

Programación Funcional

Curso 2019-20

ORDEN SUPERIOR

Funciones de orden superior

(HO functions)

- Funciones con algún argumento (o el resultado) de tipo funcional.
- Favorecen la abstracción, concisión y reutilización de código.

Funciones de primer orden: las "normales", que no son de OS

$f :: (\text{Int} \rightarrow \text{Int}) \rightarrow \text{Int}$ Es de orden superior

$f' :: (\text{Int}, \text{Int}) \rightarrow \text{Int}$ Es de primer orden

¿Cuál es el orden de la función g definida como $g\ x = f'$

Un par de funciones de primer orden...

```
-- incList xs:  sumar uno a todos los elementos de xs
-- incList [x1,...,xn] = [1+x1,...,1+xn]
incList::[Int] -> [Int]  Recomendado: declarar siempre los tipos
incList [] = []
incList (x:xs) = inc x:incList xs where inc x = x+1
```

```
-- lengthList xss: longitudes de los elementos de xss
-- lengthList [xs1,...,xsn] = [length xs1,...,length xsn]
lengthList::[[a]] -> [Int]
lengthList []      = []
lengthList (x:xs) = length x:lengthList xs
```

Se repite la misma estructura:
aplicar una función f a todos los elementos de la lista

Abstracción de OS: map

```
-- map f xs = resultado de aplicar f a todos los elementos de xs
-- map f [x1,...,xn] = [f x1,...,f xn]
map::(a -> b) -> [a] -> [b]
map f []      = []
map f (x:xs) = f x:map f xs
```

```
> map (take 2) [[1,2,3],[2],[7,6,4]]
[[1,2],[2],[7,6]]
> map not [True,False,False]
[False,True,True]
```

```
-- incList [x1,...,xn] = [1+x1,...,1+xn]
-- lengthList [xs1,...,xsn] = [length xs1,...,length xsn]
incList xs      = map inc xs
lengthList xs = map length xs
```

Cancelación de argumentos

En lugar de

```
incList xs      = map inc xs  
lengthList xs = map length xs
```

Podemos definir de modo equivalente

```
incList' = map inc  
lengthList' = map length Pointless programming
```

Tanto recursión como aplicación a argumentos quedan implícitas

Ambas definiciones son equivalentes

```
incList' [1,2,3] = (map inc) [1,2,3] = map inc [1,2,3] =  
incList [1,2,3]  
lengthList' [1,2,3] = (map length) [1,2,3] =  
map length [1,2,3] = lengthList [1,2,3]
```

Funciones de orden superior sobre listas I

filter, all, any

```
-- filter p xs = lista de elementos de xs que cumplen la propiedad p
filter :: (a -> Bool) -> [a] -> [a]
filter p []      = []
filter p (x:xs)  =
  | p x          = x:filter p xs
  | otherwise    =   filter p xs

-- all p xs = todos los elementos de xs cumplen p
-- any p xs = algún elemento de xs cumple p

all _ []          = True
all p (x:xs)      = p x && all p xs
any _ []          = False
any p (x:xs)      = p x || any p xs
```

```
> filter (> 1) [-1,3,0,4]
[3,4]
> all (> 1) [-1,3,0,4]
False
> any (> 1) [-1,3,0,4]
True
```

Funciones de orden superior sobre listas II

takeWhile, dropWhile, span, break

```
-- takeWhile p xs = mayor prefijo de xs cuyos elementos cumplen p
-- dropWhile p xs = resultado de eliminar (takeWhile p xs) de xs
takeWhile, dropWhile :: (a -> Bool) -> [a] -> [a]
span, break :: (a -> Bool) -> [a] -> ([a], [a])
takeWhile _ [] = []
takeWhile p (x:xs)
  | p x      = x : takeWhile p xs
  | otherwise = []

dropWhile _ [] = []
dropWhile p (x:xs)
  | p x      = dropWhile p xs
  | otherwise = x:xs

span p xs = (takeWhile p xs, dropWhile p xs)
break p xs = span (not.p) xs
```

- es la composición de funciones xs se podría cancelar, pero p no

Composición de funciones: .

```
-- f.g = composición de las funciones f y g
```

```
infixr 9 .
```

```
(.) :: (b -> c) -> (a -> b) -> (a -> c)
```

```
(.) f g x = f (g x)
```

```
> (head . tail) [1,2,3]  
2
```

```
> ((^ 2).(3 *)) 2  
36
```

Aplicación funcional: \$

```
-- f $ x = aplica la función f a x
```

```
infixr 0 $
```

```
($) :: (a -> b) -> a -> b
```

```
($) f x = f (x)
```

```
> head $ tail [1,2,3]  
2
```

```
> tail $ map (*2) [1,2,3]  
[4,6]
```


Más funciones de orden superior sobre listas

```
iterate, zipWith, zip
```

```
-- iterate f x = [x,f x, f (f x),...]
```

```
iterate:: (a -> a) -> a -> [a]
```

```
iterate f x = x : iterate f (f x)
```

```
iterate (* 2) 1 =[1,2,4,8,...]
```

```
take 3 (map head (iterate tail [1..]))=[1,2,3]
```

```
-- zipWith f xs ys:  combina los elementos de xs e ys mediante f
```

```
-- Si una lista es más corta, se descartan los sobrantes de la otra
```

```
-- zipWith f [x1,...,xn] [y1,...,ym]=[f x1 y1,...,f xk yk] (k=min(n,m))
```

```
zipWith :: (a -> b -> c) -> [a] -> [b] -> [c]
```

```
zipWith f (a:as) (b:bs) = f a b : zipWith f as bs
```

```
zipWith _ _ _ = []
```

```
zip [1,2,3] ['a','b'] = [(1,'a'),(2,'b')]
```

```
zipWith (+) [1,2,3] [4,5,6] = [5,7,9]
```

Define la función zip en términos de zipWith

La familia *fold*

Un patrón de recursión sobre listas muy repetido

```
sum [] = 0
sum (x:xs) = x + sum xs
```

```
and [] = True
and (x:xs) = x && and xs
```

```
or [] = False
or (x:xs) = x || or xs
```

```
length [] = 0
length (x:xs) = 1+length xs
```

```
concat [] = []
concat (xs:xss) = xs++concat xss
```

El patrón general que reconocemos es:

```
g [] = e
g (x:xs) = f x (g xs)
```

O bien, si f viene como un operador infijo \oplus :

```
g [] = e
g (x:xs) = x  $\oplus$  g xs
```

O, en términos semiformales, y suponiendo que \oplus asocia a la derecha:

```
g [x1,x2,...,xn] = x1  $\oplus$  x2  $\oplus$  ...  $\oplus$  xn  $\oplus$  e
```

Podemos abstraer en una función de OS con f/\oplus y e como parámetros

La familia *fold* (II)

$$\text{foldr } f \ e \ [x_1, x_2, \dots, x_n] = f \ x_1 \ (f \ x_2 \ (\dots (f \ x_n \ e) \dots))$$
$$\text{foldr } \oplus \ e \ [x_1, x_2, \dots, x_n] = x_1 \oplus x_2 \oplus \dots \oplus x_n \oplus e$$

```
foldr :: (a -> b -> b) -> b -> [a] -> b
foldr f e []      = e
foldr f e (x:xs) = f x (foldr f e xs)
-- 0 escribiendo f en modo infijo
foldr f e (x:xs) = x 'f' (foldr f e xs)
```

```
sum      = foldr (+) 0
product = foldr (*) 1
and      = foldr (&&) True
or       = foldr (||) False
length  = foldr f 0  where f x y = y+1
concat  = foldr (++) []
```

La familia *fold* (III)

$\text{foldl } \oplus e [x_1, x_2, \dots, x_n] = e \oplus x_1 \oplus x_2 \oplus \dots \oplus x_n$
(suponiendo que \oplus asocia por la izda)

$\text{foldl } f e [x_1, x_2, \dots, x_n] = (f (\dots (f (f e x_1) x_2) \dots) x_n)$

$\text{foldl} :: (a \rightarrow b \rightarrow a) \rightarrow a \rightarrow [b] \rightarrow a$

$\text{foldl } f e [] = e$

$\text{foldl } f e (x:xs) = \text{foldl } f (f e x) xs$

$\text{sum} = \text{foldl } (+) 0$

$\text{product} = \text{foldl } (*) 1$

$\text{and} = \text{foldl } (\&\&) \text{True}$

$\text{or} = \text{foldl } (||) \text{False}$

$\text{length} = \text{foldl } f 0$

 where $f x y = x+1$

$\text{concat} = \text{foldl } (++) []$

Más ineficiente que antes ¿Por qué?

La familia *fold* (III)

- *foldr* puede procesar listas (incluso infinitas) sin recorrerlas enteras (dependiendo de \oplus)
- *foldl* ha de recorrer la lista entera
- *foldl* presenta recursión final (*tail recursion*)
- La eficiencia comparativa es muy dependiente de cada caso

Miscelánea OS

```
-- flip f: cambia de orden los argumentos de f
flip:: (a -> b -> c) -> (b -> a -> c)
flip f x y = f y x

-- curry f:  currifica f (que es una función que se aplica a pares)
-- uncurry f:  hace que f (función currificada) se aplique a pares
curry::((a,b) -> c) -> (a -> b -> c)
curry f x y = f (x,y)

uncurry::(a -> b -> c) -> ((a,b) -> c)
uncurry f (x,y) = f x y
```

Funciones que suelen usarse en forma de aplicación parcial

```
-- id:  función identidad
id::  a -> a
id x = x

-- const x:  función constante de valor x
const::  a -> b -> a
const x y = x
```