# CPUNAME Instruction Set Architecture (ISA) Specification

## luccie-cmd

# March 18, 2025

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### 1 General overview

# 2 Registers

Register	Purpose	
R1-R32	General purpose registers	
SP	Pointer to where the stack head currently is	
IP	Pointer to where the CPU will execute instructions	
PTP	Pointer to where the MMU will translate addresses	
FLAGS	Status flags used for conditional jumps and moves	

Table 1: Register overview

See table above for register descriptions. Both SP and IP are virtual addresses meaning that they should be

#### **2.1 NOTES**

The flags register affects conditional execution such as jumps and moves and can only be effected by the ALU

# 3 Instruction encoding

### 3.1 General Instruction Format

Unlike RISC, the CPUNAME will have variable length instructions, meaning that every instruction only takes up the binary space it needs.

A typical instruction will generally look like so:

Field	Size (In bytes)	Description
Prefix	1	The instruction prefix, will change what ad-
		dressing mode(s) to use. Described in Instruc-
		tion Set subsection Prefixes
Opcode	1	The main instruction identifier. Contains
		number of sources and if the destination or
		source is a memory or register or immediate
Operands(s)	1	A register to use in the instruction
(R8/16/32/64/128)		
Operands(s)	1, 2, 4, 8 or 16	A memory location to use in the instruction.
(M8/16/32/64/128)		Number of bytes pulled from memory will de-
		pend on the instruction's prefix otherwise de-
		fault 128 bit addressing is assumed
Operands(s)	1, 2, 4, 8 or 16	An immediate value to use in the instruction.
(18/16/32/64/128)		Number of bytes will depend on the instruc-
		tion's prefix otherwise default 128 bit address-
		ing is assumed

Table 2: Instruction overview

Prefixes and operands are optional, some instructions (Described in the instruction set table where the operands are None) don't use operands. An instruction that has a prefix byte but doesn't take any operands will trigger an Invalid Prefix Opcode interrupt (#IPO). The number of operands for a instruction if specified in the table in Instruction Set, please refer to that for specific instruction encoding and arguments.

# 4 Instruction set

## 4.1 Instruction set as mnemonics

## 4.1.1 Moves

Bytes	Mnemonic	Operand 1	Operand 2	Description
00	MOVE	R8/16/32/64/128	R8/16/32/64/128	Move value at source regis-
				ter (Operand 2) to destina-
				tion register (Operand 1)
01	MOVE	R8/16/32/64/128	M8/16/32/64/128	Move value located at the
				memory address in (Operand
				2) to destination register
				(Operand 1)
02	MOVE	R8/16/32/64/128	R8/16/32/64/128+	Move value at the address
			I8/16/32/64/128	in the register of Operand 2,
				plus the immediate value of
				Operand 2 into the destina-
				tion register (Operand 1)
03	MOVE	R8/16/32/64/128	R8/16/32/64/128+	Move value at the address
			R8/16/32/64/128	in the register of Operand 2,
				plus the value of the register
				of Operand 2 into the desti-
				nation register (Operand 1)
04	MOVE	M8/16/32/64/128	R8/16/32/64/128	Move the value of register
				Operand 2 into the memory
				location specified in Operand
				1

## 4.1.2 Stack Manipulation

Bytes	Mnemonic	Operand 1	Operand 2	Description
05	PUSH	R8/16/32/64/128	None	Push value in operand 1
				on the stack and decrement
				stack pointer by N bytes
				(specified using prefix)
06	POP	R8/16/32/64/128	None	Pop value at current stack
				pointer into Operand 1 and
				increment stack pointer by N
				bytes (specified using prefix)

### 4.1.3 Control flow

Bytes	Mnemonic	Operand 1	Operand 2	Description
07	JIE	I128	None	Perform a conditional jump
				to the immediate value of
				Operand 1 if zero flag is set
08	JIL	I128	None	Perform a jump to the imme-
				diate value of Operand 1 if SF
				!= OF
09	JMP	I128	None	Perform an unconditional
				jump to the immediate value
				of Operand 1
0A	CALL	I128	None	Perform a call to the immedi-
				ate value of Operand 1, push
				the value of the return ad-
				dress (Current IP+17) first
0B	RET	None	None	Perform a return from a pro-
				cedure by popping the return
				value from the stack
0C	U2K	None	None	Performs a call to the ele-
				vated program to handle priv-
				ileged instructions, changes
				the privilege level

### 4.1.4 Arithmetic

Bytes	Mnemonic	Operand 1	Operand 2	Description
0D	TEST	R8/16/32/64/128	R8/16/32/64/128	Perform a bitwise AND but
				don't store the result only up-
				date the FLAGS
0E	XOR	R8/16/32/64/128	R8/16/32/64/128	Perform a bitwise XOR On
				the operands and store the re-
				sult in the first operand
0F	SHL	R8/16/32/64/128	I8/16/32/64/128	Perform a shift to the left by
				the number of times specified
				in Operand 2 on Operand 1

#### 4.1.5 Miscellaneous

Bytes	Mnemonic	Operand 1	Operand 2	Description
10	COMP	R8/16/32/64/128	I8/16/32/64/128	Compare the value in register
				Operand 1 against the imme-
				diate value and set flags

### 4.1.6 Prefixes

Certain prefixes can be placed before an instruction to modify its behavior. These prefixes can affect the selection of registers, immediate values, memory

addressing modes, or other instruction properties. By default, without any prefixes, all operands are 128 bits. The table below defines the available prefixes and their effects.

Bytes	Effect		
11	Changes instruction to operate		
	using 8 bit values for all operands		
12	Changes instruction to oper-		
	ate using 16 bit values for all		
	operands		
13	Changes instruction to oper-		
	ate using 32 bit values for all		
	operands		
14	Changes instruction to oper-		
	ate using 64 bit values for all		
	operands		

## 4.2 Microinstructions

### 4.2.1 Loading

Mnemonic	Arguments	Description
LDR	Register index	Load contents of the register specified in the register index onto the bus
LDI	Value	Load the provided value onto the bus

### 4.2.2 Writing

Mnemonic	Arguments	Description
STR	Register index	Store the contents of the bus into
		the register specified in register
		index

### 4.2.3 ALU

Mnemonic	Arguments	Description	
PAA	None	Prepare ALU for an addition	
LALU	Operand	Load the Nth ALU operand	
		(with N being specified in the ar-	
		gument) from the bus	
AE	None	Allow the ALU to output its re-	
		sult onto the bus	

#### 4.2.4 MMU

Mnemonic	Arguments	Description
MR	None	Load value at the address spec-
		ified in the EADDR register to
		the bus
MW	None	Write the value of the current
		bus to the address specified in
		the EADDR register

### 5 Privilige levels

The CPUNAME and versions thereof have 2 privilege levels named unprivileged and privileged. These privilege levels can be compared to X86's ring 0 and ring 3 respectively. A program running in privileged mode can execute any and all instructions and is allowed to change things like the PTP register, SP and IP via direct assignment. However an unprivileged program running is not allowed to access some instructions and registers that a privileged program can access. For example the PTP register inaccessible for a unprivileged program and will throw an the Unpriviliged Execute interrupt (#UPE). A full description of instructions and registers the unprivileged level isn't allowed to access will be provided in a list below.

#### 5.1 Instructions

Currently there are no instructions that can only be run on a privileged level.

#### 5.2 Registers

- PTP
- SP
- IP
- FLAGS
- EADDR
- TEMP

# 6 Memory

The CPUNAME can support up to 13 levels of paging, each taking up 9 bits of the virtual address with the exception of the last level which points to the physical address, this takes up 12 bits giving a total of 117 bits used for virtual addressing. The remaining 11 bits should all be the same, if for any reason these bits aren't the same when it is needed (Wether for reading or writing) it will trigger a Non Canonical interrupt (#NCC). The 12 bits are the default value and are extended by 9 bits for every paging level activated. The maximum

paging level useable is 13 by default but can be changed by setting a specific value in the paging control register (CRP) when the maximum paging level of 0 is applied, the MMU treats this as not having any paging enabled. For every new level of paging addded, 9 bits of virtual address space get added which would equate to a function like:  $2^{N\times 9}$  where N is the number of paging levels enabled.

The stack and instruction pointer (SP and IP respectively) are to be mapped in the PTP (If paging is enabled).

PTP describes a physical address stored somewhere in the RAM, the PTP is used for translating virtual addresses to physical ones via a table described below.

Bits	Name	Description
1	RE	Describes if a page can be read from, trying to
		read from a non readable page will trigger an Un
		Readable Page interrupt (#URP).
1	PR	Describes if a page is present, trying to use a non
		present page will trigger a Non Present Page in-
		terrupt (#NPP).
1	WR	Describes if a page can be written to, trying to
		write to a non writeable page will trigger an Un
		Writeable Page interrupt (#UWP)
1	PV	Describes if a page is occupied by the privileged
		program, trying to use a privileged page as an
		unprivileged program will trigger an Un Privilige
		Page interrupt (#UPP)
1	NX	Describes if a page can *NOT* be executed, tying
		to execute from a non executable page results in
		a Non Executable Page interrupt (#NEP)
40	NL	This describes the address of the next level of
		paging table, if this is the last level it contains
		the physical page number. Must be non null oth-
		erwise it will throw a Null Next Level interrupt
		(#NNL)
19	RESV	These bits are reserved for future purposes and
		should be zero at all times. Attempting to set
		these bits to any other value will result in a Non
		Zero Reserved interrupt (#NZR)

SP and IP are both virtual addresses and need to be mapped somewhere in the PTP

PTP is a physical address to somewhere in RAM and is used to The CPU fetches instructions using the virtual address specified in IP after translating it

### **7** MMU

This section describes how the CPUNAME's MMU (Memory Management Unit) works in detail

## 8 ALU

This section describes how the CPUNAME's ALU (Arithmetic Logic Unit) works in detail

# 9 I/O

The CPUNAME only support MMIO as the PIO isn't used anymore

- 9.1 MMIO
- 10 Interrupts
- 11 Pipeline