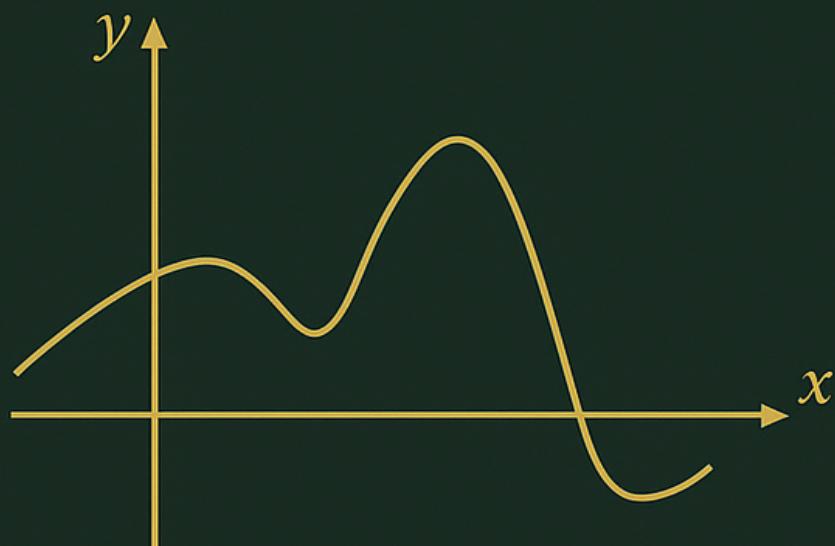


REAL ANALYSIS

A SELF-TAUGHT APPROACH



Luis Vasquez

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1 Introduction

1.1 Sources

The following notes are taken from the compilation of a few sources

- Zorich, V. A. (2004). Mathematical analysis I (R. Cooke, Trans.). Springer.
- MIT 18.100B Real Analysis, Spring 2025, available at the MIT OCW youtube channel

1.2 Purpose

This notes are taken in a way that is easy to understand math for myself (the author). No obscure proof or incomplete idea will be included, avoiding partial understanding of a certain topic. It also covers the need of having an easy-to-follow approach to *Real Analysis*, making it possible to read this through to revisit known topics without the need of rabbithole-ing at other sources. Since I will also be studying the course while taking this notes, the document as a whole will be written by hand without any AI slob nor blind copy/pasting, and since English is not my native language, typos may appear here and there.

*When I say **we**, it means you (the reader) and I. When I say **I/myself**, responsibility only lies on me (the author).*

2 Base knowledge

2.1 What is a Field

“We can verify that a set is a field by checking that multiplication is a well-defined operation, i.e., it is independent of the representative.”

For example, for arbitrary rational numbers Q :

$$\frac{m_1}{n_1} \times \frac{p_1}{q_1}$$

And evaluate an equivalent expression with different representatives of the same numbers:

$$\frac{m_2}{n_2} \times \frac{p_2}{q_2}$$

Given that:

$$\frac{m_1}{n_1} = \frac{m_2}{n_2} \quad \text{and} \quad \frac{p_1}{q_1} = \frac{p_2}{q_2}$$

We want to verify that both multiplication results are the same. To review this, we start from the tautology (intuitive truth):

$$\frac{m_1}{n_1} = \frac{m_2}{n_2} \iff m_1 \times n_2 = m_2 \times n_1 \quad (\text{A})$$

$$\frac{p_1}{q_1} = \frac{p_2}{q_2} \iff p_1 \times q_2 = p_2 \times q_1 \quad (\text{B})$$

Then, operating the multiplication using both representatives:

$$\begin{aligned} \frac{m_1}{n_1} \times \frac{p_1}{q_1} &= \frac{m_1 \cdot p_1}{n_1 \cdot q_1} \\ \frac{m_2}{n_2} \times \frac{p_2}{q_2} &= \frac{m_2 \cdot p_2}{n_2 \cdot q_2} \end{aligned}$$

Conveniently, we want to form $m_1 \times n_2$ to use the first ground truth:

$$\begin{aligned} \frac{m_1}{n_1} \times \frac{p_1}{q_1} \times n_2 &= \frac{m_1 \cdot n_2 \cdot p_1}{n_1 \cdot q_1} \\ &= \frac{m_2 \cdot n_1 \cdot p_1}{n_1 \cdot q_1} \quad (\text{replacing using A}) \\ &= \frac{m_2 \cdot p_1}{q_1} \quad (\text{simplifying } n_1) \end{aligned}$$

Applying the same logic for $p_1 \times q_2$ to use the second ground truth:

$$\begin{aligned} \frac{m_1}{n_1} \times \frac{p_1}{q_1} \times n_2 \times q_2 &= \frac{m_2 \cdot p_1 \cdot q_2}{q_1} \\ &= \frac{m_2 \cdot p_2 \cdot q_1}{q_1} \quad (\text{replacing using B}) \\ &= m_2 \cdot p_2 \quad (\text{simplifying } q_1) \end{aligned}$$

Finally, rearranging:

$$\begin{aligned}\frac{m_1}{n_1} \times \frac{p_1}{q_1} \times n_2 \times q_2 &= m_2 \cdot p_2 \\ \frac{m_1}{n_1} \times \frac{p_1}{q_1} &= \frac{m_2 \cdot p_2}{n_2 \cdot q_2} \\ \frac{m_1}{n_1} \times \frac{p_1}{q_1} &= \frac{m_2}{n_2} \times \frac{p_2}{q_2} \quad \square\end{aligned}$$

This is not a rigorous demonstration, but gives us a first step to go from the intuition of a solution (particularly for \mathbb{Q}) to a more formal procedure based on the real definition of a field.

2.2 Formal definition

Definition A **field** \mathbb{F} is a set with two operations: addition (\oplus) and multiplication (\otimes), with the following properties:

Field

2.1

- $x, y \in \mathbb{F} \implies x \oplus y \in \mathbb{F}$
- $x, y \in \mathbb{F} \implies x \oplus y = y \oplus x$
- $x, y, z \in \mathbb{F} \implies (x \oplus y) \oplus z = x \oplus (y \oplus z)$
- $\exists 0 \in \mathbb{F}$ such that $\forall x \in \mathbb{F}, x \oplus 0 = x$
- $\forall x \in \mathbb{F}, \exists (-x) \in \mathbb{F}$ such that $x \oplus (-x) = 0$
- $x, y \in \mathbb{F} \implies x \otimes y \in \mathbb{F}$
- $x, y \in \mathbb{F} \implies x \otimes y = y \otimes x$
- $x, y, z \in \mathbb{F} \implies (x \otimes y) \otimes z = x \otimes (y \otimes z)$
- $\exists 1 \in \mathbb{F}$ such that $\forall x \in \mathbb{F}, x \otimes 1 = x$
- $\forall x \in \mathbb{F} \setminus \{0\}, \exists x^{-1} \in \mathbb{F}$ such that $x \otimes x^{-1} = 1$

The first five properties correspond to the addition operation, and the last five to the multiplication operation. In order to relate both sets of properties, the following axiom is stated:

AxiomDistributive law
2.1Let $x, y, z \in \mathbb{F}$. Then

$$x \otimes (y \oplus z) = x \otimes y \oplus x \otimes z$$

TheoremZero uniqueness
2.1For any field \mathbb{F} , there exists only one zero element.

Proof | Assume 0_1 and 0_2 are zeros for a field \mathbb{F} . Then

$$\forall x \in \mathbb{F} : \begin{cases} 0_1 + x = x & (\text{A}) \\ 0_2 + x = x & (\text{B}) \end{cases}$$

For both cases, let x be 0_2 and 0_1 respectively:

$$x = 0_2 \implies (\text{A}) : 0_1 + 0_2 = 0_2 \quad (\text{C})$$

$$x = 0_1 \implies (\text{B}) : 0_2 + 0_1 = 0_1 \quad (\text{D})$$

From (C) and commutativity:

$$0_2 + 0_1 = 0_2$$

Comparing this result with (D):

$$\begin{aligned} 0_2 + 0_1 &= 0_2 \quad \text{and} \quad 0_2 + 0_1 = 0_1 \\ &\implies 0_1 = 0_2 \end{aligned}$$

□

2.3 Order

Definition A set S is **ordered** when it has an ordering “ $<$ ” such that for all $x, y \in S$, exactly one of the following properties holds:

2.2

1. $x = y$
2. $x < y$
3. $y < x$

Definition A field \mathbb{F} is **ordered** if it is also an ordered set. As a consequence, the following properties apply:

2.3

- $x, y \in \mathbb{F}, x < y \implies \forall z \in \mathbb{F}, x \oplus z < y \oplus z$
- $x, y \in \mathbb{F}, 0 < x, y \implies 0 < x \otimes y$

Theorem Given \mathbb{F} an ordered field. If $x < y$ and $0 < z$, then $x \otimes z < y \otimes z$.

2.2

Proof We prove by contradiction. Assume the opposite:

$$x \otimes z \geq y \otimes z$$

Then adding $(-x \otimes z)$ on both sides, we maintain the ordering of the expression:

$$\begin{aligned} x \otimes z \oplus (-x \otimes z) &\geq y \otimes z \oplus (-x \otimes z) \\ 0 &\geq y \otimes z \oplus (-x \otimes z) \end{aligned}$$

Now, using Axiom 2.1, we get:

$$0 \geq z \otimes (y \oplus (-x))$$

From the initial conditions, $x < y$ implies $y \oplus (-x) > 0$. Since we also have $z > 0$, we would expect the product of these two to be > 0 by Definition 2.3 (second property). Hence:

$$0 \geq z \otimes (y \oplus (-x)) \quad \text{and} \quad 0 < z \otimes (y \oplus (-x))$$

This is a contradiction. □

From this point forward, for better readability, $+$ and \cdot (or \times) will be used instead of \oplus and \otimes . They will still represent the abstraction of a field’s addition and multiplication operations, without necessarily being the familiar operations we might expect them to be.

2.4 Completeness

Definition | A set $X \subset \mathbb{F}$ (\mathbb{F} ordered field) is said to be *bounded above* (or respectively, *bounded below*) if $\exists c \in \mathbb{F}$ such that $\forall a \in \mathbb{F}, a \leq c$ (or respectively, $c \leq a$). c is called upper (or respectively lower) bound of X .

Definition | A set that is both bounded above and below, is called *bounded*.

2.5

Definition | An element $a \in X$ is called the *largest* element of X if $\forall x \in X, x \leq a$. Respectively, $a \in X$ is called the *smallest* element of X if $\forall x \in X, a \leq x$. Simplifying the notation:

$$(a = \max X) := (a \in X \wedge \forall x \in X, x \leq a)$$

$$(a = \min X) := (a \in X \wedge \forall x \in X, a \leq x)$$

These read as *maximal* and *minimal* of X . Now, given this definition, it is important to notice that not every set, not even every bounded set, has a maximal or minimal element. For example:

$$X = \{x \in \mathbb{F} \mid 0 \leq x < 1\}$$

Only has a minimal element (0), but no maximal element, since $1 \notin X$

Definition | The smallest $s \in X \subset \mathbb{F}$ that bounds X from above is called the *least upper bound* of X , and denoted $\sup X$ (read "the supremum of X ")

2.7

$$(s = \sup X) := \forall x \in X((x \leq s) \wedge (\forall s' < s \exists x' \in X(s' < x')))$$

Lets break this down by element

- $\forall x \in X$: The following definition applies to the whole set X .
- $(x \leq s)$: Given that s is an upper bound for X ...
- $(\forall s' < s) \exists x' \in X(s' < x')$:
 - $(\forall s' < s)$: Considering any arbitrary s' smaller than our upper bound s
 - $\exists x' \in X$: There will be an element x' in X , so that...
 - $(s' < x')$: It is larger than the s' , making s' to fail to be an upper bound.

So in summary, s is $\sup X$ if and only if, s is an upper bound, and no smaller number s' is an upper bound of X , because we can find an x' that is not bounded by it.

Definition Greatest Lower Bound 2.8	<p>Similarly, the greatest $i \in X \subset \mathbb{F}$ that bounds X below is called the <i>greatest lower bound of X</i>, and denoted $\inf X$ (read "the infimum of X")</p> $(i = \inf X) := \forall x \in X((i \leq x) \wedge (\forall i' < i \exists x' \in X(x' < i')))$
---	---

Thus, we have now the following definitions:

$$\begin{aligned}\sup X &:= \min\{c \in \mathbb{F} | \forall x \in X(x \leq c)\} \\ \inf X &:= \max\{c \in \mathbb{F} | \forall x \in X(c \leq x)\}\end{aligned}$$

It is important to note that the supremum and infimum of a set, as defined above, may not exist in an arbitrary ordered field \mathbb{F} . The definitions above specify what $\sup X$ and $\inf X$ mean *if they exist*, but they do not guarantee existence. We will address this issue shortly.

Theorem Uniqueness of Supremum 2.3	<p>Let $X \subset \mathbb{F}$ be a nonempty set in an ordered field \mathbb{F}. If X has a supremum, then this supremum is unique.</p>
--	---

Proof Suppose s_1 and s_2 are both suprema of X . We will show that $s_1 = s_2$. Since $s_1 = \sup X$, we know that s_1 is an upper bound of X . Since $s_2 = \sup X$, we know that s_2 is the *least* upper bound of X . Therefore:

$$s_2 \leq s_1$$

By the same reasoning (swapping the roles of s_1 and s_2), since s_2 is an upper bound and s_1 is the least upper bound:

$$s_1 \leq s_2$$

By the antisymmetry property of order in \mathbb{F} , we have $s_1 \leq s_2$ and $s_2 \leq s_1$, which implies:

$$s_1 = s_2$$

Therefore, the supremum is unique. □

The proof for the uniqueness of the infimum is analogous.

Example
2.1

Consider the ordered field \mathbb{Q} of rational numbers, and let:

$$S = \{x \in \mathbb{Q} \mid x^2 < 2\}$$

The set S is nonempty (for instance, $1 \in S$) and bounded above (for instance, $2 \in \mathbb{F}$ is an upper bound). However, S does not have a supremum in \mathbb{Q} , since if it had, it would have to equal $\sqrt{2} \notin \mathbb{Q}$, so S has no least upper bound within the rational numbers.

This shows that not every ordered field has the property that bounded sets possess suprema.

Definition
Complete Ordered
Field
2.9

An ordered field \mathbb{F} is called *complete* if every nonempty subset of \mathbb{F} that is bounded above has a supremum in \mathbb{F} .

Proposition
2.1

If \mathbb{F} is a complete ordered field, then every nonempty subset of \mathbb{F} that is bounded below has an infimum in \mathbb{F} .

Proof Let $X \subset \mathbb{F}$ be nonempty and bounded below. Define:

$$X' = \{-x \mid x \in X\}$$

1. X' is nonempty and bounded above

Since X is nonempty, $\exists x_0 \in X \Rightarrow -x_0 \in X'$ by definition of X' . Thus $X' \neq \emptyset$. Now, X is bounded below by the proposition, so there exists $c \in \mathbb{F}$ such that:

$$\begin{aligned} \forall x \in X, \quad c \leq x \\ \forall x \in X, \quad -x \leq -c \end{aligned}$$

Since $-x \in X'$ and $-c \in \mathbb{F}$, we can be sure that X' is bounded above in \mathbb{F} . Considering \mathbb{F} is complete, and we proved X' is nonempty and bounded above, X' has a supremum in \mathbb{F} .

2. $-s$ is a lower bound for X .

Let $s = \sup X'$ and $x \in X$ be arbitrary. Then $-x \in X'$ by definition of X' . Since s is an upper bound for X' :

$$\begin{aligned} \forall -x \in X', \quad -x \leq s \\ \Rightarrow -s \leq x \end{aligned}$$

Showing that $-s$ is a lower bound for X .

3: $-s$ is the greatest lower bound for X .

Let $\ell \in \mathbb{F}$ be any lower bound for X .

Since ℓ is a lower bound for X :

$$\begin{aligned} \forall x \in X, \quad \ell \leq x \\ \Rightarrow -x \leq -\ell \end{aligned}$$

Meaning that $-\ell$ is an upper bound for any $-x \in X'$. Since $s = \sup X'$, we have:

$$\begin{aligned} s &\leq -\ell \\ \Rightarrow \ell &\leq -s \end{aligned}$$

Showing that $-s$ is greater than or equal to every lower bound of X . Therefore, $-s = \inf X$. \square

Theorem
 Least Upper
 Bound Principle
 2.4

Let \mathbb{F} be a complete ordered field. Every nonempty subset of \mathbb{F} that is bounded above has a unique least upper bound in \mathbb{F} .

Proof

Let $X \subset \mathbb{F}$ be nonempty and bounded above.

Existence: Since \mathbb{F} is complete, by definition of completeness, X has a supremum in \mathbb{F} .

Uniqueness: By the Uniqueness of Supremum theorem, this supremum is unique. Therefore, X has a unique least upper bound. \square

2.5 Archimedean property

Now that we achieved completeness on an ordered field, we can now extend the definitions and properties into a theorem/property that uses the result from the *Least Upper Bound principle* (Theorem 2.4).

Theorem
 2.5

Let \mathbb{F} be a complete ordered field, then

$$\forall x \in \mathbb{F}, \exists! n \in \mathbb{N} \subset \mathbb{F} : n > x$$

Proof We first define the subset $\mathbb{N} \subset \mathbb{F}$ **inductively**, to work with a generalized version of the set of “natural numbers”

1. $1_{\mathbb{F}}$ is the multiplicative identity in \mathbb{F}
2. $n + 1_{\mathbb{F}} \in \mathbb{N}, \forall n \in \mathbb{N}$ defines the induction step over \mathbb{N} , using the addition operation from the field \mathbb{F}

We now proceed to prove by contradiction, assuming $\exists x \in \mathbb{F}, \forall n \in \mathbb{N} :$

$$n \leq x$$

Meaning that \mathbb{N} is bounded above. Now, since \mathbb{N} is not empty, and is bounded by an element of the (complete ordered) field that contains it ($x \in \mathbb{F}$), then by *L.U.B. principle* (Theorem 2.4), \mathbb{N} must have a supremum.

Let $s := \sup \mathbb{N} \in \mathbb{F}$ be that supremum of \mathbb{N} . Then, since s is the least upper bound, any element that is less than s will no longer be an upper bound for \mathbb{N} . Conveniently we take $s - 1_{\mathbb{F}}$:

$$\begin{aligned} \exists n_0 \in \mathbb{N} : s - 1_{\mathbb{F}} &< n_0 \\ &\Rightarrow s < n_0 + 1_{\mathbb{F}} \end{aligned}$$

But by our inductive definition of \mathbb{N} : $(n_0 + 1_{\mathbb{F}}) \in \mathbb{N}$. So we found out that s ($\sup \mathbb{N}$) fails to be an upper bound for an element in \mathbb{N} .

Therefore, our original assumption of x must be wrong. \square

Now, is important to notice that \mathbb{F} implies an archimedien field, but the opposite won't necessarily be true.

Example 2.2 Using an Archimedean field, that is not complete, lets take $\mathbb{F} = \mathbb{Q}$

$$x \in \mathbb{Q} \wedge n \in \mathbb{N} \subset \mathbb{Q}$$

It is true that, for any x there is a natural number n so that

$$n > x$$

But invoking the previous example (Example 2.1) we can take a subset of \mathbb{Q} that is not bounded in \mathbb{Q} , failing to properly form a complete set.

Closing this chapter, lets emphasize the fact that the construction of $\mathbb{N} \subset \mathbb{F}$ by induction was crucial to generalize the property for any complete ordered field. Next we will explain why this wasn't necessary at all.

3 \mathbb{R} as the complete ordered field

3.1 \mathbb{F} is \mathbb{R}

Turns out that there is a reason why it is called ***Real*** Analysis. So far we have worked over a generic *Complete Ordered Field* \mathbb{F} , that was also an *Archimedean* following the expected properties. We have finally reached the point were we can state that **any** set that manages to follow all these properties is actually identical to the set of real numbers \mathbb{R} . In fact, \mathbb{R} is the only set that accomplishes this.

Most Real Analysis courses don't go into proving this, but as I already mentioned in the first chapter, I want to cover all doubts and missing point existent during a course like this.

We will go through the initial proofs needed to understand the main one. At the beginning its possible that some of them look unnecessary or trivial, but I don't want to assume that much. Better to be safe than sorry.

3.1.1 Isomorphism

Our starting point will be defining what it means to have “equal fields”.

Definition Field homomorphism 3.1	Let F, F' fields. A function $\phi : F \rightarrow F'$ is a homomorphism between F and F' iif. it preserves the defined operations in both fields: $\begin{aligned}\phi(x + y) &= \phi(x) + \phi(y), \forall x, y \in F \\ \phi(x \cdot y) &= \phi(x) \cdot \phi(y), \forall x, y \in F\end{aligned}$
--	--

Definition Field isomorphism 3.2	Let F, F' fields. A function $\phi : F \rightarrow F'$ is a isomorphism iif. it is a homomorphism and also a bijection
---	--

Lemma | Let F, F' fields, $\phi : F \rightarrow F'$. If ϕ preserves order from \mathbb{F} :

Injectivity by
order preservation

3.1

$$x < y \Rightarrow \phi(x) < \phi(y), x, y \in F$$

Then ϕ is injective.

Proof | To prove injectivity, we need to show that for any two identical images of the function ϕ , we get that their arguments were the same, guaranteeing the uniqueness of the image:

Let $x, y \in \mathbb{F}$ so that

$$\phi(x) = \phi(y)$$

Using the tricotomy over F :

Case 1: $x < y$:

$$\rightarrow \phi(x) < \phi(y) (\rightarrow \leftarrow)$$

Case 2: $y < x$:

$$\rightarrow \phi(y) < \phi(x) (\rightarrow \leftarrow)$$

$$\Rightarrow x = y$$

Then, ϕ is injective. □

Lemma | Let F, F' fields, $\phi : F \rightarrow F'$. Define:

Surjectivity by
construction

3.2

$$\text{Im}_F := \{\phi(x), x \in F\}$$

Then, the bounded function $\phi' : F \rightarrow \text{Im}_F$ is surjective

Proof To prove surjectivity, we need to show that every element in $y \in \text{Im}_F$ has an $x \in F$ such that $\phi'(x) = y$
 Using the definition of our constructed set:

$$\begin{aligned} \forall y \in \text{Im}_F &\Leftrightarrow y \in \{\phi(x), x \in F\} \\ &\Rightarrow \exists x \in F / \forall y \in \text{Im}_F : y = \phi'(x) \end{aligned}$$

Finally proving that for any image of $\phi' \in \text{Im}_F$ we will have a preimage $x \in F$.
 Then ϕ' is surjective. \square

On this first part, this surjectiveness by construction might be the bit that doesn't quite seem correct. Isn't it trivial to have surjectiveness if we grab the method ϕ and limit it to its image ϕ' ?

The key here is that we defined the set of images Im_F first, and *then* proved the surjectiveness in it. The lemma confirms that the set we defined was *just right*. Not so big to leave some elements unreached, and not too small to make undefined preimages at some spots.

3.1.2 Sets construction

Now, to make it reasonable for a complete ordered field \mathbb{F} to be isomorphic to our known \mathbb{R} , we need to start shaping it using the parts that make it up. Think of \mathbb{R} as the joint set of irrational and rational numbers, which contain the other types of number as well.

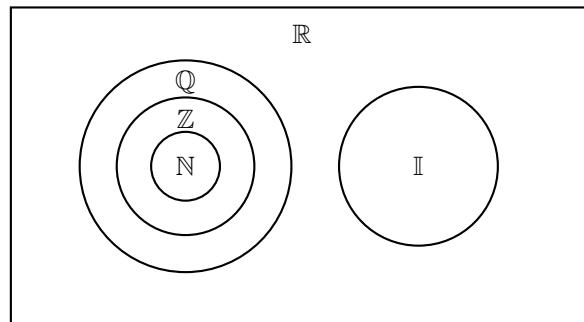


Figure 1: Hierarchy of number sets: Natural numbers (\mathbb{N}), Integers (\mathbb{Z}), Rationals (\mathbb{Q}), Irrationals (\mathbb{I}), and Real numbers (\mathbb{R}).

From this idea, we need to build $\mathbb{N}_{\mathbb{F}}$, $\mathbb{Z}_{\mathbb{F}}$, $\mathbb{Q}_{\mathbb{F}}$ and $\mathbb{I}_{\mathbb{F}}$ to finally get this \mathbb{F} we want to make some proofs on. This will all make sense in the last bit of the process, and I promise that the only tricky part will be to construct $\mathbb{I}_{\mathbb{F}}$ from $\mathbb{Q}_{\mathbb{F}}$, but we need to go step by step.

Definition
 $\mathbb{N}_{\mathbb{F}}$ by induction
3.3

Let \mathbb{F} be *the* complete ordered field. Using $1_{\mathbb{F}}$: multiplicative neutral of \mathbb{F} , we declare the following element construction:

$$\begin{aligned} 1_{\mathbb{F}} &:= 1_{\mathbb{F}} \\ 2_{\mathbb{F}} &:= 1_{\mathbb{F}} + 1_{\mathbb{F}} \\ 3_{\mathbb{F}} &:= 1_{\mathbb{F}} + 1_{\mathbb{F}} + 1_{\mathbb{F}} \\ &\vdots \\ n_{\mathbb{F}} &:= \sum_1^n 1_{\mathbb{F}} \end{aligned}$$

Finally, we define the set:

$$\mathbb{N}_{\mathbb{F}} := \{1_{\mathbb{F}}, 2_{\mathbb{F}}, 3_{\mathbb{F}}, \dots\} = \{n_{\mathbb{F}} : n \in \mathbb{N}\}$$

Definition
 $\mathbb{Z}_{\mathbb{F}}$ by extension
3.4

Within field \mathbb{F} , let the following:

1. $0_{\mathbb{F}}$: Additive identity
2. $-x$: Additive inverse

We define the set:

$$\mathbb{Z}_{\mathbb{F}} = \{0_{\mathbb{F}}\} \cup \mathbb{N}_{\mathbb{F}} \cup \{-x_{\mathbb{F}}, \forall n_{\mathbb{F}} \in \mathbb{N}_{\mathbb{F}}\}$$

Definition
 $\mathbb{Q}_{\mathbb{F}}$ by operation
3.5

Within field \mathbb{F} , we define the set:

$$\mathbb{Q}_{\mathbb{F}} := \left\{ \frac{a}{b} : a \in \mathbb{Z}_{\mathbb{F}}, b \in \mathbb{N}_{\mathbb{F}} \right\}$$

Starting similarly to Theorem 2.5 (Archimedean property) we use the inductive approach to reach $\mathbb{N}_{\mathbb{F}}$. Fancy but not complicated at all. Later, we just extend integers in \mathbb{F} to build by $\mathbb{Z}_{\mathbb{F}}$. Notice how we only used existent elements from \mathbb{F} computed using the field's operations and nothing else.