

TCR

git merge -s octopus solution cup

Ludo Pulles, Reinier Schmiermann, Pim Spelier

CONTENTS

0.1. De winnende aanpak	2
0.2. Wrong Answer	2
1. Mathematics	2
1.1. Primitive Root $O(\sqrt{m})$	2
1.2. Tonelli-Shanks algorithm	3
1.3. Numeric Integration	3
1.4. Fast Hadamard Transform	3
1.5. Tridiagonal Matrix Algorithm	3
1.6. Number of Integer Points under Line	3
1.7. Solving linear recurrences	3
1.8. Misc	3
2. Datastructures	4
2.1. Order tree	4
2.2. Segment tree $O(\log n)$	4
2.3. Binary Indexed Tree $O(\log n)$	5
2.4. Disjoint-Set / Union-Find $O(\alpha(n))$	5
2.5. Cartesian tree	5
2.6. Heap	5
2.7. Dancing Links	6
2.8. Misof Tree	6
2.9. k -d Tree	6
2.10. Sqrt Decomposition	7
2.11. Monotonic Queue	7
2.12. Line container à la 'Convex Hull Trick' $O(n \log n)$	7
2.13. Sparse Table $O(\log n)$ per query	7
3. Graph Algorithms	7
3.1. Shortest path	7
3.2. Maximum Matching	8
3.3. Cycle Detection $O(V + E)$	9
3.4. Depth first searches	9
3.5. Min Cut / Max Flow	10
3.6. Minimal Spanning Tree $O(E \log V)$	10
3.7. Euler Path $O(V + E)$ hopefully	11
3.8. Heavy-Light Decomposition	11
3.9. Centroid Decomposition	11
3.10. Least Common Ancestors, Binary Jumping	11
3.11. Miscellaneous	12
4. String algorithms	13

4.1. Trie	13
4.2. Z-algorithm $O(n)$	13
4.3. Suffix array $O(n \log n)$	13
4.4. Longest Common Subsequence $O(n^2)$	14
4.5. Levenshtein Distance $O(n^2)$	14
4.6. Knuth-Morris-Pratt algorithm $O(N + M)$	14
4.7. Aho-Corasick Algorithm $O(N + \sum_{i=1}^m S_i)$	14
4.8. eerTree	14
4.9. Suffix Automaton	14
4.10. Hashing	14
5. Geometry	15
5.1. Convex Hull $O(n \log n)$	15
5.2. Rotating Calipers $O(n)$	15
5.3. Closest points $O(n \log n)$	16
5.4. Great-Circle Distance	16
5.5. Delaunay triangulation	16
5.6. 3D Primitives	16
5.7. Polygon Centroid	18
5.8. Rectilinear Minimum Spanning Tree	18
5.9. Points and lines (CP3)	18
5.10. Polygon (CP3)	18
5.11. Triangle (CP3)	18
5.12. Circle (CP3)	18
5.13. Formulas	18
6. Miscellaneous	21
6.1. Binary search $O(\log(hi - lo))$	21
6.2. Fast Fourier Transform $O(n \log n)$	21
6.3. Minimum Assignment (Hungarian Algorithm) $O(n^3)$	21
6.4. Partial linear equation solver $O(N^3)$	21
6.5. Cycle-Finding	21
6.6. Longest Increasing Subsequence	22
6.7. Dates	22
6.8. Simplex	22
7. Combinatorics	22
8. Formulas	23
9. Game Theory	23
10. Scheduling Theory	23
11. Debugging Tips	24
11.1. Dynamic programming optimizations	24
11.2. Solution Ideas	24
Practice Contest Checklist	25

```
.bashrc
alias gg='g++ -std=c++17 -Wall -Wconversion
↳ -Wshadow'
alias g='gg -DDEBUG -g -fsanitize=address,undefined'

.vimrc

set nu rnu sw=4 ts=4 sts=4 noet ai hls shcf=-ic
sy on | colo slate

Test script (usage: ./test.sh A/B/..)
g++ -g -Wall -fsanitize=address,undefined
↳ -Wfatal-errors -std=c++17 $1.cc || exit
for i in $(ls *.in)
do
    j="$(echo $i | sed 's/.in/.ans/')"
    ./a.out < $i > output
    diff output $j || echo "!!WA on $i!!"
done

template.cc
#include <bits/stdc++.h>
using namespace std;

typedef long long ll;
typedef long double ld;
typedef pair<ll, ll> ii;
typedef vector<ll> vi;
typedef vector<vi> vvi;
typedef vector<ii> vii;

#define x first
#define y second
#define pb push_back
#define eb emplace_back
#define rep(i,a,b) for(auto i=(a); i<(b); ++i)
#define REP(i,n) rep(i,0,n)
#define all(v) (v).begin(), (v).end()
#define sz(v) ((int) (v).size())
#define rs resize
#define DBG(x) cerr << __LINE__ << ": " << #x << " = " << (x) << endl

const ld PI = acos(-1.0);

template<class T> ostream& operator<<(ostream &os,
↳ vector<T> v) {
    os << "\n[";
    for(T &x : v) os << x << ', ';
    return os << "]\n";
}

struct pairhash {
public:
    template<typename T1, typename T2>
    size_t operator() (const pair<T1, T2> &p) const {
        size_t lhs = hash<T1>() (p.x);
        size_t rhs = hash<T2>() (p.y);
        return lhs ^ (rhs<<9e3779b9+(lhs<<6)+(lhs>>2));
    }
};
```

```

    }
};

void run() {

}

signed main() {
    // DON'T MIX "scanf" and "cin"!
    ios_base::sync_with_stdio(false);
    cin.tie(NULL);
    cout << fixed << setprecision(20);
    run();
    return 0;
}

template.py

# reading input:
from sys import *
n,m = [ int(x) for x in
    ↪ stdin.readline().rstrip().split() ]
stdout.write( str(n*m)+"\n" )
# set operations:
from itertools import *
for (x,y) in product(range(3),repeat=2):
    stdout.write( str(3*x+y)+" " )
print()
for L in combinations(range(4),2):
    stdout.write( str(L)+" " )
print()
# fancy lambdas:
from functools import *
y = reduce( lambda x,y: x+y, map( lambda x: x*x,
    ↪ range(4) ), -3 )
print(y)
# formatting:
from math import *
stdout.write( "{0:.2f}\n".format(pi) )

```

0.1. De winnende aanpak.

- Slaap goed & heb een vroeg ritme!
- Drink & eet genoeg voor & tijdens de wedstrijd!
- Houd een lijst bij met info over alle problemen.
- Iedereen (incl. Ludo) moet **ALLE** opgaves goed lezen!
- Analyseer de voorbeeld test cases.
- Houd na 2 uur een pauze en overleg waar iedereen mee bezig is.
- Maak zelf (zware) test cases.
- Gebruik ll.

0.2. Wrong Answer.

- Print de oplossing om te debuggen!
- Kijk naar wellicht makkelijkere problemen.
- Bedenk zelf test cases met **randgevallen**!
- Controleer de **precisie**.

- Controleer op **overflow** (gebruik **OVERAL** ll, ld).
Kijk naar overflows in tussenantwoorden bij modulo.
- Controleer op **typo's**.
- Loop de voorbeeld test case accuraat langs.
- Controleer op off-by-one-errors (in indices of lus-grenzen)?

Detecting overflow: This GNU builtin checks for over- and underflow. Result is in res if successful:

```

bool isOverflown =
    ↪ __builtin_[add|mul|sub]_overflow(a, b, &res);

```

1. MATHEMATICS

XOR sum: $\bigoplus_{x=0}^{a-1} x = \{0, a-1, 1, a\}[a \bmod 4]$.

```

int abs(int x) { return x > 0 ? x : -x; }
int sign(int x) { return (x > 0) - (x < 0); }

```

```

// greatest common divisor
ll gcd(ll a, ll b) { while(b) a%=b, swap(a,b); return a; }
// least common multiple
ll lcm(ll a, ll b) { return a/gcd(a, b)*b; }
ll mod(ll a, ll b) { return (a%b) < 0 ? a+b : a; }

```

```

// ab % m for m <= 4e18 in O(log b)
ll mod_mul(ll a, ll b, ll m) {
    ll r = 0;
    while(b) {
        if (b & 1) r = mod(r+a,m);
        a = mod(a+a,m); b >>= 1;
    }
    return r;
}

```

```

// a^b % m for m <= 2e9 in O(log b)
ll mod_pow(ll a, ll b, ll m) {
    ll r = 1;
    while(b) {
        if (b & 1) r = (r * a) % m; // mod_mul
        a = (a * a) % m; // mod_mul
        b >>= 1;
    }
    return r;
}

```

```

// returns x, y such that ax + by = gcd(a, b)
ll egcd(ll a, ll b, ll &x, ll &y) {
    ll xx = y = 0, yy = x = 1;
    while (b) {
        x -= a / b * xx; swap(x, xx);
        y -= a / b * yy; swap(y, yy);
        a %= b; swap(a, b);
    }
    return a;
}

```

```

// Chinese Remainder Theorem: returns (u, v) s.t.
// x=u (mod v) <=> x=a (mod n) and x=b (mod m)

```

```

pair<ll, ll> crt(ll a, ll n, ll b, ll m) {
    ll s, t, d = egcd(n, m, s, t); //n,m<=1e9
    if (mod(a - b, d)) return { 0, -1 };
    return { mod(s*b*m*n + t*a*n*m, n*m)/d, n*m/d };
}

```

```

// phi[i] = #{ 0 < j <= i | gcd(i, j) = 1 } sieve
vi totient(int N) {
    vi phi(N);
    for (int i = 0; i < N; i++) phi[i] = i;
    for (int i = 2; i < N; i++) if (phi[i] == i)
        for (int j = i; j < N; j+=i) phi[j] -= phi[j]/i;
    return phi;
}

```

```

// calculate nCk % p (p prime!)
ll lucas(ll n, ll k, ll p) {
    ll ans = 1;
    while (n) {
        ll np = n % p, kp = k % p;
        if (np < kp) return 0;
        ans = mod(ans * binom(np, kp), p); // (np C kp)
        n /= p; k /= p;
    }
    return ans;
}

```

```

// returns if n is prime for n < 3e24 (>2^64)
// but use mul_mod for n > 2e9.
bool millerRabin(ll n) {
    if (n < 2 || n % 2 == 0) return n == 2;
    ll d = n - 1, ad, s = 0, r;
    for (; d % 2 == 0; d /= 2) s++;
    for (int a : { 2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37, 41 }) {
        if (n == a) return true;
        if ((ad = mod_pow(a, d, n)) == 1) continue;
        for (r = 0; r < s && ad + 1 != n; r++)
            ad = (ad * ad) % n;
        if (r == s) return false;
    }
    return true;
}

```

1.1. Primitive Root $O(\sqrt{m})$. Returns a generator of \mathbb{F}_m^* . If m not prime, replace $m-1$ by totient of m .

```

ll primitive_root(ll m) {
    vector<ll> div;
    for (ll i = 1; i*i < m; i++)
        if ((m-1) % i == 0) {
            if (i < m-1) div.pb(i);
            if ((m-1)/i < m) div.pb((m-1)/i);
        }
    rep(x, 2, m) {
        bool ok = true;
        for (ll d : div) if (mod_pow(x, d, m) == 1)

```

```

    { ok = false; break; }
    if (ok) return x;
}
return -1;
}

```

1.2. Tonelli-Shanks algorithm. Given prime p and integer $1 \leq n < p$, returns the square root r of n modulo p . There is also another solution given by $-r$ modulo p .

```

ll legendre(ll a, ll p) {
    if (a % p == 0) return 0;
    return p == 2 || mod_pow(a, (p-1)/2, p) == 1 ? 1 :
        -1;
}

ll tonelli_shanks(ll n, ll p) {
    assert(legendre(n,p) == 1);
    if (p == 2) return 1;
    ll s = 0, q = p-1, z = 2;
    while (~q & 1) s++, q >>= 1;
    if (s == 1) return mod_pow(n, (p+1)/4, p);
    while (legendre(z,p) != -1) z++;
    ll c = mod_pow(z, q, p),
        r = mod_pow(n, (q+1)/2, p),
        t = mod_pow(n, q, p),
        m = s;
    while (t != 1) {
        ll i = 1, ts = (ll)t*t % p;
        while (ts != 1) i++, ts = ((ll)ts * ts) % p;
        ll b = mod_pow(c, 1LL<<(m-i-1), p);
        r = (ll)r * b % p;
        t = (ll)t * b % p * b % p;
        c = (ll)b * b % p;
        m = i;
    }
    return r;
}

```

1.3. Numeric Integration. Numeric integration using Simpson's rule.

```

ld numint(ld (*f)(ld), ld a, ld b, ld EPS = 1e-6) {
    ld ba = b - a, m=(a+b)/2;
    return abs(ba) < EPS
        ? ba/8*(f(a)+f(b)+f(a+ba/3)*3+f(b-ba/3)*3)
        : numint(f,a,m,EPS) + numint(f,m,b,EPS);
}

```

1.4. Fast Hadamard Transform. Computes XOR-convolutions in $O(k2^k)$ on k bits.

For AND-convolution, use $(x+y, y)$, $(x-y, y)$.

For OR-convolution, use $(x, x+y)$, $(x, -x+y)$.

Note: The array size must be a power of 2.

```

void fht(vi &A, bool inv=false, int l, int r) {
    if (l+1 == r) return;
    int k = (r-l)/2;
    if (!inv) fht(A, inv, l, l+k), fht(A, inv, l+k,
        -> r);
}

```

```

rep(i,l,l+k) {
    int x = A[i], y = A[i+k];
    if (!inv) A[i] = x-y, A[i+k] = x+y;
    else A[i] = (x+y)/2, A[i+k] = (-x+y)/2;
}
if (inv) fht(A, inv, l, l+k), fht(A, inv, l+k, r);
}

```

1.5. Tridiagonal Matrix Algorithm. Solves a tridiagonal system of linear equations

$$a_i x_{i-1} + b_i x_i + c_i x_{i+1} = d_i$$

where $a_1 = c_n = 0$. Beware of numerical instability.

```

#define MAXN 5000
ld A[MAXN], B[MAXN], C[MAXN], D[MAXN], X[MAXN];
void solve(int n) {
    C[0] /= B[0]; D[0] /= B[0];
    rep(i,1,n-1) C[i] /= B[i] - A[i]*C[i-1];
    rep(i,1,n) D[i] =
        (D[i] - A[i]*D[i-1]) / (B[i] - A[i]*C[i-1]);
    X[n-1] = D[n-1];
    for (int i = n-1; i--;) X[i] = D[i] - C[i]*X[i+1];
}

```

1.6. Number of Integer Points under Line. Count the number of integer solutions to $Ax + By \leq C$, $0 \leq x \leq n$, $0 \leq y$. In other words, evaluate the sum $\sum_{x=0}^n \lfloor \frac{C-Ax}{B} + 1 \rfloor$. To count all solutions, let $n = \lfloor \frac{C}{A} \rfloor$. In any case, it must hold that $C - nA \geq 0$. Be very careful about overflows.

```

ll floor_sum(ll n, ll a, ll b, ll c) {
    if (c == 0) return 1;
    if (c < 0) return 0;
    if (a % b == 0) return
        -> (n+1)*(c/b+1)-n*(n+1)/2*a/b;
    if (a >= b) return
        -> floor_sum(n,a%b,b,c)-a/b*n*(n+1)/2;
    ll t = (c-a*n+b)/b;
    return floor_sum((c-b*t)/b,b,a,c-b*t)+t*(n+1);
}

```

1.7. Solving linear recurrences. Given some brute-forced sequence $s[0], s[1], \dots, s[2n-1]$, Berlekamp-Massey finds the shortest possible recurrence relation in $\mathcal{O}(n^2)$. After that, `lin_rec` finds $s[k]$ in $\mathcal{O}(n^2 \log k)$.

```

// Given a sequence s[0], ..., s[2n-1] finds the
// smallest linear recurrence
// of size <= n compatible with s.
vl BerlekampMassey(const vl &s, ll mod) {
    int n = sz(s), L = 0, m = 0;
    vl C(n), B(n), T;
    C[0] = B[0] = 1;
    ll b = 1;
    REP(i, n) {
        ++m;
        ll d = s[i] % mod;
        rep(j, 1, L+1) d = (d + C[j] * s[i-j]) % mod;
    }
}

```

```

if (!d) continue;
T = C;
ll coef = d * modpow(b, mod-2, mod) % mod;
rep(j,m,n) C[j] = (C[j] - coef * B[j-m]) % mod;
if (2 * L > i) continue;
L = i + 1 - L;
B = T; b = d; m = 0;
}
C.resize(L + 1);
C.erase(C.begin());
for (auto &x : C) x = (mod - x) % mod;
return C;
}

```

```

// Input: A[0,...,n-1], C[0,...,n-1] satisfying
// A[i] = \sum_{j=1}^n C[j-1] A[i-j],
// Outputs A[k]
ll lin_rec(const vl &A, const vl &C, ll k, ll mod){
    int n = sz(A);
    auto combine = [&](vl a, vl b) {
        vl res(sz(a) + sz(b) - 1, 0);
        REP(i, sz(a)) REP(j, sz(b))
            res[i+j] = (res[i+j] + a[i]*b[j]) % mod;
        for (int i = 2*n; i > n; --i) REP(j,n)
            res[i-1-j] = (res[i-1-j] + res[i]*C[j]) % mod;
        res.resize(n + 1);
        return res;
    };
    vl pol(n + 1, e(pol));
    pol[0] = e[1] = 1;
    for (++k; k; k /= 2) {
        if (k % 2) pol = combine(pol, e);
        e = combine(e, e);
    }
    ll res = 0;
    REP(i, n) res = (res + pol[i + 1] * A[i]) % mod;
    return res;
}

```

1.8. Misc.

1.8.1. Josephus problem. Last man standing out of n if every k th is killed. Zero-based, and does not kill 0 on first pass.

```

int J(int n, int k) {
    if (n == 1 || k == 1) return n-1;
    if (n < k) return (J(n-1,k)+k)%n;
    int np = n - n/k;
    return k*((J(np,k)+np-n%k*np)%np) / (k-1);
}

```

• **Prime numbers:**

1031, 32771, 1048583, 8125344, 33554467, 9982451653, 1073741827, 34359738421, 1099511627791, 35184372088891, 1125899906842679, 36028797018963971.
 $10^3 + \{-9, -3, 9, 13\}$, $10^6 + \{-17, 3, 33\}$, $10^9 + \{7, 9, 21, 33, 87\}$.

- **Generating functions:** Ordinary (ogf): $A(x) := \sum_{n=0}^{\infty} a_n x^n$.
Calculate product $c_n = \sum_{k=0}^n a_k b_{n-k}$ with FFT.
Exponential (e.g.f.): $A(x) := \sum_{n=0}^{\infty} a_n x^n / n!$,
 $c_n = \sum_{k=0}^n \binom{n}{k} a_k b_{n-k} = n! \sum_{k=0}^n \frac{a_k}{k!} \frac{b_{n-k}}{(n-k)!}$ (use FFT).
- **General linear recurrences:** If $a_n = \sum_{k=0}^{n-1} a_k b_{n-k}$, then $A(x) = \frac{a_0}{1-B(x)}$.
- **Inverse polynomial modulo x^l :** Given $A(x)$, find $B(x)$ such that $A(x)B(x) = 1 + x^l Q(x)$ for some $Q(x)$.
Step 1: Start with $B_0(x) = 1/a_0$
Step 2: $B_{k+1}(x) = (-B_k(x)^2 A(x) + 2B_k(x)) \bmod x^{2^{k+1}}$.
- **Fast subset convolution:** Given array a_i of size 2^k calculate $b_i = \sum_{j \& i = i} a_j$.

```
for (int b = 1; b < (1 << k); b <= 1)
  for (int i = 0; i < (1 << k); i++)
    if (!(i & b)) a[i | b] += a[i];
// inv: if (!(i & b)) a[i | b] -= a[i];
```
- **Primitive Roots:** It only exists when n is $2, 4, p^k, 2p^k$, where p odd prime. If g is a primitive root, all primitive roots are of the form g^k where $k, \phi(p)$ are coprime (hence there are $\phi(\phi(p))$ primitive roots).
- **Maximum number of divisors:**

$\leq N$	10^3	10^6	10^9	10^{12}	10^{18}
m	840	720720	735134400	963761198400	
$\sigma_0(m)$	32	240	1344	6270	103680

For $n = 10^{18}, m = 897612484786617600$.

2. DATASTRUCTURES

2.1. Order tree.

```
#include <bits/stdc++.h>
using namespace __gnu_pbds;
template<class TK, class TM> using order_tree =
  tree<TK, TM, greater<TK>, rb_tree_tag,
  tree_order_statistics_node_update>;
template<class TK> using order_set =
  order_tree<TK, null_type>;
```

```
vi s;
order_set<ii> t;
void update( ll k, ll v ) {
  t.erase( ii{ s[k], k } );
  s[k] = v;
  t.insert( ii{ s[k], k } );
}
```

```
signed main() {
  ll n = 4;
  s.resize(n, 0);
  rep(i, 0, n) t.insert(ii{0, i});
}
```

```
update( 2, 3 );
cout << t.find_by_order( 2 )->y << endl;
cout << t.order_of_key( ii{s[3], 3} ) << endl;
}
```

2.2. Segment tree $\mathcal{O}(\log n)$.

2.2.1. Standard segment tree.

```
typedef int S; // or define your own object
const int n = 1 << 20;
S t[2 * n];

// combine must be an associative function!
S combine(S l, S r) { return l+r; } //or max(l,r) etc

void build() {
  for (int i = n; --i; )
    t[i] = combine(t[2 * i], t[2 * i + 1]);
}

// set value v on position i
void update(int i, S v) {
  for (t[i+=n] = v; i /= 2; )
    t[i] = combine(t[2 * i], t[2 * i + 1]);
}

// sum on interval [l, r)
S query(int l, int r) {
  S resL = 0, resR = 0;
  for (l += n, r += n; l < r; l /= 2, r /= 2) {
    if (l & 1) resL = combine(resL, t[l++]);
    if (r & 1) resR = combine(t[--r], resR);
  }
  return combine(resL, resR);
}
```

2.2.2. Lazy segment tree.

Be careful: all intervals are right-closed $[\ell, r]$.

```
struct node {
  int l, r, x, lazy;
  node() {}
  node(int _l, int _r) : l(_l), r(_r), x(INT_MAX),
    lazy(0) {}
  node(int _l, int _r, int _x) : node(_l, _r) {x=_x;}
  node(node a, node b) : node(a.l, b.r) {x=min(a.x, b.x);}
  void update(int v) { x = v; }
  void range_update(int v) { lazy = v; }
  void apply() { x += lazy; lazy = 0; }
  void push(node &u) { u.lazy += lazy; }
};
```

```
struct segment_tree {
  int n;
  vector<node> arr;
  segment_tree(int n) {
    segment_tree(const vi &a) : n(sz(a)), arr(4*n) {
      mk(a, 0, 0, n-1);
    }
}
```

```
node mk(const vi &a, int i, int l, int r) {
  int m = (l+r)/2;
  return arr[i] = l > r ? node(l, r) :
    l == r ? node(l, r, a[l]) :
      node(mk(a, 2*i+1, l, m), mk(a, 2*i+2, m+1, r));
}
node update(int at, ll v, int i=0) {
  propagate(i);
  int hl = arr[i].l, hr = arr[i].r;
  if (at < hl || hr < at) return arr[i];
  if (hl == at && at == hr) {
    arr[i].update(v); return arr[i];
  }
  return arr[i] =
    node(update(at, v, 2*i+1), update(at, v, 2*i+2));
}
node query(int l, int r, int i=0) {
  propagate(i);
  int hl = arr[i].l, hr = arr[i].r;
  if (r < hl || hr < l) return node(hl, hr);
  if (l <= hl && hr <= r) return arr[i];
  return node(query(l, r, 2*i+1), query(l, r, 2*i+2));
}
node range_update(int l, int r, ll v, int i=0) {
  propagate(i);
  int hl = arr[i].l, hr = arr[i].r;
  if (r < hl || hr < l) return arr[i];
  if (l <= hl && hr <= r) {
    arr[i].range_update(v);
    propagate(i);
    return arr[i];
  }
  return arr[i] = node(range_update(l, r, v, 2*i+1),
    range_update(l, r, v, 2*i+2));
}
void propagate(int i) {
  if (arr[i].l < arr[i].r) {
    arr[i].push(arr[2*i+1]);
    arr[i].push(arr[2*i+2]);
  }
  arr[i].apply();
}
};
```

2.2.3. Persistent segment tree.

Be careful: all intervals are right-closed $[\ell, r]$, including build.

```
int segcnt = 0;
struct segment {
  int l, r, lid, rid, sum;
} S[2000000];
```

```
int build(int l, int r) {
  if (l > r) return -1;
  int id = segcnt++;
  S[id].l = l;
  S[id].r = r;
  if (l == r) S[id].lid = -1, S[id].rid = -1;
```

```

else {
    int m = (l + r) / 2;
    S[id].lid = build(l, m);
    S[id].rid = build(m + 1, r);
}
S[id].sum = 0;
return id;
}
int update(int idx, int v, int id) {
    if (id == -1) return -1;
    if (idx < S[id].l || idx > S[id].r) return id;
    int nid = segcnt++;
    S[nid].l = S[id].l;
    S[nid].r = S[id].r;
    S[nid].lid = update(idx, v, S[id].lid);
    S[nid].rid = update(idx, v, S[id].rid);
    S[nid].sum = S[id].sum + v;
    return nid;
}
int query(int id, int l, int r) {
    if (r < S[id].l || S[id].r < l) return 0;
    if (l <= S[id].l && S[id].r <= r) return S[id].sum;
    return query(S[id].lid, l, r) + query(S[id].rid, l, r);
}

```

2.3. Binary Indexed Tree $\mathcal{O}(\log n)$. Use one-based indices ($i > 0$)!

```

struct BIT {
    int n; vi A;
    BIT(int _n) : n(_n), A(_n + 1, 0) {}
    BIT(vi &v) : n(sz(v)), A(1) {
        for (auto x:v) A.pb(x);
        for (int i=1, j; j=i&-i, i<=n; i++)
            if (i+j <= n) A[i+j] += A[i];
    }
    void update(int i, ll v) { // a[i] += v
        while (i <= n) A[i] += v, i += i&-i;
    }
    ll query(int i) { // sum_{j<=i} a[j]
        ll v = 0;
        while (i) v += A[i], i -= i&-i;
        return v;
    }
};

struct rangeBIT {
    int n; BIT b1, b2;
    rangeBIT(int _n) : n(_n), b1(_n), b2(_n + 1) {}
    rangeBIT(vi &v) : n(sz(v)), b1(v), b2(sz(v) + 1) {}
    void pupdate(int i, ll v) { b1.update(i, v); }
    void rupdate(int i, int j, ll v) { // a[i,...,j] += v
        b2.update(i, v);
        b2.update(j+1, -v);
        b1.update(j+1, v*j);
        b1.update(i, (1-i)*v);
    }
    ll query(int i) { return b1.query(i) + b2.query(i)*i; }
};

```

2.4. Disjoint-Set / Union-Find $\mathcal{O}(\alpha(n))$.

```

struct dsu {
    vi par, rnk;
    dsu(int n) : par(n, -1), rnk(n, 0) {}
    int find(int i) { return
        par[i] < 0 ? i : par[i] = find(par[i]); }
    void unite(int a, int b) {
        if ((a = find(a)) == (b = find(b))) return;
        if (rnk[a] < rnk[b]) swap(a, b);
        if (rnk[a] == rnk[b]) rnk[a]++;
        par[a] += par[b]; par[b] = a;
    }
};

```

2.5. Cartesian tree.

```

struct node {
    int x, y, sz;
    node *l, *r;
    node(int _x, int _y)
        : x(_x), y(_y), sz(1), l(NULL), r(NULL) {}
};
int tsize(node* t) { return t ? t->sz : 0; }
void augment(node* t) {
    t->sz = 1 + tsize(t->l) + tsize(t->r);
}
pair<node*, node*> split(node* t, int x) {
    if (!t) return make_pair((node*)NULL, (node*)NULL);
    if (t->x < x) {
        pair<node*, node*> res = split(t->r, x);
        t->r = res.x; augment(t);
        return make_pair(t, res.y);
    }
    pair<node*, node*> res = split(t->l, x);
    t->l = res.y; augment(t);
    return make_pair(res.x, t);
}
node* merge(node* l, node* r) {
    if (!l) return r; if (!r) return l;
    if (l->y > r->y) {
        l->r = merge(l->r, r); augment(l); return l;
    }
    r->l = merge(l, r->l); augment(r); return r;
}
node* find(node* t, int x) {
    while (t) {
        if (x < t->x) t = t->l;
        else if (t->x < x) t = t->r;
        else return t;
    }
    return NULL;
}
node* insert(node* t, int x, int y) {
    if (find(t, x) != NULL) return t;
    pair<node*, node*> res = split(t, x);
    return merge(res.x, merge(new node(x, y), res.y));
}
node* erase(node* t, int x) {
    if (!t) return NULL;
    if (t->x < x) t->r = erase(t->r, x);
    else if (x < t->x) t->l = erase(t->l, x);
    else {node* old = t; t = merge(t->l, t->r); delete old;
    }
    if (t) augment(t); return t;
}
int kth(node* t, int k) {
    if (k < tsize(t->l)) return kth(t->l, k);
    else if (k == tsize(t->l)) return t->x;
}

```

```

else return kth(t->r, k - tsize(t->l) - 1);
}

```

2.6. Heap. An implementation of a binary heap.

```

#define RESIZE
#define SWP(x,y) tmp = x, x = y, y = tmp
struct int_less {
    int_less() {}
    bool operator () (const int &a, const int &b) {
        return a < b;
    }
};
template <class Compare = int_less> struct heap {
    int cap, len, *q, *loc, tmp;
    Compare _cmp;
    inline bool cmp(int i, int j) {
        return _cmp(q[i], q[j]);
    }
    inline void swp(int i, int j) {
        SWP(q[i], q[j]), SWP(loc[q[i]], loc[q[j]]);
    }
    void swim(int i) {
        while (i > 0) {
            int p = (i - 1) / 2;
            if (!cmp(i, p)) break;
            swp(i, p), i = p;
        }
    }
    void sink(int i) {
        while (true) {
            int l = 2*i + 1, r = l + 1;
            if (l >= len) break;
            int m = r >= len || cmp(l, r) ? l : r;
            if (!cmp(m, i)) break;
            swp(m, i), i = m;
        }
    }
    heap(int C=128) : len(0), cap(C), _cmp(Compare())
    {
        q = new int[C]; loc = new int[C];
        memset(loc, 255, cap << 2);
    }
    ~heap() {
        delete[] q; delete[] loc;
    }
    void push(int n, bool fix = true) {
        if (cap == len || n >= cap) {
#ifdef RESIZE
            int newcap = 2 * cap;
            while (n >= newcap) newcap *= 2;
            int *newq = new int[newcap], *newloc = new
                int[newcap];
            REP(i, cap) newq[i] = q[i], newloc[i] = loc[i];
            memset(newloc + cap, 255, (newcap - cap) << 2);
            delete[] q, delete[] loc;

```



```

    loc = newloc, q = newq, cap = newcap;
#else
    assert(false);
#endif
}
assert(loc[n] == -1);
loc[n] = len, q[len++] = n;
if (fix) swim(len-1);
}
void pop(bool fix = true) {
    assert(len > 0);
    loc[q[0]] = -1, q[0] = q[--len], loc[q[0]] = 0;
    if (fix) sink(0);
}
int top() { assert(len > 0); return q[0]; }
void heapify() {
    for (int i = len - 1; i > 0; i--)
        if (cmp(i, (i-1)/2)) swp(i, (i-1)/2);
}
void update_key(int n) {
    assert(loc[n] != -1); swim(loc[n]); sink(loc[n]);
}
bool empty() { return len == 0; }
int size() { return len; }
void clear() {
    len = 0; memset(loc, 255, cap << 2);
}
};

```

2.7. Dancing Links. An implementation of Donald Knuth's Dancing Links data structure. A linked list supporting deletion and restoration of elements.

```

template <class T>
struct dancing_links {
    struct node {
        T item;
        node *l, *r;
        node(const T &_item, node *_l=NULL, node *_r=NULL)
            : item(_item), l(_l), r(_r) {
            if (!l) l->r = this;
            if (!r) r->l = this; } };
    node *front, *back;
    dancing_links() { front = back = NULL; }
    node *push_back(const T &item) {
        back = new node(item, back, NULL);
        if (!front) front = back;
        return back; }
    node *push_front(const T &item) {
        front = new node(item, NULL, front);
        if (!back) back = front;
        return front; }
    void erase(node *n) {
        if (!n->l) front = n->r; else n->l->r = n->r;
        if (!n->r) back = n->l; else n->r->l = n->l; }
    void restore(node *n) {
        if (!n->l) front = n; else n->l->r = n;
        if (!n->r) back = n; else n->r->l = n; } };

```

2.8. Misof Tree. A simple tree data structure for inserting, erasing, and querying the n th largest element.

```

const int BITS = 15;
struct misof_tree {
    int cnt[BITS][1<<BITS];
    misof_tree() { memset(cnt, 0, sizeof(cnt)); }
    void insert(int x) {
        for (int i=0; i<BITS; cnt[i++][x]++, x >= 1); }
    void erase(int x) {
        for (int i=0; i<BITS; cnt[i++][x]--, x >= 1); }
    int nth(int n) {
        int res = 0;
        for (int i = BITS-1; i >= 0; i--)
            if (cnt[i][res <= 1] <= n)
                n -= cnt[i][res], res |= 1;
        return res;
    }
};

```

2.9. k -d Tree. A k -dimensional tree supporting fast construction, adding points, and nearest neighbor queries. NOTE: Not completely stable, occasionally segfaults.

```

#define INC(c) ((c) == K - 1 ? 0 : (c) + 1)
template <int K> struct kd_tree {
    struct pt {
        double coord[K];
        pt() {}
        pt(double c[K]) { REP(i,K) coord[i] = c[i]; }
        double dist(const pt &other) const {
            double sum = 0.0;
            REP(i,K) sum +=
                pow(coord[i]-other.coord[i], 2);
            return sqrt(sum); } };
    struct cmp {
        int c;
        cmp(int _c) : c(_c) {}
        bool operator()(const pt &a, const pt &b) {
            for (int i = 0, cc; i <= K; i++) {
                cc = i == 0 ? c : i - 1;
                if (abs(a.coord[cc] - b.coord[cc]) > EPS)
                    return a.coord[cc] < b.coord[cc];
            }
            return false; } };
    struct bb {
        pt from, to;
        bb(pt _from, pt _to) : from(_from), to(_to) {}
        double dist(const pt &p) {
            double sum = 0.0;
            REP(i,K) {
                if (p.coord[i] < from.coord[i])
                    sum += pow(from.coord[i] - p.coord[i],
                        2.0);
                else if (p.coord[i] > to.coord[i])
                    sum += pow(p.coord[i] - to.coord[i], 2.0);
            }
            return sqrt(sum); }
    bb bound(double l, int c, bool left) {

```

```

        pt nf(from.coord), nt(to.coord);
        if (left) nt.coord[c] = min(nt.coord[c], l);
        else nf.coord[c] = max(nf.coord[c], l);
        return bb(nf, nt); } };
    struct node {
        pt p; node *l, *r;
        node(pt _p, node *_l, node *_r)
            : p(_p), l(_l), r(_r) {} };
    node *root;

    // kd_tree() : root(NULL) {}
    kd_tree(vector<pt> pts) {
        root = construct(pts, 0, size(pts) - 1, 0); }
    node* construct(vector<pt> &pts, int fr, int to,
        int c) {
        if (fr > to) return NULL;
        int mid = fr + (to-fr) / 2;
        nth_element(pts.begin() + fr, pts.begin() + mid,
            pts.begin() + to + 1, cmp(c));
        return new node(pts[mid],
            construct(pts, fr, mid - 1, INC(c)),
            construct(pts, mid + 1, to, INC(c))); }
    bool contains(const pt &p) { return
        _con(p, root, 0); }
    bool _con(const pt &p, node *n, int c) {
        if (!n) return false;
        if (cmp(c)(p, n->p)) return _con(p, n->l, INC(c));
        if (cmp(c)(n->p, p)) return _con(p, n->r, INC(c));
        return true; }
    void insert(const pt &p) { _ins(p, root, 0); }
    void _ins(const pt &p, node* &n, int c) {
        if (!n) n = new node(p, NULL, NULL);
        else if (cmp(c)(p, n->p)) _ins(p, n->l, INC(c));
        else if (cmp(c)(n->p, p)) _ins(p, n->r, INC(c));
        }
    void clear() { _clr(root); root = NULL; }
    void _clr(node *n) {
        if (n) _clr(n->l), _clr(n->r), delete n; }
    pt nearest_neighbour(const pt &p, bool same=true)
        {
            assert(root);
            double mn = INFINITY, cs[K];
            REP(i,K) cs[i] = -INFINITY;
            pt from(cs);
            REP(i,K) cs[i] = INFINITY;
            pt to(cs);
            return _nn(p, root, bb(from, to), mn, 0,
                same).x;
        }
    pair<pt, bool> _nn(const pt &p, node *n, bb b,
        double &mn, int c, bool same) {
        if (!n || b.dist(p) > mn)
            return make_pair(pt(), false);
        bool found = same || p.dist(n->p) > EPS,
            l1 = true, l2 = false;
        pt resp = n->p;

```

```

if (found) mn = min(mn, p.dist(resp));
node *n1 = n->l, *n2 = n->r;
REP(i,2) {
    if (i == 1 || cmp(c)(n->p, p))
        swap(n1, n2), swap(l1, l2);
    auto res = _nn(p, n1, b.bound(n->p.coord[c],
        ↪ c, l1), mn, INC(c), same);
    if (res.y && (!found || p.dist(res.x) <
        ↪ p.dist(resp)))
        resp = res.x, found = true;
}
return make_pair(resp, found); } };
```

2.10. Sqrt Decomposition. Design principle that supports many operations in amortized \sqrt{n} per operation.

```

struct segment {
    vi arr;
    segment(vi _arr) : arr(_arr) { } };
vector<segment> T;
int K;
void rebuild() {
    int cnt = 0;
    rep(i,0,size(T))
        cnt += size(T[i].arr);
    K = static_cast<int>(ceil(sqrt(cnt)) + 1e-9);
    vi arr(cnt);
    for (int i = 0, at = 0; i < size(T); i++)
        rep(j,0,size(T[i].arr))
            arr[at++] = T[i].arr[j];
    T.clear();
    for (int i = 0; i < cnt; i += K)
        T.push_back(segment(vi(arr.begin()+i,
            ↪ arr.begin()+min(i+K,
            ↪ cnt)))); }

int split(int at) {
    int i = 0;
    while (i < size(T) && at >= size(T[i].arr))
        at -= size(T[i].arr), i++;
    if (i >= size(T)) return size(T);
    if (at == 0) return i;
    T.insert(T.begin() + i + 1,
        segment(vi(T[i].arr.begin() + at,
            ↪ T[i].arr.end())));
    T[i] = segment(vi(T[i].arr.begin(),
        ↪ T[i].arr.begin() + at));
    return i + 1; }

void insert(int at, int v) {
    vi arr; arr.push_back(v);
    T.insert(T.begin() + split(at), segment(arr)); }

void erase(int at) {
    int i = split(at); split(at + 1);
    T.erase(T.begin() + i); }
```

2.11. Monotonic Queue. A queue that supports querying for the minimum element. Useful for sliding window algorithms.

```

struct min_stack {
    stack<int> S, M;
```

```

void push(int x) {
    S.push(x);
    M.push(M.empty() ? x : min(M.top(), x)); }

int top() { return S.top(); }
int mn() { return M.top(); }
void pop() { S.pop(); M.pop(); }
bool empty() { return S.empty(); } };

struct min_queue {
    min_stack inp, outp;
    void push(int x) { inp.push(x); }
    void fix() {
        if (outp.empty()) while (!inp.empty())
            outp.push(inp.top(), inp.pop()); }
    int top() { fix(); return outp.top(); }
    int mn() {
        if (inp.empty()) return outp.mn();
        if (outp.empty()) return inp.mn();
        return min(inp.mn(), outp.mn()); }
    void pop() { fix(); outp.pop(); }
    bool empty() { return inp.empty() && outp.empty(); } };
```

2.12. Line container à la ‘Convex Hull Trick’ $\mathcal{O}(n \log n)$. Container where you can add lines of the form $y_i(x) = k_i x + m_i$ and query $\max_i y_i(x)$.

```

bool Q;
struct Line {
    mutable ll k, m, p;
    bool operator<(const Line& o) const {
        return Q ? p < o.p : k < o.k; }
};

struct LineContainer : multiset<Line> {
    // (for doubles, use inf = 1/.0, div(a,b) = a/b)
    const ll inf = LLONG_MAX;
    ll div(ll a, ll b) { // floored division
        return a / b - ((a ^ b) < 0 && a % b); }
    bool isect(iterator x, iterator y) {
        if (y == end()) { x->p = inf; return false; }
        if (x->k == y->k)
            x->p = x->m > y->m ? inf : -inf;
        else
            x->p = div(y->m - x->m, x->k - y->k);
        return x->p >= y->p; }

    void add(ll k, ll m) {
        auto z = insert({k, m, 0}), y = z++, x = y;
        while (isect(y, z)) z = erase(z);
        if (x != begin() && isect(--x, y))
            isect(x, y = erase(y));
        while ((y = x) != begin() && (--x)->p >= y->p)
            isect(x, erase(y)); }

    ll query(ll x) {
        assert(!empty());
        Q=1; auto l = *lower_bound({0,0,x}); Q=0;
        return l.k * x + l.m;
```

```

}
};

2.13. Sparse Table  $\mathcal{O}(\log n)$  per query.
struct sparse_table {
    vvi m;
    sparse_table(vi arr) {
        m.pb(arr);
        for (int k=0; (1<<(+k)) <= sz(arr); ) {
            int w = (1<<k), hw = w/2;
            m.pb(vi(sz(arr) - w + 1);
            for (int i = 0; i+w <= sz(arr); i++) {
                m[k][i] = min(m[k-1][i], m[k-1][i+hw]);
            }
        }
    }
    int query(int l, int r) { // query min in [l,r]
        int k = 31 - __builtin_clz(r-l); // k = 0;
        // while (1<<(k+1) <= r-l+1) k++;
        return min(m[k][l], m[k][r-(1<<k)+1]);
    }
};
```

3. GRAPH ALGORITHMS

3.1. Shortest path.

3.1.1. Dijkstra $\mathcal{O}(|E| \log |V|)$.

```

const ll INFY = -1;
vi dijkstra( vector<vii> G, ll s ) {
    vi d( G.size(), INFY );
    priority_queue<ii,vector<ii>,greater<ii>> Q;
    Q.emplace(0,s);
    while(!Q.empty()){
        ll c = Q.top().x, a = Q.top().y;
        Q.pop();
        if(d[a] != INFY)
            continue;
        d[a] = c;
        for(ii e : G[a])
            Q.emplace(d[a] + e.y, e.x);
    }
    return d;
}
```

3.1.2. Floyd-Warshall $\mathcal{O}(V^3)$. Be careful with negative edges! Note: $|d[i][j]|$ can grow exponentially, and $\text{INFY} + \text{negative} < \text{INFY}$.

```

const ll INF = 1LL << 61;
void floyd_warshall( vvi& d ) {
    ll n = d.size();
    REP(i,n) REP(j,n) REP(k,n)
        if(d[j][i] < INF && d[i][k] < INF) // neg edges!
            d[j][k] = max(-INF,
                min(d[j][k], d[j][i] + d[i][k]));
}
```

```
vvi d(n, vi(n, INF));
REP(i, n) d[i][i] = 0;
```

3.1.3. *Bellman Ford* $\mathcal{O}(VE)$. This is only useful if there are edges with weight $w_{ij} < 0$ in the graph.

```
const ll INF = 1LL << 61;
// G[u] = { (v,w) | edge u->v, cost w }
vi bellman_ford(vector<vii> G, ll s) {
    ll n = G.size();
    vi d(n, INF); d[s] = 0;
    REP(loops, n) REP(u, n) if(d[u] != INF)
        for(ii e : G[u]) if(d[u] + e.y < d[e.x])
            d[e.x] = d[u] + e.y;
    // detect paths of -INF length
    for(ll change = 1; change--;)
        REP(u, n) if(d[u] != INF)
            for(ii e : G[u]) if(d[e.x] != -INF)
                if(d[u] + e.y < d[e.x])
                    d[e.x] = -INF, change = 1;
    return d;
}
```

3.1.4. *IDA* algorithm*.

```
int n, cur[100], pos;
int calch() {
    int h = 0;
    rep(i, 0, n) if (cur[i] != 0) h += abs(i - cur[i]);
    return h; }
int dfs(int d, int g, int prev) {
    int h = calch();
    if (g + h > d) return g + h;
    if (h == 0) return 0;
    int mn = INT_MAX;
    rep(di, -2, 3) {
        if (di == 0) continue;
        int nxt = pos + di;
        if (nxt == prev) continue;
        if (0 <= nxt && nxt < n) {
            swap(cur[pos], cur[nxt]);
            swap(pos, nxt);
            mn = min(mn, dfs(d, g+1, nxt));
            swap(pos, nxt);
            swap(cur[pos], cur[nxt]);
        }
        if (mn == 0) break; }
    return mn; }
int idastar() {
    rep(i, 0, n) if (cur[i] == 0) pos = i;
    int d = calch();
    while (true) {
        int nd = dfs(d, 0, -1);
        if (nd == 0 || nd == INT_MAX) return d;
        d = nd; } }
```

3.2. Maximum Matching.

Matching: A set of edges without common vertices (*Maximum* is the **largest** such set, *maximal* is a set which

you cannot add more edges to without breaking the property).

Minimum Vertex Cover: A set of vertices such that each edge in the graph is incident to at least one vertex of the set.

Minimum Edge Cover: A set of edges such that every vertex is incident to at least one edge of the set.

Maximum Independent Set: A set of vertices in a graph such that no two of them are adjacent.

Minimum edge cover \iff Maximum independent set.

König's theorem: In any bipartite graph $G = (L \cup R, E)$, the number of edges in a maximum matching equals the number of vertices in a minimum vertex cover. Let U be the set of unmatched vertices in L , and Z be the set of vertices that are either in U or are connected to U by an alternating path. Then $K = (L \setminus Z) \cup (R \cap Z)$ is the minimum vertex cover.

In any bipartite graph,

$$\text{maxmatch} = \text{MVC} = V - \text{MIS}.$$

See 3.2.3.

3.2.1. *Standard bipartite matching* $\mathcal{O}(nm)$.

```
const int sizeL = 1e4, sizeR = 1e4;
bool vis[sizeR];
int par[sizeR]; // par : R -> L
vi adj[sizeL]; // adj : L -> (N -> R)

bool match(int u) {
    for (int v : adj[u]) {
        if (vis[v]) continue; vis[v] = true;
        if (par[v] == -1 || match(par[v])) {
            par[v] = u;
            return true;
        }
    }
    return false;
}

// perfect matching iff ret == sizeL == sizeR
int maxmatch() {
    fill_n(par, sizeR, -1); int ret = 0;
    for (int i = 0; i < sizeL; i++) {
        fill_n(vis, sizeR, false);
        ret += match(i);
    }
    return ret;
}
```

3.2.2. *Hopcroft-Karp bipartite matching* $\mathcal{O}(E\sqrt{V})$.

```
const ll INF = (1LL << 61LL);
struct bi_graph {
```

```
    ll n, m;
    vvi adj;
    vi L, R, d;
    queue<ll> q;
    bi_graph(ll _n, ll _m) : n(_n), m(_m),
        adj(n, L(n, -1), R(m, n), d(n+1)) {}
    ll add_edge(ll a, ll b) { adj[a].pb(b); }
    ll bfs() {
        rep(v, 0, n)
            if (L[v] == -1) d[v] = 0, q.push(v);
            else d[v] = INF;
        d[n] = INF;
        while (!q.empty()) {
            ll v = q.front(); q.pop();
            if (d[v] < d[n])
                for (ll u : adj[v]) if (d[R[u]] == INF)
                    d[R[u]] = d[v] + 1, q.push(R[u]);
        }
        return d[n] != INF;
    }
    ll dfs(ll v) {
        if (v == n) return true;
        for (ll u : adj[v])
            if (d[R[u]] == d[v] + 1 and dfs(R[u])) {
                R[u] = v; L[v] = u;
                return true;
            }
        d[v] = INF;
        return false;
    }
    ll maximum_matching() {
        ll s = 0;
        while (bfs()) rep(i, 0, n)
            s += L[i] == -1 && dfs(i);
        return s;
    }
};
```

3.2.3. *Minimum Vertex Cover in Bipartite Graphs*.

```
#include "hopcroft_karp.cpp"
vi alt;
void dfs(bi_graph &G, ll v) {
    alt[v] = 1;
    for (ll u : G.adj[v]) {
        alt[u+G.n] = 1;
        if (G.R[u] != G.n && !alt[G.R[u]])
            dfs(G, G.R[u]);
    }
}
vi mvc_bipartite(bi_graph &G) {
    vi res; G.maximum_matching();
    alt.assign(G.n + G.m, 0);
    rep(i, 0, G.n) if (G.L[i] == -1) dfs(G, i);
    rep(i, 0, G.n) if (!alt[i]) res.pb(i);
    rep(i, 0, G.n) if (alt[G.n+i]) res.pb(G.n+i);
    return res;
}
```


3.2.4. *Stable marriage*. With n men, $m \geq n$ women, n preference lists of women for each men, and for every woman j an preference of men defined by `pref[][j]` (lower is better) find for every man a women such that no pair of a men and a woman want to run off together.

```
// n = aantal mannen, m = aantal vrouwen
// voor een man i, is order[i] de prefere
vi stable(int n, int m, vvi order, vvi pref) {
    queue<int> q;
    REP(i, n) q.push(i);
    vi mas(m, -1), mak(n, -1), p(n, 0);
    while (!q.empty()) {
        int k = q.front();
        q.pop();
        int s = order[k][p[k]], k2 = mas[s];
        if (mas[s] == -1) {
            mas[s] = k;
            mak[k] = s;
        } else if (pref[k][s] < pref[k2][s]) {
            mas[s] = k;
            mak[k] = s;
            mak[k2] = -1;
            q.push(k2);
        } else {
            q.push(k);
        }
        p[k]++;
    }
    return mak;
}
```

3.3. Cycle Detection $\mathcal{O}(V + E)$.

```
bool cycle_detection(const vvi &adj) {
    int n = sz(adj); // undirected graph
    stack<int> s;
    vector<bool> vis(n, false); vi par(n, -1);
    s.push(0); vis[0] = true;
    while (!s.empty()) {
        int cur = s.top(); s.pop();
        for (int i : adj[cur]) {
            if (vis[i] && par[cur] != i) return true;
            s.push(i); par[i] = cur; vis[i] = true;
        }
    }
    return false;
}
```

3.4. Depth first searches.

3.4.1. *Topological Sort* $\mathcal{O}(V + E)$.

```
vi topo(vvi &adj) { // requires C++14
    int n=sz(adj); vector<bool> vis(n,0); vi ans;
    auto dfs = [&](int v, const auto& f)->void {
        vis[v] = true;
        for (int w : adj[v]) if (!vis[w]) f(w, f);
        ans.pb(v);
    };
    dfs(0);
    return ans;
}
```

```
REP(i, n) if (!vis[i]) dfs(i, dfs);
reverse(all(ans));
return ans;
}
```

3.4.2. *Cut Points and Bridges* $\mathcal{O}(V + E)$.

```
const int MAXN = 5000;
int low[MAXN], num[MAXN], curnum;

void dfs(vvi &adj, vi &cp, vii &bs, int u, int p) {
    low[u] = num[u] = curnum++;
    int cnt = 0; bool found = false;
    REP(i, sz(adj[u])) {
        int v = adj[u][i];
        if (num[v] == -1) {
            dfs(adj, cp, bs, v, u);
            low[u] = min(low[u], low[v]);
            cnt++;
            found = found || low[v] >= num[u];
            if (low[v] > num[u]) bs.eb(u, v);
        } else if (p != v) low[u] = min(low[u], num[v]);
    }
    if (found && (p != -1 || cnt > 1)) cp.pb(u);
}

pair<vi, vii> cut_points_and_bridges(vvi &adj) {
    int n = size(adj);
    vi cp; vii bs;
    memset(num, -1, n << 2);
    curnum = 0;
    REP(i, n) if (num[i] < 0) dfs(adj, cp, bs, i, -1);
    return make_pair(cp, bs);
}
```

3.4.3. *Strongly Connected Components* $\mathcal{O}(V + E)$.

```
vvi adj, comps;
vi tidix, lnk, cnr, st;
vector<bool> vis;
int age, ncomps;

void tarjan(int v) {
    tidix[v] = lnk[v] = ++age;
    vis[v] = true; st.pb(v);
    for (int w : adj[v]) {
        if (!tidix[w])
            tarjan(w), lnk[v] = min(lnk[v], lnk[w]);
        else if (vis[w])
            lnk[v] = min(lnk[v], tidix[w]);
    }
    if (lnk[v] != tidix[v]) return;
    comps.pb(vi());
    int w;
    do {
        vis[w = st.back()] = false; cnr[w] = ncomps;
        comps.back().pb(w);
        st.pop_back();
    } while (w != v);
    ncomps++;
}
```

```
}

void findSCC(int n) {
    age = ncomps = 0;
    vis.assign(n, false);
    tidix.assign(n, 0);
    lnk.resize(n); cnr.resize(n); comps.clear();
    REP(i, n) if (tidix[i] == 0) tarjan(i);
}
```

3.4.4. *2-SAT* $\mathcal{O}(V + E)$. Include `findSCC`.

```
void init2sat(int n) { adj.assign(2 * n, vi()); }

// (var xl = vl) ==> (var xr = vr)
void imply(int xl, bool vl, int xr, bool vr) {
    adj[2 * xl + vl].pb(2 * xr + vr);
    adj[2 * xr + !vr].pb(2 * xl + !vl);
}

void satOr(int xl, bool vl, int xr, bool vr) {
    imply(xl, !vl, xr, vr);
}

void satConst(int x, bool v) { imply(x, !v, x, v); }
void satIff(int xl, bool vl, int xr, bool vr) {
    imply(xl, vl, xr, vr); imply(xr, vr, xl, vl);
}

bool solve2sat(int n, vector<bool> &sol) {
    findSCC(2 * n);
    for (int i = 0; i < n; i++)
        if (cnr[2 * i] == cnr[2 * i + 1]) return false;
    vector<bool> seen(n, false); sol.assign(n, false);
    for (vi &comp : comps) {
        for (int v : comp) {
            if (seen[v / 2]) continue;
            seen[v / 2] = true;
            sol[v / 2] = v & 1;
        }
    }
    return true;
}
```

3.4.5. *Dominator graph*.

- A node d dominates a node n if every path from the entry node to n must go through d .
- The immediate dominator (idom) of a node n is the unique node that strictly dominates n but does not strictly dominate any other node that strictly dominates n .

```
const int N = 1e6;
vi g[N], g_rev[N], bucket[N];
int pos[N], cnt, order[N], par[N], sdom[N];
int p[N], best[N], idom[N], link[N];

void dfs(int v) {
    pos[v] = cnt;
    order[cnt++] = v;
    for (int u : g[v])
```

```

    if (pos[u] < 0) par[u] = v, dfs(u);
}

int find_best(int x) {
    if (p[x] == x) return best[x];
    int u = find_best(p[x]);
    if (pos[sdom[u]] < pos[sdom[best[x]]])
        best[x] = u;
    p[x] = p[p[x]];
    return best[x];
}

void dominators(int n, int root) {
    fill_n(pos, n, -1);
    cnt = 0;
    dfs(root);
    REP(i, n) for (int u : g[i]) g_rev[u].pb(i);
    REP(i, n) p[i] = best[i] = sdom[i] = i;

    for (int it = cnt - 1; it >= 1; it--) {
        int w = order[it];
        for (int u : g_rev[w]) {
            int t = find_best(u);
            if (pos[sdom[t]] < pos[sdom[w]])
                sdom[w] = sdom[t];
        }
        bucket[sdom[w]].pb(w);
        idom[w] = sdom[w];
        for (int u : bucket[par[w]])
            link[u] = find_best(u);
        bucket[par[w]].clear();
        p[w] = par[w];
    }

    for (int it = 1; it < cnt; it++) {
        int w = order[it];
        idom[w] = idom[link[w]];
    }
}

```

3.5. Min Cut / Max Flow.

3.5.1. Dinic's Algorithm $O(V^2E)$.

```

struct Edge { int t; ll c, f; };
struct Dinic {
    vi H, P; vvi E;
    vector<Edge> G;
    Dinic(int n) : H(n), P(n), E(n) {}

    void addEdge(int u, int v, ll c) {
        E[u].pb(G.size()); G.pb({v, c, 0LL});
        E[v].pb(G.size()); G.pb({u, 0LL, 0LL});
    }

    ll dfs(int t, int v, ll f) {
        if (v == t || !f) return f;
        for ( ; P[v] < (int) E[v].size(); P[v]++) {
            int e = E[v][P[v]], w = G[e].t;
            if (H[w] != H[v] + 1) continue;

```

```

            ll df = dfs(t, w, min(f, G[e].c - G[e].f));
            if (df > 0) {
                G[e].f += df, G[e ^ 1].f -= df;
                return df;
            }
        } return 0;
    }

    ll maxflow(int s, int t, ll f = 0) {
        while (1) {
            fill(all(H), 0); H[s] = 1;
            queue<int> q; q.push(s);
            while (!q.empty()) {
                int v = q.front(); q.pop();
                for (int w : E[v])
                    if (G[w].f < G[w].c && !H[G[w].t])
                        H[G[w].t] = H[v] + 1, q.push(G[w].t);
            }
            if (!H[t]) return f;
            fill(all(P), 0);
            while (ll df = dfs(t, s, LLONG_MAX)) f += df;
        }
    }
};

```

3.5.2. *Min-cost max-flow* $O(n^2m^2)$. Find the cheapest possible way of sending a certain amount of flow through a flow network.

```

const int maxn = 300;

struct edge { ll x, y, f, c, w; };
ll V, par[maxn], D[maxn]; vector<edge> g;
inline void addEdge(int u, int v, ll c, ll w) {
    g.pb({u, v, 0, c, w});
    g.pb({v, u, 0, 0, -w});
}

void sp(int s, int t) {
    fill_n(D, V, LLONG_MAX); D[s] = 0;
    for (int ng = g.size(), _ = V; _--; ) {
        bool ok = false;
        for (int i = 0; i < ng; i++)
            if (D[g[i].x] != LLONG_MAX && g[i].f < g[i].c
                && D[g[i].x] + g[i].w < D[g[i].y]) {
                D[g[i].y] = D[g[i].x] + g[i].w;
                par[g[i].y] = i; ok = true;
            }
        if (!ok) break;
    }
}

```

```

void minCostMaxFlow(int s, int t, ll &c, ll &f) {
    for (c = f = 0; sp(s, t), D[t] < LLONG_MAX; ) {
        ll df = LLONG_MAX, dc = 0;
        for (int v = t, e; e = par[v], v != s; v =
            <- g[e].x) df = min(df, g[e].c - g[e].f);
        for (int v = t, e; e = par[v], v != s; v =
            <- g[e].x) g[e].f += df, g[e ^ 1].f -= df, dc +=
            <- g[e].w;
    }
}

```

```

        f += df; c += dc * df;
    }
}

```

3.5.3. *Gomory-Hu Tree - All Pairs Maximum Flow*. An implementation of the Gomory-Hu Tree. The spanning tree is constructed using Gusfield's algorithm in $O(|V|^2)$ plus $|V|-1$ times the time it takes to calculate the maximum flow. If Dinic's algorithm is used to calculate the max flow, the running time is $O(|V|^3|E|)$. NOTE: Not sure if it works correctly with disconnected graphs.

```

#include "dinic.cpp"
bool same[MAXV];
pair<vii, vvi> construct_gh_tree(flow_network &g) {
    int n = g.n, v;
    vii par(n, ii(0, 0)); vvi cap(n, vi(n, -1));
    rep(s, l, n) {
        int l = 0, r = 0;
        par[s].second = g.max_flow(s, par[s].first,
            <- false);
        memset(d, 0, n * sizeof(int));
        memset(same, 0, n * sizeof(bool));
        d[q[r++] = s] = 1;
        while (l < r) {
            same[v = q[l++]] = true;
            for (int i = g.head[v]; i != -1; i =
                <- g.e[i].nxt)
                if (g.e[i].cap > 0 && d[g.e[i].v] == 0)
                    d[q[r++] = g.e[i].v] = 1;
        }
        rep(i, s+1, n)
            if (par[i].first == par[s].first && same[i])
                par[i].first = s;
        g.reset();
    }
    rep(i, 0, n) {
        int mn = INT_MAX, cur = i;
        while (true) {
            cap[cur][i] = mn;
            if (cur == 0) break;
            mn = min(mn, par[cur].second), cur =
                <- par[cur].first;
        }
        return make_pair(par, cap);
    }
    int compute_max_flow(int s, int t, const pair<vii,
        <- vvi> &gh) {
        int cur = INT_MAX, at = s;
        while (gh.second[at][t] == -1)
            cur = min(cur, gh.first[at].second),
            at = gh.first[at].first;
        return min(cur, gh.second[at][t]);
    }
}

```

3.6. Minimal Spanning Tree $O(E \log V)$.

```

struct edge { int x, y; ll w; };
ll kruskal(int n, vector<edge> edges) {
    dsu D(n);
    sort(all(edges), [] (edge a, edge b) -> bool {
        return a.w < b.w; });
}

```

```

11 ret = 0;
for (edge e : edges)
    if (D.find(e.x) != D.find(e.y))
        ret += e.w, D.unite(e.x, e.y);
return ret;
}

```

3.7. Euler Path $O(V + E)$ hopefully. Finds an Euler Path (or circuit) in a *directed* graph iff one exists.

```

const int MAXV = 1000, MAXE = 5000;
vi adj[MAXV];
int n, m, indeg[MAXV], outdeg[MAXV], res[MAXE + 1];
ii start_end() {
    int start = -1, end = -1, any = 0, c = 0;
    REP(i, n) {
        if (outdeg[i] > 0) any = i;
        if (indeg[i] + 1 == outdeg[i]) start = i, c++;
        else if (indeg[i] == outdeg[i] + 1) end = i, c++;
        else if (indeg[i] != outdeg[i]) return ii(-1, -1);
    }
    if ((start == -1) != (end == -1) || (c != 2 && c))
        return ii(-1, -1);
    if (start == -1) start = end = any;
    return ii(start, end);
}
bool euler_path() {
    ii se = start_end();
    int cur = se.first, at = m + 1;
    if (cur == -1) return false;
    stack<int> s;
    while (true) {
        if (outdeg[cur] == 0) {
            res[--at] = cur;
            if (s.empty()) break;
            cur = s.top(); s.pop();
        } else s.push(cur), cur =
            ↳ adj[cur][--outdeg[cur]];
    }
    return at == 0;
}

```

Finds an Euler *cycle* in a *undirected* graph:

```

const int MAXV = 1000;
multiset<int> adj[MAXV];
list<int> L;
list<int>::iterator euler(int at, int to,
    list<int>::iterator it) {
    if (at == to) return it;
    L.insert(it, at), --it;
    while (!adj[at].empty()) {
        int nxt = *adj[at].begin();
        adj[at].erase(adj[at].find(nxt));
        adj[nxt].erase(adj[nxt].find(at));
        if (to == -1) {
            it = euler(nxt, at, it);
            L.insert(it, at);
            --it;
        } else {
            it = euler(nxt, to, it);

```

```

        to = -1; } }
    return it; }
// usage: euler(0, -1, L.begin());

```

3.8. Heavy-Light Decomposition.

```

struct HLD {
    vvi adj; int cur_pos = 0;
    vi par, dep, hvy, head, pos;

    HLD(int n, const vvi &A) : adj(all(A)), par(n),
        dep(n, hvy(n, -1), head(n), pos(n) {
        cur_pos = 0; dfs(0); decomp(0, 0);
    }

    int dfs(int v) { // determine parent/depth/sizes
        int wei = 1, mw = 0;
        for (int c : adj[v]) if (c != par[v]) {
            par[c] = v, dep[c] = dep[v] + 1;
            int w = dfs(c);
            wei += w;
            if (w > mw) mw = w, hvy[v] = c;
        }
        return wei;
    }

    // pos: index in SegmentTree, head: root of path
    void decomp(int v, int h) {
        head[v] = h, pos[v] = cur_pos++;
        if (hvy[v] != -1) decomp(hvy[v], h);
        for (int c : adj[v])
            if (c != par[v] && c != hvy[v]) decomp(c, c);
    }

    // requires queryST(a, b) = max{A[i] | a ≤ i < b}.
    int query(int a, int b) {
        int res = 0;
        for (; head[a] != head[b]; b = par[head[b]]) {
            if (dep[head[a]] > dep[head[b]]) swap(a, b);
            res = max(res, queryST(pos[head[b]], pos[b] + 1));
        }
        if (dep[a] > dep[b]) swap(a, b);
        return max(res, queryST(pos[a], pos[b] + 1));
    }
};

```

3.9. Centroid Decomposition.

```

#define MAXV 100100
#define LGMAXV 20
int jmp[MAXV][LGMAXV],
    path[MAXV][LGMAXV],
    sz[MAXV], seph[MAXV],
    shortest[MAXV];
struct centroid_decomposition {
    int n; vvi adj;
    centroid_decomposition(int _n) : n(_n), adj(n) {}
    void add_edge(int a, int b) {
        adj[a].push_back(b); adj[b].push_back(a);
    }
    int dfs(int u, int p) {

```

```

        sz[u] = 1;
        rep(i, 0, size(adj[u]))
            if (adj[u][i] != p) sz[u] += dfs(adj[u][i],
                ↳ u);
        return sz[u];
    }
    void makepaths(int sep, int u, int p, int len) {
        jmp[u][seph[sep]] = sep, path[u][seph[sep]] =
            ↳ len;
        int bad = -1;
        rep(i, 0, size(adj[u])) {
            if (adj[u][i] == p) bad = i;
            else makepaths(sep, adj[u][i], u, len + 1);
        }
        if (p == sep)
            swap(adj[u][bad], adj[u].back()),
                ↳ adj[u].pop_back();
    }
    void separate(int h=0, int u=0) {
        dfs(u, -1); int sep = u;
        down: iter(nxt, adj[sep])
            if (sz[*nxt] < sz[sep] && sz[*nxt] > sz[u]/2)
                ↳ {
                    sep = *nxt; goto down;
                }
        seph[sep] = h, makepaths(sep, sep, -1, 0);
        rep(i, 0, size(adj[sep])) separate(h+1,
            ↳ adj[sep][i]);
    }
    void paint(int u) {
        rep(h, 0, seph[u]+1)
            shortest[jmp[u][h]] = min(shortest[jmp[u][h]],
                ↳ path[u][h]);
    }

    int closest(int u) {
        int mn = INT_MAX/2;
        rep(h, 0, seph[u]+1)
            mn = min(mn, path[u][h] +
                ↳ shortest[jmp[u][h]]);
        return mn;
    }
};

```

3.10. Least Common Ancestors, Binary Jumping.

```

const int LOGSZ = 20, SZ = 1 << LOGSZ;
int P[SZ], BP[SZ][LOGSZ];

void initLCA() { // assert P[root] == root
    rep(i, 0, SZ) BP[i][0] = P[i];
    rep(j, 1, LOGSZ) rep(i, 0, SZ)
        BP[i][j] = BP[BP[i][j-1]][j-1];
}

int LCA(int a, int b) {
    if (H[a] > H[b]) swap(a, b);
    int dh = H[b] - H[a], j = 0;
    rep(i, 0, LOGSZ) if (dh & (1 << i)) b = BP[b][i];
    while (BP[a][j] != BP[b][j]) j++;
    while (--j >= 0) if (BP[a][j] != BP[b][j])
        a = BP[a][j], b = BP[b][j];
    return a == b ? a : P[a];
}

```

3.11. Miscellaneous.

3.11.1. *Misra-Gries $D+1$ -edge coloring.* Finds a $\max_i \deg(i) + 1$ -edge coloring where there all incident edges have distinct colors. Finding a D -edge coloring is NP-hard.

```
struct Edge { int to, col, rev; };

struct MisraGries {
    int N, K=0; vvi F;
    vector<vector<Edge>> G;

    MisraGries(int n) : N(n), G(n) {}
    // add an undirected edge, NO DUPLICATES ALLOWED
    void addEdge(int u, int v) {
        G[u].pb({v, -1, (int) G[v].size()});
        G[v].pb({u, -1, (int) G[u].size()-1});
    }

    void color(int v, int i) {
        vi fan = { i };
        vector<bool> used(G[v].size());
        used[i] = true;
        for (int j = 0; j < (int) G[v].size(); j++)
            if (!used[j] && G[v][j].col >= 0 &&
                F[G[v][fan.back()].to][G[v][j].col] < 0)
                used[j] = true, fan.pb(j), j = -1;
        int c = 0; while (F[v][c] >= 0) c++;
        int d = 0; while (F[G[v][fan.back()].to][d] >= 0) d++;
        int w = v, a = d, k = 0, ccol;
        while (true) {
            swap(F[w][c], F[w][d]);
            if (F[w][c] >= 0) G[w][F[w][c]].col = c;
            if (F[w][d] >= 0) G[w][F[w][d]].col = d;
            if (F[w][a^c^d] < 0) break;
            w = G[w][F[w][a]].to;
        }
        do {
            Edge &e = G[v][fan[k]];
            ccol = F[e.to][d] < 0 ? d :
                G[v][fan[k+1]].col;
            if (e.col >= 0) F[e.to][e.col] = -1;
            F[e.to][ccol] = e.rev;
            F[v][ccol] = fan[k];
            e.col = G[e.to][e.rev].col = ccol;
            k++;
        } while (ccol != d);
    }
    // finds a K-edge-coloring
    void color() {
        REP(v, N) K = max(K, (int) G[v].size() + 1);
        F = vvi(N, vi(K, -1));
        REP(v, N) for (int i = G[v].size(); i--;)
            if (G[v][i].col < 0) color(v, i);
    }
};
```

3.11.2. *Minimum Mean Weight Cycle.* Given a strongly connected directed graph, finds the cycle of minimum mean weight. If you have a graph that is not strongly connected, run this on each strongly connected component.

```
double
↳ min_mean_cycle(vector<vector<pair<int,double>>>
↳ adj){
    int n = size(adj); double mn = INFINITY;
    vector<vector<double>> arr(n+1, vector<double>(n,
        ↳ mn));
    arr[0][0] = 0;
    rep(k,1,n+1) rep(j,0,n) iter(it,adj[j])
        arr[k][it->first] = min(arr[k][it->first],
            it->second +
            ↳ arr[k-1][j]);

    rep(k,0,n) {
        double mx = -INFINITY;
        rep(i,0,n) mx = max(mx,
            ↳ (arr[n][i]-arr[k][i])/(n-k));
        mn = min(mn, mx); }
    return mn; }
```

3.11.3. *Minimum Arborescence.* Given a weighted directed graph, finds a subset of edges of minimum total weight so that there is a unique path from the root r to each vertex. Returns a vector of size n , where the i th element is the edge for the i th vertex. The answer for the root is undefined!

$\mathcal{O}(EV)$ runtime and $\mathcal{O}(E)$ memory:

```
#include "../datastructures/union_find.cpp"
struct arborescence {
    int n; union_find uf;
    vector<vector<pair<ii,int>>> adj;
    arborescence(int _n) : n(_n), uf(n), adj(n) {}
    void add_edge(int a, int b, int c) {
        adj[b].eb(ii(a,b),c); }
    vii find_min(int r) {
        vi vis(n,-1), mn(n,INT_MAX); vii par(n);
        REP(i, n) {
            if (uf.find(i) != i) continue;
            int at = i;
            while (at != r && vis[at] == -1) {
                vis[at] = i;
                for (auto it : adj[at])
                    if (it.y < mn[at] && uf.find(it.x.x) !=
                        ↳ at)
                        mn[at] = it.y, par[at] = it.x;
                if (par[at] == ii(0,0)) return vii();
                at = uf.find(par[at].x);
            }
            if (at == r || vis[at] != i) continue;
            union_find tmp = uf;
            vi seq;
            do seq.pb(at), at = uf.find(par[at].x);
            while (at != seq.front());
            int c = uf.find(seq[0]);
        }
    }
};
```

```
for (auto it : seq) uf.unite(it, c);
for (auto &jt : adj[c]) jt.y -= mn[c];
for (auto it : seq) {
    if (it == c) continue;
    for (auto jt : adj[it])
        adj[c].eb(jt.x, jt.y - mn[it]);
    adj[it].clear();
}
vii rest = find_min(r);
if (rest.empty()) return rest;
ii use = rest[c];
rest[at = tmp.find(use.y)] = use;
for (int it : seq) if (it != at)
    rest[it] = par[it];
return rest;
}
return par; } };
```

$\mathcal{O}(V^2 \log V)$ runtime and $\mathcal{O}(E)$ memory:

```
const int oo = 0x3f3f3f3f, MAXN = 4024;

// N = #V, R = root
int N, R;
// for each node a list of pairs (predecessor,
↳ cost):
vector<pii> g[MAXN];
int pred[MAXN], label[MAXN], node[MAXN],
↳ helper[MAXN];

int get_node(int n) {
    return node[n] == n ? n :
        (node[n] = get_node(node[n]));
}

int update_node(int n) {
    int m = oo;
    for (auto ed : g[n]) m = min(m, ed.y);
    REP(j, sz(g[n])) {
        g[n][j].y -= m;
        if (g[n][j].y == 0)
            pred[n] = g[n][j].x;
    }
    return m;
}

ll cycle(vi &active, int n, int &end) {
    n = get_node(n);
    if (label[n] == 1) return false;
    if (label[n] == 2) { end = n; return 0; }

    active.pb(n);
    label[n] = 2;
    auto res = cycle(active, pred[n], end);
    if (end == n) {
        int F = find(all(active), n)-active.begin();
        vi todo(active.begin() + F, active.end());
        active.resize(F);
        vii> newg;
```

```

for (auto i: todo) node[i] = n;
for (auto i: todo) for (auto &ed : g[i])
    helper[ed.x] = get_node(ed.x) = ed.y;
for (auto i: todo) for (auto ed : g[i])
    helper[ed.x] = min(ed.y, helper[ed.x]);
for (auto i: todo) for (auto ed: g[i]) {
    if (helper[ed.x] != oo && ed.x != n) {
        newg.ed(ed.x, helper[ed.x]);
        helper[ed.x] = oo;
    }
}
g[n] = newg;
res += update_node(n);
label[n] = 0;
cend = -1;
return cycle(active, n, cend) + res;
}
if (cend == -1) {
    active.pop_back();
    label[n] = 1;
}
return res;
}

// Calculates value of minimal arborescence from R,
// assuming it exists.
// NOTE: N, R must be initialized at this point!!!
// Algo changes g!!
ll min_arbor() {
    ll res = 0;
    REP(i, N) {
        node[i] = i;
        if (i != R) res += update_node(i);
    }
    REP(i, N) label[i] = (i==R);
    REP(i, N) {
        if (label[i] == 1 || get_node(i) != i)
            continue;
        vi active;
        int cend = -1;
        res += cycle(active, i, cend);
    }
    return res;
}

```

3.11.4. *Maximum Density Subgraph*. Given (weighted) undirected graph G . Binary search density. If g is current density, construct flow network: (S, u, m) , $(u, T, m + 2g - d_u)$, $(u, v, 1)$, where m is a large constant (larger than sum of edge weights). Run floating-point max-flow. If minimum cut has empty S -component, then maximum density is smaller than g , otherwise it's larger. Distance between valid densities is at least $1/(n(n-1))$. Edge case when density is 0. This also works for weighted graphs by replacing d_u by the weighted degree, and doing more iterations (if weights are not integers).

3.11.5. *Maximum-Weight Closure*. Given a vertex-weighted directed graph G . Turn the graph into a flow network, adding weight ∞ to each edge. Add vertices S, T . For each vertex v of weight w , add edge (S, v, w) if $w \geq 0$, or edge $(v, T, -w)$ if $w < 0$. Sum of positive weights minus minimum $S - T$ cut is the answer. Vertices reachable from S are in the closure. The maximum-weight closure is the same as the complement of the minimum-weight closure on the graph with edges reversed.

3.11.6. *Maximum Weighted Independent Set in a Bipartite Graph*. This is the same as the minimum weighted vertex cover. Solve this by constructing a flow network with edges $(S, u, w(u))$ for $u \in L$, $(v, T, w(v))$ for $v \in R$ and (u, v, ∞) for $(u, v) \in E$. The minimum S, T -cut is the answer. Vertices adjacent to a cut edge are in the vertex cover.

3.11.7. *Synchronizing word problem*. A DFA has a synchronizing word (an input sequence that moves all states to the same state) iff. each pair of states has a synchronizing word. That can be checked using reverse DFS over pairs of states. Finding the shortest synchronizing word is NP-complete.

4. STRING ALGORITHMS

4.1. Trie.

```

const int SIGMA = 26;

struct trie {
    bool word; trie **adj;

    trie() : word(false), adj(new trie*[SIGMA]) {
        for (int i = 0; i < SIGMA; i++) adj[i] = NULL;
    }

    void addWord(const string &str) {
        trie *cur = this;
        for (char ch : str) {
            int i = ch - 'a';
            if (!cur->adj[i]) cur->adj[i] = new trie();
            cur = cur->adj[i];
        }
        cur->word = true;
    }

    bool isWord(const string &str) {
        trie *cur = this;
        for (char ch : str) {
            int i = ch - 'a';
            if (!cur->adj[i]) return false;
            cur = cur->adj[i];
        }
        return cur->word;
    }
};

```

4.2. Z-algorithm $\mathcal{O}(n)$.

```

// z[i] = length of longest substring starting from
// s[i] which is also a prefix of s.
vi z_function(const string &s) {
    int n = (int) s.length();
    vi z(n);
    for (int i = 1, l = 0, r = 0; i < n; ++i) {
        if (i <= r) z[i] = min(r - i + 1, z[i - l]);
        while (i + z[i] < n && s[z[i]] == s[i + z[i]])
            ++z[i];
        if (i + z[i] - 1 > r) l = i, r = i + z[i] - 1;
    }
    return z;
}

```

4.3. *Suffix array* $\mathcal{O}(n \log n)$. Lexicographically sorts the cyclic shifts of S where $p[0]$ is the index of the smallest string, etc.

```

vi sort_cyclic_shifts(const string &s) {
    const int alphabet = 256, n = sz(s);

    vi p(n), c(n), cnt(max(alphabet, n), 0);
    REP(i, n) cnt[s[i]]++;
    partial_sum(cnt.begin(), cnt.begin() + n, cnt);
    REP(i, n) p[--cnt[s[i]]] = i;
    c[p[0]] = 0;
    int cl = 1;
    rep(i, 1, n) {
        if (s[p[i]] != s[p[i-1]]) cl++;
        c[p[i]] = cl - 1;
    }

    vi pn(n), cn(n);
    for (int h = 0, l = 1; l < n; l *= 2, ++h) {
        REP(i, n) {
            pn[i] = p[i] - (1 << h);
            if (pn[i] < 0) pn[i] += n;
        }
        fill(cnt.begin(), cnt.begin() + n, 0);
        REP(i, n) cnt[c[pn[i]]]++;
        rep(i, 1, n) cnt[i] += cnt[i-1];
        for (int i = n-1; i >= 0; i--)
            p[--cnt[c[pn[i]]]] = pn[i];
        cn[p[0]] = 0;
        cl = 1;
        rep(i, 1, n) {
            if (c[p[i]] != c[p[i-1]] || c[(p[i]+l)%n]
                != c[(p[i-1]+l)%n]) cl++;
            cn[p[i]] = cl - 1;
        }
        c.swap(cn);
    }
    return p;
}

vi suffix_array(string s) {

```



```

s += '\0';
vi v = sort_cyclic_shifts(s);
v.erase(v.begin());
return v;
}

```

4.4. Longest Common Subsequence $\mathcal{O}(n^2)$. SUBSTRING: consecutive characters!!!

```

int dp[STR_SIZE][STR_SIZE]; // DP problem

int lcs(const string &w1, const string &w2) {
    int n1 = w1.size(), n2 = w2.size();
    for (int i = 0; i < n1; i++) {
        for (int j = 0; j < n2; j++) {
            if (i == 0 || j == 0) dp[i][j] = 0;
            else if (w1[i-1] == w2[j-1])
                dp[i][j] = dp[i-1][j-1] + 1;
            else dp[i][j] = max(dp[i-1][j], dp[i][j-1]);
        }
    }
    return dp[n1][n2];
}

```

```

// backtrace
string getLCS(const string &w1, const string &w2) {
    int i = w1.size(), j = w2.size(); string ret = "";
    while (i > 0 && j > 0) {
        if (w1[i-1] == w2[j-1]) ret += w1[--i], j--;
        else if (dp[i][j-1] > dp[i-1][j]) j--;
        else i--;
    }
    reverse(ret.begin(), ret.end());
    return ret;
}

```

4.5. Levenshtein Distance $\mathcal{O}(n^2)$. Minimal number of insertions, removals and edits required to transform one string in the other.

```

int dp[MAX_SIZE][MAX_SIZE]; // DP problem

int levDist(const string &w1, const string &w2) {
    int n1 = sz(w1)+1, n2 = sz(w2)+1;
    REP(i, n1) dp[i][0] = i; // removal
    REP(j, n2) dp[0][j] = j; // insertion
    rep(i, 1, n1) rep(j, 1, n2)
        dp[i][j] = min(
            1 + min(dp[i-1][j], dp[i][j-1]),
            dp[i-1][j-1] + (w1[i-1] != w2[j-1])
        );
    return dp[n1][n2];
}

```

4.6. Knuth-Morris-Pratt algorithm $\mathcal{O}(N + M)$.

```

int kmp(const string &word, const string &text) {
    int n = word.size();
    vi T(n + 1, 0);
    for (int i = 1, j = 0; i < n; ) {

```

```

        if (word[i] == word[j]) T[++i] = ++j; // match
        else if (j > 0) j = T[j]; // fallback
        else i++; // no match, keep zero
    }
    int matches = 0;
    for (int i = 0, j = 0; i < text.size(); ) {
        if (text[i] == word[j]) {
            i++;
            if (++j == n) // match at interval [i - n, i)
                matches++, j = T[j];
        } else if (j > 0) j = T[j];
        else i++;
    }
    return matches;
}

```

4.7. Aho-Corasick Algorithm $\mathcal{O}(N + \sum_{i=1}^m |S_i|)$. Dictionary substring matching as automaton. All given P must be unique!

```

const int MAXP = 100, MAXLEN = 200, SIGMA = 26,
        MAXTRIE = MAXP * MAXLEN;

```

```

int nP;
string P[MAXP], S;

```

```

int pnr[MAXTRIE], to[MAXTRIE][SIGMA],
        sLink[MAXTRIE], dLink[MAXTRIE], nnodes;

```

```

void ahoCorasick() {
    fill_n(pnr, MAXTRIE, -1);
    for (int i = 0; i < MAXTRIE; i++) fill_n(to[i],
        SIGMA, 0);
    fill_n(sLink, MAXTRIE, 0); fill_n(dLink, MAXTRIE,
        0);
    nnodes = 1;
    // STEP 1: MAKE A TREE
    for (int i = 0; i < nP; i++) {
        int cur = 0;
        for (char c : P[i]) {
            int i = c - 'a';
            if (to[cur][i] == 0) to[cur][i] = nnodes++;
            cur = to[cur][i];
        }
        pnr[cur] = i;
    }
    // STEP 2: CREATE SUFFIX_LINKS AND DICT_LINKS
    queue<int> q; q.push(0);
    while (!q.empty()) {
        int cur = q.front(); q.pop();
        for (int c = 0; c < SIGMA; c++) {
            if (to[cur][c]) {
                int sl = sLink[to[cur][c]] = cur == 0 ? 0 :
                    to[sLink[cur]][c];
                // if all strings have equal length, remove
                // this:
                dLink[to[cur][c]] = pnr[sl] >= 0 ? sl :
                    dLink[sl];
                q.push(to[cur][c]);
            }

```

```

        } else to[cur][c] = to[sLink[cur]][c];
    }
}
// STEP 3: TRAVERSE S
for (int cur = 0, i = 0, n = S.size(); i < n; i++)
    cur = to[cur][S[i] - 'a'];
for (int hit = pnr[cur] >= 0 ? cur : dLink[cur];
    hit; hit = dLink[hit]) {
    cerr << P[pnr[hit]] << " found at [" << (i + 1)
        << " - " << P[pnr[hit]].size() << ", " << i << "]"
        << endl;
}
}
}

```

4.8. eerTree. Constructs an eerTree in $\mathcal{O}(n)$, one character at a time.

```

#define MAXN 100100
#define SIGMA 26
#define BASE 'a'
char *s = new char[MAXN];
struct state {
    int len, link, to[SIGMA];
} *st = new state[MAXN+2];
struct eertree {
    int last, sz, n;
    eertree() : last(1), sz(2), n(0) {
        st[0].len = st[0].link = -1;
        st[1].len = st[1].link = 0;
    }
    int extend() {
        char c = s[n++]; int p = last;
        while (n - st[p].len - 2 < 0 || c != s[n -
            2])
            p = st[p].link;
        if (!st[p].to[c-BASE]) {
            int q = last = sz++;
            st[p].to[c-BASE] = q;
            st[q].len = st[p].len + 2;
            do { p = st[p].link;
                while (p != -1 && (n < st[p].len + 2 ||
                    c != s[n - st[p].len - 2]));
            } while (p == -1) st[q].link = 1;
            else st[q].link = st[p].to[c-BASE];
            return 1;
        }
        last = st[p].to[c-BASE];
        return 0;
    }
};

```

4.9. Suffix Automaton. Minimum automata that accepts all suffixes of a string with $\mathcal{O}(n)$ construction. The automata itself is a DAG therefore suitable for DP, examples are counting unique substrings, occurrences of substrings and suffix.

```
// TODO: Add longest common subsring
const int MAXL = 100000;
struct suffix_automaton {
    vi len, link, occur, cnt;
    vector<map<char,int> > next;
    vector<bool> isclone;
    ll *occuratleast;
    int sz, last;
    string s;
    suffix_automaton() : len(MAXL*2), link(MAXL*2),
        occur(MAXL*2), next(MAXL*2), isclone(MAXL*2) {
        clear(); }
    void clear() { sz = 1; last = len[0] = 0; link[0]
        = -1;
        next[0].clear(); isclone[0] =
            false; }
    bool issubstr(string other){
        for(int i = 0, cur = 0; i < size(other); ++i){
            if(cur == -1) return false; cur =
                next[cur][other[i]]; }
        return true; }
    void extend(char c){ int cur = sz++; len[cur] =
        len[last]+1;
        next[cur].clear(); isclone[cur] = false; int p =
            last;
        for(; p != -1 && !next[p].count(c); p = link[p])
            next[p][c] = cur;
        if(p == -1){ link[cur] = 0; }
        else{ int q = next[p][c];
            if(len[p] + 1 == len[q]){ link[cur] = q; }
            else { int clone = sz++; isclone[clone] =
                true;
                len[clone] = len[p] + 1;
                link[clone] = link[q]; next[clone] =
                    next[q];
                for(; p != -1 && next[p].count(c) &&
                    next[p][c] == q;
                    p = link[p]){
                    next[p][c] = clone; }
                link[q] = link[cur] = clone;
            } } last = cur; }
    void count(){
        cnt=vi(sz, -1); stack<ii> S; S.push(ii(0,0));
        map<char,int>::iterator i;
        while(!S.empty()){
            ii cur = S.top(); S.pop();
            if(cur.y){
                for(i = next[cur.x].begin();
                    i != next[cur.x].end();++i){
                    cnt[cur.x] += cnt[(i).y]; } }
            else if(cnt[cur.x] == -1){
                cnt[cur.x] = 1; S.push(ii(cur.x, 1));
                for(i = next[cur.x].begin();
                    i != next[cur.x].end();++i){
                    S.push(ii((i).y, 0)); } } } }
        string lexicok(ll k){
            int st=0; string res; map<char,int>::iterator i;
```

```
while(k){
    for(i = next[st].begin(); i != next[st].end();
        ++i){
        if(k <= cnt[(i).y]){ st = (i).y;
            res.push_back((i).x); k--; break;
        } else { k -= cnt[(i).y]; } } }
    return res; }
void countoccur(){
    REP(i, sz) occur[i] = 1 - isclone[i];
    vii states(sz);
    REP(i, sz) states[i] = ii(len[i],i);
    sort(states.begin(), states.end());
    for(int i = size(states)-1; i >= 0; --i){
        int v = states[i].y;
        if (link[v] != -1)
            occur[link[v]] += occur[v]; } };
```

4.10. **Hashing.** Modulus should be a large prime. Can also use multiple instances with different moduli to minimize chance of collision.

```
struct hasher {
    int b = 311, m; vi h, p;
    hasher(string s, int _m) :
        m(_m), h(sz(s)+1), p(sz(s)+1) {
        p[0] = 1; h[0] = 0;
        REP(i, sz(s)) p[i+1] = (ll)p[i] * b % m;
        REP(i, sz(s)) h[i+1] = ((ll)h[i] * b + s[i]) % m;
    }
    int hash(int l, int r) {
        return (h[r+1] + m - (ll)h[l]*p[r-l+1] % m) % m;
    }
};
```

5. GEOMETRY

```
const ld EPS = 1e-7, PI = acos(-1.0);
typedef ld NUM; // EITHER ld OR ll
typedef pair<NUM, NUM> pt;

pt operator+(pt p,pt q){ return {p.x+q.x,p.y+q.y}; }
pt operator-(pt p,pt q){ return {p.x-q.x,p.y-q.y}; }
pt operator*(pt p, NUM n) { return {p.x*n, p.y*n}; }

pt& operator+=(pt &p, pt q) { return p = p+q; }
pt& operator-=(pt &p, pt q) { return p = p-q; }
```

```
NUM operator*(pt p, pt q){ return p.x*q.x+p.y*q.y; }
NUM operator^(pt p, pt q){ return p.x*q.y-p.y*q.x; }
```

// square distance from p to line ab

```
ld distPtLineSq(pt p, pt a, pt b) {
    p -= a; b -= a;
    return ld(p^b) * (p^b) / (b*b);
}
```

// square distance from p to linesegment ab

```
ld distPtSegmentSq(pt p, pt a, pt b) {
    p -= a; b -= a;
    NUM dot = p*b, len = b*b;
```

```
if (dot <= 0) return p*p;
if (dot >= len) return (p-b)*(p-b);
return p*p - ld(dot)*dot/len;
}
// Test if p is on line segment ab
bool segmentHasPoint(pt p, pt a, pt b) {
    pt u = p-a, v = p-b;
    return abs(u^v) < EPS && u*v <= 0;
}
```

// projects p onto the line ab

```
pair<ld,ld> proj(pt p, pt a, pt b) {
    p -= a; b -= a;
    return a + b*(ld(b*p) / (b*b));
}
```

```
bool col(pt a, pt b, pt c) {
    return abs((a-b) ^ (a-c)) < EPS;
}
```

// true => 1 intersection, false => parallel or same

```
bool linesIntersect(pt a, pt b, pt c, pt d) {
    return abs((a-b) ^ (c-d)) > EPS;
}
```

```
pair<ld,ld> lineLineIntersection(pt a, pt b, pt c,
    pt d) {
    ld det = (a-b) ^ (c-d);
    assert(abs(det) > EPS);
    return ((c-d)*(a^b) - (a-b)*(c^d)) *
        (ld(1.0)/det);
}
```

// dp, dq are directions from p, q

// intersection at p + t_i dp, for 0 <= i < return

```
value
int segmentIntersection(pt p, pt dp, pt q, pt dq,
    pt &A, pt &B) {
    if (abs(dp * dp)<EPS)
        swap(p,q), swap(dp,dq); // dq=0
    if (abs(dp * dp)<EPS) {
        A = p; // dp = dq = 0
        return p == q;
    }
```

```
pt dpq = q-p;
NUM c = dp^dq, c0 = dpq^dp, c1 = dpq^dq;
if (abs(c) < EPS) { // parallel, dp > 0, dq >= 0
    if (abs(c0) > EPS) return 0; // not collinear
    NUM v0 = dpq*dp, v1 = v0 + dq*dp, dp2 = dp*dp;
    if (v1 < v0) swap(v0, v1);
```

```
v0 = max(v0, NUM(0));
v1 = min(v1, dp2);
```

```
A = p + dp * (ld(v0) / dp2);
```

```

    B = p + dp * (ld(v1) / dp2);

    return (v0 <= v1) + (v0 < v1);
}

if (c < 0) {
    c = -c; c0 = -c0; c1 = -c1;
}

A = p + dp * (ld(c1)/c);
return 0 <= min(c0,c1) && max(c0,c1) <= c;
}

// Returns TWICE the area of a polygon (for
// integers)
NUM polygonTwiceArea(const vector<pt> &p) {
    NUM area = 0;
    for (int n = sz(p), i=0, j=n-1; i<n; j = i++)
        area += p[i].x * p[j].y;
    return abs(area); // area < 0 ==> p ccw
}

bool insidePolygon(const vector<pt> &pts, pt p, bool
// strict = true) {
    int n = 0;
    for (int N = sz(pts), i = 0, j = N - 1; i < N; j =
        i++) {
        // if p is on edge of polygon
        if (segmentHasPoint(p, pts[i], pts[j])) return
            // !strict;
        // or: if(distPtSegmentSq(p, pts[i], pts[j]) <=
            // EPS) return !strict;

        // increment n if segment intersects line from p
        n += (max(pts[i].y, pts[j].y) > p.y &&
            // min(pts[i].y, pts[j].y) <= p.y &&
            // ((pts[j].x - pts[i].x)*(p-pts[i].x)) > 0) ==
            // (pts[i].y <= p.y));
    }
    return n & 1; // inside if odd number of
// intersections
}

```

5.1. Convex Hull $\mathcal{O}(n \log n)$.

```

// the convex hull consists of: { pts[ret[0]],
// pts[ret[1]], ... pts[ret.back()] }
vi convexHull(const vector<pt> &pts) {
    if (pts.empty()) return vi();
    vi ret, ord;
    int n = pts.size(), st = min_element(all(pts)) -
        pts.begin();
    rep(i, 0, n)
        if (pts[i] != pts[st]) ord.pb(i);
    sort(all(ord), [&pts,&st] (int a, int b) {
        pt p = pts[a] - pts[st], q = pts[b] - pts[st];
        return (p ^ q) != 0 ? (p ^ q) > 0 : lenSq(p) <
            lenSq(q);
    });
}

```

```

});
ord.pb(st); ret.pb(st);
for (int i : ord) {
    // use '>' to include ALL points on the
    // hull-line
    for (int s = ret.size() - 1; s > 0 &&
        ((pts[ret[s-1]] - pts[ret[s]]) ^ (pts[i] -
            pts[ret[s]])) >= 0; s--)
        ret.pop_back();
    ret.pb(i);
}
ret.pop_back();
return ret;
}

```

5.2. Rotating Calipers $\mathcal{O}(n)$. Finds the longest distance between two points in a convex hull.

```

NUM rotatingCalipers(vector<pt> &hull) {
    int n = hull.size(), a = 0, b = 1;
    if (n <= 1) return 0.0;
    while (((hull[1] - hull[0]) ^ (hull[(b + 1) % n] -
        hull[b])) > 0) b++;
    NUM ret = 0.0;
    while (a < n) {
        ret = max(ret, lenSq(hull[a], hull[b]));
        if (((hull[(a + 1) % n] - hull[a]) ^
            (hull[(b + 1) % n] - hull[b])) <= 0) a++;
        else if (++b == n) b = 0;
    }
    return ret;
}

```

5.3. Closest points $\mathcal{O}(n \log n)$.

```

int n; pt pts[maxn];

struct byY {
    bool operator()(int a, int b) const { return
        pts[a].y < pts[b].y; }
};

inline NUM dist(ii p) { return hypot(pts[p.x].x -
    pts[p.y].x, pts[p.x].y - pts[p.y].y); }

ii minpt(ii p1, ii p2) { return dist(p1) < dist(p2)
    // ? p1 : p2; }

// closest pts (by index) inside pts[l ... r], with
// sorted y values in ys
ii closest(int l, int r, vi &ys) {
    if (r - l == 2) { // don't assume 1 here.
        ys = { l, l + 1 };
        return ii(l, l + 1);
    } else if (r - l == 3) { // brute-force
        ys = { l, l + 1, l + 2 };
        sort(all(ys), byY());
        return minpt(ii(l, l + 1), minpt(ii(l, l + 2),
            ii(l + 1, l + 2)));
    }
}

```

```

}
int m = (l + r) / 2; vi yl, yr;
ii delta = minpt(closest(l, m, yl), closest(m, r,
    yr));
NUM ddelta = dist(delta), xm = .5 * (pts[m-1].x +
    pts[m].x);
merge(all(yl), all(yr), back_inserter(ys), byY());
deque<int> q;
for (int i : ys) if (abs(pts[i].x - xm) <= ddelta)
    q.pb(i);
for (int j : q) delta = minpt(delta, ii(i, j));
q.pb(i);
if (q.size() > 8) q.pop_front(); // magic from
// Introduction to Algorithms.
}
return delta;
}
}

```

5.4. Great-Circle Distance. Computes the distance between two points (given as latitude/longitude coordinates) on a sphere of radius r .

```

ld gc_distance(ld pLat, ld pLong, ld qLat, ld qLong,
// ld r) {
    pLat *= pi / 180; pLong *= pi / 180;
    qLat *= pi / 180; qLong *= pi / 180;
    return r * acos(cos(pLat)*cos(qLat)*cos(pLong -
        qLong) + sin(pLat)*sin(qLat)); }

```

5.5. Delaunay triangulation.

```

// https://cp-algorithms.com/geometry/delaunay.html
typedef long long ll;

bool ge(const ll& a, const ll& b) { return a >= b; }
bool le(const ll& a, const ll& b) { return a <= b; }
bool eq(const ll& a, const ll& b) { return a == b; }
bool gt(const ll& a, const ll& b) { return a > b; }
bool lt(const ll& a, const ll& b) { return a < b; }
int sgn(const ll& a) { return (a>0) - (a<0); }

struct pt {
    ll x, y;
    pt() {}
    pt(ll _x, ll _y) : x(_x), y(_y) {}
    pt operator-(const pt& p) const {
        return pt(x - p.x, y - p.y); }
    ll cross(const pt& p) const {
        return x*p.y - y*p.x; }
    ll cross(const pt& a, const pt& b) const {
        return (a - *this).cross(b - *this); }
    ll dot(const pt& p) const {
        return x*p.x + y*p.y; }
    ll dot(const pt& a, const pt& b) const {
        return (a - *this).dot(b - *this); }
    ll lenSq() const { return this->dot(*this); }
    bool operator==(const pt& p) const {

```

```

    return eq(x, p.x) && eq(y, p.y); }
};

const pt inf_pt = pt(1e18, 1e18);

struct Quad { // `QuadEdge` originally
    pt O; // origin
    Quad *rot = nullptr, *onext = nullptr;
    bool used = false;
    Quad* rev() const { return rot->rot; }
    Quad* lnext() const {
        return rot->rev()->onext->rot; }
    Quad* oprev() const {
        return rot->onext->rot; }
    pt dest() const { return rev()->O; }
};

Quad* make_edge(pt from, pt to) {
    Quad* e1 = new Quad, e2 = new Quad;
    Quad* e3 = new Quad, e4 = new Quad;
    e1->O = from; e2->O = to;
    e3->O = e4->O = inf_pt;
    e1->rot = e3; e2->rot = e4;
    e3->rot = e2; e4->rot = e1;
    e1->onext = e1; e2->onext = e2;
    e3->onext = e4; e4->onext = e3;
    return e1;
}

void splice(Quad* a, Quad* b) {
    swap(a->onext->rot->onext, b->onext->rot->onext);
    swap(a->onext, b->onext);
}

void delete_edge(Quad* e) {
    splice(e, e->oprev());
    splice(e->rev(), e->rev()->oprev());
    delete e->rot; delete e->rev()->rot;
    delete e; delete e->rev();
}

Quad* connect(Quad* a, Quad* b) {
    Quad* e = make_edge(a->dest(), b->O);
    splice(e, a->lnext());
    splice(e->rev(), b);
    return e;
}

bool left_of(pt p, Quad* e) {
    return gt(p.cross(e->O, e->dest()), 0); }
bool right_of(pt p, Quad* e) {
    return lt(p.cross(e->O, e->dest()), 0); }

template <class T> T det3(T a1, T a2, T a3,
    T b1, T b2, T b3, T c1, T c2, T c3) {
    return a1*(b2*c3 - c2*b3) - a2*(b1*c3 - c1*b3)
        + a3*(b1*c2 - c1*b2);
}

```

```

// Calculate directly with __int128, or with angles
bool in_circle(pt a, pt b, pt c, pt d) {
    #if defined(__LP64__) || defined(_WIN64)
        __int128 det = 0;
        det -= det3<__int128>(b.x, b.y, b.lenSq(),
            c.x, c.y, c.lenSq(), d.x, d.y, d.lenSq());
        det += det3<__int128>(a.x, a.y, a.lenSq(),
            c.x, c.y, c.lenSq(), d.x, d.y, d.lenSq());
        det -= det3<__int128>(a.x, a.y, a.lenSq(),
            b.x, b.y, b.lenSq(), d.x, d.y, d.lenSq());
        det += det3<__int128>(a.x, a.y, a.lenSq(),
            b.x, b.y, b.lenSq(), c.x, c.y, c.lenSq());
        return det > 0;
    #else
        auto ang = [](pt l, pt mid, pt r) {
            ll x = mid.dot(l, r), y = mid.cross(l, r);
            return atan2((ld) x, (ld) y);
        };
        return (ang(a, b, c) + ang(c, d, a)
            - ang(b, c, d) - ang(d, a, b)) > 1e-8;
    #endif
}

pair<Quad*, Quad*> build_tr(int l, int r,
    vector<pt>& p) {
    if (r - l + 1 == 2) {
        Quad* res = make_edge(p[l], p[r]);
        return make_pair(res, res->rev());
    }
    if (r - l + 1 == 3) {
        Quad *a = make_edge(p[l], p[l+1]);
        Quad *b = make_edge(p[l+1], p[r]);
        splice(a->rev(), b);
        int sg = sgn(p[l].cross(p[l+1], p[r]));
        if (sg == 0) return make_pair(a, b->rev());
        Quad* c = connect(b, a);
        if (sg == 1) return make_pair(a, b->rev());
        return make_pair(c->rev(), c);
    }
    int mid = (l + r) / 2;
    Quad *ldo, *ldi, *rdo, *rdi;
    tie(ldo, ldi) = build_tr(l, mid, p);
    tie(rdi, rdo) = build_tr(mid + 1, r, p);
    while (true) {
        if (left_of(rdi->O, ldi)) {
            ldi = ldi->lnext(); continue; }
        if (right_of(ldi->O, rdi)) {
            rdi = rdi->rev()->onext; continue; }
        break;
    }
    Quad* B = connect(rdi->rev(), ldi);
    auto valid = [&B](Quad* e) {
        return right_of(e->dest(), B); };
    if (ldi->O == ldo->O) ldo = B->rev();
    if (rdi->O == rdo->O) rdo = B;
    while (true) {

```

```

        Quad* lc = B->rev()->onext; // left candidate
        if (valid(lc)) {
            while (in_circle(B->dest(), B->O,
                lc->dest(), lc->onext->dest())) {
                Quad* t = lc->onext;
                delete_edge(lc);
                lc = t;
            }
        }
        Quad* rc = B->oprev(); // right candidate
        if (valid(rc)) {
            while (in_circle(B->dest(), B->O,
                rc->dest(), rc->oprev()->dest())) {
                Quad* t = rc->oprev();
                delete_edge(rc);
                rc = t;
            }
        }
        if (!valid(lc) && !valid(rc)) break;
        if (!valid(lc) || (valid(rc) && in_circle(
            lc->dest(), lc->O, rc->O, rc->dest())))
            B = connect(rc, B->rev());
        else B = connect(B->rev(), lc->rev());
    }
    return make_pair(ldo, rdo);
}

vector<tuple<pt, pt, pt>> delaunay(vector<pt> p) {
    sort(all(p), [](const pt& a, const pt& b) {
        return lt(a.x, b.x) ||
            (eq(a.x, b.x) && lt(a.y, b.y));
    });
    auto res = build_tr(0, sz(p) - 1, p);
    Quad* e = res.first;
    vector<Quad*> edges = {e};
    while (lt(e->onext->dest().cross(e->dest(), e->O), 0))
        e = e->onext;
    auto add = [&p, &e, &edges]() {
        Quad* cur = e;
        do {
            cur->used = true;
            p.pb(cur->O);
            edges.pb(cur->rev());
            cur = cur->lnext();
        } while (cur != e);
    };
    add(); p.clear();

    int kek = 0;
    while (kek < sz(edges))
        if (!(e = edges[kek++])->used) add();
    vector<tuple<pt, pt, pt>> ans;
    for (int i = 0; i < sz(p); i += 3)
        ans.pb(make_tuple(p[i], p[i + 1], p[i + 2]));
    return ans;
}

```

5.6. 3D Primitives.

```

#define P(p) const point3d &p
#define L(p0, p1) P(p0), P(p1)
#define PL(p0, p1, p2) P(p0), P(p1), P(p2)
struct point3d {
    double x, y, z;
    point3d() : x(0), y(0), z(0) {}
    point3d(double _x, double _y, double _z)
        : x(_x), y(_y), z(_z) {}
    point3d operator+(P(p)) const {
        return point3d(x + p.x, y + p.y, z + p.z); }
    point3d operator-(P(p)) const {
        return point3d(x - p.x, y - p.y, z - p.z); }
    point3d operator-() const {
        return point3d(-x, -y, -z); }
    point3d operator*(double k) const {
        return point3d(x * k, y * k, z * k); }
    point3d operator/(double k) const {
        return point3d(x / k, y / k, z / k); }
    double operator%(P(p)) const {
        return x * p.x + y * p.y + z * p.z; }
    point3d operator*(P(p)) const {
        return point3d(y*p.z - z*p.y,
                        z*p.x - x*p.z, x*p.y - y*p.x); }

    double length() const {
        return sqrt(*this % *this); }
    double distTo(P(p)) const {
        return (*this - p).length(); }
    double distTo(P(A), P(B)) const {
        // A and B must be two different points
        return ((*this - A) * (*this - B)).length() /
            ↳ A.distTo(B); }
    point3d normalize(double k = 1) const {
        // length() must not return 0
        return (*this) * (k / length()); }
    point3d getProjection(P(A), P(B)) const {
        point3d v = B - A;
        return A + v.normalize((v % (*this - A)) /
            ↳ v.length()); }
    point3d rotate(P(normal)) const {
        //normal must have length 1 and be orthogonal to
        ↳ the vector
        return (*this) * normal; }
    point3d rotate(double alpha, P(normal)) const {
        return (*this) * cos(alpha) + rotate(normal) *
            ↳ sin(alpha); }
    point3d rotatePoint(P(O), P(axe), double alpha)
        ↳ const {
        point3d Z = axe.normalize(axe % (*this - O));
        return O + Z + (*this - O - Z).rotate(alpha, O);
        ↳ }
    bool isZero() const {
        return abs(x) < EPS && abs(y) < EPS && abs(z) <
            ↳ EPS; }
    bool isOnLine(L(A, B)) const {
        return ((A - *this) * (B - *this)).isZero(); }
    bool isInSegment(L(A, B)) const {

```

```

        return isOnLine(A, B) && ((A - *this) % (B -
            ↳ *this)) < EPS; }
    bool isInSegmentStrictly(L(A, B)) const {
        return isOnLine(A, B) && ((A - *this) % (B -
            ↳ *this)) < -EPS; }
    double getAngle() const {
        return atan2(y, x); }
    double getAngle(P(u)) const {
        return atan2((*this * u).length(), *this % u); }
    bool isOnPlane(PL(A, B, C)) const {
        return
            ↳ abs((A - *this) * (B - *this) % (C - *this)) <
            ↳ EPS; } };
    int line_line_intersect(L(A, B), L(C, D), point3d
        ↳ &O) {
        if (abs((B - A) * (C - A) % (D - A)) > EPS) return
            ↳ 0;
        if ((A - B) * (C - D)).length() < EPS)
            ↳ return A.isOnLine(C, D) ? 2 : 0;
        point3d normal = ((A - B) * (C - B)).normalize();
        double s1 = (C - A) * (D - A) % normal;
        O = A + ((B - A) / (s1 + ((D - B) * (C - B) %
            ↳ normal))) * s1;
        return 1; }
    int line_plane_intersect(L(A, B), PL(C, D, E),
        ↳ point3d &O) {
        double V1 = (C - A) * (D - A) % (E - A);
        double V2 = (D - B) * (C - B) % (E - B);
        if (abs(V1 + V2) < EPS)
            ↳ return A.isOnPlane(C, D, E) ? 2 : 0;
        O = A + ((B - A) / (V1 + V2)) * V1;
        return 1; }
    bool plane_plane_intersect(P(A), P(nA), P(B), P(nB),
        ↳ point3d &P, point3d &Q) {
        point3d n = nA * nB;
        if (n.isZero()) return false;
        point3d v = n * nA;
        P = A + (n * nA) * ((B - A) % nB / (v % nB));
        Q = P + n;
        return true; }

```

5.7. Polygon Centroid.

$$C_x = \frac{1}{6A} \sum_{i=0}^{n-1} (x_i + x_{i+1})(x_i y_{i+1} - x_{i+1} y_i)$$

$$C_y = \frac{1}{6A} \sum_{i=0}^{n-1} (y_i + y_{i+1})(x_i y_{i+1} - x_{i+1} y_i)$$

$$A = \frac{1}{2} \sum_{i=0}^{n-1} (x_i y_{i+1} - x_{i+1} y_i)$$

5.8. Rectilinear Minimum Spanning Tree. Given a set of n points in the plane, and the aim is to find a minimum spanning tree connecting these n points, assuming the Manhattan

distance is used. The function candidates returns at most $4n$ edges that are a superset of the edges in a minimum spanning tree, and then one can use Kruskal's algorithm.

```

#define MAXN 100100
struct RMST {
    struct point {
        int i; ll x, y;
        point() : i(-1) {}
        ll d1() { return x + y; }
        ll d2() { return x - y; }
        ll dist(point other) {
            return abs(x - other.x) + abs(y - other.y); }
        bool operator <(const point &other) const {
            return y==other.y ? x > other.x : y < other.y;
        }
    } best[MAXN], A[MAXN], tmp[MAXN];
    int n;
    RMST() : n(0) {}
    void add_point(int x, int y) {
        A[A[n].i = n].x = x, A[n++].y = y; }
    void rec(int l, int r) {
        if (l >= r) return;
        int m = (l+r)/2;
        rec(l, m), rec(m+1, r);
        point bst;
        for(int i=l, j=m+1, k=1; i <= m || j <= r; k++){
            if(j>r || (i <= m && A[i].d1() < A[j].d1())){
                tmp[k] = A[i++];
                if (bst.i != -1 && (best[tmp[k].i].i == -1
                    ↳ || best[tmp[k].i].d2() < bst.d2()))
                    best[tmp[k].i] = bst;
            } else {
                tmp[k] = A[j++];
                if (bst.i == -1 || bst.d2() < tmp[k].d2())
                    bst = tmp[k]; } }
        rep(i, l, r+1) A[i] = tmp[i]; }
    vector<pair<ll, ii>> candidates() {
        vector<pair<ll, ii>> es;
        REP(p, 2) {
            REP(q, 2) {
                sort(A, A+n);
                REP(i, n) best[i].i = -1;
                rec(0, n-1);
                REP(i, n) {
                    if (best[A[i].i].i != -1)
                        es.pb({A[i].dist(best[A[i].i]),
                            ↳ {A[i].i, best[A[i].i].i}});
                    swap(A[i].x, A[i].y);
                    A[i].x *= -1, A[i].y *= -1; } }
                REP(i, n) A[i].x *= -1; }
            return es; } };

```

5.9. Points and lines (CP3).

```
const ld EPS = 1e-9;
```

```
ld DEG_to_RAD(ld d) { return d*PI/180.0; }
ld RAD_to_DEG(ld r) { return r*180.0/PI; }
```



```

struct point { ld x, y;
  point() { x = y = 0.0; }
  point(ld _x, ld _y) : x(_x), y(_y) {}
  // useful for sorting
  bool operator < (point other) const {
    if (fabs(x - other.x) > EPS)
      return x < other.x;
    return y < other.y; }
  // use EPS (1e-9) when testing for equality
  bool operator == (point other) const {
    return fabs(x-other.x)<EPS &&
      ↪ fabs(y-other.y)<EPS;
  }
};

ld dist(point p1, point p2) {
  // hypot(dx, dy) returns sqrt(dx * dx + dy * dy)
  return hypot(p1.x - p2.x, p1.y - p2.y);
}

// rotate p by rad RADIANS CCW w.r.t origin (0, 0)
point rotate(point p, ld rad) {
  return point(p.x*cos(rad) - p.y*sin(rad),
    p.x*sin(rad) + p.y*cos(rad));
}

// lines are (x,y) s.t. ax + by = c. AND b=0,1.
struct line { ld a, b, c; };

// gives line through p1, p2
line pointsToLine(point p1, point p2) {
  if (fabs(p1.x - p2.x) < EPS) // vertical line
    return { 1.0, 0.0, -p1.x };
  else return {
    -(ld)(p1.y - p2.y) / (p1.x - p2.x),
    1.0,
    -(ld)(1.0 * p1.x) - p1.y;
  };
}

bool areParallel(line l1, line l2) {
  return fabs(l1.a-l2.a)<EPS && fabs(l1.b-l2.b)<EPS;
}

bool areSame(line l1, line l2) {
  return areParallel(l1,l2) && fabs(l1.c-l2.c)<EPS;
}

// returns true (+ intersection) if l1,l2 intersect
bool areIntersect(line l1, line l2, point &p) {
  if (areParallel(l1, l2)) return false; // 0 or inf
  // solve two equations:
  p.x = (l2.b * l1.c - l1.b * l2.c)
    / (l2.a * l1.b - l1.a * l2.b);
  // special case: test for vertical line:
  if (fabs(l1.b) > EPS) p.y = -(l1.a * p.x + l1.c);
  else p.y = -(l2.a * p.x + l2.c);
  return true;
}

```

```

}

// name: `vec' is different from STL vector
struct vec { ld x, y;
  vec(ld _x, ld _y) : x(_x), y(_y) {} };
// convert 2 points to vector a->b
vec toVec(point a, point b) {
  return vec(b.x - a.x, b.y - a.y); }
vec scale(vec v, ld s) { return vec(v.x*s, v.y*s); }
// translate p according to v
point translate(point p, vec v) {
  return point(p.x + v.x, p.y + v.y); }

// convert point and gradient/slope to line
void pointSlopeToLine(point p, ld m, line &l) {
  l.a = -m; // always -m
  l.b = 1; // always 1
  l.c = -(l.a * p.x) + (l.b * p.y); }

void closestPoint(line l, point p, point &ans) {
  if (fabs(l.b) < EPS) { // case 1: vertical line
    ans.x = -(l.c); ans.y = p.y; return; }

  if (fabs(l.a) < EPS) { // case 2: horizontal line
    ans.x = p.x; ans.y = -(l.c); return; }
  // normal line:
  line perpendicular;
  pointSlopeToLine(p, 1 / l.a, perpendicular);
  // intersect line l with this perpendicular line
  // the intersection point is the closest point
  areIntersect(l, perpendicular, ans); }

// returns the reflection of point on a line
void reflectionPoint(line l, point p, point &ans) {
  point b;
  closestPoint(l, p, b); // similar to distToLine
  return point(2*b.x - p.x, 2*b.y - p.y); }

ld dot(vec a, vec b) { return a.x*b.x + a.y*b.y; }
ld cross(vec a,vec b){ return a.x*b.y - a.y*b.x; }
ld norm_sq(vec v) { return v.x*v.x + v.y*v.y; }

// returns the distance from p to the line defined
// by points a and b (a != b), closest point in c.
ld distToLine(point p, point a, point b, point &c) {
  // formula: c = a + u * ab
  vec ap = toVec(a, p), ab = toVec(a, b);
  ld u = dot(ap, ab) / norm_sq(ab);
  c = translate(a, scale(ab, u));
  return dist(p, c); }

// returns the distance from p to the line segment
// ab defined by points a and b (still OK if a == b)
// the closest point is stored in c byref.
ld distToLineSegment(point p, point a, point b,
  ↪ point &c) {
  vec ap = toVec(a, p), ab = toVec(a, b);
  ld u = dot(ap, ab) / norm_sq(ab);
}

```

```

if (u < 0.0) { c = point(a.x, a.y);
  return dist(p, a); } // closer to a
if (u > 1.0) { c = point(b.x, b.y);
  return dist(p, b); } // closer to b
// otherwise closest is perp to line:
return distToLine(p, a, b, c); }

// returns angle aob in rad
ld angle(point a, point o, point b) {
  vec oa = toVec(o, a), ob = toVec(o, b);
  return acos(dot(oa, ob)
    / sqrt(norm_sq(oa) * norm_sq(ob)));
}

// note: to accept collinear points, change `> 0'
// returns true if r is on the left side of line pq
bool ccw(point p, point q, point r) {
  return cross(toVec(p, q), toVec(p, r)) > 0; }

// returns true if r is on the same line as line pq
bool collinear(point p, point q, point r) {
  return fabs(cross(toVec(p,q), toVec(p,r))) < EPS;
}

5.10. Polygon (CP3). Polygons have  $P_0 = P_{n-1}$  here.

typedef vector<point> poly;

// returns the perimeter: sum of Euclidean distances
// of consecutive line segments (polygon edges)
ld perimeter(const poly &P) {
  ld result = 0.0;
  REP(i, sz(P)-1) // remember that P[0] = P[n-1]
    result += dist(P[i], P[i+1]);
  return result; }

// returns the area, which is half the determinant
ld area(const poly &P) {
  ld result = 0.0;
  REP(i, sz(P)-1)
    result += P[i].x*P[i+1].y - P[i+1].x*P[i].y;
  return result;
}

// returns true if we always make the same turn
// throughout the polygon
bool isConvex(const poly &P) {
  int n = sz(P);
  if (n <= 3) return false; // point=2; line=3
  bool isLeft = ccw(P[0], P[1], P[2]);
  rep(i, n-2) if (ccw(P[i], P[i+1],
    P[(i+2) == n ? 1 : i+2]) != isLeft)
    return false; // different sign -> concave
  return true; } // convex

// returns true if pt is in polygon P
bool inPolygon(point pt, const poly &P) {
}

```

```

    if (sz(P) == 0) return false;
    ld sum = 0; // Assume P[0] == P[n-1]
    REP(i, sz(P)-1) {
        if (ccw(pt, P[i], P[i+1]))
            sum += angle(P[i], pt, P[i+1]);
        else sum -= angle(P[i], pt, P[i+1]); }
    return fabs(fabs(sum) - 2*PI) < EPS;
}

// line segment p-q intersect with line A-B.
point lineIntersectSeg(point p, point q,
    point A, point B) {
    ld a = B.y - A.y;
    ld b = A.x - B.x;
    ld c = B.x * A.y - A.x * B.y;
    ld u = fabs(a * p.x + b * p.y + c);
    ld v = fabs(a * q.x + b * q.y + c);
    return point((p.x*v + q.x*u) / (u+v),
        (p.y*v + q.y*u) / (u+v)); }

// cuts polygon Q along the line formed by a -> b
// (note: Q[0] == Q[n-1] is assumed)
poly cutPolygon(point a, point b, const poly &Q) {
    poly P;
    REP(i, sz(Q)) {
        ld left1 = cross(toVec(a,b), toVec(a,Q[i]));
        ld left2 = 0;
        if (i != sz(Q)-1)
            left2 = cross(toVec(a, b), toVec(a, Q[i+1]));
        if (left1 > -EPS)
            P.pb(Q[i]); // Q[i] is left of ab
        if (left1 * left2 < -EPS)
            // edge Q[i]-Q[i+1] crosses line ab
            P.pb(lineIntersectSeg(Q[i], Q[i+1], a, b));
    }
    if (!P.empty() && !(P.back() == P.front()))
        P.pb(P.front()); // make P[0] == P[n-1]
    return P; }

point pivot; // sorts points by angle around pivot
bool angleCmp(point a, point b) {
    if (collinear(pivot, a, b)) // special case
        return dist(pivot, a) < dist(pivot, b);
    ld d1x = a.x - pivot.x, d1y = a.y - pivot.y;
    ld d2x = b.x - pivot.x, d2y = b.y - pivot.y;
    return (atan2(d1y, d1x) - atan2(d2y, d2x)) < 0;
}

poly CH(poly P) { // no order of P assumed!
    int i, j, n = sz(P)
    if (n <= 3) {
        // safeguard from corner case
        if (!(P[0] == P[n-1])) P.pb(P[0]);
        return P; // special case, the CH is P itself
    }

    // P0 = point with lowest Y (if tie rightmost X)
    int P0 = 0;

```

```

    rep(i, 1, n) if (P[i].y < P[P0].y
        || (P[i].y == P[P0].y && P[i].x > P[P0].x))
        P0 = i;
    // swap P[P0] with P[0]:
    point temp = P[0]; P[0] = P[P0]; P[P0] = temp;

    // second, sort points by angle w.r.t. pivot P0
    pivot = P[0];
    sort(++P.begin(), P.end(), angleCmp); // keep P[0]

    // third, the ccw tests
    poly S = { P[n-1], P[0], P[1] }; // initial S
    i = 2; // then, we check the rest
    while (i < n) { // required: N must be >= 3
        j = sz(S) - 1;
        if (ccw(S[j-1], S[j], P[i]))
            S.pb(P[i++]); // left turn, accept
        else // pop top of S when right turn
            S.pop_back();
    }
    return S;
}

5.11. Triangle (CP3).
ld perimeter(point a, point b, point c) {
    return dist(a, b) + dist(b, c) + dist(c, a); }

ld area(ld ab, ld bc, ld ca) {
    // Heron's formula
    ld s = 0.5 * (ab+bc+ca);
    return sqrt(s)*sqrt(s-ab)*sqrt(s-bc)*sqrt(s-ca);
}

ld area(point a, point b, point c) {
    return area(dist(a, b), dist(b, c), dist(c, a));
}

ld rInCircle(ld ab, ld bc, ld ca) {
    return area(ab,bc,ca)*2.0 / (ab+bc+ca);
}

```

```

ld rInCircle(point a, point b, point c) {
    return rInCircle(dist(a,b),dist(b,c),dist(c,a));
}

// assumption: the required points/lines functions
// have been written.
// Returns if there is an inCircle center
// if it returns TRUE, ctr will be the inCircle
// center and r is the same as rInCircle
int inCircle(point p1, point p2, point p3, point
    &ctr, ld &r) {
    r = rInCircle(p1, p2, p3);
    if (fabs(r) < EPS) return false;

    line l1, l2; // compute these two angle bisectors
    ld ratio = dist(p1, p2) / dist(p1, p3);
    point p = translate(p2,
        scale(toVec(p2, p3), ratio / (1 + ratio)));
    pointsToLine(p1, p, l1);

```

```

    ratio = dist(p2, p1) / dist(p2, p3);
    p = translate(p1,
        scale(toVec(p1, p3), ratio / (1 + ratio)));
    pointsToLine(p2, p, l2);
    // get their intersection point:
    areIntersect(l1, l2, ctr);
    return true;
}

ld rCircumCircle(ld ab, ld bc, ld ca) {
    return ab * bc * ca / (4.0 * area(ab, bc, ca)); }

ld rCircumCircle(point a, point b, point c) {
    return rCircumCircle(
        dist(a,b), dist(b,c), dist(c,a));
}

// assumption: the required points/lines functions
// have been written.
// Returns 1 iff there is a circumCenter center
// if this function returns 1, ctr will be the
// circumCircle center and r = rCircumCircle
bool circumCircle(point p1, point p2, point p3,
    &point &ctr, ld &r){
    ld a = p2.x - p1.x, b = p2.y - p1.y;
    ld c = p3.x - p1.x, d = p3.y - p1.y;
    ld e = a * (p1.x + p2.x) + b * (p1.y + p2.y);
    ld f = c * (p1.x + p3.x) + d * (p1.y + p3.y);
    ld g = 2.0 * (a * (p3.y-p2.y) - b * (p3.x-p2.x));
    if (fabs(g) < EPS) return false;

    ctr.x = (d*e - b*f) / g;
    ctr.y = (a*f - c*e) / g;
    r = dist(p1, ctr); // r = dist(center, p_i)
    return true;
}

// returns if pt d is inside the circumCircle
// defined by a,b,c
bool inCircumCircle(point a, point b,
    point c, point d) {
    vec va=toVec(a,d), vb=toVec(b,d), vc=toVec(c,d);
    return 0 <
        va.x * vb.y * (vc.x*vc.x + vc.y*vc.y) +
        va.y * (vb.x*vb.x + vb.y*vb.y) * vc.x +
        (va.x*va.x + va.y*va.y) * vb.x * vc.y -
        (va.x*va.x + va.y*va.y) * vb.y * vc.x -
        va.y * vb.x * (vc.x*vc.x + vc.y*vc.y) -
        va.x * (vb.x*vb.x+vb.y*vb.y) * vc.y;
}

bool canFormTriangle(ld a, ld b, ld c) {
    return a+b > c && a+c > b && b+c > a; }

```

5.12. Circle (CP3).

```

int insideCircle(point_i p, point_i c, int r) { //
↪ all integer version
    int dx = p.x - c.x, dy = p.y - c.y;
    int Euc = dx * dx + dy * dy, rSq = r * r;
    ↪ // all integer
    return Euc < rSq ? 0 : Euc == rSq ? 1 : 2; }
    ↪ //inside/border/outside

bool circle2PtsRad(point p1, point p2, double r,
↪ point &c) {
    double d2 = (p1.x - p2.x) * (p1.x - p2.x) +
                (p1.y - p2.y) * (p1.y - p2.y);
    double det = r * r / d2 - 0.25;
    if (det < 0.0) return false;
    double h = sqrt(det);
    c.x = (p1.x + p2.x) * 0.5 + (p1.y - p2.y) * h;
    c.y = (p1.y + p2.y) * 0.5 + (p2.x - p1.x) * h;
    return true; } // to get the other center,
    ↪ reverse p1 and p2

```

5.13. **Formulas.** Let $a = (a_x, a_y)$ and $b = (b_x, b_y)$ be two-dimensional vectors.

- $a \cdot b = |a||b| \cos \theta$, where θ is the angle between a and b .
- $a \times b = |a||b| \sin \theta$, where θ is the signed angle between a and b .
- $a \times b$ is equal to the area of the parallelogram with two of its sides formed by a and b . Half of that is the area of the triangle formed by a and b .
- **Euler's formula:** $V - E + F = 2$
- Side lengths a, b, c can form a triangle iff. $a + b > c$, $b + c > a$ and $a + c > b$.
- Sum of internal angles of a regular convex n -gon is $(n - 2)\pi$.
- **Law of sines:** $\frac{a}{\sin A} = \frac{b}{\sin B} = \frac{c}{\sin C}$
- **Law of cosines:** $b^2 = a^2 + c^2 - 2ac \cos B$
- Internal tangents of circles $(c_1, r_1), (c_2, r_2)$ intersect at $(c_1 r_2 + c_2 r_1) / (r_1 + r_2)$, external intersect at $(c_1 r_2 - c_2 r_1) / (r_1 + r_2)$.

6. MISCELLANEOUS

6.1. **Binary search** $\mathcal{O}(\log(hi - lo))$.

```

bool test(int n);

int search(int lo, int hi) {
    assert(test(lo) && !test(hi)); // BE CERTAIN
    while (hi - lo > 1) {
        int m = (lo + hi) / 2;
        (test(m) ? lo : hi) = m;
    }
    // assert(test(lo) && !test(hi));
    return lo;
}

```

6.2. **Fast Fourier Transform** $\mathcal{O}(n \log n)$. Given two polynomials $A(x) = a_0 + a_1x + \dots + a_{n/2}x^{n/2}$ and $B(x) = b_0 + b_1x + \dots + b_{n/2}x^{n/2}$, FFT calculates all coefficients of $C(x) = A(x) \cdot B(x) = c_0 + c_1x + \dots + c_nx^n$, with $c_i = \sum_{j=0}^i a_j b_{i-j}$.

```

typedef complex<double> cpx;
const int LOGN = 19, MAXN = 1 << LOGN;

int rev[MAXN];
cpx rt[MAXN], a[MAXN] = {}, b[MAXN] = {};

void fft(cpx *A) {
    REP(i, MAXN) if (i < rev[i]) swap(A[i],
    ↪ A[rev[i]]);
    for (int k = 1; k < MAXN; k *= 2)
        for (int i = 0; i < MAXN; i += 2*k) REP(j, k) {
            cpx t = rt[j + k] * A[i + j + k];
            A[i + j + k] = A[i + j] - t;
            A[i + j] += t;
        }
}

void multiply() { // a = convolution of a * b
    rev[0] = 0; rt[1] = cpx(1, 0);
    REP(i, MAXN) rev[i] = (rev[i/2] | (i&1)<<LOGN)/2;
    for (int k = 2; k < MAXN; k *= 2) {
        cpx z(cos(PI/k), sin(PI/k));
        rep(i, k/2, k) rt[2*i+1]=rt[i]*z;
    }
    fft(a); fft(b);
    REP(i, MAXN) a[i] *= b[i] / (double)MAXN;
    reverse(a+1, a+MAXN); fft(a);
}

```

6.3. **Minimum Assignment (Hungarian Algorithm)** $\mathcal{O}(n^3)$.

```

int a[MAXN + 1][MAXM + 1]; // matrix, 1-based
int minimum_assignment(int n, int m) { // n rows, m
    ↪ columns
    vi u(n + 1), v(m + 1), p(m + 1), way(m + 1);
    for (int i = 1; i <= n; i++) {
        p[0] = i;
        int j0 = 0;
        vi mv(m + 1, INT_MAX);
        vector<char> used(m + 1, false);
        do {
            used[j0] = true;
            int i0 = p[j0], delta = INT_MAX, j1;
            for (int j = 1; j <= m; j++)
                if (!used[j]) {
                    int cur = a[i0][j] - u[i0] - v[j];
                    if (cur < mv[j]) mv[j] = cur, way[j] = j0;
                    if (mv[j] < delta) delta = mv[j], j1 = j;
                }
            for (int j = 0; j <= m; j++) {
                if (used[j]) u[p[j]] += delta, v[j] -= delta;
                else mv[j] -= delta;
            }
        } while (j1 != j0);
        p[j0] = j1;
    }
}

```

```

    }
    j0 = j1;
    } while (p[j0] != 0);
    do {
        int j1 = way[j0]; p[j0] = p[j1]; j0 = j1;
    } while (j0);
}
// column j is assigned to row p[j]
return -v[0];
}

```

6.4. **Partial linear equation solver** $\mathcal{O}(N^3)$.

```

typedef double NUM;
const int ROWS = 200, COLS = 200;
const NUM EPS = 1e-5;

// F2: bitset<COLS+1> M[ROWS]; bitset<ROWS> vals;
NUM M[ROWS][COLS + 1], vals[COLS];
bool hasval[COLS];

bool is0(NUM a) { return -EPS < a && a < EPS; }

// finds x such that Ax = b
// A_ij is M[i][j], b_i is M[i][m]
int solveM(int n, int m) {
    // F2: vals.reset();
    int pr = 0, pc = 0;
    while (pc < m) {
        int r = pr, c;
        while (r < n && is0(M[r][pc])) r++;
        if (r == n) { pc++; continue; }

        // F2: M[pr]^=M[r]; M[r]^=M[pr]; M[pr]^=M[r];
        for (c = 0; c <= m; c++)
            swap(M[pr][c], M[r][c]);

        r = pr++; c = pc++;
        // F2: vals.set(pc, M[pr][m]);
        NUM div = M[r][c];
        for (int col = c; col <= m; col++)
            M[r][col] /= div;
        REP(row, n) {
            if (row == r) continue;
            // F2: if (M[row].test(c)) M[row] ^= M[r];
            NUM times = -M[row][c];
            for (int col = c; col <= m; col++)
                M[row][col] += times * M[r][col];
        }
    } // now M is in RREF

    for (int r = pr; r < n; r++)
        if (!is0(M[r][m])) return 0;
    // F2: return 1;
    fill_n(hasval, n, false);
    for (int col = 0, row; col < m; col++) {
        hasval[col] = !is0(M[row][col]);
        if (!hasval[col]) continue;
        for (int c = col + 1; c < m; c++) {

```

```

    if (!is0(M[row][c])) hasval[col] = false;
}
if (hasval[col]) vals[col] = M[row][m];
row++;
}
REP(i, n) if (!hasval[i]) return 2;
return 1;
}

```

6.5. Cycle-Finding.

```

ii find_cycle(int x0, int (*f)(int)) {
    int t = f(x0), h = f(t), mu = 0, lam = 1;
    while (t != h) t = f(t), h = f(f(h));
    h = x0;
    while (t != h) t = f(t), h = f(h), mu++;
    h = f(t);
    while (t != h) h = f(h), lam++;
    return ii(mu, lam); }

```

6.6. Longest Increasing Subsequence.

```

vi lis(vi arr) {
    vi seq, back(sz(arr)), ans;
    REP(i, sz(arr)) {
        int res = 0, lo = 1, hi = sz(seq);
        while (lo <= hi) {
            int mid = (lo+hi)/2;
            if (arr[seq[mid-1]] < arr[i]) res = mid, lo =
                mid + 1;
            else hi = mid - 1;
        }
        if (res < sz(seq)) seq[res] = i;
        else seq.pb(i);
        back[i] = res == 0 ? -1 : seq[res-1];
    }
    int at = seq.back();
    while (at != -1) ans.pb(at), at = back[at];
    reverse(all(ans));
    return ans;
}

```

6.7. Dates.

```

int intToDay(int jd) { return jd % 7; }
int dateToInt(int y, int m, int d) {
    return 1461 * (y + 4800 + (m - 14) / 12) / 4 +
        367 * (m - 2 - (m - 14) / 12 * 12) / 12 -
        3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +
        d - 32075; }
void intToDate(int jd, int &y, int &m, int &d) {
    int x, n, i, j;
    x = jd + 68569;
    n = 4 * x / 146097;
    x -= (146097 * n + 3) / 4;
    i = (4000 * (x + 1)) / 1461001;
    x -= 1461 * i / 4 - 31;
    j = 80 * x / 2447;
    d = x - 2447 * j / 80;
    x = j / 11;
}

```

```

m = j + 2 - 12 * x;
y = 100 * (n - 49) + i + x; }

```

6.8. Simplex.

```

typedef vector<ld> VD;
typedef vector<VD> VVD;
const ld EPS = 1e-9;
struct LPSolver {
    int m, n; vi B, N; VVD D;
    LPSolver(const VVD &A, const VD &b, const VD &c) :
        m(b.size()), n(c.size()),
        N(n + 1), B(m), D(m + 2, VD(n + 2)) {
        REP(i, m) REP(j, n) D[i][j] = A[i][j];
        REP(i, m) { B[i] = n + i; D[i][n] = -1;
            D[i][n + 1] = b[i]; }
        REP(j, n) N[j] = j, D[m][j] = -c[j];
        N[n] = -1; D[m + 1][n] = 1;
    }
    void Pivot(int r, int s) {
        double inv = 1.0 / D[r][s];
        REP(i, m+2) if (i != r) REP(j, n+2) if (j != s) {
            D[i][j] -= D[r][j] * D[i][s] * inv;
        }
        REP(j, n+2) if (j != s) D[r][j] *= inv;
        REP(i, m+2) if (i != r) D[i][s] *= -inv;
        D[r][s] = inv;
        swap(B[r], N[s]); }
    bool Simplex(int phase) {
        int x = phase == 1 ? m + 1 : m;
        while (true) {
            int s = -1;
            for (int j = 0; j <= n; j++) {
                if (phase == 2 && N[j] == -1) continue;
                if (s == -1 || D[x][j] < D[x][s] ||
                    D[x][j] == D[x][s] && N[j] < N[s]) s = j; }
            if (D[x][s] > -EPS) return true;
            int r = -1;
            REP(i, m) {
                if (D[i][s] < EPS) continue;
                if (r == -1 || D[i][n + 1] / D[i][s] < D[r][n +
                    1] /
                    D[r][s] || (D[i][n + 1] / D[i][s]) ==
                    D[r][n + 1] /
                    D[r][s]) && B[i] < B[r]) r = i; }
            if (r == -1) return false;
            Pivot(r, s); } }
    ld Solve(VD &x) {
        int r = 0;
        rep(i, 1, m) if (D[i][n+1] < D[r][n+1]) r = i;
        if (D[r][n + 1] < -EPS) {
            Pivot(r, n);
            if (!Simplex(1) || D[m + 1][n + 1] < -EPS)
                return -numeric_limits<ld>::infinity();
        }
        REP(i, m) if (B[i] == -1) {
            int s = -1;
            for (int j = 0; j <= n; j++)
                if (s == -1 || D[i][j] < D[i][s] ||
                    D[i][j] == D[i][s] && N[j] < N[s])
                    s = j;
        }
    }
}

```

```

    Pivot(i, s); }
}
if (!Simplex(2)) return
    numeric_limits<ld>::infinity();
x = VD(n);
for (int i = 0; i < m; i++) if (B[i] < n)
    x[B[i]] = D[i][n + 1];
return D[m][n + 1]; } }
// 2-phase simplex solves linear system:
// maximize c^T x
// subject to Ax <= b, x >= 0
// INPUT: A -- an m x n matrix
// b -- an m-dimensional vector
// c -- an n-dimensional vector
// x -- optimal solution (by reference)
// OUTPUT: c^T x (inf. if unbounded above, nan if
// infeasible)
// *** Example ***
// const int m = 4, n = 3;
// ld _A[m][n] = {{6,-1,0}, {-1,-5,0},
// {1,5,1}, {-1,-5,-1}};
// ld _b[m] = {10,-4,5,-5}, _c[n] = {1,-1,0};
// VVD A(m);
// VD b(_b, _b + m), c(_c, _c + n), x;
// REP(i, m) A[i] = VD(_A[i], _A[i] + n);
// LPSolver solver(A, b, c);
// ld value = solver.Solve(x);
// cerr << "VALUE: " << value << endl; // 1.29032
// cerr << "SOLUTION:"; // 1.74194 0.451613 1
// REP(i, sz(x)) cerr << " " << x[i];
// cerr << endl;

```

7. COMBINATORICS

- Catalan numbers (valid bracket seq's of length $2n$):

$$C_0 = 1, C_n = \frac{1}{n+1} \binom{2n}{n} = \sum_{i=0}^{n-1} C_i C_{n-i-1}.$$
- Stirling 1th kind ($\#\pi \in \mathfrak{S}_n$ with exactly k cycles):

$$\left[\begin{smallmatrix} n \\ k \end{smallmatrix} \right] = \left[\begin{smallmatrix} 0 \\ n \end{smallmatrix} \right] = \delta_{0n}, \left[\begin{smallmatrix} n \\ k \end{smallmatrix} \right] = (n-1) \left[\begin{smallmatrix} n-1 \\ k \end{smallmatrix} \right] + \left[\begin{smallmatrix} n-1 \\ k-1 \end{smallmatrix} \right].$$
- Stirling 2nd kind (k -partitions of $[n]$):

$$\left\{ \begin{smallmatrix} n \\ 1 \end{smallmatrix} \right\} = \left\{ \begin{smallmatrix} n \\ n \end{smallmatrix} \right\} = 1, \left\{ \begin{smallmatrix} n \\ k \end{smallmatrix} \right\} = k \left\{ \begin{smallmatrix} n-1 \\ k \end{smallmatrix} \right\} + \left\{ \begin{smallmatrix} n-1 \\ k-1 \end{smallmatrix} \right\}.$$
- Bell numbers (partitions of $[n]$):

$$B_0 = 1, B_n = \sum_{k=0}^{n-1} B_k \binom{n-1}{k} = \sum_{k=0}^n \left\{ \begin{smallmatrix} n \\ k \end{smallmatrix} \right\}.$$
- Euler ($\#\pi \in \mathfrak{S}_n$ with exactly k ascents):

$$\left\langle \begin{smallmatrix} n \\ 0 \end{smallmatrix} \right\rangle = \left\langle \begin{smallmatrix} n \\ n-1 \end{smallmatrix} \right\rangle = 1, \left\langle \begin{smallmatrix} n \\ k \end{smallmatrix} \right\rangle = (k+1) \left\langle \begin{smallmatrix} n-1 \\ k \end{smallmatrix} \right\rangle + (n-k) \left\langle \begin{smallmatrix} n-1 \\ k-1 \end{smallmatrix} \right\rangle.$$
- Euler 2nd order (nr perms of $1, 1, 2, 2, \dots, n, n$ with exactly k ascents):

$$\left\langle\!\left\langle \begin{smallmatrix} n \\ k \end{smallmatrix} \right\rangle\!\right\rangle = (k+1) \left\langle\!\left\langle \begin{smallmatrix} n-1 \\ k \end{smallmatrix} \right\rangle\!\right\rangle + (2n-k-1) \left\langle\!\left\langle \begin{smallmatrix} n-1 \\ k-1 \end{smallmatrix} \right\rangle\!\right\rangle.$$
- Rooted trees: n^{n-1} , unrooted: n^{n-2} .
- Forests of k rooted trees: $\binom{n}{k} k \cdot n^{n-k-1}$.

$$\bullet 1^2 + \dots + n^2 = \frac{n(n+1)(2n+1)}{6}, \quad 1^3 + \dots + n^3 = \frac{n^2(n+1)^2}{4}$$

- $\sum_{i=1}^n \binom{n}{i} F_i = F_{2n}$, $\sum_i \binom{n-i}{i} = F_{n+1}$
- $\sum_{k=0}^n \binom{k}{m} = \binom{n+1}{m+1}$, $x^k = \sum_{i=0}^k i! \left\{ \begin{smallmatrix} k \\ i \end{smallmatrix} \right\} \binom{x}{i} = \sum_{i=0}^k \langle i \rangle \binom{x+i}{k}$
- $a \equiv b \pmod{x, y} \Leftrightarrow a \equiv b \pmod{\text{lcm}(x, y)}$.
- $ac \equiv bc \pmod{m} \Leftrightarrow a \equiv b \pmod{m/\text{gcd}(c, m)}$.
- $\text{gcd}(n^a - 1, n^b - 1) = \text{gcd}(a, b) - 1$.
- **Möbius inversion formula:** If $f(n) = \sum_{d|n} g(d)$, then $g(n) = \sum_{d|n} \mu(d) f(n/d)$. If $f(n) = \sum_{m=1}^n g(\lfloor n/m \rfloor)$, then $g(n) = \sum_{m=1}^n \mu(m) f(\lfloor \frac{n}{m} \rfloor)$.
- **Inclusion-Exclusion:** If $g(T) = \sum_{S \subseteq T} f(S)$, then

$$f(T) = \sum_{S \subseteq T} (-1)^{|T \setminus S|} g(T).$$

Corollary: $b_n = \sum_{k=0}^n \binom{n}{k} a_k \Leftrightarrow a_n = \sum_{k=0}^n (-1)^{n-k} \binom{n}{k} b_k$.

- **The Twelffold Way:** Putting n balls into k boxes. $p(n, k)$ is # partitions of n in k parts, each > 0 . $p_k(n) = \sum_{i=0}^k p(n, k)$.

Balls Boxes	same same	distinct same	same distinct	distinct distinct
-	$p_k(n)$	$\sum_{i=0}^k \left\{ \begin{smallmatrix} n \\ i \end{smallmatrix} \right\}$	$\binom{n+k-1}{k-1}$	k^n
size ≥ 1	$p(n, k)$	$\left\{ \begin{smallmatrix} n \\ k \end{smallmatrix} \right\}$	$\binom{n-1}{k-1}$	$k! \left\{ \begin{smallmatrix} n \\ k \end{smallmatrix} \right\}$
size ≤ 1	$[n \leq k]$	$[n \leq k]$	$\binom{k}{n}$	$n! \binom{k}{n}$

8. FORMULAS

- **Legendre symbol:** $\left(\frac{a}{b}\right) = a^{(b-1)/2} \pmod{b}$, b odd prime.
- **Heron's formula:** A triangle with side lengths a, b, c has area $\sqrt{s(s-a)(s-b)(s-c)}$ where $s = \frac{a+b+c}{2}$.
- **Shoelace formula:** $A = \frac{1}{2} |\sum_{i=0}^{n-1} x_i y_{i+1} - x_{i+1} y_i|$.
- **Pick's theorem:** A polygon on an integer grid strictly containing i lattice points and having b lattice points on the boundary has area $i + \frac{b}{2} - 1$. (Nothing similar in higher dimensions)
- **Absorption probabilities** A random walk on $[0, n]$ with probability p to increase and q to decrease, starting at k has at n absorption probability $\frac{(q/p)^k - 1}{(q/p)^n - 1}$ if $q \neq p$, and k/n if $q = p$.
- A minimum Steiner tree for n vertices requires at most $n - 2$ additional Steiner vertices.
- **Lagrange polynomial** through points $(x_0, y_0), \dots, (x_k, y_k)$ is

$$L(x) = \sum_{j=0}^k y_j \prod_{\substack{0 \leq m \leq k \\ m \neq j}} \frac{x - x_m}{x_j - x_m}.$$

- **Hook length formula:** If λ is a Young diagram and $h_\lambda(i, j)$ is the hook-length of cell (i, j) , then then the number of Young tableaux $d_\lambda = n! / \prod h_\lambda(i, j)$.
- #primitive pythagorean triples with hypotenuse $< n$ approx $n/(2\pi)$.
- **Frobenius Number:** largest number which can't be expressed as a linear combination of numbers a_1, \dots, a_n with non-negative coefficients. $g(a_1, a_2) = a_1 a_2 - a_1 - a_2$, $N(a_1, a_2) = (a_1 - 1)(a_2 - 1)/2$. $g(d \cdot a_1, d \cdot a_2, a_3) = d \cdot g(a_1, a_2, a_3) + a_3(d - 1)$. An integer $x > (\max_i a_i)^2$ can be expressed in such a way iff. $x \mid \text{gcd}(a_1, \dots, a_n)$.
- **Snell's law:** $v_2 \sin \theta_1 = v_1 \sin \theta_2$ gives the shortest path between two media.
- **BEST theorem:** The number of Eulerian cycles in a *directed* graph G is:

$$t_w(G) \prod_{v \in G} (\deg v - 1)!,$$

where $t_w(G)$ is the number of arborescences ("directed spanning" tree) rooted at w : $t_w(G) = \det(q_{ij})_{i,j \neq w}$, with $q_{ij} = [i = j] \text{indeg}(i) - \# \{ (i, j) \in E \}$.

- **Burnside's Lemma:** Let a finite group G act on a set X . Denote $X^g = \{x \in X \mid gx = x\}$. For each g in G let X^g denote the set of elements in X that are fixed by g . Then the number of orbits is:

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|.$$

- **Bézout's identity:** If (x, y) is a solution to $ax + by = d$ (x, y can be found with EGCD), then all solutions are given by

$$(x + k \cdot \text{lcm}(a, b)/a, y - k \cdot \text{lcm}(a, b)/b), \quad k \in \mathbb{Z}$$

9. GAME THEORY

A game can be reduced to Nim if it is a finite impartial game. Nim and its variants include:

- **Nim:** Let $X = \bigoplus_{i=1}^n x_i$, then $(x_i)_{i=1}^n$ is a winning position iff $X \neq 0$. Find a move by picking k such that $x_k > x_k \oplus X$.
- **Misère Nim:** Regular Nim, except that the last player to move *loses*. Play regular Nim until there is only one pile of size larger than 1, reduce it to 0 or 1 such that there is an odd number of piles. The second player wins (a_1, \dots, a_n) if 1) there is a pile $a_i > 1$ and $\bigoplus_{i=1}^n a_i = 0$ or 2) all $a_i \leq 1$ and $\bigoplus_{i=1}^n a_i = 1$.
- **Staircase Nim:** Stones are moved down a staircase and only removed from the last pile. $(x_i)_{i=1}^n$ is an L -position if $(x_{2i-1})_{i=1}^{n/2}$ is (i.e. only look at odd-numbered piles).

- **Moore's Nim_k:** The player may remove from at most k piles (Nim = Nim₁). Expand the piles in base 2, do a carry-less addition in base $k+1$ (i.e. the number of ones in each column should be divisible by $k+1$).
- **Dim⁺:** The number of removed stones must be a divisor of the pile size. The Sprague-Grundy function is $k+1$ where 2^k is the largest power of 2 dividing the pile size.
- **Aliquot game:** Same as above, except the divisor should be proper (hence 1 is also a terminal state, but watch out for size 0 piles). Now the Sprague-Grundy function is just k .
- **Nim (at most half):** Write $n+1 = 2^m y$ with m maximal, then the Sprague-Grundy function of n is $(y-1)/2$.
- **Lasker's Nim:** Players may alternatively split a pile into two new non-empty piles. $g(4k+1) = 4k+1$, $g(4k+2) = 4k+2$, $g(4k+3) = 4k+4$, $g(4k+4) = 4k+3$ ($k \geq 0$).
- **Hackenbush on trees:** A tree with stalks $(x_i)_{i=1}^n$ may be replaced with a single stalk with length $\bigoplus_{i=1}^n x_i$.

10. SCHEDULING THEORY

Let p_j be the time task j takes on a machine, d_j the deadline, C_j the time it is completed, $L_j = C_j - d_j$ the lateness, $T_j = \max(L_j, 0)$ the tardiness, $U_j = 1$ iff $T_j > 0$ and else 0.

- One machine, minimise L_{\max} : do the tasks in increasing deadline
- One machine, minimise $\sum_j w_j C_j$: do the task increasing in p_j/w_j
- One machine, minimise $\sum_{j=1}^n C_j$ under the condition that all tasks can be done on time:
 - (1) Initialise $k = n, \tau = \sum_j p_j, J = [n]$
 - (2) Take $i_k \in J$ with $d_{i_k} \geq \tau$ and $p_{i_k} \geq p_\ell$ for $\ell \in J$ with $d_\ell \geq \tau$
 - (3) $\tau := \tau - p_{i_k}, k := k - 1, J := J - \{i_k\}$. If $k \neq 0$, go to step 2.
 - (4) The optimale schedule is i_1, \dots, i_n .
- One machine, minimise $\sum_j U_j$. Add all tasks in order of increasing deadline; if adding a task makes it contrary with its deadline, remove the processed task with the highest processing time.
- Two machines (all tasks have to be done on both machines, in any order), minimise C_{\max} : a greedy algorithm, when a machine is free it picks a task that hasn't been done yet on either machine and has longest processing time on the other machine.
- Two machines (all tasks have to be done first on machine 1, then machine 2), minimise C_{\max} . There is an optimal schedule with on both machines the same order of tasks. Take

- $X = \{j : p_{1j} \leq p_{2j}\}$ and Y the complement. Sort X increasing in p_{1j} and Y decreasing in p_{2j} . Then X, Y is an optimal schedule.
- Two machines (all tasks have to be done first on machine 1, then on 2, or vice versa), minimise C_{\max} : let J_{12} be the tasks that have to be done first on machine 1, then on 2 and similar J_{21} . Use the above algorithm to find S_{12}, S_{21} optimal for J_{12}, J_{21} . Then optimal is S_{12}, S_{21} for M1 and S_{21}, S_{12} for M2. (If there are tasks that have to be done on only one machine, do them in the middle.)

11. DEBUGGING TIPS

- Stack overflow? Recursive DFS on tree that is actually a long path?
- Floating-point numbers
 - Getting NaN? Make sure acos etc. are not getting values out of their range (perhaps $1+\epsilon$).
 - Rounding negative numbers?
 - Outputting in scientific notation?
- Wrong Answer?
 - Read the problem statement again!
 - Are multiple test cases being handled correctly? Try repeating the same test case many times.
 - Integer overflow?
 - Think very carefully about boundaries of all input parameters
 - Try out possible edge cases:
 - * $n = 0, n = -1, n = 1, n = 2^{31} - 1$ or $n = -2^{31}$
 - * List is empty, or contains a single element
 - * n is even, n is odd
 - * Graph is empty, or contains a single vertex
 - * Graph is a multigraph (loops or multiple edges)
 - * Polygon is concave or non-simple
 - Is initial condition wrong for small cases?
 - Are you sure the algorithm is correct?
 - Explain your solution to someone.
 - Are you using any functions that you don't completely understand? Maybe STL functions?
 - Maybe you (or someone else) should rewrite the solution?
 - Can the input line be empty?
- Run-Time Error?
 - Is it actually Memory Limit Exceeded?

11.1. Dynamic programming optimizations.

- Convex Hull
 - $dp[i] = \min_{j < i} \{dp[j] + b[j] \times a[i]\}$
 - $b[j] \geq b[j + 1]$

- optionally $a[i] \leq a[i + 1]$
 - $O(n^2)$ to $O(n)$ (see 2.12).
 - Divide & Conquer
 - $dp[i][j] = \min_{k < j} \{dp[i - 1][k] + C[k][j]\}$
 - $A[i][j] \leq A[i][j + 1]$
 - sufficient:

$$C[a][c] + C[b][d] \leq C[a][d] + C[b][c], (a \leq b \leq c \leq d) \quad (\text{QI})$$
 - $O(kn^2)$ to $O(kn \log n)$
- ```
vvi A; // A[i][j] is voor [i, j]

void divco(ll ls, ll rs, ll lt, ll rt, vi &t, vi
 ↪ &s) {
 // berekent t[_][lt, rt)
 if (lt >= rt) return;
 ll ms = ls, mt = (lt + rt) / 2;
 t[mt] = -INF;
 rep(i, ls, rs) {
 if (i >= mt) break;
 if (s[i] + A[i][mt] > t[mt]) {
 t[mt] = s[i] + A[i][mt];
 ms = i;
 }
 }
 divco(ls, ms + 1, lt, mt, t, s);
 divco(ms, rs, mt + 1, rt, t, s);
}

• Knuth
 - $dp[i][j] = \min_{i < k < j} \{dp[i][k] + dp[k][j] + C[i][j]\}$
 - $A[i][j - 1] \leq A[i][j] \leq A[i + 1][j]$
 - $O(n^3)$ to $O(n^2)$
 - sufficient: QI and $C[b][c] \leq C[a][d], a \leq b \leq c \leq d$
```

### 11.2. Solution Ideas.

- Dynamic Programming
  - Parsing CFGs: CYK Algorithm
  - Drop a parameter, recover from others
  - Swap answer and a parameter
  - When grouping: try splitting in two
  - $2^k$  trick
- Greedy
- Randomized
- Optimizations
  - Use bitset (/64)
  - Switch order of loops (cache locality)
- Process queries offline
  - Mo's algorithm
- Square-root decomposition
- Precomputation
- Efficient simulation
  - Mo's algorithm

- Sqrt decomposition
- Store  $2^k$  jump pointers
- Data structure techniques
  - Sqrt buckets
  - Store  $2^k$  jump pointers
  - $2^k$  merging trick
- Counting
  - Inclusion-exclusion principle
  - Generating functions
- Graphs
  - Can we model the problem as a graph?
  - Can we use any properties of the graph?
  - Strongly connected components
  - Cycles (or odd cycles)
  - Bipartite (no odd cycles)
    - \* Bipartite matching
    - \* Hall's marriage theorem
    - \* Stable Marriage
  - Cut vertex/bridge
  - Biconnected components
  - Degrees of vertices (odd/even)
  - Trees
    - \* Heavy-light decomposition
    - \* Centroid decomposition
    - \* Least common ancestor
    - \* Centers of the tree
  - Eulerian path/circuit
  - Chinese postman problem
  - Topological sort
  - (Min-Cost) Max Flow
  - Min Cut
    - \* Maximum Density Subgraph
  - Huffman Coding
  - Min-Cost Arborescence
  - Steiner Tree
  - Kirchoff's matrix tree theorem
  - Prüfer sequences
  - Lovász Toggle
  - Look at the DFS tree (which has no cross-edges)
  - Is the graph a DFA or NFA?
    - \* Is it the Synchronizing word problem?
- math
  - Is the function multiplicative?
  - Look for a pattern
  - Permutations
    - \* Consider the cycles of the permutation
  - Functions

- \* Sum of piecewise-linear functions is a piecewise-linear function
- \* Sum of convex (concave) functions is convex (concave)
- Modular arithmetic
  - \* Chinese Remainder Theorem
  - \* Linear Congruence
- Sieve
- System of linear equations
- Values too big to represent?
  - \* Compute using the logarithm
  - \* Divide everything by some large value
- Linear programming
  - \* Is the dual problem easier to solve?
- Can the problem be modeled as a different combinatorial problem? Does that simplify calculations?
- Logic
  - 2-SAT
  - XOR-SAT (Gauss elimination or Bipartite matching)
- Meet in the middle
- Only work with the smaller half ( $\log(n)$ )
- Strings
  - Trie (maybe over something weird, like bits)
  - Suffix array
  - Suffix automaton (+DP?)
  - Aho-Corasick
  - eerTree
  - Work with  $S + S$
- Hashing
- Euler tour, tree to array
- Segment trees
  - Lazy propagation
  - Persistent
  - Implicit
  - Segment tree of X
- Geometry
  - Minkowski sum (of convex sets)
  - Rotating calipers
  - Sweep line (horizontally or vertically?)
  - Sweep angle
  - Convex hull
- Fix a parameter (possibly the answer).
- Are there few distinct values?
- Binary search
- Sliding Window (+ Monotonic Queue)
- Computing a Convolution? Fast Fourier Transform
- Computing a 2D Convolution? FFT on each row, and then on each column

- Exact Cover (+ Algorithm X)
- Cycle-Finding
- What is the smallest set of values that identify the solution? The cycle structure of the permutation? The powers of primes in the factorization?
- Look at the complement problem
  - Minimize something instead of maximizing
- Immediately enforce necessary conditions. (All values greater than 0? Initialize them all to 1)
- Add large constant to negative numbers to make them positive
- Counting/Bucket sort

## PRACTICE CONTEST CHECKLIST

- How many operations per second? Compare to local machine.
- What is the stack size?
- How to use printf/scanf with long long/long double?
- Are `__int128` and `__float128` available?
- Does MLE give RTE or MLE as a verdict? What about stack overflow?
- What is `RAND_MAX`?
- How does the judge handle extra spaces (or missing newlines) in the output?
- Look at documentation for programming languages.
- Try different programming languages: C++, Java and Python.
- Try the submit script.
- Try local programs: `i?python[23]`, `factor`.
- Try submitting with `assert(false)` and `assert(true)`.
- Omitting `return 0;` still works?
- Look for directory with sample test cases.
- Make sure printing works.