

# Part I NachOS Problems

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**Explain how does the NachOS FS manage and find free block space? Where is this information stored on the raw disk (which sector)?**

NachOS用 PersistentBitmap \*freeMap 來記錄 free block space 。它會用一個長度為NumSectors的陣列來記錄每個block的資訊，0為尚未被使用，1則是被使用。而當需要尋找free space的時候，會去call FindAndSet()，將未被使用的block設為1，然後回傳sector的索引值。

FreeMap的資訊儲存在 `#define FreeMapSector 0`

**What is the maximum disk size can be handled by the current implementation? Explain why.**

```
const int MagicSize = sizeof(int); const int DiskSize = (MagicSize + (NumSectors * SectorSize));
```

$\text{DiskSize} = 4 + (32 \times 32) \times 128 = 131076\text{B}$

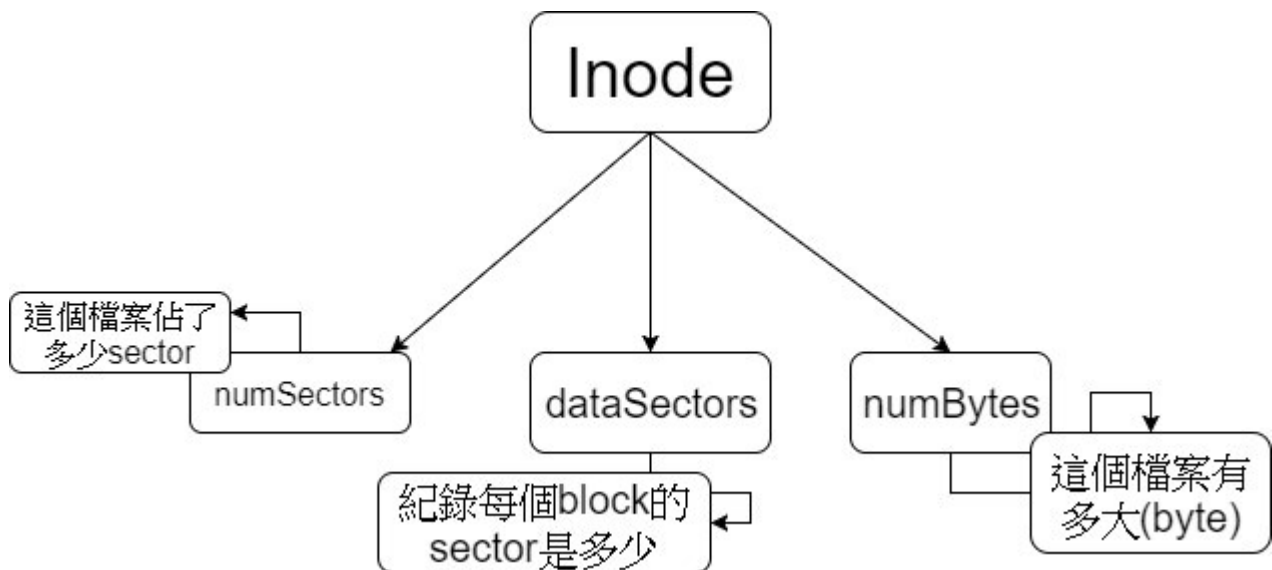
**Explain how does the NachOS FS manage the directory data structure? Where is this information stored on the raw disk (which sector)?**

當建立Directory，會有資料結構DirectoryEntry的table來記錄information。裡面的inUse用來查看是否entry被使用，sector用來尋找fileheader在disk的位置。如果需要一個資料時，會去table中尋找有沒有一樣的filename，然後將它的sector number回傳。

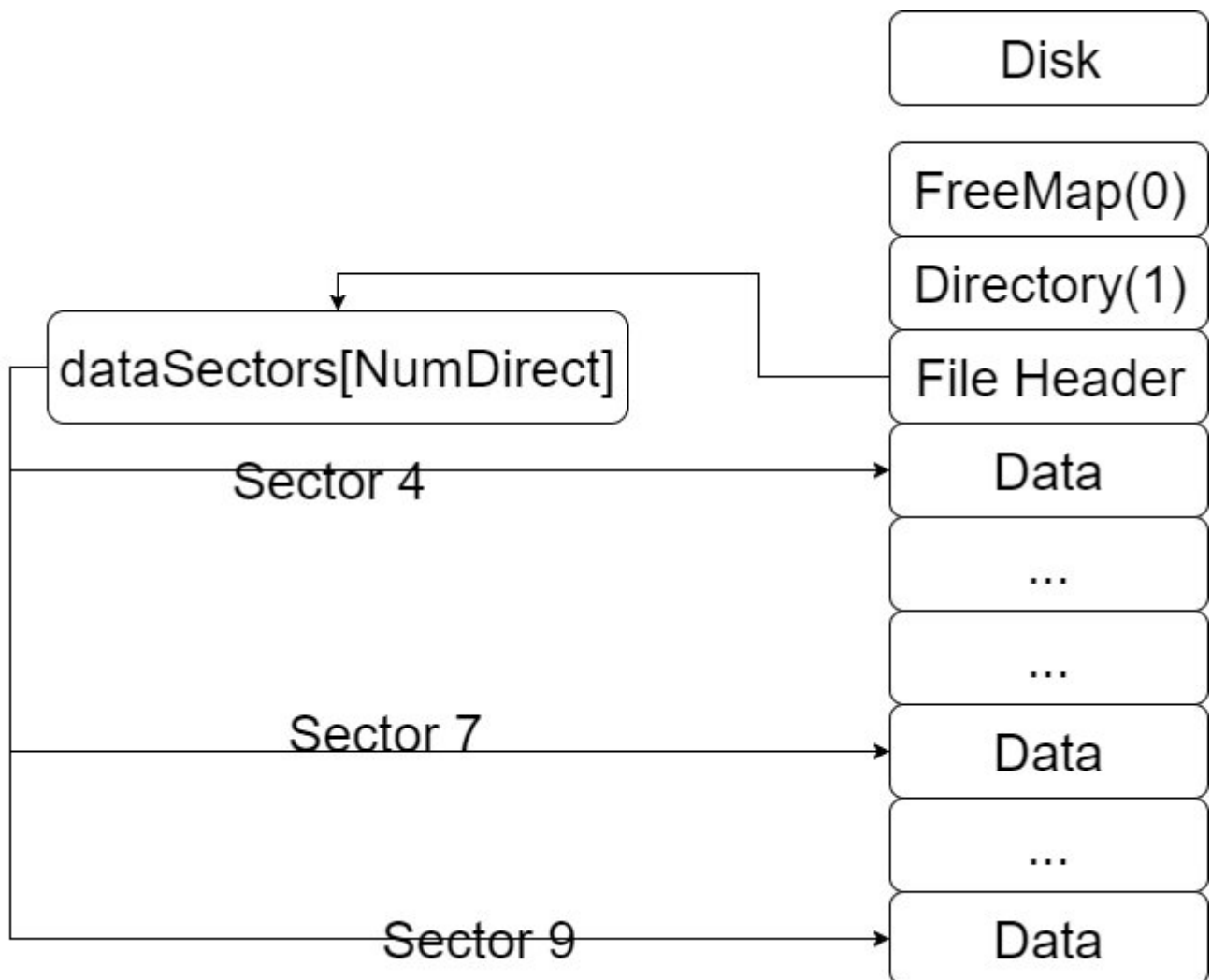
Directory information記錄在 `#define DirectorySector 1`

**Explain what information is stored in an inode, and use a figure to illustrate the disk allocation scheme of current implementation.**

**Inode Structure**



## Disk Allocation Scheme



## Why a file is limited to 4KB in the current implementation?

因為 **NumDirect=30**，而 **MaxFileSize = (NumDirect \* SectorSize) = 30 \* 128B = 3840B**，大約是4KB。

# PartII

## (1) Combine your MP1 file system call interface with NachOS FS

- `int Create(char *name, int size)`

在這裡我們多增加了size變數傳入SysCreate()，其餘皆一樣。

```
1  case SC_Create:
2      val = kernel->machine->ReadRegister(4);
3      size = kernel->machine->ReadRegister(5);
4      {
5          char *filename = &(kernel->machine->mainMemory[val]);
6          //cout << filename << endl;
7          status = SysCreate(filename,size);
8          kernel->machine->WriteRegister(2, (int) status);
9      }
10     kernel->machine->WriteRegister(PrevPCReg, kernel->machine->ReadRegister(PCReg));
11     kernel->machine->WriteRegister(PCReg, kernel->machine->ReadRegister(PCReg) + 4);
12     kernel->machine->WriteRegister(NextPCReg, kernel->machine->ReadRegister(PCReg)+4);
13     return;
14     ASSERTNOTREACHED();
15     break;
```

之後的皆與MP1實作方法一樣 `ksyscall -> interrupt -> kernel -> filesystem` 而在filesystem的Create()我們直接利用老師所提供的function。

```
1  int SysCreate(char *filename,int initialSize)
2  {
3      return kernel->interrupt->CreateFile(filename,initialSize);
4  }
```

```
1  int Interrupt::CreateFile(char *filename,int initialSize)
2  {
3      return kernel->CreateFile(filename,initialSize);
4  }
```

```
1  int Kernel::CreateFile(char *filename,int initialSize)
2  {
3      return filesystem->Create(filename,initialSize);
4  }
```

```
1  bool FileSystem::Create(char *name, int initialSize)
```

- `OpenFileId Open(char *name)`

與MP1相同

`ksyscall -> interrupt -> kernel -> filesystem`

`Open()`一樣使用老師提供的。

```

1 case SC_Open:
2     val = kernel->machine->ReadRegister(4);
3     {
4         char *filename = &(kernel->machine->mainMemory[val]);
5         status = SysOpen(filename);
6         kernel->machine->WriteRegister(2, (int) status);
7     }
8     kernel->machine->WriteRegister(PrevPCReg, kernel->machine->ReadRegister(PCReg));
9     kernel->machine->WriteRegister(PCReg, kernel->machine->ReadRegister(PCReg) + 4);
10    kernel->machine->WriteRegister(NextPCReg, kernel->machine->ReadRegister(PCReg)+4);
11    return;
12    ASSERTNOTREACHED();
13    break;

```

```

1 int SysOpen(char *filename)
2 {
3     return kernel->interrupt->Open(filename);
4 }

```

```

1 int Interrupt::Open(char *filename)
2 {
3     return kernel->Open(filename);
4 }

```

```

1 int Kernel::Open(char *filename)
2 {
3     int fileID = (int)fileSystem->Open(filename);
4     if(fileID!=0) return fileID;
5     return -1;
6 }

```

```

1 OpenFile *FileSystem::Open(char *name)

```

- int Read(char \*buf, int size, OpenFileId id)
  - int Write(char \*buf, int size, OpenFileId id)
  - int Close(OpenFileId id);
- exception -> ksyscall -> interrupt -> kernel -> filesys

```
1 case SC_Read:
2     val = kernel->machine->ReadRegister(4);
3     {
4         buffer = &(kernel->machine->mainMemory[val]);
5         size = kernel->machine->ReadRegister(5);
6         id = kernel->machine->ReadRegister(6);
7         status = SysRead(buffer,size,id);
8         kernel->machine->WriteRegister(2, (int) status);
9     }
10    kernel->machine->WriteRegister(PrevPCReg, kernel->machine->ReadRegister(PCReg));
11    kernel->machine->WriteRegister(PCReg, kernel->machine->ReadRegister(PCReg) + 4);
12    kernel->machine->WriteRegister(NextPCReg, kernel->machine->ReadRegister(PCReg)+4);
13    return;
14    ASSERTNOTREACHED();
15    break;
16 case SC_Write:
17     val = kernel->machine->ReadRegister(4);
18     {
19         buffer = &(kernel->machine->mainMemory[val]);
20         size = kernel->machine->ReadRegister(5);
21         id = kernel->machine->ReadRegister(6);
22         status = SysWrite(buffer,size,id);
23         kernel->machine->WriteRegister(2, (int) status);
24     }
25    kernel->machine->WriteRegister(PrevPCReg, kernel->machine->ReadRegister(PCReg));
26    kernel->machine->WriteRegister(PCReg, kernel->machine->ReadRegister(PCReg) + 4);
27    kernel->machine->WriteRegister(NextPCReg, kernel->machine->ReadRegister(PCReg)+4);
28    return;
29    ASSERTNOTREACHED();
30    break;
31 case SC_Close:
32     val = kernel->machine->ReadRegister(4);
33     {
34         status = SysClose(val);
35         kernel->machine->WriteRegister(2, (int) status);
36     }
37    kernel->machine->WriteRegister(PrevPCReg, kernel->machine->ReadRegister(PCReg));
38    kernel->machine->WriteRegister(PCReg, kernel->machine->ReadRegister(PCReg) + 4);
39    kernel->machine->WriteRegister(NextPCReg, kernel->machine->ReadRegister(PCReg)+4);
40    return;
41    ASSERTNOTREACHED();
42    break;
```

```

1  int SysRead(char *buffer, int size, int id)
2  {
3      return kernel->interrupt->Read(buffer,size,id);
4  }
5  int SysWrite(char *buffer, int size, int id)
6  {
7      return kernel->interrupt->Write(buffer,size,id);
8  }
9  int SysClose(int id)
10 {
11     return kernel->interrupt->Close(id);
12 }

```

```

1  int Interrupt::Read(char *buffer, int size, int id)
2  {
3      return kernel->Read(buffer, size, id);
4  }
5
6  int Interrupt::Write(char *buffer, int size, int id)
7  {
8      return kernel->Write(buffer, size, id);
9  }
10
11 int Interrupt::Close(int id)
12 {
13     return kernel->Close(id);
14 }

```

```

1  int Kernel::Write(char *buffer, int size, int id)
2  {
3      return fileSystem->Write(buffer, size, id);
4  }
5
6  int Kernel::Read(char *buffer, int size, int id)
7  {
8      return fileSystem->Read(buffer, size, id);
9  }
10
11 int Kernel::Close(int id)
12 {
13     return fileSystem->Close(id);
14 }

```

透過Read(),Write()會去call ReadAt與WriteAt去實作更底層的Read,Write。而Close則是把指標給指向NULL，然後回傳1。

```

1 | int FileSystem::Read(char *buffer, int size, int id)
2 | {
3 |     return opfile->Read(buffer, size);
4 | }
5 |
6 | int FileSystem::Write(char *buffer, int size, int id)
7 | {
8 |     return opfile->Write(buffer, size);
9 | }
10 |
11 | int FileSystem::Close(int id)
12 | {
13 |     opfile = NULL;
14 |     return 1;
15 | }

```

## (2) Enhance the FS to let it support up to 32KB file size

我們主要修改的地方為 filehdr.h/.cc

原本的 `MaxFileSize=30*128B`，於是在不改變SectorSize的情況下，我們為了要能support 32KB大小的File，我們將Sectors每32個單位綁在一起，然後再將它串起來，如此一來便能讀寫更大的File。

```

1 | #define MaxListNum 29
2 | #define SectorsPerList 32
3 | #define MaxFileSize MaxListNum * SectorsPerList * SectorSize

```

接著，我們在.h新增一個numLists紀錄有幾個list串起來。

```

1 | private:
2 |     int numLists;

```

一開始假如沒有足夠的空間可以Allocate便return FALSE，而不再繼續執行下面的動作。假如有空間可以擺放的話，便利用Lists的概念看看有多少Lists，還有最後一個Lists有多少sectors，然後以sector為單位寫回disk。

```

1  bool
2  FileHeader::Allocate(PersistentBitmap *freeMap, int fileSize)
3  {
4      char empty[128] = {0};
5      numSectors = divRoundUp(fileSize, SectorSize);
6      numBytes = fileSize;
7      numLists = divRoundUp(numSectors, SectorsPerList);
8
9      if (freeMap->NumClear() < numSectors)
10         return FALSE;          // not enough space
11
12     int SectorsRead = 0;
13     for (int i = 0; i < numLists; i++, SectorsRead += SectorsPerList) {
14         dataSectorLists[i] = freeMap->FindAndSet();
15         ASSERT(dataSectorLists[i] >= 0);
16
17         int lastSector;
18         if (SectorsRead + SectorsPerList > NumSectors)
19             lastSector = NumSectors;
20         else
21             lastSector = SectorsRead + SectorsPerList;
22
23         int *buffer = new int[SectorsPerList];
24         memset(buffer, 0, sizeof(int)*SectorsPerList);
25         for (int j = 0; j < lastSector - SectorsRead; j++) {
26             buffer[j] = freeMap->FindAndSet();
27             kernel->synchDisk->WriteSector(buffer[j], empty);
28             ASSERT(buffer[j] >= 0);
29         }
30         kernel->synchDisk->WriteSector(dataSectorLists[i], (char*) buffer);
31         delete [] buffer;
32     }
33     return true;
34 }

```

用for迴圈尋找file所涵蓋住的Lists，接著再利用for迴圈將每個Lists的sector(以及最後一個Lists剩下的sector)以sector為單位做Clear()的動作。



```

1 void
2 FileHeader::Deallocate(PersistentBitmap *freeMap)
3 {
4     int SectorsRead = 0;
5     for (int i = 0; i < numLists; i++, SectorsRead += SectorsPerList) {
6         int lastSector;
7         if (SectorsRead + SectorsPerList > numSectors)
8             lastSector = numSectors;
9         else
10            lastSector = SectorsRead + SectorsPerList;
11
12            int* buffer = new int[SectorsPerList];
13            kernel->synchDisk->ReadSector(dataSectorLists[i], (char*) buffer);
14            for (int j = 0; j < lastSector; j++) {
15                ASSERT(freeMap->Test((int) buffer[j]));
16                freeMap->Clear((int)buffer[j]);
17            }
18            delete [] buffer;
19        }
20    }
21

```

我們利用sectorID去尋找資料放在第幾個sector，之後再將sectorID / SectorsPerList與sectorID % SectorsPerList。目的是要找出資料位在第幾個list的第幾個sector。之後，將位在某個list裡的東西撈到buffer上，再找尋buffer[idInList]，這樣就能找到我們所要的資料。

```

1 int
2 FileHeader::ByteToSector(int offset)
3 {
4     //return(dataSectors[offset / SectorSize]);
5     int sectorID = offset / SectorSize;
6     int listID = sectorID / SectorsPerList, idInList = sectorID % SectorsPerList;
7     int *buffer = new int[SectorsPerList];
8     kernel->synchDisk->ReadSector(dataSectorLists[listID], (char *) buffer);
9     // get the SectorNum
10    int retVal = buffer[idInList];
11    delete [] buffer;
12    return retVal;
13 }
14

```

## Part III Implementation

### How to support subdirectory

#### 1. Entering FileSystem API

我們使用以下這個function來實作subdirectory。當NachOS收到mkdir指令時會將mkdirflag設為真，程式本身依靠這個flag來判斷該創造一個檔案(FileSystem::Create)或是資料夾(FileSystem::CreateDir)

```
1  bool FileSystem::CreateDir(char *name)
2  {
3      Directory *directory;
4      PersistentBitmap *freeMap;
5      FileHeader *hdr;
6      int sector;
7      bool success;
8
9      directory = new Directory(NumDirEntries);
10     directory->FetchFrom(directoryFile);
11
12     if (directory->Find(name) != -1)
13         success = FALSE;
14     else
15     {
16         freeMap = new PersistentBitmap(freeMapFile, NumSectors);
17         sector = freeMap->FindAndSet();
18         if (sector == -1)
19             success = FALSE;
20         else if(!directory->Add(name, sector, 'D'))
21             success = FALSE;
22         else
23         {
24             hdr = new FileHeader;
25             if (!hdr->Allocate(freeMap, DirectoryFileSize))
26                 success = FALSE;
27             else
28             {
29                 success = TRUE;
30                 hdr->WriteBack(sector);
31                 directory->WriteBack(directoryFile);
32                 freeMap->WriteBack(freeMapFile);
33             }
34         }
35     }
36     delete freeMap;
37     delete hdr;
38     delete directory;
39     return success;
40 }
```

我們在DirectoryEntry新增了type來判斷這個檔案真的是檔案或者是資料夾

## 2. Parsing Path

原本的Directory.\* 不支援subdirectory的查找，我們修改了下列function **1. int Directory::Find(char name)**

```

1  int Directory::Find(char *name)
2  {
3      name++;
4      char localName[256] = {0};
5      char localID = 0;
6      bool findNext = false;
7      while (name[0] != '\0') {
8          if (name[0] == '/') {
9              findNext = true;
10             break;
11         }
12         localName[localID++] = name[0];
13         name++;
14     }
15     int i = FindIndex(localName);
16     if (i != -1) {
17         if (findNext) {
18             OpenFile* nextDirectory = new OpenFile(table[i].sector);
19             Directory* nextDir = new Directory(NumDirEntries);
20             nextDir->FetchFrom(nextDirectory);
21             int result = nextDir->Find(name);
22             delete nextDirectory;
23             delete nextDir;
24             return result;
25         } else {
26             return table[i].sector;
27         }
28     } else {
29         return -1;
30     }
31 }

```

1. 一開始把name加一，因為我們不想把slash放進檔案名稱裡儲存
2. 程式7~14行將Path做parsing來判斷需不需要subdirectory visiting。一個字元一個字元的讀直到讀到slash或是結尾，如果有slash代表這是subdirectory，將findNext拉起來準備做recursion進入子資料夾內。
3. localName紀錄直到slash前的path，也就是第一個需要進入的資料夾。
4. 如果有子資料夾需要進入的話，就打開它並call它的Find(name)來做recursion
5. 如果不需要就直接return檔案的sector

## 2. bool Directory::Add(char name, int newSector, char inType)

```

1 bool Directory::Add(char *name, int newSector, char inType)
2 {
3     if (Find(name) != -1)
4         return FALSE;
5
6     char Path[256] = {0};
7     char File[9] = {0};
8     int len = strlen(name), slash, tmpID = 0;
9     for (int i = len - 1; i >= 0; i--) {
10         if (name[i] == '/') {
11             slash = i;
12             break;
13         }
14     }
15
16     for (int i = slash+1; i < len; i++) {
17         File[tmpID++] = name[i];
18     }
19     for (int i = 0; i < slash; i++) {
20         Path[i] = name[i];
21     }
22
23     if (Path[0] != 0) {
24         int sector = Find(Path);
25         OpenFile* nextDirectory = new OpenFile(sector);
26         Directory* nextDir = new Directory(NumDirEntries);
27         nextDir->FetchFrom(nextDirectory);
28
29         for (int i = 0; i < tableSize; i++) {
30             if (!nextDir->table[i].inUse) {
31                 nextDir->table[i].inUse = true;
32                 strncpy(nextDir->table[i].name, File, FileNameMaxLen);
33                 nextDir->table[i].sector = newSector;
34                 nextDir->table[i].type = inType;
35                 nextDir->WriteBack(nextDirectory);
36
37                 delete nextDirectory;
38                 delete nextDir;
39                 return true;
40             }
41         }
42     } else {
43         for (int i = 0; i < tableSize; i++) {
44             if (!table[i].inUse) {
45                 table[i].inUse = true;
46                 strncpy(table[i].name, File, FileNameMaxLen);
47                 table[i].sector = newSector;
48                 table[i].type = inType;
49                 return true;
50             }
51         }
52     }
53     return false;    // no space.  Fix when we have extensible files.

```

1. 和Find一樣需要Parsing，不同的是這次是從後面Parse回來直到找到第一個slash，slash後面的Name就是檔案名稱，前面的就是Path。如果沒有subdirectory，Path就會是空的。
2. 如果Path[] != 0，就call Find()來拿到該寫進的正確directory，並且寫入disk內。
3. 如果Path == 0，代表自己就是該寫進去的正確資料夾，直接寫到裡面，因為會return回去FileSystem寫入disk，所以這邊不需要寫入。

### **3. bool Directory::Remove(char name)**

```

1  bool Directory::Remove(char *name)
2  {
3      if (Find(name) == -1)
4          return false;
5
6      char Path[256] = {0};
7      char File[9] = {0};
8      int len = strlen(name), tmpID = 0, slash;
9      for (int i = len-1; i >= 0; i--) {
10         if (name[i] == '/') {
11             slash = i;
12             break;
13         }
14     }
15
16     for (int i = slash + 1; i < len; i++) {
17         File[tmpID++] = name[i];
18     }
19     for (int i = 0; i < slash; i++) {
20         Path[i] = name[i];
21     }
22
23     if (Path[0] != 0) {
24         int sector = Find(Path);
25         OpenFile* nextDirectory = new OpenFile(sector);
26         Directory* nextDir = new Directory(NumDirEntries);
27         nextDir->FetchFrom(nextDirectory);
28
29         int id = nextDir->FindIndex(File);
30         if (id == -1)
31             return false;
32
33         nextDir->table[id].inUse = false;
34         nextDir->WriteBack(nextDirectory);
35         delete nextDirectory;
36         delete nextDir;
37         return true;
38     } else {
39         int id = this->FindIndex(name);
40         if (id == -1)
41             return false;
42         this->table[id].inUse = false;
43         return true;
44     }
45 }

```

大抵上和Add()很像，只是差別在於一個是Remove一個是Add

## Support up to 64 entries in a directory

蠻簡單的，把NumdirEntries從10改成64即可

## Support recursively list the file/directory in a directory

---

```
1 void Directory::RecursiveList(int depth)
2 {
3     for (int i = 0; i < tableSize; i++) {
4         if (table[i].inUse) {
5             for (int j = 0; j < depth*8; j++)
6                 putchar(' ');
7             printf("[Entry No.%d]: %s %c\n", i, table[i].name, table[i].type);
8             if (table[i].type == 'D') {
9                 OpenFile* nextDirectory = new OpenFile(table[i].sector);
10                Directory* nextDir = new Directory(NumDirEntries);
11                nextDir->FetchFrom(nextDirectory);
12
13                nextDir->RecursiveList(depth+1);
14                delete nextDirectory;
15                delete nextDir;
16            }
17        }
18    }
19 }
```

呼叫上面的RecursiveList()，印出來時順便判斷是Directory還是File，如果是D就進去traverse一遍把東西都印出來。

## Bonus II : Remove a file or recursively remove the directory

---

我們的Recursive Remove走和正常remove一樣的flow，Directory::remove()會自己parse Path並刪除。

## Contribution

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- 陳麒懋：Part I、Part II、Report Part I、II
- 鍾昀誼：Part II、PartIII、Report PartIII