

UNIVERSIDAD EAFIT

OPERATING SYSTEMS

ST0257

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## Midterm 1: Large Datasets Processing in C++

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## Context and Approach

This project focuses on handling large datasets in C/C++ on Linux and evaluating how different design choices affect performance. The main objectives are:

- Generate and process datasets with millions of records.
- Compare using **values vs. pointers** to measure memory savings and runtime performance.
- Compare using **structs (C) vs. classes (C++)** to examine efficiency and memory layout.
- Record **execution time** and **memory usage** for each query with a custom monitor.

Each record contains: full name, date of birth, city of residence, assets, debts, ID number, and the assigned tax calendar group (A/B/C) based on the last digits of the ID.

## Queries

The required queries were:

### 1. Oldest person

- Across the entire dataset.
- Grouped by city.

### 2. Person with the most assets

- Overall.
- By city.
- By tax calendar group (A/B/C).

### 3. Tax filers

- Count people in each tax calendar group.
- Validate assignment according to ID termination rules.

Additional queries designed by the team:

1. Cities with the highest average assets.
2. Percentage of people older than 80 in each tax group.
3. Number of people in each city.

# How We Measured Performance

Two main configurations were tested:

- **By value vs. by pointer:** storing full records directly compared to storing only pointers to them.
- **struct vs. class:** using plain structs without constructors or virtual methods compared to full C++ classes with methods.

Performance was measured through:

- **Execution time** with `std::chrono`.
- **Memory usage** (RSS in KB) from `/proc/self/statm`.
- **Recorded statistics** saved by the monitor and summarized in console.

## Analysis

The results let us compare time and memory consumption across different approaches. Key points for analysis include:

- The extent of memory savings when using pointers instead of values, particularly with very large datasets.
- Whether structs offer measurable cache benefits compared to classes.
- How queries scale when the dataset grows beyond one million records.
- Which queries are more costly in terms of performance, such as full scans versus grouped lookups.

## Preliminary Results

We collected measurements for dataset creation and several queries using two main distinctions:

1. **Values vs. Pointers** — how data is stored and accessed.
2. **Structs vs. Classes** — how records are defined in C/C++.

The metrics include execution time (milliseconds) and additional memory used (KB).

## 5.1 Dataset Creation

Configuration	Time (ms)	Memory (KB)
Pointers	1806.87	105200
Values	1843.18	105112

**Table 1:** Time and memory for creating a dataset with 1,000,000 records

At the dataset creation stage, the use of structs versus classes did not show a significant difference. Both approaches required roughly the same memory footprint and initialization time. The small variations are more related to pointer vs. value than to the record type.

## 5.2 Finding Oldest Person (by City)

Configuration	Time (ms)	Memory (KB)
Pointers	0.056	0
Values	703.745	6972

**Table 2:** Performance of query: oldest person by city

Here the gap is striking: with pointers the query is almost instantaneous, while with values it takes hundreds of milliseconds and uses noticeable extra memory. In this type of traversal, structs behave similarly to classes, but classes may add small overhead due to constructors and possible alignment padding. Still, the main bottleneck here is clearly value vs. pointer.

## 5.3 Finding Person with Highest Assets (by Tax Group)

Configuration	Time (ms)	Memory (KB)
Values	92.256	0
Pointers	0.003	0

**Table 3:** Performance of query: highest assets by tax group

In this aggregation query, using pointers again proved much faster. Memory impact was negligible in both cases. When comparing structs and classes, structs offer slightly better cache locality (plain old data layout), while classes bring organizational clarity at the cost of small construction overhead.

# Notes on Early Observations

- Dataset creation showed minimal differences across all approaches — the choice between struct or class does not matter much here.
- Query performance, however, is dominated by whether pointers or values are used. Pointers give far superior execution times and avoid unnecessary memory copies.

- Structs generally help in memory layout and cache efficiency, but when queries involve heavy scanning or grouping, the pointer vs. value decision outweighs the struct vs. class effect.

In summary, structs and classes perform similarly for bulk dataset creation, but pointers versus values make a dramatic difference when running queries on large datasets.

## Critical Thinking Questions

### 7.1 Why do pointers reduce memory usage when dealing with 10 million records?

Pointers reduce memory usage because they reference a shared memory location instead of duplicating entire data structures. In scenarios with millions of records, storing full copies of the same or large objects is inefficient. By storing pointers, only the address of the data (typically 4 or 8 bytes, depending on architecture) is maintained, avoiding redundant copies. This can drastically reduce memory consumption, though the actual savings depend on the structure size and whether the pointed-to data is reused.

### 7.2 If the calendar depends on the person's ID, how can group searches be optimized?

If the calendar grouping is a deterministic function of the person's ID, it is unnecessary to store the group explicitly for each record. Instead, one can precompute or derive the mapping function between the ID and the calendar group (e.g., an enumeration such as `CalendarTaxGroup`). By indexing or hashing IDs directly to their groups, searches can be optimized, reducing both memory overhead and lookup time.

### 7.3 How does memory access differ between an array of structs and a vector of class objects?

The distinction arises primarily from data layout and object semantics:

- **Array of Structs (AoS):** Structs that are Plain Old Data (POD) types—meaning no user-defined constructors, destructors, or virtual functions—are stored contiguously in memory. This layout maximizes spatial locality and cache efficiency, which is advantageous in data-oriented design.
- **Vector of Class Objects:** While `std::vector` also stores elements contiguously, class objects often introduce overhead through constructors, destructors, and possible virtual table pointers. This makes them heavier than plain structs, potentially harming cache performance. If the class holds pointers to heap-allocated members, memory access becomes indirect and fragmented.

Thus, for performance-critical workloads requiring sequential traversal and predictable memory access, an AoS layout is typically more efficient.

## 7.4 How can `mmap` or virtual memory help when data exceeds physical RAM?

When datasets exceed available RAM, the operating system leverages virtual memory. Pages that cannot fit in RAM are stored in a swap area on disk, with the OS transparently paging them in and out. This allows programs to operate on data larger than physical memory, though with significant latency penalties due to disk access times.

The `mmap` system call provides more direct control. It can:

- Map files directly into memory space, enabling efficient I/O by allowing the OS page cache to handle reads and writes.
- Allocate large contiguous memory regions, bypassing some limitations of `malloc`.
- Enable shared memory between processes by mapping the same file or anonymous memory region into multiple address spaces.

In high-volume data processing, combining virtual memory and `mmap` can allow applications to handle datasets that far exceed the machine's physical memory capacity, though at the cost of higher latency for non-resident pages.