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Contest (1)

```
template.cpp
```

```
#include <bits/stdc++.h>
using namespace std;
#define rep(i, a, b) for(int i = a; i < (b); ++i)
#define trav(a, x) for(auto& a : x)
#define all(x) x.begin(), x.end()
#define sz(x) (int)(x).size()
typedef long long 11;
typedef long double ld;
typedef pair<int, int> pii;
typedef vector<int> vi;
const ll LLINF = ~(1LL<<63);</pre>
const ld pi = acos(-1.0);
int main() {
  cin.sync_with_stdio(0); cin.tie(0);
 cin.exceptions(cin.failbit);
.bashrc
alias c='q++ -Wall -Wshadow -Wconversion
   -Wfatal-errors -q -std=c++14
   -fsanitize=undefined,address'
xmodmap -e 'clear lock' -e 'keycode 66=less
```

.vimrc

Pre-submit:

- Write a few simple test cases, if sample is not enough.
- Are time limits close? If so, generate max cases.
- Is the memory usage fine?

- Could anything overflow?
- Make sure to submit the right file.

Wrong answer:

- Print your solution! Print debug output, as well.
- Are you clearing all datastructures between test cases?
- Can your algorithm handle the whole range of input?
- Read the full problem statement again.
- Do you handle all corner cases correctly?
- Have you understood the problem correctly?
- Any uninitialized variables?
- Any overflows?
- Confusing N and M, i and j, etc.?
- Are you sure your algorithm works?

- What special cases have you not thought of?
- Are you sure the STL functions you use work as you think?
- Add some assertions, maybe resubmit.
- Create some testcases to run your algorithm on.
- Go through the algorithm for a simple case.
- Go through this list again.
- Explain your algorithm to a team mate.
- Ask the team mate to look at your code.
- Go for a small walk, e.g. to the toilet.
- Is your output format correct? (including whitespace)
- Rewrite your solution from the start or let a team mate do it.

Runtime error:

- Have you tested all corner cases locally?
- Any uninitialized variables?
- Are you reading or writing outside the range of any vector?
- Any assertions that might fail?
- Any possible division by 0? (mod 0 for example)
- Any possible infinite recursion?
- Invalidated pointers or iterators?
- Are you using too much memory?
- Debug with resubmits (e.g. remapped signals, see Various).

Time limit exceeded:

- Do you have any possible infinite loops?
- What is the complexity of your algorithm?
- Are you copying a lot of unnecessary data? (References)
- How big is the input and output? (consider scanf)
- Avoid vector, map. (use arraysunordered_map)
- What do your team mates think about your algorithm?

Memory limit exceeded:

- What is the max amount of memory your algorithm should need?
- Are you clearing all datastructures between test cases?

Mathematics (2)

2.1 Equations

$$ax^2 + bx + c = 0 \Rightarrow x = \frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

The extremum is given by x = -b/2a.

$$ax + by = e$$

$$cx + dy = f \Rightarrow x = \frac{ed - bf}{ad - bc}$$

$$y = \frac{af - ec}{ad - bc}$$

In general, given an equation Ax = b, the solution to a variable x_i is given by

$$x_i = \frac{\det A_i'}{\det A}$$

where A'_i is A with the i'th column replaced by b.

2.2 Recurrences

If $a_n = c_1 a_{n-1} + \cdots + c_k a_{n-k}$, and r_1, \ldots, r_k are distinct roots of $x^k + c_1 x^{k-1} + \cdots + c_k$, there are d_1, \ldots, d_k s.t.

$$a_n = d_1 r_1^n + \dots + d_k r_k^n.$$

Non-distinct roots r become polynomial factors, e.g. $a_n = (d_1n + d_2)r^n$.

2.3 Trigonometry

$$\sin(v+w) = \sin v \cos w + \cos v \sin w$$
$$\cos(v+w) = \cos v \cos w - \sin v \sin w$$

$$\tan(v+w) = \frac{\tan v + \tan w}{1 - \tan v \tan w}$$
$$\sin v + \sin w = 2\sin\frac{v+w}{2}\cos\frac{v-w}{2}$$
$$\cos v + \cos w = 2\cos\frac{v+w}{2}\cos\frac{v-w}{2}$$

$$(V+W)\tan(v-w)/2 = (V-W)\tan(v+w)/2$$

where V, W are lengths of sides opposite angles v, w.

$$a\cos x + b\sin x = r\cos(x - \phi)$$

$$a\sin x + b\cos x = r\sin(x + \phi)$$

where
$$r = \sqrt{a^2 + b^2}$$
, $\phi = \operatorname{atan2}(b, a)$.

2.4 Geometry

Triangles 2.4.1

Side lengths: a, b, c

Semiperimeter: $p = \frac{a+b+c}{2}$

Area: $A = \sqrt{p(p-a)(p-b)(p-c)}$

Circumradius: $R = \frac{abc}{4A}$

Inradius: $r = \frac{A}{}$

Length of median (divides triangle into two equal-area

triangles): $m_a = \frac{1}{2}\sqrt{2b^2 + 2c^2 - a^2}$

Length of bisector (divides angles in two):

$$s_a = \sqrt{bc \left[1 - \left(\frac{a}{b+c}\right)^2\right]}$$

Law of sines: $\frac{\sin \alpha}{a} = \frac{\sin \beta}{b} = \frac{\sin \gamma}{c} = \frac{1}{2R}$ Law of cosines: $a^2 = b^2 + c^2 - 2bc \cos \alpha$

Law of tangents: $\frac{a+b}{a-b} = \frac{\tan \frac{\alpha+\beta}{2}}{\tan \frac{\alpha-\beta}{2}}$

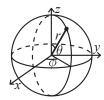
Quadrilaterals 2.4.2

With side lengths a, b, c, d, diagonals e, f, diagonals angle θ , area A and magic flux $F = b^2 + d^2 - a^2 - c^2$:

$$4A = 2ef \cdot \sin \theta = F \tan \theta = \sqrt{4e^2f^2 - F^2}$$

For cyclic quadrilaterals the sum of opposite angles is 180°, ef = ac + bd, and $A = \sqrt{(p-a)(p-b)(p-c)(p-d)}$.

Spherical coordinates 2.4.3



$$x = r \sin \theta \cos \phi \qquad r = \sqrt{x^2 + y^2 + z^2}$$

$$y = r \sin \theta \sin \phi \quad \theta = a\cos(z/\sqrt{x^2 + y^2 + z^2})$$

$$z = r \cos \theta \qquad \phi = a\tan(2y, x)$$

Derivatives/Integrals

$$\frac{d}{dx}\arcsin x = \frac{1}{\sqrt{1-x^2}} \quad \frac{d}{dx}\arccos x = -\frac{1}{\sqrt{1-x^2}}$$

$$\frac{d}{dx}\tan x = 1 + \tan^2 x \quad \frac{d}{dx}\arctan x = \frac{1}{1+x^2}$$

$$\int \tan ax = -\frac{\ln|\cos ax|}{a} \quad \int x\sin ax = \frac{\sin ax - ax\cos ax}{a^2}$$

$$\int e^{-x^2} = \frac{\sqrt{\pi}}{2}\operatorname{erf}(x) \quad \int xe^{ax}dx = \frac{e^{ax}}{a^2}(ax-1)$$

Integration by parts:

$$\int_{a}^{b} f(x)g(x)dx = [F(x)g(x)]_{a}^{b} - \int_{a}^{b} F(x)g'(x)dx$$

2.6 Sums

$$c^{a} + c^{a+1} + \dots + c^{b} = \frac{c^{b+1} - c^{a}}{c-1}, c \neq 1$$

$$1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$$

$$1^{2} + 2^{2} + 3^{2} + \dots + n^{2} = \frac{n(2n+1)(n+1)}{6}$$

$$1^{3} + 2^{3} + 3^{3} + \dots + n^{3} = \frac{n^{2}(n+1)^{2}}{4}$$

$$1^{4} + 2^{4} + 3^{4} + \dots + n^{4} = \frac{n(n+1)(2n+1)(3n^{2} + 3n - 1)}{30}$$

2.7Series

$$e^{x} = 1 + x + \frac{x^{2}}{2!} + \frac{x^{3}}{3!} + \dots, (-\infty < x < \infty)$$
$$\ln(1+x) = x - \frac{x^{2}}{2} + \frac{x^{3}}{3} - \frac{x^{4}}{4} + \dots, (-1 < x \le 1)$$

$$\sqrt{1+x} = 1 + \frac{x}{2} - \frac{x^2}{8} + \frac{2x^3}{32} - \frac{5x^4}{128} + \dots, (-1 \le x \le 1)$$

$$\sin x = x - \frac{x^3}{31} + \frac{x^5}{51} - \frac{x^7}{71} + \dots, (-\infty < x < \infty)$$

$$\cos x = 1 - \frac{x^2}{2!} + \frac{x^4}{4!} - \frac{x^6}{6!} + \dots, (-\infty < x < \infty)$$

Probability theory

Let X be a discrete random variable with probability $p_X(x)$ of assuming the value x. It will then have an expected value (mean) $\mu = \mathbb{E}(X) = \sum_x x p_X(x)$ and variance $\sigma^2 = V(X) = \mathbb{E}(X^2) - (\mathbb{E}(X))^2 = \sum_x (x - \mathbb{E}(X))^2 p_X(x)$ where σ is the standard deviation. If X is instead continuous it will have a probability density function $f_X(x)$ and the sums above will instead be integrals with $p_X(x)$ replaced by $f_X(x)$.

Expectation is linear:

$$\mathbb{E}(aX + bY) = a\mathbb{E}(X) + b\mathbb{E}(Y)$$

For independent X and Y,

$$V(aX + bY) = a^2V(X) + b^2V(Y).$$

2.8.1Discrete distributions

Binomial distribution

The number of successes in n independent yes/no experiments, each which yields success with probability p is Bin(n, p), n = 1, 2, ..., 0

$$p(k) = \binom{n}{k} p^k (1-p)^{n-k}$$

$$\mu = np, \ \sigma^2 = np(1-p)$$

Bin(n, p) is approximately Po(np) for small p.

First success distribution

The number of trials needed to get the first success in independent yes/no experiments, each wich yields success with probability p is Fs(p), $0 \le p \le 1$.

$$p(k) = p(1-p)^{k-1}, k = 1, 2, \dots$$

$$\mu = \frac{1}{p}, \sigma^2 = \frac{1-p}{p^2}$$

Poisson distribution

The number of events occurring in a fixed period of time t if these events occur with a known average rate κ and independently of the time since the last event is $Po(\lambda)$, $\lambda = t\kappa$.

$$p(k) = e^{-\lambda} \frac{\lambda^k}{k!}, k = 0, 1, 2, \dots$$
$$\mu = \lambda, \sigma^2 = \lambda$$

2.8.2 Continuous distributions

Uniform distribution

If the probability density function is constant between a and b and 0 elsewhere it is U(a,b), a < b.

$$f(x) = \begin{cases} \frac{1}{b-a} & a < x < b \\ 0 & \text{otherwise} \end{cases}$$

$$\mu = \frac{a+b}{2}, \, \sigma^2 = \frac{(b-a)^2}{12}$$

Exponential distribution

The time between events in a Poisson process is $\operatorname{Exp}(\lambda)$, $\lambda > 0$.

$$f(x) = \begin{cases} \lambda e^{-\lambda x} & x \ge 0\\ 0 & x < 0 \end{cases}$$
$$\mu = \frac{1}{\lambda}, \, \sigma^2 = \frac{1}{\lambda^2}$$

Normal distribution

Most real random values with mean μ and variance σ^2 are well described by $\mathcal{N}(\mu, \sigma^2)$, $\sigma > 0$.

$$f(x) = \frac{1}{\sqrt{2\pi\sigma^2}} e^{-\frac{(x-\mu)^2}{2\sigma^2}}$$

If $X_1 \sim \mathcal{N}(\mu_1, \sigma_1^2)$ and $X_2 \sim \mathcal{N}(\mu_2, \sigma_2^2)$ then $aX_1 + bX_2 + c \sim \mathcal{N}(\mu_1 + \mu_2 + c, a^2\sigma_1^2 + b^2\sigma_2^2)$

2.9 Markov chains

A Markov chain is a discrete random process with the property that the next state depends only on the current state. Let X_1, X_2, \ldots be a sequence of random variables generated by the Markov process. Then there is a transition matrix $\mathbf{P} = (p_{ij})$, with $p_{ij} = \Pr(X_n = i | X_{n-1} = j)$, and $\mathbf{p}^{(n)} = \mathbf{P}^n \mathbf{p}^{(0)}$ is the probability distribution for X_n (i.e., $p_i^{(n)} = \Pr(X_n = i)$), where $\mathbf{p}^{(0)}$ is the initial distribution.

 π is a stationary distribution if $\pi = \pi \mathbf{P}$. If the Markov chain is *irreducible* (it is possible to get to any state from any state), then $\pi_i = \frac{1}{\mathbb{E}(T_i)}$ where $\mathbb{E}(T_i)$ is the expected time between two visits in state i. π_j/π_i is the expected number of visits in state j between two visits in state i.

For a connected, undirected and non-bipartite graph, where the transition probability is uniform among all neighbors, π_i is proportional to node *i*'s degree.

A Markov chain is *ergodic* if the asymptotic distribution is independent of the initial distribution. A finite Markov chain is ergodic iff it is irreducible and *aperiodic* (i.e., the gcd of cycle lengths is 1). $\lim_{k\to\infty} \mathbf{P}^k = \mathbf{1}\pi$.

A Markov chain is an A-chain if the states can be partitioned into two sets **A** and **G**, such that all states in **A** are absorbing $(p_{ii} = 1)$, and all states in **G** leads to an absorbing state in **A**. The probability for absorption in state $i \in \mathbf{A}$, when the initial state is j, is $a_{ij} = p_{ij} + \sum_{k \in \mathbf{G}} a_{ik} p_{kj}$. The expected time until absorption, when the initial state is i, is $t_i = 1 + \sum_{k \in \mathbf{G}} p_{ki} t_k$.

Data structures (3)

```
OrderStatisticTree.h
#include <bits/extc++.h>
using namespace __qnu_pbds;
template < class TK, class TM> using TreeMap =

    tree<TK, TM,
</pre>
  less<TK>, rb_tree_tag,

    tree_order_statistics_node_update>;

template < class T > using Tree = TreeMap < T,</pre>
→ null type>;
void example() {
  Tree<int> t, t2; t.insert(8);
  auto it = t.insert(10).first;
  assert(it == t.lower bound(9));
  assert(t.order_of_key(10) == 1);
  assert(t.order_of_key(11) == 2);
  assert(*t.find_by_order(0) == 8);
  t.join(t2); // assuming T < T2 or T > T2, merge
  \rightarrow t2 into t
HashMap.h
#include <bits/extc++.h>
__gnu_pbds::gp_hash_table<ll, int> h({},{},{},{},
SegmentTree.h
struct Tree {
  typedef int T;
  static const T LOW = INT_MIN;
  T f(T a, T b)  { return max(a, b); } // (any
  → associative fn)
  vector<T> s; int n;
  Tree (int n = 0, T def = 0) : s(2*n, def), n(n)
  void update(int pos, T val) {
    for (s[pos += n] = val; pos > 1; pos /= 2)
      s[pos / 2] = f(s[pos & ~1], s[pos | 1]);
  T query (int b, int e) { // query [b, e)
    T ra = LOW, rb = LOW;
    for (b += n, e += n; b < e; b /= 2, e /= 2) {
      if (b % 2) ra = f(ra, s[b++]);
      if (e % 2) rb = f(s[--e], rb);
    return f(ra, rb);
```

LazySegmentTree.h

```
const int inf = 1e9;
struct Node {
 Node *1 = 0, *r = 0;
 int lo, hi, mset = inf, madd = 0, val = -inf;
 Node (int lo, int hi):lo(lo), hi(hi) {} // Large

→ interval of -inf

 Node (vi& v, int lo, int hi) : lo(lo), hi(hi) {
   if (lo + 1 < hi) {
      int mid = 10 + (hi - 10)/2;
      1 = new Node(v, lo, mid); r = new Node(v,

→ mid, hi);
      val = max(1->val, r->val);
    else val = v[lo];
 int query(int L, int R) {
   if (R <= lo || hi <= L) return -inf;</pre>
    if (L <= lo && hi <= R) return val;</pre>
    push();
    return max(l->query(L, R), r->query(L, R));
 void set(int L, int R, int x) {
   if (R <= lo | | hi <= L) return;</pre>
    if (L <= lo && hi <= R) mset = val = x, madd
    \hookrightarrow = 0;
    else {
      push(), l\rightarrow set(L, R, x), r\rightarrow set(L, R, x);
      val = max(1->val, r->val);
 void add(int L, int R, int x) {
   if (R <= lo || hi <= L) return;
    if (L <= lo && hi <= R) {
      if (mset != inf) mset += x;
      else madd += x;
      val += x;
    else {
      push(), l->add(L, R, x), r->add(L, R, x);
      val = max(1->val, r->val);
 void push() {
   if (!1) {
      int mid = 10 + (hi - 10)/2;
      1 = new Node(lo, mid); r = new Node(mid,
      \rightarrow hi);
    if (mset != inf)
      1->set(lo,hi,mset), r->set(lo,hi,mset),
      \hookrightarrow mset = inf;
    olao if (madd)
```

```
UnionFind.h
struct UF {
 vi e;
  UF (int n) : e(n, -1) {}
  bool same_set(int a, int b) { return find(a) ==
  \hookrightarrow find(b); }
  int size(int x) { return -e[find(x)]; }
  int find(int x) { return e[x] < 0 ? x : e[x] =
  \rightarrow find(e[x]); }
  void join(int a, int b) {
    a = find(a), b = find(b);
    if (a == b) return;
    if (e[a] > e[b]) swap(a, b);
    e[a] += e[b]; e[b] = a;
};
SubMatrix.h
template<class T>
struct SubMatrix
 vector<vector<T>> p;
  SubMatrix(vector<vector<T>>& v) {
    int R = sz(v), C = sz(v[0]);
    p.assign(R+1, vectorT>(C+1));
    rep(r, 0, R) rep(c, 0, C)
      p[r+1][c+1] = v[r][c] + p[r][c+1] +
       \rightarrow p[r+1][c] - p[r][c];
 T sum(int u, int l, int d, int r) {
    return p[d][r] - p[d][l] - p[u][r] + p[u][l];
} ;
```

Matrix.h

Treap.h

```
template < class T, int N > struct Matrix {
  typedef Matrix M;
  array<array<T, N>, N> d{};
 M operator* (const M& m) const {
   Ma;
    rep(i,0,N) rep(j,0,N)
      rep(k, 0, N) \ a.d[i][j] += d[i][k]*m.d[k][j];
   return a;
  vector<T> operator*(const vector<T>& vec) const
  ← {
   vector<T> ret(N);
    rep(i, 0, N) rep(j, 0, N) ret[i] += d[i][j] *
    → vec[j];
   return ret;
 M operator^(ll p) const {
    assert (p >= 0);
   M a, b(*this);
    rep(i, 0, N) \ a.d[i][i] = 1;
    while (p) {
      if (p&1) a = a*b;
      b = b*b;
      p >>= 1;
    return a;
};
```

```
bool O:
struct Line {
 mutable ll k, m, p;
 bool operator<(const Line& o) const {</pre>
    return Q ? p < o.p : k < o.k;
};
struct LineContainer : multiset<Line> {
  // (for doubles, use inf = 1/.0, div(a,b) =
  \rightarrow a/b
  const ll inf = LLONG_MAX;
 11 div(ll a, ll b) { // floored division
    return a / b - ((a ^ b) < 0 && a % b); }
 bool isect(iterator x, iterator y) {
    if (y == end()) { x->p = inf; return false; }
    if (x->k == y->k) x->p = x->m > y->m ? inf :
    → -inf;
    else x->p = div(y->m - x->m, x->k - y->k);
    return x->p >= y->p;
 void add(ll k, ll m) {
    auto z = insert(\{k, m, 0\}), y = z++, x = y;
    while (isect(y, z)) z = erase(z);
    if (x != begin() \&\& isect(--x, y)) isect(x, y)
    \Rightarrow = erase(y));
    while ((y = x) != begin() \&\& (--x) -> p >=
    isect(x, erase(y));
 11 querv(11 x) {
   assert(!empty());
    0 = 1; auto 1 = *lower_bound({0,0,x}); Q = 0;
    return 1.k * x + 1.m;
} ;
```

```
struct Node {
  Node *1 = 0, *r = 0;
  int val, y, c = 1;
 Node(int val) : val(val), y(rand()) {}
  void recalc();
};
int cnt(Node* n) { return n ? n->c : 0; }
void Node::recalc() { c = cnt(l) + cnt(r) + 1; }
template < class F > void each (Node * n, F f) {
  if (n) { each (n->1, f); f(n->val); each (n->r, f)
  \hookrightarrow f); }
pair<Node*, Node*> split(Node* n, int k) {
  if (!n) return {};
  if (cnt(n->1) >= k) { // "n->val >= v" for
  \rightarrow lower_bound(v)
    auto pa = split(n->1, k);
    n->1 = pa.second;
    n->recalc();
    return {pa.first, n};
    auto pa = split (n->r, k - cnt(n->1) - 1);
    n->r = pa.first;
    n->recalc();
    return {n, pa.second};
Node* merge(Node* 1, Node* r) {
  if (!1) return r;
  if (!r) return 1;
  if (1->y > r->y) {
   1->r = merge(1->r, r);
    1->recalc();
    return 1:
  } else {
    r->1 = merge(1, r->1);
    r->recalc();
    return r;
Node* ins(Node* t, Node* n, int pos) {
  auto pa = split(t, pos);
  return merge (merge (pa.first, n), pa.second);
// Example application: move the range [1, r) to
```

KTH

FenwickTree.h

```
struct FT {
 vector<ll> s;
 FT(int n) : s(n) {}
 void update(int pos, ll dif) { // a[pos] += dif
    for (; pos < sz(s); pos |= pos + 1) s[pos] +=

    dif;

  11 query(int pos) { // sum of values in [0,

→ pos)

   11 \text{ res} = 0;
    for (; pos > 0; pos &= pos - 1) res +=
    \rightarrow s[pos-1];
    return res;
  int lower_bound(ll sum) {// min pos st sum of
  \hookrightarrow [0, pos] >= sum
    // Returns n if no sum is \geq sum, or -1 if
    → empty sum is.
    if (sum \leq 0) return -1;
    int pos = 0;
    for (int pw = 1 << 25; pw; pw >>= 1) {
      if (pos + pw \le sz(s) \&\& s[pos + pw-1] \le

    sum)

        pos += pw, sum -= s[pos-1];
    return pos;
};
```

```
struct FT2 {
 vector<vi> vs; vector<FT> ft;
 FT2(int limx) : ys(limx) {}
 void fakeUpdate(int x, int y) {
   for (; x < sz(ys); x = x + 1)

    ys[x].push_back(y);

 void init() {
   trav(v, ys) sort(all(v)),

    ft.emplace_back(sz(v));

 int ind(int x, int y) {
   return (int) (lower_bound(all(ys[x]), y) -
    \rightarrow ys[x].begin()); }
 void update(int x, int y, ll dif) {
   for (; x < sz(ys); x = x + 1)
      ft[x].update(ind(x, y), dif);
 11 query(int x, int y) {
   11 \text{ sum} = 0;
   for (; x; x &= x - 1)
      sum += ft[x-1].query(ind(x-1, y));
   return sum;
 }
};
RMQ.h
template<class T>
struct RMO {
 vector<vector<T>> jmp;
  RMQ(const vector<T>& V) {
    int N = sz(V), on = 1, depth = 1;
   while (on < sz(V)) on *= 2, depth++;
    jmp.assign(depth, V);
   rep(i, 0, depth-1) rep(j, 0, N)
      jmp[i+1][j] = min(jmp[i][j],
      jmp[i][min(N - 1, j + (1 << i))]);
 T query(int a, int b) {
   assert (a < b); // or return inf if a == b
   int dep = 31 - __builtin_clz(b - a);
   return min(jmp[dep][a], jmp[dep][b - (1 <<</pre>
    → dep)]);
};
```

```
double gss (double a, double b, double
double r = (sqrt(5)-1)/2, eps = 1e-7;
 double x1 = b - r*(b-a), x2 = a + r*(b-a);
 double f1 = f(x1), f2 = f(x2);
 while (b-a > eps)
   if (f1 < f2) { //change to > to find maximum
      b = x2; x2 = x1; f2 = f1;
      x1 = b - r*(b-a); f1 = f(x1);
   } else {
      a = x1; x1 = x2; f1 = f2;
     x2 = a + r*(b-a); f2 = f(x2);
 return a;
Polynomial.h
struct Poly {
 vector<double> a;
 double operator()(double x) const {
    double val = 0;
    for(int i = sz(a); i--;) (val *= x) += a[i];
    return val;
 void diff() {
   rep(i,1,sz(a)) a[i-1] = i*a[i];
    a.pop_back();
 void divroot (double x0) {
    double b = a.back(), c; a.back() = 0;
    for(int i=sz(a)-1; i--;) c = a[i], a[i] =
    \rightarrow a[i+1] \times x0+b, b=c;
    a.pop_back();
};
```

GoldenSectionSearch.h.

PolyRoots.h

```
vector<double> poly_roots(Poly p, double xmin,

→ double xmax) {
 if (sz(p.a) == 2) { return {-p.a[0]/p.a[1]}; }
 vector<double> ret;
 Poly der = p_i
  der.diff();
  auto dr = poly_roots(der, xmin, xmax);
  dr.push back(xmin-1);
  dr.push_back(xmax+1);
  sort(all(dr));
  rep(i, 0, sz(dr) - 1) {
   double l = dr[i], h = dr[i+1];
   bool sign = p(1) > 0;
   if (sign ^ (p(h) > 0)) {
     rep(it, 0, 60) { // while (h - 1 > 1e-8)
        double m = (1 + h) / 2, f = p(m);
       if ((f <= 0) ^ sign) l = m;
        else h = m;
      ret.push_back((1 + h) / 2);
  return ret;
```

PolyInterpolate.h

```
typedef vector<double> vd;
vd interpolate(vd x, vd y, int n) {
  vd res(n), temp(n);
  rep(k,0,n-1) rep(i,k+1,n)
    y[i] = (y[i] - y[k]) / (x[i] - x[k]);
  double last = 0; temp[0] = 1;
  rep(k,0,n) rep(i,0,n) {
    res[i] += y[k] * temp[i];
    swap(last, temp[i]);
    temp[i] -= last * x[k];
  }
  return res;
}
```

```
vector<ll> BerlekampMassey(vector<ll> s) {
 int n = sz(s), L = 0, m = 0;
 vector<ll> C(n), B(n), T;
 C[0] = B[0] = 1;
 11 b = 1;
 rep(i, 0, n) \{ ++m;
   ll d = s[i] % mod;
   rep(j,1,L+1) d = (d + C[j] * s[i - j]) % mod;
   if (!d) continue;
   T = C; 11 coef = d * modpow(b, mod-2) % mod;
   rep(j,m,n) C[j] = (C[j] - coef * B[j - m]) %
    → mod;
   if (2 * L > i) continue;
   L = i + 1 - L; B = T; b = d; m = 0;
 C.resize(L + 1); C.erase(C.begin());
 trav(x, C) x = (mod - x) % mod;
 return C;
```

```
typedef vector<ll> Poly;
ll linearRec(Poly S, Poly tr, ll k) {
  int n = sz(S);
  auto combine = [&](Poly a, Poly b) {
    Poly res(n * 2 + 1);
    rep(i, 0, n+1) rep(j, 0, n+1)
      res[i + j] = (res[i + j] + a[i] * b[j]) %
       → mod;
    for (int i = 2 * n; i > n; --i) rep(j,0,n)
      res[i - 1 - j] = (res[i - 1 - j] + res[i] *
       \rightarrow tr[j]) % mod;
    res.resize(n + 1);
    return res;
  };
  Poly pol(n + 1), e(pol);
  pol[0] = e[1] = 1;
  for (++k; k; k /= 2) {
    if (k % 2) pol = combine(pol, e);
    e = combine(e, e);
  }
  11 \text{ res} = 0;
  rep(i, 0, n) res = (res + pol[i + 1] * S[i]) %
  → mod;
  return res;
HillClimbing.h
typedef array<double, 2> P;
double func(P p);
pair<double, P> hillClimb(P start) {
  pair<double, P> cur(func(start), start);
  for (double jmp = 1e9; jmp > 1e-20; jmp /= 2) {
    rep(j, 0, 100) rep(dx, -1, 2) rep(dy, -1, 2) {
      P p = cur.second;
      p[0] += dx * jmp;
      p[1] += dy * jmp;
      cur = min(cur, make_pair(func(p), p));
  return cur;
```

Integrate.h

```
double quad(double (*f)(double), double a, double
\rightarrow b) {
  const int n = 1000;
  double h = (b - a) / 2 / n;
  double v = f(a) + f(b);
  rep(i, 1, n*2)
   v += f(a + i*h) * (i&1 ? 4 : 2);
  return v * h / 3;
IntegrateAdaptive.h
typedef double d;
d simpson(d (*f)(d), d a, d b) {
 dc = (a+b) / 2;
  return (f(a) + 4*f(c) + f(b)) * (b-a) / 6;
d \operatorname{rec}(d (*f)(d), d a, d b, d \operatorname{eps}, d S) {
 dc = (a+b) / 2;
 d S1 = simpson(f, a, c);
 d S2 = simpson(f, c, b), T = S1 + S2;
 if (abs (T - S) <= 15*eps || b-a < 1e-10)
   return T + (T - S) / 15;
  \rightarrow eps/2, S2);
d quad(d (*f)(d), d a, d b, d eps = 1e-8) {
  return rec(f, a, b, eps, simpson(f, a, b));
Determinant.h
double det (vector<vector<double>>& a) {
  int n = sz(a); double res = 1;
  rep(i,0,n) {
   int b = i;
    rep(j,i+1,n) if (fabs(a[j][i]) >
    \rightarrow fabs(a[b][i])) b = i;
    if (i != b) swap(a[i], a[b]), res *= -1;
    res *= a[i][i];
   if (res == 0) return 0;
    rep(j, i+1, n) {
      double v = a[j][i] / a[i][i];
      if (v != 0) rep(k, i+1, n) a[j][k] -= v *
      \rightarrow a[i][k];
  return res;
```

```
IntDeterminant.h
const 11 mod = 12345;
11 det(vector<vector<11>>& a) {
  int n = sz(a); ll ans = 1;
  rep(i,0,n) {
    rep(j,i+1,n) {
      while (a[j][i] != 0) { // gcd step
        ll t = a[i][i] / a[j][i];
        if (t) rep(k,i,n)
          a[i][k] = (a[i][k] - a[j][k] * t) %
          → mod;
        swap(a[i], a[i]);
        ans *=-1;
    ans = ans * a[i][i] % mod;
    if (!ans) return 0;
  return (ans + mod) % mod;
Simplex.h
```

```
typedef double T; // long double, Rational,

    double + mod<P>...

typedef vector<T> vd;
typedef vector<vd> vvd;
const T eps = 1e-8, inf = 1/.0;
#define MP make_pair
\#define ltj(X) if (s == -1 \mid \mid MP(X[j], N[j]) <
\hookrightarrow MP(X[s],N[s])) s=j
struct LPSolver {
 int m, n;
 vi N, B;
  vvd D;
  LPSolver(const vvd& A, const vd& b, const vd&
  m(sz(b)), n(sz(c)), N(n+1), B(m), D(m+2)
    \rightarrow vd(n+2)) {
      rep(i, 0, m) rep(j, 0, n) D[i][j] = A[i][j];
      rep(i, 0, m) \{ B[i] = n+i; D[i][n] = -1;
      \rightarrow D[i][n+1] = b[i];
      rep(j, 0, n) \{ N[j] = j; D[m][j] = -c[j]; \}
      N[n] = -1; D[m+1][n] = 1;
  void pivot(int r, int s) {
    T * a = D[r].data(), inv = 1 / a[s];
    rep(i, 0, m+2) if (i != r \& \& abs(D[i][s]) >
    → eps) {
      T *b = D[i].data(), inv2 = b[s] * inv;
      rep(j, 0, n+2) b[j] -= a[j] * inv2;
      b[s] = a[s] * inv2;
    rep(j, 0, n+2) if (j!= s) D[r][j] *= inv;
    rep(i, 0, m+2) if (i != r) D[i][s] *= -inv;
    D[r][s] = inv;
    swap(B[r], N[s]);
  bool simplex(int phase) {
    int x = m + phase - 1;
    for (;;) {
      int s = -1;
      rep(j,0,n+1) if (N[j] != -phase) ltj(D[x]);
      if (D[x][s] >= -eps) return true;
      int r = -1;
      rep(i,0,m) {
        if (D[i][s] <= eps) continue;</pre>
        if (r == -1 \mid | MP(D[i][n+1] / D[i][s],
         \hookrightarrow B[i])
                       \angle MD/D[x][x+1]/D[x][a]
```

```
SolveLinear.h
typedef vector<double> vd;
const double eps = 1e-12;
int solveLinear(vector<vd>& A, vd& b, vd& x) {
  int n = sz(A), m = sz(x), rank = 0, br, bc;
  if (n) assert(sz(A[0]) == m);
  vi col(m); iota(all(col), 0);
  rep(i,0,n) {
   double v, bv = 0;
   rep(r,i,n) rep(c,i,m)
      if ((v = fabs(A[r][c])) > bv)
        br = r, bc = c, bv = v;
   if (bv <= eps) {
      rep(j, i, n) if (fabs(b[j]) > eps) return -1;
      break;
    swap(A[i], A[br]);
    swap(b[i], b[br]);
    swap(col[i], col[bc]);
    rep(j,0,n) swap(A[j][i], A[j][bc]);
   bv = 1/A[i][i];
    rep(j, i+1, n) {
      double fac = A[j][i] * bv;
      b[j] -= fac * b[i];
      rep(k,i+1,m) A[j][k] -= fac*A[i][k];
    rank++;
 x.assign(m, 0);
  for (int i = rank; i--;) {
   b[i] /= A[i][i];
   x[col[i]] = b[i];
    rep(j, 0, i) b[j] -= A[j][i] * b[i];
  return rank; // (multiple solutions if rank <</pre>
  \hookrightarrow m)
SolveLinear2.h
rep(j,0,n) if (j != i) // instead of <math>rep(j,i+1,n)
// ... then at the end:
x.assign(m, undefined);
rep(i, 0, rank) {
  rep(j,rank,m) if (fabs(A[i][j]) > eps) goto

  fail;

 x[col[i]] = b[i] / A[i][i];
fail:; }
```

```
SolveLinearBinary.h
typedef bitset<1000> bs;
int solveLinear(vector<bs>& A, vi& b, bs& x, int
\hookrightarrow m) {
 int n = sz(A), rank = 0, br;
 assert(m \le sz(x));
 vi col(m); iota(all(col), 0);
 rep(i,0,n) {
    for (br=i; br<n; ++br) if (A[br].any())</pre>
    → break;
    if (br == n) {
      rep(j,i,n) if(b[j]) return -1;
      break;
    int bc = (int)A[br]._Find_next(i-1);
    swap(A[i], A[br]);
    swap(b[i], b[br]);
    swap(col[i], col[bc]);
    rep(j, 0, n) if (A[j][i] != A[j][bc]) {
      A[j].flip(i); A[j].flip(bc);
    rep(j,i+1,n) if (A[j][i]) {
     b[j] = b[i];
      A[i] ^= A[i];
    rank++;
 x = bs();
 for (int i = rank; i--;) {
    if (!b[i]) continue;
    x[col[i]] = 1;
    rep(j,0,i) b[j] ^= A[j][i];
 return rank; // (multiple solutions if rank <</pre>
  \hookrightarrow m)
```

MatrixInverse.h

```
int matInv(vector<vector<double>>& A) {
  int n = sz(A); vi col(n);
  vector<vector<double>> tmp(n,

    vector < double > (n));

  rep(i, 0, n) tmp[i][i] = 1, col[i] = i;
 rep(i,0,n) {
    int r = i, c = i;
    rep(j,i,n) rep(k,i,n)
      if (fabs(A[j][k]) > fabs(A[r][c]))
        r = j, c = k;
    if (fabs(A[r][c]) < 1e-12) return i;
    A[i].swap(A[r]); tmp[i].swap(tmp[r]);
    rep(j,0,n)
      swap(A[j][i], A[j][c]), swap(tmp[j][i],
       \rightarrow tmp[j][c]);
    swap(col[i], col[c]);
    double v = A[i][i];
    rep(j, i+1, n) {
      double f = A[j][i] / v;
      A[j][i] = 0;
      rep(k, i+1, n) A[j][k] -= f*A[i][k];
      rep(k, 0, n) tmp[j][k] -= f*tmp[i][k];
    rep(j, i+1, n) A[i][j] /= v;
    rep(j,0,n) tmp[i][j] /= v;
    A[i][i] = 1;
  }
  for (int i = n-1; i > 0; --i) rep(j,0,i) {
    double v = A[i][i];
    rep(k, 0, n) tmp[j][k] -= v*tmp[i][k];
  }
  rep(i,0,n) rep(j,0,n) A[col[i]][col[j]] =
  \hookrightarrow tmp[i][j];
  return n;
```

Tridiagonal.h

```
typedef double T;
vector<T> tridiagonal(vector<T> diag, const

→ vector<T>& super,

    const vector<T>& sub, vector<T> b) {
 int n = sz(b); vi tr(n);
  rep(i, 0, n-1) {
   if (abs(diag[i]) < 1e-9 * abs(super[i])) { //
    \rightarrow diag[i] == 0
      b[i+1] -= b[i] * diag[i+1] / super[i];
      if (i+2 < n) b[i+2] -= b[i] * sub[i+1] /

    super[i];

      diaq[i+1] = sub[i]; tr[++i] = 1;
   } else {
      diag[i+1] -= super[i]*sub[i]/diag[i];
      b[i+1] = b[i] * sub[i] / diag[i];
 for (int i = n; i--;) {
   if (tr[i]) {
      swap(b[i], b[i-1]);
      diag[i-1] = diag[i];
      b[i] /= super[i-1];
   } else {
      b[i] /= diag[i];
      if (i) b[i-1] -= b[i] * super[i-1];
  return b;
```

```
typedef valarray<complex<double> > carray;
void fft(carray& x, carray& roots) {
 int N = sz(x);
 if (N <= 1) return;</pre>
 carray even = x[slice(0, N/2, 2)];
 carray odd = x[slice(1, N/2, 2)];
 carray rs = roots[slice(0, N/2, 2)];
  fft(even, rs);
  fft (odd, rs);
 rep(k, 0, N/2) {
   auto t = roots[k] * odd[k];
   x[k] = even[k] + t;
   x[k+N/2] = even[k] - t;
typedef vector<double> vd;
vd conv(const vd& a, const vd& b) {
 int s = sz(a) + sz(b) - 1, L =
  \rightarrow 32-_builtin_clz(s), n = 1<<L;
 if (s <= 0) return {};
 carray av(n), bv(n), roots(n);
 rep(i,0,n) roots[i] = polar(1.0, -2 * M_PI * i
  \rightarrow / n);
  copy(all(a), begin(av)); fft(av, roots);
  copy(all(b), begin(bv)); fft(bv, roots);
  roots = roots.apply(conj);
 carray cv = av * bv; fft(cv, roots);
 vd c(s); rep(i, 0, s) c[i] = cv[i].real() / n;
 return c:
```

```
typedef complex<double> cpx;
const int LOGN = 19, MAXN = 1 << LOGN;</pre>
int rev[MAXN];
cpx rt[MAXN], a[MAXN] = \{\}, b[MAXN] = \{\};
void fft(cpx *A) {
  REP(i, MAXN) if (i < rev[i]) swap(A[i],

    A[rev[i]]);
  for (int k = 1; k < MAXN; k \neq 2)
    for (int i = 0; i < MAXN; i += 2 * k) REP (i, k)
        cpx t = rt[j + k] * A[i + j + k];
        A[i + j + k] = A[i + j] - t;
        A[i + j] += t;
void multiply() { // a = convolution of a * b
  rev[0] = 0; rt[1] = cpx(1, 0);
  REP(i, MAXN) rev[i] = (rev[i/2] |
  \hookrightarrow (i&1) << LOGN) /2;
  for (int k = 2; k < MAXN; k \neq 2) {
    cpx z(cos(PI/k), sin(PI/k));
    rep(i, k/2, k) rt[2*i]=rt[i],
     \hookrightarrow rt[2*i+1]=rt[i]*z;
  fft(a); fft(b);
  REP(i, MAXN) a[i] *= b[i] / (double) MAXN;
  reverse(a+1,a+MAXN); fft(a);
```

4.1 Fourier transforms

```
const 11 mod = (119 << 23) + 1, root = 3; // =
→ 998244353
// For p < 2^30 there is also e.g. (5 << 25, 3),
// (479 << 21, 3) and (483 << 21, 5). The last
\rightarrow two are > 10^9.
typedef vector<ll> v1;
void ntt(ll* x, ll* temp, ll* roots, int N, int

    skip) {

  if (N == 1) return;
  int n2 = N/2;
  ntt(x
          , temp, roots, n2, skip*2);
  ntt(x+skip, temp, roots, n2, skip*2);
  rep(i, 0, N) temp[i] = x[i*skip];
  rep(i,0,n2) {
    11 s = temp[2*i], t = temp[2*i+1] *

→ roots[skip*i];

    x[skip*i] = (s + t) % mod; x[skip*(i+n2)] =
    \hookrightarrow (s - t) % mod;
void ntt(vl& x, bool inv = false) {
  11 e = modpow(root, (mod-1) / sz(x));
  if (inv) e = modpow(e, mod-2);
  vl roots(sz(x), 1), temp = roots;
  rep(i,1,sz(x)) roots[i] = roots[i-1] * e % mod;
  ntt(&x[0], &temp[0], &roots[0], sz(x), 1);
vl conv(vl a, vl b) {
  int s = sz(a) + sz(b) - 1; if (s \le 0) return
  int L = s > 1 ? 32 - \_builtin\_clz(s - 1) : 0,
  \hookrightarrow n = 1 << L;
  if (s <= 200) { // (factor 10 optimization for</pre>
  \rightarrow |a|, |b| = 10)
    vlc(s);
    rep(i, 0, sz(a)) rep(j, 0, sz(b))
      c[i + j] = (c[i + j] + a[i] * b[j]) % mod;
    return c;
  a.resize(n); ntt(a);
  b.resize(n); ntt(b);
  v1 c(n); 11 d = modpow(n, mod-2);
  rep(i, 0, n) c[i] = a[i] * b[i] % mod * d % mod;
  ntt(c, true); c.resize(s); return c;
```

```
void FST(vi& a, bool inv) {
  for (int n = sz(a), step = 1; step < n; step *=</pre>

→ 2) {

    for (int i = 0; i < n; i += 2 * step)</pre>
    \hookrightarrow rep(j,i,i+step) {
      int &u = a[j], &v = a[j + step]; tie(u, v)
        inv ? pii(v - u, u) : pii(v, u + v); //
        inv ? pii(v, u - v) : pii(u + v, u); //
         \hookrightarrow OR
        pii(u + v, u - v);

→ XOR

 if (inv) trav(x, a) x = sz(a); // XOR only
vi conv(vi a, vi b) {
 FST(a, 0); FST(b, 0);
 rep(i, 0, sz(a)) a[i] *= b[i];
 FST(a, 1); return a;
```

Number theory (5)

5.1 Modular arithmetic

```
const 11 mod = 17; // change to something else
struct Mod {
  11 x;
  Mod(ll xx) : x(xx) \{ \}
  Mod operator + (Mod b) { return Mod((x + b.x) %
  \rightarrow mod); }
  Mod operator-(Mod b) { return Mod((x - b.x +
  \rightarrow mod) % mod); }
  Mod operator* (Mod b) { return Mod((x * b.x) %
  \rightarrow mod); }
  Mod operator/(Mod b) { return *this *

    invert(b); }

  Mod invert (Mod a) {
    ll x, y, q = \text{euclid}(a.x, \text{mod}, x, y);
    assert(g == 1); return Mod((x + mod) % mod);
  Mod operator (11 e) {
    if (!e) return Mod(1);
    Mod r = *this ^ (e / 2); r = r * r;
    return e&1 ? *this * r : r;
};
ModInverse.h
const 11 mod = 1000000007, LIM = 200000;
11* inv = new 11[LIM] - 1; inv[1] = 1;
rep(i, 2, LIM) inv[i] = mod - (mod / i) * inv[mod %]

    il % mod;

ModPow.h
const 11 mod = 1000000007; // faster if const
11 modpow(ll a, ll e) {
if (e == 0) return 1;
11 \times = modpow(a * a % mod, e >> 1);
 return e & 1 ? x * a % mod : x;
```

FastSubsetTransform.h

Modular Arithmetic.h

ModSum.h

ModSqrt.h

```
typedef unsigned long long ull;
ull sumsq(ull to) { return to / 2 * ((to-1) | 1);
ull divsum(ull to, ull c, ull k, ull m) {
 ull res = k / m * sumsq(to) + c / m * to;
 k %= m; c %= m;
 if (k) {
   ull to2 = (to * k + c) / m;
   res += to * to2;
   res -= divsum(to2, m-1 - c, m, k) + to2;
  return res;
11 modsum(ull to, ll c, ll k, ll m) {
 c = ((c % m) + m) % m;
 k = ((k % m) + m) % m;
 return to * c + k * sumsq(to) - m * divsum(to,
  \hookrightarrow c, k, m);
ModMulLL.h
typedef unsigned long long ull;
const int bits = 10;
// if all numbers are less than 2 k, set bits =
const ull po = 1 << bits;</pre>
ull mod_mul(ull a, ull b, ull &c) {
 ull x = a * (b & (po - 1)) % c;
 while ((b >>= bits) > 0) {
   a = (a << bits) % c;
   x += (a * (b & (po - 1))) % c;
  return x % c;
ull mod_pow(ull a, ull b, ull mod) {
 if (b == 0) return 1;
 ull res = mod pow(a, b / 2, mod);
  res = mod_mul(res, res, mod);
 if (b & 1) return mod_mul(res, a, mod);
 return res:
```

```
11 sqrt(ll a, ll p) {
 a \% = p; if (a < 0) a += p;
 if (a == 0) return 0;
 assert (modpow(a, (p-1)/2, p) == 1);
 if (p % 4 == 3) return modpow(a, (p+1)/4, p);
  // a^{(n+3)/8} \text{ or } 2^{(n+3)/8} * 2^{(n-1)/4} \text{ works if}
  → p % 8 == 5
 11 s = p - 1;
 int r = 0;
 while (s % 2 == 0)
   ++r, s /= 2;
 11 n = 2; // find a non-square mod p
 while (modpow(n, (p - 1) / 2, p) != p - 1) ++n;
 11 x = modpow(a, (s + 1) / 2, p);
 11 b = modpow(a, s, p);
 11 q = modpow(n, s, p);
 for (;;) {
    11 t = b;
    int m = 0;
    for (; m < r; ++m) {
      if (t == 1) break;
      t = t * t % p;
    if (m == 0) return x;
    11 \text{ qs} = \text{modpow}(q, 1 << (r - m - 1), p);
    q = qs * qs % p;
    x = x * qs % p;
    b = b * q % p;
```

5.2 Primality

eratosthenes.h

```
12
const int MAX_PR = 5000000;
bitset<MAX PR> isprime;
vi eratosthenes_sieve(int lim) {
  isprime.set(); isprime[0] = isprime[1] = 0;
  for (int i = 4; i < lim; i += 2) isprime[i] =</pre>
  for (int i = 3; i*i < lim; i += 2) if
  for (int j = i*i; j < lim; j += i*2)</pre>
    \rightarrow isprime[i] = 0;
  vi pr;
  rep(i,2,lim) if (isprime[i]) pr.push_back(i);
  return pr;
MillerRabin.h
bool prime(ull p) {
  if (p == 2) return true;
  if (p == 1 || p % 2 == 0) return false;
 ull s = p - 1;
  while (s % 2 == 0) s /= 2;
 rep(i, 0, 15) {
   ull a = rand() % (p - 1) + 1, tmp = s;
    ull mod = mod_pow(a, tmp, p);
    while (tmp != p - 1 \&\& mod != 1 \&\& mod != p -
    mod = mod mul(mod, mod, p);
     tmp *= 2;
```

if (mod != p - 1 && tmp % 2 == 0) **return**

false;

return true;

factor.h

```
vector<ull> pr;
ull f(ull a, ull n, ull &has) {
  return (mod_mul(a, a, n) + has) % n;
vector<ull> factor(ull d) {
  vector<ull> res;
  for (int i = 0; i < sz(pr) && pr[i]*pr[i] <= d;</pre>
  \hookrightarrow i++)
    if (d % pr[i] == 0) {
      while (d % pr[i] == 0) d /= pr[i];
      res.push_back(pr[i]);
  //d is now a product of at most 2 primes.
 if (d > 1) {
    if (prime(d))
      res.push back(d);
    else while (true)
      ull has = rand() % 2321 + 47;
      ull x = 2, y = 2, c = 1;
      for (; c==1; c = \_gcd((y > x ? y - x : x -
      \rightarrow y), d)) {
        x = f(x, d, has);
        y = f(f(y, d, has), d, has);
      if (c != d) {
        res.push_back(c); d /= c;
        if (d != c) res.push_back(d);
        break;
  return res;
void init(int bits) {//how many bits do we use?
  vi p = eratosthenes_sieve(1 << ((bits + 2) /</pre>

→ 3));
  pr.assign(all(p));
```

5.3 Divisibility

euclid.h

```
11 gcd(ll a, ll b) { return __gcd(a, b); }
ll euclid(ll a, ll b, ll &x, ll &y) {
 if (b) { ll d = euclid(b, a % b, y, x);
   return y = a/b * x, d; }
 return x = 1, y = 0, a;
Euclid.java
static BigInteger[] euclid(BigInteger a,
→ BigInteger b)
 BigInteger x = BigInteger.ONE, yy = x;
 BigInteger y = BigInteger.ZERO, xx = y;
 while (b.signum() != 0) {
   BigInteger q = a.divide(b), t = b;
   b = a.mod(b); a = t;
   t = xx; xx = x.subtract(q.multiply(xx)); x =
   t = yy; yy = y.subtract(q.multiply(yy)); y =
 return new BigInteger[]{x, y, a};
```

5.3.1 Bézout's identity

For $a \neq b \neq 0$, then d = gcd(a, b) is the smallest positive integer for which there are integer solutions to

$$ax + by = d$$

If (x, y) is one solution, then all solutions are given by

$$\left(x + \frac{kb}{\gcd(a,b)}, y - \frac{ka}{\gcd(a,b)}\right), \quad k \in \mathbb{Z}$$

phiFunction.h

5.4 Fractions

ContinuedFractions.h

```
typedef double d; // for N ~ 1e7; long double for
\hookrightarrow N ~ 1e9
pair<11, 11> approximate(d x, 11 N) {
  11 LP = 0, LQ = 1, P = 1, Q = 0, inf = 0
  \hookrightarrow LLONG_MAX; d y = x;
  for (;;) {
    ll lim = min(P ? (N-LP) / P : inf, Q ? (N-LQ)
     \hookrightarrow / \bigcirc : inf),
        a = (ll) floor(y), b = min(a, lim),
       NP = b * P + LP, NQ = b * Q + LQ;
    if (a > b) {
      // If b > a/2, we have a semi-convergent
       → that gives us a
      // better approximation; if b = a/2, we
       → *may* have one.
      // Return {P, Q} here for a more canonical
       → approximation.
      return (abs (x - (d) NP / (d) NQ) < abs <math>(x - (d) NQ)
       \rightarrow (d) P / (d)Q)) ?
         make_pair(NP, NQ) : make_pair(P, Q);
    if (abs(y = 1/(y - (d)a)) > 3*N) {
      return {NP, NQ};
    LP = P; P = NP;
    LQ = Q; Q = NQ;
```

FracBinarySearch.h

```
struct Frac { ll p, q; };
template < class F>
Frac fracBS(F f, ll N) {
 bool dir = 1, A = 1, B = 1;
  Frac lo{0, 1}, hi{1, 1}; // Set hi to 1/0 to
  \rightarrow search (0, N]
  assert(!f(lo)); assert(f(hi));
  while (A | | B) {
    11 adv = 0, step = 1; // move hi if dir, else
    for (int si = 0; step; (step *= 2) >>= si) {
      adv += step;
      Frac mid{lo.p * adv + hi.p, lo.g * adv +

    hi.q};
      if (abs(mid.p) > N || mid.q > N || dir ==
      \rightarrow !f(mid)) {
        adv -= step; si = 2;
    hi.p += lo.p * adv;
    hi.q += lo.q * adv;
    dir = !dir;
    swap(lo, hi);
    A = B; B = !!adv;
  return dir ? hi : lo;
```

5.5 Chinese remainder theorem

chinese.h

5.6 Pythagorean Triples

The Pythagorean triples are uniquely generated by

$$a = k \cdot (m^2 - n^2), b = k \cdot (2mn), c = k \cdot (m^2 + n^2),$$

with m > n > 0, k > 0, $m \perp n$, and either m or n even.

5.7 Primes

p=962592769 is such that $2^{21}\mid p-1$, which may be useful. For hashing use 970592641 (31-bit number), 31443539979727 (45-bit), 3006703054056749 (52-bit). There are 78498 primes less than 1000000.

Primitive roots exist modulo any prime power p^a , except for p=2, a>2, and there are $\phi(\phi(p^a))$ many. For p=2, a>2, the group $\mathbb{Z}_{2^a}^{\times}$ is instead isomorphic to $\mathbb{Z}_2 \times \mathbb{Z}_{2^{a-2}}$.

5.8 Estimates

$$\sum_{d|n} d = O(n \log \log n).$$

The number of divisors of n is at most around 100 for n < 5e4, 500 for n < 1e7, 2000 for n < 1e10, 200 000 for n < 1e19.

Combinatorial (6)

6.1 Permutations

6.1.1 Factorial

n	1 2 3	4	5 6	7	8	9	10	
$\overline{n!}$	1 2 6	24 1	20 72	0 5040	40320	362880	3628800	-
						5 16		
n!	4.0e7	4.8e	8 6.2e	9 8.7e	10 1.3	e12 2.1e	13 3.6e14 0 171	
n	20	25	30	40	50 1	100 15	0 171	
n!	2e18	2e25	3e32	8e47 :	3e64 9e	$e157 \ 6e2$	$62 > DBL_N$	1AX

IntPerm.h

6.1.2 Cycles

Let $g_S(n)$ be the number of *n*-permutations whose cycle lengths all belong to the set S. Then

$$\sum_{n=0}^{\infty} g_S(n) \frac{x^n}{n!} = \exp\left(\sum_{n \in S} \frac{x^n}{n}\right)$$

6.1.3 Derangements

Permutations of a set such that none of the elements appear in their original position.

$$D(n) = (n-1)(D(n-1) + D(n-2)) = nD(n-1) + (-1)^n = \left\lfloor \frac{n!}{e} \right\rfloor$$

6.1.4 Burnside's lemma

Given a group G of symmetries and a set X, the number of elements of X up to symmetry equals

$$\frac{1}{|G|} \sum_{g \in G} |X^g|,$$

where X^g are the elements fixed by g (g.x = x).

If f(n) counts "configurations" (of some sort) of length n, we can ignore rotational symmetry using $G = \mathbb{Z}_n$ to get

$$g(n) = \frac{1}{n} \sum_{k=0}^{n-1} f(\gcd(n,k)) = \frac{1}{n} \sum_{k|n} f(k)\phi(n/k).$$

6.2 Partitions and subsets

Partition function 6.2.1

Number of ways of writing n as a sum of positive integers, disregarding the order of the summands.

$$p(0) = 1, \ p(n) = \sum_{k \in \mathbb{Z} \setminus \{0\}} (-1)^{k+1} p(n - k(3k - 1)/2)$$

$$p(n) \sim 0.145/n \cdot \exp(2.56\sqrt{n})$$

Binomials

binomialModPrime.h

```
ll chooseModP(ll n, ll m, int p, vi& fact, vi&
→ invfact) {
 11 c = 1;
  while (n || m) {
   11 a = n \% p, b = m \% p;
    if (a < b) return 0;</pre>
    c = c * fact[a] % p * invfact[b] % p *

    invfact[a - b] % p;

   n /= p; m /= p;
  return c;
```

multinomial.h

General purpose numbers 6.3

6.3.1Stirling numbers of the first kind

Number of permutations on n items with k cycles.

$$c(n,k) = c(n-1,k-1) + (n-1)c(n-1,k), \ c(0,0) = 1$$

$$\sum_{k=0}^{n} c(n,k)x^{k} = x(x+1)\dots(x+n-1)$$

c(8, k) = 8, 0, 5040, 13068, 13132, 6769, 1960, 322, 28, 1 $c(n,2) = 0, 0, 1, 3, 11, 50, 274, 1764, 13068, 109584, \dots$

6.3.2 Eulerian numbers

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. k j:s s.t. $\pi(j) > \pi(j+1), k+1 \text{ } j:s \text{ s.t. } \pi(j) \geq j, k \text{ } j:s \text{ s.t.}$ $\pi(j) > j$.

$$E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)$$

$$E(n,0) = E(n,n-1) = 1$$

$$E(n,k) = \sum_{j=0}^{k} (-1)^{j} \binom{n+1}{j} (k+1-j)^{n}$$

Stirling numbers of the second 6.3.3 kind

Partitions of n distinct elements into exactly kgroups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k)$$

$$S(n,1) = S(n,n) = 1$$

$$S(n,k) = \frac{1}{k!} \sum_{j=0}^{k} (-1)^{k-j} \binom{k}{j} j^{n}$$

6.3.4 Bell numbers

Total number of partitions of n distinct elements. B(n) = $1, 1, 2, 5, 15, 52, 203, 877, 4140, 21147, \ldots$ For p prime,

$$B(p^m + n) \equiv mB(n) + B(n+1) \pmod{p}$$

Catalan numbers 6.3.5

$$C_n = \frac{1}{n+1} {2n \choose n} = {2n \choose n} - {2n \choose n+1} = \frac{(2n)!}{(n+1)!n!}$$

$$C_0 = 1, \ C_{n+1} = \frac{2(2n+1)}{n+2}C_n, \ C_{n+1} = \sum_{i=1}^{n} C_i C_{n-i}$$

 $C_n = 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, \dots$

- sub-diagonal monotone paths in an $n \times n$ grid.
- strings with n pairs of parenthesis, correctly nested.
- binary trees with with n+1 leaves (0 or 2 children).
- ordered trees with n+1 vertices.
- ways a convex polygon with n+2 sides can be cut into triangles by connecting vertices with straight lines.
- permutations of [n] with no 3-term increasing subseq.

Graph (7)

Fundamentals

bellmanFord.h

TopoSort.h

vector<pii> outs; // (dest, edge index)

struct V {

int nins = 0;

16

```
typedef 11 T; // or whatever
struct Edge { int src, dest; T weight; };
struct Node { T dist; int prev; };
struct Graph { vector<Node> nodes; vector<Edge>

    edges; };

const T inf = numeric_limits<T>::max();
bool bellmanFord2 (Graph& q, int start node)
  trav(n, g.nodes) { n.dist = inf; n.prev = -1; }
  g.nodes[start node].dist = 0;
  rep(i,0,sz(q.nodes)) trav(e, q.edges) {
   Node& cur = g.nodes[e.src];
    Node& dest = q.nodes[e.dest];
   if (cur.dist == inf) continue;
    T ndist = cur.dist + (cur.dist == -inf ? 0 :
    if (ndist < dest.dist) {</pre>
      dest.prev = e.src;
      dest.dist = (i \ge sz(g.nodes) - 1 ? -inf :

→ ndist);
  }
  bool ret = 0;
  rep(i,0,sz(q.nodes)) trav(e, q.edges) {
   if (g.nodes[e.src].dist == -inf)
      q.nodes[e.dest].dist = -inf, ret = 1;
  return ret;
FloydWarshall.h
const 11 inf = 1LL << 62;</pre>
void floydWarshall(vector<vector<ll>>>& m) {
  int n = sz(m);
  rep(i, 0, n) m[i][i] = min(m[i][i], {});
  rep(k, 0, n) rep(i, 0, n) rep(j, 0, n)
   if (m[i][k] != inf && m[k][j] != inf) {
      auto newDist = max(m[i][k] + m[k][j],
      \rightarrow -inf);
      m[i][j] = min(m[i][j], newDist);
  rep(k, 0, n) if (m[k][k] < 0) rep(i, 0, n)
  \rightarrow rep(j,0,n)
   if (m[i][k] != inf && m[k][j] != inf) m[i][j]
    \Rightarrow = -inf;
```

```
template<class E, class I>
bool topo_sort(const E &edges, I &idx, int n) {
 vi indeq(n);
  rep(i,0,n)
   trav(e, edges[i])
      indeg[e]++;
  queue<int> q; // use priority queue for lexic.
  \hookrightarrow smallest ans.
  rep(i,0,n) if (indeg[i] == 0) q.push(-i);
  int nr = 0;
 while (q.size() > 0) {
   int i = -q.front(); // top() for priority

→ queue

    idx[i] = nr++;
   q.pop();
   trav(e, edges[i])
      if (--indeg[e] == 0) q.push(-e);
 return nr == n;
```

7.2 Euler walk

```
};
vi euler_walk(vector<V>& nodes, int nedges, int
\hookrightarrow src=0) {
 int c = 0;
 trav(n, nodes) c += abs(n.nins - sz(n.outs));
 if (c > 2) return {};
 vector<vector<pii>::iterator> its;
  trav(n, nodes)
    its.push_back(n.outs.begin());
 vector<bool> eu(nedges);
 vi ret, s = \{src\};
 while(!s.emptv()) {
    int x = s.back();
    auto& it = its[x], end = nodes[x].outs.end();
    while(it != end && eu[it->second]) ++it;
    if(it == end) { ret.push_back(x);

    s.pop_back(); }

    else { s.push_back(it->first); eu[it->second]
    \hookrightarrow = true; }
 if(sz(ret) != nedges+1)
    ret.clear(); // No Eulerian cycles/paths.
 // else, non-cycle if ret.front() != ret.back()
 reverse(all(ret));
 return ret;
```

7.3 Network flow

EulerWalk.h

PushRelabel.h

```
typedef ll Flow;
struct Edge {
 int dest, back;
 Flow f, c;
} ;
struct PushRelabel {
 vector<vector<Edge>> g;
 vector<Flow> ec;
 vector<Edge*> cur;
 vector<vi> hs; vi H;
  PushRelabel(int n) : g(n), ec(n), cur(n),
  \rightarrow hs(2*n), H(n) {}
  void add_edge(int s, int t, Flow cap, Flow
  \hookrightarrow rcap=0) {
   if (s == t) return;
   Edge a = \{t, sz(q[t]), 0, cap\};
   Edge b = \{s, sz(g[s]), 0, rcap\};
   g[s].push_back(a);
   g[t].push_back(b);
 void add_flow(Edge& e, Flow f) {
   Edge &back = g[e.dest][e.back];
   if (!ec[e.dest] && f)

→ hs[H[e.dest]].push_back(e.dest);
   e.f += f; e.c -= f; ec[e.dest] += f;
   back.f -= f; back.c += f; ec[back.dest] -= f;
  Flow maxflow(int s, int t) {
   int v = sz(g); H[s] = v; ec[t] = 1;
   vi co(2*v); co[0] = v-1;
   rep(i, 0, v) cur[i] = g[i].data();
   trav(e, g[s]) add_flow(e, e.c);
    for (int hi = 0;;) {
      while (hs[hi].empty()) if (!hi--) return
      \rightarrow -ec[s];
      int u = hs[hi].back(); hs[hi].pop_back();
      while (ec[u] > 0) // discharge u
        if (cur[u] == g[u].data() + sz(g[u])) {
          H[u] = 1e9;
          trav(e, g[u]) if (e.c && H[u] >
          \rightarrow H[e.dest]+1)
            H[u] = H[e.dest]+1, cur[u] = &e;
          if (++co[H[u]], !--co[hi] && hi < v)</pre>
            rep(i, 0, v) if (hi < H[i] && H[i] < v)
              --co[H[i]], H[i] = v + 1;
          hi = H[u];
        } else if (cur[u]->c && H[u] ==
```

 $U \left[\alpha u x \left[u \right] - \lambda d \alpha c + 1 + 1 \right)$

MinCostMaxFlow.h

```
#include <bits/extc++.h>
const ll INF = numeric_limits<ll>::max() / 4;
typedef vector<ll> VL;
struct MCMF {
  int N;
 vector<vi> ed, red;
 vector<VL> cap, flow, cost;
 vi seen;
 VL dist, pi;
 vector<pii> par;
 MCMF (int N) :
   N(N), ed(N), red(N), cap(N, VL(N)),

→ flow(cap), cost(cap),
   seen(N), dist(N), pi(N), par(N) {}
 void addEdge(int from, int to, ll cap, ll cost)
  ← {
   this->cap[from][to] = cap;
   this->cost[from][to] = cost;
   ed[from].push_back(to);
   red[to].push_back(from);
 void path(int s) {
   fill(all(seen), 0);
   fill(all(dist), INF);
   dist[s] = 0; ll di;
    __gnu_pbds::priority_queue<pair<11, int>> q;
   vector<decltype(g)::point_iterator> its(N);
   q.push({0, s});
   auto relax = [&](int i, ll cap, ll cost, int

    dir) {

     ll val = di - pi[i] + cost;
     if (cap && val < dist[i]) {
        dist[i] = val;
       par[i] = \{s, dir\};
        if (its[i] == q.end()) its[i] =
        \rightarrow q.push(\{-dist[i], i\});
        else q.modify(its[i], {-dist[i], i});
   };
   while (!q.empty()) {
     s = q.top().second; q.pop();
      seen[s] = 1; di = dist[s] + pi[s];
      trav(i, ed[s]) if (!seen[i])
       rolay/i gan[gl[i] - flow[gl[i]
```

```
EdmondsKarp.h
template<class T> T
→ edmondsKarp(vector<unordered_map<int, T>>&
assert(source != sink);
 T flow = 0;
 vi par(sz(graph)), q = par;
 for (;;) {
   fill(all(par), -1);
   par[source] = 0;
   int ptr = 1;
   q[0] = source;
   rep(i,0,ptr) {
     int x = q[i];
     trav(e, graph[x]) {
       if (par[e.first] == -1 \&\& e.second > 0) {
         par[e.first] = x;
         q[ptr++] = e.first;
         if (e.first == sink) goto out;
   return flow;
out:
   T inc = numeric_limits<T>::max();
   for (int y = sink; y != source; y = par[y])
     inc = min(inc, graph[par[y]][y]);
   flow += inc;
   for (int y = sink; y != source; y = par[y]) {
     int p = par[y];
     if ((graph[p][y] -= inc) <= 0)

    graph[p].erase(y);
     graph[y][p] += inc;
 }
MinCut.h
GlobalMinCut.h
```

```
pair<int, vi> GetMinCut(vector<vi>& weights) {
  int N = sz(weights);
  vi used(N), cut, best_cut;
 int best_weight = -1;
 for (int phase = N-1; phase >= 0; phase--) {
    vi w = weights[0], added = used;
   int prev, k = 0;
    rep(i,0,phase){
      prev = k;
      k = -1;
      rep(j,1,N)
        if (!added[j] && (k == -1 || w[j] >
        \hookrightarrow w[k])) k = \dot{\gamma};
      if (i == phase-1) {
        rep(j,0,N) weights[prev][j] +=
        → weights[k][j];
        rep(j,0,N) weights[j][prev] =

    weights[prev][j];

        used[k] = true;
        cut.push_back(k);
        if (best_weight == -1 \mid \mid w[k] <

    best_weight) {

          best_cut = cut;
          best_weight = w[k];
      } else {
        rep(j,0,N)
          w[j] += weights[k][j];
        added[k] = true;
  return {best_weight, best_cut};
```

7.4 Matching

```
bool dfs(int a, int layer, const vector<vi>& q,

    vi & btoa,

      vi& A, vi& B) {
 if (A[a] != layer) return 0;
 A[a] = -1;
 trav(b, q[a]) if (B[b] == layer + 1) {
   B[b] = -1;
   if (btoa[b] == -1 || dfs(btoa[b], layer+2, q,
    \hookrightarrow btoa, A, B))
      return btoa[b] = a, 1;
 return 0;
int hopcroftKarp(const vector<vi>& q, vi& btoa) {
 int res = 0;
 vi A(g.size()), B(btoa.size()), cur, next;
 for (;;) {
   fill(all(A), 0);
   fill(all(B), -1);
   cur.clear();
    trav(a, btoa) if(a !=-1) A[a] = -1;
    rep(a, 0, sz(q)) if(A[a] == 0)

    cur.push_back(a);

    for (int lay = 1;; lay += 2) {
      bool islast = 0;
      next.clear();
      trav(a, cur) trav(b, g[a]) {
        if (btoa[b] == -1) {
          B[b] = lay;
          islast = 1;
        else if (btoa[b] != a && B[b] == -1) {
          B[b] = lay;
          next.push_back(btoa[b]);
      if (islast) break;
      if (next.empty()) return res;
      trav(a, next) A[a] = lay+1;
      cur.swap(next);
    rep(a, 0, sz(q)) {
      if(dfs(a, 0, g, btoa, A, B))
        ++res;
```

```
DFSMatching.h
vi match;
vector<bool> seen;
bool find(int j, const vector<vi>& g) {
  if (match[j] == -1) return 1;
  seen[j] = 1; int di = match[j];
  trav(e, q[di])
    if (!seen[e] && find(e, g)) {
      match[e] = di;
      return 1;
  return 0;
int dfs_matching(const vector<vi>& q, int n, int
\hookrightarrow m) {
 match.assign(m, -1);
  rep(i,0,n) {
    seen.assign(m, 0);
    trav(j,g[i])
      if (find(j, g)) {
        match[j] = i;
        break;
  return m - (int) count (all (match), -1);
```

```
typedef vector<double> vd;
bool zero(double x) { return fabs(x) < 1e-10; }</pre>
double MinCostMatching(const vector<vd>& cost,
int n = sz(cost), mated = 0;
 vd dist(n), u(n), v(n);
 vi dad(n), seen(n);
  rep(i,0,n) {
   u[i] = cost[i][0];
    rep(j,1,n) u[i] = min(u[i], cost[i][j]);
  rep(j,0,n) {
   v[j] = cost[0][j] - u[0];
   rep(i,1,n) v[j] = min(v[j], cost[i][j] -
    \hookrightarrow u[i]);
 L = R = vi(n, -1);
  rep(i, 0, n) rep(j, 0, n) {
   if (R[\dot{j}] != -1) continue;
   if (zero(cost[i][j] - u[i] - v[j])) {
     L[i] = i;
      R[j] = i;
      mated++;
      break;
  for (; mated < n; mated++) { // until solution</pre>
  \hookrightarrow is feasible
   int s = 0;
   while (L[s] !=-1) s++;
    fill(all(dad), -1);
    fill(all(seen), 0);
    rep(k,0,n)
      dist[k] = cost[s][k] - u[s] - v[k];
    int j = 0;
    for (;;) {
      j = -1;
      rep(k,0,n){
        if (seen[k]) continue;
        if (j == -1 \mid \mid dist[k] < dist[j]) j = k;
      seen[j] = 1;
      int i = R[j];
      if (i == -1) break;
      rep(k,0,n) {
        if (seen[k]) continue;
        auto new_dist = dist[j] + cost[i][k] -
           11 [ i ] _ 17 [ le ] •
```

```
GeneralMatching.h
vector<pii> generalMatching(int N, vector<pii>&
\rightarrow ed) {
 vector<vector<ll>> mat(N, vector<ll>(N)), A;
 trav(pa, ed) {
    int a = pa.first, b = pa.second, r = rand() %

→ mod;

    mat[a][b] = r, mat[b][a] = (mod - r) % mod;
  int r = matInv(A = mat), M = 2*N - r, fi, fj;
  assert (r % 2 == 0);
  if (M != N) do {
    mat.resize(M, vector<11>(M));
    rep(i,0,N) {
      mat[i].resize(M);
      rep(j, N, M) {
        int r = rand() % mod;
        mat[i][j] = r, mat[j][i] = (mod - r) %

→ mod;

  } while (matInv(A = mat) != M);
  vi has(M, 1); vector<pii> ret;
  rep(it, 0, M/2) {
    rep(i,0,M) if (has[i])
      rep(j,i+1,M) if (A[i][j] && mat[i][j]) {
        fi = i; fj = j; goto done;
    } assert(0); done:
    if (fj < N) ret.emplace_back(fi, fj);</pre>
    has[fi] = has[fj] = 0;
    rep(sw, 0, 2) {
      11 a = modpow(A[fi][fj], mod-2);
      rep(i,0,M) if (has[i] && A[i][fj]) {
        ll b = A[i][fj] * a % mod;
        rep(j, 0, M) A[i][j] = (A[i][j] - A[fi][j]
         \leftrightarrow * b) % mod;
      swap(fi,fj);
  return ret;
MinimumVertexCover.h
```

```
vi cover(vector<vi>& q, int n, int m) {
 int res = dfs_matching(g, n, m);
 seen.assign(m, false);
 vector<bool> lfound(n, true);
 trav(it, match) if (it !=-1) lfound[it] =

    false;

 vi q, cover;
 rep(i,0,n) if (lfound[i]) q.push_back(i);
 while (!q.empty()) {
   int i = q.back(); q.pop_back();
   lfound[i] = 1;
   trav(e, q[i]) if (!seen[e] && match[e] != -1)
     seen[e] = true;
     q.push_back(match[e]);
 rep(i,0,n) if (!lfound[i]) cover.push_back(i);
 rep(i,0,m) if (seen[i]) cover.push_back(n+i);
 assert(sz(cover) == res);
 return cover;
```

7.5 DFS algorithms

```
SCC.h
vi val, comp, z, cont;
int Time, ncomps;
template < class G, class F> int dfs (int j, G& g, F
→ f) {
 int low = val[j] = ++Time, x; z.push_back(j);
 trav(e,q[j]) if (comp[e] < 0)
   low = min(low, val[e] ?: dfs(e,g,f));
 if (low == val[j]) {
    do {
      x = z.back(); z.pop_back();
      comp[x] = ncomps;
      cont.push_back(x);
   } while (x != j);
    f(cont); cont.clear();
   ncomps++;
  return val[j] = low;
template < class G, class F > void scc(G& g, F f) {
  int n = sz(q);
 val.assign(n, 0); comp.assign(n, -1);
 Time = ncomps = 0;
  rep(i, 0, n) if (comp[i] < 0) dfs(i, q, f);
```

```
vi num, st;
vector<vector<pii>>> ed;
int Time;
template<class F>
int dfs(int at, int par, F f) {
 int me = num[at] = ++Time, e, y, top = me;
 trav(pa, ed[at]) if (pa.second != par) {
   tie(y, e) = pa;
   if (num[y]) {
     top = min(top, num[y]);
     if (num[y] < me)
        st.push_back(e);
    } else {
      int si = sz(st);
      int up = dfs(y, e, f);
      top = min(top, up);
      if (up == me) {
        st.push_back(e);
        f(vi(st.begin() + si, st.end()));
        st.resize(si);
      else if (up < me) st.push_back(e);</pre>
      else { /* e is a bridge */ }
 return top;
template<class F>
void bicomps(F f) {
 num.assign(sz(ed), 0);
 rep(i, 0, sz(ed)) if (!num[i]) dfs(i, -1, f);
```

```
struct TwoSat {
 int N;
 vector<vi> gr;
 vi values; // 0 = false, 1 = true
 TwoSat(int n = 0) : N(n), gr(2*n) {}
 int add_var() { // (optional)
    gr.emplace_back();
    gr.emplace_back();
    return N++;
 void either(int f, int j) {
    f = \max(2 * f, -1 - 2 * f);
    j = \max(2*j, -1-2*j);
    gr[f^1].push_back(j);
    gr[j^1].push_back(f);
  void set_value(int x) { either(x, x); }
 void at_most_one(const vi& li) { // (optional)
    if (sz(li) <= 1) return;</pre>
    int cur = ~li[0];
    rep(i,2,sz(li)) {
      int next = add_var();
      either(cur, ~li[i]);
      either(cur, next);
      either(~li[i], next);
      cur = ~next;
    either(cur, ~li[1]);
 vi val, comp, z; int time = 0;
 int dfs(int i) {
    int low = val[i] = ++time, x; z.push_back(i);
    trav(e, qr[i]) if (!comp[e])
      low = min(low, val[e] ?: dfs(e));
    ++time;
    if (low == val[i]) do {
      x = z.back(); z.pop_back();
      comp[x] = time;
      if (values[x >> 1] == -1)
        values[x>>1] = !(x&1);
    } while (x != i);
    return val[i] = low;
 bool solve()
    values.assign(N, -1);
    \text{val} = \text{agaign} (2 + \mathbf{N} - 0) \cdot \text{gamp} = \text{val} \cdot
```

7.6 Heuristics

7.7 Trees

```
TreePower.h
vector<vi> treeJump(vi& P) {
  int on = 1, d = 1;
  while (on < sz(P)) on *= 2, d++;
  vector<vi> jmp(d, P);
  rep(i,1,d) rep(j,0,sz(P))
    jmp[i][j] = jmp[i-1][jmp[i-1][j]];
  return jmp;
int jmp(vector<vi>& tbl, int nod, int steps) {
  rep(i, 0, sz(tbl))
   if(steps&(1<<i)) nod = tbl[i][nod];
  return nod;
int lca(vector<vi>& tbl, vi& depth, int a, int b)
← {
  if (depth[a] < depth[b]) swap(a, b);</pre>
  a = jmp(tbl, a, depth[a] - depth[b]);
  if (a == b) return a;
  for (int i = sz(tbl); i--;) {
   int c = tbl[i][a], d = tbl[i][b];
   if (c != d) a = c, b = d;
 return tbl[0][a];
```

```
LCA.h
typedef vector<pii> vpi;
typedef vector<vpi> graph;
struct LCA {
 vi time;
 vector<ll> dist;
 RMQ<pii> rmq;
 LCA(graph\&C): time(sz(C), -99), dist(sz(C)),
  \rightarrow rmq(dfs(C)) {}
 vpi dfs(graph& C) {
   vector<tuple<int, int, int, ll>> q(1);
   vpi ret;
   int T = 0, v, p, d; ll di;
   while (!q.empty()) {
     tie(v, p, d, di) = q.back();
     q.pop_back();
      if (d) ret.emplace_back(d, p);
      time[v] = T++;
      dist[v] = di;
      trav(e, C[v]) if (e.first != p)
        q.emplace_back(e.first, v, d+1, di +

    e.second);
   return ret;
 int query(int a, int b) {
   if (a == b) return a;
   a = time[a], b = time[b];
   return rmq.query(min(a, b), max(a,
    11 distance(int a, int b) {
   int lca = query(a, b);
   return dist[a] + dist[b] - 2 * dist[lca];
 }
} ;
```

CompressTree.h

```
vpi compressTree(LCA& lca, const vi& subset) {
 static vi rev; rev.resize(sz(lca.dist));
 vi li = subset, &T = lca.time;
 auto cmp = [&](int a, int b) { return T[a] <</pre>
  \hookrightarrow T[b]; };
 sort(all(li), cmp);
 int m = sz(li)-1;
 rep(i,0,m) {
   int a = li[i], b = li[i+1];
   li.push_back(lca.query(a, b));
 sort(all(li), cmp);
 li.erase(unique(all(li)), li.end());
 rep(i, 0, sz(li)) rev[li[i]] = i;
 vpi ret = {pii(0, li[0])};
 rep(i, 0, sz(li) - 1) {
   int a = li[i], b = li[i+1];
    ret.emplace_back(rev[lca.query(a, b)], b);
 }
 return ret;
```

HLD.h

```
typedef vector<pii> vpi;
struct Node {
 int d, par, val, chain = -1, pos = -1;
};
struct Chain {
 int par, val;
 vector<int> nodes;
 Tree tree;
} ;
struct HLD {
 typedef int T;
  const T LOW = -(1 << 29);
 void f(T& a, T b) { a = max(a, b); }
 vector<Node> V;
 vector<Chain> C;
 HLD(vector<vpi>& g) : V(sz(g)) {
   dfs(0, -1, g, 0);
   trav(c, C) {
     c.tree = {sz(c.nodes), 0};
     for (int ni : c.nodes)
       c.tree.update(V[ni].pos, V[ni].val);
 void update(int node, T val) {
   Node& n = V[node]; n.val = val;
   if (n.chain != -1)
    int pard(Node& nod) {
   if (nod.par == -1) return -1;
   return V[nod.chain == -1 ? nod.par :
    // query all *edges* between n1, n2
 pair<T, int> query(int i1, int i2) {
   T ans = LOW;
   while(i1 != i2) {
     Node n1 = V[i1], n2 = V[i2];
     if (n1.chain != -1 && n1.chain == n2.chain)
       int lo = n1.pos, hi = n2.pos;
       if (lo > hi) swap(lo, hi);
       f(ans, C[n1.chain].tree.query(lo, hi));
       i1 - i2 - C[n1 \text{ abainl nodes [bile]}]
```

```
struct Node { // Splay tree. Root's pp contains

    tree's parent.

 Node *p = 0, *pp = 0, *c[2];
 bool flip = 0;
 Node() { c[0] = c[1] = 0; fix(); }
 void fix() {
   if (c[0]) c[0]->p = this;
   if (c[1]) c[1]->p = this;
   // (+ update sum of subtree elements etc. if
    → wanted)
 void push_flip() {
   if (!flip) return;
   flip = 0; swap(c[0], c[1]);
   if (c[0]) c[0]->flip ^= 1;
   if (c[1]) c[1]->flip ^= 1;
  int up() { return p ? p->c[1] == this : -1; }
 void rot(int i, int b) {
   int h = i \hat{b};
   Node *x = c[i], *y = b == 2 ? x : x -> c[h], *z
    \Rightarrow = b ? \forall : x;
   if ((y->p = p)) p->c[up()] = y;
   c[i] = z -> c[i ^ 1];
   if (b < 2) {
     x->c[h] = y->c[h ^ 1];
      z - c[h ^1] = b ? x : this;
   y - c[i ^1] = b ? this : x;
   fix(); x->fix(); y->fix();
   if (p) p->fix();
   swap(pp, y->pp);
 void splay() {
   for (push_flip(); p; ) {
     if (p->p) p->p->push_flip();
     p->push_flip(); push_flip();
     int c1 = up(), c2 = p->up();
     if (c2 == -1) p->rot(c1, 2);
     else p->p->rot(c2, c1 != c2);
   }
 Node* first() {
   push_flip();
   return c[0] ? c[0]->first() : (splay(),
    → this);
struct LinkCut
 vector<Node> node;
 Tinl(Cut (int N)) \cdot nodo(N) ()
```

```
MatrixTree.h
```

Geometry (8)

8.1 Geometric primitives

Point.h

```
template<class T>
struct Point {
 typedef Point P;
 T x, y;
  explicit Point (T x=0, T y=0) : x(x), y(y) {}
 bool operator<(P p) const { return tie(x,y) <</pre>
  \hookrightarrow tie(p.x,p.y); }
 bool operator==(P p) const { return
  \rightarrow tie(x,y) == tie(p.x,p.y); }
 P operator+(P p) const { return P(x+p.x,
  \rightarrow y+p.y); }
 P operator-(P p) const { return P(x-p.x,
  \hookrightarrow y-p.y); }
 P operator*(T d) const { return P(x*d, y*d); }
 P operator/(T d) const { return P(x/d, y/d); }
 T dot(P p) const { return x*p.x + y*p.y; }
 T cross(P p) const { return x*p.y - y*p.x; }
 T cross(P a, P b) const { return
  T dist2() const { return x*x + y*y; }
 double dist() const { return

    sqrt((double)dist2()); }

  // angle to x-axis in interval [-pi, pi]
  double angle() const { return atan2(y, x); }
 P unit() const { return *this/dist(); } //
  → makes dist()=1
 P perp() const { return P(-y, x); } // rotates
  → +90 degrees
 P normal() const { return perp().unit(); }
  // returns point rotated 'a' radians ccw around

→ the origin

 P rotate (double a) const {
    \rightarrow P(x*cos(a)-y*sin(a),x*sin(a)+y*cos(a)); }
};
lineDistance.h
```

```
SegmentDistance.h
typedef Point < double > P;
double segDist(P& s, P& e, P& p) {
 if (s==e) return (p-s).dist();
  auto d = (e-s).dist2(), t =
  \rightarrow min(d, max(.0, (p-s).dot(e-s)));
 return ((p-s)*d-(e-s)*t).dist()/d;
SegmentIntersection.h
template < class P>
int segmentIntersection (const P& s1, const P& e1,
    const P& s2, const P& e2, P& r1, P& r2) {
 if (e1==s1) {
   if (e2==s2) {
      if (e1==e2) { r1 = e1; return 1; } //all

→ equal

      else return 0; //different point segments
    } else return

    segmentIntersection(s2,e2,s1,e1,r1,r2);//s

  //segment directions and separation
  P v1 = e1-s1, v2 = e2-s2, d = s2-s1;
  auto a = v1.cross(v2), a1 = v1.cross(d), a2 =
  \rightarrow v2.cross(d):
 if (a == 0) { //if parallel
    auto b1=s1.dot(v1), c1=e1.dot(v1),
         b2=s2.dot(v1), c2=e2.dot(v1);
   if (a1 || a2 ||
    \rightarrow max(b1,min(b2,c2))>min(c1,max(b2,c2)))
      return 0;
    r1 = min(b2,c2) < b1 ? s1 : (b2 < c2 ? s2 : e2);
    r2 = max(b2,c2)>c1 ? e1 : (b2>c2 ? s2 : e2);
   return 2-(r1==r2);
 if (a < 0) { a = -a; a1 = -a1; a2 = -a2; }
 if (0<a1 || a<-a1 || 0<a2 || a<-a2)
   return 0;
 r1 = s1-v1*a2/a;
  return 1;
SegmentIntersectionQ.h
```

```
template<class P>
bool segmentIntersectionQ(P s1, P e1, P s2, P e2)
← {
 if (e1 == s1) {
    if (e2 == s2) return e1 == e2;
    swap(s1,s2); swap(e1,e2);
 }
  P v1 = e1-s1, v2 = e2-s2, d = s2-s1;
  auto a = v1.cross(v2), a1 = d.cross(v1), a2 =
  \rightarrow d.cross(v2);
  if (a == 0) { // parallel
    auto b1 = s1.dot(v1), c1 = e1.dot(v1),
         b2 = s2.dot(v1), c2 = e2.dot(v1);
    return !a1 && max(b1,min(b2,c2)) <=</pre>
    \rightarrow min(c1, max(b2, c2));
  if (a < 0) { a = -a; a1 = -a1; a2 = -a2; }
  return (0 <= a1 && a1 <= a && 0 <= a2 && a2 <=
  → a);
lineIntersection.h
template<class P>
int lineIntersection (const P& s1, const P& e1,
const P& s2,
    const P& e2, P& r) {
  if ((e1-s1).cross(e2-s2)) { //if not parallell
    \leftrightarrow s2-(e2-s2) * (e1-s1).cross(s2-s1) / (e1-s1).cross(e2-s2);
    return 1;
  } else
    return - ((e1-s1).cross(s2-s1) == 0 | | s2==e2);
sideOf.h
template<class P>
int sideOf(const P& s, const P& e, const P& p) {
  auto a = (e-s).cross(p-s);
  return (a > 0) - (a < 0);
template < class P>
int sideOf(const P& s, const P& e, const P& p,

→ double eps) {
 auto a = (e-s).cross(p-s);
  double 1 = (e-s).dist()*eps;
  return (a > 1) - (a < -1);
```

```
onSegment.h
template < class P>
bool onSegment (const P& s, const P& e, const P&
→ p) {
  P ds = p-s, de = p-e;
  return ds.cross(de) == 0 && ds.dot(de) <= 0;
linearTransformation.h
typedef Point < double > P;
P linearTransformation(const P& p0, const P& p1,
    const P& q0, const P& q1, const P& r) {
  P dp = p1-p0, dq = q1-q0, num(dp.cross(dq),

    dp.dot(dq));
  return q0 + P((r-p0).cross(num),
  \rightarrow (r-p0).dot(num))/dp.dist2();
Angle.h
```

```
struct Angle {
  int x, y;
  int t;
  Angle(int x, int y, int t=0) : x(x), y(y), t(t)
  Angle operator-(Angle b) const { return {x-b.x,
  \hookrightarrow y-b.y, t}; }
  int quad() const {
    assert(x || y);
    if (v < 0) return (x >= 0) + 2;
    if (y > 0) return (x <= 0);
    return (x <= 0) * 2;
  Angle t90() const { return \{-y, x, t + (quad())\}
  \hookrightarrow == 3)}; }
  Angle t180() const { return \{-x, -y, t +
  \rightarrow (quad() >= 2)}; }
  Angle t360() const { return {x, y, t + 1}; }
bool operator<(Angle a, Angle b) {</pre>
  // add a.dist2() and b.dist2() to also compare

→ distances

  return make_tuple(a.t, a.quad(), a.y * (11)b.x)
          make_tuple(b.t, b.quad(), a.x *
          \hookrightarrow (11)b.y);
// Given two points, this calculates the smallest
→ angle between
// them, i.e., the angle that covers the defined
→ line segment.
pair < Angle, Angle > segment Angles (Angle a, Angle
\rightarrow b) {
  if (b < a) swap(a, b);
  return (b < a.t180() ?
          make_pair(a, b) : make_pair(b,
           \rightarrow a.t360());
Angle operator+(Angle a, Angle b) { // point a +
\hookrightarrow vector b
  Angle r(a.x + b.x, a.y + b.y, a.t);
  if (a.t180() < r) r.t--;
  return r.t180() < a ? r.t360() : r;</pre>
Angle angleDiff(Angle a, Angle b) { // angle b -
→ angle a
  int tu = b.t - a.t; a.t = b.t;
  return {a.x*b.x + a.y*b.y, a.x*b.y - a.y*b.x,
  \rightarrow tu - (b < a) };
```

8.2 Circles

MinimumEnclosingCircle.h

```
CircleIntersection.h
typedef Point < double > P;
bool circleIntersection (P a, P b, double r1,

→ double r2,

    pair<P, P>* out) {
  P delta = b - a;
  assert (delta.x || delta.y || r1 != r2);
  if (!delta.x && !delta.y) return false;
  double r = r1 + r2, d2 = delta.dist2();
  double p = (d2 + r1*r1 - r2*r2) / (2.0 * d2);
  double h2 = r1*r1 - p*p*d2;
  if (d2 > r*r \mid \mid h2 < 0) return false;
  P mid = a + delta*p, per = delta.perp() *
  \rightarrow sqrt(h2 / d2);
  *out = {mid + per, mid - per};
  return true;
circleTangents.h
template < class P>
pair < P, P > circleTangents (const P &p, const P &c,
→ double r) {
 P a = p-c;
  double x = r*r/a.dist2(), y = sqrt(x-x*x);
  return make_pair(c+a*x+a.perp()*y,
  \hookrightarrow c+a*x-a.perp()*y);
circumcircle.h
typedef Point < double > P;
double ccRadius (const P& A, const P& B, const P&

→ C) {
  return (B-A).dist() * (C-B).dist() * (A-C).dist() /
      abs ((B-A) \cdot cross(C-A))/2;
P ccCenter(const P& A, const P& B, const P& C) {
 P b = C-A, c = B-A;
  return A +
  \rightarrow (b*c.dist2()-c*b.dist2()).perp()/b.cross(c)/2;
```

```
pair<double, P> mec2(vector<P>& S, P a, P b, int
\hookrightarrow n) {
  double hi = INFINITY, lo = -hi;
 rep(i,0,n) {
    auto si = (b-a).cross(S[i]-a);
    if (si == 0) continue;
    P m = ccCenter(a, b, S[i]);
    auto cr = (b-a).cross(m-a);
    if (si < 0) hi = min(hi, cr);
    else lo = max(lo, cr);
  double v = (0 < lo ? lo : hi < 0 ? hi : 0);
  P c = (a + b) / 2 + (b - a) .perp() * v / (b - a)
  \rightarrow a).dist2();
  return { (a - c).dist2(), c};
pair<double, P> mec(vector<P>& S, P a, int n) {
  random_shuffle(S.begin(), S.begin() + n);
  P b = S[0], c = (a + b) / 2;
  double r = (a - c).dist2();
  rep(i,1,n) if ((S[i] - c).dist2() > r * (1 + c)
  \rightarrow 1e-8)) {
   tie(r,c) = (n == sz(S) ?
      mec(S, S[i], i) : mec2(S, a, S[i], i));
  return {r, c};
pair<double, P> enclosingCircle(vector<P> S) {
  assert(!S.empty()); auto r = mec(S, S[0],
  \hookrightarrow sz(S));
  return {sqrt(r.first), r.second};
```

8.3 Polygons

insidePolygon.h

```
template<class It, class P>
bool insidePolygon (It begin, It end, const P& p,
    bool strict = true) {
  int n = 0; //number of isects with line from p
  \leftrightarrow to (inf,p.v)
  for (It i = begin, j = end-1; i != end; j =
  //if p is on edge of polygon
    if (onSegment(*i, *j, p)) return !strict;
    //or: if (segDist(*i, *j, p) \leq epsilon)
    → return !strict;
    //increment n if segment intersects line from
    n += (max(i->y, j->y) > p.y \&\& min(i->y, j->y)
    ((*j-*i).cross(p-*i) > 0) == (i->y <=

    p.y));
  return n&1; //inside if odd number of

    intersections

PolygonArea.h
template<class T>
T polygonArea2(vector<Point<T>>& v) {
  T = v.back().cross(v[0]);
  rep(i, 0, sz(v) -1) a += v[i].cross(v[i+1]);
  return a;
PolygonCenter.h
typedef Point < double > P;
Point < double > polygonCenter (vector < P > & v) {
  auto i = v.begin(), end = v.end(), j = end-1;
  Point<double> res{0,0}; double A = 0;
  for (; i != end; j=i++) {
    res = res + (*i + *j) * j \rightarrow cross(*i);
    A += j->cross(*i);
  return res / A / 3;
PolygonCut.h
```

```
typedef Point < double > P;
vector<P> polygonCut(const vector<P>& poly, P s,
  vector<P> res:
  rep(i, 0, sz(poly)) {
    P cur = poly[i], prev = i ? poly[i-1] :
    → poly.back();
    bool side = s.cross(e, cur) < 0;</pre>
    if (side != (s.cross(e, prev) < 0)) {</pre>
      res.emplace back();
      lineIntersection(s, e, cur, prev,

→ res.back());
    if (side)
      res.push_back(cur);
  return res:
ConvexHull.h
typedef Point<ll> P;
pair<vi, vi> ulHull(const vector<P>& S) {
 vi Q(sz(S)), U, L;
 iota(all(Q), 0);
 sort(all(Q), [&S](int a, int b) { return S[a] <</pre>
  \hookrightarrow S[b]; \});
 trav(it, Q) {
#define ADDP(C, cmp) while (sz(C) > 1 \&\&
\hookrightarrow S[C[sz(C)-2]].cross(\
 S[it], S[C.back()]) cmp 0) C.pop_back();
ADDP(U, <=); ADDP(L, >=);
 return {U, L};
vi convexHull(const vector<P>& S) {
 vi u, l; tie(u, l) = ulHull(S);
 if (sz(S) <= 1) return u;
 if (S[u[0]] == S[u[1]]) return {0};
 l.insert(l.end(), u.rbegin()+1, u.rend()-1);
 return 1:
PolygonDiameter.h
```

```
vector<pii> antipodal(const vector<P>& S, vi& U,

    vi& L) {

  vector<pii> ret;
  int i = 0, j = sz(L) - 1;
  while (i < sz(U) - 1 || j > 0) {
    ret.emplace_back(U[i], L[j]);
    if (j == 0 \mid | (i != sz(U) - 1 \&\& (S[L[j]] - 1))

→ S[L[i-1]])

          .cross(S[U[i+1]] - S[U[i]]) > 0)) ++i;
    else -- i;
  return ret;
pii polygonDiameter(const vector<P>& S) {
  vi U, L; tie(U, L) = ulHull(S);
  pair<ll, pii> ans;
  trav(x, antipodal(S, U, L))
    ans = max(ans, {(S[x.first] - 
    \rightarrow S[x.second]).dist2(), x});
  return ans.second;
PointInsideHull.h
typedef Point<ll> P;
int insideHull2 (const vector<P>& H, int L, int R,

→ const P& p) {
  int len = R - L;
  if (len == 2) {
    int sa = sideOf(H[0], H[L], p);
    int sb = sideOf(H[L], H[L+1], p);
    int sc = sideOf(H[L+1], H[0], p);
    if (sa < 0 || sb < 0 || sc < 0) return 0;
    if (sb==0 || (sa==0 && L == 1) || (sc == 0 &&
    \hookrightarrow R == sz(H)))
      return 1;
    return 2;
  int mid = L + len / 2;
  if (sideOf(H[0], H[mid], p) >= 0)
    return insideHull2(H, mid, R, p);
  return insideHull2(H, L, mid+1, p);
int insideHull(const vector<P>& hull, const P& p)
  if (sz(hull) < 3) return onSegment(hull[0],</pre>
  → hull.back(), p);
  else return insideHull2(hull, 1, sz(hull), p);
```

LineHullIntersection.h

```
11 sqn(11 a) { return (a > 0) - (a < 0); }</pre>
typedef Point<ll> P;
struct HullIntersection {
 int N;
 vector<P> p;
 vector<pair<P, int>> a;
 HullIntersection(const vector<P>& ps) :
  \rightarrow N(sz(ps)), p(ps) {
   p.insert(p.end(), all(ps));
   int b = 0;
   rep(i,1,N) if (P\{p[i].y,p[i].x) < P\{p[b].y,
    \rightarrow p[b].x}) b = i;
   rep(i,0,N) {
      int f = (i + b) % N;
      a.emplace_back(p[f+1] - p[f], f);
 }
 int qd(P p) {
   return (p.y < 0) ? (p.x >= 0) + 2
         : (p.x \le 0) * (1 + (p.y \le 0));
 int bs(P dir) {
   int lo = -1, hi = N;
   while (hi - lo > 1) {
      int mid = (lo + hi) / 2;
      if (make_pair(qd(dir), dir.y *

    a[mid].first.x) <</pre>
        make_pair(qd(a[mid].first), dir.x *

    a[mid].first.y))

        hi = mid;
      else lo = mid;
    return a[hi%N].second;
 }
 bool isign (P a, P b, int x, int y, int s) {
   return sgn(a.cross(p[x], b)) *
    \rightarrow sgn(a.cross(p[y], b)) == s;
 int bs2(int lo, int hi, P a, P b) {
   int L = lo;
   if (hi < lo) hi += N;
   while (hi - lo > 1) {
      int mid = (lo + hi) / 2;
      if (isign(a, b, mid, L, -1)) hi = mid;
      else lo = mid;
    roturn lo.
```

| 8.4 Misc. Point Set Problems

closestPair.h

```
template < class It >
bool it_less(const It& i, const It& j) { return
template < class It >
bool y_it_less(const It& i,const It& j) {return
\hookrightarrow i->y < j->y; }
template<class It, class IIt> /* IIt =
→ vector<It>::iterator */
double cp_sub(IIt ya, IIt yaend, IIt xa, It &i1,
typedef typename

    iterator_traits<It>::value_type P;

  int n = yaend-ya, split = n/2;
 if(n <= 3) { // base case
    double a = (*xa[1] - *xa[0]).dist(), b = 1e50,
    \rightarrow c = 1e50;
   if (n==3) b= (*xa[2]-*xa[0]).dist(),
    \hookrightarrow c=(*xa[2]-*xa[1]).dist();
   if(a \le b) \{ i1 = xa[1];
      if(a <= c) return i2 = xa[0], a;
      else return i2 = xa[2], c;
   } else { i1 = xa[2];
      if(b <= c) return i2 = xa[0], b;
      else return i2 = xa[1], c;
  vector<It> ly, ry, stripy;
  P splitp = *xa[split];
  double splitx = splitp.x;
  for(IIt i = ya; i != yaend; ++i) { // Divide
   if(*i != xa[split] && (**i-splitp).dist2() <</pre>
    → 1e-12)
      return i1 = *i, i2 = xa[split], 0;// nasty

→ special case!

    if (**i < splitp) ly.push_back(*i);</pre>
    else ry.push_back(*i);
  } // assert((signed)lefty.size() == split)
  It j1, j2; // Conquer
  double a = cp_sub(ly.begin(), ly.end(), xa, i1,
  double b = cp_sub(ry.begin(), ry.end(),
  \rightarrow xa+split, j1, j2);
  if (b < a) a = b, i1 = 1, i2 = 12;
  double a2 = a*a;
  for(IIt i = ya; i != yaend; ++i) { // Create

    strip (y-sorted)

   double x = (*i) -> x;
   if(x >= splitx-a && x <= splitx+a)</pre>

    stripy.push_back(*i);

  for(IIt i = stripy.begin(); i != stripy.end();
```

kdTree.h

```
typedef long long T;
typedef Point<T> P;
const T INF = numeric_limits<T>::max();
bool on_x(const P& a, const P& b) { return a.x <
\rightarrow b.x; }
bool on_y(const P& a, const P& b) { return a.y <
\rightarrow b.y; }
struct Node {
 P pt; // if this is a leaf, the single point in
 T x0 = INF, x1 = -INF, y0 = INF, y1 = -INF; //
  → bounds
 Node *first = 0, *second = 0;
 T distance (const P& p) { // min squared

→ distance to a point

   T x = (p.x < x0 ? x0 : p.x > x1 ? x1 : p.x);
   T y = (p.y < y0 ? y0 : p.y > y1 ? y1 : p.y);
   return (P(x,y) - p).dist2();
 Node (vector < P > \&\& vp) : pt (vp[0]) {
   for (P p : vp) {
     x0 = min(x0, p.x); x1 = max(x1, p.x);
     y0 = min(y0, p.y); y1 = max(y1, p.y);
   if (vp.size() > 1) {
      // split on x if the box is wider than high
      sort (all (vp), x1 - x0 >= y1 - y0 ? on_x :
      \rightarrow on_y);
      // divide by taking half the array for each
      // best performance with many duplicates in

    → the middle)

     int half = sz(vp)/2;
      first = new Node({vp.begin(), vp.begin() +
      \rightarrow half});
      second = new Node({vp.begin() + half,
      \rightarrow vp.end()});
} ;
struct KDTree
 Node* root;
 KDTree (const vector < P > & vp) : root (new
  naireT Do goarch (Nodo unodo gongt Do n) (
```

```
DelaunayTriangulation.h
template < class P, class F>
void delaunay (vector < P > & ps, F trifun) {
  if (sz(ps) == 3) \{ int d = (ps[0].cross(ps[1],
  \rightarrow ps[2]) < 0);
    trifun(0,1+d,2-d); }
  vector<P3> p3;
  trav(p, ps) p3.emplace_back(p.x, p.y,
   \rightarrow p.dist2());
  if (sz(ps) > 3) trav(t, hull3d(p3)) if
   \hookrightarrow ((p3[t.b]-p3[t.a]).
      cross(p3[t.c]-p3[t.a]).dot(P3(0,0,1)) < 0)
    trifun(t.a, t.c, t.b);
8.5
       3D
PolyhedronVolume.h
template<class V, class L>
double signed_poly_volume(const V& p, const L&

    trilist) {
 double v = 0;
 trav(i, trilist) v +=

    p[i.a].cross(p[i.b]).dot(p[i.c]);
  return v / 6;
```

Point3D.h

```
template < class T > struct Point3D {
  typedef Point3D P;
  typedef const P& R;
  T x, y, z;
  explicit Point3D (T x=0, T y=0, T z=0) : x(x),
  \hookrightarrow \forall (\forall), z(z) \{ \}
  bool operator<(R p) const {</pre>
    return tie(x, y, z) < tie(p.x, p.y, p.z); }
  bool operator==(R p) const {
    return tie(x, y, z) == tie(p.x, p.y, p.z); }
  P operator+(R p) const { return P(x+p.x, y+p.y,
  \hookrightarrow z+p.z); }
  P operator-(R p) const { return P(x-p.x, y-p.y,
  \hookrightarrow z-p.z); }
  P operator* (T d) const { return P(x*d, y*d,
  P operator/(T d) const { return P(x/d, y/d,
  \rightarrow z/d); }
  T dot(R p) const { return x*p.x + y*p.y +
  \hookrightarrow z * p.z;
  P cross(R p) const {
    return P(y*p.z - z*p.y, z*p.x - x*p.z, x*p.y
    \rightarrow - y*p.x);
  T dist2() const { return x*x + y*y + z*z; }
  double dist() const { return

    sqrt((double)dist2()); }

  //Azimuthal angle (longitude) to x-axis in

    interval [-pi, pi]

  double phi() const { return atan2(y, x); }
  //Zenith angle (latitude) to the z-axis in

    interval [0, pi]

  double theta() const { return
  \rightarrow atan2(sqrt(x*x+y*y),z); }
  P unit() const { return *this/(T)dist(); }
  \rightarrow //makes dist()=1
  //returns unit vector normal to *this and p
  P normal(P p) const { return cross(p).unit(); }
  //returns point rotated 'angle' radians ccw

→ around axis

  P rotate (double angle, P axis) const {
    double s = sin(angle), c = cos(angle); P u =

    axis.unit();
    return u*dot(u)*(1-c) + (*this)*c -

    cross(u) *s;

} ;
```

```
typedef Point3D<double> P3;
struct PR {
 void ins(int x) { (a == -1 ? a : b) = x; }
 void rem(int x) { (a == x ? a : b) = -1; }
 int cnt() { return (a !=-1) + (b !=-1); }
 int a, b;
struct F { P3 q; int a, b, c; };
vector<F> hull3d(const vector<P3>& A) {
 assert (sz(A) >= 4);
 vector<vector<PR>> E(sz(A), vector<PR>(sz(A),
  \hookrightarrow {-1, -1}));
#define E(x,y) E[f.x][f.y]
 vector<F> FS:
 auto mf = [&](int i, int j, int k, int l) {
   P3 q = (A[j] - A[i]).cross((A[k] - A[i]));
   if (q.dot(A[1]) > q.dot(A[i]))
      q = q * -1;
   F f{q, i, j, k};
   E(a,b).ins(k); E(a,c).ins(j); E(b,c).ins(i);
   FS.push_back(f);
 rep(i, 0, 4) rep(j, i+1, 4) rep(k, j+1, 4)
   mf(i, j, k, 6 - i - j - k);
 rep(i,4,sz(A)) {
   rep(j, 0, sz(FS)) {
      F f = FS[i];
      if(f.q.dot(A[i]) > f.q.dot(A[f.a])) {
        E(a,b).rem(f.c);
        E(a,c).rem(f.b);
        E(b,c).rem(f.a);
        swap(FS[j--], FS.back());
        FS.pop_back();
      }
   int nw = sz(FS);
   rep(j,0,nw) {
     F f = FS[i];
#define C(a, b, c) if (E(a,b).cnt() != 2) mf(f.a,
\hookrightarrow f.b, i, f.c);
      C(a, b, c); C(a, c, b); C(b, c, a);
 trav(it, FS) if ((A[it.b] - A[it.a]).cross(
   A[it.c] - A[it.a]).dot(it.q) <= 0) swap(it.c,
    \rightarrow it.b);
 return FS:
```

```
sphericalDistance.h
double sphericalDistance(double f1, double t1,
    double f2, double t2, double radius) {
  double dx = \sin(t2) \cdot \cos(f2) - \sin(t1) \cdot \cos(f1);
  double dy = sin(t2) * sin(f2) - sin(t1) * sin(f1);
  double dz = cos(t2) - cos(t1);
  double d = sqrt(dx*dx + dy*dy + dz*dz);
  return radius*2*asin(d/2);
Strings (9)
KMP.h
vi pi(const string& s) {
  vi p(sz(s));
  rep(i,1,sz(s)) {
    int q = p[i-1];
    while (g \&\& s[i] != s[g]) g = p[g-1];
    p[i] = q + (s[i] == s[q]);
  return p;
vi match (const string& s, const string& pat) {
  vi p = pi(pat + '\0' + s), res;
  rep(i,sz(p)-sz(s),sz(p))
    if (p[i] == sz(pat)) res.push_back(i - 2 *
    \hookrightarrow sz(pat));
  return res;
Manacher.h
void manacher(const string& s) {
  int n = sz(s);
  vi p[2] = {vi(n+1), vi(n)};
  rep(z, 0, 2) for (int i=0, l=0, r=0; i < n; i++) {
    int t = r-i+!z;
    if (i<r) p[z][i] = min(t, p[z][1+t]);
    int L = i-p[z][i], R = i+p[z][i]-!z;
    while (L>=1 && R+1<n && s[L-1] == s[R+1])
      p[z][i]++, L--, R++;
    if (R>r) l=L, r=R;
} }
MinRotation.h
```

```
int min_rotation(string s) {
  int a=0, N=sz(s); s += s;
  rep(b,0,N) rep(i,0,N) {
    if (a+i == b || s[a+i] < s[b+i]) {b += max(0, cond); break;}
    if (s[a+i] > s[b+i]) { a = b; break; }
  }
  return a;
}
```

```
typedef pair<ll, int> pli;
void count_sort(vector<pli> &b, int bits) { //

→ (optional)

 //this is just 3 times faster than stl sort for
  \rightarrow N=10^6
 int mask = (1 << bits) - 1;</pre>
 rep(it, 0, 2) {
   int move = it * bits;
   vi q(1 << bits), w(sz(q) + 1);
   rep(i, 0, sz(b))
      q[(b[i].first >> move) & mask]++;
    partial_sum(q.begin(), q.end(), w.begin() +
    \hookrightarrow 1);
    vector<pli> res(b.size());
   rep(i,0,sz(b))
     res[w[(b[i].first >> move) \& mask]++] =
      \rightarrow b[i];
    swap(b, res);
struct SuffixArray
 vi a;
 string s;
 SuffixArray(const string& _s) : s(_s + '\0') {
   int N = sz(s);
   vector<pli> b(N);
   a.resize(N);
   rep(i,0,N) {
     b[i].first = s[i];
      b[i].second = i;
   }
    int q = 8;
    while ((1 << q) < N) q++;
    for (int moc = 0;; moc++) {
      count_sort(b, q); // sort(all(b)) can be

→ used as well

      a[b[0].second] = 0;
      rep(i,1,N)
        a[b[i].second] = a[b[i - 1].second] +
          (b[i - 1].first != b[i].first);
      if ((1 << moc) >= N) break;
      rep(i,0,N) {
       b[i].first = (ll)a[i] << q;
        if (i + (1 << moc) < N)
          b[i].first += a[i + (1 << moc)];
        b[i].second = i;
    rep(i, 0, sz(a)) a[i] = b[i].second;
```

SuffixTree.h

```
struct SuffixTree {
  enum { N = 200010, ALPHA = 26 }; // N ~

→ 2*maxlen+10

  int toi(char c) { return c - 'a'; }
 string a; // v = cur node, q = cur position
  int
  \rightarrow t[N][ALPHA], l[N], r[N], p[N], s[N], v=0, q=0, m=2;
 void ukkadd(int i, int c) { suff:
   if (r[v]<=a) {
      if (t[v][c]==-1) { t[v][c]=m; l[m]=i;
        p[m++]=v; v=s[v]; q=r[v]; goto suff; }
      v=t[v][c]; q=l[v];
   if (q==-1 || c==toi(a[q])) q++; else {
      l[m+1]=i; p[m+1]=m; l[m]=l[v]; r[m]=q;
      p[m] = p[v]; t[m][c] = m+1;
      \rightarrow t[m][toi(a[q])]=v;
      l[v]=q; p[v]=m; t[p[m]][toi(a[l[m]])]=m;
      v=s[p[m]]; q=l[m];
      while (q<r[m]) { v=t[v][toi(a[q])];
      \rightarrow q+=r[v]-l[v]; }
      if (q==r[m]) s[m]=v; else s[m]=m+2;
      q=r[v]-(q-r[m]); m+=2; goto suff;
 SuffixTree(string a) : a(a) {
    fill(r,r+N,sz(a));
   memset(s, 0, sizeof s);
   memset (t, -1, sizeof t);
    fill(t[1],t[1]+ALPHA,0);
    s[0] = 1; 1[0] = 1[1] = -1; r[0] = r[1] =
    \rightarrow p[0] = p[1] = 0;
    rep(i, 0, sz(a)) ukkadd(i, toi(a[i]));
  // example: find longest common substring (uses
  \hookrightarrow ALPHA = 28)
  pii best;
  int lcs(int node, int i1, int i2, int olen) {
   if (l[node] <= i1 && i1 < r[node]) return 1;</pre>
   if (1[node] <= i2 && i2 < r[node]) return 2;</pre>
   int mask = 0, len = node ? olen + (r[node] -
    \rightarrow l[node]) : 0;
    rep(c, 0, ALPHA) if (t[node][c] != -1)
      mask |= lcs(t[node][c], i1, i2, len);
    if (mask == 3)
      best = max(best, {len, r[node] - len});
    return mask;
  static nii ICC(string a string t) [
```

```
Hashing.h
```

```
// Arithmetic mod 2^64-1. 2x slower than mod 2^64

→ and more

// code, but works on evil test data (e.g.
→ Thue-Morse, where
// ABBA... and BAAB... of length 2^10 hash the
\rightarrow same mod 2^64).
// "typedef ull H;" instead if you think test
→ data is random,
// or work mod 10^9+7 if the Birthday paradox is
→ not a problem.
struct H {
  typedef uint64_t ull;
 ull x; H(ull x=0) : x(x) {}
#define OP(O,A,B) H operator O(H \circ) { ull r = x;

→ asm \

  (A "addg %%rdx, %0\n adcg $0, %0" : "+a"(r) :
\rightarrow B); return r; }
 OP(+,,"d"(o.x)) OP(*,"mul %1\n", "r"(o.x) :
  \rightarrow "rdx")
  H operator-(H o) { return *this + ~o.x; }
  ull get() const { return x + !~x; }
  bool operator==(H o) const { return get() ==
  → o.get(); }
  bool operator<(H o) const { return get() <</pre>
  → o.get(); }
};
static const H C = (11)1e11+3; // (order ~ 3e9;
struct HashInterval {
  vector<H> ha, pw;
  HashInterval(string& str) : ha(sz(str)+1),
  \rightarrow pw(ha) {
   pw[0] = 1;
    rep(i, 0, sz(str))
     ha[i+1] = ha[i] * C + str[i],
      pw[i+1] = pw[i] * C;
 H hashInterval(int a, int b) { // hash [a, b)
    return ha[b] - ha[a] * pw[b - a];
};
vector<H> getHashes(string& str, int length) {
 if (sz(str) < length) return {};</pre>
 H h = 0, pw = 1;
  rep(i,0,length)
   h = h * C + str[i], pw = pw * C;
  vector<H> ret = {h};
  rep(i,length,sz(str)) {
    ret.push\_back(h = h * C + str[i] - pw *
       atr[i_lonath]).
```

```
AhoCorasick.h
```

```
struct AhoCorasick {
  enum {alpha = 26, first = 'A'};
 struct Node {
   // (nmatches is optional)
   int back, next[alpha], start = -1, end = -1,

    nmatches = 0;

   Node (int v) { memset (next, v, sizeof (next));
 };
 vector<Node> N;
 vector<int> backp;
 void insert(string& s, int j) {
   assert(!s.empty());
   int n = 0;
   trav(c, s) {
     int& m = N[n].next[c - first];
     if (m == -1) { n = m = sz(N);
      \rightarrow N.emplace_back(-1); }
      else n = m;
   if (N[n].end == -1) N[n].start = j;
   backp.push_back(N[n].end);
   N[n].end = j;
   N[n].nmatches++;
 AhoCorasick(vector<string>& pat) {
   N.emplace\_back(-1);
   rep(i, 0, sz(pat)) insert(pat[i], i);
   N[0].back = sz(N);
   N.emplace_back(0);
   queue<int> q;
    for (q.push(0); !q.empty(); q.pop()) {
      int n = q.front(), prev = N[n].back;
      rep(i,0,alpha) {
        int &ed = N[n].next[i], y =
        → N[prev].next[i];
        if (ed == -1) ed = y;
        else {
         N[ed].back = y;
          (N[ed].end == -1 ? N[ed].end:

→ backp[N[ed].start])
            = N[y].end;
         N[ed].nmatches += N[y].nmatches;
          q.push(ed);
 vi find(string word) {
   int n = 0;
   mi root // 11 count - 0.
```

Various (10)

10.1 Intervals

IntervalContainer.h

```
set<pii>::iterator addInterval(set<pii>& is, int
\hookrightarrow L, int R) {
 if (L == R) return is.end();
  auto it = is.lower_bound({L, R}), before = it;
  while (it != is.end() && it->first <= R) {</pre>
    R = max(R, it->second);
   before = it = is.erase(it);
  if (it != is.begin() && (--it)->second >= L) {
   L = min(L, it->first);
   R = max(R, it->second);
    is.erase(it);
  return is.insert(before, {L,R});
void removeInterval(set<pii>& is, int L, int R) {
 if (L == R) return;
  auto it = addInterval(is, L, R);
  auto r2 = it->second;
  if (it->first == L) is.erase(it);
  else (int&)it->second = L;
  if (R != r2) is.emplace(R, r2);
```

IntervalCover.h

```
template < class T>
vi cover(pair<T, T> G, vector<pair<T, T>> I) {
  vi S(sz(I)), R;
  iota(all(S), 0);
  sort(all(S), [&](int a, int b) { return I[a] <
  \hookrightarrow I[b]; });
  T cur = G.first;
  int at = 0;
 while (cur < G.second) { // (A)</pre>
    pair<T, int> mx = make_pair(cur, -1);
    while (at < sz(I) && I[S[at]].first <= cur) {
      mx = max(mx, make_pair(I[S[at]].second,
      \rightarrow S[at]));
      at++;
    if (mx.second == -1) return {};
    cur = mx.first;
    R.push_back (mx.second);
  return R;
```

ConstantIntervals.h

```
template < class F, class G, class T>
void rec(int from, int to, F f, G g, int& i, T&
\rightarrow p, T q) {
 if (p == q) return;
 if (from == to) {
    g(i, to, p);
    i = to; p = q;
 } else {
    int mid = (from + to) >> 1;
    rec(from, mid, f, q, i, p, f(mid));
    rec(mid+1, to, f, g, i, p, q);
template<class F, class G>
void constantIntervals(int from, int to, F f, G
\rightarrow g) {
 if (to <= from) return;</pre>
 int i = from; auto p = f(i), q = f(to-1);
 rec(from, to-1, f, g, i, p, q);
  q(i, to, q);
```

| 10.2 Misc. algorithms

```
TernarySearch.h

template < class F >
int ternSearch(int a, int b, F f) {
   assert(a <= b);
   while (b - a >= 5) {
     int mid = (a + b) / 2;
     if (f(mid) < f(mid+1)) // (A)
        a = mid;
     else
        b = mid+1;
   }
   rep(i,a+1,b+1) if (f(a) < f(i)) a = i; // (B)
   return a;
}</pre>
```

Karatsuba.h

LIS.h

```
template < class I > vi lis(vector < I > S) {
  vi prev(sz(S));
  typedef pair<I, int> p;
  vector res;
  rep(i, 0, sz(S))
    p el { S[i], i };
    //S[i]+1 for non-decreasing
    auto it = lower_bound(all(res), p { S[i], 0
    if (it == res.end()) res.push_back(el), it =
    → --res.end();
    *it = el;
    prev[i] = it = res.begin() ?0:(it-1) -> second;
  int L = sz(res), cur = res.back().second;
  vi ans(L);
  while (L--) ans[L] = cur, cur = prev[cur];
  return ans;
LCS.h
```

```
template < class T > T lcs (const T &X, const T &Y) {
   int a = sz(X), b = sz(Y);
   vector < vi > dp(a+1, vi(b+1));
   rep(i,1,a+1) rep(j,1,b+1)
      dp[i][j] = X[i-1] == Y[j-1] ? dp[i-1][j-1]+1 :
        max(dp[i][j-1],dp[i-1][j]);
   int len = dp[a][b];
   T ans(len,0);
   while (a && b)
      if (X[a-1] == Y[b-1]) ans[--len] = X[--a], --b;
      else if (dp[a][b-1] > dp[a-1][b]) --b;
      else --a;
   return ans;
}
```

10.3 Dynamic programming

```
DivideAndConquerDP.h
struct DP { // Modify at will:
  int lo(int ind) { return 0; }
  int hi(int ind) { return ind; }
  11 f(int ind, int k) { return dp[ind][k]; }
  void store(int ind, int k, ll v) { res[ind] =
  \rightarrow pii(k, v); }
  void rec(int L, int R, int LO, int HI) {
    if (L >= R) return;
    int mid = (L + R) \gg 1;
    pair<11, int> best(LLONG_MAX, LO);
    rep(k, max(LO, lo(mid)), min(HI, hi(mid)))
      best = min(best, make_pair(f(mid, k), k));
    store (mid, best.second, best.first);
    rec(L, mid, LO, best.second+1);
    rec(mid+1, R, best.second, HI);
  void solve(int L, int R) { rec(L, R, INT_MIN,
  → INT_MAX); }
};
```

10.4 Debugging tricks

KnuthDP.h

- signal(SIGSEGV, [](int) { _Exit(0); }); converts segfaults into Wrong Answers. Similarly one can catch SIGABRT (assertion failures) and SIGFPE (zero divisions). _GLIBCXX_DEBUG violations generate SIGABRT (or SIGSEGV on gcc 5.4.0 apparently).
- feenableexcept (29); kills the program on NaNs (1), 0-divs (4), infinities (8) and denormals (16).

10.5 Optimization tricks

10.5.1 Bit hacks

• x & -x is the least bit in x.

- for (int x = m; x;) { --x &= m; ... } loops over all subset masks of m (except m itself).
- c = x&-x, r = x+c; (((r^x) >> 2)/c) | r is the next number after x with the same number of bits set.
- rep(b,0,K) rep(i,0,(1 << K)) if (i & 1
 << b) D[i] += D[i^(1 << b)]; computes all
 sums of subsets.</pre>

10.5.2 Pragmas

- #pragma GCC optimize ("Ofast") will make GCC auto-vectorize for loops and optimizes floating points better (assumes associativity and turns off denormals).
- #pragma GCC target ("avx,avx2") can double performance of vectorized code, but causes crashes on old machines.
- #pragma GCC optimize ("trapv") kills the program on integer overflows (but is really slow).

// Either globally or in a single class:

${\bf BumpAllocator.h}$

```
static char buf[450 << 201;
void* operator new(size_t s) {
 static size_t i = sizeof buf;
 assert(s < i);
 return (void*) &buf[i -= s];
void operator delete(void*) {}
SmallPtr.h
template < class T> struct ptr {
 unsigned ind;
 ptr(T* p = 0) : ind(p ? unsigned((char*)p -
  \rightarrow buf) : 0) {
   assert(ind < sizeof buf);</pre>
 T& operator*() const { return *(T*)(buf + ind);
 T* operator->() const { return &**this; }
 T& operator[](int a) const { return
  explicit operator bool() const { return ind; }
```

```
{\bf BumpAllocator STL.h}
```

```
char buf[450 << 20] alignas(16);
size_t buf_ind = sizeof buf;

template < class T > struct small {
   typedef T value_type;
   small() {}
   template < class U > small(const U&) {}
   T* allocate(size_t n) {
      buf_ind -= n * sizeof(T);
      buf_ind &= 0 - alignof(T);
      return (T*) (buf + buf_ind);
   }
   void deallocate(T*, size_t) {}
};
```

Unrolling.h

SIMD.h

```
#pragma GCC target ("avx2") // or sse4.1
#include "immintrin.h"
typedef __m256i mi;
#define L(x) _mm256_loadu_si256((mi*)&(x))
// High-level/specific methods:
// load(u)?_si256, store(u)?_si256,
→ setzero_si256, _mm_malloc
// blendv_(epi8|ps|pd) (z?y:x), movemask_epi8
// i32gather_epi32(addr, x, 4): map addr[] over
\rightarrow 32-b parts of x
// sad_epu8: sum of absolute differences of u8,
 → outputs 4xi64
// maddubs_epi16: dot product of unsigned i7's,
 → outputs 16xi15
// madd_epi16: dot product of signed i16's,
 → outputs 8xi32
// extractf128_si256(, i) (256->128),
 // permute2f128_si256(x,x,1) swaps 128-bit lanes
// shuffle_epi32(x, 3*64+2*16+1*4+0) == x for
→ each lane
// shuffle_epi8(x, y) takes a vector instead of

→ an imm

// Methods that work with most data types (append
→ e.g. _epi32):
// set1, blend (i8?x:y), add, adds (sat.), mullo,
 → sub, and/or,
// and not, abs, min, max, sign(1,x), cmp(gt/eq),
 → unpack(lo/hi)
int sumi32(mi m) { union {int v[8]; mi m;} u; u.m
 \hookrightarrow = m;
    int ret = 0; rep(i,0,8) ret += u.v[i]; return
     → ret; }
mi zero() { return _mm256_setzero_si256(); }
mi one() { return _mm256_set1_epi32(-1); }
bool all_zero(mi m) { return

    _mm256_testz_si256(m, m); }

bool all_one(mi m) { return _mm256_testc_si256(m,

    one());
}
ll example_filteredDotProduct(int n, short* a,
 ⇔ short* b) {
    int i = 0; 11 r = 0;
    mi zero = _mm256_setzero_si256(), acc = zero;
    while (i + 16 \le n) {
        mi \ va = L(a[i]), \ vb = L(b[i]); \ i += 16;
        m = mm^2 = mm^
```

Techniques (A)

techniques.txt

Combinatorics

159 lines

Recursion Divide and conquer Finding interesting points in N log N Algorithm analysis Master theorem Amortized time complexity Greedy algorithm Scheduling Max contiquous subvector sum Invariants Huffman encoding Graph theory Dynamic graphs (extra book-keeping) Breadth first search Depth first search * Normal trees / DFS trees Dijkstra's algorithm MST: Prim's algorithm Bellman-Ford Konig's theorem and vertex cover Min-cost max flow Lovasz toggle Matrix tree theorem Maximal matching, general graphs Hopcroft-Karp Hall's marriage theorem Graphical sequences Floyd-Warshall Euler cycles Flow networks * Augmenting paths * Edmonds-Karp Bipartite matching Min. path cover Topological sorting Strongly connected components Cut vertices, cut-edges and biconnected components Edge coloring * Trees Vertex coloring * Bipartite graphs (=> trees) * 3^n (special case of set cover) Diameter and centroid K'th shortest path Shortest cycle Dynamic programming Knapsack Coin change Longest common subsequence Longest increasing subsequence Number of paths in a dag Shortest path in a dag Dynprog over intervals Dynprog over subsets Dynprog over probabilities Dynprog over trees 3^n set cover Divide and conquer Knuth optimization Convex hull optimizations RMQ (sparse table a.k.a 2^k-jumps) Bitonic cycle Log partitioning (loop over most restricted)

Computation of binomial coefficients Pigeon-hole principle Inclusion/exclusion Catalan number Pick's theorem Number theory Integer parts Divisibility Euclidean algorithm Modular arithmetic * Modular multiplication * Modular inverses * Modular exponentiation by squaring Chinese remainder theorem Fermat's little theorem Euler's theorem Phi function Frobenius number Quadratic reciprocity Pollard-Rho Miller-Rabin Hensel lifting Vieta root jumping Game theory Combinatorial games Game trees Mini-may Nim Games on graphs Games on graphs with loops Grundy numbers Bipartite games without repetition General games without repetition Alpha-beta pruning Probability theory Optimization Binary search Ternary search Unimodality and convex functions Binary search on derivative Numerical methods Numeric integration Newton's method Root-finding with binary/ternary search Golden section search Matrices Gaussian elimination Exponentiation by squaring Sorting Radix sort Geometry Coordinates and vectors * Cross product * Scalar product Convex hull Polygon cut Closest pair Coordinate-compression Quadtrees All segment-segment intersection Sweeping Discretization (convert to events and sweep) Angle sweeping Line sweeping Discrete second derivatives Strings Longest common substring Palindrome subsequences

Knuth-Morris-Pratt Tries Rolling polynomial hashes Suffix arrav Suffix tree Aho-Corasick Manacher's algorithm Letter position lists Combinatorial search Meet in the middle Brute-force with pruning Best-first (A*) Bidirectional search Iterative deepening DFS / A* Data structures LCA (2^k-jumps in trees in general) Pull/push-technique on trees Heavy-light decomposition Centroid decomposition Lazy propagation Self-balancing trees Convex hull trick (wcipeg.com/wiki/Convex_hull_trick) Monotone queues / monotone stacks / sliding queues Sliding queue using 2 stacks Persistent segment tree

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