MICRO II PS#2 (Q2&Q3)

2. Consider the linear program

$$v_P(b) = \max c_x x + c_y y$$

subject to $x \le 1$
 $x \ge 0$

- (a) Rewrite this in the canonical form.
- (b) Rewrite this in the standard form.
- (a) Sing y doesn't have the nonnegativity constraint, we can express y as tradifference of two numeroative variables.

$$V_{P(b)} = MAX \quad (C_{X} \quad C_{Y} - C_{Y}) \begin{bmatrix} X \\ Y_{1} \\ Y_{2} \end{bmatrix}$$

$$f.T. \quad (1 \quad 0 \quad 0) \begin{bmatrix} X \\ Y_{1} \\ Y_{2} \end{bmatrix} \leq 1$$

$$\begin{bmatrix} X \\ Y_{1} \\ Y_{2} \end{bmatrix} \geqslant 0$$

(10)
$$V_{PLD}) = MAX Cxx + Cyy_1 - Cyy_2 + 0.5$$

S.T. $x + 5 = 1$
 $x, y_1, y_2, 5 > 0.$

$$\Rightarrow V_{P(b)} = MAX \quad (CX \quad Cy - Cy \quad O) \begin{pmatrix} x_1 \\ y_2 \\ y_2 \end{pmatrix} = 1$$

$$\begin{pmatrix} x_1 \\ y_2 \\ y_2 \end{pmatrix} = 0$$

3. Consider the primal problem

$$v_p(b) = \max x_1 + 2x_2$$
 s.t. $x_1 + x_2 \le 4$
$$x_1 + 3x_2 \le b$$

- (a) Write down the dual.
- (b) For b = 1, plot the constraint sets for both problems, and solve them.
- (c) Describe $v_P(b)$, and compute $\partial v_P(b)$ on the range $0 \le b \le 14$.

(a) We want to first rewrite the problem with nonnegativity constraints.

$$VP(b) = MAX (X_1-X_2)+2(X_3-X_4) \qquad \underline{Note:} \quad \chi_1, \chi_2 \text{ here are not}$$

$$\int_{1.7} (X_1-X_2)+(X_3-X_4) \leq \psi \qquad \text{the problem}$$

$$(X_1-X_2)+3(X_3-X_4) \leq b$$

$$X_{1,2}X_{2,1}X_{3,1}X_{4,2} \geq 0.$$

$$\Rightarrow Vp(b) = MAX \left(1 - 1 \ 2 - 2 \right) \begin{bmatrix} \chi_1 \\ \chi_2 \\ \chi_3 \\ \chi_4 \end{bmatrix}$$

$$\begin{cases} 1 - 1 \ 1 - 1 \\ 1 - 1 \ 3 - 3 \end{cases} \begin{bmatrix} \chi_1 \\ \chi_2 \\ \chi_3 \\ \chi_4 \end{bmatrix} \leq \begin{bmatrix} 4 \\ b \end{bmatrix}$$

$$A \left(\begin{array}{c} \chi_1 \\ \chi_2 \\ \chi_3 \\ \chi_4 \end{array} \right) \geq 0$$

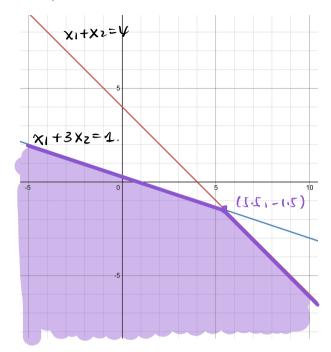
Henu, the anal problem is

$$V_{D(c)} = MIN \left(y_{1} \quad y_{2} \right] \left(\begin{array}{c} 4 \\ b \end{array} \right)$$

$$S.T. \quad \left(y_{1} \quad y_{2} \right] \left(\begin{array}{c} 1 & -1 & 1 & -1 \\ 1 & -1 & 3 & -3 \end{array} \right) \geqslant \left(1 & -1 & 2 & -2 \right)$$

$$\left[y_{1} \quad y_{2} \right] \geqslant 0.$$

(b) First consider the primal problem



$$v_p(b) = \max x_1 + 2x_2$$

s.t. $x_1 + x_2 \le 4$
 $x_1 + 3x_2 \le 4$

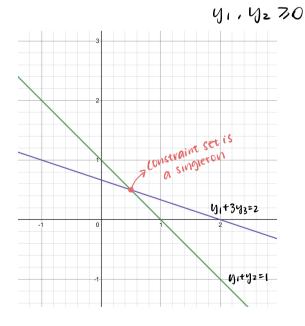
The shaded region is the lonstraint set and by vertex theorem, the solution is $(5.5, -1.5)_{2}$ which gives Vp(b) = 2.5

Next consider the dual problem:

MIN
$$491+92$$

SIT. $91+92>1$
 $-91-92>-1$ $\Rightarrow 91+92=1$

$$y_1 + 3y_2 \ge Z$$
 => $y_1 + 3y_3 = 2$ - $y_1 - 3y_2 \ge -2$



Since the constraint set is a (ingleton, the solution to the dual problem is (0.5,0.5), which gives Vo(b)= 2.5.

$$\Rightarrow V_D(c) = V_P(b)$$

(1)

Shadow Prices

Theorem. If either $v_P(b)$ or $v_D(c)$ is finite, then

- 1. $v_P(b) = v_D(c)$,
- 2. both programs have optimal solutions, and
- 3. $\partial v_D(b)$ is the set of optimal solutions to the primal, and $\partial v_P(b)$ is the set of optimal solutions to the dual.

By computation above, we know that $V_P(b)$ and $V_O(c)$ are both finite when be(0,18). and $V_P(b) = V_P(c)$. The solution for the clual publishm is $\{0.5, 0.5\}$. Hence, $V_P(b) = V_P(c) = 410.5$) + 0.5b = 2+0.5b.

By the shadow price theorem, $\partial V_P(4) = (0.5, 0.5)$ for be [0.14]

You could also solve the primal for an arbitrary b_1 and then obtain $\begin{bmatrix} \dot{x}_1 \\ \dot{x}_2 \end{bmatrix} = \begin{bmatrix} 12-b \\ 2 \\ b-4 \\ 2 \end{bmatrix}$

and plug it back into $Vp(b) = x_1 + 2x_2 = 2 + \frac{b}{2}$ $\frac{d Vp(b)}{d b} = \frac{1}{2}$