Intérprete con estrategia Call By Name (CBN)





Machete de interpretación y semántica denotacional



$\overline{\Gamma,x} = \widehat{\langle M,\Gamma'\rangle}, \underline{\Lambda} \vdash x \hookrightarrow \overline{V} \quad x \notin \mathsf{D}(\Delta)$ $\Gamma' \vdash M \hookrightarrow V$

$$\begin{array}{c|c} \hline \Gamma \vdash M \hookrightarrow \langle x, M', \Gamma' \rangle & \Gamma', x = \langle N, \Gamma \rangle \vdash M' \hookrightarrow V \\ \hline \hline \Gamma \vdash \lambda x. M \hookrightarrow \langle x, M, \Gamma \rangle \\ \hline \end{array}$$

 $\Gamma \vdash MN \hookrightarrow V$

$$\begin{array}{c} \mathbf{1} \ \ \, \Gamma \\ \Gamma \vdash M \hookrightarrow \mathsf{True} \quad \Gamma \vdash N_1 \hookrightarrow V \end{array}$$

$$\Gamma \vdash M \hookrightarrow_{\mathsf{False}} \Gamma \vdash N_2 \hookrightarrow V$$

$$: \mathcal{F} M \hookrightarrow_{\mathsf{False}} N_1 \hookrightarrow_{\mathsf{False}} N_2 \hookrightarrow_{\mathsf{False}}$$

$$\overline{\Gamma} \vdash$$
 if M then N_1 else $N_2 \hookrightarrow \overline{V}$ $\overline{\Gamma} \vdash$ $\overline{\Gamma}$

$$\begin{array}{c} \Gamma \vdash M \hookrightarrow \text{False } \Gamma \vdash N_2 \hookrightarrow V \\ \hline \Gamma \vdash \text{if } M \text{ then } N_1 \text{ else } N_2 \hookrightarrow V \\ \hline \end{array}$$

$$\Gamma, x = \langle \mu x. M, \Gamma \rangle \vdash M \hookrightarrow V$$

$$x = \langle \mu x. M, \Gamma' \rangle \vdash M$$
$$\Gamma \vdash \mu x. M \hookrightarrow V$$

Intérprete con estrategia Call By Value (CBV)



CBN y CBV

Extensión de los intérpretes con números naturales

 $\Gamma \vdash N \hookrightarrow W \quad \Gamma \vdash M \hookrightarrow \langle x, M', \Gamma' \rangle \quad \Gamma', x = W \vdash M' \hookrightarrow V$

 $\Gamma \vdash MN \hookrightarrow V$

 $\overline{\Gamma,x} = \langle \mu y.M, \Gamma' \rangle, \Delta \vdash x \hookrightarrow \overline{V} \ x \notin \mathsf{D}(\Delta)$

 $\overline{\Gamma, x = V, \Delta \vdash x \hookrightarrow V} \ x \notin \mathsf{D}(\Delta)$

 $\Gamma' \vdash \mu y.M \to V$

$$\Gamma \vdash M \hookrightarrow \text{zero}$$

 $\Gamma \vdash \text{pred}(M) \hookrightarrow \text{zero}$

$$M\hookrightarrow \text{succ}(V)$$

$$\Gamma \vdash isZero(M) \hookrightarrow True$$

Γ ⊢ M → zero

$$\Gamma \vdash M \hookrightarrow succ(V)$$

 $\Gamma \vdash isZero(M) \hookrightarrow False$

 $\Gamma \vdash \operatorname{succ}(M) \hookrightarrow \operatorname{succ}(V)$ **V** → M → J

> $\overline{\Gamma} dash$ if M then N_1 else $N_2 \hookrightarrow \overline{V}$ $\Gamma \vdash M \hookrightarrow$ False $\Gamma \vdash N_2 \hookrightarrow V$

 $\overline{\Gamma \vdash \text{if } M \text{ then } N_1 \text{ else } N_2 \hookrightarrow V} \qquad \overline{\Gamma \vdash \text{if } M \text{ then } \Gamma, x = \langle \mu x. M, \Gamma \rangle \vdash M \hookrightarrow V}$

 $\Gamma \vdash M \hookrightarrow \mathsf{True} \ \Gamma \vdash N_1 \hookrightarrow V$

 $\Gamma \vdash \lambda x.M \hookrightarrow \langle x, M, \Gamma \rangle$

 $\Gamma \vdash \mu x.M \hookrightarrow V$

 $\Gamma \vdash M \hookrightarrow succ(V)$ Γ ⊢ pred(M)
 ∨ V

Semántica denotacional del Cálculo Lambda (sin error) 🥠

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