

Logistics

- TA Office Hours:
 - Mon & Wed 18:30-19:30pm
 - In person (Uris 465) or by Zoom (same link as before).
- All problem sets are due before Friday section.
- Section's material https://github.com/luciacasal/-Cornell_ECON6100.

1 Walrasian equilibrium: pure exchange economy ¹

1.1 Important concepts:

1. **Setup.** Consider an exchange economy:

- (a) L goods, I consumers
- (b) endowment $e^i = (e_1^i, \dots, e_L^i) \in \mathbb{R}_+^L$ and aggregate endowments $e = \sum_i e^i$
- (c) prices $p = (p_1, \dots, p_L) \in \mathbb{R}_+^L / \{0\}$
- (d) demand function of consumer i is $d_i(p, e^i) \in \mathbb{R}_+^L$, excess demand is $z_i(p, e) = d_i(p, e^i) - e^i \in \mathbb{R}^L$, with aggregate excess demand $z(p, e) = \sum_i z_i(p, e^i) \in \mathbb{R}^L$.

2. **Definition 1. Walrasian equilibrium:** (or competitive equilibrium) is a price vector p^* such that (1) $z(p^*, e) \leq 0$; (2) $p^* z(p^*, e) = 0$.

3. **Theorem 1. Existence of Walrasian Equilibrium.** If excess demand function $z(p, e)$ satisfies:

- (a) A1: [HOD0] $z(p, e)$ is HOD in p .
- (b) A2: [Walras' Law] $\sum_l p_l z_l(p, e) = 0$
- (c) A3: [Continuity] $z(p, e)$ is continuous in p .

then there is a Walrasian equilibrium price p^* .

¹These notes borrow from Jaden Chen's notes from 2020.

1.2 Questions

1. (Varian 17.4) There are two consumers A and B with the following utility functions and endowments

$$\begin{aligned} u^A(x, y) &= \alpha \log x + (1 - \alpha) \log y & e^A &= (0, 1) \\ u^B(x, y) &= \min\{x, y\} & e^B &= (1, 0) \end{aligned}$$

- (a) Illustrate this situation in an Edgeworth box diagram for $\alpha = 1/2$. What are the market clearing prices and equilibrium allocation?
- (b) For $\alpha \in (0, 1)$, calculate the market clearing prices and equilibrium allocations. How do things change with α ?
2. (Modified from Q2009 June) Consider an exchange economy with L goods and N consumers. Each consumer's utility function is of the form $u_n(x_1, \dots, x_L) = \sum_l v_n(x_l)$, where each v_n is strictly concave, strictly increasing, differentiable and satisfies an Inada condition at the origin. Suppose that each consumer has a strictly positive endowment $\omega_n = (\omega_{n1}, \dots, \omega_{nL}) \gg 0$.
- (a) If $\sum_n \omega_{n1} = \sum_n \omega_{n2} = \dots = \sum_n \omega_{nL}$, what are equilibrium prices?
- (b) Show that if $\sum_n \omega_{n1} > \sum_n \omega_{n2} > \dots > \sum_n \omega_{nL}$ then for the competitive equilibrium price vector p^* , we have $p_1^* < p_2^* < \dots < p_L^*$.

2 General Equilibrium

2.1 Important concepts:

1. Definition 1: Private-ownership economy. We move from pure exchange economy to a private-ownership economy by introducing a set of producers.
 - (a) L commodities, aggregate endowments $\bar{\omega} \in \mathbb{R}^L$.
 - (b) Firms: $j = 1, \dots, J$ with nonempty and closed production set $Y_j \in \mathbb{R}^L$
 - (c) Consumers: $i = 1, \dots, I$ with consumption set $X_i \subset \mathbb{R}^L$, preferences \succeq_i and endowment $\omega_i \in X_i$. Consumer i 's ownership of firm j is $\theta_{ij} \in [0, 1]$ and $\sum_i \theta_{ij} = 1$.
2. Definition 2. Walrasian equilibrium for a private-ownership economy: denotes an allocation (x^*, y^*) and a price vector p^* such that:
 - (a) (Profit maximization): For all j , $p^* \cdot y_j^* \geq p^* \cdot y_j$ for all $y_j \in Y_j$
 - (b) (Utility maximization): For all i , x_i^* is \succeq_i -*maximal* in the budget set:

$$\left\{ x_i \in X_i : p^* \cdot x_i \leq p^* \cdot \omega_i + \sum_i \theta_{ij} p^* \cdot y_j^* \right\}$$

- (c) (Markets clear): $\sum_i x_i^* = \bar{\omega} + \sum_j y_j^*$

Q1)

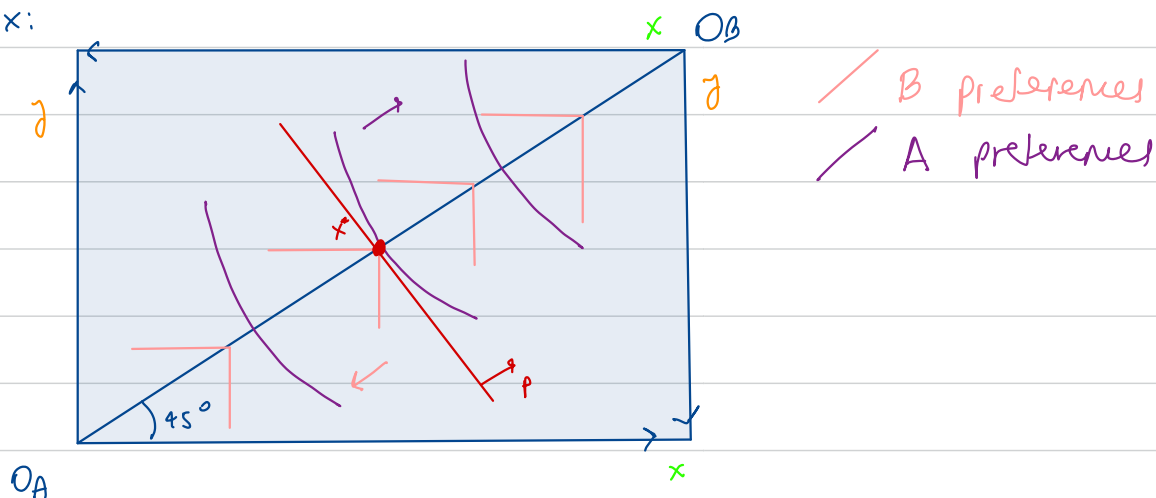
2 consumers, A and B. with utility functions:

$$u^A(x, z) = \alpha \log x + (1-\alpha) \log z \quad e^A = (0, 1)$$

$$u^B(x, z) = \min\{x, z\} \quad e^B = (1, 0)$$

a) $\alpha = 1/2$. Edgeworth box. Mkt clearing prices and eq. allocation.

Edgeworth box:



normalize $p_x = p$ $p_z = 1$

consumer A

$$\begin{cases} \max & \alpha \log x_A + (1-\alpha) \log z_A \\ \text{st} & p \cdot x_A + 1 \cdot z_A \leq p \cdot 0 + 1 \cdot 1 \\ & x_A, z_A \geq 0 \end{cases}$$

\therefore Inada cond \rightarrow interior solution \rightarrow ignore non-neg

\therefore concave \rightarrow SOC \checkmark (FOC sufficient)

\therefore monotone preferences \rightarrow Walras Law $\rightarrow p_x x_A + p_z z_A = 1$

We know in a competitive equilibrium

$$\frac{MgU_{x_A}}{MgU_{z_A}} = \frac{p_x}{p_z} = p$$

$$\rightarrow \frac{\frac{\alpha}{x_A}}{\frac{1-\alpha}{z_A}} = p$$

$$\frac{y_A}{x_A} \cdot \frac{\alpha}{1-\alpha} = p$$

$$y_A = p x_A \cdot \frac{1-\alpha}{\alpha}$$

From budget constraint:

$$p \cdot x_A + y_A = 1$$

$$p x_A + p x_A \cdot \frac{1-\alpha}{\alpha} = 1$$

$$p x_A \left(\frac{1}{\alpha} \right) = 1$$

①

$$x_A^* = \frac{\alpha}{p}$$

②

$$y_A = 1 - \alpha$$

For $\alpha = 1/2$

$$x_A = \frac{1}{2p}$$

$$y_A = 1/2$$

consumer B

$$\max \min \{x_B, y_B\}$$

$$\text{s.t. } p x_B + y_B \leq 1 \cdot p$$

$$x_B, y_B \geq 0$$

we know $p_x, p_y > 0$ because:

Claim: if $\exists i$ such that z_i is strongly monotone, then equilibrium price $p^* > 0$. (If not, the agent would like to consume ∞ amounts of the good when price is 0, and market clearing condition would be violated).

$$\Rightarrow x_B = y_B \quad (\text{see graph})$$

$$\therefore \text{from budget constraint } p x_B + x_B = p$$

③

$$y_B = x_B = \frac{p}{1+p}$$

Now we use market clearing to find prices:

$$\begin{cases} x_A + x_B = 1 = e x_A + e x_B \\ y_A + y_B = 1 = e y_A + e y_B \end{cases}$$

Substitute ① ② ③:

- $\frac{\alpha}{p} + \frac{p}{1+p} = 1$

- $(1-\alpha) + \frac{p}{1+p} = 1 \longrightarrow \frac{p}{1+p} = \alpha \longrightarrow p = \frac{\alpha}{1-\alpha}$ *

For $\alpha = 1/2 \longrightarrow p = 1 \longrightarrow$

$x_A = 1/2$	$y_A = 1/2$
$x_B = 1/2$	$y_B = 1/2$

b) Done in part a. If $\alpha \nearrow \rightarrow p \nearrow$ α is like the marginal utility of the log x...

claim:

* If Walras Law holds, only need to solve for $N-1$ equations in the market clearing.

Proof. take $N=2$

if WL holds, $p(x^A + x^B) + 1(y^A + y^B) = p e x + 1 e y$

If $\underbrace{x^A + x^B = e x}_{\text{market clearing for } x}$, then we have

$\underbrace{y^A + y^B = e y}_{\text{mkt clearing for } y}.$

Q2) (Modified from P 2009 June)

L goods
 N consumers

each n :

$$\begin{aligned} \max \quad & \sum_l v_n(x_l) \\ \text{st} \quad & \sum_l p_l x_l^n \leq \sum_l w_l^n \end{aligned}$$

for each consumer n :

$U_n(x_1, \dots, x_L) = \sum_l v_n(x_l)$; with v_n strictly concave (SOC ✓)
 " increasing (Walras Law)
 differentiable (calculus)
 budget constraint (interior sol)

suppose strictly positive endowment $w_n = (w_{n1}, \dots, w_{nL}) \gg 0$

a) If $\sum_n w_{n1} = \sum_n w_{n2} = \dots = \sum_n w_{nL}$, what are eq. prices?

from FOC: for all n and l, m :

$$\frac{v_n'(x_l^n)}{v_n'(x_m^n)} = \frac{p_l}{p_m}$$

if $\exists l, m$ such that $p_l > p_m$, since v_n is strictly concave, we have $x_m^n > x_l^n$.

$$\Rightarrow \sum_n x_m^n > \sum_n x_l^n \quad \forall n$$

$$\therefore \text{in equilibrium} \quad \begin{cases} \sum_n x_l^n = \sum_n w_l^n \\ \sum_n x_m^n = \sum_n w_m^n \end{cases} \quad [\text{market clearing}]$$

then from ①: $\sum_n w_m^n > \sum_n w_l^n$ contradiction!
 with

$$\Rightarrow p_l^* = p_m^* \quad \forall m, l$$

b) Show that if $\sum_n w_{n1} > \sum_n w_{n2} > \dots > \sum_n w_{nL}$, then for the competitive equilibrium price vector p^* , we have $p_1^* < p_2^* < \dots < p_L^*$

notice that if there are m and l such that

$$\sum_n w_{nl} > \sum_n w_{nm}$$

but $p_l \geq p_m$

from a) we must have $x_l^n \leq x_m^n \quad \forall n$

$$\underbrace{\sum_n x_l^n}_{\parallel} \leq \underbrace{\sum_n x_m^n}_{\parallel}$$

$$\sum_n w_{nl} < \sum_n w_{nm}$$

contradiction!

market clearing