



Radix Sort

4	2	7	6	5	6	1
---	---	---	---	---	---	---

1	2	4	5	6	6	7
---	---	---	---	---	---	---

Radix Sort

4	1	0	0
2	0	1	0
7	1	1	1
6	1	1	0
5	1	0	1
6	1	1	0
1	0	0	1

Radix Sort

4	1	0	<i>0</i>
2	0	1	<i>0</i>
7	1	1	<i>1</i>
6	1	1	<i>0</i>
5	1	0	<i>1</i>
6	1	1	<i>0</i>
1	0	0	<i>1</i>

Sort by least significant bit

Radix Sort

4	1	0	<i>0</i>
2	0	1	<i>0</i>
6	1	1	<i>0</i>
6	1	1	<i>0</i>
7	1	1	<i>1</i>
5	1	0	<i>1</i>
1	0	0	<i>1</i>

Sort by least significant bit

Radix Sort

4
2
6
6
7
5
1

1	0	0
0	1	0
1	1	0
1	1	0
1	1	1
1	0	1
0	0	1

Sort by second least significant bit

Radix Sort

4
5
1
2
6
6
7

1	0	0
1	0	1
0	0	1
0	1	0
1	1	0
1	1	0
1	1	1

Sort by second least significant bit

Radix Sort

4	1	0	0
5	1	0	1
1	0	0	1
2	0	1	0
6	1	1	0
6	1	1	0
7	1	1	1

Sort by third least significant bit

Radix Sort

1	0	0	1
2	0	1	0
4	1	0	0
5	1	0	1
6	1	1	0
7	1	1	1

Sort by third least significant bit

Parallel Radix Sort

4	1	0	<i>0</i>
2	0	1	<i>0</i>
7	1	1	<i>1</i>
6	1	1	<i>0</i>
5	1	0	<i>1</i>
6	1	1	<i>0</i>
1	0	0	<i>1</i>

Sort by least significant bit

Parallel Radix Sort

4
2
7
6
5
6
1

0
0
1
0
1
0
1

Extracting least significant bit can be done in parallel $O(1)$ when $p=n$

Parallel Radix Sort

4
2
7
6
5
6
1

!	
0	1
0	1
1	0
0	1
1	0
0	1
1	0
!	

Using not (!) we can
“select” the ones with 0.
 $O(1)$ when $p=n$



Parallel Radix Sort



Parallel Radix Sort

4
2
7
6
5
6
1

!	
0	1
0	1
1	0
0	1
1	0
0	1
1	0
!	

**We can calculate the new positions using scan (partial sums)
 $O(\log(N))$ when $p=n$**

Parallel Radix Sort

New positions for bits with value 0

4
2
7
6
5
6
1

0	1	1
0	1	2
1	0	2
0	1	3
1	0	3
0	1	4
1	0	4

We can calculate the new positions using scan (partial sums)
 $O(\log(N))$ when $p=n$

MAX



Parallel Radix Sort



Parallel Radix Sort

New positions for bits with value 1

4
2
7
6
5
6
1

0	0	4
0	0	4
1	1	5
0	1	5
1	2	6
0	2	6
1	3	7

We can calculate the new positions using scan (partial sums) and adding the **MAX** new position from the 0 bits.

$O(\log(N))$ when $p=n$



Parallel Radix Sort



Parallel Radix Sort

New positions for all



4
2
7
6
5
6
1

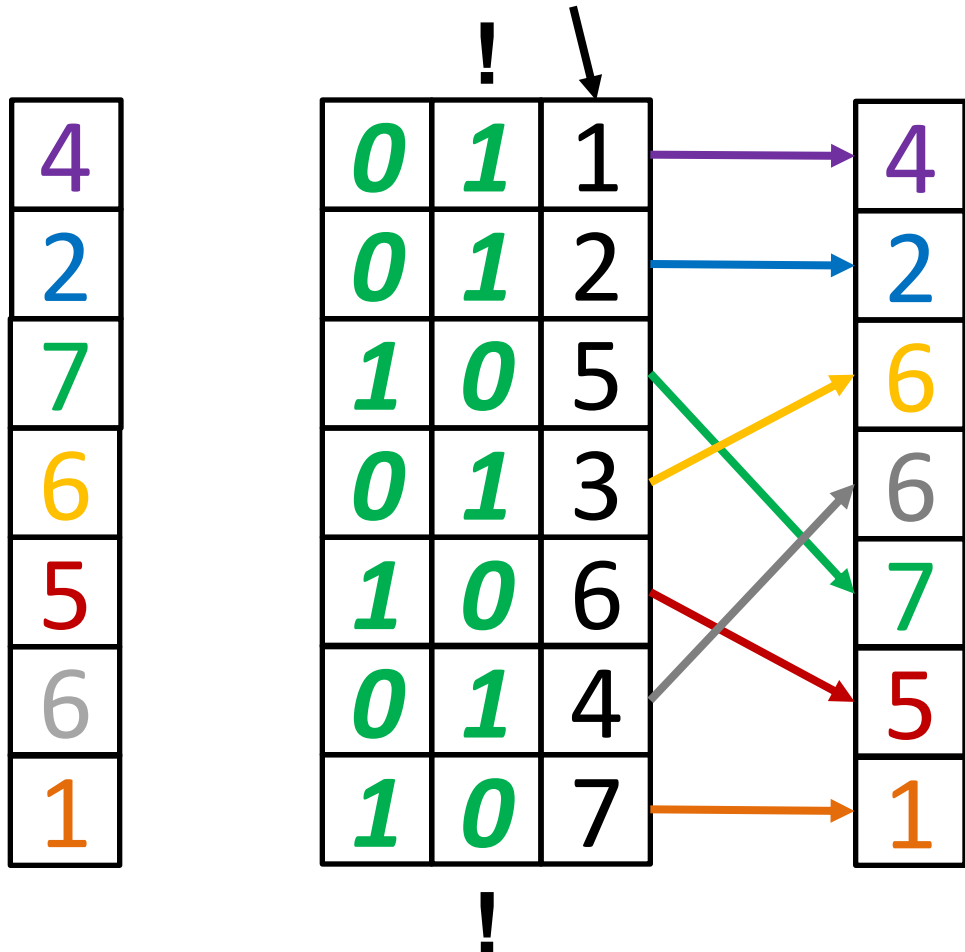
0	1	1
0	1	2
1	0	5
0	1	3
1	0	6
0	1	4
1	0	7



This step does not need to exit in the code. Smart uses of masks permits one to skip it. However it is a good way to visualize our result.

Parallel Radix Sort

New positions for all



Move the values in the new positions.
 $O(1)$ when $p=n$

Radix Sort

4	1	0	<i>0</i>
2	0	1	<i>0</i>
6	1	1	<i>0</i>
6	1	1	<i>0</i>
7	1	1	<i>1</i>
5	1	0	<i>1</i>
1	0	0	<i>1</i>

The elements are now in correct positions as expected after the first step of sequential radix sort.



Parallel Radix Sort



**Repeat all operations with second,
then third, and so on, least
significant bits.**