

## Sum of Elements (medium)

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### Problem Statement

Given an array, find the sum of all numbers between the K1'th and K2'th smallest elements of that array.

Example 1:

```
Input: [1, 3, 12, 5, 15, 11], and K1=3, K2=6
Output: 23
Explanation: The 3rd smallest number is 5 and 6th smallest number 15. The sum of numbers coming between 5 and 15 is 23 (11+12).
```

Example 2:

```
Input: [3, 5, 8, 7], and K1=1, K2=4
Output: 12
Explanation: The sum of the numbers between the 1st smallest number (3) and the 4th smallest number (8) is 12 (5+7).
```

### Try it yourself

Try solving this question here:

JavaPython3JS C++

```
1 def find_sum_of_elements(nums, k1, k2):
2     # TODO: Write your code here
3     return -1
4
5
6 def main():
7
8     print("Sum of all numbers between k1 and k2 smallest numbers: " +
9           str(find_sum_of_elements([1, 3, 12, 5, 15, 11], 3, 6)))
10    print("Sum of all numbers between k1 and k2 smallest numbers: " +
11          str(find_sum_of_elements([3, 5, 8, 7], 1, 4)))
12
13
14    main()
15
```

RunSaveReset

### Solution

This problem follows the [Top 'K' Numbers](#) pattern, and shares similarities with [Kth Smallest Number](#).

We can find the sum of all numbers coming between the K1'th and K2'th smallest numbers in the following steps:

1. First, insert all numbers in a min-heap.
2. Remove the first `K1` smallest numbers from the min-heap.
3. Now take the next `K2-K1-1` numbers out of the heap and add them. This sum will be our required output.

### Code

Here is what our algorithm will look like:

JavaPython3C++JS

```
1 from heapq import *
2
3
4 def find_sum_of_elements(nums, k1, k2):
5     minHeap = []
6     # insert all numbers to the min heap
7     for num in nums:
8         heappush(minHeap, num)
9
10    # remove k1 small numbers from the min heap
11    for _ in range(k1):
12        heappop(minHeap)
13
14    elementSum = 0
15    # sum next k2-k1-1 numbers
16    for _ in range(k2 - k1 - 1):
17        elementSum += heappop(minHeap)
18
19    return elementSum
20
21
22 def main():
23
24    print("Sum of all numbers between k1 and k2 smallest numbers: " +
25          str(find_sum_of_elements([1, 3, 12, 5, 15, 11], 3, 6)))
26    print("Sum of all numbers between k1 and k2 smallest numbers: " +
27          str(find_sum_of_elements([3, 5, 8, 7], 1, 4)))
28
29
30    main()
31
```

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#### Time complexity

Since we need to put all the numbers in a min-heap, the time complexity of the above algorithm will be  $O(N * \log N)$  where 'N' is the total input numbers.

#### Space complexity

The space complexity will be  $O(N)$ , as we need to store all the 'N' numbers in the heap.

### Alternate Solution

We can iterate the array and use a max-heap to keep track of the top K2 numbers. We can, then, add the top `K2-K1-1` numbers in the max-heap to find the sum of all numbers coming between the K1'th and K2'th smallest numbers. Here is what the algorithm will look like:

JavaPython3C++JS

```
1 from heapq import *
2
3
4 def find_sum_of_elements(nums, k1, k2):
5     maxHeap = []
6     # keep smallest k2 numbers in the max heap
7     for i in range(len(nums)):
8         if i < k2 - 1:
9             heappush(maxHeap, -nums[i])
10        elif nums[i] < -maxHeap[0]:
11            heappop(maxHeap) # as we are interested only in the smallest k2 numbers
12            heappush(maxHeap, -nums[i])
13
14    # get the sum of numbers between k1 and k2 indices
15    # these numbers will be at the top of the max heap
16    elementSum = 0
17    for _ in range(k2 - k1 - 1):
18        elementSum += -heappop(maxHeap)
19
20    return elementSum
21
22
23 def main():
24
25    print("Sum of all numbers between k1 and k2 smallest numbers: " +
26          str(find_sum_of_elements([1, 3, 12, 5, 15, 11], 3, 6)))
27    print("Sum of all numbers between k1 and k2 smallest numbers: " +
28          str(find_sum_of_elements([3, 5, 8, 7], 1, 4)))
29
30
31    main()
```

RunSaveReset

#### Time complexity

Since we need to put only the top K2 numbers in the max-heap at any time, the time complexity of the above algorithm will be  $O(N * \log K2)$ .

#### Space complexity

The space complexity will be  $O(K2)$ , as we need to store the smallest 'K2' numbers in the heap.