

3) a)

$$T(n) = T(n/2) + n$$

$$T(1) = 1$$

$$T(n) = T(n/2) + n$$

$$T(n/2) = T(n/2^2) + n/2$$

$$T(n/2^2) = T(n/2^3) + n/2^2$$

↓

$$T(n) = \underbrace{T(n/2^k)}_{T(1)} + n/2^{k-1}$$

$$\frac{n}{2^k} = 1$$

$$n = 2^k$$

$$k = \log_2 n$$

$$T(n) = 1 + \sum_{i=0}^{\log n} \frac{n}{2^i}$$

$$T(n) = 1 + n \sum_{i=0}^{\log n} \frac{1}{2^i}$$

$\sum_{i=0}^{\log n} \frac{1}{2^i} \rightarrow$ Somatório nunca será > 2 .

$$T(n) = 1 + n \cdot \sum_{i=0}^{\log n} \frac{1}{2^i}$$

→ Domina o Somatório

$$\text{Complexidade} = O(n)$$