

Dijkstra's Shortest Path Algorithm

The shortest path between u and v is denoted $\delta(u, v)$, if there is no path, then $\delta(u, v) = \infty$

1 Can a shortest path contain a cycle?

A directed cycle is:

- Positive: if its edge weights sum up to a positive number
- Negative: if its edge weights sum up to a negative number

If there is a positive cycle in the graph, it will not be contained in any shortest path between u and v so we can assume that the shortest paths we find contain no positive cycles.

However if there is a negative cycle between u and v , then $\delta(u, v) = -\infty$ so we shall assume that the graphs we consider do not contain negative cycles.

2 Single-Source Shortest Paths

- Aim: to describe an algorithm that solves the single-source shortest paths problem, i.e. an algorithm that finds the shortest path from a specific source vertex
- This is a generalization of BFS
- So the output of the algorithm should be two arrays d, π where for each vertex v :
 - $d(v) = \delta(s, v)$
 - $\pi(v)$ is the predecessor of v

3 Relaxation

- Assume that the weight on every edge is non-negative
- We do not directly compute the entry $d(v) = \delta(s, v)$
- Instead, at every step, $d(v)$ is an estimate for $\delta(s, v)$
 - Initially, $d(v) = \infty$, and it always remains $d(v) \geq \delta(s, v)$
 - $d(v)$ is updated (i.e. it decreases) as shorter paths are found
 - At the end of the algorithm we have $d(v) = \delta(s, v)$

Listing 1 Initialise-Single-Source(G, s)

```

1 for each vertex  $v \in V(G)$  do
2    $d(v) = \infty$ 
3    $\pi(v) = \text{NIL}$ 
4  $d(s) = 0$ 
```

The process of relaxing an edge (u, v) :

- Test whether we can improve the shortest path from s to v that we found so far, by going through u
- If yes, then update $d(v)$ and $\pi(v)$
 - Decrease the estimate $d(v)$
 - Update the predecessor $\pi(v)$ to u
- The algorithm first calls initialise-single-source and then it repeatedly relaxes the appropriate edges (according to the weight function w)

Listing 2 Relax(u, v, w)

```
1 if  $d(v) > d(u) + w(u, v)$  then  
2    $d(v) = d(u) + w(u, v)$   
3    $\pi(v) = u$ 
```

4 Dijkstra's Algorithm