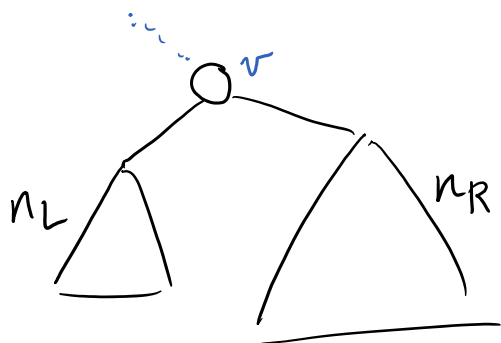


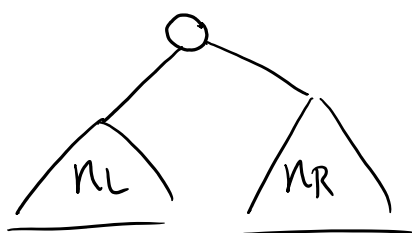
Size-Balanced Property for a node



$$\text{MAX}(n_L, n_R) \leq 2 \times \text{MIN}(n_L, n_R) + 1$$

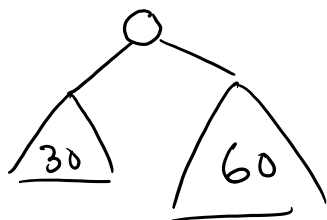
A tree is size-balanced iff all nodes are size-balanced. [NOT JUST GLOBAL ROOT].

Ex SUPPOSE $N=91$



$$n_L + n_R = 90$$

?
 $60 \leq 2 \times 30 + 1$
 \checkmark 61



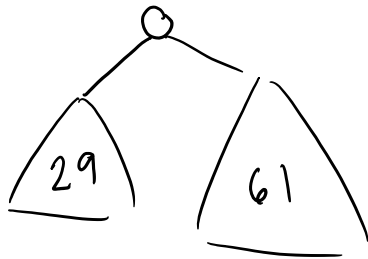
ok?
 (a) YES
 (b) NO

?
 $61 \leq 58 + 1$



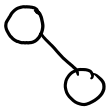
1. 7 (a) YES

✓ ?
 $6 \leq 58 + 1$
 59
~~X~~

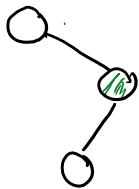


ok? (a) YES
 (b) NO

$1 \leq 2 \times 0 + 1$
 $1 \leq 1$ ✓

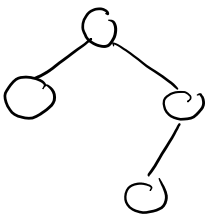


ok? (a) yes
 (b) No



?
 $2 \leq 2 \times 0 + 1$
 $2 \leq 1$ ~~X~~

ok? (a) yes
 (b) NO

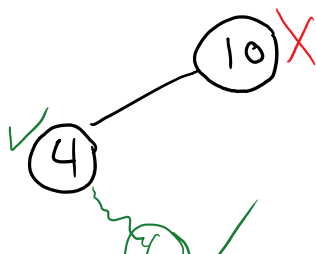


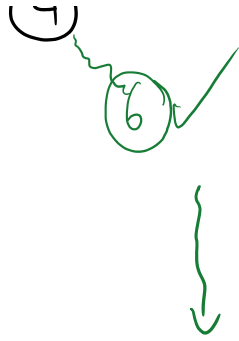
ok? (a) yes
 (b) No

Maintaining size balanced property...

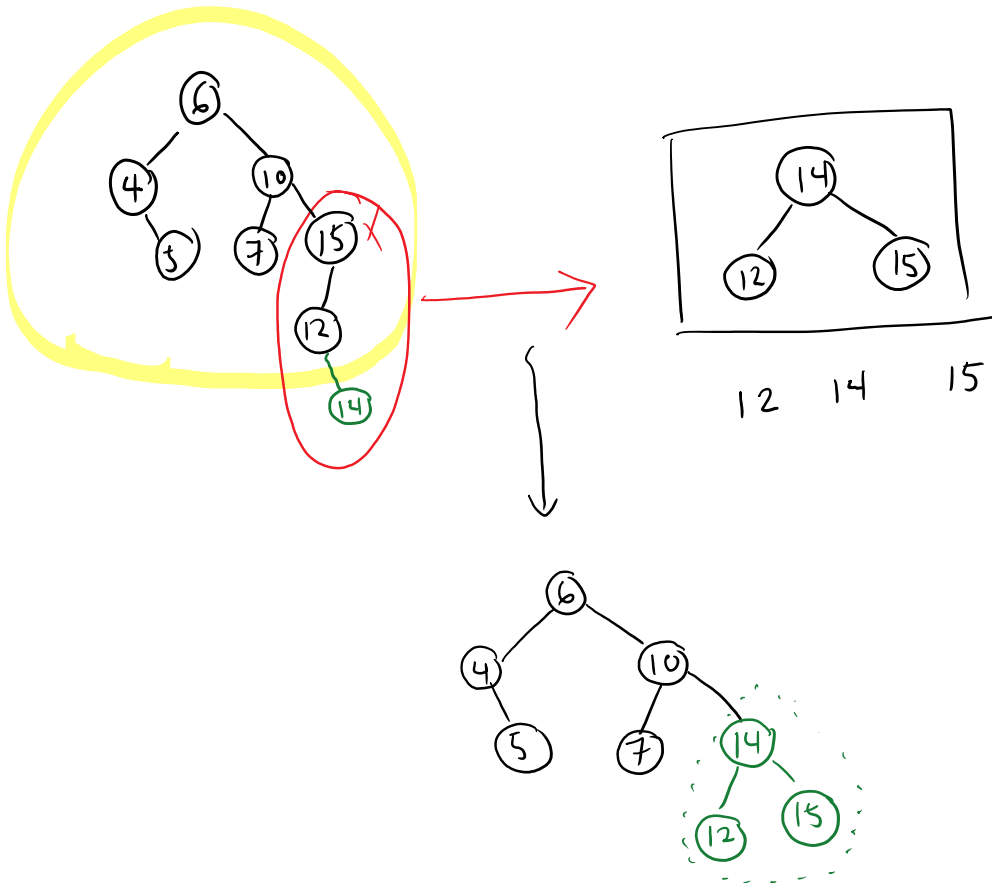
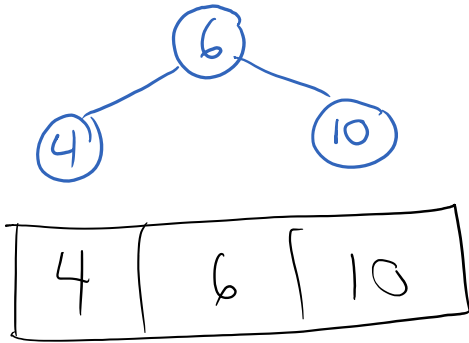
insertion seq:

$10 - 4 - 6$



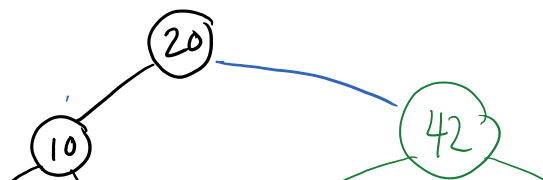
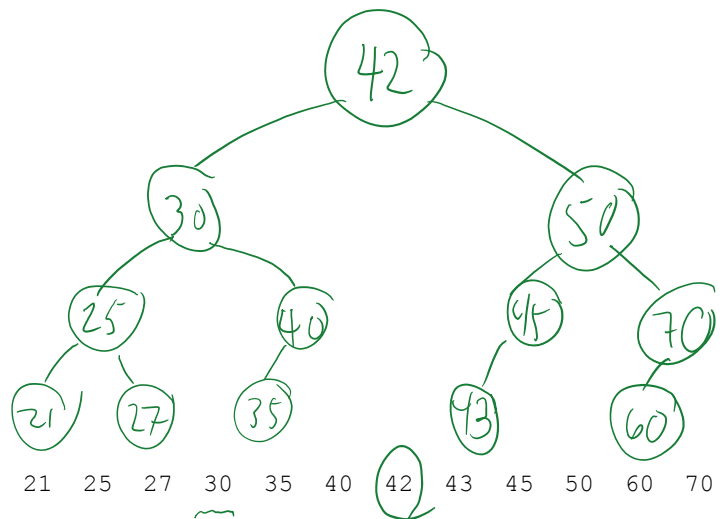
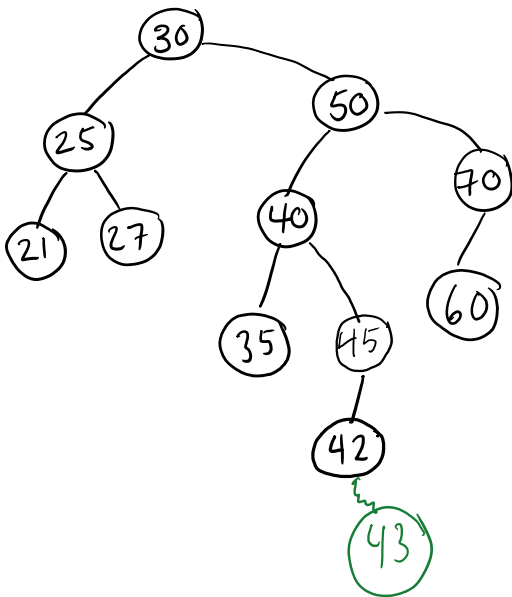
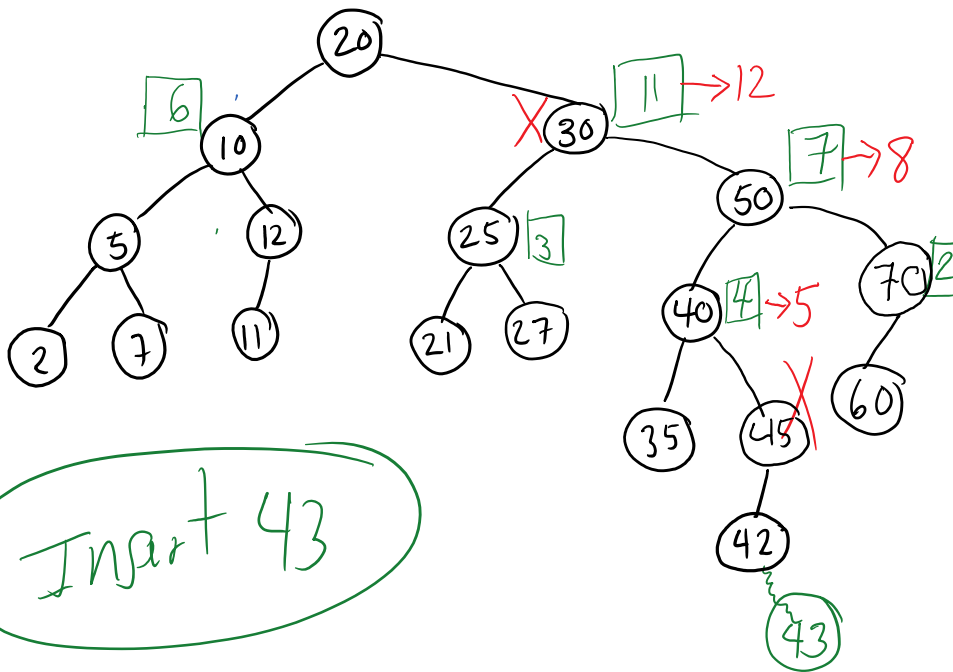


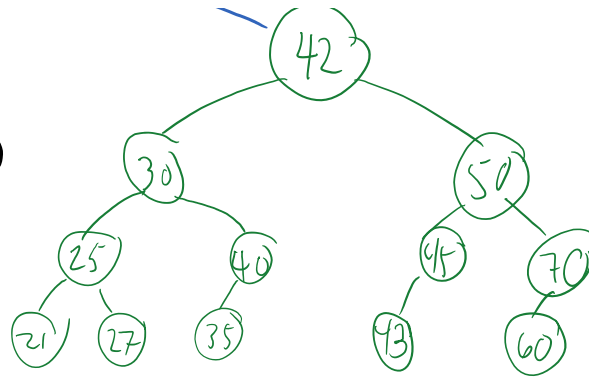
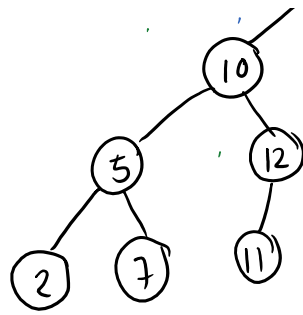
more insertions :

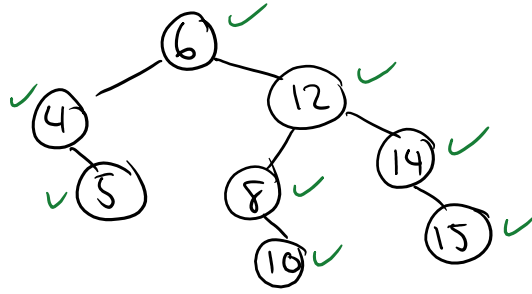


BIGGER EXAMPLE

SIZE-BALANCED?







MAX HEIGHT OF A SIZE-BALANCED
TREE W/ N NODES?

$$\approx O(\log N)$$

$$\left[\approx \log_{3/2}(n) \right]$$

Intuition of Amortized Runtime Claim...

Claim: a sequence of N insertions into an initially empty Size-Balanced BST takes

$O(N \log N)$ time

Thus, "on average" each insertion takes $O(\log N)$ time.

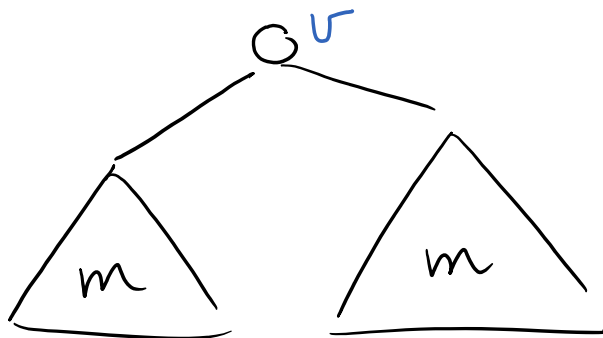
BUT: Any **particular** insertion may **not** be $O(\log N)$

Worst case for one insertion:

$\Theta(N)$

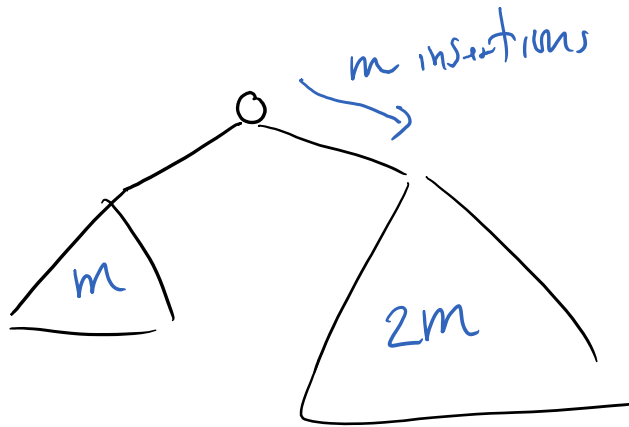
(Not A PROOF...)

consider this config:



$$N = 2m + 1$$

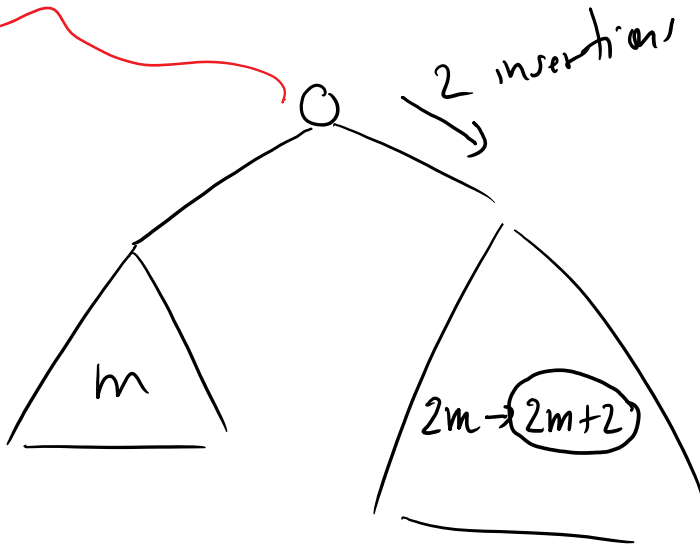
GAME: force v to be re-balanced as soon as possible via seq. of insertions



m insertions into
Right-side...
still ok..

two more!

\$ $3(m+2)$
 $3m+6$



$$N = m + 2m + 2 + 1$$

$$(2m+2) \stackrel{?}{\leq} 2 \times m + 1$$

NOPE!

Rebalance!

$$N_{\text{start}} = 2m + 1$$

$$N_{\text{final}} = 3m + 3$$

Total Work?

$$9m+2 \text{ insertions } O(\underline{(2m+2) \log(N_f)})$$

$$\underline{2m+2} \text{ insertions } \underline{O((2m+2) \log(N_f))}$$

$$+ O(N \log N)$$

Rebalance

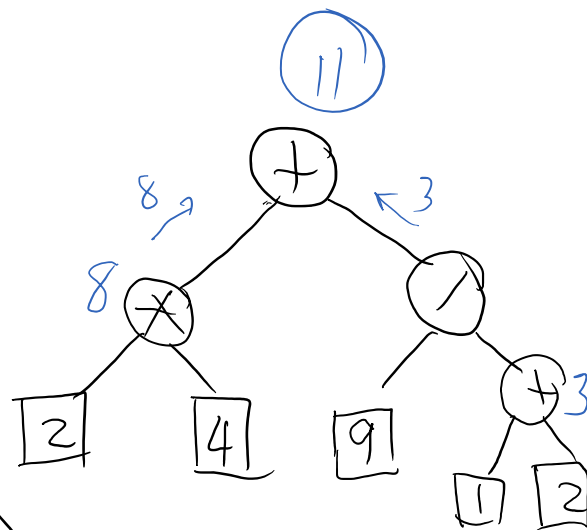
$$O(N)$$

$$\text{Total: } N \log N + N = \Theta(N \log N)$$

$2 + 2$ | $2 \ 2 +$ | $+ 2 \ 2$

$+ \quad \cancel{*} \ 2 \ 4 \quad / \ 9 + 1 \ 2$
 $- \quad \times \quad \sim \quad - \quad - \quad \leftarrow \quad -$

$\langle \text{operator} \rangle \quad \langle \text{exprA} \rangle \quad \langle \text{exprB} \rangle$



$$(2 \times 4) + (9 / (1 + 2))$$