

*CUDD*

# Colorado University Decision Diagram Package

Systems Design Laboratory (2021/2022)

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# Outline

- 1 Introduction
- 2 Basic Architecture
- 3 Basic Functions
- 4 Example: Half-Adder
- 5 Variable reordering
- 6 Converting BDDs to ZDDs and Vice Versa

- CUDD is the **C**olorado **U**niversity **D**ecision **D**iagram package.
- It is a C/C++ library for creating the following different types of decision diagrams:
  - ▶ BDD: **B**inary **D**ecision **D**iagrams.
  - ▶ ZDD: **Z**ero-Suppressed **BDD**s.
  - ▶ ADD: **A**lgebraic **D**ecision **D**iagrams.
- The slide, source codes, and all materials related to this lesson are available here:  
<https://github.com/luigicapogrosso/SDL>

# Getting CUDD

- The CUDD package is **available via anonymous FTP** from `vlsi.colorado.edu`.
- You can download the CUDD package from the server using an FTP client such as FileZilla or you can use the `ftp` command from the command line.
- Alternatively, you can download the latest version of CUDD **directly from the SDL GitHub repository**, so:

```
$ git clone  
https://github.com/luigicapogrosso/SDL.git
```

# Getting CUDD (cont'd)

- The library is tested using GCC (9.4.0) and GNU Make (9.4.0). To build the library from sources in a clean way, it is preferable that you set up a build subdirectory, say:

```
$ cd SDL/lecture_01/cudd-3.0.0
$ export CUDD_INSTALL_DIRECTORY=$HOME/<path>
$ mkdir objdir && cd objdir
$ ../configure --prefix=$CUDD_INSTALL_DIRECTORY
$ make && make install
```

# Including and linking the CUDD library

- To build an application that uses the CUDD package, you should **add, in your source codes, the following line:**

- ▶ `#include "cudd.h"`
- ▶ `#include "util.h"`

- To **compile** and link a C program that uses CUDD:

```
$ gcc -o main main.c -lcudd -lutil
```

- Or, you can refer to the following **Makefile**:

```
https://github.com/luigicapogrosso/SDL/blob/master/lecture\_01/code/Makefile
```

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# Garbage Collection

- CUDD has a **built-in garbage collection** system.
- When a BDD is not used anymore, its memory can be reclaimed.
- To facilitate the garbage collector, **we need to “reference” and “dereference”** each node in our BDD:
  - ▶ `Cudd_Ref (DdNode*)` to reference a node.
  - ▶ `Cudd_RecursiveDeref (DdNode*)` to dereference a node and all its descendants.



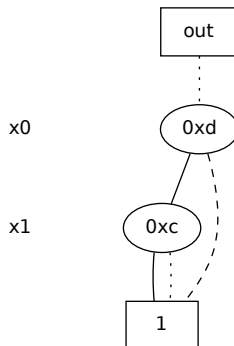
# Complemented arcs

- Each node of a BDD can be:
  - ▶ **A variable** with two children.
  - ▶ **A leaf** with a constant value.
- The two children of a node are referred to as the “*then*” child and the “*else*” child.
- To assign a value to a BDD, we follow “*then*” and “*else*” children until we reach a leaf:
  - ▶ **The value of our assignment is the value of the leaf we reach.**
- **However:** “*else*” children can be **complemented**:
  - ▶ When an “*else*” child is complemented, then we take the **complement of the value of the leaf**:
    - ★ *i.e.*, if the value of the leaf is 1 and we have traversed an odd number of complement arcs, the value of our assignment is 0.

# Complemented arcs: example

- $out = x_0\bar{x}_1$
- “*then*” arcs are solid.
- Normal “*else*” arcs are dashed.
- Complemented “*else*” arcs are dotted.
- The  $out$  arc is complemented:

$$\begin{aligned}\overline{out} &= \bar{x}_0 + x_1 \\ &= \bar{x}_0 + x_0x_1\end{aligned}$$



# The DdManager

- The `DdManager` is the **key data structure of CUDD**:
  - ▶ It must be created before calling any other CUDD function.
  - ▶ It needs to be passed to almost every CUDD function.
- To initialize the `DdManager`, we use the following function:

```
DdManager * Cudd_Init(  
    unsigned int numVars,      // initial number of BDD variables (i.e., subtables)  
    unsigned int numVarsZ,    // initial number of ZDD variables (i.e., subtables)  
    unsigned int numSlots,    // initial size of the unique tables  
    unsigned int cacheSize,   // initial size of the cache  
    unsigned long maxMemory    // target maximum memory occupation.(0 means unlimited)  
);
```

# The DdManager: C code

```
#include<stdio.h>
#include"cudd.h"

int main()
{
    DdManager* manager = Cudd_Init(0, 0,
        CUDD_UNIQUE_SLOTS, CUDD_CACHE_SLOTS, 0);
    if(manager == NULL)
    {
        printf("Error when initializing CUDD.\n");
        return 1;
    }
    ...

    return 0;
}
```

# The DdNode

- The DdNode is the **core building block of BDDs**:

```
struct DdNode {  
    DdHalfWord index;      // Index of the variable represented by this node  
    DdHalfWord ref;        // reference count  
    DdNode *next;          // next pointer for unique table  
    union {  
        CUDD_VALUE_TYPE value; // for constant nodes  
        DdChildren kids;       // for internal nodes  
    } type;  
};
```

- `index` is a unique index for the variable represented by this node.
  - ▶ It is permanent: if we reorder variables, the `idx` remains the same.
- `ref` stores the reference count for this node.
  - ▶ It is incremented by `Cudd_Ref()` and decremented by `Cudd_Recursive_Deref()`.

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- **Common manipulations** of BDDs can be accomplished by **calling operators on variables**.
- The CUDD package includes Boolean functions that can be used for BDD operations such as: *NOT*, *AND*, *NAND*, *OR*, *NOR*, *Exclusive-OR*, *XNOR*, and etc.

# BDD for the NOT Boolean function

- For the **NOT Boolean function**, we use `Cudd_Not()`.
- The truth table for a NOT:

$x_1$	$f$
0	1
1	0

- Exercise: write the code to build the BDD for the function  $f = \neg x_1$ .



# BDD for the AND Boolean function

- For the **AND Boolean function**, we use `Cudd_bddAnd()`.
- The truth table for an AND:

$x_1$	$x_2$	$f$
0	0	0
0	1	0
1	0	0
1	1	1

- Exercise: write the code to build the BDD for the function  $f = x_1 \wedge x_2$ .

# BDD for the NAND Boolean function

- For the **NAND Boolean function**, we use `Cudd_bddNand()`.
- The truth table for a NAND:

$x_1$	$x_2$	$f$
0	0	1
0	1	1
1	0	1
1	1	0

- Exercise: write the code to build the BDD for the function  $f = \neg(x_1 \wedge x_2)$ .

# BDD for the OR Boolean function

- For the **OR Boolean function**, we use `Cudd_bddOr()`.
- The truth table for a logic OR:

$x_1$	$x_2$	$f$
0	0	0
0	1	1
1	0	1
1	1	1

- Exercise: write the code to build the BDD for the function  $f = x_1 \vee x_2$ .

# BDD for the NOR Boolean function

- For the **NOR Boolean function**, we use `Cudd_bddNor()`.
- The truth table for a NOR:

$x_1$	$x_2$	$f$
0	0	1
0	1	0
1	0	0
1	1	0

- Exercise: write the code to build the BDD for the function  $f = \neg(x_1 \vee x_2)$ .

# BDD for Exclusive-OR Boolean function

- For the **Exclusive-OR Boolean function**, we use `Cudd_bddXor()`.
- The truth table for an Exclusive-OR:

$x_1$	$x_2$	$f$
0	0	0
0	1	1
1	0	1
1	1	0

- Exercise: write the code to build the BDD for the function  $f = x_1 \oplus x_2$ .

# BDD for the XNOR Boolean function

- For the **XNOR Boolean function**, we use `Cudd_bddXnor()`.
- The truth table for an XNOR:

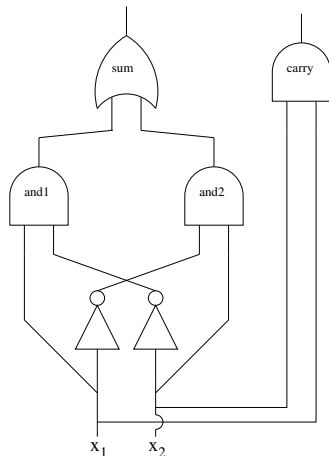
$x_1$	$x_2$	$f$
0	0	1
0	1	0
1	0	0
1	1	1

- Exercise: write the code to build the BDD for the function  $f = \neg(x_1 \oplus x_2)$ .

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# The Half-Adder circuit



This is the schematic of a **half-adder circuit** that we want to compile into an OBDD. It has the following truth table:

$x_1$	$x_2$	sum	carry
0	0	0	0
0	1	1	0
1	0	1	0
1	1	0	1



# Create the BDD for `sum`

```
DdNode* x1 = Cudd_bddIthVar(manager, 0);
DdNode* x2 = Cudd_bddIthVar(manager, 1);

DdNode* and1;
and1 = Cudd_bddAnd(manager, x1, Cudd_Not(x2));
Cudd_Ref(and1);

DdNode* and2;
and2 = Cudd_bddAnd(manager, Cudd_Not(x1), x2);
Cudd_Ref(and2);

DdNode* sum;
sum = Cudd_bddOr(manager, and1, and2);
Cudd_Ref(sum);

Cudd_RecursiveDeref(manager, and1);
Cudd_RecursiveDeref(manager, and2);
```

- **Exercise:** write the code for `carry`.

# Restricting the BDD

- **Restricting a BDD** means **assigning a truth value to some of the variables**.
- The `Cudd_bddRestrict()` function returns the restricted BDD.

```
DdNode * Cudd_bddRestrict(  
    DdManager * manager,    // DD manager  
    DdNode * BDD,           // The BDD to restrict  
    DdNode * restrictBy)    // The BDD to restrict by.
```

- BDD is the original BDD to restrict.
- `restrictBy` is the truth assignment of the variables.

# Print the truth table

```
DdNode *restrictBy;
restrictBy = Cudd_bddAnd(manager, x1, Cudd_Not(x2));
Cudd_Ref(restrictBy);

DdNode *testSum;
testSum = Cudd_bddRestrict(manager, sum, restrictBy);
Cudd_Ref(testSum);
DdNode *testCarry;
testCarry = Cudd_bddRestrict(manager, carry, restrictBy);
Cudd_Ref(testCarry);

printf("x1 = 1, x2 = 0: sum = %d, carry = %d\n",
       1 - Cudd_IsComplement(testSum),
       1 - Cudd_IsComplement(testCarry));

Cudd_RecursiveDeref(manager, restrictBy);
Cudd_RecursiveDeref(manager, testSum);
Cudd_RecursiveDeref(manager, testCarry);
```

- Exercise: Write the code for the complete truth table.

# Print the BDD with `graphviz`

- The function `Cudd_DumpDot()` dumps the BDD to a file in **GraphViz** format.
- The `.dot` file can be converted to a PDF by the command `dot:`

```
$ dot -O -Tpdf half_adder.dot
```

# Print the BDD: C code

```
char *inputNames[2];
inputNames[0] = "x1";
inputNames[1] = "x2";
char *outputNames[2];
outputNames[0] = "sum";
outputNames[1] = "carry";

DdNode *outputs[2];
outputs[0] = sum;
Cudd_Ref(outputs[0]);
outputs[1] = carry;
Cudd_Ref(outputs[1]);

FILE *f = fopen("half_adder.dot", "w");

Cudd_DumpDot(manager, 2, outputs, inputNames, outputNames, f);

Cudd_RecursiveDeref(manager, outputs[0]);
Cudd_RecursiveDeref(manager, outputs[1]);
fclose(f);
```

# Variable reordering

- The **order** of variables can have a **tremendous effects on the size of BDDs**.
- CUDD provides a rich set of tools for reordering BDDs:
  - ▶ Automatic reordering (using heuristics) when the number of nodes in the BDD passes a certain threshold.
  - ▶ Manual reordering using different heuristics.
  - ▶ Manual reordering with a user-specified variable order.
- The function `Cudd_ShuffleHeap()` is used to define the variable order:

```
int Cudd_ShuffleHeap(  
    DdManager * manager,    // DD manager  
    int * permutation        // required variable permutation  
)
```

# Exercise: play with the variable order!

- Create the BDD for the function  $x_1x_2 + x_3x_4 + x_5x_6$ .
- Try the following variable orders and compare the results:
  - ▶  $x_1 < x_2 < x_3 < x_4 < x_5 < x_6$
  - ▶  $x_1 < x_3 < x_5 < x_2 < x_4 < x_6$

## HINTS

- `int Cudd_ReadPerm(manager, x2->index)` returns the position of variable `x2` in the order.
- `int Cudd_ReadNodeCount(manager)` returns the number of nodes in the BDD.

# Converting BDDs to ZDDs

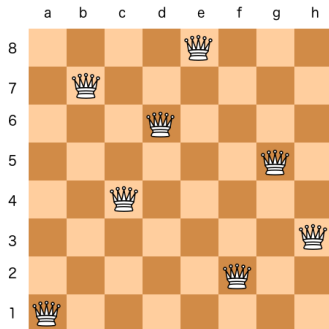
- Many applications **first build a set of BDDs and then derive ZDDs** from the BDDs.
- These applications should create the manager with 0 ZDD variables and create the BDDs.
- Then they should call `Cudd_zddVarsFromBddVars()` to create the necessary ZDD variables—whose number is likely to be known once the BDDs are available.



# Converting BDDs to ZDDs (cont'd)

- The simplest conversion from BDDs to ZDDs is a simple change of representation, which preserves the functions.
- Simply put, given a BDD for  $f$ , a ZDD for  $f$  is requested. In this case the correspondence between the BDD variables and ZDD variables are *one-to-one*.
- Hence, `Cudd_zddVarsFromBddVars()` should be called with the *multiplicity* parameter equal to 1.
- The **conversion can then be performed by calling:**  
`Cudd_zddPortFromBdd()`.
- The **inverse transformation is performed by calling:**  
`Cudd_zddPortToBdd()`.

# The N-Queens problem



- The **N-Queens problem** is the problem of placing  $N$  non-attacking queens on an  $N \times N$  chessboard.
- Our implementation of these benchmarks are based on the description of Kunkle10. We construct an BDD row-by-row that represents whether the row is in a legal state.
- On the accumulated BDD we then count the number of satisfying assignments.