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Modular Arithmetic

1 Basic modular arithmetic

1.1 Definition

if a,b are integers and m is a positive integer then a is congruent to b modulo m iff m—(a-b). Notation $a \equiv b \pmod{m}$ Example: $8 \equiv 5 \pmod{3}$ because $3 \mid (8-5); -5 \equiv 9 \pmod{7}$ because $7 \mid (-5-9)$

1.2 Lemma

If a,b,m are integers and m > 0 then $a \equiv b(mod m)$ iff rem(a, m) = rem(b, m)

1.3 Lemma

Let m be a positive integer and let $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$ Then

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a + c \equiv b + d \pmod{m} and ac \equiv bd \pmod{m}
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2 Linear congruences

Congruences work a lot like equations, but there are some differences:

- if $ac \equiv bc \pmod{m}$ and $c \not\equiv 0 \pmod{m}$ it is possible that $a \not\equiv b \pmod{m}$. For example $2 \cdot 3 \equiv 4 \cdot 3 \pmod{6}$, but $2 \neq 4 \pmod{6}$
- If $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$, it is possible that $a^c \not\equiv d^d \pmod{m}$
- A congruence of the form $ax \equiv b \pmod{m}$ is called a **linear congruence**
- Such congruences often appear in applications of number theory
- Solving such a congruence means finding all c (by default, in the range $\{0, 1, ..., m-1\}$) such that $ac \equiv b \pmod{m}$, such c might not be unique

2.1 Example

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Solve 3x + 1 \equiv 5 \pmod{7}
Subtract 1 from both sides, get 3x \equiv 4 \pmod{7}
Now multiply both sides by 5, have 15x = 20 \pmod{7}
Since 15 \equiv 1 \pmod{7} and 20 \equiv 6 \pmod{7}, have x \equiv 6 \pmod{7}. So x = 6
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3 Multiplicative inverses

- An easy way to solve equation ax = b is to multiply both parts by a^{-1}
- We cannot do this within integers, but we often can when working modulo m
- Call \bar{a} the (multiplicative) inverse of a modulo m if $\bar{a}a \equiv 1 \pmod{m}$
- Multiplicative inverses do not always exist, e.g. $2\bar{a} \not\equiv 1 \pmod{4}$ for any \bar{a}

3.1 Theorem

If gcd(a,m)=1 then the inverse of a modulo m exists, and is unique (that is, there is a unique $0 \le \overline{a} < m$ with $\overline{a}a \equiv 1 \pmod{m}$)

3.2 Proof

We show existence, and leave uniqueness as an exercise Since gcd(a,m)=1, we have sa+tm=1 for some s,t. Then $sa\equiv 1 \pmod m$ Let s=qm+r where $0 \le r < m$, then $ra=(s-qm)a\equiv sa\equiv 1 \pmod m$ so r is the required inverse

• note that s can be found by using euclid's algorithm. Then r is easy to find

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4 The Chinese remainder theorem

4.1 Theorem

Let $m_1, ..., m_n$ be pairwise relatively prime positive integers and $a_1, ..., a_n$ arbitrary integers. Then the system

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x \equiv a_1 \pmod{m_1}
x \equiv a_2 \pmod{m_2}
\vdots
x \equiv a_n \pmod{m_n}
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has a unique solution modulo $m = m_1 m_2 \cdots m_n$. That is, there is a unique solution x with $- \le x < m$ and every other solution is congruent to x modulo m.

4.2 Proof

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Let M_k = m/m_k
We have gcd(M_k, m_k). Hence M_k y_k \equiv q(\bmod m_k) for some y_k
Let x = a_1 M_1 y_1 + a_2 M_2 y_2 + ... + a_n M_n y_n
We show that x \equiv a_k (\bmod m_k) for all k.
Notice that M_j \equiv 0 (\bmod m_k) if j \neq k. Hence x \equiv a_k M_k y_k \equiv a_k (\bmod m_k)
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5 Computer arithmetic with large numbers

- Suppose $m_1, m_2, ..., m_n$ are all pairwise relatively prime (and all ≥ 2)
- By the chinese remainder theorem, any number $0 \le a < m$ can all be uniquely represented by the n-tuple $(a_1, a_2, ..., a_n)$ where $a_i = rem(a, m_i)$ for all i
- Example: let $m_1 = 3$ and $m_2 = 4$. Then the numbers < 12 are represented as

- To perform arithmetic with large numbers, choose the moduli $m_1, ..., m_n$ so that $m = m_1 \cdots m_n >$ the result of the operations you want to carry out
- Then arithmetic can be performed with representations of numbers
- Example: compute 2 · 5. Instead, multiply (2,2) and (2,1) component wise. 1st component modulo 3 and 2nd modulo 4. Get (1,2) which represents 10
- Advantages: can work with very large numbers and can compute in parallel
- Particularly good choices for m_i : numbers of the form $2^p 1$

6 Fermat's Little theorem

6.1 Theorem

If p is a prime and a is not a multiple of p then $a^{p-1} \equiv 1 \pmod{p}$ Furthermore, for every integer a, $a^p \equiv a \pmod{p}$ MCS - DMLA - Term 2 Sam Robbins

6.2 Example

- We know how to find inverses modulo prime p (via Euclid's algorithm)
- The above theorem gives an alternate approach: $a^{p-2} \cdot a \equiv 1 \pmod{p}$ hence $rem(a^{p-2}, p)$ is the required inverse

Find the multiplicative inverse of 6 modulo 17

Solution: we need to compute $rem(6^{15}, 17)$, which can be done as follows.

$$6^{2} \equiv 36 \equiv 2 \pmod{17}$$

$$6^{4} \equiv (6^{2})^{2} \equiv 2^{2} \equiv 4 \pmod{17}$$

$$6^{8} \equiv (6^{4})^{2} \equiv 4^{2} \equiv 16 \pmod{17}$$

$$6^{15} \equiv 6^{8} \cdot 6^{4} \cdot 6^{2} \cdot 6 \equiv 16 \cdot 4 \cdot 2 \cdot 6 \equiv 3 \pmod{17}$$

Therefore, $rem(6^16, 17) = 3$ is the required inverse. Indeed $3 \cdot 6 \equiv 1 \pmod{17}$

7 Euler's theorem

Recall Euler's ϕ -function. $\phi(n)$ is the number of integers $1 \le a \le n$ that are relatively prime with n. Euler's theorem generalises Fermat's little theorem to non-prime moduli

7.1 Theorem

If n is a positive integer and gcd(a, n) = 1 then $a^{\phi(n)} \equiv 1 \pmod{n}$

7.2 Method

- If n is a prime then $\phi(n) = n 1$, so this is indeed a generalisation
- If gcd(a,n)=1, then, as we proved, the inverse of a modulo n exists and can be found using Euclid's algorithm
- By Euler's Theorem, it can also be found as $rem(a^{\phi(n)-1}, n)$
- Can a have a multiplicative inverse modulo n is gcd(a, n) > 1? No
 - If the inverse \bar{a} exists we have $\bar{a}a \equiv 1 \pmod{n}$, i.e. $\bar{a}a 1 = kn$ for some k
 - Rewrite as $\overline{a}a + (-k)n = 1$, it follows that gcd(a, n) = 1

8 Computing Euler's ϕ -function

8.1 Lemma

If m_1 and m_2 are relatively prime then $\phi(m_1 \cdot m_2) = \phi(m_1) \cdot \phi(m_2)$. If p is prime then $\phi(p^k) = p^k - p^{k-1}$

8.2 Proof

By the chinese remainder theorem, there is a 1 to 1 correspondence between

- numbers x with $0 \le x < m_1 m_2$ and
- pairs $(1_2, a_2)$ such that $0 \le a_i < m_i$ and $x \equiv a_i \pmod{m_i}$ for i = 1, 2

Since $a_i = rem(x, m_i)$ we have $gcd(x, m_i) = gcd(a_i, m_i)$ for i=1,2

We have $gcd(x, m_1, m_2) = gcd(x, m_1) \cdot gcd(x, m_2) = gcd(a_1, m_1) \cdot gcd(a_2, m_2)$ (the first equality holds because $gcd(m_1, m_2) = 1$) In particular, $gcd(x, m_1m_2) = 1$ iff $gcd(a_1, m_1) = gcd(a_2, m_2) = 1$. This immediately implies $\phi(m_1cdotm_2) = \phi(m_1) \cdot \phi(m_2)$

8.3 Example

$$\phi(75) = \phi(3 \cdot 5^2) = \phi(3) \cdot \phi(5^2) = (3^1 - 3^0) \cdot (5^2 - 5) = 40$$