TCR

git diff solution (Jens Heuseveldt, Ludo Pulles, Pim Spelier)

Contents

At the start of a contest, type this in a terminal:

```
printf "set nu sw=4 ts=4 sts=4 noet ai hls shellcmdflag=-ic\nsy on | colo slate" > .vimrc
printf "\nalias gsubmit='g++ -Wall -Wshadow -std=c++11'" >> .bashrc
printf "\nalias g11='gsubmit -DLOCAL -g'" >> .bashrc
. .bashrc
mkdir contest; cd contest
```

template.cpp

```cpp
#include<bits/stdc++.h>
using namespace std;

// Order statistics tree (if supported by judge!):
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;

template<class TK, class TM>
using order_tree = tree<TK, TM, less<TK>, rb_tree_tag,
  tree_order_statistics_node_update>;
// iterator  find_by_order(int r) (zero based)
// int    order_of_key(TK v)
template<class TV> using order_set = order_tree<TV,
  null_type>;

#define x first
#define y second
#define pb push_back
#define eb emplace_back
#define rep(i,a,b) for(auto i=(a);i!=(b); ++i)
#define all(v) (v).begin(), (v).end()
#define rs resize

typedef long long ll;
typedef pair<int, int> pii;
typedef vector<int> vi;
typedef vector<vi> vvi;
template<class T> using min_queue = priority_queue<T,
  vector<T>, greater<T>>;

const int INF = 2147483647; // (1 << 30) - 1 + (1 << 30)
const ll LLINF = (1LL << 62) - 1 + (1LL << 62); // =
  9.223.372.036.854.775.807
const double PI = acos(-1.0);

#ifdef LOCAL
#define DBG(x) cerr << __LINE__ << ": " << #x << " = " << (x)
  << endl
#else
#define DBG(x)
const bool LOCAL = false;
#endif

void Log() { if(LOCAL) cerr << "\n\n"; }
template<class T, class... S>
void Log(T t, S... s) { if(LOCAL) cerr << t << "\t",
  Log(s...); }

// lambda-expression: [] (args) -> retType { body }
int main() {
  ios_base::sync_with_stdio(false); // fast IO
  cin.tie(NULL); // fast IO
  cerr << boolalpha; // print true/false
  (cout << fixed).precision(10); // adjust precision

  return 0;
}
```

Prime numbers: *982451653*, *81253449*, $10^3+\{-9,-3,9,13\}$, $10^6+\{-17,3,33\}$, $10^9+\{7,9,21,33,87\}$

0.1. De winnende aanpak.

- Goed slapen & een vroeg ritme hebben
- Genoeg drinken & eten voor en tijdens de wedstrijd
- Een lijst van alle problemen met info waar het over gaat, en wie het goed kan oplossen
- Ludo moet **ALLE** opgaves **goed** lezen
- Test de kleine voorbeeldgevallen
- Houd na 2 uur een pauze en overleg waar iedereen mee bezig is
- Maak zelf wat test-cases
- Typ de dingen uit de TCR, die je zeker nodig hebt, alvast in
- Als iemand niks te doen heeft, kan hij nodige dingen uit de TCR typen.
- We moeten ook een voorbeeld test-case voor TCR algoritmes hebben om te testen of het goed overgetypt is
- Bij geometrie moeten we om kunnen gaan met meerdere input manieren (voor bv. lijnen)
- Gebruik veel long long's

0.2. Wrong Answer.

(1) Print de oplossing om te debuggen! Kijk ook naar andere (mogelijk makkelijkere) problemen.
(2) Bedenk zelf test-cases met **randgevallen**!
(3) Controleer op **overflow** (gebruik **OVERAL** long long, long double).
Kijk naar overflows in tussenantwoorden bij modulo.
(4) Controleer de **precisie**.
(5) Controleer op **typo's**.
(6) Loop de voorbeeldinput accuraat langs.
(7) Controller op off-by-one-errors (in indices of lus-grenzen)?

0.3. Detecting overflow. These are GNU builtins, detect both over- and underflow. Returns a boolean upon failure, otherwise the result is present in `ref`. Follow the template:

```
bool isOverflown = __builtin_[add|mul|sub]_overflow(a, b, \&res);
```

0.4. Covering problems.

Minimum edge cover $\iff$ *Maximum independent set*

Matching: A set of edges without common vertices *(Maximum is the* ***largest*** *such set, maximal is a set which you cannot add more edges to without breaking the property).*

Minimum Vertex Cover: A set vertices (cover) such that each edge in the graph is incident to at least one vertex of the set.

Minimum Edge Cover: A set of edges (cover) such that every vertex is incident to at least one edge of the set.

Maximum Independent Set: A set of vertices in a graph such that no two of them are adjacent.

König's theorem: In any bipartite graph, the number of edges in a maximum matching equals the number of vertices in a minimum vertex cover

0.5. Game theory. A game can be reduced to Nim if it is a finite impartial game. Nim and its variants include:

Nim: Let $X = \bigoplus_{i=1}^n x_i$, then $(x_i)_{i=1}^n$ is a winning position iff $X \neq 0$. Find a move by picking k such that $x_k > x_k \oplus X$.

Misère Nim: Regular Nim, except that the last player to move *loses*. Play regular Nim until there is only one pile of size larger than 1, reduce it to 0 or 1 such that there is an odd number of piles.

Staircase Nim: Stones are moved down a staircase and only removed from the last pile. $(x_i)_{i=1}^n$ is an L-position if $(x_{2i-1})_{i=1}^{n/2}$ is (i.e. only look at odd-numbered piles).

Moore's Nim$_k$: The player may remove from at most k piles (Nim = Nim$_1$). Expand the piles in base 2, do a carry-less addition in base $k+1$ (i.e. the number of ones in each column should be divisible by $k+1$).

Dim$^+$: The number of removed stones must be a divisor of the pile size. The Sprague-Grundy function is $k+1$ where 2^k is the largest power of 2 dividing the pile size.

Aliquot game: Same as above, except the divisor should be proper (hence 1 is also a terminal state, but watch out for size 0 piles). Now the Sprague-Grundy function is just k.

Nim (at most half): Write $n+1 = 2^m y$ with m maximal, then the Sprague-Grundy function of n is $(y-1)/2$.

Lasker's Nim: Players may alternatively split a pile into two new non-empty piles. $g(4k+1) = 4k+1$, $g(4k+2) = 4k+2$, $g(4k+3) = 4k+4$, $g(4k+4) = 4k+3$ $(k \geq 0)$.

Hackenbush on trees: A tree with stalks $(x_i)_{i=1}^n$ may be replaced with a single stalk with length $\bigoplus_{i=1}^n x_i$.

A useful identity: $\bigoplus_{x=0}^{a-1} x = \{0, a-1, 1, a\}[a \bmod 4]$.

1. Math

```
int abs(int x) { return x > 0 ? x : -x; }
int sign(int x) { return (x > 0) - (x < 0); }

// greatest common divisor
ll gcd(ll a, ll b) { while (b) a %= b, swap(a, b); return a;
↪  };
// least common multiple
ll lcm(ll a, ll b) { return a / gcd(a, b) * b; }
ll mod(ll a, ll b) { return (a %= b) < 0 ? a + b : a; }

// safe multiplication (ab % m) for m <= 4e18 in O(log b)
ll mulmod(ll a, ll b, ll m) {
  ll r = 0;
  while (b) {
    if (b & 1) r = (r + a) % m; a = (a + a) % m; b >>= 1;
  }
  return r;
}

// safe exponentation (a^b % m) for m <= 2e9 in O(log b)
ll powmod(ll a, ll b, ll m) {
  ll r = 1;
  while (b) {
    if (b & 1) r = (r * a) % m; // r = mulmod(r, a, m);
    a = (a * a) % m; // a = mulmod(a, a, m);
    b >>= 1;
  }
  return r;
}

// returns x, y such that ax + by = gcd(a, b)
ll egcd(ll a, ll b, ll &x, ll &y) {
  ll xx = y = 0, yy = x = 1;
  while (b) {
    x -= a / b * xx; swap(x, xx);
    y -= a / b * yy; swap(y, yy);
    a %= b; swap(a, b);
  }
  return a;
}

// Chinese remainder theorem
const pll NO_SOLUTION(0, -1);
// Returns (u, v) such that x = u % v <=> x = a % n and x = b
↪  % m
pll crt(ll a, ll n, ll b, ll m) {
  ll s, t, d = egcd(n, m, s, t), nm = n * m;
  if (mod(a - b, d)) return NO_SOLUTION;
  return pll(mod(s * b * n + t * a * m, nm) / d, nm / d);
  /* when n, m > 10^6, avoid overflow:
  return pll(mod(mulmod(mulmod(s, b, nm), n, nm)
               + mulmod(mulmod(t, a, nm), m, nm), nm) / d, nm
↪  / d); */
}

// phi[i] = #{ 0 < j <= i | gcd(i, j) = 1 }
vi totient(int N) {
  vi phi(N);
  for (int i = 0; i < N; i++) phi[i] = i;
  for (int i = 2; i < N; i++)
    if (phi[i] == i)
      for (int j = i; j < N; j += i) phi[j] -= phi[j] / i;
  return phi;
}

// calculate nCk % p (p prime!)
ll lucas(ll n, ll k, ll p) {
  ll ans = 1;
  while (n) {
    ll np = n % p, kp = k % p;
    if (np < kp) return 0;
    ans = mod(ans * binom(np, kp), p); // (np C kp)
    n /= p; k /= p;
  }
  return ans;
}

// returns if n is prime for n < 3e24 ( > 2^64)
bool millerRabin(ll n){
  if (n < 2 || n % 2 == 0) return n == 2;
  ll d = n - 1, ad, s = 0, r;
  for (; d % 2 == 0; d /= 2) s++;
  for (int a : { 2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37,
↪  41 }) {
    if (n == a) return true;
    if ((ad = powmod(a, d, n)) == 1) continue;
    for (r = 0; r < s && ad + 1 != n; r++)
      ad = mulmod(ad, ad, n);
    if (r == s) return false;
  }
  return true;
}
```

1.1. Primitive Root.

```
#include "mod_pow.cpp"
ll primitive_root(ll m) {
  vector<ll> div;
  for (ll i = 1; i*i <= m-1; i++) {
    if ((m-1) % i == 0) {
      if (i < m) div.push_back(i);
      if (m/i < m) div.push_back(m/i); } }
  rep(x,2,m) {
    bool ok = true;
    iter(it,div) if (mod_pow<ll>(x, *it, m) == 1) {
      ok = false; break; }
    if (ok) return x; }
```

```
  return -1; }
// vim: cc=60 ts=2 sts=2 sw=2:
```

1.2. **Tonelli-Shanks algorithm.** Given prime p and integer $1 \leq n < p$, returns the square root r of n modulo p. There is also another solution given by $-r$ modulo p.

```
#include "mod_pow.cpp"
ll legendre(ll a, ll p) {
  if (a % p == 0) return 0;
  if (p == 2) return 1;
  return mod_pow(a, (p-1)/2, p) == 1 ? 1 : -1; }
ll tonelli_shanks(ll n, ll p) {
  assert(legendre(n,p) == 1);
  if (p == 2) return 1;
  ll s = 0, q = p-1, z = 2;
  while (~q & 1) s++, q >>= 1;
  if (s == 1) return mod_pow(n, (p+1)/4, p);
  while (legendre(z,p) != -1) z++;
  ll c = mod_pow(z, q, p),
    r = mod_pow(n, (q+1)/2, p),
    t = mod_pow(n, q, p),
    m = s;
  while (t != 1) {
    ll i = 1, ts = (ll)t*t % p;
    while (ts != 1) i++, ts = ((ll)ts * ts) % p;
    ll b = mod_pow(c, 1LL<<(m-i-1), p);
    r = (ll)r * b % p;
    t = (ll)t * b % p * b % p;
    c = (ll)b * b % p;
    m = i; }
  return r; }
// vim: cc=60 ts=2 sts=2 sw=2:
```

1.3. **Numeric Integration.** Numeric integration using Simpson's rule.

```
double integrate(double (*f)(double), double a, double b,
    double delta = 1e-6) {
  if (abs(a - b) < delta)
    return (b-a)/8 *
      (f(a) + 3*f((2*a+b)/3) + 3*f((a+2*b)/3) + f(b));
  return integrate(f, a,
      (a+b)/2, delta) + integrate(f, (a+b)/2, b, delta); }
// vim: cc=60 ts=2 sts=2 sw=2:
```

1.4. **Fast Hadamard Transform.** Computes the Hadamard transform of the given array. Can be used to compute the `XOR`-convolution of arrays, exactly like with FFT. For `AND`-convolution, use $(x + y, y)$ and $(x - y, y)$. For `OR`-convolution, use $(x, x + y)$ and $(x, -x + y)$. **Note**: Size of array must be a power of 2.

```
void fht(vi &arr, bool inv=false, int l=0, int r=-1) {
  if (r == -1) { fht(arr,inv,0,size(arr)); return; }
  if (l+1 == r) return;
  int k = (r-l)/2;
  if (!inv) fht(arr, inv, l, l+k), fht(arr, inv, l+k, r);
  rep(i,l,l+k) { int x = arr[i], y = arr[i+k];
    if (!inv) arr[i] = x-y, arr[i+k] = x+y;
    else arr[i] = (x+y)/2, arr[i+k] = (-x+y)/2; }
  if (inv) fht(arr, inv, l, l+k), fht(arr, inv, l+k, r); }
// vim: cc=60 ts=2 sts=2 sw=2:
```

1.5. **Tridiagonal Matrix Algorithm.** Solves a tridiagonal system of linear equations $a_i x_{i-1} + b_i x_i + c_i x_{i+1} = d_i$ where $a_1 = c_n = 0$. Beware of numerical instability.

```
#define MAXN 5000
long double A[MAXN], B[MAXN], C[MAXN], D[MAXN], X[MAXN];
void solve(int n) {
  C[0] /= B[0]; D[0] /= B[0];
  rep(i,1,n-1) C[i] /= B[i] - A[i]*C[i-1];
  rep(i,1,n)
    D[i] = (D[i] - A[i] * D[i-1]) / (B[i] - A[i] * C[i-1]);
  X[n-1] = D[n-1];
  for (int i = n-2; i>=0; i--)
    X[i] = D[i] - C[i] * X[i+1]; }
// vim: cc=60 ts=2 sts=2 sw=2:
```

1.6. **Mertens Function.** Mertens function is $M(n) = \sum_{i=1}^{n} \mu(i)$. Let $L \approx (n \log\log n)^{2/3}$ and the algorithm runs in $O(n^{2/3})$.

```
#define L 9000000
int mob[L], mer[L];
unordered_map<ll,ll> mem;
ll M(ll n) {
  if (n < L) return mer[n];
  if (mem.find(n) != mem.end()) return mem[n];
  ll ans = 0, done = 1;
  for (ll i = 2; i*i <= n; i++) ans += M(n/i), done = i;
  for (ll i = 1; i*i <= n; i++)
    ans += mer[i] * (n/i - max(done, n/(i+1)));
  return mem[n] = 1 - ans; }
void sieve() {
  for (int i = 1; i < L; i++) mer[i] = mob[i] = 1;
  for (int i = 2; i < L; i++) {
    if (mer[i]) {
      mob[i] = -1;
      for (int j = i+i; j < L; j += i)
        mer[j] = 0, mob[j] = (j/i)%i == 0 ? 0 : -mob[j/i]; }
    mer[i] = mob[i] + mer[i-1]; } }
// vim: cc=60 ts=2 sts=2 sw=2:
```

1.7. **Summatory Phi.** The summatory phi function $\Phi(n) = \sum_{i=1}^{n} \phi(i)$. Let $L \approx (n \log\log n)^{2/3}$ and the algorithm runs in $O(n^{2/3})$.

```
#define N 10000000
ll sp[N];
unordered_map<ll,ll> mem;
ll sumphi(ll n) {
  if (n < N) return sp[n];
  if (mem.find(n) != mem.end()) return mem[n];
  ll ans = 0, done = 1;
  for (ll i = 2; i*i <= n; i++) ans += sumphi(n/i), done = i;
  for (ll i = 1; i*i <= n; i++)
    ans += sp[i] * (n/i - max(done, n/(i+1)));
  return mem[n] = n*(n+1)/2 - ans; }
void sieve() {
  for (int i = 1; i < N; i++) sp[i] = i;
  for (int i = 2; i < N; i++) {
    if (sp[i] == i) {
      sp[i] = i-1;
      for (int j = i+i; j < N; j += i) sp[j] -= sp[j] / i; }
    sp[i] += sp[i-1]; } }
// vim: cc=60 ts=2 sts=2 sw=2:
```

1.8. **Josephus problem.** Last man standing out of n if every kth is killed. Zero-based, and does not kill 0 on first pass.

```
int J(int n, int k) {
  if (n == 1) return 0;
  if (k == 1) return n-1;
  if (n < k) return (J(n-1,k)+k)%n;
  int np = n - n/k;
  return k*((J(np,k)+np-n%k%np)%np) / (k-1); }
// vim: cc=60 ts=2 sts=2 sw=2:
```

1.9. **Number of Integer Points under Line.** Count the number of integer solutions to $Ax + By \leq C$, $0 \leq x \leq n$, $0 \leq y$. In other words, evaluate the sum $\sum_{x=0}^{n} \left\lfloor \frac{C-Ax}{B} + 1 \right\rfloor$. To count all solutions, let $n = \left\lfloor \frac{c}{a} \right\rfloor$. In any case, it must hold that $C - nA \geq 0$. Be very careful about overflows.

```
ll floor_sum(ll n, ll a, ll b, ll c) {
  if (c == 0) return 1;
  if (c < 0) return 0;
  if (a % b == 0) return (n+1)*(c/b+1)-n*(n+1)/2*a/b;
  if (a >= b) return floor_sum(n,a%b,b,c)-a/b*n*(n+1)/2;
  ll t = (c-a*n+b)/b;
  return floor_sum((c-b*t)/b,b,a,c-b*t)+t*(n+1); }
// vim: cc=60 ts=2 sts=2 sw=2:
```

1.10. **Numbers and Sequences.** Some random prime numbers: 1031, 32771, 1048583, 33554467, 1073741827, 34359738421, 1099511627791, 35184372088891, 1125899906842679, 36028797018963971.

More random prime numbers: $10^3 + \{-9, -3, 9, 13\}$, $10^6 + \{-17, 3, 33\}$, $10^9 + \{7, 9, 21, 33, 87\}$.

Some maximal divisor counts:

840	32
720 720	240
735 134 400	1 344
963 761 198 400	6 720
866 421 317 361 600	26 880
897 612 484 786 617 600	103 680

2. Datastructures

2.1. **Standard segment tree** $\mathcal{O}(\log n)$.

```
typedef /* Tree element */ S;
const int n = 1 << 20; S t[2 * n];

// required axiom: associativity
S combine(S l, S r) { return l + r; } // sum segment tree
S combine(S l, S r) { return max(l, r); } // max segment tree

void build() { for (int i = n; --i; ) t[i] = combine(t[2 * i],
 ↪  t[2 * i + 1]); }

// set value v on position i
```

```
void update(int i, S v) { for (t[i += n] = v; i /= 2; ) t[i] =
↪  combine(t[2 * i], t[2 * i + 1]);}
}

// sum on interval [l, r)
S query(int l, int r) {
  S resL, resR;
  for (l += n, r += n; l < r; l /= 2, r /= 2) {
    if (l & 1) resL = combine(resL, t[l++]);
    if (r & 1) resR = combine(t[--r], resR);
  }
  return combine(resL, resR);
}
```

2.2. Binary Indexed Tree $\mathcal{O}(\log n)$. Use one-based indices ($i > 0$)!

```
int bit[MAXN + 1];

// arr[i] += v
void update(int i, int v) {
  while (i <= MAXN) bit[i] += v, i += i & -i;
}

// returns sum of arr[i], where i: [1, i]
int query(int i) {
  int v = 0; while (i) v += bit[i], i -= i & -i; return v;
}
```

2.3. Disjoint-Set / Union-Find $\mathcal{O}(\alpha(n))$.

```
int par[MAXN], rnk[MAXN];

void uf_init(int n) {
  fill_n(par, n, -1); fill_n(rnk, n, 0);
}

int uf_find(int v) { return par[v] < 0 ? v : par[v] =
↪  uf_find(par[v]); }

void uf_union(int a, int b) {
  if ((a = uf_find(a)) == (b = uf_find(b))) return;
  if (rnk[a] < rnk[b]) swap(a, b);
  if (rnk[a] == rnk[b]) rnk[a]++;
  par[b] = a;
}
```

3. Graph Algorithms

3.1. Maximum matching $\mathcal{O}(nm)$. This problem could be solved with a flow algorithm like Dinic's algorithm which runs in $\mathcal{O}(\sqrt{V}E)$, too.

```
const int sizeL = 1e4, sizeR = 1e4;

bool vis[sizeR];
int par[sizeR]; // par : R -> L
vi adj[sizeL]; // adj : L -> (N -> R)

bool match(int u) {
  for (int v : adj[u]) {
    if (vis[v]) continue; vis[v] = true;
    if (par[v] == -1 || match(par[v])) {
      par[v] = u;
      return true;
    }
  }
  return false;
}

// perfect matching iff ret == sizeL == sizeR
int maxmatch() {
  fill_n(par, sizeR, -1); int ret = 0;
  for (int i = 0; i < sizeL; i++) {
    fill_n(vis, sizeR, false);
    ret += match(i);
  }
  return ret;
}
```

3.2. Strongly Connected Components $\mathcal{O}(V + E)$.

```
vvi adj, comps; vi tidx, lnk, cnr, st; vector<bool> vis; int
↪  age, ncomps;

void tarjan(int v) {
  tidx[v] = lnk[v] = ++age; vis[v] = true; st.pb(v);

  for (int w : adj[v]) {
    if (!tidx[w]) tarjan(w), lnk[v] = min(lnk[v], lnk[w]);
    else if (vis[w]) lnk[v] = min(lnk[v], tidx[w]);
  }

  if (lnk[v] != tidx[v]) return;

  comps.pb(vi()); int w;
  do {
    vis[w = st.back()] = false; cnr[w] = ncomps;
↪  comps.back().pb(w);
    st.pop_back();
  } while (w != v);
  ncomps++;
}

void findSCC(int n) {
  age = ncomps = 0; vis.assign(n, false); tidx.assign(n, 0);
↪  lnk.resize(n);
  cnr.resize(n); comps.clear();

  for (int i = 0; i < n; i++)
    if (tidx[i] == 0) tarjan(i);
}
```

3.2.1. *2-SAT* $\mathcal{O}(V + E)$. Include `findSCC`.

```
void init2sat(int n) { adj.assign(2 * n, vi()); }

// vl, vr = true -> variable l, variable r should be negated.
void imply(int xl, bool vl, int xr, bool vr) {
  adj[2 * xl + vl].pb(2 * xr + vr); adj[2 * xr +!vr].pb(2 * xl
↪  +!vl); }

void satOr(int xl, bool vl, int xr, bool vr) { imply(xl, !vl,
↪  xr, vr); }
void satConst(int x, bool v) { imply(x, !v, x, v); }
void satIff(int xl, bool vl, int xr, bool vr) {
  imply(xl, vl, xr, vr); imply(xr, vr, xl, vl);}

bool solve2sat(int n, vector<bool> &sol) {
  findSCC(2 * n);
  for (int i = 0; i < n; i++)
    if (cnr[2 * i] == cnr[2 * i + 1]) return false;
  vector<bool> seen(n, false); sol.assign(n, false);
  for (vi &comp : comps) {
    for (int v : comp) {
      if (seen[v / 2]) continue;
      seen[v / 2] = true; sol[v / 2] = v & 1;
    }
  }
  return true;
}
```

3.3. Cycle Detection $\mathcal{O}(V + E)$.

```
vvi adj; // assumes bidirected graph, adjust accordingly

bool cycle_detection() {
  stack<int> s; vector<bool> vis(MAXN, false); vi par(MAXN,
↪  -1); s.push(0);
  vis[0] = true;
  while(!s.empty()) {
    int cur = s.top(); s.pop();
    for(int i : adj[cur]) {
      if(vis[i] && par[cur] != i) return true;
      s.push(i); par[i] = cur; vis[i] = true;
    }
  }
  return false;}
```

3.4. Shortest path.

3.4.1. *Dijkstra* $\mathcal{O}(E + V \log V)$.

3.4.2. *Floyd-Warshall* $\mathcal{O}(V^3)$.

```
int n = 100; ll d[MAXN][MAXN];
for (int i = 0; i < n; i++) fill_n(d[i], n, 1e18);
// set direct distances from i to j in d[i][j] (and d[j][i])
for (int i = 0; i < n; i++)
  for (int j = 0; j < n; j++)
    for (int k = 0; k < n; k++)
      d[j][k] = min(d[j][k], d[j][i] + d[i][k]);
```

3.4.3. *Bellman Ford* $\mathcal{O}(VE)$. This is only useful if there are edges with weight $w_{ij} < 0$ in the graph.

```
vector< pair<pii, ll> > edges; // ((from, to), weight)
vector<ll> dist;

// when undirected, add back edges
bool bellman_ford(int V, int source) {
  dist.assign(V, 1e18); dist[source] = 0;

  bool updated = true; int loops = 0;
  while (updated && loops < n) {
    updated = false;
    for (auto e : edges) {
      int alt = dist[e.x.x] + e.y;
      if (alt < dist[e.x.y]) {
        dist[e.x.y] = alt; updated = true;
      }
    }
  }
  return loops < n; // loops >= n: negative cycles
}
```

3.5. Max-flow min-cut.

3.5.1. *Dinic's Algorithm* $\mathcal{O}(V^2E)$.

```
struct edge {
  int to, rev; ll cap, flow;
  edge(int t, int r, ll c) : to(t), rev(r), cap(c), flow(0) {}
};

int s, t, level[MAXN]; // s = source, t = sink
vector<edge> g[MAXN];

void add_edge(int fr, int to, ll cap) {
  g[fr].pb(edge(to, g[to].size(), cap)); g[to].pb(edge(fr,
↪  g[fr].size() - 1, 0));
}

bool dinic_bfs() {
  fill_n(level, MAXN, 0); level[s] = 1;

  queue<int> q; q.push(s);
  while (!q.empty()) {
    int cur = q.front(); q.pop();
    for (edge e : g[cur]) {
      if (level[e.to] == 0 && e.flow < e.cap) {
        level[e.to] = level[cur] + 1; q.push(e.to);
      }
    }
  }
  return level[t] != 0;
}

ll dinic_dfs(int cur, ll maxf) {
  if (cur == t) return maxf;

  ll f = 0; bool isSat = true;
  for (edge &e : g[cur]) {
    if (level[e.to] != level[cur] + 1 || e.flow >= e.cap)
      continue;
    ll df = dinic_dfs(e.to, min(maxf - f, e.cap - e.flow));
    f += df; e.flow += df; g[e.to][e.rev].flow -= df; isSat &=
↪  e.flow == e.cap;
    if (maxf == f) break;
  }
  if (isSat) level[cur] = 0;
  return f;
}

ll dinic_maxflow() {
  ll f = 0;
  while (dinic_bfs()) f += dinic_dfs(s, LLINF);
  return f;
}
```

3.6. Min-cost max-flow.

Find the cheapest possible way of sending a certain amount of flow through a flow network.

```
struct edge {
  // to, rev, flow, capacity, weight
  int t, r; ll f, c, w;
  edge(int _t, int _r, ll _c, ll _w) : t(_t), r(_r), f(0),
↪  c(_c), w(_w) {}
};

int n, par[MAXN]; vector<edge> adj[MAXN]; ll dist[MAXN];

bool findPath(int s, int t) {
  fill_n(dist, n, LLINF); fill_n(par, n, -1);

  priority_queue< pii, vector<pii>, greater<pii> > q;
  q.push(pii(dist[s] = 0, s));

  while (!q.empty()) {
    int d = q.top().x, v = q.top().y; q.pop();
    if (d > dist[v]) continue;

    for (edge e : adj[v]) {
      if (e.f < e.c && d + e.w < dist[e.t]) {
        q.push(pii(dist[e.t] = d + e.w, e.t)); par[e.t] = e.r;
      }
    }
  }
  return dist[t] < INF;
}

pair<ll, ll> minCostMaxFlow(int s, int t) {
  ll cost = 0, flow = 0;
  while (findPath(s, t)) {
    ll f = INF, c = 0; int cur = t;
    while (cur != s) {
      const edge &rev = adj[cur][par[cur]], &e =
↪  adj[rev.t][rev.r];
      f = min(f, e.c - e.f); cur = rev.t;
    }
    cur = t;
    while (cur != s) {
      edge &rev = adj[cur][par[cur]], &e = adj[rev.t][rev.r];
      c += e.w; e.f += f; rev.f -= f; cur = rev.t;
    }
    cost += f * c; flow += f;
  }
  return pair<ll, ll>(cost, flow);
}

inline void addEdge(int from, int to, ll cap, ll weight) {
  adj[from].pb(edge(to, adj[to].size(), cap, weight));
  adj[to].pb(edge(from, adj[from].size() - 1, 0, -weight));
}
```

3.7. Minimal Spanning Tree.

3.7.1. *Kruskal* $\mathcal{O}(E \log V)$.

4. String algorithms

4.1. Trie.

```
const int SIGMA = 26;

struct trie {
  bool word; trie **adj;

  trie() : word(false), adj(new trie*[SIGMA]) {
    for (int i = 0; i < SIGMA; i++) adj[i] = NULL;
  }

  void addWord(const string &str) {
    trie *cur = this;
    for (char ch : str) {
      int i = ch - 'a';
      if (!cur->adj[i]) cur->adj[i] = new trie();
      cur = cur->adj[i];
    }
    cur->word = true;
  }

  bool isWord(const string &str) {
    trie *cur = this;
    for (char ch : str) {
      int i = ch - 'a';
      if (!cur->adj[i]) return false;
      cur = cur->adj[i];
    }
    return cur->word;
  }
};
```

4.2. Z-algorithm $\mathcal{O}(n)$.

```
// z[i] = length of longest substring starting from s[i] which
↪ is also a prefix of s.
vi z_function(const string &s) {
  int n = (int) s.length();
  vi z(n);
  for (int i = 1, l = 0, r = 0; i < n; ++i) {
    if (i <= r) z[i] = min (r - i + 1, z[i - l]);
    while (i + z[i] < n && s[z[i]] == s[i + z[i]]) ++z[i];
    if (i + z[i] - 1 > r) l = i, r = i + z[i] - 1;
  }
  return z;
}
```

4.3. Suffix array $\mathcal{O}(n \log^2 n)$.

This creates an array $P[0], P[1], \ldots, P[n-1]$ such that the suffix $S[i \ldots n]$ is the $P[i]^{th}$ suffix of S when lexicographically sorted.

```
typedef pair<pii, int> tii;

const int maxlogn = 17, int maxn = 1 << maxlogn;

tii make_triple(int a, int b, int c) { return tii(pii(a, b),
↪ c); }

int p[maxlogn + 1][maxn]; tii L[maxn];

int suffixArray(string S) {
  int N = S.size(), stp = 1, cnt = 1;
  for (int i = 0; i < N; i++) p[0][i] = S[i];
  for (; cnt < N; stp++, cnt <<= 1) {
    for (int i = 0; i < N; i++)
      L[i] = tii(pii(p[stp-1][i], i + cnt < N ? p[stp-1][i +
↪ cnt] : -1), i);
    sort(L, L + N);
    for (int i = 0; i < N; i++)
      p[stp][L[i].y] = i > 0 && L[i].x == L[i-1].x ?
↪ p[stp][L[i-1].y] : i;
  }
  return stp - 1; // result is in p[stp - 1][0 .. (N - 1)]
}
```

4.4. Longest Common Subsequence $\mathcal{O}(n^2)$.

SUBSTRING: *consecutive characters* !!!

```
int dp[STR_SIZE][STR_SIZE]; // DP problem

int lcs(const string &w1, const string &w2) {
  int n1 = w1.size(), n2 = w2.size();
  for (int i = 0; i < n1; i++) {
    for (int j = 0; j < n2; j++) {
      if (i == 0 || j == 0) dp[i][j] = 0;
      else if (w1[i - 1] == w2[j - 1]) dp[i][j] = dp[i - 1][j
↪ - 1] + 1;
      else dp[i][j] = max(dp[i - 1][j], dp[i][j - 1]);
    }
  }
  return dp[n1][n2];
}

// backtrace
string getLCS(const string &w1, const string &w2) {
  int i = w1.size(), j = w2.size(); string ret = "";
  while (i > 0 && j > 0) {
    if (w1[i - 1] == w2[j - 1]) ret += w1[--i], j--;
    else if (dp[i][j - 1] > dp[i - 1][j]) j--;
    else i--;
  }
  reverse(ret.begin(), ret.end());
  return ret;
}
```

4.5. Levenshtein Distance $\mathcal{O}(n^2)$.

Also known as the 'Edit distance'.

```
int dp[MAX_SIZE][MAX_SIZE]; // DP problem

int levDist(const string &w1, const string &w2) {
  int n1 = w1.size(), n2 = w2.size();
  for (int i = 0; i <= n1; i++) dp[i][0] = i; // removal
  for (int j = 0; j <= n2; j++) dp[0][j] = j; // insertion
  for (int i = 1; i <= n1; i++)
    for (int j = 1; j <= n2; j++)
      dp[i][j] = min(
        1 + min(dp[i - 1][j], dp[i][j - 1]),
        dp[i - 1][j - 1] + (w1[i - 1] != w2[j - 1])
      );
  return dp[n1][n2];
}
```

4.6. Knuth-Morris-Pratt algorithm $\mathcal{O}(N + M)$.

```
int kmp_search(const string &word, const string &text) {
  int n = word.size();
  vi T(n + 1, 0);
  for (int i = 1, j = 0; i < n; ) {
    if (word[i] == word[j]) T[++i] = ++j; // match
    else if (j > 0) j = T[j]; // fallback
    else i++; // no match, keep zero
  }
  int matches = 0;
  for (int i = 0, j = 0; i < text.size(); ) {
    if (text[i] == word[j]) {
      i++;
      if (++j == n) { // match at interval [i - n, i)
        matches++; j = T[j];
      }
    } else if (j > 0) j = T[j];
    else i++;
  }
  return matches;
}
```

4.7. Aho-Corasick Algorithm $\mathcal{O}(N + \sum_{i=1}^{m} |S_i|)$.

All given P must be unique!

```
const int MAXP = 100, MAXLEN = 200, SIGMA = 26, MAXTRIE = MAXP
↪ * MAXLEN;
int nP;
string P[MAXP], S;

int pnr[MAXTRIE], to[MAXTRIE][SIGMA], sLink[MAXTRIE],
↪ dLink[MAXTRIE], nnodes;

void ahoCorasick() {
  fill_n(pnr, MAXTRIE, -1);
  for (int i = 0; i < MAXTRIE; i++) fill_n(to[i], SIGMA, 0);
  fill_n(sLink, MAXTRIE, 0); fill_n(dLink, MAXTRIE, 0);
  nnodes = 1;
  // STEP 1: MAKE A TREE
  for (int i = 0; i < nP; i++) {
    int cur = 0;
    for (char c : P[i]) {
      int i = c - 'a';
      if (to[cur][i] == 0) to[cur][i] = nnodes++;
      cur = to[cur][i];
    }
    pnr[cur] = i;
  }
  // STEP 2: CREATE SUFFIX_LINKS AND DICT_LINKS
  queue<int> q; q.push(0);
  while (!q.empty()) {
    int cur = q.front(); q.pop();
    for (int c = 0; c < SIGMA; c++) {
      if (to[cur][c]) {
        int sl = sLink[to[cur][c]] = cur == 0 ? 0 :
↪ to[sLink[cur]][c];
        // if all strings have equal length, remove this:
        dLink[to[cur][c]] = pnr[sl] >= 0 ? sl : dLink[sl];
        q.push(to[cur][c]);
      } else to[cur][c] = to[sLink[cur]][c];
    }
  }
  // STEP 3: TRAVERSE S
  for (int cur = 0, i = 0, n = S.size(); i < n; i++) {
    cur = to[cur][S[i] - 'a'];
    for (int hit = pnr[cur] >= 0 ? cur : dLink[cur]; hit; hit
↪ = dLink[hit]) {
      cerr << P[pnr[hit]] << " found at [" << (i + 1 -
↪ P[pnr[hit]].size()) << ", " << i << "]" << endl;
    }
  }
}
```

5. Geometry

```
const double EPS = 1e-7, PI = acos(-1.0);

typedef long long NUM; // EITHER double OR long long
typedef pair<NUM, NUM> pt;
#define x first
#define y second
```

```cpp
pt operator+(pt p, pt q) { return pt(p.x + q.x, p.y + q.y); }
pt operator-(pt p, pt q) { return pt(p.x - q.x, p.y - q.y); }

pt& operator+=(pt &p, pt q) { return p = p + q; }
pt& operator-=(pt &p, pt q) { return p = p - q; }

pt operator*(pt p, NUM l) { return pt(p.x * l, p.y * l); }
pt operator/(pt p, NUM l) { return pt(p.x / l, p.y / l); }

NUM operator*(pt p, pt q) { return p.x * q.x + p.y * q.y; }
NUM operator^(pt p, pt q) { return p.x * q.y - p.y * q.x; }

istream& operator>>(istream &in, pt &p) { return in >> p.x >> p.y; }
ostream& operator<<(ostream &out, pt p) { return out << '(' << p.x << ", " << p.y << ')'; }

NUM lenSq(pt p) { return p * p; }
NUM lenSq(pt p, pt q) { return lenSq(p - q); }
double len(pt p) { return hypot(p.x, p.y); } // more overflow safe
double len(pt p, pt q) { return len(p - q); }

typedef pt frac;
typedef pair<double, double> vec;
vec getvec(pt p, pt dp, frac t) { return vec(p.x + 1. * dp.x * t.x / t.y, p.y + 1. * dp.y * t.x / t.y); }

// square distance from pt a to line bc
frac distPtLineSq(pt a, pt b, pt c) {
  a -= b, c -= b;
  return frac((a ^ c) * (a ^ c), c * c);
}

// square distance from pt a to linesegment bc
frac distPtSegmentSq(pt a, pt b, pt c) {
  a -= b; c -= b;
  NUM dot = a * c, len = c * c;
  if (dot <= 0) return frac(a * a, 1);
  if (dot >= len) return frac((a - c) * (a - c), 1);
  return frac(a * a * len - dot * dot, len);
}

// projects pt a onto linesegment bc
frac proj(pt a, pt b, pt c) { return frac((a - b) * (c - b), (c - b) * (c - b)); }
vec projv(pt a, pt b, pt c) { return getvec(b, c - b, proj(a, b, c)); }

bool collinear(pt a, pt b, pt c) { return ((a - b) ^ (a - c)) == 0; }

bool pointOnSegment(pt a, pt b, pt c) {
  NUM dot = (a - b) * (c - b), len = (c - b) * (c - b);
  return collinear(a, b, c) && 0 <= dot && dot <= len;
}

// true => 1 intersection, false => parallel, so 0 or \infty solutions
bool linesIntersect(pt a, pt b, pt c, pt d) { return ((a - b) ^ (c - d)) != 0; }
vec lineLineIntersection(pt a, pt b, pt c, pt d) {
  double det = (a - b) ^ (c - d); pt ret = (c - d) * (a ^ b) - (a - b) * (c ^ d);
  return vec(ret.x / det, ret.y / det);
}

// dp, dq are directions from p, q
// intersection at p + t_i dp, for 0 <= i < return value
int segmentIntersection(pt p, pt dp, pt q, pt dq, frac &t0, frac &t1){
  if (dp * dp == 0) swap(p, q), swap(dp, dq); // dq = 0
  if (dp * dp == 0) { t0 = t1 = frac(0, 1); return p == q; } // dp = dq = 0
  pt dpq = (q - p); NUM c = dp ^ dq, c0 = dpq ^ dp, c1 = dpq ^ dq;
  if (c == 0) { // parallel, dp > 0, dq >= 0
    if (c0 != 0) return 0; // not collinear
    NUM v0 = dpq * dp, v1 = v0 + dq * dp, dp2 = dp * dp;
    if (v1 < v0) swap(v0, v1);
    t0 = frac(v0 = max(v0, (NUM) 0), dp2);
    t1 = frac(v1 = min(v1, dp2), dp2);
    return (v0 <= v1) + (v0 < v1);
  } else if (c < 0) c = -c, c0 = -c0, c1 = -c1;
  t0 = t1 = frac(c1, c);
  return 0 <= min(c0, c1) && max(c0, c1) <= c;
}

// Returns TWICE the area of a polygon to keep it an integer
NUM polygonTwiceArea(const vector<pt> &pts) {
  NUM area = 0;
  for (int N = pts.size(), i = 0, j = N - 1; i < N; j = i++)
    area += pts[i] ^ pts[j];
  return abs(area); // area < 0 <=> pts ccw
}

bool pointInPolygon(pt p, const vector<pt> &pts) {
  double sum = 0;
  for (int N = pts.size(), i = 0, j = N - 1; i < N; j = i++) {
    if (pointOnSegment(p, pts[i], pts[j])) return true; // boundary
    double angle = acos((pts[i] - p) * (pts[j] - p) / len(pts[i], p) / len(pts[j], p));
    sum += ((pts[i] - p) ^ (pts[j] - p)) < 0 ? angle : -angle;}
  return abs(abs(sum) - 2 * PI) < EPS;
}
```

5.1. **Convex Hull** $\mathcal{O}(n \log n)$.

```cpp
// points are given by: pts[ret[0]], pts[ret[1]], ... pts[ret[ret.size()-1]]
vi convexHull(const vector<pt> &pts) {
  if (pts.empty()) return vi();
  vi ret;
  // find one outer point:
  int fsti = 0, n = pts.size(); pt fstpt = pts[0];
  for(int i = n; i--; ) if (pts[i] < fstpt) fstpt = pts[fsti = i];
  ret.pb(fsti); pt refr = pts[fsti];
  vi ord; // index into pts
  for (int i = n; i--; ) if (pts[i] != refr) ord.pb(i);
  sort(ord.begin(), ord.end(), [&pts, &refr] (int a, int b) -> bool {
    NUM cross = (pts[a] - refr) ^ (pts[b] - refr);
    return cross != 0 ? cross > 0 : lenSq(refr, pts[a]) < lenSq(refr, pts[b]);
  });
  for (int i : ord) {
    // NOTE: > INCLUDES points on the hull-line, >= EXCLUDES
    while (ret.size() > 1 &&
        ((pts[ret[ret.size()-2]]-pts[ret.back()]) ^ (pts[i]-pts[ret.back()])) >= 0)
      ret.pop_back();
    ret.pb(i);
  }
  return ret;
}
```

5.2. **Rotating Calipers** $\mathcal{O}(n)$. Finds the longest distance between two points in a convex hull.

```cpp
NUM rotatingCalipers(vector<pt> &hull) {
  int n = hull.size(), a = 0, b = 1;
  if (n <= 1) return 0.0;
  while (((hull[1] - hull[0]) ^ (hull[(b + 1) % n] - hull[b])) > 0) b++;
  NUM ret = 0.0;
  while (a < n) {
    ret = max(ret, lenSq(hull[a], hull[b]));
    if (((hull[(a + 1) % n] - hull[a % n]) ^ (hull[(b + 1) % n] - hull[b])) <= 0) a++;
    else if (++b == n) b = 0;
  }
  return ret;
}
```

5.3. **Closest points** $\mathcal{O}(n \log n)$.

```cpp
int n;pt pts[maxn];

struct byY {
  bool operator()(int a, int b) const { return pts[a].y < pts[b].y; }
};

inline NUM dist(pii p) {
```

```
  return hypot(pts[p.x].x - pts[p.y].x, pts[p.x].y - pts[p.y].y);
}

pii minpt(pii p1, pii p2) { return (dist(p1) < dist(p2)) ? p1 : p2;}

// closest pts (by index) inside pts[l ... r], with sorted y values in ys
pii closest(int l, int r, vi &ys) {
  if (r - l == 2) { // don't assume 1 here.
    ys = { l, l + 1 };
    return pii(l, l + 1);
  } else if (r - l == 3) { // brute-force
    ys = { l, l + 1, l + 2 };
    sort(ys.begin(), ys.end(), byY());
    return minpt(pii(l, l + 1), minpt(pii(l, l + 2), pii(l + 1, l + 2)));
  }
  int m = (l + r) / 2; vi yl, yr;
  pii delta = minpt(closest(l, m, yl), closest(m, r, yr));
  NUM ddelta = dist(delta), xm = .5 * (pts[m-1].x + pts[m].x);
  merge(yl.begin(), yl.end(), yr.begin(), yr.end(), back_inserter(ys), byY());
  deque<int> q;
  for (int i : ys) {
    if (abs(pts[i].x - xm) <= ddelta) {
      for (int j : q) delta = minpt(delta, pii(i, j));
      q.pb(i);
      if (q.size() > 8) q.pop_front(); // magic from Introduction to Algorithms.
    }
  }
  return delta;
}
```

5.4. **Great-Circle Distance.** Computes the distance between two points (given as latitude/longitude coordinates) on a sphere of radius r.

```
double gc_distance(double pLat, double pLong,
         double qLat, double qLong, double r) {
  pLat *= pi / 180; pLong *= pi / 180;
  qLat *= pi / 180; qLong *= pi / 180;
  return r * acos(cos(pLat) * cos(qLat) * cos(pLong - qLong) +
          sin(pLat) * sin(qLat)); }
// vim: cc=60 ts=2 sts=2 sw=2:
```

5.5. **3D Primitives.**

```
#define P(p) const point3d &p
#define L(p0, p1) P(p0), P(p1)
#define PL(p0, p1, p2) P(p0), P(p1), P(p2)
struct point3d {
  double x, y, z;
  point3d() : x(0), y(0), z(0) {}
  point3d(double _x, double _y, double _z)
    : x(_x), y(_y), z(_z) {}
  point3d operator+(P(p)) const {
    return point3d(x + p.x, y + p.y, z + p.z); }
  point3d operator-(P(p)) const {
    return point3d(x - p.x, y - p.y, z - p.z); }
  point3d operator-() const {
    return point3d(-x, -y, -z); }
  point3d operator*(double k) const {
    return point3d(x * k, y * k, z * k); }
  point3d operator/(double k) const {
    return point3d(x / k, y / k, z / k); }
  double operator%(P(p)) const {
    return x * p.x + y * p.y + z * p.z; }
  point3d operator*(P(p)) const {
    return point3d(y*p.z - z*p.y,
                   z*p.x - x*p.z, x*p.y - y*p.x); }
  double length() const {
    return sqrt(*this % *this); }
  double distTo(P(p)) const {
    return (*this - p).length(); }
  double distTo(P(A), P(B)) const {
    // A and B must be two different points
    return ((*this - A) * (*this - B)).length() / A.distTo(B);}
  point3d normalize(double k = 1) const {
    // length() must not return 0
    return (*this) * (k / length()); }
  point3d getProjection(P(A), P(B)) const {
    point3d v = B - A;
    return A + v.normalize((v % (*this - A)) / v.length()); }
  point3d rotate(P(normal)) const {
    //normal must have length 1 and be orthogonal to the vector
    return (*this) * normal; }
  point3d rotate(double alpha, P(normal)) const {
    return (*this) * cos(alpha) + rotate(normal) * sin(alpha);}
  point3d rotatePoint(P(O), P(axe), double alpha) const{
    point3d Z = axe.normalize(axe % (*this - O));
    return O + Z + (*this - O - Z).rotate(alpha, O); }
  bool isZero() const {
    return abs(x) < EPS && abs(y) < EPS && abs(z) < EPS; }
  bool isOnLine(L(A, B)) const {
    return ((A - *this) * (B - *this)).isZero(); }
  bool isInSegment(L(A, B)) const {
    return isOnLine(A, B) && ((A - *this) % (B - *this))<EPS;}
  bool isInSegmentStrictly(L(A, B)) const {
    return isOnLine(A, B) && ((A - *this) % (B - *this))<-EPS;}
  double getAngle() const {
    return atan2(y, x); }
  double getAngle(P(u)) const {
    return atan2((*this * u).length(), *this % u); }
  bool isOnPlane(PL(A, B, C)) const {
    return
      abs((A - *this) * (B - *this) % (C - *this)) < EPS; } };
int line_line_intersect(L(A, B), L(C, D), point3d &O){
  if (abs((B - A) * (C - A) % (D - A)) > EPS) return 0;
  if (((A - B) * (C - D)).length() < EPS)
    return A.isOnLine(C, D) ? 2 : 0;
  point3d normal = ((A - B) * (C - B)).normalize();
  double s1 = (C - A) * (D - A) % normal;
  O = A + ((B - A) / (s1 + ((D - B) * (C - B) % normal))) * s1;
  return 1; }
int line_plane_intersect(L(A, B), PL(C, D, E), point3d & O) {
  double V1 = (C - A) * (D - A) % (E - A);
  double V2 = (D - B) * (C - B) % (E - B);
  if (abs(V1 + V2) < EPS)
    return A.isOnPlane(C, D, E) ? 2 : 0;
  O = A + ((B - A) / (V1 + V2)) * V1;
  return 1; }
bool plane_plane_intersect(P(A), P(nA), P(B), P(nB),
    point3d &P, point3d &Q) {
  point3d n = nA * nB;
  if (n.isZero()) return false;
  point3d v = n * nA;
  P = A + (n * nA) * ((B - A) % nB / (v % nB));
  Q = P + n;
  return true; }
// vim: cc=60 ts=2 sts=2 sw=2:
```

5.6. **Polygon Centroid.**

$$C_x = \frac{1}{6A}\sum_{i=0}^{n-1}(x_i + x_{i+1})(x_i y_{i+1} - x_{i+1} y_i)$$

$$C_Y = \frac{1}{6A}\sum_{i=0}^{n-1}(y_i + y_{i+1})(x_i y_{i+1} - x_{i+1} y_i)$$

$$A = \frac{1}{2}\sum_{i=0}^{n-1}(x_i y_{i+1} - x_{i+1} y_i)$$

5.7. **Rectilinear Minimum Spanning Tree.** Given a set of n points in the plane, and the aim is to find a minimum spanning tree connecting these n points, assuming the Manhattan distance is used. The function `candidates` returns at most $4n$ edges that are a superset of the edges in a minimum spanning tree, and then one can use Kruskal's algorithm.

```
#define MAXN 100100
struct RMST {
  struct point {
    int i; ll x, y;
    point() : i(-1) { }
    ll d1() { return x + y; }
    ll d2() { return x - y; }
    ll dist(point other) {
      return abs(x - other.x) + abs(y - other.y); }
    bool operator <(const point &other) const {
      return y == other.y ? x > other.x : y < other.y; }
  } best[MAXN], arr[MAXN], tmp[MAXN];
  int n;
  RMST() : n(0) {}
  void add_point(int x, int y) {
```

```
    arr[arr[n].i = n].x = x, arr[n++].y = y; }
  void rec(int l, int r) {
    if (l >= r) return;
    int m = (l+r)/2;
    rec(l,m), rec(m+1,r);
    point bst;
    for (int i = l, j = m+1, k = l; i <= m || j <= r; k++) {
      if (j > r || (i <= m && arr[i].d1() < arr[j].d1())) {
        tmp[k] = arr[i++];
        if (bst.i != -1 && (best[tmp[k].i].i == -1
                        || best[tmp[k].i].d2() < bst.d2()))
          best[tmp[k].i] = bst;
      } else {
        tmp[k] = arr[j++];
        if (bst.i == -1 || bst.d2() < tmp[k].d2())
          bst = tmp[k]; } }
    rep(i,l,r+1) arr[i] = tmp[i]; }
  vector<pair<ll,ii> > candidates() {
    vector<pair<ll, ii> > es;
    rep(p,0,2) {
      rep(q,0,2) {
        sort(arr, arr+n);
        rep(i,0,n) best[i].i = -1;
        rec(0,n-1);
        rep(i,0,n) {
          if(best[arr[i].i].i != -1)
            es.push_back({arr[i].dist(best[arr[i].i]),
                          {arr[i].i, best[arr[i].i].i}});
          swap(arr[i].x, arr[i].y);
          arr[i].x *= -1, arr[i].y *= -1; } }
      rep(i,0,n) arr[i].x *= -1; }
    return es; } };
// vim: cc=60 ts=2 sts=2 sw=2:
```

5.8. **Formulas.** Let $a = (a_x, a_y)$ and $b = (b_x, b_y)$ be two-dimensional vectors.

- $a \cdot b = |a||b| \cos\theta$, where θ is the angle between a and b.
- $a \times b = |a||b| \sin\theta$, where θ is the signed angle between a and b.
- $a \times b$ is equal to the area of the parallelogram with two of its sides formed by a and b. Half of that is the area of the triangle formed by a and b.
- **Euler's formula:** $V - E + F = 2$
- Side lengths a, b, c can form a triangle iff. $a + b > c$, $b + c > a$ and $a + c > b$.
- Sum of internal angles of a regular convex n-gon is $(n - 2)\pi$.
- **Law of sines:** $\frac{a}{\sin A} = \frac{b}{\sin B} = \frac{c}{\sin C}$
- **Law of cosines:** $b^2 = a^2 + c^2 - 2ac \cos B$
- Internal tangents of circles $(c_1, r_1), (c_2, r_2)$ intersect at $(c_1 r_2 + c_2 r_1)/(r_1 + r_2)$, external intersect at $(c_1 r_2 - c_2 r_1)/(r_1 + r_2)$.

6. Miscellaneous

6.1. **Binary search** $\mathcal{O}(\log(hi - lo))$.

```
bool test(int n);

int search(int lo, int hi) {
  // assert(test(lo) && !test(hi));
  while (hi - lo > 1) {
    int m = (lo + hi) / 2;
    (test(m) ? lo : hi) = m;
  }
  // assert(test(lo) && !test(hi));
  return lo;
}
```

6.2. **Fast Fourier Transform** $\mathcal{O}(n \log n)$. Given two polynomials $A(x) = a_0 + a_1 x + \cdots + a_{n/2} x^{n/2}$ and $B(x) = b_0 + b_1 x + \cdots + b_{n/2} x^{n/2}$, FFT calculates all coefficients of $C(x) = A(x) \cdot B(x) = c_0 + c_1 x + \ldots c_n x^n$, with $c_i = \sum_{j=0}^{i} a_j b_{i-j}$.

```
typedef complex<double> cpx;
const int logmaxn = 20, maxn = 1 << logmaxn;

cpx a[maxn] = {}, b[maxn] = {}, c[maxn];

void fft(cpx *src, cpx *dest) {
  for (int i = 0, rep = 0; i < maxn; i++, rep = 0) {
    for (int j = i, k = logmaxn; k--; j >>= 1) rep = (rep << 1) | (j & 1);
    dest[rep] = src[i];
  }
  for (int s = 1, m = 1; m <= maxn; s++, m *= 2) {
    cpx r = exp(cpx(0, 2.0 * PI / m));
    for (int k = 0; k < maxn; k += m) {
      cpx cr(1.0, 0.0);
      for (int j = 0; j < m / 2; j++) {
        cpx t = cr * dest[k + j + m / 2]; dest[k + j + m / 2] = dest[k + j] - t;
        dest[k + j] += t; cr *= r;
      }
    }
  }
}

void multiply() {
  fft(a, c); fft(b, a);
  for (int i = 0; i < maxn; i++) b[i] = conj(a[i] * c[i]);
  fft(b, c);
  for (int i = 0; i < maxn; i++) c[i] = conj(c[i]) / (1.0 * maxn);
}
```

6.3. **Minimum Assignment (Hungarian Algorithm)** $\mathcal{O}(n^3)$.

```
int a[MAXN + 1][MAXM + 1]; // matrix, 1-based

int minimum_assignment(int n, int m) { // n rows, m columns
  vi u(n + 1), v(m + 1), p(m + 1), way(m + 1);

  for (int i = 1; i <= n; i++) {
    p[0] = i;
    int j0 = 0;
    vi minv(m + 1, INF);
    vector<char> used(m + 1, false);
    do {
      used[j0] = true;
      int i0 = p[j0], delta = INF, j1;
      for (int j = 1; j <= m; j++)
        if (!used[j]) {
          int cur = a[i0][j] - u[i0] - v[j];
          if (cur < minv[j]) minv[j] = cur, way[j] = j0;
          if (minv[j] < delta) delta = minv[j], j1 = j;
        }
      for (int j = 0; j <= m; j++) {
        if(used[j]) u[p[j]] += delta, v[j] -= delta;
        else minv[j] -= delta;
      }
      j0 = j1;
    } while (p[j0] != 0);
    do {
      int j1 = way[j0]; p[j0] = p[j1]; j0 = j1;
    } while (j0);
  }

  // column j is assigned to row p[j]
  // for (int j = 1; j <= m; ++ j) ans[p[j]] = j;
  return -v[0];
}
```

6.4. **Partial linear equation solver** $\mathcal{O}(N^3)$.

```
typedef double NUM;

#define MAXN 110
#define EPS 1e-5

NUM mat[MAXN][MAXN + 1], vals[MAXN]; bool hasval[MAXN];

bool is_zero(NUM a) { return -EPS < a && a < EPS; }
bool eq(NUM a, NUM b) { return is_zero(a - b); }

int solvemat(int n){ //mat[i][j] contains the matrix A, mat[i][n] contains b
  int pivrow = 0, pivcol = 0;
  while (pivcol < n) {
    int r = pivrow, c;
    while (r < n && is_zero(mat[r][pivcol])) r++;
    if (r == n) { pivcol++; continue; }

    for (c = 0; c <= n; c++) swap(mat[pivrow][c], mat[r][c]);

    r = pivrow++; c = pivcol++;
    NUM div = mat[r][c];
    for (int col = c; col <= n; col++) mat[r][col] /= div;
    for (int row = 0; row < n; row++) {
      if (row == r) continue;
      NUM times = -mat[row][c];
      for (int col = c; col <= n; col++) mat[row][col] += times * mat[r][col];
    }
  } // now mat is in RREF
```

```
  for (int r = pivrow; r < n; r++)
    if (!is_zero(mat[r][n])) return 0;

  fill_n(hasval, n, false);
  for (int col = 0, row; col < n; col++) {
    hasval[col] = !is_zero(mat[row][col]);
    if (!hasval[col]) continue;
    for (int c = col + 1; c < n; c++) {
      if (!is_zero(mat[row][c])) hasval[col] = false;
    }
    if (hasval[col]) vals[col] = mat[row][n];
    row++;
  }

  for (int i = 0; i < n; i++)
    if (!hasval[i]) return 2;
  return 1;
}
```

6.5. Cycle-Finding.

```
ii find_cycle(int x0, int (*f)(int)) {
  int t = f(x0), h = f(t), mu = 0, lam = 1;
  while (t != h) t = f(t), h = f(f(h));
  h = x0;
  while (t != h) t = f(t), h = f(h), mu++;
  h = f(t);
  while (t != h) h = f(h), lam++;
  return ii(mu, lam); }
// vim: cc=60 ts=2 sts=2 sw=2:
```

6.6. Longest Increasing Subsequence.

```
vi lis(vi arr) {
  vi seq, back(size(arr)), ans;
  rep(i,0,size(arr)) {
    int res = 0, lo = 1, hi = size(seq);
    while (lo <= hi) {
      int mid = (lo+hi)/2;
      if (arr[seq[mid-1]] < arr[i]) res = mid, lo = mid + 1;
      else hi = mid - 1; }
    if (res < size(seq)) seq[res] = i;
    else seq.push_back(i);
    back[i] = res == 0 ? -1 : seq[res-1]; }
  int at = seq.back();
  while (at != -1) ans.push_back(at), at = back[at];
  reverse(ans.begin(), ans.end());
  return ans; }
// vim: cc=60 ts=2 sts=2 sw=2:
```

6.7. Dates.

```
int intToDay(int jd) { return jd % 7; }
int dateToInt(int y, int m, int d) {
  return 1461 * (y + 4800 + (m - 14) / 12) / 4 +
    367 * (m - 2 - (m - 14) / 12 * 12) / 12 -
    3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +
    d - 32075; }
void intToDate(int jd, int &y, int &m, int &d) {
  int x, n, i, j;
  x = jd + 68569;
  n = 4 * x / 146097;
  x -= (146097 * n + 3) / 4;
  i = (4000 * (x + 1)) / 1461001;
  x -= 1461 * i / 4 - 31;
  j = 80 * x / 2447;
  d = x - 2447 * j / 80;
  x = j / 11;
  m = j + 2 - 12 * x;
  y = 100 * (n - 49) + i + x; }
// vim: cc=60 ts=2 sts=2 sw=2:
```

6.8. Simplex.

```
typedef long double DOUBLE;
typedef vector<DOUBLE> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
const DOUBLE EPS = 1e-9;
struct LPSolver {
 int m, n;
 VI B, N;
 VVD D;
 LPSolver(const VVD &A, const VD &b, const VD &c) :
  m(b.size()), n(c.size()),
  N(n + 1), B(m), D(m + 2, VD(n + 2)) {
  for (int i = 0; i < m; i++) for (int j = 0; j < n; j++)
    D[i][j] = A[i][j];
  for (int i = 0; i < m; i++) { B[i] = n + i; D[i][n] = -1;
    D[i][n + 1] = b[i]; }
  for (int j = 0; j < n; j++) { N[j] = j; D[m][j] = -c[j]; }
  N[n] = -1; D[m + 1][n] = 1; }
 void Pivot(int r, int s) {
  double inv = 1.0 / D[r][s];
  for (int i = 0; i < m + 2; i++) if (i != r)
   for (int j = 0; j < n + 2; j++) if (j != s)
    D[i][j] -= D[r][j] * D[i][s] * inv;
  for (int j = 0; j < n + 2; j++) if (j != s) D[r][j] *= inv;
  for (int i = 0; i < m + 2; i++) if (i != r) D[i][s] *= -inv;
  D[r][s] = inv;
  swap(B[r], N[s]); }
 bool Simplex(int phase) {
  int x = phase == 1 ? m + 1 : m;
  while (true) {
   int s = -1;
   for (int j = 0; j <= n; j++) {
    if (phase == 2 && N[j] == -1) continue;
    if (s == -1 || D[x][j] < D[x][s] ||
        D[x][j] == D[x][s] && N[j] < N[s]) s = j; }
   if (D[x][s] > -EPS) return true;
   int r = -1;
   for (int i = 0; i < m; i++) {
    if (D[i][s] < EPS) continue;
    if (r == -1 || D[i][n + 1] / D[i][s] < D[r][n + 1] /
        D[r][s] || (D[i][n + 1] / D[i][s]) == (D[r][n + 1] /
        D[r][s]) && B[i] < B[r]) r = i; }
   if (r == -1) return false;
   Pivot(r, s); } }
 DOUBLE Solve(VD &x) {
  int r = 0;
  for (int i = 1; i < m; i++) if (D[i][n + 1] < D[r][n + 1])
    r = i;
  if (D[r][n + 1] < -EPS) {
   Pivot(r, n);
   if (!Simplex(1) || D[m + 1][n + 1] < -EPS)
     return -numeric_limits<DOUBLE>::infinity();
   for (int i = 0; i < m; i++) if (B[i] == -1) {
    int s = -1;
    for (int j = 0; j <= n; j++)
     if (s == -1 || D[i][j] < D[i][s] ||
         D[i][j] == D[i][s] && N[j] < N[s])
       s = j;
    Pivot(i, s); } }
  if (!Simplex(2)) return numeric_limits<DOUBLE>::infinity();
  x = VD(n);
  for (int i = 0; i < m; i++) if (B[i] < n)
    x[B[i]] = D[i][n + 1];
  return D[m][n + 1]; } };
// Two-phase simplex algorithm for solving linear programs
// of the form
//     maximize     c^T x
//     subject to   Ax <= b
//                  x >= 0
// INPUT: A -- an m x n matrix
//        b -- an m-dimensional vector
//        c -- an n-dimensional vector
//        x -- a vector where the optimal solution will be
//             stored
// OUTPUT: value of the optimal solution (infinity if
//                  unbounded above, nan if infeasible)
// To use this code, create an LPSolver object with A, b,
// and c as arguments.  Then, call Solve(x).
// #include <iostream>
// #include <iomanip>
// #include <vector>
// #include <cmath>
// #include <limits>
// using namespace std;
// int main() {
//   const int m = 4;
//   const int n = 3;
//   DOUBLE _A[m][n] = {
//     { 6, -1, 0 },
//     { -1, -5, 0 },
//     { 1, 5, 1 },
//     { -1, -5, -1 }
//   };
//   DOUBLE _b[m] = { 10, -4, 5, -5 };
//   DOUBLE _c[n] = { 1, -1, 0 };
//   VVD A(m);
//   VD b(_b, _b + m);
//   VD c(_c, _c + n);
```

```
//   for (int i = 0; i < m; i++) A[i] = VD(_A[i], _A[i] + n);
//   LPSolver solver(A, b, c);
//   VD x;
//   DOUBLE value = solver.Solve(x);
//   cerr << "VALUE: " << value << endl; // VALUE: 1.29032
//   cerr << "SOLUTION:"; // SOLUTION: 1.74194 0.451613 1
//   for (size_t i = 0; i < x.size(); i++) cerr << " " <<
 ↪ x[i];
//   cerr << endl;
//   return 0;
// }
// vim: cc=60 ts=2 sts=2 sw=2:
```

7. Geometry (CP3)

7.1. Points and lines.

```
#define INF 1e9
#define EPS 1e-9
#define PI acos(-1.0) // important constant; alternative
 ↪ #define PI (2.0 * acos(0.0))

double DEG_to_RAD(double d) { return d * PI / 180.0; }

double RAD_to_DEG(double r) { return r * 180.0 / PI; }

struct point { double x, y;   // only used if more precision
 ↪ is needed
  point() { x = y = 0.0; }                         // default
 ↪ constructor
  point(double _x, double _y) : x(_x), y(_y) {}        //
 ↪ user-defined
  bool operator < (point other) const { // override less than
 ↪ operator
    if (fabs(x - other.x) > EPS)                 // useful for
 ↪ sorting
      return x < other.x;          // first criteria , by
 ↪ x-coordinate
    return y < other.y; }          // second criteria, by
 ↪ y-coordinate
  // use EPS (1e-9) when testing equality of two floating
 ↪ points
  bool operator == (point other) const {
   return (fabs(x - other.x) < EPS && (fabs(y - other.y) <
 ↪ EPS)); } };

double dist(point p1, point p2) {                // Euclidean
 ↪ distance
                      // hypot(dx, dy) returns sqrt(dx * dx +
 ↪ dy * dy)
  return hypot(p1.x - p2.x, p1.y - p2.y); }           //
 ↪ return double

// rotate p by theta degrees CCW w.r.t origin (0, 0)
point rotate(point p, double theta) {
  double rad = DEG_to_RAD(theta);    // multiply theta with PI
 ↪ / 180.0
  return point(p.x * cos(rad) - p.y * sin(rad),
               p.x * sin(rad) + p.y * cos(rad)); }

struct line { double a, b, c; };          // a way to
 ↪ represent a line

// the answer is stored in the third parameter (pass by
 ↪ reference)
void pointsToLine(point p1, point p2, line &l) {
  if (fabs(p1.x - p2.x) < EPS) {              // vertical line
 ↪ is fine
    l.a = 1.0;   l.b = 0.0;   l.c = -p1.x;           //
 ↪ default values
  } else {
    l.a = -(double)(p1.y - p2.y) / (p1.x - p2.x);
    l.b = 1.0;              // IMPORTANT: we fix the value of
 ↪ b to 1.0
    l.c = -(double)(l.a * p1.x) - p1.y;
} }

bool areParallel(line l1, line l2) {       // check
 ↪ coefficients a & b
  return (fabs(l1.a-l2.a) < EPS) && (fabs(l1.b-l2.b) < EPS); }

bool areSame(line l1, line l2) {           // also check
 ↪ coefficient c
  return areParallel(l1 ,l2) && (fabs(l1.c - l2.c) < EPS); }

// returns true (+ intersection point) if two lines are
 ↪ intersect
bool areIntersect(line l1, line l2, point &p) {
  if (areParallel(l1, l2)) return false;           // no
 ↪ intersection
  // solve system of 2 linear algebraic equations with 2
 ↪ unknowns
  p.x = (l2.b * l1.c - l1.b * l2.c) / (l2.a * l1.b - l1.a *
 ↪ l2.b);
  // special case: test for vertical line to avoid division by
 ↪ zero
  if (fabs(l1.b) > EPS) p.y = -(l1.a * p.x + l1.c);
  else                  p.y = -(l2.a * p.x + l2.c);
  return true; }

struct vec { double x, y;  // name: `vec' is different from
 ↪ STL vector
  vec(double _x, double _y) : x(_x), y(_y) {} };

vec toVec(point a, point b) {       // convert 2 points to
 ↪ vector a->b
  return vec(b.x - a.x, b.y - a.y); }

vec scale(vec v, double s) {        // nonnegative s = [<1 ..
 ↪ 1 .. >1]
  return vec(v.x * s, v.y * s); }                //
 ↪ shorter.same.longer

point translate(point p, vec v) {        // translate p
 ↪ according to v
  return point(p.x + v.x , p.y + v.y); }

// convert point and gradient/slope to line
void pointSlopeToLine(point p, double m, line &l) {
  l.a = -m;                                                //
 ↪ always -m
  l.b = 1;                                                 //
 ↪ always 1
  l.c = -((l.a * p.x) + (l.b * p.y)); }                //
 ↪ compute this

void closestPoint(line l, point p, point &ans) {
  line perpendicular;         // perpendicular to l and pass
 ↪ through p
  if (fabs(l.b) < EPS) {              // special case 1:
 ↪ vertical line
    ans.x = -(l.c);   ans.y = p.y;      return; }

  if (fabs(l.a) < EPS) {             // special case 2:
 ↪ horizontal line
    ans.x = p.x;      ans.y = -(l.c);   return; }

  pointSlopeToLine(p, 1 / l.a, perpendicular);          //
 ↪ normal line
  // intersect line l with this perpendicular line
  // the intersection point is the closest point
  areIntersect(l, perpendicular, ans); }

// returns the reflection of point on a line
void reflectionPoint(line l, point p, point &ans) {
  point b;
  closestPoint(l, p, b);                       // similar to
 ↪ distToLine
  vec v = toVec(p, b);                                // create a
 ↪ vector
  ans = translate(translate(p, v), v); }         // translate
 ↪ p twice

double dot(vec a, vec b) { return (a.x * b.x + a.y * b.y); }

double norm_sq(vec v) { return v.x * v.x + v.y * v.y; }

// returns the distance from p to the line defined by
// two points a and b (a and b must be different)
// the closest point is stored in the 4th parameter (byref)
double distToLine(point p, point a, point b, point &c) {
  // formula: c = a + u * ab
  vec ap = toVec(a, p), ab = toVec(a, b);
  double u = dot(ap, ab) / norm_sq(ab);
```

```
  c = translate(a, scale(ab, u));                  //
↪  translate a to c
  return dist(p, c); }           // Euclidean distance between
↪  p and c

// returns the distance from p to the line segment ab defined
↪  by
// two points a and b (still OK if a == b)
// the closest point is stored in the 4th parameter (byref)
double distToLineSegment(point p, point a, point b, point &c)
↪  {
  vec ap = toVec(a, p), ab = toVec(a, b);
  double u = dot(ap, ab) / norm_sq(ab);
  if (u < 0.0) { c = point(a.x, a.y);                   //
↪  closer to a
    return dist(p, a); }         // Euclidean distance between
↪  p and a
  if (u > 1.0) { c = point(b.x, b.y);                   //
↪  closer to b
    return dist(p, b); }         // Euclidean distance between
↪  p and b
  return distToLine(p, a, b, c); }          // run distToLine
↪  as above

double angle(point a, point o, point b) {  // returns angle
↪  aob in rad
  vec oa = toVec(o, a), ob = toVec(o, b);
  return acos(dot(oa, ob) / sqrt(norm_sq(oa) * norm_sq(ob)));
↪  }

double cross(vec a, vec b) { return a.x * b.y - a.y * b.x; }

// note: to accept collinear points, we have to change the `>
↪  0'
// returns true if point r is on the left side of line pq
bool ccw(point p, point q, point r) {
  return cross(toVec(p, q), toVec(p, r)) > 0; }

// returns true if point r is on the same line as the line pq
bool collinear(point p, point q, point r) {
  return fabs(cross(toVec(p, q), toVec(p, r))) < EPS; }
```

7.2. Polygon.

```
// returns the perimeter, which is the sum of Euclidian
↪  distances
// of consecutive line segments (polygon edges)
double perimeter(const vector<point> &P) {
  double result = 0.0;
  for (int i = 0; i < (int)P.size()-1; i++)  // remember that
↪  P[0] = P[n-1]
    result += dist(P[i], P[i+1]);
  return result; }

// returns the area, which is half the determinant
double area(const vector<point> &P) {
  double result = 0.0, x1, y1, x2, y2;
  for (int i = 0; i < (int)P.size()-1; i++) {
    x1 = P[i].x; x2 = P[i+1].x;
    y1 = P[i].y; y2 = P[i+1].y;
    result += (x1 * y2 - x2 * y1);
  }
  return fabs(result) / 2.0; }

// returns true if we always make the same turn while
↪  examining
// all the edges of the polygon one by one
bool isConvex(const vector<point> &P) {
  int sz = (int)P.size();
  if (sz <= 3) return false;   // a point/sz=2 or a line/sz=3
↪  is not convex
  bool isLeft = ccw(P[0], P[1], P[2]);               //
↪  remember one result
  for (int i = 1; i < sz-1; i++)            // then compare
↪  with the others
    if (ccw(P[i], P[i+1], P[(i+2) == sz ? 1 : i+2]) != isLeft)
      return false;            // different sign -> this
↪  polygon is concave
  return true; }                                 // this
↪  polygon is convex

// returns true if point p is in either convex/concave polygon
↪  P
bool inPolygon(point pt, const vector<point> &P) {
  if ((int)P.size() == 0) return false;
  double sum = 0;    // assume the first vertex is equal to
↪  the last vertex
  for (int i = 0; i < (int)P.size()-1; i++) {
    if (ccw(pt, P[i], P[i+1]))
        sum += angle(P[i], pt, P[i+1]);                   //
↪  left turn/ccw
    else sum -= angle(P[i], pt, P[i+1]); }                //
↪  right turn/cw
  return fabs(fabs(sum) - 2*PI) < EPS; }

// line segment p-q intersect with line A-B.
point lineIntersectSeg(point p, point q, point A, point B) {
  double a = B.y - A.y;
  double b = A.x - B.x;
  double c = B.x * A.y - A.x * B.y;
  double u = fabs(a * p.x + b * p.y + c);
  double v = fabs(a * q.x + b * q.y + c);
  return point((p.x * v + q.x * u) / (u+v), (p.y * v + q.y *
↪  u) / (u+v)); }

// cuts polygon Q along the line formed by point a -> point b
// (note: the last point must be the same as the first point)
vector<point> cutPolygon(point a, point b, const vector<point>
↪  &Q) {
  vector<point> P;
  for (int i = 0; i < (int)Q.size(); i++) {
    double left1 = cross(toVec(a, b), toVec(a, Q[i])), left2 =
↪  0;
    if (i != (int)Q.size()-1) left2 = cross(toVec(a, b),
↪  toVec(a, Q[i+1]));
    if (left1 > -EPS) P.push_back(Q[i]);       // Q[i] is on
↪  the left of ab
    if (left1 * left2 < -EPS)        // edge (Q[i], Q[i+1])
↪  crosses line ab
      P.push_back(lineIntersectSeg(Q[i], Q[i+1], a, b));
  }
  if (!P.empty() && !(P.back() == P.front()))
    P.push_back(P.front());        // make P's first point =
↪  P's last point
  return P; }

point pivot;
bool angleCmp(point a, point b) {                  //
↪  angle-sorting function
  if (collinear(pivot, a, b))                                 //
↪  special case
    return dist(pivot, a) < dist(pivot, b);    // check which
↪  one is closer
  double d1x = a.x - pivot.x, d1y = a.y - pivot.y;
  double d2x = b.x - pivot.x, d2y = b.y - pivot.y;
  return (atan2(d1y, d1x) - atan2(d2y, d2x)) < 0; }   //
↪  compare two angles

vector<point> CH(vector<point> P) {   // the content of P may
↪  be reshuffled
  int i, j, n = (int)P.size();
  if (n <= 3) {
    if (!(P[0] == P[n-1])) P.push_back(P[0]); // safeguard
↪  from corner case
    return P;                              // special case, the
↪  CH is P itself
  }

  // first, find P0 = point with lowest Y and if tie:
↪  rightmost X
  int P0 = 0;
  for (i = 1; i < n; i++)
    if (P[i].y < P[P0].y || (P[i].y == P[P0].y && P[i].x >
↪  P[P0].x))
      P0 = i;

  point temp = P[0]; P[0] = P[P0]; P[P0] = temp;    // swap
↪  P[P0] with P[0]

  // second, sort points by angle w.r.t. pivot P0
  pivot = P[0];                       // use this global variable
↪  as reference
  sort(++P.begin(), P.end(), angleCmp);              // we do
↪  not sort P[0]
```

```cpp
  // third, the ccw tests
  vector<point> S;
  S.push_back(P[n-1]); S.push_back(P[0]); S.push_back(P[1]); // initial S
  i = 2;                                                      // then, we check the rest
  while (i < n) {           // note: N must be >= 3 for this method to work
    j = (int)S.size()-1;
    if (ccw(S[j-1], S[j], P[i])) S.push_back(P[i++]);  // left turn, accept
    else S.pop_back(); }   // or pop the top of S until we have a left turn
  return S; }                                         // return the result
```

7.3. Triangle.

```cpp
double perimeter(double ab, double bc, double ca) {
  return ab + bc + ca; }

double perimeter(point a, point b, point c) {
  return dist(a, b) + dist(b, c) + dist(c, a); }

double area(double ab, double bc, double ca) {
  // Heron's formula, split sqrt(a * b) into sqrt(a) * sqrt(b); in implementation
  double s = 0.5 * perimeter(ab, bc, ca);
  return sqrt(s) * sqrt(s - ab) * sqrt(s - bc) * sqrt(s - ca); }

double area(point a, point b, point c) {
  return area(dist(a, b), dist(b, c), dist(c, a)); }

double rInCircle(double ab, double bc, double ca) {
  return area(ab, bc, ca) / (0.5 * perimeter(ab, bc, ca)); }

double rInCircle(point a, point b, point c) {
  return rInCircle(dist(a, b), dist(b, c), dist(c, a)); }

// assumption: the required points/lines functions have been written
// returns 1 if there is an inCircle center, returns 0 otherwise
// if this function returns 1, ctr will be the inCircle center
// and r is the same as rInCircle
int inCircle(point p1, point p2, point p3, point &ctr, double &r) {
  r = rInCircle(p1, p2, p3);
  if (fabs(r) < EPS) return 0;                      // no inCircle center

  line l1, l2;                    // compute these two angle bisectors
  double ratio = dist(p1, p2) / dist(p1, p3);
  point p = translate(p2, scale(toVec(p2, p3), ratio / (1 + ratio)));
  pointsToLine(p1, p, l1);

  ratio = dist(p2, p1) / dist(p2, p3);
  p = translate(p1, scale(toVec(p1, p3), ratio / (1 + ratio)));
  pointsToLine(p2, p, l2);

  areIntersect(l1, l2, ctr);           // get their intersection point
  return 1; }

double rCircumCircle(double ab, double bc, double ca) {
  return ab * bc * ca / (4.0 * area(ab, bc, ca)); }

double rCircumCircle(point a, point b, point c) {
  return rCircumCircle(dist(a, b), dist(b, c), dist(c, a)); }

// assumption: the required points/lines functions have been written
// returns 1 if there is a circumCenter center, returns 0 otherwise
// if this function returns 1, ctr will be the circumCircle center
// and r is the same as rCircumCircle
int circumCircle(point p1, point p2, point p3, point &ctr, double &r){
  double a = p2.x - p1.x, b = p2.y - p1.y;
  double c = p3.x - p1.x, d = p3.y - p1.y;
  double e = a * (p1.x + p2.x) + b * (p1.y + p2.y);
  double f = c * (p1.x + p3.x) + d * (p1.y + p3.y);
  double g = 2.0 * (a * (p3.y - p2.y) - b * (p3.x - p2.x));
  if (fabs(g) < EPS) return 0;

  ctr.x = (d*e - b*f) / g;
  ctr.y = (a*f - c*e) / g;
  r = dist(p1, ctr);  // r = distance from center to 1 of the 3 points
  return 1; }

// returns true if point d is inside the circumCircle defined by a,b,c
int inCircumCircle(point a, point b, point c, point d) {
  return (a.x - d.x) * (b.y - d.y) * ((c.x - d.x) * (c.x - d.x) + (c.y - d.y) * (c.y - d.y)) +
         (a.y - d.y) * ((b.x - d.x) * (b.x - d.x) + (b.y - d.y) * (b.y - d.y)) * (c.x - d.x) +
         ((a.x - d.x) * (a.x - d.x) + (a.y - d.y) * (a.y - d.y)) * (b.x - d.x) * (c.y - d.y) -
         ((a.x - d.x) * (a.x - d.x) + (a.y - d.y) * (a.y - d.y)) * (b.y - d.y) * (c.x - d.x) -
         (a.y - d.y) * (b.x - d.x) * ((c.x - d.x) * (c.x - d.x) + (c.y - d.y) * (c.y - d.y)) -
         (a.x - d.x) * ((b.x - d.x) * (b.x - d.x) + (b.y - d.y) * (b.y - d.y)) * (c.y - d.y) > 0 ? 1 : 0;
}

bool canFormTriangle(double a, double b, double c) {
  return (a + b > c) && (a + c > b) && (b + c > a); }
```

7.4. Circle.

```cpp
int insideCircle(point_i p, point_i c, int r) { // all integer version
  int dx = p.x - c.x, dy = p.y - c.y;
  int Euc = dx * dx + dy * dy, rSq = r * r;             // all integer
  return Euc < rSq ? 0 : Euc == rSq ? 1 : 2; } //inside/border/outside

bool circle2PtsRad(point p1, point p2, double r, point &c) {
  double d2 = (p1.x - p2.x) * (p1.x - p2.x) +
              (p1.y - p2.y) * (p1.y - p2.y);
  double det = r * r / d2 - 0.25;
  if (det < 0.0) return false;
  double h = sqrt(det);
  c.x = (p1.x + p2.x) * 0.5 + (p1.y - p2.y) * h;
  c.y = (p1.y + p2.y) * 0.5 + (p2.x - p1.x) * h;
  return true; }         // to get the other center, reverse p1 and p2
```

8. Combinatorics

Catalan	$C_0 = 1,\ C_n = \frac{1}{n+1}\binom{2n}{n} = \sum_{i=0}^{n-1} C_i C_{n-i-1} = \frac{4n-2}{n+1}$
Stirling 1st kind	$\left[{0 \atop 0}\right] = 1,\ \left[{n \atop 0}\right] = \left[{0 \atop n}\right] = 0,\ \left[{n \atop k}\right] = (n-1)\left[{n-1 \atop k}\right] + \left[{n \atop k}\right.$
Stirling 2nd kind	$\left\{{n \atop 1}\right\} = \left\{{n \atop n}\right\} = 1,\ \left\{{n \atop k}\right\} = k\left\{{n-1 \atop k}\right\} + \left\{{n-1 \atop k-1}\right\}$
Euler	$\left\langle{n \atop 0}\right\rangle = \left\langle{n \atop n-1}\right\rangle = 1,\ \left\langle{n \atop k}\right\rangle = (k+1)\left\langle{n-1 \atop k}\right\rangle + (n-k)$
Euler 2nd Order	$\left\langle\!\left\langle{n \atop k}\right\rangle\!\right\rangle = (k+1)\left\langle\!\left\langle{n-1 \atop k}\right\rangle\!\right\rangle + (2n-k-1)\left\langle\!\left\langle{n-1 \atop k-1}\right\rangle\!\right\rangle$
Bell	$B_1 = 1,\ B_n = \sum_{k=0}^{n-1} B_k \binom{n-1}{k} = \sum_{k=0}^{n} \left\{{n \atop k}\right\}$

#labeled rooted trees	n^{n-1}
#labeled unrooted trees	n^{n-2}
#forests of k rooted trees	$\frac{k}{n}\binom{n}{k}n^{n-k}$
$\sum_{i=1}^{n} i^2 = n(n+1)(2n+1)/6$	$\sum_{i=1}^{n} i^3 = n^2(n+1)^2/4$
$!n = n\times!(n-1) + (-1)^n$	$!n = (n-1)(!(n-1)+!($
$\sum_{i=1}^{n}\binom{n}{i}F_i = F_{2n}$	$\sum_i \binom{n-i}{i} = F_{n+1}$
$\sum_{k=0}^{n}\binom{k}{m} = \binom{n+1}{m+1}$	$x^k = \sum_{i=0}^{k} i!\left\{{k \atop i}\right\}\binom{x}{i} = \sum$
$a \equiv b \pmod{x, y} \Rightarrow a \equiv b \pmod{\mathrm{lcm}(x, y)}$	$\sum_{d\mid n}\phi(d) = n$
$ac \equiv bc \pmod{m} \Rightarrow a \equiv b \pmod{\frac{m}{\gcd(c,m)}}$	$(\sum_{d\mid n}\sigma_0(d))^2 = \sum_{d\mid n}\sigma$
p prime $\Leftrightarrow (p-1)! \equiv -1 \pmod{p}$	$\gcd(n^a - 1, n^b - 1) = n^{\mathrm{g}}$
$\sigma_x(n) = \prod_{i=0}^{r}\frac{p_i^{(a_i+1)x}-1}{p_i^x - 1}$	$\sigma_0(n) = \prod_{i=0}^{r}(a_i + 1)$
$\sum_{k=0}^{m}(-1)^k\binom{n}{k} = (-1)^m\binom{n-1}{m}$	
$2^{\omega(n)} = O(\sqrt{n})$	$\sum_{i=1}^{n} 2^{\omega(i)} = O(n\log n)$
$d = v_i t + \frac{1}{2}at^2$	$v_f^2 = v_i^2 + 2ad$
$v_f = v_i + at$	$d = \frac{v_i + v_f}{2}t$

8.1. **The Twelvefold Way.** Putting n balls into k boxes.

Balls Boxes	same same	distinct same	same distinct	distinct distinct	Remarks
-	$\mathrm{p}_k(n)$	$\sum_{i=0}^{k} \left\{ {n \atop i} \right\}$	$\binom{n+k-1}{k-1}$	k^n	$\mathrm{p}_k(n)$: #partitions of n into $\leq k$ positive parts
size ≥ 1	$\mathrm{p}(n,k)$	$\left\{ {n \atop k} \right\}$	$\binom{n-1}{k-1}$	$k!\left\{ {n \atop k} \right\}$	$\mathrm{p}(n,k)$: #partitions of n into k positive parts
size ≤ 1	$[n \leq k]$	$[n \leq k]$	$\binom{k}{n}$	$n!\binom{k}{n}$	$[cond]$: 1 if $cond = true$, else 0

9. Useful Information

10. Misc

10.1. Debugging Tips.

- Stack overflow? Recursive DFS on tree that is actually a long path?
- Floating-point numbers
  - Getting `NaN`? Make sure `acos` etc. are not getting values out of their range (perhaps `1+eps`).
  - Rounding negative numbers?
  - Outputting in scientific notation?
- Wrong Answer?
  - Read the problem statement again!
  - Are multiple test cases being handled correctly? Try repeating the same test case many times.
  - Integer overflow?
  - Think very carefully about boundaries of all input parameters
  - Try out possible edge cases:
    - $n = 0, n = -1, n = 1, n = 2^{31} - 1$ or $n = -2^{31}$
    - List is empty, or contains a single element
    - n is even, n is odd
    - Graph is empty, or contains a single vertex
    - Graph is a multigraph (loops or multiple edges)
    - Polygon is concave or non-simple
  - Is initial condition wrong for small cases?
  - Are you sure the algorithm is correct?
  - Explain your solution to someone.
  - Are you using any functions that you don't completely understand? Maybe STL functions?
  - Maybe you (or someone else) should rewrite the solution?
  - Can the input line be empty?
- Run-Time Error?
  - Is it actually Memory Limit Exceeded?

10.2. Solution Ideas.

- Dynamic Programming
  - Parsing CFGs: CYK Algorithm
  - Drop a parameter, recover from others
  - Swap answer and a parameter
  - When grouping: try splitting in two
  - 2^k trick
  - When optimizing
    - Convex hull optimization
      - $\text{dp}[i] = \min_{j<i}\{\text{dp}[j] + b[j] \times a[i]\}$
      - $b[j] \geq b[j+1]$
      - optionally $a[i] \leq a[i+1]$
      - $O(n^2)$ to $O(n)$
    - Divide and conquer optimization
      - $\text{dp}[i][j] = \min_{k<j}\{\text{dp}[i-1][k] + C[k][j]\}$
      - $A[i][j] \leq A[i][j+1]$
      - $O(kn^2)$ to $O(kn \log n)$
      - sufficient: $C[a][c] + C[b][d] \leq C[a][d] + C[b][c]$, $a \leq b \leq c \leq d$ (QI)
    - Knuth optimization
      - $\text{dp}[i][j] = \min_{i<k<j}\{\text{dp}[i][k] + \text{dp}[k][j] + C[i][j]\}$
      - $A[i][j-1] \leq A[i][j] \leq A[i+1][j]$
      - $O(n^3)$ to $O(n^2)$
      - sufficient: QI and $C[b][c] \leq C[a][d]$, $a \leq b \leq c \leq d$
- Greedy
- Randomized
- Optimizations
  - Use bitset (/64)
  - Switch order of loops (cache locality)
- Process queries offline
  - Mo's algorithm
- Square-root decomposition
- Precomputation
- Efficient simulation
  - Mo's algorithm
  - Sqrt decomposition
  - Store 2^k jump pointers
- Data structure techniques
  - Sqrt buckets
  - Store 2^k jump pointers
  - 2^k merging trick
- Counting
  - Inclusion-exclusion principle
  - Generating functions
- Graphs
  - Can we model the problem as a graph?
  - Can we use any properties of the graph?
  - Strongly connected components
  - Cycles (or odd cycles)
  - Bipartite (no odd cycles)
    - Bipartite matching
    - Hall's marriage theorem
    - Stable Marriage
  - Cut vertex/bridge
  - Biconnected components
  - Degrees of vertices (odd/even)
  - Trees
    - Heavy-light decomposition
    - Centroid decomposition
    - Least common ancestor
    - Centers of the tree
  - Eulerian path/circuit
  - Chinese postman problem
  - Topological sort
  - (Min-Cost) Max Flow
  - Min Cut
    - Maximum Density Subgraph
  - Huffman Coding
  - Min-Cost Arborescence
  - Steiner Tree
  - Kirchoff's matrix tree theorem
  - Prüfer sequences
  - Lovász Toggle
  - Look at the DFS tree (which has no cross-edges)
  - Is the graph a DFA or NFA?
    - Is it the Synchronizing word problem?
- math
  - Is the function multiplicative?
  - Look for a pattern
  - Permutations
    - Consider the cycles of the permutation
  - Functions
    - Sum of piecewise-linear functions is a piecewise-linear function
    - Sum of convex (concave) functions is convex (concave)
  - Modular arithmetic
    - Chinese Remainder Theorem
    - Linear Congruence
  - Sieve
  - System of linear equations
  - Values too big to represent?
    - Compute using the logarithm
    - Divide everything by some large value
  - Linear programming
    - Is the dual problem easier to solve?
  - Can the problem be modeled as a different combinatorial problem? Does that simplify calculations?
- Logic
  - 2-SAT
  - XOR-SAT (Gauss elimination or Bipartite matching)
- Meet in the middle
- Only work with the smaller half ($\log(n)$)
- Strings
  - Trie (maybe over something weird, like bits)
  - Suffix array
  - Suffix automaton (+DP?)
  - Aho-Corasick
  - eerTree
  - Work with $S + S$
- Hashing
- Euler tour, tree to array
- Segment trees
  - Lazy propagation
  - Persistent
  - Implicit
  - Segment tree of X
- Geometry
  - Minkowski sum (of convex sets)
  - Rotating calipers
  - Sweep line (horizontally or vertically?)
  - Sweep angle
  - Convex hull
- Fix a parameter (possibly the answer).
- Are there few distinct values?
- Binary search
- Sliding Window (+ Monotonic Queue)
- Computing a Convolution? Fast Fourier Transform
- Computing a 2D Convolution? FFT on each row, and then on each column
- Exact Cover (+ Algorithm X)
- Cycle-Finding
- What is the smallest set of values that identify the solution? The cycle structure of the permutation? The powers of primes in the factorization?
- Look at the complement problem

  - Minimize something instead of maximizing
- Immediately enforce necessary conditions. (All values greater than 0? Initialize them all to 1)
- Add large constant to negative numbers to make them positive
- Counting/Bucket sort

11. Formulas

- **Legendre symbol:** $\left(\frac{a}{b}\right) = a^{(b-1)/2} \pmod{b}$, b odd prime.
- **Heron's formula:** A triangle with side lengths a, b, c has area $\sqrt{s(s-a)(s-b)(s-c)}$ where $s = \frac{a+b+c}{2}$.
- **Pick's theorem:** A polygon on an integer grid strictly containing i lattice points and having b lattice points on the boundary has area $i + \frac{b}{2} - 1$. (Nothing similar in higher dimensions)
- **Euler's totient:** The number of integers less than n that are coprime to n are $n\prod_{p|n}\left(1-\frac{1}{p}\right)$ where each p is a distinct prime factor of n.
- **König's theorem:** In any bipartite graph $G = (L \cup R, E)$, the number of edges in a maximum matching is equal to the number of vertices in a minimum vertex cover. Let U be the set of unmatched vertices in L, and Z be the set of vertices that are either in U or are connected to U by an alternating path. Then $K = (L \setminus Z) \cup (R \cap Z)$ is the minimum vertex cover.
- A minumum Steiner tree for n vertices requires at most $n-2$ additional Steiner vertices.
- The number of vertices of a graph is equal to its minimum vertex cover number plus the size of a maximum independent set.
- **Lagrange polynomial** through points $(x_0, y_0), \ldots, (x_k, y_k)$ is $L(x) = \sum_{j=0}^{k} y_j \prod_{\substack{0 \le m \le k \\ m \ne j}} \frac{x - x_m}{x_j - x_m}$
- **Hook length formula:** If λ is a Young diagram and $h_\lambda(i,j)$ is the hook-length of cell (i,j), then then the number of Young tableux $d_\lambda = n!/\prod h_\lambda(i,j)$.
- **Möbius inversion formula:** If $f(n) = \sum_{d|n} g(d)$, then $g(n) = \sum_{d|n} \mu(d) f(n/d)$. If $f(n) = \sum_{m=1}^{n} g(\lfloor n/m \rfloor)$, then $g(n) = \sum_{m=1}^{n} \mu(m) f(\lfloor \frac{n}{m} \rfloor)$.
- #primitive pythagorean triples with hypotenuse $< n$ approx $n/(2\pi)$.
- **Frobenius Number:** largest number which can't be expressed as a linear combination of numbers $a_1, \ldots, a_n$ with non-negative coefficients. $g(a_1, a_2) = a_1 a_2 - a_1 - a_2$, $N(a_1, a_2) = (a_1 - 1)(a_2 - 1)/2$. $g(d \cdot a_1, d \cdot a_2, a_3) = d \cdot g(a_1, a_2, a_3) + a_3(d-1)$. An integer $x > (\max_i a_i)^2$ can be expressed in such a way iff. $x \mid \gcd(a_1, \ldots, a_n)$.

11.1. Physics.

- **Snell's law:** $\frac{\sin\theta_1}{v_1} = \frac{\sin\theta_2}{v_2}$

11.2. Markov Chains.

A Markov Chain can be represented as a weighted directed graph of states, where the weight of an edge represents the probability of transitioning over that edge in one timestep. Let $P^{(m)} = (p_{ij}^{(m)})$ be the probability matrix of transitioning from state i to state j in m timesteps, and note that $P^{(1)}$ is the adjacency matrix of the graph. **Chapman-Kolmogorov:** $p_{ij}^{(m+n)} = \sum_k p_{ik}^{(m)} p_{kj}^{(n)}$. It follows that $P^{(m+n)} = P^{(m)} P^{(n)}$ and $P^{(m)} = P^m$. If $p^{(0)}$ is the initial probability distribution (a vector), then $p^{(0)} P^{(m)}$ is the probability distribution after m timesteps.

The return times of a state i is $R_i = \{m \mid p_{ii}^{(m)} > 0\}$, and i is *aperiodic* if $\gcd(R_i) = 1$. A MC is aperiodic if any of its vertices is aperiodic. A MC is *irreducible* if the corresponding graph is strongly connected.

A distribution π is stationary if $\pi P = \pi$. If MC is irreducible then $\pi_i = 1/\mathbb{E}[T_i]$, where T_i is the expected time between two visits at i. π_j/π_i is the expected number of visits at j in between two consecutive visits at i. A MC is *ergodic* if $\lim_{m\to\infty} p^{(0)} P^m = \pi$. A MC is ergodic iff. it is irreducible and aperiodic.

A MC for a random walk in an undirected weighted graph (unweighted graph can be made weighted by adding 1-weights) has $p_{uv} = w_{uv}/\sum_x w_{ux}$. If the graph is connected, then $\pi_u = \sum_x w_{ux}/\sum_v \sum_x w_{vx}$. Such a random walk is aperiodic iff. the graph is not bipartite.

An *absorbing* MC is of the form $P = \begin{pmatrix} Q & R \\ 0 & I_r \end{pmatrix}$. Let $N = \sum_{m=0}^{\infty} Q^m = (I_t - Q)^{-1}$. Then, if starting in state i, the expected number of steps till absorption is the i-th entry in $N1$. If starting in state i, the probability of being absorbed in state j is the (i,j)-th entry of NR.

Many problems on MC can be formulated in terms of a system of recurrence relations, and then solved using Gaussian elimination.

11.3. Burnside's Lemma.

Let G be a finite group that acts on a set X. For each g in G let X^g denote the set of elements in X that are fixed by g. Then the number of orbits

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$

$$Z(S_n) = \frac{1}{n} \sum_{l=1}^{n} a_l Z(S_{n-l})$$

11.4. Bézout's identity.

If (x, y) is any solution to $ax + by = d$ (e.g. found by the Extended Euclidean Algorithm), then all solutions are given by

$$\left(x + k\frac{b}{\gcd(a,b)}, y - k\frac{a}{\gcd(a,b)}\right)$$

11.5. Misc.

11.5.1. *Determinants and PM.*

$$det(A) = \sum_{\sigma \in S_n} \operatorname{sgn}(\sigma) \prod_{i=1}^{n} a_{i,\sigma(i)}$$

$$perm(A) = \sum_{\sigma \in S_n} \prod_{i=1}^{n} a_{i,\sigma(i)}$$

$$pf(A) = \frac{1}{2^n n!} \sum_{\sigma \in S_{2n}} \operatorname{sgn}(\sigma) \prod_{i=1}^{n} a_{\sigma(2i-1),\sigma(2i)}$$

$$= \sum_{M \in \mathrm{PM}(n)} \operatorname{sgn}(M) \prod_{(i,j) \in M} a_{i,j}$$

11.5.2. *BEST Theorem.*

Count directed Eulerian cycles. Number of OST given by Kirchoff's Theorem (remove r/c with root) $\#\mathrm{OST}(G,r) \cdot \prod_v (d_v - 1)!$

11.5.3. *Primitive Roots.*

Only exists when n is $2, 4, p^k, 2p^k$, where p odd prime. Assume n prime. Number of primitive roots $\phi(\phi(n))$ Let g be primitive root. All primitive roots are of the form g^k where $k, \phi(p)$ are coprime.
k-roots: $g^{i \cdot \phi(n)/k}$ for $0 \le i < k$

11.5.4. *Sum of primes.*

For any multiplicative f:

$$S(n,p) = S(n,p-1) - f(p) \cdot (S(n/p,p-1) - S(p-1,p-1))$$

11.5.5. *Floor.*

$$\lfloor \lfloor x/y \rfloor / z \rfloor = \lfloor x/(yz) \rfloor$$

$$x\%y = x - y \lfloor x/y \rfloor$$

Practice Contest Checklist

- How many operations per second? Compare to local machine.
- What is the stack size?
- How to use printf/scanf with long long/long double?
- Are `__int128` and `__float128` available?
- Does MLE give RTE or MLE as a verdict? What about stack overflow?
- What is `RAND_MAX`?
- How does the judge handle extra spaces (or missing newlines) in the output?
- Look at documentation for programming languages.
- Try different programming languages: C++, Java and Python.
- Try the submit script.
- Try local programs: i?python[23], factor.
- Try submitting with `assert(false)` and `assert(true)`.
- Return-value from `main`.
- Look for directory with sample test cases.
- Make sure printing works.
- Remove this page from the notebook.