

Exercício Prático 06

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Parte 1

1. O que é um arquivo fonte?

A. um arquivo de texto que contém instruções de linguagem de programação.

2. O que é um registrador?

A. parte do sistema de computador que mantém o controle dos parâmetros do sistema.

3. Qual o caracter que, na linguagem assembly do SPIM, inicia um comentário?

A. #

4. Quantos bits há em cada instrução de máquina MIPS?

C. 32

5. O que é o contador de programa?

D. parte do processador que contém o endereço da próxima instrução de máquina para ser obtida.

6. Ao executarmos uma instrução, quanto será adicionado ao contador de programa?

C. 4

7. O que é uma diretiva, tal como a diretiva .text?

D. uma declaração que diz o montador algo sobre o que o programador quer, mas não corresponde diretamente a uma instrução de máquina.

8. O que é um endereço simbólico?

D. um nome usado no código-fonte em linguagem assembly para um local na memória.

9. Em qual endereço o simulador SPIM coloca a primeira instrução de máquina quando ele está sendo executado?

B. 0x00400000

10. Algumas instruções de máquina possuem uma constante como um dos operandos. Como é chamado tal operando?

A. operando imediato

11. Como é chamada uma operação lógica executada entre bits de cada coluna dos operandos para produzir um bit de resultado para cada coluna?

B. operação bitwise

12. Quando uma operação é de fato executada, como estão os operandos na ALU?

D. Cada um dos registradores deve possuir 32 bit.

13. Dezesesseis bits de dados de uma instrução de ori são usados como um operando imediato. Durante execução, o que deve ser feito primeiro?

B. Os dados são estendidos em zero à esquerda por 16 bits.

14. Qual das instruções seguintes armazenam no registrador \$5 um padrão de bits que representa positivo 48?

C. ori \$5,\$0,48

15. A instrução de ori pode armazenar o complemento de dois de um número em um registrador?

B. Sim.

16. Qual das instruções seguintes limpa todos os bits no registrador \$8 com exceção do byte de baixa ordem que fica inalterado?

D. andi \$8,\$8,0xFF

17. Qual é o resultado de um ou exclusivo de padrão sobre ele mesmo?

A. Todos os bits em zero.

18. Todas as instruções de máquina têm os mesmos campos?

A. Não. Diferentes de instruções de máquina possuem campos diferentes.

Parte 2

Implementar em MIPS/MARS os seguintes programas

programa 1 (add, addi, sub, lógicas) {

```
a = 2;
b = 3;
c = 4;
d = 5;
x = (a+b) - (c+d);
y = a - b + x;
b = x - y;
```

}

The screenshot displays the MARS MIPS simulator interface. The top window shows the assembly code for the program, which implements the given logic in MIPS assembly. The code includes comments for each step, such as initializing registers \$a0 through \$a5 with values 2 through 5, calculating x = (a+b) - (c+d), y = a - b + x, and finally updating b = x - y. The bottom window shows the memory state, with the Text Segment and Data Segment visible. The Text Segment contains the compiled MIPS instructions, and the Data Segment shows the memory addresses and values for the registers. The registers \$a0 through \$a5 are initialized with values 2 through 5, respectively. The registers \$t0 through \$t5 are initialized with values 0 through 5, respectively. The registers \$s0 through \$s5 are initialized with values 0 through 5, respectively. The registers \$f0 through \$f5 are initialized with values 0.0 through 5.0, respectively. The registers \$d0 through \$d5 are initialized with values 0.0 through 5.0, respectively. The registers \$e0 through \$e5 are initialized with values 0.0 through 5.0, respectively. The registers \$o0 through \$o5 are initialized with values 0.0 through 5.0, respectively. The registers \$l0 through \$l5 are initialized with values 0.0 through 5.0, respectively. The registers \$q0 through \$q5 are initialized with values 0.0 through 5.0, respectively. The registers \$r0 through \$r5 are initialized with values 0.0 through 5.0, respectively. The registers \$t0 through \$t5 are initialized with values 0 through 5, respectively. The registers \$s0 through \$s5 are initialized with values 0 through 5, respectively. The registers \$f0 through \$f5 are initialized with values 0.0 through 5.0, respectively. The registers \$d0 through \$d5 are initialized with values 0.0 through 5.0, respectively. The registers \$e0 through \$e5 are initialized with values 0.0 through 5.0, respectively. The registers \$o0 through \$o5 are initialized with values 0.0 through 5.0, respectively. The registers \$l0 through \$l5 are initialized with values 0.0 through 5.0, respectively. The registers \$q0 through \$q5 are initialized with values 0.0 through 5.0, respectively. The registers \$r0 through \$r5 are initialized with values 0.0 through 5.0, respectively.

```
1
2 # a -> $a0
3 # b -> $a1
4 # c -> $a2
5 # d -> $a3
6 # x -> $a4
7 # y -> $a5
8
9 .text
10 # inicio
11 addi $a0,$zero, 2 # a = 2
12 addi $a1,$zero, 3 # b = 3
13 addi $a2,$zero, 4 # c = 4
14 addi $a3,$zero, 5 # d = 5
15 add $t0, $a0, $a1 # a + b
16 add $t1, $a2, $a3 # c + d
17 sub $a4, $t0, $t1 # (a + b) - (c + d)
18 sub $t0, $a0, $a1 # a - b
19 add $a5, $t0, $a5 # (a - b) + 5
20 sub $a1, $a4, $a5 # b = x - y
21 # fim
22
```

| Name | Number | Value |
|--------|--------|------------|
| \$zero | 0 | 0x00000000 |
| \$at | 1 | 0x00000000 |
| \$v0 | 2 | 0x00000000 |
| \$v1 | 3 | 0x00000000 |
| \$a0 | 4 | 0x00000000 |
| \$a1 | 5 | 0x00000000 |
| \$a2 | 6 | 0x00000000 |
| \$a3 | 7 | 0x00000000 |
| \$t0 | 8 | 0xffffffff |
| \$t1 | 9 | 0x00000009 |
| \$t2 | 10 | 0x00000000 |
| \$t3 | 11 | 0x00000000 |
| \$t4 | 12 | 0x00000000 |
| \$t5 | 13 | 0x00000000 |
| \$t6 | 14 | 0x00000000 |
| \$t7 | 15 | 0x00000000 |
| \$s0 | 16 | 0x00000002 |
| \$s1 | 17 | 0xffffffff |
| \$s2 | 18 | 0x00000004 |
| \$s3 | 19 | 0x00000005 |
| \$s4 | 20 | 0xffffffff |
| \$s5 | 21 | 0xffffffff |
| \$s6 | 22 | 0x00000000 |
| \$s7 | 23 | 0x00000000 |
| \$t8 | 24 | 0x00000000 |
| \$t9 | 25 | 0x00000000 |
| \$t0 | 26 | 0x00000000 |
| \$t1 | 27 | 0x00000000 |
| \$gp | 28 | 0x10000000 |
| \$sp | 29 | 0x7fffffc0 |
| \$fp | 30 | 0x00000000 |
| \$ra | 31 | 0x00000000 |
| pc | | 0x00400028 |
| hi | | 0x00000000 |
| lo | | 0x00000000 |

Mars Messages Run I/O

Clear -- program is finished running (dropped off bottom) --

```
//programa 2 (add, addi, sub, lógicas) {
    x = 1;
    y = 5*x + 15;
}
```

The screenshot shows the Mars MIPS simulator interface. The top window displays the assembly code for 'mips1.asm'. The code implements Program 2, which calculates $y = 5x + 15$ for $x = 1$. The code uses the `addi` instruction to set `$s4` to 1, then a loop of `add` instructions to calculate $5x$ (adding `$s4` to `$t0` four times), and finally `addi` to add 15 to `$t0` and store the result in `$s5`.

The bottom window shows the 'Text Segment' and 'Data Segment' memory views. The 'Text Segment' table lists the instructions being executed, and the 'Data Segment' table shows the memory layout, with `$s5` at address 0x10010014 containing the value 0x00000015.

| Name | Number | Value |
|-------------|-----------|-------------------|
| \$zero | 0 | 0x00000000 |
| \$at | 1 | 0x00000000 |
| \$v0 | 2 | 0x00000000 |
| \$v1 | 3 | 0x00000000 |
| \$a0 | 4 | 0x00000000 |
| \$a1 | 5 | 0x00000000 |
| \$a2 | 6 | 0x00000000 |
| \$a3 | 7 | 0x00000000 |
| \$t0 | 8 | 0x00000005 |
| \$t1 | 9 | 0x00000000 |
| \$t2 | 10 | 0x00000000 |
| \$t3 | 11 | 0x00000000 |
| \$t4 | 12 | 0x00000000 |
| \$t5 | 13 | 0x00000000 |
| \$t6 | 14 | 0x00000000 |
| \$t7 | 15 | 0x00000000 |
| \$s0 | 16 | 0x00000000 |
| \$s1 | 17 | 0x00000000 |
| \$s2 | 18 | 0x00000000 |
| \$s3 | 19 | 0x00000000 |
| \$s4 | 20 | 0x00000001 |
| \$s5 | 21 | 0x00000015 |
| \$s6 | 22 | 0x00000000 |
| \$s7 | 23 | 0x00000000 |
| \$s8 | 24 | 0x00000000 |
| \$t9 | 25 | 0x00000000 |
| \$k0 | 26 | 0x00000000 |
| \$k1 | 27 | 0x00000000 |
| \$gp | 28 | 0x10008000 |
| \$sp | 29 | 0x7fffffc0 |
| \$fp | 30 | 0x00000000 |
| \$ra | 31 | 0x00000000 |
| pc | | 0x00400018 |
| hi | | 0x00000000 |
| lo | | 0x00000000 |

```
// programa 3 (add, addi, sub, lógicas) {
    x = 3;
    y = 4;
    z = ( 15*x + 67*y)*4
}
```

```
1 .text
2
3 addi $s0, $zero, 3    # x = 3
4 addi $s1, $zero, 4    # y = 4
5
6 # 15 * x
7 addi $t0, $s0, $s0    # t0 = 2 * x
8 addi $t0, $t0, $t0    # t0 = 4 * x
9 addi $t0, $t0, $t0    # t0 = 8 * x
10 addi $t0, $t0, $t0    # t0 = 16 * x
11 sub $t0, $t0, $s0     # t0 = 15 * x
12
13 # 67 * y
14 addi $t1, $s1, $s1    # t1 = 2 * y
15 addi $t1, $t1, $t1    # t1 = 4 * y
16 addi $t1, $t1, $t1    # t1 = 8 * y
17 addi $t1, $t1, $t1    # t1 = 16 * y
18 addi $t1, $t1, $t1    # t1 = 32 * y
19 addi $t1, $t1, $t1    # t1 = 64 * y
20 add $t1, $t1, $s1     # t1 = 67 * y
21
22 # (15 * x + 67 * y) * 4
23 add $t2, $t0, $t1     # t2 = (15 * x + 67 * y)
24 addi $s2, $t2, $t2    # s2 = 2 * (15 * x + 67 * y)
25 addi $t2, $t2, $t2    # t2 = 4 * (15 * x + 67 * y)
26 # Fim
27
```



```
// programa 5 {
    x = 100000;
    y = 200000;
    z = x + y;
}
```

```
1 .text
2 ori $s0, $zero, 0x186A
3 sll $s0, $s0, 4
4 ori $s1, $zero, 0x30D4
5 sll $s1, $s1, 4
6 add $s2, $s0, $s1
7 |
```

The screenshot shows the Mars MIPS simulator interface. The main window displays assembly code for a program. Below the code, there are two panels: 'Text Segment' and 'Data Segment'. The 'Text Segment' panel shows the assembly code with addresses, codes, and sources. The 'Data Segment' panel shows memory addresses and values. On the right side, there is a table of registers and their values.

| Name | Number | Value |
|-------------|-----------|-------------------|
| \$zero | 0 | 0x00000000 |
| \$at | 1 | 0x00000000 |
| \$v0 | 2 | 0x00000000 |
| \$v1 | 3 | 0x00000000 |
| \$a0 | 4 | 0x00000000 |
| \$a1 | 5 | 0x00000000 |
| \$a2 | 6 | 0x00000000 |
| \$a3 | 7 | 0x00000000 |
| \$t0 | 8 | 0x00000000 |
| \$t1 | 9 | 0x00000000 |
| \$t2 | 10 | 0x00000000 |
| \$t3 | 11 | 0x00000000 |
| \$t4 | 12 | 0x00000000 |
| \$t5 | 13 | 0x00000000 |
| \$t6 | 14 | 0x00000000 |
| \$t7 | 15 | 0x00000000 |
| \$s0 | 16 | 0x000186A0 |
| \$s1 | 17 | 0x00030D40 |
| \$s2 | 18 | 0x000493E0 |
| \$s3 | 19 | 0x00000000 |
| \$s4 | 20 | 0x00000000 |
| \$s5 | 21 | 0x00000000 |
| \$s6 | 22 | 0x00000000 |
| \$s7 | 23 | 0x00000000 |
| \$s8 | 24 | 0x00000000 |
| \$s9 | 25 | 0x00000000 |
| \$k0 | 26 | 0x00000000 |
| \$k1 | 27 | 0x00000000 |
| \$gp | 28 | 0x10008000 |
| \$sp | 29 | 0x7fffffc |
| \$fp | 30 | 0x00000000 |
| \$ra | 31 | 0x00000000 |
| pc | | 0x00400014 |
| hi | | 0x00000000 |
| lo | | 0x00000000 |

```
// programa 6 {
    x = o maior inteiro possível;
    y = 300000;
    z = x - 4y
}
```

```
mips1.asm
1 # carregar o maior valor possível (32767) em $s0
2 ori $s0, $zero, 0x7FFF # x = 32767 (0x7FFF)
3
4 # carregar o valor de y (300000) em $s1 usando deslocamento e soma
5 ori $s1, $zero, 0x493E # carrega 4 em $s1
6
7 # calcular 4 * y (multiplicar y por 4)
8 sll $s2, $s1, 4 # desloca y para a esquerda por 2 bits (multiplica por 4). Resultado em $s2 = 1200000
9
10 # Calcular z = x - 4y
11 sub $s3, $s0, $s2 # z = x - 4y, o resultado será armazenado em $s3
12
```


// programa 8

Inicialmente escreva um programa que faça:

\$8 = 0x12345678.

A partir do registrador \$8 acima, usando apenas instruções lógicas (or, ori, and, andi, xor, xori) e instruções de deslocamento (sll, srl e sra), você deverá obter os seguintes valores nos respectivos registradores:

\$9 = 0x12

\$10 = 0x34

\$11 = 0x56

\$12 = 0x78

The screenshot displays the MIPS assembler interface with the following components:

- Assembly Code (mips1.asm):**

```
1 # atribuir $8 = 0x12345678
2 ori $8, $zero, 0x1234
3 sll $8, $8, 16
4 ori $8, $8, 0x5678
5
6 # obtem $9 = 12
7 srl $9, $8, 24 # $9 = 12
8
9 # obtem $10 = 34
10 srl $10, $8, 16
11 andi $10, $10, 0x00FF
12
13 # obtem $11 = 56
14 srl $11, $8, 8
15 andi $11, $11, 0x00FF
16
17 # obtem $12 = 78
18 andi $12, $8, 0x00FF
19
```
- Text Segment Table:**

| Bkpt | Address | Code | Basic | Source |
|------|------------|------------|--------------------------|------------------------------|
| 0 | 0x00400000 | 0x34081234 | ori \$8,\$0,0x00001234 | 2: ori \$8, \$zero, 0x1234 |
| 1 | 0x00400004 | 0x00084400 | sll \$8,\$8,0x00000010 | 3: sll \$8,\$8, 16 |
| 2 | 0x00400008 | 0x35085678 | ori \$8,\$8,0x00005678 | 4: ori \$8, \$8, 0x5678 |
| 3 | 0x0040000c | 0x00084e02 | srl \$9,\$8,0x00000018 | 7: srl \$9,\$8,24 # \$9 = 12 |
| 4 | 0x00400010 | 0x00085402 | srl \$10,\$8,0x00000010 | 10: srl \$10, \$8, 16 |
| 5 | 0x00400014 | 0x314a00ff | andi \$10,\$10,0x0000... | 11: andi \$10, \$10, 0x00FF |
| 6 | 0x00400018 | 0x00085a02 | srl \$11,\$8,0x00000008 | 14: srl \$11, \$8, 8 |
| 7 | 0x0040001c | 0x316b00ff | andi \$11,\$11,0x0000... | 15: andi \$11, \$11, 0x00FF |
| 8 | 0x00400020 | 0x310c00ff | andi \$12,\$8,0x000000ff | 18: andi \$12, \$8, 0x00FF |
- Data Segment Table:**

| Address | Value (+0) | Value (+4) | Value (+8) | Value (+c) | Value (+10) | Value (+14) | Value (+18) | Value (+1c) |
|------------|------------|------------|------------|------------|-------------|-------------|-------------|-------------|
| 0x10010000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
| 0x10010020 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
| 0x10010040 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
- Register Values Table:**

| Name | Number | Value |
|--------|--------|------------|
| \$zero | 0 | 0x00000000 |
| \$at | 1 | 0x00000000 |
| \$v0 | 2 | 0x00000000 |
| \$v1 | 3 | 0x00000000 |
| \$a0 | 4 | 0x00000000 |
| \$a1 | 5 | 0x00000000 |
| \$a2 | 6 | 0x00000000 |
| \$a3 | 7 | 0x00000000 |
| \$t0 | 8 | 0x12345678 |
| \$t1 | 9 | 0x00000012 |
| \$t2 | 10 | 0x00000034 |
| \$t3 | 11 | 0x00000056 |
| \$t4 | 12 | 0x00000078 |
| \$t5 | 13 | 0x00000000 |
| \$t6 | 14 | 0x00000000 |
| \$t7 | 15 | 0x00000000 |
| \$t8 | 16 | 0x00000000 |
| \$t9 | 17 | 0x00000000 |
| \$s0 | 18 | 0x00000000 |
| \$s1 | 19 | 0x00000000 |
| \$s2 | 20 | 0x00000000 |
| \$s3 | 21 | 0x00000000 |
| \$s4 | 22 | 0x00000000 |
| \$s5 | 23 | 0x00000000 |
| \$s6 | 24 | 0x00000000 |
| \$s7 | 25 | 0x00000000 |
| \$s8 | 26 | 0x00000000 |
| \$s9 | 27 | 0x00000000 |
| \$gp | 28 | 0x10008000 |
| \$sp | 29 | 0x7fffffc |
| \$fp | 30 | 0x00000000 |
| \$ra | 31 | 0x00000000 |
| pc | | 0x00400024 |
| hi | | 0x00000000 |
| lo | | 0x00000000 |

Para os programas a seguir use instruções de Memória (lw e sw)

// programa 9

Considere a memória inicial da seguinte forma:

.text

.data

x1: .word 15

x2: .word 25

x3: .word 13

x4: .word 17

soma: .word -1

Escrever um programa que leia todos os números, calcule e substitua o valor da variável soma por este valor.

The screenshot displays the MARS MIPS simulator interface. The top window shows the assembly code for `mips1.asm`. The code declares variables `x1`, `x2`, `x3`, `x4`, and `soma` in the `.data` segment. It then loads these values into registers `$s0`, `$s1`, `$s2`, and `$s3` in the `.text` segment. The program calculates the sum of `x1`, `x2`, `x3`, and `x4` and stores the result in `soma`.

```
1 .data
2 x1: .word 15      # Declara a variável x1 com o valor 15
3 x2: .word 25      # Declara a variável x2 com o valor 25
4 x3: .word 13      # Declara a variável x3 com o valor 13
5 x4: .word 17      # Declara a variável x4 com o valor 17
6 soma: .word -1    # Inicializa a variável soma com valor -1
7
8 .text
9 lui $t0, 0x1001
10 lw $s0, 0($t0)   # x1 = MEM[0]
11 lw $s1, 4($t0)   # x2 = MEM[1]
12 lw $s2, 8($t0)   # x3 = MEM[2]
13 lw $s3, 12($t0)  # x4 = MEM[3]
14 add $t1, $s0,$s1 # x1 + x2
15 add $t1, $t1, $s2 # x1 + x2 + x3
16 add $t1, $t1, $s3 # x1 + x2 + x3 + x4
17 sw $t1, 16($t0)  # MEM[soma] = x1 + x2 + x3 + x4
18
```

The bottom window shows the `Text Segment` and `Data Segment`. The `Text Segment` table lists the instructions and their addresses. The `Data Segment` table shows the memory layout, with the value `0x00000046` stored at address `0x10010000`, which corresponds to the `soma` variable.

| Name | Number | Value |
|--------|--------|------------|
| \$zero | 0 | 0x00000000 |
| \$at | 1 | 0x00000000 |
| \$v0 | 2 | 0x00000000 |
| \$v1 | 3 | 0x00000000 |
| \$a0 | 4 | 0x00000000 |
| \$a1 | 5 | 0x00000000 |
| \$a2 | 6 | 0x00000000 |
| \$a3 | 7 | 0x00000000 |
| \$t0 | 8 | 0x10010000 |
| \$t1 | 9 | 0x00000046 |
| \$t2 | 10 | 0x00000000 |
| \$t3 | 11 | 0x00000000 |
| \$t4 | 12 | 0x00000000 |
| \$t5 | 13 | 0x00000000 |
| \$t6 | 14 | 0x00000000 |
| \$t7 | 15 | 0x00000000 |
| \$s0 | 16 | 0x0000000f |
| \$s1 | 17 | 0x00000019 |
| \$s2 | 18 | 0x0000000d |
| \$s3 | 19 | 0x00000011 |
| \$s4 | 20 | 0x00000000 |
| \$s5 | 21 | 0x00000000 |
| \$s6 | 22 | 0x00000000 |
| \$s7 | 23 | 0x00000000 |
| \$s8 | 24 | 0x00000000 |
| \$s9 | 25 | 0x00000000 |
| \$k0 | 26 | 0x00000000 |
| \$k1 | 27 | 0x00000000 |
| \$gp | 28 | 0x10008000 |
| \$sp | 29 | 0x7fffffc0 |
| \$fp | 30 | 0x00000000 |
| \$ra | 31 | 0x00000000 |
| pc | | 0x00400024 |
| hi | | 0x00000000 |
| lo | | 0x00000000 |

// programa 10

Considere o seguinte programa: $y = 127x - 65z + 1$

Faça um programa que calcule o valor de y conhecendo os valores de x e z . Os valores de x e z

estão armazenados na memória e, na posição imediatamente a seguir, o valor de y deverá ser escrito, ou seja:

.data

x: .word 5

z: .word 7

y: .word 0 # esse valor deverá ser sobrescrito após a execução do programa.

The screenshot displays the MARS MIPS simulator interface. The top pane shows the assembly code for 'mips1.asm'. The bottom-left pane shows the 'Text Segment' with columns for Bkpt, Address, Code, Basic, and Source. The bottom-right pane shows the 'Data Segment' with columns for Address, Value (+0), Value (+4), Value (+8), Value (+c), Value (+10), Value (+14), Value (+18), and Value (+1c). The 'Name' column on the right lists registers and memory locations.

Assembly Code (mips1.asm):

```
1 .data
2 x: .word 5
3 z: .word 7
4 y: .word 0 # esse valor deverá ser sobrescrito após a execução do programa.
5
6 # y = 127x - 65z + 1
7 # s0 -> x
8 # s1 -> y
9 # s2 -> z
10
11 .text
12 lui $t0, 0x1001
13 lw $s0, 0($t0) # MEM[0] = x
14 lw $s2, 4($t0) # MEM[1] = z
15
16 # calcular 127 * x
17 ori $t1, $zero, 127
18 mult $t1, $s0
19 mflo $t2 # t2 = 127 * x
20
21 # calcular 65 * z
22 ori $t3, $zero, 65
23 mult $t3, $s2
24 mflo $t4 # t2 = 127 * x
25
26 # fazer 127x - 65z
27 sub $t5, $t2, $t4
28 # adicionar mais 1
29 addi $t5, $t5, 1
30
31 # retornar valor de t2 pra memoria
32 sw $t5, 8($t0)
33
```

Text Segment:

| Bkpt | Address | Code | Basic | Source |
|------|------------|------------|-------------------------------|-----------------------------------|
| | 0x00400000 | 0x3c081001 | lui \$t0, 0x1001 | 12: lui \$t0, 0x1001 |
| | 0x00400004 | 0x8d100000 | lw \$s0, 0(\$t0) # MEM[0] = x | 13: lw \$s0, 0(\$t0) # MEM[0] = x |
| | 0x00400008 | 0x8d120004 | lw \$s2, 4(\$t0) # MEM[1] = z | 14: lw \$s2, 4(\$t0) # MEM[1] = z |
| | 0x0040000c | 0x3409007f | ori \$t1, \$zero, 127 | 17: ori \$t1, \$zero, 127 |
| | 0x00400010 | 0x01300018 | mult \$t1, \$s0 | 18: mult \$t1, \$s0 |
| | 0x00400014 | 0x0005012 | mflo \$t2 | 19: mflo \$t2 # t2 = 127 * x |
| | 0x00400018 | 0x340b0041 | ori \$t3, \$zero, 65 | 22: ori \$t3, \$zero, 65 |
| | 0x0040001c | 0x01720018 | mult \$t3, \$s2 | 23: mult \$t3, \$s2 |
| | 0x00400020 | 0x0006012 | mflo \$t4 | 24: mflo \$t4 # t2 = 127 * x |
| | 0x00400024 | 0x014c6822 | sub \$t5, \$t2, \$t4 | 27: sub \$t5, \$t2, \$t4 |
| | 0x00400028 | 0x21ad0001 | addi \$t5, \$t5, 1 | 29: addi \$t5, \$t5, 1 |
| | 0x0040002c | 0xad0d0008 | sw \$t5, 8(\$t0) | 32: sw \$t5, 8(\$t0) |

Data Segment:

| Address | Value (+0) | Value (+4) | Value (+8) | Value (+c) | Value (+10) | Value (+14) | Value (+18) | Value (+1c) |
|------------|------------|------------|------------|------------|-------------|-------------|-------------|-------------|
| 0x10010000 | 0x00000005 | 0x00000007 | 0x000000b5 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
| 0x10010020 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
| 0x10010040 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |

Register File:

| Name | Number | Value |
|--------|--------|------------|
| \$zero | 0 | 0x00000000 |
| \$at | 1 | 0x00000000 |
| \$v0 | 2 | 0x00000000 |
| \$v1 | 3 | 0x00000000 |
| \$a0 | 4 | 0x00000000 |
| \$a1 | 5 | 0x00000000 |
| \$a2 | 6 | 0x00000000 |
| \$a3 | 7 | 0x00000000 |
| \$t0 | 8 | 0x10010000 |
| \$t1 | 9 | 0x0000007f |
| \$t2 | 10 | 0x0000027b |
| \$t3 | 11 | 0x00000041 |
| \$t4 | 12 | 0x000001c7 |
| \$t5 | 13 | 0x000000b5 |
| \$t6 | 14 | 0x00000000 |
| \$t7 | 15 | 0x00000000 |
| \$s0 | 16 | 0x00000005 |
| \$s1 | 17 | 0x00000000 |
| \$s2 | 18 | 0x00000007 |
| \$s3 | 19 | 0x00000000 |
| \$s4 | 20 | 0x00000000 |
| \$s5 | 21 | 0x00000000 |
| \$s6 | 22 | 0x00000000 |
| \$s7 | 23 | 0x00000000 |
| \$s8 | 24 | 0x00000000 |
| \$s9 | 25 | 0x00000000 |
| \$k0 | 26 | 0x00000000 |
| \$k1 | 27 | 0x00000000 |
| \$gp | 28 | 0x10008000 |
| \$sp | 29 | 0x7fffffc0 |
| \$fp | 30 | 0x00000000 |
| \$ra | 31 | 0x00000000 |
| pc | | 0x00400030 |
| hi | | 0x00000000 |
| lo | | 0x000001c7 |

// programa 11

Considere o seguinte programa: $y = x - z + 300000$

Faça um programa que calcule o valor de y conhecendo os valores de x e z. Os valores de x e z

estão armazenados na memória e, na posição imediatamente a seguir, o valor de y deverá ser escrito, ou seja:

.data

x: .word 100000

z: .word 200000

y: .word 0 # esse valor deverá ser sobrescrito após a execução do programa.

The screenshot displays a MIPS assembler interface with the following components:

- Assembly Code (mips1.asm):**

```
1 .data
2 x: .word 100000
3 z: .word 200000
4 y: .word 0 # esse valor deverá ser sobrescrito após a execução do programa
5
6 # y = x - z + 300000
7 .text
8 # colocando 300000
9 ori $t1, $zero, 0x493E
10 sll $t1, $t1, 4
11
12 lui $t0, 0x1001
13 lw $s0, 0($t0) # MEM[0] = x
14 lw $s1, 4($t0) # MEM[1] = z
15 sub $t2, $s0, $s1
16 add $t2, $t2, $t1
17 sw $t2, 8($t0)
```
- Text Segment Table:**

| Bkpt | Address | Code | Basic | Source |
|------|------------|------------|-------------------------|-----------------------------------|
| | 0x00400000 | 0x3409493e | ori \$9,\$0,0x0000493e | 9: ori \$t1, \$zero, 0x493E |
| | 0x00400004 | 0x00094900 | sll \$9,\$9,0x00000004 | 10: sll \$t1, \$t1, 4 |
| | 0x00400008 | 0x3c081001 | lui \$8,0x00001001 | 12: lui \$t0,0x1001 |
| | 0x0040000c | 0x8d100000 | lw \$16,0x00000000(\$8) | 13: lw \$s0, 0(\$t0) # MEM[0] = x |
| | 0x00400010 | 0x8d110004 | lw \$17,0x00000004(\$8) | 14: lw \$s1, 4(\$t0) # MEM[1] = z |
| | 0x00400014 | 0x02115022 | sub \$10,\$16,\$17 | 15: sub \$t2, \$s0, \$s1 |
| | 0x00400018 | 0x01495020 | add \$10,\$10,\$9 | 16: add \$t2, \$t2, \$t1 |
| | 0x0040001c | 0xad0a0008 | sw \$10,0x00000008(\$8) | 17: sw \$t2, 8(\$t0) |
- Data Segment Table:**

| Address | Value (+0) | Value (+4) | Value (+8) | Value (+c) | Value (+10) | Value (+14) | Value (+18) | Value (+1c) |
|------------|------------|------------|------------|------------|-------------|-------------|-------------|-------------|
| 0x10010000 | 0x000186a0 | 0x00030d40 | 0x00030d40 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
| 0x10010020 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
| 0x10010040 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
- Registers Table:**

| Name | Number | Value |
|--------|--------|------------|
| \$zero | 0 | 0x00000000 |
| \$at | 1 | 0x00000000 |
| \$v0 | 2 | 0x00000000 |
| \$v1 | 3 | 0x00000000 |
| \$a0 | 4 | 0x00000000 |
| \$a1 | 5 | 0x00000000 |
| \$a2 | 6 | 0x00000000 |
| \$a3 | 7 | 0x00000000 |
| \$t0 | 8 | 0x10010000 |
| \$t1 | 9 | 0x000493e0 |
| \$t2 | 10 | 0x00030d40 |
| \$t3 | 11 | 0x00000000 |
| \$t4 | 12 | 0x00000000 |
| \$t5 | 13 | 0x00000000 |
| \$t6 | 14 | 0x00000000 |
| \$t7 | 15 | 0x00000000 |
| \$s0 | 16 | 0x000186a0 |
| \$s1 | 17 | 0x00030d40 |
| \$s2 | 18 | 0x00000000 |
| \$s3 | 19 | 0x00000000 |
| \$s4 | 20 | 0x00000000 |
| \$s5 | 21 | 0x00000000 |
| \$s6 | 22 | 0x00000000 |
| \$s7 | 23 | 0x00000000 |
| \$s8 | 24 | 0x00000000 |
| \$t9 | 25 | 0x00000000 |
| \$k0 | 26 | 0x00000000 |
| \$k1 | 27 | 0x00000000 |
| \$gp | 28 | 0x10008000 |
| \$sp | 29 | 0x7fffffc |
| \$fp | 30 | 0x00000000 |
| \$ra | 31 | 0x00000000 |
| pc | | 0x00400020 |
| hi | | 0x00000000 |
| lo | | 0x00000000 |

// programa 12

Considere a seguinte situação:

```
int ***x;
```

onde x contem um ponteiro para um ponteiro para um ponteiro para um inteiro.

Nessa situação, considere que a posição inicial de memória contenha o inteiro em questão.

Coloque todos os outros valores em registradores, use os endereços de memória que quiser dentro do espaço de endereçamento do Mips.

Resumo do problema:

$k = \text{MEM} [\text{MEM} [\text{MEM} [x]]]$.

Crie um programa que implemente a estrutura de dados acima, leia o valor de K, o multiplique por 2 e o reescreva no local correto conhecendo-se apenas o valor de x.

```
1  .data
2  ppp_x: .word 0x10010004
3  pp_x:  .word 0x10010008
4  p_x:   .word 0x1001000C
5  x:     .word 10
6
7  .text
8  lui $t0, 0x1001
9  lw $t1, 0($t0) # ponteiro de ponteiro de ponteiro
10 lw $t1, 0($t1) # ponteiro de ponteiro
11 lw $t1, 0($t1) # ponteiro
12 lw $s0, 0($t1) # valor
13 sll $s0, $s0, 1
14 sw $s0, 0($t1) # MEM[t1] = k
15
```

| Name | Number | Value |
|--------|--------|------------|
| \$zero | 0 | 0x00000000 |
| \$at | 1 | 0x00000000 |
| \$v0 | 2 | 0x00000000 |
| \$v1 | 3 | 0x00000000 |
| \$a0 | 4 | 0x00000000 |
| \$a1 | 5 | 0x00000000 |
| \$a2 | 6 | 0x00000000 |
| \$a3 | 7 | 0x00000000 |
| \$t0 | 8 | 0x10010000 |
| \$t1 | 9 | 0x1001000C |
| \$t2 | 10 | 0x00000000 |
| \$t3 | 11 | 0x00000000 |
| \$t4 | 12 | 0x00000000 |
| \$t5 | 13 | 0x00000000 |
| \$t6 | 14 | 0x00000000 |
| \$t7 | 15 | 0x00000000 |
| \$s0 | 16 | 0x00000014 |
| \$s1 | 17 | 0x00000000 |
| \$s2 | 18 | 0x00000000 |
| \$s3 | 19 | 0x00000000 |
| \$s4 | 20 | 0x00000000 |
| \$s5 | 21 | 0x00000000 |
| \$s6 | 22 | 0x00000000 |
| \$s7 | 23 | 0x00000000 |
| \$s8 | 24 | 0x00000000 |
| \$s9 | 25 | 0x00000000 |
| \$k0 | 26 | 0x00000000 |
| \$k1 | 27 | 0x00000000 |
| \$gp | 28 | 0x10008000 |
| \$sp | 29 | 0x7ffffc00 |
| \$fp | 30 | 0x00000000 |
| \$ra | 31 | 0x00000000 |
| pc | | 0x0040001C |
| hi | | 0x00000000 |
| lo | | 0x00000000 |

Para os programas a seguir use instruções de desvio (beq, bne, j)

// programa 13:

Escreva um programa que leia um valor A da memória, identifique se o número é negativo ou não e encontre o seu módulo. O valor deverá ser reescrito sobre A.

The screenshot displays the Mips1.asm editor and its execution interface. The editor shows the following assembly code:

```
1 .text
2 .globl main
3 main:
4     lui $t0, 0x1001
5     lw $s0, 0($t0) # x = A
6     srl $t1, $s0, 31 # armazenar o bit do sinal do numero ( em complemento de dois )
7     beq $t1, $zero, fim # se nao for negativo, ir para o fim
8
9 modulo:
10    sub $s0, $zero, $s0 # 0 - (-4) = 4
11    sw $s0, 0($t0) # retorna o valor pra memoria
12
13
14 fim:
15 .data
16 A: .word -54
17
```

The execution interface shows the registers and data segment. The registers table is as follows:

| Name | Number | Value |
|--------|--------|------------|
| \$zero | 0 | 0x00000000 |
| \$at | 1 | 0x00000000 |
| \$v0 | 2 | 0x00000000 |
| \$v1 | 3 | 0x00000000 |
| \$a0 | 4 | 0x00000000 |
| \$a1 | 5 | 0x00000000 |
| \$a2 | 6 | 0x00000000 |
| \$a3 | 7 | 0x00000000 |
| \$t0 | 8 | 0x10010000 |
| \$t1 | 9 | 0x00000001 |
| \$t2 | 10 | 0x00000000 |
| \$t3 | 11 | 0x00000000 |
| \$t4 | 12 | 0x00000000 |
| \$t5 | 13 | 0x00000000 |
| \$t6 | 14 | 0x00000000 |
| \$t7 | 15 | 0x00000000 |
| \$s0 | 16 | 0x00000036 |
| \$s1 | 17 | 0x00000000 |
| \$s2 | 18 | 0x00000000 |
| \$s3 | 19 | 0x00000000 |
| \$s4 | 20 | 0x00000000 |
| \$s5 | 21 | 0x00000000 |
| \$s6 | 22 | 0x00000000 |
| \$s7 | 23 | 0x00000000 |
| \$s8 | 24 | 0x00000000 |
| \$s9 | 25 | 0x00000000 |
| \$k0 | 26 | 0x00000000 |
| \$k1 | 27 | 0x00000000 |
| \$gp | 28 | 0x10000000 |
| \$sp | 29 | 0x7ffffc00 |
| \$fp | 30 | 0x00000000 |
| \$ra | 31 | 0x00000000 |
| \$pc | | 0x00400018 |
| \$hi | | 0x00000000 |
| \$lo | | 0x00000000 |

The data segment table is as follows:

| Address | Value (+0) | Value (+4) | Value (+8) | Value (+c) | Value (+10) | Value (+14) | Value (+18) | Value (+1c) |
|------------|------------|------------|------------|------------|-------------|-------------|-------------|-------------|
| 0x10010000 | 0x00000036 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
| 0x10010020 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
| 0x10010040 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |

The execution interface also shows the Text Segment and Data Segment tables. The Text Segment table is as follows:

| Bkpt | Address | Code | Basic | Source |
|------|------------|------------|-------------------------|---|
| | 0x00400000 | 0x3c081001 | lui \$8,0x00001001 | 4: lui \$t0, 0x1001 |
| | 0x00400004 | 0x8d100000 | lw \$16,0x00000000(\$8) | 5: lw \$s0, 0(\$t0) # x = A |
| | 0x00400008 | 0x00104fc2 | srl \$9,\$16,0x0000001f | 6: srl \$t1, \$s0, 31 # armazenar o bit do sinal do numero (em ... |
| | 0x0040000c | 0x11200002 | beq \$9,\$0,0x00000002 | 7: beq \$t1, \$zero, fim # se nao for negativo, ir para o fim |
| | 0x00400010 | 0x00108022 | sub \$16,\$0,\$16 | 10: sub \$s0, \$zero, \$s0 # 0 - (-4) = 4 |
| | 0x00400014 | 0xad100000 | sw \$16,0x00000000(\$8) | 11: sw \$s0, 0(\$t0) # retorna o valor pra memoria |

The Data Segment table is as follows:

| Address | Value (+0) | Value (+4) | Value (+8) | Value (+c) | Value (+10) | Value (+14) | Value (+18) | Value (+1c) |
|------------|------------|------------|------------|------------|-------------|-------------|-------------|-------------|
| 0x10010000 | 0x00000036 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
| 0x10010020 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
| 0x10010040 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |

The execution interface also shows the Registers table and the Data Segment table. The Registers table is as follows:

| Name | Number | Value |
|--------|--------|------------|
| \$zero | 0 | 0x00000000 |
| \$at | 1 | 0x00000000 |
| \$v0 | 2 | 0x00000000 |
| \$v1 | 3 | 0x00000000 |
| \$a0 | 4 | 0x00000000 |
| \$a1 | 5 | 0x00000000 |
| \$a2 | 6 | 0x00000000 |
| \$a3 | 7 | 0x00000000 |
| \$t0 | 8 | 0x10010000 |
| \$t1 | 9 | 0x00000001 |
| \$t2 | 10 | 0x00000000 |
| \$t3 | 11 | 0x00000000 |
| \$t4 | 12 | 0x00000000 |
| \$t5 | 13 | 0x00000000 |
| \$t6 | 14 | 0x00000000 |
| \$t7 | 15 | 0x00000000 |
| \$s0 | 16 | 0x00000036 |
| \$s1 | 17 | 0x00000000 |
| \$s2 | 18 | 0x00000000 |
| \$s3 | 19 | 0x00000000 |
| \$s4 | 20 | 0x00000000 |
| \$s5 | 21 | 0x00000000 |
| \$s6 | 22 | 0x00000000 |
| \$s7 | 23 | 0x00000000 |
| \$s8 | 24 | 0x00000000 |
| \$s9 | 25 | 0x00000000 |
| \$k0 | 26 | 0x00000000 |
| \$k1 | 27 | 0x00000000 |
| \$gp | 28 | 0x10000000 |
| \$sp | 29 | 0x7ffffc00 |
| \$fp | 30 | 0x00000000 |
| \$ra | 31 | 0x00000000 |
| \$pc | | 0x00400018 |
| \$hi | | 0x00000000 |
| \$lo | | 0x00000000 |

The Data Segment table is as follows:

| Address | Value (+0) | Value (+4) | Value (+8) | Value (+c) | Value (+10) | Value (+14) | Value (+18) | Value (+1c) |
|------------|------------|------------|------------|------------|-------------|-------------|-------------|-------------|
| 0x10010000 | 0x00000036 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
| 0x10010020 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
| 0x10010040 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |

The execution interface also shows the Registers table and the Data Segment table. The Registers table is as follows:

| Name | Number | Value |
|--------|--------|------------|
| \$zero | 0 | 0x00000000 |
| \$at | 1 | 0x00000000 |
| \$v0 | 2 | 0x00000000 |
| \$v1 | 3 | 0x00000000 |
| \$a0 | 4 | 0x00000000 |
| \$a1 | 5 | 0x00000000 |
| \$a2 | 6 | 0x00000000 |
| \$a3 | 7 | 0x00000000 |
| \$t0 | 8 | 0x10010000 |
| \$t1 | 9 | 0x00000001 |
| \$t2 | 10 | 0x00000000 |
| \$t3 | 11 | 0x00000000 |
| \$t4 | 12 | 0x00000000 |
| \$t5 | 13 | 0x00000000 |
| \$t6 | 14 | 0x00000000 |
| \$t7 | 15 | 0x00000000 |
| \$s0 | 16 | 0x00000036 |
| \$s1 | 17 | 0x00000000 |
| \$s2 | 18 | 0x00000000 |
| \$s3 | 19 | 0x00000000 |
| \$s4 | 20 | 0x00000000 |
| \$s5 | 21 | 0x00000000 |
| \$s6 | 22 | 0x00000000 |
| \$s7 | 23 | 0x00000000 |
| \$s8 | 24 | 0x00000000 |
| \$s9 | 25 | 0x00000000 |
| \$k0 | 26 | 0x00000000 |
| \$k1 | 27 | 0x00000000 |
| \$gp | 28 | 0x10000000 |
| \$sp | 29 | 0x7ffffc00 |
| \$fp | 30 | 0x00000000 |
| \$ra | 31 | 0x00000000 |
| \$pc | | 0x00400018 |
| \$hi | | 0x00000000 |
| \$lo | | 0x00000000 |

The Data Segment table is as follows:

| Address | Value (+0) | Value (+4) | Value (+8) | Value (+c) | Value (+10) | Value (+14) | Value (+18) | Value (+1c) |
|------------|------------|------------|------------|------------|-------------|-------------|-------------|-------------|
| 0x10010000 | 0x00000036 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
| 0x10010020 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
| 0x10010040 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |

The execution interface also shows the Registers table and the Data Segment table. The Registers table is as follows:

| Name | Number | Value |
|--------|--------|------------|
| \$zero | 0 | 0x00000000 |
| \$at | 1 | 0x00000000 |
| \$v0 | 2 | 0x00000000 |
| \$v1 | 3 | 0x00000000 |
| \$a0 | 4 | 0x00000000 |
| \$a1 | 5 | 0x00000000 |
| \$a2 | 6 | 0x00000000 |
| \$a3 | 7 | 0x00000000 |
| \$t0 | 8 | 0x10010000 |
| \$t1 | 9 | 0x00000001 |
| \$t2 | 10 | 0x00000000 |
| \$t3 | 11 | 0x00000000 |
| \$t4 | 12 | 0x00000000 |
| \$t5 | 13 | 0x00000000 |
| \$t6 | 14 | 0x00000000 |
| \$t7 | 15 | 0x00000000 |
| \$s0 | 16 | 0x00000036 |
| \$s1 | 17 | 0x00000000 |
| \$s2 | 18 | 0x00000000 |
| \$s3 | 19 | 0x00000000 |
| \$s4 | 20 | 0x00000000 |
| \$s5 | 21 | 0x00000000 |
| \$s6 | 22 | 0x00000000 |
| \$s7 | 23 | 0x00000000 |
| \$s8 | 24 | 0x00000000 |
| \$s9 | 25 | 0x00000000 |
| \$k0 | 26 | 0x00000000 |
| \$k1 | 27 | 0x00000000 |
| \$gp | 28 | 0x10000000 |
| \$sp | 29 | 0x7ffffc00 |
| \$fp | 30 | 0x00000000 |
| \$ra | 31 | 0x00000000 |
| \$pc | | 0x00400018 |
| \$hi | | 0x00000000 |
| \$lo | | 0x00000000 |

The Data Segment table is as follows:

| Address | Value (+0) | Value (+4) | Value (+8) | Value (+c) | Value (+10) | Value (+14) | Value (+18) | Value (+1c) |
|------------|------------|------------|------------|------------|-------------|-------------|-------------|-------------|
| 0x10010000 | 0x00000036 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
| 0x10010020 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
| 0x10010040 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |

The execution interface also shows the Registers table and the Data Segment table. The Registers table is as follows:

| Name | Number | Value |
|--------|--------|------------|
| \$zero | 0 | 0x00000000 |
| \$at | 1 | 0x00000000 |
| \$v0 | 2 | 0x00000000 |
| \$v1 | 3 | 0x00000000 |
| \$a0 | 4 | 0x00000000 |
| \$a1 | 5 | 0x00000000 |
| \$a2 | 6 | 0x00000000 |
| \$a3 | 7 | 0x00000000 |
| \$t0 | 8 | 0x10010000 |
| \$t1 | 9 | 0x00000001 |
| \$t2 | 10 | 0x00000000 |
| \$t3 | 11 | 0x00000000 |
| \$t4 | 12 | 0x00000000 |
| \$t5 | 13 | 0x00000000 |
| \$t6 | 14 | 0x00000000 |
| \$t7 | 15 | 0x00000000 |
| \$s0 | 16 | 0x00000036 |
| \$s1 | 17 | 0x00000000 |
| \$s2 | 18 | 0x00000000 |
| \$s3 | 19 | 0x00000000 |
| \$s4 | 20 | 0x00000000 |
| \$s5 | 21 | 0x00000000 |
| \$s6 | 22 | 0x00000000 |
| \$s7 | 23 | 0x00000000 |
| \$s8 | 24 | 0x00000000 |
| \$s9 | 25 | 0x00000000 |
| \$k0 | 26 | 0x00000000 |
| \$k1 | 27 | 0x00000000 |
| \$gp | 28 | 0x10000000 |
| \$sp | 29 | 0x7ffffc00 |
| \$fp | 30 | 0x00000000 |
| \$ra | 31 | 0x00000000 |
| \$pc | | 0x00400018 |
| \$hi | | 0x00000000 |
| \$lo | | 0x00000000 |

The Data Segment table is as follows:

| Address | Value (+0) | Value (+4) | Value (+8) | Value (+c) | Value (+10) | Value (+14) | Value (+18) | Value (+1c) |
|------------|------------|------------|------------|------------|-------------|-------------|-------------|-------------|
| 0x10010000 | 0x00000036 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
| 0x10010020 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
| 0x10010040 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |

The execution interface also shows the Registers table and the Data Segment table. The Registers table is as follows:

| Name | Number | Value |
|--------|--------|------------|
| \$zero | 0 | 0x00000000 |
| \$at | 1 | 0x00000000 |
| \$v0 | 2 | 0x00000000 |
| \$v1 | 3 | 0x00000000 |
| \$a0 | 4 | 0x00000000 |
| \$a1 | 5 | 0x00000000 |
| \$a2 | 6 | 0x00000000 |
| \$a3 | 7 | 0x00000000 |
| \$t0 | 8 | 0x10010000 |
| \$t1 | 9 | 0x00000001 |
| \$t2 | 10 | 0x00000000 |
| \$t3 | 11 | 0x00000000 |
| \$t4 | 12 | 0x00000000 |
| \$t5 | 13 | 0x00000000 |
| \$t6 | 14 | 0x00000000 |
| \$t7 | 15 | 0x00000000 |
| \$s0 | 16 | 0x00000036 |
| \$s1 | 17 | 0x00000000 |
| \$s2 | 18 | 0x00000000 |
| \$s3 | 19 | 0x00000000 |
| \$s4 | 20 | 0x00000000 |
| \$s5 | 21 | 0x00000000 |
| \$s6 | 22 | 0x00000000 |
| \$s7 | 23 | 0x00000000 |
| \$s8 | 24 | 0x00000000 |
| \$s9 | 25 | 0x00000000 |
| \$k0 | 26 | 0x00000000 |
| \$k1 | 27 | 0x00000000 |
| \$gp | 28 | 0x10000000 |
| \$sp | 29 | 0x7ffffc00 |
| \$fp | 30 | 0x00000000 |
| \$ra | 31 | 0x00000000 |
| \$pc | | 0x00400018 |
| \$hi | | 0x00000000 |
| \$lo | | 0x00000000 |

The Data Segment table is as follows:

| Address | Value (+0) | Value (+4) | Value (+8) | Value (+c) | Value (+10) | Value (+14) | Value (+18) | Value (+1c) |
|------------|------------|------------|------------|------------|-------------|-------------|-------------|-------------|
| 0x10010000 | 0x00000036 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
| 0x10010020 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
| 0x10010040 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |

The execution interface also shows the Registers table and the Data Segment table. The Registers table is as follows:

| Name | Number | Value |
|--------|--------|------------|
| \$zero | 0 | 0x00000000 |
| \$at | 1 | 0x00000000 |
| \$v0 | 2 | 0x00000000 |
| \$v1 | 3 | 0x00000000 |
| \$a0 | 4 | 0x00000000 |
| \$a1 | 5 | 0x00000000 |
| \$a2 | 6 | 0x00000000 |
| \$a3 | 7 | 0x00000000 |
| \$t0 | 8 | 0x10010000 |
| \$t1 | 9 | 0x00000001 |
| \$t2 | 10 | 0x00000000 |
| \$t3 | 11 | 0x00000000 |
| \$t4 | 12 | 0x00000000 |
| \$t5 | 13 | 0x00000000 |
| \$t6 | 14 | 0x00000000 |
| \$t7 | 15 | 0x00000000 |
| \$s0 | 16 | 0x00000036 |
| | | |

// programa 14:

Escreva um programa que leia um valor A da memória, identifique se o número é par ou não. Um valor deverá ser escrito na segunda posição livre da memória (0 para par e 1 para ímpar).

The screenshot displays a MIPS assembler and simulator interface. The main window shows the assembly code for a program named `mips1.asm`. The code is as follows:

```
1 .data
2 A: .word 2
3 resp: .word -1
4
5 .text
6 lui $t0, 0x1001
7 lw $s0, 0($t0)
8 andi $t1, $s0, 0x1
9 sw $t1, 4($t0) # armazena 0 se for par e 1 se for impar
10
```

The interface includes a "Text Segment" window showing the disassembled code with addresses, codes, and sources. The "Data Segment" window shows the memory layout with addresses and values. The "Registers" window on the right shows the state of the MIPS registers, with `$t1` highlighted as containing `0x00000000`.

| Name | Number | Value |
|-------------|----------|-------------------|
| \$zero | 0 | 0x00000000 |
| \$at | 1 | 0x00000000 |
| \$v0 | 2 | 0x00000000 |
| \$v1 | 3 | 0x00000000 |
| \$a0 | 4 | 0x00000000 |
| \$a1 | 5 | 0x00000000 |
| \$a2 | 6 | 0x00000000 |
| \$a3 | 7 | 0x00000000 |
| \$t0 | 8 | 0x10010000 |
| \$t1 | 9 | 0x00000000 |
| \$t2 | 10 | 0x00000000 |
| \$t3 | 11 | 0x00000000 |
| \$t4 | 12 | 0x00000000 |
| \$t5 | 13 | 0x00000000 |
| \$t6 | 14 | 0x00000000 |
| \$t7 | 15 | 0x00000000 |
| \$s0 | 16 | 0x00000002 |
| \$s1 | 17 | 0x00000000 |
| \$s2 | 18 | 0x00000000 |
| \$s3 | 19 | 0x00000000 |
| \$s4 | 20 | 0x00000000 |
| \$s5 | 21 | 0x00000000 |
| \$s6 | 22 | 0x00000000 |
| \$s7 | 23 | 0x00000000 |
| \$s8 | 24 | 0x00000000 |
| \$s9 | 25 | 0x00000000 |
| \$k0 | 26 | 0x00000000 |
| \$k1 | 27 | 0x00000000 |
| \$gp | 28 | 0x10008000 |
| \$sp | 29 | 0x7ffffc00 |
| \$fp | 30 | 0x00000000 |
| \$ra | 31 | 0x00000000 |
| pc | | 0x00400010 |
| hi | | 0x00000000 |
| lo | | 0x00000000 |

// programa 15:

Escrever um programa que crie um vetor de 100 elementos na memória onde $\text{vetor}[i] = 2*i + 1$.

Após a última posição do vetor criado, escrever a soma de todos os valores armazenados do vetor.

Use o MARS para verificar a quantidade de instruções conforme o tipo (ULA, Desvios, Mem ou Outras)

The image shows the MARS MIPS assembler interface. The main window displays the assembly code for 'mips1.asm'. The code is as follows:

```
1  # $t0 => vet[0]
2  # $s0 -> i
3  # $s1 -> acumulador
4  # $t1 -> TAM
5
6  .text
7  .globl main
8  main:
9      lui $t0, 0x1001
10     ori $s0, $zero, 0      # i = 0
11     ori $s1, $zero, 0      # acumulador = 0
12     ori $t1, $zero, 100    # TAM = 100
13
14     vetor:
15         sll $t2, $s0, 1      # t2 = 2*i
16         addi $t2, $t2, 1      # t2 = 2*i + 1
17         sw $t2, 0($t0)        # vet[i] = 2*i + 1
18         add $s1, $s1, $t2      # acumulador = acumulador + (2*i + 1)
19         addi $t0, $t0, 4      # endereço a ser guardado, aumenta 4 para cada iteração
20         addi $s0, $s0, 1      # i = i + 1
21         bne $s0, $t1, vetor
22
23     fim:
24         sw $s1, 0($t0)        # resultado da soma de todos os números é armazenado
25
```

The interface also shows the memory layout. The Text Segment table is as follows:

| Bkpt | Address | Code | Basic | Source |
|------|------------|------------|-----------------------|--|
| | 0x00400000 | 0x3c081001 | lui \$t0, 0x0001001 | 9: lui \$t0, 0x1001 |
| | 0x00400004 | 0x34100000 | ori \$t2, \$zero, 0 | 10: ori \$s0, \$zero, 0 # i = 0 |
| | 0x00400008 | 0x34100000 | ori \$t2, \$zero, 0 | 11: ori \$s1, \$zero, 0 # acumulador = 0 |
| | 0x0040000c | 0x34090064 | ori \$t2, \$zero, 100 | 12: ori \$t1, \$zero, 100 # TAM = 100 |
| | 0x00400010 | 0x00105040 | sll \$t2, \$s0, 1 | 15: sll \$t2, \$s0, 1 # t2 = 2*i |
| | 0x00400014 | 0x214a0001 | addi \$t2, \$t2, 1 | 16: addi \$t2, \$t2, 1 # t2 = 2*i + 1 |
| | 0x00400018 | 0xad0a0000 | sw \$t2, 0(\$t0) | 17: sw \$t2, 0(\$t0) # vet[i] = 2*i + 1 |
| | 0x0040001c | 0x022a8820 | add \$s1, \$s1, \$t2 | 18: add \$s1, \$s1, \$t2 # acumulador = acumulador + (2*i + 1) |
| | 0x00400020 | 0x21080004 | addi \$t0, \$t0, 4 | 19: addi \$t0, \$t0, 4 # endereço a ser guardado, aumenta 4 para ... |
| | 0x00400024 | 0x22100001 | addi \$s0, \$s0, 1 | 20: addi \$s0, \$s0, 1 # i = i + 1 |
| | 0x00400028 | 0x1609ffff | bne \$s0, \$t1, vetor | 21: bne \$s0, \$t1, vetor |
| | 0x0040002c | 0xad110000 | sw \$s1, 0(\$t0) | 24: sw \$s1, 0(\$t0) # resultado da soma de todos os números é ... |

The Data Segment table is as follows:

| Address | Value (+0) | Value (+4) | Value (+8) | Value (+c) | Value (+10) | Value (+14) | Value (+18) | Value (+1c) |
|------------|------------|------------|------------|------------|-------------|-------------|-------------|-------------|
| 0x10010160 | 0x000000b1 | 0x000000b3 | 0x000000b5 | 0x000000b7 | 0x000000b9 | 0x000000bb | 0x000000bd | 0x000000bf |
| 0x10010180 | 0x000000c1 | 0x000000c3 | 0x000000c5 | 0x000000c7 | 0x000000c9 | 0x000000cb | 0x000000cd | 0x000000cf |
| 0x100101a0 | 0x000000d1 | 0x000000d3 | 0x000000d5 | 0x000000d7 | 0x000000d9 | 0x000000db | 0x000000dd | 0x000000df |

The right panel shows the register values:

| Name | Number | Value |
|--------|--------|------------|
| \$zero | 0 | 0x00000000 |
| \$at | 1 | 0x00000000 |
| \$v0 | 2 | 0x00000000 |
| \$v1 | 3 | 0x00000000 |
| \$a0 | 4 | 0x00000000 |
| \$a1 | 5 | 0x00000000 |
| \$a2 | 6 | 0x00000000 |
| \$a3 | 7 | 0x00000000 |
| \$t0 | 8 | 0x10010190 |
| \$t1 | 9 | 0x00000064 |
| \$t2 | 10 | 0x000000c7 |
| \$t3 | 11 | 0x00000000 |
| \$t4 | 12 | 0x00000000 |
| \$t5 | 13 | 0x00000000 |
| \$t6 | 14 | 0x00000000 |
| \$t7 | 15 | 0x00000000 |
| \$s0 | 16 | 0x00000064 |
| \$s1 | 17 | 0x00002710 |
| \$s2 | 18 | 0x00000000 |
| \$s3 | 19 | 0x00000000 |
| \$s4 | 20 | 0x00000000 |
| \$s5 | 21 | 0x00000000 |
| \$s6 | 22 | 0x00000000 |
| \$s7 | 23 | 0x00000000 |
| \$s8 | 24 | 0x00000000 |
| \$s9 | 25 | 0x00000000 |
| \$k0 | 26 | 0x00000000 |
| \$k1 | 27 | 0x00000000 |
| \$gp | 28 | 0x10008000 |
| \$sp | 29 | 0x7ffffc |
| \$fp | 30 | 0x00000000 |
| \$ra | 31 | 0x00000000 |
| pc | | 0x00400030 |
| hi | | 0x00000000 |
| lo | | 0x00000000 |

// programa 16

Escreva um programa que avalie a expressão: $(x*y)/z$.

Use $x = 1600000$ ($=0x186A00$), $y = 80000$ ($=0x13880$), e $z = 400000$ ($=0x61A80$). Inicializar os registradores com os valores acima.

mips1.asm

```
1  # $s0 -> x
2  # $s1 -> y
3  # $s2 -> z
4
5  .text
6  .globl main
7  main:
8      ori $t0, $zero, 0x186A
9      sll $s0, $t0, 8      # x = 1.600.000
10     ori $t0, $zero, 0x1388
11     sll $s1, $t0, 4      # y = 80.000
12     ori $t0, $zero, 0x6A8
13     sll $s2, $t0, 4      # z = 400.000
14
15     div $s0, $s2        # lo = 1.600.000 / 400.000
16     mflo $t0            # t0 = 4
17     mult $t0, $s1       # lo = 4 * 80.000
18     mflo $s2            # s2 = lo
19
```

Text Segment

| Bkpt | Address | Code | Basic | Source |
|------|------------|------------|-------------------------|---|
| | 0x00400000 | 0x3408186a | ori \$s0,\$0,0x0000186a | 8: ori \$t0, \$zero, 0x186A |
| | 0x00400004 | 0x00088200 | sll \$t6,\$8,0x00000008 | 9: sll \$s0, \$t0, 8 # x = 1.600.000 |
| | 0x00400008 | 0x34081388 | ori \$s0,\$0,0x00001388 | 10: ori \$t0, \$zero, 0x1388 |
| | 0x0040000c | 0x00088900 | sll \$t7,\$8,0x00000004 | 11: sll \$s1, \$t0, 4 # y = 80.000 |
| | 0x00400010 | 0x340806a8 | ori \$s0,\$0,0x000006a8 | 12: ori \$t0, \$zero, 0x6A8 |
| | 0x00400014 | 0x00089100 | sll \$t8,\$8,0x00000004 | 13: sll \$s2, \$t0, 4 # z = 400.000 |
| | 0x00400018 | 0x0212001a | div \$t6,\$t8 | 15: div \$s0, \$s2 # lo = 1.600.000 / 400.000 |
| | 0x0040001c | 0x00004012 | mflo \$8 | 16: mflo \$t0 # t0 = 4 |
| | 0x00400020 | 0x01110018 | mult \$8,\$t7 | 17: mult \$t0, \$s1 # lo = 4 * 80.000 |
| | 0x00400024 | 0x00009012 | mflo \$t8 | 18: mflo \$s2 # s2 = lo |

Data Segment

| Address | Value (+0) | Value (+4) | Value (+8) | Value (+c) | Value (+10) | Value (+14) | Value (+18) | Value (+1c) |
|------------|------------|------------|------------|------------|-------------|-------------|-------------|-------------|
| 0x10010000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
| 0x10010020 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
| 0x10010040 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |

0x10010000 (.data)

☒ Hexadecimal Addresses ☒ Hexadecimal Values ☐ ASCII

| Name | Number | Value |
|--------|--------|-------------|
| \$zero | 0 | 0x00000000 |
| \$at | 1 | 0x00000000 |
| \$v0 | 2 | 0x00000000 |
| \$v1 | 3 | 0x00000000 |
| \$a0 | 4 | 0x00000000 |
| \$a1 | 5 | 0x00000000 |
| \$a2 | 6 | 0x00000000 |
| \$a3 | 7 | 0x00000000 |
| \$t0 | 8 | 0x00000003a |
| \$t1 | 9 | 0x00000000 |
| \$t2 | 10 | 0x00000000 |
| \$t3 | 11 | 0x00000000 |
| \$t4 | 12 | 0x00000000 |
| \$t5 | 13 | 0x00000000 |
| \$t6 | 14 | 0x00000000 |
| \$t7 | 15 | 0x00000000 |
| \$s0 | 16 | 0x00186a00 |
| \$s1 | 17 | 0x00013880 |
| \$s2 | 18 | 0x0046cd00 |
| \$s3 | 19 | 0x00000000 |
| \$s4 | 20 | 0x00000000 |
| \$s5 | 21 | 0x00000000 |
| \$s6 | 22 | 0x00000000 |
| \$s7 | 23 | 0x00000000 |
| \$s8 | 24 | 0x00000000 |
| \$s9 | 25 | 0x00000000 |
| \$k0 | 26 | 0x00000000 |
| \$k1 | 27 | 0x00000000 |
| \$gp | 28 | 0x10008000 |
| \$sp | 29 | 0x7ffffc |
| \$fp | 30 | 0x00000000 |
| \$ra | 31 | 0x00000000 |
| pc | | 0x00400028 |
| hi | | 0x00000000 |
| lo | | 0x0046cd00 |

// programa 17

Para a expressão a seguir, escreva um programa que calcule o valor de k:

$k = x * y$ (Você deverá realizar a multiplicação através de somas!)

O valor de x deve ser lido da primeira posição livre da memória e o valor de y deverá lido da segunda posição livre. O valor de k, após calculado, deverá ainda ser escrito na terceira posição livre da memória.

The screenshot displays a MIPS assembler interface with the following components:

- Assembly Code (mips1.asm):**

```
1  # $s0 -> endereço base
2  # $s1 -> x
3  # $s2 -> k
4  # $s3 -> i
5
6  .text
7  .globl main
8  main:
9      lui $t0, 0x1001          # Carrega o endereço base em $t0
10     lw $s0, 0($t0)           # x = MEM[$t0]
11     lw $s1, 4($t0)           # y = MEM[$t0 + 4]
12     ori $s3, $zero, 0        # i = 0
13
14     multiplicacao:
15         add $s2, $s2, $s0      # k = k + x
16         addi $s3, $s3, 1       # i = i + 1
17         bne $s3, $s1, multiplicacao # Se i != y, volta para "multiplicacao"
18
19     resultado:
20         sw $s2, 8($t0)         # Salva o resultado em MEM[$t0 + 8]
21
22     .data
23     x: .word 20
24     y: .word 10
25     k: .word -1              # Será sobrescrito no final
26
```
- Text Segment Table:**

| Bkpt | Address | Code | Basic | Source |
|------|------------|-------------------------------|-------|---|
| 9 | 0x00400000 | lui \$t0, 0x1001 | 9: | lui \$t0, 0x1001 # Carrega o endereço base em \$t0 |
| 10 | 0x00400004 | lw \$s0, 0(\$t0) | 10: | lw \$s0, 0(\$t0) # x = MEM[\$t0] |
| 11 | 0x00400008 | lw \$s1, 4(\$t0) | 11: | lw \$s1, 4(\$t0) # y = MEM[\$t0 + 4] |
| 12 | 0x0040000c | ori \$s3, \$zero, 0 | 12: | ori \$s3, \$zero, 0 # i = 0 |
| 15 | 0x00400010 | add \$s2, \$s2, \$s0 | 15: | add \$s2, \$s2, \$s0 # k = k + x |
| 16 | 0x00400014 | addi \$s3, \$s3, 1 | 16: | addi \$s3, \$s3, 1 # i = i + 1 |
| 17 | 0x00400018 | bne \$s3, \$s1, multiplicacao | 17: | bne \$s3, \$s1, multiplicacao # Se i != y, volta para "multiplicacao" |
| 20 | 0x0040001c | sw \$s2, 8(\$t0) | 20: | sw \$s2, 8(\$t0) # Salva o resultado em MEM[\$t0 + 8] |
- Data Segment Table:**

| Address | Value (+0) | Value (+4) | Value (+8) | Value (+c) | Value (+10) | Value (+14) | Value (+18) | Value (+1c) |
|------------|------------|------------|------------|------------|-------------|-------------|-------------|-------------|
| 0x10010000 | 0x00000014 | 0x0000000a | 0x000000c8 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
| 0x10010020 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
| 0x10010040 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
- Registers Table:**

| Name | Number | Value |
|--------|--------|------------|
| \$zero | 0 | 0x00000000 |
| \$at | 1 | 0x00000000 |
| \$v0 | 2 | 0x00000000 |
| \$v1 | 3 | 0x00000000 |
| \$a0 | 4 | 0x00000000 |
| \$a1 | 5 | 0x00000000 |
| \$a2 | 6 | 0x00000000 |
| \$a3 | 7 | 0x00000000 |
| \$t0 | 8 | 0x10010000 |
| \$t1 | 9 | 0x00000000 |
| \$t2 | 10 | 0x00000000 |
| \$t3 | 11 | 0x00000000 |
| \$t4 | 12 | 0x00000000 |
| \$t5 | 13 | 0x00000000 |
| \$t6 | 14 | 0x00000000 |
| \$t7 | 15 | 0x00000000 |
| \$s0 | 16 | 0x00000014 |
| \$s1 | 17 | 0x0000000a |
| \$s2 | 18 | 0x000000c8 |
| \$s3 | 19 | 0x0000000a |
| \$s4 | 20 | 0x00000000 |
| \$s5 | 21 | 0x00000000 |
| \$s6 | 22 | 0x00000000 |
| \$s7 | 23 | 0x00000000 |
| \$s8 | 24 | 0x00000000 |
| \$s9 | 25 | 0x00000000 |
| \$k0 | 26 | 0x00000000 |
| \$k1 | 27 | 0x00000000 |
| \$gp | 28 | 0x10008000 |
| \$sp | 29 | 0x7fffffc |
| \$fp | 30 | 0x00000000 |
| \$ra | 31 | 0x00000000 |
| pc | | 0x00400020 |
| hi | | 0x00000000 |
| lo | | 0x00000000 |

// programa 18

Para a expressão a seguir, escreva um programa que calcule o valor de k:

$k = x^y$

Obs: Você poderá utilizar o exercício anterior.

O valor de x deve ser lido da primeira posição livre da memória e o valor de y deverá lido da segunda posição livre. O valor de k, após calculado, deverá ainda ser escrito na terceira posição livre da memória.

Dê um valor para x e y (dê valores pequenos !!) e use o MARS para verificar a quantidade de instruções conforme o tipo (ULA, Desvios, Mem ou Outras)

mips1.asm

```
1 .globl main
2 main:
3     # Carregar valores de x e y da memória
4     lui $t0, 0x1001      # Carregar base do endereço
5     lw $t1, 0($t0)       # Carregar x em $t1
6     lw $t2, 4($t0)       # Carregar y em $t2
7     ori $t3, $zero, 1    # Inicializar $t3 (resultado inicial = 1)
8     ori $t4, $zero, 0    # Inicializar contador de expoente
9
10    potencia_loop:
11        beq $t4, $t2, resultado_k # Se o contador for igual a y, terminou a exponenciação
12        mul $t3, $t3, $t1        # Multiplica o resultado atual por x
13        addi $t4, $t4, 1         # Incrementa o contador
14        j potencia_loop         # Volta para multiplicar novamente
15
16    resultado_k:
17        sw $t3, 8($t0)          # Armazena o resultado (k) na posição de memória 3
18        # Fim do programa
19        li $v0, 10              # Sair do programa
20        syscall
21
22    .data
23    x: .word 3                  # Exemplo de valor de x
24    y: .word 4                  # Exemplo de valor de y
25    z: .word -1                 # Para armazenar o resultado k
26
```

Text Segment

| Bkpt | Address | Code | Basic | Source |
|------|------------|------------|--------------------------|---|
| | 0x00400000 | 0x3c081001 | lui \$8,0x00001001 | 4: lui \$t0, 0x1001 # Carregar base do endereço |
| | 0x00400004 | 0x8d090000 | lw \$9,0x00000000(\$8) | 5: lw \$t1, 0(\$t0) # Carregar x em \$t1 |
| | 0x00400008 | 0x8d0a0004 | lw \$10,0x00000004(\$8) | 6: lw \$t2, 4(\$t0) # Carregar y em \$t2 |
| | 0x0040000c | 0x340b0001 | ori \$11,\$0,0x00000001 | 7: ori \$t3, \$zero, 1 # Inicializar \$t3 (resultado inicial = 1) |
| | 0x00400010 | 0x340c0000 | ori \$12,\$0,0x00000000 | 8: ori \$t4, \$zero, 0 # Inicializar contador de expoente |
| | 0x00400014 | 0x118a0003 | beq \$12,\$10,0x00000003 | 11: beq \$t4, \$t2, resultado_k # Se o contador for igual a y, ter... |
| | 0x00400018 | 0x71695802 | mul \$11,\$11,\$9 | 12: mul \$t3, \$t3, \$t1 # Multiplica o resultado atual por x |
| | 0x0040001c | 0x218c0001 | addi \$12,\$12,0x0000... | 13: addi \$t4, \$t4, 1 # Incrementa o contador |
| | 0x00400020 | 0x08100005 | j 0x00400014 | 14: j potencia_loop # Volta para multiplicar novamente |
| | 0x00400024 | 0xad0b0008 | sw \$11,0x00000008(\$8) | 17: sw \$t3, 8(\$t0) # Armazena o resultado (k) na posição de ... |
| | 0x00400028 | 0x2402000a | addiu \$2,\$0,0x0000000a | 19: li \$v0, 10 # Sair do programa |
| | 0x0040002c | 0x0000000c | syscall | 20: syscall |

Data Segment

| Address | Value (+0) | Value (+4) | Value (+8) | Value (+c) | Value (+10) | Value (+14) | Value (+18) | Value (+1c) |
|------------|------------|------------|------------|------------|-------------|-------------|-------------|-------------|
| 0x10010000 | 0x00000003 | 0x00000004 | 0x00000051 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
| 0x10010020 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
| 0x10010040 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |

0x10010000 (.data) [X] Hexadecimal Addresses [X] Hexadecimal Values [] ASCII

| Name | Number | Value |
|--------|--------|------------|
| \$zero | 0 | 0x00000000 |
| \$at | 1 | 0x00000000 |
| \$v0 | 2 | 0x0000000a |
| \$v1 | 3 | 0x00000000 |
| \$a0 | 4 | 0x00000000 |
| \$a1 | 5 | 0x00000000 |
| \$a2 | 6 | 0x00000000 |
| \$a3 | 7 | 0x00000000 |
| \$t0 | 8 | 0x10010000 |
| \$t1 | 9 | 0x00000003 |
| \$t2 | 10 | 0x00000004 |
| \$t3 | 11 | 0x00000051 |
| \$t4 | 12 | 0x00000004 |
| \$t5 | 13 | 0x00000000 |
| \$t6 | 14 | 0x00000000 |
| \$t7 | 15 | 0x00000000 |
| \$s0 | 16 | 0x00000000 |
| \$s1 | 17 | 0x00000000 |
| \$s2 | 18 | 0x00000000 |
| \$s3 | 19 | 0x00000000 |
| \$s4 | 20 | 0x00000000 |
| \$s5 | 21 | 0x00000000 |
| \$s6 | 22 | 0x00000000 |
| \$s7 | 23 | 0x00000000 |
| \$t8 | 24 | 0x00000000 |
| \$t9 | 25 | 0x00000000 |
| \$k0 | 26 | 0x00000000 |
| \$k1 | 27 | 0x00000000 |
| \$gp | 28 | 0x10008000 |
| \$sp | 29 | 0x7ffffc |
| \$fp | 30 | 0x00000000 |
| \$ra | 31 | 0x00000000 |
| pc | | 0x00400030 |
| hi | | 0x00000000 |
| lo | | 0x00000051 |

Desafio:

Todos viram durante a parte aritmética que podemos utilizar a ULA e registradores para multiplicar dois números através de 3 algoritmos.

Você deverá escrever um programa que leia dois números da memória (primeira posição e segunda posição) os multiplique e coloque o resultado na terceira posição a memória.

Procure usar a versão 3 do algoritmo de multiplicação, pode ser mais simples !!

Atenção que, ao multiplicarmos dois números de 32 bits a resposta poderá ser um número de 64 bits, assim a resposta deverá estar contida em dois registradores temporários, um armazenará a parte superior do número e outro a parte inferior, portanto duas posições de memória serão escritas (a terceira e a quarta).

mips1.asm

```
1 .globl main
2 main:
3     # Carregar valores de x e y da memória
4     lui $t0, 0x1001 # Carregar base do endereço
5     lw $t1, 0($t0) # Carregar x em $t1
6     lw $t2, 4($t0) # Carregar y em $t2
7     ori $t3, $zero, 0 # Inicializar resultado inferior (parte baixa) em $t3
8     ori $t4, $zero, 0 # Inicializar resultado superior (parte alta) em $t4
9     ori $t5, $zero, 0 # Inicializar contador (bit de multiplicação) em $t5
10
11 # Algoritmo de multiplicação por shift-and-add
12 multiplicacao_loop:
13     beq $t5, 32, resultado_k # Se o contador de bits chegar a 32, terminou a multiplicação
14
15     andi $t6, $t2, 1 # Verifica o bit menos significativo de y (multiplicador)
16     beq $t6, $zero, shift_step # Se o bit é 0, apenas desloca
17
18     # Se o bit de y for 1, soma o valor de x (movido) ao resultado inferior
19     add $t3, $t3, $t1 # Soma o valor de x ao resultado inferior
20     # Verifica se houve overflow no resultado inferior (soma maior que 32 bits)
21     bltz $t3, carry_bit # Se houve overflow, ajuste no resultado superior
22     j shift_step
23
24 carry_bit:
25     addi $t4, $t4, 1 # Ajusta a parte superior (overflow)
26
27 shift_step:
28     srl $t1, $t1, 1 # Desloca x para a esquerda (multiplica por 2)
29     srl $t2, $t2, 1 # Desloca y para a direita (divide por 2)
30     addi $t5, $t5, 1 # Incrementa o contador de bits
31     j multiplicacao_loop # Continua o loop de multiplicação
32
33 resultado_k:
34     # Armazena a parte inferior do resultado
35     sw $t3, 0($t0) # Armazena a parte inferior (k baixa) na posição de memória 3
36     # Armazena a parte superior do resultado
37     sw $t4, 12($t0) # Armazena a parte superior (k alta) na posição de memória 4
38
39 # Fim do programa
40 li $v0, 10 # Sair do programa
41 syscall
42
43 .data
44 x: .word 3 # Exemplo de valor de x
45 y: .word 5 # Exemplo de valor de y
46 z1: .word -1 # Para armazenar a parte baixa do resultado k
47 z2: .word -1 # Para armazenar a parte alta do resultado k
48
```

Text Segment

| Bkpt | Address | Code | Basic | Source |
|------|------------|------------|--------------------------|---|
| | 0x00400000 | 0x3c081001 | lui \$8,0x00001001 | 4: lui \$t0, 0x1001 # Carregar base do endereço |
| | 0x00400004 | 0x8d090000 | lw \$9,0x00000000(\$8) | 5: lw \$t1, 0(\$t0) # Carregar x em \$t1 |
| | 0x00400008 | 0x8d0a0004 | lw \$10,0x00000004(\$8) | 6: lw \$t2, 4(\$t0) # Carregar y em \$t2 |
| | 0x0040000c | 0x340b0000 | ori \$11,\$0,0x00000000 | 7: ori \$t3, \$zero, 0 # Inicializar resultado inferior (parte b... |
| | 0x00400010 | 0x340c0000 | ori \$12,\$0,0x00000000 | 8: ori \$t4, \$zero, 0 # Inicializar resultado superior (parte a... |
| | 0x00400014 | 0x340d0000 | ori \$13,\$0,0x00000000 | 9: ori \$t5, \$zero, 0 # Inicializar contador (bit de multiplica... |
| | 0x00400018 | 0x20010020 | addi \$1,\$0,0x00000020 | 13: beq \$t5, 32, resultado_k # Se o contador de bits chegar a 32... |
| | 0x0040001c | 0x102d000a | beq \$1,\$13,0x0000000a | |
| | 0x00400020 | 0x314e0001 | andi \$14,\$10,0x0000... | 15: andi \$t6, \$t2, 1 # Verifica o bit menos significativo de y... |
| | 0x00400024 | 0x11c00004 | beq \$14,\$0,0x00000004 | 16: beq \$t6, \$zero, shift_step # Se o bit é 0, apenas desloca |
| | 0x00400028 | 0x016e5820 | add \$11,\$11,\$9 | 19: add \$t3, \$t3, \$t1 # Soma o valor de x ao resultado inferior |
| | 0x0040002c | 0x05600001 | bltz \$11,0x00000001 | 21: bltz \$t3, carry_bit # Se houve overflow, ajuste no resultado ... |
| | 0x00400030 | 0x0810000e | j 0x00400038 | 22: j shift_step |
| | 0x00400034 | 0x218c0001 | addi \$12,\$12,0x0000... | 25: addi \$t4, \$t4, 1 # Ajusta a parte superior (overflow) |
| | 0x00400038 | 0x00094840 | srl \$9,\$9,0x00000001 | 28: srl \$t1, \$t1, 1 # Desloca x para a esquerda (multiplica p... |
| | 0x0040003c | 0x000a5042 | srl \$10,\$10,0x00000001 | 29: srl \$t2, \$t2, 1 # Desloca y para a direita (divide por 2) |
| | 0x00400040 | 0x218d0001 | addi \$19,\$13,0x0000... | 30: addi \$t5, \$t5, 1 # Incrementa o contador de bits |
| | 0x00400044 | 0x0810000e | j 0x00400018 | 31: j multiplicacao_loop # Continua o loop de multiplicação |
| | 0x00400048 | 0xad0b0008 | sw \$11,0x00000008(\$8) | 35: sw \$t3, 0(\$t0) # Armazena a parte inferior (k baixa) na ... |
| | 0x0040004c | 0xad0c000c | sw \$12,0x0000000c(\$8) | 37: sw \$t4, 12(\$t0) # Armazena a parte superior (k alta) na p... |
| | 0x00400050 | 0x2402000a | addiu \$2,\$0,0x0000000a | 40: li \$v0, 10 # Sair do programa |
| | 0x00400054 | 0x0000000c | syscall | 41: syscall |

Data Segment

| Address | Value (+0) | Value (+4) | Value (+8) | Value (+c) | Value (+10) | Value (+14) | Value (+18) | Value (+1c) |
|------------|------------|------------|------------|------------|-------------|-------------|-------------|-------------|
| 0x10010000 | 0x00000003 | 0x00000005 | 0x0000000f | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |
| 0x10010020 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 | 0x00000000 |

| Name | Number | Value |
|--------|--------|------------|
| \$zero | 0 | 0x00000000 |
| \$at | 1 | 0x00000020 |
| \$v0 | 2 | 0x0000000a |
| \$v1 | 3 | 0x00000000 |
| \$a0 | 4 | 0x00000000 |
| \$a1 | 5 | 0x00000000 |
| \$a2 | 6 | 0x00000000 |
| \$a3 | 7 | 0x00000000 |
| \$t0 | 8 | 0x10010000 |
| \$t1 | 9 | 0x00000000 |
| \$t2 | 10 | 0x00000000 |
| \$t3 | 11 | 0x0000000f |
| \$t4 | 12 | 0x00000000 |
| \$t5 | 13 | 0x00000020 |
| \$t6 | 14 | 0x00000000 |
| \$t7 | 15 | 0x00000000 |
| \$a0 | 16 | 0x00000000 |
| \$s1 | 17 | 0x00000000 |
| \$s2 | 18 | 0x00000000 |
| \$s3 | 19 | 0x00000000 |
| \$s4 | 20 | 0x00000000 |
| \$s5 | 21 | 0x00000000 |
| \$s6 | 22 | 0x00000000 |
| \$s7 | 23 | 0x00000000 |
| \$t8 | 24 | 0x00000000 |
| \$t9 | 25 | 0x00000000 |
| \$k0 | 26 | 0x00000000 |
| \$k1 | 27 | 0x00000000 |
| \$gp | 28 | 0x10008000 |
| \$sp | 29 | 0x7ffffcfc |
| \$fp | 30 | 0x00000000 |
| \$ra | 31 | 0x00000000 |
| \$PC | | 0x00400058 |
| \$hi | | 0x00000000 |
| \$lo | | 0x00000000 |

Parte 3

- 1. Se tivermos 2 inteiros, cada um com 32 bits, quantos bits podemos esperar para o produto?**
C. 64
- 2. Quais os registradores que armazenam os resultados na multiplicação?**
B. hi e lo
- 3. Qual a operação usada para multiplicar inteiros em comp. de dois?**
A. mult
- 4. Qual instrução move os bits menos significativos da multiplicação para o reg. 8?**
C. mflo \$8
- 5. Se tivermos dois inteiros, cada um com 32 bits, quantos bits deveremos estar preparados para receber no quociente?**
B. 32
- 6. Após a instrução div, qual registrador possui o quociente?**
A. lo
- 7. Qual a inst. Usada para dividir dois inteiros em comp. de dois?**
D. div
- 8. Faça um arithmetic shift right de dois no seguinte padrão de bits: 1001 1011**
B. 0010 0110
- 9. Qual o efeito de um arithmetic shift right de uma posição?**
A. Se o inteiro for unsigned, o shift o divide por 2. Se o inteiro for signed, o shift o divide por 2.
- 10. Qual sequencia de instruções avalia $3x+7$, onde x é iniciado no reg. \$8 e o resultado armazenado em \$9?**
A. ori \$3,\$0,3
 mult \$8,\$3
 mflo \$9
 addi \$9,\$9,7

Parte 4

// programa 19

Escrever um programa que leia dois números da memória, a primeira e segunda posições respectivamente (os coloque em \$s0 e \$s1) e determine a quantidade de bits significantes de cada um. Coloque as respostas em \$t0 e \$t1, a partir desse resultado faça a multiplicação. Caso o número de bits significantes de ambos seja menor do que 32 a resposta deverá estar apenas em \$s2, caso contrário a resposta estará em \$s2 e \$s3 (LO e HI respectivamente). Para os exercícios a seguir, considere as variáveis com números abaixo de 16 bits, salvo se mencionado ao contrário.

```
1  # inicio
2
3  .data
4  x: .word 3
5  y: .word 4
6  .text
7  .globl main
8  main:
9      #t0 -> 0x10010000 (first position)
10     #x -> s0
11     #y -> s1
12     #k -> s2
13     ori $t2, $0, 0x1001    #t2 = 0x1001
14     sll $t2, $t2, 16       #t2 = 0x10010000
15     lw $s0, 0($t2)         #s0 = MEM[$t0]
16     lw $s1, 4($t2)         #s1 = MEM[$t0+4]
17     or $t3, $s0, $0        #t3 = x
18     or $t4, $s1, $0        #t4 = y
19     or $t0, $0, $0         #t0 = 0 (contador de x)
20     or $t1, $0, $0         #t1 = 0 (contador de y)
21 if:  beq $t3, $0, if2      #if(t3==0) (goto fim1)
22     addi $t0, $t0, 1       #t0 = t0 + 1
23     srl $t3, $t3, 1        #t3 >> 1
24     j if
25 if2: beq $t4, $0, fim      #if(t4==0) (goto fim2)
26     addi $t1, $t1, 1       #t1 = t1 + 1
27     srl $t4, $t4, 1        #t4 >> 1
28     j if2
29 fim: mult $t0, $t1         #t0 * t1
30     mflo $t5               #t5 = t0
31     slti $t6, $t5, 32      #if(t5<32) (t6=1) else (t6 = 0)
32     bne $t6, 1, maior     #if(t6!=1) (goto maior)
33     or $s2, $0, $t5        #s2 = t5
34     mflo $s2               #s2 = t0
35     j fim2
36 maior: mfhi $s2            #s2 = hi
37     mflo $s3               #s3 = lo
38 fim2:
39
40 # fim
```

| Bkpt | Address | Code | Basic | Source |
|------|------------|------------|---------------------------|---|
| | 0x00400000 | 0x349e1001 | ori \$t0, \$0, 0x10010000 | 13: ori \$t2, \$0, 0x1001 #t2 = 0x1001 |
| | 0x00400004 | 0x000a5400 | sll \$t2, \$t2, 16 | 14: sll \$t2, \$t2, 16 #t2 = 0x10010000 |
| | 0x00400008 | 0x8d500000 | lw \$s0, 0(\$t2) | 15: lw \$s0, 0(\$t2) #s0 = MEM[\$t0] |
| | 0x0040000c | 0x8d510004 | lw \$s1, 4(\$t2) | 16: lw \$s1, 4(\$t2) #s1 = MEM[\$t0+4] |
| | 0x00400010 | 0x02005825 | or \$t3, \$s0, \$0 | 17: or \$t3, \$s0, \$0 #t3 = x |
| | 0x00400014 | 0x02206025 | or \$t4, \$s1, \$0 | 18: or \$t4, \$s1, \$0 #t4 = y |
| | 0x00400018 | 0x00004025 | or \$t0, \$0, \$0 | 19: or \$t0, \$0, \$0 #t0 = 0 (contador de x) |
| | 0x0040001c | 0x00004825 | or \$t1, \$0, \$0 | 20: or \$t1, \$0, \$0 #t1 = 0 (contador de y) |
| | 0x00400020 | 0x11600003 | beq \$t3, \$0, if2 | 21: if: beq \$t3, \$0, if2 #if(t3==0) (goto fim1) |
| | 0x00400024 | 0x21080001 | addi \$t0, \$t0, 1 | 22: addi \$t0, \$t0, 1 #t0 = t0 + 1 |
| | 0x00400028 | 0x000b5842 | srl \$t3, \$t3, 1 | 23: srl \$t3, \$t3, 1 #t3 >> 1 |
| | 0x0040002c | 0x08100008 | j if | 24: j if #goto if |
| | 0x00400030 | 0x11800003 | beq \$t4, \$0, fim | 25: if2: beq \$t4, \$0, fim #if(t4==0) (goto fim2) |
| | 0x00400034 | 0x21290001 | addi \$t1, \$t1, 1 | 26: addi \$t1, \$t1, 1 #t1 = t1 + 1 |
| | 0x00400038 | 0x000c6042 | srl \$t4, \$t4, 1 | 27: srl \$t4, \$t4, 1 #t4 >> 1 |
| | 0x0040003c | 0x0810000c | j if2 | 28: j if2 #goto if |
| | 0x00400040 | 0x01090018 | mult \$t0, \$t1 | 29: fim: mult \$t0, \$t1 #t0 * t1 |
| | 0x00400044 | 0x00006812 | mflo \$t5 | 30: mflo \$t5 #t5 = t0 |
| | 0x00400048 | 0x29ae0020 | slti \$t6, \$t5, 32 | 31: slti \$t6, \$t5, 32 #if(t5<32) (t6=1) else (t6 = 0) |
| | 0x0040004c | 0x20010001 | addi \$t6, \$t6, 1 | 32: bne \$t6, 1, maior #if(t6!=1) (goto maior) |
| | 0x00400050 | 0x142e0003 | bne \$t6, 1, maior | |
| | 0x00400054 | 0x000d9025 | or \$s2, \$0, \$t5 | 33: or \$s2, \$0, \$t5 #s2 = t5 |
| | 0x00400058 | 0x00009012 | mflo \$s2 | 34: mflo \$s2 #s2 = t0 |
| | 0x0040005c | 0x0810001a | j fim2 | 35: j fim2 #goto fim |
| | 0x00400060 | 0x00009010 | mfhi \$s2 | 36: maior: mfhi \$s2 #s2 = hi |
| | 0x00400064 | 0x00009812 | mflo \$s3 | 37: mflo \$s3 #s3 = lo |

| Name | Number | Value |
|--------|--------|------------|
| \$zero | 0 | 0x00000000 |
| \$at | 1 | 0x00000001 |
| \$v0 | 2 | 0x00000000 |
| \$v1 | 3 | 0x00000000 |
| \$a0 | 4 | 0x00000000 |
| \$a1 | 5 | 0x00000000 |
| \$a2 | 6 | 0x00000000 |
| \$a3 | 7 | 0x00000000 |
| \$t0 | 8 | 0x00000002 |
| \$t1 | 9 | 0x00000003 |
| \$t2 | 10 | 0x10010000 |
| \$t3 | 11 | 0x00000000 |
| \$t4 | 12 | 0x00000000 |
| \$t5 | 13 | 0x00000006 |
| \$t6 | 14 | 0x00000001 |
| \$t7 | 15 | 0x00000000 |
| \$s0 | 16 | 0x00000003 |
| \$s1 | 17 | 0x00000004 |
| \$s2 | 18 | 0x00000006 |
| \$s3 | 19 | 0x00000000 |
| \$s4 | 20 | 0x00000000 |
| \$s5 | 21 | 0x00000000 |
| \$s6 | 22 | 0x00000000 |
| \$s7 | 23 | 0x00000000 |
| \$s8 | 24 | 0x00000000 |
| \$s9 | 25 | 0x00000000 |
| \$k0 | 26 | 0x00000000 |
| \$k1 | 27 | 0x00000000 |
| \$gp | 28 | 0x10008000 |
| \$sp | 29 | 0x7ffffcfc |
| \$fp | 30 | 0x00000000 |
| \$ra | 31 | 0x00000000 |
| \$pc | | 0x00400068 |
| \$hi | | 0x00000000 |
| \$lo | | 0x00000006 |

```
// programa 20
```

$$y = \begin{cases} x^4 + x^3 - 2x^2 & \text{se } x \text{ for par} \\ x^5 - x^3 + 1 & \text{se } x \text{ for impar} \end{cases}$$

Os valores de x devem ser lidos da primeira posição livre da memória e o valor de y deverá ser escrito na segunda posição livre.

```

1  # inicio
2
3  .data
4  x: .word 3
5  .text
6  .globl main
7  main:
8      ori $t0, $0, 0x1001      #t0 -> 0x10010000 (first position)
9      sll $t0, $t0, 16        #t0 << 16
10
11     lw $s0, 0($t0)           #s0 = MEM[$t0]
12     andi $t1, $s0, 1         #t1 = s0 & 1
13     mult $s0, $s0            #s0 * s0
14     mflo $t2                 #t2 = x^2
15     mult $t2, $s0            #t2 * s0
16     mflo $t3                 #t3 = x^4
17     mult $t3, $s0            #t3 * s0
18     mflo $t4                 #t4 = x^8
19     mult $t4, $s0            #t4 * s0
20     mflo $t5                 #t5 = x^16
21     bne $t1, $0, impar      #if(t1!=0){ goto impar }
22     add $t2, $t2, $t2        #t2 = t2 + t2
23     add $s1, $t4, $t3        #s1 = t4 + t3
24     sub $s1, $s1, $t2        #s1 = s1 - t2
25     j fim                   #goto fim
26 impar: sub $s1, $t5, $t3     #s1 = t5 - t3
27     addi $s1, $s1, 1         #s1 = s1 + 1
28 fim:   sw $s1, 4($t0)        #MEM[4+$t0] = s1
29
30 # fim
31

```

[illegible]

```
// programa 21
```

$$y = \begin{cases} x^3 + 1 & \text{se } x > 0 \\ x^4 - 1 & \text{se } x \leq 0 \end{cases}$$

Os valores de x devem ser lidos da primeira posição livre da memória e o valor de y deverá ser escrito na segunda posição livre.

[illegible]