

EDAN40

examination

4 hp

12th April 2023

8:00 - 13:00

WRITE ONLY ON ONE SIDE OF THE PAPER - the exams will be scanned in and only the front/odd pages will be read.

DO NOT WRITE WITH OTHER COLOUR THAN BLACK OR DARK BLUE - lightly coloured text may disappear during scanning

PUT YOUR ID AND PAGE NUMBER ON EACH PAGE YOU SUBMIT - make sure that the amount of pages is equal to the amount you note on the front information page

WRITE CLEARLY - if we cannot read you we cannot (properly) grade you.

PRELIMINARY MAX AMOUNT OF POINTS: 6

Exam

1. Point-free notation

Rewrite the following two definitions into a point-free form (i.e., $f = \dots$, $g = \dots$), using neither lambda-expressions nor list comprehensions nor enumeration nor where clause nor let clause:

```
f x y = (42 - y) * x
g x y = y x
```

2. Type derivation

Find the types of the following expressions:

```
(( $ ) $)
((. ) .)
((: ) :)
((==) ==)
((||) ||)
```

3. Proving program properties

The Functor class is defined as follows:

```
class Functor f where
  fmap :: (a -> b) -> f a -> f b
```

It is mandatory that all instances of Functor should obey:

```
fmap id      = id
fmap (p . q) = (fmap p) . (fmap q)
```

Assume the following definition of Maybe types as a functor instance:

```
instance Functor Maybe where
  fmap f (Just x) = Just (f x)
  fmap f Nothing  = Nothing
```

Is this a correct definition of a functor instance? Why or why not? **Prove your claim.**

4. Evaluation

Explain what a *thunk* is.

5. Monadic computations

What is the type of `e` defined below? Motivate your answer.

```
e k = do
  x <- k
  Nothing
  return False
```

6. Types and type classes

- Define a tree data structure so that the trees are ternary (i.e., each node has either three children or is a leaf) and store strings in leaves.
- Generalize your definition so that your ternary trees can contain objects of an arbitrary predetermined type in a leaf.
- Assuming your polymorphic trees type is denoted by `Tree3 a` (you may use your own name used above), write all necessary code so that the following function is correct:

```
myLength :: Tree3 String -> Tree3 Integer
myLength = fmap length
```

and yields a ternary tree with leaves containing lengths of the strings placed in the respective leaves of the argument tree.

Make sure that the following works as well:

```
myReverse :: Tree3 String -> Tree3 String
myReverse = fmap reverse
```

Good Luck!

```

{-A list of selected functions from the Haskell modules:
Prelude
Data.List
Data.Maybe
Data.Char -}

-----
-- standard type classes
class Show a where
  show :: a -> String

class Eq a where
  (==), (/=) :: a -> a -> Bool

class (Eq a) => Ord a where
  (<), (<=), (>=), (>) :: a -> a -> Bool
  max, min :: a -> a -> a

class (Eq a, Show a) => Num a where
  (+), (-), (*) :: a -> a -> a
  negate :: a -> a
  abs, signum :: a -> a
  fromInteger :: Integer -> a

class (Num a, Ord a) => Real a where
  toRational :: a -> Rational

class (Real a, Enum a) => Integral a where
  quot, rem :: a -> a -> a
  div, mod :: a -> a -> a
  toInteger :: a -> Integer

class (Num a) => Fractional a where
  (/) :: a -> a -> a
  fromRational :: Rational -> a

class (Fractional a) => Floating a where
  exp, log, sqrt :: a -> a
  sin, cos, tan :: a -> a

class (Real a, Fractional a) => RealFrac a where
  truncate, round :: (Integral b) => a -> b
  ceiling, floor :: (Integral b) => a -> b

-----
-- numerical functions
even, odd :: (Integral a) => a -> Bool
even n = n `rem` 2 == 0
odd = not . even

-----
-- monadic functions
sequence :: Monad m => [m a] -> m [a]
sequence = foldr mcons (return [])

where mcons p q = do x <- p; xs <- q; return (x:xs)

-----
sequence_ :: Monad m => [m a] -> m ()
sequence_ xs = do sequence xs; return ()

-----
-- functions on functions
id :: a -> a
id x = x

const :: a -> b -> a
const x _ = x

(.) :: (b -> c) -> (a -> b) -> a -> c
f . g = \x -> f (g x)

flip :: (a -> b -> c) -> b -> a -> c
flip f x y = f y x

($) :: (a -> b) -> a -> b
f $ x = f x

-----
-- functions on Booleans
data Bool = False | True

(&&), (||) :: Bool -> Bool -> Bool
True && x = x
False && _ = False
True || _ = True
False || x = x

not :: Bool -> Bool
not True = False
not False = True

-----
-- functions on Maybe
data Maybe a = Nothing | Just a

isJust :: Maybe a -> Bool
isJust (Just a) = True
isJust Nothing = False

isNothing :: Maybe a -> Bool
isNothing = not . isJust

fromJust :: (Just a) -> a
fromJust (Just a) = a

maybetoList :: Maybe a -> [a]
maybetoList Nothing = []
maybetoList (Just a) = [a]

-----

```

```

listToMaybe :: [a] -> Maybe a
listToMaybe [] = Nothing
listToMaybe (a:_) = Just a

-- a hidden goodie
instance Monad [] where
  return x = [x]
  xs >=> f = concat (map f xs)

-- functions on pairs
fst :: (a, b) -> a
fst (x, y) = x
snd :: (a, b) -> b
snd (x, y) = y
curry :: ((a, b) -> c) -> a -> b -> c
curry f x y = f (x, y)
uncurry :: (a -> b -> c) -> (a, b) -> c
uncurry f p = f (fst p) (snd p)

-- functions on lists
map :: (a -> b) -> [a] -> [b]
map f xs = [ f x | x <- xs ]
(++) :: [a] -> [a] -> [a]
xs ++ ys = foldr (:) ys xs
filter :: (a -> Bool) -> [a] -> [a]
filter p xs = [ x | x <- xs, p x ]
concat :: [[a]] -> [a]
concat xss = foldr (++) [] xss
concatMap :: (a -> [b]) -> [a] -> [b]
concatMap f = concat . map f
head, last :: [a] -> a
head (x:_) = x
last [x] = x
last (_:xs) = last xs
tail, init :: [a] -> [a]
tail (_:xs) = xs
init [x] = []
init (x:xs) = x : init xs

```

```

null :: [a] -> Bool
null [] = True
null (_:_) = False
length :: [a] -> Int
length [] = 0
length (_:_) = 1 + length l
(x:_) !! 0 = x
(x:_) !! (n-1) = xs !! (n-1)
foldr :: (a -> b -> b) -> b -> [a] -> b
foldr f z [] = z
foldr f z (x:xs) = f x (foldr f z xs)
foldl :: (a -> b -> a) -> a -> [b] -> a
foldl f z [] = z
foldl f z (x:xs) = foldl f (f z x) xs
iterate :: (a -> a) -> a -> [a]
iterate f x = x : iterate f (f x)
repeat :: a -> [a]
repeat x = xs where xs = x:xs
replicate :: Int -> a -> [a]
replicate n x = take n (repeat x)
cycle :: [a] -> [a]
cycle [] = error "Prelude.cycle: empty list"
cycle xs = xs' where xs' = xs++xs'
take, drop :: Int -> [a] -> [a]
take n _ | n <= 0 = []
take _ [] = []
take n (x:xs) = x : take (n-1) xs
drop n xs | n <= 0 = xs
drop _ [] = []
drop n (x:xs) = drop (n-1) xs
splitAt :: Int -> [a] -> ([a], [a])
splitAt n xs = (take n xs, drop n xs)
takeWhile, dropWhile :: (a -> Bool) -> [a] -> [a]
takeWhile p [] = []
takeWhile p (x:xs) = x : takeWhile p xs
dropWhile p [] = []
dropWhile p (x:xs) = x : dropWhile p xs
dropWhile p [] = []
dropWhile p (x:xs) = dropWhile p xs'
| otherwise = xs

```

```

lines, words      :: String -> [String]
-- lines "apa\bepea\ncepa\n" == ["apa", "bepea", "cepa"]
-- words "apa bepea\ncepa" == ["apa", "bepea", "cepa"]

unlines, unwords  :: [String] -> String
-- unlines ["apa", "bepea", "cepa"] == "apa\bepea\ncepa"
-- unwords ["apa", "bepea", "cepa"] == "apa bepea cepa"

and, or            :: [Bool] -> Bool
and               = foldr (&) True
or                = foldr (||) False

any, all           :: (a -> Bool) -> [a] -> Bool
any p             = or . map p
all p             = and . map p

elem, notElem     :: (Eq a) => a -> [a] -> Bool
elem x            = any (== x)
notElem x         = all (/= x)

lookup            :: (Eq a) => a -> [(a,b)] -> Maybe b
lookup key []     = Nothing
lookup key ((x,y):xys)
  | key == x      = Just y
  | otherwise     = lookup key xys

sum, product      :: (Num a) => [a] -> a
sum               = foldl (+) 0
product           = foldl (*) 1

maximum, minimum  :: (Ord a) => [a] -> a
maximum []        = error "Prelude.maximum: empty list"
maximum xs        = foldl max xs
minimum []        = error "Prelude.minimum: empty list"
minimum xs        = foldl min xs

zip              :: [a] -> [b] -> [(a,b)]
zip              = zipWith (,)

zipWith          :: (a -> b -> c) -> [a] -> [b] -> [c]
zipWith z (a:as) (b:bs)
  = z a b : zipWith z as bs
zipWith _ _ _    = []

unzip            :: [(a,b)] -> ([a], [b])
unzip            = foldr (\(a,b) ~> (as,bs) -> (a:as,b:bs)) ([], [])

nub              :: [a] -> [a]
nub []           = []
nub (x:xs)       = x : nub [ y | y <- xs, x /= y ]

delete           :: Eq a => a -> [a] -> [a]
delete y []      = []

```

```

delete y (x:xs)  = if x == y then xs else x : delete y xs

(\\)             :: Eq a => [a] -> [a] -> [a]
(\\)             = foldl (flip delete)

union            :: Eq a => [a] -> [a] -> [a]
union xs ys     = xs ++ ( ys \\ xs )

intersect        :: Eq a => [a] -> [a] -> [a]
intersect xs ys = [ x | x <- xs, x `elem` ys ]

intersperse     :: a -> [a] -> [a]
-- intersperse 0 [1,2,3,4] == [1,0,2,0,3,0,4]

transpose       :: [[a]] -> [[a]]
-- transpose [[1,2,3],[4,5,6]] == [[1,4],[2,5],[3,6]]

partition       :: (a -> Bool) -> [a] -> ([a], [a])
partition p xs  = (filter p xs, filter (not . p) xs)

group           :: Eq a => [a] -> [[a]]
-- group "aapaabbeee" == ["aa", "p", "aa", "bbb", "eeee"]

isPrefixOf, isSuffixOf
  :: Eq a => [a] -> [a] -> Bool
isPrefixOf [] _ = True
isPrefixOf _ [] = False
isPrefixOf (x:xs) (y:ys) = x == y && isPrefixOf xs ys

isSuffixOf x y  = reverse x `isPrefixOf` reverse y

sort            :: [a] -> [a]
sort            = foldr insert []
insert          :: Ord a => a -> [a] -> [a]
insert x []     = [x]
insert x (y:xs) = if x <= y then x:y:xs else y:insert x xs

-- functions on Char

type String = [Char]

toupper, tolower :: Char -> Char
-- toupper 'a' == 'A'
-- tolower 'Z' == 'z'

digitToInt      :: Char -> Int
-- digitToInt '8' == 8

intToDigit      :: Int -> Char
-- intToDigit 3 == '3'

ord             :: Char -> Int
chr             :: Int -> Char

```

230412

N40-0001-DTN

(1)

Ableton

7. $f = \text{flip } \$ (*) . (42-)$

$g = \text{flip } \underline{\underline{(\$)}}$
10

0.9

2. a) $(1.\$)\ :: (a \rightarrow b) \rightarrow a \rightarrow b$ } renaming to
 $(2.\$)\ :: (c \rightarrow d) \rightarrow c \rightarrow d$ } distinguish between them

$$a := c \rightarrow d$$

$$b := c \rightarrow d$$

Substitute and remove first argument:

$$(\$)\ \$\ :: (c \rightarrow d) \rightarrow c \rightarrow d$$

Beautify expression:

Answer: $((\$)\ \$)\ :: (a \rightarrow b) \rightarrow a \rightarrow b$

0.2

- b) $((.2) . 1) \rightarrow$ renaming to distinguish between the functions

$$(.1)\ :: (b \rightarrow c) \rightarrow (a \rightarrow b) \rightarrow a \rightarrow c$$

$$(.2)\ :: (e \rightarrow f) \rightarrow (d \rightarrow e) \rightarrow d \rightarrow f$$

$$b := e \rightarrow f$$

$$c := (d \rightarrow e) \rightarrow d \rightarrow f$$

Substitute and remove first argument:

$$((.) .) \ :: (a \rightarrow e \rightarrow f) \rightarrow a \rightarrow (d \rightarrow e) \rightarrow d \rightarrow f$$

Beautify expression:

Answer: $((.) .) \ :: (a \rightarrow c \rightarrow d) \rightarrow a \rightarrow (b \rightarrow c) \rightarrow b \rightarrow d$

0.2

2. c) $((:1) : 2) \rightarrow$ Renaming to distinguish between them.

$$(:1) : a \rightarrow [a] \rightarrow [a]$$

$$(:2) : b \rightarrow [b] \rightarrow [b]$$

$$a := b \rightarrow [b] \rightarrow [b]$$

Substitute and remove first argument:

$$((:)) :: [(b \rightarrow [b] \rightarrow [b])] \rightarrow [(b \rightarrow [b] \rightarrow [b])]$$

0.2

d) $((==1) ==2) \rightarrow$ Renaming for distinction

$$(==1) :: \text{Eq } a \Rightarrow a \rightarrow a \rightarrow \text{Bool}$$

$$(==2) :: \text{Eq } b \Rightarrow b \rightarrow b \rightarrow \text{Bool}$$

$((==) ==) : \text{Type error because } (b \rightarrow b \rightarrow \text{Bool})$
is not of instance Eq.

D 0.2

$$e) ((!!) !!) :: \text{Bool} \rightarrow \text{Bool} \rightarrow \text{Bool} \rightarrow \text{Bool}$$

0

3. Law 1: $fmap\ id = id$

$$1. fmap\ f\ (Just\ x) = Just\ (f\ x)$$

$$fmap\ id\ (Just\ x) = Just\ (id\ x) = Just\ (x) = id\ Just\ x$$

$$2. fmap\ f\ Nothing = Nothing$$

$$fmap\ id\ Nothing = Nothing = id\ Nothing$$

Law 1 holds.

Law 2: $fmap\ (p.\ q) = (fmap\ p) . (fmap\ q)$

$$1. fmap\ f\ (Just\ x) = Just\ (f\ x)$$

$$fmap\ (p.\ q)\ (Just\ x) = Just\ ((p.\ q)\ x) = Just\ (p\ (q\ x))$$

$$\begin{aligned} & ((fmap\ p) . (fmap\ q))\ (Just\ x) = (fmap\ p)\ (fmap\ q\ (Just\ x)) \\ & = (fmap\ p)\ (Just\ (q\ x)) = Just\ (p\ (q\ x)) \end{aligned}$$

$$2. fmap\ f\ Nothing = Nothing$$

$$fmap\ (p.\ q)\ Nothing = Nothing$$

$$\begin{aligned} & ((fmap\ p) . (fmap\ q))\ Nothing = (fmap\ p)\ (fmap\ q\ Nothing) \\ & = fmap\ p\ Nothing = Nothing \end{aligned}$$

Law 2 holds.

(1)

Answer: Yes the definition is correct since both laws hold

4. Thunk is related to the concept of lazy evaluation where expressions are not evaluated until being used.

0.2

5. $e :: \text{Maybe } a \rightarrow \text{Maybe Bool}$

Because the expression in `do` has a `Nothing`, we can extract that the Monad being used is `Maybe`, which has values `Nothing` and `Just`.

Because a monad always returns a monad and we can see that the function returns `False`, we can therefore conclude that the output type is `Maybe Bool`.

1

6.

Ableton

- data Tree3 String = Leaf String | Node ~~String~~
(Tree3 String) (Tree3 String) (Tree3 String)
- data Tree3 a = Leaf a | Node ~~a~~ (Tree3 a)
(Tree3 a) (Tree3 a)
deriving Show
- instance Functor (Tree3 a) where
fmap f (Leaf x) = Leaf (f x)
fmap f (Node ~~x~~ x y z) = Node ~~(f x)~~ (fmap f x)
(fmap f x) (fmap f z)

0.9

EDAN40

examination

4 hp

2nd June 2022

14:00 - 19:00

WRITE ONLY ON ONE SIDE OF THE PAPER - the exams will be scanned in and only the front/odd pages will be read.

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PUT YOUR ID AND PAGE NUMBER ON EACH PAGE YOU SUBMIT - make sure that the amount of pages is equal to the amount you note on the front information page

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PRELIMINARY MAX AMOUNT OF POINTS: 6

Exam

1. Type derivation (1p)

- (a) Assume that the type of reduce is

```
reduce :: a -> a
```

Find the type of

```
prepare = reduce . words . map toLower . filter
                                              (not . flip
                                              elem ".,;:;!#%&|")
```

- (b) Given that

```
map2 :: (a -> b, c -> d) -> (a, c) -> (b, d)
```

find the destination type b of the following function:

```
rulesCompile :: [(String, [String])] -> b
rulesCompile = (map . map2) (words . map toLower, map words)
```

- (c) Given that

```
transformationApply :: Eq a => a -> ([a] -> [a]) -> [a] -> ([a], [a])
                                                         -> Maybe [a]
```

```
orElse :: Maybe a -> Maybe a -> Maybe a
```

find the type of

```
foldr1 orElse (map (transformationApply wildcard f) pats)
```

2. Proving program properties (2p)

The Functor class is defined as follows:

```
class Functor f where
  fmap :: (a -> b) -> f a -> f b
```

It is mandatory that all instances of Functor should obey:

```
fmap id      = id
fmap (p . q) = (fmap p) . (fmap q)
```

Let Either be defined as follows:

```
data Either a b = Left a | Right b
```

Assume the following definition of Either types as a functor instance:

```
instance Functor (Either a) where
  fmap f (Right x) = Right (f x)
  fmap f (Left x)  = Left x
```



```

{-A list of selected functions from the Haskell modules:
Prelude
Data.List
Data.Maybe
Data.Char -}

-----
-- standard type classes
class Show a where
  show :: a -> String

class Eq a where
  (==), (/=) :: a -> a -> Bool

class (Eq a) => Ord a where
  (<), (<=), (>=), (>) :: a -> a -> Bool
  max, min :: a -> a -> a

class (Eq a, Show a) => Num a where
  (+), (-), (*) :: a -> a -> a
  negate :: a -> a
  abs, signum :: a -> a
  fromInteger :: Integer -> a

class (Num a, Ord a) => Real a where
  toRational :: a -> Rational

class (Real a, Enum a) => Integral a where
  quot, rem :: a -> a -> a
  div, mod :: a -> a -> a
  toInteger :: a -> Integer

class (Num a) => Fractional a where
  (/) :: a -> a -> a
  fromRational :: Rational -> a

class (Fractional a) => Floating a where
  exp, log, sqrt :: a -> a
  sin, cos, tan :: a -> a

class (Real a, Fractional a) => RealFrac a where
  truncate, round :: (Integral b) => a -> b
  ceiling, floor :: (Integral b) => a -> b

-----
-- numerical functions
even, odd :: (Integral a) => a -> Bool
even n = n `rem` 2 == 0
odd = not . even

-----
-- monadic functions
sequence :: Monad m => [m a] -> m [a]
sequence = foldr mcons (return [])

where mcons p q = do x <- p; xs <- q; return (x:xs)

sequence_ :: Monad m => [m a] -> m ()
sequence_ xs = do sequence xs; return ()

-----
-- functions on functions
id :: a -> a
id x = x

const :: a -> b -> a
const x _ = x

(.,) :: (b -> c) -> (a -> b) -> a -> c
f . g = \x -> f (g x)

flip :: (a -> b -> c) -> b -> a -> c
flip f x y = f y x

($) :: (a -> b) -> a -> b
f $ x = f x

-----
-- functions on Booleans
data Bool = False | True

(&&), (||) :: Bool -> Bool -> Bool
True && x = x
False && _ = False
True || _ = True
False || x = x

not :: Bool -> Bool
not True = False
not False = True

-----
-- functions on Maybe
data Maybe a = Nothing | Just a

isJust :: Maybe a -> Bool
isJust (Just a) = True
isJust Nothing = False

isNothing :: Maybe a -> Bool
isNothing = not . isJust

fromJust :: Maybe a -> a
fromJust (Just a) = a

maybeToList :: Maybe a -> [a]
maybeToList Nothing = []
maybeToList (Just a) = [a]

```

```

listToMaybe :: [a] -> Maybe a
listToMaybe [] = Nothing
listToMaybe (a:_) = Just a

-- a hidden goodie
instance Monad [] where
  return x = [x]
  xs >=> f = concat (map f xs)

-- functions on pairs
fst :: (a, b) -> a
fst (x, y) = x
snd :: (a, b) -> b
snd (x, y) = y
curry :: ((a, b) -> c) -> a -> b -> c
curry f x y = f (x, y)
uncurry :: (a -> b -> c) -> (a, b) -> c
uncurry f p = f (fst p) (snd p)

-- functions on lists
map :: (a -> b) -> [a] -> [b]
map f xs = [ f x | x <- xs ]
(+++) :: [a] -> [a] -> [a]
xs +++ ys = foldr (:) ys xs
filter :: (a -> Bool) -> [a] -> [a]
filter p xs = [ x | x <- xs, p x ]
concat :: [[a]] -> [a]
concat xss = foldr (++) [] xss
concatMap :: (a -> [b]) -> [a] -> [b]
concatMap f = concat . map f
head, last :: [a] -> a
head (x:_) = x
last [x] = x
last (_:xs) = last xs
tail, init :: [a] -> [a]
tail (_:xs) = xs
init [x] = []
init (x:xs) = x : init xs

```

```

null :: [a] -> Bool
null [] = True
null (_:_) = False
length :: [a] -> Int
length [] = 0
length (_:_) = 1 + length l
(!!) :: [a] -> Int -> a
(x:_) !! 0 = x
(_:xs) !! n = xs !! (n-1)
foldr :: (a -> b -> b) -> b -> [a] -> b
foldr f z [] = z
foldr f z (x:xs) = f x (foldr f z xs)
foldl :: (a -> b -> a) -> a -> [b] -> a
foldl f z [] = z
foldl f z (x:xs) = foldl f (f z x) xs
iterate :: (a -> a) -> a -> [a]
iterate f x = x : iterate f (f x)
repeat :: a -> [a]
repeat x = xs where xs = x:xs
replicate :: Int -> a -> [a]
replicate n x = take n (repeat x)
cycle :: [a] -> [a]
cycle [] = error "Prelude.cycle: empty list"
cycle xs = xs' where xs' = xs++xs'
take, drop :: Int -> [a] -> [a]
take n _ | n <= 0 = []
take _ [] = []
take n (x:xs) = x : take (n-1) xs
drop n xs | n <= 0 = xs
drop _ [] = []
drop n (_:xs) = drop (n-1) xs
splitAt :: Int -> [a] -> ([a], [a])
splitAt n xs = (take n xs, drop n xs)
takeWhile, dropWhile :: (a -> Bool) -> [a] -> [a]
takeWhile p [] = []
takeWhile p (x:xs) = x : takeWhile p xs
  | p x
  | otherwise = []
dropWhile p [] = []
dropWhile p xs@(x:xs') = dropWhile p xs'
  | p x
  | otherwise = xs

```

```

lines, words :: String -> [String]
-- lines "apa\bepa\ncepa\n" == ["apa", "bepa", "cepa"]
-- words "apa bepa\n cepa" == ["apa", "bepa", "cepa"]

unlines, unwords :: [String] -> String
-- unlines ["apa", "bepa", "cepa"] == "apa\bepa\ncepa\n"
-- unwords ["apa", "bepa", "cepa"] == "apa bepa cepa"

and, or :: [Bool] -> Bool
and = foldr (&&) True
or = foldr (||) False

any, all :: (a -> Bool) -> [a] -> Bool
any p = or . map p
all p = and . map p

elem, notElem :: (Eq a) => a -> [a] -> Bool
elem x = any (== x)
notElem x = all (/= x)

lookup :: (Eq a) => a -> [(a,b)] -> Maybe b
lookup key [] = Nothing
lookup key ((x,y):xys)
    | key == x = Just y
    | otherwise = lookup key xys

sum, product :: (Num a) => [a] -> a
sum = foldl (+) 0
product = foldl (*) 1

maximum, minimum :: (Ord a) => [a] -> a
maximum [] = error "Prelude.maximum: empty list"
maximum xs = foldl1 max xs
minimum [] = error "Prelude.minimum: empty list"
minimum xs = foldl1 min xs

zip :: [a] -> [b] -> [(a,b)]
zip = zipWith (,)

zipWith :: (a -> b -> c) -> [a] -> [b] -> [c]
zipWith z (a:as) (b:bs)
    = z a b : zipWith z as bs
zipWith _ _ _ = []

unzip :: [(a,b)] -> ([a],[b])
unzip = foldr \((a,b) ~ (as,bs)) -> (a:as,b:bs) ([],[])

nub :: [a] -> [a]
nub [] = []
nub (x:xs) = x : nub [ y | y <- xs, x /= y ]

delete :: Eq a => a -> [a] -> [a]
delete y [] = []

```

```

delete y (x:xs) = if x == y then xs else x : delete y xs

(\\) :: Eq a => [a] -> [a] -> [a]
(\\) = foldl (flip delete)

union :: Eq a => [a] -> [a] -> [a]
union xs ys = xs ++ ( ys \\ xs )

intersect :: Eq a => [a] -> [a] -> [a]
intersect xs ys = [ x | x <- xs, x `elem` ys ]

intersperse :: Int -> [a] -> [a]
intersperse 0 [1,2,3,4] == [1,0,2,0,3,0,4]

transpose :: [[a]] -> [[a]]
transpose [[1,2,3],[4,5,6]] == [[1,4],[2,5],[3,6]]

partition :: (a -> Bool) -> [a] -> ([a],[a])
partition p xs = (filter p xs, filter (not . p) xs)

group :: Eq a => [a] -> [[a]]
-- group "aapaaabbbccc" == ["aa","p","aa","bbb","ccc"]

isPrefixOf :: [a] -> [a] -> Bool
isPrefixOf [] = True
isPrefixOf [ ] = False
isPrefixOf (x:xs) (y:ys) = x == y && isPrefixOf xs ys

isSuffixOf :: [a] -> [a] -> Bool
isSuffixOf xs y = reverse x `isPrefixOf` reverse y

sort :: [a] -> [a]
sort = foldr insert []

insert :: [a] -> [a] -> [a]
insert x [] = [x]
insert x (y:xs) = if x <= y then x:y:xs else y:insert x xs

-- functions on Char

type String = [Char]

toupper, tolower :: Char -> Char
-- toupper 'a' == 'A'
-- tolower 'Z' == 'z'

digitToInt :: Char -> Int
-- digitToInt '8' == 8

intToDigit :: Int -> Char
-- intToDigit 3 == '3'

ord :: Char -> Int
chr :: Int -> Char

```

Q1 (a). For simplicity, denote

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- $t1 \ u = \text{elem } u \text{ ".,; * ! \% \& |"}$
- $t2 = \text{flip elem ".,; * ! \% \& |"}$
- $t3 = \text{filter (not . } t2)$
- $t4 = \text{map } \underline{\text{toLower . } t3}$

$\therefore \text{elem} :: (\text{Eq } a) \Rightarrow a \rightarrow [a] \rightarrow \text{Bool}$

\therefore For $t1 \ u$, since type of a is Char , u has the type Char .

$\therefore \text{flip} :: (a \rightarrow b \rightarrow c) \rightarrow b \rightarrow a \rightarrow c$

$\text{flip elem} :: (\text{Eq } a) \Rightarrow [a] \rightarrow a \rightarrow \text{Bool}$

$\therefore t2 :: \text{Char} \rightarrow \text{Bool} \quad (\text{Char is in class Eq})$

$\therefore \text{not . } t2 :: \text{Char} \rightarrow \text{Bool}$

$\therefore t3 :: [\text{Char}] \rightarrow [\text{Char}]$

$\therefore \text{toLower} :: \text{Char} \rightarrow \text{Char}$

$\therefore t4 :: [\text{Char}] \rightarrow [\text{Char}]$, or $\text{String} \rightarrow \text{String}$.

$\therefore \text{words} :: \text{String} \rightarrow [\text{String}]$

$\therefore \text{words . } t4 :: \text{String} \rightarrow [\text{String}]$

$\therefore \text{prepare} :: \text{String} \rightarrow [\text{String}]$.

Q1 (b). $\therefore \text{words . map toLower} :: \text{String} \rightarrow [\text{String}]$

$\text{map words} :: [\text{String}] \rightarrow [[\text{String}]]$.

$\text{map . map2} :: (a \rightarrow b, c \rightarrow d) \rightarrow [(a, c)] \rightarrow [(b, d)]$

$\therefore \text{rulesCompile} :: ((\text{String}, [\text{String}]) \rightarrow [([\text{String}], [\text{String}])])$

$\vdash :: \Gamma \vdash [\text{String}] \vdash \Gamma \vdash [\text{String}] \vdash$

Q1. (c). For simplicity, denote

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$t = \text{transformationApply wildcard } f \ x$

(assume wildcard, f , x are fixed so that t does not depend on them)

Then $t :: \text{Eq } a \Rightarrow ([a], [a]) \rightarrow \text{Maybe } [a]$.

So, $\text{map } t \text{ pats} :: \text{Eq } a \Rightarrow [\text{Maybe } [a]]$

$\therefore \text{foldr1} :: (a \rightarrow b \rightarrow b) \rightarrow [a] \rightarrow b$.

$\therefore \text{foldr1 orElse (map } t \text{ pats)} :: \text{Eq } a \Rightarrow \text{Maybe } [a]$.

Q2. Proof 1: $\text{fmap id} = \text{id}$

Suppose $x = \text{Left } y$, then

$$\begin{aligned}\text{fmap id } x &= \text{fmap id } (\text{Left } y) \\ &= \text{Left } y \\ &= x\end{aligned}$$

Suppose $x = \text{Right } y$, then

$$\begin{aligned}\text{fmap id } x &= \text{fmap id } (\text{Right } y) \\ &= \text{Right } (\text{id } y) \\ &= \text{Right } y \\ &= x.\end{aligned}$$

Hence $\text{fmap id} = \text{id}$.

Proof 2: $\text{fmap } (p \cdot q) = (\text{fmap } p) \cdot (\text{fmap } q)$. We let $f' = p \cdot q$.

Suppose $x = \text{Left } y$, then

$$\text{fmap } f' (\text{Left } y) = \text{Left } y$$

Also, we have

$$\begin{aligned}&((\text{fmap } p) \cdot (\text{fmap } q))(\text{Left } y) \\ &= \text{fmap } p (\text{fmap } q (\text{Left } y)) \\ &= \text{fmap } p (\text{Left } y) \\ &= \text{Left } y.\end{aligned}$$

Suppose $x = \text{Right } y$, then.

$$\text{fmap } f' (\text{Right } y) = \text{Right } (f' y)$$

Also, we have

$$\begin{aligned}&((\text{fmap } p) \cdot (\text{fmap } q)) (\text{Right } y) \\ &= \text{fmap } p (\text{fmap } q (\text{Right } y)) \\ &= \text{fmap } p (\text{Right } (q y)) \\ &= \text{Right } (p (q y)) \\ &= \text{Right } ((p \cdot q) y) = \text{Right } (f' y)\end{aligned}$$

Q2. cont'd) Since the definition of the functor instance obeys the two laws suggested, the functor instance is correct. —

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Q3. - Spark is the potential of thread creation.

- It occurs in Haskell when users call 'par' function for parallel programming
- Spark is good for its lazy evaluation. Since it only evaluates the first argument in 'par' function to WHNF, the program does not require to evaluate the exact value of the first argument. This is helpful when exact value is hard to be evaluated, say, an infinite list results in the argument

Q4. (a).

paste reg im1 im2 = \pos → if (reg pos) then (im1 pos) else (im2 pos)

~~reg pos = im1 pos~~
~~otherwise = im2 pos~~

Q4 (b). Define ~~inRange :: Region~~ such that

~~inRange (x, y) = (inUnit x) && (inUnit y)~~
~~where inUnit z = (z >= 0) && (z <= 1)~~

lift0 val = const val

lift1 func im1 = func . im1

lift2 func im1 im2 = \p → func (im1 p) (im2 p)

Q4(c). The operator (-) is in class Num, which is the child of Eq and Show classes. Hence before declaring the operator (-), we have to first declare

- show :: Image a → String (in Show Class; instance)
- (==) :: Image a → Image a → Bool (in Eq Class instance)

Then we can declare (-) :: Image a → Image a → Image a in Num Class instance. which was made use of lift2 function in Q4(b).

Q5.

Rewrite the function as

$$e\ k = k \gg= \lambda x \rightarrow \text{Nothing} \gg= \lambda y \rightarrow \text{return } 42$$

Recall $(\gg=) :: (\text{Monad } m) \Rightarrow m\ a \rightarrow (a \rightarrow m\ b) \rightarrow m\ b$ - Focus on $\text{Nothing} \gg= \lambda y \rightarrow \text{return } 42$ $\therefore \text{Nothing} :: \text{Maybe } a$ $\therefore \text{return } 42 :: (\text{Num } b) \Rightarrow \text{Maybe } b$ So the type of this subfunction is $(\text{Num } b) \Rightarrow \text{Maybe } b$.However since the first argument is Nothing , the function of the 2nd argument is not conducted, hence it returns Nothing .Hence $e\ k = k \gg= \lambda x \rightarrow \text{Nothing}$.- Hence $k :: \text{Maybe } a$ As a result, $e :: (\text{Num } b) \Rightarrow \text{Maybe } a \rightarrow \text{Maybe } b$, and it returns Nothing for all $k :: \text{Maybe } a$.

EDAN40

examination

4 hp

18th August 2022

8:00 - 13:00

WRITE ONLY ON ONE SIDE OF THE PAPER - the exams will be scanned in and only the front/odd pages will be read.

DO NOT WRITE WITH OTHER COLOUR THAN BLACK OR DARK BLUE - lightly coloured text may disappear during scanning

PUT YOUR ID AND PAGE NUMBER ON EACH PAGE YOU SUBMIT - make sure that the amount of pages is equal to the amount you note on the front information page

WRITE CLEARLY - if we cannot read you we cannot (properly) grade you.

PRELIMINARY MAX AMOUNT OF POINTS: 6

Exam

1. Proving program properties (1.5p)

Given the following function:

```
foo :: [a] -> [a]
foo []    = []
foo (x:xs) = foo xs ++ [x]
```

prove that the following holds:

```
foo (foo xs) = xs
```

2. Type derivation (1p)

Find types of: a) $(: (.))$, b) $((: (.) .)$, c) $(. (:))$, d) $((. (:) .)$.

3. Types and type classes (1.5p)

- (0.3p) Define a tree data structure so that the trees are ternary (i.e., each node has either three children or is a leaf) and store strings in each node.
- (0.2p) Generalize your definition so that your ternary trees can contain objects of an arbitrary predetermined type in a node.
- (1p) Assuming your polymorphic trees type is denoted by `Tree3 a`, write all necessary code so that the following function is correct:

```
myLength :: Tree3 String -> Tree3 Integer
myLength = fmap length
```

and yields a ternary tree with nodes containing lengths of the strings placed in the respective nodes of the argument tree.

Make sure that the following works as well:

```
myReverse :: Tree3 String -> Tree3 String
myReverse = fmap reverse
```

4. Do notation (1p)

Given the following function:

```
f x y = do
  a <- x
  b <- y
  return (a*b)
```

What is the type of f?
 What is the value of f [1,2,3] [2,4,8] ?
 What is the value of f (Just 5) Nothing ?
 Is the expression fmap (+2) (Just 5) correct?
 What is the type of expression return 5?

5. Memoization (1p)

Consider the following two versions of similarity score computations. The difference is in the expression defining value for simEntry i j .

- (a) (0.1p) Which of the versions is much faster than the other?
 (b) (0.9p) Why?

VERSION 1:

```
similScore :: String -> String -> Int
similScore xs ys = simScore (length xs) (length ys)
  where
    simScore i j = simTable!!i!!j
    simTable = [[ simEntry i j | j<-[0..]] | i<-[0..] ]
    simEntry :: Int -> Int -> Int
    simEntry 0 0 = 0
    simEntry i 0 = (i * scoreSpace)
    simEntry 0 j = (scoreSpace * j)
    simEntry i j = maximum [((simScore (i-1) (j-1)) + (score x y)),
                             ((simScore (i-1) j) + (score x '-')),
                             ((simScore i (j-1)) + (score '-' y))]
```

where
 x = xs!!(i-1)
 y = ys!!(j-1)

VERSION 2:

```
similScore :: String -> String -> Int
similScore xs ys = simScore (length xs) (length ys)
  where
    simScore i j = simTable!!i!!j
    simTable = [[ simEntry i j | j<-[0..]] | i<-[0..] ]
    simEntry :: Int -> Int -> Int
    simEntry 0 0 = 0
    simEntry i 0 = (i * scoreSpace)
    simEntry 0 j = (scoreSpace * j)
    simEntry i j = maximum [((simEntry (i-1) (j-1)) + (score x y)),
                             ((simEntry (i-1) j) + (score x '-')),
                             ((simEntry i (j-1)) + (score '-' y))]
```

where
 x = xs!!(i-1)
 y = ys!!(j-1)

Good Luck!

```

{-A list of selected functions from the Haskell modules:
Prelude
Data.List
Data.Maybe
Data.Char -}

-----
-- standard type classes
class Show a where
  show :: a -> String

class Eq a where
  (==), (/=) :: a -> a -> Bool

class (Eq a) => Ord a where
  (<), (<=), (>=), (>) :: a -> a -> Bool
  max, min :: a -> a -> a

class (Eq a, Show a) => Num a where
  (+), (-), (*) :: a -> a -> a
  negate :: a -> a
  abs, signum :: a -> a
  fromInteger :: Integer -> a

class (Num a, Ord a) => Real a where
  toRational :: a -> Rational

class (Real a, Enum a) => Integral a where
  quot, rem :: a -> a -> a
  div, mod :: a -> a -> a
  toInteger :: a -> Integer

class (Num a) => Fractional a where
  (/) :: a -> a -> a
  fromRational :: Rational -> a

class (Fractional a) => Floating a where
  exp, log, sqrt :: a -> a
  sin, cos, tan :: a -> a

class (Real a, Fractional a) => RealFrac a where
  truncate, round :: (Integral b) => a -> b
  ceiling, floor :: (Integral b) => a -> b

-----
-- numerical functions
even, odd :: (Integral a) => a -> Bool
even n = n `rem` 2 == 0
odd = not . even

-----
-- monadic functions
sequence :: Monad m => [m a] -> m [a]
sequence = foldr mcons (return [])

```

```

                                where mcons p q = do x <- p; xs <- q; return (x:xs)

sequence_ :: Monad m => [m a] -> m ()
sequence_ xs = do sequence xs; return ()

-----
-- functions on functions
id :: a -> a
id x = x

const :: a -> b -> a
const x _ = x

(.) :: (b -> c) -> (a -> b) -> a -> c
f . g = \x -> f (g x)

flip :: (a -> b -> c) -> b -> a -> c
flip f x y = f y x

($) :: (a -> b) -> a -> b
f $ x = f x

-----
-- functions on Booleans
data Bool = False | True

(&&), (||) :: Bool -> Bool -> Bool
True && x = x
False && _ = False
True || _ = True
False || x = x

not :: Bool -> Bool
not True = False
not False = True

-----
-- functions on Maybe
data Maybe a = Nothing | Just a

isJust :: Maybe a -> Bool
isJust (Just a) = True
isJust Nothing = False

isNothing :: Maybe a -> Bool
isNothing = not . isJust

fromJust :: Maybe a -> a
fromJust (Just a) = a

maybeToJust :: Maybe a -> [a]
maybeToJust Nothing = []
maybeToJust (Just a) = [a]

```

```

listToMaybe :: [a] -> Maybe a
listToMaybe [] = Nothing
listToMaybe (a:_) = Just a

-- a hidden goodie
instance Monad [] where
  return x = [x]
  xs >= f = concat (map f xs)

-- functions on pairs
fst :: (a, b) -> a
fst (x, y) = x
snd :: (a, b) -> b
snd (x, y) = y
curry :: ((a, b) -> c) -> a -> b -> c
curry f x y = f (x, y)
uncurry :: (a -> b -> c) -> (a, b) -> c
uncurry f p = f (fst p) (snd p)

-- functions on lists
map :: (a -> b) -> [a] -> [b]
map f xs = [ f x | x <- xs ]
map f xs ++ ys = foldr (:) ys xs
foldr (:) ys xs
filter p xs :: (a -> Bool) -> [a] -> [a]
filter p xs = [ x | x <- xs, p x ]
concat :: [[a]] -> [a]
concat xss = foldr (++) [] xss
concatMap :: (a -> [b]) -> [a] -> [b]
concatMap f = concat . map f
head, last :: [a] -> a
head (x:_) = x
last [x] = x
last (_:xs) = last xs
tail, init :: [a] -> [a]
tail (_:xs) = xs
init [x] = []
init (x:xs) = x : init xs

```

```

null :: [a] -> Bool
null [] = True
null (_:_) = False
length :: [a] -> Int
length [] = 0
length (_:_) = 1 + length l
(!!) :: [a] -> Int -> a
(x:_) !! 0 = x
(_:xs) !! n = xs !! (n-1)
foldr f z [] = z
foldr f z (x:xs) = f x (foldr f z xs)
foldl f z [] = z
foldl f z (x:xs) = foldl f (f z x) xs
iterate f x :: (a -> a) -> a -> [a]
iterate f x = x : iterate f (f x)
repeat x :: a -> [a]
repeat x = xs where xs = x:xs
replicate n x :: Int -> a -> [a]
replicate n x = take n (repeat x)
cycle [] = []
cycle xs = xs' where xs' = xs++xs'
take, drop :: Int -> [a] -> [a]
take n [] = []
take n _ | n <= 0 = []
take n (x:xs) = x : take (n-1) xs
drop n xs | n <= 0 = xs
drop n [] = []
drop n (_:xs) = drop (n-1) xs
splitAt n xs :: Int -> [a] -> ([a], [a])
splitAt n xs = (take n xs, drop n xs)
takeWhile, dropWhile :: (a -> Bool) -> [a] -> [a]
takeWhile p [] = []
takeWhile p (x:xs) = x : takeWhile p xs
dropWhile p [] = []
dropWhile p (x:xs) = x : dropWhile p xs
dropWhile p [] = []
dropWhile p xs@(x:xs') = dropWhile p xs'

```

```

lines, words      :: String -> [String]
-- lines "apa\bepea\ncepa\n" == ["apa","bepea","cepa"]
-- words "apa bepea\ncepa" == ["apa","bepea","cepa"]

unlines, unwords :: [String] -> String
-- unlines ["apa","bepea","cepa"] == "apa\bepea\ncepa"
-- unwords ["apa","bepea","cepa"] == "apa bepea cepa"

and, or           :: [Bool] -> Bool
and               = foldr (&&) True
or                = foldr (||) False

any, all          :: (a -> Bool) -> [a] -> Bool
any p             = or . map p
all p             = and . map p

elem, notElem    :: (Eq a) => a -> [a] -> Bool
elem x            = any (== x)
notElem x         = all (/= x)

lookup key []    :: (Eq a) => a -> [(a,b)] -> Maybe b
lookup key ((x,y):xys)
  | key == x     = Just y
  | otherwise    = lookup key xys

sum, product     :: (Num a) => [a] -> a
sum              = foldl (+) 0
product          = foldl (*) 1

maximum, minimum :: (Ord a) => [a] -> a
maximum []       = error "Prelude.maximum: empty list"
maximum xs       = foldl1 max xs
minimum []       = error "Prelude.minimum: empty list"
minimum xs       = foldl1 min xs

zip             :: [a] -> [b] -> [(a,b)]
zip             = zipWith (,)

zipWith         :: (a -> b -> c) -> [a] -> [b] -> [c]
zipWith z (a:as) (b:bs)
  = z a b : zipWith z as bs
zipWith _ _ _   = []

unzip           :: [(a,b)] -> ([a],[b])
unzip           = foldr \((a,b) ~(as,bs)) -> (a:as,b:bs) ([],[])

nub            :: [a] -> [a]
nub []          = []
nub (x:xs)      = x : nub [ y | y <- xs, x /= y ]

delete         :: Eq a => a -> [a] -> [a]
delete y []     = []

```

```

delete y (x:xs) = if x == y then xs else x : delete y xs

(\\)            :: Eq a => [a] -> [a] -> [a]
(\\)            = foldl (flip delete)

union           :: Eq a => [a] -> [a] -> [a]
union xs ys     = xs ++ ( ys \\ xs )

intersect       :: Eq a => [a] -> [a] -> [a]
intersect xs ys = [ x | x <- xs, x `elem` ys ]

intersperse    :: a -> [a] -> [a]
-- intersperse 0 [1,2,3,4] == [1,0,2,0,3,0,4]
:: a -> [a] -> [a]

transpose      :: [[a]] -> [[a]]
-- transpose [[1,2,3],[4,5,6]] == [[1,4],[2,5],[3,6]]

partition      :: (a -> Bool) -> [a] -> ([a],[a])
partition p xs  = (filter p xs, filter (not . p) xs)

group          :: Eq a => [a] -> [[a]]
-- group "aapabbbee" == ["aa","p","aa","bbb","eee"]

isPrefixOf, isSuffixOf :: Eq a => [a] -> [a] -> Bool
isPrefixOf [] _      = True
isPrefixOf _ []      = False
isPrefixOf (x:xs) (y:ys) = x == y && isPrefixOf xs ys

isSuffixOf x y      = reverse x `isPrefixOf` reverse y

sort            :: (Ord a) => [a] -> [a]
sort            = foldr insert []

insert         :: (Ord a) => a -> [a] -> [a]
insert x []      = [x]
insert x (y:xs)  = if x <= y then x:y:xs else y:insert x xs

-----
-- functions on Char

type String = [Char]

toUpper, toLower :: Char -> Char
-- toUpper 'a' == 'A'
-- toLower 'Z' == 'z'

digitToInt :: Char -> Int
-- digitToInt '8' == 8

intToDigit :: Int -> Char
-- intToDigit 3 == '3'

ord :: Char -> Int
chr :: Int -> Char

```

220018

(1)

(\$)

N40-0017-SHA

$$\text{foo} :: [a] \rightarrow [a]$$

$$\text{foo} [] = []$$

$$\text{foo} (x:xs) = \text{foo} xs ++ [x]$$

Induction:Base case: $xs = []$

$$\text{foo} (\text{foo} xs) = \text{foo} (\text{foo} []) = \text{foo} [] = [] = xs \quad 0.3$$

Hypothesis: Assume that for xs , $\text{foo} (\text{foo} xs) = xs$ Now we need to prove that if the hypothesis holds, that implies that the property holds for $(x:xs)$

Answer:

Induction step:

$$\begin{aligned} \text{foo} (\text{foo} (x:xs)) &= \text{foo} (\text{foo} xs ++ [x]) \quad \text{this step needs additional proof!} \\ &\quad \text{= xs through hypothesis} \\ &= \text{foo} [x] ++ \underbrace{\text{foo} (\text{foo} xs)}_{= xs} = \text{foo} [x] ++ xs = ([] ++ [x]) ++ xs = \\ &= [x] ++ xs = (x:xs) \end{aligned}$$

Because we have showed that if

$$\text{foo} (\text{foo} xs) = xs \implies \text{foo} (\text{foo} (x:xs))$$

and that the base case holds, we have shown that

$$\underline{\text{foo} (\text{foo} xs) = xs} \quad 0.6$$

(2)

$$a) (.) :: (b \rightarrow c) \rightarrow (a \rightarrow b) \rightarrow a \rightarrow c$$

$$(:) :: a \rightarrow [a] \rightarrow [a]$$

(\$)

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Answer: The second argument of $(:)$ must be a list, so $(:) \text{ is not a list, } (: (.)) \text{ will result in an error.}$

0.25

b) ...

$$\text{Answer} = ((:)) :: (a1 \rightarrow a2) \rightarrow a1 \rightarrow [a2] \rightarrow [a2]$$

0.25

c) ...

$$\text{Answer: } (.(:)) :: ([a] \rightarrow [a] \rightarrow c) \rightarrow a \rightarrow c$$

0.25

d)

$$\text{Answer: } ((.:)) :: [(b \rightarrow c) \rightarrow (a \rightarrow b) \rightarrow a \rightarrow c] \rightarrow$$

$$\rightarrow [(b \rightarrow c) \rightarrow (a \rightarrow b) \rightarrow a \rightarrow c]$$

0.25

3.

(3)

a) data Tree = Leaf String | Branch String Tree Tree Tree

(b)

N40-0017-SHA

0.3

0.2

b) data Tree3 a = Leaf a | Branch a (Tree3 a) (Tree3 a) (Tree3 a)

c) instance Functor Tree3 where

fmap f (Leaf a) = Leaf (f a)

fmap f (Branch a x y z) = Branch (f a) (fmap f x) (fmap f y) (fmap f z)

Answer: When we make Tree3 a functor and implement fmap both myLength and myReverse will work as intended.

1

$$9. \quad \frac{f}{f} \times \frac{g}{g} = do$$

$$a \leftarrow x$$

$$c \leftarrow \frac{g}{g}$$

$$do\ m.\ (a * b)$$

(4)

(8)

N40-0017-SHA

$$a) \quad f :: (Morad\ m, Num\ b) \Rightarrow m\ b \rightarrow m\ b \rightarrow m\ b$$

0.2

b) The value of $f\ [2,2,3]\ [2,4,8]$ is:

$$[2, 8, 8, 4, 8, 16, 6, 12, 24]$$

0.2

c) The value of f (Just 5) Nothing is:

Nothing

0.2

d) The expression is correct, it evaluates to (Just 7)

0.2

$$e) \quad (return\ 5) :: (Morad\ m, Num\ a) \Rightarrow m\ a$$

0.2

5. Memoization

(5)

(S)
N90-0017-SHA

a) The first version is much faster

b) The first version uses the tabularized values from SumScore and makes use of previous computations to reduce the execution time.

The second version, on the other hand, doesn't use previous computations and instead just recursively tries to compute the value for each i . It has more computations and is slower.