Mathematical Logic Part One

Question: How do we formalize the definitions and reasoning we use in our proofs?

Where We're Going

- Propositional Logic (Today)
 - Basic logical connectives.
 - Truth tables.
 - Logical equivalences.
- First-Order Logic (Wednesday/Friday)
 - Reasoning about properties of multiple objects.

Propositional Logic

A *proposition* is a statement that is, by itself, either true or false.

Propositional Logic

- Propositional logic is a mathematical system for reasoning about propositions and how they relate to one another.
- Every statement in propositional logic consists of propositional variables combined via propositional connectives.
 - Each variable represents some proposition, so each variable has value true or false.
 - Connectives encode how propositions are related.

Propositional Connectives

• Logical NOT: $\neg p$

- $\neg p$ is true if and only if p is false.
- Also called *logical negation*.

Logical AND: p ∧ q

- $p \land q$ is true if and only if both p and q are true.
- Also called *logical conjunction*.

Logical OR: p v q

- p v q is true if and only if at least one of p or q are true (inclusive OR)
- Also called *logical disjunction*.

Truth Tables

- A *truth table* is a table showing the truth value of a propositional logic formula as a function of its inputs.
- Useful for several reasons:
 - They give a formal definition of what a connective "means."
 - They give us a mechanical way to evaluate a complex propositional formula.

Truth Table for XOR

- Recall that our OR connective is inclusive.
- The truth table at right defines an exclusive or called XOR.
- We also could have expressed XOR using just the connectives we already had.

p	q	p XOR q
F	F	F
F	Т	Т
Т	F	Т
Т	Т	F

Which expresses XOR?

$$(A)$$
 $(p \land q) \lor (p \lor q)$

(B)
$$(p \land q) \lor \neg (p \lor q)$$

(C)
$$(p \lor q) \land \neg (p \land q)$$

(D)
$$(p \land q) \land (p \lor q)$$

Answer at **PollEv.com/cs103** or text **CS103** to **22333** once to join, then **A**, **B**, or **C**.

Mathematical Implication

Truth Table for $p \rightarrow q$ (implies)

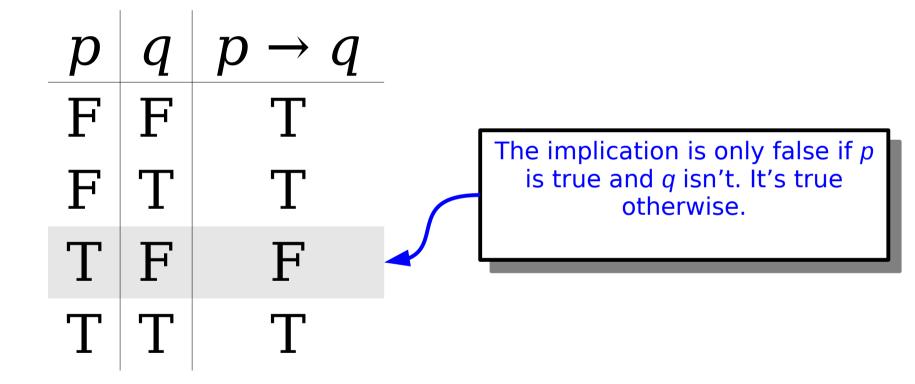
What is the correct truth table for implication? Enter your guess as a list of four values to fill in the rightmost column of the table. (ex: F, T, ?, F)

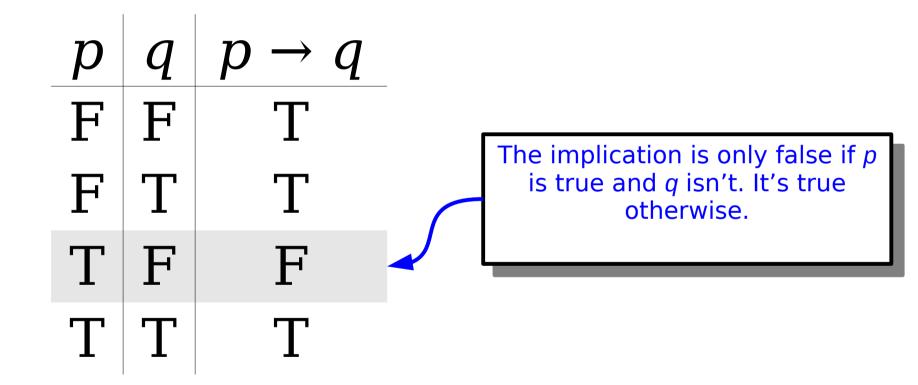
p	q	$p \rightarrow q$
F	F	
F	Т	
Т	F	
Т	Т	

Answer at **PollEv.com/cs103** or text **CS103** to **22333** once to join, then your response.

p	q	$p \rightarrow q$
F	F	T
F	T	T
T	F	F
T	T	T

p	q	$p \rightarrow q$	
F	F	T • 1	Bad bracket, didn't get A
F	T	T • 1	Bad bracket, got A
T	F	F • 1	Perfect bracket, didn't get A
T	T	T • 1	Perfect bracket, got A





You will need to commit this table to memory. (Consider a tattoo on your forearm.)
We're going to be using it *a lot* over the rest of the week.

Why This Truth Table?

- The truth values of the → are the way they are because they're *defined* that way.
- Are there other ways we could write a proposition that has the same truth table as →, using the other connectives? (like what we did for XOR)
 - Yep!
 - Try to think of some
 - What's the truth table for $\neg(p \land \neg q)$?

The Biconditional Connective

The Biconditional Connective

- The biconditional connective

 is used to represent a two-directional implication.
- Specifically, $p \leftrightarrow q$ means both that $p \rightarrow q$ and that $q \rightarrow p$.
- Based on that, what should its truth table look like?
- Take a guess, and talk it over with your neighbor!

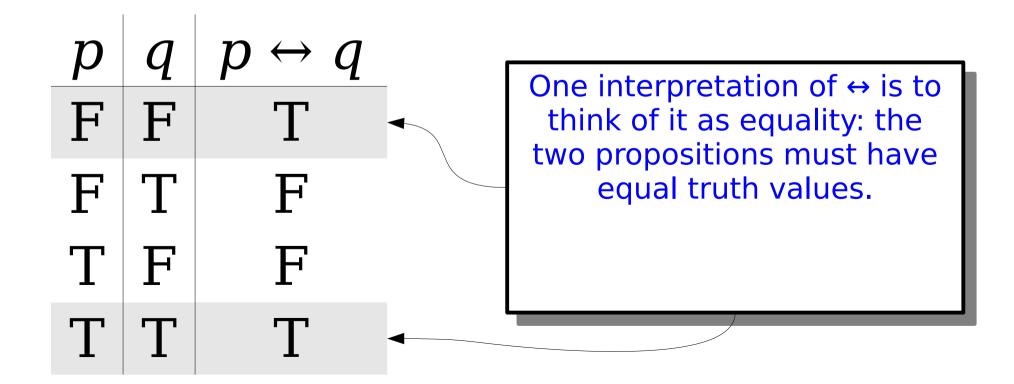
Biconditionals

- The **biconditional** connective $p \leftrightarrow q$ is read "p if and only if q."
- Here's its truth table:

p	\boldsymbol{q}	$p \leftrightarrow q$
F	F	\mathbf{T}
F	T	F
T	F	F
T	Т	T

Biconditionals

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True and False

- In addition to variables and connectives, we have constants: true and false.
 - The symbol \top is a value that is always true.
 - The symbol \bot is value that is always false.

- These are often called connectives, though they don't connect anything.
 - (Or rather, they connect zero things.)

Proof by Contradiction

- Suppose you want to prove *p* is true using a proof by contradiction.
- The setup looks like this:
 - Assume *p* is false.
 - Derive something that we know is false.
 - Conclude that p is true.
- In propositional logic:

$$(\neg p \rightarrow \bot) \rightarrow p$$

How do we parse this statement?

$$\neg x \rightarrow y \lor z \rightarrow x \lor y \land z$$

Operator precedence for propositional logic:

- All operators are right-associative.
- We can use parentheses to disambiguate.

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∧ ∨ → ↔

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$$(\neg x) \rightarrow y \lor z \rightarrow x \lor (y \land z)$$

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How do we parse this statement?

$$(\neg x) \rightarrow y \lor z \rightarrow x \lor (y \land z)$$

Operator precedence for propositional logic:

∧
 ∨
 →
 ↔

- All operators are right-associative.
- We can use parentheses to disambiguate.

How do we parse this statement?

$$(\neg x) \to (y \lor z) \to (x \lor (y \land z))$$

Operator precedence for propositional logic:

- All operators are right-associative.
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How do we parse this statement?

$$(\neg x) \rightarrow (y \lor z) \rightarrow (x \lor (y \land z))$$

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Operator precedence for propositional logic:

- All operators are right-associative.
- We can use parentheses to disambiguate.

- The main points to remember:
 - ¬ binds to whatever immediately follows it.
 - Λ and V bind more tightly than \rightarrow .
- We will commonly write expressions like $p \land q$ $\rightarrow r$ without adding parentheses.
- For more complex expressions,
 - you should add parens (the TAs thank you!)
 - we'll try to add parentheses (if I ever don't and you're confused, ok to ask!)

The Big Table

Connective	Read As	C++ Version	Fancy Name
_	"not"	!	Negation
٨	"and"	&&	Conjunction
V	"or"		Disjunction
\rightarrow	"implies"	see PS2!	Implication
\leftrightarrow	"if and only if"	see PS2!	Biconditional
Т	"true"	true	Truth
Т	"false"	false	Falsity

Recap So Far

- A *propositional variable* is a variable that is either true or false.
- The *propositional connectives* are
 - Negation: $\neg p$
 - Conjunction: $p \land q$
 - Disjunction: p v q
 - Implication: $p \rightarrow q$
 - Biconditional: $p \leftrightarrow q$
 - True: T
 - False: ⊥

Translating into Propositional Logic

a: I will be in the path of totality.

b: I will see a total solar eclipse.

a: I will be in the path of totality.

b: I will see a total solar eclipse.

"I won't see a total solar eclipse if I'm not in the path of totality."

a: I will be in the path of totality.

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"I won't see a total solar eclipse if I'm not in the path of totality."

$$\neg a \rightarrow \neg b$$

translates to

$$q \rightarrow p$$

It does *not* translate to

$$p \rightarrow q$$

"p, but q"

translates to

 $p \land q$

a: I will be in the path of totality.

b: I will see a total solar eclipse.

c: There is a total solar eclipse today.

a: I will be in the path of totality.

b: I will see a total solar eclipse.

c: There is a total solar eclipse today.

"If I will be in the path of totality, but there's no solar eclipse today, I won't see a total solar eclipse." Which is equivalent to the sentence at left?

(A)
$$a \wedge \neg c \wedge \neg b$$

(B)
$$(a \land \neg c) \rightarrow \neg b$$

(C)
$$a \rightarrow (\neg c \land \neg b)$$

(D)
$$(\neg a \lor c) \lor \neg b$$

a: I will be in the path of totality.

b: I will see a total solar eclipse.

c: There is a total solar eclipse today.

"If I will be in the path of totality, but there's no solar eclipse today, I won't see a total solar eclipse."

$$a \wedge \neg c \rightarrow \neg b$$

Propositional Equivalences

Quick Question:

What would I have to show you to convince you that the statement $p \land q$ is false?

Quick Question:

What would I have to show you to convince you that the statement $p \lor q$ is false?

De Morgan's Laws

Using truth tables, we concluded that

$$\neg(p \land q)$$

is equivalent to

$$\neg p \lor \neg q$$

We also saw that

$$\neg (p \lor q)$$

is equivalent to

$$\neg p \land \neg q$$

These two equivalences are called *De Morgan's Laws*.

De Morgan's Laws in Code

• **Pro tip:** Don't write this:

```
if (!(p() && q()) {
    /* ... */
}
```

• Write this instead:

```
if (!p() || !q()) {
    /* ... */
}
```

(This even short-circuits correctly!)

Logical Equivalence

- Because $\neg(p \land q)$ and $\neg p \lor \neg q$ have the same truth tables, we say that they're **equivalent** to one another.
- We denote this by writing

$$\neg (p \land q) \equiv \neg p \lor \neg q$$

- The \equiv symbol is not a connective.
 - The statement $\neg(p \land q) \leftrightarrow (\neg p \lor \neg q)$ is a propositional formula. If you plug in different values of p and q, it will evaluate to a truth value. It just happens to evaluate to true every time.
 - The statement $\neg(p \land q) \equiv \neg p \lor \neg q$ means "these two formulas have exactly the same truth table."
- In other words, the notation $\phi \equiv \psi$ means " ϕ and ψ always have the same truth values, regardless of how the variables are assigned."

An Important Equivalence

• Earlier, we talked about the truth table for $p \rightarrow q$. We chose it so that

$$p \rightarrow q \equiv \neg (p \land \neg q)$$

• Later on, this equivalence will be incredibly useful:

$$\neg (p \to q) \equiv p \land \neg q$$

Another Important Equivalence

Here's a useful equivalence. Start with

$$p \to q \equiv \neg (p \land \neg q)$$

• By De Morgan's laws:

$$p \rightarrow q \equiv \neg (p \land \neg q)$$

$$\equiv \neg p \lor \neg \neg q$$

$$\equiv \neg p \lor q$$

• Thus $p \rightarrow q \equiv \neg p \lor q$

Another Important Equivalence

Here's a useful equivalence. Start with

$$p \to q \equiv \neg (p \land \neg q)$$

• By De Morgan's laws:

$$p \rightarrow q \equiv \neg (p \land \neg q)$$

$$\equiv \neg p \lor \neg \neg q$$

$$\equiv \neg p \lor \neg q = \neg p \lor q \text{ is true. If } p \text{ is true, then } q \text{ has to be true for the whole expression to be true.}$$
• Thus $p \rightarrow q \equiv \neg p \lor q$

One Last Equivalence

The Contrapositive

The contrapositive of the statement

$$p \rightarrow q$$

is the statement

$$\neg q \rightarrow \neg p$$

• These are logically equivalent, which is why proof by contrapositive works:

$$p \rightarrow q \equiv \neg q \rightarrow \neg p$$

 Suppose we want to prove the following statement:

"If x + y = 16, then $x \ge 8$ or $y \ge 8$ "

"If
$$x + y = 16$$
, then $x \ge 8$ or $y \ge 8$ "

$$x + y = 16 \rightarrow x \ge 8 \ \forall \ y \ge 8$$

"If
$$x + y = 16$$
, then $x \ge 8$ or $y \ge 8$ "

$$x + y = 16 \rightarrow x \ge 8 \ \text{v} \ y \ge 8$$

"If
$$x + y = 16$$
, then $x \ge 8$ or $y \ge 8$ "

$$\neg(x \ge 8 \ \lor \ y \ge 8) \to \neg(x + y = 16)$$

"If
$$x + y = 16$$
, then $x \ge 8$ or $y \ge 8$ "

$$\neg(x \ge 8 \ \lor \ y \ge 8) \to \neg(x + y = 16)$$

"If
$$x + y = 16$$
, then $x \ge 8$ or $y \ge 8$ "

$$\neg(x \ge 8 \ \lor \ y \ge 8) \to \neg(x + y = 16)$$

"If
$$x + y = 16$$
, then $x \ge 8$ or $y \ge 8$ "

$$\neg(x \ge 8 \ \lor \ y \ge 8) \rightarrow x + y \ne 16$$

"If
$$x + y = 16$$
, then $x \ge 8$ or $y \ge 8$ "

$$\neg(x \ge 8 \ \lor \ y \ge 8) \rightarrow x + y \ne 16$$

"If
$$x + y = 16$$
, then $x \ge 8$ or $y \ge 8$ "

$$\neg(x \ge 8 \lor y \ge 8) \rightarrow x + y \ne 16$$

"If
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, then $x \ge 8$ or $y \ge 8$ "

$$\neg(x \ge 8) \land \neg(y \ge 8) \rightarrow x + y \ne 16$$

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"If
$$x + y = 16$$
, then $x \ge 8$ or $y \ge 8$ "

$$\neg (x \ge 8) \land \neg (y \ge 8) \to x + y \ne 16$$

"If
$$x + y = 16$$
, then $x \ge 8$ or $y \ge 8$ "

$$x < 8 \land \neg (y \ge 8) \to x + y \ne 16$$

"If
$$x + y = 16$$
, then $x \ge 8$ or $y \ge 8$ "

$$x < 8 \land \neg (y \ge 8) \to x + y \ne 16$$

"If
$$x + y = 16$$
, then $x \ge 8$ or $y \ge 8$ "

$$x < 8 \land \neg (y \ge 8) \to x + y \ne 16$$

"If
$$x + y = 16$$
, then $x \ge 8$ or $y \ge 8$ "

$$x < 8 \land y < 8 \rightarrow x + y \neq 16$$

"If
$$x + y = 16$$
, then $x \ge 8$ or $y \ge 8$ "

$$x < 8 \land y < 8 \rightarrow x + y \neq 16$$

 Suppose we want to prove the following statement:

"If
$$x + y = 16$$
, then $x \ge 8$ or $y \ge 8$ "

$$x < 8 \land y < 8 \rightarrow x + y \neq 16$$

"If x < 8 and y < 8, then $x + y \ne 16$ "

Theorem: If x + y = 16, then $x \ge 8$ or $y \ge 8$.

Proof: By contrapositive. We will prove that if x < 8 and y < 8, then $x + y \ne 16$. Let x and y be arbitrary numbers such that x < 8 and y < 8.

Note that

$$x + y < 8 + y$$

 $< 8 + 8$
 $= 16.$

This means that x + y < 16, so $x + y \ne 16$, which is what we needed to show.

Why This Matters

- Propositional logic is a tool for reasoning about how various statements affect one another.
- To better understand how to prove a result, it often helps to translate what you're trying to prove into propositional logic first.
- That said, propositional logic isn't expressive enough to capture all statements. For that, we need something more powerful.

Next Time

- First-Order Logic
 - Reasoning about groups of objects.
- First-Order Translations
 - Expressing yourself in symbolic math!