

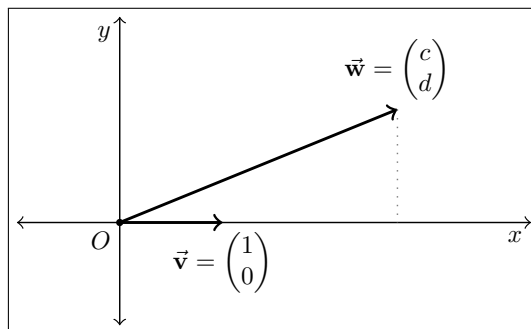
1 Inner products

We describe points in 2D-space by their x - and y -coordinate. For example: $\begin{pmatrix} a \\ b \end{pmatrix}$ is the point with x -coordinate a and y -coordinate b . If we have two points $\vec{v} = \begin{pmatrix} a \\ b \end{pmatrix}$ and $\vec{w} = \begin{pmatrix} c \\ d \end{pmatrix}$, the “inner product” of \vec{v} and \vec{w} is:

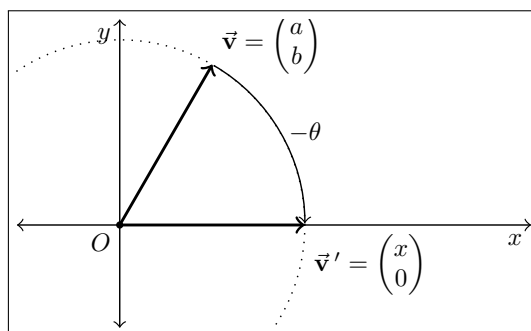
$$\vec{v} \cdot \vec{w} = a \times c + b \times d.$$

In the exercises below, we will see why this definition is nice and meaningful.

- (a) Simplify the inner product in the special case when the first vector is $\vec{v} = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$. Then: $\begin{pmatrix} 1 \\ 0 \end{pmatrix} \cdot \begin{pmatrix} c \\ d \end{pmatrix} = \dots$



- (b) How does this relate to \vec{w} ? If $\vec{v} \cdot \vec{w} = 0$, what is the angle between \vec{v} and \vec{w} ?
- (c) Now look at the inner product with $\vec{v} = \begin{pmatrix} a \\ 0 \end{pmatrix}$. Compute: $\frac{\vec{v} \cdot \vec{w}}{\vec{v} \cdot \vec{v}} \times \vec{v} = \begin{pmatrix} \dots \\ 0 \end{pmatrix}$.¹
- (d) We can rotate a point $\vec{v} = \begin{pmatrix} a \\ b \end{pmatrix}$ counter-clockwise by an angle of θ as follows: $\vec{v}' = \begin{pmatrix} a \cos(\theta) - b \sin(\theta) \\ a \sin(\theta) + b \cos(\theta) \end{pmatrix}$.
- (1) Verify the equation for $\theta = 90^\circ$. Then: $\vec{v}' = \begin{pmatrix} \dots \\ \dots \end{pmatrix}$.
- (2) Write \vec{w}' for the rotation of \vec{w} by an angle of θ . Show: $\vec{v} \cdot \vec{w} = \vec{v}' \cdot \vec{w}'$.²



- (e) Describe the angle θ (using a, b and $\sin / \cos / \tan$) that rotates $\vec{v} = \begin{pmatrix} a \\ b \end{pmatrix}$ to $\vec{v}' = \begin{pmatrix} x \\ 0 \end{pmatrix}$ for some $x > 0$. What is x ?
- (f) The length of a vector $\vec{v} = \begin{pmatrix} a \\ b \end{pmatrix}$ is given by $\|\vec{v}\| = \sqrt{a^2 + b^2}$. Prove: $\|\vec{v}\|^2 = \vec{v} \cdot \vec{v}$.
- (g) Suppose the angle between \vec{v} and \vec{w} is θ . Prove (using the previous exercises):³

$$\vec{v} \cdot \vec{w} = \cos(\theta) \times \|\vec{v}\| \times \|\vec{w}\|.$$

¹Note: this fraction is a number. If z is some number, then $z \times \begin{pmatrix} a \\ b \end{pmatrix} = \begin{pmatrix} za \\ zb \end{pmatrix}$.

²Hint: recall $\cos^2(\theta) + \sin^2(\theta) = 1$.

³Hint: use (5) to rotate \vec{w} with the same angle that rotates \vec{v} onto the x -axis.

2 Parallelograms

Recall that a basis for a lattice gives a tiling of space.

For (a), (b) and (c), assume \vec{v} is on the x -axis. First, we relate the area of the parallelogram to the basis.

- Show: the parallelogram $O, \vec{v}, \vec{v} + \vec{w}, \vec{w}$, and the rectangle O, A, \vec{w}, C have equal area.⁴
- Show: replacing \vec{w} by $\vec{w} + z\vec{v}$ does not change the area for any z (the area is independent of c).
- What is the area in terms of x and d ?

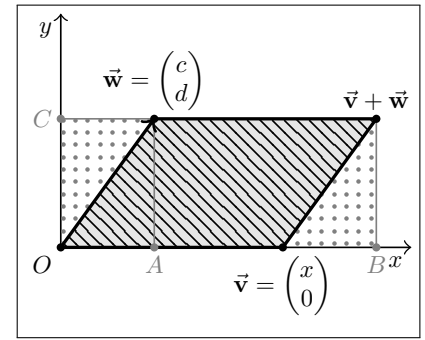


Figure 1: One tile

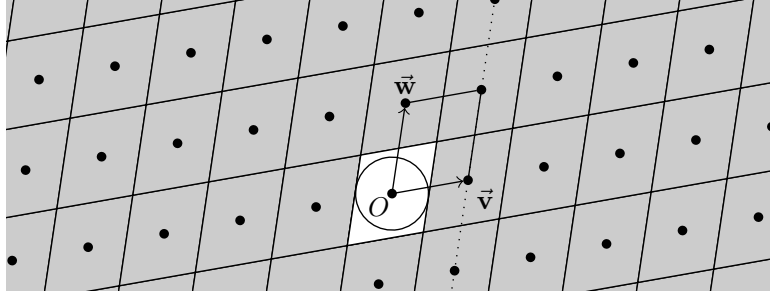


Figure 2: Tiling of 2D-space by the lattice basis $[\vec{v}, \vec{w}]$

Now consider the general case of $\vec{v} = \begin{pmatrix} a \\ b \end{pmatrix}$ and $\vec{w} = \begin{pmatrix} c \\ d \end{pmatrix}$.

- Using (b), slide \vec{v} to a point on the x -axis. Show: the parallelogram has area $ad - bc$.⁵
- How does the area change when you:
 - swap \vec{v} and \vec{w} ?
 - replace \vec{w} by $\vec{w} + \vec{v}$?
 - multiply \vec{v} and \vec{w} by 2?
- Conclude from (1) and (2) that the area stays the same during lattice reduction.
- Show: the circle in Figure 3 has radius $r = \frac{1}{2}\|\vec{v}\|$.
Show: the circle in Figure 2 has radius $r' = \frac{1}{2}\|\vec{v}\| \times \sin(\theta) \leq r$, where θ is the angle between \vec{v} and \vec{w} .

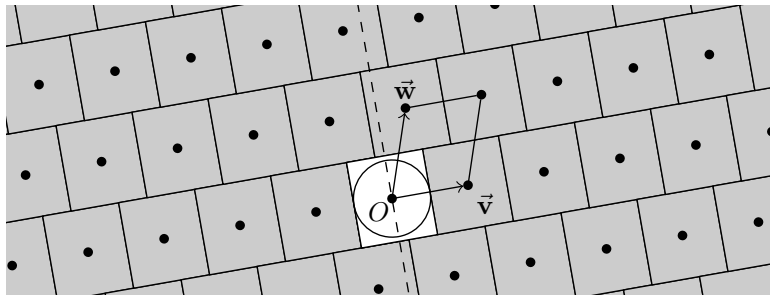


Figure 3: A different tiling of 2D-space, using a rectangular tile.

Encryption with lattices. With the secret key (the good basis $[\vec{v}, \vec{w}]$) only you can decrypt ciphertexts sent by others! Namely, a ciphertext is $\vec{t} = \vec{m} + \vec{e}$, where $\vec{m} = a\vec{v} + b\vec{w}$ (a, b integer) corresponds to a message, and \vec{e} is a short vector. With the secret key, you can find the tile containing \vec{t} and recover \vec{m} , if $\|\vec{e}\|$ is small enough.

- Show: decrypting with the tile in Fig. 2 works if $\|\vec{e}\| \leq r'$, and with the tile in Fig. 3, it works if $\|\vec{e}\| \leq r$. Which of the two tilings works best at decryption?

⁴Hint: go from one to the other by adding a triangle and removing a triangle (of same area).

⁵Hint: first determine for which value of z , the point $\vec{v} - z\vec{w}$ is on the x -axis. Then determine the x -coordinate of this point.

3 Gram–Schmidt and projections

Instead of 2D-space (\mathbb{R}^2), in this exercise we will look at n -dimensional space (\mathbb{R}^n). Now, vectors are described by n coordinates:

$$\vec{v} = \begin{pmatrix} v_1 \\ v_2 \\ \dots \\ v_n \end{pmatrix}$$

We still have the inner product, which is: $\vec{v} \cdot \vec{w} = v_1 \times w_1 + v_2 \times w_2 + \dots v_n \times w_n$.

- (a) Prove the following, where $\vec{u}, \vec{v}, \vec{w}$ are vectors and z is any number:

$$\vec{u} \cdot (\vec{v} + \vec{w}) = (\vec{u} \cdot \vec{v}) + (\vec{u} \cdot \vec{w}), \quad \vec{v} \cdot (z \times \vec{w}) = z \times (\vec{v} \cdot \vec{w}).$$

- (b) There is no simple “division” operation with vectors, because you cannot express \vec{w} in terms of \vec{v} whenever it is not pointing in the same direction. Still, using (c) from sheet 1, we can make a vector \vec{w} *orthogonal* to \vec{v} by computing the following:

$$\pi_{\vec{v}}(\vec{w}) = \vec{w} - \frac{\vec{v} \cdot \vec{w}}{\vec{v} \cdot \vec{v}} \times \vec{v}.$$

Show: $\pi_{\vec{v}}(\vec{w})$ satisfies $\vec{v} \cdot \pi_{\vec{v}}(\vec{w}) = 0$ (use (a)).

- (c) Consider the following process, called “Gram–Schmidt orthogonalization”:

$$\vec{b}_1^* = \vec{b}_1, \quad \vec{b}_2^* = \pi_{\vec{b}_1^*}(\vec{b}_2), \quad \vec{b}_3^* = \pi_{\vec{b}_2^*}(\pi_{\vec{b}_1^*}(\vec{b}_3)), \quad \dots$$

Show: \vec{b}_3^* is orthogonal to both \vec{b}_1^* and \vec{b}_2^* . (in general show: $\vec{b}_i^* \cdot \vec{b}_j^* = 0$ for all $1 \leq i < j \leq n$)

- (d) For finding short basis vectors, we want to reduce \vec{b}_3 by removing (integer) multiples of \vec{b}_1 and \vec{b}_2 that make \vec{b}_3 as short as possible. Considering the tiling in Figure 3, we want that \vec{b}_3 (when ignoring the third dimension) is in the tile containing the origin O , by removing some multiples of $\vec{b}_1 = \vec{v}$ and $\vec{b}_2 = \vec{w}$. This tile containing O is given by the points \vec{x} for which $-\frac{1}{2} \leq \frac{\vec{b}_1^* \cdot \vec{x}}{\vec{b}_1^* \cdot \vec{b}_1^*} \leq \frac{1}{2}$ and similarly for \vec{b}_2^* .

Write $\lceil x \rceil$ for rounding a number to its nearest integer. This means $-\frac{1}{2} \leq x - \lceil x \rceil \leq \frac{1}{2}$. Given \vec{b}_3 , show:

$$\vec{x} = \vec{b}_3 - \left\lceil \frac{\vec{b}_2^* \cdot \vec{b}_3}{\vec{b}_2^* \cdot \vec{b}_2^*} \right\rceil \vec{b}_2,$$

satisfies $-\frac{1}{2} \leq \frac{\vec{b}_2^* \cdot \vec{x}}{\vec{b}_2^* \cdot \vec{b}_2^*} \leq \frac{1}{2}$.

Given \vec{x} as above, show:

$$\vec{y} = \vec{x} - \left\lceil \frac{\vec{b}_1^* \cdot \vec{x}}{\vec{b}_1^* \cdot \vec{b}_1^*} \right\rceil \vec{b}_1,$$

is in the tile containing O , that is it satisfies: $-\frac{1}{2} \leq \frac{\vec{b}_1^* \cdot \vec{y}}{\vec{b}_1^* \cdot \vec{b}_1^*} \leq \frac{1}{2}$, and $-\frac{1}{2} \leq \frac{\vec{b}_2^* \cdot \vec{y}}{\vec{b}_2^* \cdot \vec{b}_2^*} \leq \frac{1}{2}$.

The process is called “size-reduction” and forms the basis of lattice reduction!

In general, to size-reduce \vec{b}_n , we update the value of \vec{b}_n $n-1$ times by computing the following first for $i = n-1$, then $i = n-2, \dots$ down to $i = 1$:

$$\vec{b}_n := \vec{b}_n - \left\lceil \frac{\vec{b}_i^* \cdot \vec{b}_n}{\vec{b}_i^* \cdot \vec{b}_i^*} \right\rceil \vec{b}_i,$$

where $x := y$ means “put the value of y in the slot of x ”.

4 Lagrange reduction (dimension 2)

Lattice reduction in dimension 2 is easy. Namely, we can perform the following algorithm:

1. Perform size-reduction (Ex 3(d)) on \vec{b}_2 :

$$\vec{b}'_2 = \vec{b}_2 - \left\lceil \frac{\vec{b}_1 \cdot \vec{b}_2}{\vec{b}_1 \cdot \vec{b}_1} \right\rceil \vec{b}_1, \quad (1)$$

(and replace \vec{b}_2 by \vec{b}'_2).

2. If $\|\vec{b}_2\| \geq \|\vec{b}_1\|$, the basis $[\vec{b}_1, \vec{b}_2]$ is reduced, so stop.
3. Otherwise, swap \vec{b}_1 and \vec{b}_2 , then go back to step 1.

In Figure 4, \vec{b}_2 becomes shorter than \vec{b}_1 after size-reduction. Then, by size-reducing \vec{b}_1 with respect to \vec{b}_2 one may possibly shorten \vec{b}_1 further.

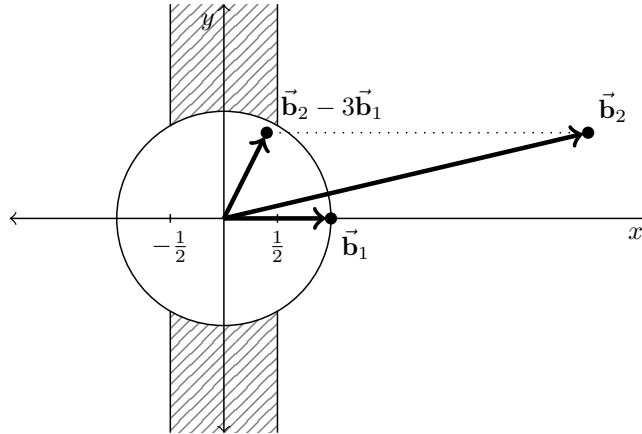


Figure 4: Example of one iteration of Lagrange reduction.

- (a) Show that in each iteration (size reduction & swap) the length of \vec{b}_1 stays the same, or decreases.
- (b) Conclude that the number of iterations is finite.⁶
- (c) When the process terminates, show that \vec{b}_2 must be in the striped region.
- (d) Show: the area of the parallelogram spanned by \vec{b}_1, \vec{b}_2 is the same as the area of the rectangle spanned by \vec{b}_1 and \vec{b}_2^* . Show: this area is

$$\|\vec{b}_1\| \times \|\vec{b}_2^*\|.$$

- (e) Given a lattice in dimension 2, show that we can draw circles around each lattice point, with radius $\frac{1}{2}\|\vec{b}_1\|$ if \vec{b}_1 is one of the shortest vectors. Any lattice gives a so-called “sphere packing”.
- (f) What is the densest (largest radius) sphere packing for a lattice in dimension 2, when you require that $\|\vec{b}_1\| \times \|\vec{b}_2^*\| = 1$?

Hint: (1) look at the ratio $\frac{\left(\frac{1}{2}\|\vec{b}_1\|\right)^2}{\|\vec{b}_1\| \times \|\vec{b}_2^*\|} = \frac{\|\vec{b}_1\|}{4 \times \|\vec{b}_2^*\|}$, and assume \vec{b}_1 lies on the x -axis.

- (2) What is the possible value of \vec{b}_2 lying in the striped region that maximizes this ratio?
- (3) What lattice is this?

⁶Hint: show if $\|\vec{b}_1\|$ stays the same in one iteration that the algorithm stops and the shortest vectors are found.