

# Decision Problems

## 1 Optimisation vs Decision Problems

- A major variation of optimization problems: decision problems
- Answer is not a value but YES/NO
- Every optimisation -problem has a decision counterpart
- An optimisation problem has a fast algorithm iff the corresponding decision problem has a fast algorithm

## 2 Decision problems and encodings

The standard way to define a decision problem is to describe a generic instance and a yes-no question about each instance

**Reachability:**

- Instance: A finite directed graph  $G$  and vertices  $s$  and  $t$
- Question: Is there a path in  $G$  from  $s$  to  $t$ ?

To input problems to a computer, each instance must be encoded as a string of symbols over some alphabet. We need an encoding scheme

To ensure that encoding the problem does not change its essential nature, an encoding scheme must be concise

- Represent numbers efficiently
- Not add unnecessary information

## 3 Languages

An alphabet,  $\Sigma$  is a (finite) set of symbols

A string over  $\Sigma$  is a finite sequence of symbols from  $\Sigma$

A language over  $\Sigma$  is any set of strings over  $\Sigma$

For a problem  $\Pi$  and an encoding scheme  $e$  with alphabet  $\Sigma$ , the set of all strings corresponding to instances with answer yes is denoted  $\mathcal{L}(\Pi, e)$  and is called the language associated with  $\Pi$  and  $e$

For decision problems, we just want to decide whether the (encoding of a) given instance belongs to the alphabet  $\mathcal{L}(\Pi, e)$

## 4 Complexity of Problems

- The problems encountered so far in this course have all proved to be tractable (since we have found fast algorithms for all of them)
- There are many problems however, which cannot be quickly solved in practice, i.e., which are intractable
- There are many difficulties:
  - What do we mean by tractable and intractable
  - Can we define these notions formally
  - Can we prove that a problem is one but not the other
- One technique that will prove very useful: showing that one decision problem can be transformed(reduced) into another

## 5 Complexity measures

Every decidable problems has a set of algorithms that solved it. The properties of this set of algorithms:

- The difficulty of constructing the algorithm
- The length of the shortest possible algorithm
  - Static complexity measure
  - Useful for classifying the complexity of strings, called Kolmogorov complexity
- The efficiency of the most possible algorithm?
  - Dynamic complexity measure
  - A numerical function that measures the maximum resources used by an algorithm to compute the answer to a given instance

## 6 Dynamic complexity measures

- The most critical resources are often time/space
- By considering Turing Machines as our model of computation

### Definition: Time complexity

The time complexity of a Turing Machine  $T$  is the function  $Time_T$  such that  $Time_T(x)$  is the number of steps taken by the computation  $T(x)$

### Definition: Space complexity

The space complexity of a Turing Machine  $T$  is the function  $Space_T$  such that  $Space_T(x)$  is the number of distinct tape cells visited during the computation  $T(x)$

## 7 Equivalence

**Theorem 1** An  $n$ -vertex graph  $G$  has an independent set of size  $k$  iff  $G$  has a vertex cover of size  $n-k$

**Corollary** - The decision problems independent set and vertex cover are equivalent

## 8 Clique

- Instance - A graph  $G$  and an integer  $k$
- Question: Does  $G$  have a clique of size at least  $k$  - i.e. a set of at least  $k$  vertices that are all adjacent to one another

## 9 Complement

### Definition: Complement

The complement of a graph  $G$  has the same vertex set, and there is an edge between two vertices  $u$  and  $v$  in the complement iff there is no edge from  $u$  to  $v$  in  $G$

**Theorem 2** A graph  $G$  has an independent set of size  $k$  iff its complement  $\bar{G}$  has a clique of size  $k$

**Corollary** - The decision problems independent set and clique are equivalent

## 10 Set cover

- Instance - A set  $U$  of  $n$  elements, a collection of subsets  $S_1, S_2, \dots, S_t$  whose union equals  $U$ , and an integer  $k$
- Question: Does there exist a collection of at most  $k$  of these sets whose union is equal to all of  $U$

Vertex cover can be solved using an algorithm for SET COVER