

## Logistics

- TA Office Hours:
  - Mon & Wed 18:30-19:30pm (in person (Uris 465) or by Zoom).
- All problem sets are due before Friday section.
- Section's material [https://github.com/luciacasal/-Cornell\\_ECON6100](https://github.com/luciacasal/-Cornell_ECON6100).

# 1 General Equilibrium and Welfare Theorems<sup>1</sup>

## 1.1 Important concepts:

1. Definition 1: Private-ownership economy. We move from pure exchange economy to a private-ownership economy by introducing a set of producers.
  - (a)  $L$  commodities, aggregate endowments  $\bar{\omega} \in \mathbb{R}^L$ .
  - (b) Firms:  $j = 1, \dots, J$  with nonempty and closed production set  $Y_j \in \mathbb{R}^L$
  - (c) Consumers:  $i = 1, \dots, I$  with consumption set  $X_i \subset \mathbb{R}^L$ , preferences  $\succeq_i$  and endowment  $\omega_i \in X_i$ . Consumer  $i$ 's ownership of firm  $j$  is  $\theta_{ij} \in [0, 1]$  and  $\sum_i \theta_{ij} = 1$ .
2. Definition 2. Walrasian equilibrium for a private-ownership economy: denotes an allocation  $(x^*, y^*)$  and a price vector  $p^*$  such that:
  - (a) (Profit maximization): For all  $j$ ,  $p^* \cdot y_j^* \geq p^* \cdot y_j$  for all  $y_j \in Y_j$
  - (b) (Utility maximization): For all  $i$ ,  $x_i^*$  is  $\succeq_i$ -maximal in the budget set:
 
$$\left\{ x_i \in X_i : p^* \cdot x_i \leq p^* \cdot \omega_i + \sum_i \theta_{ij} p^* \cdot y_j^* \right\}$$
  - (c) (Markets clear):  $\sum_i x_i^* = \bar{\omega} + \sum_j y_j^*$
3. Definition 3: [Walrasian Quasi-equilibrium] A Walrasian quasi-equilibrium for a private-ownership economy denotes an allocation  $(x^*, y^*)$  and a price vector  $p^*$  such that:
  - (a) (Profit Maximization): For all  $j$ ,  $p^* \cdot y_j^* \geq p^* \cdot y_j$  for all  $y_j \in Y_j$ .
  - (b) (“Expenditure Minimization”) <sup>2</sup>: For all  $i$ , we have  $p^* \cdot x_i^* \leq p^* \cdot \omega_i + P_i \theta_{ij} p^* \cdot y_j^*$  and for all  $x'_i \in X_i$  such that  $x'_i \succeq_i x_i^*$  we have  $p^* \cdot p^* \cdot x'_i \geq p^* \cdot \omega_i + P_i \theta_{ij} p^* \cdot y_j^*$ .
  - (c) (Markets Clear):  $\sum_i x_i^* = \bar{\omega} + \sum_j y_j^*$

<sup>1</sup>These notes borrow from Jaden Chen's notes from 2020.

<sup>2</sup>when  $\succeq_i$  is LNS, this condition is equivalent to  $x_i^*$  minimizes expenditure on  $\{x_i : x_i \succeq_i x_i^*\}$

4. Theorem 1: [Existence of Walrasian equilibrium in private-ownership Economy] Suppose:

- (a) For all  $i$ ,  $X_i$  is compact and convex,  $\succeq_i$  is continuous and convex.
- (b) For all  $j$ ,  $Y_j$  is compact and convex,  $0 \in Y_j$  ;

Then there is a Walrasian equilibrium.

5. Theorem 2: [First Welfare Theorem] Suppose  $\succeq_i$ 's are locally nonsatiated for all  $i$  and  $(p^*, x^*, y^*)$  is Walrasian equilibrium with  $p^* \geq 0$ , then  $(x^*, y^*)$  is Pareto optimal.
6. Theorem 3: [Second Welfare Theorem] Suppose every  $Y_j$  is convex and every preference relation  $\succeq_i$  is convex and locally nonsatiated. Then for every Pareto optimal allocation  $(x^*, y^*)$ , there is a way to reallocate endowments and shareholdings, such that  $(p^*, x^*, y^*)$  is Walrasian quasi-equilibrium where  $p^* \neq 0$ .

## 1.2 Questions:

1. Consider a private ownership economy with two individuals, Mr. 1 and Mr. 2 and two goods  $x$  and  $y$ . Consumers cannot consume in negative quantities. Mr. 1 and Mr. 2 have the following utility functions:

$$\begin{aligned} u_1 &= x_1 - \gamma y_2 \\ u_2 &= \frac{1}{2} \log x_2 + \frac{1}{2} \log y_2 \end{aligned}$$

where  $\gamma \in [0, 1)$ . Each consumer has endowment of 1 unit of each good. Let good  $x$  be the numeraire good and denote the price of good  $y$  by  $p$ .

- (a) Find the CE allocation and price for this economy.
  - (b) For what values of  $\gamma$  is the CE Pareto optimal?
  - (c) Does First Welfare Theorem hold here? Why?
2. (Welfare Economics Problem 2)
3. In a two-person, two good exchange economy, individuals A and B have the following preferences:

$$\begin{aligned} u^A(x, y) &= \min\{x, 2y\} \\ u^B(x, y) &= \min\{3x, y\} \end{aligned}$$

where  $x$  and  $y$  are two goods. The aggregate endowment is  $(12, 12)$ . Answer the following using an Edgeworth box:

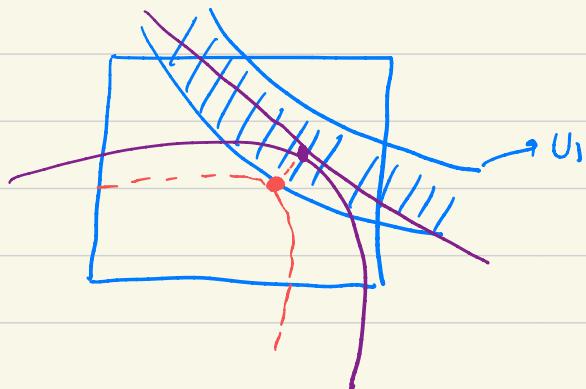
- (a) Find the set of Pareto optima.
- (b) Assume A has all of good  $x$  and B all of good  $y$ . Describe the competitive equilibria.
- (c) Which endowment allocations have competitive equilibria with positive prices for both goods?

- It is better to assume preferences not utility functions. (General Rule).
- It is sometimes more convenient to assume on the utility functions

1st WT: most important thing is  $\succsim$  is LNS.

If not then the CE may not be P.O.

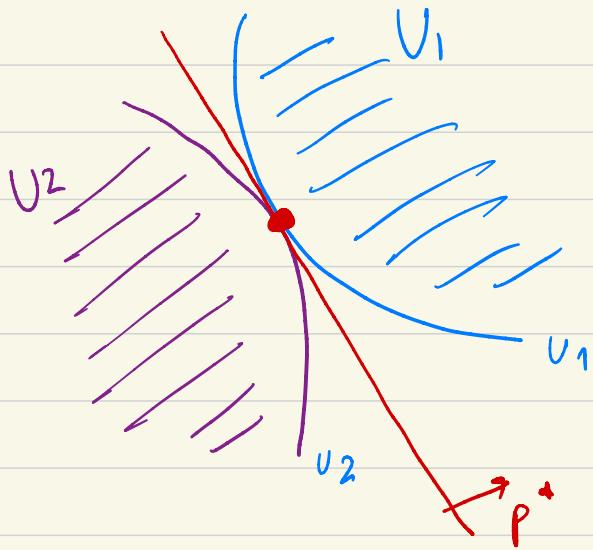
Counterexample:



- Is a CE but not
    - corner 1 is the same
    - corner 2 is better off
- ↓  
Not P.O!

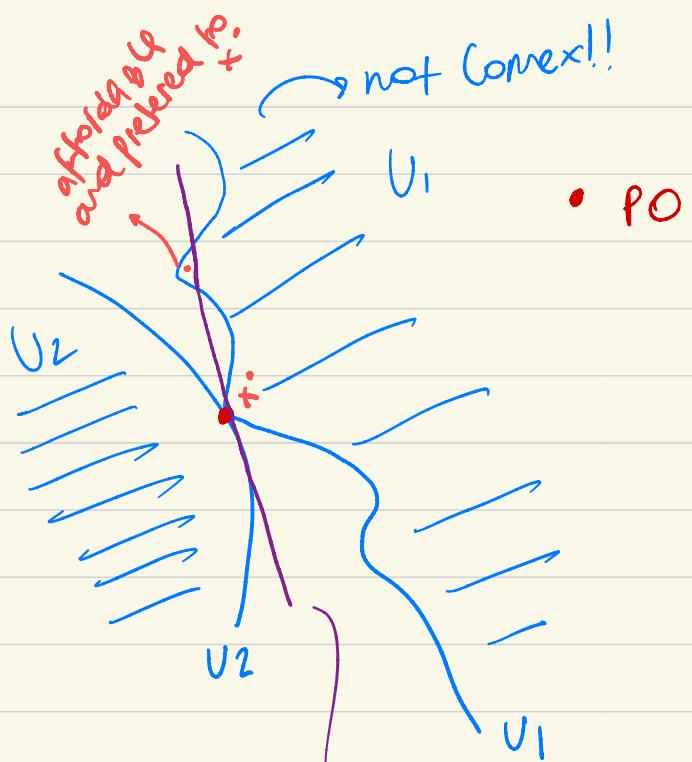
CE is not PO.

2nd WT:  $\succsim$  is LNS. Most important thing is  $\succsim$  is CONVEX because we need to use the Separated Hyperplane theorem!!



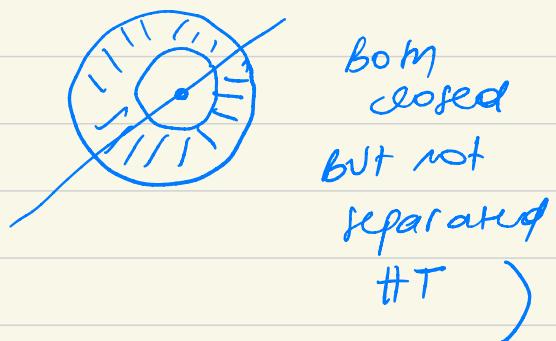
- PO,  
upper contour set of  
person 1 and 2  
are convex  $\rightsquigarrow$   
SHT  
we can find a  
plane that  
separates the  
convex sets

If not convex, it is not able to find supporting planes!



You are going to meet  
either with  $U_1$  or  $U_2$ ,  
so  $PO$  is not a CE!!

(If they are closed it's not enough



## Problems

- sound free exchange ..
- Consider a private ownership economy with two individuals, Mr. 1 and Mr. 2 and two goods  $x$  and  $y$ . Consumers cannot consume in negative quantities. Mr. 1 and Mr. 2 has the following utility functions:

$$u_1 = x_1 - \gamma y_2$$

$$u_2 = \frac{1}{2} \log x_2 + \frac{1}{2} \log y_2$$

where  $\gamma \in [0, 1]$ . Each consumer has endowment of 1 unit of each good. Let good  $x$  be the numeraire good and denote the price of good  $y$  by  $p$ .

- Find the CE allocation and price for this economy.
- For what values of  $\gamma$  is the CE Pareto optimal?
- Does First Welfare Theorem hold here? Why?

$$\left\{ \begin{array}{l} u_1 = x_1 - \gamma y_2 \quad e_1 = (1, 1) \\ \qquad \qquad \qquad e_2 = (1, 1) \\ u_2 = \frac{1}{2} \log x_2 + \frac{1}{2} \log y_2 \end{array} \right. \rightarrow \begin{array}{l} \text{releasing in both } x_2, y_2 \\ \text{if } p_2 = 0 \rightarrow \text{good 1 } \infty \end{array}$$

as long as there is a CE  $\rightarrow$  we can guarantee  $p > 0$

market  
can not  
clear!  
because we  
have a finite  
endowment.

Super  
Important!

- as long as there is 1 consumer who has strong marshale preferences and  $\exists$  CE  $\rightarrow$  eq prices are  $> 0$  !!
- as long as 1 consumer whose preferences are marshale  $\rightarrow$  price vector  $\geq 0$   
 $\exists$  CE but it can not be negative.

$$(a) \begin{cases} \max x_1 - \gamma j_2 & \text{Linear Utility Function!} \\ \text{st } px_1 + j_1 \leq p+1 & [\text{normalize } px=p; pj=1] \\ x_1, j_1 \geq 0 & \text{Walras Law holds in eq} \\ & \text{after we let } j=0. \end{cases}$$

$$\Rightarrow x_1 = \frac{p+1}{p}; j_1 = 0$$

$$\begin{cases} \max \frac{1}{2} \log x_2 + \frac{1}{2} \log j_2 \\ \text{st } px_2 + j_2 \leq p+1 (\lambda) \quad \text{sue (randa)} \\ & \text{and ramperd} \\ & \therefore \text{drop } x_2 j_2 \geq 0 \end{cases}$$

$$\begin{cases} [x_2] \quad \frac{1}{2x_2} = p\lambda \\ [j_2] \quad \frac{1}{2j_2} = \lambda \end{cases} \quad \left| \begin{array}{l} x_2 = \frac{1}{2p\lambda} \\ j_2 = \frac{1}{2\lambda} \end{array} \right. \quad \therefore \text{smiley convex} \quad \therefore \text{SOC} \checkmark$$

$\therefore \text{LNS} \quad \therefore \text{plug in the budget constraint and get}$

$$\frac{1}{2\lambda} + \frac{1}{2\lambda} = p+1 \rightarrow \lambda = \frac{1}{p+1}$$

$$\therefore x_2 = \frac{p+1}{2p} \quad j_2 = \frac{p+1}{2}$$

Mkt clearing

$$\begin{cases} x_1 + x_2 = 2 \\ y_1 + y_2 = 2 \end{cases} \Rightarrow p = 3$$

$$\therefore \begin{cases} x_1^* = 4/3 & y_1^* = 0 \\ x_2^* = 2/3 & y_2^* = 2 \end{cases}$$

(b) Value of  $\gamma$  st CE is Pareto Optimal.

Method 1: use definition of PO.

We want  $(x_1^*, x_2^*, y_1^*, y_2^*)$  to be PO

At  $(x_2^*, y_2^*)$ ,  $v_2 = \frac{1}{2} \log 2/3 + \frac{1}{2} \log 2 = \bar{v}$

$$\begin{cases} \max_{x_1, x_2, y_2} x_1 - \gamma y_2 \\ \text{st } \frac{1}{2} \log x_2 + \frac{1}{2} \log y_2 \geq \bar{v} \\ 0 \leq x_1 + x_2 \leq 2 \\ 0 \leq y_2 \leq 2 \end{cases}$$

allocate consumer of  
 $x, y$  st consumer  
2 gets  $\bar{v}$   
already  
and consumer  
1 gets the  
max possible given  
that

We want to find the range of  $\gamma$  such that

$(x_1^*, x_2^*, j_2^*)$  solves the problem

$$\Rightarrow \gamma \leq 1/3 \quad \therefore \gamma \in [0, 1/3]$$

How to solve? By K-T

Method 2: graphically by Edgeworth Box in the eq.

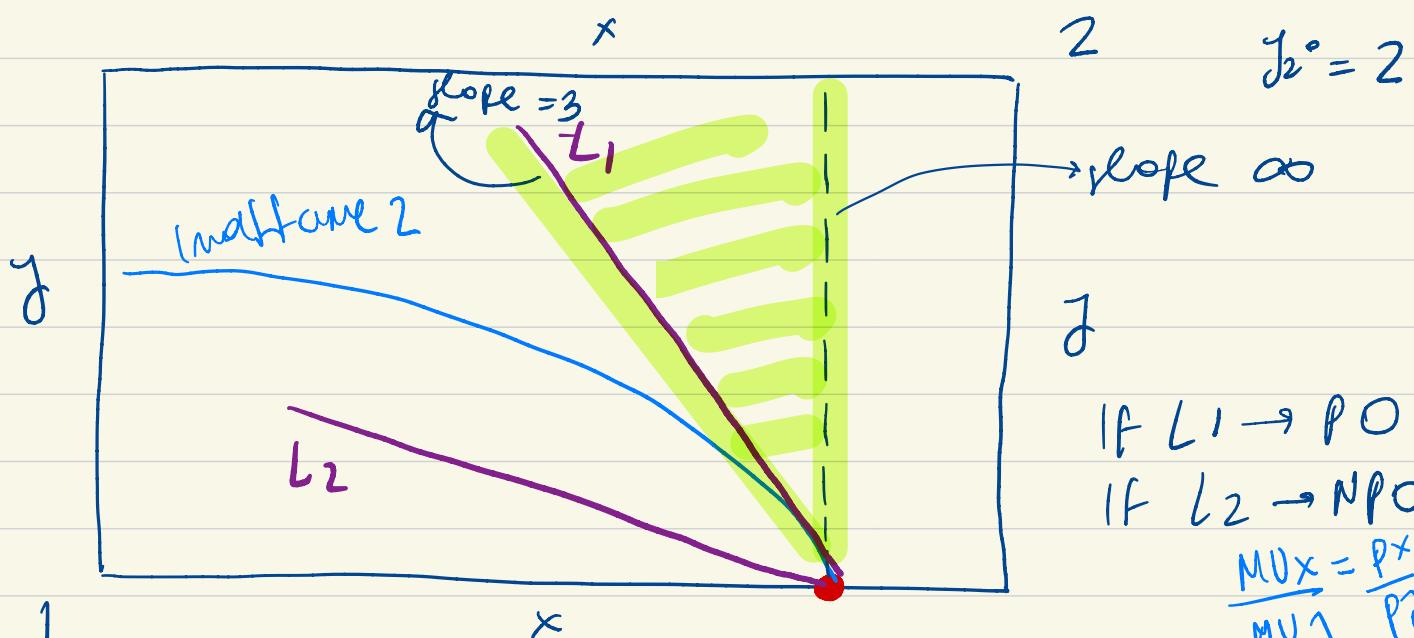
$$U_1 = x_1 - \gamma j_2 = x_1 - \gamma (2 - j_1) = x_1 + \gamma j_1 - 2\gamma$$

$2 - j_1$  by MC

(covers 1 <sup>↑</sup> area)  
very in eq.)

$$U_2 = \frac{1}{2} \log(x_2) + \frac{1}{2} \log(j_2)$$

Eg.  $x_1^* = 1/3, j_1^* = 0$   
 $x_2^* = 2/3$



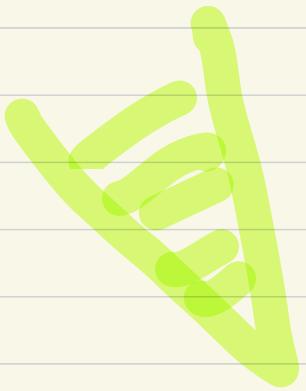
$$\begin{aligned} x_1 + \gamma j_1 - 2\gamma &= j \\ x_1 - 2\gamma &= -\gamma j \\ \text{slope: } -\frac{j}{\gamma} &= 3 \end{aligned}$$

concave 1 indifference

$$\frac{1}{\gamma} = 3$$

$$-\frac{1}{f} \in [-3; -\infty) \rightarrow f \in (\frac{1}{3}, +\infty)$$

within this range corner 1  
we have PO.



$$\therefore \frac{1}{f} \in [3, \infty] \rightarrow f \in [0, \frac{1}{3}]$$

c) If welfare theorem holds? One assumption at the very beginning is violated  
No,  
because we have externality.

assumption everybody  
only depends on  
his or her own  
consumption  
we need NO EXTERNALITIES!!

It should hold for any  $f$  if we say it  
holds.

If person 1 a shapley tunner  $\rightarrow$  he would care  
more about  $j_1$  to give  
corner 2 less  $j_2$   
 $\rightarrow$  but here you take  $j_2$   
as given when  
tuning demand.  
ex ante he does not take  
me  
Ex post  $\rightarrow$  Yes!

[Be explicit about your assumptions]

## Problem 2

(a) If  $\gtrsim$  are LNS, at a CE pure vector, expenditure on  $U(U_1, U_2)$  is minimized at agg endowment.

2 consumers:  $x, y$

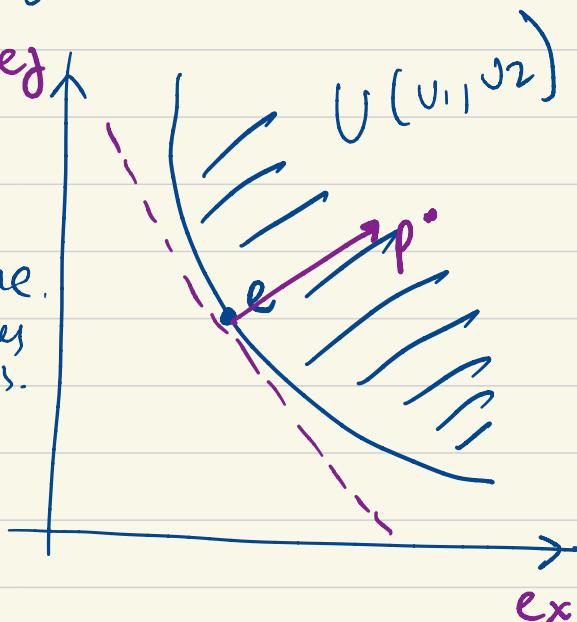
2 currencies: 1, 2

I will assume  $\gtrsim$  are LNS and monotone.

monotone implies LNS.

If  $\exists e' \text{ s.t. } e' \in U(U_1, U_2)$  and

$$p^* e' < p^* e \quad \therefore e' \in U(U_1, U_2),$$



By definition there are  $c_1', c_2'$  such that:

$$\left\{ \begin{array}{l} c_1' + c_2' \leq e' \\ u_1(c_1') \geq u_1(c_1^*) \\ u_2(c_2') \geq u_2(c_2^*) \end{array} \right.$$

[some write =]

is not mat  
re endow  
the endow  
num num  
exhaust

$\therefore \gtrsim$  is monotone, we have  $p^* \geq 0$ .

$$\begin{aligned} \text{So } p^* c_1 + p^* c_2 &\leq p^* e' \\ &< p^* e \end{aligned}$$

prices are the  
same  $p^*$   
[monotone]

$$\therefore p^* c_1 + p^* c_2 < p^* e_1 + p^* e_2$$

So there is some  $i \in \{1, 2\}$  such that  $p^* c_i' < p^* e_i$

But notice that  $u_i(c_i^*) \geq u_i(c_i')$ , violating LNS.

(b) utility smly quasi-concave,  $U(u_1, u_2)$  smly convex set

Assume that  $\gtrsim$  is lower-semi continuous. ( $\gtrsim$  is cont)

If  $e^1, e'' \in U(u_1, u_2)$ . WTS  $\alpha e^1 + (1-\alpha)e'' \in \text{Int } U(u_1, u_2)$

for  $\alpha \in (0, 1)$  (ie  $\exists$  open set  $O \subset U(u_1, u_2)$  st  $\alpha e^1 + (1-\alpha)e'' \in O$ )

Suppose  $c_i^1 \gtrsim c_i^*$  and  $c_i'' \gtrsim c_i^*$  where  $c_i^1$  and  $c_i''$  are allocations of  $e^1$  and  $e''$ .

$\therefore \gtrsim$  is smly convex ( $U(\cdot)$  is smly anal concave)

$\therefore \alpha c_i^1 + (1-\alpha)c_i'' \gtrsim c_i^*$ . Define  $P(c_i^*)$  as the smallest upper contour set of  $c_i^*$

$$P(c_i^*) = \{c_i : c_i \gtrsim_i c_i^*\}$$

$\therefore \alpha c_i^1 + (1-\alpha)c_i'' \in P(c_i^*)$

Define  $P = P(c_1^*) + P(c_2^*) = \left\{ c : c = c_1 + c_2 ; c_1 \in P(c_1^*) \right. \\ \left. c_2 \in P(c_2^*) \right\}$

$\therefore \gtrsim$  is lower-semi continuous.

$\therefore P(c_i^*)$  is open (since strict lower-semi-continuity)

$\therefore P$  is open  $\subset U(u_1, u_2)$  (sum of two open sets)

Since  $\alpha c_i^1 + (1-\alpha)c_i'' \in P(c_i^*) \forall i$ , then

$$\alpha(c_1^1 + c_2^1) + (1-\alpha)(c_1'' + c_2'') \in P$$

$$\therefore \alpha e^1 + (1-\alpha) e'' \in P$$

$$\therefore \alpha e^1 + (1-\alpha) e'' \in \text{Int } U(v_1, v_2)$$