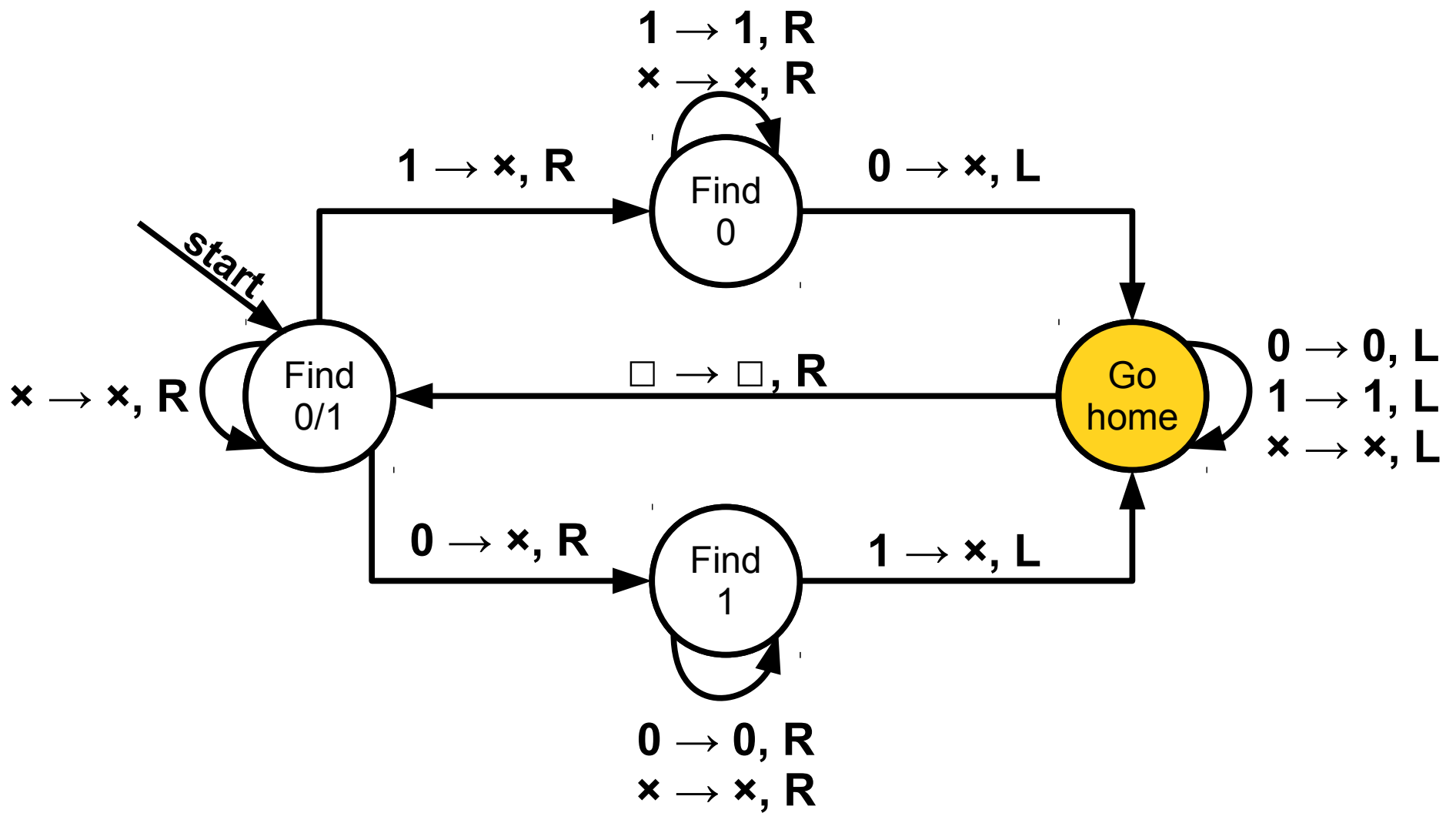


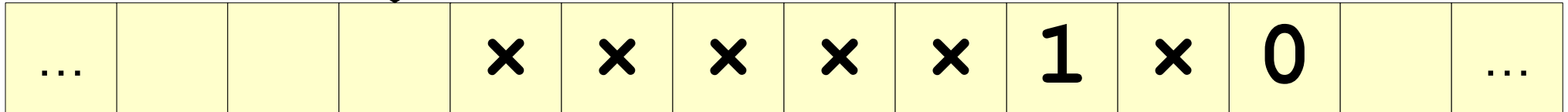
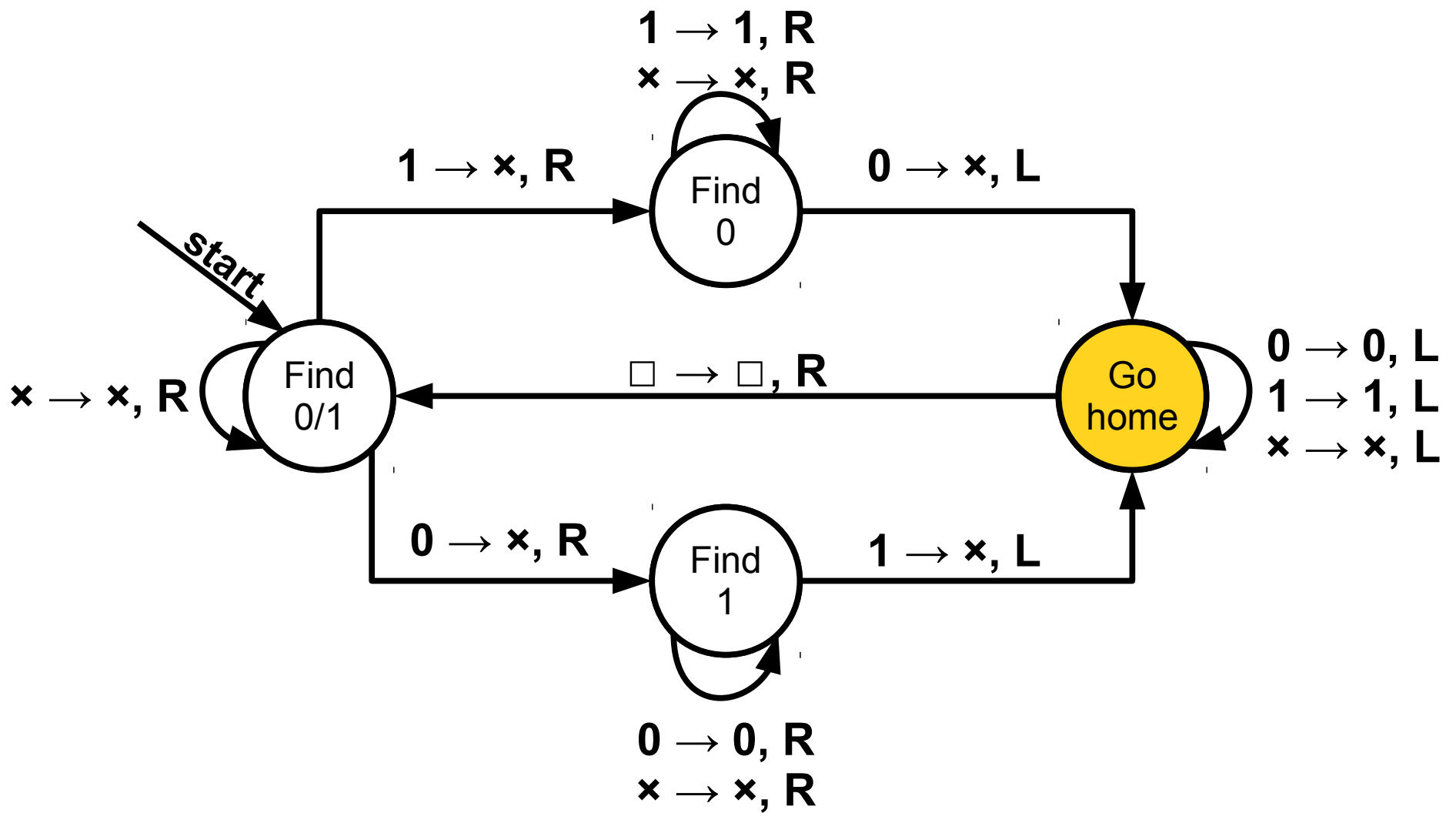


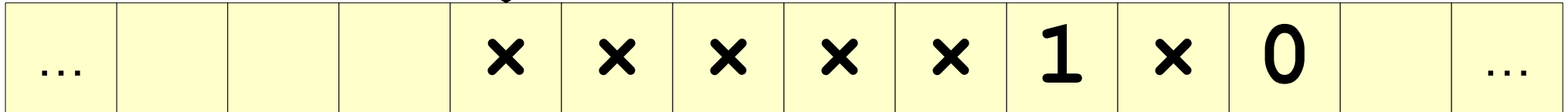
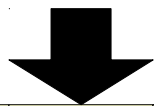
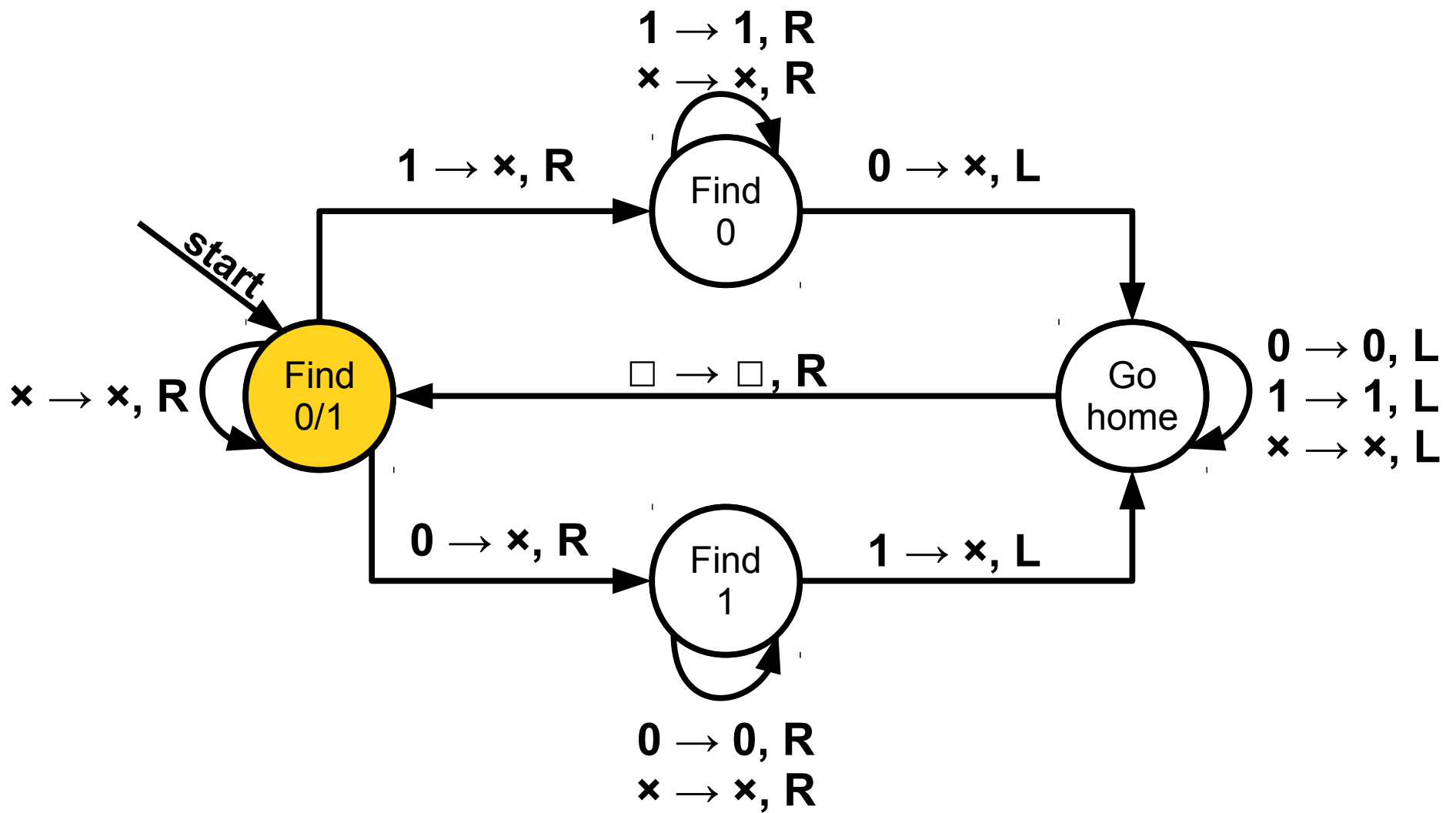


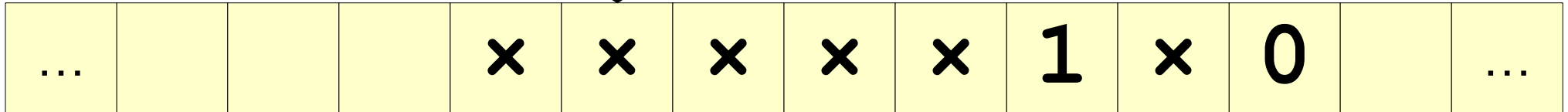
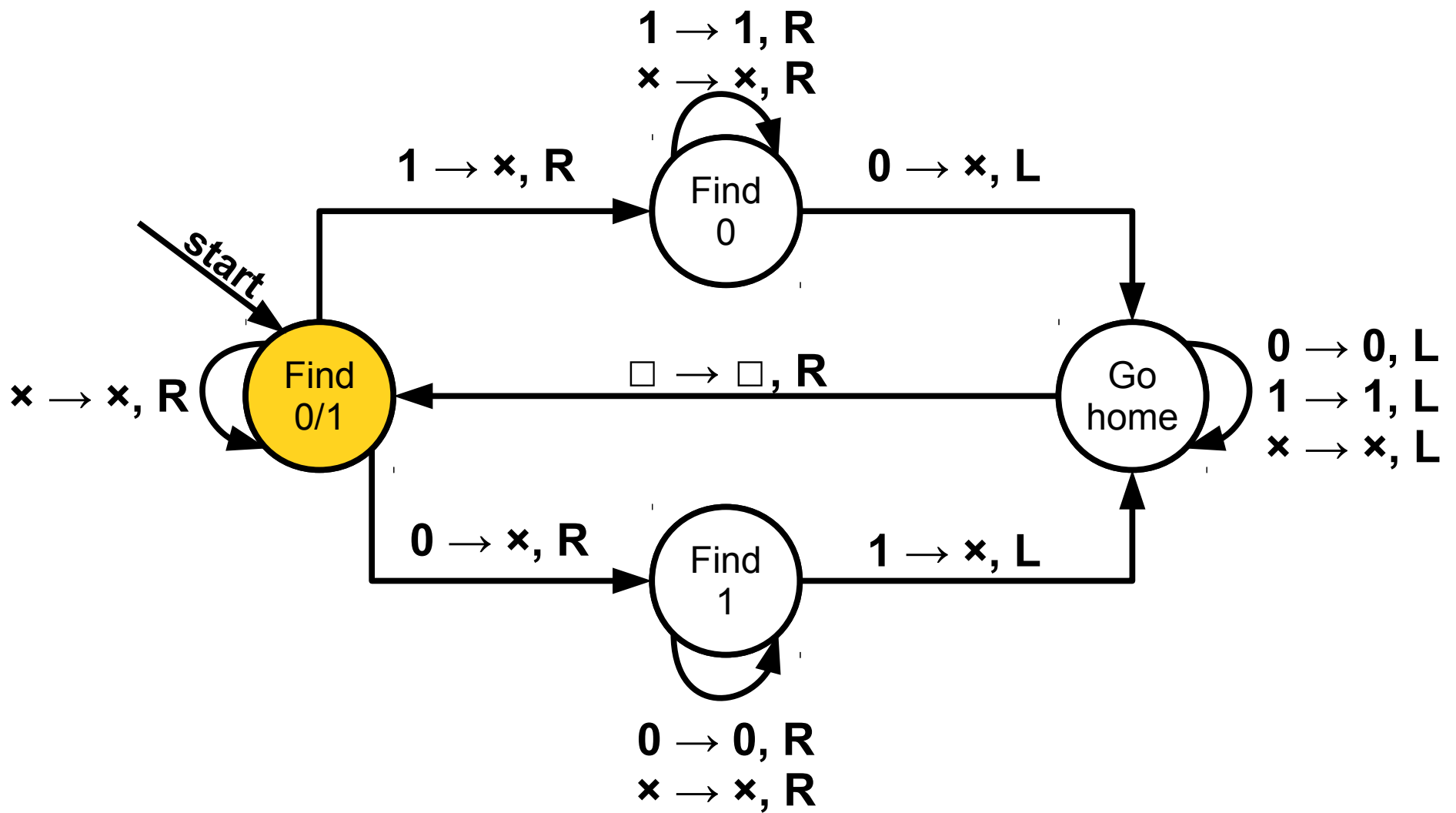


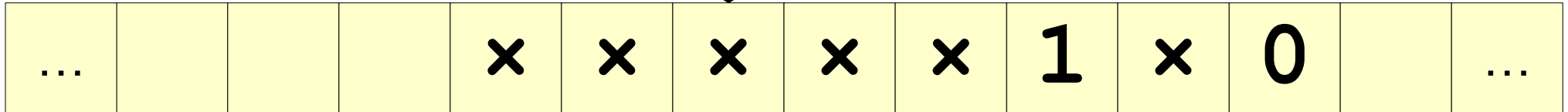
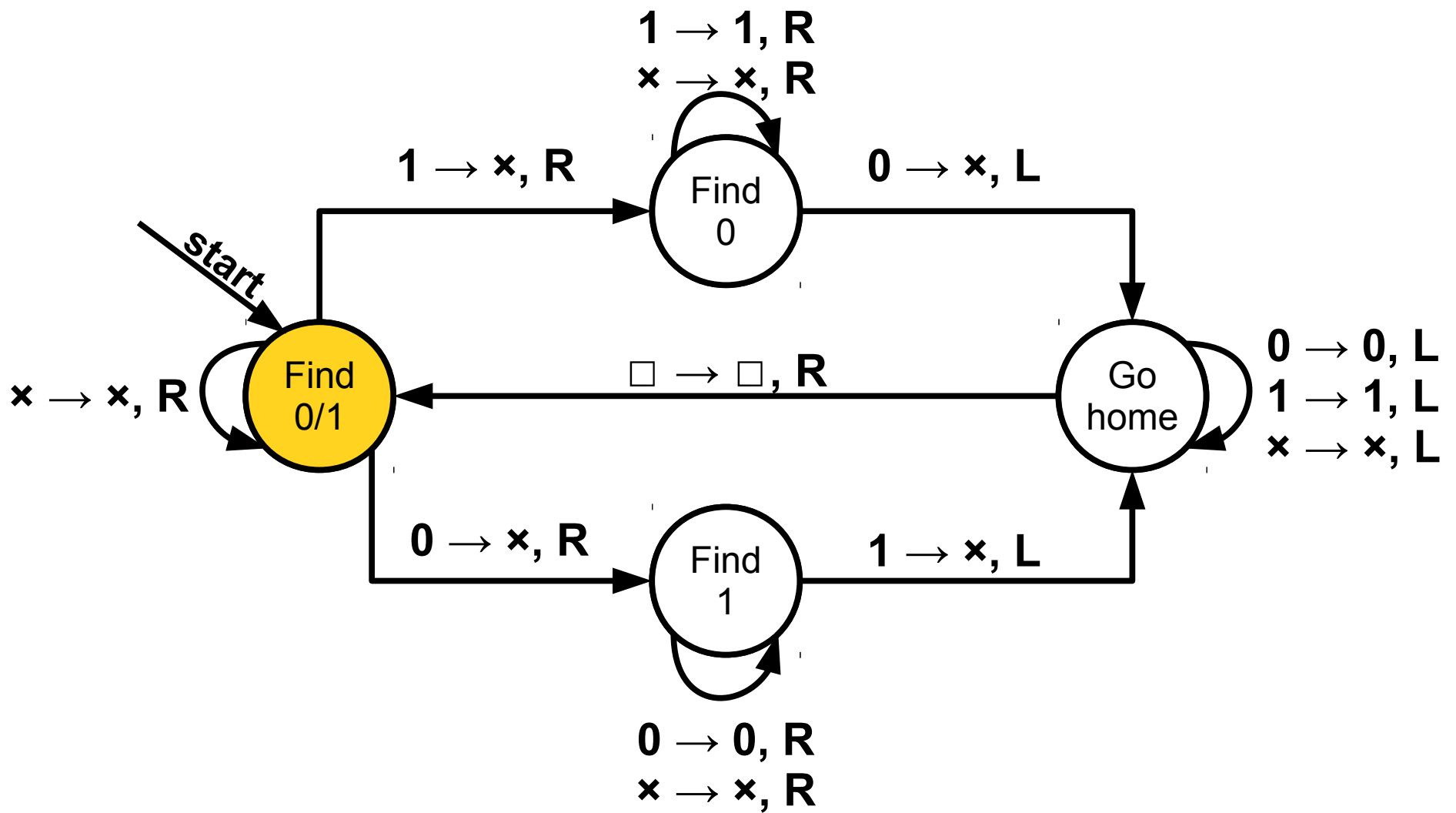


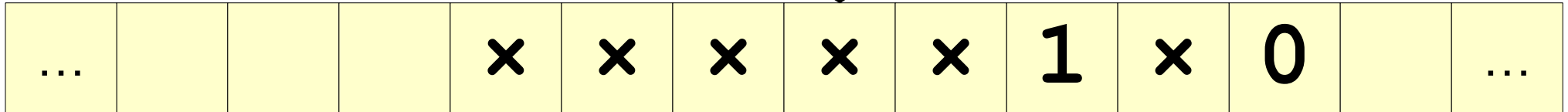
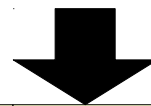
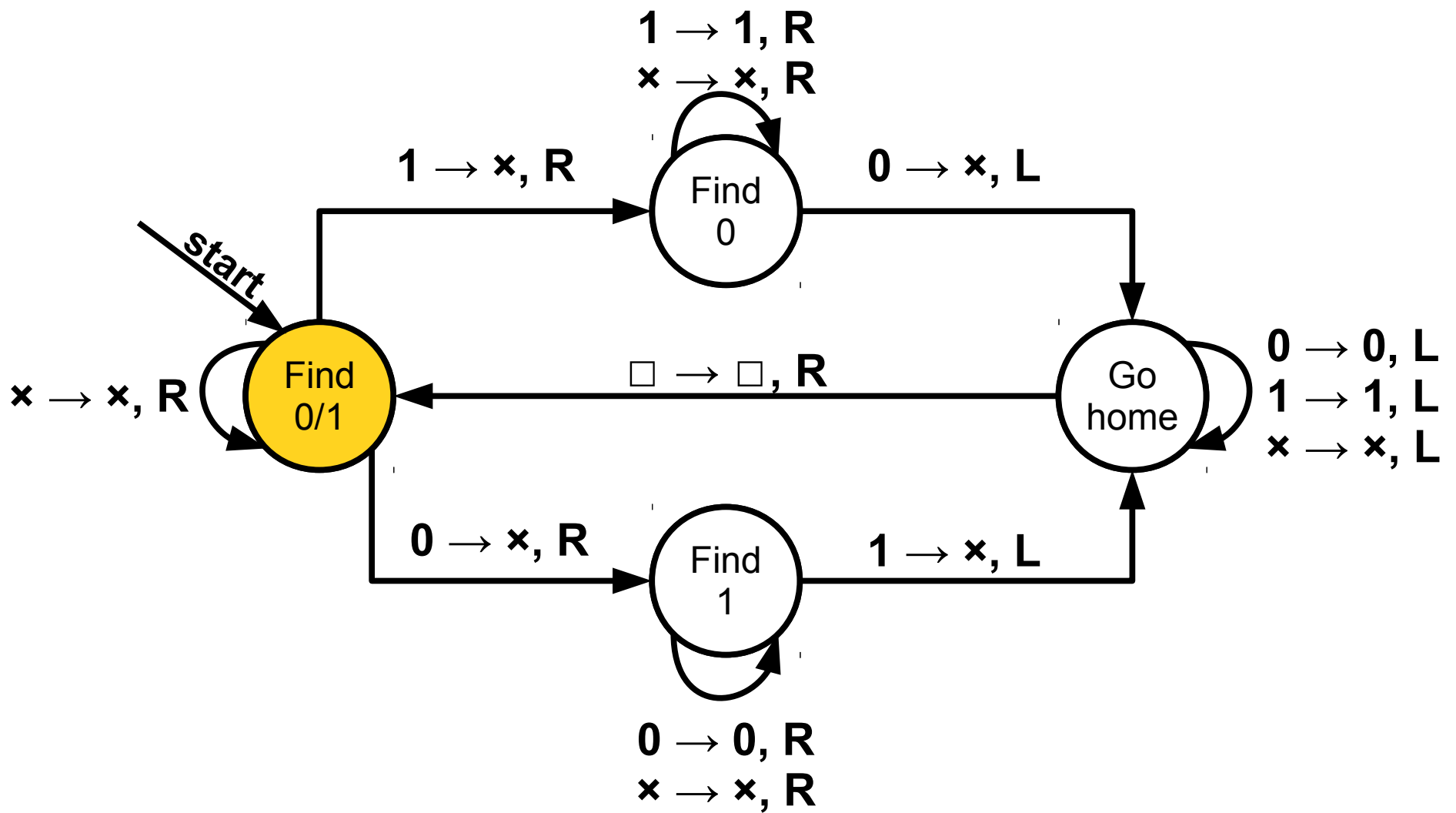


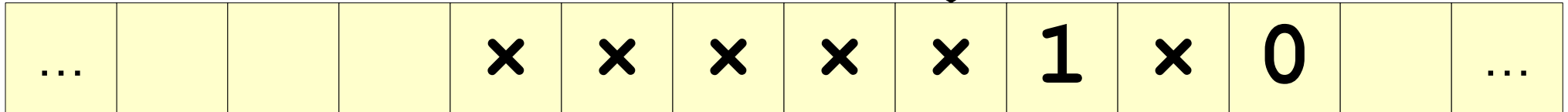
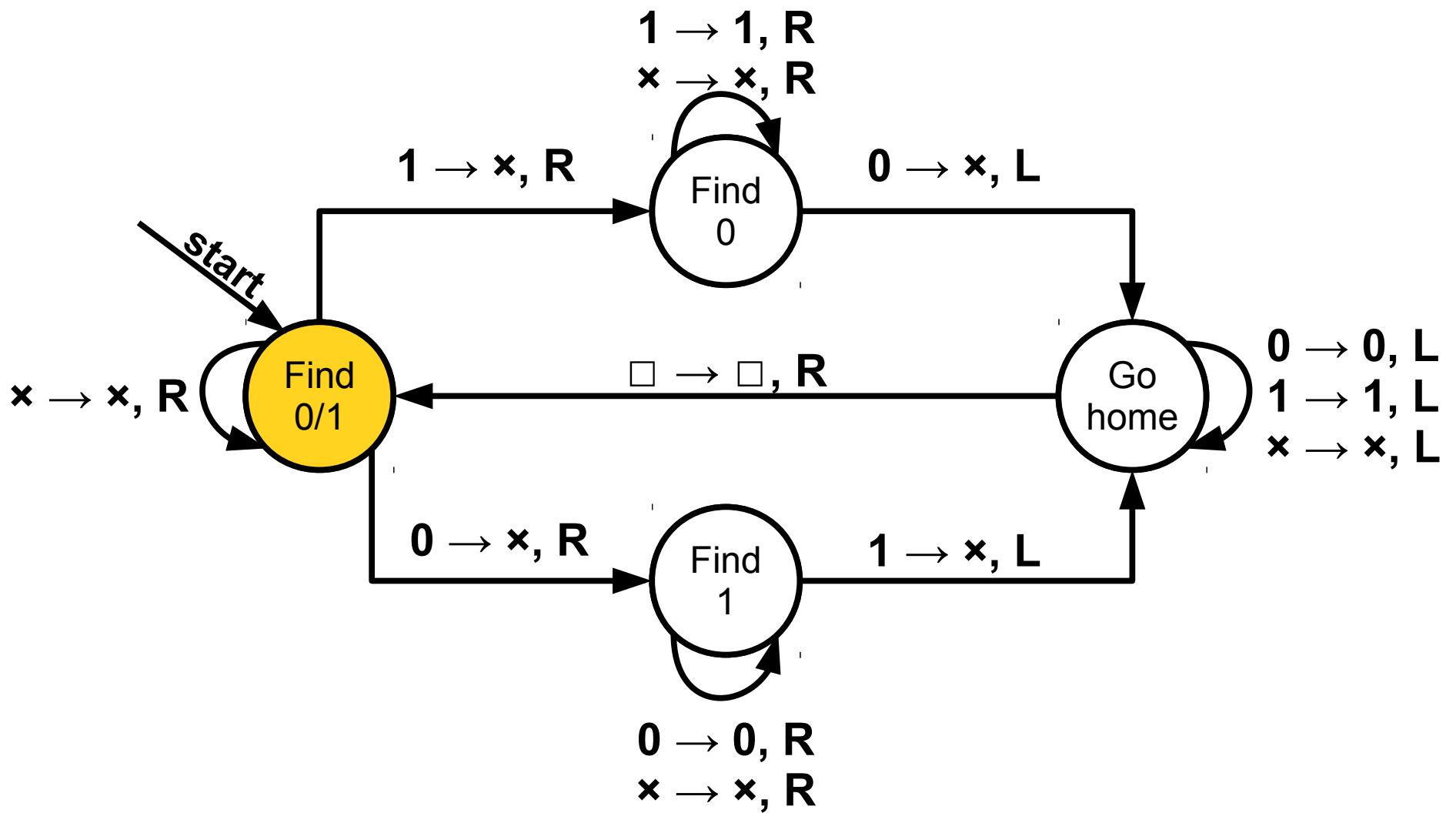


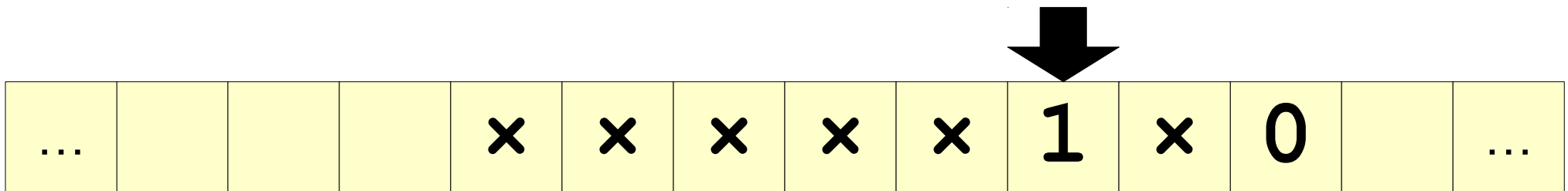
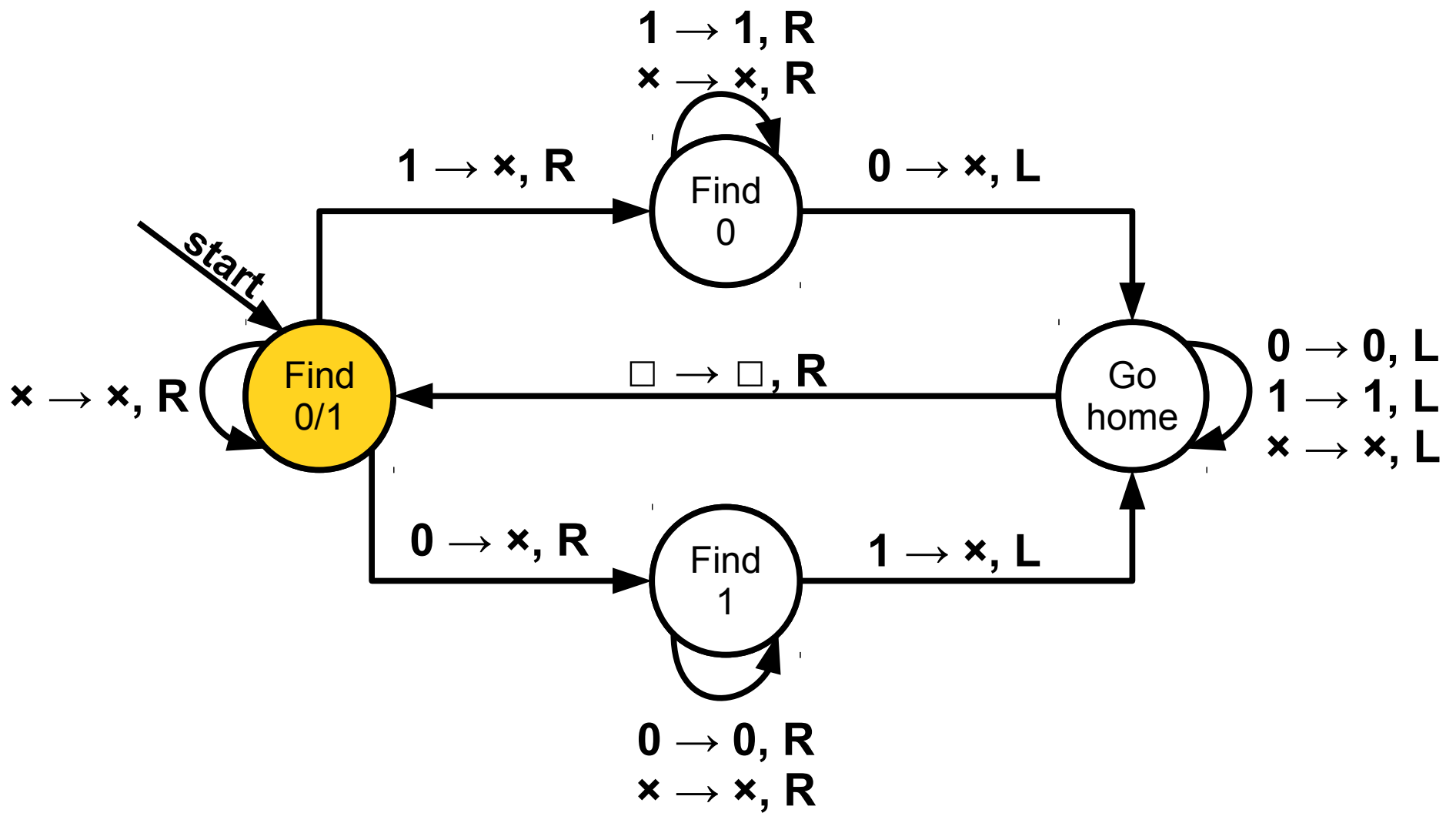


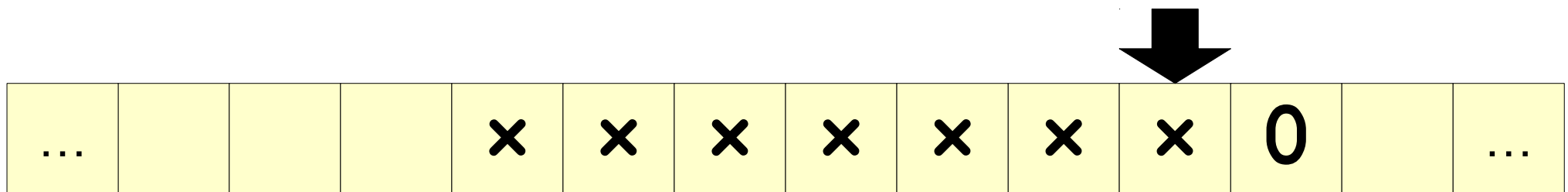
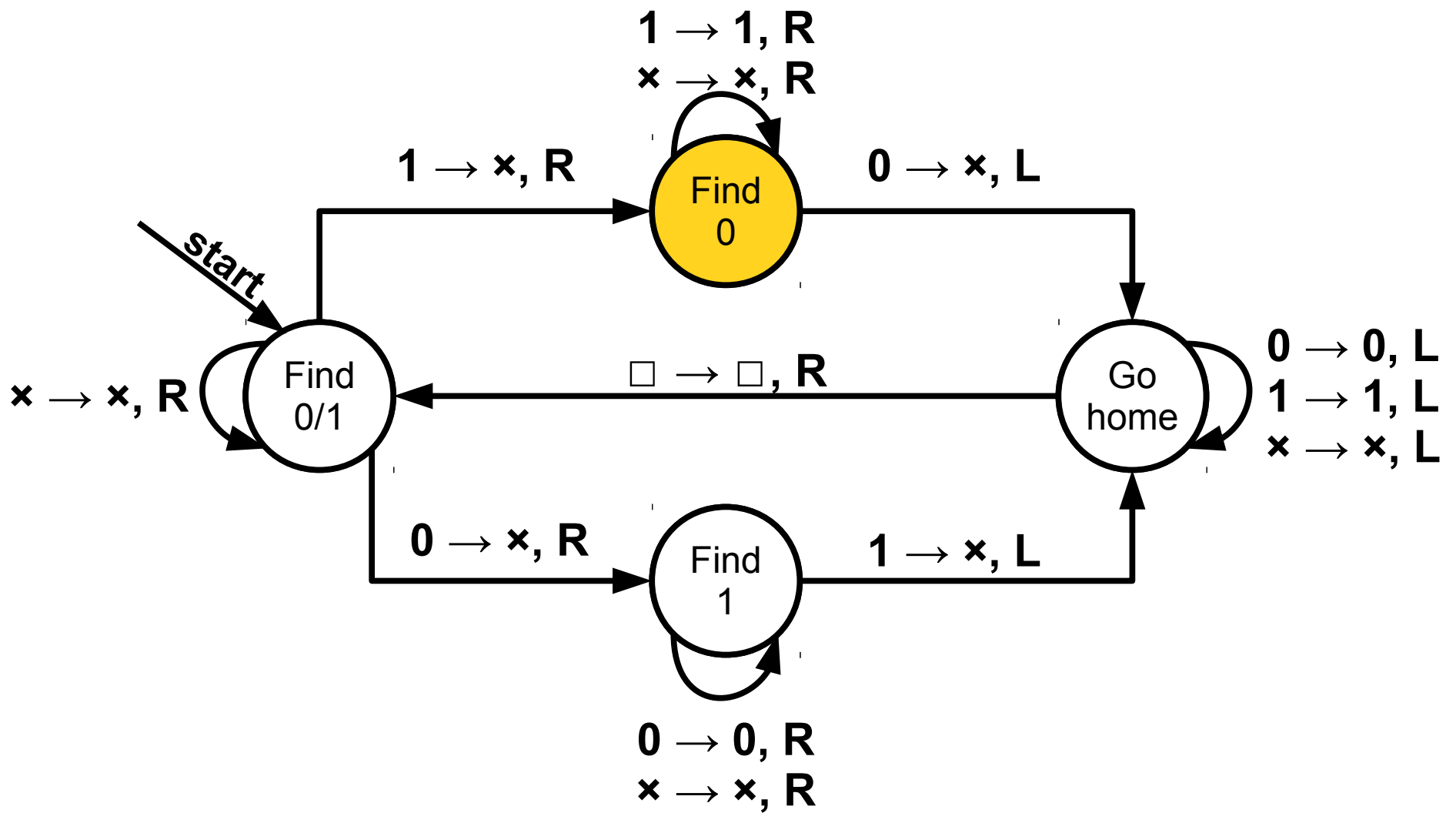


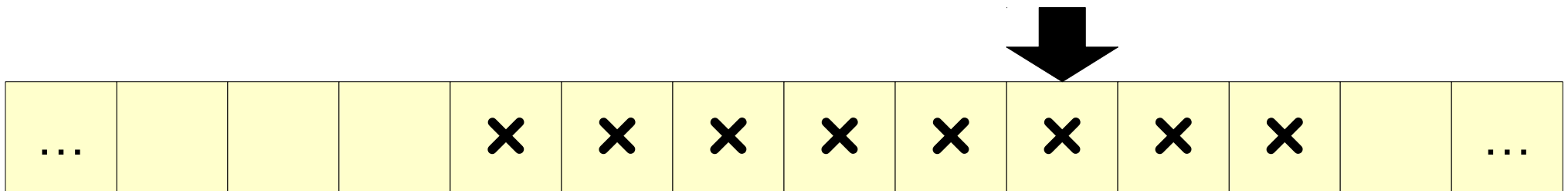
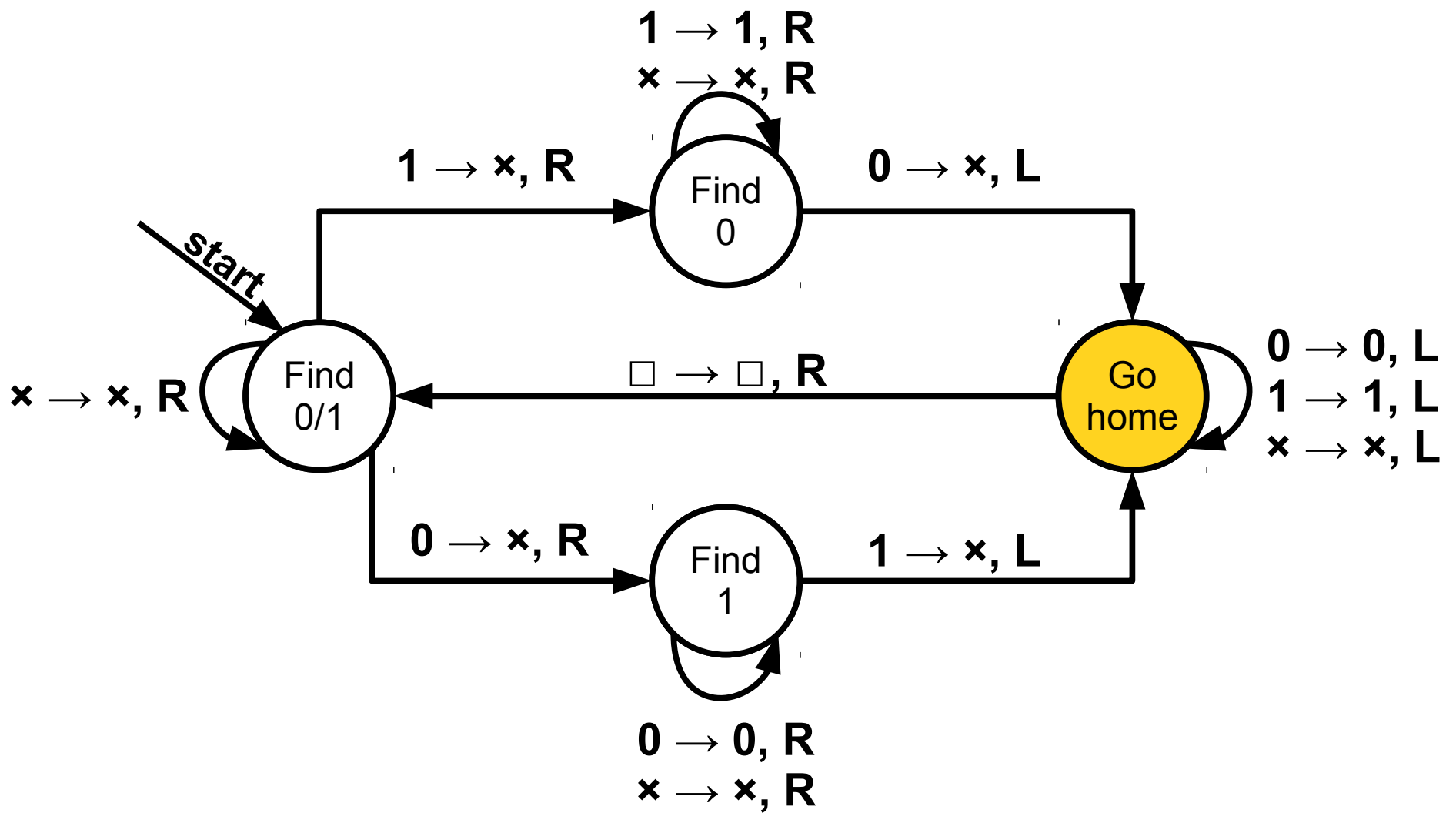


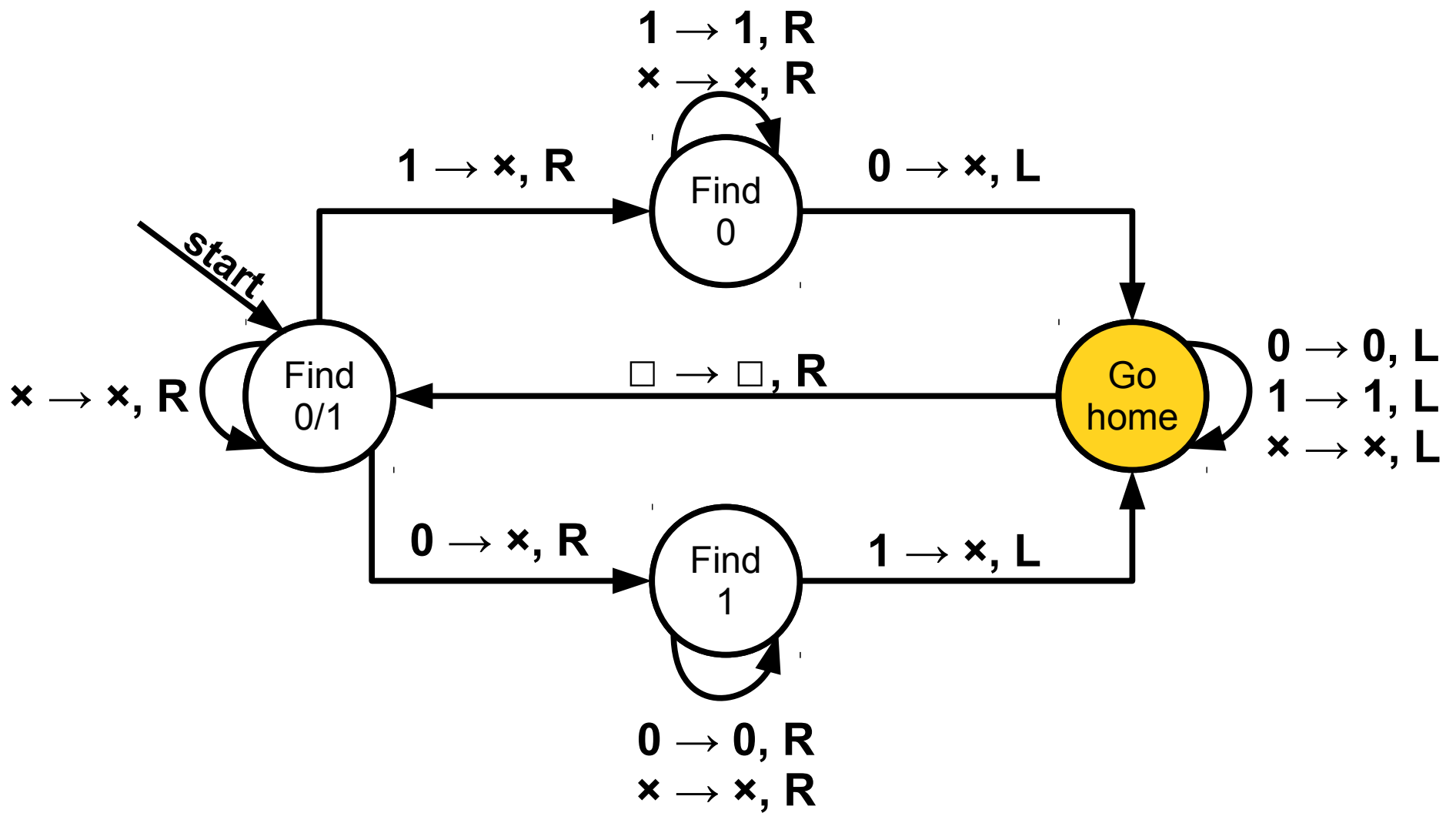


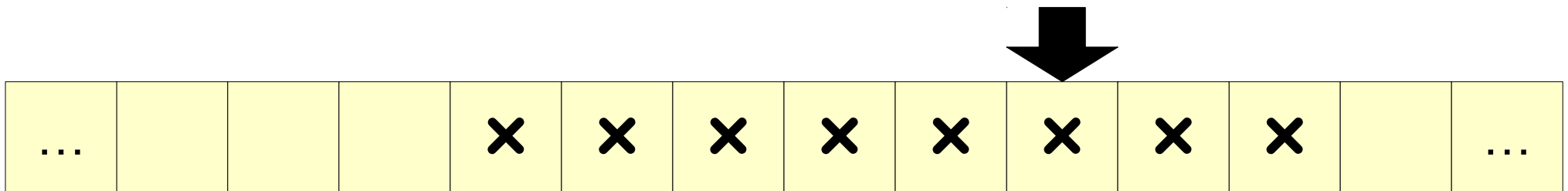
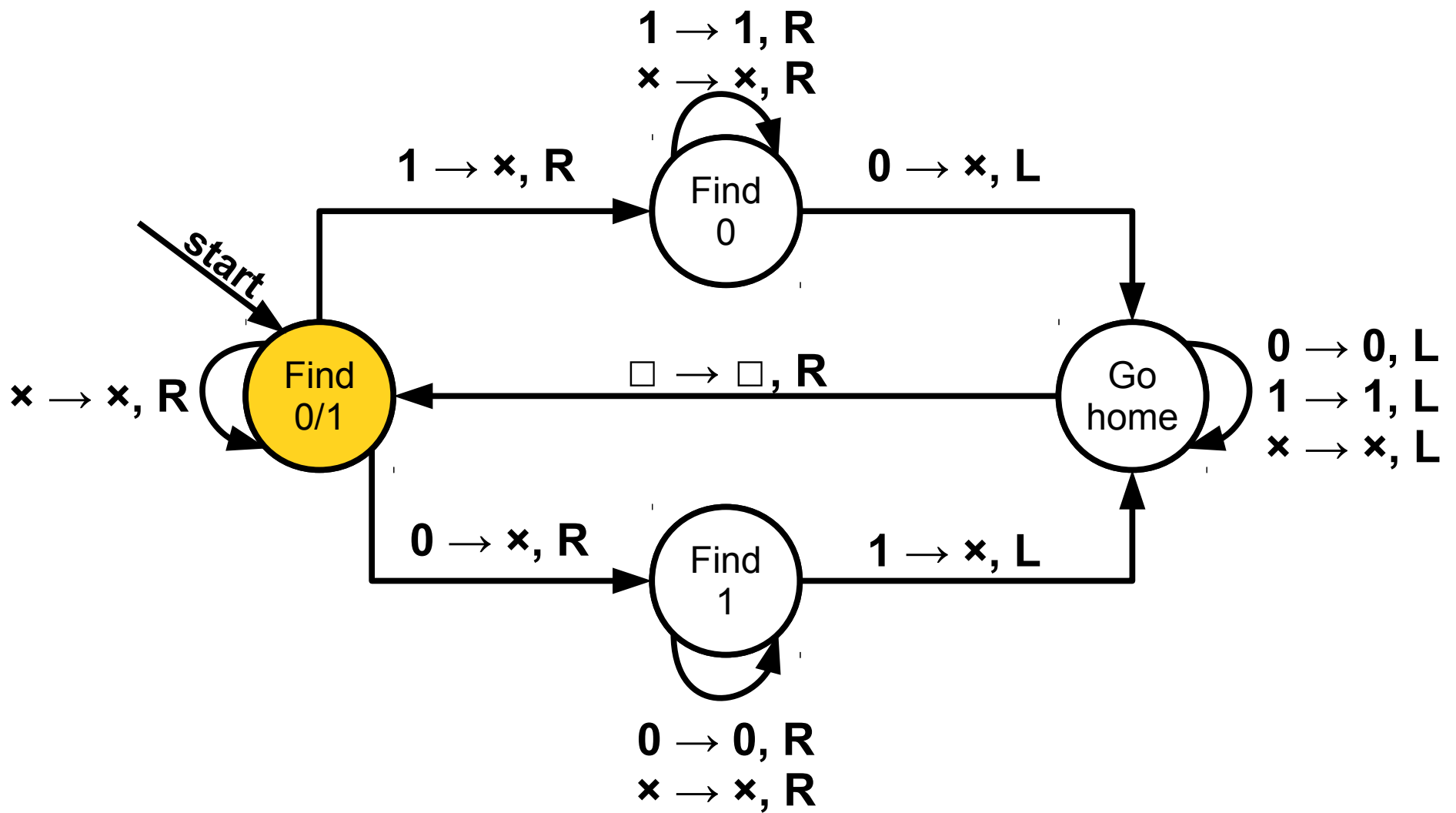


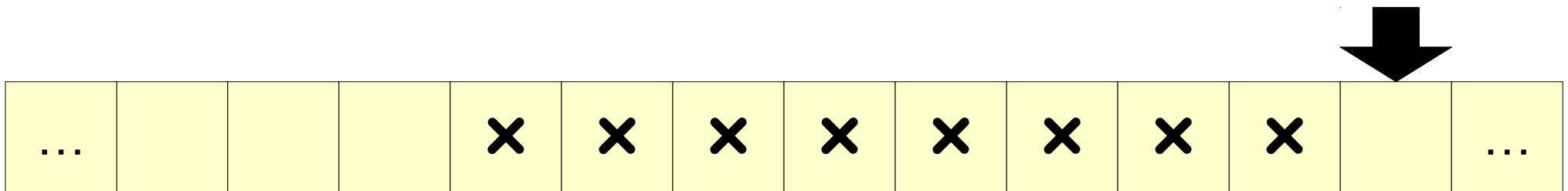
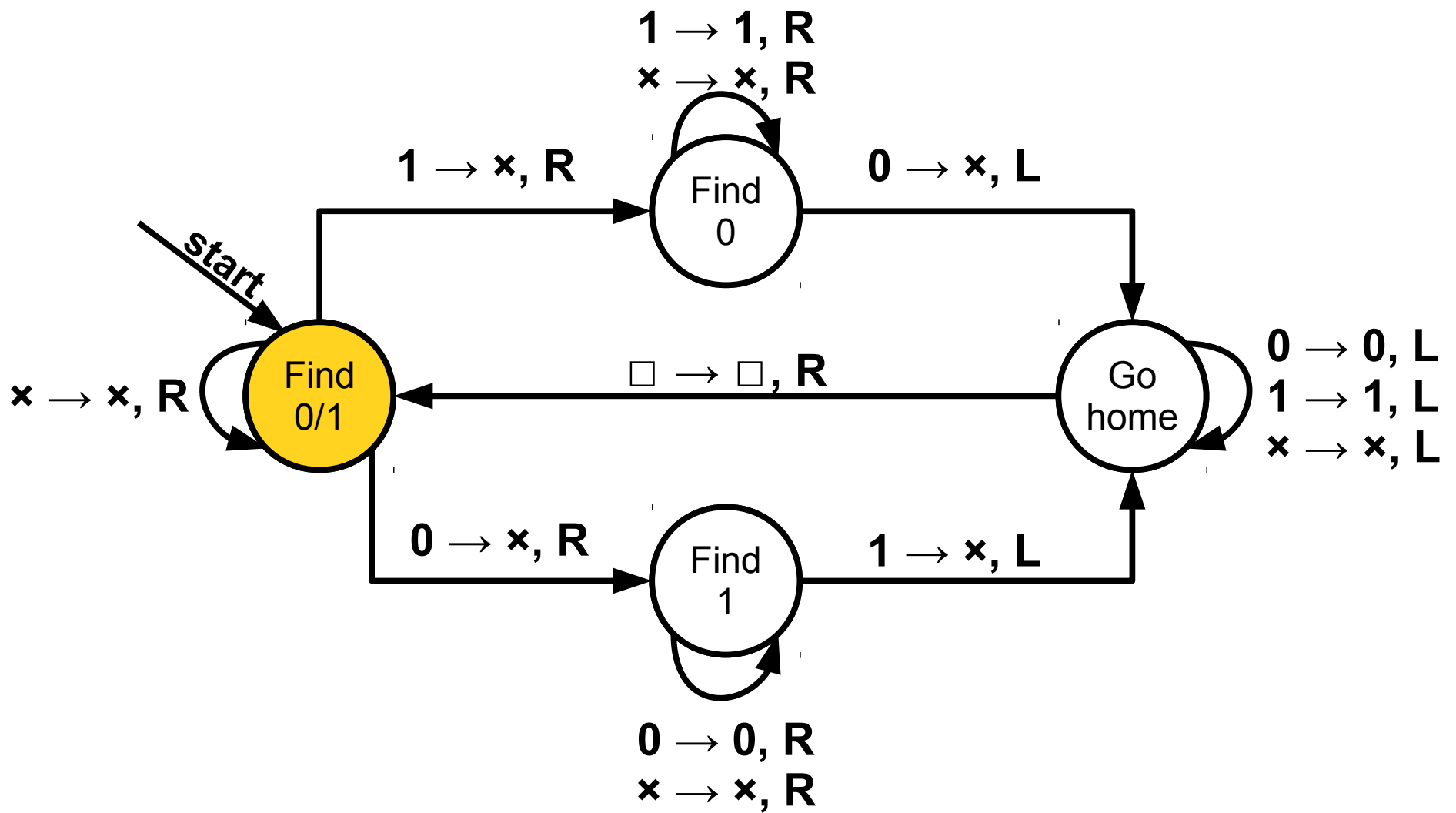


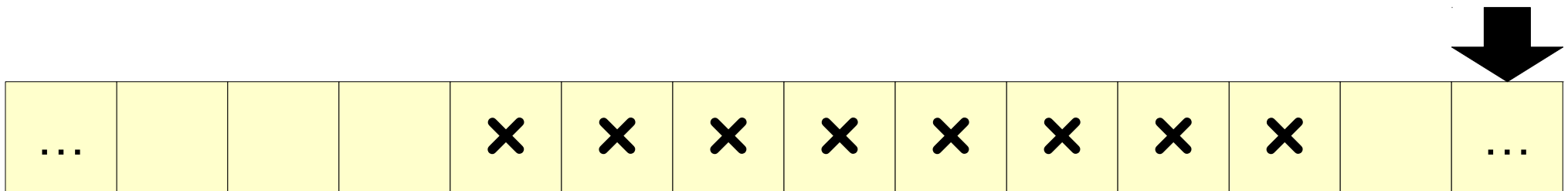
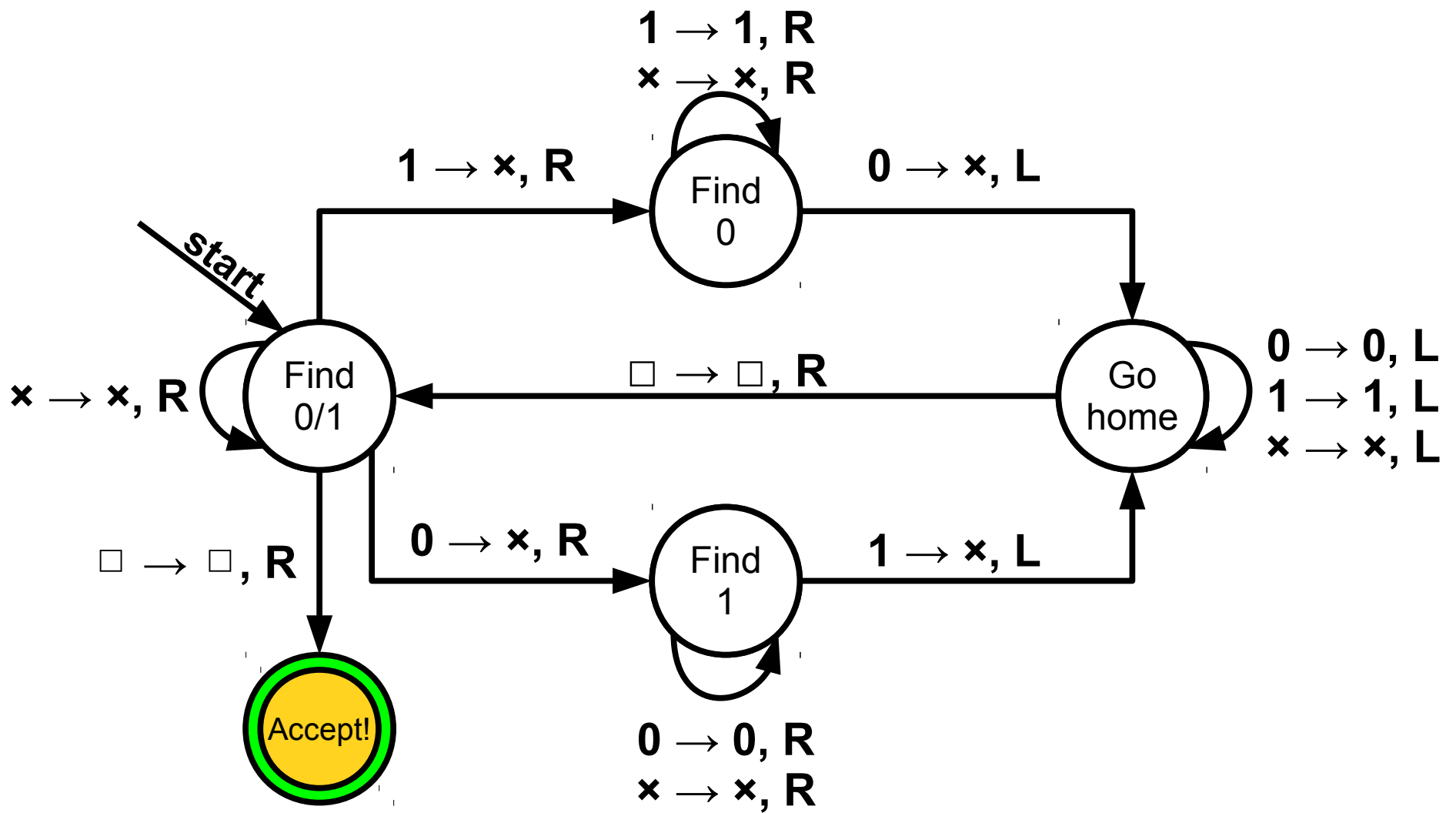


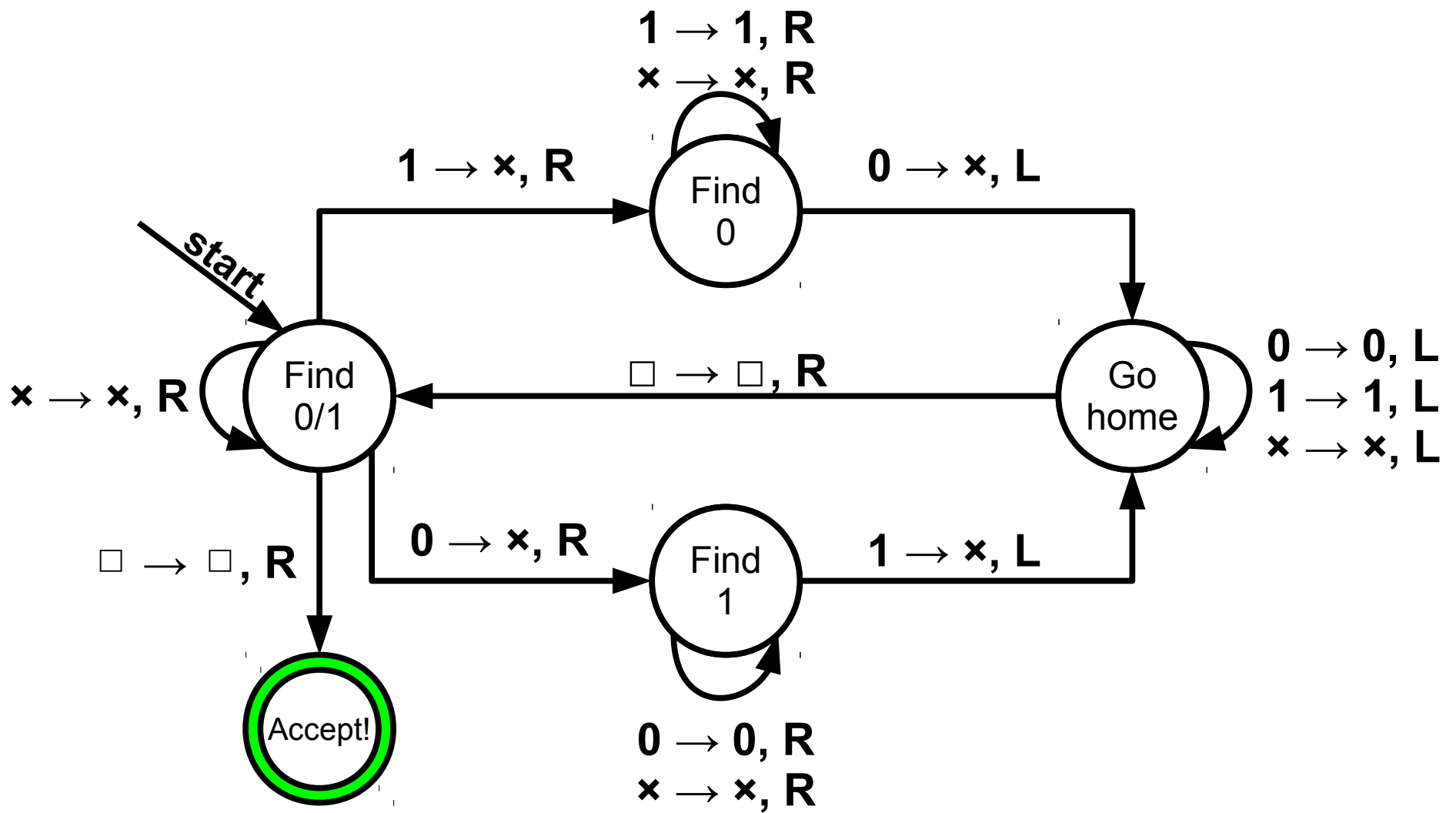


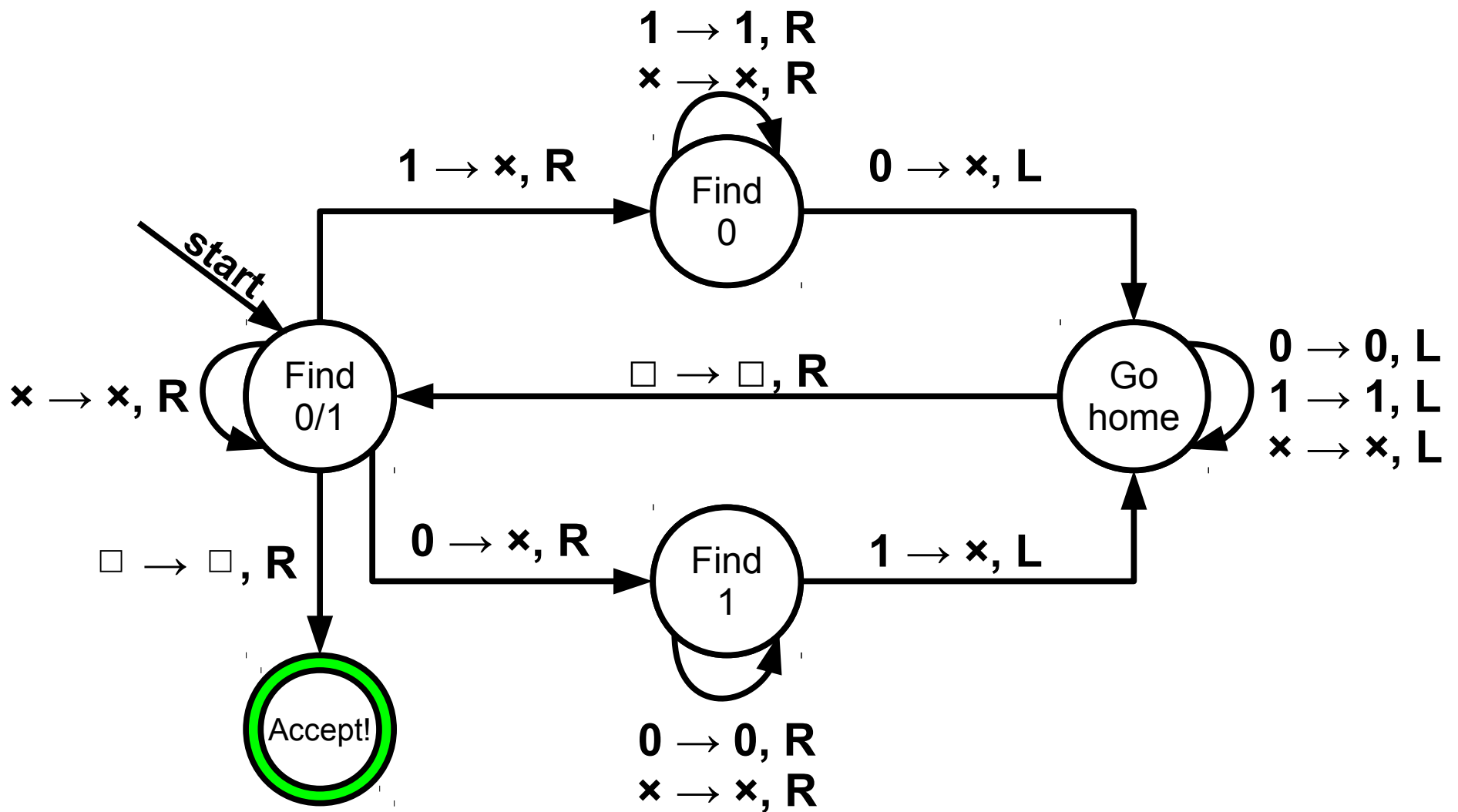




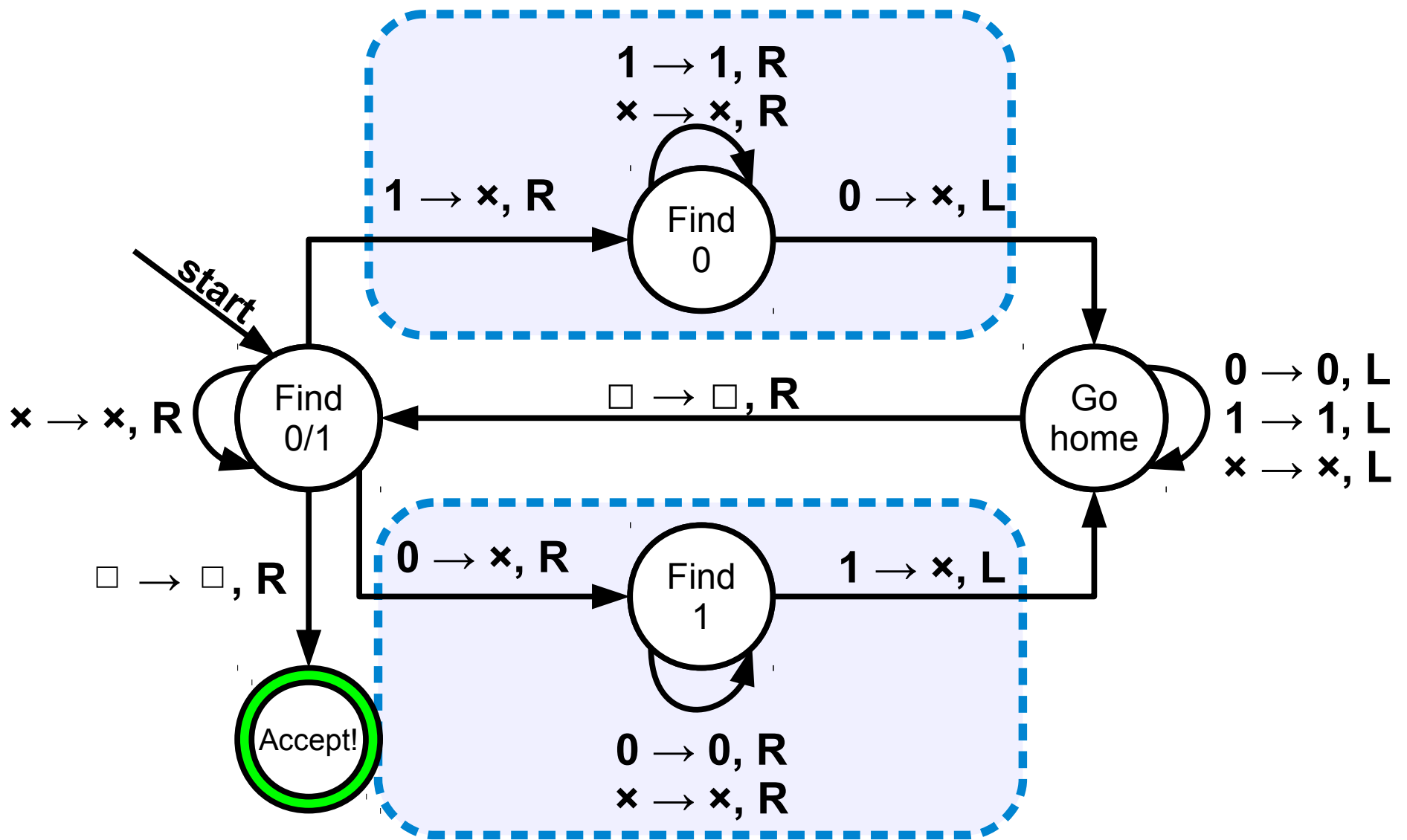








Remember that all missing transitions implicitly reject.



Constant Storage

- Sometimes, a TM needs to remember some additional information that can't be put on the tape.
- In this case, you can use similar techniques from DFAs and introduce extra states into the TM's finite-state control.
- The finite-state control can only remember one of finitely many things, but that might be all that you need!

Time-Out for Announcements!

Problem Set Seven

- Problem Set Seven is due this Friday at 2:30PM.
 - As always, if you have questions, feel free to stop by office hours or ask on Piazza!
- Problem Set Eight will go out on Friday.

Back to CS103!

Another TM Design

- We just designed a TM for this language over $\Sigma = \{0, 1\}$:

$$L = \{ w \in \Sigma^* \mid w \text{ has the same number of } 0\text{s and } 1\text{s} \}$$

- Let's do a quick review of how it worked.

A Leap of Faith



...			0	0	0	1	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

A Leap of Faith



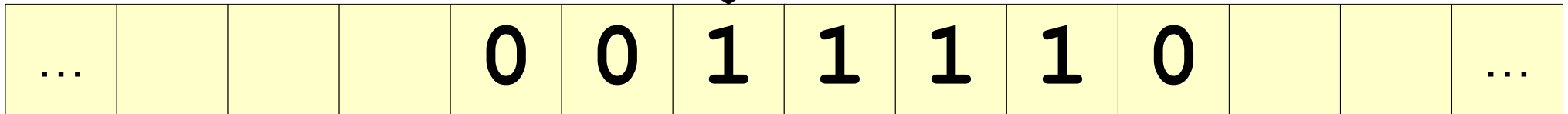
...				0	0	1	1	1	1	0			...
-----	--	--	--	---	---	---	---	---	---	---	--	--	-----

A Leap of Faith



...				0	0	1	1	1	1	0			...
-----	--	--	--	---	---	---	---	---	---	---	--	--	-----

A Leap of Faith



A Leap of Faith



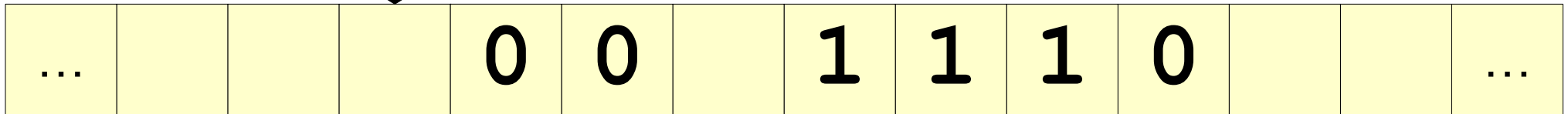
...				0	0		1	1	1	0			...
-----	--	--	--	---	---	--	---	---	---	---	--	--	-----

A Leap of Faith



...				0	0		1	1	1	0			...
-----	--	--	--	---	---	--	---	---	---	---	--	--	-----

A Leap of Faith

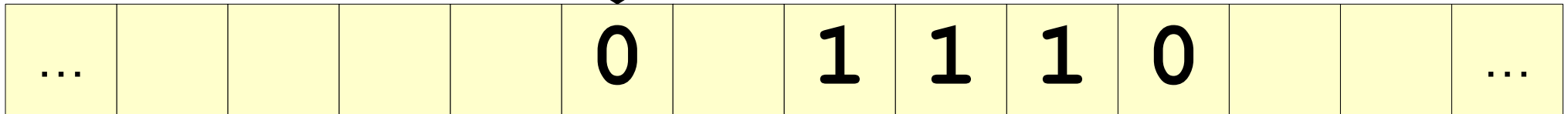


A Leap of Faith

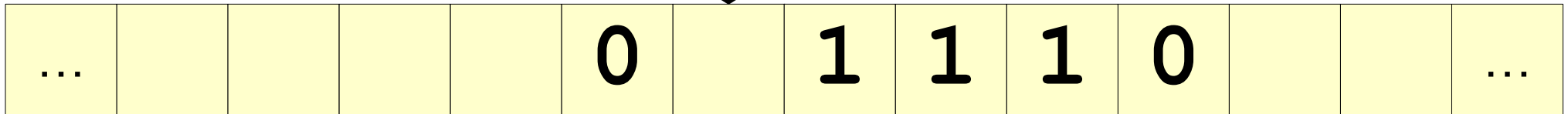


...				0	0		1	1	1	0			...
-----	--	--	--	---	---	--	---	---	---	---	--	--	-----

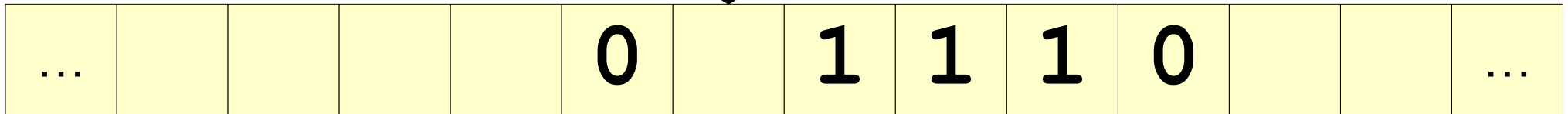
A Leap of Faith



A Leap of Faith



A Leap of Faith



How do we know that
this blank isn't one of
the infinitely many
blanks after our input
string?

The Solution



...			0	0	0	1	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

The Solution



...			×	0	0	1	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

The Solution



...			×	0	0	1	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

The Solution



...			×	0	0	1	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

The Solution



...			×	0	0	×	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

The Solution



...			x	0	0	x	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

The Solution



...			x	0	0	x	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

The Solution



...			×	0	0	×	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

The Solution



...			x	0	0	x	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

The Solution



...			x	0	0	x	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

The Solution



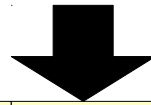
...			x	x	0	x	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

The Solution



...			x	x	0	x	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

The Solution



...			x	x	0	x	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

The Solution



...			x	x	0	x	x	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

The Solution



...			x	x	0	x	x	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

The Solution



...			x	x	0	x	x	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

The Solution



...			x	x	0	x	x	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

The Solution



...			x	x	0	x	x	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

The Solution



...			x	x	0	x	x	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

A Different Idea

A Different Strategy



...			0	0	0	1	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

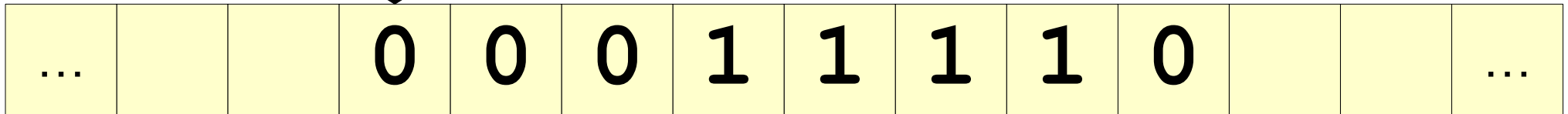
A Different Strategy



...			0	0	0	1	1	1	1	0			...
-----	--	--	---	---	---	---	---	---	---	---	--	--	-----

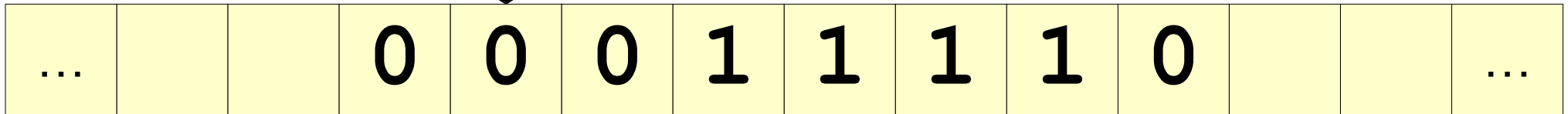
Could we sort
the characters
of this string?

A Different Strategy



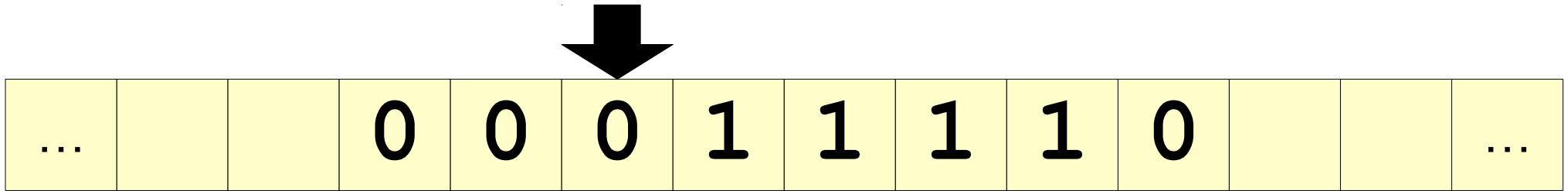
Observation 1: A string of 0s and 1s is sorted if it matches the regex 0^*1^* .

A Different Strategy



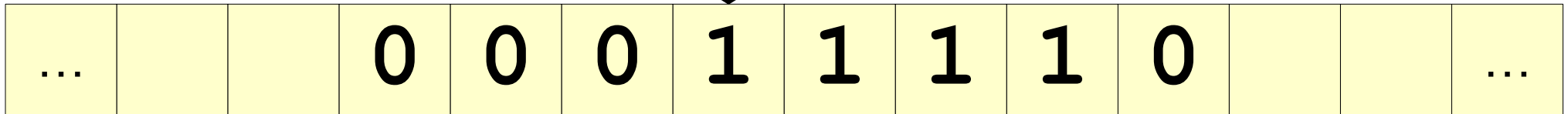
Observation 1: A string of 0s and 1s is sorted if it matches the regex 0^*1^* .

A Different Strategy



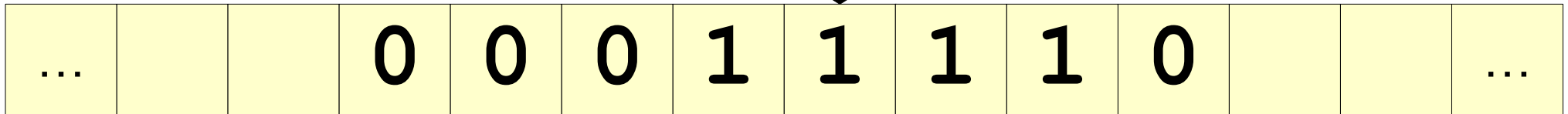
Observation 1: A string
of 0s and 1s is sorted
if it matches the
regex 0^*1^* .

A Different Strategy



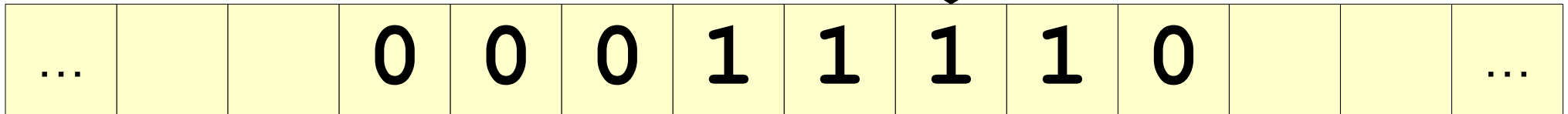
Observation 1: A string of 0s and 1s is sorted if it matches the regex 0^*1^* .

A Different Strategy



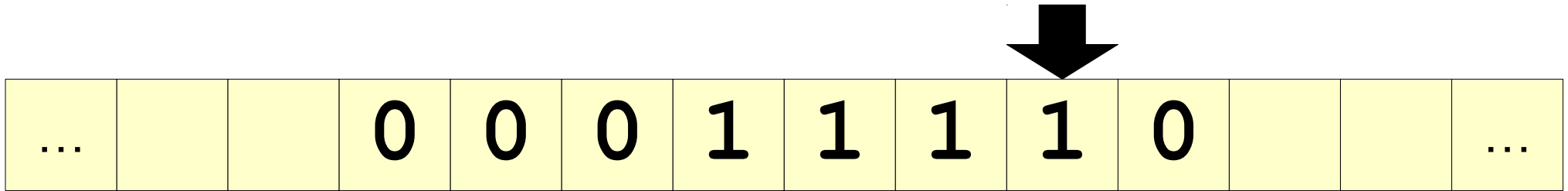
Observation 1: A string of 0s and 1s is sorted if it matches the regex 0^*1^* .

A Different Strategy



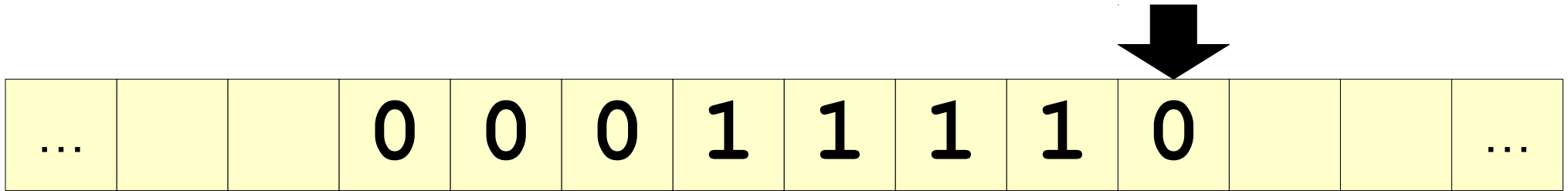
Observation 1: A string
of 0s and 1s is sorted
if it matches the
regex 0^*1^* .

A Different Strategy



Observation 1: A string of 0s and 1s is sorted if it matches the regex 0^*1^* .

A Different Strategy



Observation 1: A string of 0s and 1s is sorted if it matches the regex 0^*1^* .

A Different Strategy



Observation 1: A string of 0s and 1s is sorted if it matches the regex 0^*1^* .

A Different Strategy



Observation 2: A string of 0s and 1s is not sorted if it contains 10 as a substring.

A Different Strategy



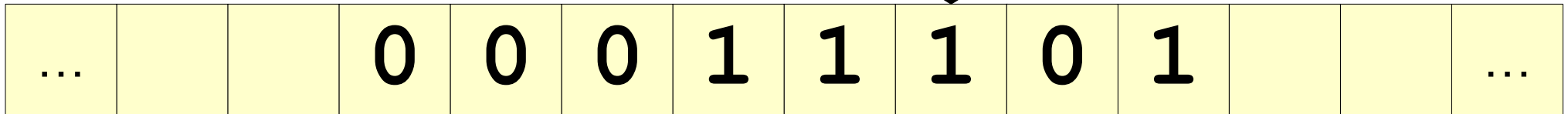
Observation 2: A string of 0s and 1s is not sorted if it contains 10 as a substring.

A Different Strategy



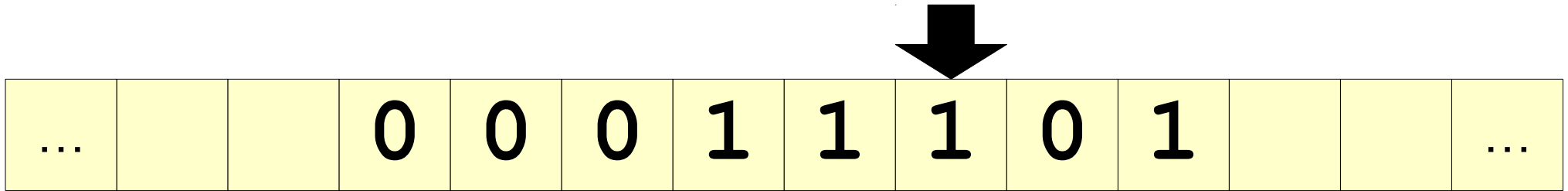
Observation 2: A string of 0s and 1s is not sorted if it contains 10 as a substring.

A Different Strategy



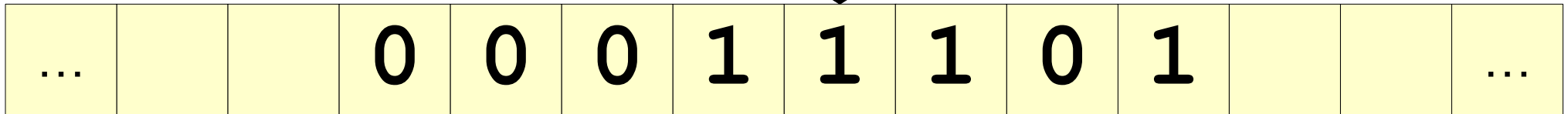
Observation 2: A string of 0s and 1s is not sorted if it contains 10 as a substring.

A Different Strategy



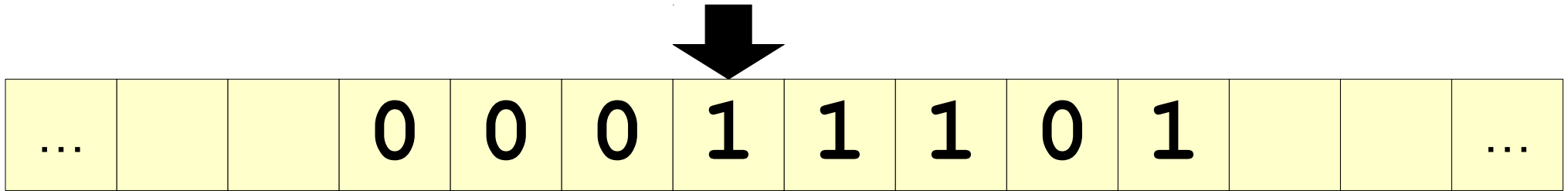
Idea: Repeatedly find
a copy of 10 and
replace it with 01.

A Different Strategy



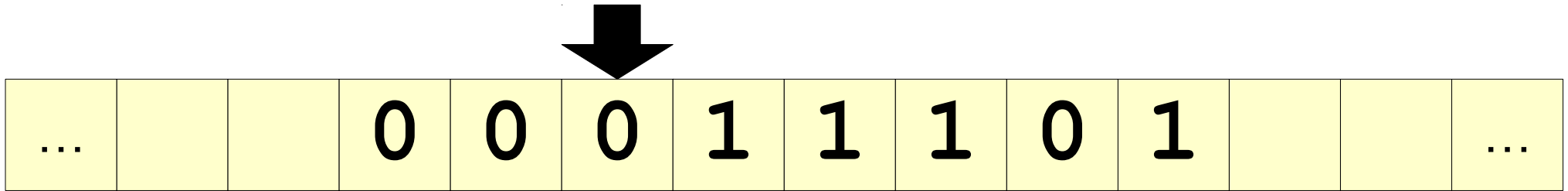
Idea: Repeatedly find
a copy of 10 and
replace it with 01.

A Different Strategy



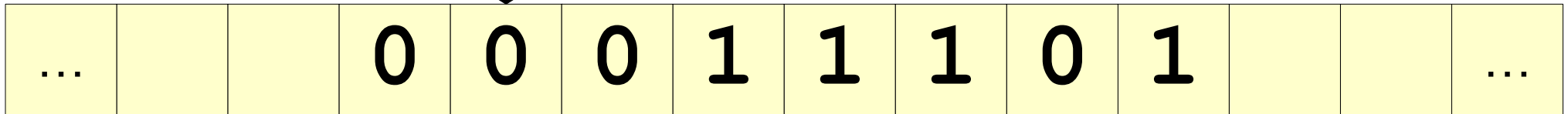
Idea: Repeatedly find
a copy of 10 and
replace it with 01.

A Different Strategy



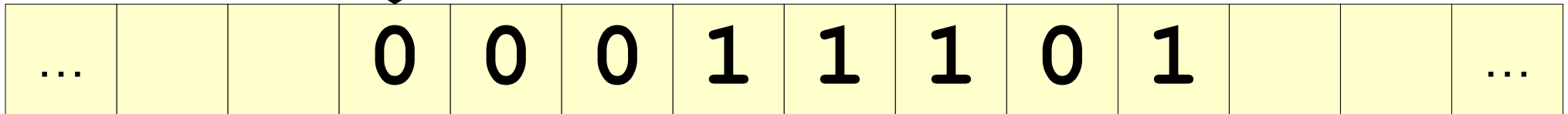
Idea: Repeatedly find
a copy of 10 and
replace it with 01.

A Different Strategy



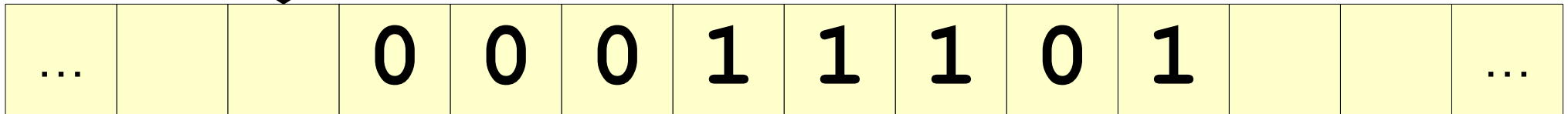
Idea: Repeatedly find
a copy of 10 and
replace it with 01.

A Different Strategy



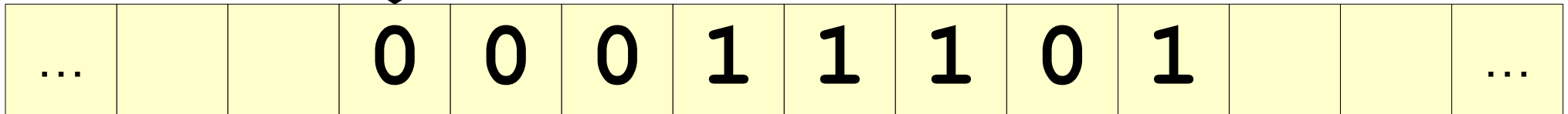
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A Different Strategy



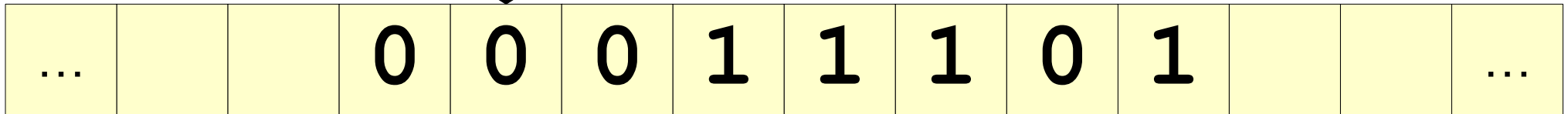
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A Different Strategy



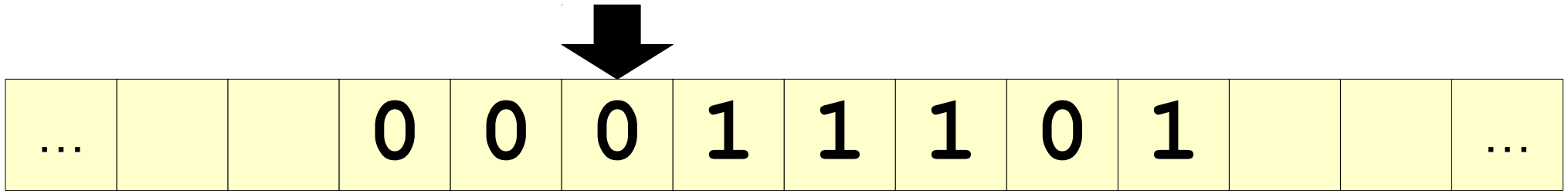
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A Different Strategy



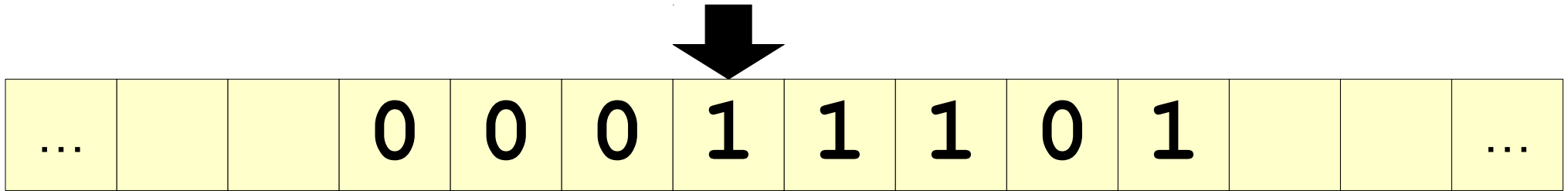
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A Different Strategy



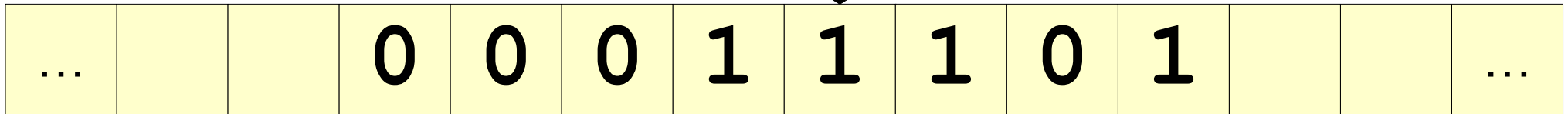
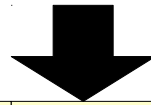
Idea: Repeatedly find
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A Different Strategy



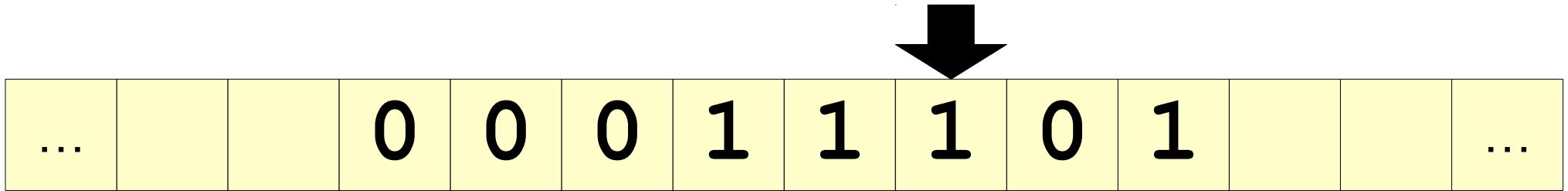
Idea: Repeatedly find
a copy of 10 and
replace it with 01.

A Different Strategy



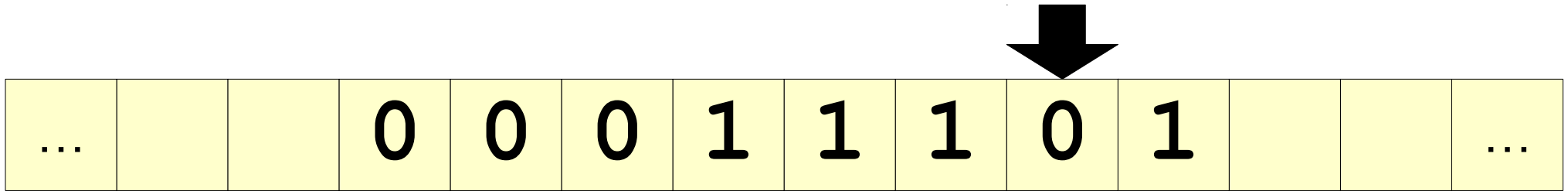
Idea: Repeatedly find
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A Different Strategy



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a copy of 10 and
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A Different Strategy



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a copy of 10 and
replace it with 01.

A Different Strategy



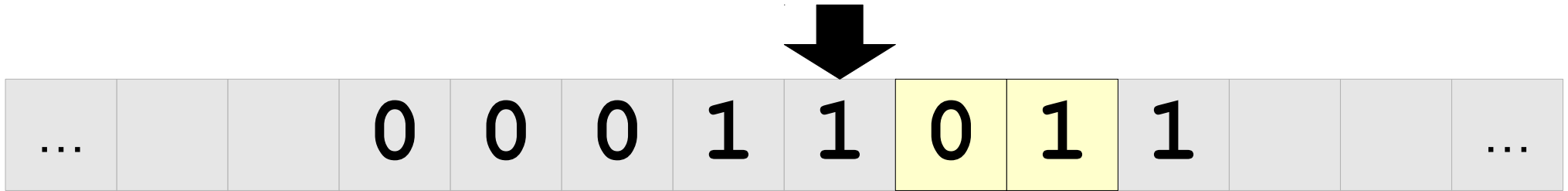
Idea: Repeatedly find
a copy of 10 and
replace it with 01.

A Different Strategy



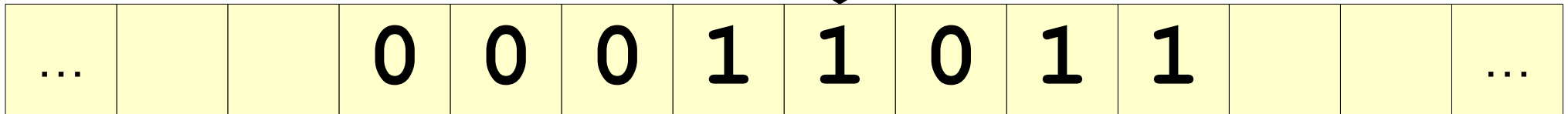
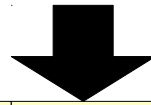
Idea: Repeatedly find
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replace it with 01.

A Different Strategy



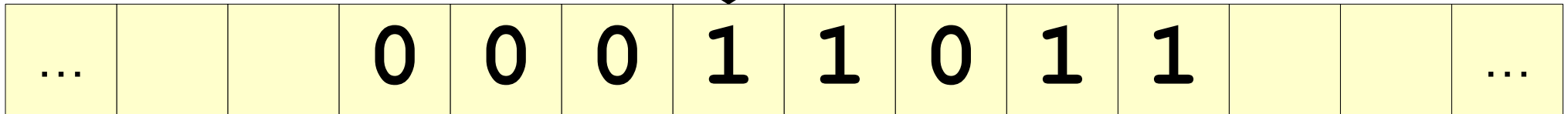
Idea: Repeatedly find
a copy of 10 and
replace it with 01.

A Different Strategy



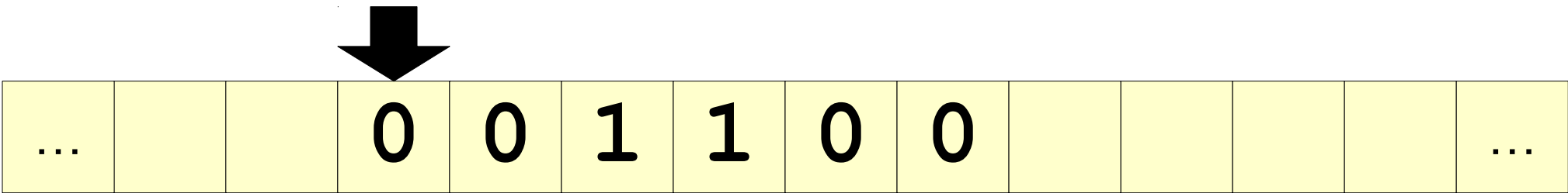
Idea: Repeatedly find
a copy of 10 and
replace it with 01.

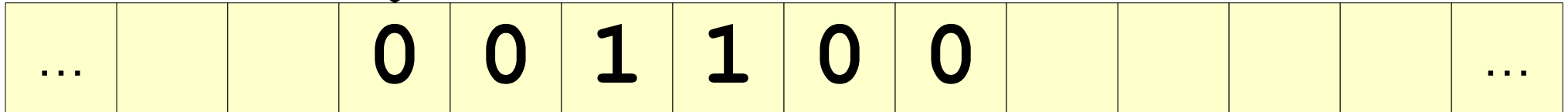
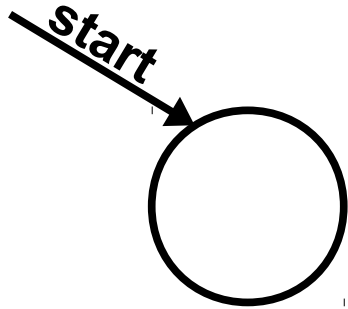
A Different Strategy

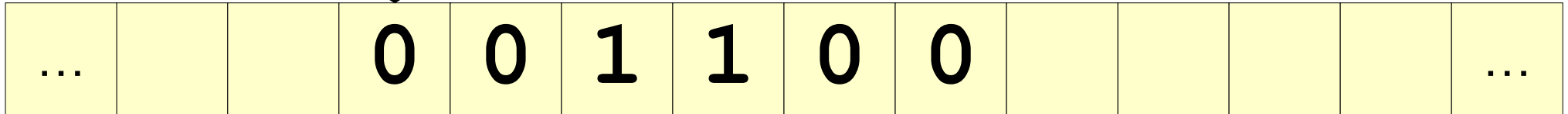
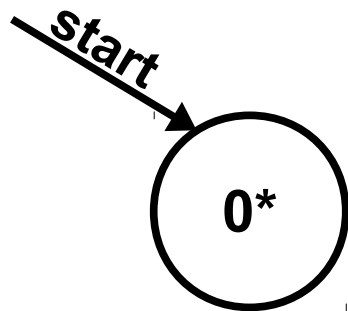


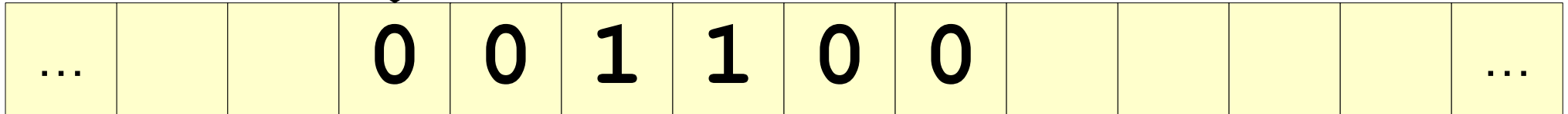
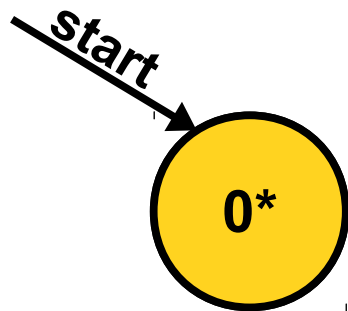
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a copy of 10 and
replace it with 01.

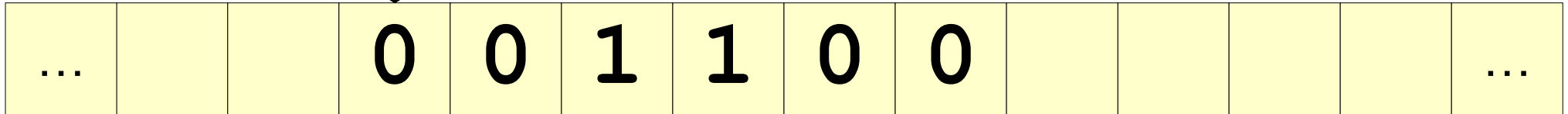
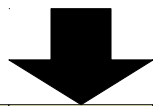
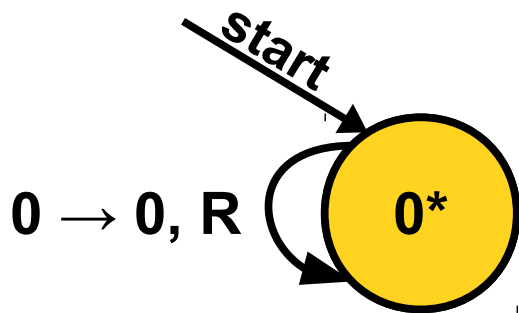
Let's Build It!

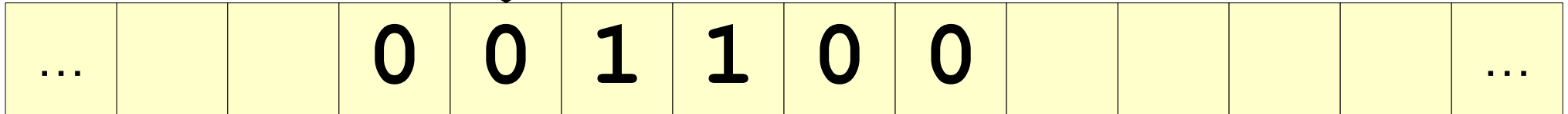
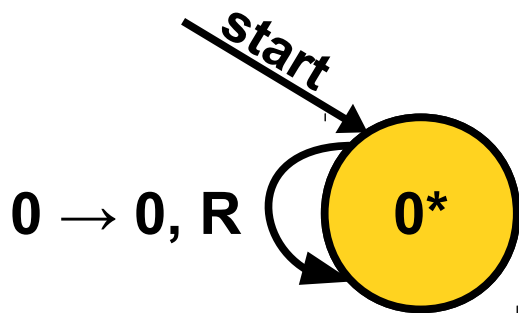


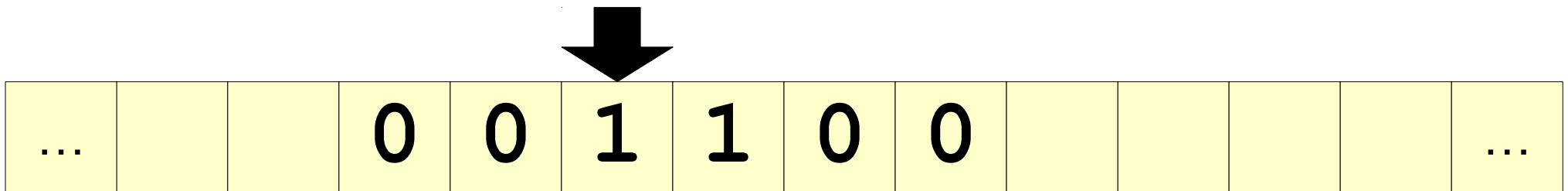
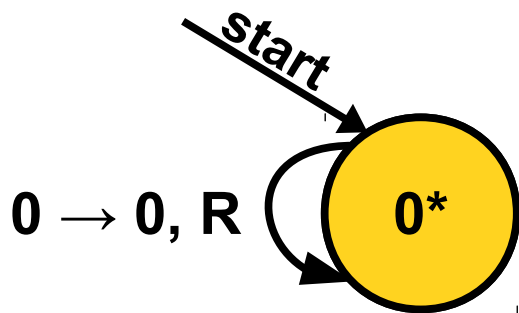


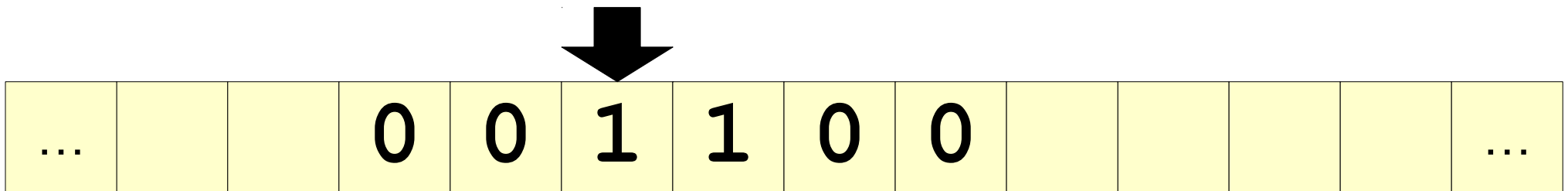
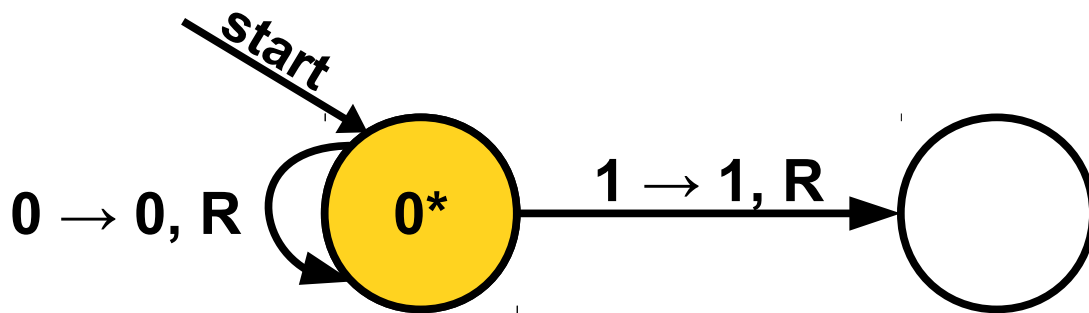


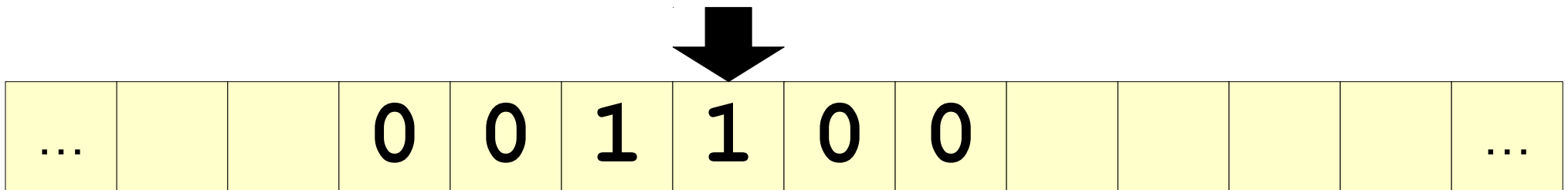
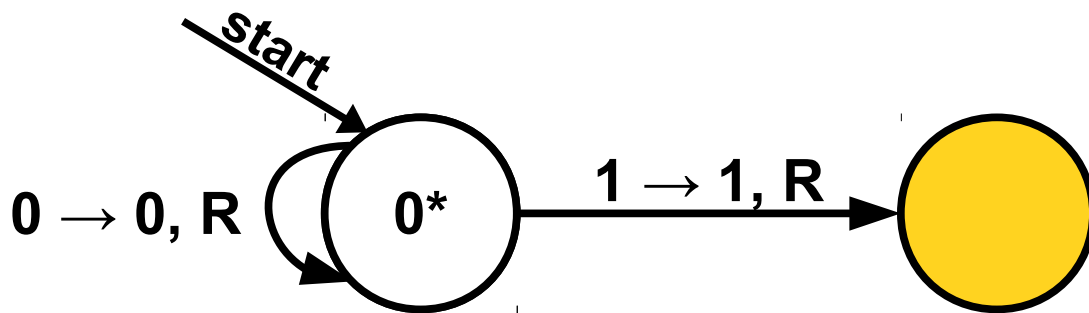


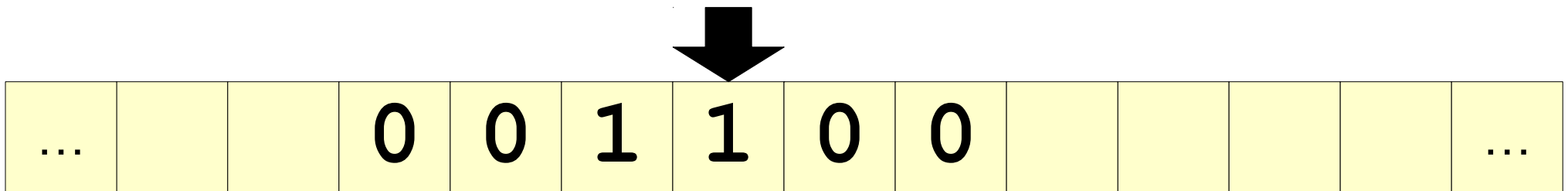
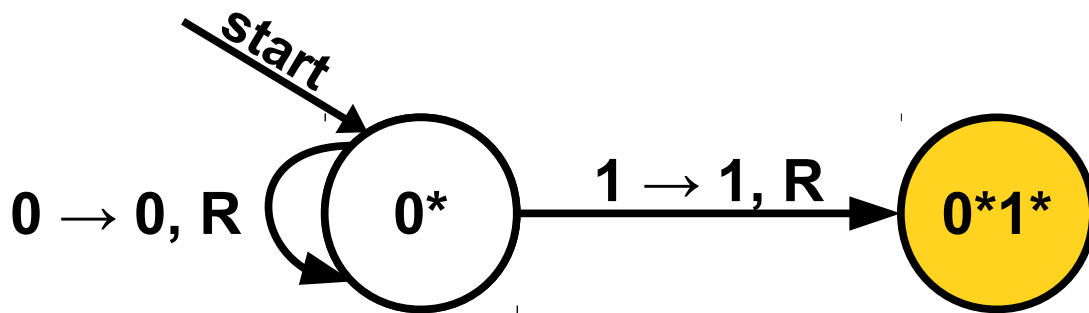


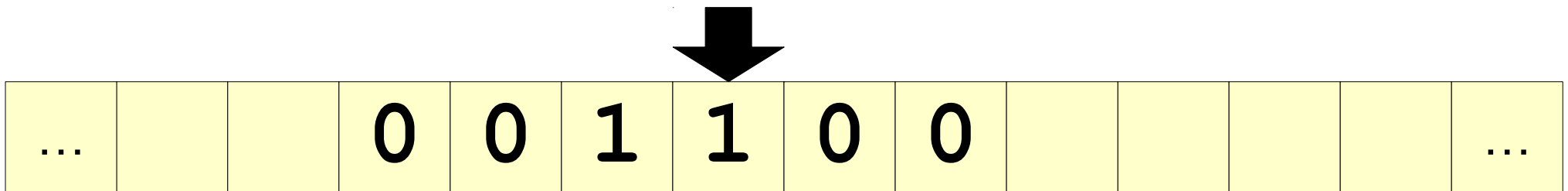
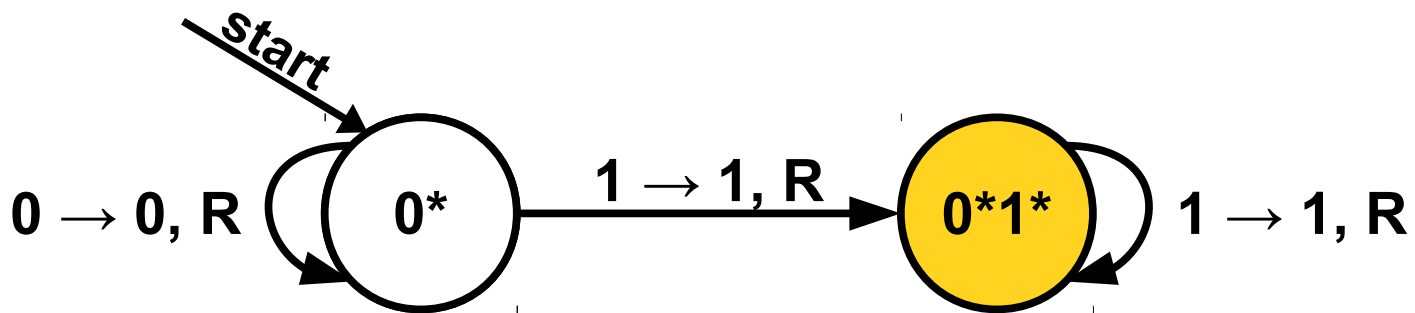


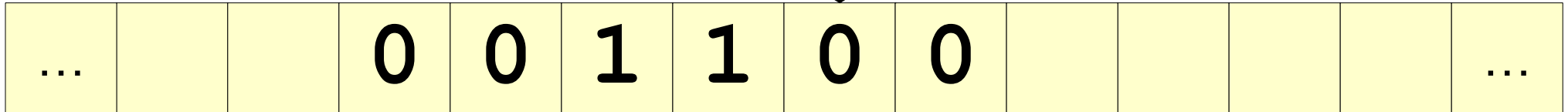
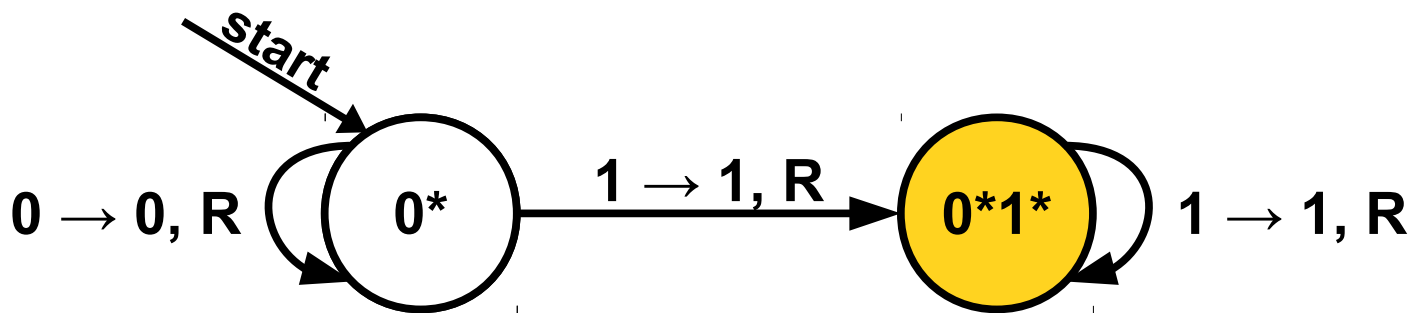


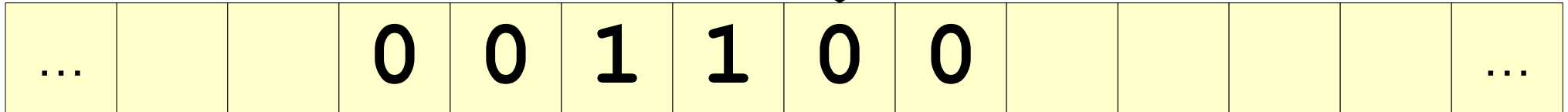
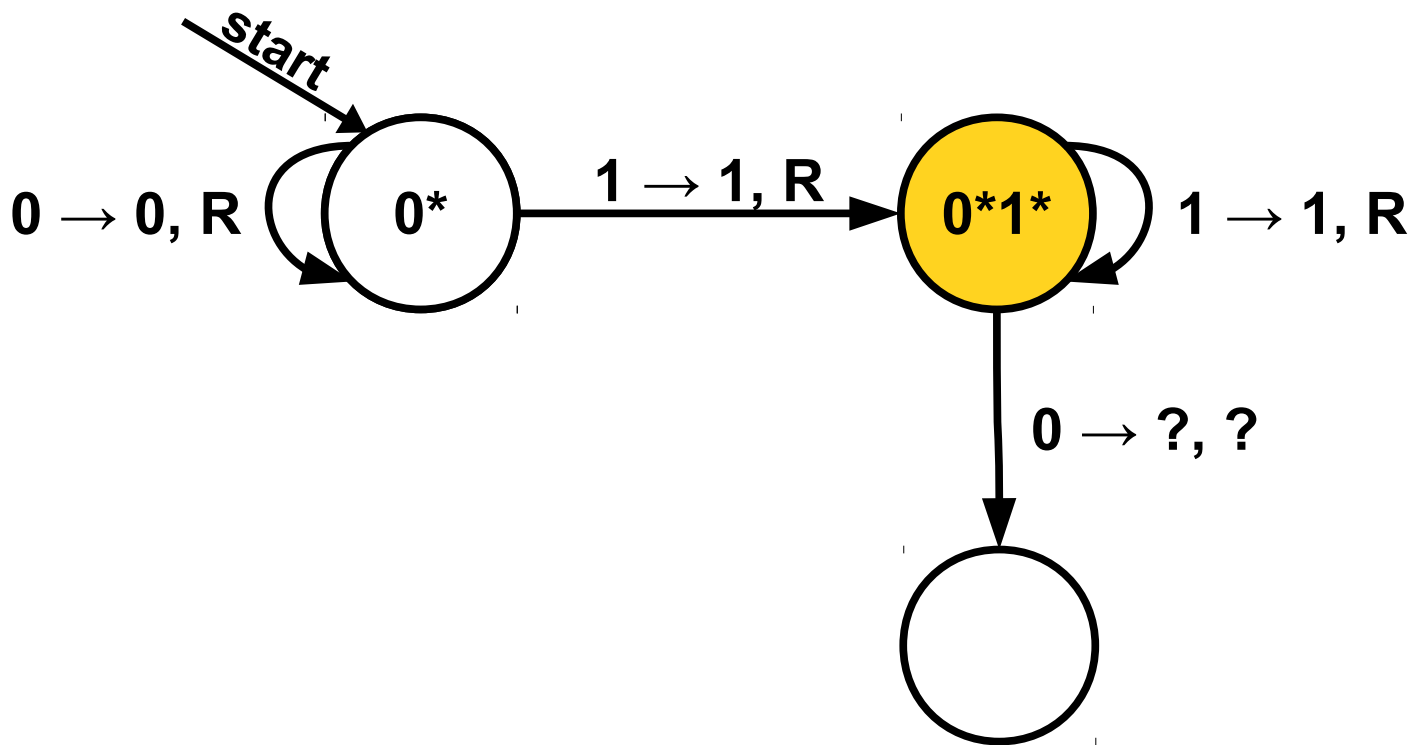


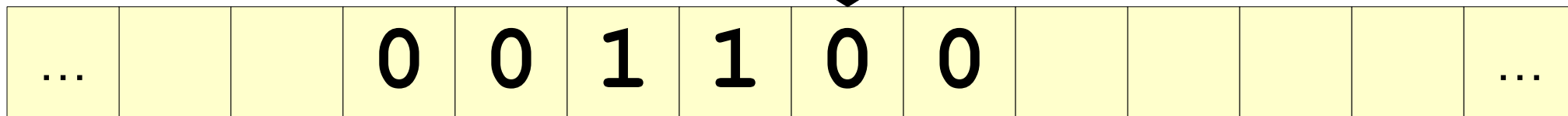
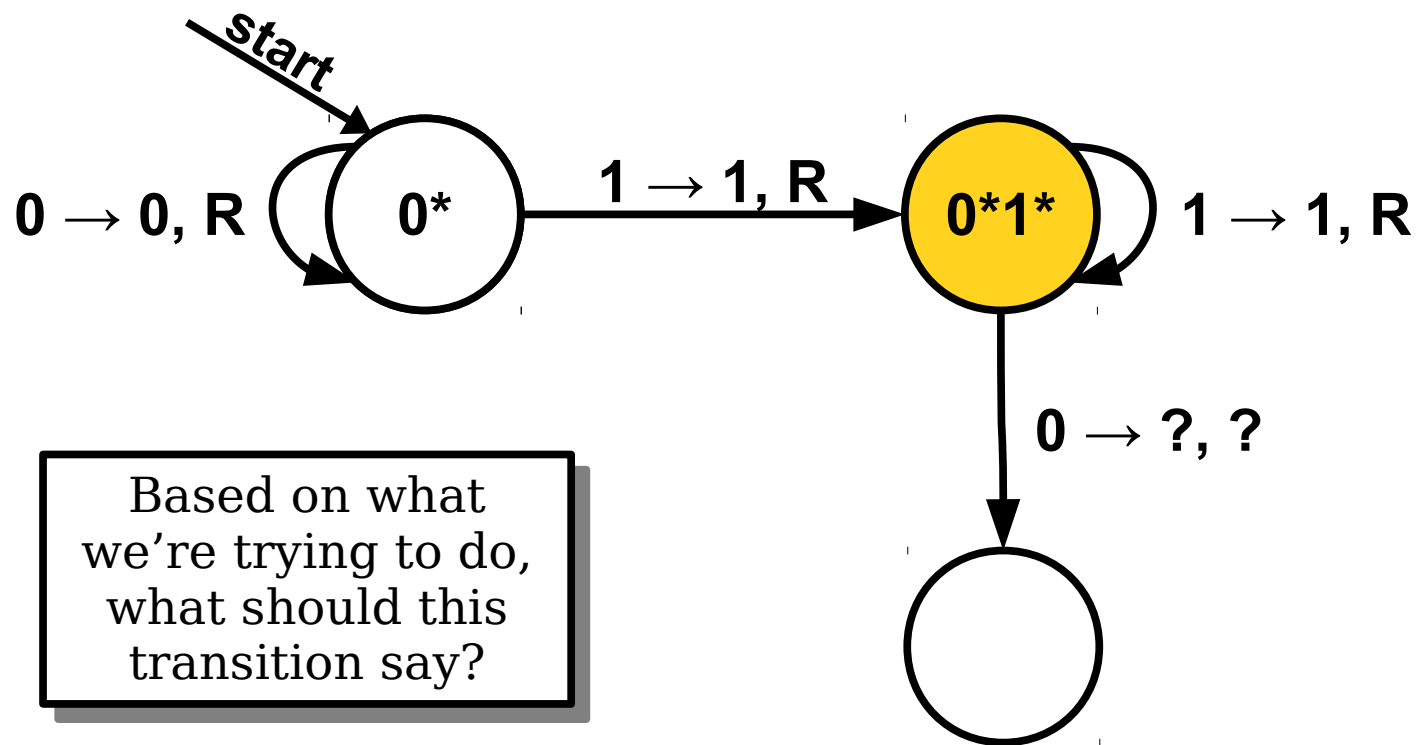


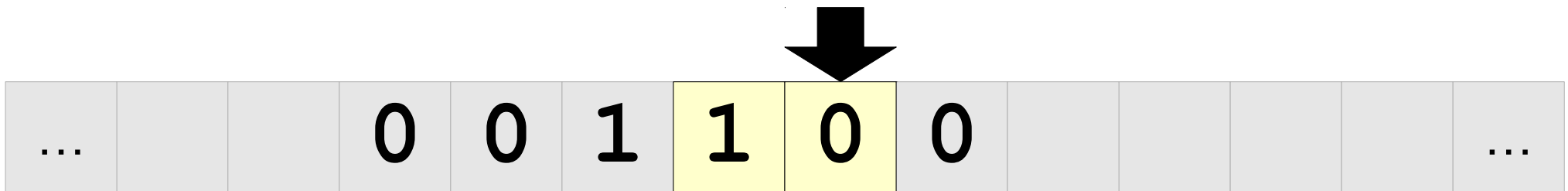
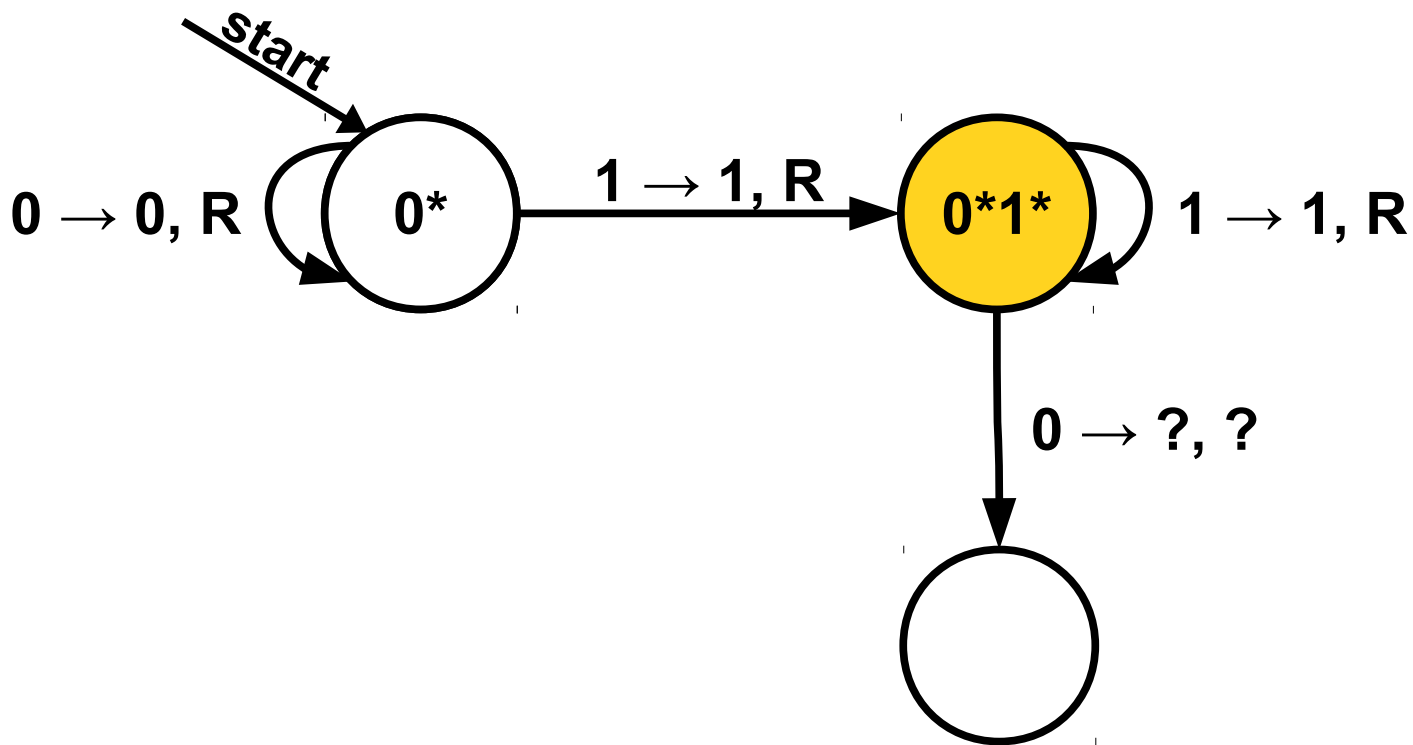


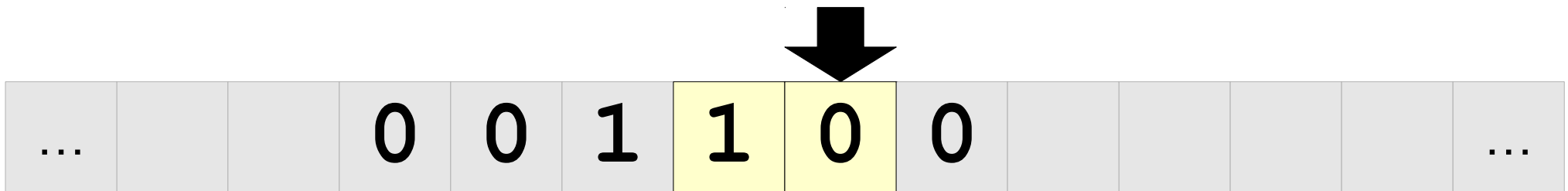
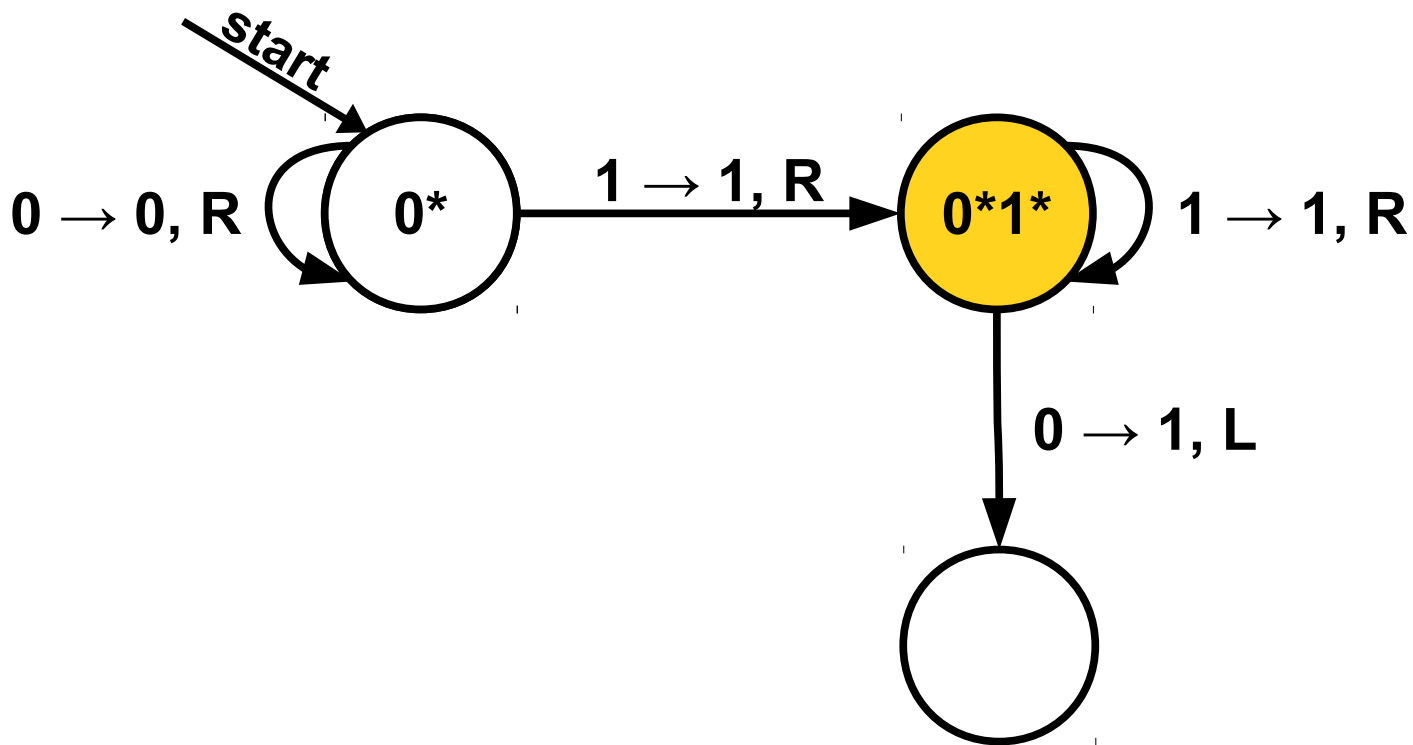


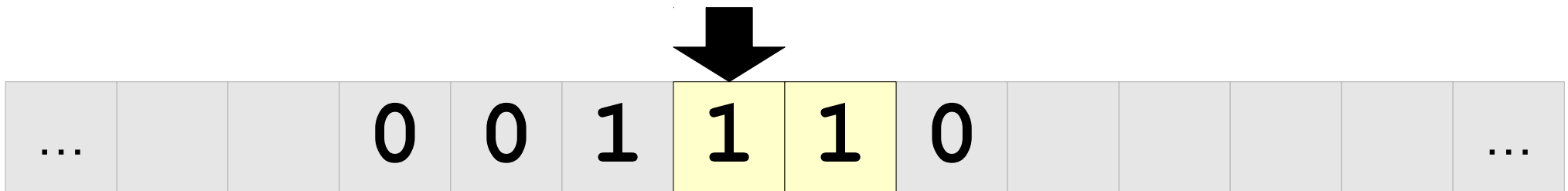
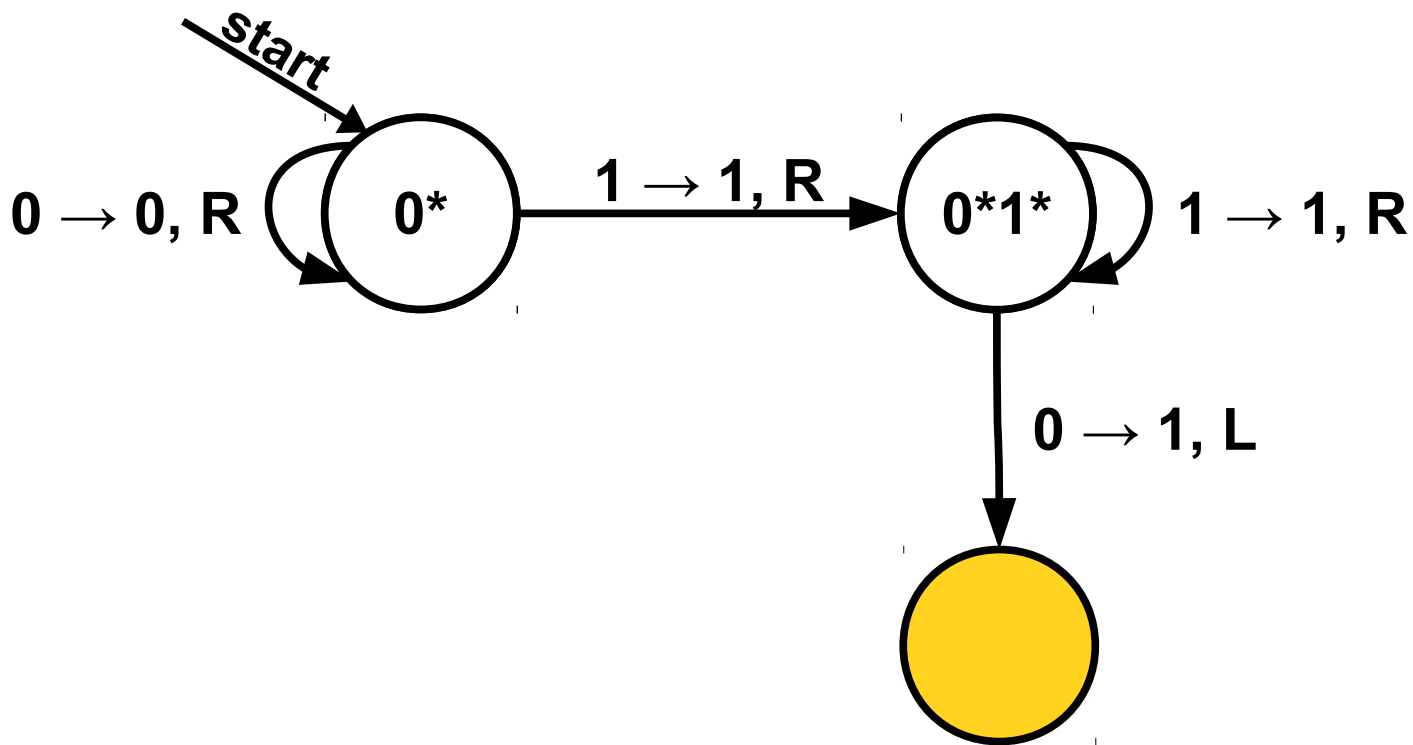


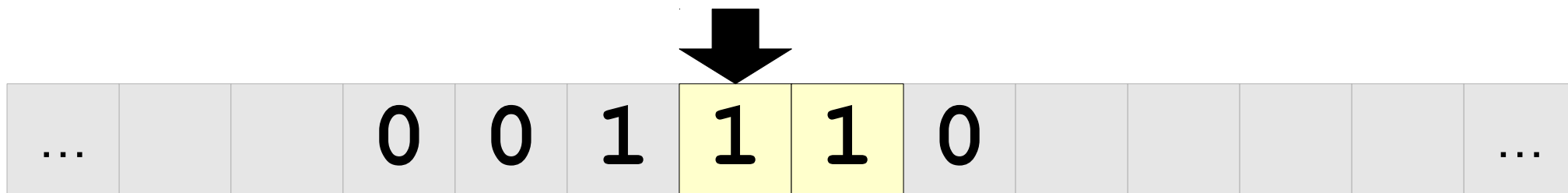
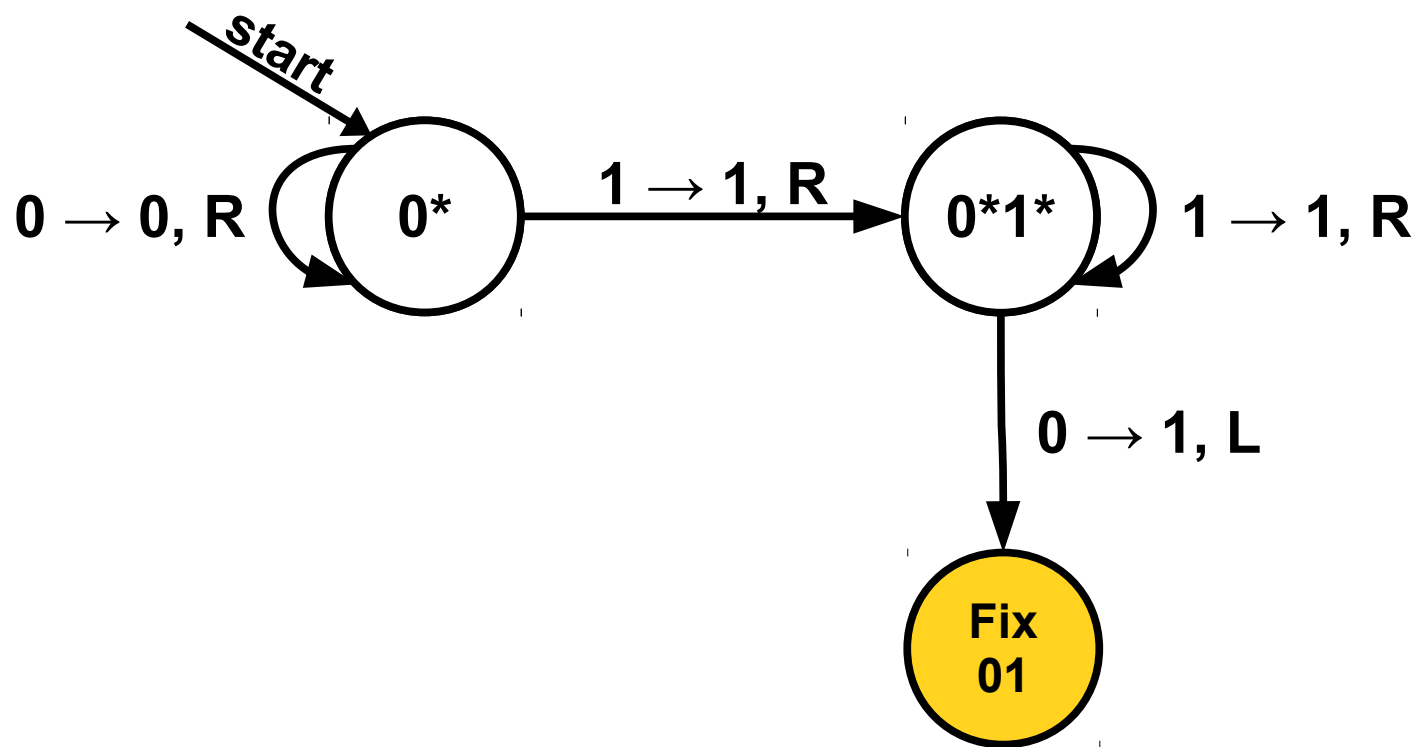


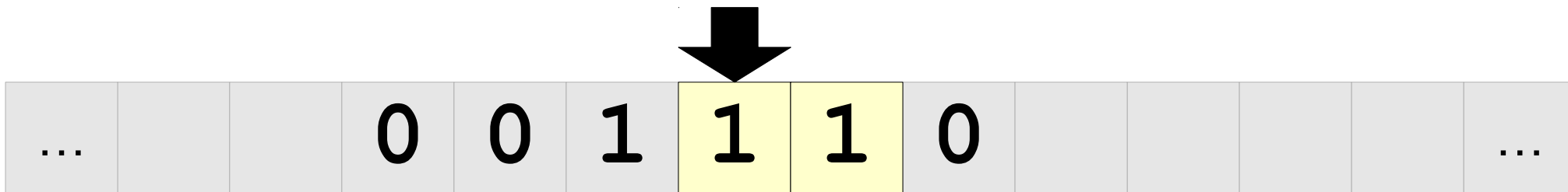
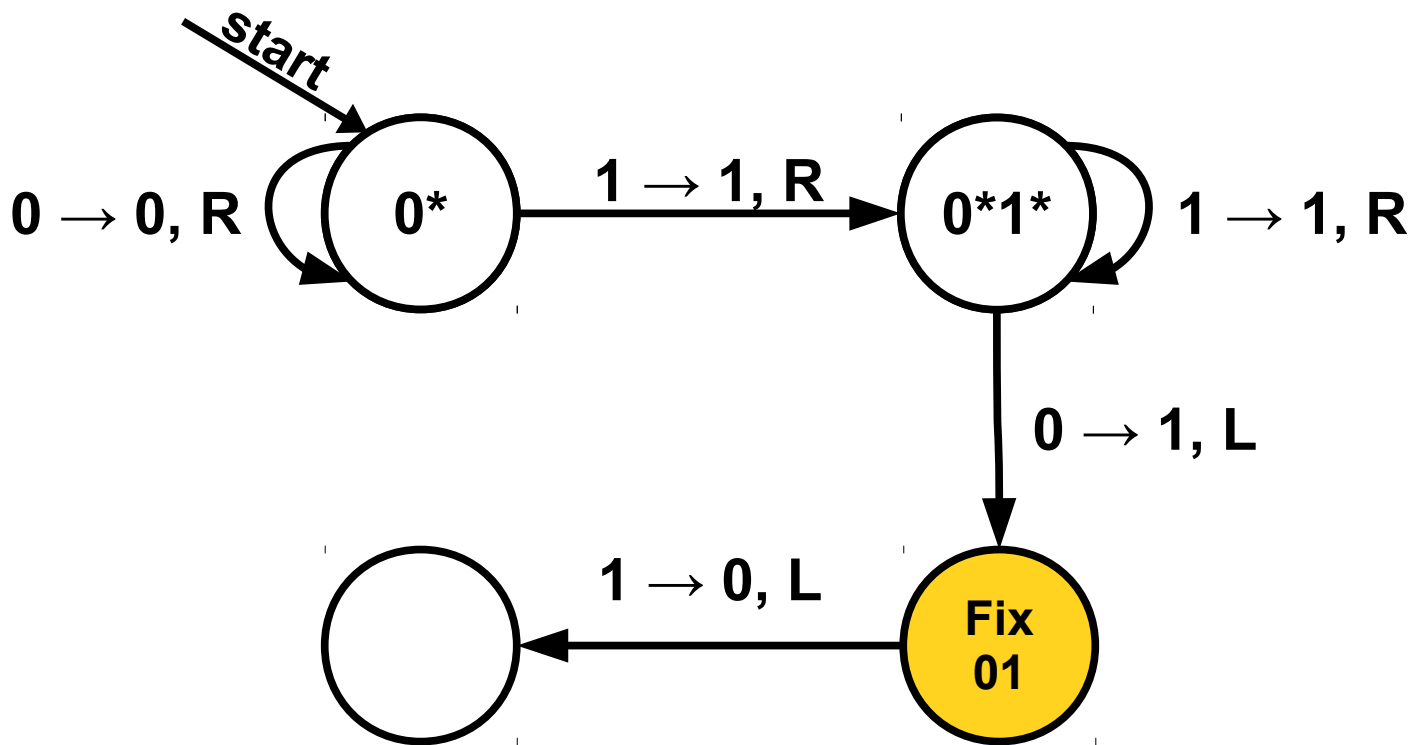


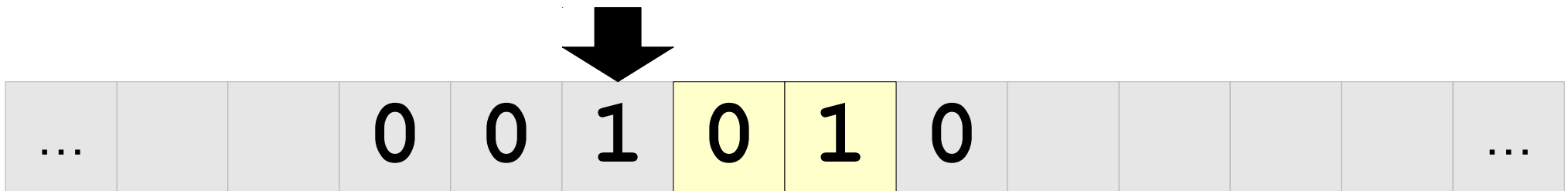
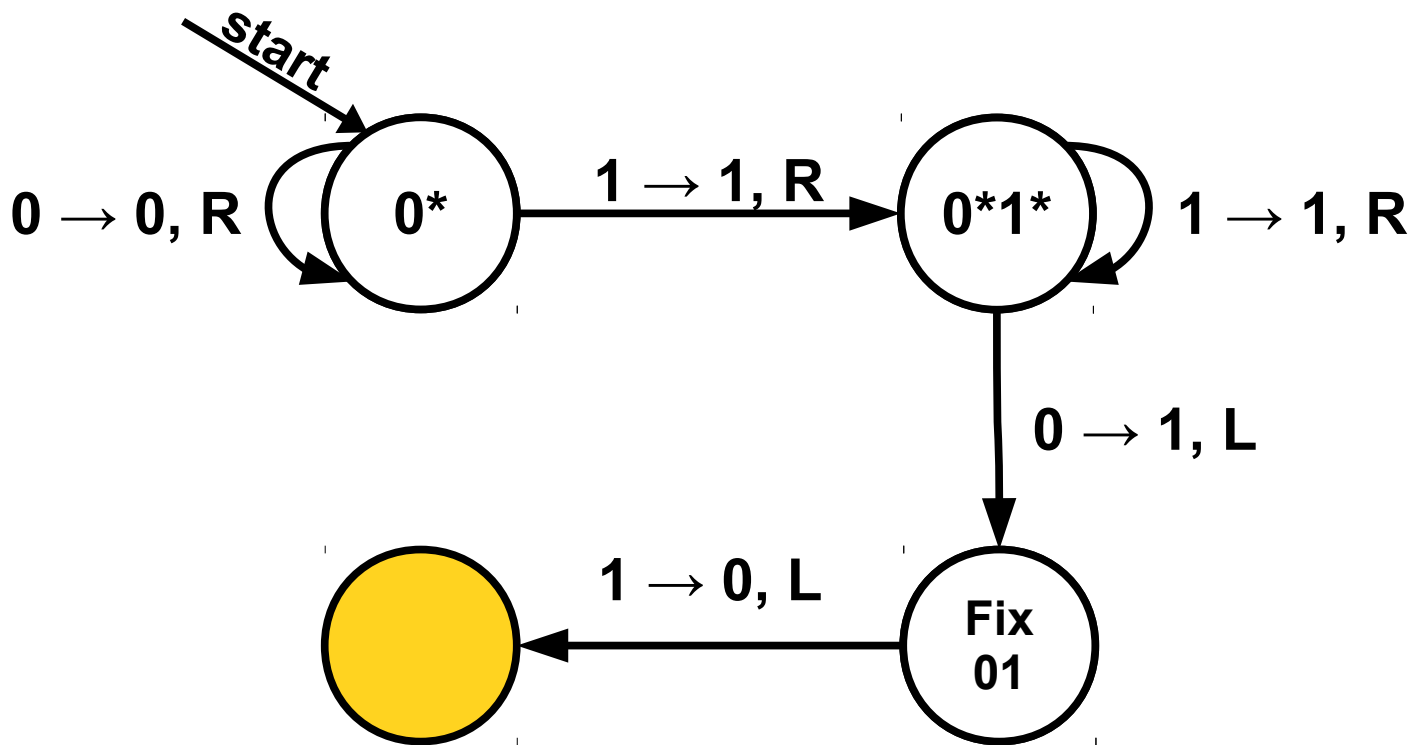


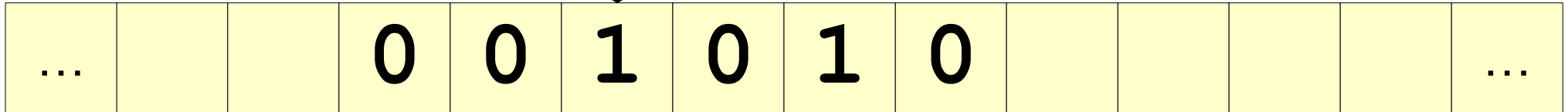
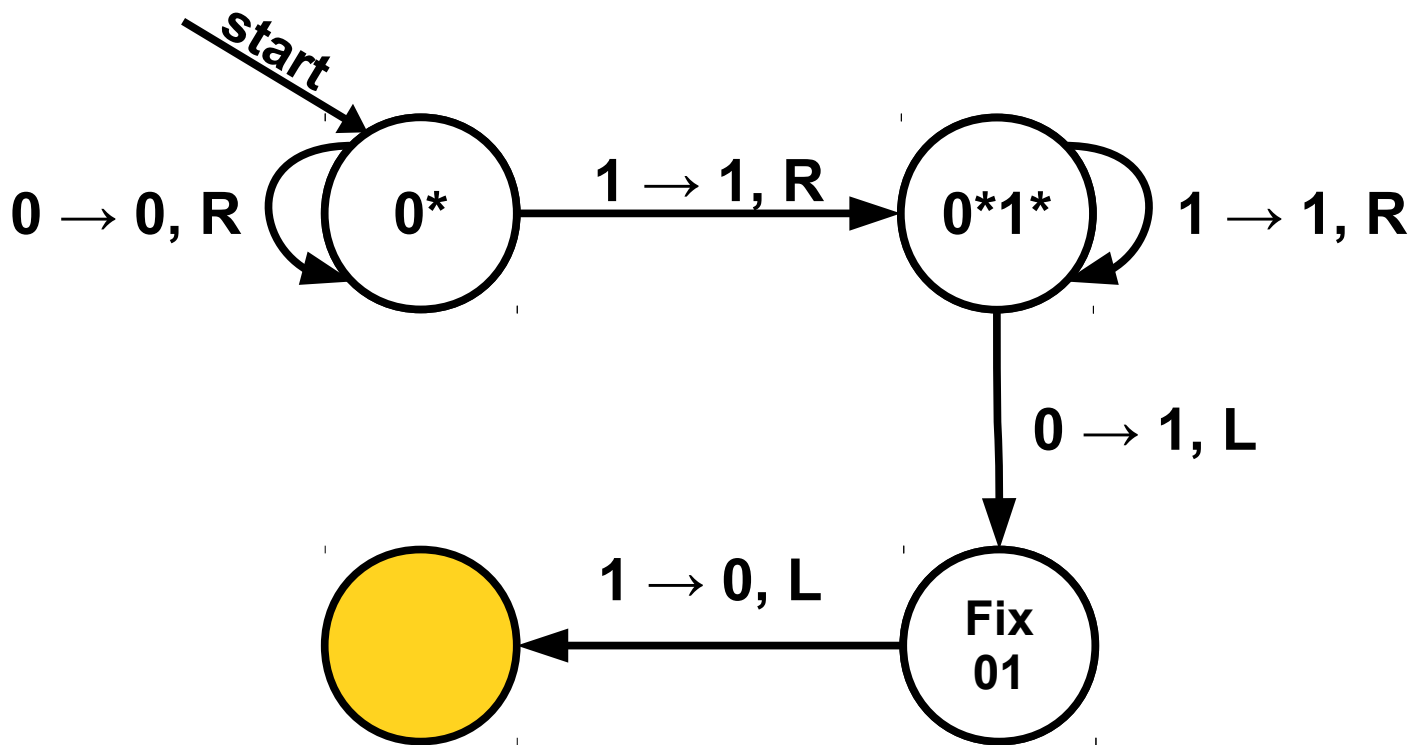


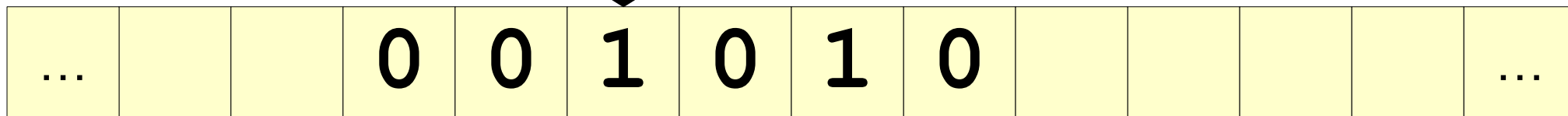
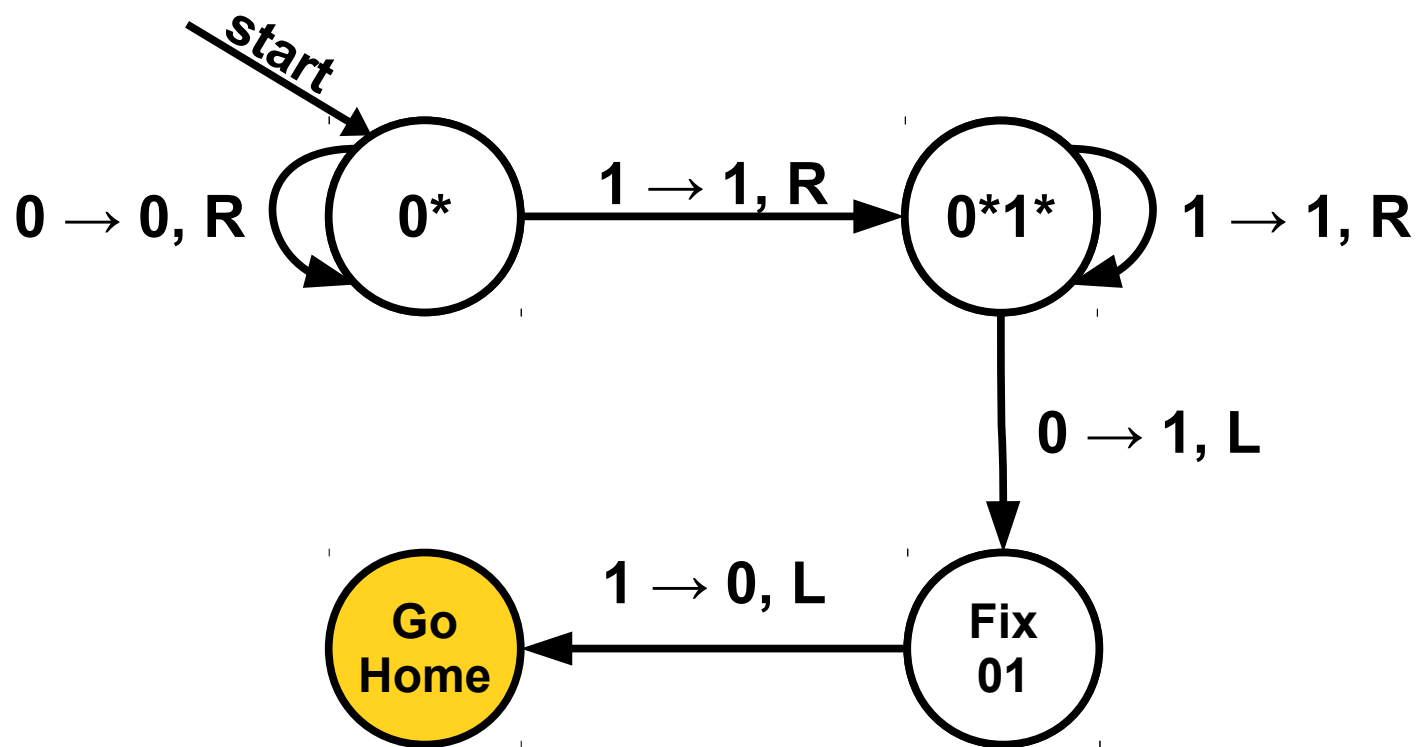


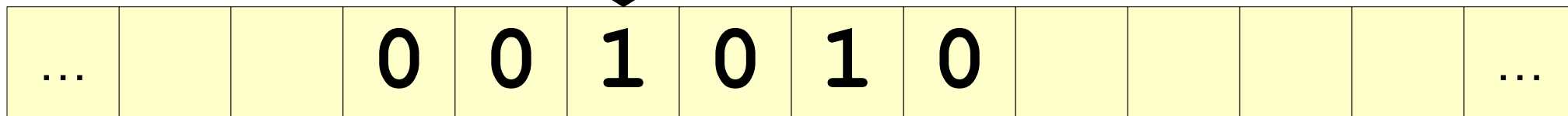
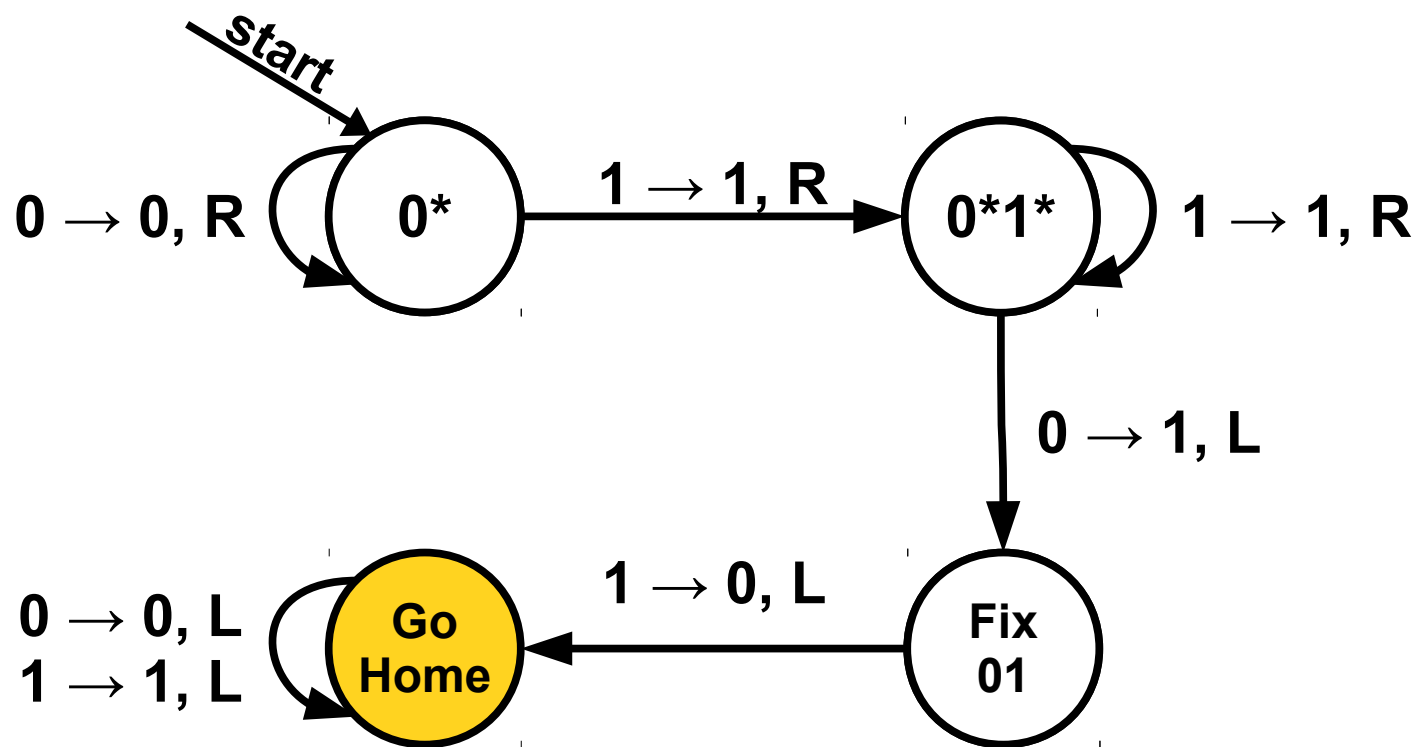


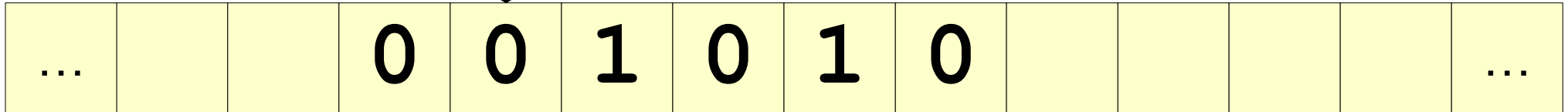
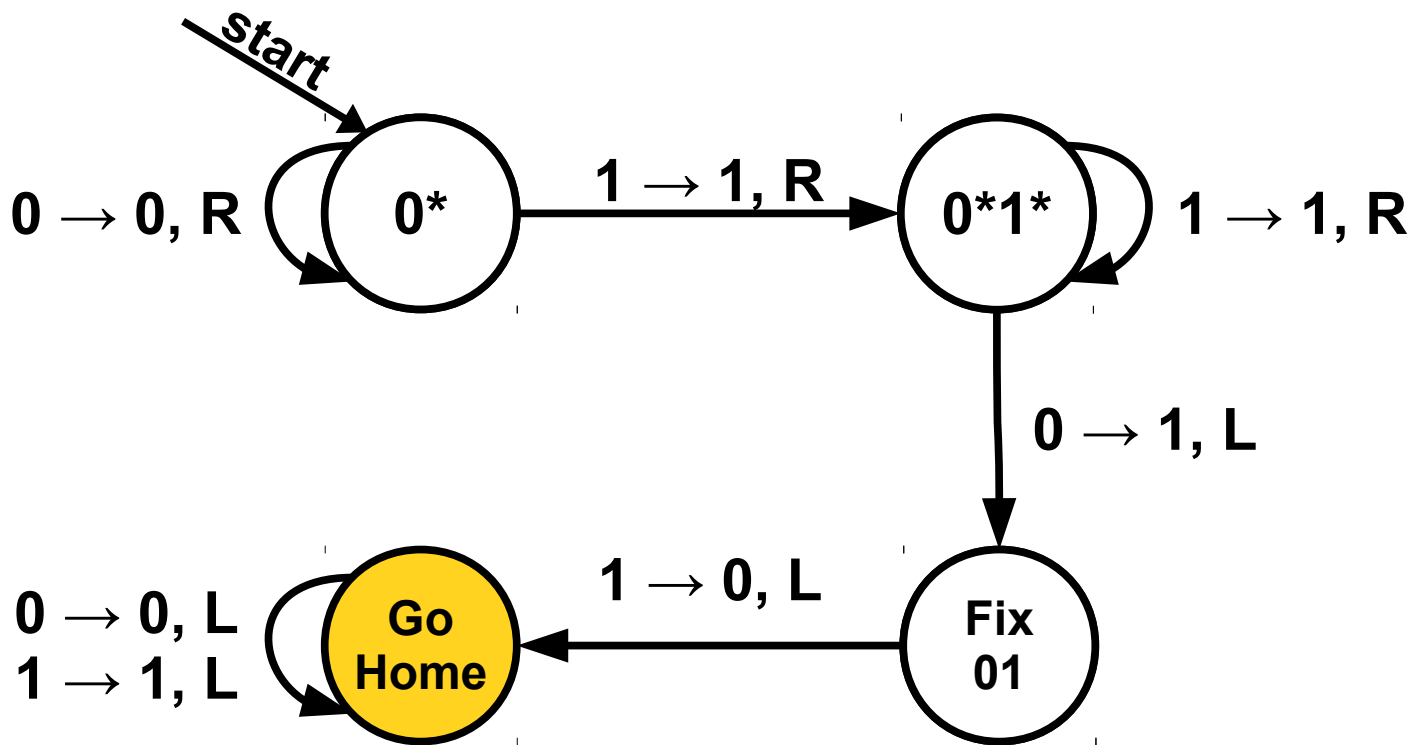


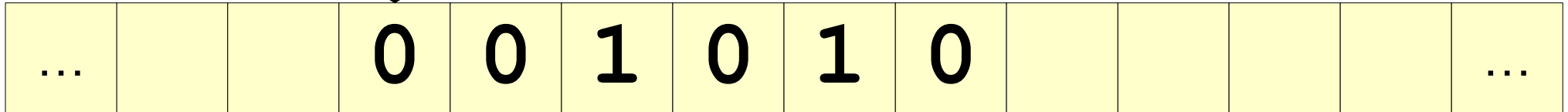
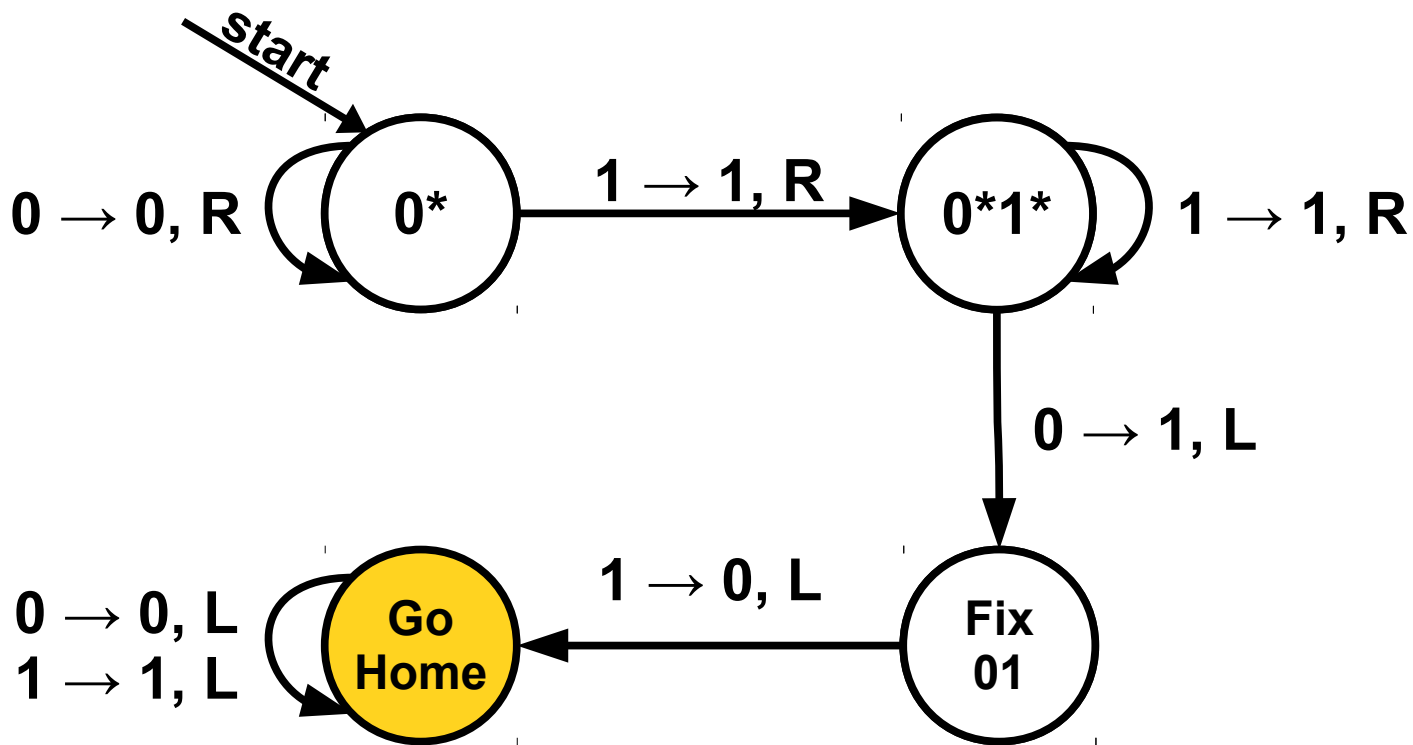


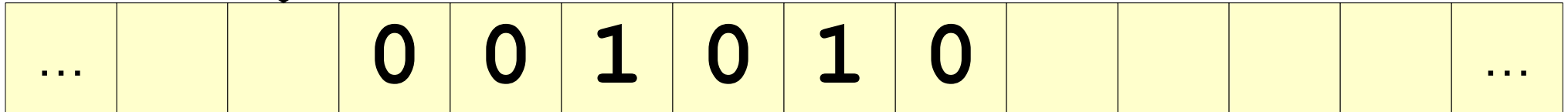
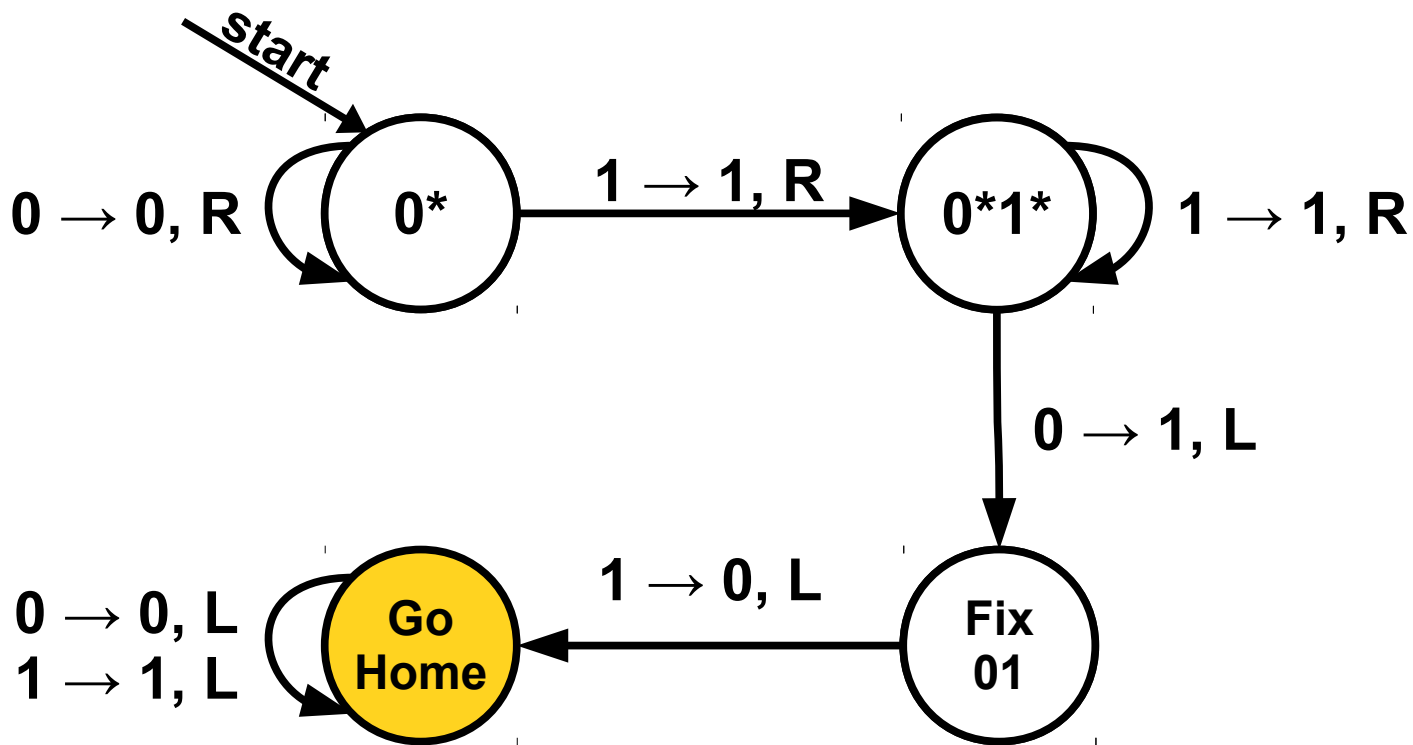


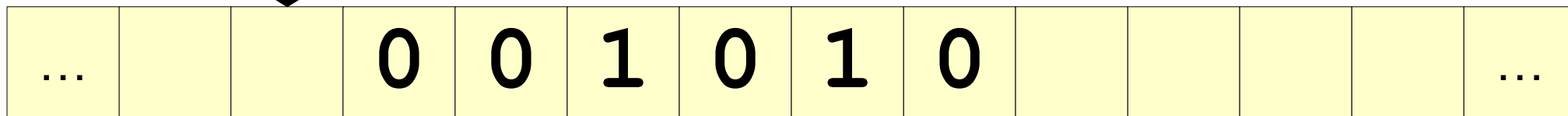
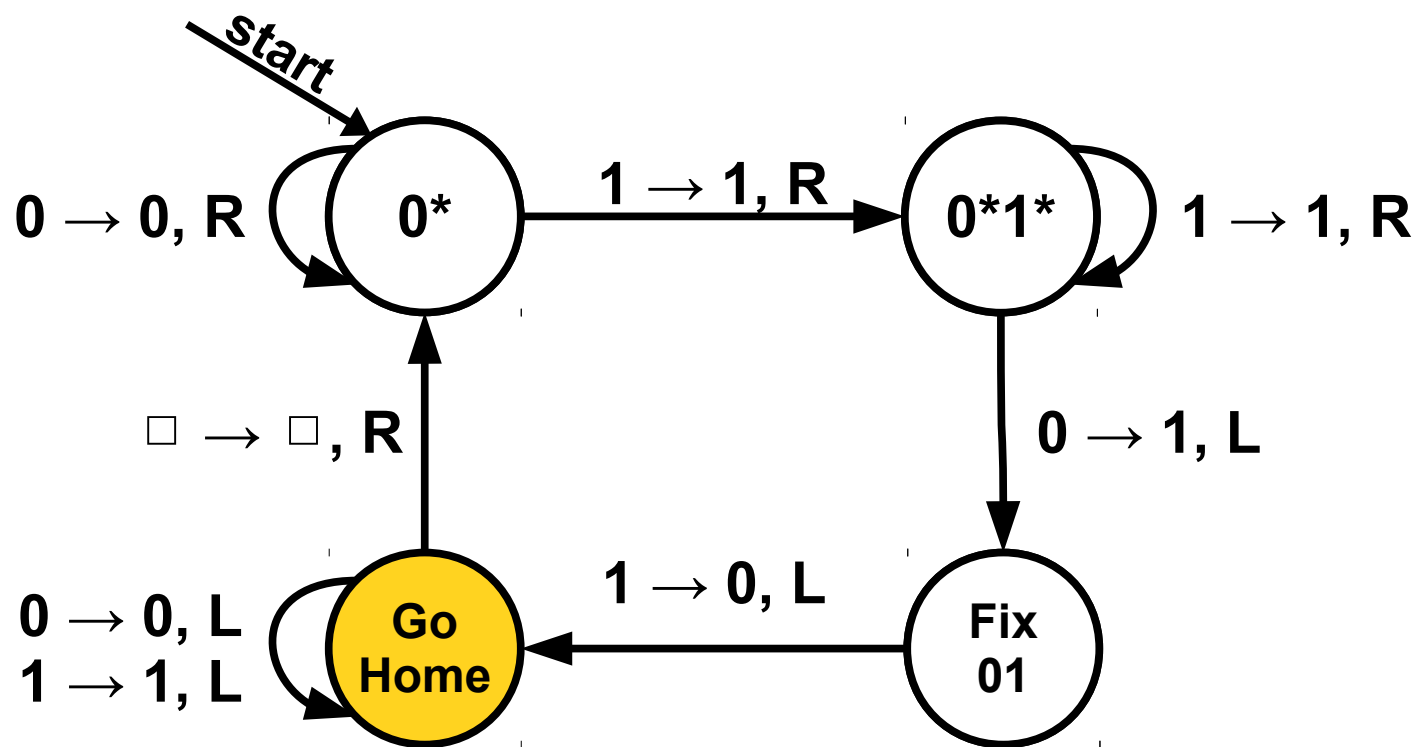


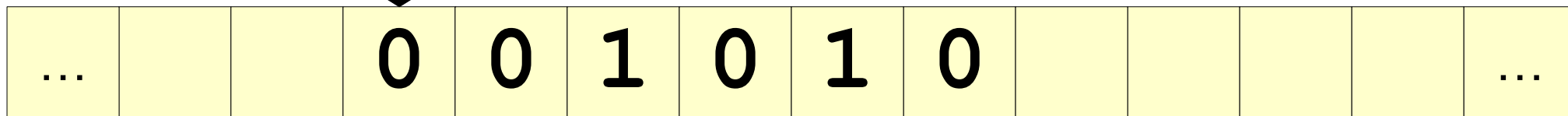
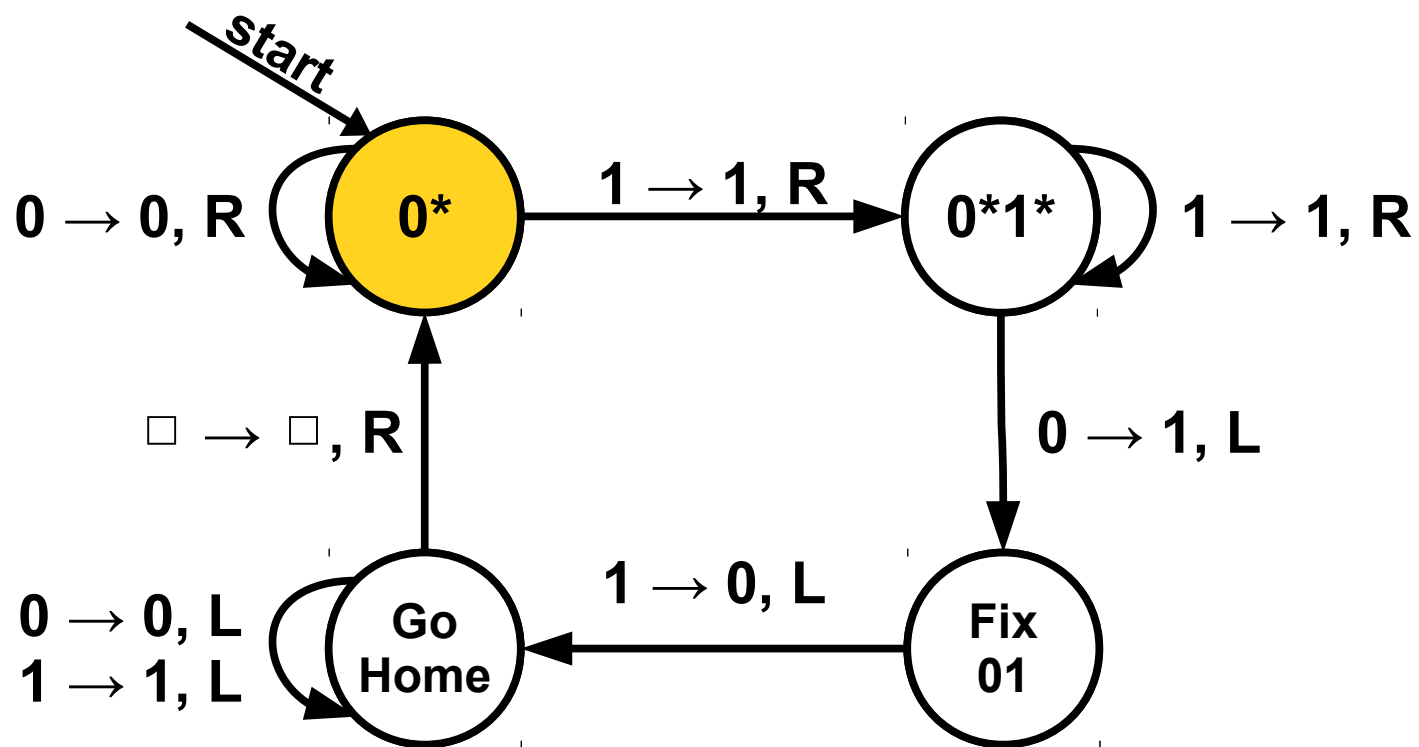


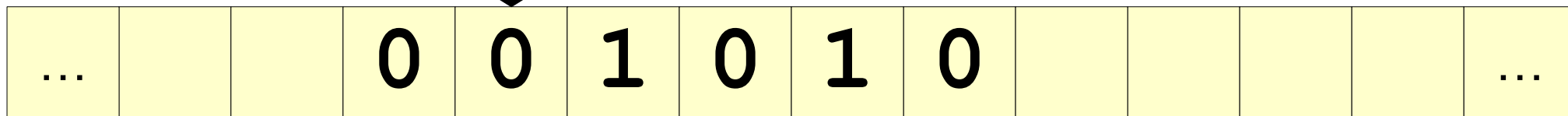
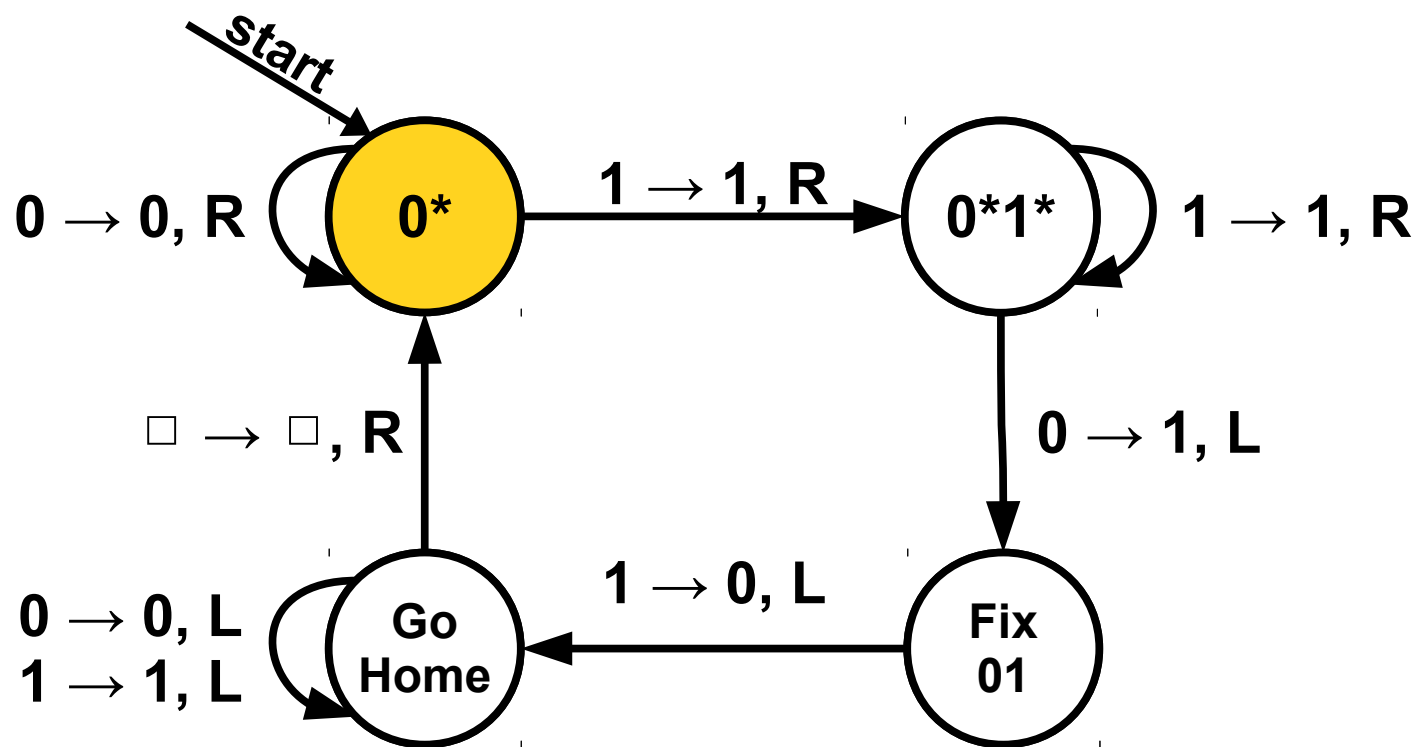


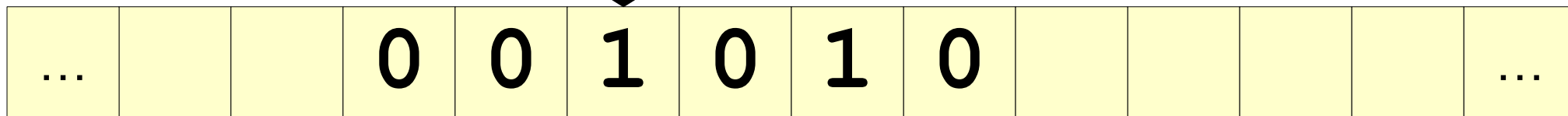
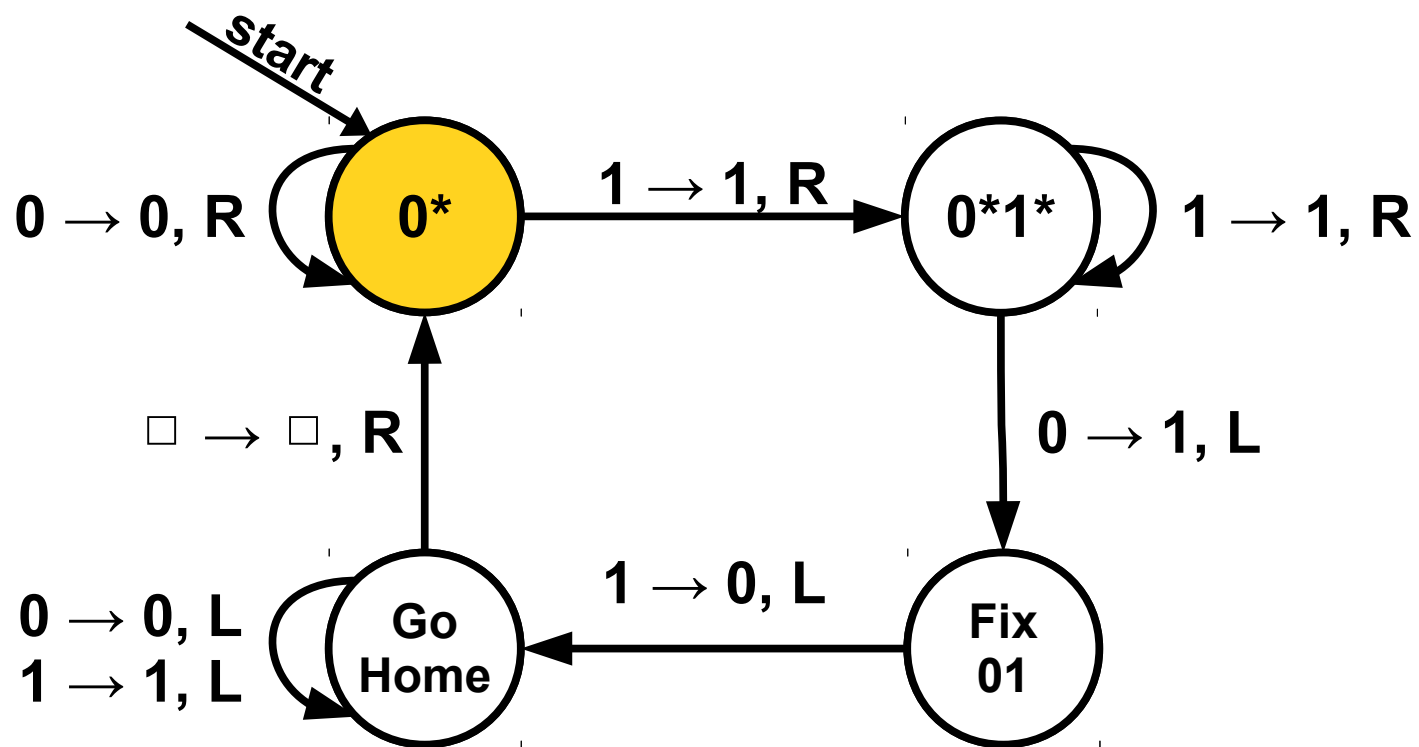


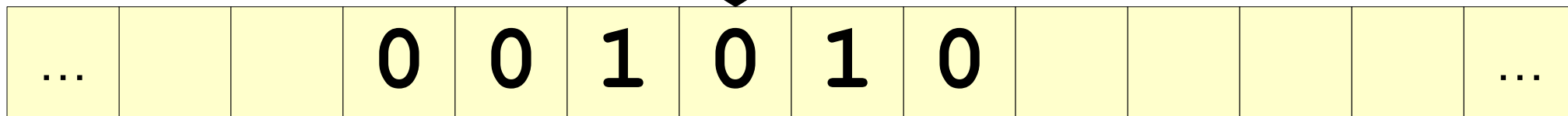
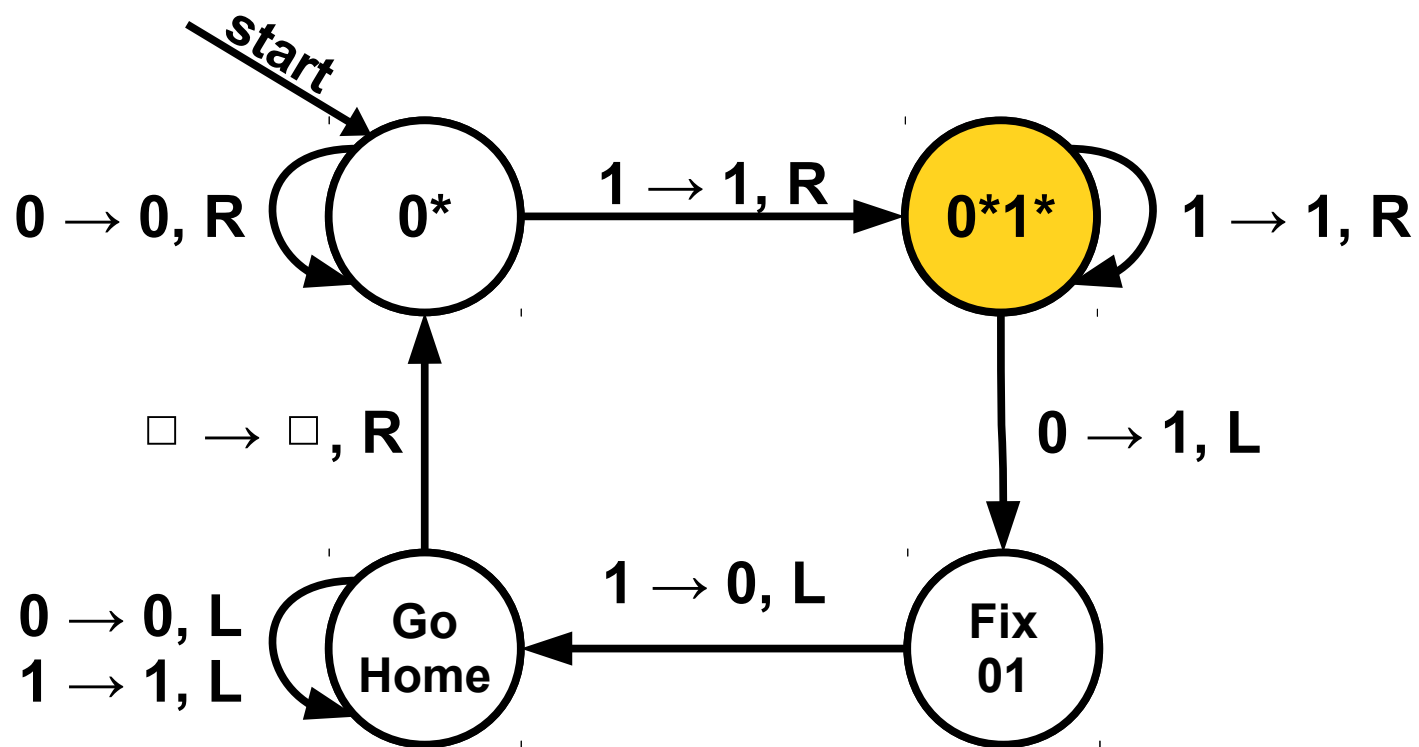


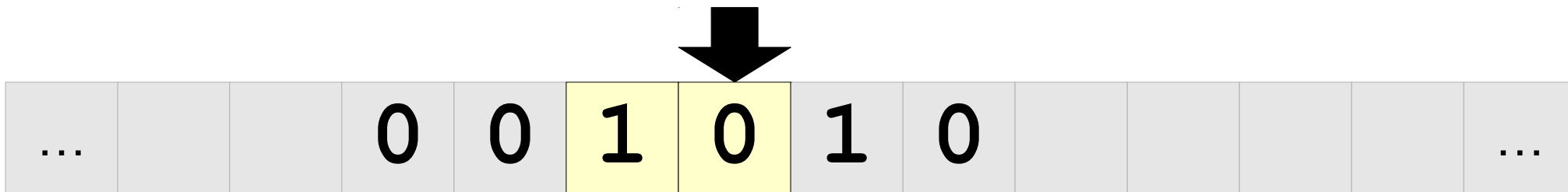
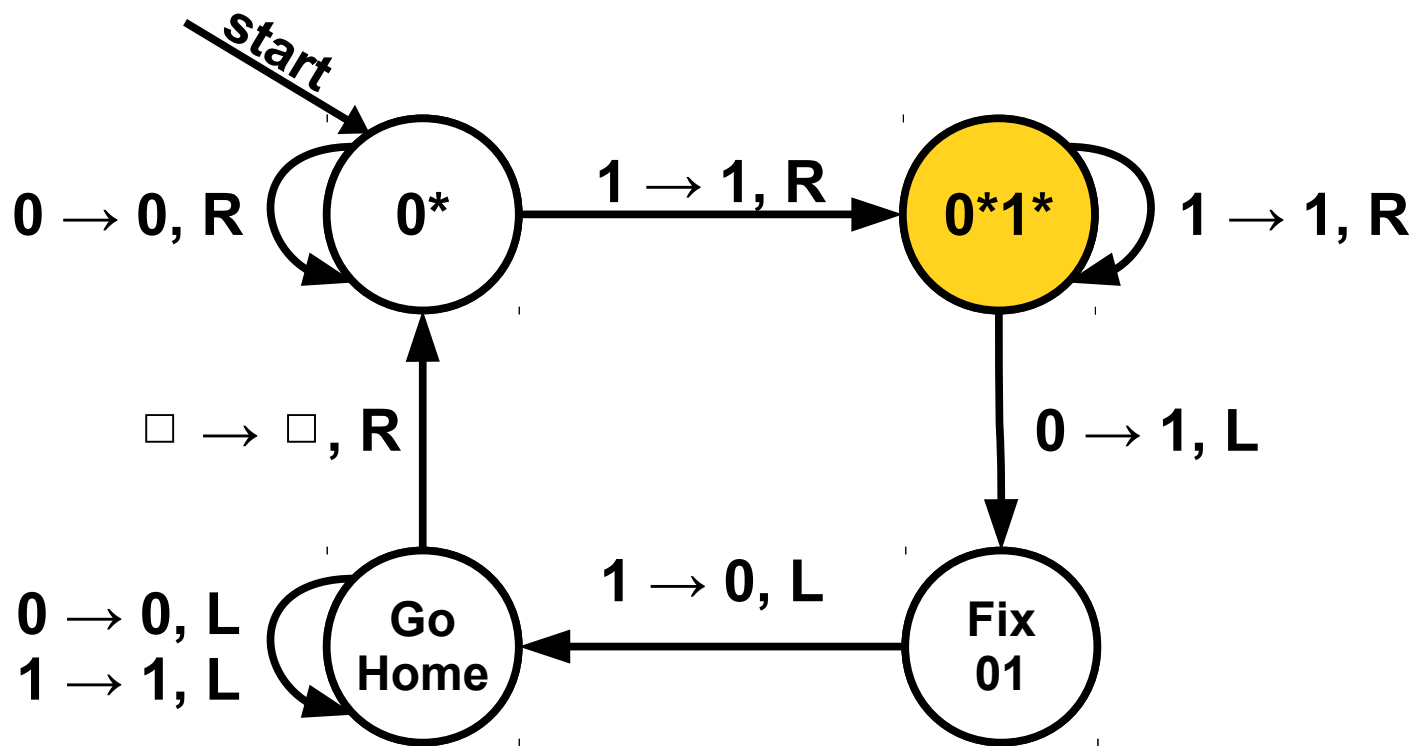


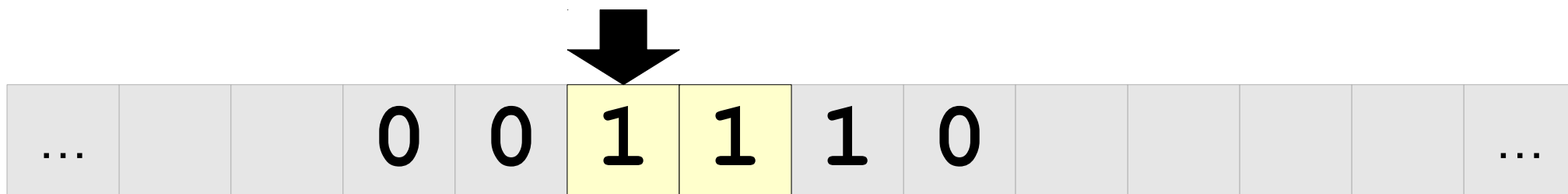
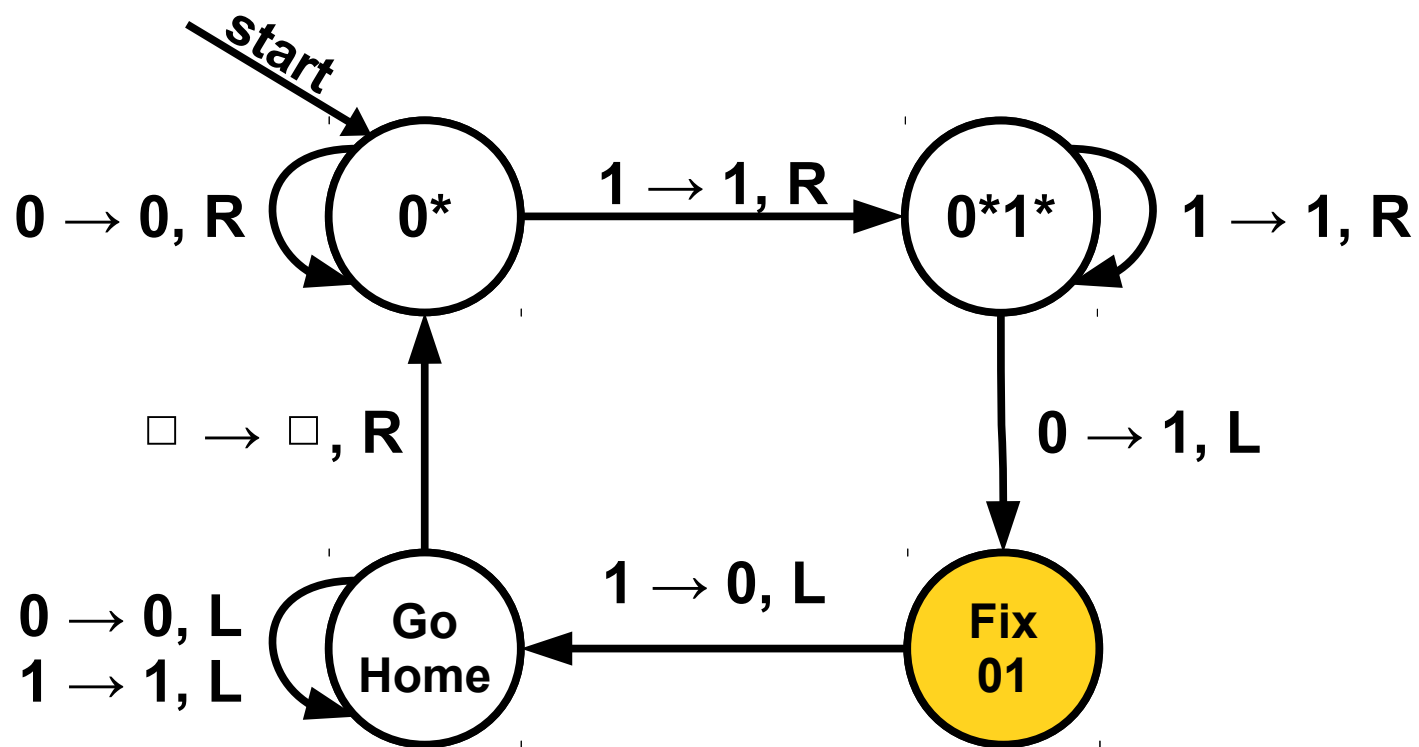


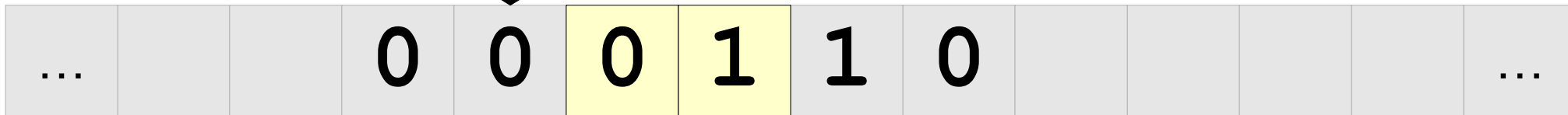
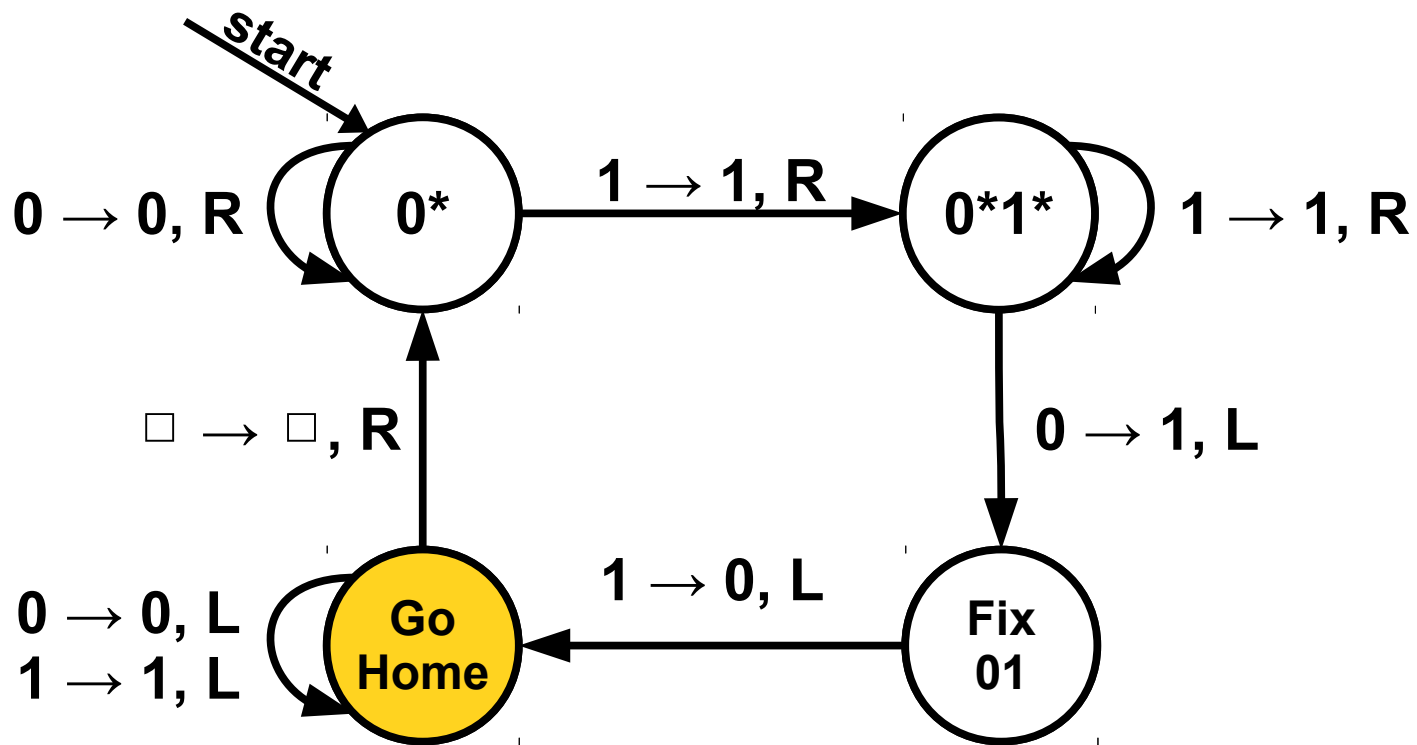


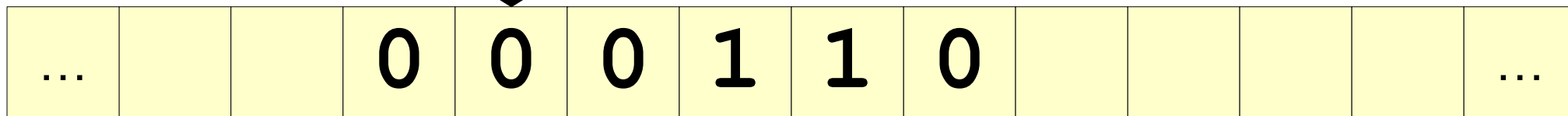
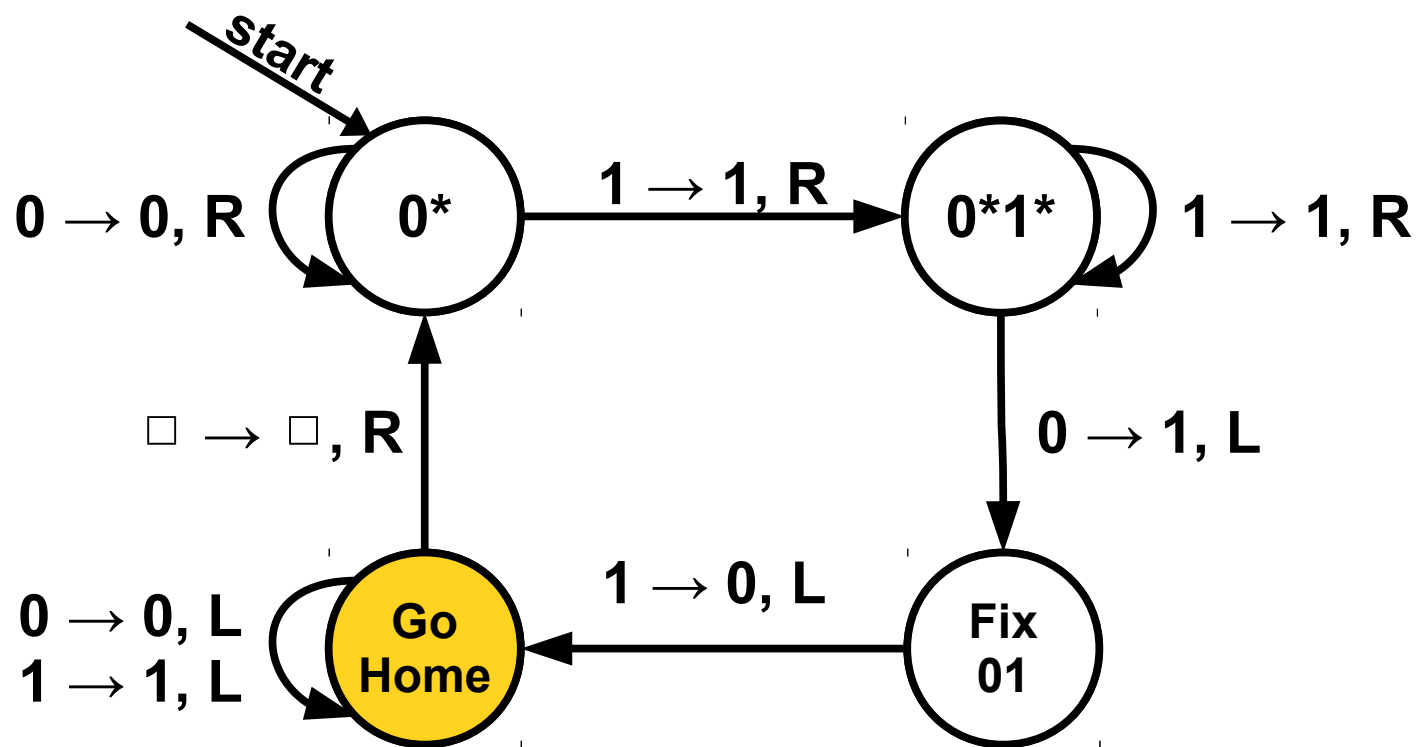


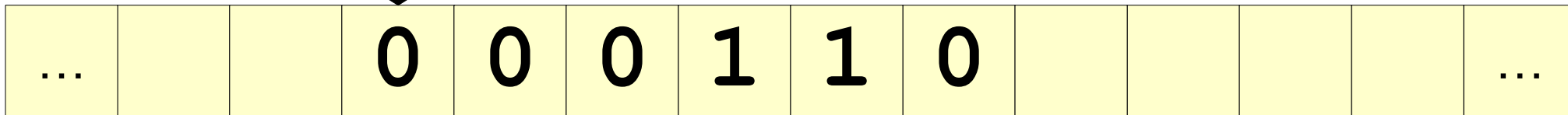
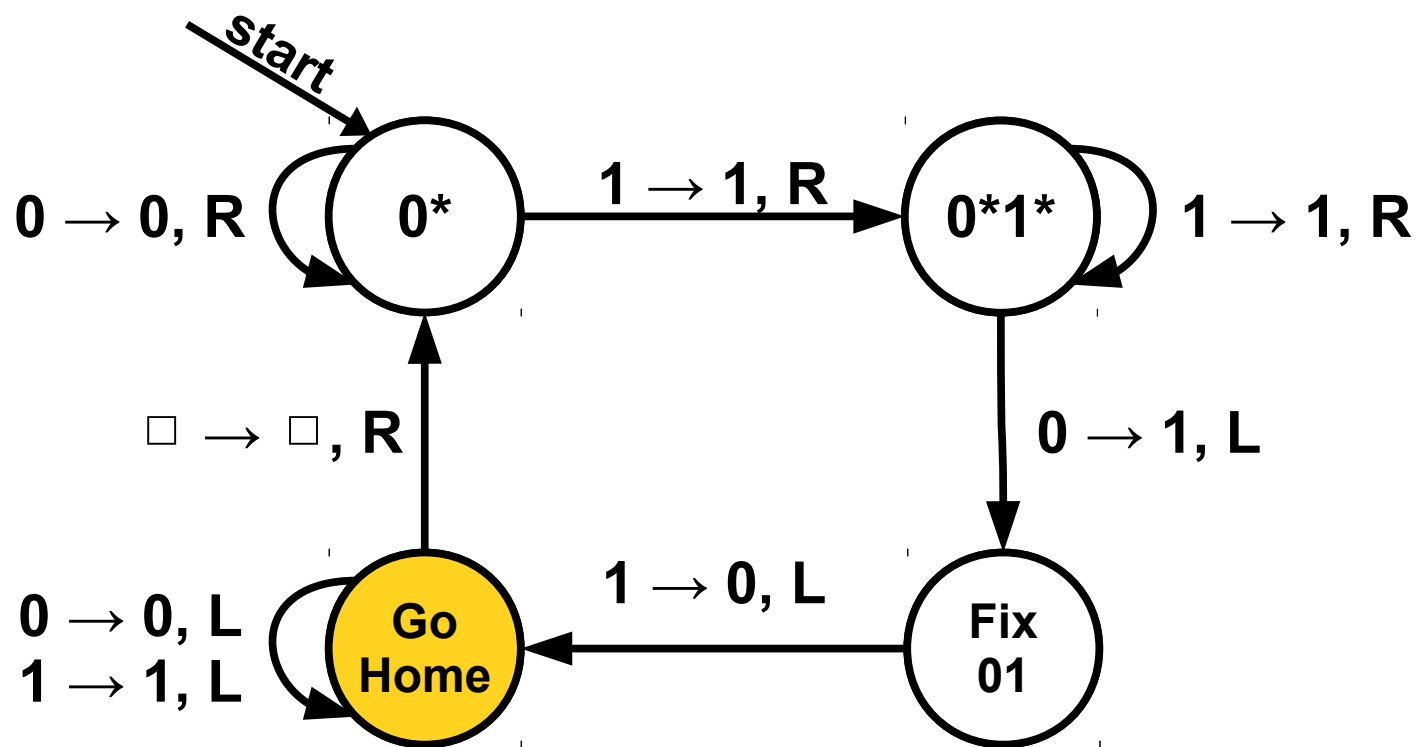


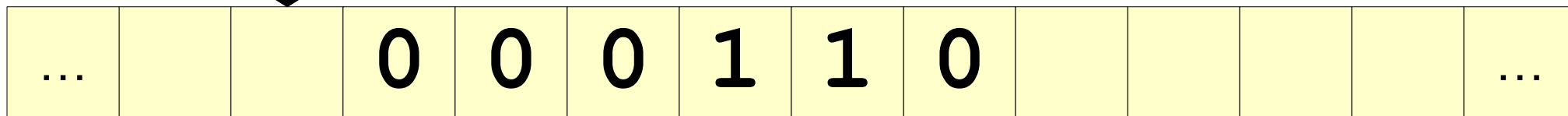
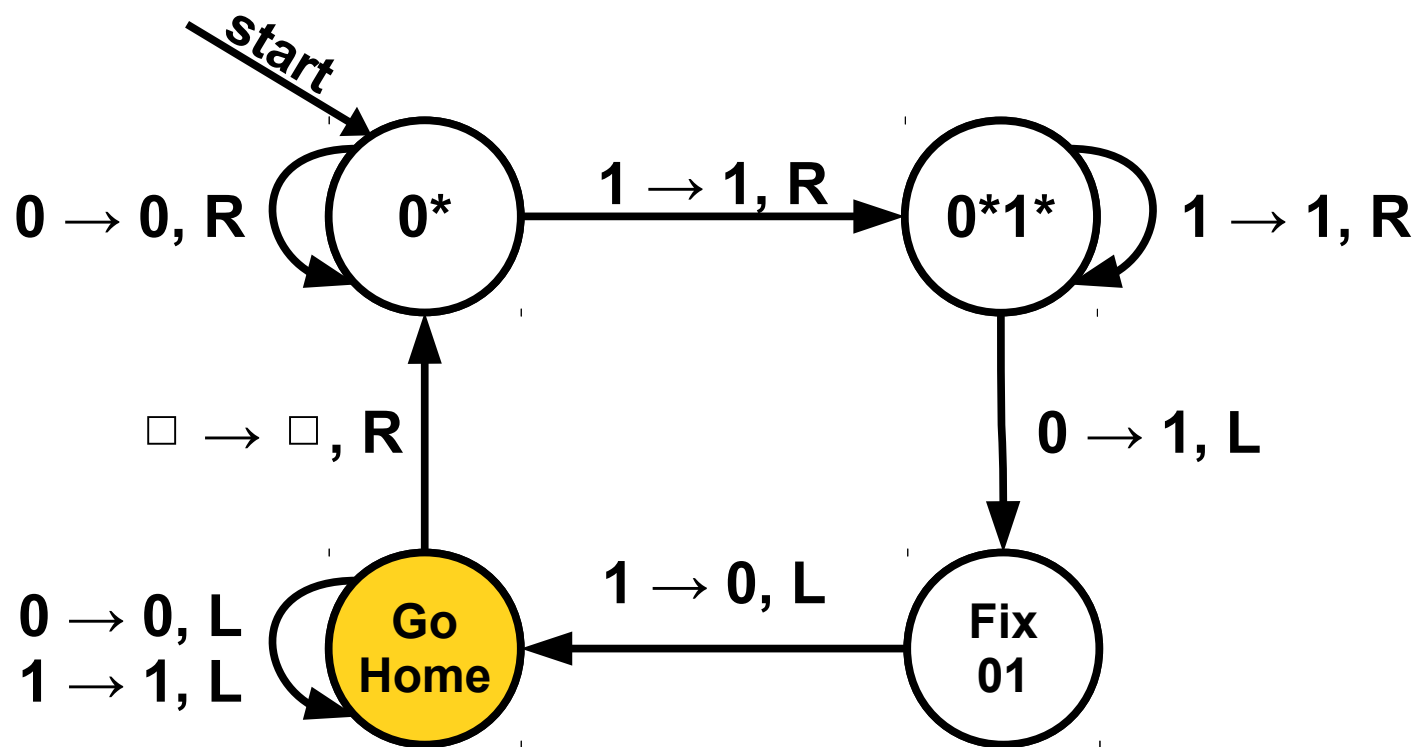


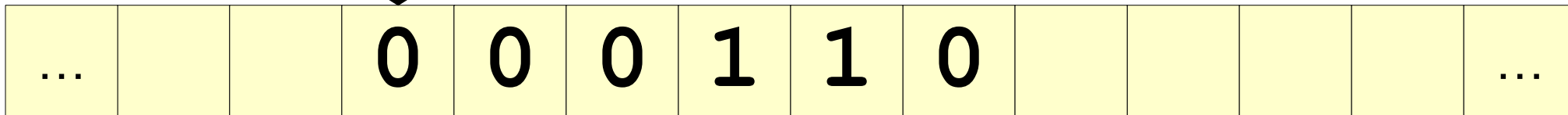
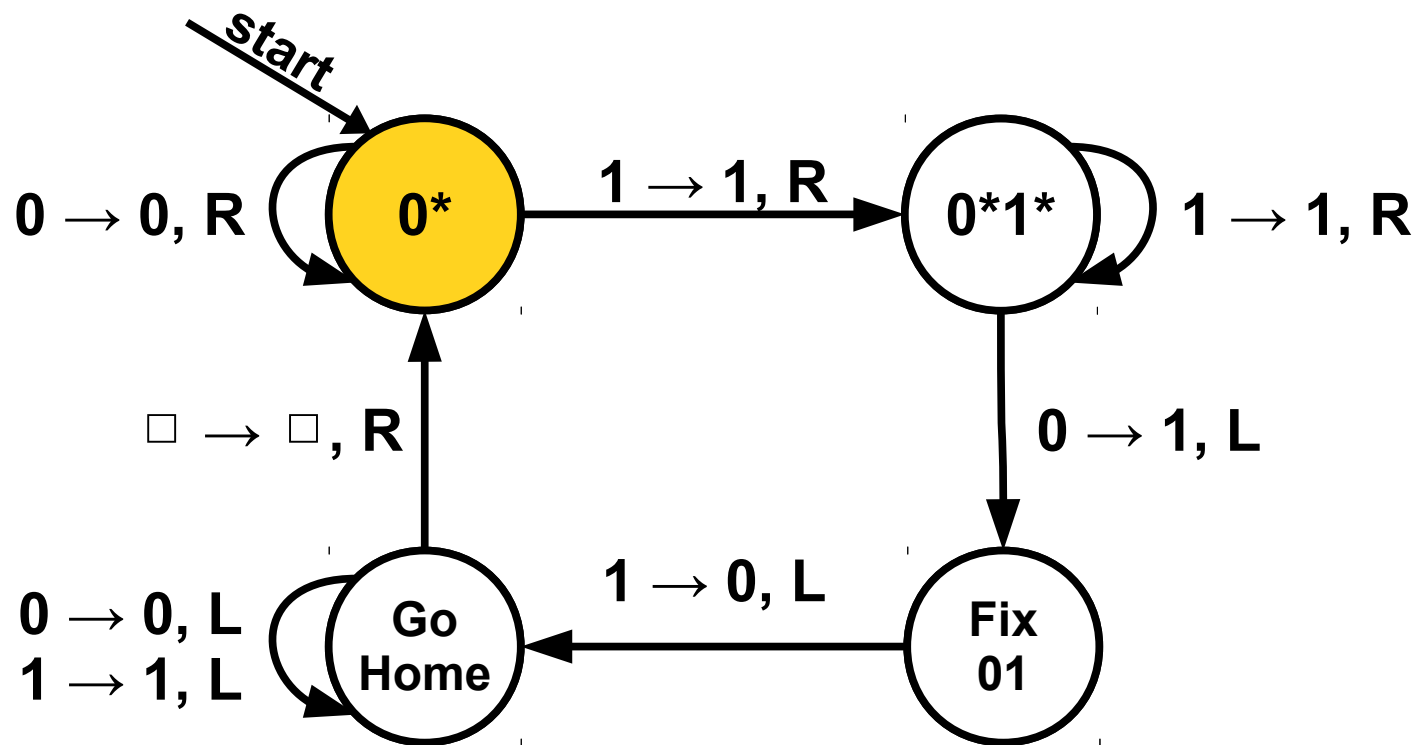


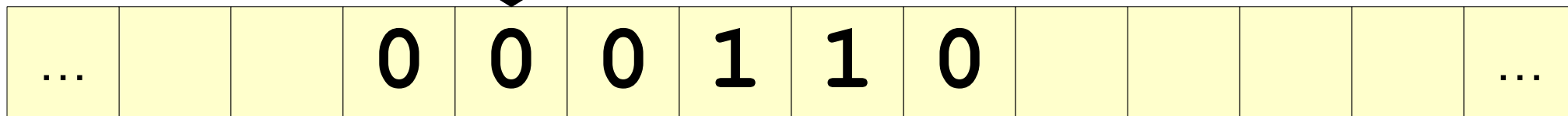
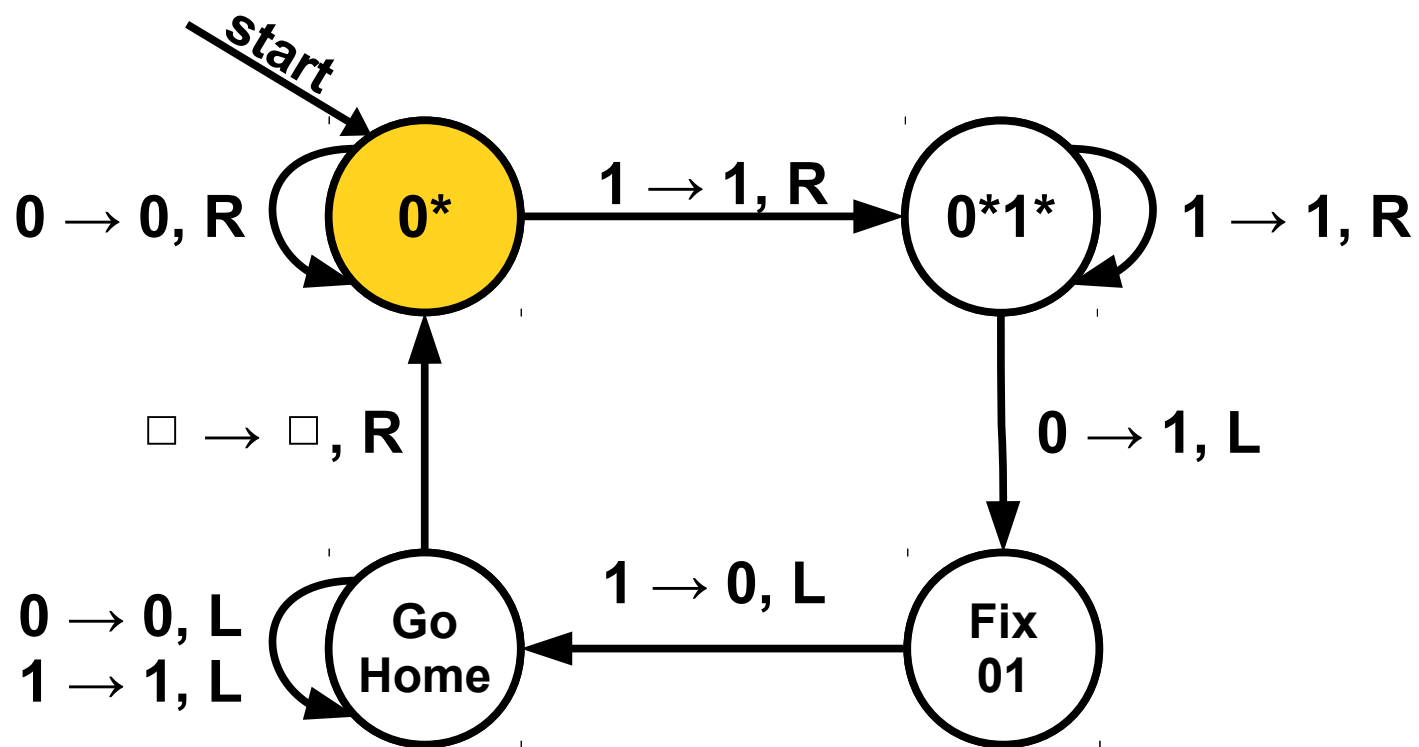


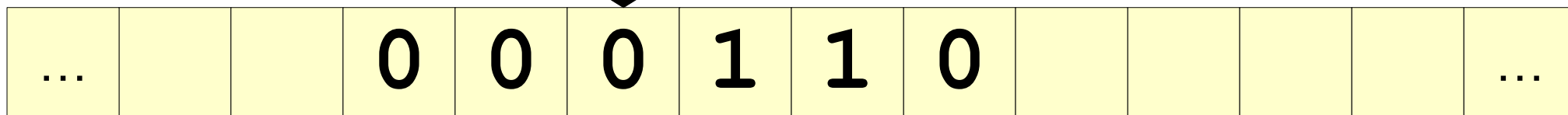
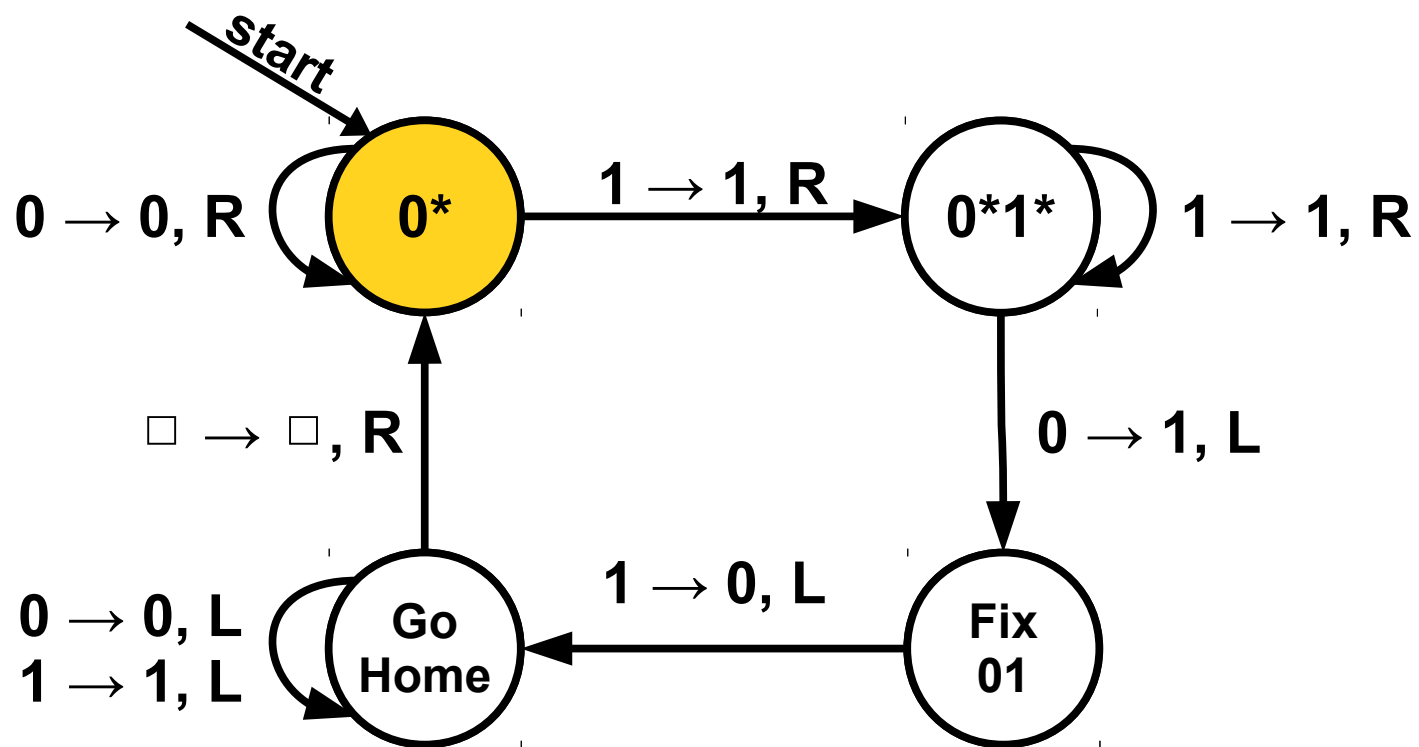


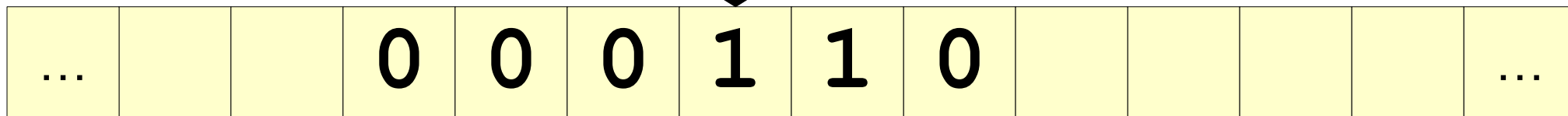
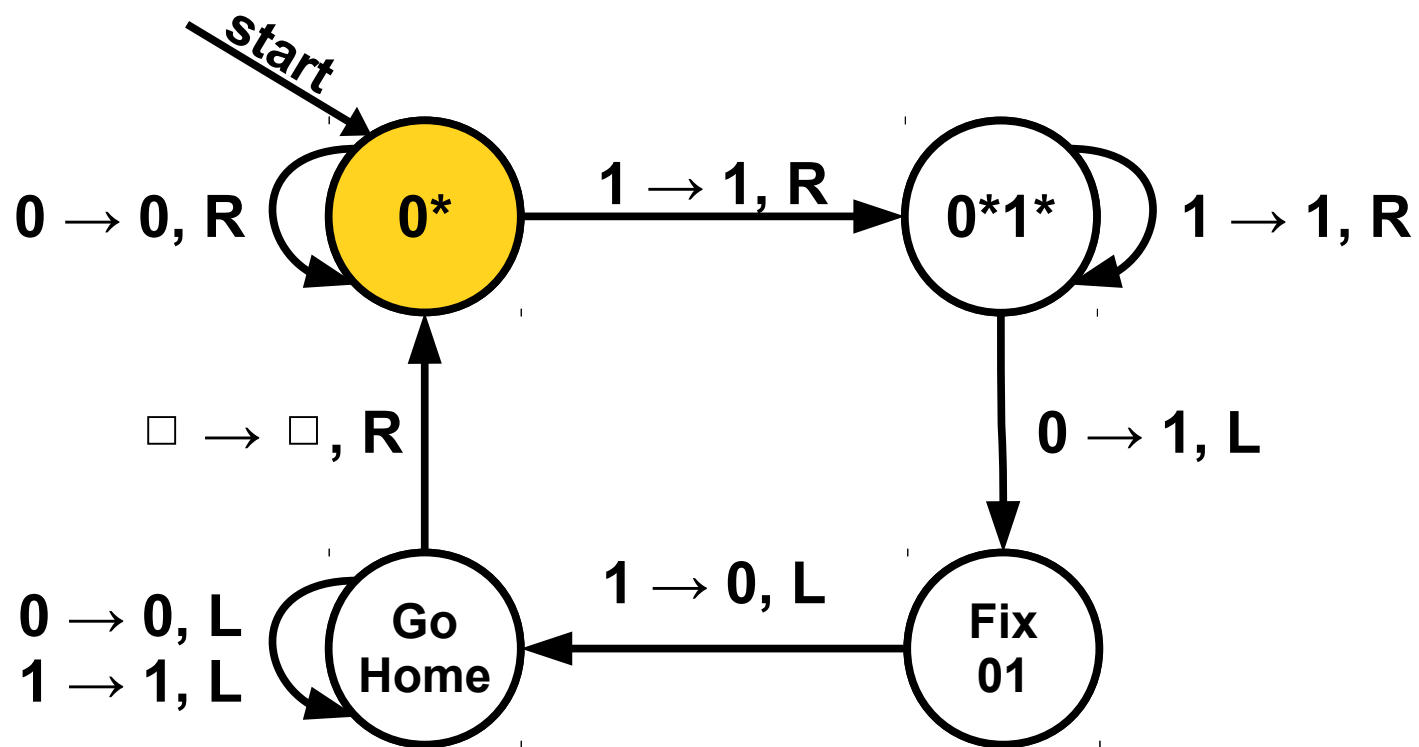


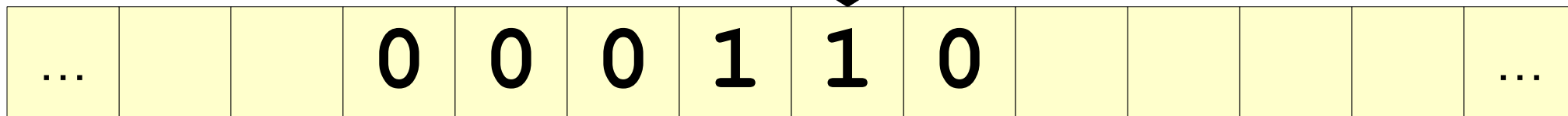
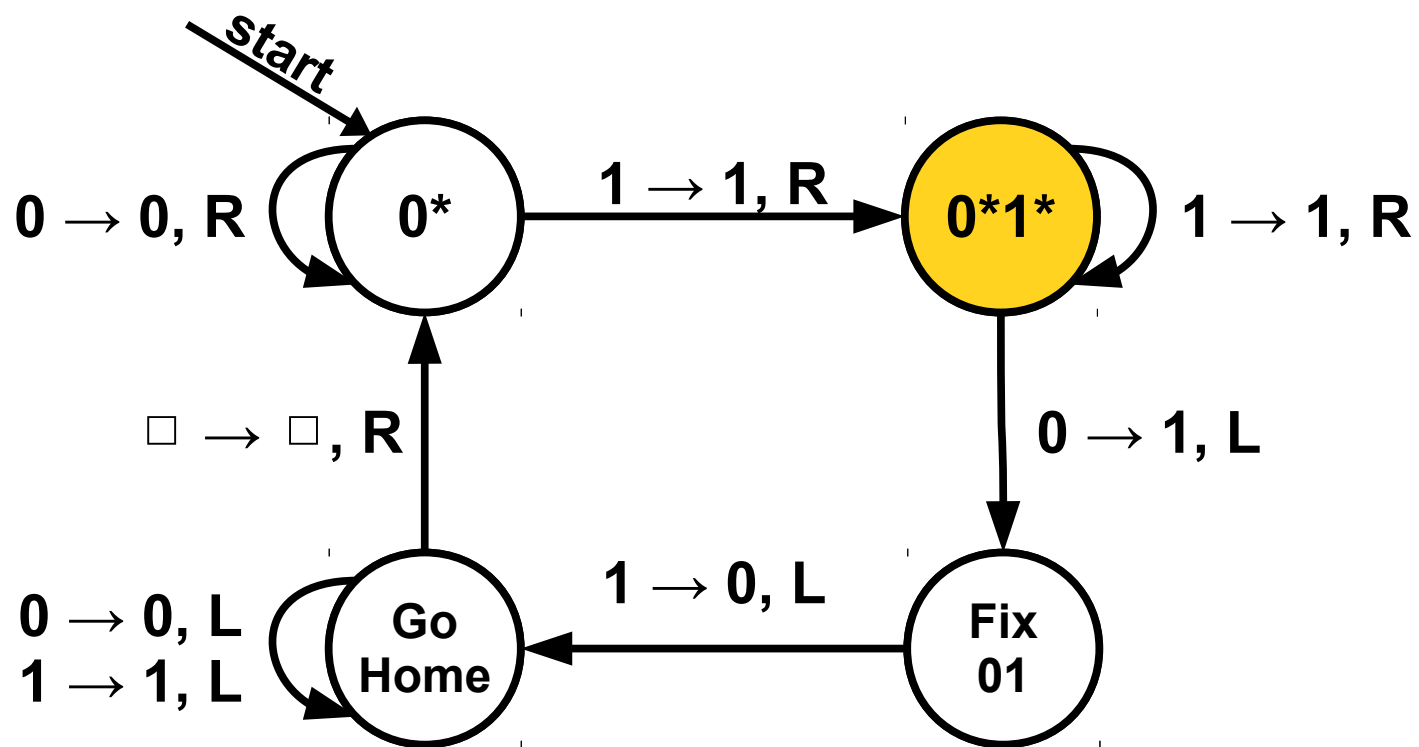


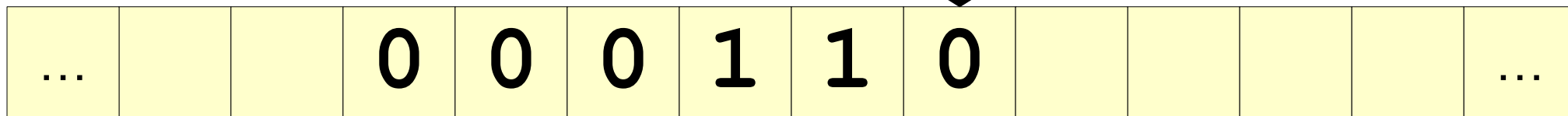
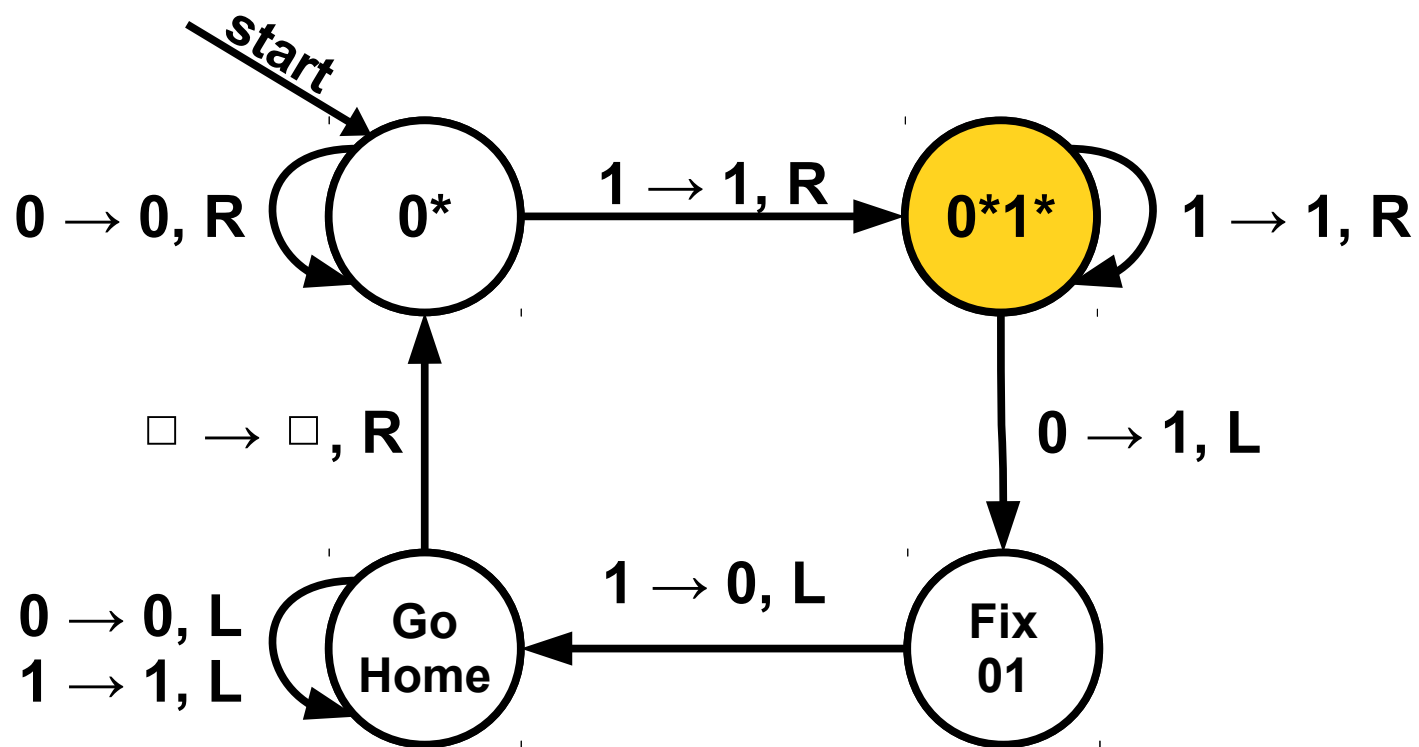


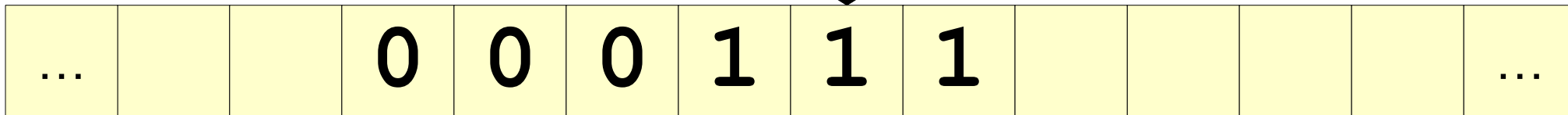
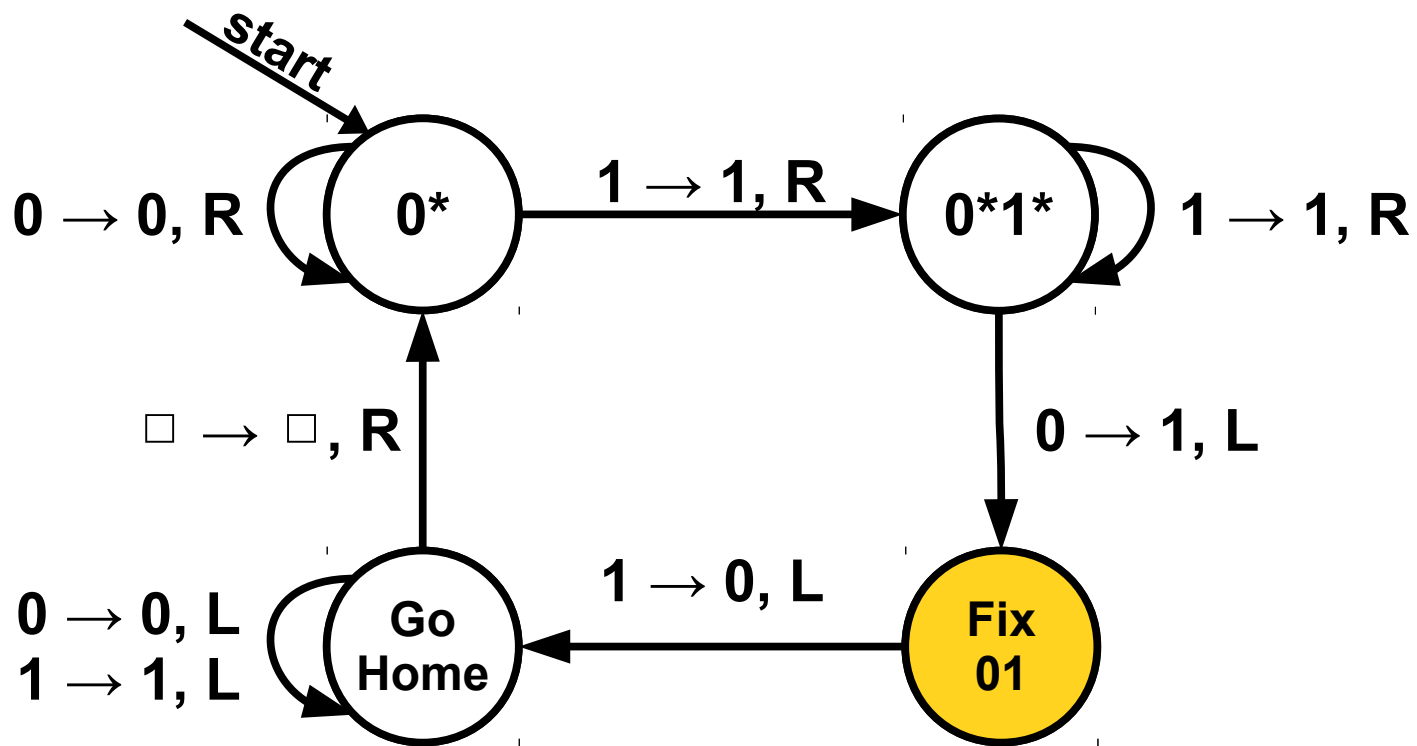


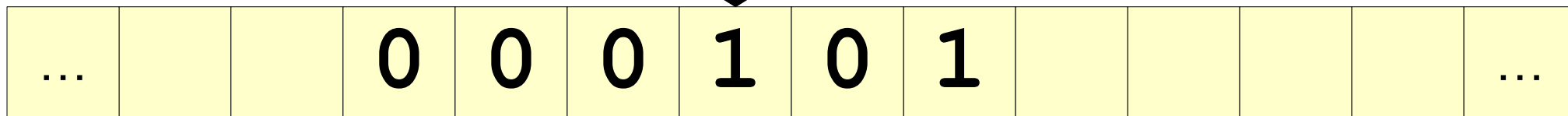
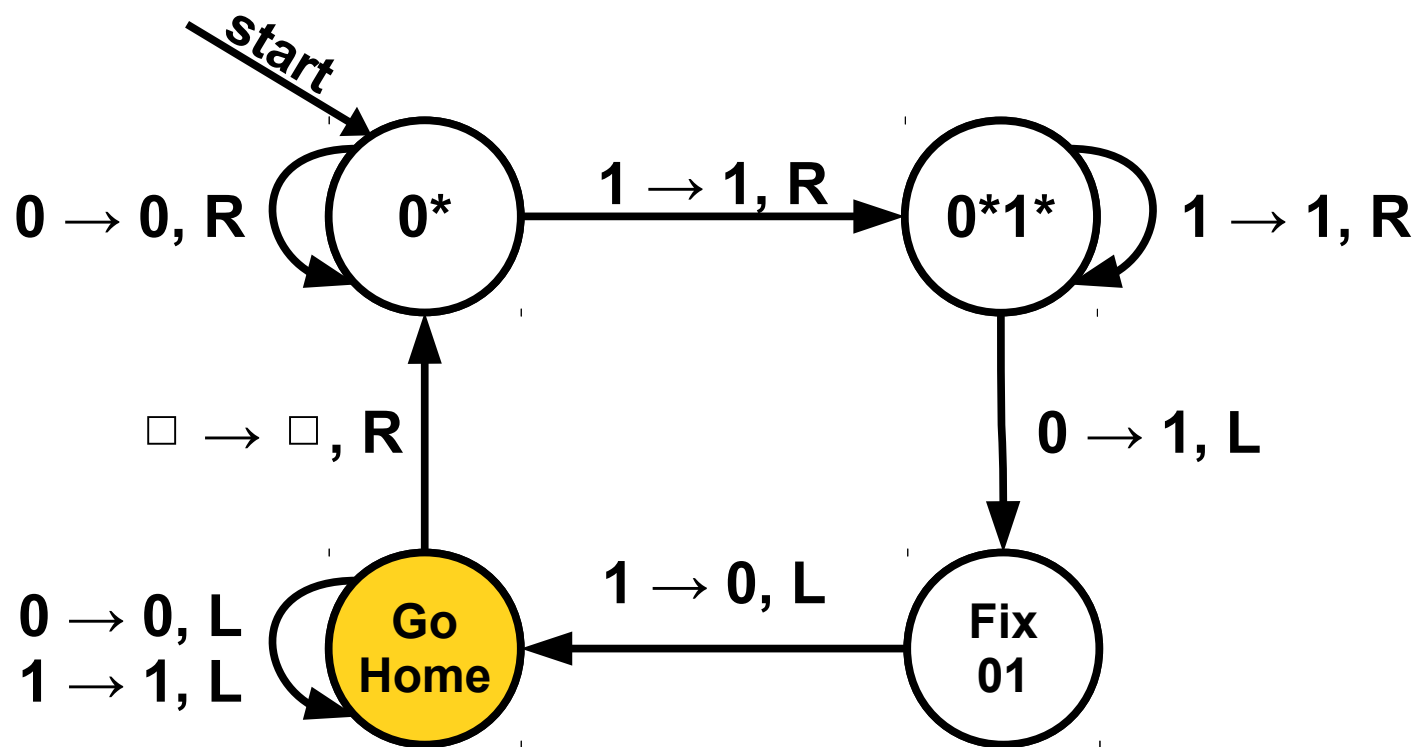


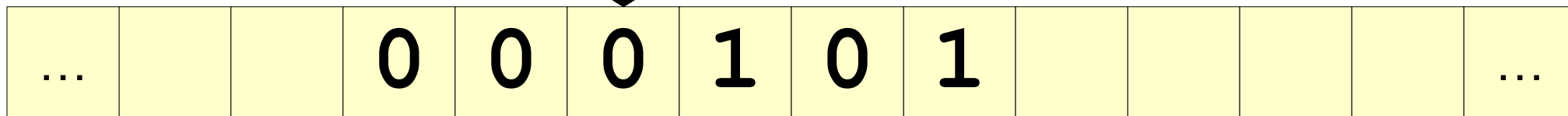
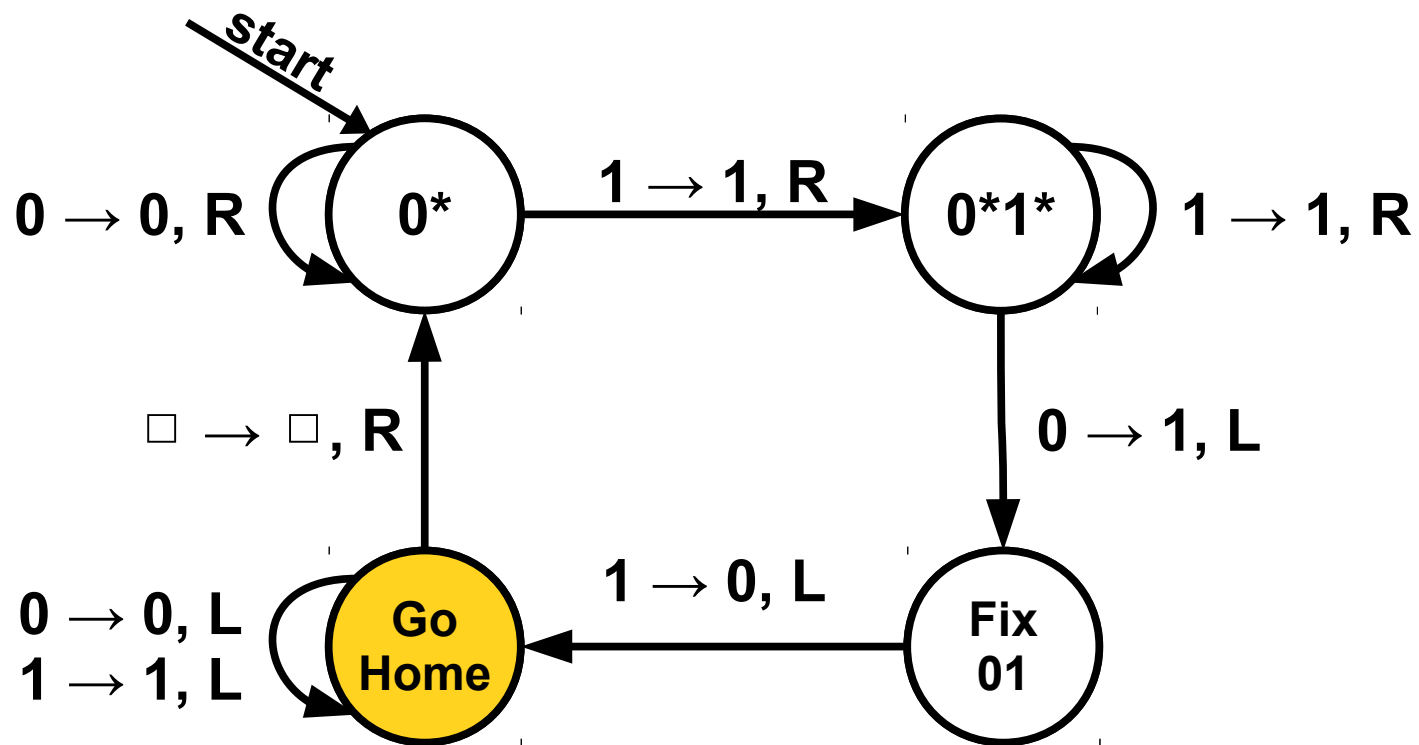


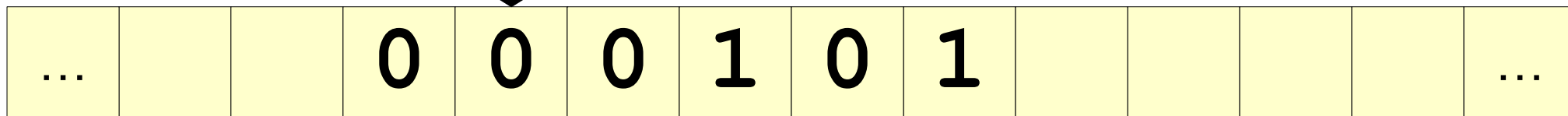
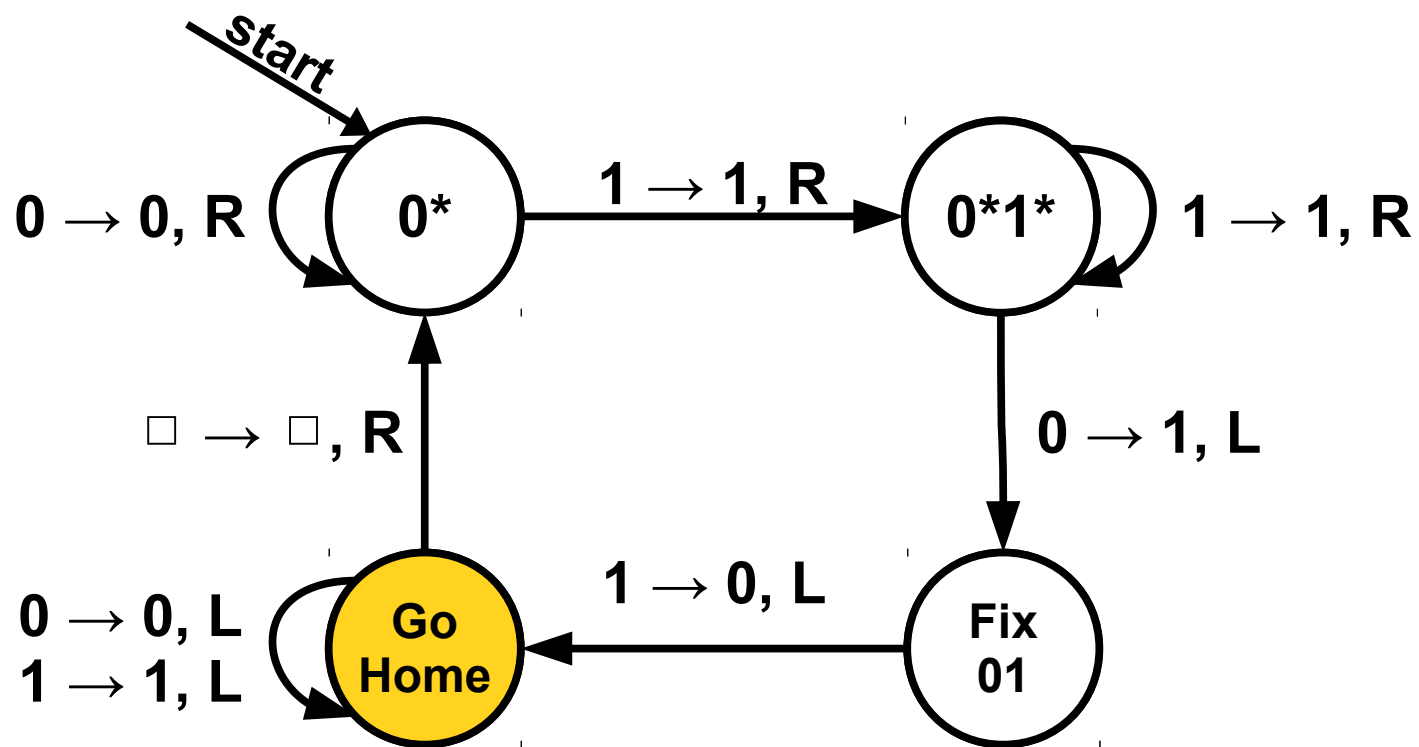


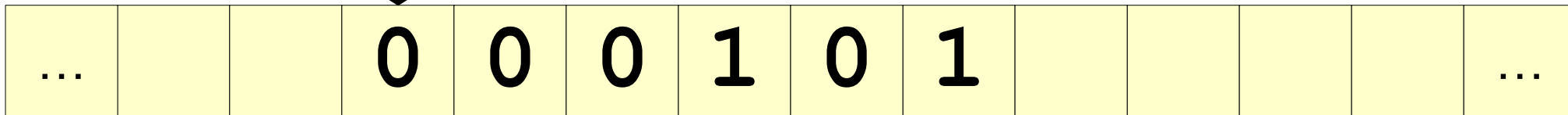
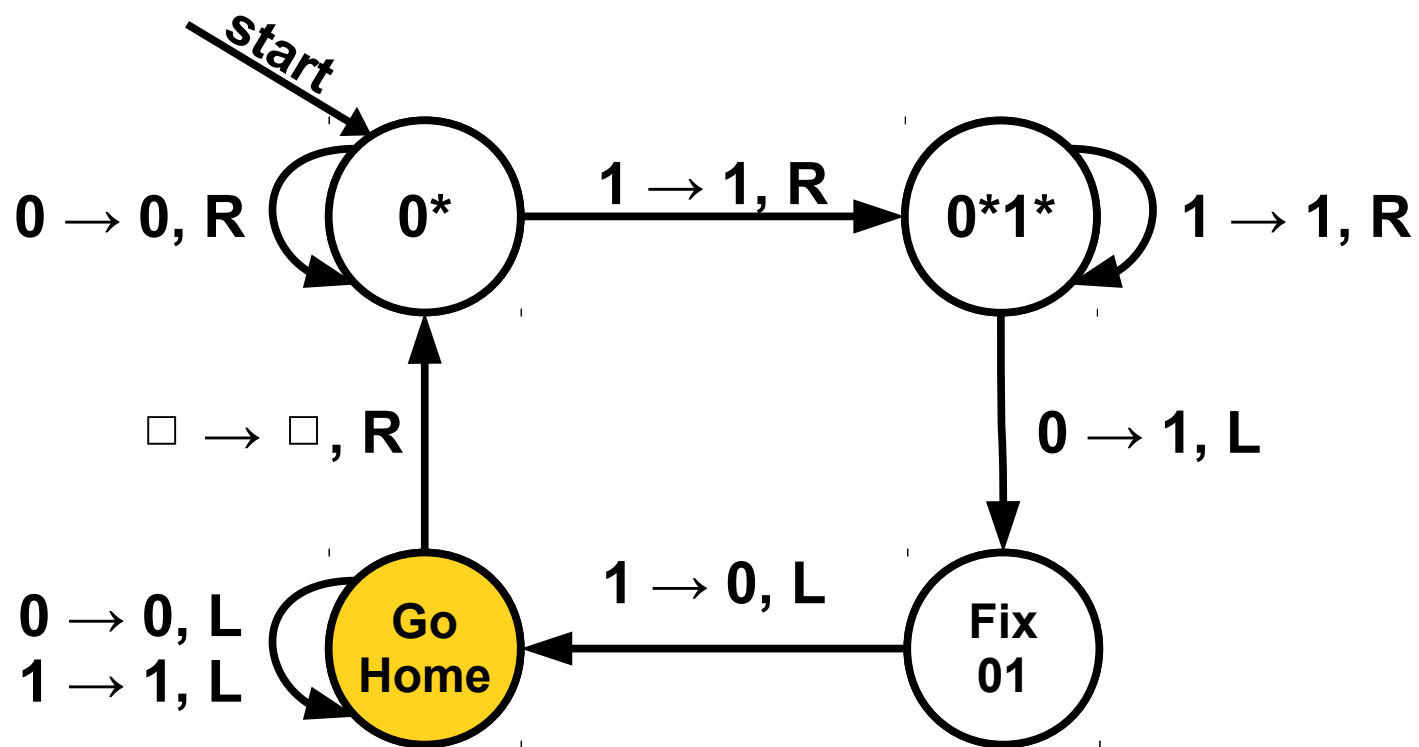


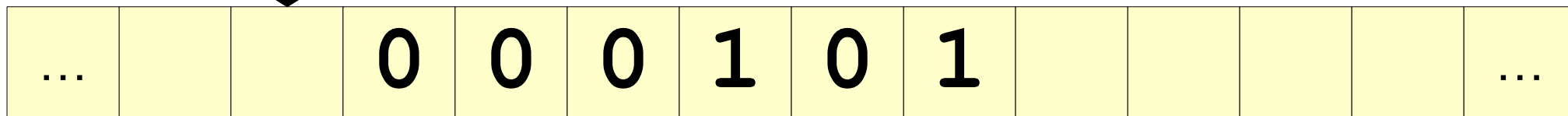
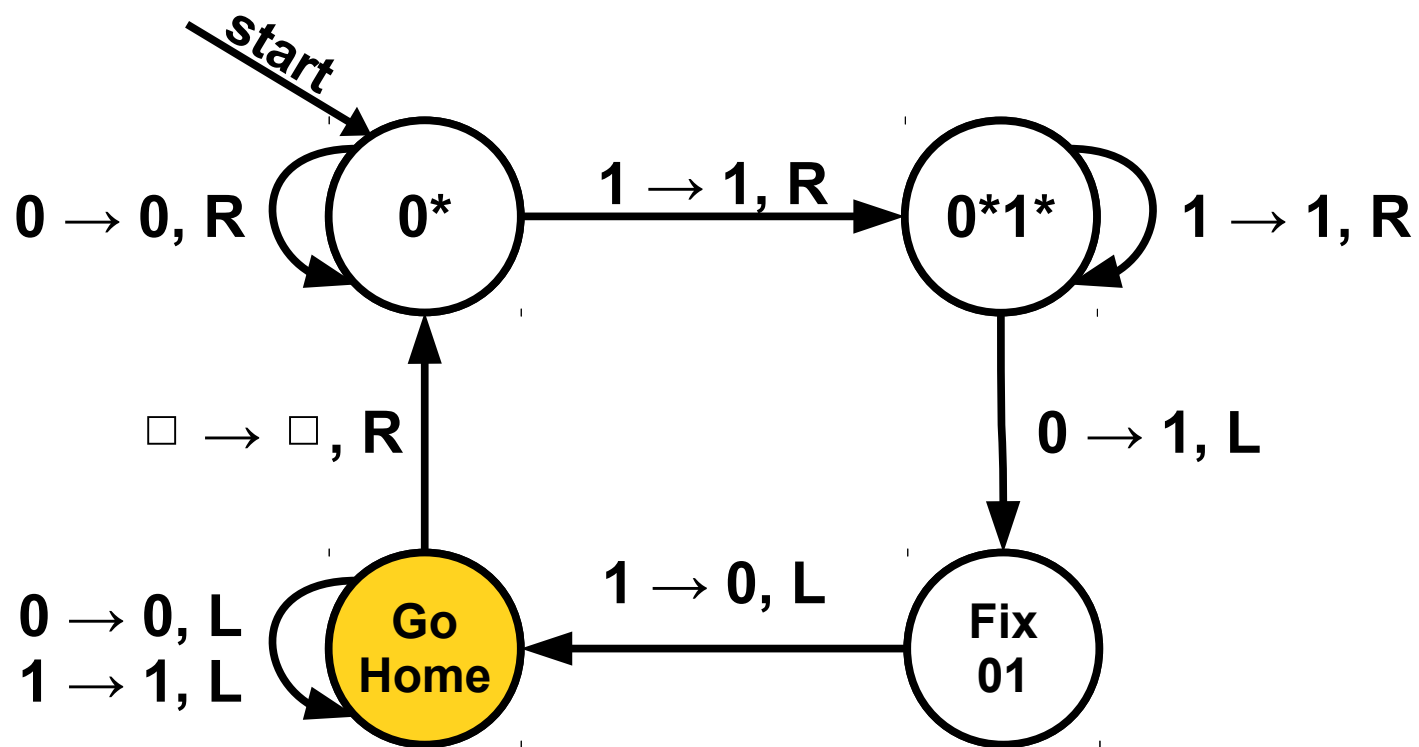


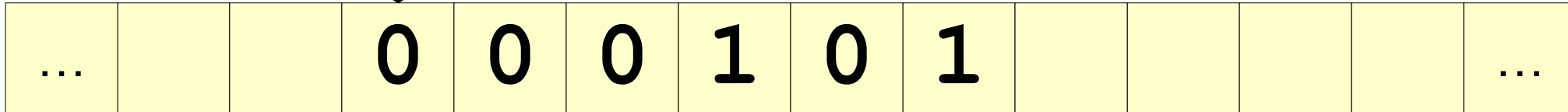
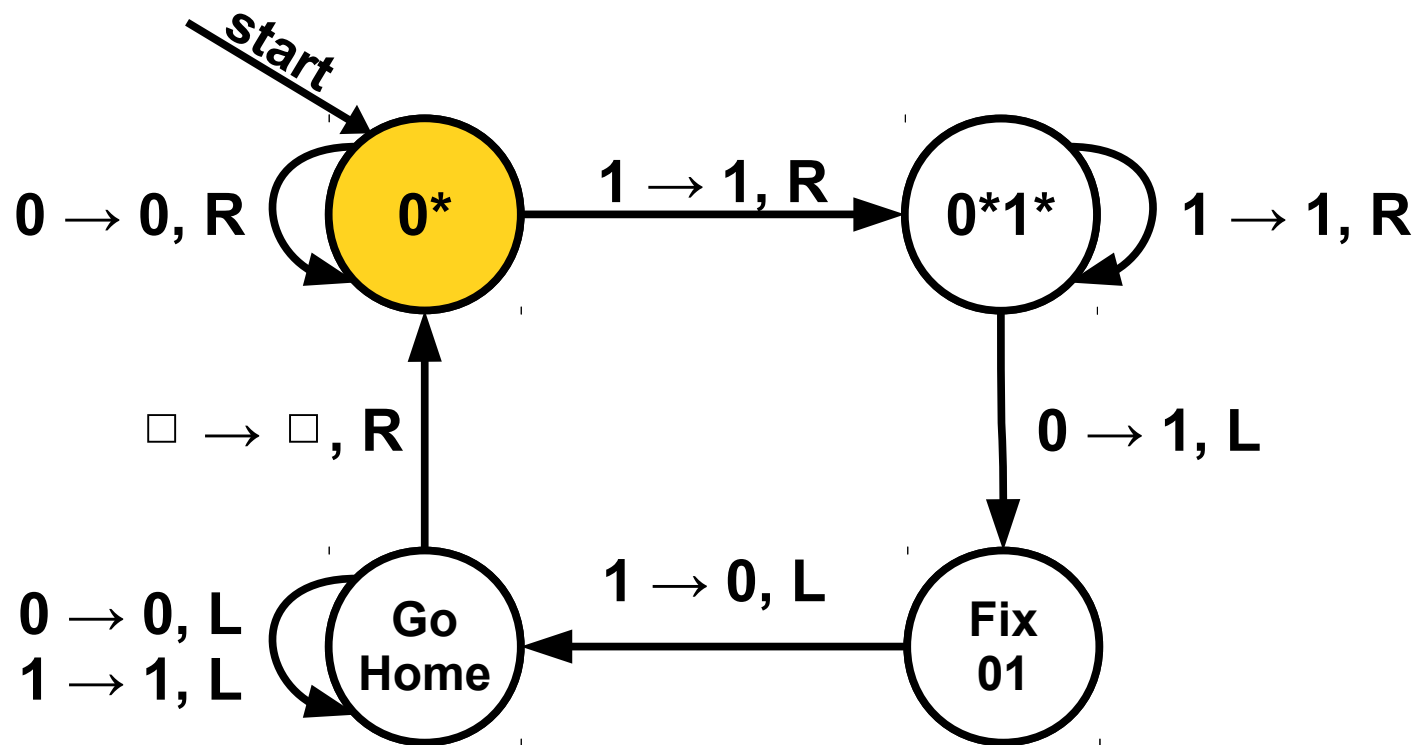


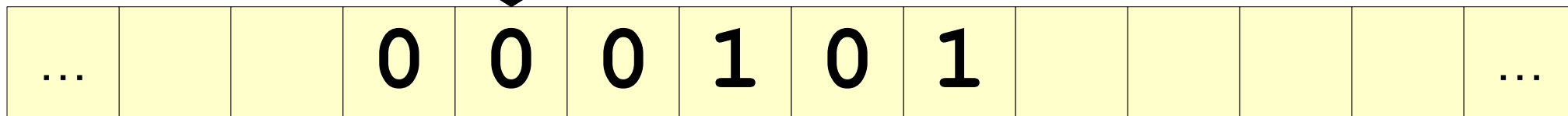
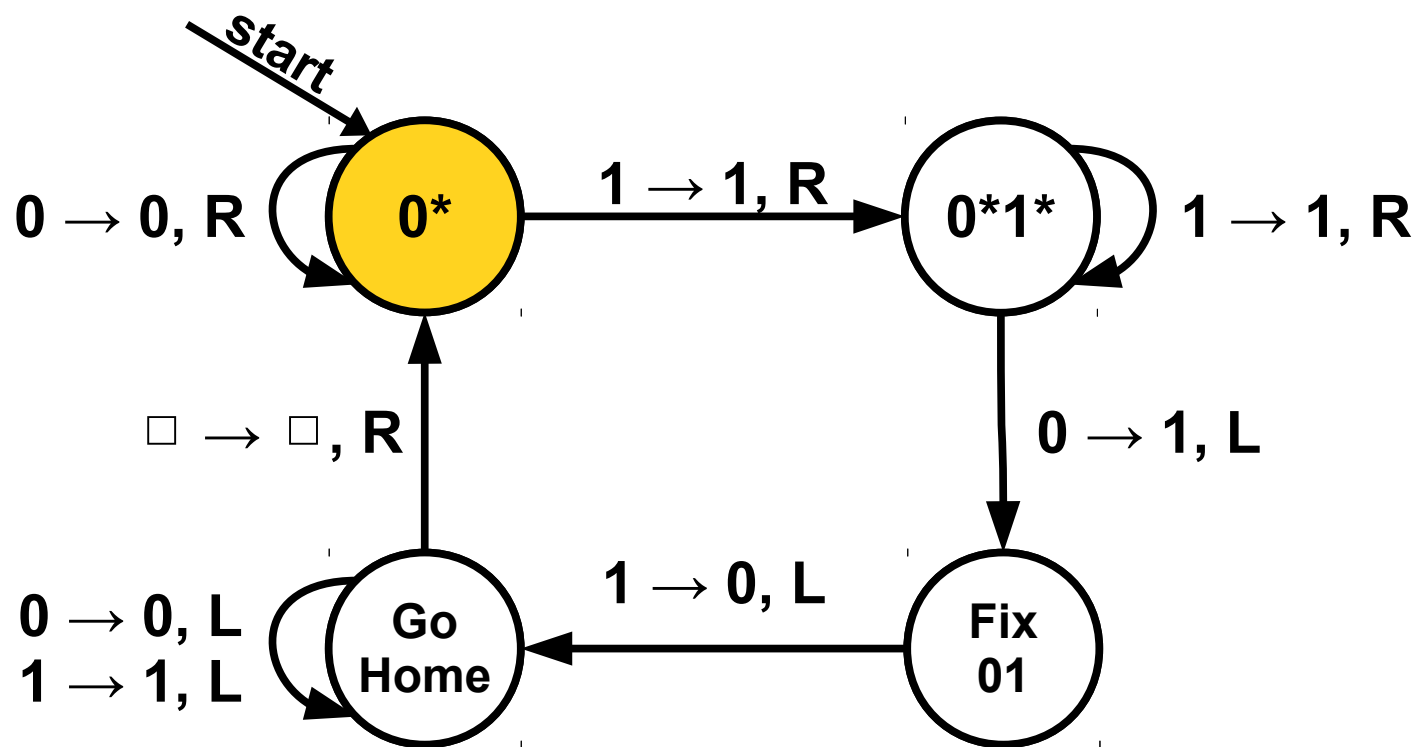


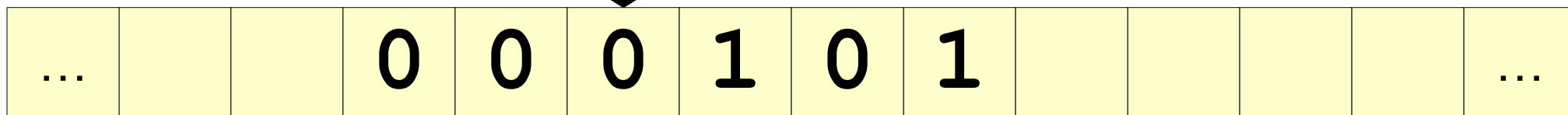
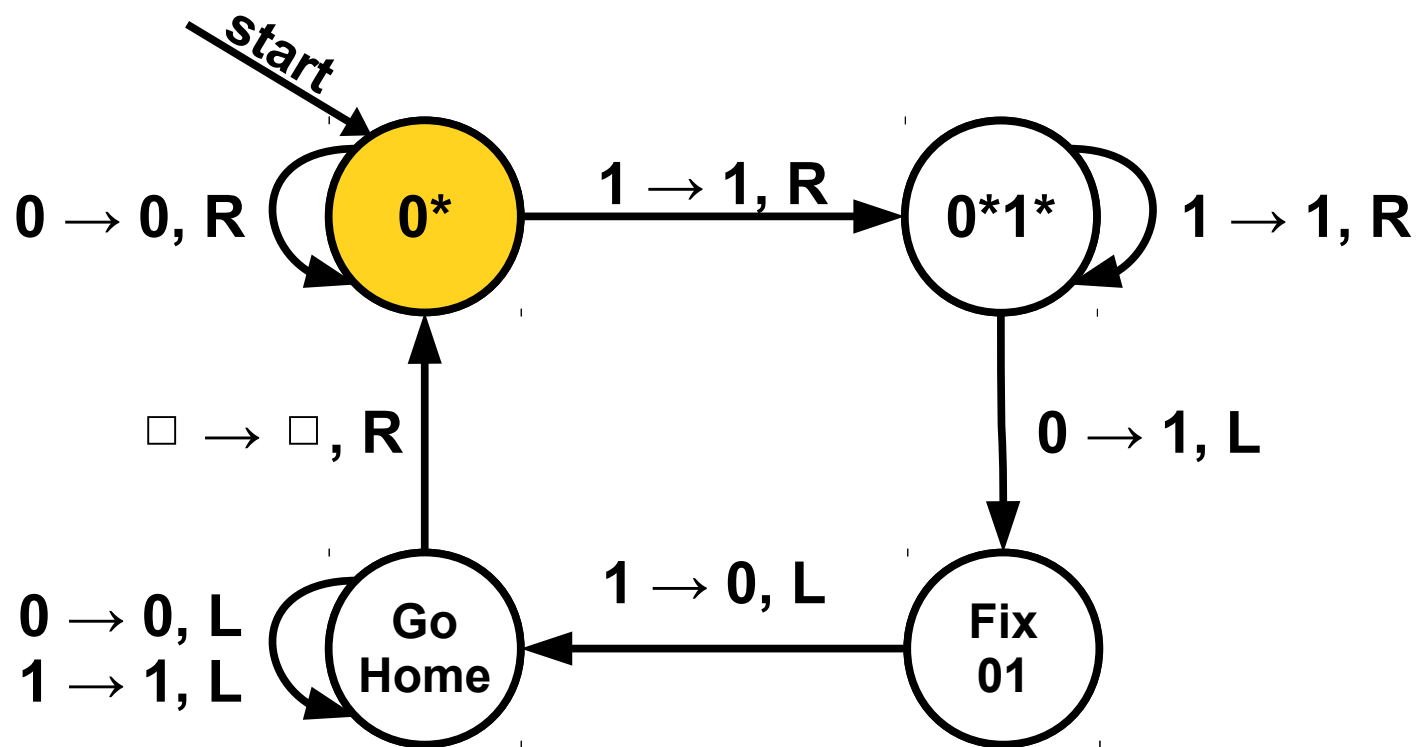


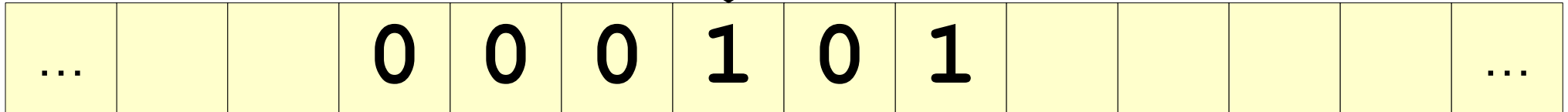
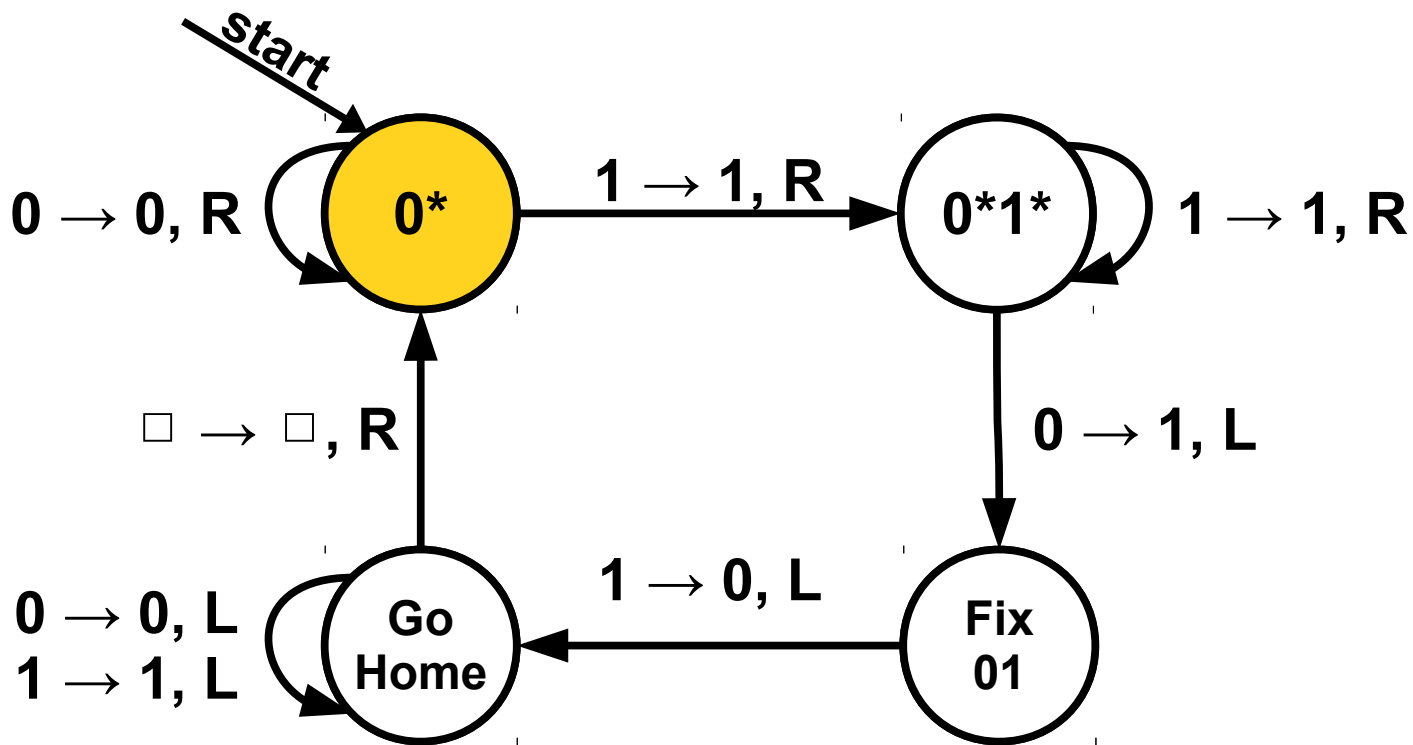


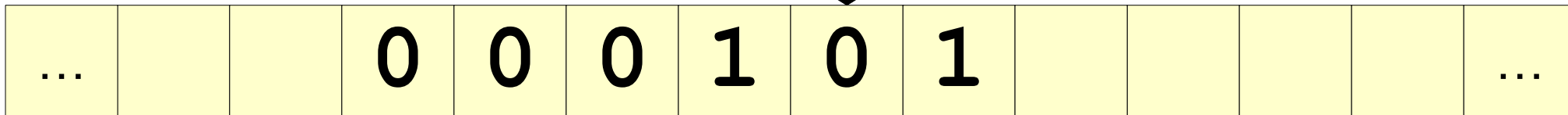
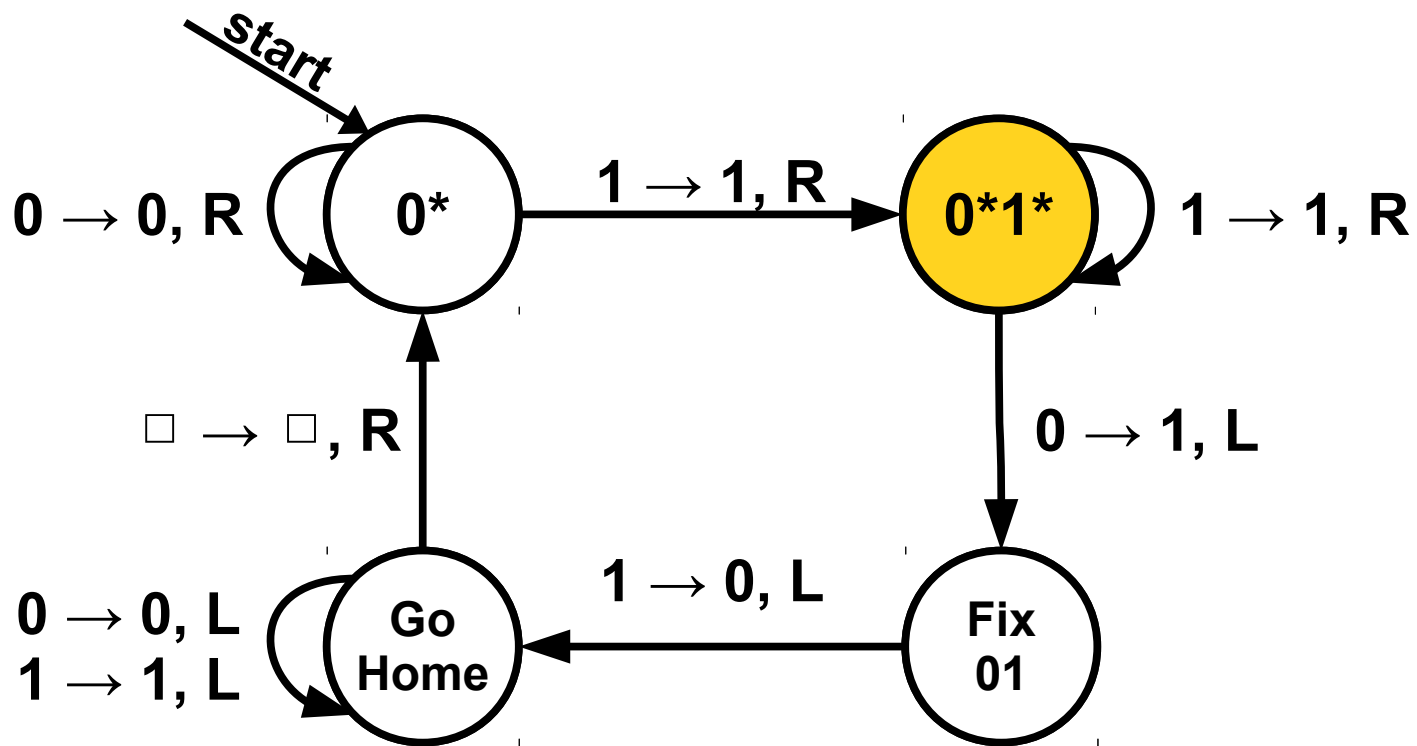


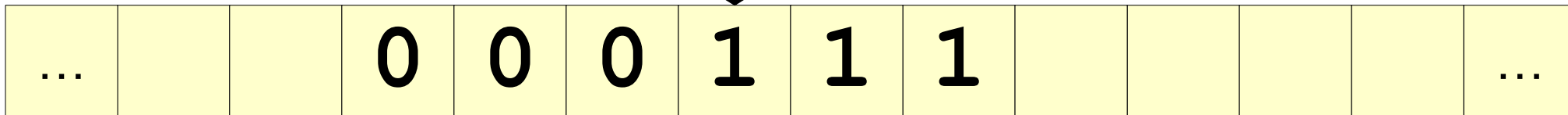
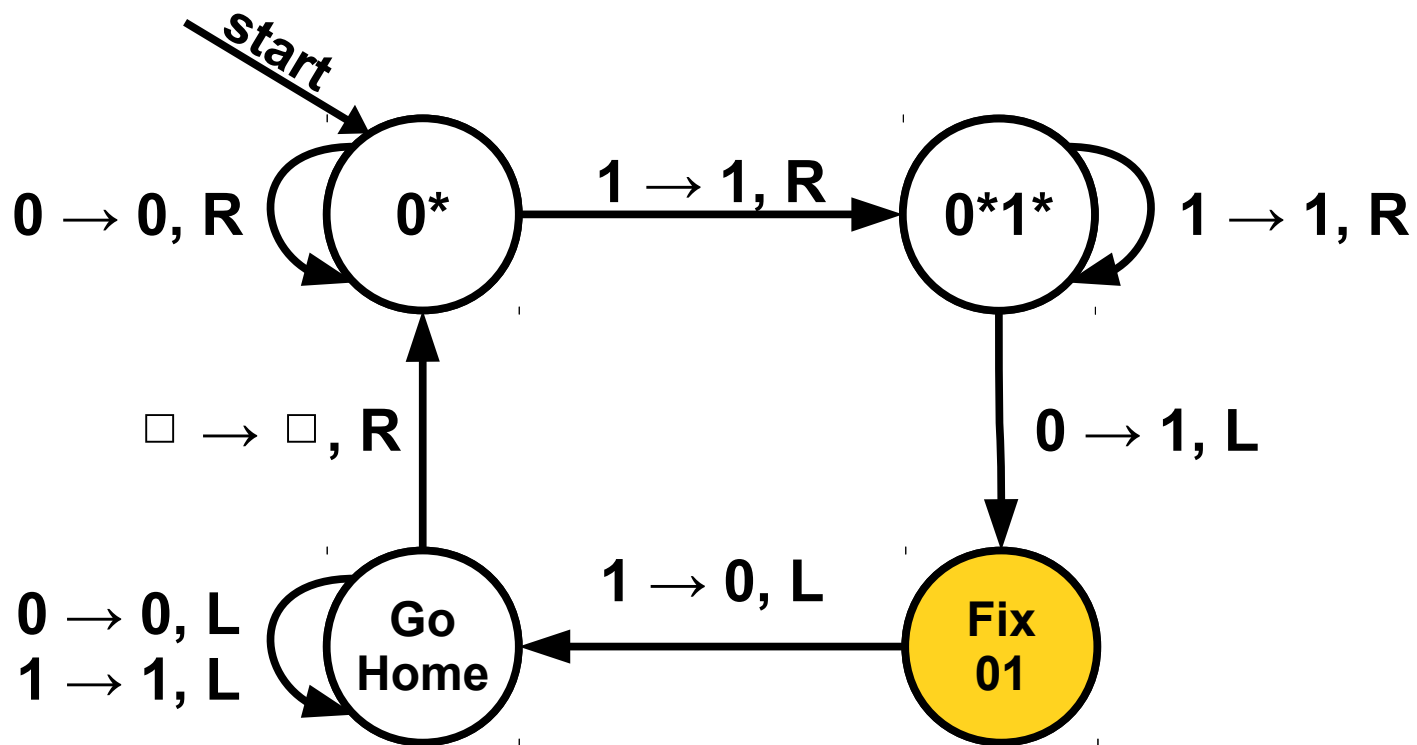


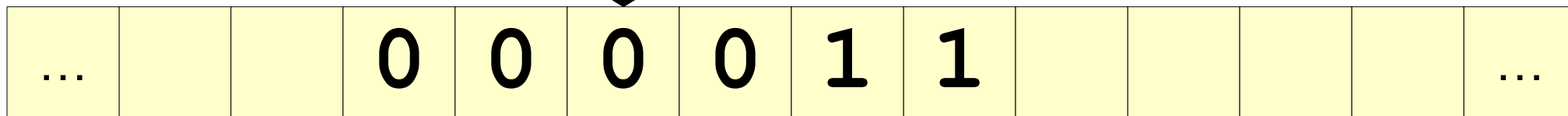
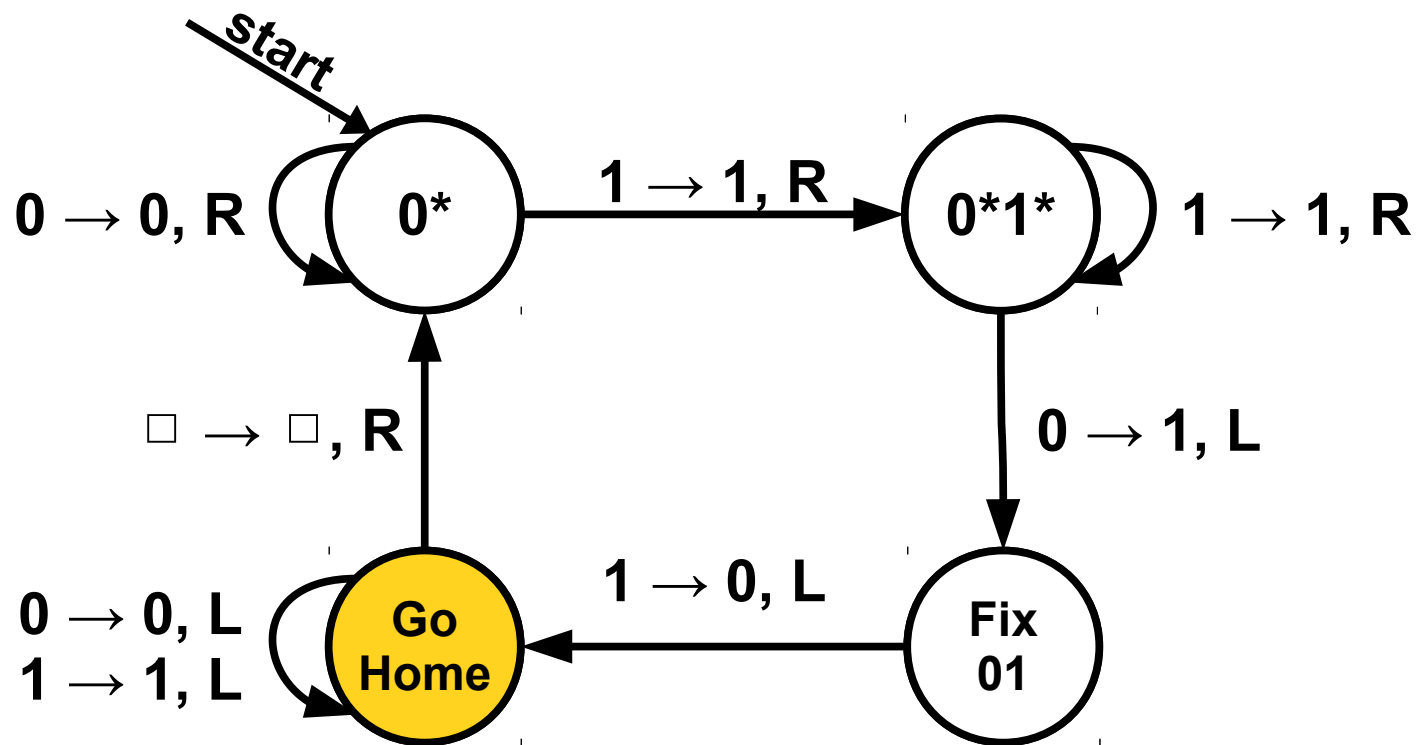


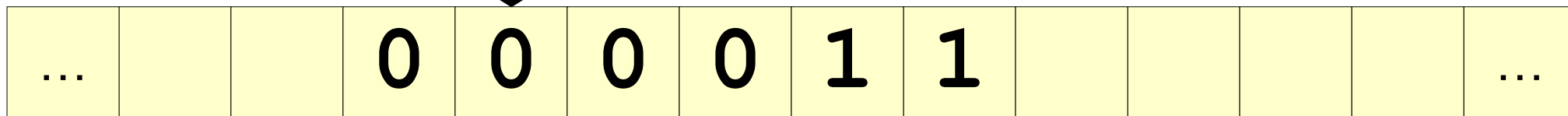
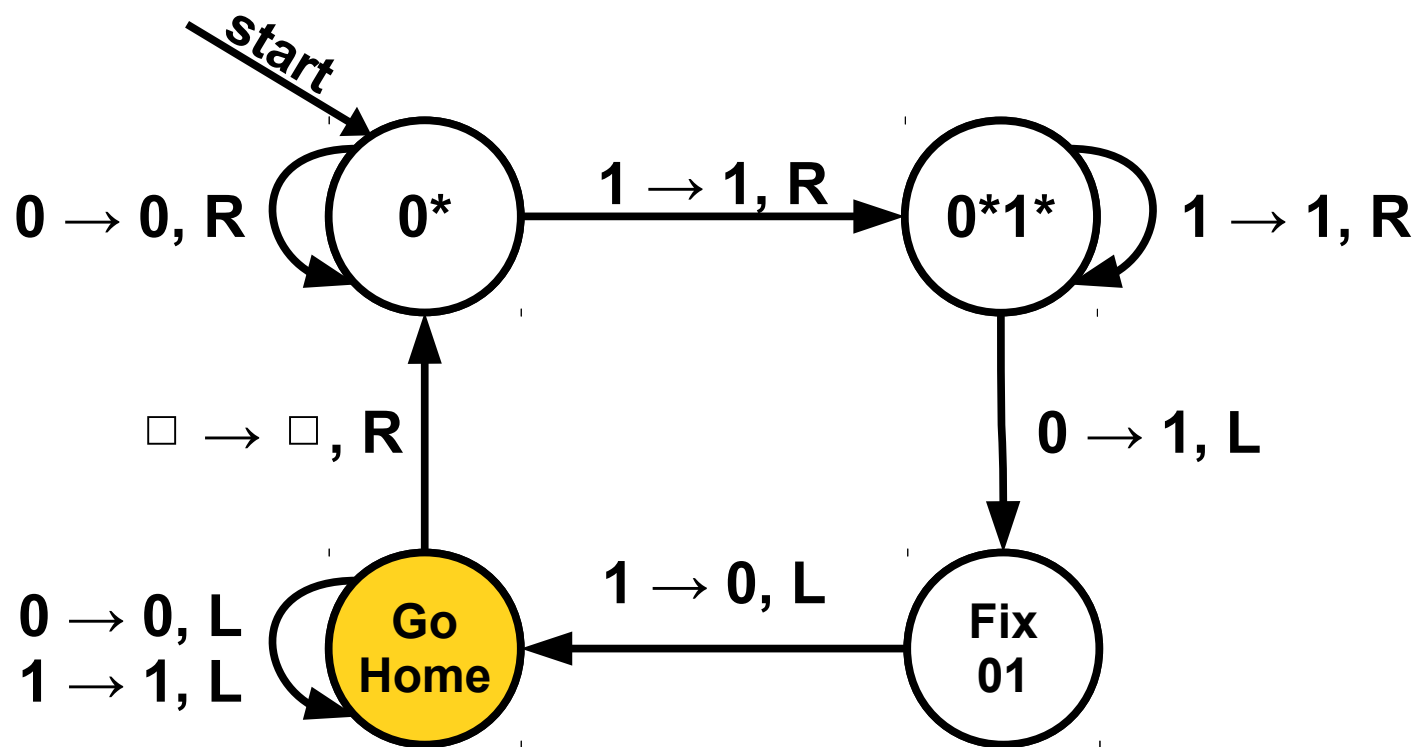


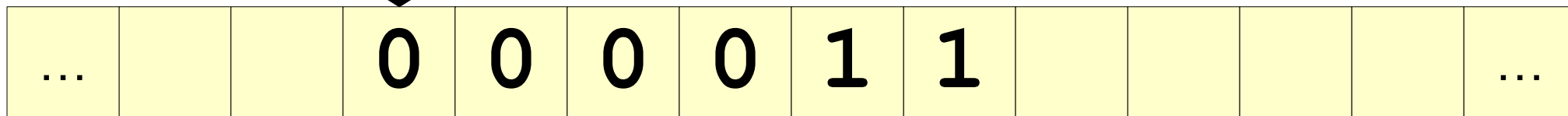
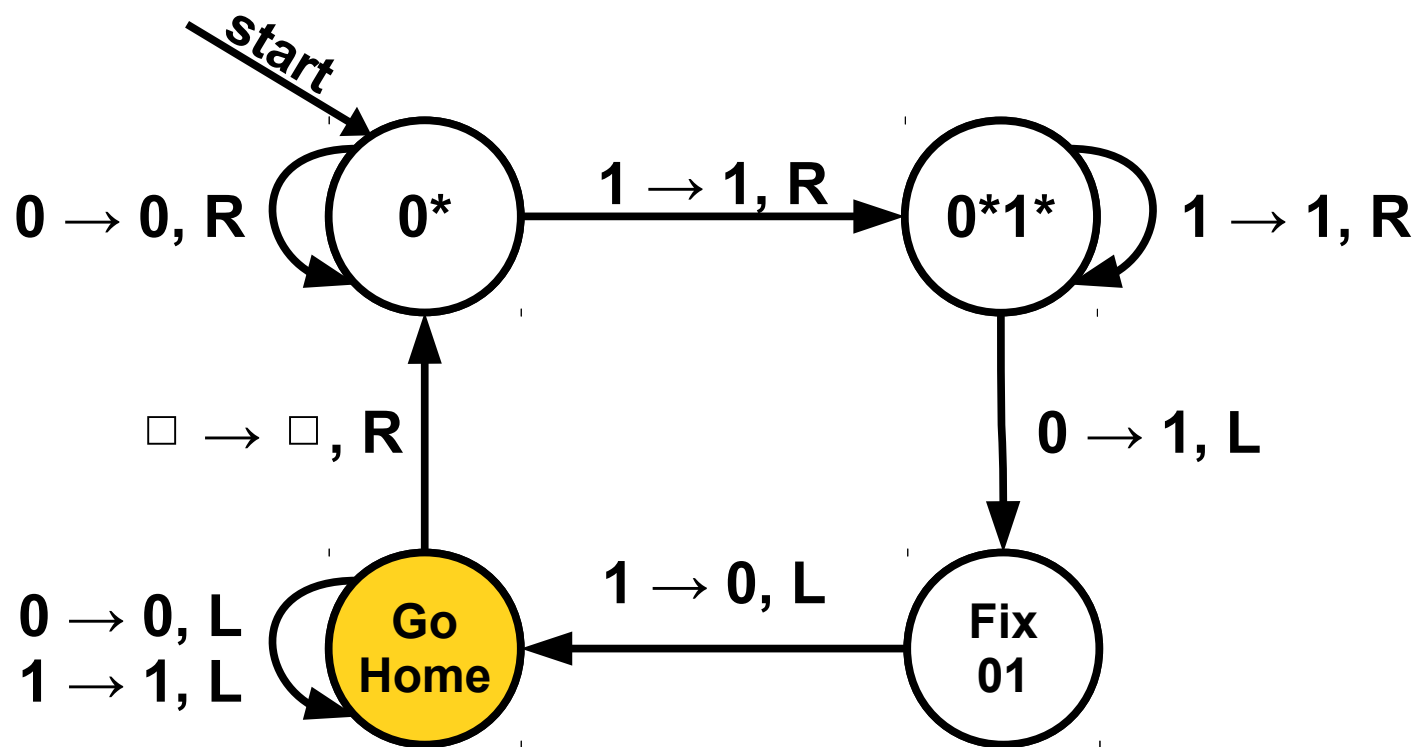


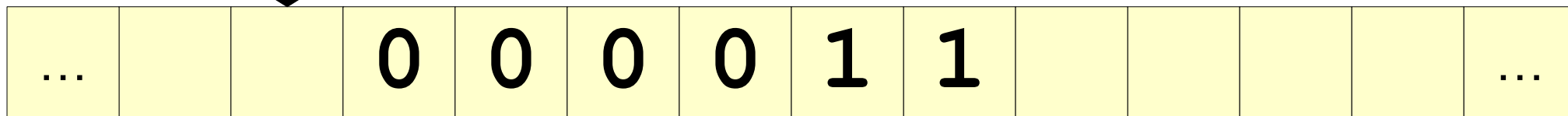
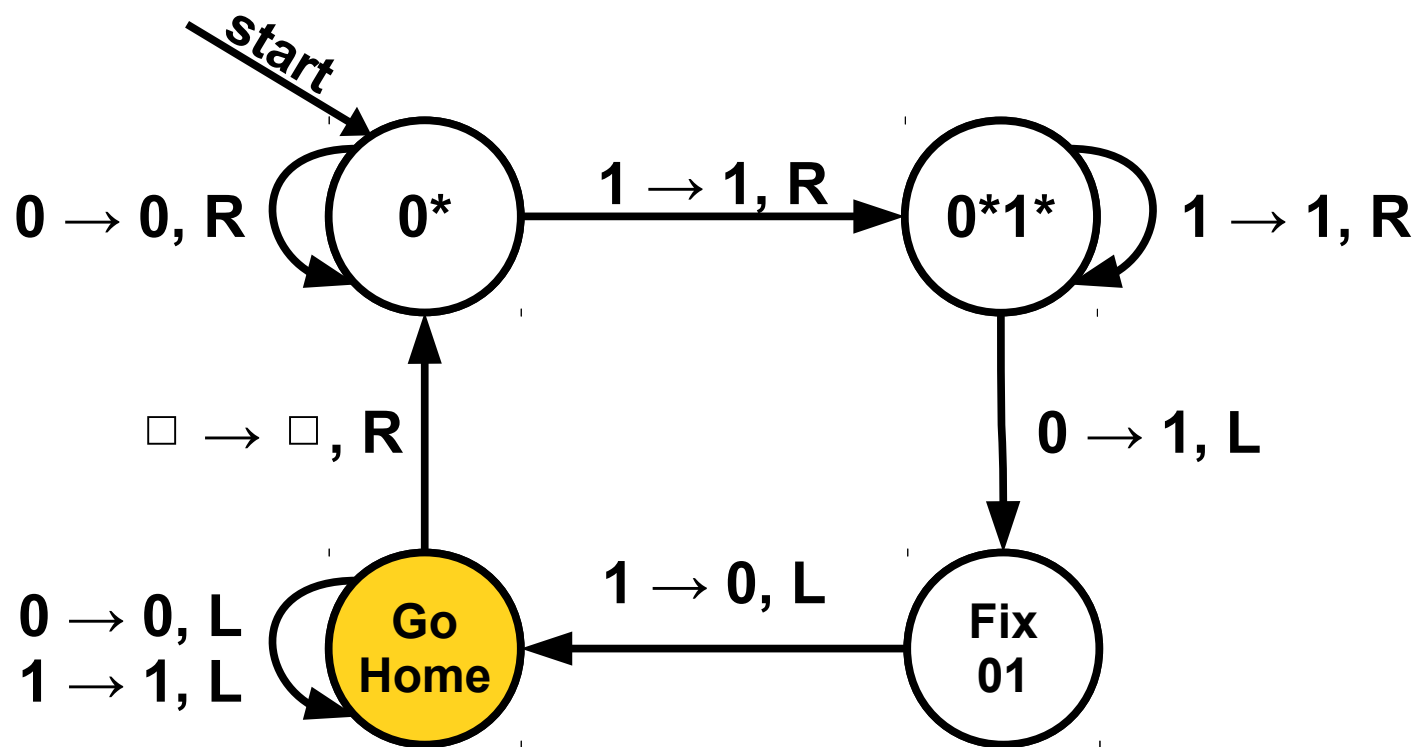


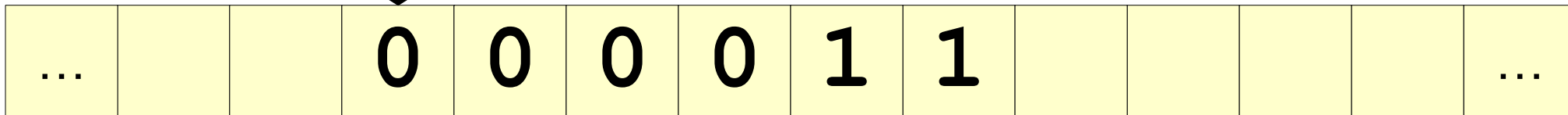
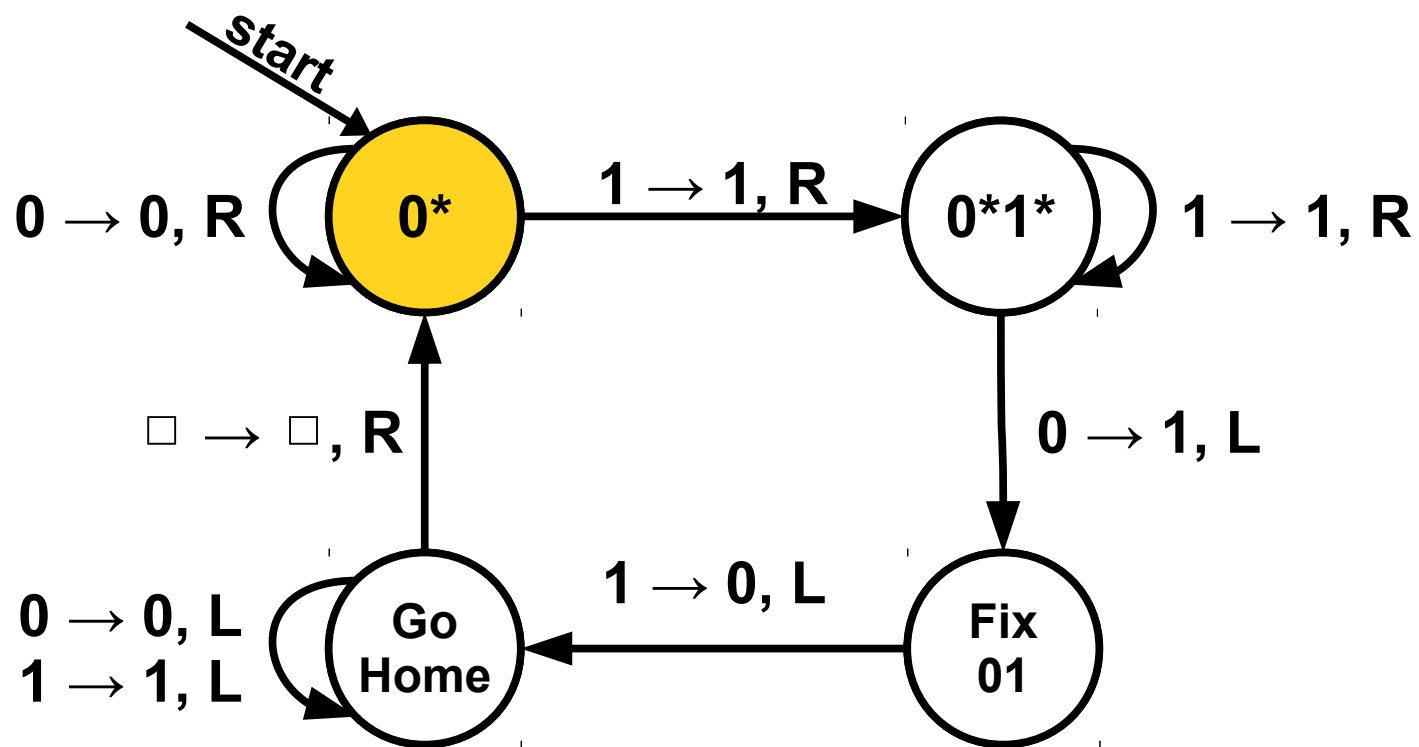


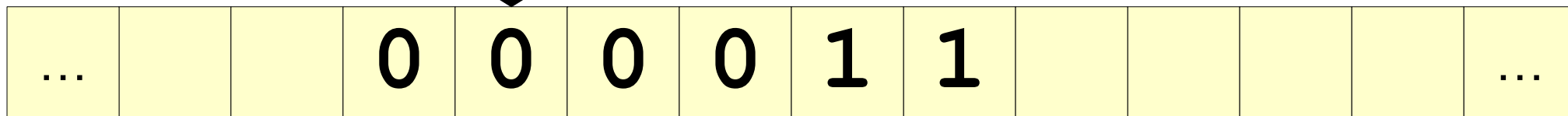
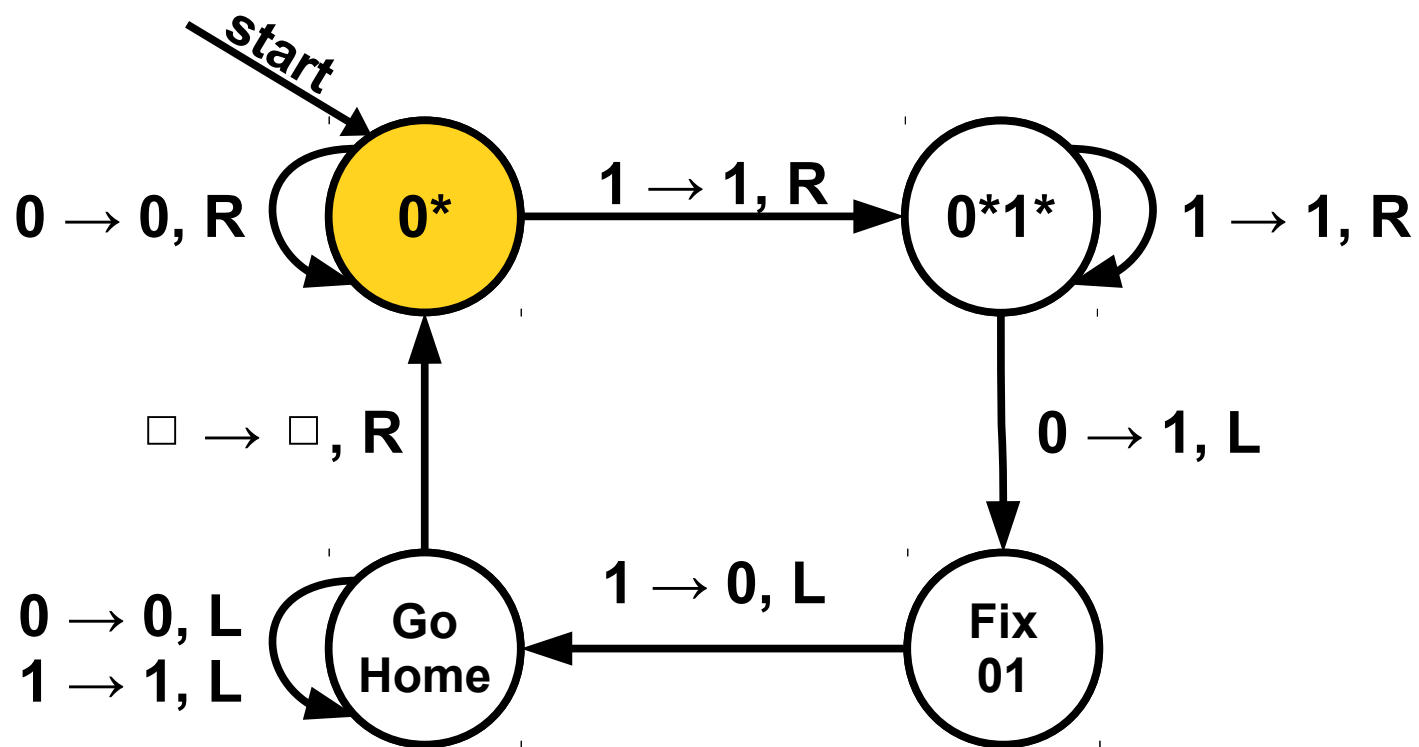


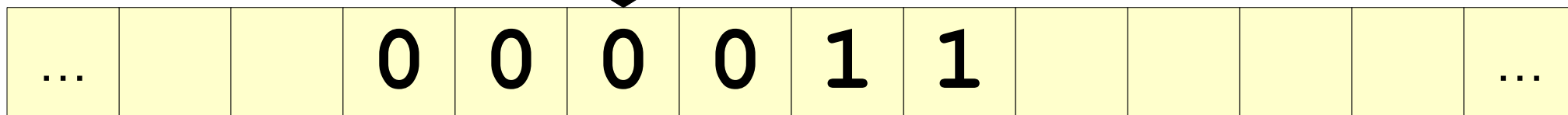
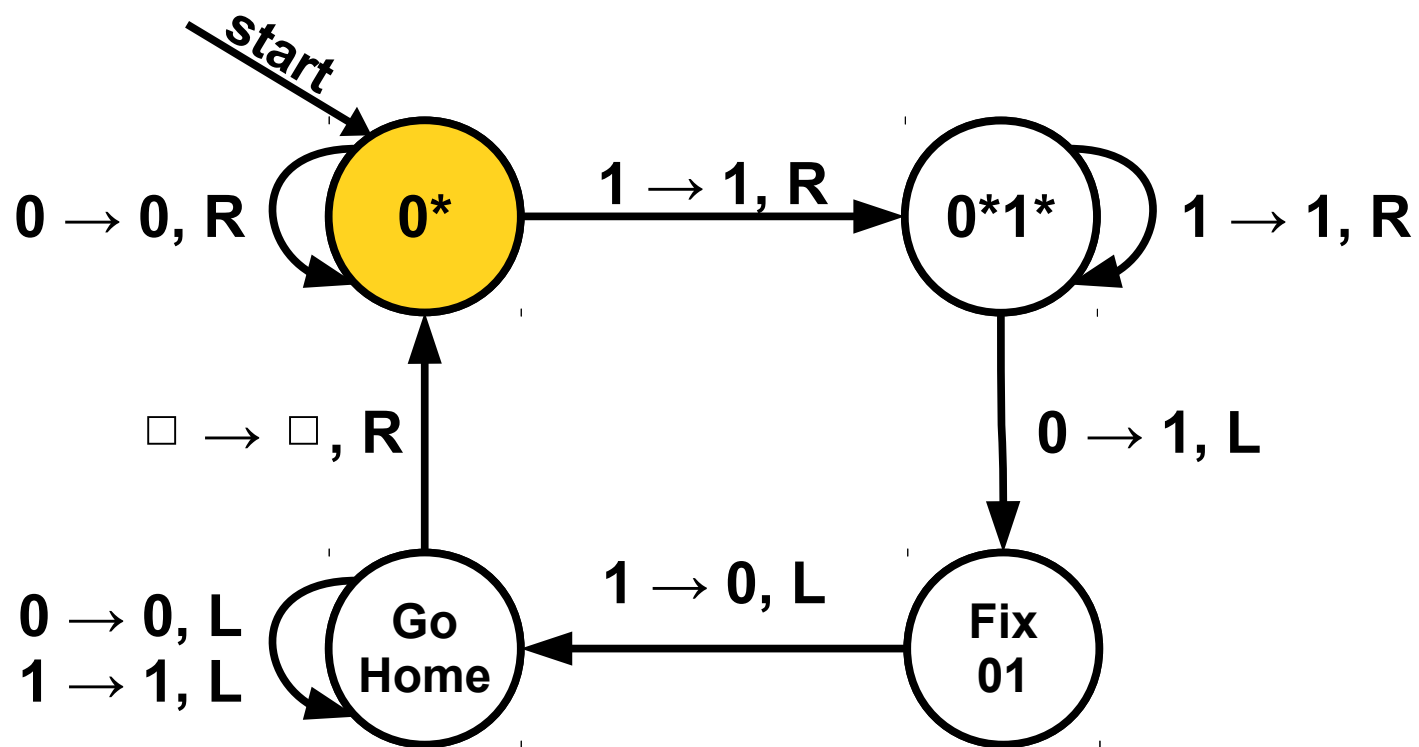


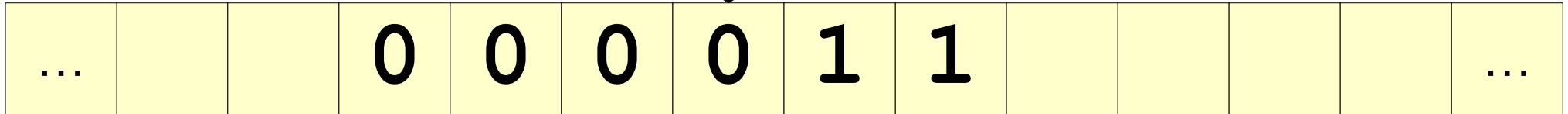
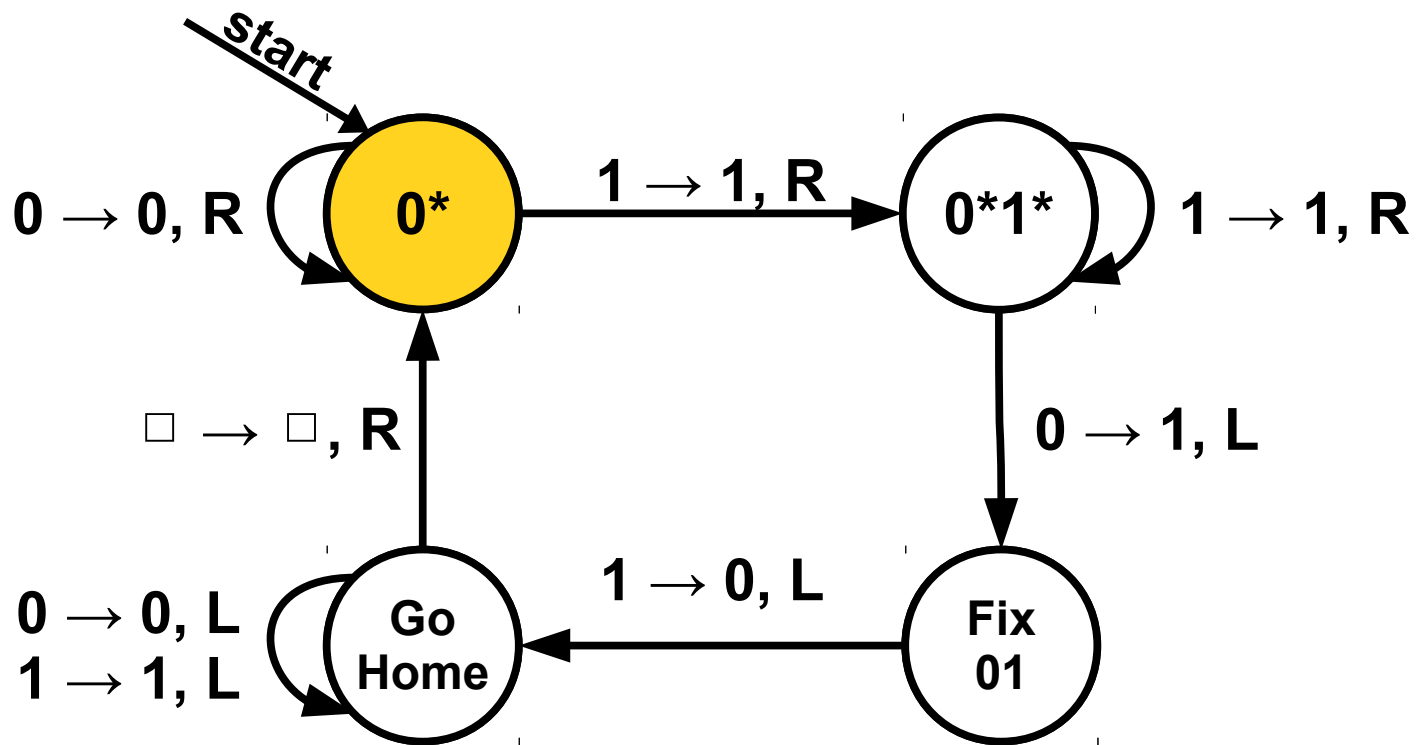


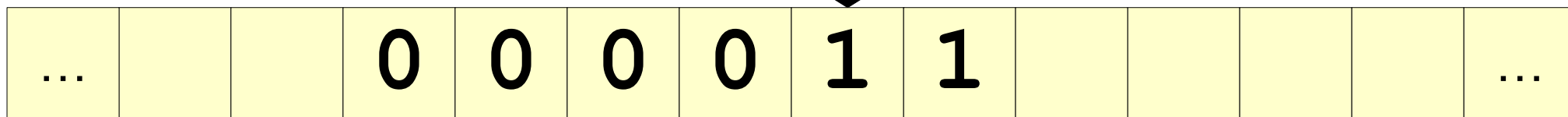
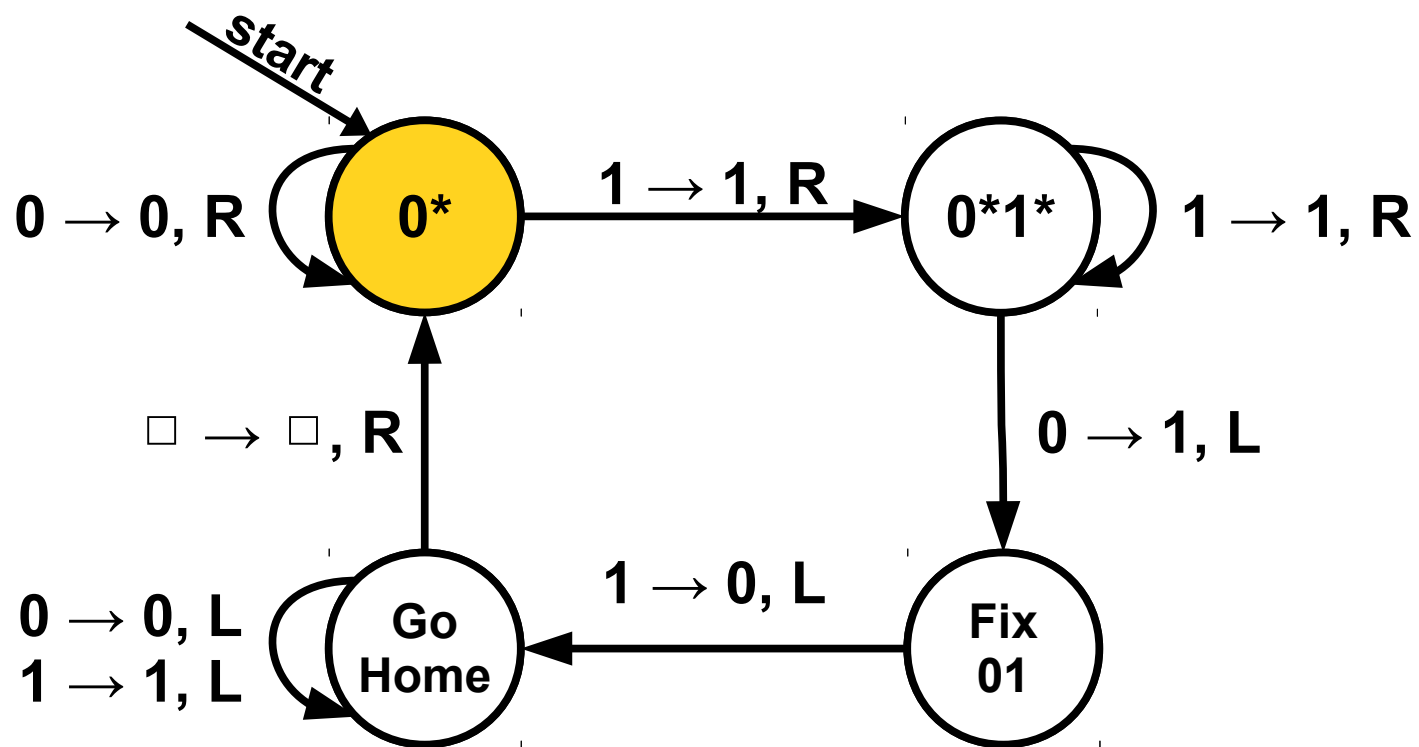


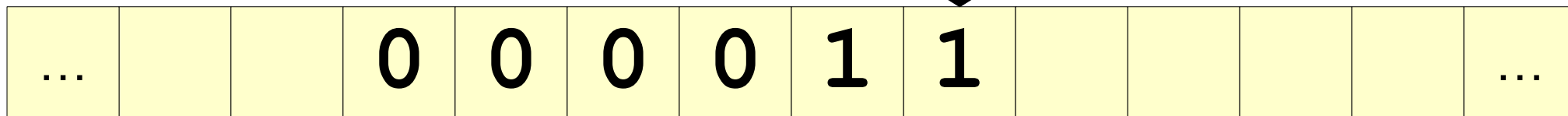
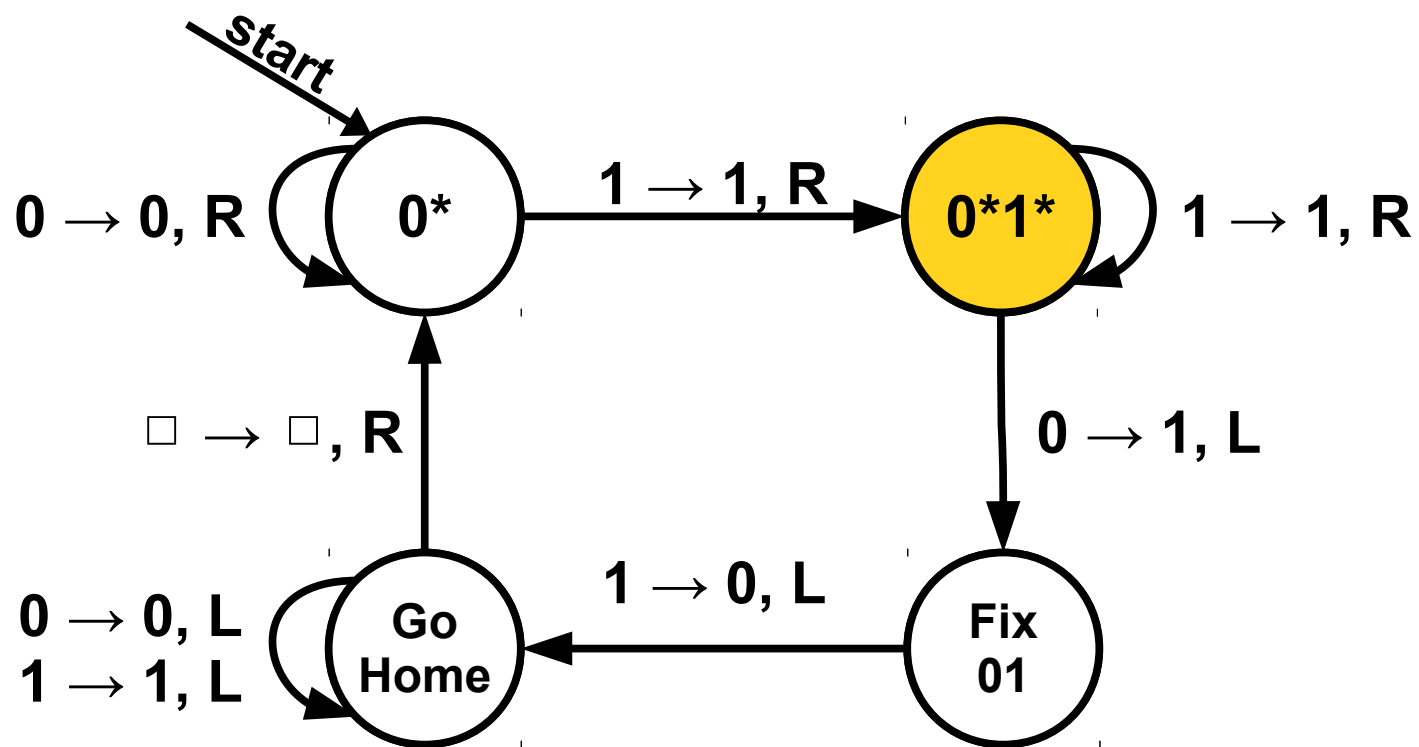


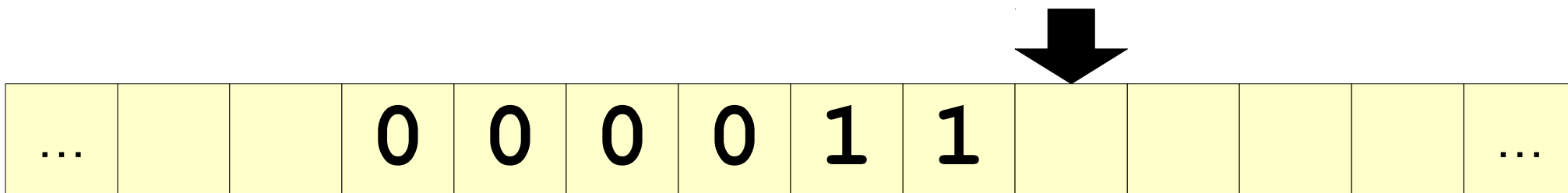
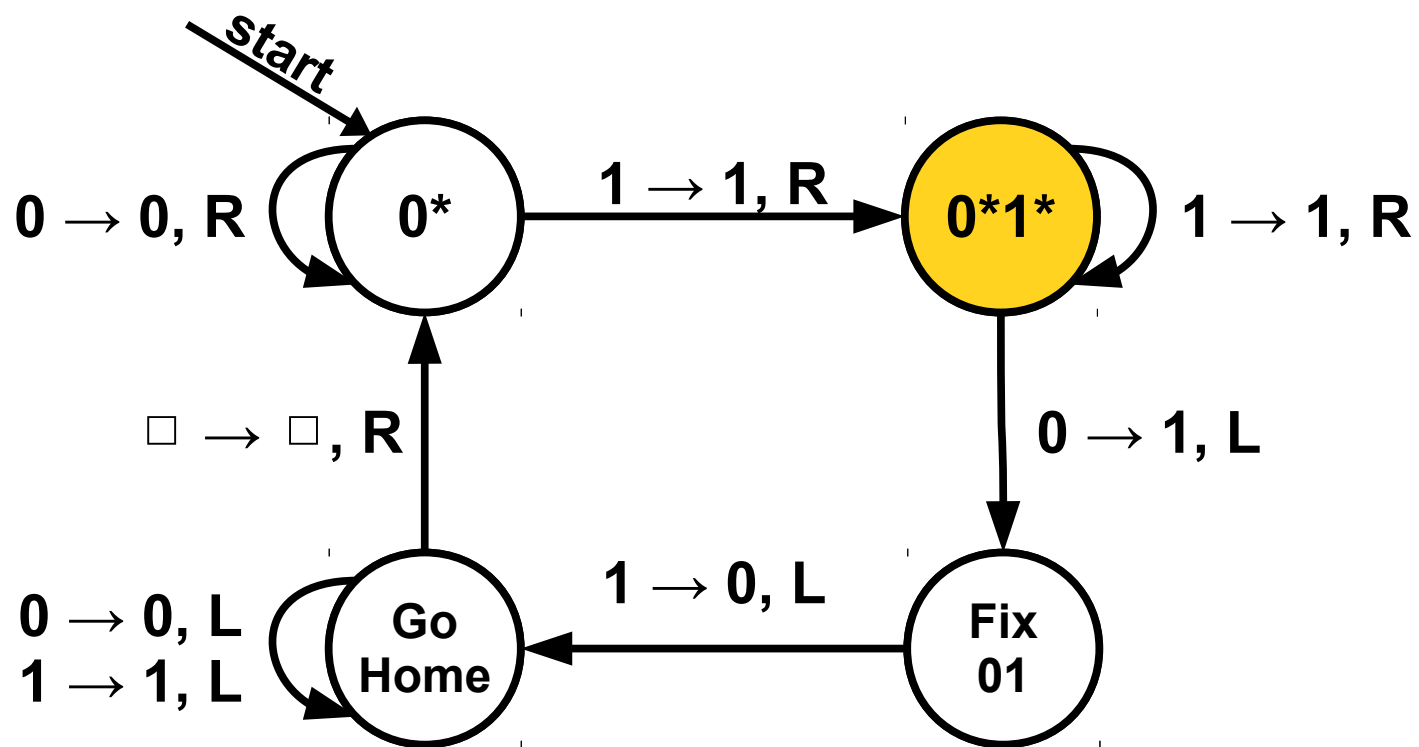




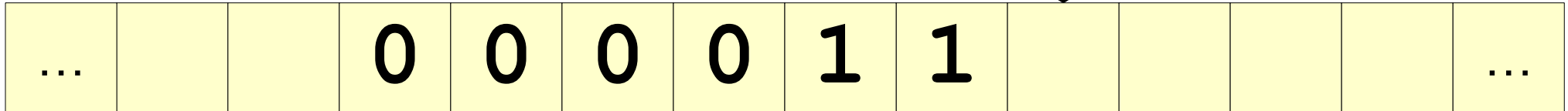
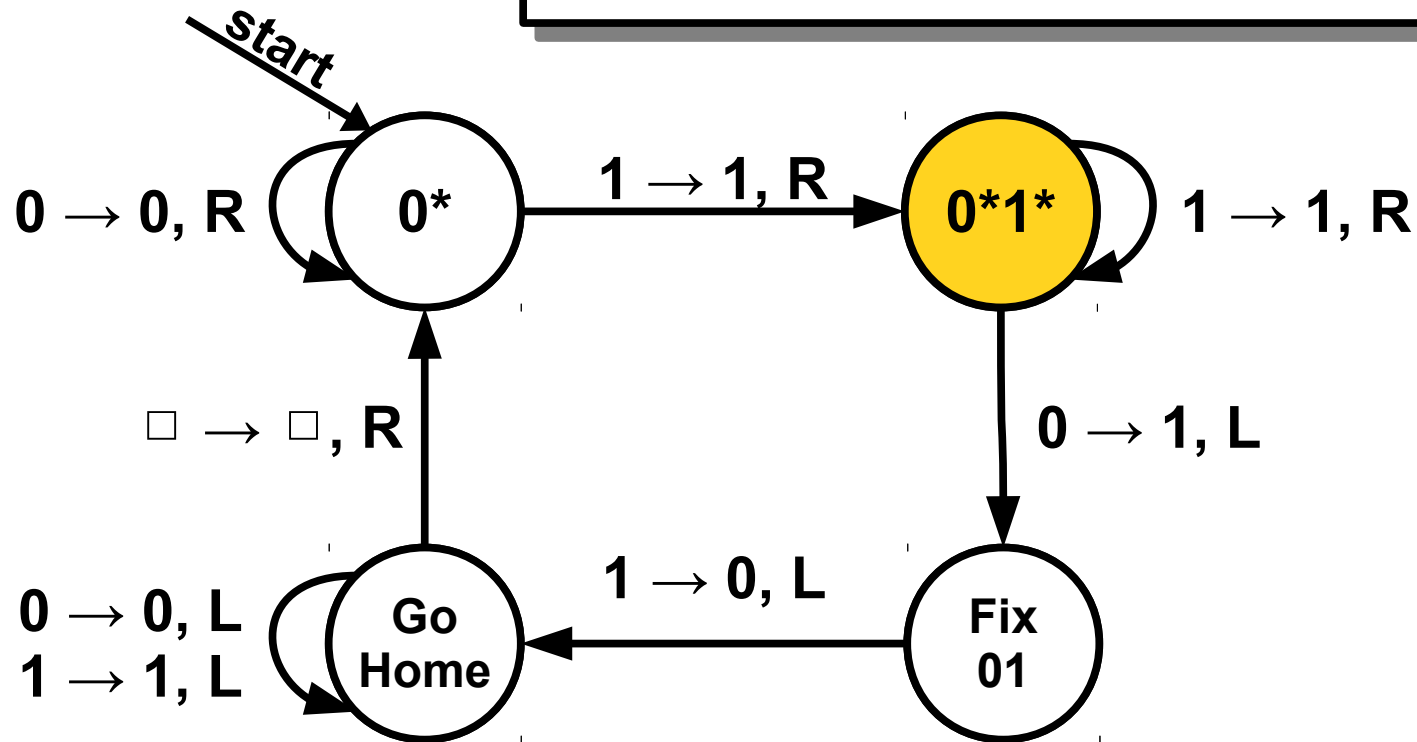


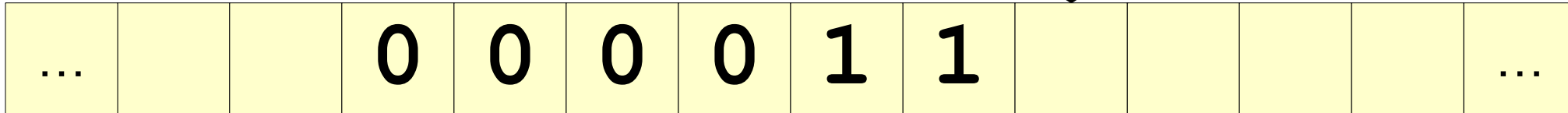
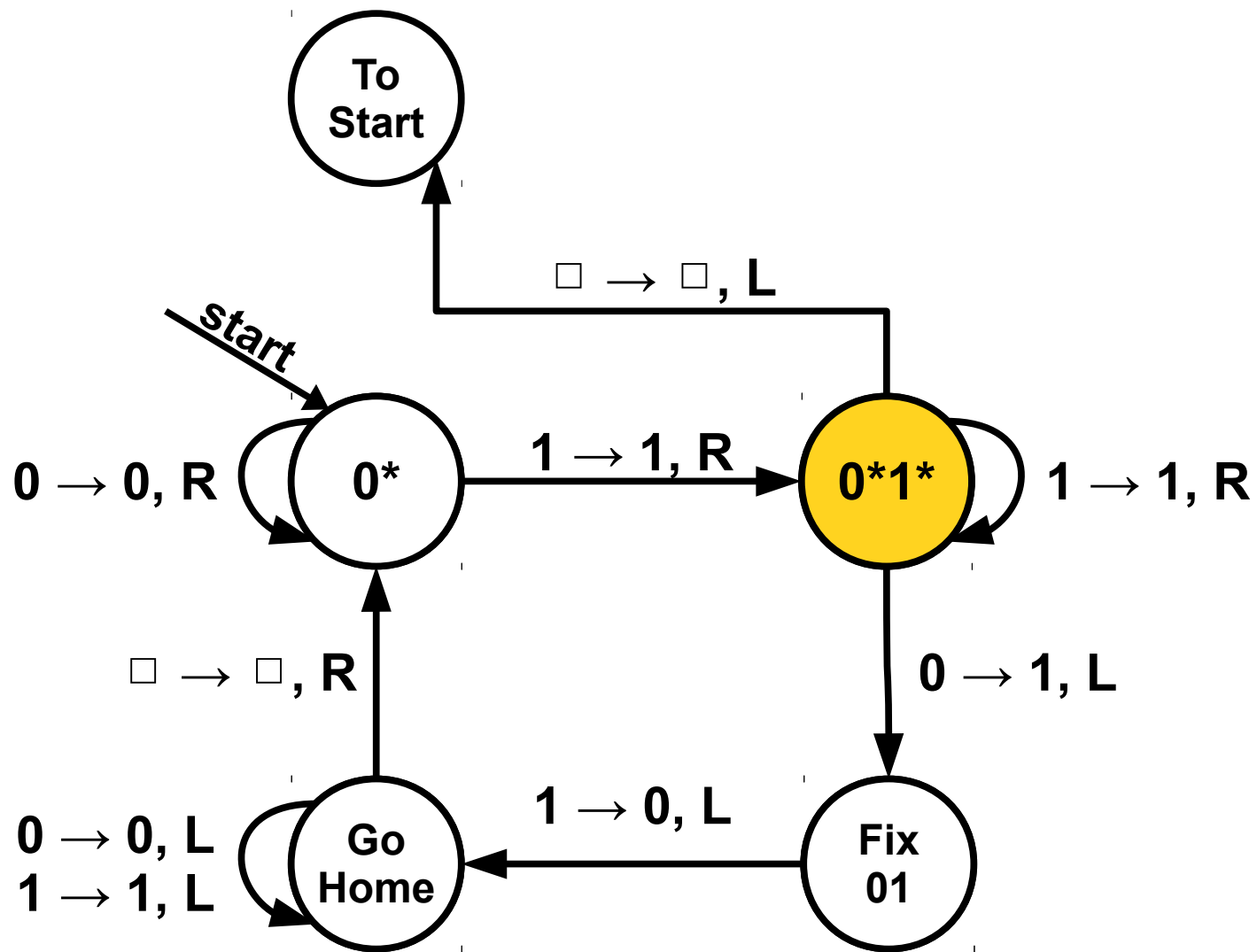


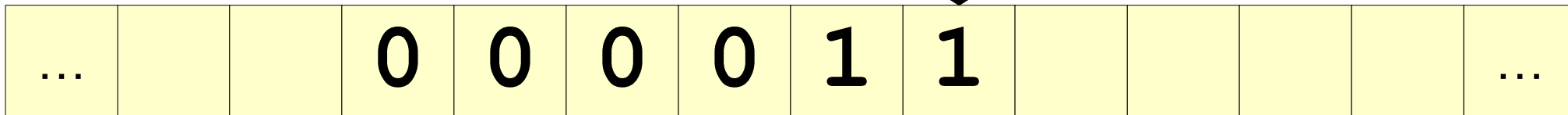
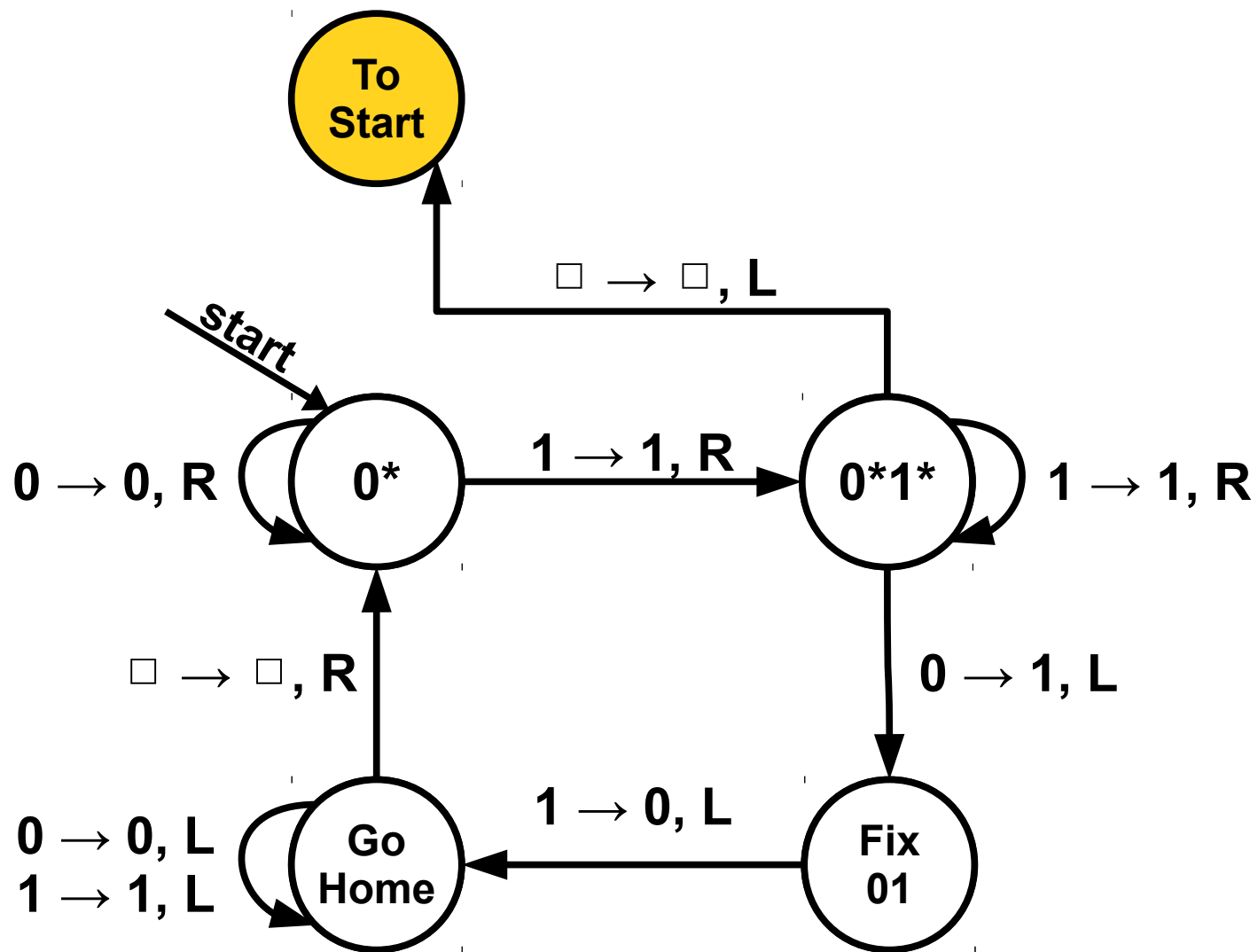


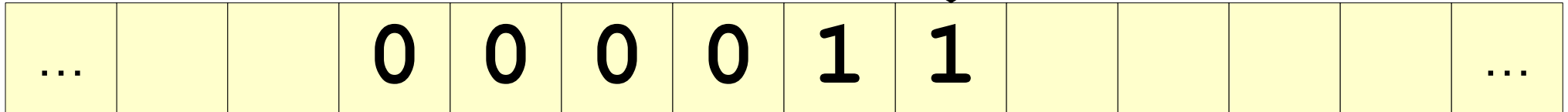
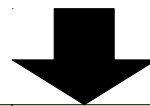
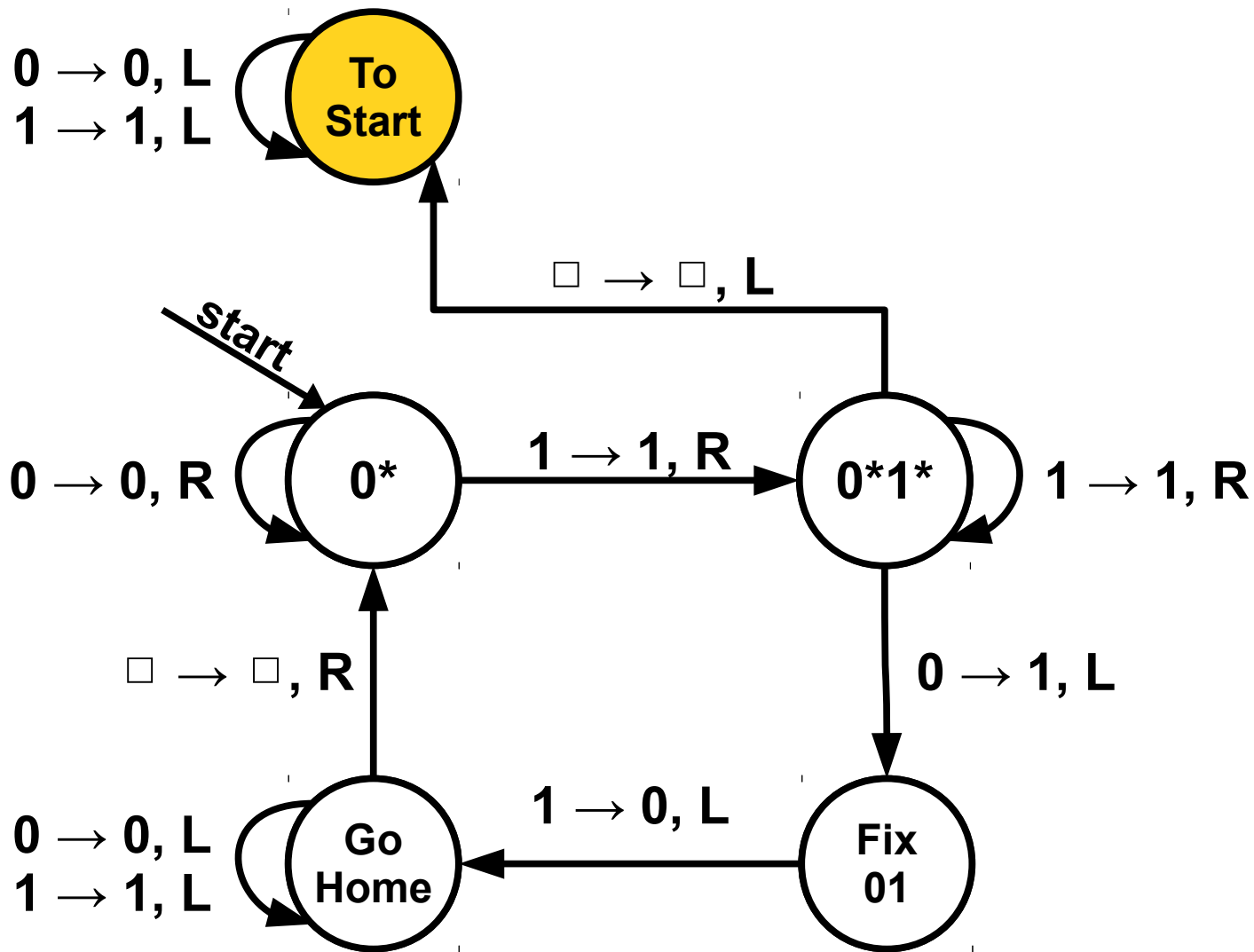


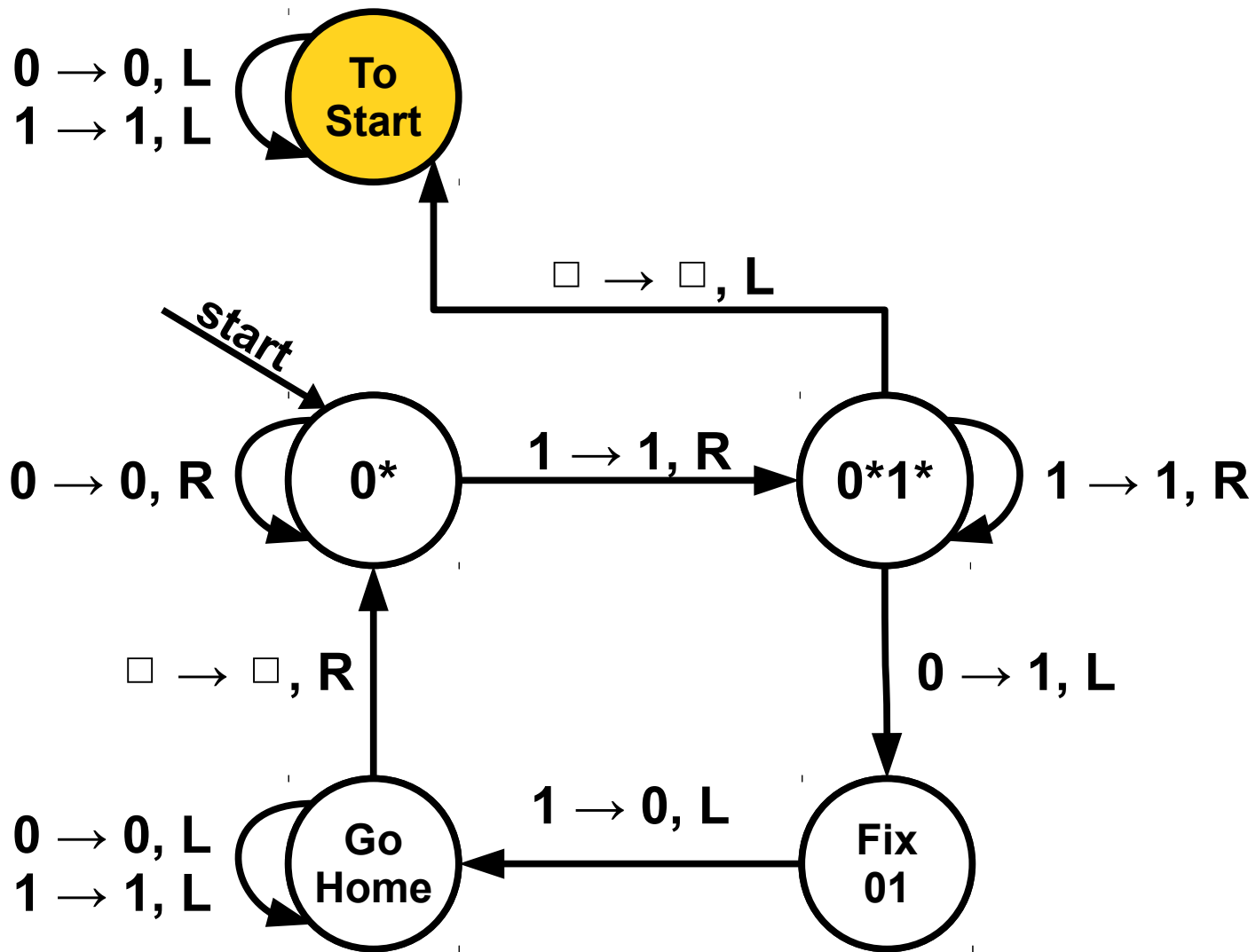
Our ultimate goal here was to sort everything so we could hand it off to the machine to check for $0^n 1^n$. Let's rewind the tape head back to the start.

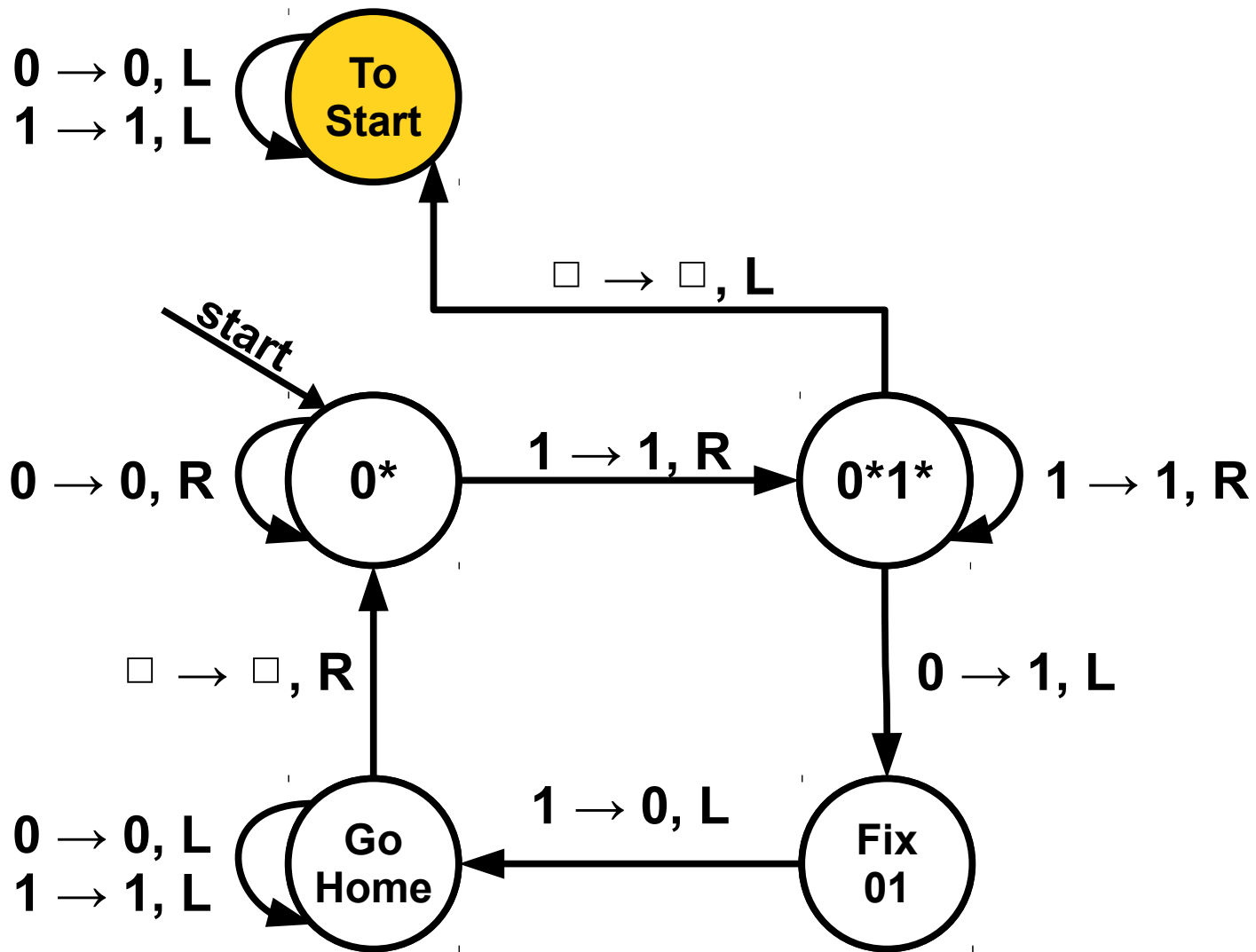


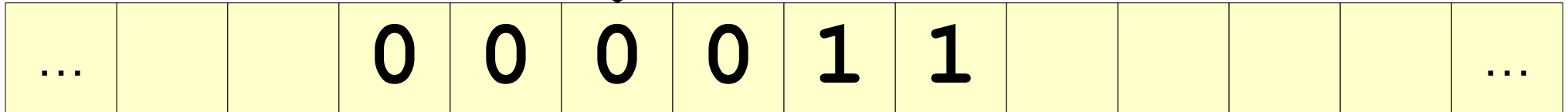
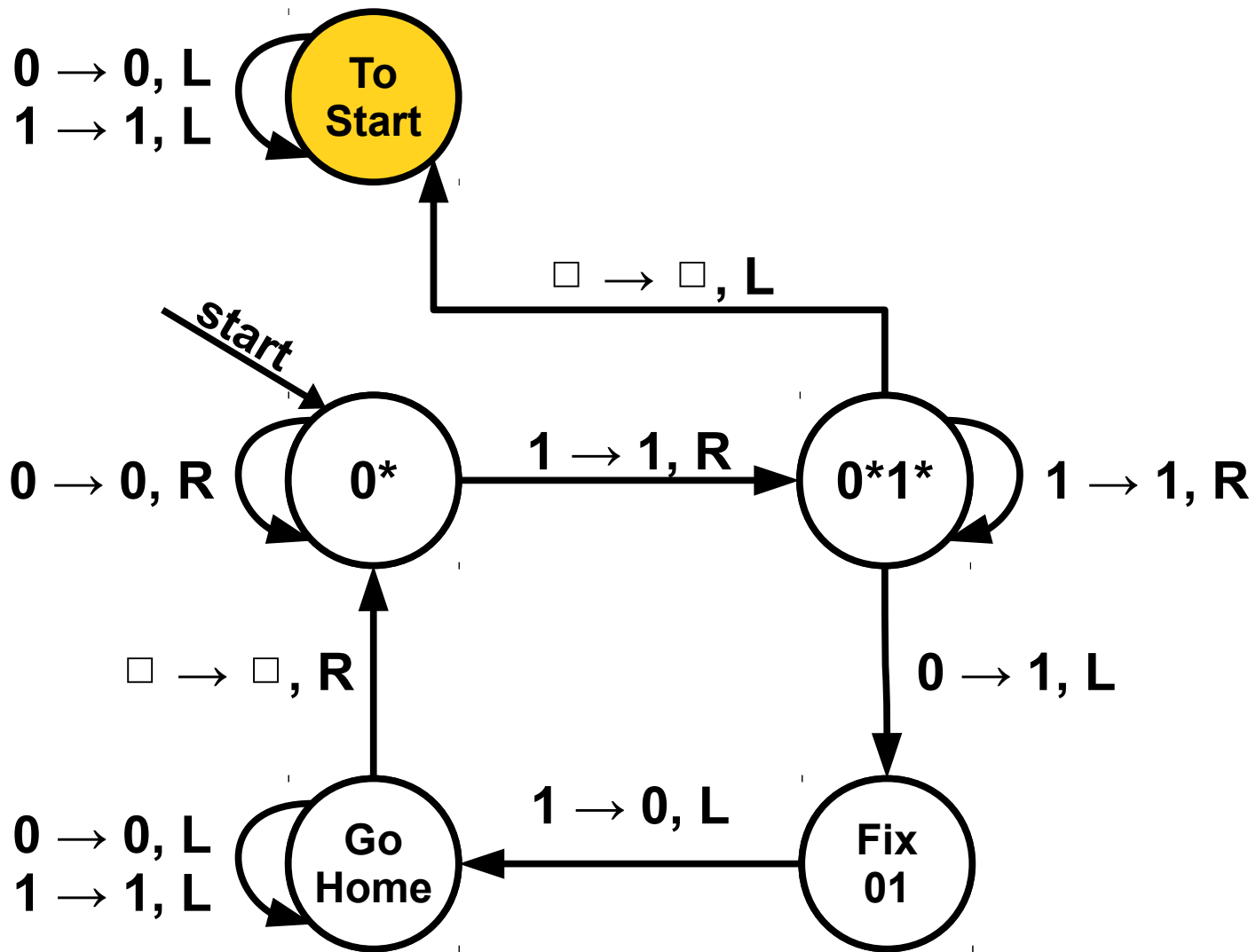


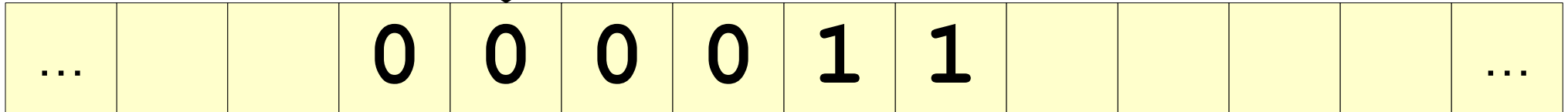
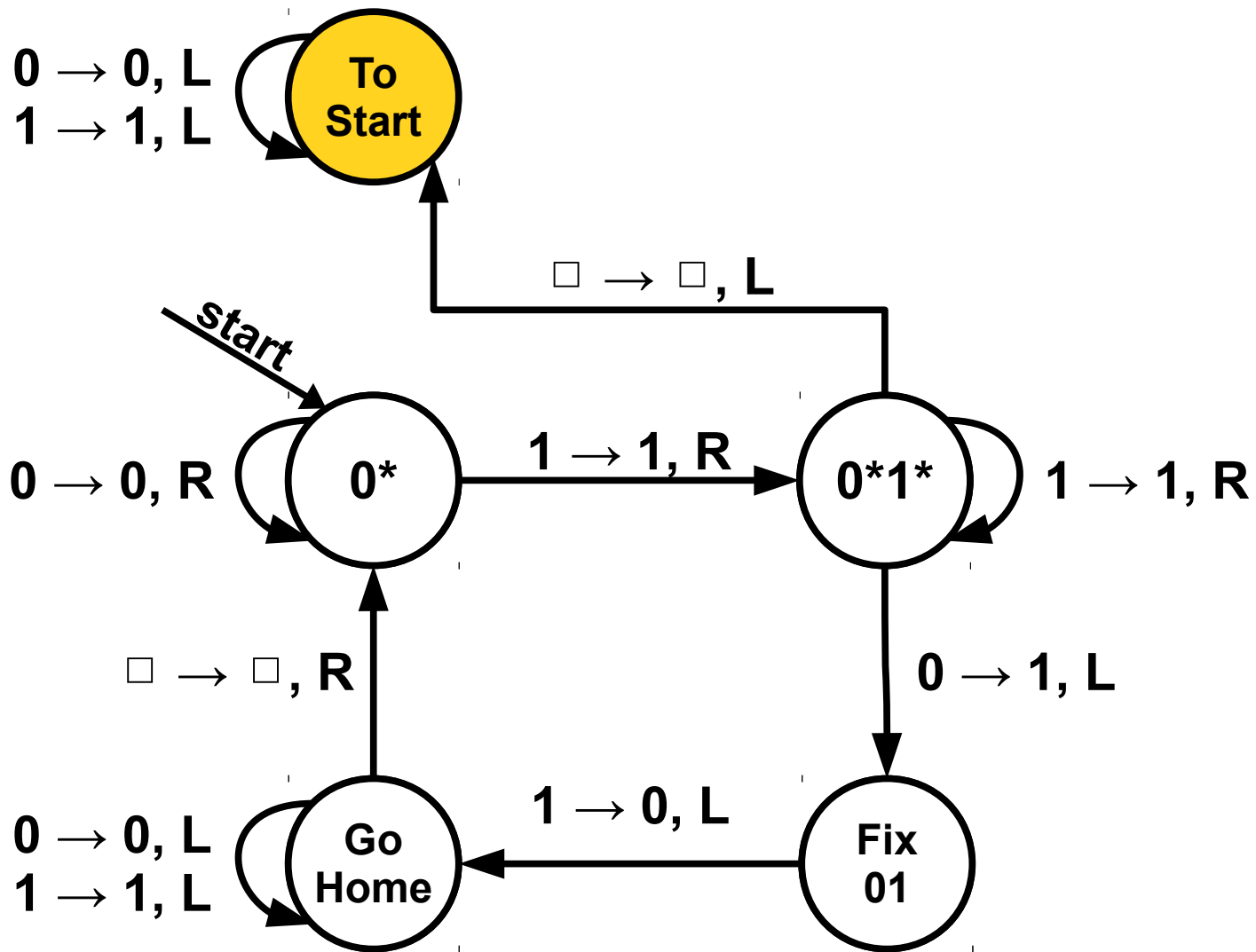


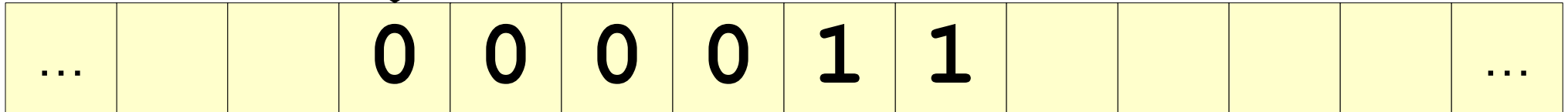
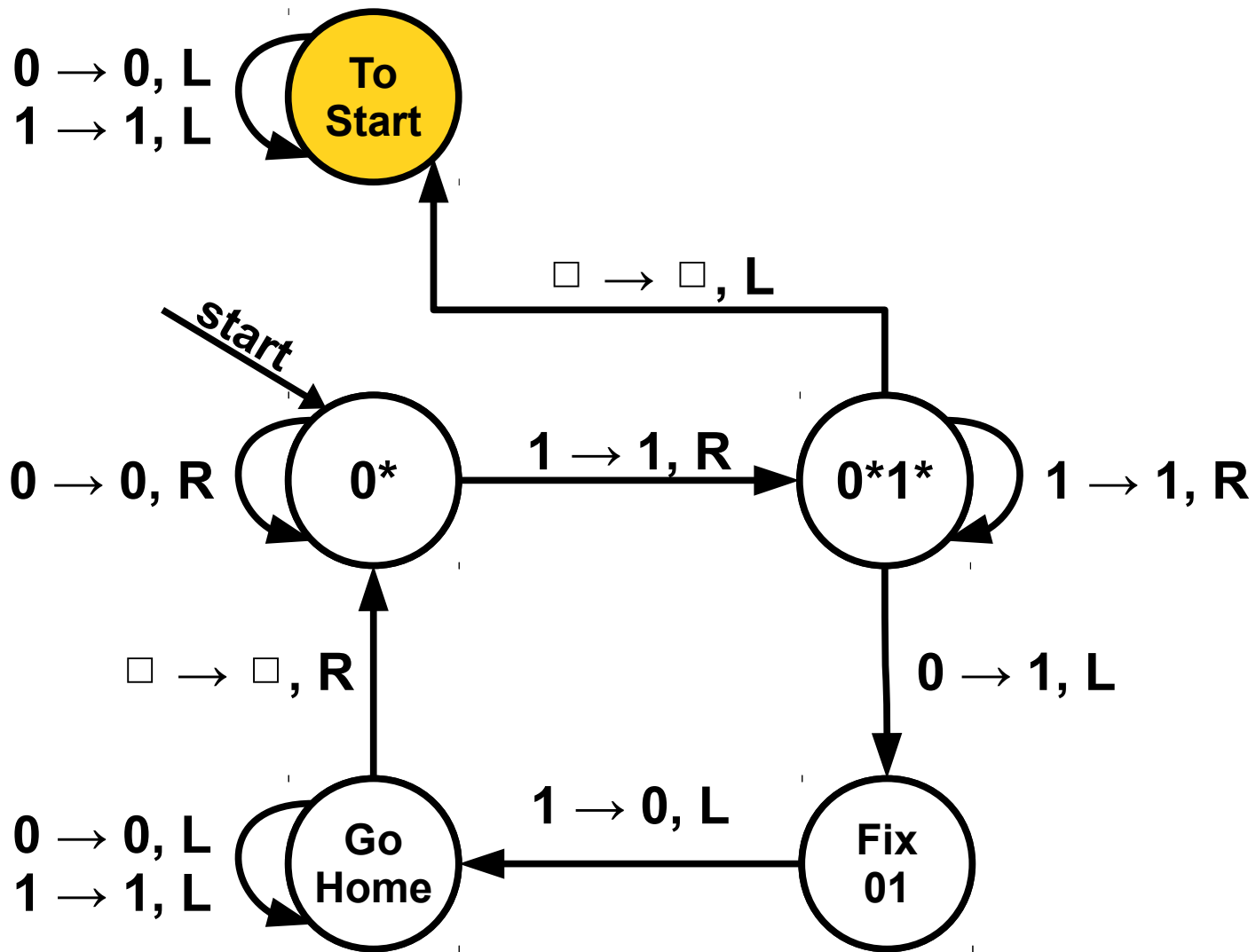


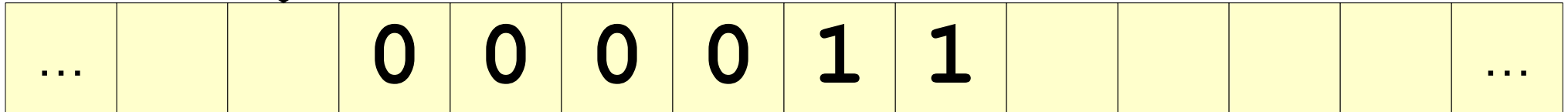
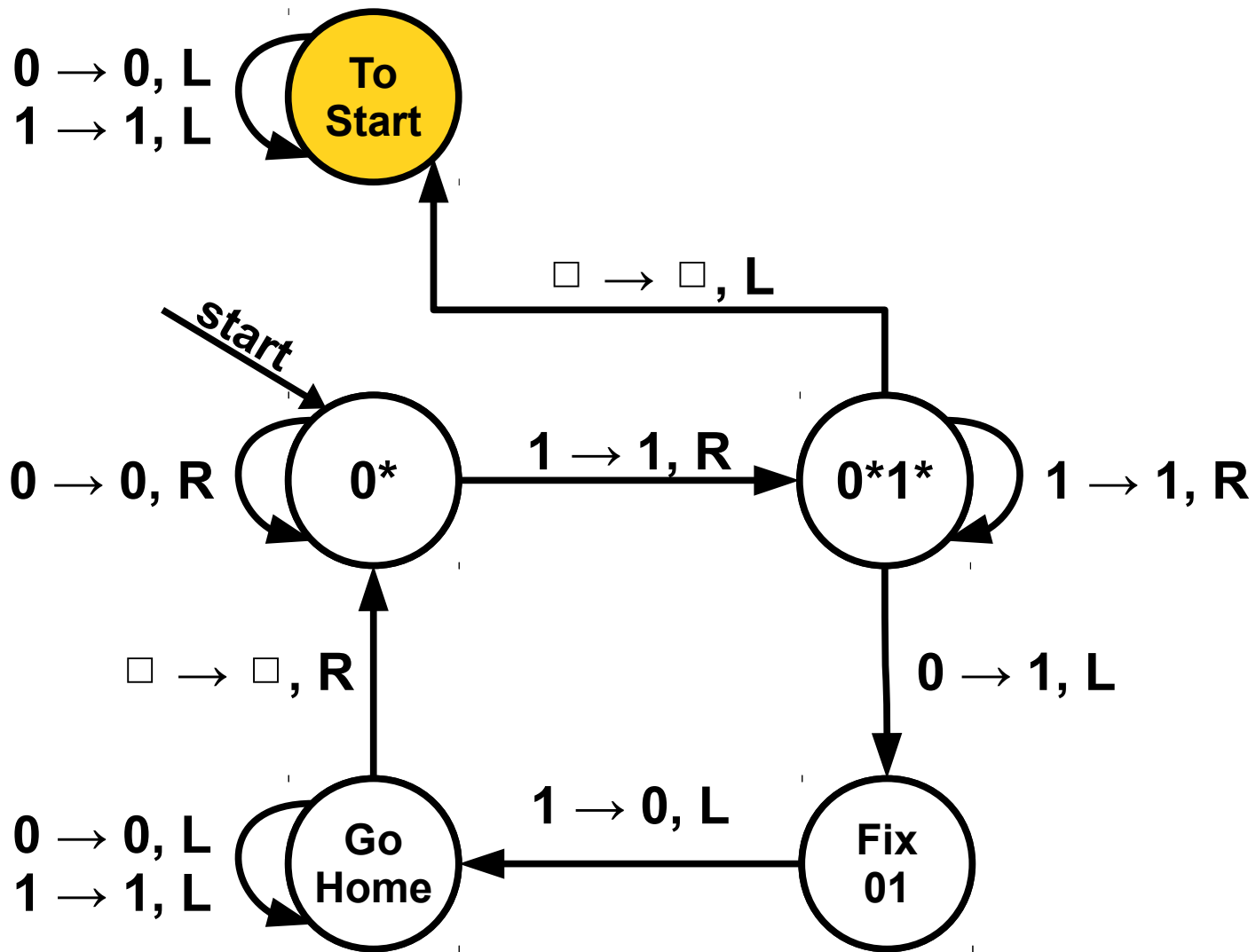


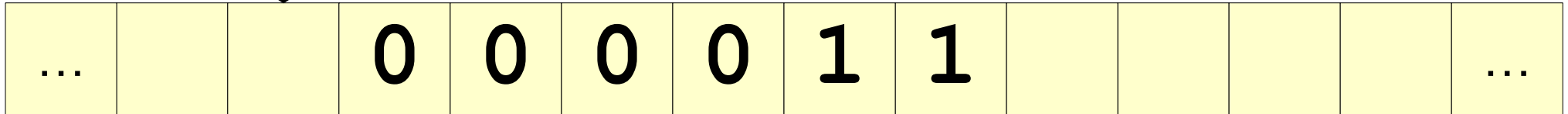
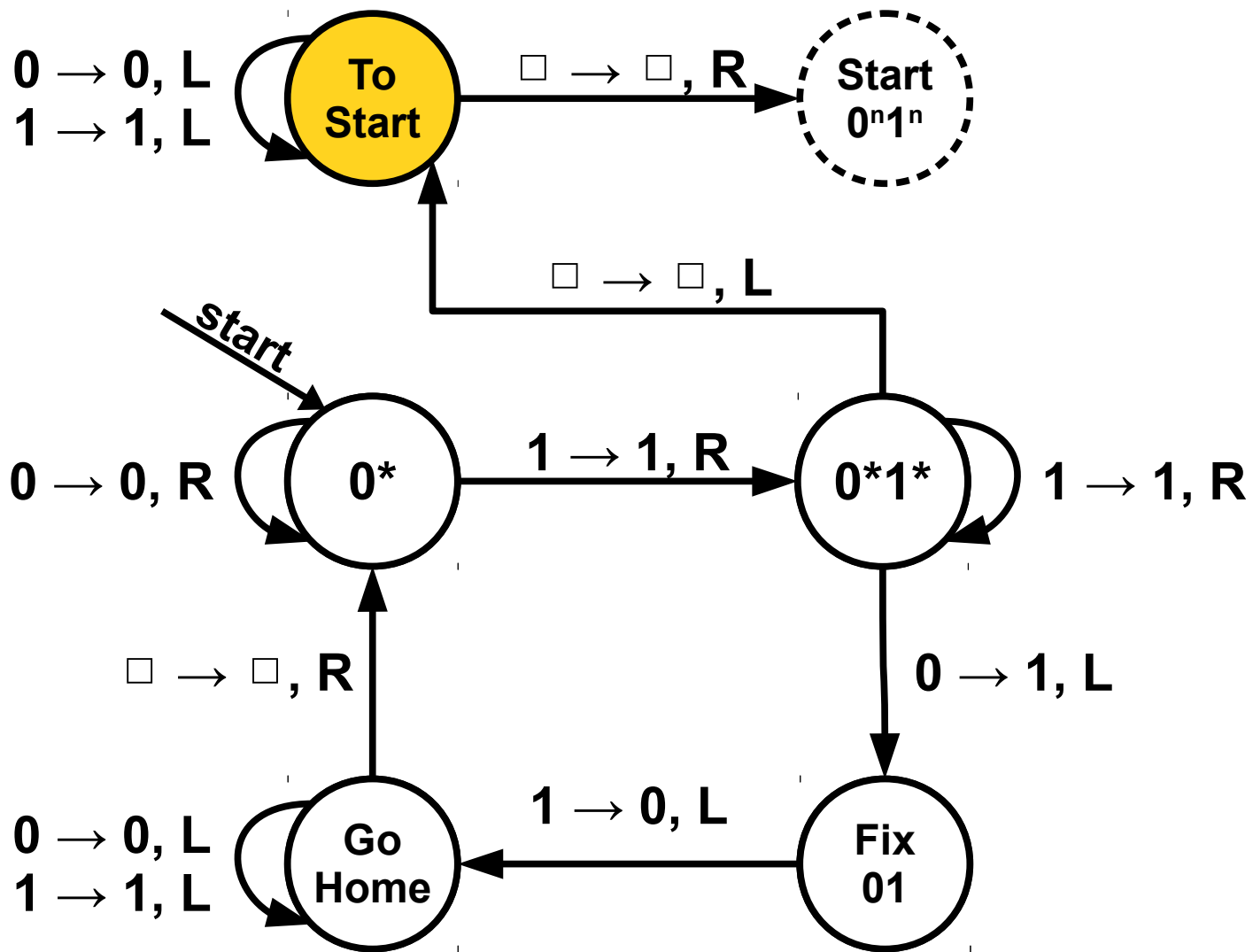


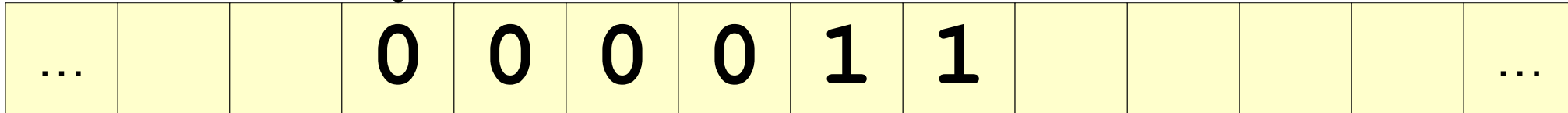
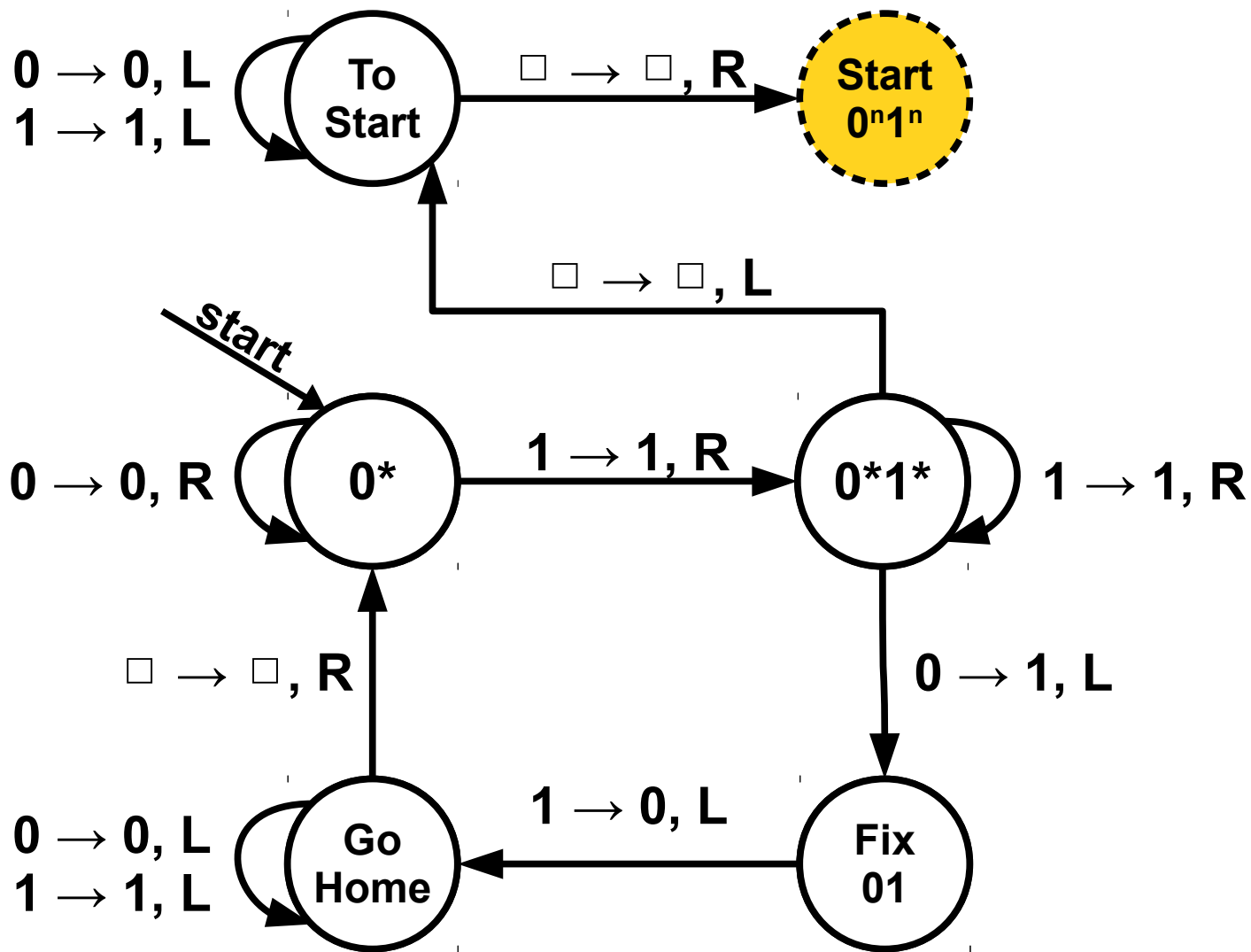


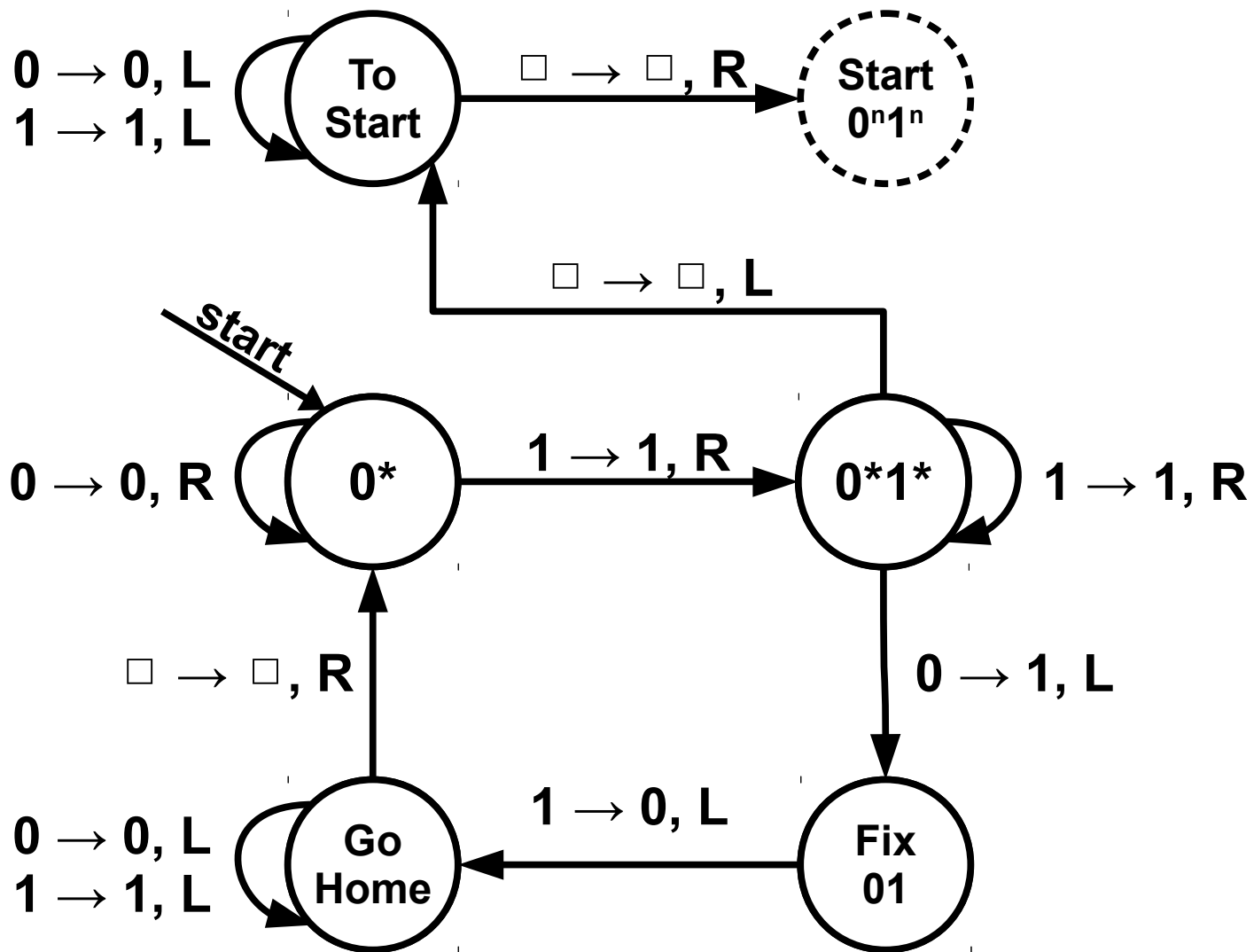


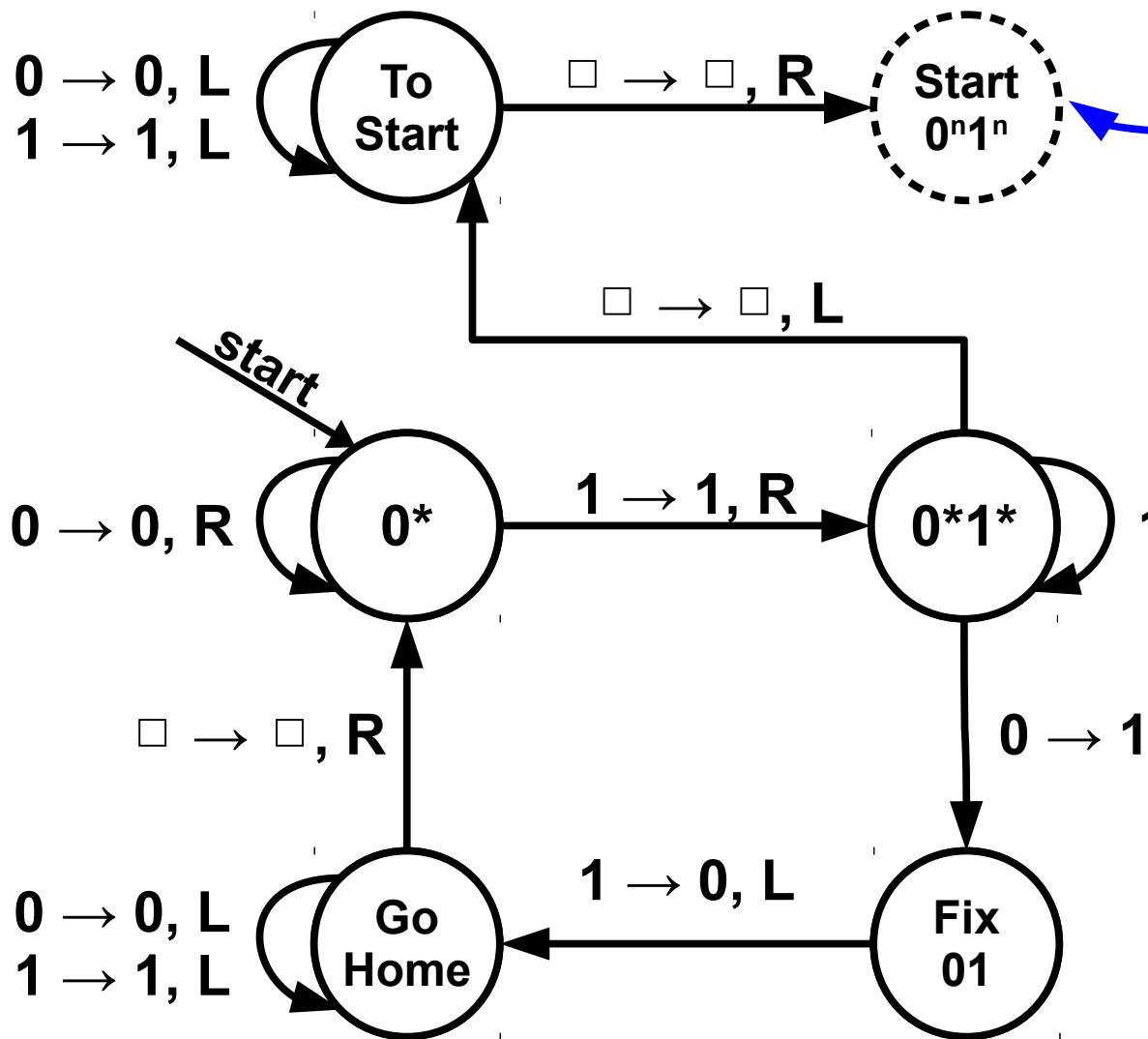




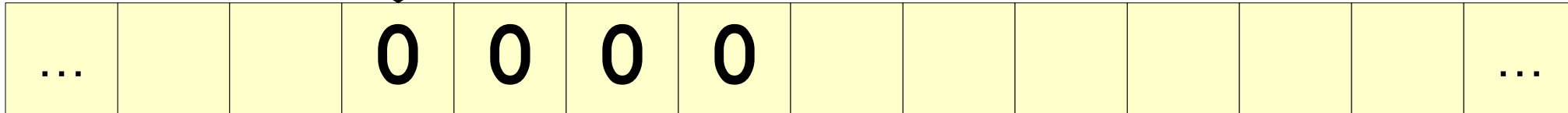
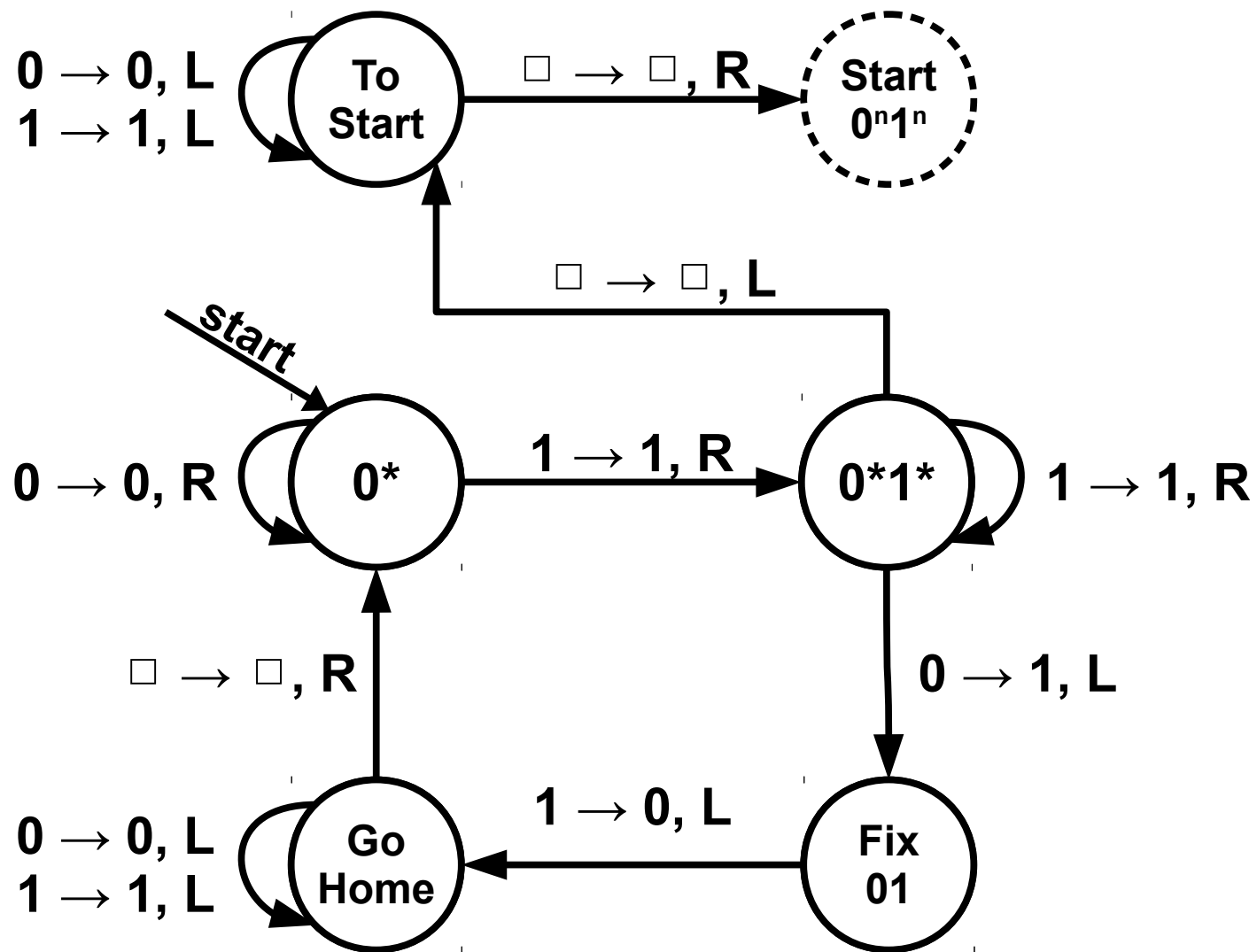


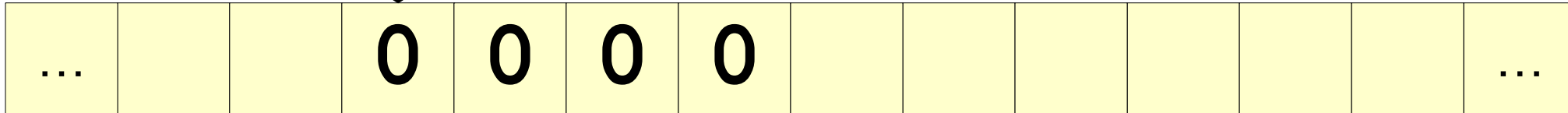
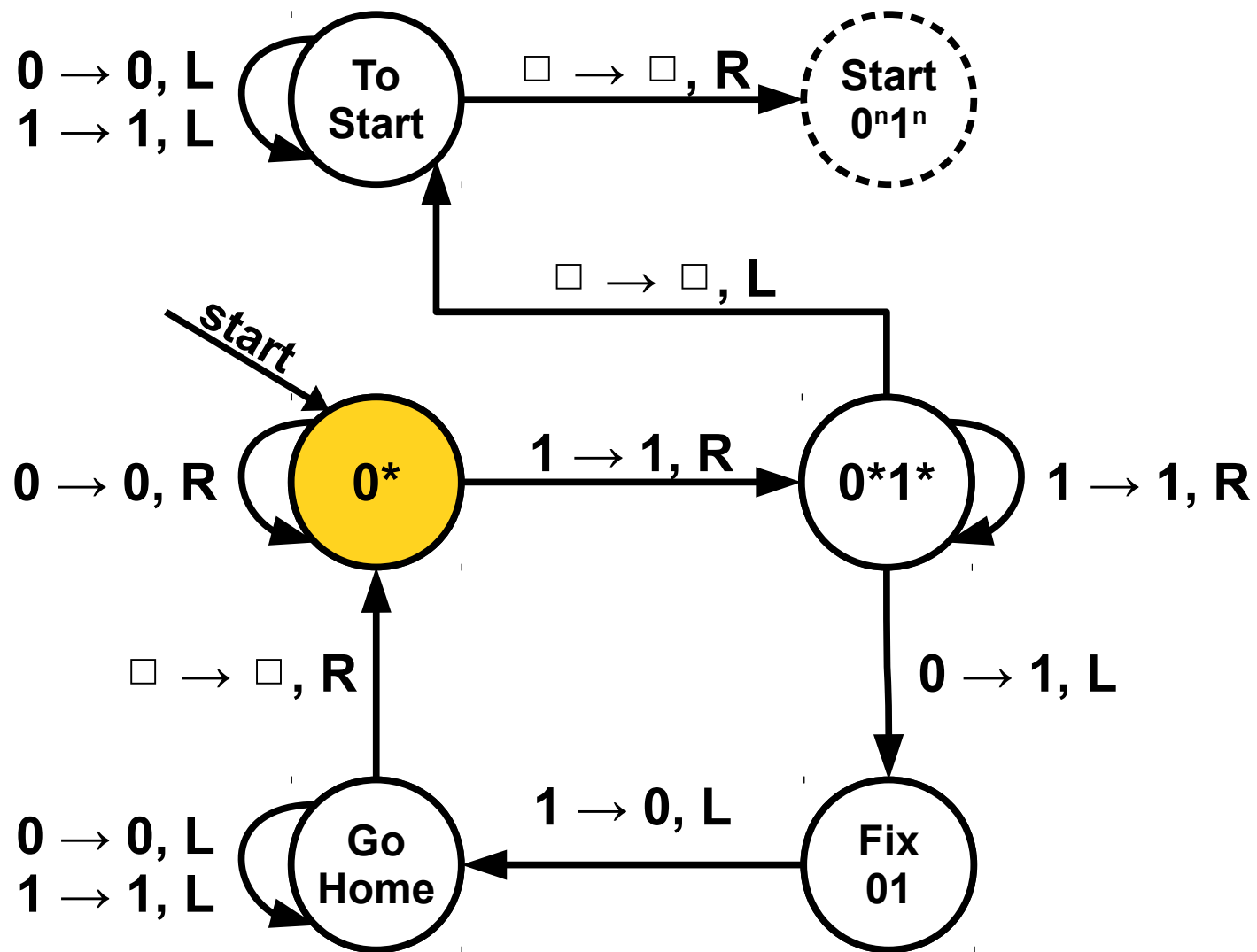


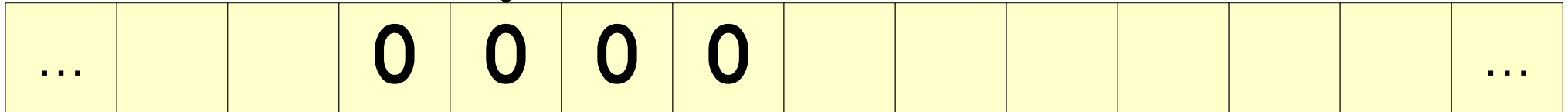
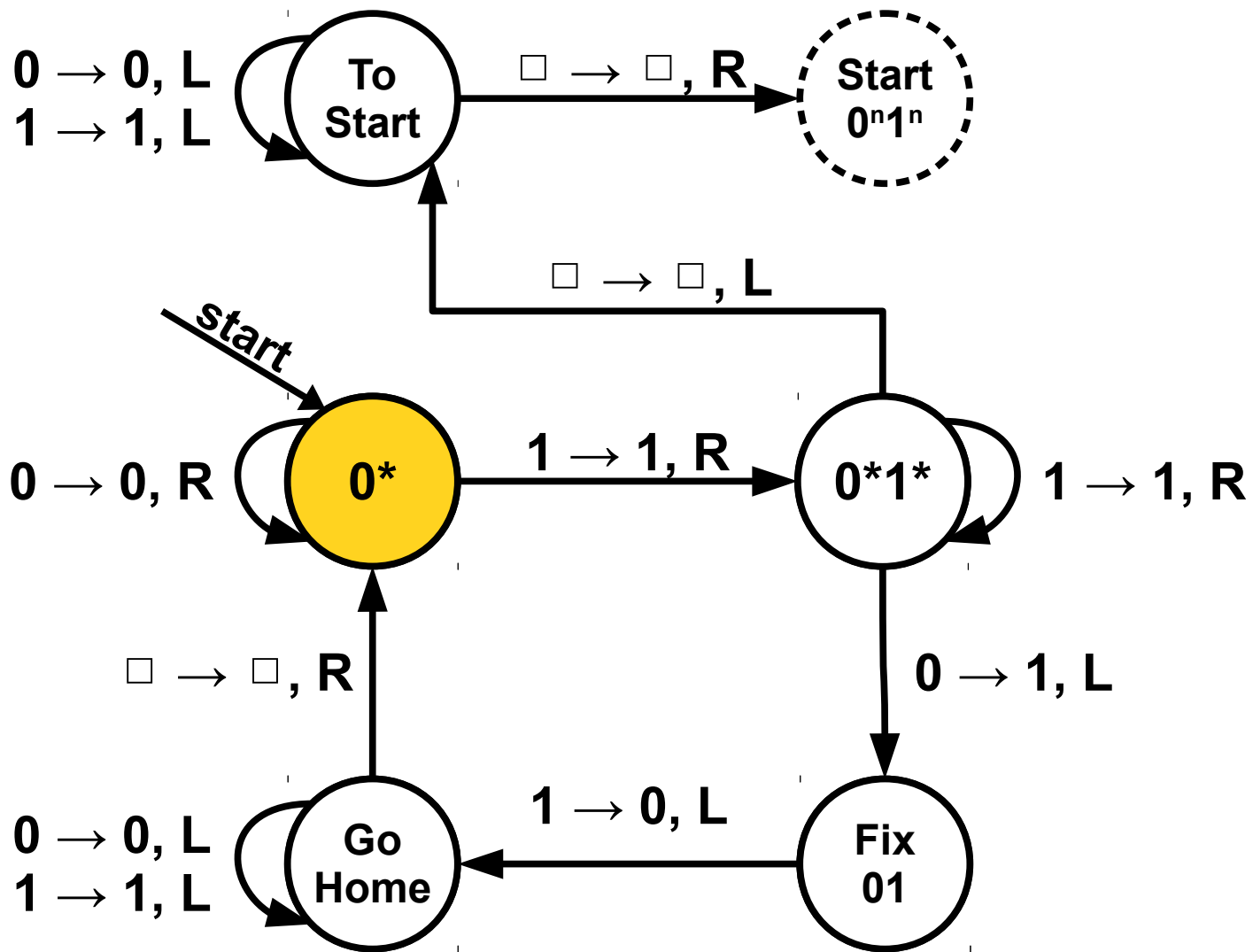


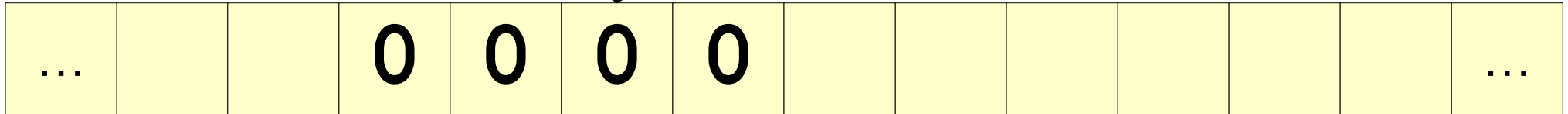
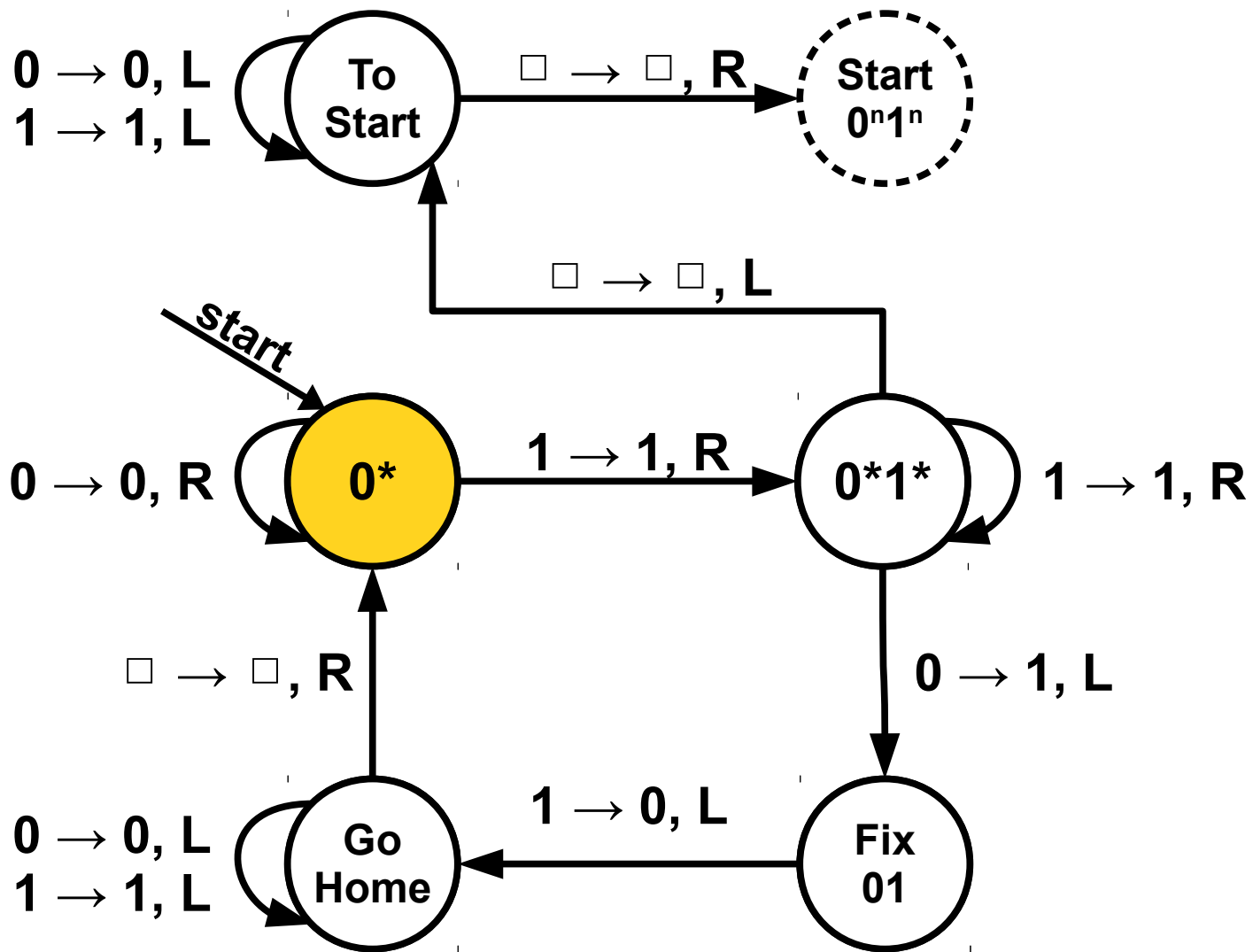


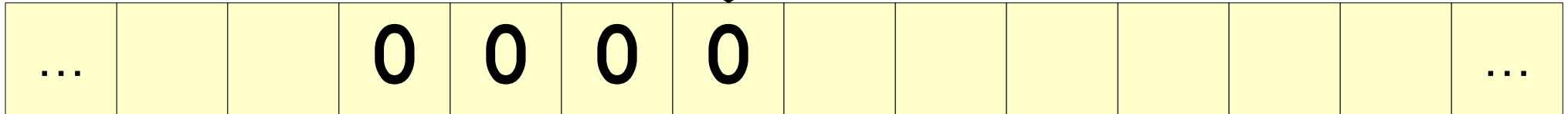
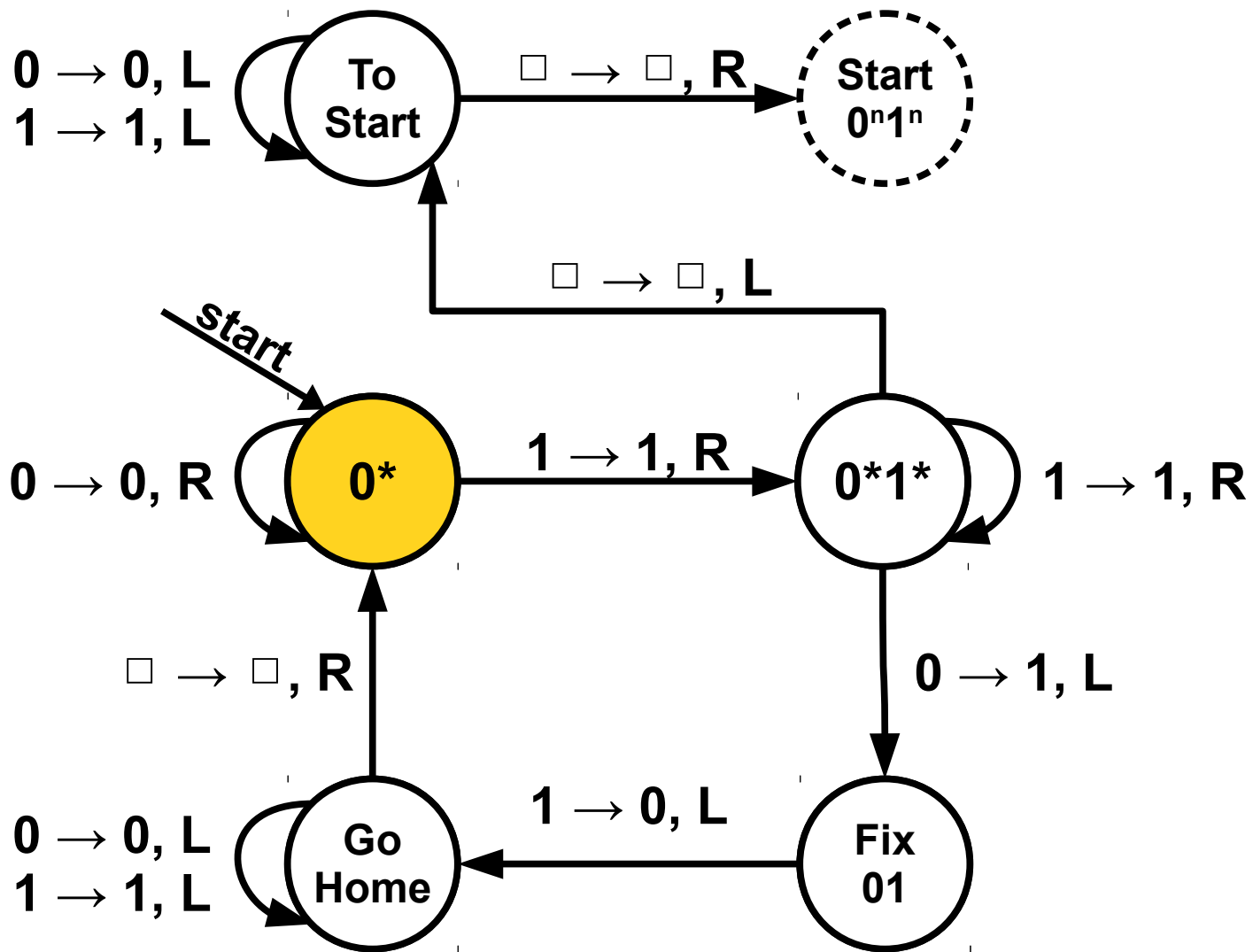
This is just a placeholder. Imagine snapping in the entire TM for $0^n 1^n$ into this diagram, putting the start state in the dashed area.

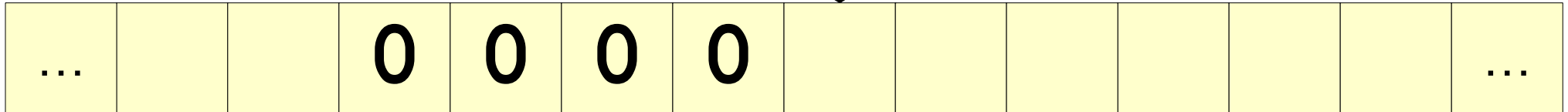
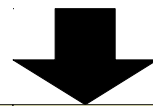
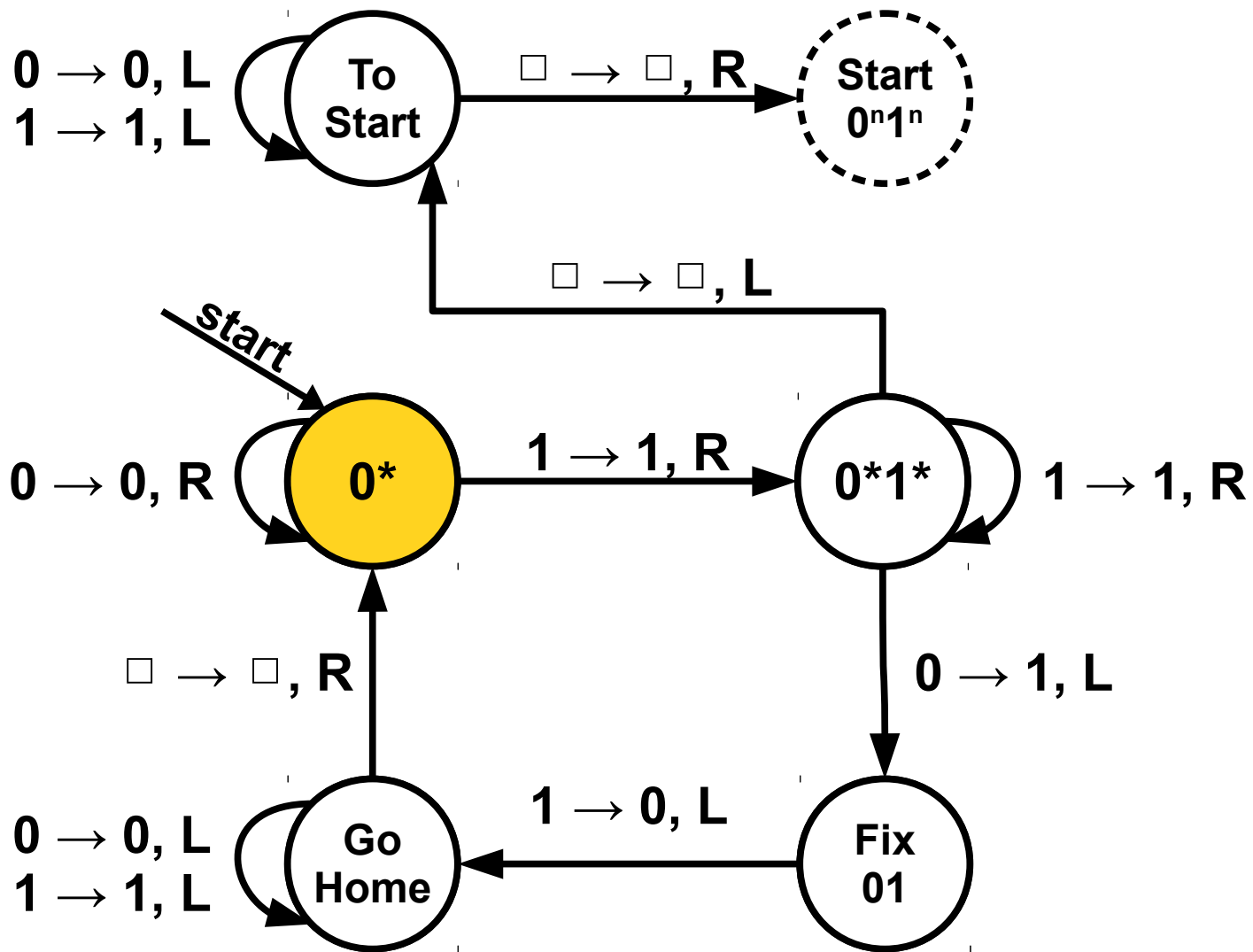


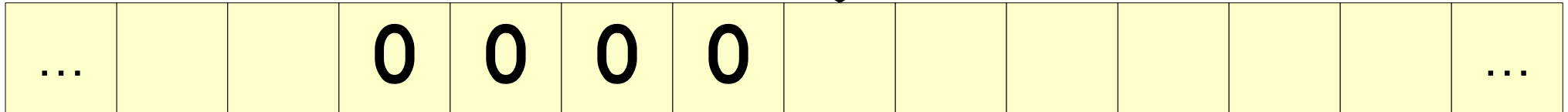
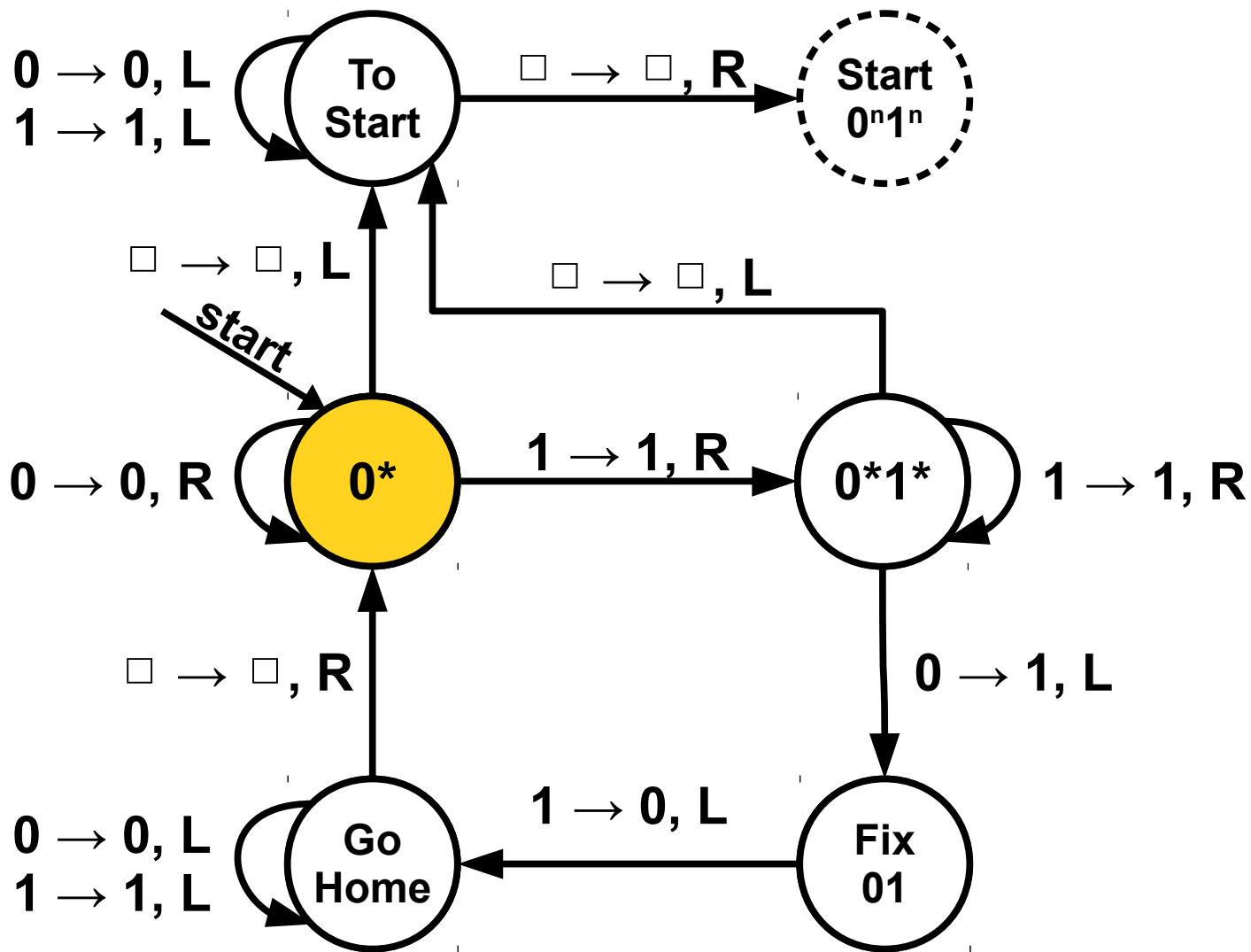


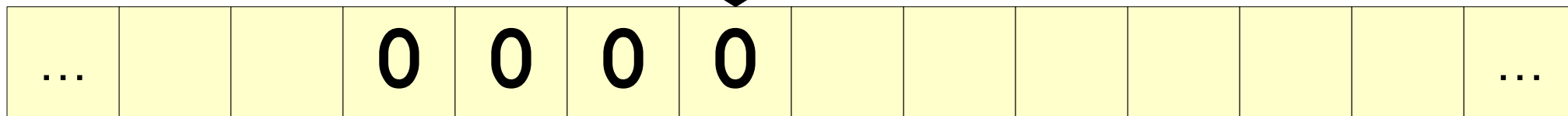
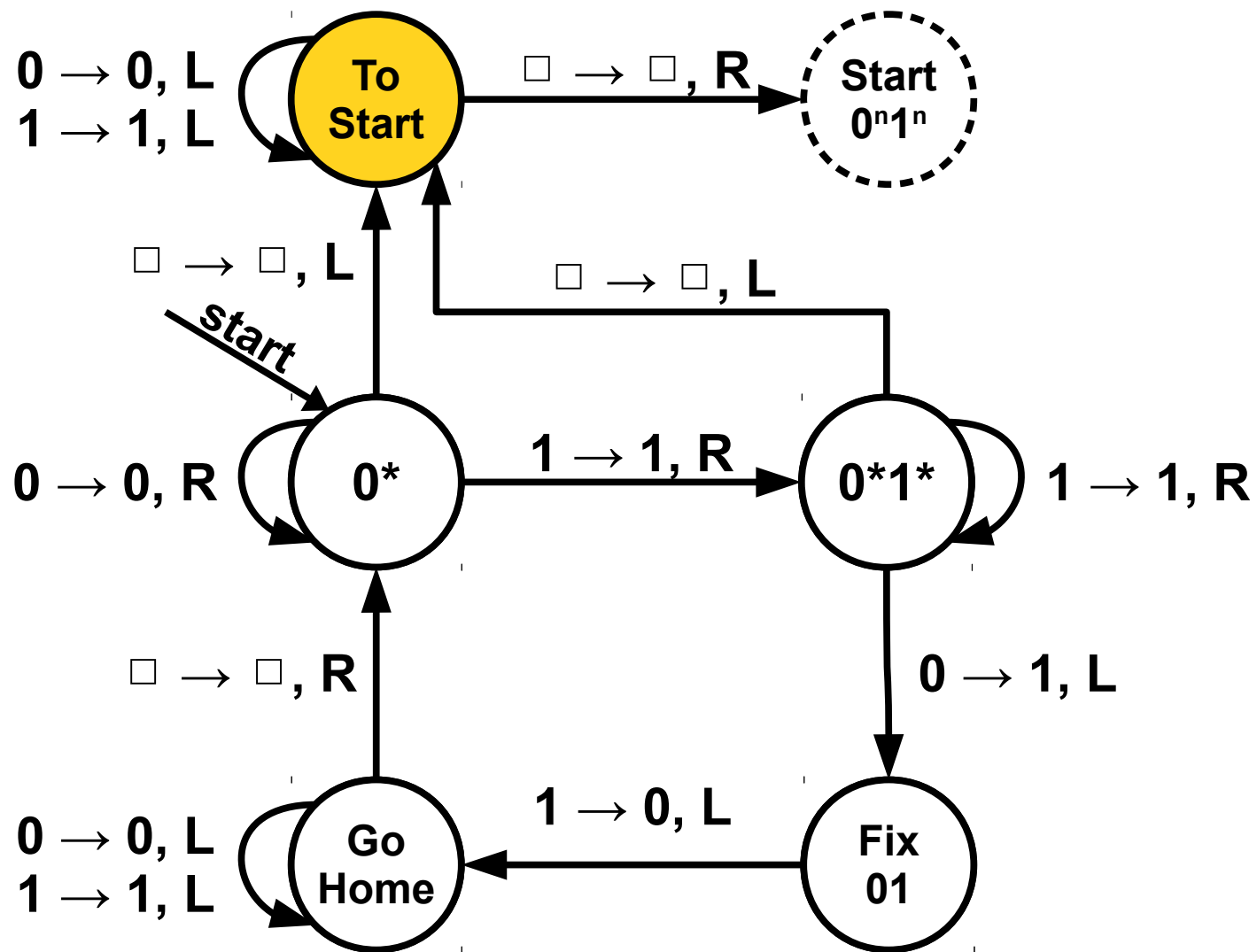


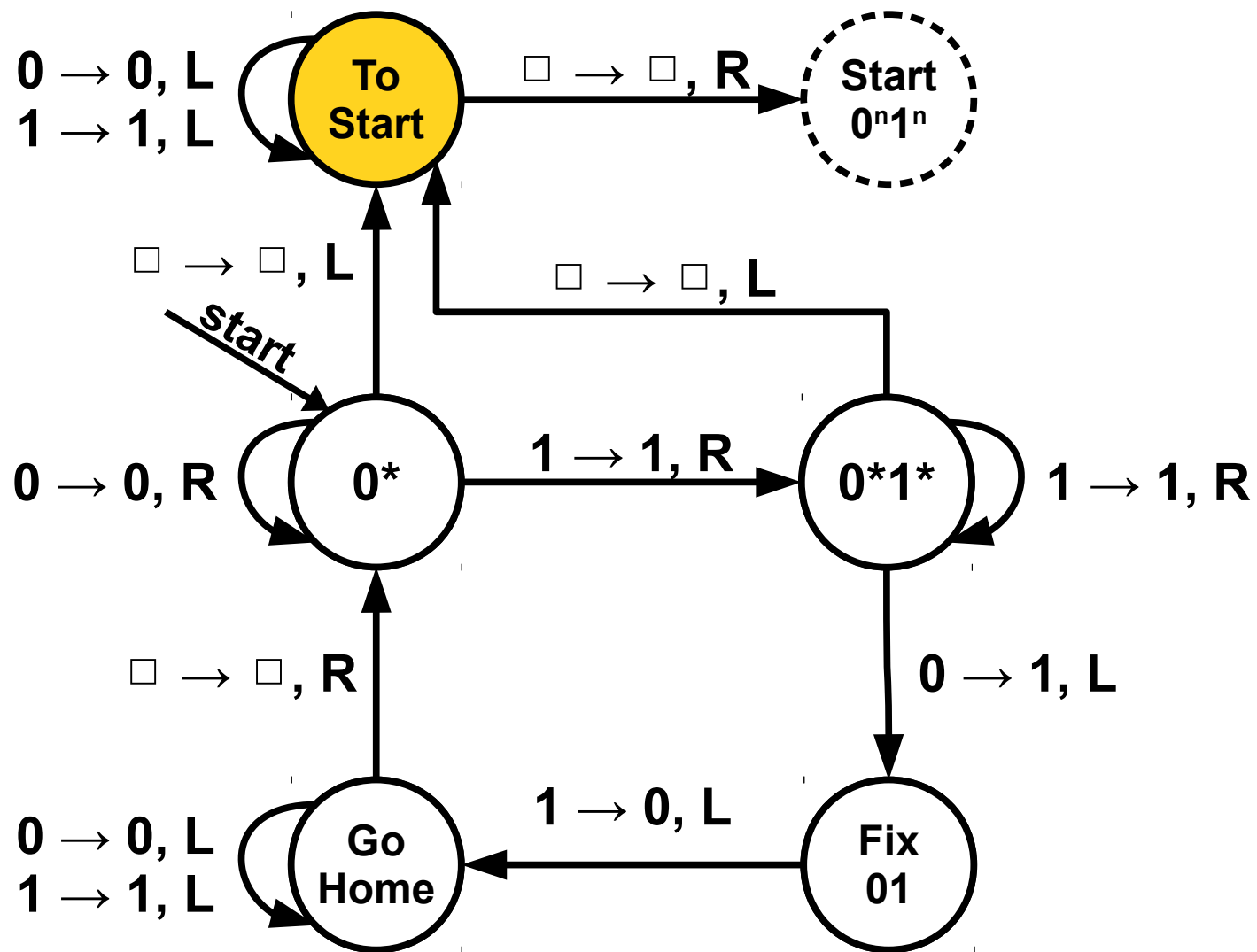


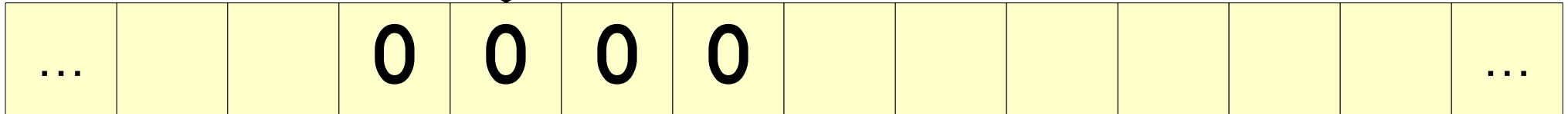
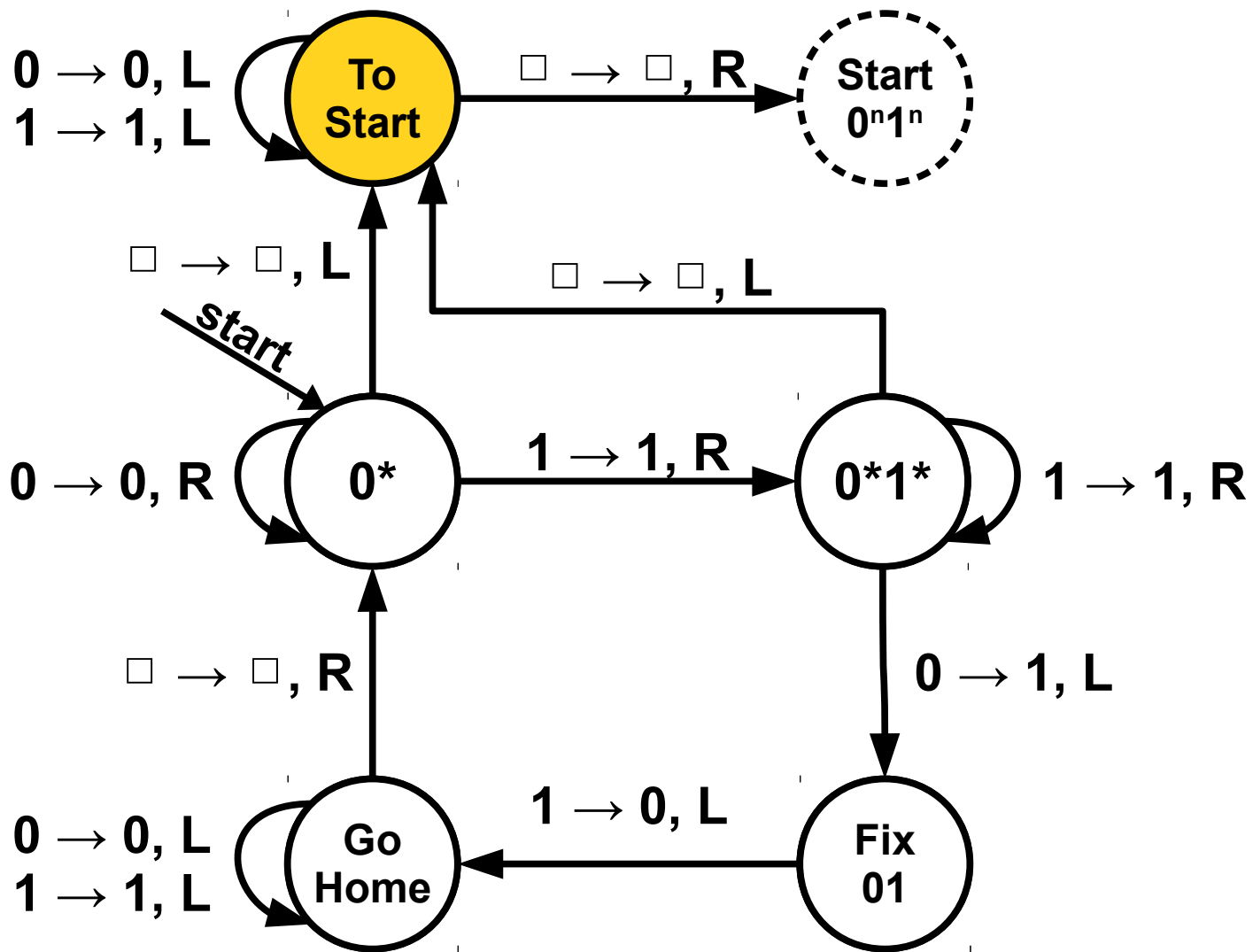


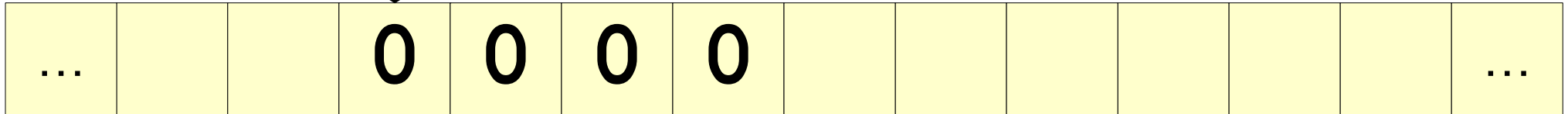
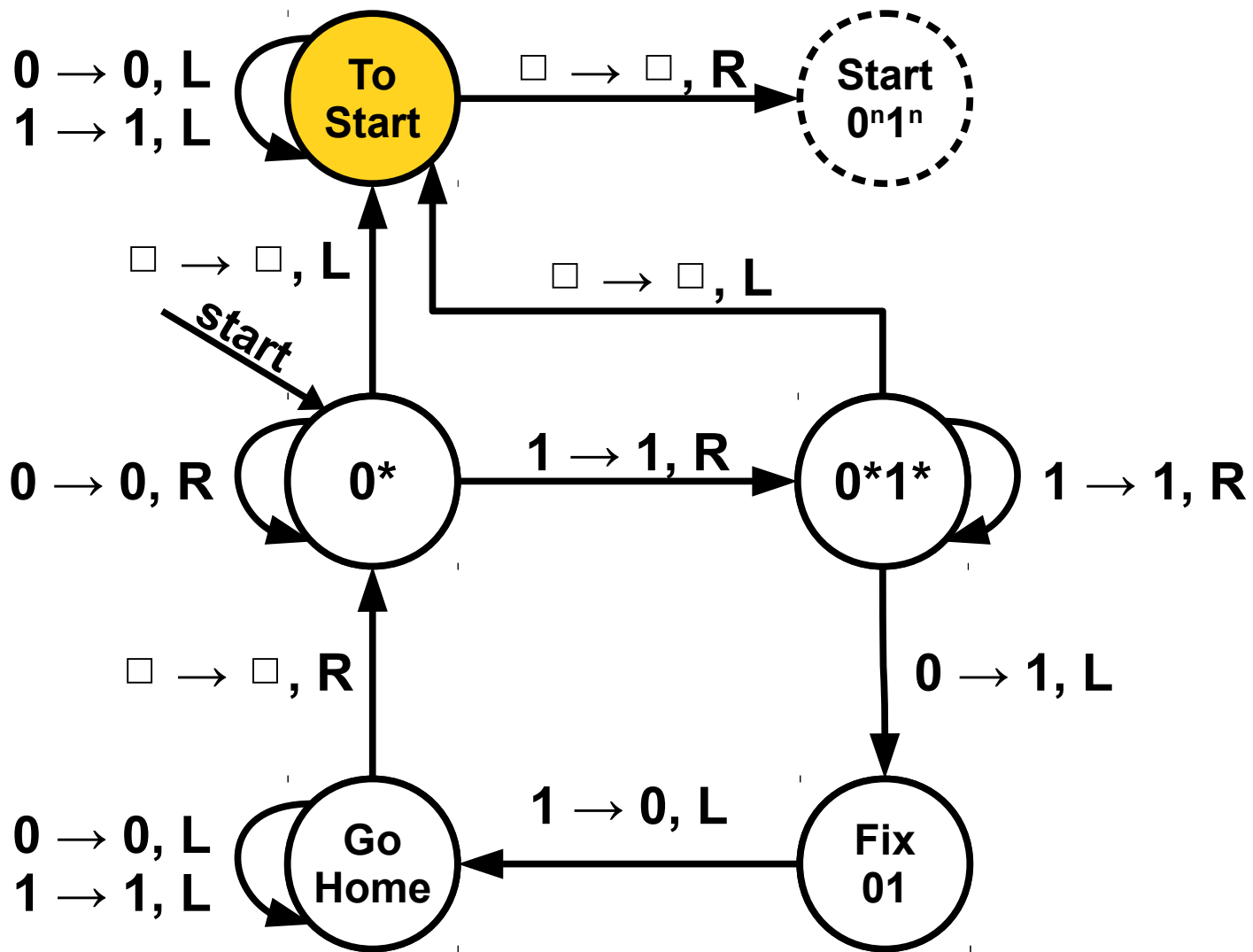


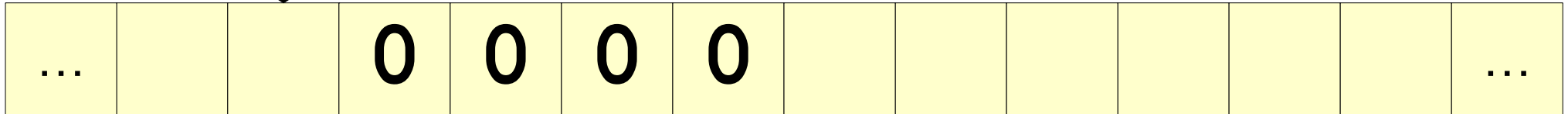
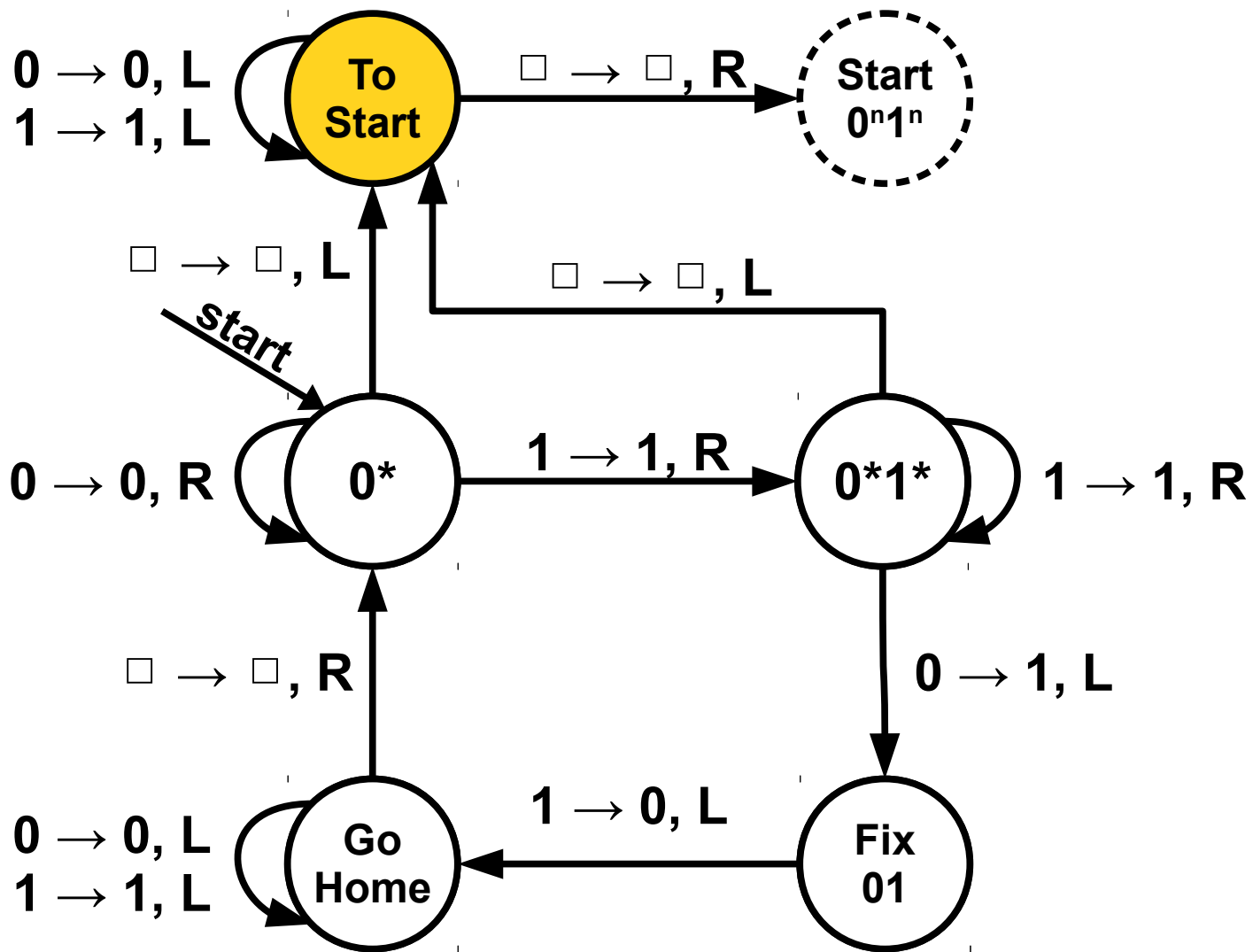


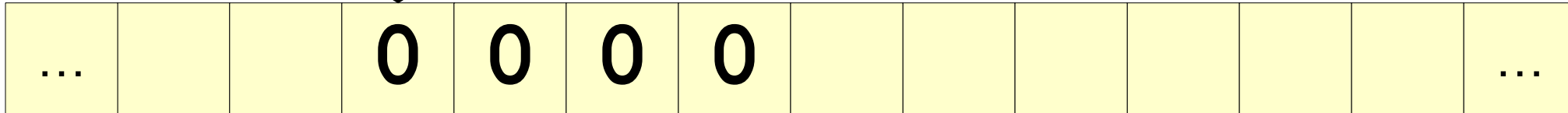
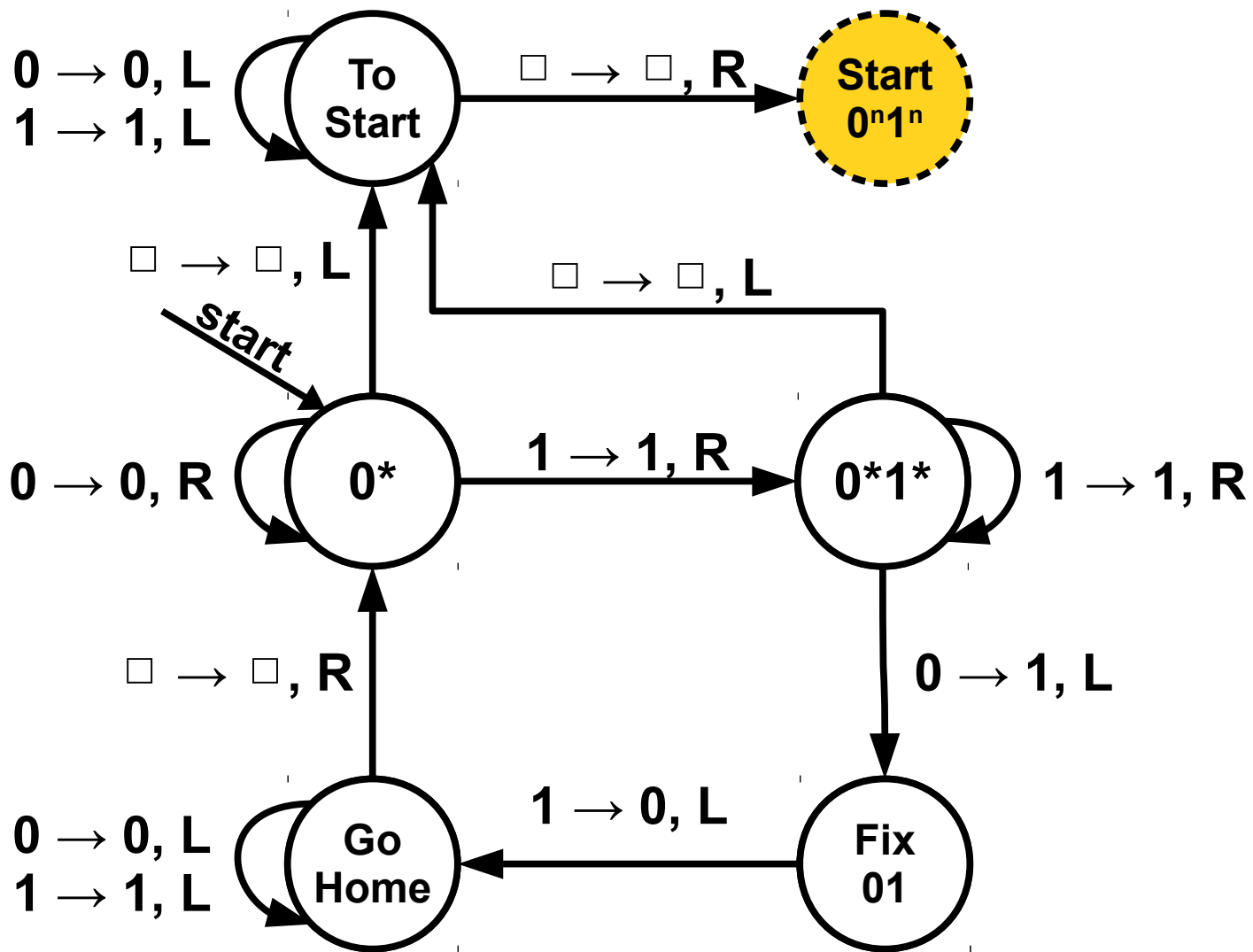


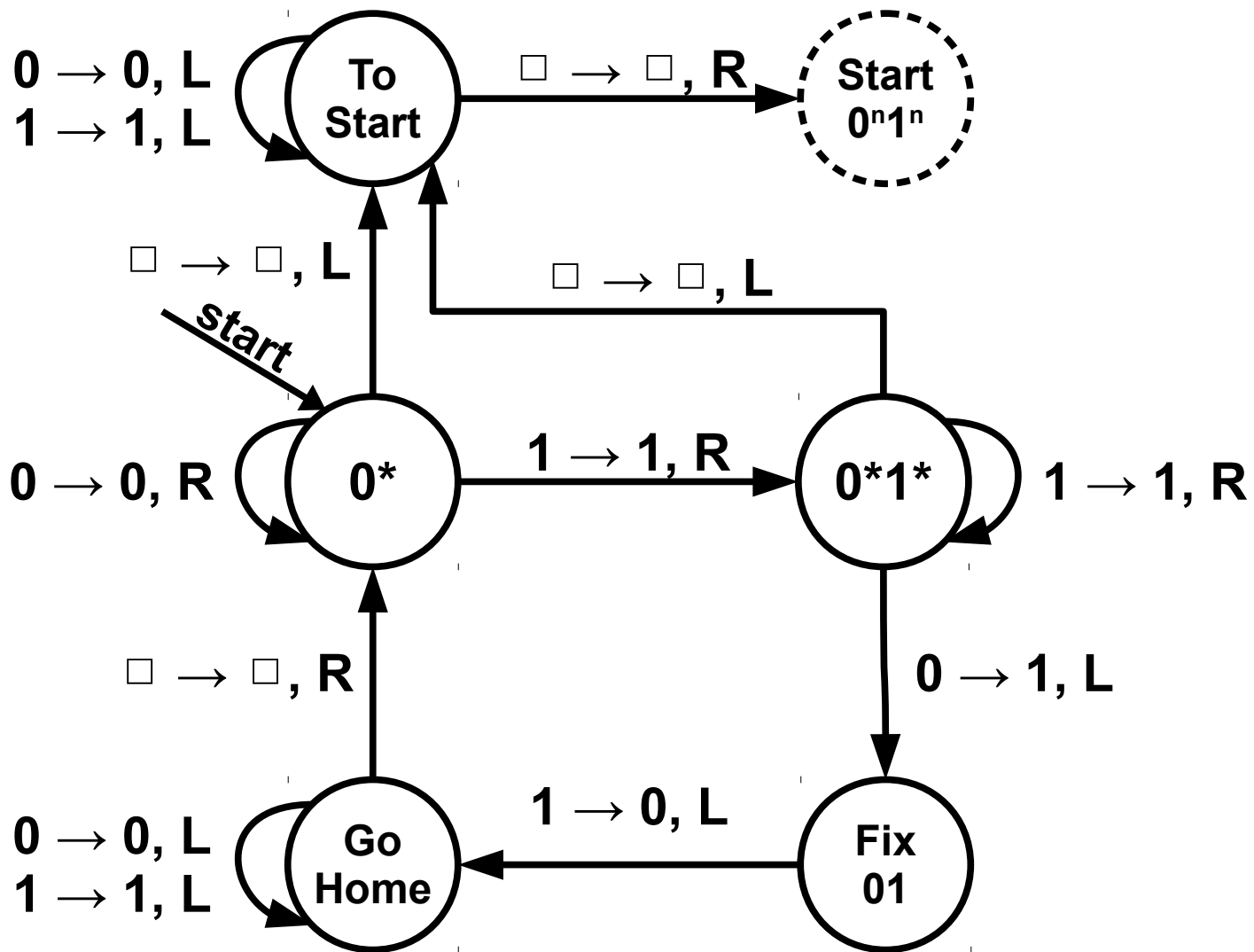


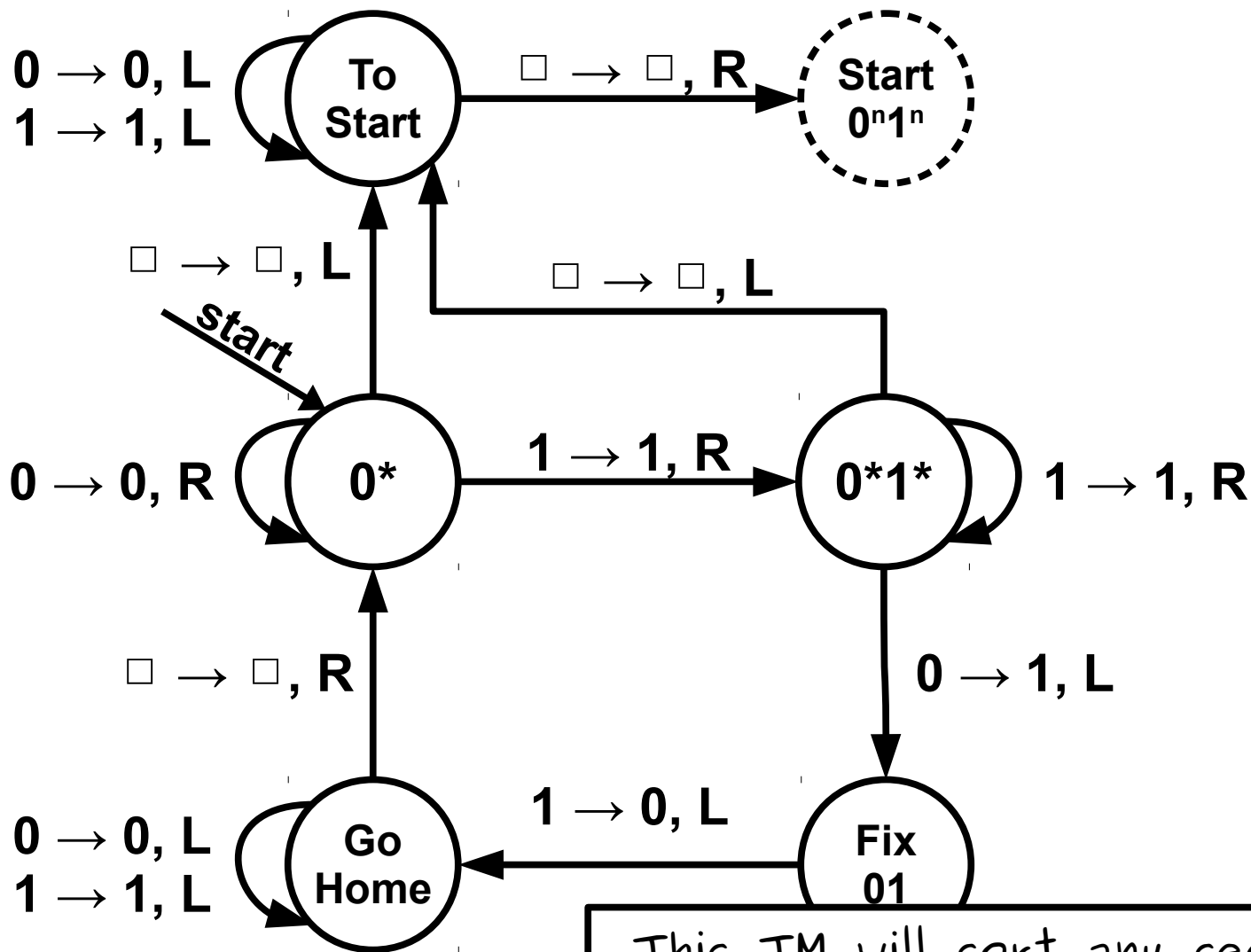












This TM will sort any sequence of 0s and 1s, but it might take a while.

Fun problem: design a TM that sorts a string of 0s and 1s, but does so while taking way fewer steps than this machine.

TM Subroutines

- A ***TM subroutine*** is a Turing machine that, instead of accepting or rejecting an input, does some sort of processing job.
- TM subroutines let us compose larger TMs out of smaller TMs, just as you'd write a larger program using lots of smaller helper functions.
- Here, we saw a TM subroutine that sorts a sequence of 0s and 1s into ascending order.

TM Subroutines

- Typically, when a subroutine is done running, you have it enter a state marked “done” with a dashed line around it.
- When we're composing multiple subroutines together – which we'll do in a bit – the idea is that we'll snap in some real state for the “done” state.

What other subroutines can we make?

Next Time

- ***Arithmetic Subroutines***
 - What else can we do with TMs?
- ***The Church-Turing Thesis***
 - A scientific hypothesis about computation.
- ***R and RE Languages***
 - What does it mean to solve a problem?