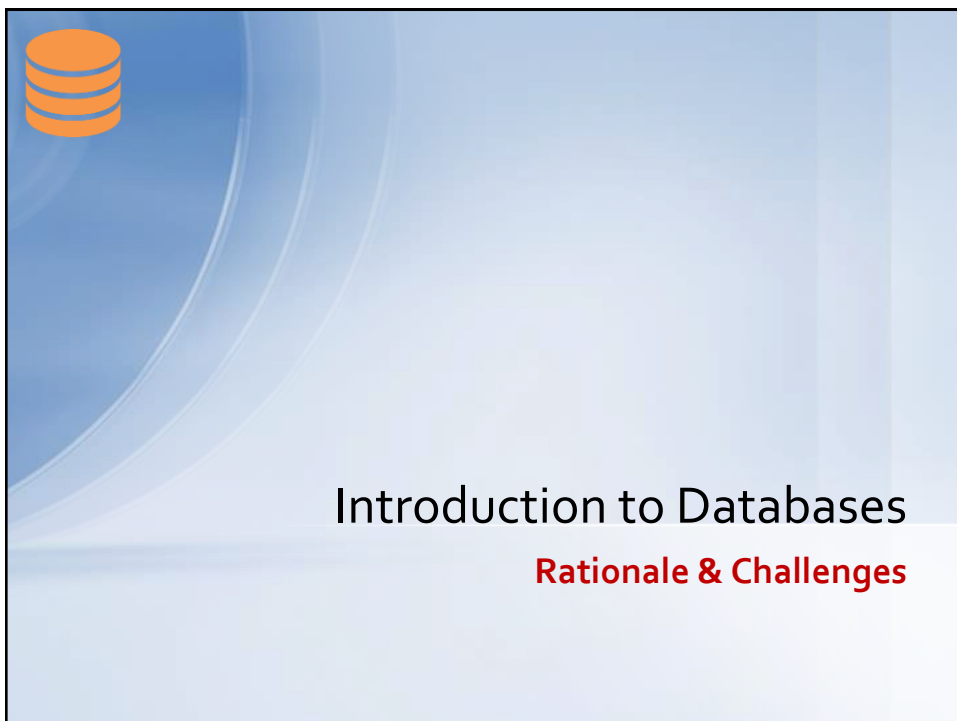



1



2



Where can we start!?


- **Definition:**
 - A **database** is a collection of **structured data** stored on **persistent storage**.
- **Environment and Contextual**
 - **Available** to a large number of end users.
 - As in “Access my data when I need it ...”
 - No delay, no access denial, no loss of data
 - **Correctness**.
 - As in “What’s my balance?” gets the same value until it is changed.
- **What’s different from, say, programming?**
 - Data **persistence**;
 - Data **sharing**;
 - Data **independence**.

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Example Database Instance (From Elmasri)

- A lot of databases are built with a particular data modelling language called **relational**.
- This is an example of a relational database:
 - It is built by set of **relations** (e.g., tables).
 - A group of tables are sometimes grouped in a **schema**
- There are two aspects immediately visible:
 - **Structures**;
 - **Data** (fitting the given structure).

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STUDENT

Name	Student_number	Class	Major
Smith	17	1	CS
Brown	8	2	CS

COURSE

Course_name	Course_number	Credit_hours	Department
Intro to Computer Science	CS1310	4	CS
Data Structures	CS3320	4	CS
Discrete Mathematics	MATH2410	3	MATH
Database	CS3380	3	CS

SECTION

Section_identifier	Course_number	Semester	Year	Instructor
85	MATH2410	Fall	07	King
92	CS1310	Fall	07	Anderson
102	CS3320	Spring	08	Knuth
112	MATH2410	Fall	08	Chang
119	CS1310	Fall	08	Anderson
135	CS3380	Fall	08	Stone

GRADE REPORT

Student_number	Section_identifier	Grade
17	112	B
17	119	C
8	85	A
8	92	A
8	102	B
8	135	A

PREREQUISITE

Course_number	Prerequisite_number
CS3380	CS3320
CS3380	MATH2410
CS3320	CS1310

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Database schema
Structures built by a relational data modelling

STUDENT	Name	StudentNumber	Class	Major
	Smith	17	1	CS
	Brown	8	2	CS

COURSE	CourseName	CourseNumber	CreditHours	Department
	Intro to Computer Science	CS1310	4	CS
	Data Structures	CS3320	4	CS
	Discrete Mathematics	MATH2410	3	MATH
	Database	CS3380	3	CS

SECTION	SectionIdentifier	CourseNumber	Semester	Year	Instructor
	85	MATH2410	Fall	98	King
	92	CS1310	Fall	98	Anderson
	102	CS3320	Spring	99	Knuth
	112	MATH2410	Fall	99	Chang
	119	CS1310	Fall	99	Anderson
	135	CS3380	Fall	99	Stone

GRADE_REPORT	StudentNumber	SectionIdentifier	Grade
	17	112	B
	17	119	C
	8	85	A
	8	92	A
	8	102	B
	8	135	A

PREREQUISITE	CourseNumber	PrerequisiteNumber
	CS3380	CS3320
	CS3380	MATH2410
	CS3320	CS1310

THE END!



Database instance
Data fitting respective relations

STUDENT	Name	StudentNumber	Class	Major
	Smith	17	1	CS
	Brown	8	2	CS


COURSE	CourseName	CourseNumber	CreditHours	Department
	Intro to Computer Science	CS1310	4	CS
	Data Structures	CS3320	4	CS
	Discrete Mathematics	MATH2410	3	MATH
	Database	CS3380	3	CS

SECTION	SectionIdentifier	CourseNumber	Semester	Year	Instructor
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	119	CS1310	Fall	99	Anderson
	135	CS3380	Fall	99	Stone

GRADE REPORT	StudentNumber	SectionIdentifier	Grade
	17	112	B
	17	119	C
	8	85	A
	8	92	A
	8	102	B
	8	135	A

PREREQUISITE	CourseNumber	PrerequisiteNumber
	CS3380	CS3320
	CS3380	MATH2410
	CS3320	CS1310

THE END!



Example Database Structure and Data (From Elmasri)

STUDENT

Name	Student_number	Class	Major
Smith	17	1	CS
Brown	8	2	CS


COURSE

Course_name	Course_number	Credit_hours	Department
Intro to Computer Science	CS1310	4	CS
Data Structures	CS3320	4	CS
Discrete Mathematics	MATH2410	3	MATH
Database	CS3380	3	CS

- Note:**
 - Table structure composition (e.g., attribute name and data type (assumed));
 - Table's data;
 - Are there any other rules which the data must follow?
 - Primary key, limiting possible values in a data type, foreign key, etc.

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For what do we use a database? How do we use a database?

- Typical operations over a database are:**
 - accepting structured data (ADD, AMEND and PURGE/DELETE).
- How do we use databases?**
 - computing functions on the present data;
 - sorting the structured data; and
 - providing responses to questions (querying).

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How is usage managed and controlled?

- **A Database Management System (DBMS)**
 - Is a computer-based application tool set for: defining, creating, manipulating, controlling, and managing databases.
- **DBMS Requirements:**
End users, developers, administrators demand **efficiency** in terms of allocation of computational resources and storage space.
 - **Reliability** (e.g., if the system falls, can it relive itself)
 - **Openness** (e.g., in term of data connectivity)
 - **Scalability** (e.g., storage and throughput)

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DBMS Functionality

- **DBMS also facilitate and insulates data access by system's end users (including "computer" naive *end users*).**
- **Three important functionalities include:**
 - *Query Processing*
 - *Transaction Management*
 - *In some context TP subsumes QP*
 - *Storage management*
- **Two main areas of concern:**
 - *Consistency* (e.g., returns the same results under invariant states), and
 - *Efficiency* (e.g., computational time and space) of DBMS activities.

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DBMS Advantages

- **Some advantages:**
 - Execute a query in an efficient and opportunistic approach – **query optimisation**;
 - Iron out sharing inconsistencies between conflicting queries and transactions – **concurrency control**;
 - Undoing the effects of incomplete transactions under a system's instructions or in rectifying a system failure – **recovery**.
 - **Note:** system here is most general: OS, DBMS, hard disk, etc.


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Query Processing (and Optimisation)

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Query Processing & Optimisation


- We now briefly describe the structure and functionality of a *query processor* -- a process that **transforms** and **executes** a user-defined query against a database state and returns those instances that satisfy the query.
- A **declarative query** is a data access expression that indicates the structure and the property of the result (e.g. **what you want**). In contrast to a procedural query that specifies structure, property and an exact computational program(e.g. **how to get**).
 - Efficiency of a **declarative relational query processor** (for languages such as SQL) is a major project.
 - **Query optimisation** is the tuning of queries and storage structures to favour the performance of *some* frequently used query.
- **What is a data access path?**
 - A procedure to traverse and access required data items through the use of available look-up methods associated to the data file organisation.

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QP Example:

Return all student's name and grade in the database unit.

STUDENT	Name	StudentNumber	Class	Major
✓	Smith	17	1	CS
✓	Brown	8	2	CS

COURSE	CourseName	CourseNumber	CreditHours	Department
✓	Intro to Computer Science	CS1310	4	CS
	Data Structures	CS3320	4	CS
	Discrete Mathematics	MATH2410	3	MATH
	Database	CS3380	3	CS

SECTION	SectionIdentifier	CourseNumber	Semester	Year	Instructor
	85	MATH2410	Fall	98	King
	92	CS1310	Fall	98	Anderson
	102	CS3320	Spring	99	Knuth
	112	MATH2410	Fall	99	Chang
	119	CS1310	Fall	99	Anderson
	135	CS3380	Fall	99	Stone

GRADE_REPORT	StudentNumber	SectionIdentifier	Grade
✓	17	112	B
	17	119	C
	8	85	A
	8	92	A
	8	102	B
	8	135	A

PREREQUISITE	CourseNumber	PrerequisiteNumber
	CS3380	CS3320
	CS3380	MATH2410
	CS3320	CS1310

Step One: Identify the attributes to output
i.e. Student / Name & Grade_report / Grade.

Step Two: Identify the relationships between structures of interest to the query
i.e. a) Student / StudentNumber to Grade_report / StudentNumber,
b) Course / CourseNumber to Section / CourseNumber,
c) Section / SectionIdentifier to Grade_report / SectionIdentifier.

Step Three: Apply any restrictions
i.e. Course / CourseName is equal to 'Database'.


Step Four: Compute the output.
i.e. Student 'Brown' with a grade of 'A'.

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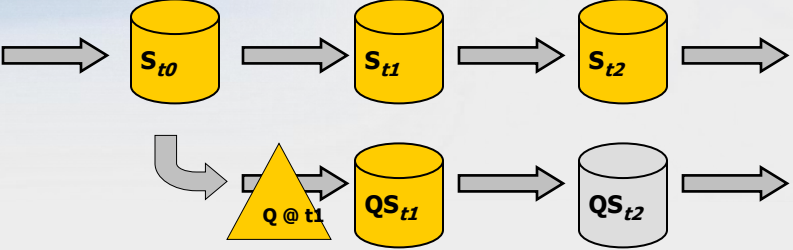
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QP while the database is available (to other users)



At time t1, query Q's state (i.e. result QS) is related to the database state at t0.

Database states at t0 and t1 are equivalent (i.e. $St_0 = St_1$).

Database states at t1 and t2 are not necessary equivalent (i.e. $St_2 \neq St_1$).


At time t2, query Q's state (i.e. result QS at t2) is NOT related to the database sate at t2.

Query Q states at t1 and t2 are equivalent (i.e. $QSt_1 = QSt_2$).

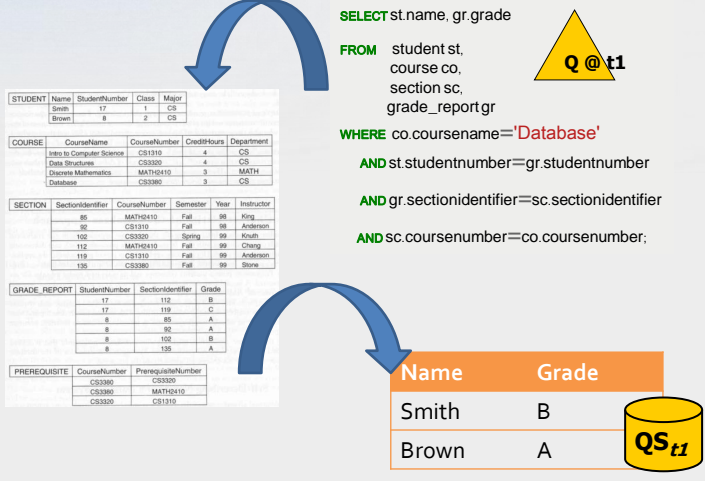
The end!

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QP Example – another view



SELECT st.name, gr.grade

FROM student st,
course co,
section sc,
grade_report gr

WHERE co.coursename='Database'

AND st.studentnumber=gr.studentnumber

AND gr.sectionidentifier=sc.sectionidentifier

AND sc.coursenumber=co.coursenumber;

STUDENT	Name	StudentNumber	Class	Major
Smith		17	1	CS
Brown		8	2	CS

COURSE	CourseName	CourseNumber	CreditHours	Department
Intro to Computer Science		CS1310	4	CS
Data Structures		CS3320	4	CS
Discrete Mathematics		MATH2410	3	MATH
Database		CS3380	3	CS

SECTION	SectionIdentifier	CourseNumber	Semester	Year	Instructor
85	MATH2410		Fall	98	King
92	CS1310		Fall	98	Anderson
102	CS3320		Spring	99	Krauth
112	MATH2410		Fall	99	Cheng
119	CS1310		Fall	99	Anderson
135	CS3380		Fall	99	Stone

GRADE REPORT	StudentNumber	SectionIdentifier	Grade
	17	112	B
	17	119	C
	8	85	A
	8	92	A
	8	102	B
	8	135	A

PREREQUISITE	CourseNumber	PrerequisiteNumber
	CS3380	CS3320
	CS3380	MATH2410
	CS3320	CS1310

Name	Grade
Smith	B
Brown	A


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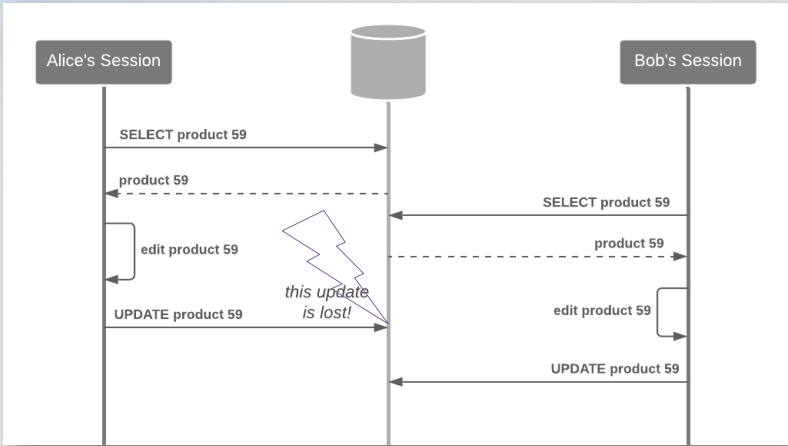


Transaction Processing and Recovery

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Sharing Issues – e.g., lost update type




```
sequenceDiagram
    participant Alice as Alice's Session
    participant Bob as Bob's Session
    participant DB as Database

    Alice->>DB: SELECT product 59
    DB-->Alice: product 59
    Note over Alice: edit product 59
    Alice->>DB: UPDATE product 59
    Bob->>DB: SELECT product 59
    DB-->Bob: product 59
    Note over Bob: edit product 59
    Bob->>DB: UPDATE product 59
    Note over Alice,DB: this update is lost!
```

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Transaction Modelling

- An unrestricted sharing over data is problematic:
 - The example on the right is called *dirty read* as T2 reads what T1 has written. If Commit of T1 does not pass then T2 is wrong.
- An explicit sequence of data manipulation commands over a database are abstracted as a single logical update (i.e a **transaction**).
 - A transaction is **valid** if **all** its sub-parts (e.g. updates, deletes) have succeeded. A valid transaction can **commit**, otherwise we have to **rollback** the transaction.
- Operationally this complicates our scenario when allowing for data sharing! The DBMS typically insulates this from users (programmers and end users alike!).


$$S = \begin{bmatrix} T1 & T2 \\ R(A) & \\ W(A) & \\ & R(A) \\ & W(A) \\ & R(B) \\ & W(B) \\ & Com. \\ R(B) & \\ W(B) & \\ Com. & \end{bmatrix}$$

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TP Principles - ACID


- Transaction's ACID principles for short-lived transactions are (over and above valid transactions):
 - Atomicity
 - The transaction executes completely or not at all,
 - Consistency
 - The transaction preserves the internal consistency of the database.
 - Isolation
 - The transaction executes as if it were running alone, with no other transactions.
 - Durability
 - The transaction's result will not be lost in the future.

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Transaction Modeling: an Example

- “Ensure that the employees' salary is upped by 25% and each salesperson commission is augmented by Euro 250.”

Remark SQL command sequence

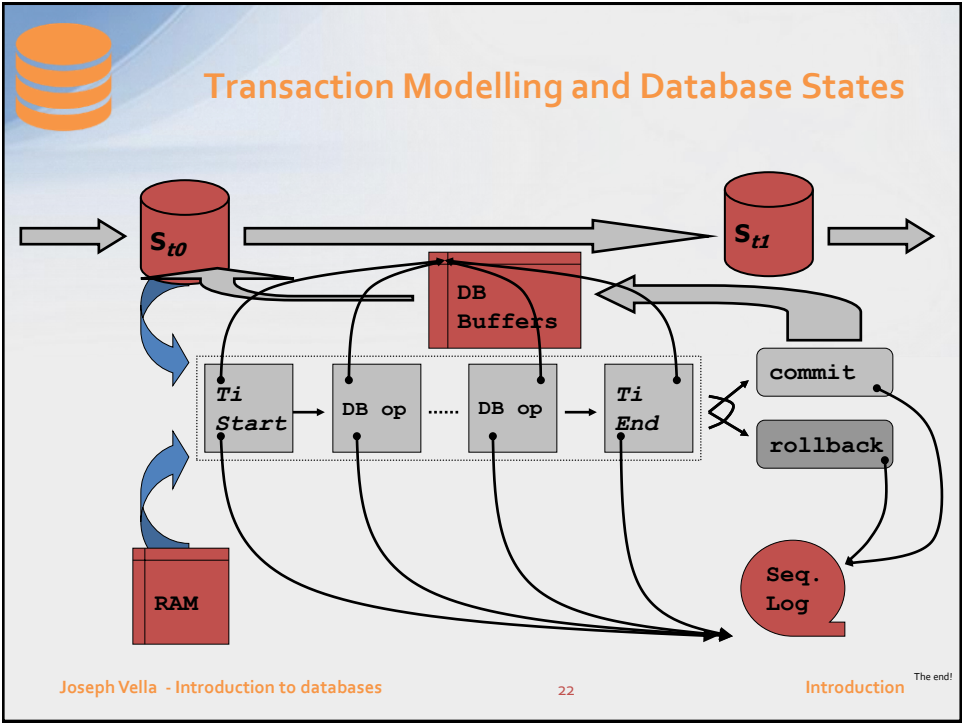
```
start transaction;  
  update emp  
    set sal = sal * 1.25;  
remark implicit mechanism - on error rollback  
  update emp  
    set comm = comm + 250  
  where comm is not null;  
remark implicit mechanism - on error rollback  
commit;  
end transaction;
```

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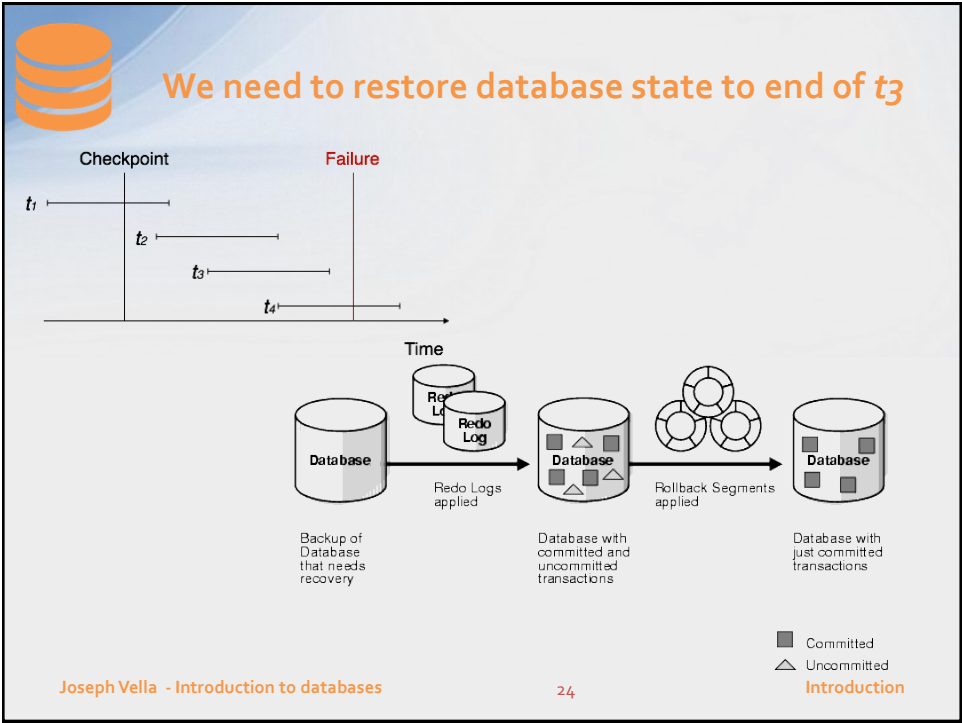
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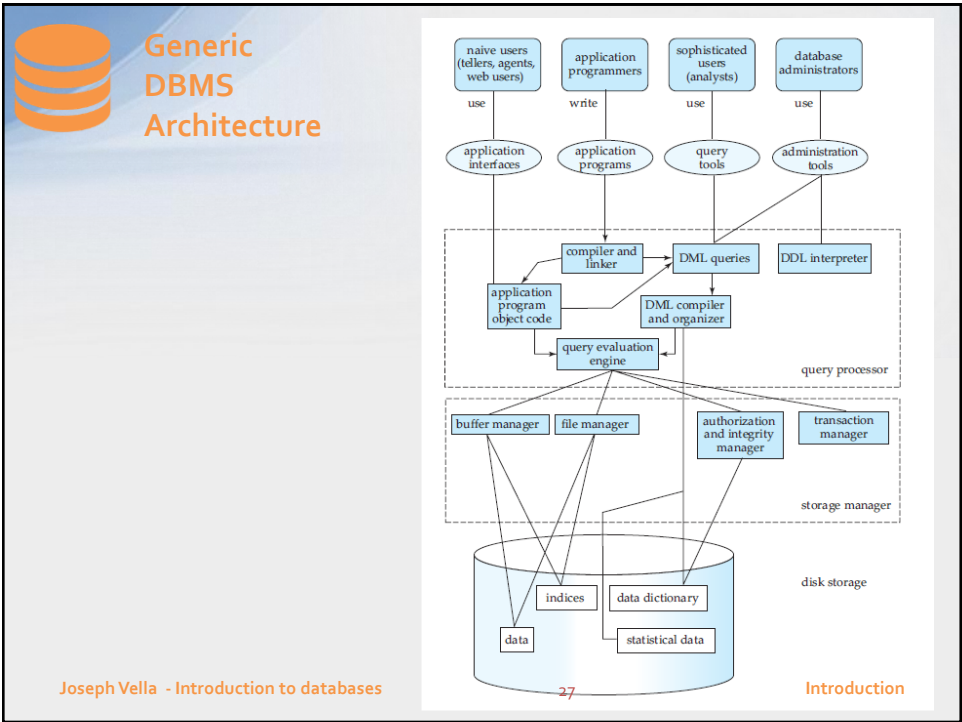
A presentation slide with a light blue background and a white footer bar. In the top left corner, there is an orange icon of a database cylinder. The title "DBMS components" is written in orange in the top right. A list of components is shown in red bullet points. The footer contains the text "Joseph Vella - Introduction to databases", "26", and "Introduction".

DBMS components


- Procedural and Declarative database programming languages;
- Database browsers;
- Administration and Data Dictionary tools;
- Database connectivity protocols (e.g. ODBC);
- Performance tuning tools;
- Import / Export Data utilities;
- CASE (and BPR) tools for conversion of a conceptual database design into a set of data modelling constructs;
- Forms and Report generators;

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Available DBMS (a biased list 😊)



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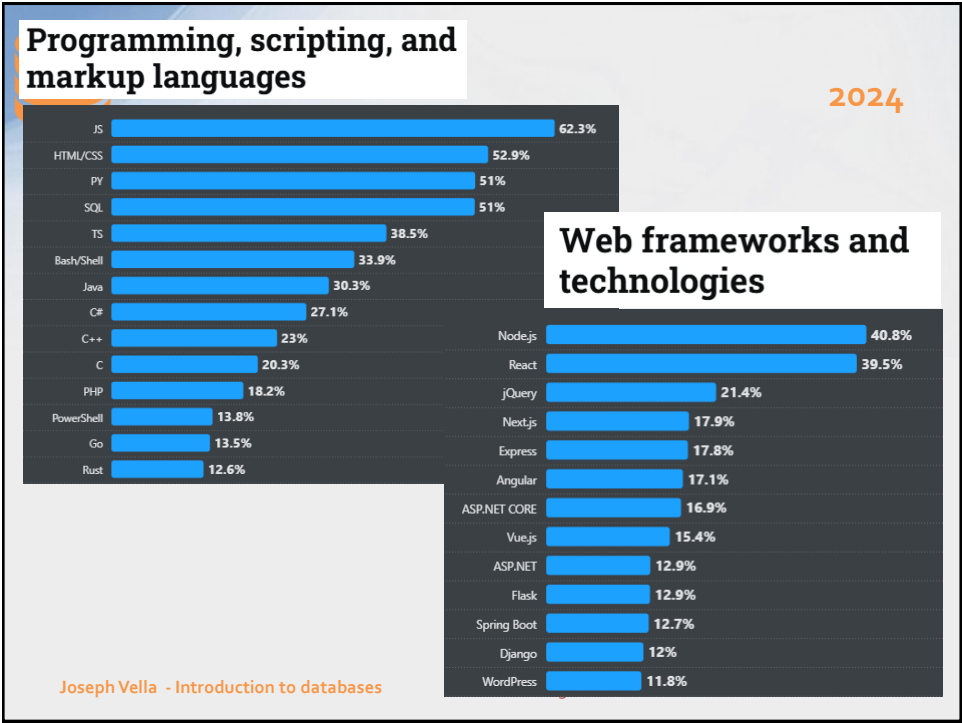


Available DBMS (a biased collective list) Stackoverflow 2024 ()

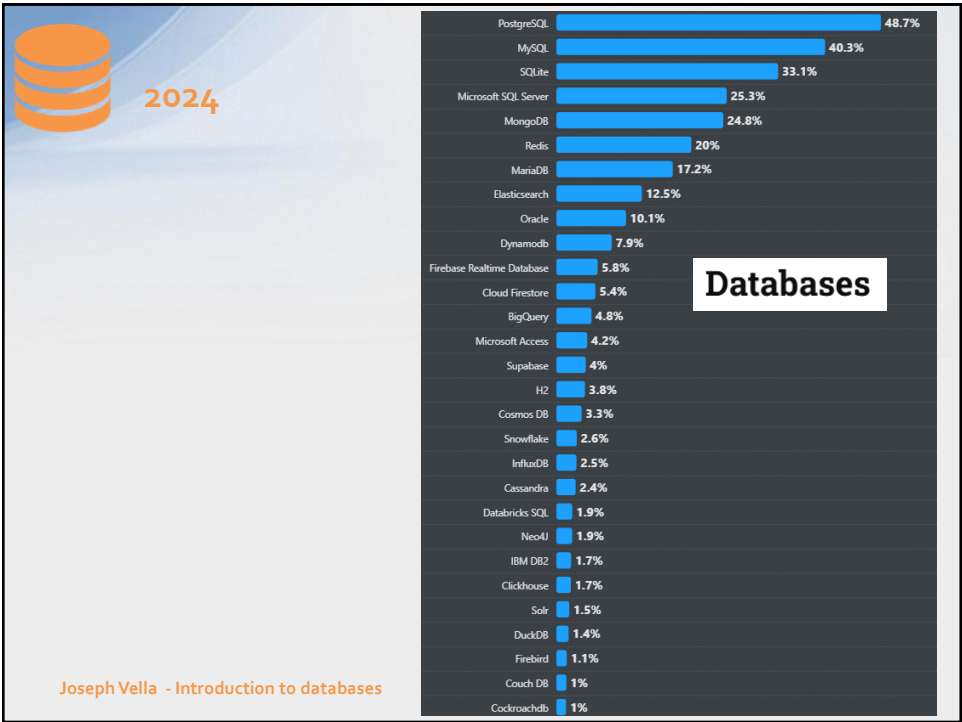
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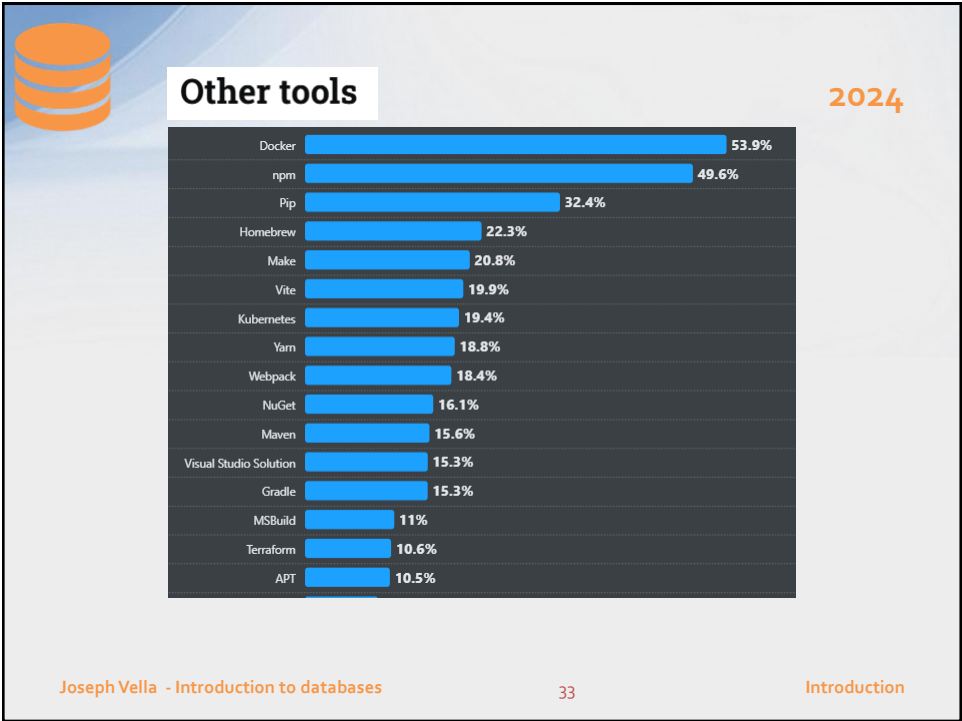
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Old copies

Just to keep the older copies

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Introducti3g

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Available DBMS (a biased collective list)

Stackoverflow 2022 (June)

PostgreSQL	46.48%
MySQL	45.68%
SQLite	30.83%
Microsoft SQL Server	28.77%
MongoDB	28.29%

JavaScript	67.9%
HTML/CSS	54.93%
SQL	52.64%
Python	43.51%
TypeScript	40.08%

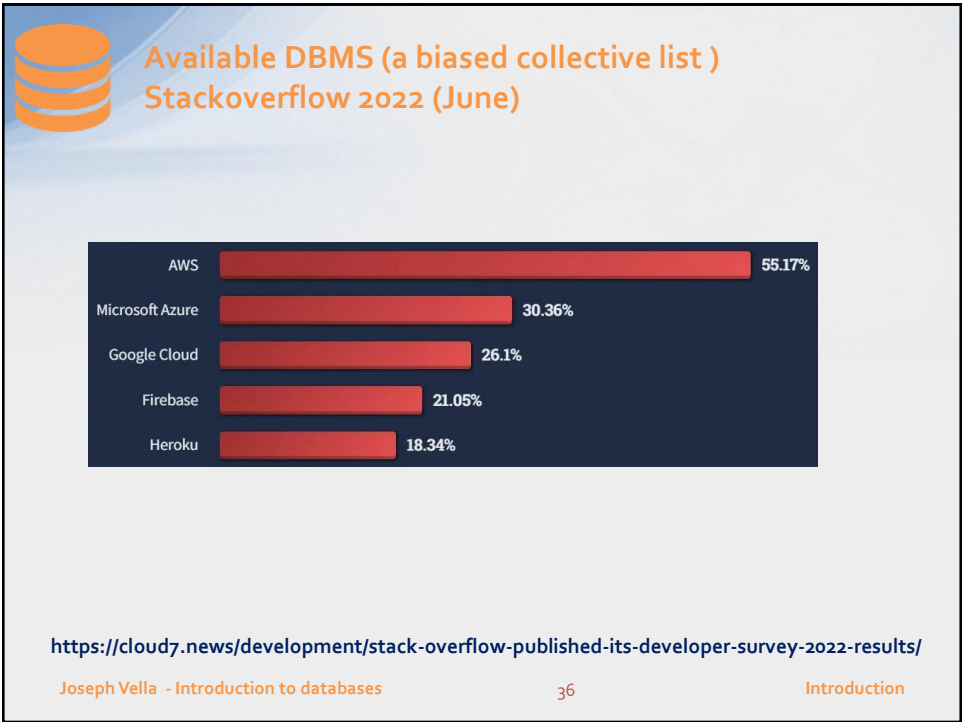
<https://cloud7.news/development/stack-overflow-published-its-developer-survey-2022-results/>

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 A presentation slide with a light blue background and a white horizontal bar at the bottom. In the top left corner, there is an orange icon of a database cylinder. The title 'Keys ideas (for this unit)' is written in orange. Below the title is a bulleted list in blue and red. The bottom of the slide has a white bar with three small pieces of text: 'Joseph Vella - Introduction to databases', '39', and 'Introduction'.

Keys ideas (for this unit)

- **Data Modelling**
 - Build a database design to requirements (including data requirements).
- **Schema and data management**
- **Structured Query Language (SQL):**
 - Data definition language (DDL)
 - Data manipulation (DML)
 - *Query modelling as in read only operations;*
 - *Transaction modelling as in read & write operations.*
 - Rule definitions
 - *Integrity constraints*
 - *(and database triggers)*
- **Elements of good design (e.g., data normalisation)**
- **Latching databases to programming exercises:**
 - Intra: stored procedures;
 - Inter: latching to programming languages.

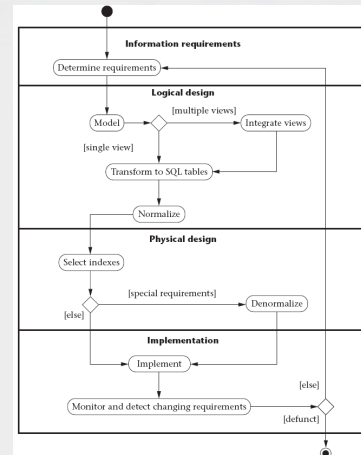
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Database life cycle (i)

- Get the requirements and extract database related information ...
 - Do not shy from certain topics! e.g., transaction processing, platforms, data security
- It is best to start with the conceptual schema and continue with adding and integrating external views as required.
 - These have to be supported from the requirements.
 - **Note:** external views are indispensable for application development (e.g. forms and reports) and data security basis.
- When global and external views are available one translates them into a DBMS specific data model constructs.



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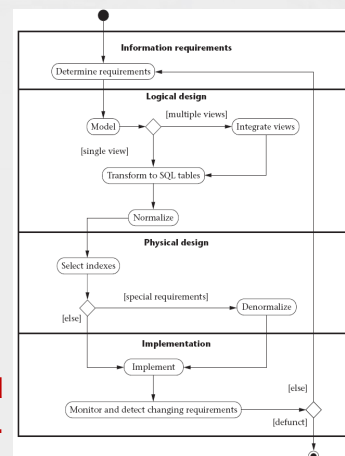
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Database life cycle (ii)

- Once DBMS specific artifacts are introduced one can start to introduce physical objects to address *performance*, *reliability* and *connectivity* issues.
 - Again it's best to reconcile with requirements.
 - Physical artifacts include:
 - Variety of indexes;
 - Data placement policies;
 - DBMS set-up parameters.
- Once the database is active, DBMS tools are used for monitoring performance and comparing this with expected guidelines.




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Database systems: areas of interest for IT people (i)

- **Design of database schemas**
 - How to build a useful database?
 - What descriptions and relationships need to be depicted in a schema?
 - How can we integrate with other established databases?
 - How to capture and quality check external data?

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Database systems: areas of interest for IT people (ii)

- **Database programming**
 - How to access and control database 'connections'?
 - How to express pattern matching for data retrieval (query specifications)?
 - How to express state change instructions over the database (transaction specification) and control their execution?
 - How best to live with DBMS facilities (data persistence, integrity constraints, transaction processing)?
 - How to align database programming with front-end programming (e.g. GUI at client presentation)?

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Database systems: areas of interest for IT people (iii)

- **DBMS implementation**
 - What's available (e.g. indexing methods)?
 - What's configurable?
 - e.g. can we overrule default B-tree with a hash index?
 - Are there any indications on performance issues?
 - e.g. in 95% of queries submitted their response time must be within 3 s.
 - e.g. 95% of submitted transaction have to commit.
 - What are the tradeoffs between various components (e.g. storage, authorization, transaction and query processing)?


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Guiding Principles

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Database systems: guiding principles (i)

ServiceCard

Serial Numb.	Service Type	Client Number	Client Name	Hours	Hrs to Charge
110	SERVICE	25010	Jeffrey	3	0
120	SERVICE	25150	Jeremy	2	2
130	REPAIR	25060	JJ	2	2
140	INSTALL	25090	Jaz	3	1
150	MA	25090	Jaz	2	0

- Consider the following data extracts.

ServiceCard


Serial Numb.	Service Type	Client Number	Hours	Hrs to Charge
110	SERVICE	25010	3	0
120	SERVICE	25150	2	2
130	REPAIR	25060	2	2
140	INSTALL	25090	3	1
150	MA	25090	2	0

ClientDetails

Client Numb.	Client Name
25010	Jeffrey
25150	Jeremy
25060	JJ
25090	Jaz

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Database systems: guiding principles (i)

- Consider the previous slide with data extracts:
 - Not in 3nf
 - And
 - Fixing it by decomposing into two tables.
- Data redundancy (i.e. at Conceptual level)**
 - Is the need to ensure, to a high degree, that data in a database is not duplicated.
 - If duplication (i.e. redundancy) exists then inconsistency arising from it has to be addressed (e.g. extra coding, obscures data meaning).
- We want to minimize data redundancy at logical level.**

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Data Independence: Motivation & Definitions

- The **separation** of physical data items (e.g. files), logical data items (relational database), and application programs that interact with the database minimises their inter dependence.
 - The concept of data independence conveys our wish to make artefact changes without unduly having to *re-map* (or *recompile*) all the higher artefacts that depend on the artefact.
 - Lowest is the physical level , middle is the logical and higher is the AP.
- **Logical data independence**
 - changes to the schema do not necessitate changes in the application programs. For example: adding new data types - programmed access should still be consistent.
- **Physical data independence**
 - changes to the physical schema are insulated from conceptual and application programs that use the database. Typical physical changes include: index re-structuring; records re-ordering in a data file.

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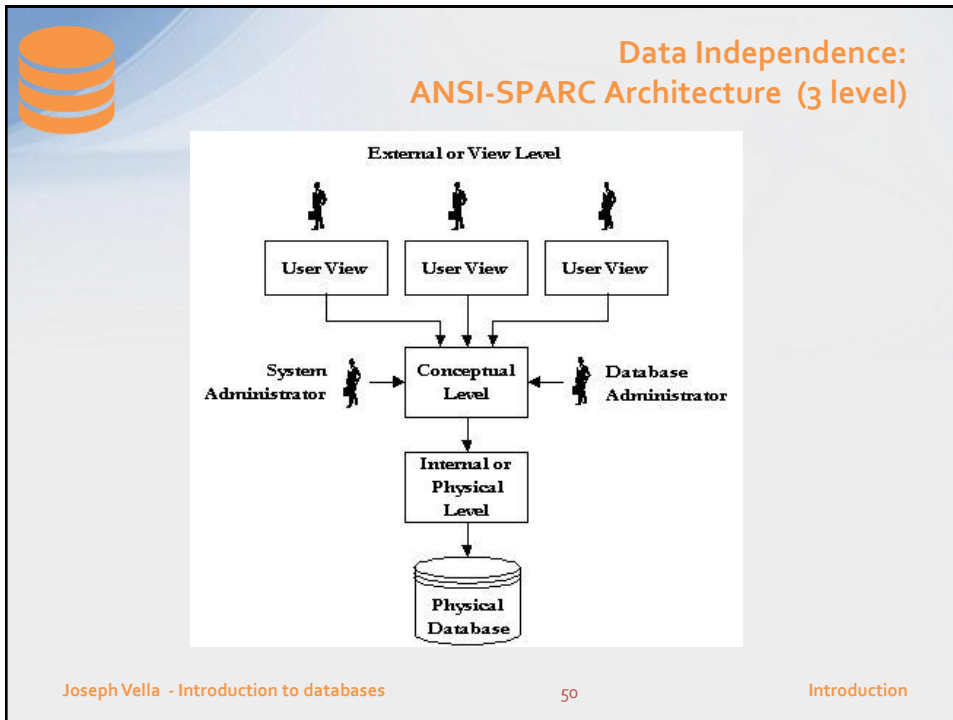
Data Independence: ANSI-SPARC Architecture (3 level)

The classic representation of a database and DBMS data architecture is the **ANSI/X3/SPARC** proposal (1978). This architecture comprises three levels of abstraction.

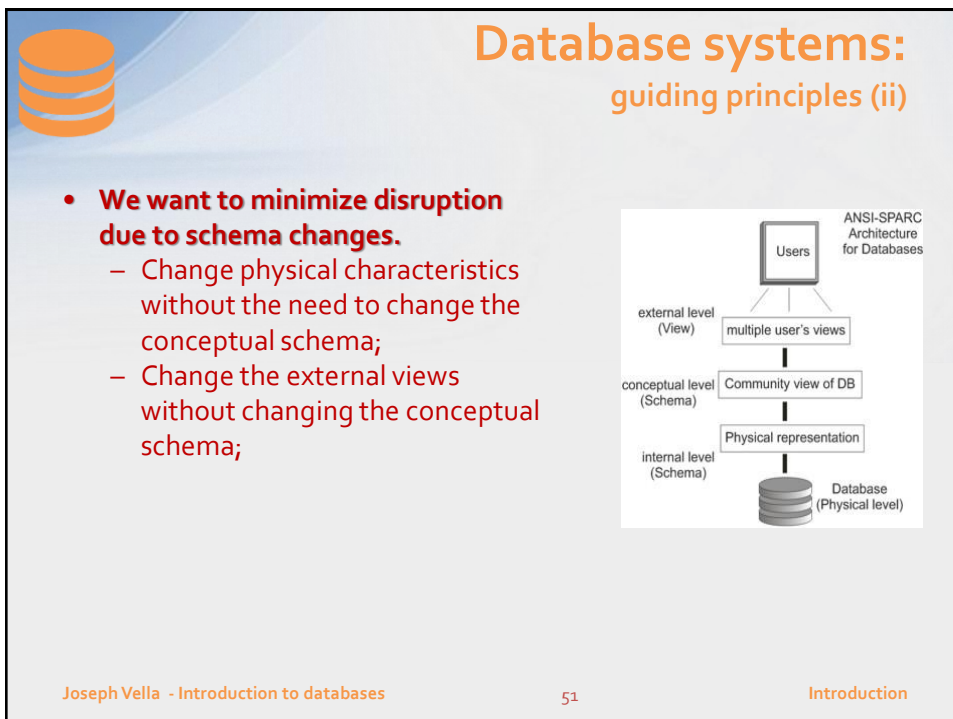
- **Internal Schema**
 - This is the physical storage model containing information such as devices, file locations, structures, indexing systems and access methods.
- **Conceptual Schema**
 - A single and consistent model of the database of interest.
- **External Schemas**
 - A view of the conceptual model applicable users or processes.
 - There are usually many of these and view overlapping is allowed.

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Database systems: guiding principles (iii)

- **Data security**
 - We said earlier that Data is indispensable:
 - **Protect from 'abuse' – e.g. security breaches, errors, contention, h/w failures (e.g. hard disks fail).**

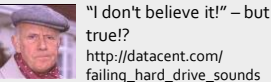




Write order from CPU for data "ABCD"

Parity data creation

A~D~Parity ABCD
E~H~Parity EFGH
I~L~Parity IJKL
M~P~Parity MNOP



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