



Topic 4

Assembly Programming - Procedure Calling

Procedure (Function) Calling

- Used to improve reusability and manageability
- Steps of procedure operation
 1. Place parameters in parameter registers
 2. Transfer control to procedure
 3. Acquire storage for procedure in stack
 4. Perform procedure's operations
 5. Place results in result register(s) for caller
 6. Release storage
 7. Return to the place before procedure calling

Procedure Call Instructions

- Procedure call operations: jump and link

`jal ProcedureLabel` (J-type)

- $\$ra = PC+4$; Address of following instruction put in $\$ra$
- $PC = PC_{31...28} : 26\text{-bit address} : 00$
- Jumps to target address

- Procedure return operations: jump register

`jr $ra` (R-type)

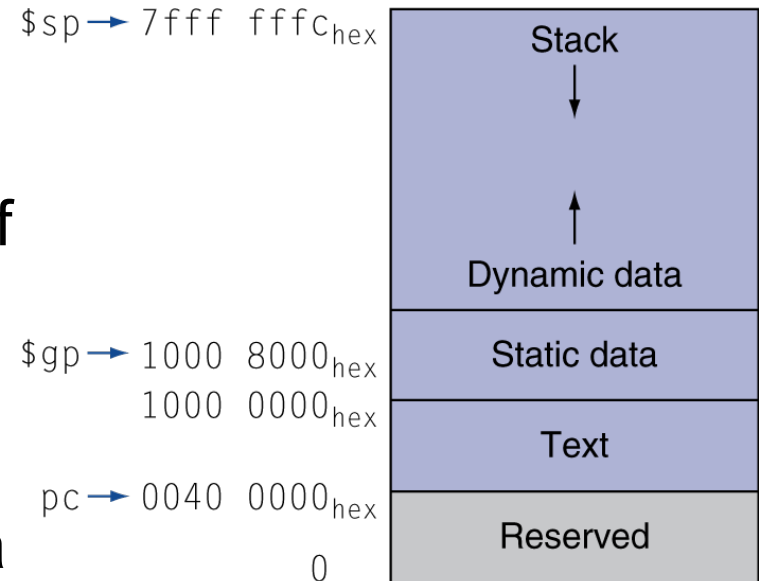
- $PC = \$ra$; Copies $\$ra$ to program counter
- Can also be used for computed jumps (to any other register)
 - e.g., for case/switch statements

Register Usage

- \$zero: constant 0 (reg 0)
- \$at: Assembler Temporary (reg 1)
- \$v0, \$v1: result values (reg's 2 and 3)
- \$a0 – \$a3: arguments (reg's 4 – 7)
- \$t0 – \$t9: temporaries (reg's 8 – 15)
 - Can be overwritten by callee
- \$s0 – \$s7: saved (reg's 16 – 23)
 - Must be saved/restored by callee
- \$t8, \$t9: temporaries (reg's 24 and 25)
- \$k0, \$k1: reserved for OS kernel (reg's 26 and 27)
- \$gp: global pointer for static data (reg 28)
- \$sp: stack pointer (reg 29)
- \$fp: frame pointer (reg 30)
- \$ra: return address (reg 31)

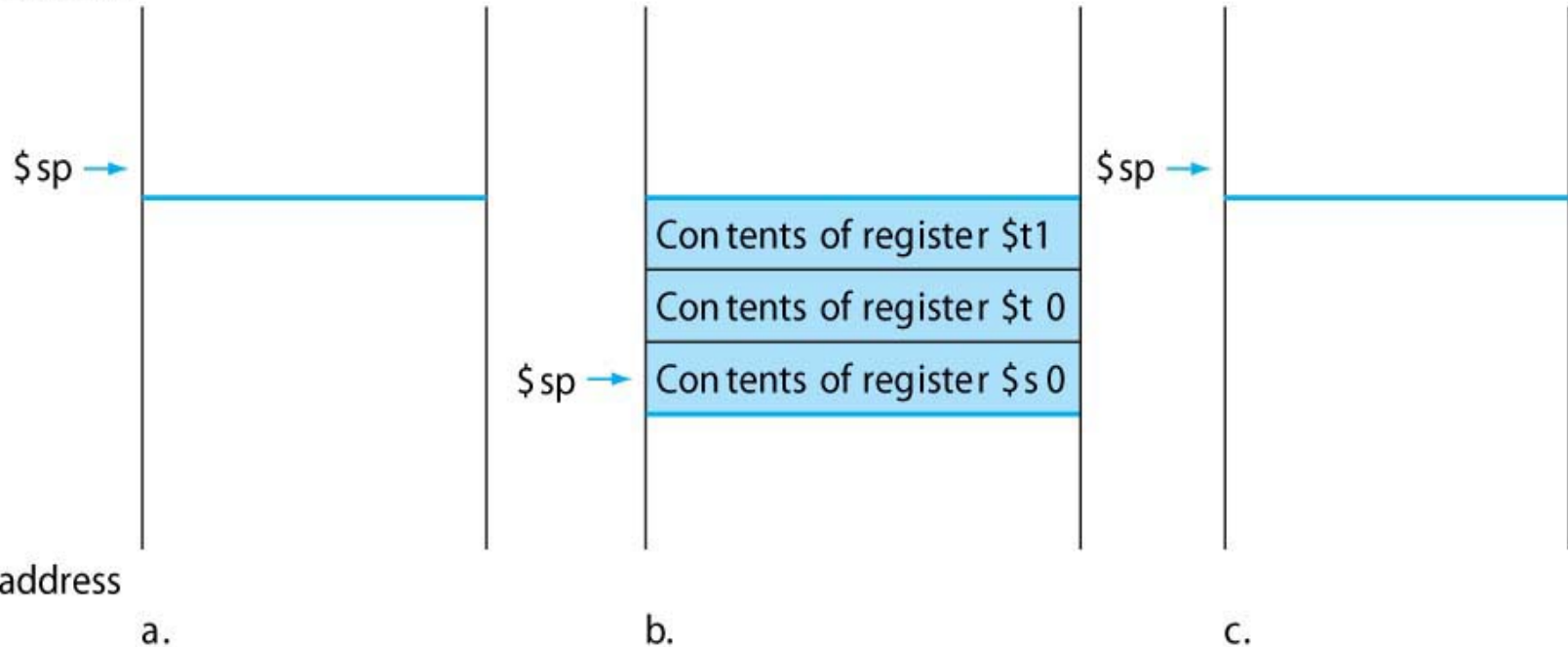
Memory Layout

- Text: program code
 - PC initialized to 0x00400000
- Static data: global/static variables
 - \$gp initialized to the middle of this segment, 0x10008000 allowing \pm offset
- Dynamic data: heap
 - E.g., malloc in C, new in Java
- Stack: storage for temporary variable in functions
 - \$sp initialized to 0x7ffffffc, growing towards low address



Uses of Stack in Function Call

High address



Before calling

During procedure

- For storing important registers
- For temporary variables

After calling

- Important registers restored
- Temporary variables destroyed

Leaf Procedure

- Procedures that don't call other procedures

- C code:

```
int leaf_example (int g, h, i, j)
{ int f;
  f = (g + h) - (i + j);
  return f;}
```

- Assumptions:

- Arguments g, ..., j in \$a0, ..., \$a3
- f in \$s0 (need to save \$s0 before it's overwritten)

Result in \$v0



Leaf Procedure Example

■ MIPS code

leaf_example:

addi \$sp, \$sp, -12 #create spaces on stack

sw \$t1, 8(\$sp)

sw \$t0, 4(\$sp)

sw \$s0, 0(\$sp)

add \$t0, \$a0, \$a1

add \$t1, \$a2, \$a3

sub \$s0, \$t0, \$t1

add \$v0, \$s0, \$zero

lw \$s0, 0(\$sp)

lw \$t0, 4(\$sp)

lw \$t1, 8(\$sp)

addi \$sp, \$sp, 12

jr \$ra

#store data on stack

Unnecessary

#restore data from stack

#destroy spaces on stack

#return from procedure

String Copy Example

- C code:

- Assuming null-terminated string

```
void strcpy (char x[], char y[])  
{ int i;  
  i = 0;  
  while ((x[i]=y[i])!='\0')  
    i += 1;  
}
```

- Base addresses of x, y in \$a0, \$a1
- i in \$s0

String Copy Example

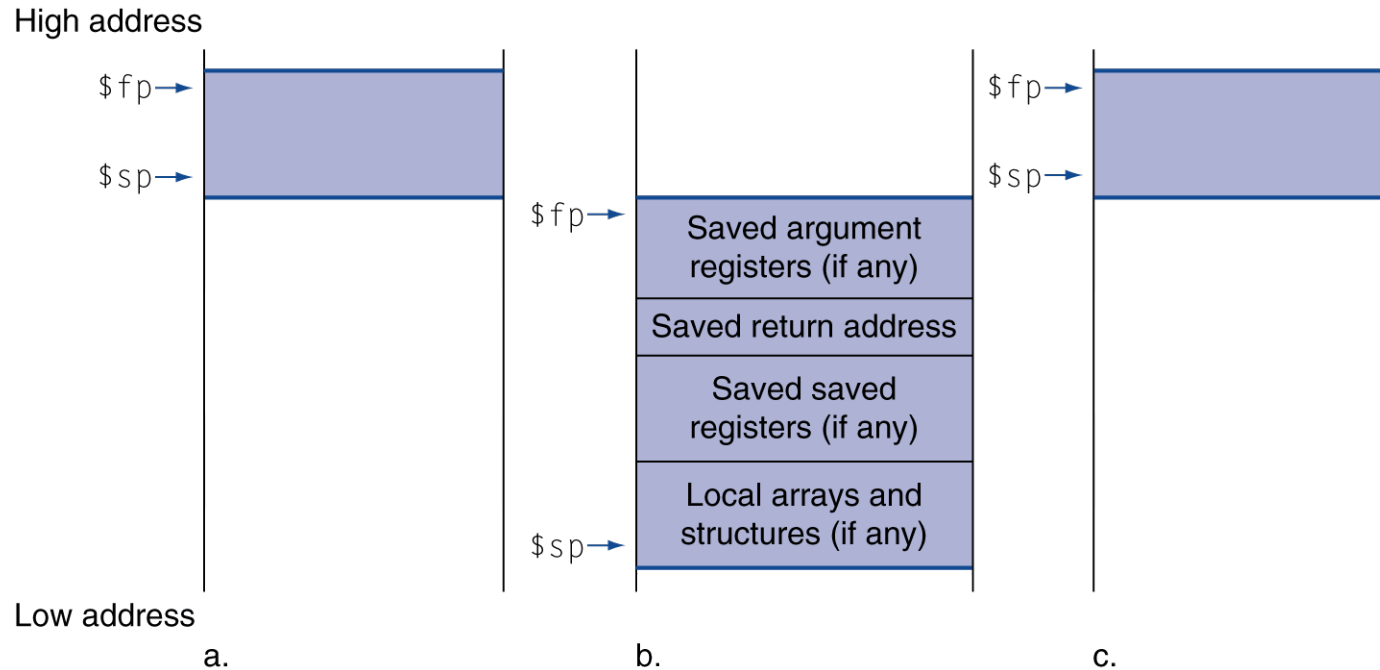
■ MIPS code:

strcpy:		
	addi \$sp, \$sp, -4	# adjust stack for 1 item
	sw \$s0, 0(\$sp)	# save \$s0
	add \$s0, \$zero, \$zero	# i = 0
L1:	add \$t1, \$s0, \$a1	# addr of y[i] in \$t1
	lbu \$t2, 0(\$t1)	# \$t2 = y[i]
	add \$t3, \$s0, \$a0	# addr of x[i] in \$t3
	sb \$t2, 0(\$t3)	# x[i] = y[i]
	beq \$t2, \$zero, L2	# exit loop if y[i] == 0
	addi \$s0, \$s0, 1	# i = i + 1
	j L1	# next iteration of loop
L2:	lw \$s0, 0(\$sp)	# restore saved \$s0
	addi \$sp, \$sp, 4	# pop 1 item from stack
	jr \$ra	# and return

Non-Leaf Procedures

- Procedures that call other procedures
- For nested call, caller needs to save on the stack:
 - Its return address
 - Any arguments and temporaries needed after the call
- Restore from the stack after the call

Local Data on the Stack



- Procedure frame (activation record)
 - Saved registers
 - Local data allocated by procedure
- Two pointers manages stack
 - $\$sp$ manages frames
 - $\$fp$ manages elements in each frame

Non-Leaf Procedure Example

- C code:

```
int fact (int n)
{
    if (n < 1) return f;
    else return n * fact(n - 1);
}
```

- Argument n in $\$a0$
- Result in $\$v0$

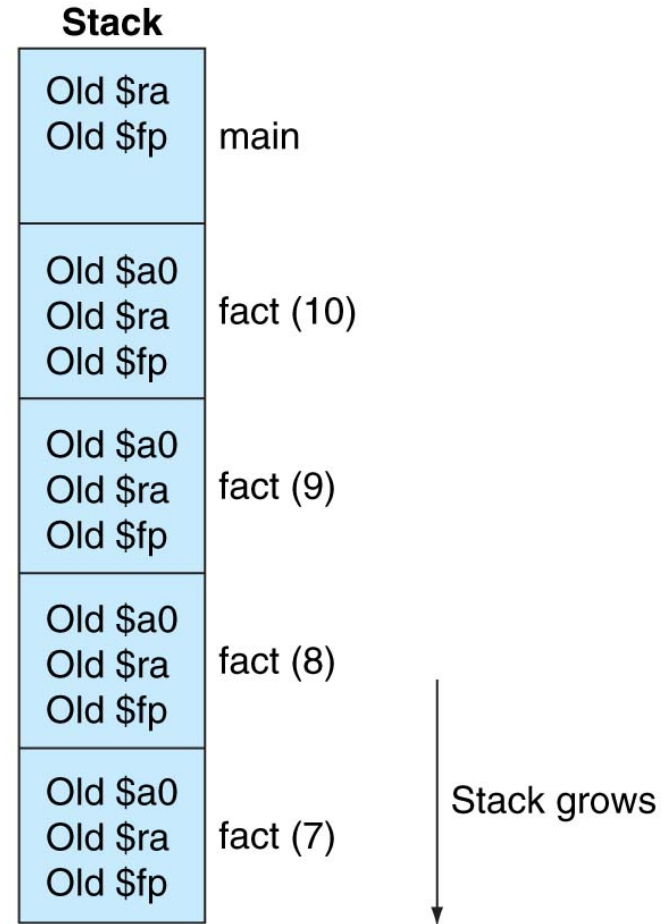
Non-Leaf Procedure Example

- MIPS code:

fact:		
addi	\$sp, \$sp, -8	# adjust stack for 2 items
sw	\$ra, 4(\$sp)	# save return address
sw	\$a0, 0(\$sp)	# save argument
slti	\$t0, \$a0, 1	# test for n < 1
beq	\$t0, \$zero, L1	
addi	\$v0, \$zero, 1	# if so, result is 1
addi	\$sp, \$sp, 8	# release stack
jr	\$ra	# and return
L1:	addi \$a0, \$a0, -1	# else decrement n
	jal fact	# recursive call
lw	\$a0, 0(\$sp)	# restore original n
lw	\$ra, 4(\$sp)	# and return address
addi	\$sp, \$sp, 8	# pop 2 items from stack
mul	\$v0, \$a0, \$v0	# multiply to get result
jr	\$ra	# and return

Usage of Stack Frames

- `fact (int n)` is a function, can be called recursively
- Note: `$fp` wasn't used in previous example



Procedure Calling Convention

- Three places in procedure calling when conventions apply
 - Immediately before the procedure is called
 - In procedure, but before it starts executing
 - Immediately before the procedure finishes

Procedure Calling Convention

- Before the procedure is called
 1. Pass arguments to \$a0-\$a3
 - more arguments on stack
 2. Save registers that should be saved by caller,
 - such as \$a0-\$a3 (non-leaf procedure), \$t0-\$t9 (if necessary)
 3. jal

Procedure Calling Convention

- Before procedure starts executing
 1. Allocate memory of frame's size
 - by moving `$sp` downwards for frame's size
 2. Save registers that should be saved by the procedure in the frame, before they are overwritten
 - `$s0-$s7` (if to be used), `$fp` (if used), `$ra` (non-leaf procedure),
 3. Establish `$fp` (if desired), $\text{\$fp} = \text{\$sp} + \text{frame's size} - 4$

Procedure Calling Convention

- Before procedure finishes
 1. If necessary, place procedure result to \$v0, \$v1
 2. Restore registers saved by the procedure
 - Pop from frame
 3. Destroy stack frame by moving \$sp upword
 4. `jr $ra`

C Sort Example

- Illustrates use of assembly instructions for a C bubble sort function

- Swap procedure (leaf)

```
void swap(int v[], int k)
{
    int temp;
    temp = v[k];
    v[k] = v[k+1];
    v[k+1] = temp;
}
```

- v in \$a0, k in \$a1, temp in \$t0

The Procedure Swap

swap:	sll \$t1, \$a1, 2	# \$t1 = k * 4
	add \$t1, \$a0, \$t1	# \$t1 = v+(k*4)
		# (address of v[k])
	lw \$t0, 0(\$t1)	# \$t0 (temp) = v[k]
	lw \$t2, 4(\$t1)	# \$t2 = v[k+1]
	sw \$t2, 0(\$t1)	# v[k] = \$t2 (v[k+1])
	sw \$t0, 4(\$t1)	# v[k+1] = \$t0 (temp)
	jr \$ra	# return to calling routine

The Sort Procedure in C

- Non-leaf procedure (calls swap)

```
void sort (int v[], int n)
{
    int i, j;
    for (i = 0; i < n; i += 1) {
        for (j = i - 1;
             j >= 0 && v[j] > v[j + 1];
             j -= 1) {
            swap(v, j);
        }
    }
}
```

- v in \$a0, n in \$a1, i in \$s0, j in \$s1

The Full Procedure

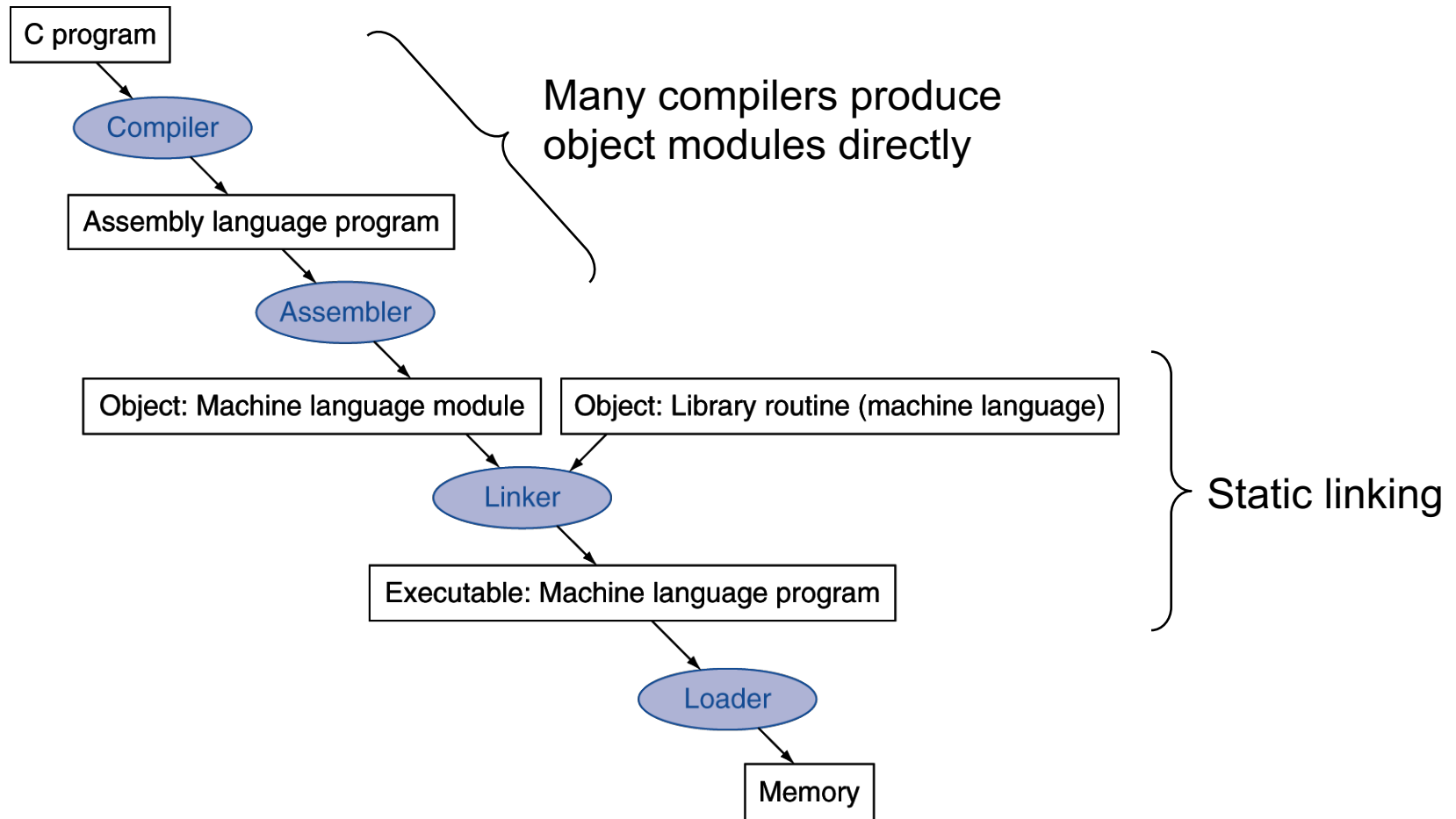
sort:	addi \$sp,\$sp, -20	# make room on stack for 5 # registers
	sw \$ra, 16(\$sp)	# save \$ra on stack
	sw \$s3, 12(\$sp)	# save \$s3 on stack
	sw \$s2, 8(\$sp)	# save \$s2 on stack
	sw \$s1, 4(\$sp)	# save \$s1 on stack
	sw \$s0, 0(\$sp)	# save \$s0 on stack
	...	# procedure body
	...	
exit1:	lw \$s0, 0(\$sp)	# restore \$s0 from stack
	lw \$s1, 4(\$sp)	# restore \$s1 from stack
	lw \$s2, 8(\$sp)	# restore \$s2 from stack
	lw \$s3, 12(\$sp)	# restore \$s3 from stack
	lw \$ra, 16(\$sp)	# restore \$ra from stack
	addi \$sp,\$sp, 20	# restore stack pointer
	jr \$ra	# return to calling routine

The Procedure Body

	move \$s2, \$a0	# save \$a0 into \$s2	Move params
	move \$s3, \$a1	# save \$a1 into \$s3	
for1tst:	move \$s0, \$zero	# i = 0	Outer loop
	slt \$t0, \$s0, \$s3	# \$t0 = 0 if \$s0 ≥ \$s3 (i ≥ n)	
for2tst:	beq \$t0, \$zero, exit1	# go to exit1 if \$s0 ≥ \$s3 (i ≥ n)	
	addi \$s1, \$s0, -1	# j = i - 1	
	slti \$t0, \$s1, 0	# \$t0 = 1 if \$s1 < 0 (j < 0)	
	bne \$t0, \$zero, exit2	# go to exit2 if \$s1 < 0 (j < 0)	
	sll \$t1, \$s1, 2	# \$t1 = j * 4	Inner loop
	add \$t2, \$s2, \$t1	# \$t2 = v + (j * 4)	
	lw \$t3, 0(\$t2)	# \$t3 = v[j]	
	lw \$t4, 4(\$t2)	# \$t4 = v[j + 1]	
	slt \$t0, \$t4, \$t3	# \$t0 = 0 if \$t4 ≥ \$t3	
	beq \$t0, \$zero, exit2	# go to exit2 if \$t4 ≥ \$t3	
	move \$a0, \$s2	# 1st param of swap is v (old \$a0)	Pass params & call
	move \$a1, \$s1	# 2nd param of swap is j	
	jal swap	# call swap procedure	
	addi \$s1, \$s1, -1	# j -= 1	Inner loop
	j for2tst	# jump to test of inner loop	
exit2:	addi \$s0, \$s0, 1	# i += 1	Outer loop
	j for1tst	# jump to test of outer loop	



Translation and Startup



Producing an Object Module

- Assembler (or compiler) translates program into machine instructions
- Provides information for building a complete program from the pieces
 - *Header*: described contents of object module
 - *Text segment*: translated instructions
 - *Static data segment*: data allocated for the life of the program
 - *Relocation info*: for contents that depend on absolute location of loaded program
 - *Symbol table*: global definitions and external refs
 - *Debug info*: for associating with source code

Example of Object Modules

```
int x[100], y[100];
Procedure_A(int m)
{ m = X[0];
  ...
  Procedure_B(...);
  ...
}
```



```
lw $a0, offset1($gp)
jal Procedure_B
...
```

```
Procedure_B(int n)
{ Y[0] = n;
  ...
  Procedure_A(...);
  ...
}
```



```
sw $a1, offset2($gp)
jal Procedure_A
...
```

- m and n are parameters to the C functions
- Array X and Y are global variables
- By default \$gp = 1000 8000_{hex}

Example of Object Modules

Object file header			
	Name	Procedure A	
	Text size	100 _{hex}	
	Data size	20 _{hex}	
Text segment	Address	Instruction	
	0	lw \$a0, 0(\$gp)	
	4	jal 0	
	
Data segment	0	(X)	
	
Relocation information	Address	Instruction type	Dependency
	0	lw	X
	4	jal	B
Symbol table	Label	Address	
	X	---	
	B	---	

Object file header			
	Name	Procedure B	
	Text size	200 _{hex}	
	Data size	30 _{hex}	
Text segment	Address	Instruction	
	0	sw \$a1, 0(\$gp)	
	4	jal 0	
	
Data segment	0	(Y)	
	
Relocation information	Address	Instruction type	Dependency
	0	sw	Y
	4	jal	A
Symbol table	Label	Address	
	Y	---	
	A	---	

Linking Object Modules

- Produces an executable image
 1. Merges segments
 2. Resolve labels (determine their addresses)
 3. Patch location-dependent and external references

Example of Linked Objects

Executable File Header		
	Text size	300 _{hex}
	Data size	50 _{hex}
Text Segment	Address	Instruction
	0040 0000 _{hex}	lw \$a0, 8000 _{hex} (\$gp)
	0040 0004 _{hex}	jal 40 0100 _{hex}

	0040 0100 _{hex}	sw \$a1, 8020 _{hex} (\$gp)
	0040 0104 _{hex}	jal 40 0000 _{hex}

Data Segment	Address	
	1000 0000 _{hex}	(X)

	1000 0020 _{hex}	(Y)

Loading a Program

- Load from image file on disk into memory
 1. Read header to determine segment sizes
 2. Create virtual address space
 3. Copy text and initialized data into memory
 4. Set up arguments on stack
 5. Initialize registers (including `$sp`, `$fp`, `$gp`)
 6. Jump to startup routine
 - Copies arguments to `$a0`, ... and calls `main`
 - When `main` returns, do `exit` syscall

Dynamic Linking

- Only link/load library procedure when it is called
 - Requires procedure code to be relocatable
 - Avoids big executable caused by static linking of all referenced libraries
 - Some of them may be never used
 - Automatically picks up new library versions

MIPS R2000 Assembly Language

- Appendix B.10